ESP32-S2

ESP-IDF 编程指南

Release v4.2.5

乐鑫信息科技

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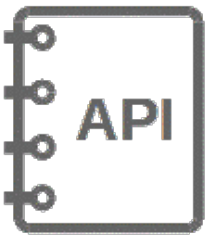
这里是乐鑫 IoT 开发框架 (esp-idf) 的文档中心。ESP-IDF 是 ESP32 和 ESP32-S2 系列芯片的官方开发框架。

本文档仅包含针对 ESP32-S2 芯片的 ESP-IDF 使用。

ESP-IDF 是 ESP32 芯片的官方开发框架。

		
快速入门	API 参考	H/W 参考
		
API 指南	贡献代码	相关资源

Chapter 1

快速入门

本文档旨在指导用户搭建 ESP32-S2 硬件开发的软件环境，通过一个简单的示例展示如何使用 ESP-IDF (Espressif IoT Development Framework) 配置菜单，并编译、下载固件至 ESP32-S2 开发板等步骤。

注解： 这是 ESP-IDF 稳定版本 v4.2.5 的文档，还有其他版本的文档*ESP-IDF 版本简介* 供参考。

1.1 概述

ESP32-S2 SoC 芯片支持以下功能：

- 2.4 GHz Wi-Fi
- 高性能单核
- 运行 RISC-V 或 FSM 内核的超低功耗协处理器
- 多种外设
- 内置安全硬件
- USB OTG 接口

ESP32-S2 采用 40 nm 工艺制成，具有最佳的功耗性能、射频性能、稳定性、通用性和可靠性，适用于各种应用场景和不同功耗需求。

乐鑫为用户提供完整的软、硬件资源，进行 ESP32-S2 硬件设备的开发。其中，乐鑫的软件开发环境 ESP-IDF 旨在协助用户快速开发物联网 (IoT) 应用，可满足用户对 Wi-Fi、蓝牙、低功耗等方面的要求。

1.2 准备工作

硬件：

- 一款 **ESP32-S2** 开发板
- **USB 数据线**（A 转 Micro-B）
- 电脑（Windows、Linux 或 mac OS）

软件：

您可以选择下载并手动安装以下软件：

- 设置 **工具链**，用于编译 ESP32-S2 代码；
- **编译构建工具**——CMake 和 Ninja 编译构建工具，用于编译 ESP32-S2 **应用程序**；
- 获取 **ESP-IDF** 软件开发框架。该框架已经基本包含 ESP32-S2 使用的 API（软件库和源代码）和运行 **工具链**的脚本；

或者，您也可以通过以下集成开发环境 (IDE) 中的官方插件完成安装流程：

- Eclipse 插件 (安装)

- VS Code 插件 (安装)

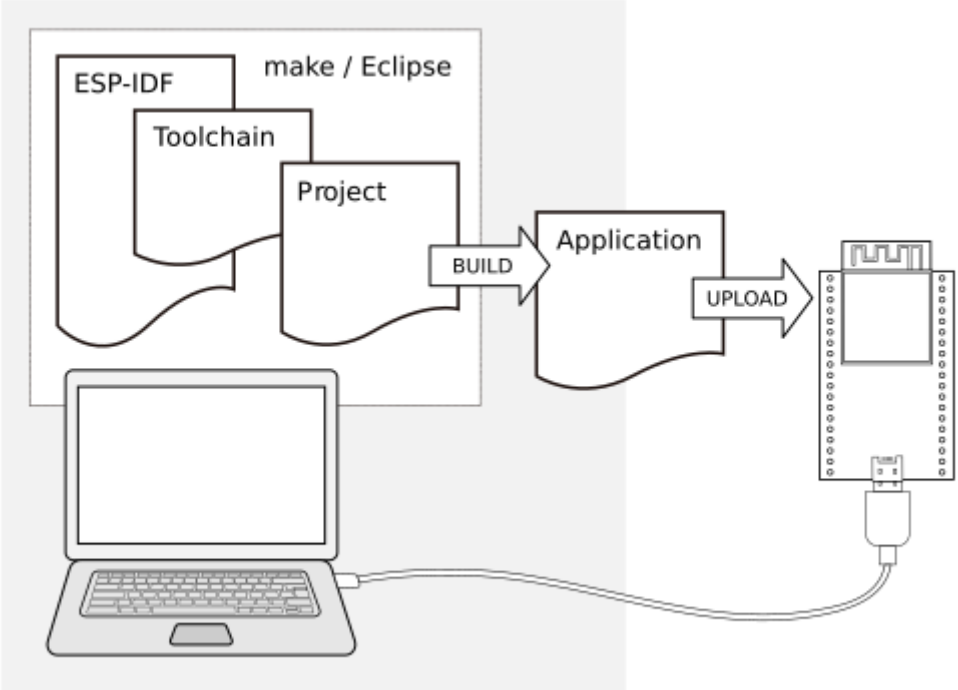


图 1: ESP32-S2 应用程序开发

1.3 开发板简介

请点击下方连接，了解有关具体开发板的详细信息。

1.3.1 ESP32-S2-Saola-1

本指南介绍了乐鑫一款基于 ESP32-S2 的小型开发板 ESP32-S2-Saola-1。

本指南包括如下内容：

- 入门指南：简要介绍了 ESP32-S2-Saola-1 和硬件、软件设置指南。
- 硬件参考：详细介绍了 ESP32-S2-Saola-1 的硬件。
- 硬件版本：介绍硬件历史版本和已知问题，并提供链接至历史版本开发板的入门指南（如有）。
- 相关文档：列出了相关文档的链接。

入门指南

本节介绍了如何快速上手 ESP32-S2-Saola-1。开头部分介绍了 ESP32-S2-Saola-1，开始开发应用 小节介绍了怎样在 ESP32-S2-Saola-1 上安装模组、设置及烧录固件。

概述 ESP32-S2-Saola-1 是乐鑫一款基于 ESP32-S2 的小型开发板。板上的绝大部分管脚均已引出，开发人员可根据实际需求，轻松通过跳线连接多种外围器件，或将开发板插在面包板上使用。

为了更好地满足不同用户需求，ESP32-S2-Saola-1 支持以下模组：

- ESP32-S2-WROVER
- ESP32-S2-WROVER-I

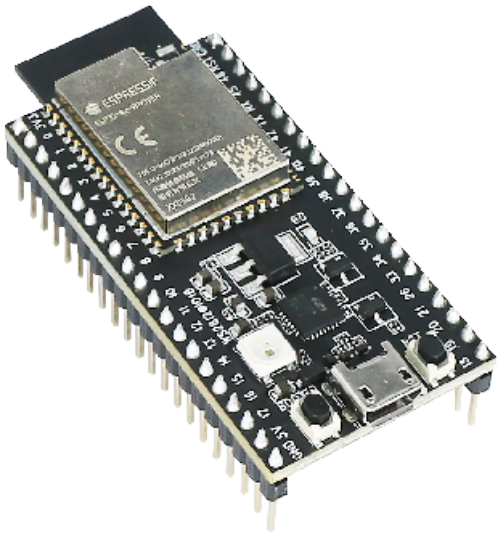


图 2: ESP32-S2-Saola-1

- ESP32-S2-WROOM
- ESP32-S2-WROOM-I

本指南以搭载 ESP32-S2-WROVER 模组的 ESP32-S2-Saola-1 为例。

内含组件和包装

零售订单 如购买样品，每个 ESP32-S2-Saola-1 底板将以防静电袋或零售商选择的其他方式包装。

零售订单请前往 https://www.espressif.com/zh-hans/company/contact/buy-a-sample。

批量订单 如批量购买，ESP32-S2-Saola-1 烧录底板将以大纸板箱包装。

批量订单请参考 乐鑫产品订购信息 (PDF)。

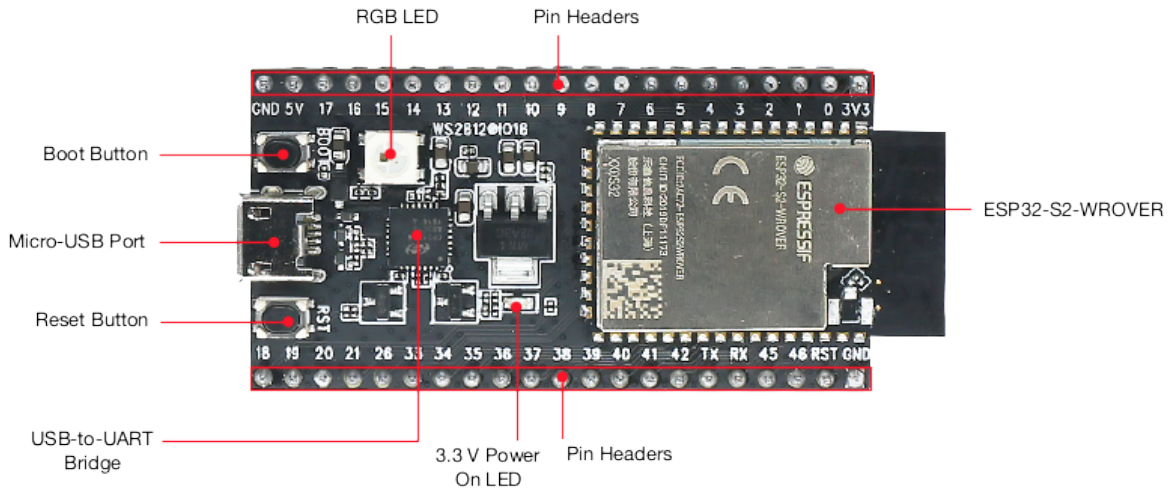


图 3: ESP32-S2-Saola-1 - 正面

组件介绍 以下按照顺时针的顺序依次介绍开发板上的主要组件。

主要组件	介绍
ESP32-S2-WROVER	ESP32-S2-WROVER 集成 ESP32-S2，是通用型 Wi-Fi MCU 模组，功能强大。该模组采用 PCB 板载天线，配置了 4 MB SPI flash 和 2 MB SPI PSRAM。
Pin Headers（排针）	所有可用 GPIO 管脚（除 Flash 和 PSRAM 的 SPI 总线）均已引出至开发板的排针。用户可对 ESP32-S2 芯片编程，使能 SPI、I2S、UART、I2C、触摸传感器、PWM 等多种功能。
3.3 V Power On LED（3.3 V 电源指示灯）	开发板连接 USB 电源后，该指示灯亮起。
USB-to-UART Bridge（USB 转 UART 桥接器）	单芯片 USB 至 UART 桥接器，可提供高达 3 Mbps 的传输速率。
Reset Button（Reset 键）	复位按键。
Micro-USB Port（Micro-USB 接口）	USB 接口。可用作开发板的供电电源或 PC 和 ESP32-S2 芯片的通信接口。
Boot Button（Boot 键）	下载按键。按住 **Boot** 键的同时按一下 **Reset** 键进入“固件下载”模式，通过串口下载固件。
RGB LED	可寻址 RGB 发光二极管 (WS2812)，由 GPIO18 驱动。

开始开发应用 通电前，请确保 ESP32-S2-Saola-1 完好无损。

必备硬件

- ESP32-S2-Saola-1
- USB 2.0 数据线（标准 A 型转 Micro-B 型）
- 电脑（Windows、Linux 或 macOS）

注解： 请确保使用适当的 USB 数据线。部分数据线仅可用于充电，无法用于数据传输和编程。

软件设置 请前往*快速入门*，在*详细安装步骤* 一节查看如何快速设置开发环境，将应用程序烧录至 ESP32-S2-Saola-1。

注解： ESP32-S2 系列芯片仅支持 ESP-IDF master 分支或 v4.2 以上版本。

硬件参考

功能框图 ESP32-S2-Saola-1 的主要组件和连接方式如下图所示。

电源选项 您可从以下三种供电方式中任选其一给 ESP32-S2-Saola-1 供电：

- Micro-USB 接口供电（默认）
- 5V 和 GND 排针供电
- 3V3 和 GND 排针供电

建议选择第一种供电方式：Micro-USB 接口供电。

排针 下表列出了开发板两侧排针（J2 和 J3）的 **名称**和 **功能**，排针的名称如图*ESP32-S2-Saola-1 - 正面* 所示，排针的序号与 ESP32-S2-Saola-1 原理图 (PDF) 一致。

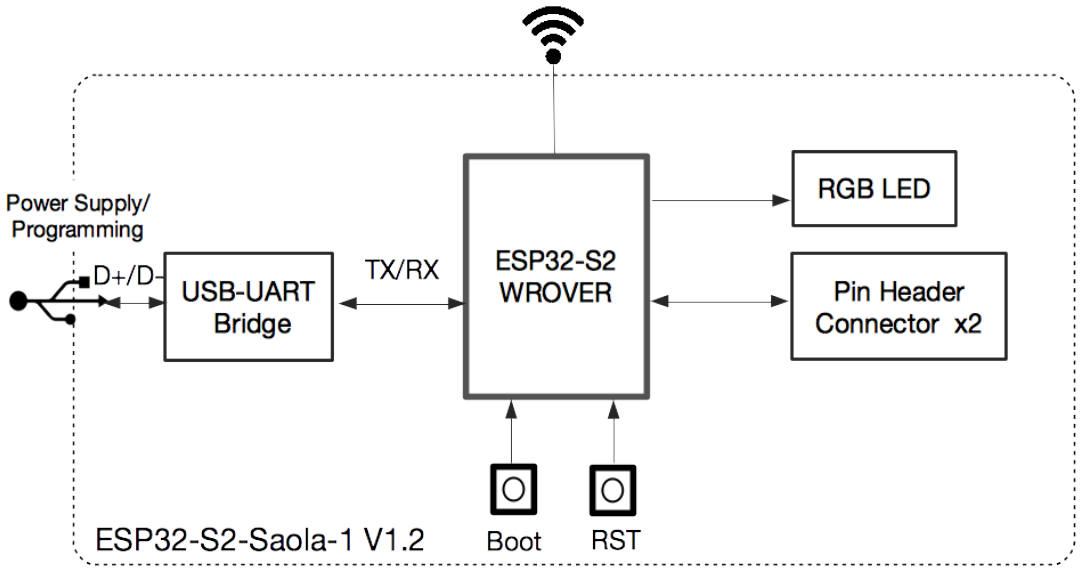


图 4: ESP32-S2-Saola-1 （点击放大）

J2

序号	名称	类型[1]	功能
1	3V3	P	3.3 V 电源
2	IO0	I/O	GPIO0, 启动
3	IO1	I/O	GPIO1, ADC1_CH0, TOUCH_CH1
4	IO2	I/O	GPIO2, ADC1_CH1, TOUCH_CH2
5	IO3	I/O	GPIO3, ADC1_CH2, TOUCH_CH3
6	IO4	I/O	GPIO4, ADC1_CH3, TOUCH_CH4
7	IO5	I/O	GPIO5, ADC1_CH4, TOUCH_CH5
8	IO6	I/O	GPIO6, ADC1_CH5, TOUCH_CH6
9	IO7	I/O	GPIO7, ADC1_CH6, TOUCH_CH7
10	IO8	I/O	GPIO8, ADC1_CH7, TOUCH_CH8
11	IO9	I/O	GPIO9, ADC1_CH8, TOUCH_CH9
12	IO10	I/O	GPIO10, ADC1_CH9, TOUCH_CH10
13	IO11	I/O	GPIO11, ADC2_CH0, TOUCH_CH11
14	IO12	I/O	GPIO12, ADC2_CH1, TOUCH_CH12
15	IO13	I/O	GPIO13, ADC2_CH2, TOUCH_CH13
16	IO14	I/O	GPIO14, ADC2_CH3, TOUCH_CH14
17	IO15	I/O	GPIO15, ADC2_CH4, XTAL_32K_P
18	IO16	I/O	GPIO16, ADC2_CH5, XTAL_32K_N
19	IO17	I/O	GPIO17, ADC2_CH6, DAC_1
20	5V0	P	5 V 电源
21	GND	G	接地

[1] P：电源；I：输入；O：输出；T：可设置为高阻。

J3

序号	名称	类型	功能
1	GND	G	接地
2	RST	I	CHIP_PU, 复位
3	IO46	I	GPIO46
4	IO45	I/O	GPIO45
5	IO44	I/O	GPIO44, U0RXD
6	IO43	I/O	GPIO43, U0TXD
7	IO42	I/O	GPIO42, MTMS
8	IO41	I/O	GPIO41, MTDI
9	IO40	I/O	GPIO40, MTDO
10	IO39	I/O	GPIO39, MTCK
11	IO38	I/O	GPIO38
12	IO37	I/O	GPIO37
13	IO36	I/O	GPIO36
14	IO35	I/O	GPIO35
16	IO34	I/O	GPIO34
17	IO33	I/O	GPIO33
17	IO26	I/O	GPIO26
18	IO21	I/O	GPIO21
19	IO20	I/O	GPIO20, ADC2_CH9, USB_D+
20	IO19	I/O	GPIO19, ADC2_CH8, USB_D-
21	IO18	I/O	GPIO18, ADC2_CH7, DAC_2, RGB LED

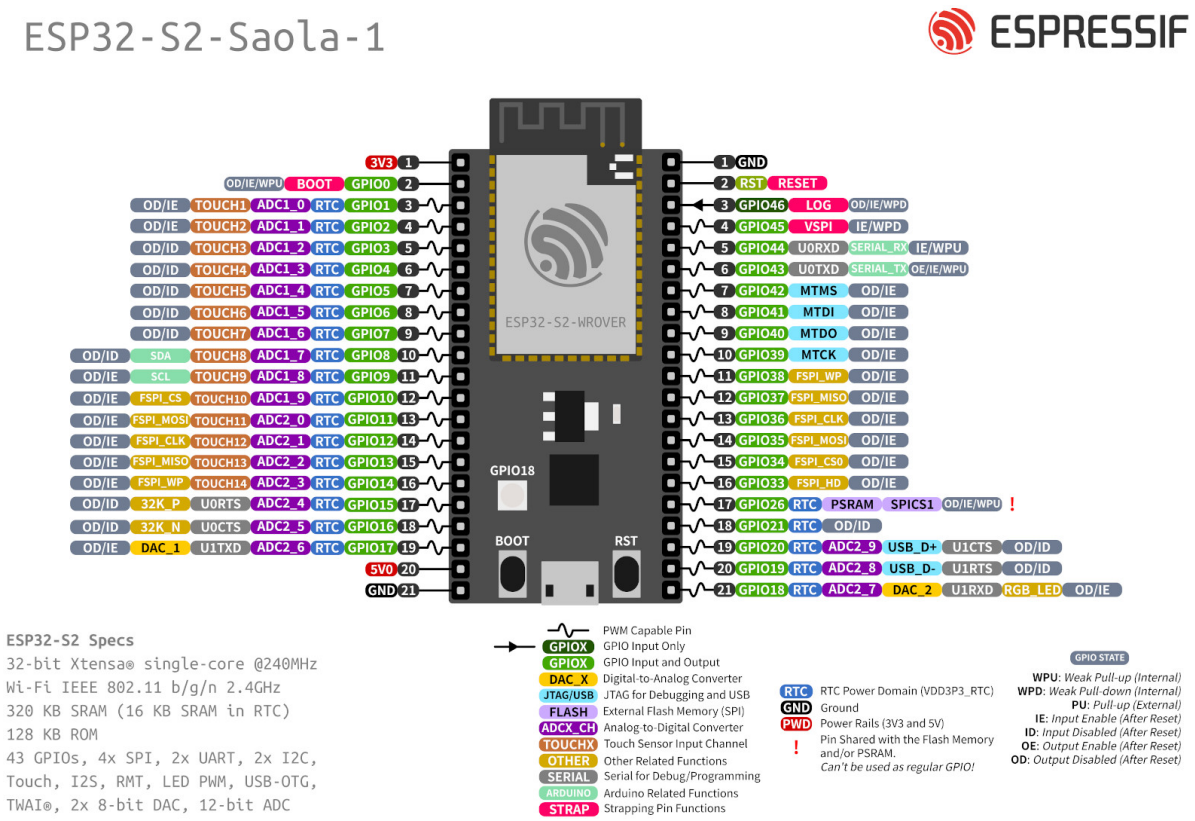


图 5: ESP32-S2-Saola-1 管脚布局（点击放大）

管脚布局

硬件版本

无历史版本。

相关文档

- ESP32-S2-Saola-1 原理图 (PDF)
- ESP32-S2-Saola-1 尺寸图 (PDF)
- ESP32-S2 技术规格书 (PDF)
- ESP32-S2-WROVER & ESP32-S2-WROVER-I 技术规格书 (PDF)
- ESP32-S2-WROOM & ESP32-S2-WROOM-I 技术规格书 (PDF)
- 乐鑫产品选型工具

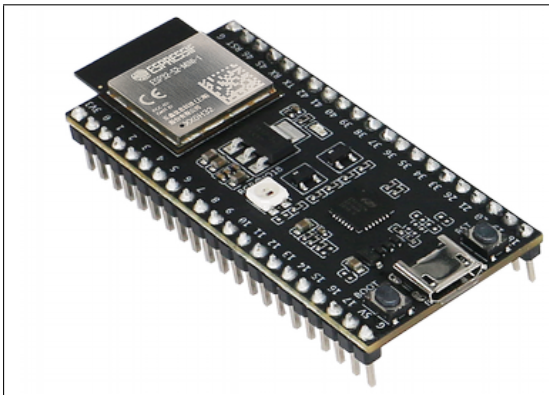
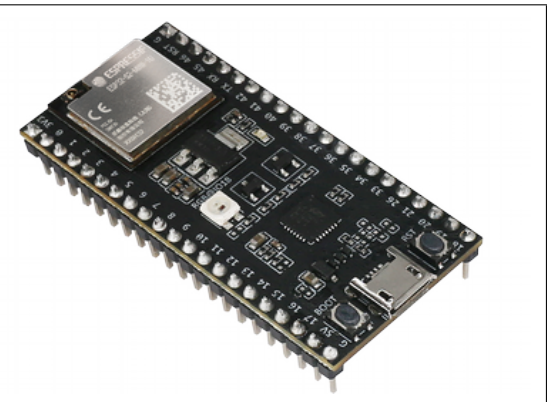
有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

1.3.2 ESP32-S2-DevKitM-1(U)

本指南介绍了乐鑫的小型开发板 ESP32-S2-DevKitM-1(U)。

ESP32-S2-DevKitM-1(U) 是一款基于 ESP32-S2FH4 芯片（ESP32-S2 系列）的通用型开发板。该款开发板具有丰富的外设和优化的引脚布局，令产品开发更快捷。

ESP32-S2-DevKitM-1 搭载的是 ESP32-S2-MINI-1 模组（PCB 板载天线），ESP32-S2-DevKitM-1U 搭载的是 ESP32-S2-MINI-1U 模组（外部天线连接器）。

	
ESP32-S2-DevKitM-1	ESP32-S2-DevKitM-1U

本指南包括如下内容：

- 入门指南: 简要介绍了 ESP32-S2-DevKitM-1(U) 和硬件、软件设置指南。
- 硬件参考: 详细介绍了 ESP32-S2-DevKitM-1(U) 的硬件。
- 硬件版本：介绍硬件历史版本和已知问题，并提供链接至历史版本开发板的入门指南（如有）。
- 相关文档: 列出了相关文档的链接。

入门指南

本节介绍了如何快速上手 ESP32-S2-DevKitM-1(U)。开头部分介绍了 ESP32-S2-DevKitM-1(U)，开始开发应用 小节介绍了怎样在 ESP32-S2-DevKitM-1(U) 上烧录固件及相关准备工作。

概述 ESP32-S2-DevKitM-1(U) 是乐鑫一款搭载 ESP32-S2-MINI-1 或 ESP32-S2-MINI-1U 模组的入门级开发板。板上模组大部分管脚均已引出至两侧排针，开发人员可根据实际需求，轻松通过跳线连接多种外围设备，同时也可将开发板插在面包板上使用。

内含组件和包装

零售订单 如购买样品，每个 ESP32-S2-DevKitM-1(U) 底板将以防静电袋或零售商选择的其他方式包装。零售订单请前往 https://www.espressif.com/zh-hans/company/contact/buy-a-sample。

批量订单 如批量购买，ESP32-S2-DevKitM-1(U) 烧录底板将以大纸板箱包装。批量订单请参考 乐鑫产品订购信息 (PDF)。

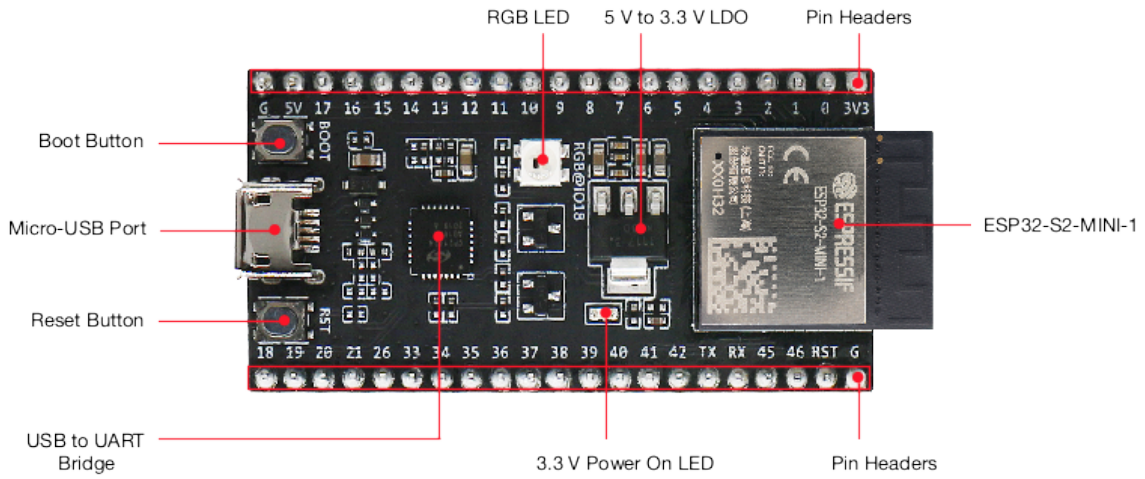


图 6: ESP32-S2-DevKitM-1 - 正面

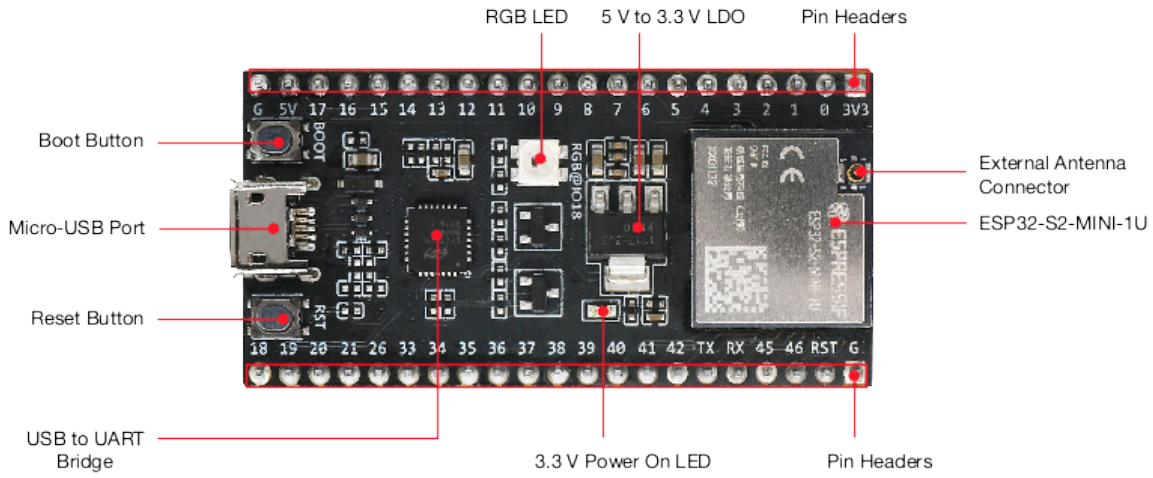


图 7: ESP32-S2-DevKitM-1U - 正面

组件介绍 以下按照顺时针的顺序依次介绍开发板上的主要组件。

主要组件	介绍
ESP32-S2-MINI-1 或 ESP32-S2-MINI-1U	ESP32-S2-MINI-1 和 ESP32-S2-MINI-1U 是集成 ESP32-S2FH4 的通用型 Wi-Fi MCU 模组，ESP32-S2-MINI-1 采用 PCB 板载天线，ESP32-S2-MINI-1U 采用外部天线连接器。两款模组均配置了 4 MB SPI flash。
Pin Headers（排针）	所有可用 GPIO 管脚（除 flash 的 SPI 总线）均已引出至开发板的排针。用户可对 ESP32-S2FH4 芯片编程，使能 SPI、I2S、UART、I2C、触摸传感器、PWM 等多种功能。请查看排针 获取更多信息。
3.3 V Power On LED（3.3 V 电源指示灯）	开发板连接 USB 电源后，该指示灯亮起。
USB-to-UART Bridge（USB 转 UART 桥接器）	单芯片 USB 至 UART 桥接器，可提供高达 3 Mbps 的传输速率。
Reset Button（Reset 键）	复位按键。
Micro-USB（Micro-USB 接口）	USB 接口。可用作开发板的供电电源或 PC 和 ESP32-S2FH4 芯片的通信接口。
Boot Button（Boot 键）	下载按键。按住 **Boot** 键的同时按一下 **Reset** 键进入“固件下载”模式，通过串口下载固件。
RGB LED	可寻址 RGB 发光二极管，由 GPIO18 驱动。
5 V to 3.3 V LDO（5 V 转 3.3 V LDO）	电源转换器，输入 5 V，输出 3.3 V。
External Antenna Connector（外部天线连接器）	仅 **ESP32-S2-MINI-1U** 模组带有外部天线连接器。连接器尺寸，请参考《ESP32-S2-MINI-1 & ESP32-S2-MINI-1U 技术规格书》的外部天线连接器尺寸章节。

开始开发应用 通电前，请确保 ESP32-S2-DevKitM-1(U) 完好无损。

必备硬件

- ESP32-S2-DevKitM-1(U)
  - 如使用 ESP32-S2-DevKitM-1U，还需准备天线
- USB 2.0 数据线（标准 A 型转 Micro-B 型）
- 电脑（Windows、Linux 或 macOS）

注解： 请确保使用适当的 USB 数据线。部分数据线仅可用于充电，无法用于数据传输和编程。

软件设置 请前往快速入门，在详细安装步骤 一节查看如何快速设置开发环境，将应用程序烧录至 ESP32-S2-DevKitM-1(U)。

注解： ESP32-S2 系列芯片仅支持 ESP-IDF master 分支或 v4.2 以上版本。

硬件参考

功能框图 ESP32-S2-DevKitM-1(U) 的主要组件和连接方式如下图所示。

电源选项 您可从以下三种供电方式中任选其一给 ESP32-S2-DevKitM-1(U) 供电：

- Micro-USB 接口供电（默认）
- 5V 和 GND 排针供电
- 3V3 和 GND 排针供电

建议选择第一种供电方式：Micro-USB 接口供电。

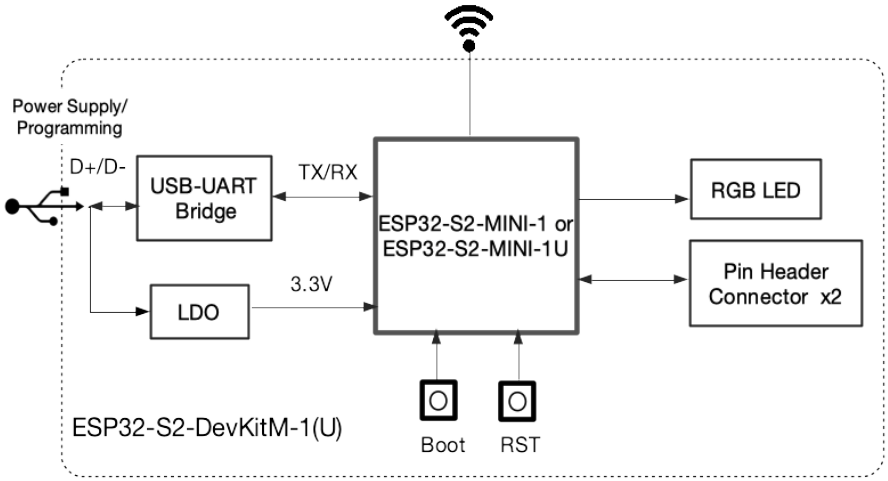


图 8: ESP32-S2-DevKitM-1(U) （点击放大）

排针 下表列出了开发板两侧排针（J1 和 J3）的 **名称**和 **功能**，排针的名称如图*ESP32-S2-DevKitM-1 - 正面* 所示，排针的序号与 ESP32-S2-DevKitM-1(U) 原理图 (PDF) 一致。

J1

序号	名称	类型[1]	功能
1	3V3	P	3.3 V 电源
2	0	I/O/T	RTC_GPIO0, GPIO0
3	1	I/O/T	RTC_GPIO1, GPIO1, TOUCH1, ADC1_CH0
4	2	I/O/T	RTC_GPIO2, GPIO2, TOUCH2, ADC1_CH1
5	3	I/O/T	RTC_GPIO3, GPIO3, TOUCH3, ADC1_CH2
6	4	I/O/T	RTC_GPIO4, GPIO4, TOUCH4, ADC1_CH3
7	5	I/O/T	RTC_GPIO5, GPIO5, TOUCH5, ADC1_CH4
8	6	I/O/T	RTC_GPIO6, GPIO6, TOUCH6, ADC1_CH5
9	7	I/O/T	RTC_GPIO7, GPIO7, TOUCH7, ADC1_CH6
10	8	I/O/T	RTC_GPIO8, GPIO8, TOUCH8, ADC1_CH7
11	9	I/O/T	RTC_GPIO9, GPIO9, TOUCH9, ADC1_CH8, FSPIHD
12	10	I/O/T	RTC_GPIO10, GPIO10, TOUCH10, ADC1_CH9, FSPICS0, FSPIIO4
13	11	I/O/T	RTC_GPIO11, GPIO11, TOUCH11, ADC2_CH0, FSPID, FSPIIO5
14	12	I/O/T	RTC_GPIO12, GPIO12, TOUCH12, ADC2_CH1, FSPICLK, FSPIIO6
15	13	I/O/T	RTC_GPIO13, GPIO13, TOUCH13, ADC2_CH2, FSPIQ, FSPIIO7
16	14	I/O/T	RTC_GPIO14, GPIO14, TOUCH14, ADC2_CH3, FSPIWP, FSPIDQS
17	15	I/O/T	RTC_GPIO15, GPIO15, U0RTS, ADC2_CH4, XTAL_32K_P
18	16	I/O/T	RTC_GPIO16, GPIO16, U0CTS, ADC2_CH5, XTAL_32K_N
19	17	I/O/T	RTC_GPIO17, GPIO17, U1TXD, ADC2_CH6, DAC_1
20	5V	P	5 V 电源
21	G	G	接地

P：电源；I：输入；O：输出；T：可设置为高阻。

J3

序号	名称	类型	功能
1	G	G	接地
2	RST	I	CHIP_PU
3	46	I	GPIO46
4	45	I/O/T	GPIO45
5	RX	I/O/T	U0RXD, GPIO44, CLK_OUT2
6	TX	I/O/T	U0TXD, GPIO43, CLK_OUT1
7	42	I/O/T	MTMS, GPIO42
8	41	I/O/T	MTDI, GPIO41, CLK_OUT1
9	40	I/O/T	MTDO, GPIO40, CLK_OUT2
10	39	I/O/T	MTCK, GPIO39, CLK_OUT3
11	38	I/O/T	GPIO38, FSPIWP
12	37	I/O/T	SPIDQS, GPIO37, FSPIQ
13	36	I/O/T	SPIIO7, GPIO36, FSPICLK
14	35	I/O/T	SPIIO6, GPIO35, FSPID
15	34	I/O/T	SPIIO5, GPIO34, FSPICS0
16	33	I/O/T	SPIIO4, GPIO33, FSPIHD
17	26	I/O/T	SPICS1, GPIO26
18	21	I/O/T	RTC_GPIO21, GPIO21
19	20	I/O/T	RTC_GPIO20, GPIO20, U1CTS, ADC2_CH9, CLK_OUT1, USB_D+
20	19	I/O/T	RTC_GPIO19, GPIO19, U1RTS, ADC2_CH8, CLK_OUT2, USB_D-
21	18	I/O/T	RTC_GPIO18, GPIO18, U1RXD, ADC2_CH7, DAC_2, CLK_OUT3, RGB LED

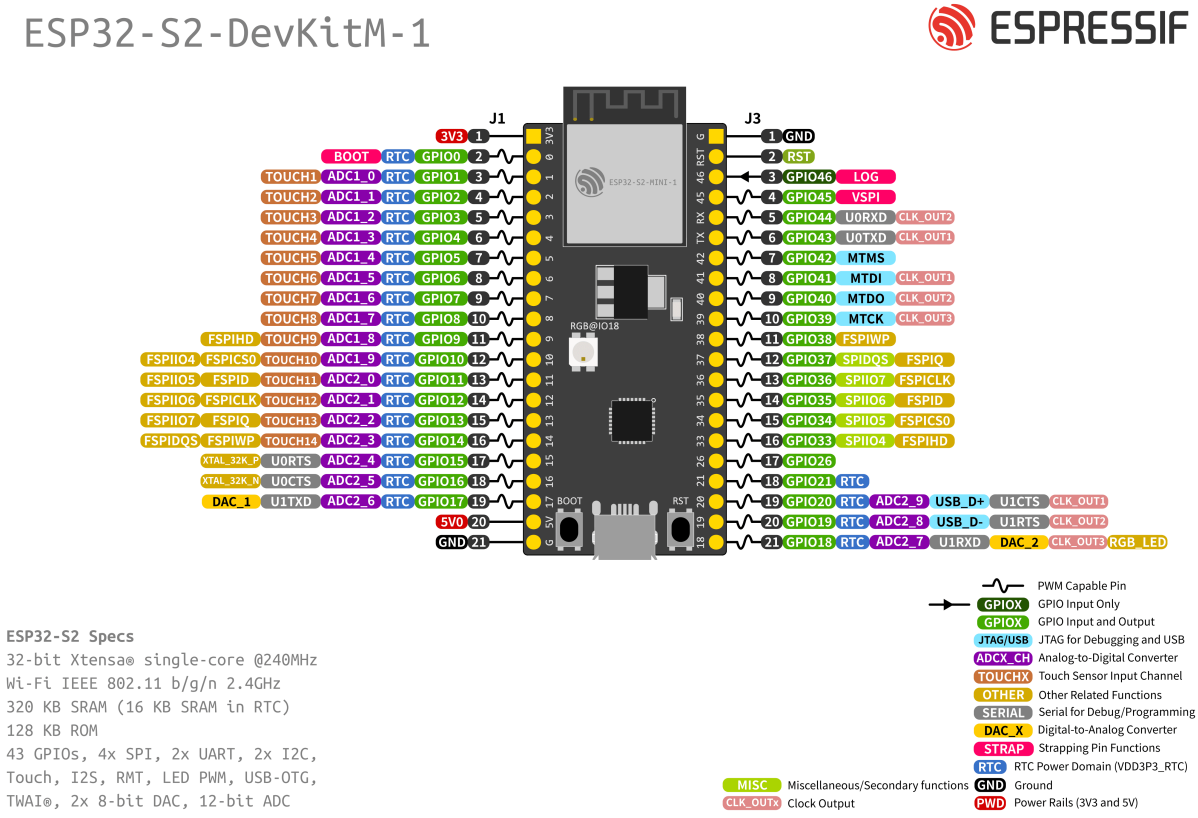


图 9: ESP32-S2-DevKitM-1(U) 管脚布局（点击放大）

管脚布局

硬件版本

无历史版本。

相关文档

- ESP32-S2-DevKitM-1(U) 原理图 (PDF)
- ESP32-S2-DevKitM-1(U) PCB 布局 (PDF)
- ESP32-S2-DevKitM-1(U) 尺寸图 (PDF)
- ESP32-S2 系列技术规格书 (PDF)
- ESP32-S2-MINI-1 & ESP32-S2-MINI-1U 技术规格书 (PDF)
- 乐鑫产品订购信息 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

1.3.3 ESP32-S2-DevKitC-1

本指南将帮助您快速上手 ESP32-S2-DevKitC-1，并提供该款开发板的详细信息。

ESP32-S2-DevKitC-1 是一款入门级开发板，使用带有 4 MB SPI flash 的 ESP32-S2-SOLO（板载 PCB 天线）或 ESP32-S2-SOLO-U（外部天线连接器）模组。该款开发板具备完整的 Wi-Fi 功能。

板上模组大部分管脚均已引出至两侧排针，开发人员可根据实际需求，轻松通过跳线连接多种外围设备，同时也可将开发板插在面包板上使用。

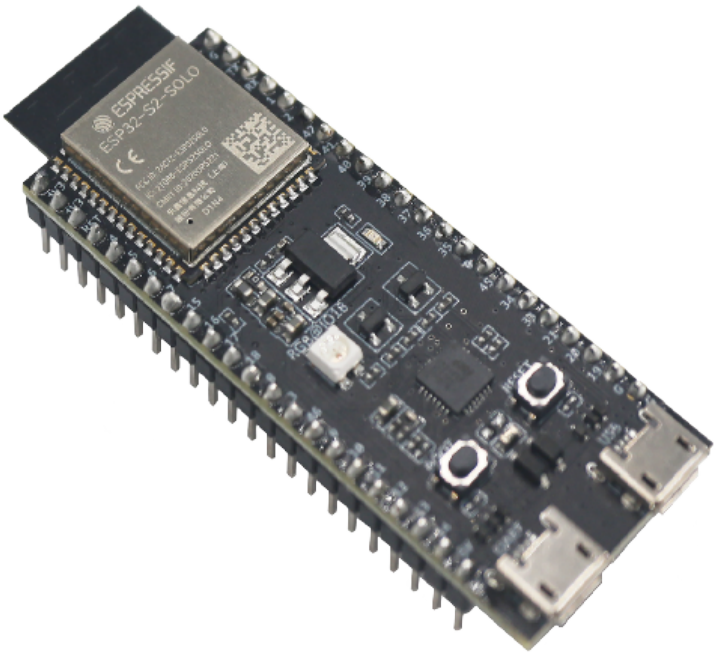


图 10: ESP32-S2-DevKitC-1（板载 ESP32-S2-SOLO 模组）

本指南包括如下内容：

- 入门指南：简要介绍了 ESP32-S2-DevKitC-1 和硬件、软件设置指南。
- 硬件参考：详细介绍了 ESP32-S2-DevKitC-1 的硬件。
- 硬件版本：介绍硬件历史版本和已知问题，并提供链接至历史版本开发板的入门指南（如有）。
- 相关文档：列出了相关文档的链接。

入门指南

本小节将简要介绍 ESP32-S2-DevKitC-1，说明如何在 ESP32-S2-DevKitC-1 上烧录固件及相关准备工作。

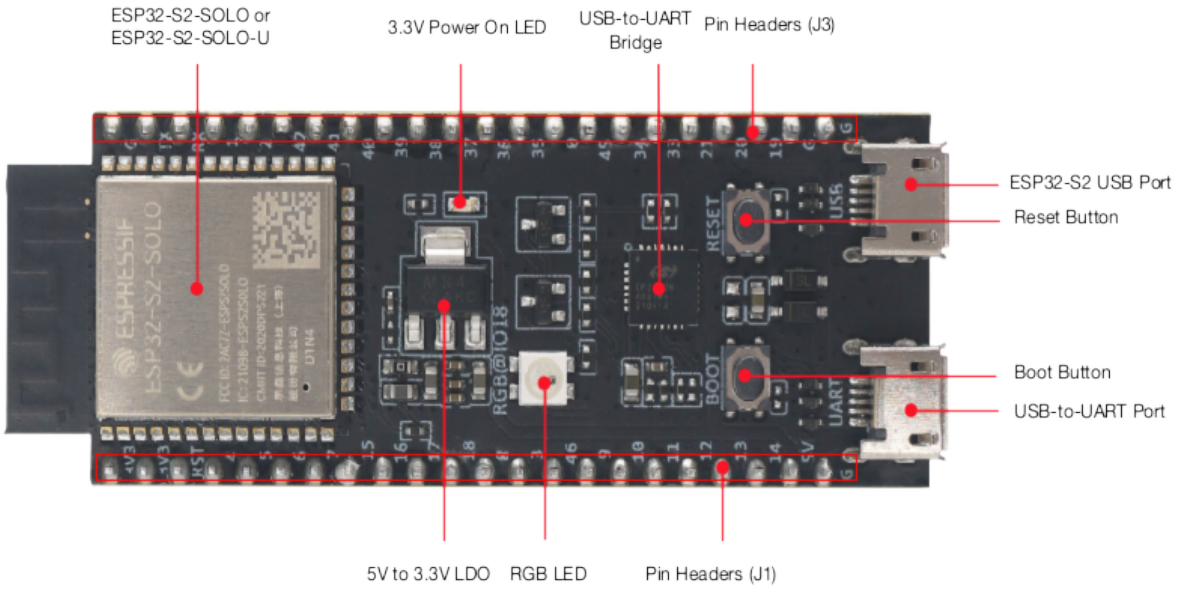


图 11: ESP32-S2-DevKitC-1 - 正面

组件介绍 以下按照顺时针的顺序依次介绍开发板上的主要组件。

主要组件	介绍
ESP32-S2-SOLO 或 ESP32-S2-SOLO-U	ESP32-S2-SOLO 和 ESP32-S2-SOLO-U 是两款通用型 Wi-Fi 模组。ESP32-S2-SOLO 采用 PCB 板载天线，ESP32-S2-SOLO-U 采用连接器连接外部天线。开发板上的 ESP32-S2-SOLO 或 ESP32-S2-SOLO-U 模组可配置 4 MB flash，也可配置 4 MB flash 加 2 MB PSRAM（芯片内置）。
3.3 V Power On LED（3.3 V 电源指示灯）	开发板连接 USB 电源后，该指示灯亮起。
USB-to-UART Bridge（USB 转 UART 桥接器）	单芯片 USB 转 UART 桥接器，可提供高达 3 Mbps 的传输速率。
Pin Headers（排针）	所有可用 GPIO 管脚（除 flash 的 SPI 总线）均已引出至开发板的排针。请查看排针 获取更多信息。
ESP32-S2 USB Port（ESP32-S2 USB 接口）	ESP32-S2 USB OTG 接口，支持全速 USB 1.1 标准。该接口可用作开发板的供电接口，可烧录固件至芯片，也可通过 USB 协议与芯片通信。
Reset Button（Reset 键）	复位按键。
Boot Button（Boot 键）	下载按键。按住 **Boot** 键的同时按一下 **Reset** 键进入“固件下载”模式，通过串口下载固件。
USB-to-UART Port（USB 转 UART 接口）	Micro-USB 接口，可用作开发板的供电接口，可烧录固件至芯片，也可作为通信接口，通过板载 USB 转 UART 桥接器与 ESP32-S2 芯片通信。
RGB LED	可寻址 RGB 发光二极管，由 GPIO18 驱动。
5 V to 3.3 V LDO（5 V 转 3.3 V LDO）	电源转换器，输入 5 V，输出 3.3 V。

开始开发应用 通电前，请确保 ESP32-S2-DevKitC-1 完好无损。

必备硬件

- ESP32-S2-DevKitC-1
- USB 2.0 数据线（标准 A 型转 Micro-B 型）
- 电脑（Windows、Linux 或 macOS）

注解： 请确保使用适当的 USB 数据线。部分数据线仅可用于充电，无法用于数据传输和编程。

硬件设置 通过 **USB 转 UART 接口**连接开发板与电脑。目前有关 **ESP32-S2 USB 接口**连接的文档尚不完善。在后续步骤中，默认使用 **USB 转 UART 接口**。

软件设置 请前往 ESP-IDF 快速入门，在 详细安装步骤 小节查看如何快速设置开发环境，将应用程序烧录至 ESP32-S2-DevKitC-1。

内含组件和包装

零售订单 如购买样品，每个 ESP32-S2-DevKitC-1 将以防静电袋或零售商选择的其他方式包装。

零售订单请前往 https://www.espressif.com/zh-hans/company/contact/buy-a-sample。

批量订单 如批量购买，ESP32-S2-DevKitC-1 将以大纸板箱包装。

批量订单请前往 https://www.espressif.com/zh-hans/contact-us/sales-questions。

硬件参考

功能框图 ESP32-S2-DevKitC-1 的主要组件和连接方式如下图所示。

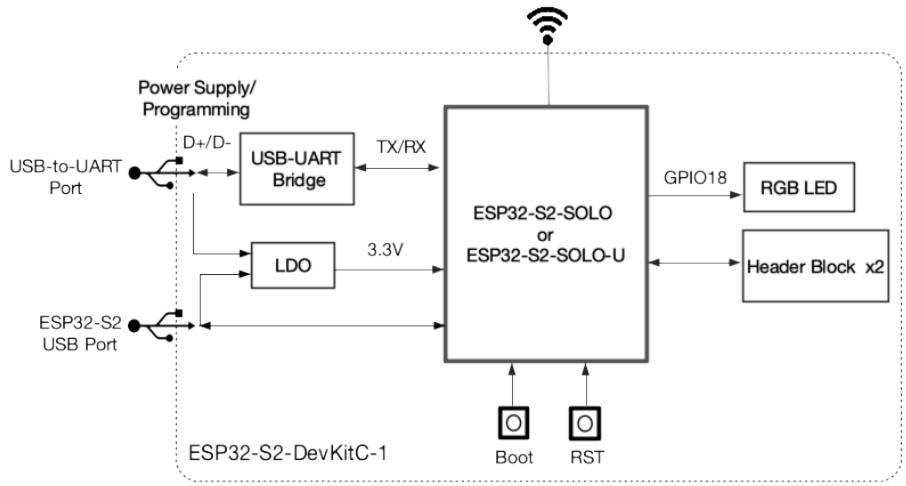


图 12: ESP32-S2-DevKitC-1（点击放大）

电源选项 您可从以下三种供电方式中任选其一给 ESP32-S2-DevKitC-1 供电：

- USB 转 UART 接口供电或 ESP32-S2 USB 接口供电（选择其一或同时供电），默认供电方式（推荐）
- 5V 和 G (GND) 排针供电
- 3V3 和 G (GND) 排针供电

排针 下表列出了开发板两侧排针（J1 和 J3）的 **名称**和 **功能**，排针的名称如图*ESP32-S2-DevKitC-1 - 正面* 所示，排针的序号与 ESP32-S2-DevKitC-1 原理图 (PDF) 一致。

J1

序号	名称	类型[1]	功能
1	3V3	P	3.3 V 电源
2	3V3	P	3.3 V 电源
3	RST	I	CHIP_PU
4	4	I/O/T	RTC_GPIO4, GPIO4, TOUCH4, ADC1_CH3
5	5	I/O/T	RTC_GPIO5, GPIO5, TOUCH5, ADC1_CH4
6	6	I/O/T	RTC_GPIO6, GPIO6, TOUCH6, ADC1_CH5
7	7	I/O/T	RTC_GPIO7, GPIO7, TOUCH7, ADC1_CH6
8	15	I/O/T	RTC_GPIO15, GPIO15, U0RTS, ADC2_CH4, XTAL_32K_P
9	16	I/O/T	RTC_GPIO16, GPIO16, U0CTS, ADC2_CH5, XTAL_32K_N
10	17	I/O/T	RTC_GPIO17, GPIO17, U1TXD, ADC2_CH6, DAC_1
11	18	I/O/T	RTC_GPIO18, GPIO18, U1RXD, ADC2_CH7, DAC_2, CLK_OUT3, RGB LED
12	8	I/O/T	RTC_GPIO8, GPIO8, TOUCH8, ADC1_CH7
13	3	I/O/T	RTC_GPIO3, GPIO3, TOUCH3, ADC1_CH2
14	46	I	GPIO46
15	9	I/O/T	RTC_GPIO9, GPIO9, TOUCH9, ADC1_CH8, FSPIHD
16	10	I/O/T	RTC_GPIO10, GPIO10, TOUCH10, ADC1_CH9, FSPICS0, FSPIIO4
17	11	I/O/T	RTC_GPIO11, GPIO11, TOUCH11, ADC2_CH0, FSPID, FSPIIO5
18	12	I/O/T	RTC_GPIO12, GPIO12, TOUCH12, ADC2_CH1, FSPICLK, FSPIIO6
19	13	I/O/T	RTC_GPIO13, GPIO13, TOUCH13, ADC2_CH2, FSPIQ, FSPIIO7
20	14	I/O/T	RTC_GPIO14, GPIO14, TOUCH14, ADC2_CH3, FSPIWP, FSPIDQS
21	5V	P	5 V 电源
22	G	G	接地

J3

序号	名称	类型	功能
1	G	G	接地
2	TX	I/O/T	U0TXD, GPIO43, CLK_OUT1
3	RX	I/O/T	U0RXD, GPIO44, CLK_OUT2
4	1	I/O/T	RTC_GPIO1, GPIO1, TOUCH1, ADC1_CH0
5	2	I/O/T	RTC_GPIO2, GPIO2, TOUCH2, ADC1_CH1
6	42	I/O/T	MTMS, GPIO42
7	41	I/O/T	MTDI, GPIO41, CLK_OUT1
8	40	I/O/T	MTDO, GPIO40, CLK_OUT2
9	39	I/O/T	MTCK, GPIO39, CLK_OUT3
10	38	I/O/T	GPIO38, FSPIWP
11	37	I/O/T	SPIDQS, GPIO37, FSPIQ
12	36	I/O/T	SPIIO7, GPIO36, FSPICLK
13	35	I/O/T	SPIIO6, GPIO35, FSPID
14	0	I/O/T	RTC_GPIO0, GPIO0
15	45	I/O/T	GPIO45
16	34	I/O/T	SPIIO5, GPIO34, FSPICS0
17	33	I/O/T	SPIIO4, GPIO33, FSPIHD
18	21	I/O/T	RTC_GPIO21, GPIO21
19	20	I/O/T	RTC_GPIO20, GPIO20, U1CTS, ADC2_CH9, CLK_OUT1, USB_D+
20	19	I/O/T	RTC_GPIO19, GPIO19, U1RTS, ADC2_CH8, CLK_OUT2, USB_D-
21	G	G	接地
22	G	G	接地

管脚布局

P：电源；I：输入；O：输出；T：可设置为高阻。

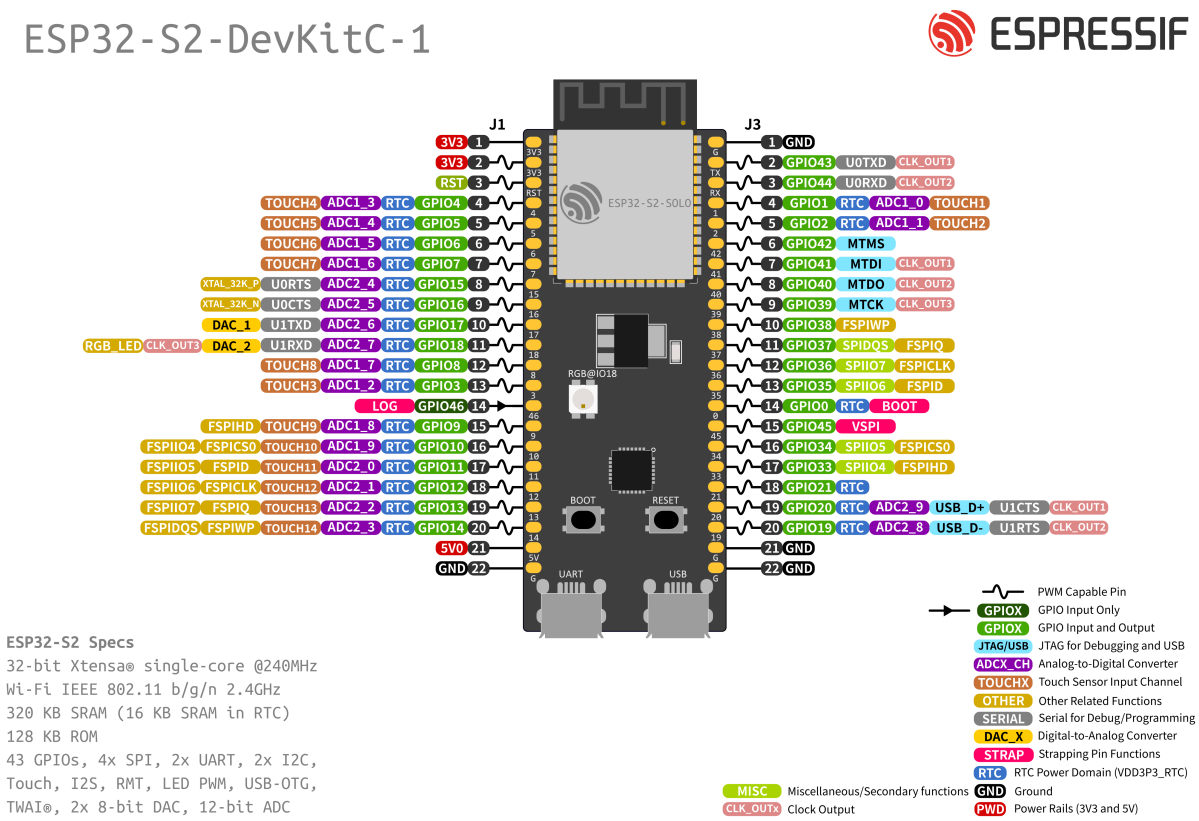


图 13: ESP32-S2-DevKitC-1 管脚布局（点击放大）

硬件版本

无历史版本。

相关文档

- ESP32-S2 系列芯片规格书 (PDF)
- ESP32-S2-SOLO & ESP32-S2-SOLO-U 模组技术规格书 (PDF)
- ESP32-S2-DevKitC-1 原理图 (PDF)
- ESP32-S2-DevKitC-1 PCB 布局图 (PDF)
- ESP32-S2-DevKitC-1 尺寸图 (PDF)
- ESP32-S2-DevKitC-1 尺寸图源文件 (DXF) - 可使用 Autodesk Viewer 查看

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

1.3.4 ESP32-S2-Kaluga-1 套件 v1.3

更早版本：*ESP32-S2-Kaluga-1 套件 v1.2*

ESP32-S2-Kaluga-1 v1.3 是一款来自乐鑫的开发套件，主要可用于以下目的：

- 展示 ESP32-S2 芯片的人机交互功能
- 为用户提供基于 ESP32-S2 的人机交互应用开发工具

ESP32-S2 的功能强大，应用场景非常丰富。对于初学者来说，可能的用例包括：

- **智能家居**：从最简单的智能照明、智能门锁、智能插座，到更复杂的视频流设备、安防摄像头、OTT 设备和家用电器等
- **电池供电设备**：Wi-Fi mesh 传感器网络、Wi-Fi 网络玩具、可穿戴设备、健康管理设备等

- **工业自动化设备**：无线控制与机器人技术、智能照明、HVAC 控制设备等
- **零售和餐饮业**：POS 机和服务机器人

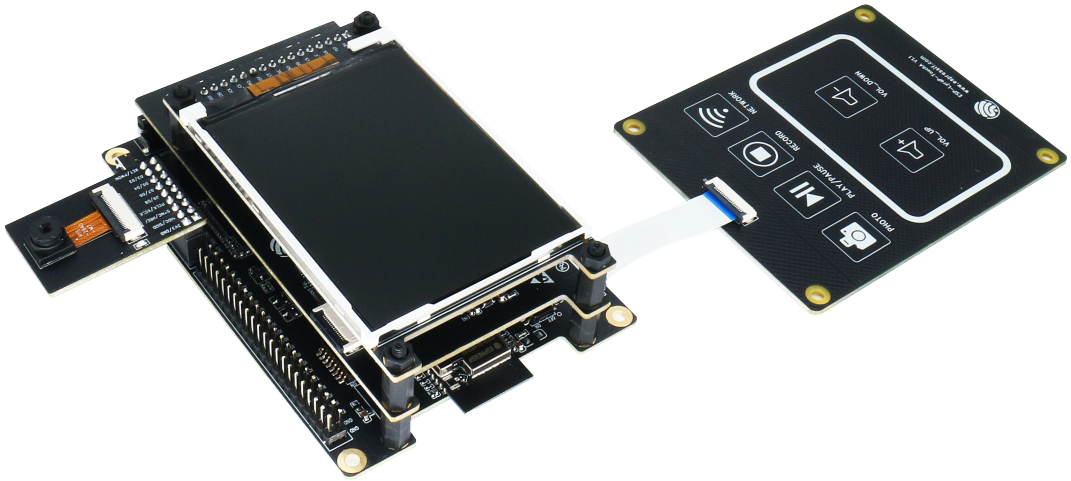


图 14: ESP32-S2-Kaluga-1-Kit 概述（点击放大）

ESP32-S2-Kaluga-1 套件包括以下几个开发板：

- 主板：*ESP32-S2-Kaluga-1*
- 扩展板：
  - *ESP-LyraT-8311A v1.3* - 音频播放器
  - *ESP-LyraP-TouchA v1.1* - 触摸板
  - *ESP-LyraP-LCD32 v1.2* - 3.2” LCD 屏
  - *ESP-LyraP-CAM v1.1* - 摄像头

由于 ESP32-S2 的管脚复用，部分扩展板的兼容性有所限制，具体请见扩展板的兼容性。

本文档主要介绍 **ESP32-S2-Kaluga-1 主板**及其与扩展板的交互。更多有关具体扩展板的信息，请点击相应的链接。

本指南包括：

- 快速入门：提供 ESP32-S2-Kaluga-1 的简要概述及必须了解的硬件和软件信息。
- 硬件参考：提供 ESP32-S2-Kaluga-1 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

快速入门

本节介绍如何开始使用 ESP32-S2-Kaluga-1，主要包括三大部分：首先，介绍一些关于 ESP32-S2-Kaluga-1 的基本信息，然后在应用程序开发 章节介绍如何进行硬件初始化，最后介绍如何为 ESP32-S2-Kaluga-1 烧录固件。

概述 ESP32-S2-Kaluga-1 主板是整个套件的核心。该主板集成了 ESP32-S2-WROVER 模组，并配备连接至各个扩展板的连接器。ESP32-S2-Kaluga-1 是人机交互接口原型设计的关键工具。

ESP32-S2-Kaluga-1 主板配备了多个连接器，可用于连接相应扩展板：

- 扩展板连接器，用于连接 ESP-LyraT-8311A、ESP-LyraP-LCD32
- 摄像头连接器，用于连接 ESP-LyraP-CAM
- 触摸 FPC 连接器，用于连接 ESP-LyraP-TouchA
- LCD FPC 连接器（尚无可用官方配套扩展板）
- I2C FPC 连接器（尚无可用官方配套扩展板）

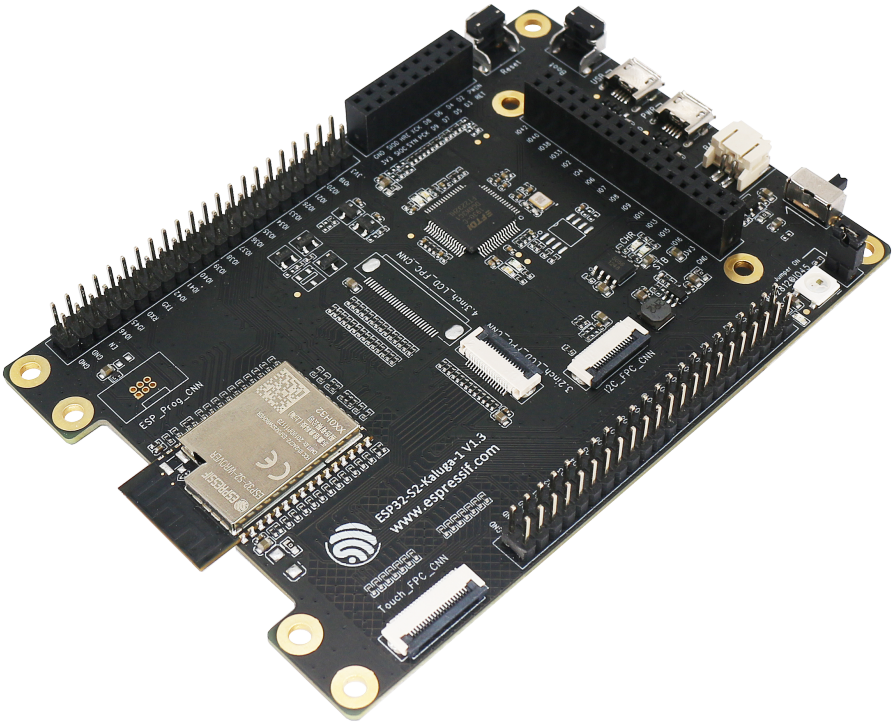


图 15: ESP32-S2-Kaluga-1（点击放大）

所有四个扩展板都经过特别设计，以支持以下功能：

- **触摸板控制**
  - 带有 6 个触摸按钮
  - 支持最大 5 mm 亚克力板
  - 支持湿手操作
  - 支持防水功能。ESP32-S2 可以配置为在多个触摸板同时被水复盖时自动禁用所有触摸板功能，并在去除水滴后重新启用触摸板
- **音频播放**
  - 连接扬声器，以播放音频
  - 配合触控板使用，可控制音频播放和调节音量
- **LCD 显示屏**
  - LCD 接口（8 位并行 RGB、8080 和 6800 接口）
- **摄像头图像采集**
  - 支持 OV2640 和 OV3660 摄像头模块
  - 8-bit DVP 图像传感器接口（ESP32-S2 还支持 16 位 DVP 图像传感器，但需要您自行进行二次开发）
  - 支持高达 40 MHz 时钟频率
  - 优化 DMA 传输带宽，便于传输高分辨率图像

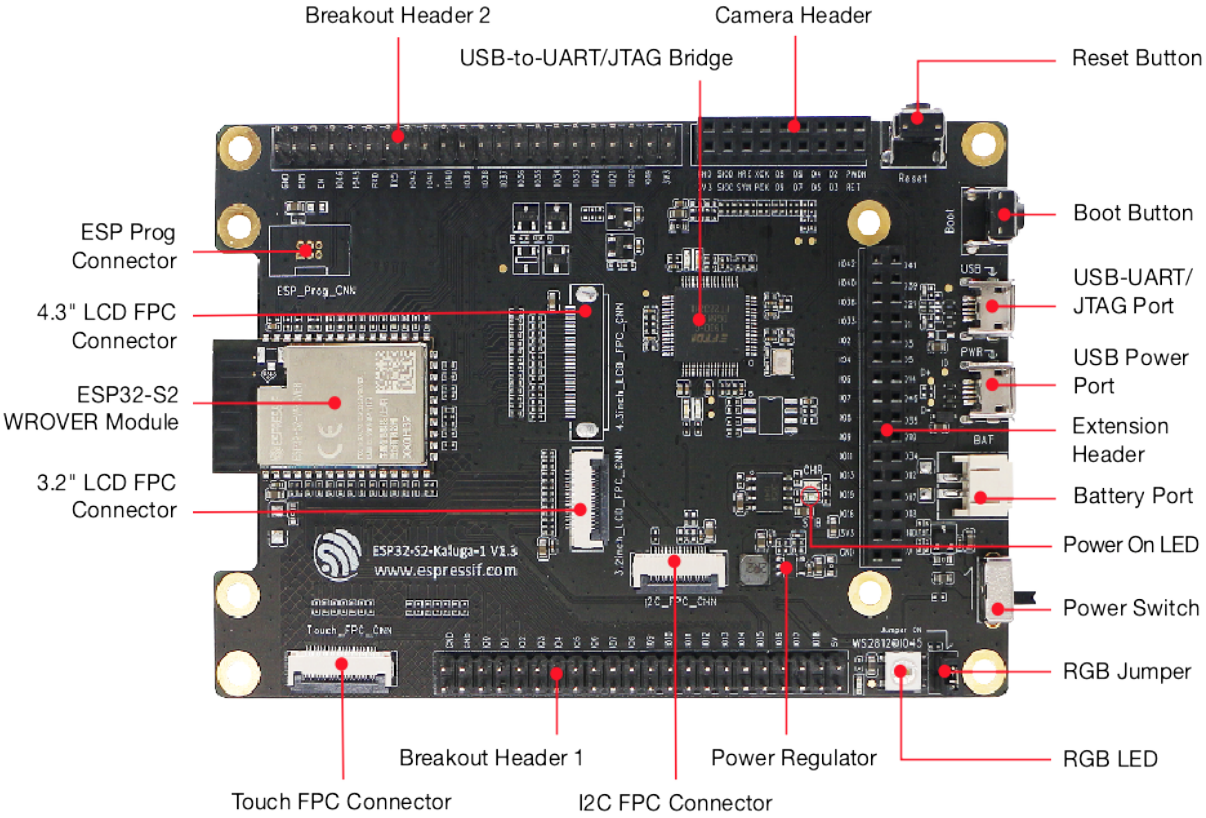


图 16: ESP32-S2-Kaluga-1 - 正面（点击放大）

组件描述 下表将从左边的 ESP32-S2 模组开始，以顺时针顺序介绍上图中的主要组件。

保留表示该功能可用，但当前版本的套件并未启用该功能。

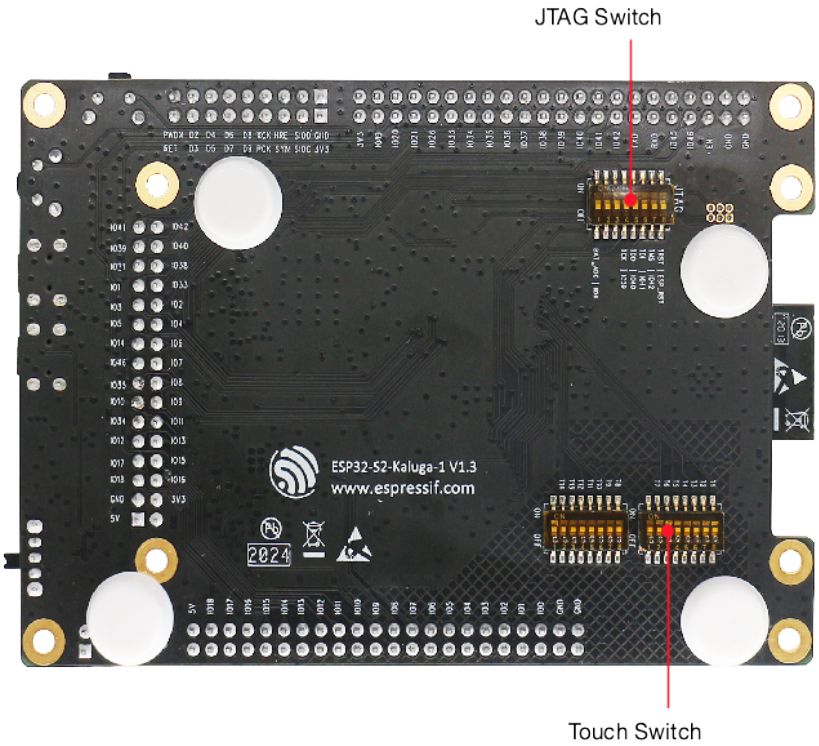


图 17: ESP32-S2-Kaluga-1 - 反面（点击放大）

主要组件	描述
ESP32-S2-WROVER 模组	集成 ESP32-S2 芯片，可提供 Wi-Fi 连接、数据处理和灵活的数据存储功能。
4.3” LCD FPC 连接器	（保留）可使用 FPC 线连接 4.3” LCD 扩展板。
ESP Prog 连接器	（保留）用于连接乐鑫固件烧录设备 (ESP-Prog)。
JTAG 开关	切换到 ON 方向，启用 ESP32-S2 和 FT2232 之间的连接。此时，可通过 USB-UART/JTAG 端口进行 JTAG 调试，详见*JTAG 调试*。
引出管脚排针 2	ESP32-S2-WROVER 模组的部分 GPIO 直接引出至该开发板（详见开发板上的标记）。
USB-to-UART/JTAG 桥接器	FT2232 适配器开发板，允许在 USB 端口使用 UART/JTAG 协议通信。
摄像头连接器	用于连接摄像头扩展板，比如 ESP-LyraP-CAM。
扩展板连接器	用于连接带有配套连接器的扩展板。
Reset 复位按钮	用于重启系统。
Boot 按钮	按下 **Boot** 键并保持，同时按一下 **Reset** 键，进入 “固件下载” 模式，通过串口下载固件。
USB-UART/JTAG 端口	PC 和 ESP32-S2 模组之间的通信接口（UART 或 JTAG）。
USB 电源端口	为开发板供电。
电池端口	2 针连接器，用于连接外部电池。
电源 LED 指示灯	当 USB 电源或外部电源供电电压正常，则 LED 亮起。
电源开关	打开可为系统供电。
RGB 跳线	如需使用 RGB LED，需在该位置增加一个跳线。
RGB LED 指示灯	可编程 RGB LED 指示灯，受控于 GPIO45。在使用前需要安装 RGB 跳线。
调压器	5 V 转 3.3 V 调压器。
I2C FPC 连接器	（保留）可通过 FPC 线连接其他 I2C 扩展板。
引出管脚排针 1	ESP32-S2-WROVER 模组的部分 GPIO 直接引出至该开发板（详见开发板上的标记）。
触摸 FPC 连接器	可通过 FPC 线连接 ESP-LyraP-TouchA 扩展板。
触摸开关	切换到 OFF 方向，配置 GPIO1 到 GPIO14 连接触摸传感器；切换到 ON 方向，配置 GPIO1 到 GPIO14 用于其他目的。
3.2” LCD FPC 连接器	可通过 FPC 线连接 3.2” LCD 扩展板，比如 ESP-LyraP-LCD32。

应用程序开发 ESP32-S2-Kaluga-1 上电前，请首先确认开发板完好无损。

硬件准备

- ESP32-S2-Kaluga-1
- 两根 USB 2.0 电缆（标准 A 转 Micro-B）
  - 电源选项
  - 用于 UART/JTAG 通信
- PC（Windows、Linux 或 macOS）
- 您选择的任何扩展板

硬件设置

1. 连接您选择的扩展板（更多信息，请见对应拓展板的用户指南）
2. 插入两根 USB 电缆
3. 打开 **电源开关**时，**电源 LED 指示灯**应点亮。

软件设置 请前往快速入门，在详细安装步骤 一节查看如何快速设置开发环境。

您还可以点击 这里，获取有关 ESP32-S2-Kaluga-1 套件编程指南与应用示例的更多内容。

内容和包装

零售订单 每一个零售 ESP32-S2-Kaluga-1 开发套件均有独立包装。

内含以下部分：

- **主板**
  - ESP32-S2-Kaluga-1
- **扩展板：**
  - ESP-LyraT-8311A
  - ESP-LyraP-CAM
  - ESP-LyraP-TouchA
  - ESP-LyraP-LCD32
- **连接器**
  - 20 针 FPC 线（用于连接 ESP32-S2-Kaluga-1 主板至 ESP-LyraP-TouchA 扩展板）
- **紧固件**
  - 安装螺栓 (x 8)
  - 螺丝 (x 4)
  - 螺母 (x 4)

零售购买，请前往 https://www.espressif.com/zh-hans/contact-us/get-samples。

批发订单 ESP32-S2-Kaluga-1 开发套件的批发包装为纸板箱。

批量订货，请参考 乐鑫产品订购信息 (PDF)。

硬件参考

功能框图 ESP32-S2-Kaluga-1 的主要组件和连接方式如下图所示。

电源选项 开发板可任一选用以下四种供电方式：

- Micro USB 端口供电（默认）
- 通过 2 针电池连接器使用外部电池供电
- 5V / GND 管脚供电
- 3V3 / GND 管脚供电

图 18: ESP32-S2-Kaluga-1 - 包装

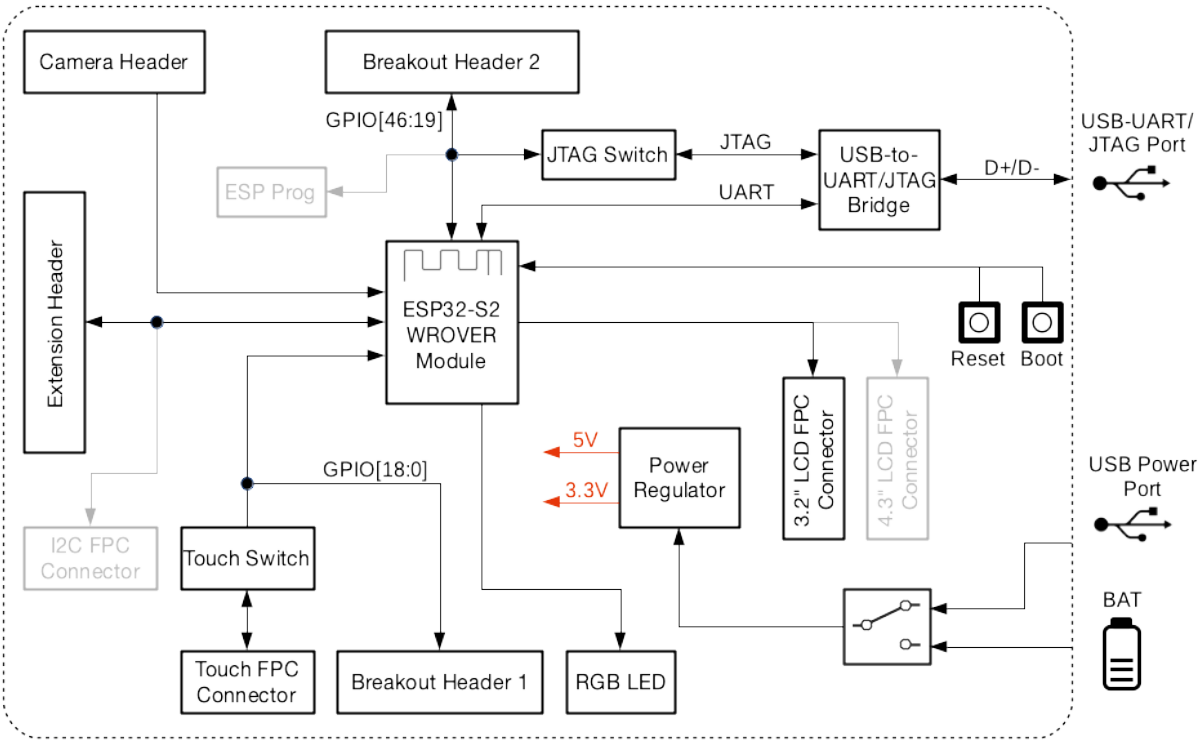


图 19: ESP32-S2-Kaluga-1 功能框图

扩展板的兼容性 如需同时使用多块扩展板，请首先查看以下兼容性信息：

扩展板组合	复用接口或管脚	无法运行原因	解决方案
8311A v1.3 + CAM v1.1	I2S 控制器	ESP32-S2 仅有 1 个 I2S 接口，但这两个开发板均需使用 ESP32-S2 的 I2S 接口进行通信（ESP-LyraT-8311A 使用标准模式；ESP-LyraP-CAM 使用 Camera 协议）。	采用分时复用；或另外选择一款可以通过其他 GPIOs 或 DAC 连接的音频扩展板。
TouchA v1.1 + LCD32 v1.2	IO11、IO6	由于管脚 IO11 复用，导致无法触发触摸动作；ESP-LyraP-LCD32 则由于其 BLCT 管脚已与 IO6 断开，因此不受影响。	不要初始化 ESP-LyraP-TouchA 扩展板的 IO11 (NETWORK) 管脚；或者配置 ESP-LyraP-LCD32 扩展板的 BLCT 管脚为 *-1*（相当于不使用 BLCT）。
8311A v1.3 + LCD32 v1.2	IO6	配置 ESP-LyraP-LCD32 扩展板的 BK 管脚为 *-1*（相当于不使用 BK）。	ESP32-S2-Kaluga-1 的 BLCT 管脚将从 IO6 断开。
TouchA v1.1 + 8311A v1.3	ESP-LyraT-8311A 的 BT_ADC 管脚	ESP-LyraT-8311A 在初始化 6 个按钮时需要使用 BT_ADC 管脚，而 ESP-LyraP-TouchA 在完成触摸动作时也需要使用 BT_ADC 管脚。	如需使用 ESP-LyraT-8311A 的 6 个按钮，则不要初始化 ESP-LyraP-TouchA 的 IO6 (PHOTO) 管脚。
TouchA v1.1 + CAM v1.1	IO1、IO2、IO3	由于管脚复用无法同时使用。	不要初始化 ESP-LyraP-TouchA 的 IO1 (VOL_UP)、IO2 (PLAY) 和 IO3 (VOL_DOWN)。
TouchA v1.1 + LCD32 v1.2 + CAM v1.1	IO1、IO2、IO3、IO11	由于管脚复用无法同时使用。	不要初始化 ESP-LyraP-TouchA 的 IO1 (VOL_UP)、IO2 (PLAY)、IO3 (VOL_DOWN) 和 IO11 (NETWORK)。
TouchA v1.1 + LCD32 v1.2 + 8311A v1.3	IO6、IO11	如果使用 ESP-LyraT-8311A 的 BT_ADC 管脚初始化开发板的 6 个按钮，其他扩展板则无法使用 IO6 和 IO11。	不要初始化 ESP-LyraP-TouchA 的 IO11 (NETWORK)。此外，如果需要使用 BT_ADC，则不要初始化 IO6 (PHOTO)。

硬件修订历史

ESP32-S2-Kaluga-1 Kit v1.3

- 以下管脚已重新分配，以解决固件烧录问题：
  - Camera D2：GPIO36
  - Camera D3：GPIO37
  - AU_I2S1_SDI：GPIO34
  - AU_WAKE_INT：GPIO46
- RGB 已移动至开发板边缘
- 所有 dip 开关均移动至开发板的反面，从而便利用户操作

ESP32-S2-Kaluga-1 Kit v1.2 *首次发布*

相关文档

ESP32-S2-Kaluga-1 套件 v1.2 最新版本：*ESP32-S2-Kaluga-1 套件 v1.3*

ESP32-S2-Kaluga-1 v1.2 是一款来自乐鑫的开发套件，主要可用于以下目的：

- 展示 ESP32-S2 芯片的人机交互功能
- 为用户提供基于 ESP32-S2 的人机交互应用开发工具

ESP32-S2 的功能强大，应用场景非常丰富。对于初学者来说，可能的用例包括：

- **智能家居**：从最简单的智能照明、智能门锁、智能插座，到更复杂的视频流设备、安防摄像头、OTT 设备和家用电器等
- **电池供电设备**：Wi-Fi mesh 传感器网络、Wi-Fi 网络玩具、可穿戴设备、健康管理设备等
- **工业自动化设备**：无线控制与机器人技术、智能照明、HVAC 控制设备等
- **零售和餐饮业**：POS 机和服务机器人

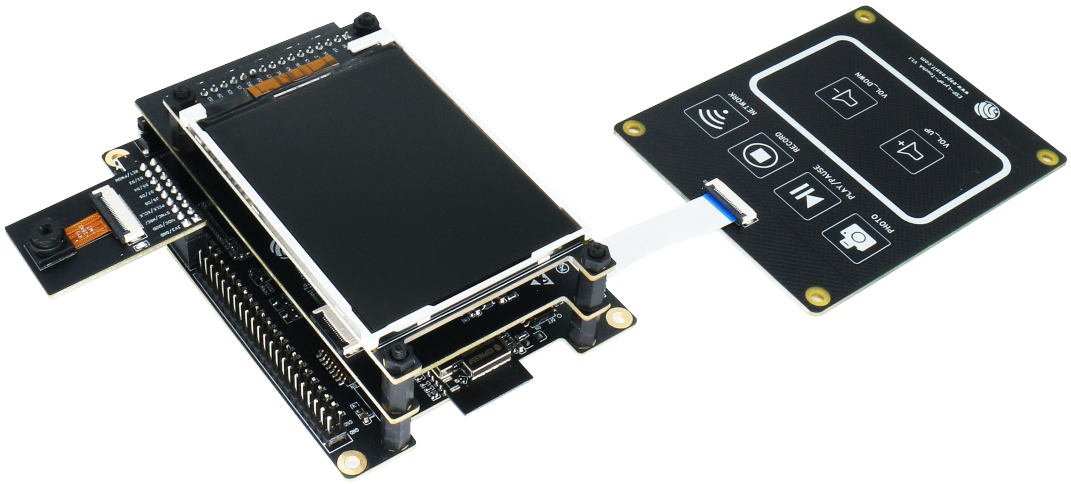


图 20: ESP32-S2-Kaluga-1-Kit 概述（点击放大）

ESP32-S2-Kaluga-1 套件包括以下几个开发板：

- 主板：*ESP32-S2-Kaluga-1*
- 扩展板：
  - *ESP-LyraT-8311A v1.2* - 音频播放器
  - *ESP-LyraP-TouchA v1.1* - 触摸板
  - *ESP-LyraP-LCD32 v1.1* - 3.2” LCD 屏
  - *ESP-LyraP-CAM v1.0* - 摄像头

由于 ESP32-S2 的管脚复用，部分扩展板的兼容性有所限制，具体请见扩展板的兼容性。

本文档主要介绍 **ESP32-S2-Kaluga-1 主板**及其与扩展板的交互。更多有关具体扩展板的信息，请点击相应的链接。

本指南包括：

- 快速入门：提供 ESP32-S2-Kaluga-1 的简要概述及必须了解的硬件和软件信息。
- 硬件参考：提供 ESP32-S2-Kaluga-1 的详细硬件信息。
- 硬件修订历史：提供该开发板的“修订历史”、“已知问题”以及此开发板之前版本的用户指南链接。
- 相关文档：提供相关文档的链接。

快速入门 本节介绍如何开始使用 ESP32-S2-Kaluga-1，主要包括三大部分：首先，介绍一些关于 ESP32-S2-Kaluga-1 的基本信息，然后在应用程序开发 章节介绍如何进行硬件初始化，最后介绍如何为 ESP32-S2-Kaluga-1 烧录固件。

概述 ESP32-S2-Kaluga-1 主板是整个套件的核心。该主板集成了 ESP32-S2-WROVER 模组，并配备连接至各个扩展板的连接器。ESP32-S2-Kaluga-1 是人机交互接口原型设计的关键工具。

ESP32-S2-Kaluga-1 主板配备了多个连接器，可用于连接相应扩展板：

- 扩展板连接器，用于连接 ESP-LyraT-8311A、ESP-LyraP-LCD32
- 摄像头连接器，用于连接 ESP-LyraP-CAM
- 触摸 FPC 连接器，用于连接 ESP-LyraP-TouchA

- LCD FPC 连接器（尚无可用官方配套扩展板）
- I2C FPC 连接器（尚无可用官方配套扩展板）

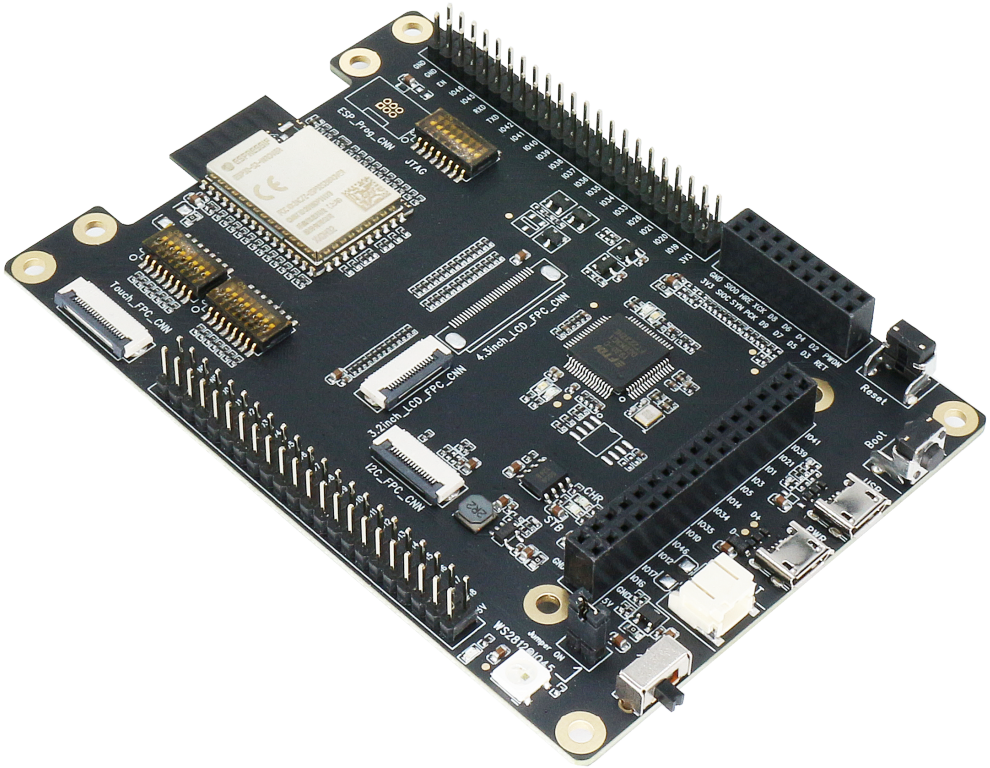


图 21: ESP32-S2-Kaluga-1（点击放大）

所有四个扩展板都经过特别设计，以支持以下功能：

- **触摸板控制**
  - 带有 6 个触摸按钮
  - 支持最大 5 mm 亚克力板
  - 支持湿手操作
  - 支持防水功能。ESP32-S2 可以配置为在多个触摸板同时被水复盖时自动禁用所有触摸板功能，并在去除水滴后重新启用触摸板
- **音频播放**
  - 连接扬声器，以播放音频
  - 配合触控板使用，可控制音频播放和调节音量
- **LCD 显示屏**
  - LCD 接口（8 位并行 RGB、8080 和 6800 接口）
- **摄像头图像采集**
  - 支持 OV2640 和 OV3660 摄像头模块
  - 8-bit DVP 图像传感器接口（ESP32-S2 还支持 16 位 DVP 图像传感器，但需要您自行进行二次开发）
  - 支持高达 40 MHz 时钟频率
  - 优化 DMA 传输带宽，便于传输高分辨率图像

组件描述 下表将从左边的 ESP32-S2 模组开始，以顺时针顺序介绍上图中的主要组件。

保留表示该功能可用，但当前版本的套件并未启用该功能。

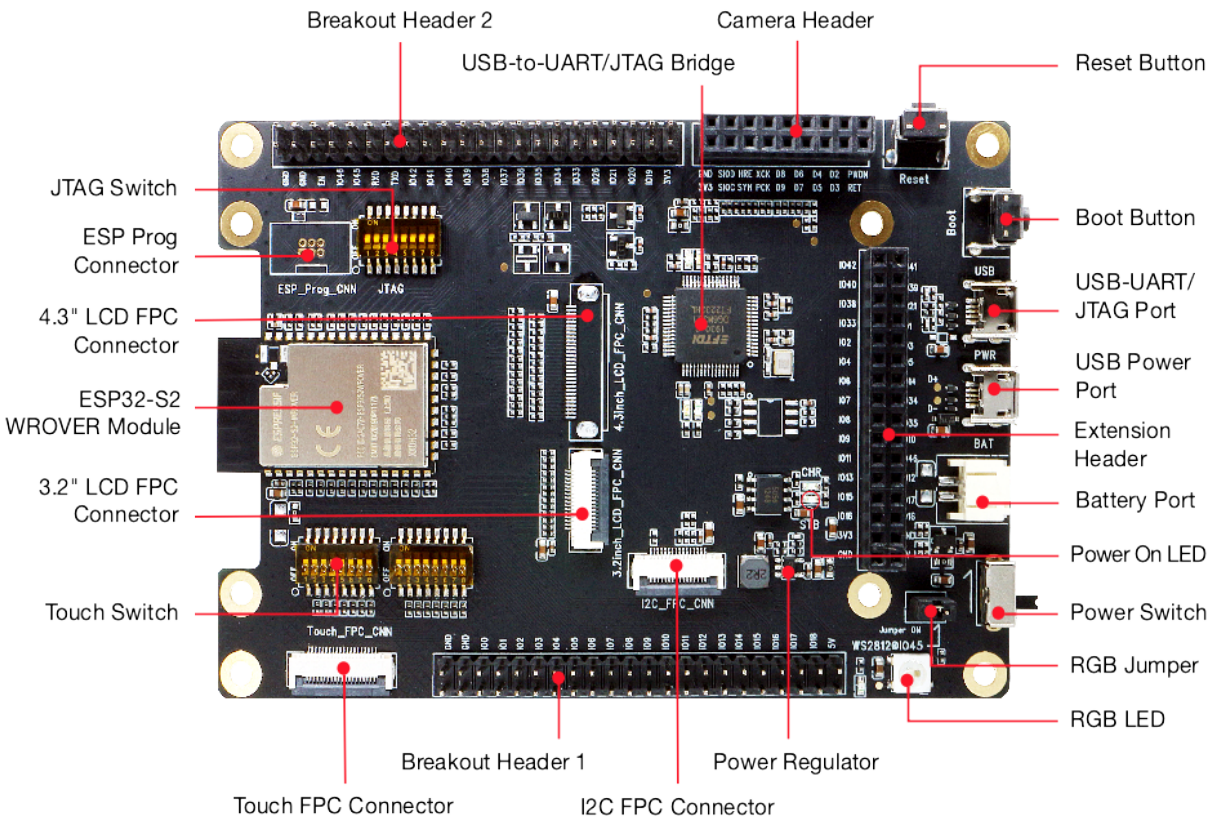


图 22: ESP32-S2-Kaluga-1 - 正面（点击放大）

主要组件	描述
ESP32-S2-WROVER 模组	集成 ESP32-S2 芯片，可提供 Wi-Fi 连接、数据处理和灵活的数据存储功能。
4.3” LCD FPC 连接器	（保留）可使用 FPC 线连接 4.3” LCD 扩展板。
ESP Prog 连接器	（保留）用于连接乐鑫固件烧录设备 (ESP-Prog)。
JTAG 开关	切换到 ON 方向，启用 ESP32-S2 和 FT2232 之间的连接。此时，可通过 USB-UART/JTAG 端口进行 JTAG 调试，详见*JTAG 调试*。
引出管脚排针 2	ESP32-S2-WROVER 模组的部分 GPIO 直接引出至该开发板（详见开发板上的标记）。
USB-to-UART/JTAG 桥接器	FT2232 适配器开发板，允许在 USB 端口使用 UART/JTAG 协议通信。
摄像头连接器	用于连接摄像头扩展板，比如 ESP-LyraP-CAM。
扩展板连接器	用于连接带有配套连接器的扩展板。
Reset 复位按钮	用于重启系统。
Boot 按钮	按下 **Boot** 键并保持，同时按一下 **Reset** 键，进入“固件下载”模式，通过串口下载固件。
USB-UART/JTAG 端口	PC 和 ESP32-S2 模组之间的通信接口（UART 或 JTAG）。
USB 电源端口	为开发板供电。
电池端口	2 针连接器，用于连接外部电池。
电源 LED 指示灯	当 USB 电源或外部电源供电电压正常，则 LED 亮起。
电源开关	打开可为系统供电。
RGB 跳线	如需使用 RGB LED，需在该位置增加一个跳线。
RGB LED 指示灯	可编程 RGB LED 指示灯，受控于 GPIO45。在使用前需要安装 RGB 跳线。
调压器	5 V 转 3.3 V 调压器。
I2C FPC 连接器	（保留）可通过 FPC 线连接其他 I2C 扩展板。
引出管脚排针 1	ESP32-S2-WROVER 模组的部分 GPIO 直接引出至该开发板（详见开发板上的标记）。
触摸 FPC 连接器	可通过 FPC 线连接 ESP-LyraP-TouchA 扩展板。
触摸开关	切换到 OFF 方向，配置 GPIO1 到 GPIO14 连接触摸传感器；切换到 ON 方向，配置 GPIO1 到 GPIO14 用于其他目的。
3.2” LCD FPC 连接器	可通过 FPC 线连接 3.2” LCD 扩展板，比如 ESP-LyraP-LCD32。

应用程序开发 ESP32-S2-Kaluga-1 上电前，请首先确认开发板完好无损。

硬件准备

- ESP32-S2-Kaluga-1
- 两根 USB 2.0 电缆（标准 A 转 Micro-B）
  - 电源选项
  - 用于 UART/JTAG 通信
- PC（Windows、Linux 或 macOS）
- 您选择的任何扩展板

硬件设置

1. 连接您选择的扩展板（更多信息，请见对应拓展板的用户指南）
2. 插入两根 USB 电缆
3. 打开 **电源开关**时，**电源 LED 指示灯**应点亮。

软件设置 请前往快速入门，在详细安装步骤 一节查看如何快速设置开发环境。

您还可以点击 这里，获取有关 ESP32-S2-Kaluga-1 套件编程指南与应用示例的更多内容。

内容和包装

零售订单 每一个零售 ESP32-S2-Kaluga-1 开发套件均有独立包装，内含以下部分：

- **主板**
  - ESP32-S2-Kaluga-1
- **扩展板：**
  - ESP-LyraT-8311A
  - ESP-LyraP-CAM
  - ESP-LyraP-TouchA
  - ESP-LyraP-LCD32
- **连接器**
  - 20 针 FPC 线（用于连接 ESP32-S2-Kaluga-1 主板至 ESP-LyraP-TouchA 扩展板）
- **紧固件**
  - 安装螺栓 (x 8)
  - 螺丝 (x 4)
  - 螺母 (x 4)

零售购买，请前往 https://www.espressif.com/zh-hans/contact-us/get-samples。

批发订单 ESP32-S2-Kaluga-1 开发套件的批发包装为纸板箱。

批量订货，请参考 乐鑫产品订购信息 (PDF)。

硬件参考

功能框图 ESP32-S2-Kaluga-1 的主要组件和连接方式如下图所示。

电源选项 开发板可任一选用以下四种供电方式：

- Micro USB 端口供电（默认）
- 通过 2 针电池连接器使用外部电池供电
- 5V / GND 管脚供电
- 3V3 / GND 管脚供电

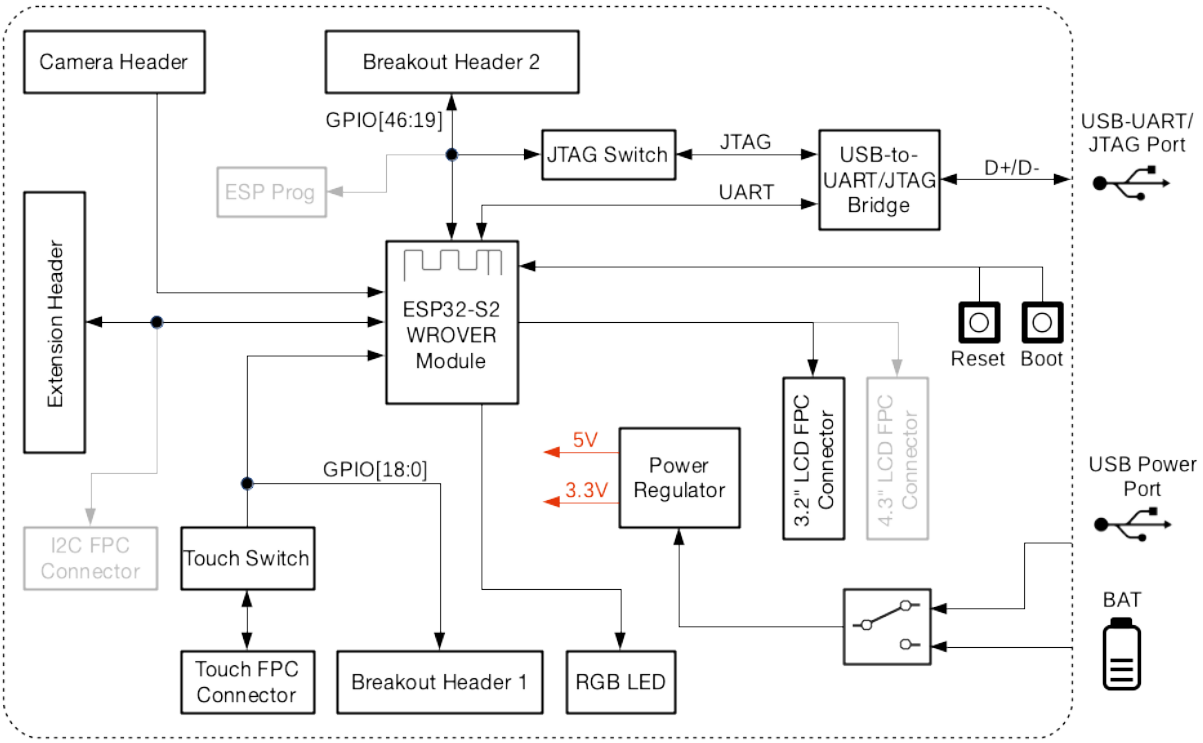


图 23: ESP32-S2-Kaluga-1 功能框图

扩展板的兼容性 如需同时使用多块扩展板，请首先查看以下兼容性信息：

扩展板组合	复用接口或管脚	无法运行原因	解决方案
8311A v1.2 + CAM v1.0	I2S 控制、IO46	ESP32-S2 仅有 1 个 I2S 接口，但这两个开发板均需使用 ESP32-S2 的 I2S 接口进行通信（ESP-LyraT-8311A 使用标准模式；ESP-LyraP-CAM 使用 Camera 协议）。如两个扩展板同时复用 IO46，ESP-LyraP-CAM 的正常使用将受到干扰。	暂无解决方法。
TouchA v1.1 + LCD32 v1.1	IO11、IO6	ESP-LyraP-TouchA 因管脚 IO11 复用，导致无法触发触摸动作；ESP-LyraP-LCD32 因 BK (BLCT) 管脚连接至 IO6 管脚复用，因此也无法使用。	不要初始化 ESP-LyraP-TouchA 扩展板的 IO11 (NETWORK) 和 IO6 (PHOTO) 管脚。
8311A v1.2 + LCD32 v1.1	IO6	这两款扩展板可以同时使用，但由于 ESP32-S2-Kaluga-1 的 BK (BLCT) 管脚已连接至 IO6，因此，ESP-LyraT-8311A 的 BT_ADC 管脚和 6 个按钮均无法使用。	用户也可通过以下配置使用 ESP-LyraT-8311A 的 BT_ADC 管脚：移除 ESP-LyraP-LCD32 扩展板上的 R39，将 R41 换为 100 欧，并将 BLCT_L 开关打开。注意，此配置将导致用户无法通过软件控制显示屏的背光亮度。
TouchA v1.1 + 8311A v1.2	ESP-LyraT-8311A 的 BT_ADC 管脚	这两款扩展板可以同时使用。然而，当 ESP-LyraT-8311A 的 BT_ADC 管脚用于初始化扩展板的 6 个按钮时，ESP-LyraP-TouchA 无法成功触发。	如果计划使用 ESP-LyraT-8311A 的 BT_ADC 管脚，请不要初始化 ESP-LyraP-TouchA 扩展板的 IO6 管脚 (PHOTO)。
TouchA v1.1 + CAM v1.0	IO1、IO2、IO3	由于管脚复用无法同时使用。	不要初始化 ESP-LyraP-TouchA 的 IO1 (VOL_UP)、IO2 (PLAY) 和 IO3 (VOL_DOWN)。
TouchA v1.1 + LCD32 v1.1 + CAM v1.0	IO1、IO2、IO3、IO6、IO11	由于管脚复用无法同时使用。	**解决方案 1**：不要初始化 ESP-LyraP-TouchA 扩展板的 IO1 (VOL_UP)、IO2 (PLAY)、IO3 (VOL_DOWN)、IO6 (PHOTO) 和 IO11 (NETWORK)。**解决方案 2**：用户也可通过以下配置正常初始化 IO6 (PHOTO)：移除 ESP-LyraP-LCD32 扩展板上的 R39，将 R41 换为 100 欧，并将 BLCT_L 开关打开。注意，此配置将导致用户无法通过软件控制显示屏的背光亮度。
TouchA v1.1 + LCD32 v1.1 + 8311A v1.2	IO6、IO11	IO11 管脚复用导致无法同时使用；IO6 管脚复用导致 ESP-LyraT-8311A 的 BT_ADC 管脚无法使用，因此无法初始化该扩展板的 6 个按钮。	**解决方法 1**：不要初始化 ESP-LyraP-TouchA 扩展板的 IO6 (PHOTO) 和 IO11 (NETWORK)。注意，此时 ESP-LyraT-8311A 的 6 个按钮依然无法使用。**解决方法 2**：移除 ESP-LyraP-LCD32 扩展板上的 R39，将 R41 换为 100 欧，并将 BLCT_L 开关打开。不要初始化 ESP-LyraP-TouchA 的 IO11 (NETWORK)。如果希望使用 ESP-LyraT-8311A 的 6 个按钮，则也不要初始化 IO6 (PHOTO)。

已知问题

问题硬件	描述	主要原因	解决方法
ESP-LyraP-CAM v1.0、管脚 IO45、管脚 IO46	当 ESP-LyraP-CAM v1.0 连接至主板时，可能导致主板无法烧录固件。	开发板上电时，strapping 管脚 IO45 和 IO46 的上电时序错误，导致开发板无法正常启动。	主板烧录固件时，不应连接该扩展板。
ESP-LyraP-CAM v1.0、管脚 IO45、管脚 IO46	使用 Reset 复位按键重启开发板可能无法达到期望结果。	开发板上电时，strapping 管脚 IO45 和 IO46 的上电时序错误，导致开发板无法正常启动。	v1.2 暂无解决方法。该问题已经在 ESP32-S2-Kaluga-1 V1.3 中进行了修复。
ESP-LyraT-8311A v1.2、管脚 IO46	当 ESP-LyraT-8311A v1.2 连接至主板时，可能导致主板无法烧录固件。	开发板上电时，strapping 管脚 IO46 的上电时序错误，导致开发板无法正常启动。	主板烧录固件时，不应连接该扩展板。
ESP-LyraT-8311A v1.2、管脚 IO46	使用 Reset 复位按键重启开发板可能无法达到期望结果。	开发板上电时，strapping 管脚 IO46 的上电时序错误，导致开发板无法正常启动。	v1.2 暂无解决方法。该问题已经在 ESP32-S2-Kaluga-1 V1.3 中进行了修复。

硬件修订历史 尚无版本升级历史。

相关文档

ESP-LyraP-CAM v1.0 本用户指南可提供 ESP-LyraP-CAM 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 *主板*，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraP-CAM v1.0 扩展板正在搭配*ESP32-S2-Kaluga-1 套件 v1.2* 销售。

ESP-LyraP-CAM 可为您的主板增加摄像头功能。

本指南包括如下内容：

- 概述：提供为了使用 ESP-LyraP-CAM 而必须了解的硬件和软件信息。
- 硬件参考：提供 ESP-LyraP-CAM 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

概述 ESP-LyraP-CAM 扩展板可为您的主板增加一个摄像头。

组件描述

主要组件	描述
主板摄像头排针	连接至主板摄像头连接器
电源 LED 指示灯	如果电源供电电压正常，则红色 LED 亮起
摄像头模块连接器	硬件支持 OV2640 和 OV3660 摄像头模块；目前，ESP-LyraP-CAM 扩展板默认提供 OV2640 摄像头模块
电源调节器	LDO 调压器（3.3 V 至 2.8 V 和 1.5 V）

应用程序开发 ESP-LyraP-CAM 上电前，请首先确认开发板完好无损。

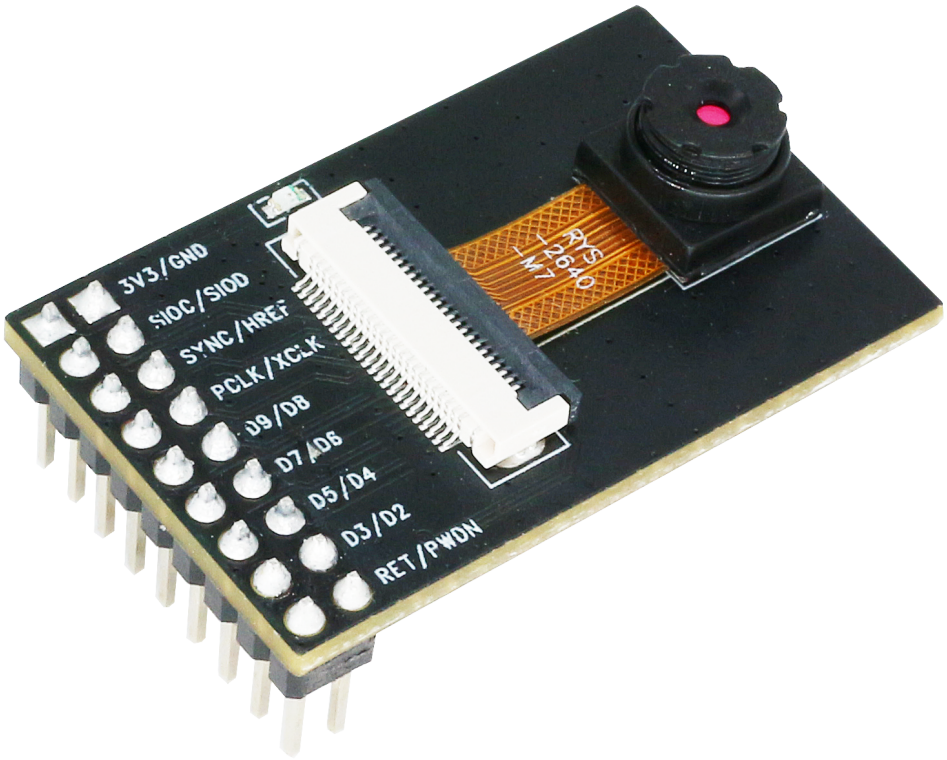


图 24: ESP-LyraP-CAM

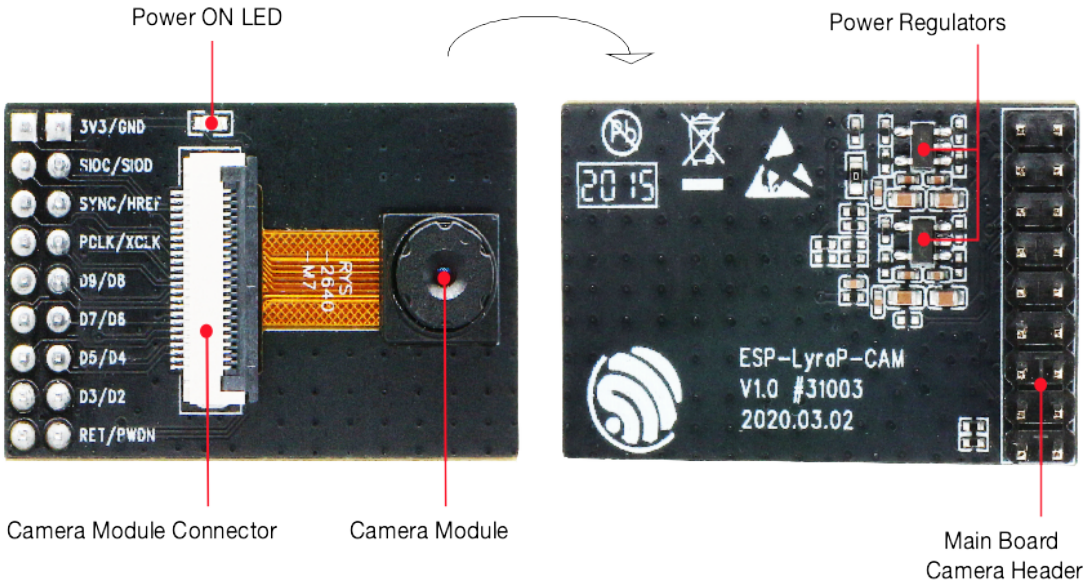


图 25: ESP-LyraP-CAM - 正面和反面

硬件准备

- 带有摄像头扩展板连接器（排母）的主板（例如 ESP32-S2-Kaluga-1）
- ESP-LyraP-CAM 扩展板
- PC（Windows、Linux 或 macOS）

硬件设置 将 ESP-LyraP-CAM 扩展板插入主板的连接头排母中。

软件设置 请前往 ESP32-S2-Kaluga-1 开发套件用户指南的软件设置 章节。

硬件参考

功能框图 ESP-LyraP-CAM 的主要组件和连接方式如下图所示。

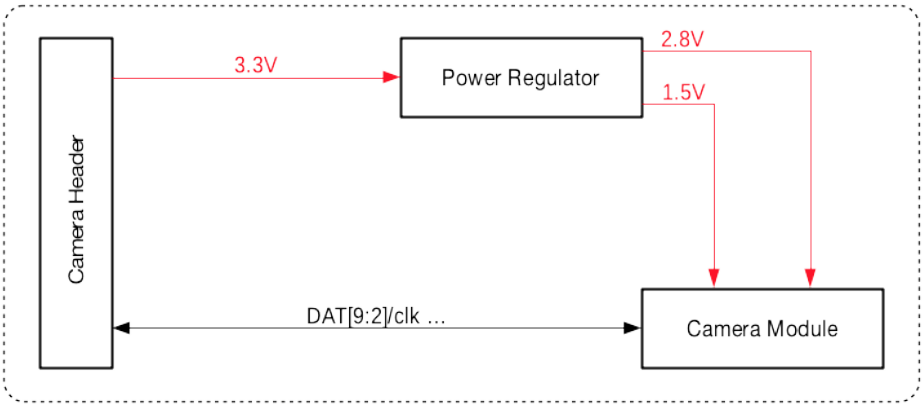


图 26: ESP-LyraP-CAM 功能框图

硬件修订历史 尚无版本升级历史。

相关文档

- ESP-LyraP-CAM 原理图 (PDF)
- ESP-LyraP-CAM PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

ESP-LyraP-LCD32 v1.1 本用户指南可提供 ESP-LyraP-LCD32 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 *主板*，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraP-CAM v1.1 扩展板正在搭配*ESP32-S2-Kaluga-1 套件 v1.2* 销售。

ESP-LyraP-LCD32 可为您的主板增加 LCD 图像显示功能。

本指南包括如下内容：

- 概述：提供为了使用 ESP-LyraP-LCD32 而必须了解的硬件和软件信息。
- 硬件参考：提供 ESP-LyraP-LCD32 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

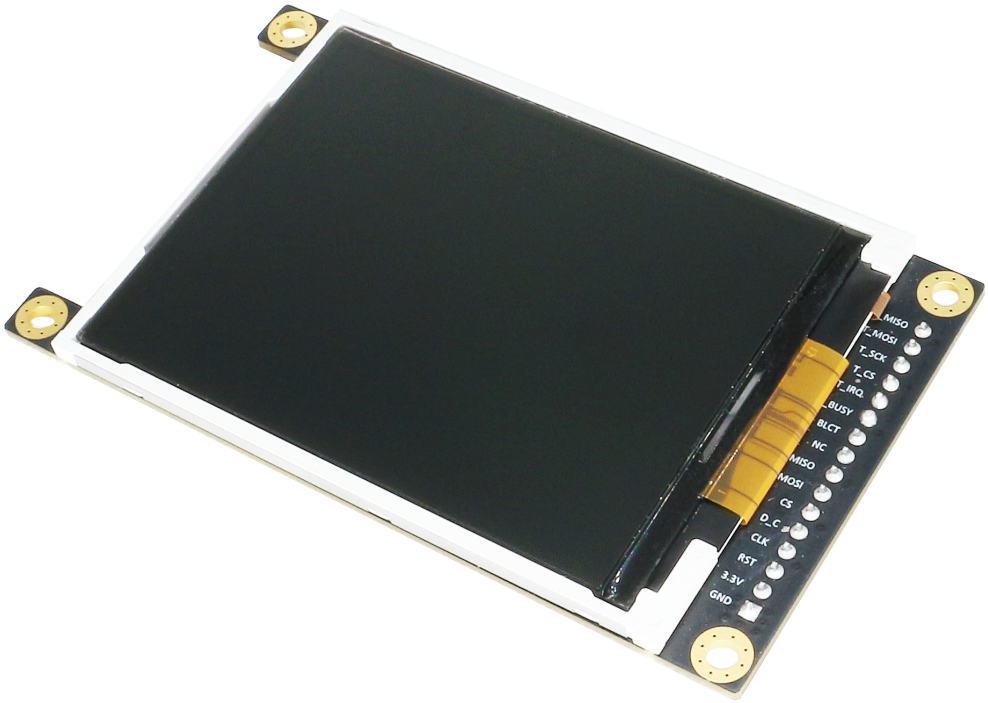


图 27: ESP-LyraP-LCD32（点击放大）

概述 ESP-LyraP-LCD32 可为您的主板增加了一块 3.2” LCD 图形显示屏（320 x 240 分辨率）。该显示屏通过 SPI 总线连接到 ESP32-S2。

组件描述 在下面的组件描述中，**保留**表示该功能可用，但当前版本的套件并未启用该功能。

主要组件	描述
扩展板排针	连接器排针，用于插入主板上的排母
LCD 显示屏	本版本支持 3.2” 的 SPI LCD 显示模块（320 x 240 分辨率）；显示器驱动（控制器）为 Sitronix ST7789V
触摸屏开关	暂不支持触摸屏，因此请注意保持关闭，确保相关管脚复用不受影响。
主板 3.2” LCD FPC 连接器	（保留）连接到主板的 3.2” LCD FPC 连接器
控制开关	打开将 Reset/Backlight_control/CS 设置为默认高电平或低电平；关闭允许释放这些管脚用作它用。

应用程序开发 ESP-LyraP-LCD32 上电前，请首先确认开发板完好无损。

硬件准备

- 带有摄像头扩展板连接器（排母）的主板（例如 ESP32-S2-Kaluga-1、ESP-LyraT-8311A）
- ESP-LyraP-LCD32 扩展板
- 4 x 螺栓，用于保证安装稳定
- PC（Windows、Linux 或 macOS）

硬件设置 请按照以下步骤将 ESP-LyraP-LCD32 安装到带有排母的主板上：

1. 先将 4 个螺栓固定到主板的相应位置上

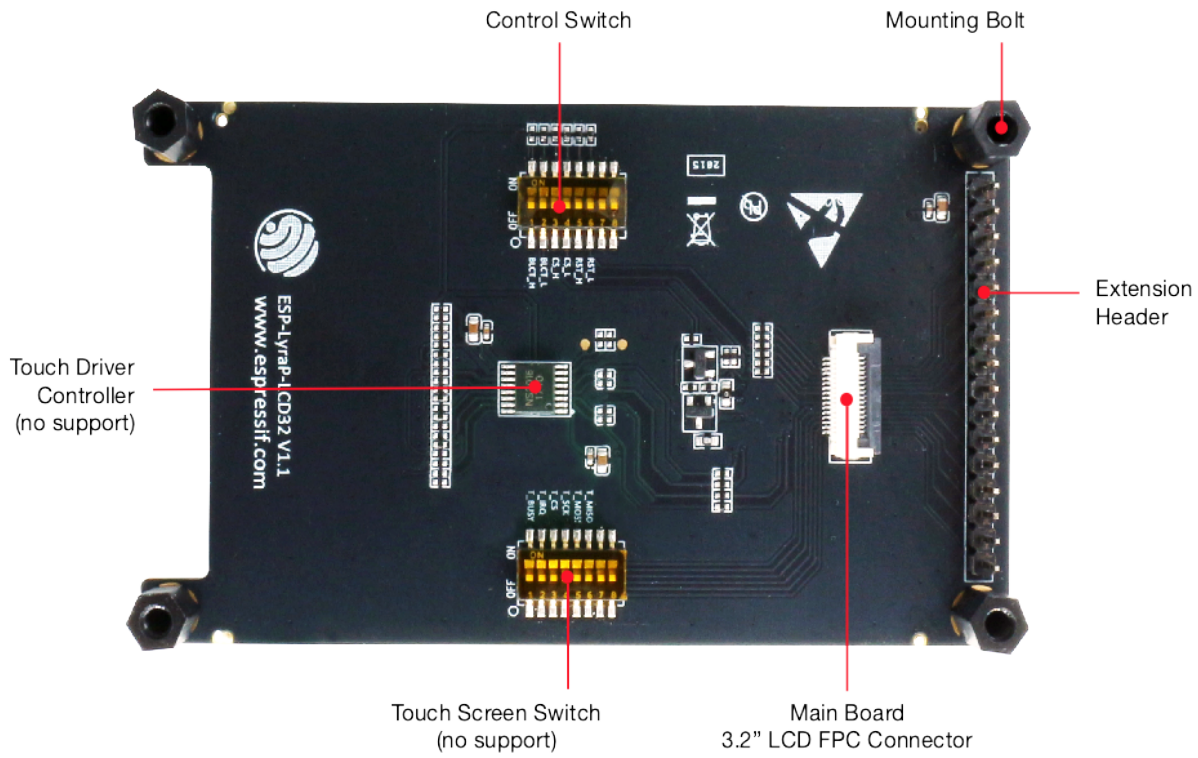


图 28: ESP-LyraP-LCD32 - 正面（点击放大）

2. 对齐 ESP-LyraP-LCD32 与主板和螺栓的位置，并小心插入

软件设置 请前往 ESP32-S2-Kaluga-1 开发套件用户指南的软件设置 章节。

硬件参考

功能框图 ESP-LyraP-LCD32 的主要组件和连接方式如下图所示。

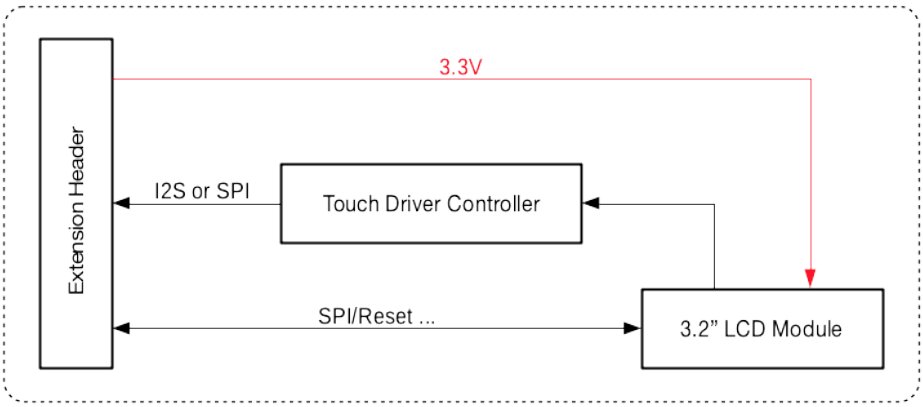


图 29: ESP-LyraP-LCD32 功能框图

硬件修订历史 尚无版本升级历史。

相关文档

- ESP-LyraP-LCD32 原理图 (PDF)
- ESP-LyraP-LCD32 PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

ESP-LyraP-TouchA v1.1 本用户指南可提供 ESP-LyraP-TouchA 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 *主板*，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraP-TouchA v1.1 扩展板正在搭配以下套件销售：

- *ESP32-S2-Kaluga-1 套件 v1.3*
- *ESP32-S2-Kaluga-1 套件 v1.2*

ESP-LyraP-TouchA 可为您的主板增加触摸按键功能。

图 30: ESP-LyraP-TouchA

本指南包括如下内容：

- 概述：提供为了使用 ESP-LyraT-8311A 而必须了解的硬件和软件信息。
- 硬件参考：提供 ESP-LyraP-TouchA 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

概述 ESP-LyraP-TouchA 共有 6 个触摸按钮，主要用于音频应用，但也可以根据实际需要用作它用。

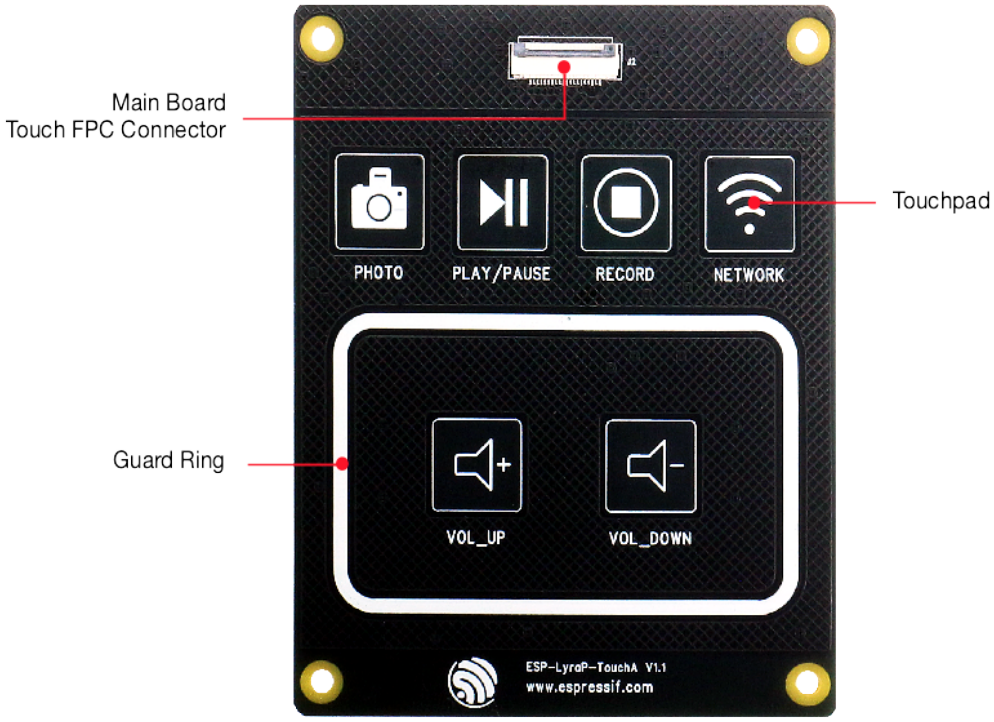


图 31: ESP-LyraP-TouchA

组件描述

主要组件	描述
主板触摸 FPC 连接器	用于将子板连接到主板的触摸 FPC 连接器。
触摸板	电容式触摸电极。
保护环	连接至触摸传感器，可在开发板遇水时触发中断保护（遇水电路保护）。此时，传感器阵列也将遇水，绝大多数（或全部）触摸板将由于大量误触而无法使用。在接收到此中断后，用户可自行裁决是否通过软件禁用所有触摸传感器。

应用程序开发 ESP-LyraP-TouchA 上电前，请首先确认开发板完好无损。

硬件准备

- 带有触摸 FPC 扩展板连接器的主板（例如 ESP32-S2-Kaluga-1）
- ESP-LyraP-TouchA 扩展板
- FPC 线
- PC（Windows、Linux 或 macOS）

硬件设置 使用 FPC 连接两个 FPC 连接器。

软件设置 请前往 ESP32-S2-Kaluga-1 开发套件用户指南的软件设置 章节。

硬件参考

功能框图 ESP-LyraP-TouchA 的主要组件和连接方式如下图所示。

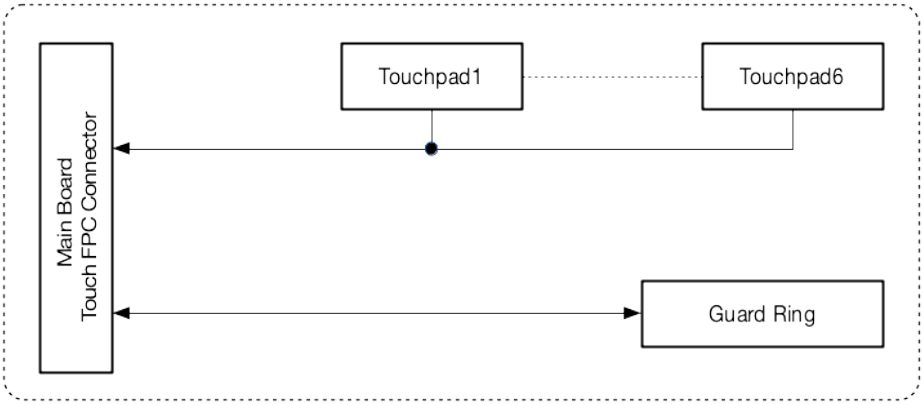


图 32: ESP-LyraP-TouchA-v1.1 功能框图

硬件修订历史 尚无版本升级历史。

相关文档

- ESP-LyraP-TouchA 原理图 (PDF)
- ESP-LyraP-TouchA PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

ESP-LyraT-8311A v1.2 本用户指南可提供 ESP-LyraT-8311A 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 *主板*，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraT-8311A v1.2 扩展板正在搭配*ESP32-S2-Kaluga-1 套件 v1.2* 销售。

ESP-LyraT-8311A 扩展板可为您的主板增加音频处理功能。

- 音频播放/录音
- 音频信号处理
- 支持可编程按钮，可实现轻松控制

ESP-LyraT-8311A 扩展板有多种使用方式。该应用程序包括语音用户界面、语音控制、语音授权、录音和播放等功能。

本指南包括如下内容：

- *概述*：提供为了使用 ESP-LyraT-8311A 而必须了解的硬件和软件信息。
- *硬件参考*：提供 ESP-LyraT-8311A 的详细硬件信息。
- *硬件修订历史*：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- *相关文档*：提供相关文档的链接。

概述 ESP-LyraT-8311A 主要用于音频应用，但也可根据实际需求用作它用。

组件描述 下表将从图片右上角开始，以顺时针顺序介绍上图中的主要组件。

保留表示该功能可用，但当前版本的套件并未启用该功能。

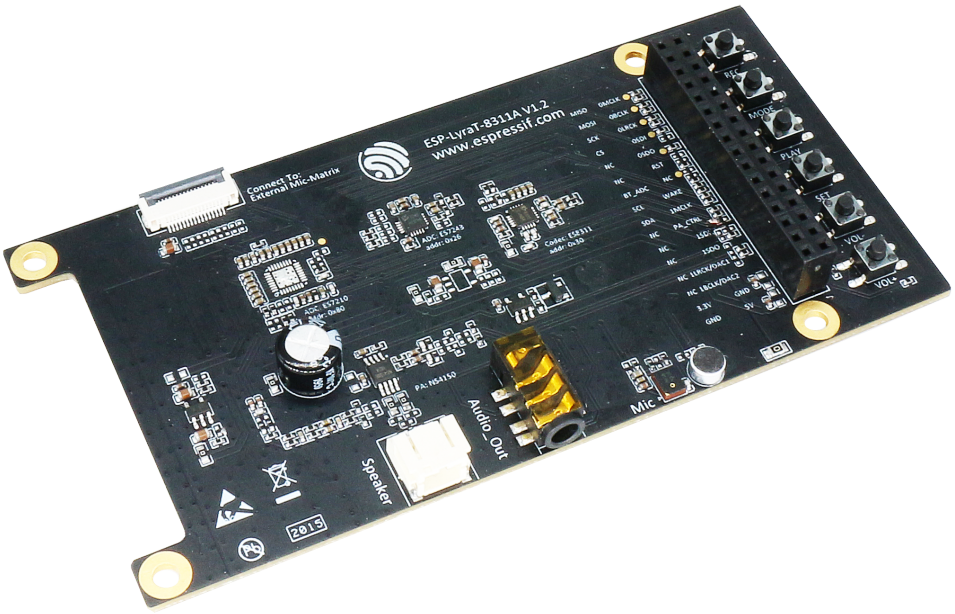


图 33: ESP-LyraT-8311A（点击放大）

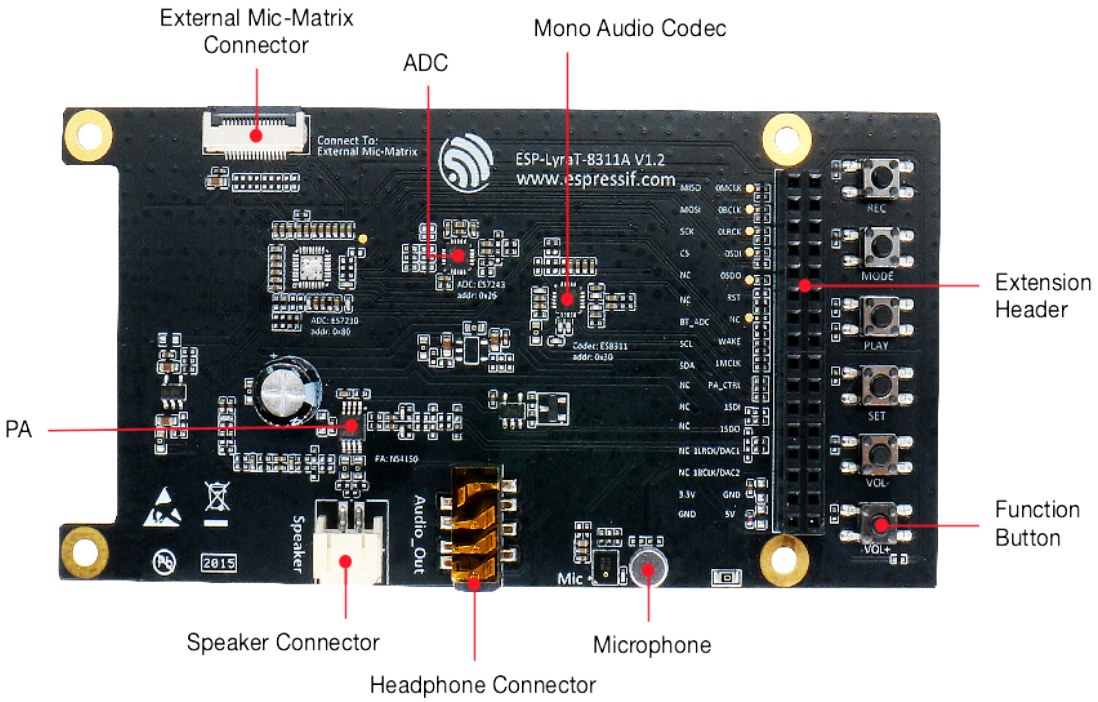


图 34: ESP-LyraT-8311A - 正面（点击放大）

主要组件	描述
扩展板排针	反面的排针用于于主板上的排母相连；排母用于连接其他带有排针的主板
功能按钮	带有 6 个可编程按钮
麦克风	支持驻极体和 MEMS 麦克风；此扩展板默认带有驻极体麦克风
耳机接口	6.3 mm (1/8”) 立体声耳机接口
扬声器连接器	2 针连接器，用于连接外部扬声器
PA	3 W 音频信号放大器，配合外部扬声器使用
外部麦克风矩阵连接器	（保留）FPC 连接器，用于连接外部麦克风矩阵（麦克风开发板）
ADC	（保留）高性能 ADC/ES7243，包括 1 个麦克风通道、1 个声学回声消除 (AEC) 功能通道
单声道音频编解器	ES8311 音频 ADC 和 DAC，可转换麦克风拾音的模拟信号；或转换数字信号，使其可通过扬声器或耳机进行播放

应用程序开发 ESP-LyraT-8311A 上电前，请首先确认开发板完好无损。

硬件准备

- 带有连接器（排母）的主板（例如 ESP32-S2-Kaluga-1）
- ESP-LyraT-8311A 扩展板
- 4 x 螺栓，用于保证安装稳定
- PC（Windows、Linux 或 macOS）

硬件设置 请按照以下步骤将 ESP-LyraT-8311A 安装到带有排母的主板上：

1. 先将 4 个螺栓固定到主板的相应位置上
2. 对齐 ESP-LyraT-8311A 与主板和螺栓的位置，并小心插入

软件设置 请根据您的具体应用，参考以下部分：

- ESP-ADF（乐鑫音频开发框架）的用户，请前往 ESP-ADF 入门指南。
- ESP32-IDF（乐鑫 IoT 开发框架）的用户，请前往 ESP32-S2-Kaluga-1 开发套件用户指南软件设置 章节。

硬件参考

功能框图 ESP-LyraT-8311A 的主要组件和连接方式如下图所示。

硬件修订历史 尚无版本升级历史。

相关文档

- ESP-LyraT-8311A 原理图 (PDF)
- ESP-LyraT-8311A PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

- ESP32-S2-WROVER 技术规格书 (PDF)
- 乐鑫产品订购信息 (PDF)
- *JTAG 调试*
- ESP32-S2-Kaluga-1 原理图 (PDF)
- ESP32-S2-Kaluga-1 PCB 布局图 (PDF)
- ESP32-S2-Kaluga-1 管脚映射 (Excel)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

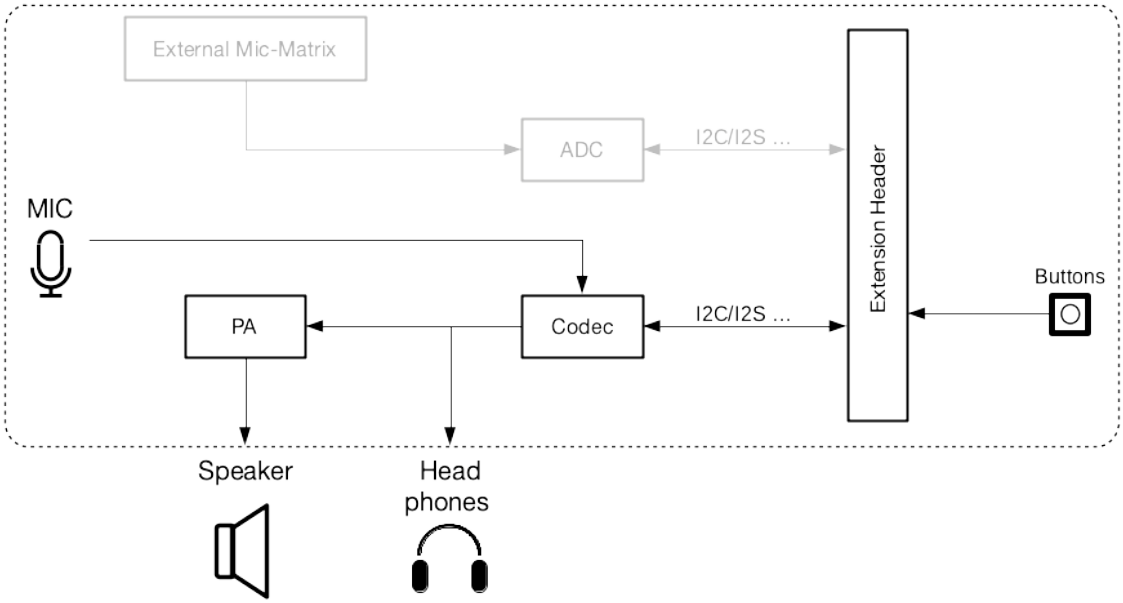


图 35: ESP-LyraT-8311A 功能框图

ESP-LyraP-CAM v1.1 本用户指南可提供 ESP-LyraP-CAM 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 *主板*，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraP-CAM v1.1 扩展板正在搭配*ESP32-S2-Kaluga-1 套件 v1.3* 销售。

ESP-LyraP-CAM 可为您的主板增加摄像头功能。

本指南包括如下内容：

- 概述：提供为了使用 ESP-LyraP-CAM 而必须了解的硬件和软件信息。
- 硬件参考：提供 ESP-LyraP-CAM 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

概述 ESP-LyraP-CAM 扩展板可为您的主板增加一个摄像头。

组件描述

主要组件	描述
主板摄像头排针	连接至主板摄像头连接器
电源 LED 指示灯	如果电源供电电压正常，则红色 LED 亮起
摄像头模块连接器	硬件支持 OV2640 和 OV3660 摄像头模块；目前，ESP-LyraP-CAM 扩展板默认提供 OV2640 摄像头模块
电源调节器	LDO 调压器（3.3 V 至 2.8 V 和 1.5 V）

应用程序开发 ESP-LyraP-CAM 上电前，请首先确认开发板完好无损。

硬件准备

- 带有摄像头扩展板连接器（排母）的主板（例如 ESP32-S2-Kaluga-1）
- ESP-LyraP-CAM 扩展板
- PC（Windows、Linux 或 macOS）

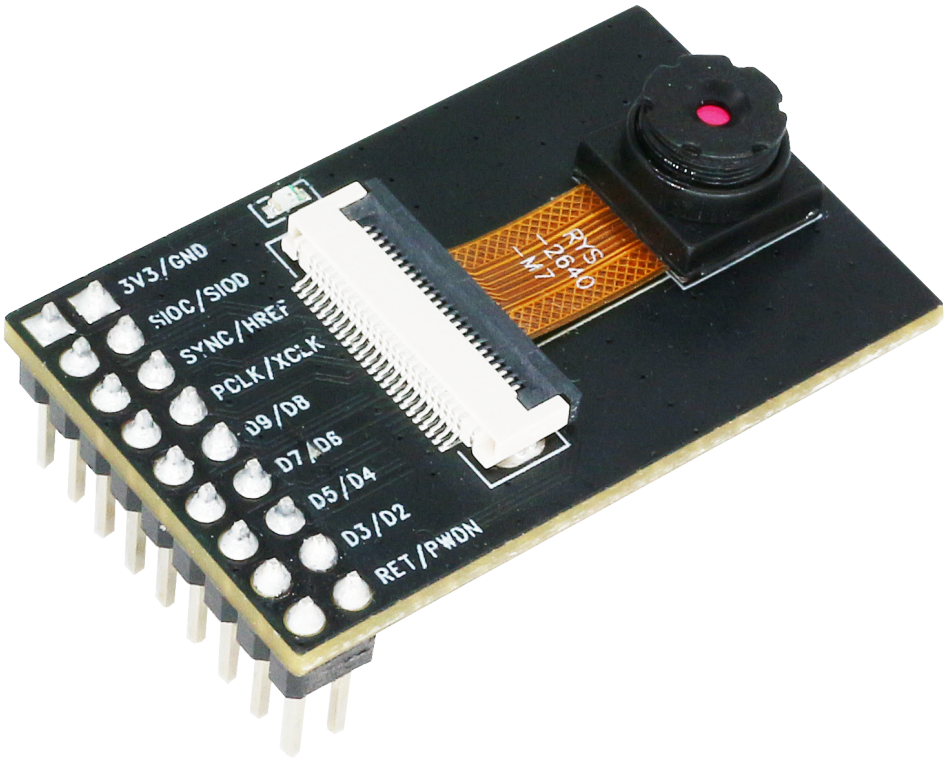


图 36: ESP-LyraP-CAM

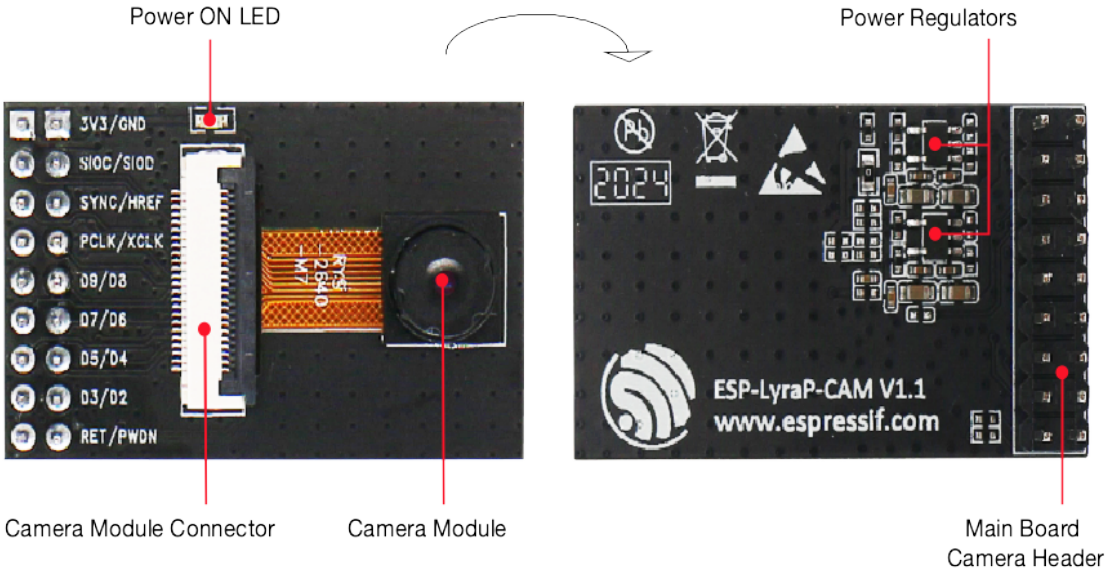


图 37: ESP-LyraP-CAM - 正面和反面

硬件设置 将 ESP-LyraP-CAM 扩展板插入主板的连接头排母中。

软件设置 请前往 ESP32-S2-Kaluga-1 开发套件用户指南的软件设置 章节。

硬件参考

功能框图 ESP-LyraP-CAM 的主要组件和连接方式如下图所示。

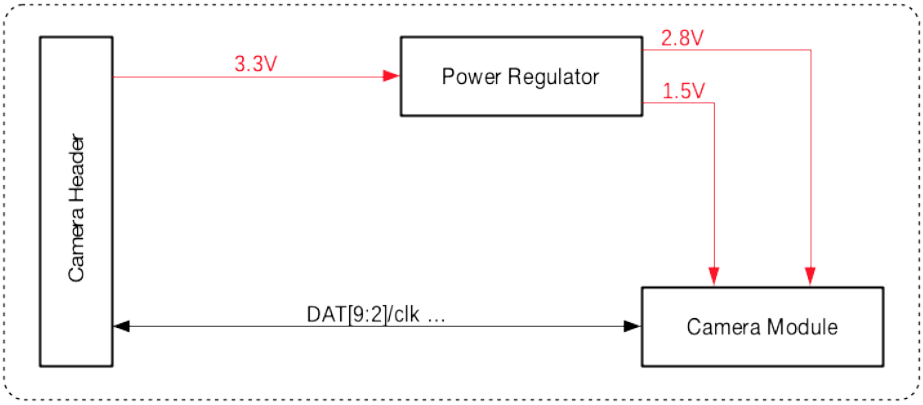


图 38: ESP-LyraP-CAM 功能框图

硬件修订历史

ESP-LyraP-CAM v1.1

- 仅更新丝印
- 无实际硬件升级

ESP-LyraP-CAM v1.0 首次发布

相关文档

- ESP-LyraP-CAM 原理图 (PDF)
- ESP-LyraP-CAM PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

ESP-LyraP-LCD32 v1.2 本用户指南可提供 ESP-LyraP-LCD32 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 *主板*，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraP-LCD32 v1.2 扩展板正在搭配*ESP32-S2-Kaluga-1 套件 v1.3* 销售。

ESP-LyraP-LCD32 可为您的主板增加 LCD 图像显示功能。

本指南包括如下内容：

- 概述：提供为了使用 ESP-LyraP-LCD32 而必须了解的硬件和软件信息。
- 硬件参考：提供 ESP-LyraP-LCD32 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

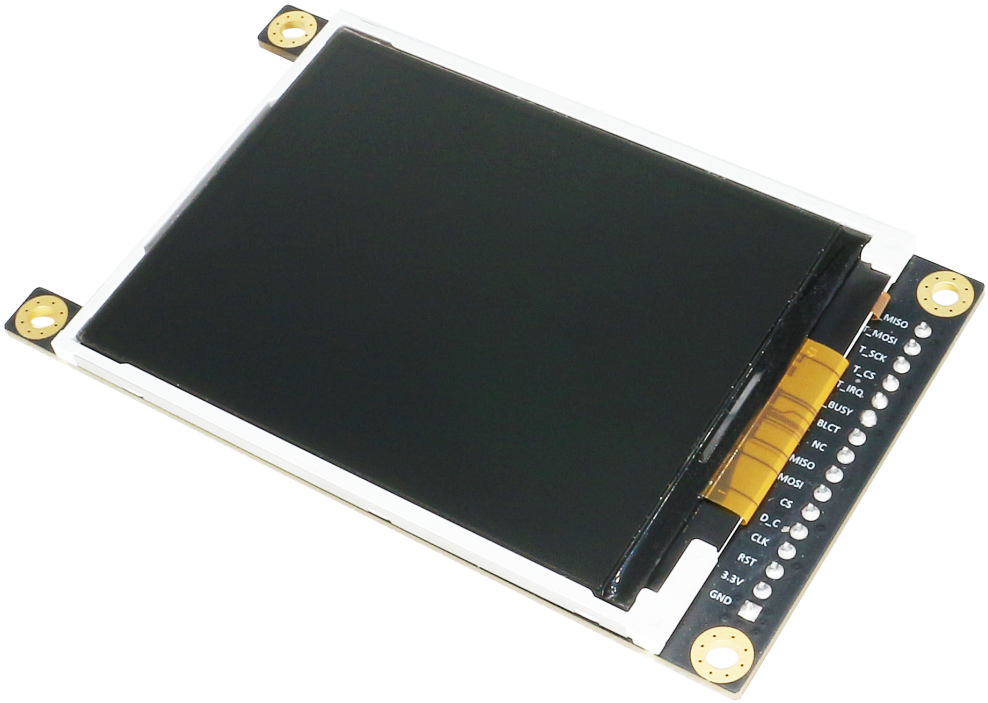


图 39: ESP-LyraP-LCD32（点击放大）

概述 ESP-LyraP-LCD32 可为您的主板增加了一块 3.2” LCD 图形显示屏（320 x 240 分辨率）。该显示屏通过 SPI 总线连接到 ESP32-S2。

组件描述 在下面的组件描述中，**保留**表示该功能可用，但当前版本的套件并未启用该功能。

主要组件	描述
扩展板排针	连接器排针，用于插入主板上的排母
LCD 显示屏	本版本支持 3.2” 的 SPI LCD 显示模块（320 x 240 分辨率）；显示器驱动（控制器）为 Sitronix ST7789V 或 Ilitek ILI9341
触摸屏开关	暂不支持触摸屏，因此请注意保持关闭，确保相关管脚复用不受影响。
主板 3.2” LCD FPC 连接器	（保留）连接到主板的 3.2” LCD FPC 连接器
控制开关	打开将 Reset/Backlight_control/CS 设置为默认高电平或低电平；关闭允许释放这些管脚用作它用。

应用程序开发 ESP-LyraP-LCD32 上电前，请首先确认开发板完好无损。

硬件准备

- 带有摄像头扩展板连接器（排母）的主板（例如 ESP32-S2-Kaluga-1、ESP-LyraT-8311A）
- ESP-LyraP-LCD32 扩展板
- 4 x 螺栓，用于保证安装稳定
- PC（Windows、Linux 或 macOS）

硬件设置 请按照以下步骤将 ESP-LyraP-LCD32 安装到带有排母的主板上：

1. 先将 4 个螺栓固定到主板的相应位置上

图 40: ESP-LyraP-LCD32 - 正面（点击放大）

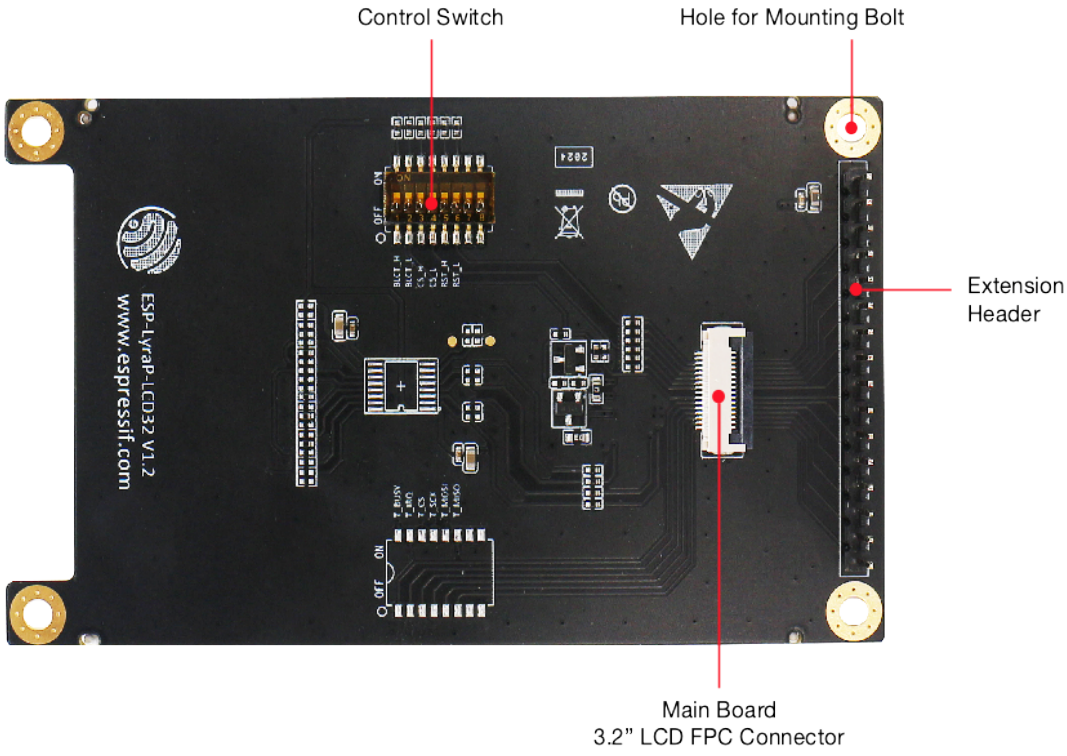


图 41: ESP-LyraP-LCD32 - 反面（点击放大）

2. 对齐 ESP-LyraP-LCD32 与主板和螺栓的位置，并小心插入

软件设置 请前往 ESP32-S2-Kaluga-1 开发套件用户指南的软件设置 章节。

硬件参考

功能框图 ESP-LyraP-LCD32 的主要组件和连接方式如下图所示。

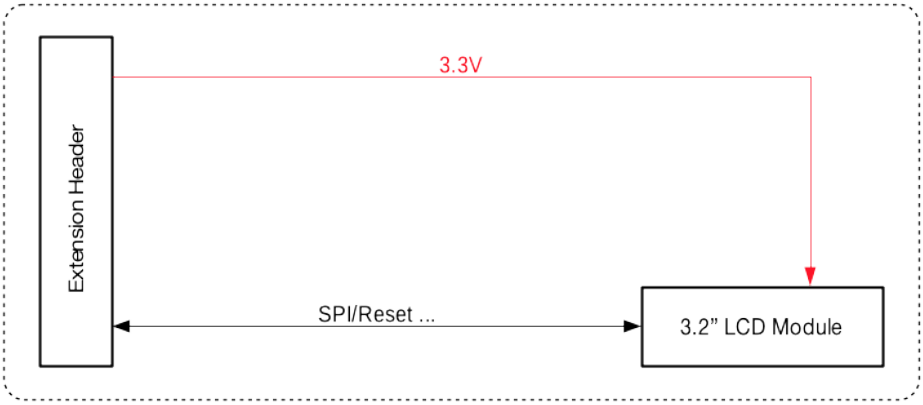


图 42: ESP-LyraP-LCD32 功能框图

硬件修订历史

ESP-LyraP-LCD32 v1.2

- LCD 背光默认打开 (ON)，无法通过 MCU 实现控制
- 移除 Touch 驱动及相关开关，以避免管脚复用带来的影响

ESP-LyraP-LCD32 v1.1 首次发布

相关文档

- ESP-LyraP-LCD32 原理图 (PDF)
- ESP-LyraP-LCD32 PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

ESP-LyraT-8311A v1.3 本用户指南可提供 ESP-LyraT-8311A 扩展板的相关信息。

本扩展板通常仅与乐鑫其他开发板一起销售（即 主板，比如 ESP32-S2-Kaluga-1），不可单独购买。

目前，ESP-LyraT-8311A v1.3 扩展板正在搭配*ESP32-S2-Kaluga-1 套件 v1.3* 销售。

ESP-LyraT-8311A 扩展板可为您的主板增加音频处理功能。

- 音频播放/录音
- 音频信号处理
- 支持可编程按钮，可实现轻松控制

ESP-LyraT-8311A 扩展板有多种使用方式。该应用程序包括语音用户界面、语音控制、语音授权、录音和播放等功能。

本指南包括如下内容：

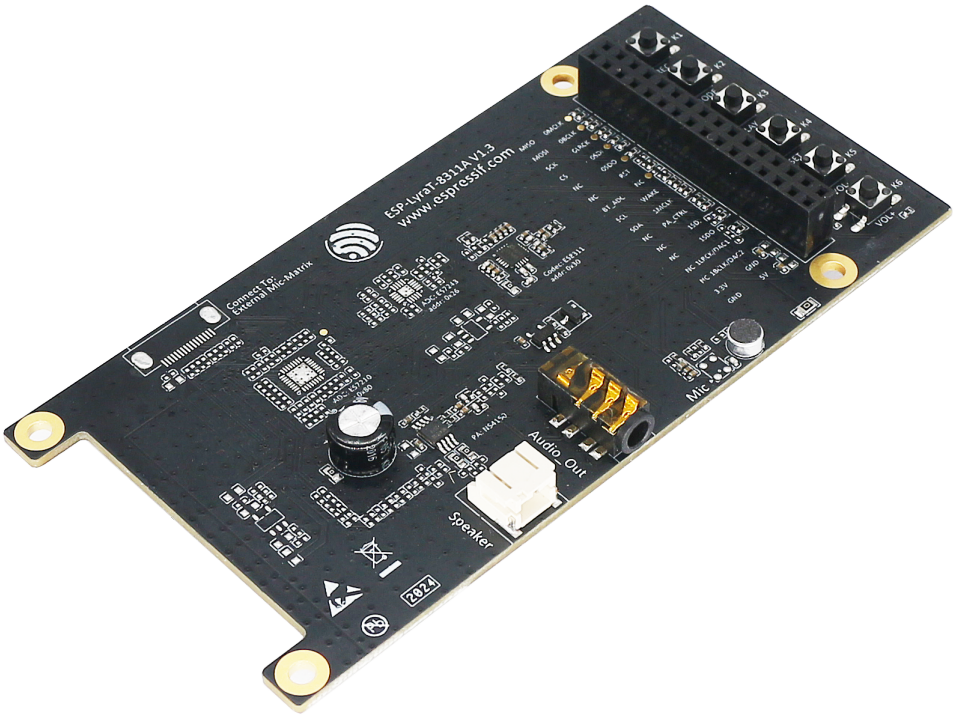


图 43: ESP-LyraT-8311A (click to enlarge)

- 概述：提供为了使用 ESP-LyraT-8311A 而必须了解的硬件和软件信息。
- 硬件参考：提供 ESP-LyraT-8311A 的详细硬件信息。
- 硬件修订历史：提供该开发版的“修订历史”、“已知问题”以及此前版本开发板的用户指南链接。
- 相关文档：提供相关文档的链接。

概述 ESP-LyraT-8311A 主要用于音频应用，但也可根据实际需求用作它用。

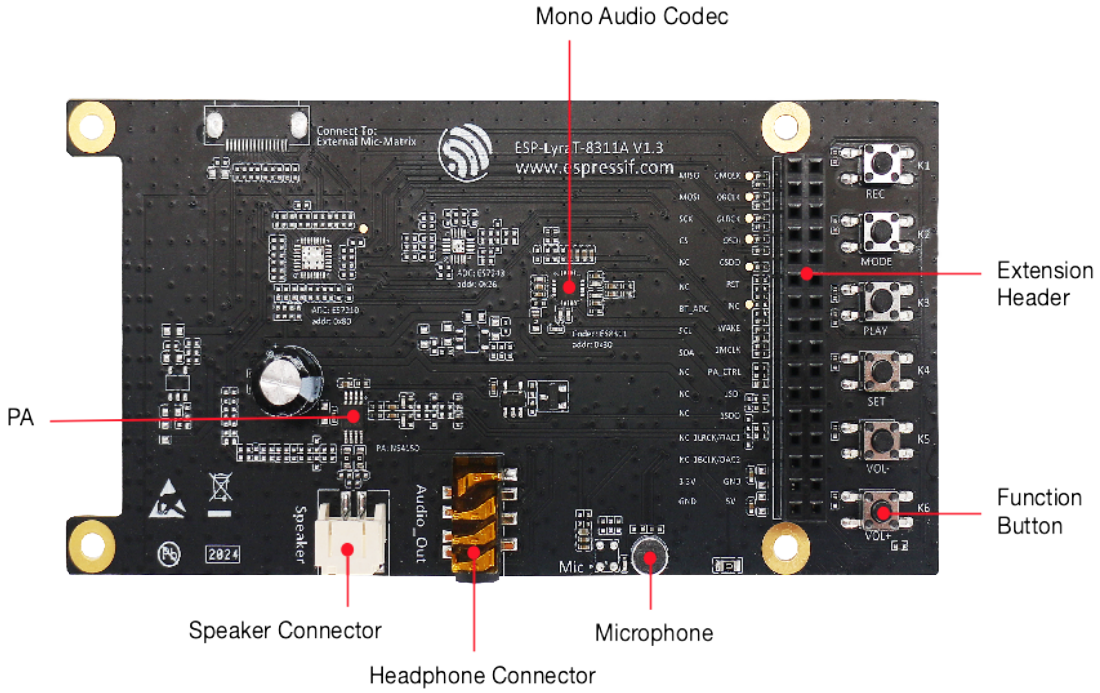


图 44: ESP-LyraT-8311A - 正面（点击放大）

组件描述 下表将从图片右上角开始，以顺时针顺序介绍上图中的主要组件。

保留表示该功能可用，但当前版本的套件并未启用该功能。

主要组件	描述
扩展板排针	反面的排针用于于主板上的排母相连；排母用于连接其他带有排针的主板
功能按钮	带有 6 个可编程按钮
麦克风	支持驻极体和 MEMS 麦克风；此扩展板默认带有驻极体麦克风
耳机接口	6.3 mm (1/8”) 立体声耳机接口
扬声器连接器	2 针连接器，用于连接外部扬声器
PA	3 W 音频信号放大器，配合外部扬声器使用
外部麦克风矩阵连接器	（保留）FPC 连接器，用于连接外部麦克风矩阵（麦克风开发板）
ADC	（保留）高性能 ADC/ES7243，包括 1 个麦克风通道、1 个声学回声消除 (AEC) 功能通道
单声道音频编解器	ES8311 音频 ADC 和 DAC，可转换麦克风拾音的模拟信号；或转换数字信号，使其可通过扬声器或耳机进行播放

应用程序开发 ESP-LyraT-8311A 上电前，请首先确认开发板完好无损。

硬件准备

- 带有连接器（排母）的主板（例如 ESP32-S2-Kaluga-1）
- ESP-LyraT-8311A 扩展板
- 4 x 螺栓，用于保证安装稳定
- PC（Windows、Linux 或 macOS）

硬件设置 请按照以下步骤将 ESP-LyraT-8311A 安装到带有排母的主板上：

1. 先将 4 个螺栓固定到主板的相应位置上
2. 对齐 ESP-LyraT-8311A 与主板和螺栓的位置，并小心插入

软件设置 请根据您的具体应用，参考以下部分：

- ESP-ADF（乐鑫音频开发框架）的用户，请前往 ESP-ADF 入门指南。
- ESP32-IDF（乐鑫 IoT 开发框架）的用户，请前往 ESP32-S2-Kaluga-1 开发套件用户指南软件设置 章节。

硬件参考

功能框图 ESP-LyraT-8311A 的主要组件和连接方式如下图所示。

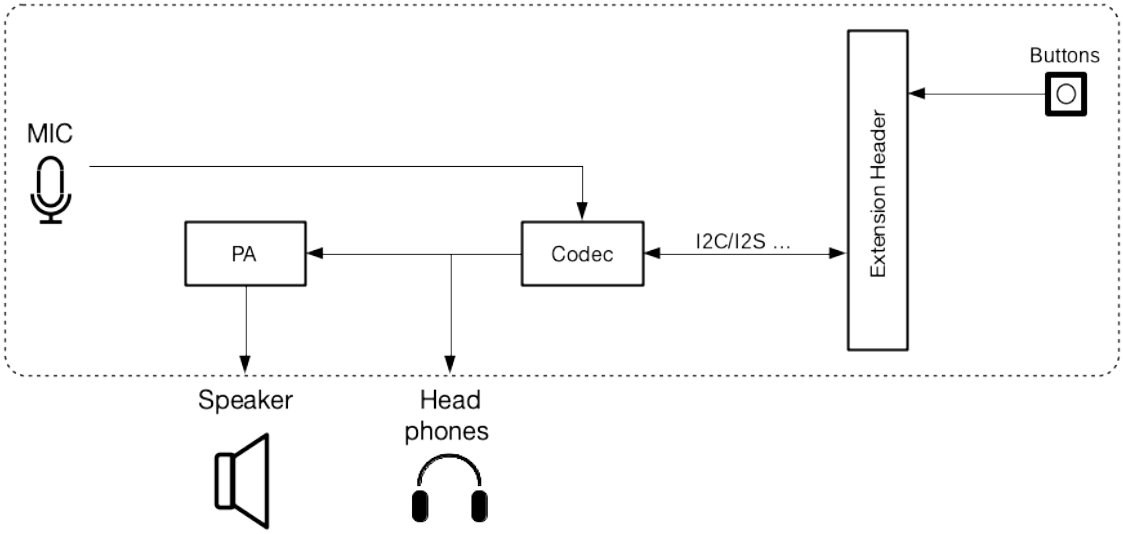


图 45: ESP-LyraT-8311A 功能框图

硬件修订历史

ESP-LyraT-8311A v1.3

- 移除 ADC/ES7243 和 ADC/ES7210，相关功能由单声道音频编解器提供。

ESP-LyraT-8311A v1.2 首次发布

相关文档

- ESP-LyraT-8311A 原理图 (PDF)
- ESP-LyraT-8311A PCB 布局图 (PDF)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

- ESP32-S2-WROVER 技术规格书 (PDF)
- 乐鑫产品订购信息 (PDF)
- *JTAG* 调试
- ESP32-S2-Kaluga-1 原理图 (PDF)
- ESP32-S2-Kaluga-1 PCB 布局图 (PDF)
- ESP32-S2-Kaluga-1 管脚映射 (Excel)

有关本开发板的更多设计文档，请联系我们的商务部门 sales@espressif.com。

1.4 详细安装步骤

请根据下方详细步骤，完成安装过程。

1.4.1 设置开发环境

- 第一步：安装准备 （*Windows* 、*Linux* 和*macOS*）
- 第二步：获取 *ESP-IDF*
- 第三步：设置工具
- 第四步：设置环境变量

1.4.2 创建您的第一个工程

- 第五步：开始创建工程
- 第六步：连接设备
- 第七步：配置
- 第八步：编译工程
- 第九步：烧录到设备
- 第十步：监视器

1.5 第一步：安装准备

在正式开始创建工程前，请先完成工具的安装，具体步骤见下：

1.5.1 Windows 平台工具链的标准设置

概述

ESP-IDF 需要安装一些必备工具，才能围绕 ESP32-S2 构建固件，包括 Python、Git、交叉编译器、CMake 和 Ninja 编译工具等。

在本入门指南中，我们通过 **命令提示符**进行有关操作。不过，您在安装 ESP-IDF 后还可以使用*Eclipse* 或其他支持 CMake 的图形化工具 IDE。

注解: 限定条件：Python 或 ESP-IDF 的安装路径中一定不能包含空格或括号。与此同时，除非操作系统配置为支持 Unicode UTF-8，否则 Python 或 ESP-IDF 的安装路径中也不能包括特殊字符（非 ASCII 码字符）

系统管理员可以通过如下方式将操作系统配置为支持 Unicode UTF-8：控制面板-更改日期、时间或数字格式-管理选项卡-更改系统地域-勾选选项“Beta：使用 Unicode UTF-8 支持全球语言”-点击确定-重启电脑。

ESP-IDF 工具安装器

安装 ESP-IDF 必备工具最简易的方式是从 https://dl.espressif.com/dl/esp-idf/?idf=4.2 中下载 ESP-IDF 工具安装器。

在线安装与离线安装的区别 在线安装程序非常小，可以安装 ESP-IDF 的所有版本。在安装过程中，安装程序只下载必要的依赖文件，包括 Git For Windows 安装器。在线安装程序会将下载的文件存储在缓存目录 `%userprofile%/espressif` 中。

离线安装程序不需要任何网络连接。安装程序中包含了所有需要的依赖文件，包括 Git For Windows 安装器。

安装内容 安装程序会安装以下组件：

- 内置的 Python
- 交叉编译器
- OpenOCD
- CMake 和 Ninja 编译工具
- ESP-IDF

安装程序允许将程序下载到现有的 ESP-IDF 目录。推荐将 ESP-IDF 下载到 `%userprofile%\Desktop\esp-idf` 目录下，其中 `%userprofile%` 代表家目录。

启动 ESP-IDF 环境 安装结束时，如果勾选了 `Run ESP-IDF PowerShell Environment` 或 `Run ESP-IDF Command Prompt (cmd.exe)`，安装程序会在选定的提示符窗口启动 ESP-IDF。

`Run ESP-IDF PowerShell Environment`:

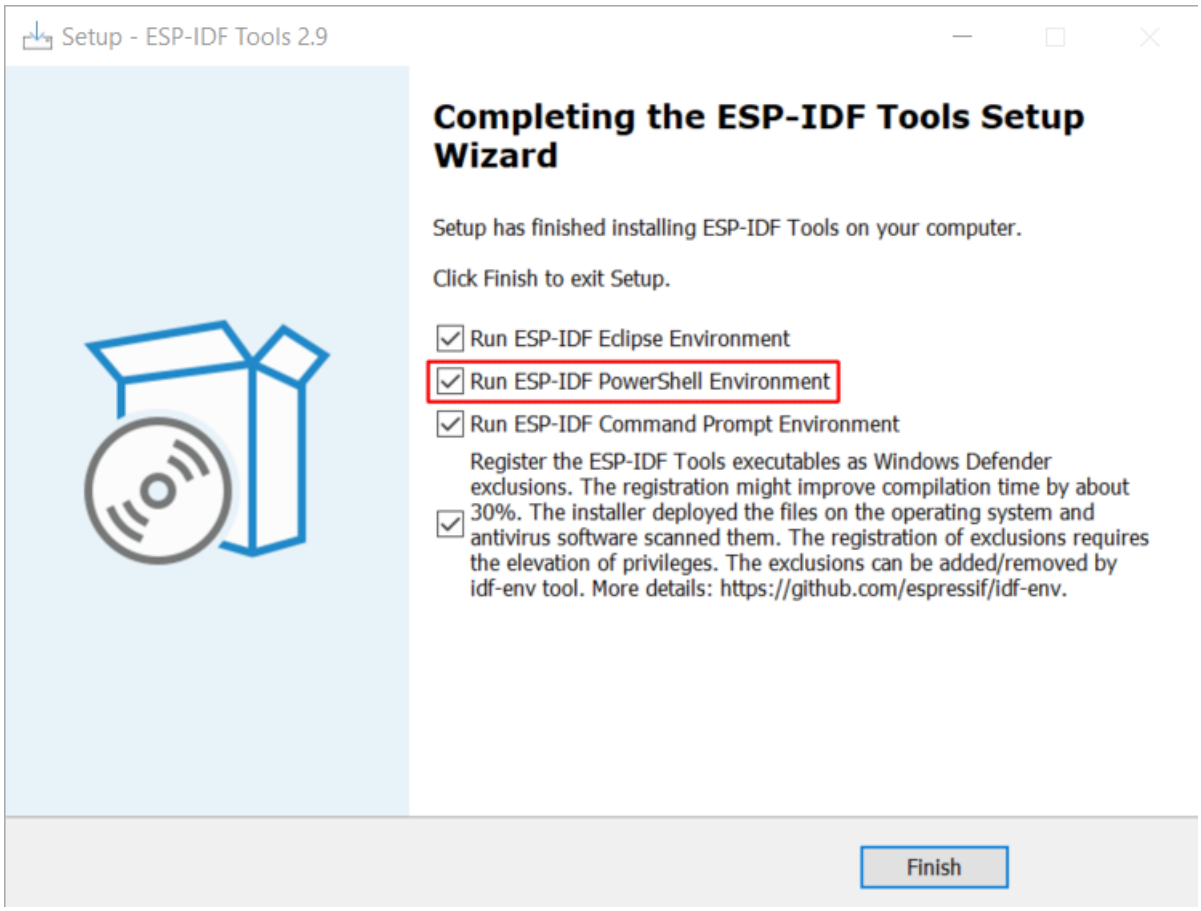


图 46: 完成 ESP-IDF 工具安装向导时运行 Run ESP-IDF PowerShell Environment

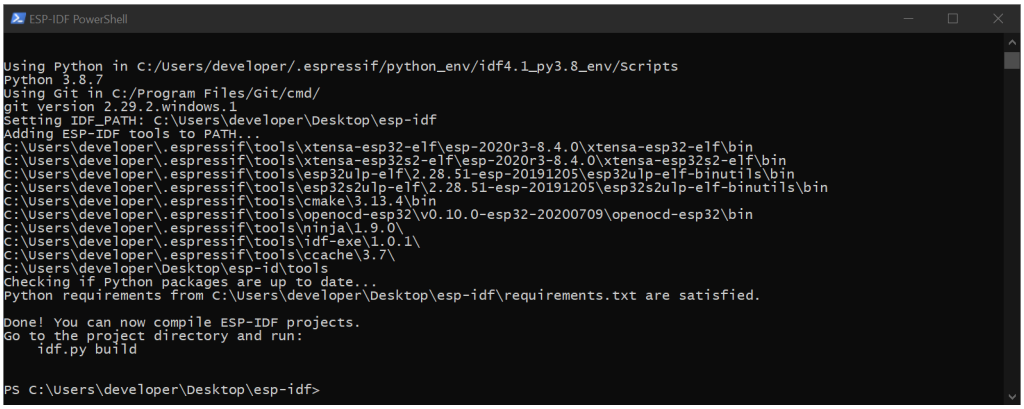


图 47: ESP-IDF PowerShell

`Run ESP-IDF Command Prompt (cmd.exe):`

使用命令提示符

在后续步骤中，我们将使用 Windows 的命令提示符进行操作。

ESP-IDF 工具安装器可在“开始”菜单中，创建一个打开 ESP-IDF 命令提示符窗口的快捷方式。本快捷方式可以打开 Windows 命令提示符（即 cmd.exe），并运行 `export.bat` 脚本以设置各环境变量（比如 `PATH`，`IDF_PATH` 等）。此外，您可还以通过 Windows 命令提示符使用各种已经安装的工具。

注意，本快捷方式仅适用 ESP-IDF 工具安装器中指定的 ESP-IDF 路径。如果您的电脑上存在多个 ESP-IDF 路径（比如您需要不同版本的 ESP-IDF），您有以下两种解决方法：

1. 为 ESP-IDF 工具安装器创建的快捷方式创建一个副本，并将新快捷方式的 ESP-IDF 工作路径指定为您希望使用的 ESP-IDF 路径。
2. 或者，您可以运行 `cmd.exe`，并切换至您希望使用的 ESP-IDF 目录，然后运行 `export.bat`。注意，这种方法要求 `PATH` 中存在 Python 和 Git。如果您在使用时遇到有关“找不到 Python 或 Git”的错误信息，请使用第一种方法。

后续步骤

当 ESP-IDF 工具安装器安装成功后，开发环境设置也到此结束。后续开发步骤，请前往第五步：开始创建工程查看。

相关文档

想要自定义安装流程的高阶用户可参照：

在 Windows 环境下更新 ESP-IDF 工具

使用脚本安装 ESP-IDF 工具 请从 Windows “命令提示符”窗口，切换至 ESP-IDF 的安装目录。然后运行：

```
install.bat
```

对于 Powershell，请切换至 ESP-IDF 的安装目录。然后运行：

```
install.ps1
```

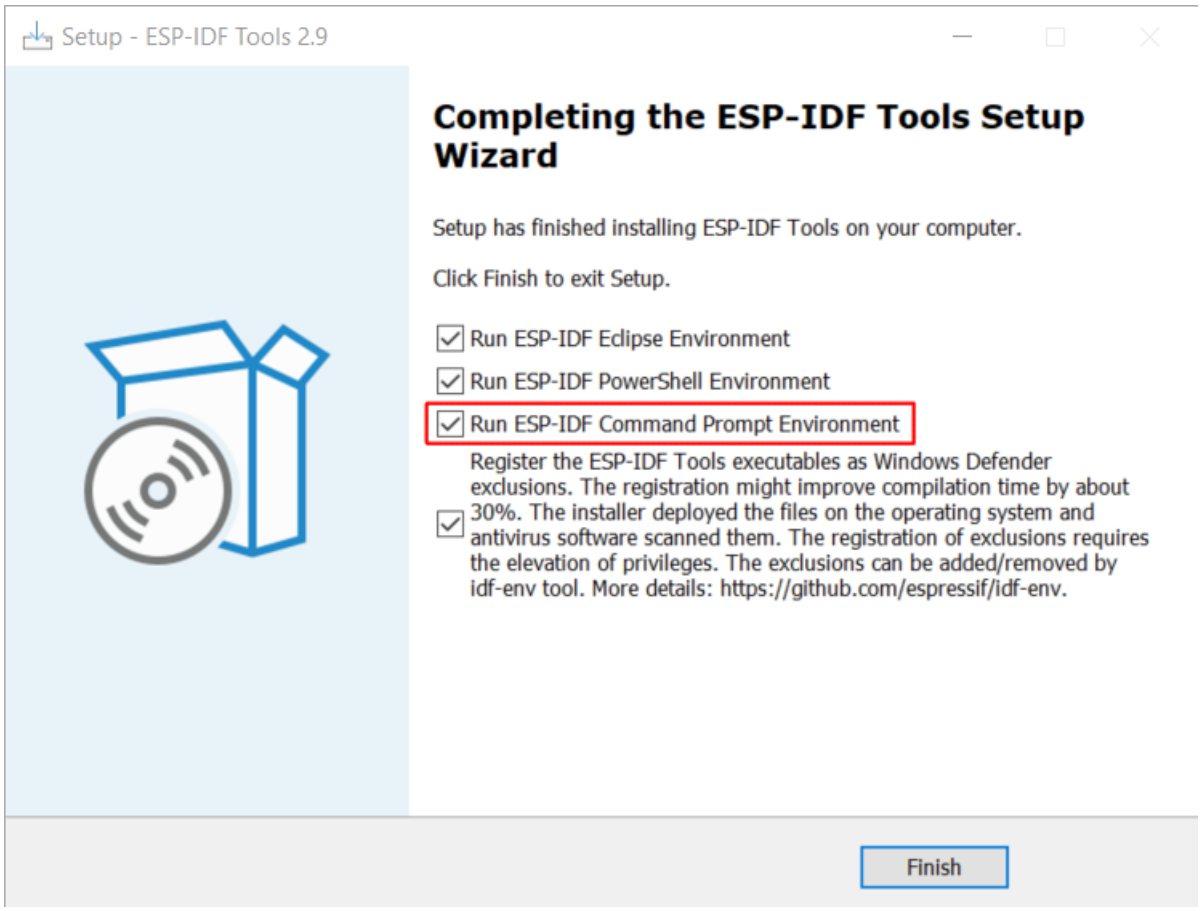


图 48: 完成 ESP-IDF 工具安装向导时运行 Run ESP-IDF Command Prompt (cmd.exe)

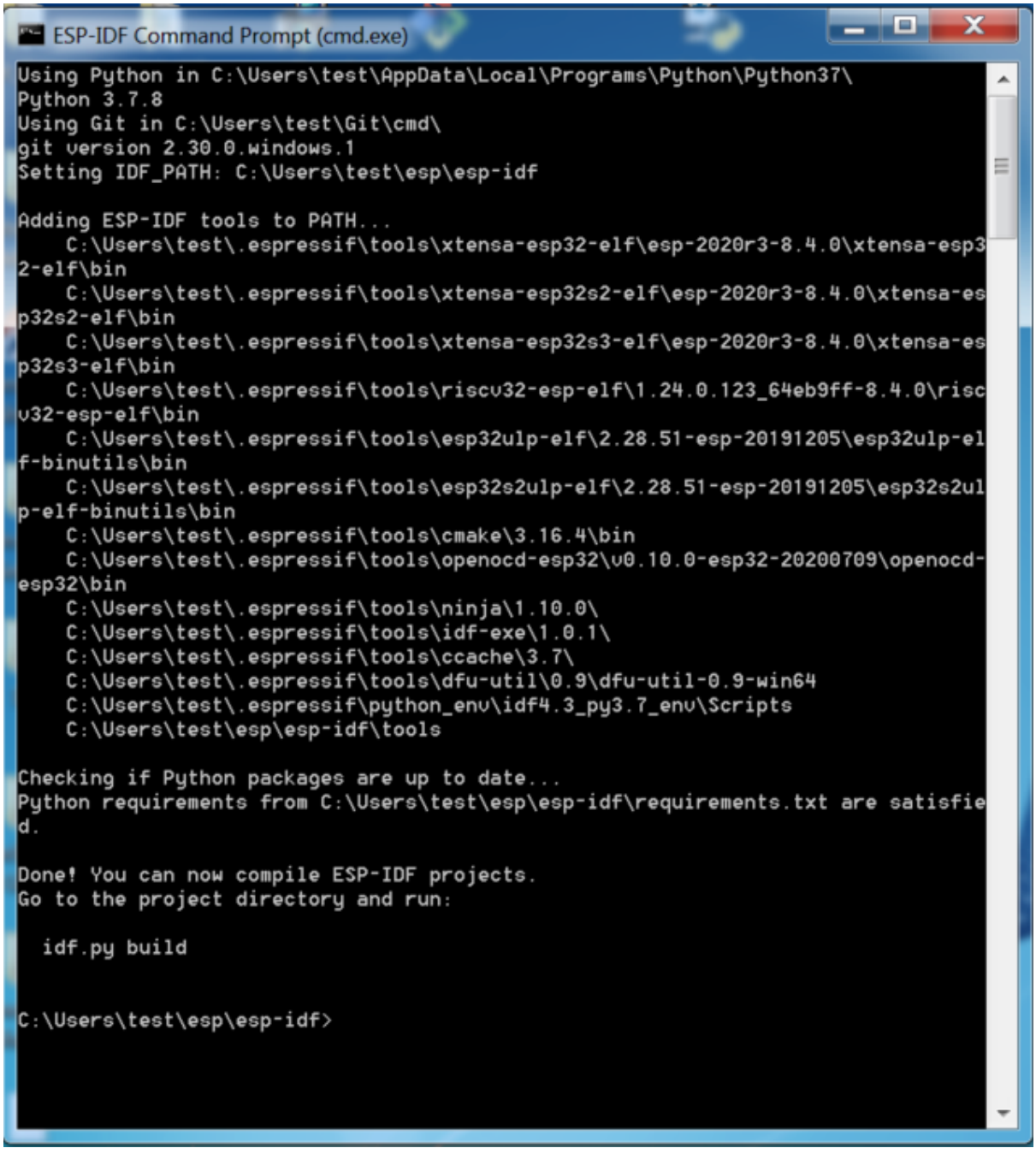


图 49: ESP-IDF 命令提示符窗口

该命令可下载安装 ESP-IDF 所需的工具。如您已经安装了某个版本的工具，则该命令将无效。

该工具的下载安装位置由 ESP-IDF 工具安装器的设置决定，默认情况下为：`C:\Users\username\.espressif`。

使用“导出脚本”将 ESP-IDF 工具添加至 PATH ESP-IDF 工具安装器将在“开始菜单”为“ESP-IDF 命令提示符”创建快捷方式。点击该快捷方式可打开 Windows 命令提示符窗口，您可在该窗口使用所有已安装的工具。

有些情况下，您正在使用的 ESP-IDF 版本可能并未创建命令提示符快捷方式，此时您可以根据下方步骤将 ESP-IDF 工具添加至 PATH。

首先，请打开需要使用 ESP-IDF 的命令提示符窗口，切换至 ESP-IDF 的安装路径，然后执行 `export.bat`：

```
cd %userprofile%\esp\esp-idf
export.bat
```

对于 Powershell 用户，请同样切换至 ESP-IDF 的安装路径，然后执行 `export.ps1`：

```
cd ~/esp/esp-idf
export.ps1
```

运行完成后，您就可以通过命令提示符使用 ESP-IDF 工具了。

1.5.2 Linux 平台工具链的标准设置

安装准备

编译 ESP-IDF 需要以下软件包：

- CentOS 7:

```
sudo yum -y update && sudo yum install git wget flex bison gperf python3␣
↪python3-pip python3-setuptools cmake ninja-build ccache dfu-util
```

目前仍然支持 CentOS 7，但为了更好的用户体验，建议使用 CentOS 8。

- Ubuntu 和 Debian:

```
sudo apt-get install git wget flex bison gperf python3 python3-pip python3-
↪setuptools cmake ninja-build ccache libffi-dev libssl-dev dfu-util
```

- Arch:

```
sudo pacman -S --needed gcc git make flex bison gperf python-pip python-
↪pyserial cmake ninja ccache dfu-util
```

注解： 使用 ESP-IDF 需要 CMake 3.5 或以上版本。较早版本的 Linux 可能需要升级才能向后移植仓库，或安装“cmake3”软件包，而不是安装“cmake”。

其他提示

权限问题 /dev/ttyUSB0 使用某些 Linux 版本向 ESP32-S2 烧录固件时，可能会出现 `Failed to open port /dev/ttyUSB0` 错误消息。此时可以将用户添加至*Linux Dialout 组*。

修复 Ubuntu 16.04 损坏的 pip

`python3-pip` 包可能已损坏无法升级。需使用脚本 get-pip.py 手动删除并安装该包:

```
apt remove python3-pip python3-virtualenv; rm -r ~/.local
rm -r ~/.espressif/python_env && python get-pip.py
```

停用 Python 2

Python 2 已经 结束生命周期，ESP-IDF 很快将不再支持 Python 2。请安装 Python 3.6 或以上版本。可参考上面列出的目前主流 Linux 发行版的安装说明。

后续步骤

继续设置开发环境，请前往*第二步：获取 ESP-IDF* 章节。

1.5.3 MacOS 平台工具链的标准设置

安装准备

ESP-IDF 将使用 macOS 上默认安装的 Python 版本。

- 安装 pip:

  ```
  sudo easy_install pip
  ```

- 安装 CMake 和 Ninja 编译工具：
  - 若有 HomeBrew，您可以运行:

    ```
    brew install cmake ninja dfu-util
    ```

  - 若有 MacPorts，您可以运行:

    ```
    sudo port install cmake ninja dfu-util
    ```

  - 若以上均不适用，请访问 CMake 和 Ninja 主页，查询有关 macOS 平台的下载安装问题。
- 强烈建议同时安装 ccache 以获得更快的编译速度。如有 HomeBrew，可通过 MacPorts 上的 `brew install ccache` 或 `sudo port install ccache` 完成安装。

注解： 如您在上述任何步骤中遇到以下错误:

```
``xcrun: error: invalid active developer path (/Library/Developer/
↪CommandLineTools), missing xcrun at:/Library/Developer/CommandLineTools/usr/bin/
↪xcrun``
```

则必须安装 XCode 命令行工具，具体可运行 `xcode-select --install`。

安装 Python 3 Catalina 10.15 发布说明 中表示不推荐使用 Python 2.7 版本，在未来的 macOS 版本中也不会默认包含 Python 2.7。执行以下命令来检查您当前使用的 Python 版本:

```
python --version
```

如果输出结果是 `Python 2.7.17`，则代表您的默认解析器是 Python 2.7。这时需要您运行以下命令检查电脑上是否已经安装过 Python 3:

```
python3 --version
```

如果运行上述命令出现错误，则代表电脑上没有安装 Python 3。

请根据以下步骤安装 Python 3：

- 使用 HomeBrew 进行安装的方法如下：

```
brew install python3
```

- 使用 MacPorts 进行安装的方法如下：

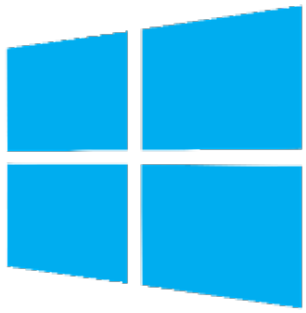
```
sudo port install python38
```

停用 Python 2

Python 2 已经 结束生命周期，ESP-IDF 很快将不再支持 Python 2。请安装 Python 3.6 或以上版本。可参考上面列出的 macOS 安装说明。

后续步骤

继续设置开发环境，请前往第二步：获取 *ESP-IDF* 章节。

		
Windows	Linux	mac OS

1.6 第二步：获取 ESP-IDF

在围绕 ESP32-S2 构建应用程序之前，请先获取乐鑫提供的软件库文件 ESP-IDF 仓库。

获取 ESP-IDF 的本地副本：打开终端，切换到您要保存 ESP-IDF 的工作目录，使用 `git clone` 命令克隆远程仓库。针对不同操作系统的详细步骤，请见下文。

注解： 在本文档中，Linux 和 macOS 操作系统中 ESP-IDF 的默认安装路径为 `~/esp`；Windows 操作系统的默认路径为 `%userprofile%\esp`。您也可以将 ESP-IDF 安装在任何其他路径下，但请注意在使用命令行时进行相应替换。注意，ESP-IDF 不支持带有空格的路径。

1.6.1 Linux 和 macOS 操作系统

打开终端后运行以下命令：

```
cd ~/esp
git clone -b v4.2.5 --recursive https://github.com/espressif/esp-idf.git
```

ESP-IDF 将下载至 `~/esp/esp-idf`。

请前往*ESP-IDF 版本简介*，查看 ESP-IDF 不同版本的具体适用场景。

1.6.2 Windows 操作系统

除了安装必要工具外，第一步中介绍的*ESP-IDF 工具安装器* 也能同时下载 ESP-IDF 本地副本。

请前往*ESP-IDF 版本简介*，查看 ESP-IDF 不同版本的具体适用场景。

除了使用 ESP-IDF 工具安装器，您也可以参考指南 手动下载 ESP-IDF。

1.7 第三步：设置工具

除了 ESP-IDF 本身，您还需要安装 ESP-IDF 使用的各种工具，比如编译器、调试器、Python 包等。

1.7.1 Windows 操作系统

请根据第一步中对 Windows (*ESP-IDF 工具安装器*) 的介绍，安装所有必需工具。

除了使用 ESP-IDF 工具安装器，您也可以通过 **命令提示符**窗口手动安装这些工具。具体步骤见下：

```
cd %userprofile%\esp\esp-idf
install.bat esp32s2
```

或使用 Windows PowerShell

```
cd ~/esp/esp-idf
./install.ps1 esp32s2
```

1.7.2 Linux 和 macOS 操作系统

```
cd ~/esp/esp-idf
./install.sh esp32s2
```

或使用 Fish shell

```
cd ~/esp/esp-idf
./install.fish esp32s2
```

注解： 通过一次性指定多个目标，可为多个目标芯片同时安装工具，如运行 `./install.sh esp32,esp32c3,esp32s3`。通过运行 `./install.sh` 或 `./install.sh all` 可一次性为所有支持的目标芯片安装工具。

1.7.3 下载工具备选方案

ESP-IDF 工具安装器会下载 Github 发布版本中附带的一些工具，如果访问 Github 较为缓慢，则可以设置一个环境变量，实现优先选择 Espressif 的下载服务器进行 Github 资源下载。

注解： 该设置只影响从 Github 发布版本中下载的单个工具，它并不会改变访问任何 Git 仓库的 URL。

Windows 操作系统

如果希望在运行 ESP-IDF 工具安装器或在使用命令行安装工具时优先选择 Espressif 下载服务器，可通过以下方式设置：打开系统控制面板，然后点击高级设置，添加一个新的环境变量（类型为用户或系统都可以，名称为 `IDF_GITHUB_ASSETS`，值为 `dl.espressif.com/github_assets`），最后点击确定。

如果在添加新的环境变量前命令行窗口或 ESP-IDF 工具安装器窗口已经打开，请关闭这些窗口后重新打开。

当设置好这个新的环境变量后，ESP-IDF 工具安装器以及命令行安装程序将会优先选择 Espressif 下载服务器。

Linux 和 macOS 操作系统

要在安装工具时优先选择 Espressif 下载服务器，请在运行 `install.sh` 时使用以下命令：

```
cd ~/esp/esp-idf
export IDF_GITHUB_ASSETS="dl.espressif.com/github_assets"
./install.sh
```

1.7.4 自定义工具安装路径

本步骤中介绍的脚本将 ESP-IDF 所需的编译工具默认安装在用户根文件夹中，即 Linux 和 macOS 系统中的 `$HOME/.espressif` 和 Windows 系统的 `%USERPROFILE%\.espressif`。您也可以选择将工具安装到其他目录中，但请在运行安装脚本前，重新设置环境变量 `IDF_TOOLS_PATH`。注意，请确保您的用户已经具备了读写该路径的权限。

如果修改了 `IDF_TOOLS_PATH` 变量，请确保该变量在每次执行“安装脚本”（`install.bat`、`install.ps1` 或 `install.sh`）和导出脚本（`export.bat`、`export.ps1` 或 `export.sh`）时均保持一致。

1.8 第四步：设置环境变量

此时，您刚刚安装的工具尚未添加至 PATH 环境变量，无法通过“命令窗口”使用这些工具。因此，必须设置一些环境变量，这可以通过 ESP-IDF 提供的另一个脚本完成。

1.8.1 Windows 操作系统

Windows 安装器（*ESP-IDF 工具安装器*）可在“开始”菜单创建一个“ESP-IDF Command Prompt”快捷方式。该快捷方式可以打开命令提示符窗口，并设置所有环境变量。您可以点击该快捷方式，然后继续下一步。

此外，如果您希望在当下命令提示符窗口使用 ESP-IDF，请使用下方代码：

```
%userprofile%\esp\esp-idf\export.bat
```

或使用 Windows PowerShell

```
.$HOME/esp/esp-idf/export.ps1
```

1.8.2 Linux 和 macOS 操作系统

请在您需要运行 ESP-IDF 的“命令提示符”窗口运行以下命令：

```
. $HOME/esp/esp-idf/export.sh
```

注意，命令开始的 "." 与路径之间应有一个空格！

如果您需要经常运行 ESP-IDF，您可以为执行 `export.sh` 创建一个别名，具体步骤如下：

1. 复制并粘贴以下命令到 shell 配置文件中（`.profile`，`.bashrc`，`.zprofile` 等）

```
alias get_idf='. $HOME/esp/esp-idf/export.sh'
```

2. 通过重启终端窗口或运行 `source [path to profile]`，如 `source ~/.bashrc` 来刷新配置文件。

现在您可以在任何终端窗口中运行 `get_idf` 来设置或刷新 esp-idf 环境。

这里不建议您直接将 `export.sh` 添加到 shell 的配置文件。因为这会导致在每个终端会话中都激活 IDF 虚拟环境（包括无需使用 IDF 的情况），从而破坏使用虚拟环境的目的，并可能影响其他软件的使用。

1.9 第五步：开始创建工程

现在，您可以开始准备开发 ESP32-S2 应用程序了。您可以从 ESP-IDF 中 examples 目录下的 get-started/hello_world 工程开始。

将 get-started/hello_world 复制至您本地的 `~/esp` 目录下：

1.9.1 Linux 和 macOS 操作系统

```
cd ~/esp
cp -r $IDF_PATH/examples/get-started/hello_world .
```

1.9.2 Windows 操作系统

```
cd %userprofile%\esp
xcopy /e /i %IDF_PATH%\examples\get-started\hello_world hello_world
```

ESP-IDF 的 examples 目录下有一系列示例工程，都可以按照上面的方法进行创建。您可以按照上述方法复制并运行其中的任何示例，也可以直接编译示例，无需进行复制。

重要： ESP-IDF 编译系统不支持带有空格的路径。

1.10 第六步：连接设备

现在，请将您的 ESP32-S2 开发板连接到 PC，并查看开发板使用的串口。

通常，串口在不同操作系统下显示的名称有所不同：

- **Windows 操作系统：** `COM1` 等
- **Linux 操作系统：** 以 `/dev/tty` 开始
- **macOS 操作系统：** 以 `/dev/cu.` 开始

有关如何查看串口名称的详细信息，请见*与 ESP32-S2 创建串口连接*。

注解： 请记住串口名，您会在下面的步骤中用到。

1.11 第七步：配置

请进入第五步：开始创建工程 中提到的 hello_world 目录，并运行工程配置工具 menuconfig。

1.11.1 Linux 和 macOS 操作系统

```
cd ~/esp/hello_world
idf.py set-target esp32s2
idf.py menuconfig
```

1.11.2 Windows 操作系统

```
cd %userprofile%\esp\hello_world
idf.py set-target esp32s2
idf.py menuconfig
```

打开一个新项目后，应首先设置“目标”芯片 idf.py set-target esp32s2。注意，此操作将清除并初始化项目之前的编译和配置（如有）。您也可以直接将“目标”配置为环境变量（则可跳过该步骤）。更多信息，请见选择目标芯片。

如果之前的步骤都正确，则会显示下面的菜单：

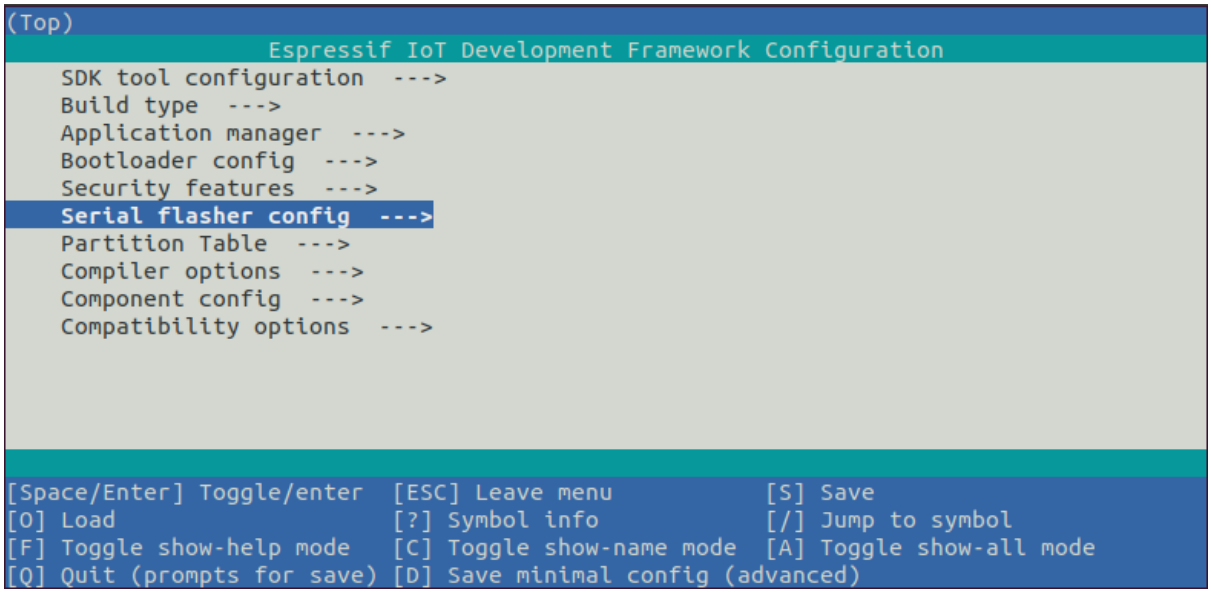


图 50: 工程配置—主窗口

您可以通过此菜单设置项目的具体变量，包括 Wi-Fi 网络名称、密码和处理器速度等. hello_world 示例项目会以默认配置运行，因此可以跳过使用 menuconfig 进行项目配置这一步骤。

注解： 您终端窗口中显示出的菜单颜色可能会与上图不同。您可以通过选项 --style 来改变外观。更多信息，请运行 idf.py menuconfig --help 命令。

1.12 第八步：编译工程

请使用以下命令，编译烧录工程：

```
idf.py build
```

运行以上命令可以编译应用程序和所有 ESP-IDF 组件，接着生成 bootloader、分区表和应用程序二进制文件。

```
$ idf.py build
Running cmake in directory /path/to/hello_world/build
Executing "cmake -G Ninja --warn-uninitialized /path/to/hello_world"...
Warn about uninitialized values.
-- Found Git: /usr/bin/git (found version "2.17.0")
-- Building empty aws_iot component due to configuration
-- Component names: ...
-- Component paths: ...

... (more lines of build system output)

[527/527] Generating hello-world.bin
esptool.py v2.3.1

Project build complete. To flash, run this command:
../../../components/esptool_py/esptool/esptool.py -p (PORT) -b 921600 write_flash -
↪-flash_mode dio --flash_size detect --flash_freq 40m 0x10000 build/hello-world.
↪bin  build 0x1000 build/bootloader/bootloader.bin 0x8000 build/partition_table/
↪partition-table.bin
or run 'idf.py -p PORT flash'
```

如果一切正常，编译完成后将生成.bin 文件。

1.13 第九步：烧录到设备

请使用以下命令，将刚刚生成的二进制文件烧录（bootloader.bin, partition-table.bin 和 hello-world.bin）至您的 ESP32-S2 开发板：

```
idf.py -p PORT [-b BAUD] flash
```

请将 PORT 替换为 ESP32-S2 开发板的串口名称，具体可见*第六步：连接设备*。

您还可以将 BAUD 替换为您希望的烧录波特率。默认波特率为 `460800`。

更多有关 idf.py 参数的详情，请见*idf.py*。

注解： 勾选 `flash` 选项将自动编译并烧录工程，因此无需再运行 `idf.py build`。

1.13.1 烧录过程中可能遇到的问题

如果在运行给定命令时出现如 “连接失败” 这样的错误，原因之一则可能是运行 `esptool.py` 出现错误。`esptool.py` 是编译系统调用的程序，用于重置芯片、与 ROM 引导加载器交互以及烧录固件的工具。解决该问题的一个简单的方法就是按照以下步骤进行手动复位。如果问题仍未解决，请参考 Troubleshooting 获取更多信息。

`esptool.py` 通过使 USB 转串口转接器芯片（如 FTDI 或 CP210x）的 DTR 和 RTS 控制线生效来自动复位 ESP32-S2（请参考*与 ESP32-S2 创建串口连接* 获取更多详细信息）。DTR 和 RTS 控制线又连接到 ESP32-S2 的 `GPIO0` 和 `CHIP_PU` (EN) 管脚上，因此 DTR 和 RTS 的电压水平变化会使 ESP32-S2 进入固件下载模式。相关示例可查看 ESP32 DevKitC 开发板的 原理图。

一般来说，使用官方的 esp-idf 开发板不会出现问题。但是，`esptool.py` 在以下情况下不能自动重置硬件。

- 您的硬件没有连接到 GPIO0 和 CIHP_PU 的 DTR 和 RTS 控制线。
- DTR 和 RTS 控制线的配置方式不同
- 根本没有这样的串行控制线路

根据您硬件的种类，也可以将您 ESP32-S2 开发板手动设置成固件下载模式（复位）。- 对于 Espressif 的开发板，您可以参考对应开发板的入门指南或用户指南。例如，可以通过按住 **Boot** 按钮 (GPIO0) 再按住 **EN** 按钮 (CHIP_PU) 来手动复位 esp-idf 开发板。- 对于其他类型的硬件，可以尝试将 GPIO0 拉低。

1.13.2 常规操作

在烧录过程中，您会看到类似如下的输出日志：

```
...
esptool.py --chip esp32s2 -p /dev/ttyUSB0 -b 460800 --before=default_reset --
↪after=hard_reset write_flash --flash_mode dio --flash_freq 40m --flash_size 2MB␣
↪0x8000 partition_table/partition-table.bin 0x1000 bootloader/bootloader.bin␣
↪0x10000 hello-world.bin
esptool.py v3.0-dev
Serial port /dev/ttyUSB0
Connecting....
Chip is ESP32-S2
Features: WiFi
Crystal is 40MHz
MAC: 18:fe:34:72:50:e3
Uploading stub...
Running stub...
Stub running...
Changing baud rate to 460800
Changed.
Configuring flash size...
Compressed 3072 bytes to 103...
Writing at 0x00008000... (100 %)
Wrote 3072 bytes (103 compressed) at 0x00008000 in 0.0 seconds (effective 3851.6␣
↪kbit/s)...
Hash of data verified.
Compressed 22592 bytes to 13483...
Writing at 0x00001000... (100 %)
Wrote 22592 bytes (13483 compressed) at 0x00001000 in 0.3 seconds (effective 595.1␣
↪kbit/s)...
Hash of data verified.
Compressed 140048 bytes to 70298...
Writing at 0x00010000... (20 %)
Writing at 0x00014000... (40 %)
Writing at 0x00018000... (60 %)
Writing at 0x0001c000... (80 %)
Writing at 0x00020000... (100 %)
Wrote 140048 bytes (70298 compressed) at 0x00010000 in 1.7 seconds (effective 662.
↪5 kbit/s)...
Hash of data verified.

Leaving...
Hard resetting via RTS pin...
Done
```

如果一切顺利，烧录完成后，开发板将会复位，应用程序“hello_world”开始运行。

如果您希望使用 Eclipse 或是 VS Code IDE，而非 idf.py，请参考*Eclipse 指南*，以及*VS Code 指南*。

1.14 第十步：监视器

您可以使用 `idf.py -p PORT monitor` 命令，监视“hello_world”的运行情况。注意，不要忘记将 PORT 替换为您的串口名称。

运行该命令后，*IDF 监视器* 应用程序将启动:

```
$ idf.py -p /dev/ttyUSB0 monitor
Running idf_monitor in directory [...]/esp/hello_world/build
Executing "python [...]/esp-idf/tools/idf_monitor.py -b 115200 [...]/esp/hello_
↪world/build/hello-world.elf"...
--- idf_monitor on /dev/ttyUSB0 115200 ---
--- Quit: Ctrl+] | Menu: Ctrl+T | Help: Ctrl+T followed by Ctrl+H ---
ets Jun  8 2016 00:22:57

rst:0x1 (POWERON_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
ets Jun  8 2016 00:22:57
...
```

此时，您就可以在启动日志和诊断日志之后，看到打印的“Hello world!”了。

```
...
Hello world!
Restarting in 10 seconds...
This is esp32 chip with 2 CPU cores, WiFi/BT/BLE, silicon revision 1, 2MB external␣
↪flash
Restarting in 9 seconds...
Restarting in 8 seconds...
Restarting in 7 seconds...
```

您可使用快捷键 `Ctrl+]`，退出 IDF 监视器。

注解： 您也可以运行以下命令，一次性执行构建、烧录和监视过程：

```
idf.py -p PORT flash monitor
```

此外，

- 请前往*IDF 监视器*，了解更多使用 IDF 监视器的快捷键和其他详情。
- 请前往*idf.py*，查看更多 `idf.py` 命令和选项。

恭喜，您已完成 ESP32-S2 的入门学习！

现在，您可以尝试一些其他 examples，或者直接开发自己的应用程序。

重要： 一些示例程序不支持 ESP32-S2，因为 ESP32-S2 中不包含所需的硬件。

在编译示例程序前请查看 README 文件中 `Supported Targets` 表格。如果表格中包含 ESP32-S2，或者不存在这个表格，那么即表示 ESP32-S2 支持这个示例程序。

1.15 更新 ESP-IDF

乐鑫会不定期推出更新版本的 ESP-IDF，修复 bug 或提供新的功能。因此，您在使用时，也应注意更新您本地的版本。最简单的方法是：直接删除您本地的 `esp-idf` 文件夹，然后按照*第二步：获取 ESP-IDF* 中的指示，重新完成克隆。

此外，您可以仅更新变更部分。具体方式，请前往*更新* 章节查看。

注意，更新完成后，请再次运行安装脚本，以防新版 ESP-IDF 所需的工具也有所更新。具体请参考*第三步：设置工具*。

一旦重新安装好工具，请使用导出脚本更新环境，具体请参考第四步：设置环境变量。

1.16 相关文档

1.16.1 与 ESP32-S2 创建串口连接

本章节主要介绍如何创建 ESP32-S2 和 PC 之间的串口连接。

连接 ESP32-S2 和 PC

用 USB 线将 ESP32-S2 开发板连接到 PC。如果设备驱动程序没有自动安装，请先确认 ESP32-S2 开发板上的 USB 转串口芯片（或外部转串口适配器）型号，然后在网上搜索驱动程序，并进行手动安装。

以下是乐鑫 ESP32-S2 开发板驱动程序的链接：

- CP210x: CP210x USB 至 UART 桥 VCP 驱动程序
- FTDI: FTDI 虚拟 COM 端口驱动程序

以上驱动仅供参考，请参考开发板用户指南，查看开发板具体使用的 USB 转串口芯片。一般情况下，当 ESP32-S2 开发板与 PC 连接时，对应驱动程序应该已经被打包在操作系统中，并已经自动安装。

在 Windows 上查看端口

检查 Windows 设备管理器中的 COM 端口列表。断开 ESP32-S2 与 PC 的连接，然后重新连接，查看哪个端口从列表中消失，然后再次出现。

以下为 ESP32 DevKitC 和 ESP32 WROVER KIT 串口：

在 Linux 和 macOS 上查看端口

查看 ESP32-S2 开发板（或外部转串口适配器）的串口设备名称，请运行两次以下命令。首先，断开开发板或适配器，第一次运行命令；然后，连接开发板或适配器，第二次运行命令。其中，第二次运行命令后出现的端口即是 ESP32-S2 对应的串口：

Linux

```
ls /dev/tty*
```

macOS

```
ls /dev/cu.*
```

注解： 对于 macOS 用户：若你没有看到串口，请检查你是否已按照《入门指南》安装了适用于你特定开发板的 USB/串口驱动程序。对于 macOS High Sierra (10.13) 的用户，你可能还需要手动允许驱动程序的加载，具体可打开 系统偏好设置 -> 安全和隐私 -> 通用，检查是否有信息显示：“来自开发人员的系统软件…”，其中开发人员的名称为 Silicon Labs 或 FTDI。

在 Linux 中添加用户到 dialout

当前登录用户应当可以通过 USB 对串口进行读写操作。在多数 Linux 版本中，你都可以通过以下命令，将用户添加到 dialout 组，来获许读写权限：

```
sudo usermod -a -G dialout $USER
```

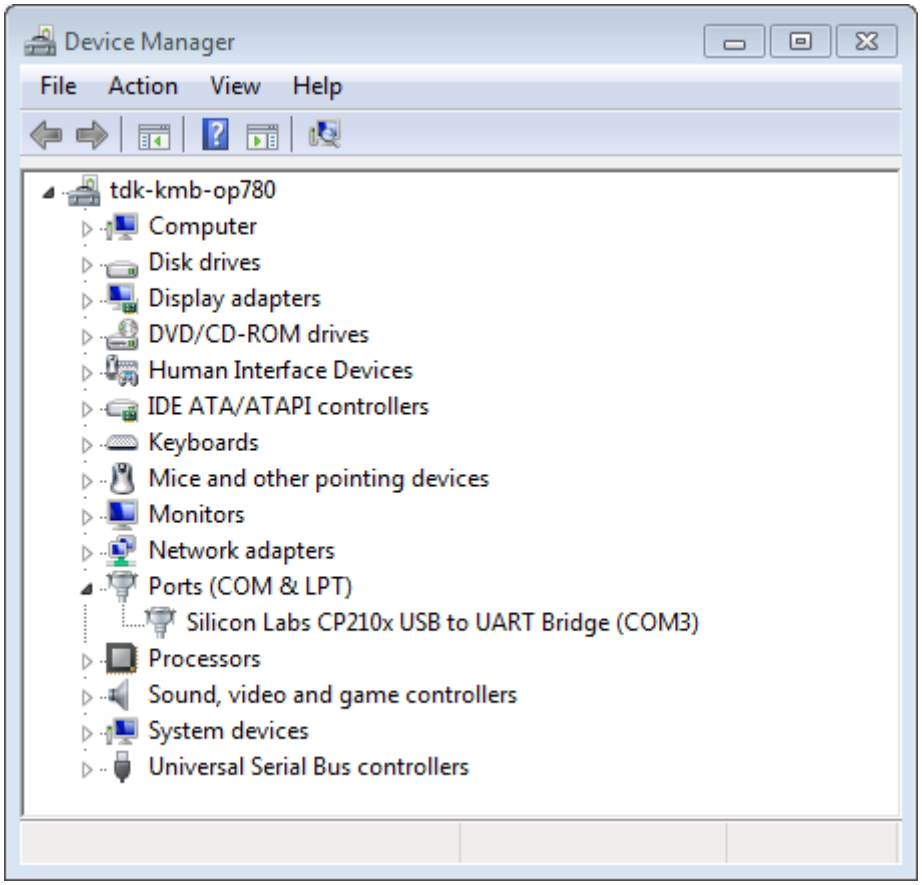


图 51: 设备管理器中 ESP32-DevKitC 的 USB 至 UART 桥

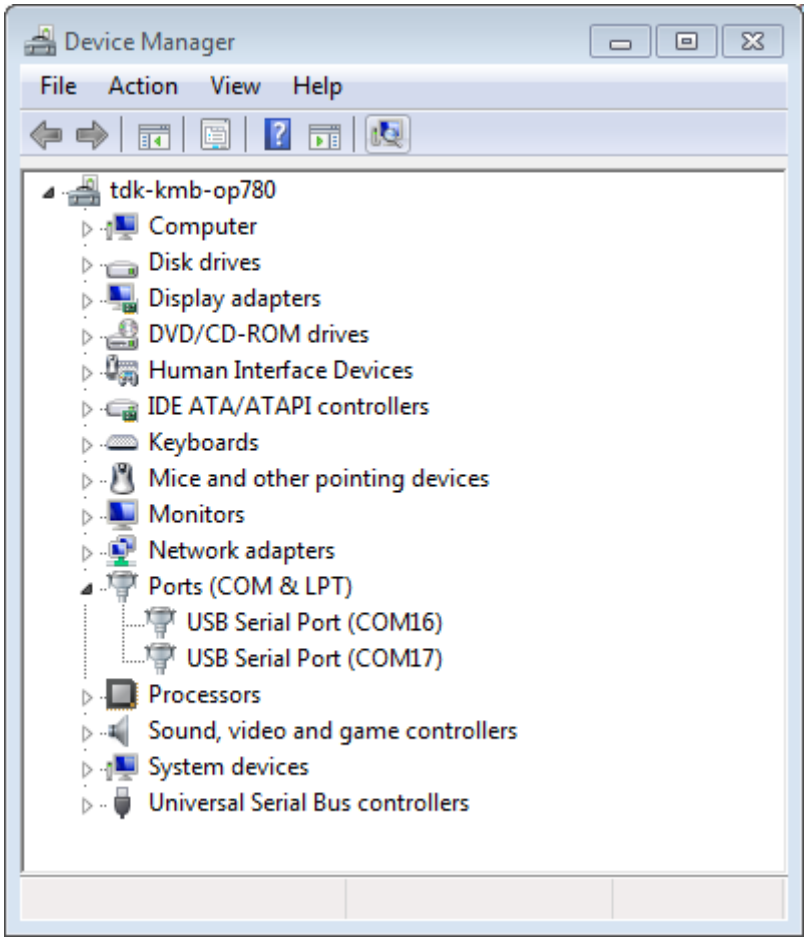


图 52: Windows 设备管理器中 ESP-WROVER-KIT 的两个 USB 串行端口

在 Arch Linux 中，需要通过以下命令将用户添加到 uucp 组中:

```
sudo usermod -a -G uucp $USER
```

请重新登录，确保串口读写权限可以生效。

确认串口连接

现在，请使用串口终端程序，验证串口连接是否可用。在本示例中，我们将使用 PuTTY SSH Client，PuTTY SSH Client 既可用于 Windows 也可用于 Linux。你也可以使用其他串口程序并设置如下的通信参数。

运行终端，配置串口：波特率 = 115200，数据位 = 8，停止位 = 1，奇偶校验 = N。以下截屏分别展示了在 Windows 和 Linux 中配置串口和上述通信参数（如 115200-8-1-N）。注意，这里一定要选择在上述步骤中确认的串口进行配置。

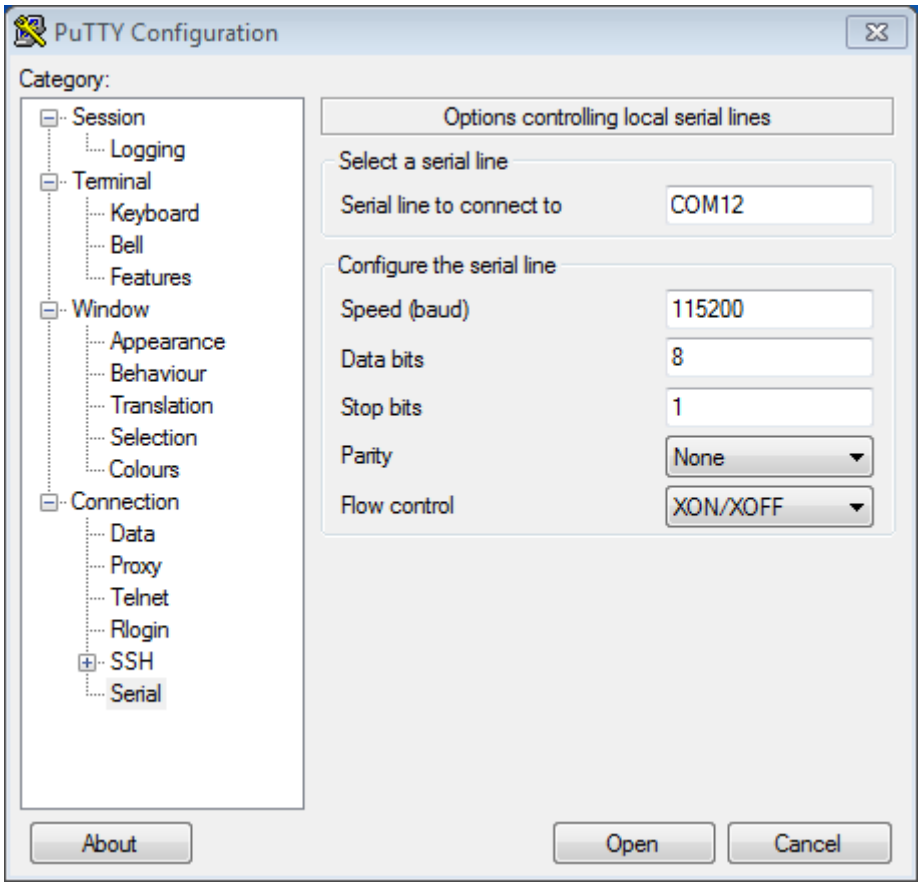


图 53: 在 Windows 操作系统中使用 PuTTY 设置串口通信参数

然后，请检查 ESP32-S2 是否有打印日志。如有，请在终端打开串口进行查看。这里，日志内容取决于加载到 ESP32-S2 的应用程序，下图即为一个示例。

```
ets Jun  8 2016 00:22:57

rst:0x5 (DEEPSLEEP_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
ets Jun  8 2016 00:22:57

rst:0x7 (TG0WDT_SYS_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
configsip: 0, SPIWP:0x00
clk_drv:0x00,q_drv:0x00,d_drv:0x00,cs0_drv:0x00,hd_drv:0x00,wp_drv:0x00
mode:DIO, clock div:2
load:0x3fff0008,len:8
```

(下页继续)

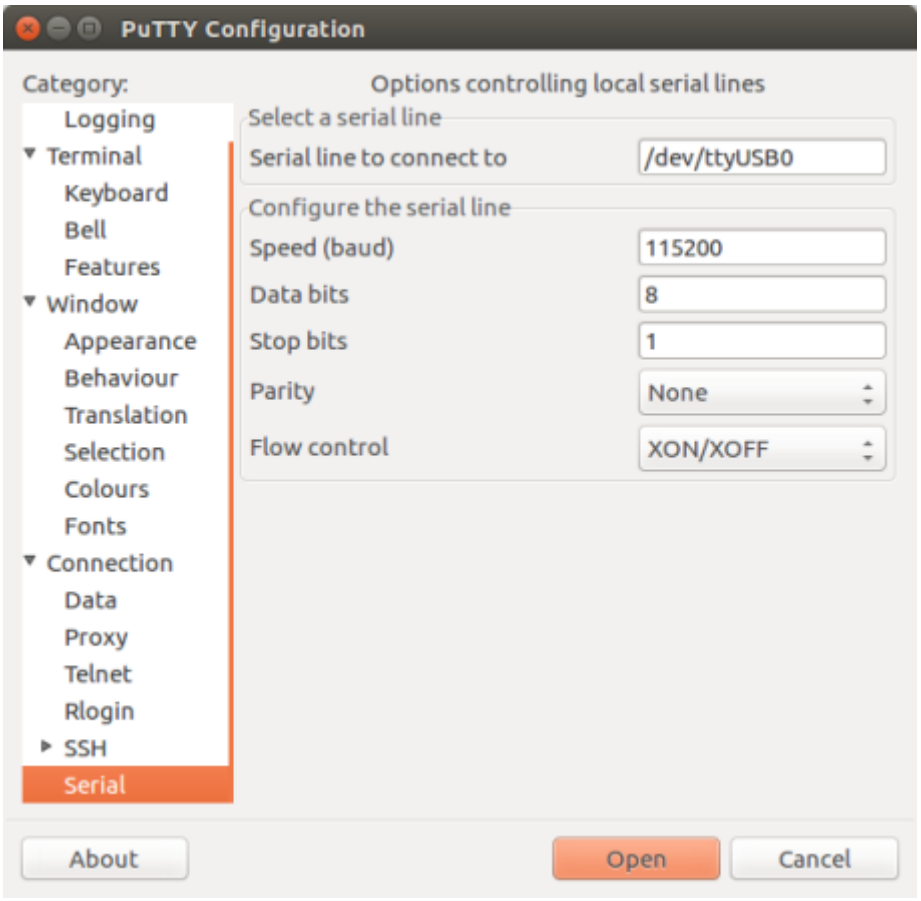


图 54: 在 Linux 操作系统中使用 PuTTY 设置串口通信参数

(续上页)

```
load:0x3fff0010,len:3464
load:0x40078000,len:7828
load:0x40080000,len:252
entry 0x40080034
I (44) boot: ESP-IDF v2.0-rc1-401-gf9fba35 2nd stage bootloader
I (45) boot: compile time 18:48:10

...
```

如果打印出的日志是可读的（而不是乱码），则表示串口连接正常。此时，你可以继续进行安装，并最终将应用程序上载到 ESP32-S2。

注解： 在某些串口接线方式下，在 ESP32-S2 启动并开始打印串口日志前，需要在终端程序中禁用串口 RTS & DTR 引脚。该问题仅存在于将 RTS & DTR 引脚直接连接到 EN & GPIO0 引脚上的情况，绝大多数开发板（包括乐鑫所有的开发板）都没有这个问题。更多详细信息，参见 esptool 文档。

注解： 请在验证完串口通信正常后，关闭串口终端。下一步，我们将使用另一个应用程序将新的固件上传到 ESP32-S2。此时，如果串口被占用则无法成功。

如你在安装 ESP32-S2 硬件开发的软件环境时，从第六步：连接设备 跳转到了这里，请从第七步：配置 继续阅读。

1.16.2 Eclipse IDE 创建和烧录指南

ESP-IDF V4.0 默认采用基于 CMake 的构建系统。

针对 CMake 构建系统，ESP-IDF 有一款新的 Eclipse 插件。具体操作指南，请见 ESP-IDF Eclipse 插件。

注解： ESP-IDF Eclipse 插件 中使用的是 macOS 截图，但安装指南对 Windows、Linux 和 macOS 均适用。

1.16.3 VS Code IDE 快速入门

我们支持 VS Code，并且致力于为所有与 ESP-IDF 相关的操作提供完善的端到端支持，包括构建、烧录、监控、调试、追踪、core-dump、以及系统追踪查看器等操作。

快速安装指南

推荐您从 VS Code 插件市场 中下载 ESP-IDF VS Code 插件，或根据 快速安装指南 安装 ESP-IDF VS Code 插件。

查看 ESP-IDF VS Code 插件教程 *<https://github.com/espressif/vscode-esp-idf-extension/blob/master/docs/tutorial/toc.md>* 了解如何使用所有功能。

支持如下功能

- **安装程序**：帮助您迅速安装 ESP-IDF 及其相关工具链。
- **构建**：通过一键构建和多目标构建，轻松构建并部署您的应用程序。
- **烧录**：UART 和 JTAG 均可完成烧录。
- **监控**：内置终端带有监控功能，您可以在 VS Code 中启用 IDF 监控命令，操作方法和传统终端一样。

- 调试 <*https://github.com/espressif/vscode-esp-idf-extension/blob/master/docs/tutorial/debugging.md*>：提供立即可用的硬件调试功能，同时支持事后剖析调试如 core-dump 功能，分析 bug 更加方便。
- **GUI 菜单配置**：提供简化的用户界面，用于配置您的芯片。
- **应用程序追踪 & 堆追踪**：支持从应用程序中收集跟踪，并提供简化的用户界面分析跟踪。
- **系统视图查看器**：读取并显示 *.svdat* 文件到用户追踪界面，同时支持多个内核追踪视图。
- **IDF 二进制大小分析**：为分析二进制文件大小提供用户界面。
- Rainmaker Cloud：我们有内置的 Rainmaker Cloud 支持，您可以轻松编辑/读取连接的物联网设备的状态。
- **代码覆盖**：我们有内置的代码覆盖支持，将用颜色突出显示已经覆盖的行。我们也会在 IDE 内部直接渲染现有的 HTML 报告。

Bugs 问题 & 功能请求

如果您在使用 VS Code 或其某些功能上遇到问题，建议您在 论坛 或是 github 上提出您的问题，我们开发团队会对问题进行解答。

我们也欢迎您提出新的功能需求，正是由于用户要求新功能或是建议对现有功能进行改善，才成就我们今天所具备的大多数功能。欢迎您在 github 上提出功能请求。

1.16.4 IDF 监视器

IDF 监视器是一个串行终端程序，用于收发目标设备串口的串行数据，IDF 监视器同时还兼具 IDF 的其他特性。

在 IDF 中调用以下目标函数可以启用此监视器：

- **若使用 CMake 编译系统，则请调用**：`idf.py monitor`
- **若使用传统 GNU Make 编译系统，请调用**：`make monitor`

操作快捷键

为了方便与 IDF 监视器进行交互，请使用表中给出的快捷键。

快捷键	操作	描述
Ctrl+]	退出监视器程序	
Ctrl+T	菜单退出键	按下如下给出的任意键之一，并按指示操作。
• Ctrl+T	将菜单字符发送至远程	
• Ctrl+]	将 exit 字符发送至远程	
• Ctrl+P	重置目标设备，进入 Bootloader，通过 RTS 线暂停应用程序	重置目标设备，通过 RTS 线（如已连接）进入 Bootloader，此时开发板不运行任何程序。等待其他设备启动时可以使用此操作。
• Ctrl+R	通过 RTS 线重置目标设备	重置设备，并通过 RTS 线（如已连接）重新启动应用程序。
• Ctrl+F	编译并烧录此项目	暂停 idf_monitor，运行 `flash` 目标，然后恢复 idf_monitor。任何改动的源文件都会被重新编译，然后重新烧录。如果 idf_monitor 是以参数 `-E` 启动的，则会运行目标 `encrypted-flash`。
• Ctrl+A (或者 A)	仅编译及烧录应用程序	暂停 idf_monitor，运行 `app-flash` 目标，然后恢复 idf_monitor。这与 `flash` 类似，但只有主应用程序被编译并被重新烧录。如果 idf_monitor 是以参数 `-E` 启动的，则会运行目标 `encrypted-flash`。
• Ctrl+Y	停止/恢复在屏幕上打印日志输出	激活时，会丢弃所有传入的串行数据。允许在不退出监视器的情况下快速暂停和检查日志输出。
• Ctrl+L	停止/恢复向文件写入日志输出	在工程目录下创建一个文件，用于写入日志输出。可使用快捷键停止/恢复该功能（退出 IDF 监视器也会终止该功能）
• Ctrl+H (或者 H)	显示所有快捷键	
• Ctrl+X (或者 X)	退出监视器程序	

除了 `Ctrl-]` 和 `Ctrl-T`，其他快捷键信号会通过串口发送到目标设备。

兼具 IDF 特性

自动解码地址 ESP-IDF 输出形式为 `0x4_____` 的十六进制代码地址后，IDF 监视器将使用 addr2line 查找该地址在源代码中的位置和对应的函数名。

ESP-IDF 应用程序发生 crash 和 panic 事件时，将产生如下的寄存器转储和回溯:

```
  Guru Meditation Error of type StoreProhibited occurred on core  0. Exception was␣
↪unhandled.
  Register dump:
  PC      : 0x400f360d  PS      : 0x00060330  A0      : 0x800dbf56  A1      :␣
↪0x3ffb7e00
  A2      : 0x3ffb136c  A3      : 0x00000005  A4      : 0x00000000  A5      :␣
↪0x00000000
  A6      : 0x00000000  A7      : 0x00000080  A8      : 0x00000000  A9      :␣
↪0x3ffb7dd0
  A10     : 0x00000003  A11     : 0x00060f23  A12     : 0x00060f20  A13     :␣
↪0x3ffba6d0
```

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```
  A14      : 0x00000047  A15      : 0x0000000f  SAR      : 0x00000019  EXCCAUSE:␣
↪0x0000001d
  EXCVADDR: 0x00000000  LBEG     : 0x4000c46c  LEND     : 0x4000c477  LCOUNT  :␣
↪0x00000000

  Backtrace: 0x400f360d:0x3ffb7e00 0x400dbf56:0x3ffb7e20 0x400dbf5e:0x3ffb7e40␣
↪0x400dbf82:0x3ffb7e60 0x400d071d:0x3ffb7e90

IDF 监视器为寄存器转储补充如下信息::

  Guru Meditation Error of type StoreProhibited occurred on core  0. Exception was␣
↪unhandled.
  Register dump:
  PC       : 0x400f360d  PS       : 0x00060330  A0       : 0x800dbf56  A1       :␣
↪0x3ffb7e00
  0x400f360d: do_something_to_crash at /home/gus/esp/32/idf/examples/get-started/
↪hello_world/main/./hello_world_main.c:57
  (inlined by) inner_dont_crash at /home/gus/esp/32/idf/examples/get-started/hello_
↪world/main/./hello_world_main.c:52
  A2       : 0x3ffb136c  A3       : 0x00000005  A4       : 0x00000000  A5       :␣
↪0x00000000
  A6       : 0x00000000  A7       : 0x00000080  A8       : 0x00000000  A9       :␣
↪0x3ffb7dd0
  A10      : 0x00000003  A11      : 0x00060f23  A12      : 0x00060f20  A13      :␣
↪0x3ffba6d0
  A14      : 0x00000047  A15      : 0x0000000f  SAR      : 0x00000019  EXCCAUSE:␣
↪0x0000001d
  EXCVADDR: 0x00000000  LBEG     : 0x4000c46c  LEND     : 0x4000c477  LCOUNT  :␣
↪0x00000000

  Backtrace: 0x400f360d:0x3ffb7e00 0x400dbf56:0x3ffb7e20 0x400dbf5e:0x3ffb7e40␣
↪0x400dbf82:0x3ffb7e60 0x400d071d:0x3ffb7e90
  0x400f360d: do_something_to_crash at /home/gus/esp/32/idf/examples/get-started/
↪hello_world/main/./hello_world_main.c:57
  (inlined by) inner_dont_crash at /home/gus/esp/32/idf/examples/get-started/hello_
↪world/main/./hello_world_main.c:52
  0x400dbf56: still_dont_crash at /home/gus/esp/32/idf/examples/get-started/hello_
↪world/main/./hello_world_main.c:47
  0x400dbf5e: dont_crash at /home/gus/esp/32/idf/examples/get-started/hello_world/
↪main/./hello_world_main.c:42
  0x400dbf82: app_main at /home/gus/esp/32/idf/examples/get-started/hello_world/
↪main/./hello_world_main.c:33
  0x400d071d: main_task at /home/gus/esp/32/idf/components/esp32s2/./cpu_start.
↪c:254
```

IDF 监视器在后台运行以下命令，解码各地址:

```
xtensa-esp32s2-elf-addr2line -pfiaC -e build/PROJECT.elf ADDRESS
```

配置 GDBStub 以启用 GDB 默认情况下，如果 ESP-IDF 应用程序发生 crash 事件，panic 处理器将在串口上打印相关寄存器和堆栈转储（类似上述情况），然后重置开发板。

或者选择配置 panic 处理器以运行 GDBStub，GDBStub 工具可以与 GDB 项目调试器进行通信，允许读取内存、检查调用堆栈帧和变量等。GDBStub 虽然没有 JTAG 通用，但不需要使用特殊硬件。

如需启用 GDBStub，请运行 `idf.py menuconfig`（适用于 CMake 编译系统），并将*CONFIG_ESP_SYSTEM_PANIC* 选项设置为 `Invoke GDBStub`。

在这种情况下，如果 panic 处理器被触发，只要 IDF 监视器监控到 GDBStub 已经加载，panic 处理器就会自动暂停串行监控并使用必要的参数运行 GDB。GDB 退出后，通过 RTS 串口线复位开发板。如果未连接 RTS 串口线，请按复位键，手动复位开发板。

IDF 监控器在后台运行如下命令:

```
xtensa-esp32s2-elf-gdb -ex "set serial baud BAUD" -ex "target remote PORT" -ex␣
↪interrupt build/PROJECT.elf :idf_target:`Hello NAME chip`
```

输出筛选 可以调用 `idf.py monitor --print-filter="xyz"` 启动 IDF 监视器，其中，`--print-filter` 是输出筛选的参数。参数默认值为空字符串，可打印任何内容。

若需对打印内容设置限制，可指定 `<tag>:<log_level>` 等选项，其中 `<tag>` 是标签字符串，`<log_level>` 是 {N, E, W, I, D, V, *} 集合中的一个字母，指的是日志级别。

例如，`PRINT_FILTER="tag1:W"` 只匹配并打印 `ESP_LOGW("tag1", ...)` 所写的输出，或者写在较低日志详细度级别的输出，即 `ESP_LOGE("tag1", ...)`。请勿指定 `<log_level>` 或使用详细级别默认值 `*`。

注解: 编译时，可以使用主日志在日志库 中禁用不需要的输出。也可以使用 IDF 监视器筛选输出来调整筛选设置，且无需重新编译应用程序。

应用程序标签不能包含空格、星号 `*`、冒号 `:`，以便兼容输出筛选功能。

如果应用程序输出的最后一行后面没有回车，可能会影响输出筛选功能，即，监视器开始打印该行，但后来发现该行不应该被写入。这是一个已知问题，可以通过添加回车来避免此问题（特别是在没有输出紧跟其后的情况下）。

筛选规则示例

- `*` 可用于匹配任何类型标签。但 `PRINT_FILTER="*:I tag1:E"` 打印关于 `tag1` 的输出时会报错，这是因为 `tag1` 规则比 `*` 规则的优先级高。
- 默认规则（空）等价于 `*:V`，因为在详细级别或更低级别匹配任意标签即意味匹配所有内容。
- `"*:N"` 不仅抑制了日志功能的输出，也抑制了 `printf` 的打印输出。为了避免这一问题，请使用 `*:E` 或更高的冗余级别。
- 规则 `"tag1:V"`、`"tag1:v"`、`"tag1:"`、`"tag1:*"` 和 `"tag1"` 等同。
- 规则 `"tag1:W tag1:E"` 等同于 `"tag1:E"`，这是因为后续出现的具有相同名称的标签会覆盖掉前一个标签。
- 规则 `"tag1:I tag2:W"` 仅在 Info 详细度级别或更低级别打印 `tag1`，在 Warning 详细度级别或更低级别打印 `tag2`。
- 规则 `"tag1:I tag2:W tag3:N"` 在本质上等同于上一规则，这是因为 `tag3:N` 指定 `tag3` 不打印。
- `tag3:N` 在规则 `"tag1:I tag2:W tag3:N *:V"` 中更有意义，这是因为如果没有 `tag3:N`，`tag3` 信息就可能打印出来了；`tag1` 和 `tag2` 错误信息会打印在指定的详细度级别（或更低级别），并默认打印所有内容。

高级筛选规则示例 如下日志是在没有设置任何筛选选项的情况下获得的:

```
load:0x40078000,len:13564
entry 0x40078d4c
E (31) esp_image: image at 0x30000 has invalid magic byte
W (31) esp_image: image at 0x30000 has invalid SPI mode 255
E (39) boot: Factory app partition is not bootable
I (568) cpu_start: Pro cpu up.
I (569) heap_init: Initializing. RAM available for dynamic allocation:
I (603) cpu_start: Pro cpu start user code
D (309) light_driver: [light_init, 74]:status: 1, mode: 2
D (318) vfs: esp_vfs_register_fd_range is successful for range <54; 64) and VFS ID␣
↪1
I (328) wifi: wifi driver task: 3ffdbf84, prio:23, stack:4096, core=0
```

`PRINT_FILTER="wifi esp_image:E light_driver:I"` 筛选选项捕获的输出如下所示:

```
E (31) esp_image: image at 0x30000 has invalid magic byte
I (328) wifi: wifi driver task: 3ffdbf84, prio:23, stack:4096, core=0
```

`PRINT_FILTER="light_driver:D esp_image:N boot:N cpu_start:N vfs:N wifi:N *:V"` 选项的输出如下:

```
load:0x40078000,len:13564
entry 0x40078d4c
I (569) heap_init: Initializing. RAM available for dynamic allocation:
D (309) light_driver: [light_init, 74]:status: 1, mode: 2
```

IDF 监视器已知问题

Windows 环境下已知问题

- 若在 Windows 环境下，出现 “winpty: command not found” 错误，请运行 `pacman -S winpty` 进行修复。
- 由于 Windows 控制台限制，有些箭头键及其他一些特殊键无法在 GDB 中使用。
- 偶然情况下，`idf.py` 或 `make` 退出时，可能会在 IDF 监视器恢复之前暂停 30 秒。
- GDB 运行时，可能会暂停一段时间，然后才开始与 GDBStub 进行通信。

1.16.5 工具链的自定义设置

除了从乐鑫官网（请见*第三步：设置工具*）下载二进制工具链外，您还可以自行编译工具链。

如无特殊需求，建议直接使用我们提供的预编译二进制工具链。不过，您可以在以下情况考虑自行编译工具链:

- 需要定制工具链编译配置
- 需要使用其他 GCC 版本（如 4.8.5）
- 需要破解 gcc、newlib 或 libstdc++
- 有相关兴趣或时间充裕
- 不信任从网站下载的 bin 文件

如需自行编译工具链，请查看以下文档:

从零开始设置 Windows 环境下的工具链

除了使用*ESP-IDF 工具安装器*，用户也可以手动设置 Windows 环境下的工具链，这也是本文的主要内容。手动安装工具可以更好地控制安装流程，同时也方便高阶用户进行自定义安装。

使用 ESP-IDF 工具安装器对工具链及其他工具进行快速标准设置，请参照*Windows 平台工具链的标准设置*。

注解: 基于 GNU Make 的构建系统要求 Windows 兼容 MSYS2 Unix，基于 CMake 的构建系统则无此要求。

获取 ESP-IDF

注解: 较早版本 ESP-IDF 使用了 **MSYS2 bash 终端**命令行。目前，基于 CMake 的编译系统可使用常见的 **Windows 命令窗口**，即本指南中使用的终端。

请注意，如果您使用基于 bash 的终端或 PowerShell 终端，一些命令语法将与下面描述有所不同。

打开命令提示符，运行以下命令:

```
mkdir %userprofile%\esp
cd %userprofile%\esp
git clone -b v4.2.5 --recursive https://github.com/espressif/esp-idf.git
```

ESP-IDF 将下载至 `%userprofile%\esp\esp-idf`。

请前往*ESP-IDF 版本简介*，查看 ESP-IDF 不同版本的具体适用场景。

注解： `git clone` 命令的 `-b v4.2.5` 选项告诉 git 从 ESP-IDF 仓库中克隆与此版本的文档对应的分支。

注解： 作为备份，还可以从 Releases page 下载此稳定版本的 zip 文件。不要下载由 GitHub 自动生成的”源代码”的 zip 文件，它们不适用于 ESP-IDF。

注解： 在克隆远程仓库时，不要忘记加上 `--recursive` 选项。否则，请接着运行以下命令，获取所有子模块

```
cd esp-idf
git submodule update --init
```

工具

cmake 工具 下载最新发布的 Windows 平台稳定版 CMake，并运行安装器。

当安装器询问“安装选项”时，选择“Add CMake to the system PATH for all users”（为所有用户的系统路径添加 CMake）或“Add CMake to the system PATH for the current user”（为当前用户的系统路径添加 CMake）。

Ninja 编译工具

注解： 目前，Ninja 仅提供支持 64 位 Windows 版本的 bin 文件。您也可以配合其他编译工具在 32 位 Windows 版本中使用 CMake 和 `idf.py`，比如 mingw-make。但是目前暂无关于此工具的说明文档。

从（下载页面）下载最新发布的 Windows 平台稳定版 ninja。

适用于 Windows 平台的 Ninja 下载文件是一个.zip 文件，包含一个 `ninja.exe` 文件。您需要将该文件解压到目录，并*添加到您的路径*（或者选择您路径中的已有目录）。

Python 下载并运行适用于 Windows 安装器的最新版 Python。

Python 安装器的“自定义”菜单可为您提供一系列选项，最后一项为“Add python.exe to Path”（添加 python.exe 到路径中）。请将该选项更改到“Will be installed”（将会安装）。

Python 安装完成后，从 Windows 开始菜单中打开“命令提示符”窗口，并运行以下命令:

```
pip install --user pyserial
```

工具链设置 下载预编译的 Windows 工具链:

https://dl.espressif.com/dl/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-win32.zip

解压压缩包文件到 `C:\Program Files`（或其他位置）。压缩包文件包含一个 `xtensa-esp32s2-elf` 目录。

然后，请将该目录下的 bin 子目录添加到您的路径。例如，C:\Program Files\xtensa-esp32s2-elf\bin。

注解： 如果您已安装 MSYS2 环境（适用“GNU Make”编译系统），则可以跳过下载那一步，直接添加目录 C:\msys32\opt\xtensa-esp32s2-elf\bin 到路径，因为 MSYS2 环境已包含工具链。

添加目录到路径 在 Windows 环境下，向 Path 环境变量增加任何新目录，请：

打开系统“控制面板”，找到环境变量对话框（Windows 10 用户请前往“高级系统设置”）。

双击 Path 变量（选择“用户”或“系统路径”，具体取决于您是否希望其他用户路径中也存在该目录)。在最后数值那里新添 ;<new value>。

后续步骤 继续设置开发环境，请前往第三步：设置工具 章节。

从零开始设置 Linux 环境下的工具链

除了从乐鑫官网直接下载已编译好的二进制工具链外，您还可以按照本文介绍，从头开始设置自己的工具链。如需快速使用已编译好的二进制工具链，可回到*Linux 平台工具链的标准设置* 章节。

注解： 设置自己的工具链可以解决 Y2K38 问题（time_t 从 32 位扩展到 64 位）。

安装准备 编译 ESP-IDF 需要以下软件包：

- CentOS 7:

```
sudo yum -y update && sudo yum install git wget ncurses-devel flex bison gperf␣
↪python3 python3-pip cmake ninja-build ccache dfu-util
```

目前仍然支持 CentOS 7，但为了更好的用户体验，建议使用 CentOS 8。

- Ubuntu 和 Debian:

```
sudo apt-get install git wget libncurses-dev flex bison gperf python3 python3-␣
↪pip python3-setuptools python3-serial python3-cryptography python3-future␣
↪python3-pyparsing python3-pyelftools cmake ninja-build ccache libffi-dev␣
↪libssl-dev dfu-util
```

- Arch:

```
sudo pacman -Sy --needed gcc git make ncurses flex bison gperf python-pyserial␣
↪python-cryptography python-future python-pyparsing python-pyelftools cmake␣
↪ninja ccache dfu-util
```

注解： 使用 ESP-IDF 需要 CMake 3.5 或以上版本。较早的 Linux 发行版可能需要升级自身的软件源仓库，或开启 backports 套件库，或安装“cmake3”软件包（不是安装“cmake”）。

从源代码编译工具链 安装依赖项：

- CentOS 7:

```
sudo yum install gawk gperf grep gettext ncurses-devel python3 python3-devel␣
↪automake bison flex texinfo help2man libtool make
```

- Ubuntu pre-16.04:

```
sudo apt-get install gawk gperf grep gettext libncurses-dev python python-dev␣
↪automake bison flex texinfo help2man libtool make
```

- Ubuntu 16.04 或以上

```
sudo apt-get install gawk gperf grep gettext python python-dev automake bison␣
↪flex texinfo help2man libtool libtool-bin make
```

- Debian 9:

```
sudo apt-get install gawk gperf grep gettext libncurses-dev python python-dev␣
↪automake bison flex texinfo help2man libtool libtool-bin make
```

- Arch:

```
sudo pacman -Sy --needed python-pip
```

创建工作目录，并进入该目录:

```
mkdir -p ~/esp
cd ~/esp
```

下载并编译 crosstool-NG：

```
git clone https://github.com/espressif/crosstool-NG.git
cd crosstool-NG
git checkout esp-2020r3
git submodule update --init
./bootstrap && ./configure --enable-local && make
```

注解： 在设置支持 64 位 time_t 的工具链时，您需要将 crosstool-NG/samples/xtensa-esp32-elf/crosstool.config 文件中第 33 和 43 行的可选参数 --enable-newlib-long-time_t 删除。

编译工具链:

```
./ct-ng xtensa-esp32s2-elf
./ct-ng build
chmod -R u+w builds/xtensa-esp32s2-elf
```

编译得到的工具链会被保存至 ~/esp/crosstool-NG/builds/xtensa-esp32s2-elf。

添加工具链到 PATH 环境变量 需要将自定义工具链复制到一个二进制目录中，并将其添加到 PATH 中。例如，您可以将编译好的工具链复制到 ~/esp/xtensa-esp32s2-elf/ 目录中。

为了正常使用工具链，您需要更新 ~/.profile 文件中 PATH 环境变量。此外，您还可以在 ~/.profile 文件中增加以下代码。这样，所有终端窗口均可以使用 xtensa-esp32s2-elf:

```
export PATH="$HOME/esp/xtensa-esp32s2-elf/bin:$PATH"
```

注解： 如果您已将 /bin/bash 设置为登录 shell，且同时存在 .bash_profile 和 .profile 两个文件，则请更新 .bash_profile。在 CentOS 环境下，alias 需要添加到 .bashrc 文件中。

退出并重新登录以使 .profile 的更改生效。运行以下命令来检查 PATH 设置是否正确:

```
printenv PATH
```

此时您需要检查输出结果的开头中是否包含类似如下的工具链路径:

```
$ printenv PATH
/home/user-name/esp/xtensa-esp32s2-elf/bin:/home/user-name/bin:/home/user-name/.
↪local/bin:/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr/bin:/sbin:/bin:/usr/
↪games:/usr/local/games:/snap/bin
```

注意这里的 /home/user-name 应该替换成您安装的主路径。

停用 Python 2 Python 2 已经 结束生命周期，ESP-IDF 很快将不再支持 Python 2。请安装 Python 3.6 或以上版本。可参考上面列出的目前主流 Linux 发行版的安装说明。

后续步骤 继续设置开发环境，请前往*第二步：获取 ESP-IDF* 章节。

从零开始设置 macOS 环境下的工具链

软件包管理器 从零开始设置工具链，您需要安装 MacPorts 或 homebrew 软件包管理器。或者，您也可以直接下载预编译的工具链。

MacPorts 需要完整的 XCode 软件，而 homebrew 只需要安装 XCode 命令行工具即可。

请参考工具链自定义设置 章节，查看可能需要从头开始设置工具链的情况。

安装准备

- 安装 pip:

```
sudo easy_install pip
```

- 安装 pyserial:

```
pip install --user pyserial
```

- 安装 CMake 和 Ninja 编译工具：
  - 若有 HomeBrew，您可以运行:

```
brew install cmake ninja dfu-util
```

  - 若有 MacPorts，您可以运行:

```
sudo port install cmake ninja dfu-util
```

从源代码编译工具链 安装依赖项：

- 对于 MacPorts:

```
sudo port install gsed gawk binutils gperf grep gettext wget libtool autoconf␣
↪automake make
```

- 对于 homebrew:

```
brew install gnu-sed gawk binutils gperftools gettext wget help2man libtool␣
↪autoconf automake make
```

创建一个文件系统镜像（区分大小写）：

```
hdiutil create ~/esp/crosstool.dmg -volname "ctng" -size 10g -fs "Case-sensitive␣
↪HFS+"
```

挂载:

```
hdiutil mount ~/esp/crosstool.dmg
```

创建指向您工作目录的符号链接:

```
mkdir -p ~/esp
ln -s /Volumes/ctng ~/esp/ctng-volume
```

前往新创建的目录

```
cd ~/esp/ctng-volume
```

下载并编译 crosstool-NG

```
git clone https://github.com/espressif/crosstool-NG.git
cd crosstool-NG
git checkout esp-2020r3
git submodule update --init
./bootstrap && ./configure --enable-local && make
```

编译工具链:

```
./ct-ng xtensa-esp32s2-elf
./ct-ng build
chmod -R u+w builds/xtensa-esp32s2-elf
```

编译得到的工具链会被保存到 ~/esp/ctng-volume/crosstool-NG/builds/xtensa-esp32s2-elf。使用工具链前，请将 ~/esp/ctng-volume/crosstool-NG/builds/xtensa-esp32s2-elf/bin 添加至 PATH 环境变量。

停用 Python 2 Python 2 已经 结束生命周期，ESP-IDF 很快将不再支持 Python 2。请安装 Python 3.6 或以上版本。可参考上面列出的 macOS 安装说明。

后续步骤 继续设置开发环境，请前往第二步：获取 *ESP-IDF* 章节。

Chapter 2

API 参考

2.1 连网 API

2.1.1 Wi-Fi

Wi-Fi 库

概述 Wi-Fi 库支持配置及监控 ESP32-S2 Wi-Fi 连网功能。

支持配置：

- 基站模式（即 STA 模式或 Wi-Fi 客户端模式），此时 ESP32-S2 连接到接入点 (AP)。
- AP 模式（即 Soft-AP 模式或接入点模式），此时基站连接到 ESP32-S2。
- AP-STA 共存模式（ESP32-S2 既是接入点，同时又作为基站连接到另外一个接入点）。
- 上述模式的各种安全模式（WPA、WPA2 及 WEP 等）。
- 扫描接入点（包括主动扫描及被动扫描）。
- 使用混杂模式监控 IEEE802.11 Wi-Fi 数据包。

应用示例 ESP-IDF 示例项目的 wifi 目录下包含以下应用程序：

- Wi-Fi 示例代码；
- 另外一个简单的应用程序 esp-idf-template，演示了如何将 ESP32-S2 模组连接到 AP。

API 参考

Header File

- esp_wifi/include/esp_wifi.h

Functions

esp_err_t **esp_wifi_init**(**const** *wifi_init_config_t* **config*)

Initialize WiFi Allocate resource for WiFi driver, such as WiFi control structure, RX/TX buffer, WiFi NVS structure etc. This WiFi also starts WiFi task.

Attention 1. This API must be called before all other WiFi API can be called

Attention 2. Always use WIFI_INIT_CONFIG_DEFAULT macro to initialize the configuration to default values, this can guarantee all the fields get correct value when more fields are added into *wifi_init_config_t* in future release. If you want to set your own initial values, overwrite the default values which are set by WIFI_INIT_CONFIG_DEFAULT. Please be notified that the field 'magic' of *wifi_init_config_t* should always be WIFI_INIT_CONFIG_MAGIC!

Return

- ESP_OK: succeed
- ESP_ERR_NO_MEM: out of memory
- others: refer to error code esp_err.h

Parameters

- `config`: pointer to WiFi initialized configuration structure; can point to a temporary variable.

esp_err_t **esp_wifi_deinit** (void)

Deinit WiFi Free all resource allocated in esp_wifi_init and stop WiFi task.

Attention 1. This API should be called if you want to remove WiFi driver from the system

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

esp_err_t **esp_wifi_set_mode** (*wifi_mode_t mode*)

Set the WiFi operating mode.

Set the WiFi operating mode as station, soft-AP or station+soft-AP, The default mode is soft-AP mode.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- others: refer to error code in esp_err.h

Parameters

- `mode`: WiFi operating mode

esp_err_t **esp_wifi_get_mode** (*wifi_mode_t *mode*)

Get current operating mode of WiFi.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters

- `[out] mode`: store current WiFi mode

esp_err_t **esp_wifi_start** (void)

Start WiFi according to current configuration If mode is WIFI_MODE_STA, it create station control block and start station If mode is WIFI_MODE_AP, it create soft-AP control block and start soft-AP If mode is WIFI_MODE_APSTA, it create soft-AP and station control block and start soft-AP and station.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_NO_MEM: out of memory
- ESP_ERR_WIFI_CONN: WiFi internal error, station or soft-AP control block wrong
- ESP_FAIL: other WiFi internal errors

esp_err_t **esp_wifi_stop** (void)

Stop WiFi If mode is WIFI_MODE_STA, it stop station and free station control block If mode is WIFI_MODE_AP, it stop soft-AP and free soft-AP control block If mode is WIFI_MODE_APSTA, it stop station/soft-AP and free station/soft-AP control block.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

esp_err_t **esp_wifi_restore** (void)

Restore WiFi stack persistent settings to default values.

This function will reset settings made using the following APIs:

- esp_wifi_set_bandwidth,
- esp_wifi_set_protocol,

- esp_wifi_set_config related
- esp_wifi_set_mode

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

esp_err_t **esp_wifi_connect**(void)

Connect the ESP32 WiFi station to the AP.

Attention 1. This API only impact WIFI_MODE_STA or WIFI_MODE_APSTA mode

Attention 2. If the ESP32 is connected to an AP, call esp_wifi_disconnect to disconnect.

Attention 3. The scanning triggered by esp_wifi_scan_start() will not be effective until connection between ESP32 and the AP is established. If ESP32 is scanning and connecting at the same time, ESP32 will abort scanning and return a warning message and error number ESP_ERR_WIFI_STATE. If you want to do reconnection after ESP32 received disconnect event, remember to add the maximum retry time, otherwise the called scan will not work. This is especially true when the AP doesn't exist, and you still try reconnection after ESP32 received disconnect event with the reason code WIFI_REASON_NO_AP_FOUND.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start
- ESP_ERR_WIFI_CONN: WiFi internal error, station or soft-AP control block wrong
- ESP_ERR_WIFI_SSID: SSID of AP which station connects is invalid

esp_err_t **esp_wifi_disconnect**(void)

Disconnect the ESP32 WiFi station from the AP.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi was not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi was not started by esp_wifi_start
- ESP_FAIL: other WiFi internal errors

esp_err_t **esp_wifi_clear_fast_connect**(void)

Currently this API is just an stub API.

Return

- ESP_OK: succeed
- others: fail

esp_err_t **esp_wifi_deauth_sta**(uint16_t *aid*)

deauthenticate all stations or associated id equals to aid

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi was not started by esp_wifi_start
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_WIFI_MODE: WiFi mode is wrong

Parameters

- `aid`: when aid is 0, deauthenticate all stations, otherwise deauthenticate station whose associated id is aid

esp_err_t **esp_wifi_scan_start**(**const** *wifi_scan_config_t* **config*, bool *block*)

Scan all available APs.

Attention If this API is called, the found APs are stored in WiFi driver dynamic allocated memory and the will be freed in esp_wifi_scan_get_ap_records, so generally, call esp_wifi_scan_get_ap_records to cause the memory to be freed once the scan is done

Attention The values of maximum active scan time and passive scan time per channel are limited to 1500 milliseconds. Values above 1500ms may cause station to disconnect from AP and are not recommended.

Return

- ESP_OK: succeed

- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi was not started by esp_wifi_start
- ESP_ERR_WIFI_TIMEOUT: blocking scan is timeout
- ESP_ERR_WIFI_STATE: wifi still connecting when invoke esp_wifi_scan_start
- others: refer to error code in esp_err.h

Parameters

- `config`: configuration of scanning
- `block`: if block is true, this API will block the caller until the scan is done, otherwise it will return immediately

esp_err_t **esp_wifi_scan_stop**(void)

Stop the scan in process.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start

esp_err_t **esp_wifi_scan_get_ap_num**(uint16_t **number*)

Get number of APs found in last scan.

Attention This API can only be called when the scan is completed, otherwise it may get wrong value.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start
- ESP_ERR_INVALID_ARG: invalid argument

Parameters

- `[out] number`: store number of APIs found in last scan

esp_err_t **esp_wifi_scan_get_ap_records**(uint16_t **number*, *wifi_ap_record_t *ap_records*)

Get AP list found in last scan.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_NO_MEM: out of memory

Parameters

- `[inout] number`: As input param, it stores max AP number ap_records can hold. As output param, it receives the actual AP number this API returns.
- `ap_records`: *wifi_ap_record_t* array to hold the found APs

esp_err_t **esp_wifi_sta_get_ap_info**(*wifi_ap_record_t *ap_info*)

Get information of AP which the ESP32 station is associated with.

Attention When the obtained country information is empty, it means that the AP does not carry country information

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_CONN: The station interface don’ t initialized
- ESP_ERR_WIFI_NOT_CONNECT: The station is in disconnect status

Parameters

- `ap_info`: the *wifi_ap_record_t* to hold AP information sta can get the connected ap’ s phy mode info through the struct member phy_11b, phy_11g, phy_11n, phy_lr in the *wifi_ap_record_t* struct. For example, phy_11b = 1 imply that ap support 802.11b mode

esp_err_t **esp_wifi_set_ps**(*wifi_ps_type_t type*)

Set current WiFi power save type.

Attention Default power save type is WIFI_PS_MIN_MODEM.

Return ESP_OK: succeed

Parameters
- `type`: power save type

esp_err_t **esp_wifi_get_ps** (*wifi_ps_type_t *type*)
Get current WiFi power save type.

Attention Default power save type is WIFI_PS_MIN_MODEM.
Return ESP_OK: succeed
Parameters
- `[out] type`: store current power save type

esp_err_t **esp_wifi_set_protocol** (*wifi_interface_t ifx*, uint8_t *protocol_bitmap*)
Set protocol type of specified interface The default protocol is (WIFI_PROTOCOL_11B|WIFI_PROTOCOL_11G|WIFI_PROT

Attention Support 802.11b or 802.11bg or 802.11bgn or LR mode
Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_IF: invalid interface
- others: refer to error codes in esp_err.h

Parameters
- `ifx`: interfaces
- `protocol_bitmap`: WiFi protocol bitmap

esp_err_t **esp_wifi_get_protocol** (*wifi_interface_t ifx*, uint8_t **protocol_bitmap*)
Get the current protocol bitmap of the specified interface.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_IF: invalid interface
- ESP_ERR_INVALID_ARG: invalid argument
- others: refer to error codes in esp_err.h

Parameters
- `ifx`: interface
- `[out] protocol_bitmap`: store current WiFi protocol bitmap of interface ifx

esp_err_t **esp_wifi_set_bandwidth** (*wifi_interface_t ifx*, *wifi_bandwidth_t bw*)
Set the bandwidth of ESP32 specified interface.

Attention 1. API return false if try to configure an interface that is not enabled
Attention 2. WIFI_BW_HT40 is supported only when the interface support 11N
Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_IF: invalid interface
- ESP_ERR_INVALID_ARG: invalid argument
- others: refer to error codes in esp_err.h

Parameters
- `ifx`: interface to be configured
- `bw`: bandwidth

esp_err_t **esp_wifi_get_bandwidth** (*wifi_interface_t ifx*, *wifi_bandwidth_t *bw*)
Get the bandwidth of ESP32 specified interface.

Attention 1. API return false if try to get a interface that is not enable
Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_IF: invalid interface
- ESP_ERR_INVALID_ARG: invalid argument

Parameters
- `ifx`: interface to be configured

- `[out] bw`: store bandwidth of interface ifx

esp_err_t **esp_wifi_set_channel**(uint8_t *primary*, *wifi_second_chan_t second*)

Set primary/secondary channel of ESP32.

Attention 1. This API should be called after esp_wifi_start()

Attention 2. When ESP32 is in STA mode, this API should not be called when STA is scanning or connecting to an external AP

Attention 3. When ESP32 is in softAP mode, this API should not be called when softAP has connected to external STAs

Attention 4. When ESP32 is in STA+softAP mode, this API should not be called when in the scenarios described above

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_IF: invalid interface
- ESP_ERR_INVALID_ARG: invalid argument

Parameters

- `primary`: for HT20, primary is the channel number, for HT40, primary is the primary channel
- `second`: for HT20, second is ignored, for HT40, second is the second channel

esp_err_t **esp_wifi_get_channel**(uint8_t **primary*, *wifi_second_chan_t *second*)

Get the primary/secondary channel of ESP32.

Attention 1. API return false if try to get a interface that is not enable

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters

- `primary`: store current primary channel
- `[out] second`: store current second channel

esp_err_t **esp_wifi_set_country**(**const** *wifi_country_t *country*)

configure country info

Attention 1. The default country is CHINA {.cc=" CN" , .schan=1, .nchan=13, .policy=WIFI_COUNTRY_POLICY_AUTO}.

Attention 2. The third octect of country code string is one of the following: ' ', 'O', 'I', 'X', otherwise it is considered as ' '.

Attention 3. When the country policy is WIFI_COUNTRY_POLICY_AUTO, the country info of the AP to which the station is connected is used. E.g. if the configured country info is {.cc=" US" , .schan=1, .nchan=11} and the country info of the AP to which the station is connected is {.cc=" JP" , .schan=1, .nchan=14} then the country info that will be used is {.cc=" JP" , .schan=1, .nchan=14}. If the station disconnected from the AP the country info is set back to the country info of the station automatically, {.cc=" US" , .schan=1, .nchan=11} in the example.

Attention 4. When the country policy is WIFI_COUNTRY_POLICY_MANUAL, then the configured country info is used always.

Attention 5. When the country info is changed because of configuration or because the station connects to a different external AP, the country IE in probe response/beacon of the soft-AP is also changed.

Attention 6. The country configuration is stored into flash.

Attention 7. This API doesn' t validate the per-country rules, it' s up to the user to fill in all fields according to local regulations.

Attention 8. When this API is called, the PHY init data will switch to the PHY init data type corresponding to the country info.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters

- `country`: the configured country info

esp_err_t **esp_wifi_get_country**(*wifi_country_t* **country*)
get the current country info

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters
- `country`: country info

esp_err_t **esp_wifi_set_mac**(*wifi_interface_t ifx*, **const** uint8_t *mac*[6])
Set MAC address of the ESP32 WiFi station or the soft-AP interface.

Attention 1. This API can only be called when the interface is disabled
Attention 2. ESP32 soft-AP and station have different MAC addresses, do not set them to be the same.
Attention 3. The bit 0 of the first byte of ESP32 MAC address can not be 1. For example, the MAC address can set to be "1a:XX:XX:XX:XX:XX", but can not be "15:XX:XX:XX:XX:XX".

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_WIFI_IF: invalid interface
- ESP_ERR_WIFI_MAC: invalid mac address
- ESP_ERR_WIFI_MODE: WiFi mode is wrong
- others: refer to error codes in esp_err.h

Parameters
- `ifx`: interface
- `mac`: the MAC address

esp_err_t **esp_wifi_get_mac**(*wifi_interface_t ifx*, uint8_t *mac*[6])
Get mac of specified interface.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_WIFI_IF: invalid interface

Parameters
- `ifx`: interface
- `[out] mac`: store mac of the interface ifx

esp_err_t **esp_wifi_set_promiscuous_rx_cb**(*wifi_promiscuous_cb_t cb*)
Register the RX callback function in the promiscuous mode.

Each time a packet is received, the registered callback function will be called.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters
- `cb`: callback

esp_err_t **esp_wifi_set_promiscuous**(bool *en*)
Enable the promiscuous mode.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters
- `en`: false - disable, true - enable

esp_err_t **esp_wifi_get_promiscuous**(bool **en*)
Get the promiscuous mode.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters
- `[out] en`: store the current status of promiscuous mode

esp_err_t **esp_wifi_set_promiscuous_filter**(**const** *wifi_promiscuous_filter_t *filter*)
Enable the promiscuous mode packet type filter.

Note The default filter is to filter all packets except WIFI_PKT_MISC

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters
- `filter`: the packet type filtered in promiscuous mode.

esp_err_t **esp_wifi_get_promiscuous_filter**(*wifi_promiscuous_filter_t *filter*)
Get the promiscuous filter.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters
- `[out] filter`: store the current status of promiscuous filter

esp_err_t **esp_wifi_set_promiscuous_ctrl_filter**(**const** *wifi_promiscuous_filter_t *filter*)
Enable subtype filter of the control packet in promiscuous mode.

Note The default filter is to filter none control packet.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters
- `filter`: the subtype of the control packet filtered in promiscuous mode.

esp_err_t **esp_wifi_get_promiscuous_ctrl_filter**(*wifi_promiscuous_filter_t *filter*)
Get the subtype filter of the control packet in promiscuous mode.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: invalid argument

Parameters
- `[out] filter`: store the current status of subtype filter of the control packet in promiscuous mode

esp_err_t **esp_wifi_set_config**(*wifi_interface_t interface*, *wifi_config_t *conf*)
Set the configuration of the ESP32 STA or AP.

Attention 1. This API can be called only when specified interface is enabled, otherwise, API fail

Attention 2. For station configuration, bssid_set needs to be 0; and it needs to be 1 only when users need to check the MAC address of the AP.

Attention 3. ESP32 is limited to only one channel, so when in the soft-AP+station mode, the soft-AP will adjust its channel automatically to be the same as the channel of the ESP32 station.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_WIFI_IF: invalid interface
- ESP_ERR_WIFI_MODE: invalid mode
- ESP_ERR_WIFI_PASSWORD: invalid password

- ESP_ERR_WIFI_NVS: WiFi internal NVS error
- others: refer to the erro code in esp_err.h

Parameters

- `interface`: interface
- `conf`: station or soft-AP configuration

esp_err_t **esp_wifi_get_config**(*wifi_interface_t interface*, *wifi_config_t *conf*)

Get configuration of specified interface.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_WIFI_IF: invalid interface

Parameters

- `interface`: interface
- `[out] conf`: station or soft-AP configuration

esp_err_t **esp_wifi_ap_get_sta_list**(*wifi_sta_list_t *sta*)

Get STAs associated with soft-AP.

Attention SSC only API

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_WIFI_MODE: WiFi mode is wrong
- ESP_ERR_WIFI_CONN: WiFi internal error, the station/soft-AP control block is invalid

Parameters

- `[out] sta`: station list ap can get the connected sta's phy mode info through the struct member phy_11b, phy_11g, phy_11n, phy_lr in the *wifi_sta_info_t* struct. For example, phy_11b = 1 imply that sta support 802.11b mode

esp_err_t **esp_wifi_ap_get_sta_aid**(**const** uint8_t *mac*[6], uint16_t *aid*)

Get AID of STA connected with soft-AP.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument
- ESP_ERR_NOT_FOUND: Requested resource not found
- ESP_ERR_WIFI_MODE: WiFi mode is wrong
- ESP_ERR_WIFI_CONN: WiFi internal error, the station/soft-AP control block is invalid

Parameters

- `mac`: STA's mac address
- `[out] aid`: Store the AID corresponding to STA mac

esp_err_t **esp_wifi_set_storage**(*wifi_storage_t storage*)

Set the WiFi API configuration storage type.

Attention 1. The default value is WIFI_STORAGE_FLASH

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_INVALID_ARG: invalid argument

Parameters

- `storage`: : storage type

esp_err_t **esp_wifi_set_vendor_ie**(bool *enable*, *wifi_vendor_ie_type_t type*, *wifi_vendor_ie_id_t idx*, **const** void *vnd_ie*)

Set 802.11 Vendor-Specific Information Element.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init()
- ESP_ERR_INVALID_ARG: Invalid argument, including if first byte of vnd_ie is not WIFI_VENDOR_IE_ELEMENT_ID (0xDD) or second byte is an invalid length.
- ESP_ERR_NO_MEM: Out of memory

Parameters

- `enable`: If true, specified IE is enabled. If false, specified IE is removed.
- `type`: Information Element type. Determines the frame type to associate with the IE.
- `idx`: Index to set or clear. Each IE type can be associated with up to two elements (indices 0 & 1).
- `vnd_ie`: Pointer to vendor specific element data. First 6 bytes should be a header with fields matching *vendor_ie_data_t*. If enable is false, this argument is ignored and can be NULL. Data does not need to remain valid after the function returns.

esp_err_t **esp_wifi_set_vendor_ie_cb** (*esp_vendor_ie_cb_t cb*, void **ctx*)

Register Vendor-Specific Information Element monitoring callback.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters

- `cb`: Callback function
- `ctx`: Context argument, passed to callback function.

esp_err_t **esp_wifi_set_max_tx_power** (int8_t *power*)

Set maximum transmitting power after WiFi start.

Attention 1. Maximum power before wifi startup is limited by PHY init data bin.

Attention 2. The value set by this API will be mapped to the max_tx_power of the structure *wifi_country_t* variable.

Attention 3. Mapping Table {Power, max_tx_power} = {{8, 2}, {20, 5}, {28, 7}, {34, 8}, {44, 11}, {52, 13}, {56, 14}, {60, 15}, {66, 16}, {72, 18}, {80, 20}}.

Attention 4. Param power unit is 0.25dBm, range is [8, 84] corresponding to 2dBm - 20dBm.

Attention 5. Relationship between set value and actual value. As follows: {set value range, actual value} = {{[8, 19],8}, {[20, 27],20}, {[28, 33],28}, {[34, 43],34}, {[44, 51],44}, {[52, 55],52}, {[56, 59],56}, {[60, 65],60}, {[66, 71],66}, {[72, 79],72}, {[80, 84],80}}.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start
- ESP_ERR_WIFI_ARG: invalid argument, e.g. parameter is out of range

Parameters

- `power`: Maximum WiFi transmitting power.

esp_err_t **esp_wifi_get_max_tx_power** (int8_t **power*)

Get maximum transmiting power after WiFi start.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start
- ESP_ERR_WIFI_ARG: invalid argument

Parameters

- `power`: Maximum WiFi transmitting power, unit is 0.25dBm.

esp_err_t **esp_wifi_set_event_mask** (uint32_t *mask*)

Set mask to enable or disable some WiFi events.

Attention 1. Mask can be created by logical OR of various WIFI_EVENT_MASK_ constants. Events which have corresponding bit set in the mask will not be delivered to the system event handler.

Attention 2. Default WiFi event mask is WIFI_EVENT_MASK_AP_PROBEREQRECVED.

Attention 3. There may be lots of stations sending probe request data around. Don't unmask this event unless you need to receive probe request data.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters
- `mask`: WiFi event mask.

esp_err_t **esp_wifi_get_event_mask**(uint32_t **mask*)
Get mask of WiFi events.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: invalid argument

Parameters
- `mask`: WiFi event mask.

esp_err_t **esp_wifi_80211_tx**(*wifi_interface_t ifx*, **const** void **buffer*, int *len*, bool *en_sys_seq*)
Send raw ieee80211 data.

Attention Currently only support for sending beacon/probe request/probe response/action and non-QoS data frame

Return
- ESP_OK: success
- ESP_ERR_WIFI_IF: Invalid interface
- ESP_ERR_INVALID_ARG: Invalid parameter
- ESP_ERR_WIFI_NO_MEM: out of memory

Parameters
- `ifx`: interface if the Wi-Fi mode is Station, the ifx should be WIFI_IF_STA. If the Wi-Fi mode is SoftAP, the ifx should be WIFI_IF_AP. If the Wi-Fi mode is Station+SoftAP, the ifx should be WIFI_IF_STA or WIFI_IF_AP. If the ifx is wrong, the API returns ESP_ERR_WIFI_IF.
- `buffer`: raw ieee80211 buffer
- `len`: the length of raw buffer, the len must be <= 1500 Bytes and >= 24 Bytes
- `en_sys_seq`: indicate whether use the internal sequence number. If en_sys_seq is false, the sequence in raw buffer is unchanged, otherwise it will be overwritten by WiFi driver with the system sequence number. Generally, if esp_wifi_80211_tx is called before the Wi-Fi connection has been set up, both en_sys_seq==true and en_sys_seq==false are fine. However, if the API is called after the Wi-Fi connection has been set up, en_sys_seq must be true, otherwise ESP_ERR_WIFI_ARG is returned.

esp_err_t **esp_wifi_set_csi_rx_cb**(*wifi_csi_cb_t cb*, void **ctx*)
Register the RX callback function of CSI data.

Each time a CSI data is received, the callback function will be called.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init

Parameters
- `cb`: callback
- `ctx`: context argument, passed to callback function

esp_err_t **esp_wifi_set_csi_config**(**const** *wifi_csi_config_t* **config*)
Set CSI data configuration.

return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start or promiscuous mode is not enabled
- ESP_ERR_INVALID_ARG: invalid argument

Parameters
- `config`: configuration

esp_err_t **esp_wifi_set_csi**(bool *en*)
Enable or disable CSI.

return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start or promiscuous mode is not enabled
- ESP_ERR_INVALID_ARG: invalid argument

Parameters
- `en`: true - enable, false - disable

esp_err_t **esp_wifi_set_ant_gpio**(**const** *wifi_ant_gpio_config_t* **config*)
Set antenna GPIO configuration.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: Invalid argument, e.g. parameter is NULL, invalid GPIO number etc

Parameters
- `config`: Antenna GPIO configuration.

esp_err_t **esp_wifi_get_ant_gpio**(*wifi_ant_gpio_config_t* **config*)
Get current antenna GPIO configuration.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: invalid argument, e.g. parameter is NULL

Parameters
- `config`: Antenna GPIO configuration.

esp_err_t **esp_wifi_set_ant**(**const** *wifi_ant_config_t* **config*)
Set antenna configuration.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: Invalid argument, e.g. parameter is NULL, invalid antenna mode or invalid GPIO number

Parameters
- `config`: Antenna configuration.

esp_err_t **esp_wifi_get_ant**(*wifi_ant_config_t* **config*)
Get current antenna configuration.

Return
- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: invalid argument, e.g. parameter is NULL

Parameters
- `config`: Antenna configuration.

int64_t **esp_wifi_get_tsf_time**(*wifi_interface_t interface*)
Get the TSF time In Station mode or SoftAP+Station mode if station is not connected or station doesn't receive at least one beacon after connected, will return 0.

Attention Enabling power save may cause the return value inaccurate, except WiFi modem sleep
Return 0 or the TSF time
Parameters
- `interface`: The interface whose tsf_time is to be retrieved.

esp_err_t **esp_wifi_set_inactive_time** (*wifi_interface_t ifx*, uint16_t *sec*)

Set the inactive time of the ESP32 STA or AP.

Attention 1. For Station, If the station does not receive a beacon frame from the connected SoftAP during the inactive time, disconnect from SoftAP. Default 6s.

Attention 2. For SoftAP, If the softAP doesn't receive any data from the connected STA during inactive time, the softAP will force deauth the STA. Default is 300s.

Attention 3. The inactive time configuration is not stored into flash

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_NOT_STARTED: WiFi is not started by esp_wifi_start
- ESP_ERR_WIFI_ARG: invalid argument, For Station, if sec is less than 3. For SoftAP, if sec is less than 10.

Parameters

- `ifx`: interface to be configured.
- `sec`: Inactive time. Unit seconds.

esp_err_t **esp_wifi_get_inactive_time** (*wifi_interface_t ifx*, uint16_t **sec*)

Get inactive time of specified interface.

Return

- ESP_OK: succeed
- ESP_ERR_WIFI_NOT_INIT: WiFi is not initialized by esp_wifi_init
- ESP_ERR_WIFI_ARG: invalid argument

Parameters

- `ifx`: Interface to be configured.
- `sec`: Inactive time. Unit seconds.

esp_err_t **esp_wifi_statis_dump** (uint32_t *modules*)

Dump WiFi statistics.

Return

- ESP_OK: succeed
- others: failed

Parameters

- `modules`: statistic modules to be dumped

esp_err_t **esp_wifi_disable_pmf_config** (*wifi_interface_t ifx*)

Disable PMF configuration for specified interface.

Attention This API should be called after esp_wifi_set_config() and before esp_wifi_start().

Return

- ESP_OK: succeed
- others: failed

Parameters

- `ifx`: Interface to be configured.

Structures

struct wifi_init_config_t

WiFi stack configuration parameters passed to esp_wifi_init call.

Public Members

system_event_handler_t **event_handler**

WiFi event handler

wifi_osi_funcs_t ***osi_funcs**

WiFi OS functions

wpa_crypto_funcs_t **wpa_crypto_funcs**
WiFi station crypto functions when connect

int **static_rx_buf_num**
WiFi static RX buffer number

int **dynamic_rx_buf_num**
WiFi dynamic RX buffer number

int **tx_buf_type**
WiFi TX buffer type

int **static_tx_buf_num**
WiFi static TX buffer number

int **dynamic_tx_buf_num**
WiFi dynamic TX buffer number

int **cache_tx_buf_num**
WiFi TX cache buffer number

int **csi_enable**
WiFi channel state information enable flag

int **ampdu_rx_enable**
WiFi AMPDU RX feature enable flag

int **ampdu_tx_enable**
WiFi AMPDU TX feature enable flag

int **nvs_enable**
WiFi NVS flash enable flag

int **nano_enable**
Nano option for printf/scan family enable flag

int **rx_ba_win**
WiFi Block Ack RX window size

int **wifi_task_core_id**
WiFi Task Core ID

int **beacon_max_len**
WiFi softAP maximum length of the beacon

int **mgmt_sbuf_num**
WiFi management short buffer number, the minimum value is 6, the maximum value is 32

uint64_t **feature_caps**
Enables additional WiFi features and capabilities

int **magic**
WiFi init magic number, it should be the last field

Macros

ESP_ERR_WIFI_NOT_INIT
WiFi driver was not installed by esp_wifi_init

ESP_ERR_WIFI_NOT_STARTED
WiFi driver was not started by esp_wifi_start

ESP_ERR_WIFI_NOT_STOPPED
WiFi driver was not stopped by esp_wifi_stop

ESP_ERR_WIFI_IF
WiFi interface error

ESP_ERR_WIFI_MODE
WiFi mode error

ESP_ERR_WIFI_STATE
WiFi internal state error

ESP_ERR_WIFI_CONN
WiFi internal control block of station or soft-AP error

ESP_ERR_WIFI_NVS
WiFi internal NVS module error

ESP_ERR_WIFI_MAC
MAC address is invalid

ESP_ERR_WIFI_SSID
SSID is invalid

ESP_ERR_WIFI_PASSWORD
Password is invalid

ESP_ERR_WIFI_TIMEOUT
Timeout error

ESP_ERR_WIFI_WAKE_FAIL
WiFi is in sleep state(RF closed) and wakeup fail

ESP_ERR_WIFI_WOULD_BLOCK
The caller would block

ESP_ERR_WIFI_NOT_CONNECT
Station still in disconnect status

ESP_ERR_WIFI_POST
Failed to post the event to WiFi task

ESP_ERR_WIFI_INIT_STATE
Invalid WiFi state when init/deinit is called

ESP_ERR_WIFI_STOP_STATE
Returned when WiFi is stopping

ESP_ERR_WIFI_NOT_ASSOC
The WiFi connection is not associated

ESP_ERR_WIFI_TX_DISALLOW
The WiFi TX is disallowed

WIFI_STATIC_TX_BUFFER_NUM

WIFI_CACHE_TX_BUFFER_NUM

WIFI_DYNAMIC_TX_BUFFER_NUM

WIFI_CSI_ENABLED

WIFI_AMPDU_RX_ENABLED

WIFI_AMPDU_TX_ENABLED

WIFI_NVS_ENABLED

WIFI_NANO_FORMAT_ENABLED

WIFI_INIT_CONFIG_MAGIC

WIFI_DEFAULT_RX_BA_WIN

WIFI_TASK_CORE_ID

WIFI_SOFTAP_BEACON_MAX_LEN

WIFI_MGMT_SBUF_NUM

CONFIG_FEATURE_WPA3_SAE_BIT

CONFIG_FEATURE_CACHE_TX_BUF_BIT

WIFI_INIT_CONFIG_DEFAULT()

Type Definitions

typedef void (***wifi_promiscuous_cb_t**)(void *buf, *wifi_promiscuous_pkt_type_t* type)

The RX callback function in the promiscuous mode. Each time a packet is received, the callback function will be called.

Parameters

- `buf`: Data received. Type of data in buffer (*wifi_promiscuous_pkt_t* or *wifi_pkt_rx_ctrl_t*) indicated by 'type' parameter.
- `type`: promiscuous packet type.

typedef void (***esp_vendor_ie_cb_t**)(void *ctx, *wifi_vendor_ie_type_t* type, **const** uint8_t sa[6], **const** *vendor_ie_data_t* *vnd_ie, int rssi)

Function signature for received Vendor-Specific Information Element callback.

Parameters

- `ctx`: Context argument, as passed to esp_wifi_set_vendor_ie_cb() when registering callback.
- `type`: Information element type, based on frame type received.
- `sa`: Source 802.11 address.
- `vnd_ie`: Pointer to the vendor specific element data received.
- `rssi`: Received signal strength indication.

typedef void (***wifi_csi_cb_t**)(void *ctx, *wifi_csi_info_t* *data)

The RX callback function of Channel State Information(CSI) data.

Each time a CSI data is received, the callback function will be called.

Parameters

- `ctx`: context argument, passed to esp_wifi_set_csi_rx_cb() when registering callback function.
- `data`: CSI data received. The memory that it points to will be deallocated after callback function returns.

Header File

- esp_wifi/include/esp_wifi_types.h

Unions

union wifi_config_t

#include <esp_wifi_types.h> Configuration data for ESP32 AP or STA.

The usage of this union (for ap or sta configuration) is determined by the accompanying interface argument passed to esp_wifi_set_config() or esp_wifi_get_config()

Public Members

wifi_ap_config_t **ap**

configuration of AP

wifi_sta_config_t **sta**

configuration of STA

Structures

struct wifi_country_t

Structure describing WiFi country-based regional restrictions.

Public Members

char **cc**[3]
country code string

uint8_t **schan**
start channel

uint8_t **nchan**
total channel number

int8_t **max_tx_power**
This field is used for getting WiFi maximum transmitting power, call esp_wifi_set_max_tx_power to set the maximum transmitting power.

wifi_country_policy_t **policy**
country policy

struct wifi_active_scan_time_t
Range of active scan times per channel.

Public Members

uint32_t **min**
minimum active scan time per channel, units: millisecond

uint32_t **max**
maximum active scan time per channel, units: millisecond, values above 1500ms may cause station to disconnect from AP and are not recommended.

struct wifi_scan_time_t
Aggregate of active & passive scan time per channel.

Public Members

wifi_active_scan_time_t **active**
active scan time per channel, units: millisecond.

uint32_t **passive**
passive scan time per channel, units: millisecond, values above 1500ms may cause station to disconnect from AP and are not recommended.

struct wifi_scan_config_t
Parameters for an SSID scan.

Public Members

uint8_t ***ssid**
SSID of AP

uint8_t ***bssid**
MAC address of AP

uint8_t **channel**
channel, scan the specific channel

bool **show_hidden**
enable to scan AP whose SSID is hidden

wifi_scan_type_t **scan_type**
scan type, active or passive

wifi_scan_time_t **scan_time**
scan time per channel

struct wifi_ap_record_t
Description of a WiFi AP.

Public Members

uint8_t **bssid**[6]
MAC address of AP

uint8_t **ssid**[33]
SSID of AP

uint8_t **primary**
channel of AP

wifi_second_chan_t **second**
secondary channel of AP

int8_t **rssi**
signal strength of AP

wifi_auth_mode_t **authmode**
authmode of AP

wifi_cipher_type_t **pairwise_cipher**
pairwise cipher of AP

wifi_cipher_type_t **group_cipher**
group cipher of AP

wifi_ant_t **ant**
antenna used to receive beacon from AP

uint32_t **phy_11b** : 1
bit: 0 flag to identify if 11b mode is enabled or not

uint32_t **phy_11g** : 1
bit: 1 flag to identify if 11g mode is enabled or not

uint32_t **phy_11n** : 1
bit: 2 flag to identify if 11n mode is enabled or not

uint32_t **phy_lr** : 1
bit: 3 flag to identify if low rate is enabled or not

uint32_t **wps** : 1
bit: 4 flag to identify if WPS is supported or not

uint32_t **reserved** : 27
bit: 5..31 reserved

wifi_country_t **country**
country information of AP

struct wifi_scan_threshold_t
Structure describing parameters for a WiFi fast scan.

Public Members

int8_t **rssi**
The minimum rssi to accept in the fast scan mode

wifi_auth_mode_t **authmode**
The weakest authmode to accept in the fast scan mode

struct wifi_pmf_config_t
Configuration structure for Protected Management Frame

Public Members

bool **capable**
Deprecated variable. Device will always connect in PMF mode if other device also advertizes PMF capability.

bool **required**
Advertizes that Protected Management Frame is required. Device will not associate to non-PMF capable devices.

struct wifi_ap_config_t
Soft-AP configuration settings for the ESP32.

Public Members

uint8_t **ssid**[32]
SSID of ESP32 soft-AP. If ssid_len field is 0, this must be a Null terminated string. Otherwise, length is set according to ssid_len.

uint8_t **password**[64]
Password of ESP32 soft-AP.

uint8_t **ssid_len**
Optional length of SSID field.

uint8_t **channel**
Channel of ESP32 soft-AP

wifi_auth_mode_t **authmode**
Auth mode of ESP32 soft-AP. Do not support AUTH_WEP in soft-AP mode

uint8_t **ssid_hidden**
Broadcast SSID or not, default 0, broadcast the SSID

uint8_t **max_connection**
Max number of stations allowed to connect in, default 4, max 10

uint16_t **beacon_interval**
Beacon interval which should be multiples of 100. Unit: TU(time unit, 1 TU = 1024 us). Range: 100 ~ 60000. Default value: 100

struct wifi_sta_config_t
STA configuration settings for the ESP32.

Public Members

uint8_t **ssid**[32]
SSID of target AP.

uint8_t **password**[64]
Password of target AP.

wifi_scan_method_t **scan_method**
do all channel scan or fast scan

bool **bssid_set**
whether set MAC address of target AP or not. Generally, station_config.bssid_set needs to be 0; and it needs to be 1 only when users need to check the MAC address of the AP.

uint8_t **bssid**[6]
MAC address of target AP

uint8_t **channel**
channel of target AP. Set to 1~13 to scan starting from the specified channel before connecting to AP. If the channel of AP is unknown, set it to 0.

uint16_t **listen_interval**
Listen interval for ESP32 station to receive beacon when WIFI_PS_MAX_MODEM is set. Units: AP beacon intervals. Defaults to 3 if set to 0.

wifi_sort_method_t **sort_method**
sort the connect AP in the list by rssi or security mode

wifi_scan_threshold_t **threshold**
When sort_method is set, only APs which have an auth mode that is more secure than the selected auth mode and a signal stronger than the minimum RSSI will be used.

wifi_pmf_config_t **pmf_cfg**
Configuration for Protected Management Frame. Will be advertized in RSN Capabilities in RSN IE.

struct wifi_sta_info_t
Description of STA associated with AP.

Public Members

uint8_t **mac**[6]
mac address

int8_t **rssi**
current average rssi of sta connected

uint32_t **phy_11b** : 1
bit: 0 flag to identify if 11b mode is enabled or not

uint32_t **phy_11g** : 1
bit: 1 flag to identify if 11g mode is enabled or not

uint32_t **phy_11n** : 1
bit: 2 flag to identify if 11n mode is enabled or not

uint32_t **phy_lr** : 1
bit: 3 flag to identify if low rate is enabled or not

uint32_t **is_mesh_child** : 1
bit: 4 flag to identify mesh child

uint32_t **reserved** : 27
bit: 5..31 reserved

struct wifi_sta_list_t
List of stations associated with the ESP32 Soft-AP.

Public Members

wifi_sta_info_t **sta**[**ESP_WIFI_MAX_CONN_NUM**]
station list

int **num**
number of stations in the list (other entries are invalid)

struct vendor_ie_data_t
Vendor Information Element header.

The first bytes of the Information Element will match this header. Payload follows.

Public Members

uint8_t **element_id**
Should be set to WIFI_VENDOR_IE_ELEMENT_ID (0xDD)

uint8_t **length**
Length of all bytes in the element data following this field. Minimum 4.

uint8_t **vendor_oui**[3]
Vendor identifier (OUI).

uint8_t **vendor_oui_type**
Vendor-specific OUI type.

uint8_t **payload**[0]
Payload. Length is equal to value in ‘length’ field, minus 4.

struct wifi_pkt_rx_ctrl_t
Received packet radio metadata header, this is the common header at the beginning of all promiscuous mode RX callback buffers.

Public Members

signed **rssi** : 8
Received Signal Strength Indicator(RSSI) of packet. unit: dBm

unsigned **rate** : 5
PHY rate encoding of the packet. Only valid for non HT(11bg) packet

unsigned **__pad0__** : 1
reserved

unsigned **sig_mode** : 2
0: non HT(11bg) packet; 1: HT(11n) packet; 3: VHT(11ac) packet

unsigned **__pad1__** : 16
reserved

unsigned **mcs** : 7
Modulation Coding Scheme. If is HT(11n) packet, shows the modulation, range from 0 to 76(MSC0 ~ MCS76)

unsigned **cwb** : 1
Channel Bandwidth of the packet. 0: 20MHz; 1: 40MHz

unsigned **__pad2__** : 16
reserved

unsigned **smoothing** : 1
reserved

unsigned **not_sounding** : 1
reserved

unsigned **__pad3__** : 1
reserved

unsigned **aggregation** : 1
Aggregation. 0: MPDU packet; 1: AMPDU packet

unsigned **stbc** : 2
Space Time Block Code(STBC). 0: non STBC packet; 1: STBC packet

unsigned **fec_coding** : 1
Flag is set for 11n packets which are LDPC

unsigned **sgi** : 1
Short Guide Interval(SGI). 0: Long GI; 1: Short GI

unsigned **__pad4__** : 8
reserved

unsigned **ampdu_cnt** : 8
ampdu cnt

unsigned **channel** : 4
primary channel on which this packet is received

unsigned **secondary_channel** : 4
secondary channel on which this packet is received. 0: none; 1: above; 2: below

unsigned **__pad5__** : 8
reserved

unsigned **timestamp** : 32
timestamp. The local time when this packet is received. It is precise only if modem sleep or light sleep is not enabled. unit: microsecond

unsigned **__pad6__** : 32
reserved

unsigned **__pad7__** : 32
reserved

unsigned **__pad8__** : 31
reserved

unsigned **ant** : 1
antenna number from which this packet is received. 0: WiFi antenna 0; 1: WiFi antenna 1

signed **noise_floor** : 8
noise floor of Radio Frequency Module(RF). unit: 0.25dBm

unsigned **__pad9__** : 24
reserved

unsigned **sig_len** : 12
length of packet including Frame Check Sequence(FCS)

unsigned **__pad10__** : 12
reserved

unsigned **rx_state** : 8
state of the packet. 0: no error; others: error numbers which are not public

struct wifi_promiscuous_pkt_t
Payload passed to ‘buf’ parameter of promiscuous mode RX callback.

Public Members

wifi_pkt_rx_ctrl_t **rx_ctrl**
metadata header

uint8_t **payload**[0]
Data or management payload. Length of payload is described by rx_ctrl.sig_len. Type of content determined by packet type argument of callback.

struct wifi_promiscuous_filter_t
Mask for filtering different packet types in promiscuous mode.

Public Members

uint32_t **filter_mask**
OR of one or more filter values WIFI_PROMIS_FILTER_*

struct wifi_csi_config_t
Channel state information(CSI) configuration type.

Public Members

bool **lltf_en**
enable to receive legacy long training field(lltf) data. Default enabled

bool **htltf_en**
enable to receive HT long training field(htltf) data. Default enabled

bool **stbc_htltf2_en**
enable to receive space time block code HT long training field(stbc-htltf2) data. Default enabled

bool **ltf_merge_en**
enable to generate htlft data by averaging lltf and ht_ltf data when receiving HT packet. Otherwise, use ht_ltf data directly. Default enabled

bool **channel_filter_en**
enable to turn on channel filter to smooth adjacent sub-carrier. Disable it to keep independence of adjacent sub-carrier. Default enabled

bool **manu_scale**
manually scale the CSI data by left shifting or automatically scale the CSI data. If set true, please set the shift bits. false: automatically. true: manually. Default false

uint8_t **shift**
manually left shift bits of the scale of the CSI data. The range of the left shift bits is 0~15

struct wifi_csi_info_t
CSI data type.

Public Members

wifi_pkt_rx_ctrl_t **rx_ctrl**
received packet radio metadata header of the CSI data

uint8_t **mac**[6]
source MAC address of the CSI data

bool **first_word_invalid**
first four bytes of the CSI data is invalid or not

int8_t ***buf**
buffer of CSI data

uint16_t **len**
length of CSI data

struct wifi_ant_gpio_t
WiFi GPIO configuration for antenna selection.

Public Members

uint8_t **gpio_select** : 1
Whether this GPIO is connected to external antenna switch

uint8_t **gpio_num** : 7
The GPIO number that connects to external antenna switch

struct wifi_ant_gpio_config_t
WiFi GPIOs configuration for antenna selection.

Public Members

wifi_ant_gpio_t **gpio_cfg**[4]
The configurations of GPIOs that connect to external antenna switch

struct wifi_ant_config_t
WiFi antenna configuration.

Public Members

wifi_ant_mode_t **rx_ant_mode**
WiFi antenna mode for receiving

wifi_ant_t **rx_ant_default**
Default antenna mode for receiving, it’ s ignored if rx_ant_mode is not WIFI_ANT_MODE_AUTO

wifi_ant_mode_t **tx_ant_mode**
WiFi antenna mode for transmission, it can be set to WIFI_ANT_MODE_AUTO only if rx_ant_mode is set to WIFI_ANT_MODE_AUTO

uint8_t **enabled_ant0** : 4
Index (in antenna GPIO configuration) of enabled WIFI_ANT_MODE_ANT0

uint8_t **enabled_ant1** : 4
Index (in antenna GPIO configuration) of enabled WIFI_ANT_MODE_ANT1

struct wifi_event_sta_scan_done_t
Argument structure for WIFI_EVENT_SCAN_DONE event

Public Members

uint32_t **status**
status of scanning APs: 0 —success, 1 - failure

uint8_t **number**
number of scan results

uint8_t **scan_id**
scan sequence number, used for block scan

struct wifi_event_sta_connected_t
Argument structure for WIFI_EVENT_STA_CONNECTED event

Public Members

uint8_t **ssid**[32]
SSID of connected AP

uint8_t **ssid_len**
SSID length of connected AP

uint8_t **bssid**[6]
BSSID of connected AP

uint8_t **channel**
channel of connected AP

wifi_auth_mode_t **authmode**
authentication mode used by AP

struct wifi_event_sta_disconnected_t
Argument structure for WIFI_EVENT_STA_DISCONNECTED event

Public Members

uint8_t **ssid**[32]
SSID of disconnected AP

uint8_t **ssid_len**
SSID length of disconnected AP

uint8_t **bssid**[6]
BSSID of disconnected AP

uint8_t **reason**
reason of disconnection

struct wifi_event_sta_authmode_change_t
Argument structure for WIFI_EVENT_STA_AUTHMODE_CHANGE event

Public Members

wifi_auth_mode_t **old_mode**
the old auth mode of AP

wifi_auth_mode_t **new_mode**
the new auth mode of AP

struct wifi_event_sta_wps_er_pin_t
Argument structure for WIFI_EVENT_STA_WPS_ER_PIN event

Public Members

uint8_t **pin_code**[8]
PIN code of station in enrollee mode

struct wifi_event_sta_wps_er_success_t
Argument structure for WIFI_EVENT_STA_WPS_ER_SUCCESS event

Public Members

uint8_t **ap_cred_cnt**
Number of AP credentials received

uint8_t **ssid**[**MAX_SSID_LEN**]
SSID of AP

uint8_t **passphrase**[**MAX_PASSPHRASE_LEN**]
Passphrase for the AP

struct *wifi_event_sta_wps_er_success_t*::**[anonymous]** **ap_cred**[**MAX_WPS_AP_CRED**]
All AP credentials received from WPS handshake

struct wifi_event_ap_staconnected_t
Argument structure for WIFI_EVENT_AP_STACONNECTED event

Public Members

uint8_t **mac**[6]
MAC address of the station connected to ESP32 soft-AP

uint8_t **aid**
the aid that ESP32 soft-AP gives to the station connected to

bool **is_mesh_child**
flag to identify mesh child

struct wifi_event_ap_stadisconnected_t
Argument structure for WIFI_EVENT_AP_STADISCONNECTED event

Public Members

uint8_t **mac**[6]
MAC address of the station disconnects to ESP32 soft-AP

uint8_t **aid**
the aid that ESP32 soft-AP gave to the station disconnects to

bool **is_mesh_child**
flag to identify mesh child

struct wifi_event_ap_probe_req_rx_t
Argument structure for WIFI_EVENT_AP_PROBEREQRECVED event

Public Members

int **rssi**
Received probe request signal strength

uint8_t **mac**[6]
MAC address of the station which send probe request

Macros

WIFI_IF_STA
WIFI_IF_AP

WIFI_PROTOCOL_11B

WIFI_PROTOCOL_11G

WIFI_PROTOCOL_11N

WIFI_PROTOCOL_LR

ESP_WIFI_MAX_CONN_NUM
max number of stations which can connect to ESP32 soft-AP

WIFI_VENDOR_IE_ELEMENT_ID

WIFI_PROMIS_FILTER_MASK_ALL
filter all packets

WIFI_PROMIS_FILTER_MASK_MGMT
filter the packets with type of WIFI_PKT_MGMT

WIFI_PROMIS_FILTER_MASK_CTRL
filter the packets with type of WIFI_PKT_CTRL

WIFI_PROMIS_FILTER_MASK_DATA
filter the packets with type of WIFI_PKT_DATA

WIFI_PROMIS_FILTER_MASK_MISC
filter the packets with type of WIFI_PKT_MISC

WIFI_PROMIS_FILTER_MASK_DATA_MPDU
filter the MPDU which is a kind of WIFI_PKT_DATA

WIFI_PROMIS_FILTER_MASK_DATA_AMPDU
filter the AMPDU which is a kind of WIFI_PKT_DATA

WIFI_PROMIS_FILTER_MASK_FCSFAIL
filter the FCS failed packets, do not open it in general

WIFI_PROMIS_CTRL_FILTER_MASK_ALL
filter all control packets

WIFI_PROMIS_CTRL_FILTER_MASK_WRAPPER
filter the control packets with subtype of Control Wrapper

WIFI_PROMIS_CTRL_FILTER_MASK_BAR
filter the control packets with subtype of Block Ack Request

WIFI_PROMIS_CTRL_FILTER_MASK_BA
filter the control packets with subtype of Block Ack

WIFI_PROMIS_CTRL_FILTER_MASK_PSPOLL
filter the control packets with subtype of PS-Poll

WIFI_PROMIS_CTRL_FILTER_MASK_RTS
filter the control packets with subtype of RTS

WIFI_PROMIS_CTRL_FILTER_MASK_CTS
filter the control packets with subtype of CTS

WIFI_PROMIS_CTRL_FILTER_MASK_ACK
filter the control packets with subtype of ACK

WIFI_PROMIS_CTRL_FILTER_MASK_CFEND
filter the control packets with subtype of CF-END

WIFI_PROMIS_CTRL_FILTER_MASK_CFENDACK
filter the control packets with subtype of CF-END+CF-ACK

WIFI_EVENT_MASK_ALL
mask all WiFi events

WIFI_EVENT_MASK_NONE
mask none of the WiFi events

WIFI_EVENT_MASK_AP_PROBEREQRECVED
mask SYSTEM_EVENT_AP_PROBEREQRECVED event

MAX_SSID_LEN

MAX_PASSPHRASE_LEN

MAX_WPS_AP_CRED

WIFI_STATIS_BUFFER

WIFI_STATIS_RXTX

WIFI_STATIS_HW

WIFI_STATIS_DIAG

WIFI_STATIS_ALL

Type Definitions

typedef esp_interface_t **wifi_interface_t**

Enumerations

enum wifi_mode_t

Values:

WIFI_MODE_NULL = 0
null mode

WIFI_MODE_STA
WiFi station mode

WIFI_MODE_AP
WiFi soft-AP mode

WIFI_MODE_APSTA
WiFi station + soft-AP mode

WIFI_MODE_MAX

enum wifi_country_policy_t

Values:

WIFI_COUNTRY_POLICY_AUTO
Country policy is auto, use the country info of AP to which the station is connected

WIFI_COUNTRY_POLICY_MANUAL
Country policy is manual, always use the configured country info

enum wifi_auth_mode_t

Values:

WIFI_AUTH_OPEN = 0
authenticate mode : open

WIFI_AUTH_WEP
authenticate mode : WEP

WIFI_AUTH_WPA_PSK
authenticate mode : WPA_PSK

WIFI_AUTH_WPA2_PSK
authenticate mode : WPA2_PSK

WIFI_AUTH_WPA_WPA2_PSK
authenticate mode : WPA_WPA2_PSK

WIFI_AUTH_WPA2_ENTERPRISE
authenticate mode : WPA2_ENTERPRISE

WIFI_AUTH_WPA3_PSK
authenticate mode : WPA3_PSK

WIFI_AUTH_WPA2_WPA3_PSK
authenticate mode : WPA2_WPA3_PSK

WIFI_AUTH_MAX

enum wifi_err_reason_t

Values:

WIFI_REASON_UNSPECIFIED = 1

WIFI_REASON_AUTH_EXPIRE = 2

WIFI_REASON_AUTH_LEAVE = 3

WIFI_REASON_ASSOC_EXPIRE = 4

WIFI_REASON_ASSOC_TOOMANY = 5

WIFI_REASON_NOT_AUTHED = 6

WIFI_REASON_NOT_ASSOCED = 7

WIFI_REASON_ASSOC_LEAVE = 8

WIFI_REASON_ASSOC_NOT_AUTHED = 9

WIFI_REASON_DISASSOC_PWRCAP_BAD = 10

WIFI_REASON_DISASSOC_SUPCHAN_BAD = 11

WIFI_REASON_IE_INVALID = 13

WIFI_REASON_MIC_FAILURE = 14

WIFI_REASON_4WAY_HANDSHAKE_TIMEOUT = 15

WIFI_REASON_GROUP_KEY_UPDATE_TIMEOUT = 16

WIFI_REASON_IE_IN_4WAY_DIFFERS = 17

WIFI_REASON_GROUP_CIPHER_INVALID = 18

WIFI_REASON_PAIRWISE_CIPHER_INVALID = 19

WIFI_REASON_AKMP_INVALID = 20

WIFI_REASON_UNSUPP_RSN_IE_VERSION = 21

WIFI_REASON_INVALID_RSN_IE_CAP = 22

WIFI_REASON_802_1X_AUTH_FAILED = 23

WIFI_REASON_CIPHER_SUITE_REJECTED = 24

WIFI_REASON_INVALID_PMKID = 53

WIFI_REASON_BEACON_TIMEOUT = 200

WIFI_REASON_NO_AP_FOUND = 201

WIFI_REASON_AUTH_FAIL = 202

WIFI_REASON_ASSOC_FAIL = 203

WIFI_REASON_HANDSHAKE_TIMEOUT = 204

WIFI_REASON_CONNECTION_FAIL = 205

WIFI_REASON_AP_TSF_RESET = 206

WIFI_REASON_ASSOC_COMEBACK_TIME_TOO_LONG = 208

enum wifi_second_chan_t

Values:

WIFI_SECOND_CHAN_NONE = 0
the channel width is HT20

WIFI_SECOND_CHAN_ABOVE
the channel width is HT40 and the secondary channel is above the primary channel

WIFI_SECOND_CHAN_BELOW
the channel width is HT40 and the secondary channel is below the primary channel

enum wifi_scan_type_t

Values:

WIFI_SCAN_TYPE_ACTIVE = 0
active scan

WIFI_SCAN_TYPE_PASSIVE
passive scan

enum wifi_cipher_type_t

Values:

WIFI_CIPHER_TYPE_NONE = 0
the cipher type is none

WIFI_CIPHER_TYPE_WEP40
the cipher type is WEP40

WIFI_CIPHER_TYPE_WEP104
the cipher type is WEP104

WIFI_CIPHER_TYPE_TKIP
the cipher type is TKIP

WIFI_CIPHER_TYPE_CCMP
the cipher type is CCMP

WIFI_CIPHER_TYPE_TKIP_CCMP
the cipher type is TKIP and CCMP

WIFI_CIPHER_TYPE_AES_CMAC128
the cipher type is AES-CMAC-128

WIFI_CIPHER_TYPE_UNKNOWN
the cipher type is unknown

enum wifi_ant_t
WiFi antenna.

Values:

WIFI_ANT_ANT0
WiFi antenna 0

WIFI_ANT_ANT1
WiFi antenna 1

WIFI_ANT_MAX
Invalid WiFi antenna

enum wifi_scan_method_t
Values:

WIFI_FAST_SCAN = 0
Do fast scan, scan will end after find SSID match AP

WIFI_ALL_CHANNEL_SCAN
All channel scan, scan will end after scan all the channel

enum wifi_sort_method_t
Values:

WIFI_CONNECT_AP_BY_SIGNAL = 0
Sort match AP in scan list by RSSI

WIFI_CONNECT_AP_BY_SECURITY
Sort match AP in scan list by security mode

enum wifi_ps_type_t
Values:

WIFI_PS_NONE
No power save

WIFI_PS_MIN_MODEM
Minimum modem power saving. In this mode, station wakes up to receive beacon every DTIM period

WIFI_PS_MAX_MODEM
Maximum modem power saving. In this mode, interval to receive beacons is determined by the listen_interval parameter in *wifi_sta_config_t*

enum wifi_bandwidth_t

Values:

WIFI_BW_HT20 = 1

WIFI_BW_HT40

enum wifi_storage_t

Values:

WIFI_STORAGE_FLASH
all configuration will store in both memory and flash

WIFI_STORAGE_RAM
all configuration will only store in the memory

enum wifi_vendor_ie_type_t
Vendor Information Element type.

Determines the frame type that the IE will be associated with.

Values:

WIFI_VND_IE_TYPE_BEACON

WIFI_VND_IE_TYPE_PROBE_REQ

WIFI_VND_IE_TYPE_PROBE_RESP

WIFI_VND_IE_TYPE_ASSOC_REQ

WIFI_VND_IE_TYPE_ASSOC_RESP

enum wifi_vendor_ie_id_t
Vendor Information Element index.

Each IE type can have up to two associated vendor ID elements.

Values:

WIFI_VND_IE_ID_0

WIFI_VND_IE_ID_1

enum wifi_promiscuous_pkt_type_t
Promiscuous frame type.

Passed to promiscuous mode RX callback to indicate the type of parameter in the buffer.

Values:

WIFI_PKT_MGMT
Management frame, indicates ‘buf’ argument is *wifi_promiscuous_pkt_t*

WIFI_PKT_CTRL
Control frame, indicates ‘buf’ argument is *wifi_promiscuous_pkt_t*

WIFI_PKT_DATA
Data frame, indiciates ‘buf’ argument is *wifi_promiscuous_pkt_t*

WIFI_PKT_MISC
Other type, such as MIMO etc. ‘buf’ argument is *wifi_promiscuous_pkt_t* but the payload is zero length.

enum wifi_ant_mode_t
WiFi antenna mode.

Values:

WIFI_ANT_MODE_ANT0
Enable WiFi antenna 0 only

WIFI_ANT_MODE_ANT1
Enable WiFi antenna 1 only

WIFI_ANT_MODE_AUTO
Enable WiFi antenna 0 and 1, automatically select an antenna

WIFI_ANT_MODE_MAX
Invalid WiFi enabled antenna

enum wifi_phy_rate_t
WiFi PHY rate encodings.

Values:

WIFI_PHY_RATE_1M_L = 0x00
1 Mbps with long preamble

WIFI_PHY_RATE_2M_L = 0x01
2 Mbps with long preamble

WIFI_PHY_RATE_5M_L = 0x02
5.5 Mbps with long preamble

WIFI_PHY_RATE_11M_L = 0x03
11 Mbps with long preamble

WIFI_PHY_RATE_2M_S = 0x05
2 Mbps with short preamble

WIFI_PHY_RATE_5M_S = 0x06
5.5 Mbps with short preamble

WIFI_PHY_RATE_11M_S = 0x07
11 Mbps with short preamble

WIFI_PHY_RATE_48M = 0x08
48 Mbps

WIFI_PHY_RATE_24M = 0x09
24 Mbps

WIFI_PHY_RATE_12M = 0x0A
12 Mbps

WIFI_PHY_RATE_6M = 0x0B
6 Mbps

WIFI_PHY_RATE_54M = 0x0C
54 Mbps

WIFI_PHY_RATE_36M = 0x0D
36 Mbps

WIFI_PHY_RATE_18M = 0x0E
18 Mbps

WIFI_PHY_RATE_9M = 0x0F
9 Mbps

WIFI_PHY_RATE_MCS0_LGI = 0x10
MCS0 with long GI, 6.5 Mbps for 20MHz, 13.5 Mbps for 40MHz

WIFI_PHY_RATE_MCS1_LGI = 0x11
MCS1 with long GI, 13 Mbps for 20MHz, 27 Mbps for 40MHz

WIFI_PHY_RATE_MCS2_LGI = 0x12
MCS2 with long GI, 19.5 Mbps for 20MHz, 40.5 Mbps for 40MHz

WIFI_PHY_RATE_MCS3_LGI = 0x13
MCS3 with long GI, 26 Mbps for 20MHz, 54 Mbps for 40MHz

WIFI_PHY_RATE_MCS4_LGI = 0x14
MCS4 with long GI, 39 Mbps for 20MHz, 81 Mbps for 40MHz

WIFI_PHY_RATE_MCS5_LGI = 0x15
MCS5 with long GI, 52 Mbps for 20MHz, 108 Mbps for 40MHz

WIFI_PHY_RATE_MCS6_LGI = 0x16
MCS6 with long GI, 58.5 Mbps for 20MHz, 121.5 Mbps for 40MHz

WIFI_PHY_RATE_MCS7_LGI = 0x17
MCS7 with long GI, 65 Mbps for 20MHz, 135 Mbps for 40MHz

WIFI_PHY_RATE_MCS0_SGI = 0x18
MCS0 with short GI, 7.2 Mbps for 20MHz, 15 Mbps for 40MHz

WIFI_PHY_RATE_MCS1_SGI = 0x19
MCS1 with short GI, 14.4 Mbps for 20MHz, 30 Mbps for 40MHz

WIFI_PHY_RATE_MCS2_SGI = 0x1A
MCS2 with short GI, 21.7 Mbps for 20MHz, 45 Mbps for 40MHz

WIFI_PHY_RATE_MCS3_SGI = 0x1B
MCS3 with short GI, 28.9 Mbps for 20MHz, 60 Mbps for 40MHz

WIFI_PHY_RATE_MCS4_SGI = 0x1C
MCS4 with short GI, 43.3 Mbps for 20MHz, 90 Mbps for 40MHz

WIFI_PHY_RATE_MCS5_SGI = 0x1D
MCS5 with short GI, 57.8 Mbps for 20MHz, 120 Mbps for 40MHz

WIFI_PHY_RATE_MCS6_SGI = 0x1E
MCS6 with short GI, 65 Mbps for 20MHz, 135 Mbps for 40MHz

WIFI_PHY_RATE_MCS7_SGI = 0x1F
MCS7 with short GI, 72.2 Mbps for 20MHz, 150 Mbps for 40MHz

WIFI_PHY_RATE_LORA_250K = 0x29
250 Kbps

WIFI_PHY_RATE_LORA_500K = 0x2A
500 Kbps

WIFI_PHY_RATE_MAX

enum wifi_event_t
WiFi event declarations

Values:

WIFI_EVENT_WIFI_READY = 0
ESP32 WiFi ready

WIFI_EVENT_SCAN_DONE
ESP32 finish scanning AP

WIFI_EVENT_STA_START
ESP32 station start

WIFI_EVENT_STA_STOP
ESP32 station stop

WIFI_EVENT_STA_CONNECTED
ESP32 station connected to AP

WIFI_EVENT_STA_DISCONNECTED
ESP32 station disconnected from AP

WIFI_EVENT_STA_AUTHMODE_CHANGE
the auth mode of AP connected by ESP32 station changed

WIFI_EVENT_STA_BEACON_TIMEOUT
ESP32 station beacon timeout

WIFI_EVENT_STA_WPS_ER_SUCCESS
ESP32 station wps succeeds in enrollee mode

WIFI_EVENT_STA_WPS_ER_FAILED
ESP32 station wps fails in enrollee mode

WIFI_EVENT_STA_WPS_ER_TIMEOUT
ESP32 station wps timeout in enrollee mode

WIFI_EVENT_STA_WPS_ER_PIN
ESP32 station wps pin code in enrollee mode

WIFI_EVENT_STA_WPS_ER_PBC_OVERLAP
ESP32 station wps overlap in enrollee mode

WIFI_EVENT_AP_START
ESP32 soft-AP start

WIFI_EVENT_AP_STOP
ESP32 soft-AP stop

WIFI_EVENT_AP_STACONNECTED
a station connected to ESP32 soft-AP

WIFI_EVENT_AP_STADISCONNECTED
a station disconnected from ESP32 soft-AP

WIFI_EVENT_AP_PROBEREQRECVED
Receive probe request packet in soft-AP interface

WIFI_EVENT_MAX
Invalid WiFi event ID

enum wifi_event_sta_wps_fail_reason_t
Argument structure for WIFI_EVENT_STA_WPS_ER_FAILED event

Values:

WPS_FAIL_REASON_NORMAL = 0
ESP32 WPS normal fail reason

WPS_FAIL_REASON_RECV_M2D
ESP32 WPS receive M2D frame

WPS_FAIL_REASON_MAX

SmartConfig

The SmartConfig™ is a provisioning technology developed by TI to connect a new Wi-Fi device to a Wi-Fi network. It uses a mobile app to broadcast the network credentials from a smartphone, or a tablet, to an un-provisioned Wi-Fi device.

The advantage of this technology is that the device does not need to directly know SSID or password of an Access Point (AP). This information is provided using the smartphone. This is particularly important to headless device and systems, due to their lack of a user interface.

If you are looking for other options to provision your ESP32-S2 devices, check 配网 *API*.

Application Example Connect ESP32-S2 to target AP using SmartConfig: wifi/smart_config.

API Reference

Header File

- esp_wifi/include/esp_smartconfig.h

Functions

const char ***esp_smartconfig_get_version**(void)

Get the version of SmartConfig.

Return

- SmartConfig version const char.

esp_err_t **esp_smartconfig_start**(**const** *smartconfig_start_config_t* **config*)

Start SmartConfig, config ESP device to connect AP. You need to broadcast information by phone APP. Device sniffer special packets from the air that containing SSID and password of target AP.

Attention 1. This API can be called in station or softAP-station mode.

Attention 2. Can not call esp_smartconfig_start twice before it finish, please call esp_smartconfig_stop first.

Return

- ESP_OK: succeed
- others: fail

Parameters

- `config`: pointer to smartconfig start configure structure

esp_err_t **esp_smartconfig_stop**(void)

Stop SmartConfig, free the buffer taken by esp_smartconfig_start.

Attention Whether connect to AP succeed or not, this API should be called to free memory taken by smartconfig_start.

Return

- ESP_OK: succeed
- others: fail

esp_err_t **esp_esptouch_set_timeout**(uint8_t *time_s*)

Set timeout of SmartConfig process.

Attention Timing starts from SC_STATUS_FIND_CHANNEL status. SmartConfig will restart if timeout.

Return

- ESP_OK: succeed
- others: fail

Parameters

- `time_s`: range 15s~255s, offset:45s.

esp_err_t **esp_smartconfig_set_type**(*smartconfig_type_t type*)

Set protocol type of SmartConfig.

Attention If users need to set the SmartConfig type, please set it before calling esp_smartconfig_start.

Return

- ESP_OK: succeed
- others: fail

Parameters

- `type`: Choose from the smartconfig_type_t.

esp_err_t **esp_smartconfig_fast_mode**(bool *enable*)

Set mode of SmartConfig. default normal mode.

Attention 1. Please call it before API esp_smartconfig_start.

Attention 2. Fast mode have corresponding APP(phone).

Attention 3. Two mode is compatible.

Return

- ESP_OK: succeed
- others: fail

Parameters

- `enable`: false-disable(default); true-enable;

Structures

struct smartconfig_event_got_ssid_pswd_t

Argument structure for SC_EVENT_GOT_SSID_PSWD event

Public Members

uint8_t **ssid**[32]
SSID of the AP. Null terminated string.

uint8_t **password**[64]
Password of the AP. Null terminated string.

bool **bssid_set**
whether set MAC address of target AP or not.

uint8_t **bssid**[6]
MAC address of target AP.

smartconfig_type_t **type**
Type of smartconfig(ESPTouch or AirKiss).

uint8_t **token**
Token from cellphone which is used to send ACK to cellphone.

uint8_t **cellphone_ip**[4]
IP address of cellphone.

struct smartconfig_start_config_t

Configure structure for esp_smartconfig_start

Public Members

bool **enable_log**
Enable smartconfig logs.

Macros

SMARTCONFIG_START_CONFIG_DEFAULT()

Enumerations

enum smartconfig_type_t

Values:

SC_TYPE_ESPTOUCH = 0
protocol: ESPTouch

SC_TYPE_AIRKISS
protocol: AirKiss

SC_TYPE_ESPTOUCH_AIRKISS
protocol: ESPTouch and AirKiss

enum smartconfig_event_t

Smartconfig event declarations

Values:

SC_EVENT_SCAN_DONE
ESP32 station smartconfig has finished to scan for APs

SC_EVENT_FOUND_CHANNEL
ESP32 station smartconfig has found the channel of the target AP

SC_EVENT_GOT_SSID_PSWD
ESP32 station smartconfig got the SSID and password

SC_EVENT_SEND_ACK_DONE
ESP32 station smartconfig has sent ACK to cellphone

ESP-NOW

概述 ESP-NOW 是一种由乐鑫公司定义的无连接 Wi-Fi 通信协议。在 ESP-NOW 中，应用程序数据被封装在各个供应商的动作帧中，然后在无连接的情况下，从一个 Wi-Fi 设备传输到另一个 Wi-Fi 设备。CTR 与 CBC-MAC 协议 (CCMP) 可用来保护动作帧的安全。ESP-NOW 广泛应用于智能照明、远程控制、传感器等领域。

帧格式 ESP-NOW 使用各个供应商的动作帧传输数据，默认比特率为 1 Mbps。各个供应商的动作帧格式为：

MAC 报头	分类代码	组织标识符	随机值	供应商特定内容	FCS
24 字节	1 字节	3 字节	4 字节	7~257 字节	4 字节

- 分类代码：分类代码字段可用于指示各个供应商的类别（比如 127）。
- 组织标识符：组织标识符包含一个唯一标识符 (比如 0x18fe34)，为乐鑫指定的 MAC 地址的前三个字节。
- 随机值：防止重放攻击。
- 供应商特定内容：供应商特定内容包含供应商特定字段，如下所示：

元素 ID	长度	组织标识符	类型	版本	正文
1 字节	1 字节	3 字节	1 字节	1 字节	0 ~ 250 字节

- 元素 ID：元素 ID 字段可用于指示特定于供应商的元素。
- 长度：长度是组织标识符、类型、版本和正文的总长度。
- 组织标识符：组织标识符包含一个唯一标识符 (比如 0x18fe34)，为乐鑫指定的 MAC 地址的前三个字节。
- 类型：类型字段设置为 4，代表 ESP-NOW。
- 版本：版本字段设置为 ESP-NOW 的版本。
- 正文：正文包含 ESP-NOW 数据。

由于 ESP-NOW 是无连接的，因此 MAC 报头与标准帧略有不同。FrameControl 字段的 FromDS 和 ToDS 位均为 0。第一个地址字段用于配置目标地址。第二个地址字段用于配置源地址。第三个地址字段用于配置广播地址 (0xff:0xff:0xff:0xff:0xff:0xff)。

安全

ESP-NOW 采用 CCMP 方法保护供应商特定动作帧的安全，具体可参考 IEEE Std. 802.11-2012。Wi-Fi 设备维护一个初始

- PMK 可使用 AES-128 算法加密 LMK。请调用*esp_now_set_pmk()* 设置 PMK。如果未设置 PMK，将使用默认 PMK。
- LMK 可通过 CCMP 方法对供应商特定的动作帧进行加密，最多拥有 6 个不同的 LMK。如果未设置配对设备的 LMK，则动作帧不进行加密。

目前，不支持加密组播供应商特定的动作帧。

初始化和反初始化 调用*esp_now_init()* 初始化 ESP-NOW，调用*esp_now_deinit()* 反初始化 ESP-NOW。ESP-NOW 数据必须在 Wi-Fi 启动后传输，因此建议在初始化 ESP-NOW 之前启动 Wi-Fi，并

在反初始化 ESP-NOW 之后停止 Wi-Fi。当调用 *esp_now_deinit()* 时，配对设备的所有信息都将被删除。

添加配对设备 在将数据发送到其他设备之前，请先调用 esp_now_add_peer() 将其添加到配对设备列表中。如果启用了加密，则必须设置 LMK。ESP-NOW 数据可以从 Station 或 Softap 接口发送。确保在发送 ESP-NOW 数据之前已启用该接口。

配对设备的最大数量是 20，其中加密设备的数量不超过 16，默认值是 6。

在发送广播数据之前必须添加具有广播 MAC 地址的设备。配对设备的信道范围是从 0 ~ 14。如果信道设置为 0，数据将在当前信道上发送。否则，必须使用本地设备所在的通道。

发送 ESP-NOW 数据 调用 *esp_now_send()* 发送 ESP-NOW 数据，调用 *esp_now_register_send_cb()* 注册发送回调函数。如果 MAC 层成功接收到数据，则该函数将返回 *ESP_NOW_SEND_SUCCESS* 事件。否则，它将返回 *ESP_NOW_SEND_FAIL*。ESP-NOW 数据发送失败可能有几种原因，比如目标设备不存在、设备的信道不相同、动作帧在传输过程中丢失等。应用层并不一定可以总能接收到数据。如果需要，应用层可在接收 ESP-NOW 数据时发回一个应答 (ACK) 数据。如果接收 ACK 数据超时，则将重新传输 ESP-NOW 数据。可以为 ESP-NOW 数据设置序列号，从而删除重复的数据。

如果有大量 ESP-NOW 数据要发送，则调用 *esp_now_send()* 一次性发送不大于 250 字节的数据。请注意，两个 ESP-NOW 数据包的发送间隔太短可能导致回调函数返回混乱。因此，建议在等到上一次回调函数返回 ACK 后再发送下一个 ESP-NOW 数据。发送回调函数从高优先级的 Wi-Fi 任务中运行。因此，不要在回调函数中执行冗长的操作。相反，将必要的数据发布到队列，并交给优先级较低的任务处理。

接收 ESP-NOW 数据 调用 *esp_now_register_recv_cb()* 注册接收回调函数。当接收 ESP-NOW 数据时，需要调用接收回调函数。接收回调函数也在 Wi-Fi 任务任务中运行。因此，不要在回调函数中执行冗长的操作。相反，将必要的数据发布到队列，并交给优先级较低的任务处理。

API 参考

Header File

- esp_wifi/include/esp_now.h

Functions

esp_err_t **esp_now_init** (void)

Initialize ESPNOW function.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_INTERNAL : Internal error

esp_err_t **esp_now_deinit** (void)

De-initialize ESPNOW function.

Return

- ESP_OK : succeed

esp_err_t **esp_now_get_version** (uint32_t **version*)

Get the version of ESPNOW.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_ARG : invalid argument

Parameters

- version: ESPNOW version

esp_err_t **esp_now_register_recv_cb**(*esp_now_recv_cb_t cb*)
Register callback function of receiving ESPNOW data.

Return
- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_INTERNAL : internal error

Parameters
- `cb`: callback function of receiving ESPNOW data

esp_err_t **esp_now_unregister_recv_cb**(void)
Unregister callback function of receiving ESPNOW data.

Return
- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized

esp_err_t **esp_now_register_send_cb**(*esp_now_send_cb_t cb*)
Register callback function of sending ESPNOW data.

Return
- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_INTERNAL : internal error

Parameters
- `cb`: callback function of sending ESPNOW data

esp_err_t **esp_now_unregister_send_cb**(void)
Unregister callback function of sending ESPNOW data.

Return
- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized

esp_err_t **esp_now_send**(**const** uint8_t **peer_addr*, **const** uint8_t **data*, size_t *len*)
Send ESPNOW data.

Attention 1. If peer_addr is not NULL, send data to the peer whose MAC address matches peer_addr
Attention 2. If peer_addr is NULL, send data to all of the peers that are added to the peer list
Attention 3. The maximum length of data must be less than ESP_NOW_MAX_DATA_LEN
Attention 4. The buffer pointed to by data argument does not need to be valid after esp_now_send returns

Return
- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument
- ESP_ERR_ESPNOW_INTERNAL : internal error
- ESP_ERR_ESPNOW_NO_MEM : out of memory, when this happens, you can delay a while before sending the next data
- ESP_ERR_ESPNOW_NOT_FOUND : peer is not found
- ESP_ERR_ESPNOW_IF : current WiFi interface doesn’ t match that of peer

Parameters
- `peer_addr`: peer MAC address
- `data`: data to send
- `len`: length of data

esp_err_t **esp_now_add_peer**(**const** *esp_now_peer_info_t* **peer*)
Add a peer to peer list.

Return
- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument
- ESP_ERR_ESPNOW_FULL : peer list is full
- ESP_ERR_ESPNOW_NO_MEM : out of memory

- ESP_ERR_ESPNOW_EXIST : peer has existed

Parameters

- `peer`: peer information

esp_err_t **esp_now_del_peer**(**const** uint8_t *peer_addr)

Delete a peer from peer list.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument
- ESP_ERR_ESPNOW_NOT_FOUND : peer is not found

Parameters

- `peer_addr`: peer MAC address

esp_err_t **esp_now_mod_peer**(**const** *esp_now_peer_info_t* *peer)

Modify a peer.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument
- ESP_ERR_ESPNOW_FULL : peer list is full

Parameters

- `peer`: peer information

esp_err_t **esp_now_get_peer**(**const** uint8_t *peer_addr, *esp_now_peer_info_t* *peer)

Get a peer whose MAC address matches peer_addr from peer list.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument
- ESP_ERR_ESPNOW_NOT_FOUND : peer is not found

Parameters

- `peer_addr`: peer MAC address
- `peer`: peer information

esp_err_t **esp_now_fetch_peer**(bool *from_head*, *esp_now_peer_info_t* *peer)

Fetch a peer from peer list. Only return the peer which address is unicast, for the multicast/broadcast address, the function will ignore and try to find the next in the peer list.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument
- ESP_ERR_ESPNOW_NOT_FOUND : peer is not found

Parameters

- `from_head`: fetch from head of list or not
- `peer`: peer information

bool **esp_now_is_peer_exist**(**const** uint8_t *peer_addr)

Peer exists or not.

Return

- true : peer exists
- false : peer not exists

Parameters

- `peer_addr`: peer MAC address

esp_err_t **esp_now_get_peer_num**(*esp_now_peer_num_t* *num)

Get the number of peers.

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument

Parameters

- num: number of peers

esp_err_t **esp_now_set_pmk**(**const** uint8_t **pmk*)

Set the primary master key.

Attention 1. primary master key is used to encrypt local master key

Return

- ESP_OK : succeed
- ESP_ERR_ESPNOW_NOT_INIT : ESPNOW is not initialized
- ESP_ERR_ESPNOW_ARG : invalid argument

Parameters

- pmk: primary master key

Structures

struct esp_now_peer_info

ESPNOW peer information parameters.

Public Members

uint8_t **peer_addr**[**ESP_NOW_ETH_ALEN**]

ESPNOW peer MAC address that is also the MAC address of station or softap

uint8_t **lmk**[**ESP_NOW_KEY_LEN**]

ESPNOW peer local master key that is used to encrypt data

uint8_t **channel**

Wi-Fi channel that peer uses to send/receive ESPNOW data. If the value is 0, use the current channel which station or softap is on. Otherwise, it must be set as the channel that station or softap is on.

wifi_interface_t **ifidx**

Wi-Fi interface that peer uses to send/receive ESPNOW data

bool **encrypt**

ESPNOW data that this peer sends/receives is encrypted or not

void ***priv**

ESPNOW peer private data

struct esp_now_peer_num

Number of ESPNOW peers which exist currently.

Public Members

int **total_num**

Total number of ESPNOW peers, maximum value is ESP_NOW_MAX_TOTAL_PEER_NUM

int **encrypt_num**

Number of encrypted ESPNOW peers, maximum value is ESP_NOW_MAX_ENCRYPT_PEER_NUM

Macros

ESP_ERR_ESPNOW_BASE

ESPNOW error number base.

ESP_ERR_ESPNOW_NOT_INIT

ESPNOW is not initialized.

ESP_ERR_ESPNOW_ARG

Invalid argument

ESP_ERR_ESPNOW_NO_MEM
Out of memory

ESP_ERR_ESPNOW_FULL
ESPNOW peer list is full

ESP_ERR_ESPNOW_NOT_FOUND
ESPNOW peer is not found

ESP_ERR_ESPNOW_INTERNAL
Internal error

ESP_ERR_ESPNOW_EXIST
ESPNOW peer has existed

ESP_ERR_ESPNOW_IF
Interface error

ESP_NOW_ETH_ALEN
Length of ESPNOW peer MAC address

ESP_NOW_KEY_LEN
Length of ESPNOW peer local master key

ESP_NOW_MAX_TOTAL_PEER_NUM
Maximum number of ESPNOW total peers

ESP_NOW_MAX_ENCRYPT_PEER_NUM
Maximum number of ESPNOW encrypted peers

ESP_NOW_MAX_DATA_LEN
Maximum length of ESPNOW data which is sent very time

Type Definitions

typedef struct *esp_now_peer_info* **esp_now_peer_info_t**
ESPNOW peer information parameters.

typedef struct *esp_now_peer_num* **esp_now_peer_num_t**
Number of ESPNOW peers which exist currently.

typedef void (***esp_now_recv_cb_t**)(**const** uint8_t *mac_addr, **const** uint8_t *data, int data_len)
Callback function of receiving ESPNOW data.

Parameters
- `mac_addr`: peer MAC address
- `data`: received data
- `data_len`: length of received data

typedef void (***esp_now_send_cb_t**)(**const** uint8_t *mac_addr, *esp_now_send_status_t* status)
Callback function of sending ESPNOW data.

Parameters
- `mac_addr`: peer MAC address
- `status`: status of sending ESPNOW data (succeed or fail)

Enumerations

enum esp_now_send_status_t
Status of sending ESPNOW data .

Values:

ESP_NOW_SEND_SUCCESS = 0
Send ESPNOW data successfully

ESP_NOW_SEND_FAIL
Send ESPNOW data fail

ESP-MESH 编程指南

这是 ESP-MESH 的编程指南，包括 API 参考和编码示例。本指南分为以下部分：

1. *ESP-MESH 编程模型*
2. *编写 ESP-MESH 应用程序*
3. *自组网*
4. *应用实例*
5. *API 参考*

有关 ESP-MESH 协议的文档，请见*ESP-MESH API 指南*。有关 ESP-MESH 开发框架的更多内容，请见 ESP-MESH 开发框架。

ESP-MESH 编程模型

软件栈 ESP-MESH 软件栈基于 Wi-Fi 驱动程序和 FreeRTOS 构建，某些情况下（如根节点）也会使用 LwIP 软件栈。下图展示了 ESP-MESH 软件栈。

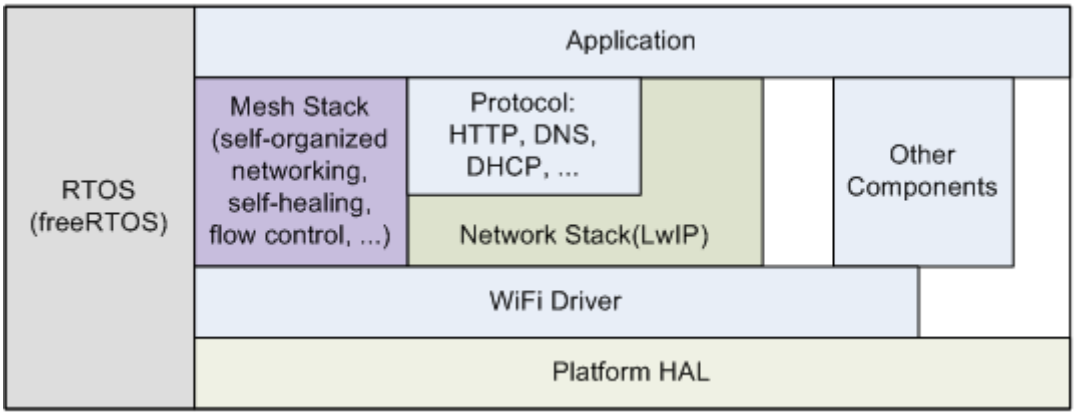

图 1: ESP-MESH 软件栈

系统事件 应用程序可通过 **ESP-MESH 事件**与 ESP-MESH 交互。由于 ESP-MESH 构建在 Wi-Fi 软件栈之上，因此也可以通过 **Wi-Fi 事件任务**与 Wi-Fi 驱动程序进行交互。下图展示了 ESP-MESH 应用程序中各种系统事件的接口。

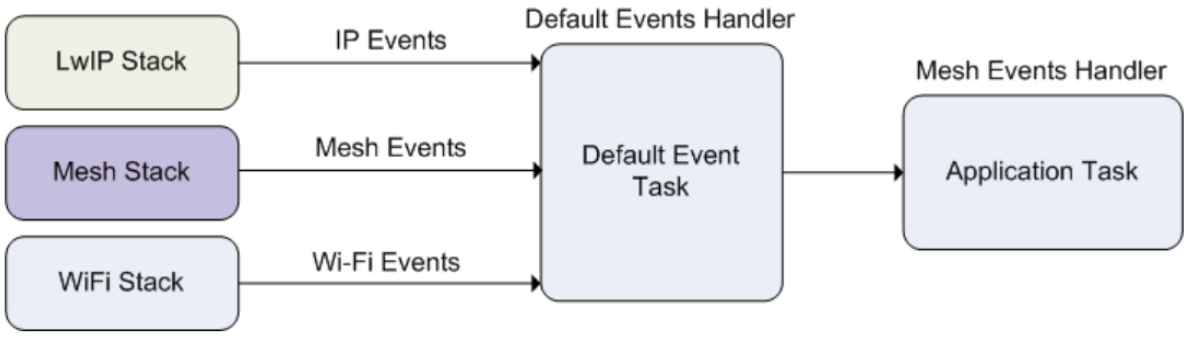

图 2: ESP-MESH 系统事件交付

mesh_event_id_t 定义了所有可能的 ESP-MESH 事件，并且可以指示父节点和子节点的连接或断开等事件。应用程序如需使用 ESP-MESH 事件，则必须通过*esp_event_handler_register()* 将 **Mesh 事件处理程序**注册在默认事件任务中。注册完成后，ESP-MESH 事件将包含与应用程序所有相关事件相关的处理程序。

Mesh 事件的典型应用场景包括：使用*MESH_EVENT_PARENT_CONNECTED* 和*MESH_EVENT_CHILD_CONNECTED* 事件来指示节点何时可以分别开始传输上行和下行的数据。同样，也可以使用 IP_EVENT_STA_GOT_IP 和 IP_EVENT_STA_LOST_IP 事件来指示根节点何时可以向外部 IP 网络传输数据。

警告： 在自组网模式下使用 ESP-MESH 时，用户必须确保不得调用 Wi-Fi API。原因在于：自组网模式将在内部调用 Wi-Fi API 实现连接/断开/扫描等操作。**此时，如果外部应用程序调用 Wi-Fi API（包括来自回调函数和 Wi-Fi 事件处理程序的调用）都可能会干扰 ESP-MESH 的自组网行为**。因此，用户不应该在*esp_mesh_start()* 和*esp_mesh_stop()* 之间调用 Wi-Fi API。

LwIP & ESP-MESH 应用程序无需通过 LwIP 层便可直接访问 ESP-MESH 软件栈，LwIP 层仅在根节点和外部 IP 网络的数据发送与接收时会用到。但是，由于每个节点都有可能成为根节点（由于自动根节点选择机制的存在），每个节点仍必须初始化 LwIP 软件栈。

每个节点都需要通过调用 tcpip_adapter_init() **初始化 LwIP 软件栈**。为了防止非根节点访问 LwIP，应用程序应该在 LwIP 初始化完成后停止以下服务：

- SoftAP 接口上的 DHCP 服务器服务。
- Station 接口上的 DHCP 客户端服务。

下方代码片段展示如何为 ESP-MESH 应用程序进行 LwIP 初始化。

```
/*  tcpip 初始化 */
tcpip_adapter_init();
/*
* 对于 MESH
* 默认情况下，在 SoftAP 接口上停止 DHCP 服务器
* 默认情况下，在 Station 接口上停止 DHCP 客户端
*/
ESP_ERROR_CHECK(tcpip_adapter_dhcps_stop(TCPIP_ADAPTER_IF_AP));
ESP_ERROR_CHECK(tcpip_adapter_dhcpc_stop(TCPIP_ADAPTER_IF_STA));
```

注解： ESP-MESH 的根节点必须与路由器连接。因此，当一个节点成为根节点时，**该节点对应的处理程序必须启动 DHCP 客户端服务并立即获取 IP 地址**。这样做将允许其他节点开始向/从外部 IP 网络发送/接收数据包。但是，如果使用静态 IP 设置，则不需要执行此步骤。

编写 ESP-MESH 应用程序 ESP-MESH 在正常启动前必须先初始化 LwIP 和 Wi-Fi 软件栈。下方代码展示了 ESP-MESH 在开始自身初始化前必须完成的步骤。

```
tcpip_adapter_init();
/*
* 对于 MESH
* 默认情况下，在 SoftAP 接口上停止 DHCP 服务器
* 默认情况下，在 Station 接口上停止 DHCP 客户端
*/
ESP_ERROR_CHECK(tcpip_adapter_dhcps_stop(TCPIP_ADAPTER_IF_AP));
ESP_ERROR_CHECK(tcpip_adapter_dhcpc_stop(TCPIP_ADAPTER_IF_STA));

/* 事件初始化 */
ESP_ERROR_CHECK(esp_event_loop_create_default());

/*Wi-Fi 初始化 */
```

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```
wifi_init_config_t config = WIFI_INIT_CONFIG_DEFAULT();
ESP_ERROR_CHECK(esp_wifi_init(&config));
/* 注册 IP 事件处理程序 */
ESP_ERROR_CHECK(esp_event_handler_register(IP_EVENT, IP_EVENT_STA_GOT_IP, &ip_
↪event_handler, NULL));
ESP_ERROR_CHECK(esp_wifi_set_storage(WIFI_STORAGE_FLASH));
ESP_ERROR_CHECK(esp_wifi_start());
```

在完成 LwIP 和 Wi-Fi 的初始化后，需完成以下三个步骤以启动并运行 ESP-MESH。

1. *初始化 Mesh*
2. *配置 ESP-MESH 网络*
3. *启动 Mesh*

初始化 Mesh 下方代码片段展示如何初始化 ESP-MESH。

```
/*Mesh 初始化 */
ESP_ERROR_CHECK(esp_mesh_init());
/* 注册 mesh 事件处理程序 */
ESP_ERROR_CHECK(esp_event_handler_register(MESH_EVENT, ESP_EVENT_ANY_ID, &mesh_
↪event_handler, NULL));
```

配置 ESP-MESH 网络 ESP-MESH 可通过*esp_mesh_set_config()* 进行配置，并使用*mesh_cfg_t* 结构体传递参数。该结构体包含以下 ESP-MESH 的配置参数：

参数	描述
Channel（信道）	1 到 14 信道
Mesh ID	ESP-MESH 网络的 ID，见*mesh_addr_t*。
Router（路由器）	路由器配置，见*mesh_router_t*。
Mesh AP	Mesh AP 配置，见*mesh_ap_cfg_t*
Crypto Functions（加密函数）	Mesh IE 的加密函数，见 mesh_crypto_funcs_t。

下方代码片段展示如何配置 ESP-MESH。

```
/* 默认启用 MESH IE 加密 */
mesh_cfg_t cfg = MESH_INIT_CONFIG_DEFAULT();
/* Mesh ID */
memcpy((uint8_t *) &cfg.mesh_id, MESH_ID, 6);
/* 信道（需与路由器信道匹配）*/
cfg.channel = CONFIG_MESH_CHANNEL;
/* 路由器 */
cfg.router.ssid_len = strlen(CONFIG_MESH_ROUTER_SSID);
memcpy((uint8_t *) &cfg.router.ssid, CONFIG_MESH_ROUTER_SSID, cfg.router.ssid_len);
memcpy((uint8_t *) &cfg.router.password, CONFIG_MESH_ROUTER_PASSWD,
    strlen(CONFIG_MESH_ROUTER_PASSWD));
/* Mesh softAP */
cfg.mesh_ap.max_connection = CONFIG_MESH_AP_CONNECTIONS;
memcpy((uint8_t *) &cfg.mesh_ap.password, CONFIG_MESH_AP_PASSWD,
    strlen(CONFIG_MESH_AP_PASSWD));
ESP_ERROR_CHECK(esp_mesh_set_config(&cfg));
```

启动 Mesh 下方代码片段展示如何启动 ESP-MESH。

```
/* 启动 Mesh */
ESP_ERROR_CHECK(esp_mesh_start());
```

启动 ESP-MESH 后，应用程序应检查 ESP-MESH 事件，以确定它是何时连接到网络的。连接后，应用程序可使用*esp_mesh_send()* 和*esp_mesh_recv()* 在 ESP-MESH 网络中发送、接收数据包。

自组网 自组网是 ESP-MESH 的功能之一，允许节点自动扫描/选择/连接/重新连接到其他节点和路由器。此功能允许 ESP-MESH 网络具有很高的自主性，可适应变化的动态网络拓扑结构和环境。启用自组网功能后，ESP-MESH 网络中的节点能够自主完成以下操作：

- 选择或选举根节点（见建立网络 中的 **自动根节点选择**）
- 选择首选的父节点（见建立网络 中的 **父节点选择**）
- 网络断开时自动重新连接（见管理网络 中的 **中间父节点失败**）

启用自组网功能后，ESP-MESH 软件栈将内部调用 Wi-Fi API。因此，**在启用自组网功能时，应用层不得调用 Wi-Fi API，否则会干扰 ESP-MESH 的工作**。

开关自组网 应用程序可以在运行时通过调用*esp_mesh_set_self_organized()* 函数，启用或禁用自组网功能。该函数具有以下两个参数：

- `bool enable` 指定启用或禁用自组网功能。
- `bool select_parent` 指定在启用自组网功能时是否应选择新的父节点。根据节点类型和节点当前状态，选择新的父节点具有不同的作用。在禁用自组网功能时，此参数不使用。

禁用自组网 下方代码片段展示了如何禁用自组网功能。

```
//禁用自组网
esp_mesh_set_self_organized(false, false);
```

ESP-MESH 将在禁用自组网时尝试维护节点的当前 Wi-Fi 状态。

- 如果节点先前已连接到其他节点，则将保持连接。
- 如果节点先前已断开连接并且正在扫描父节点或路由器，则将停止扫描。
- 如果节点以前尝试重新连接到父节点或路由器，则将停止重新连接。

启用自组网 ESP-MESH 将尝试在启用自组网时保持节点的当前 Wi-Fi 状态。但是，根据节点类型以及是否选择了新的父节点，节点的 Wi-Fi 状态可能会发生变化。下表显示了启用自组网的效果。

是否选择父节点	是否为根结点	作用
N	N	已连接到父节点的节点将保持连接。
		之前扫描父节点的节点将停止扫描。调用*esp_mesh_connect()* 重新启动。
	Y	已连接到路由器的根节点将保持连接。
		从路由器断开的根结点需调用*esp_mesh_connect()* 进行重连。
Y	N	没有父节点的节点将自动选择首选父节点并连接。
		已连接到父节点的节点将断开连接，重新选择首选父节点并进行重连。
	Y	根结点在连接至父节点前必须放弃“根结点”的角色。因此，根节点将断开与路由器和所有子节点的连接，选择首选父节点并进行连接。

下方代码片段展示了如何启用自组网功能。

```
//启用自组网，并选择一个新的父节点
esp_mesh_set_self_organized(true, true);

...

//启用自组网并手动重新连接
esp_mesh_set_self_organized(true, false);
esp_mesh_connect();
```

调用 Wi-Fi API 在有些情况下，应用程序可能希望在使用 ESP-MESH 期间调用 Wi-Fi API。例如，应用程序可能需要手动扫描邻近的接入点 (AP)。**但在应用程序调用任何 Wi-Fi API 之前，必须先禁用自组网**。否则，ESP-MESH 软件栈可能会同时调用 Wi-Fi API，进而影响应用程序的正常调用。

应用程序不应在*esp_mesh_set_self_organized()* 之间调用 Wi-Fi API。下方代码片段展示了应用程序如何在 ESP-MESH 运行期间安全地调用*esp_wifi_scan_start()*。

```
//禁用自组网
esp_mesh_set_self_organized(0, 0);

//停止任何正在进行的扫描
esp_wifi_scan_stop();
//手动启动扫描运行完成时自动停止
esp_wifi_scan_start();

//进程扫描结果

...

//如果仍为连接状态，则重新启用自组网
esp_mesh_set_self_organized(1, 0);

...

//如果不为根节点且未连接，则重新启用自组网
esp_mesh_set_self_organized(1, 1);

...

//如果为根节点且未连接，则重新启用
esp_mesh_set_self_organized(1, 0);   //不选择新的父节点
esp_mesh_connect();                  //手动重新连接到路由器
```

应用实例 ESP-IDF 包含以下 ESP-MESH 示例项目：

内部通信示例 展示了如何搭建 ESP-MESH 网络，并让根节点向网络中的每个节点发送数据包。

手动连网示例 展示了如何在禁用自组网功能的情况下使用 ESP-MESH。此示例展示了如何对节点进行编程，以手动扫描潜在父节点的列表，并根据自定义标准选择父节点。

API 参考

Header File

- esp_wifi/include/esp_mesh.h

Functions

esp_err_t **esp_mesh_init**(void)

Mesh initialization.

- Check whether Wi-Fi is started.
- Initialize mesh global variables with default values.

Attention This API shall be called after Wi-Fi is started.

Return

- ESP_OK
- ESP_FAIL

esp_err_t **esp_mesh_deinit**(void)

Mesh de-initialization.

- Release resources and stop the mesh

Return

- ESP_OK
- ESP_FAIL

esp_err_t **esp_mesh_start**(void)

Start mesh.

- Initialize mesh IE.
- Start mesh network management service.
- Create TX and RX queues according to the configuration.
- Register mesh packets receive callback.

Attention This API shall be called after mesh initialization and configuration.

Return

- ESP_OK
- ESP_FAIL
- ESP_ERR_MESH_NOT_INIT
- ESP_ERR_MESH_NOT_CONFIG
- ESP_ERR_MESH_NO_MEMORY

esp_err_t **esp_mesh_stop**(void)

Stop mesh.

- Deinitialize mesh IE.
- Disconnect with current parent.
- Disassociate all currently associated children.
- Stop mesh network management service.
- Unregister mesh packets receive callback.
- Delete TX and RX queues.
- Release resources.
- Restore Wi-Fi softAP to default settings if Wi-Fi dual mode is enabled.
- Set Wi-Fi Power Save type to WIFI_PS_NONE.

Return

- ESP_OK
- ESP_FAIL

esp_err_t **esp_mesh_send**(**const** *mesh_addr_t* *to, **const** *mesh_data_t* *data, int *flag*, **const** *mesh_opt_t opt*[], int *opt_count*)

Send a packet over the mesh network.

- Send a packet to any device in the mesh network.
- Send a packet to external IP network.

Attention This API is not reentrant.

Return

- ESP_OK
- ESP_FAIL
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_START
- ESP_ERR_MESH_DISCONNECTED
- ESP_ERR_MESH_OPT_UNKNOWN
- ESP_ERR_MESH_EXCEED_MTU
- ESP_ERR_MESH_NO_MEMORY
- ESP_ERR_MESH_TIMEOUT
- ESP_ERR_MESH_QUEUE_FULL
- ESP_ERR_MESH_NO_ROUTE_FOUND
- ESP_ERR_MESH_DISCARD

Parameters

- `[in] to`: the address of the final destination of the packet
  - If the packet is to the root, set this parameter to NULL.
  - If the packet is to an external IP network, set this parameter to the IPv4:PORT combination. This packet will be delivered to the root firstly, then the root will forward this packet to the final

IP server address.

- `[in] data`: pointer to a sending mesh packet
  - Field size should not exceed MESH_MPS. Note that the size of one mesh packet should not exceed MESH_MTU.
  - Field proto should be set to data protocol in use (default is MESH_PROTO_BIN for binary).
  - Field tos should be set to transmission tos (type of service) in use (default is MESH_TOS_P2P for point-to-point reliable).
- `[in] flag`: bitmap for data sent
  - Speed up the route search
    * If the packet is to the root and "to" parameter is NULL, set this parameter to 0.
    * If the packet is to an internal device, MESH_DATA_P2P should be set.
    * If the packet is to the root ("to" parameter isn' t NULL) or to external IP network, MESH_DATA_TODS should be set.
    * If the packet is from the root to an internal device, MESH_DATA_FROMDS should be set.
  - Specify whether this API is block or non-block, block by default
    * If needs non-blocking, MESH_DATA_NONBLOCK should be set. Otherwise, may use esp_mesh_send_block_time() to specify a blocking time.
  - In the situation of the root change, MESH_DATA_DROP identifies this packet can be dropped by the new root for upstream data to external IP network, we try our best to avoid data loss caused by the root change, but there is a risk that the new root is running out of memory because most of memory is occupied by the pending data which isn' t read out in time by esp_mesh_recv_toDS(). Generally, we suggest esp_mesh_recv_toDS() is called after a connection with IP network is created. Thus data outgoing to external IP network via socket is just from reading esp_mesh_recv_toDS() which avoids unnecessary memory copy.
- `[in] opt`: options
  - In case of sending a packet to a certain group, MESH_OPT_SEND_GROUP is a good choice. In this option, the value field should be set to the target receiver addresses in this group.
  - Root sends a packet to an internal device, this packet is from external IP network in case the receiver device responds this packet, MESH_OPT_RECV_DS_ADDR is required to attach the target DS address.
- `[in] opt_count`: option count
  - Currently, this API only takes one option, so opt_count is only supported to be 1.

esp_err_t **esp_mesh_send_block_time** (uint32_t *time_ms*)

Set blocking time of esp_mesh_send()

Attention This API shall be called before mesh is started.

Return

- ESP_OK

Parameters

- `[in] time_ms`: blocking time of esp_mesh_send(), unit:ms

esp_err_t **esp_mesh_recv** (*mesh_addr_t* **from*, *mesh_data_t* **data*, int *timeout_ms*, int **flag*, *mesh_opt_t* *opt*[], int *opt_count*)

Receive a packet targeted to self over the mesh network.

flag could be MESH_DATA_FROMDS or MESH_DATA_TODS.

Attention Mesh RX queue should be checked regularly to avoid running out of memory.

- Use esp_mesh_get_rx_pending() to check the number of packets available in the queue waiting to be received by applications.

Return

- ESP_OK
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_START
- ESP_ERR_MESH_TIMEOUT
- ESP_ERR_MESH_DISCARD

Parameters

- `[out] from`: the address of the original source of the packet
- `[out] data`: pointer to the received mesh packet
  - Field proto is the data protocol in use. Should follow it to parse the received data.

  - Field tos is the transmission tos (type of service) in use.
- `[in] timeout_ms`: wait time if a packet isn't immediately available (0:no wait, portMAX_DELAY:wait forever)
- `[out] flag`: bitmap for data received
  - MESH_DATA_FROMDS represents data from external IP network
  - MESH_DATA_TODS represents data directed upward within the mesh network

Parameters

- `[out] opt`: options desired to receive
  - MESH_OPT_RECV_DS_ADDR attaches the DS address
- `[in] opt_count`: option count desired to receive
  - Currently, this API only takes one option, so opt_count is only supported to be 1.

esp_err_t **esp_mesh_recv_toDS**(*mesh_addr_t* *from, *mesh_addr_t* *to, *mesh_data_t* *data, int *timeout_ms*, int *flag, *mesh_opt_t* opt[], int *opt_count*)

Receive a packet targeted to external IP network.

- Root uses this API to receive packets destined to external IP network
- Root forwards the received packets to the final destination via socket.
- If no socket connection is ready to send out the received packets and this esp_mesh_recv_toDS() hasn't been called by applications, packets from the whole mesh network will be pending in toDS queue.

Use esp_mesh_get_rx_pending() to check the number of packets available in the queue waiting to be received by applications in case of running out of memory in the root.

Using esp_mesh_set_xon_qsize() users may configure the RX queue size, default:32. If this size is too large, and esp_mesh_recv_toDS() isn't called in time, there is a risk that a great deal of memory is occupied by the pending packets. If this size is too small, it will impact the efficiency on upstream. How to decide this value depends on the specific application scenarios.

flag could be MESH_DATA_TODS.

Attention This API is only called by the root.

Return

- ESP_OK
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_START
- ESP_ERR_MESH_TIMEOUT
- ESP_ERR_MESH_DISCARD
- ESP_ERR_MESH_RECV_RELEASE

Parameters

- `[out] from`: the address of the original source of the packet
- `[out] to`: the address contains remote IP address and port (IPv4:PORT)
- `[out] data`: pointer to the received packet
  - Contain the protocol and applications should follow it to parse the data.
- `[in] timeout_ms`: wait time if a packet isn't immediately available (0:no wait, portMAX_DELAY:wait forever)
- `[out] flag`: bitmap for data received
  - MESH_DATA_TODS represents the received data target to external IP network. Root shall forward this data to external IP network via the association with router.

Parameters

- `[out] opt`: options desired to receive
- `[in] opt_count`: option count desired to receive

esp_err_t **esp_mesh_set_config**(**const** *mesh_cfg_t* *config)

Set mesh stack configuration.

- Use MESH_INIT_CONFIG_DEFAULT() to initialize the default values, mesh IE is encrypted by default.
- Mesh network is established on a fixed channel (1-14).
- Mesh event callback is mandatory.
- Mesh ID is an identifier of an MBSS. Nodes with the same mesh ID can communicate with each other.
- Regarding to the router configuration, if the router is hidden, BSSID field is mandatory.

If BSSID field isn't set and there exists more than one router with same SSID, there is a risk that more roots than one connected with different BSSID will appear. It means more than one mesh network is established with the same mesh ID.

Root conflict function could eliminate redundant roots connected with the same BSSID, but couldn't handle roots connected with different BSSID. Because users might have such requirements of setting up routers with same SSID for the future replacement. But in that case, if the above situations happen, please make sure applications implement forward functions on the root to guarantee devices in different mesh networks can communicate with each other. max_connection of mesh softAP is limited by the max number of Wi-Fi softAP supported (max:10).

Attention This API shall be called before mesh is started after mesh is initialized.
Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_ALLOWED

Parameters
- `[in] config`: pointer to mesh stack configuration

esp_err_t **esp_mesh_get_config**(*mesh_cfg_t* **config*)
Get mesh stack configuration.

Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT

Parameters
- `[out] config`: pointer to mesh stack configuration

esp_err_t **esp_mesh_set_router**(**const** *mesh_router_t* **router*)
Get router configuration.

Attention This API is used to dynamically modify the router configuration after mesh is configured.
Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT

Parameters
- `[in] router`: pointer to router configuration

esp_err_t **esp_mesh_get_router**(*mesh_router_t* **router*)
Get router configuration.

Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT

Parameters
- `[out] router`: pointer to router configuration

esp_err_t **esp_mesh_set_id**(**const** *mesh_addr_t* **id*)
Set mesh network ID.

Attention This API is used to dynamically modify the mesh network ID.
Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT: invalid argument

Parameters
- `[in] id`: pointer to mesh network ID

esp_err_t **esp_mesh_get_id**(*mesh_addr_t* **id*)
Get mesh network ID.

Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT

Parameters

- `[out] id`: pointer to mesh network ID

esp_err_t **esp_mesh_set_type**(*mesh_type_t type*)

Designate device type over the mesh network.

- MESH_IDLE: designates a device as a self-organized node for a mesh network
- MESH_ROOT: designates the root node for a mesh network
- MESH_LEAF: designates a device as a standalone Wi-Fi station that connects to a parent
- MESH_STA: designates a device as a standalone Wi-Fi station that connects to a router

Return

- ESP_OK
- ESP_ERR_MESH_NOT_ALLOWED

Parameters

- `[in] type`: device type

mesh_type_t **esp_mesh_get_type**(void)

Get device type over mesh network.

Attention This API shall be called after having received the event MESH_EVENT_PARENT_CONNECTED.

Return mesh type

esp_err_t **esp_mesh_set_max_layer**(int *max_layer*)

Set network max layer value.

- for tree topology, the max is 25.
- for chain topology, the max is 1000.
- Network max layer limits the max hop count.

Attention This API shall be called before mesh is started.

Return

- ESP_OK
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_ALLOWED

Parameters

- `[in] max_layer`: max layer value

int **esp_mesh_get_max_layer**(void)

Get max layer value.

Return max layer value

esp_err_t **esp_mesh_set_ap_password**(**const** uint8_t **pwd*, int *len*)

Set mesh softAP password.

Attention This API shall be called before mesh is started.

Return

- ESP_OK
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_ALLOWED

Parameters

- `[in] pwd`: pointer to the password
- `[in] len`: password length

esp_err_t **esp_mesh_set_ap_authmode**(*wifi_auth_mode_t authmode*)

Set mesh softAP authentication mode.

Attention This API shall be called before mesh is started.

Return

- ESP_OK
- ESP_ERR_MESH_ARGUMENT
- ESP_ERR_MESH_NOT_ALLOWED

Parameters

- `[in] authmode`: authentication mode

wifi_auth_mode_t **esp_mesh_get_ap_authmode**(void)
Get mesh softAP authentication mode.

Return authentication mode

esp_err_t **esp_mesh_set_ap_connections**(int *connections*)
Set mesh max connection value.

- Set mesh softAP max connection = mesh max connection + non-mesh max connection

Attention This API shall be called before mesh is started.
Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT

Parameters
- `[in] connections`: the number of max connections

int **esp_mesh_get_ap_connections**(void)
Get mesh max connection configuration.

Return the number of mesh max connections

int **esp_mesh_get_non_mesh_connections**(void)
Get non-mesh max connection configuration.

Return the number of non-mesh max connections

int **esp_mesh_get_layer**(void)
Get current layer value over the mesh network.

Attention This API shall be called after having received the event MESH_EVENT_PARENT_CONNECTED.
Return layer value

esp_err_t **esp_mesh_get_parent_bssid**(*mesh_addr_t* *bssid)
Get the parent BSSID.

Attention This API shall be called after having received the event MESH_EVENT_PARENT_CONNECTED.
Return
- ESP_OK
- ESP_FAIL

Parameters
- `[out] bssid`: pointer to parent BSSID

bool **esp_mesh_is_root**(void)
Return whether the device is the root node of the network.

Return true/false

esp_err_t **esp_mesh_set_self_organized**(bool *enable*, bool *select_parent*)
Enable/disable self-organized networking.

- Self-organized networking has three main functions: select the root node; find a preferred parent; initiate reconnection if a disconnection is detected.
- Self-organized networking is enabled by default.
- If self-organized is disabled, users should set a parent for the device via esp_mesh_set_parent().

Attention This API is used to dynamically modify whether to enable the self organizing.
Return
- ESP_OK
- ESP_FAIL

Parameters
- `[in] enable`: enable or disable self-organized networking
- `[in] select_parent`: Only valid when self-organized networking is enabled.
  - if select_parent is set to true, the root will give up its mesh root status and search for a new parent like other non-root devices.

bool **esp_mesh_get_self_organized**(void)
Return whether enable self-organized networking or not.

Return true/false

esp_err_t **esp_mesh_waive_root**(**const** *mesh_vote_t* *vote, int *reason*)
Cause the root device to give up (waive) its mesh root status.

- A device is elected root primarily based on RSSI from the external router.
- If external router conditions change, users can call this API to perform a root switch.
- In this API, users could specify a desired root address to replace itself or specify an attempts value to ask current root to initiate a new round of voting. During the voting, a better root candidate would be expected to find to replace the current one.
- If no desired root candidate, the vote will try a specified number of attempts (at least 15). If no better root candidate is found, keep the current one. If a better candidate is found, the new better one will send a root switch request to the current root, current root will respond with a root switch acknowledgment.
- After that, the new candidate will connect to the router to be a new root, the previous root will disconnect with the router and choose another parent instead.

Root switch is completed with minimal disruption to the whole mesh network.

Attention This API is only called by the root.
Return
- ESP_OK
- ESP_ERR_MESH_QUEUE_FULL
- ESP_ERR_MESH_DISCARD
- ESP_FAIL

Parameters
- `[in] vote`: vote configuration
  - If this parameter is set NULL, the vote will perform the default 15 times.
  - Field percentage threshold is 0.9 by default.
  - Field is_rc_specified shall be false.
  - Field attempts shall be at least 15 times.
- `[in] reason`: only accept MESH_VOTE_REASON_ROOT_INITIATED for now

esp_err_t **esp_mesh_set_vote_percentage**(float *percentage*)
Set vote percentage threshold for approval of being a root (default:0.9)

- During the networking, only obtaining vote percentage reaches this threshold, the device could be a root.

Attention This API shall be called before mesh is started.
Return
- ESP_OK
- ESP_FAIL

Parameters
- `[in] percentage`: vote percentage threshold

float **esp_mesh_get_vote_percentage**(void)
Get vote percentage threshold for approval of being a root.

Return percentage threshold

esp_err_t **esp_mesh_set_ap_assoc_expire**(int *seconds*)
Set mesh softAP associate expired time (default:10 seconds)

- If mesh softAP hasn't received any data from an associated child within this time, mesh softAP will take this child inactive and disassociate it.
- If mesh softAP is encrypted, this value should be set a greater value, such as 30 seconds.

Return
- ESP_OK
- ESP_FAIL

Parameters
- `[in] seconds`: the expired time

int **esp_mesh_get_ap_assoc_expire**(void)
Get mesh softAP associate expired time.

Return seconds

int **esp_mesh_get_total_node_num**(void)
Get total number of devices in current network (including the root)

Attention The returned value might be incorrect when the network is changing.
Return total number of devices (including the root)

int **esp_mesh_get_routing_table_size**(void)
Get the number of devices in this device's sub-network (including self)

Return the number of devices over this device's sub-network (including self)

esp_err_t **esp_mesh_get_routing_table**(*mesh_addr_t* *mac, int len, int *size)
Get routing table of this device's sub-network (including itself)

Return
- ESP_OK
- ESP_ERR_MESH_ARGUMENT

Parameters
- `[out] mac`: pointer to routing table
- `[in] len`: routing table size(in bytes)
- `[out] size`: pointer to the number of devices in routing table (including itself)

esp_err_t **esp_mesh_post_toDS_state**(bool *reachable*)
Post the toDS state to the mesh stack.

Attention This API is only for the root.
Return
- ESP_OK
- ESP_FAIL

Parameters
- `[in] reachable`: this state represents whether the root is able to access external IP network

esp_err_t **esp_mesh_get_tx_pending**(*mesh_tx_pending_t* *pending)
Return the number of packets pending in the queue waiting to be sent by the mesh stack.

Return
- ESP_OK
- ESP_FAIL

Parameters
- `[out] pending`: pointer to the TX pending

esp_err_t **esp_mesh_get_rx_pending**(*mesh_rx_pending_t* *pending)
Return the number of packets available in the queue waiting to be received by applications.

Return
- ESP_OK
- ESP_FAIL

Parameters
- `[out] pending`: pointer to the RX pending

int **esp_mesh_available_txupQ_num**(**const** *mesh_addr_t* *addr, uint32_t *xseqno_in)
Return the number of packets could be accepted from the specified address.

Return the number of upQ for a certain address
Parameters
- `[in] addr`: self address or an associate children address
- `[out] xseqno_in`: sequence number of the last received packet from the specified address

esp_err_t **esp_mesh_set_xon_qsize**(int *qsize*)
Set the number of RX queue for the node, the average number of window allocated to one of its child node is:

wnd = xon_qsize / (2 * max_connection + 1). However, the window of each child node is not strictly equal to the average value, it is affected by the traffic also.

Attention This API shall be called before mesh is started.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] qsize`: default:32 (min:16)

int **esp_mesh_get_xon_qsize**(void)

Get queue size.

Return the number of queue

esp_err_t **esp_mesh_allow_root_conflicts**(bool *allowed*)

Set whether allow more than one root existing in one network.

Return

- ESP_OK
- ESP_WIFI_ERR_NOT_INIT
- ESP_WIFI_ERR_NOT_START

Parameters

- `[in] allowed`: allow or not

bool **esp_mesh_is_root_conflicts_allowed**(void)

Check whether allow more than one root to exist in one network.

Return true/false

esp_err_t **esp_mesh_set_group_id**(**const** *mesh_addr_t* *addr, int *num*)

Set group ID addresses.

Return

- ESP_OK
- ESP_MESH_ERR_ARGUMENT

Parameters

- `[in] addr`: pointer to new group ID addresses
- `[in] num`: the number of group ID addresses

esp_err_t **esp_mesh_delete_group_id**(**const** *mesh_addr_t* *addr, int *num*)

Delete group ID addresses.

Return

- ESP_OK
- ESP_MESH_ERR_ARGUMENT

Parameters

- `[in] addr`: pointer to deleted group ID address
- `[in] num`: the number of group ID addresses

int **esp_mesh_get_group_num**(void)

Get the number of group ID addresses.

Return the number of group ID addresses

esp_err_t **esp_mesh_get_group_list**(*mesh_addr_t* *addr, int *num*)

Get group ID addresses.

Return

- ESP_OK
- ESP_MESH_ERR_ARGUMENT

Parameters

- `[out] addr`: pointer to group ID addresses
- `[in] num`: the number of group ID addresses

bool **esp_mesh_is_my_group**(**const** *mesh_addr_t* *addr)
Check whether the specified group address is my group.

Return true/false

esp_err_t **esp_mesh_set_capacity_num**(int *num*)
Set mesh network capacity (max:1000, default:300)

Attention This API shall be called before mesh is started.
Return
- ESP_OK
- ESP_ERR_MESH_NOT_ALLOWED
- ESP_MESH_ERR_ARGUMENT

Parameters
- `[in] num`: mesh network capacity

int **esp_mesh_get_capacity_num**(void)
Get mesh network capacity.

Return mesh network capacity

esp_err_t **esp_mesh_set_ie_crypto_funcs**(**const** mesh_crypto_funcs_t *crypto_funcs)
Set mesh IE crypto functions.

Attention This API can be called at any time after mesh is initialized.
Return
- ESP_OK

Parameters
- `[in] crypto_funcs`: crypto functions for mesh IE
  - If crypto_funcs is set to NULL, mesh IE is no longer encrypted.

esp_err_t **esp_mesh_set_ie_crypto_key**(**const** char *key, int *len*)
Set mesh IE crypto key.

Attention This API can be called at any time after mesh is initialized.
Return
- ESP_OK
- ESP_MESH_ERR_ARGUMENT

Parameters
- `[in] key`: ASCII crypto key
- `[in] len`: length in bytes, range:8~64

esp_err_t **esp_mesh_get_ie_crypto_key**(char *key, int *len*)
Get mesh IE crypto key.

Return
- ESP_OK
- ESP_MESH_ERR_ARGUMENT

Parameters
- `[out] key`: ASCII crypto key
- `[in] len`: length in bytes, range:8~64

esp_err_t **esp_mesh_set_root_healing_delay**(int *delay_ms*)
Set delay time before starting root healing.

Return
- ESP_OK

Parameters
- `[in] delay_ms`: delay time in milliseconds

int **esp_mesh_get_root_healing_delay**(void)
Get delay time before network starts root healing.

Return delay time in milliseconds

esp_err_t **esp_mesh_fix_root**(bool *enable*)
Enable network Fixed Root Setting.

- Enabling fixed root disables automatic election of the root node via voting.
- All devices in the network shall use the same Fixed Root Setting (enabled or disabled).
- If Fixed Root is enabled, users should make sure a root node is designated for the network.

Return
- ESP_OK

Parameters
- `[in] enable`: enable or not

bool **esp_mesh_is_root_fixed**(void)
Check whether network Fixed Root Setting is enabled.

- Enable/disable network Fixed Root Setting by API esp_mesh_fix_root().
- Network Fixed Root Setting also changes with the "flag" value in parent networking IE.

Return true/false

esp_err_t **esp_mesh_set_parent**(**const** *wifi_config_t* **parent*, **const** *mesh_addr_t* **parent_mesh_id*, *mesh_type_t* *my_type*, int *my_layer*)
Set a specified parent for the device.

Attention This API can be called at any time after mesh is configured.

Return
- ESP_OK
- ESP_ERR_ARGUMENT
- ESP_ERR_MESH_NOT_CONFIG

Parameters
- `[in] parent`: parent configuration, the SSID and the channel of the parent are mandatory.
  - If the BSSID is set, make sure that the SSID and BSSID represent the same parent, otherwise the device will never find this specified parent.
- `[in] parent_mesh_id`: parent mesh ID,
  - If this value is not set, the original mesh ID is used.
- `[in] my_type`: mesh type
  - MESH_STA is not supported.
  - If the parent set for the device is the same as the router in the network configuration, then my_type shall set MESH_ROOT and my_layer shall set MESH_ROOT_LAYER.
- `[in] my_layer`: mesh layer
  - my_layer of the device may change after joining the network.
  - If my_type is set MESH_NODE, my_layer shall be greater than MESH_ROOT_LAYER.
  - If my_type is set MESH_LEAF, the device becomes a standalone Wi-Fi station and no longer has the ability to extend the network.

esp_err_t **esp_mesh_scan_get_ap_ie_len**(int **len*)
Get mesh networking IE length of one AP.

Return
- ESP_OK
- ESP_ERR_WIFI_NOT_INIT
- ESP_ERR_WIFI_ARG
- ESP_ERR_WIFI_FAIL

Parameters
- `[out] len`: mesh networking IE length

esp_err_t **esp_mesh_scan_get_ap_record**(*wifi_ap_record_t* **ap_record*, void **buffer*)
Get AP record.

Attention Different from esp_wifi_scan_get_ap_records(), this API only gets one of APs scanned each time. See "manual_networking" example.

Return
- ESP_OK
- ESP_ERR_WIFI_NOT_INIT

- ESP_ERR_WIFI_ARG
- ESP_ERR_WIFI_FAIL

Parameters

- `[out] ap_record`: pointer to one AP record
- `[out] buffer`: pointer to the mesh networking IE of this AP

esp_err_t **esp_mesh_flush_upstream_packets**(void)

Flush upstream packets pending in to_parent queue and to_parent_p2p queue.

Return

- ESP_OK

esp_err_t **esp_mesh_get_subnet_nodes_num**(**const** *mesh_addr_t* *child_mac, int *nodes_num)

Get the number of nodes in the subnet of a specific child.

Return

- ESP_OK
- ESP_ERR_MESH_NOT_START
- ESP_ERR_MESH_ARGUMENT

Parameters

- `[in] child_mac`: an associated child address of this device
- `[out] nodes_num`: pointer to the number of nodes in the subnet of a specific child

esp_err_t **esp_mesh_get_subnet_nodes_list**(**const** *mesh_addr_t* *child_mac, *mesh_addr_t* *nodes, int *nodes_num*)

Get nodes in the subnet of a specific child.

Return

- ESP_OK
- ESP_ERR_MESH_NOT_START
- ESP_ERR_MESH_ARGUMENT

Parameters

- `[in] child_mac`: an associated child address of this device
- `[out] nodes`: pointer to nodes in the subnet of a specific child
- `[in] nodes_num`: the number of nodes in the subnet of a specific child

esp_err_t **esp_mesh_disconnect**(void)

Disconnect from current parent.

Return

- ESP_OK

esp_err_t **esp_mesh_connect**(void)

Connect to current parent.

Return

- ESP_OK

esp_err_t **esp_mesh_flush_scan_result**(void)

Flush scan result.

Return

- ESP_OK

esp_err_t **esp_mesh_switch_channel**(**const** uint8_t *new_bssid, int csa_newchan, int csa_count)

Cause the root device to add Channel Switch Announcement Element (CSA IE) to beacon.

- Set the new channel
- Set how many beacons with CSA IE will be sent before changing a new channel
- Enable the channel switch function

Attention This API is only called by the root.

Return

- ESP_OK

Parameters

- `[in] new_bssid`: the new router BSSID if the router changes

- `[in] csa_newchan`: the new channel number to which the whole network is moving
- `[in] csa_count`: channel switch period(beacon count), unit is based on beacon interval of its softAP, the default value is 15.

esp_err_t **esp_mesh_get_router_bssid**(uint8_t **router_bssid*)
Get the router BSSID.

Return
- ESP_OK
- ESP_ERR_WIFI_NOT_INIT
- ESP_ERR_WIFI_ARG

Parameters
- `[out] router_bssid`: pointer to the router BSSID

int64_t **esp_mesh_get_tsf_time**(void)
Get the TSF time.

Return the TSF time

esp_err_t **esp_mesh_set_topology**(*esp_mesh_topology_t topo*)
Set mesh topology. The default value is MESH_TOPO_TREE.

- MESH_TOPO_CHAIN supports up to 1000 layers

Attention This API shall be called before mesh is started.
Return
- ESP_OK
- ESP_MESH_ERR_ARGUMENT
- ESP_ERR_MESH_NOT_ALLOWED

Parameters
- `[in] topo`: MESH_TOPO_TREE or MESH_TOPO_CHAIN

esp_mesh_topology_t **esp_mesh_get_topology**(void)
Get mesh topology.

Return MESH_TOPO_TREE or MESH_TOPO_CHAIN

esp_err_t **esp_mesh_enable_ps**(void)
Enable mesh Power Save function.

Attention This API shall be called before mesh is started.
Return
- ESP_OK
- ESP_ERR_WIFI_NOT_INIT
- ESP_ERR_MESH_NOT_ALLOWED

esp_err_t **esp_mesh_disable_ps**(void)
Disable mesh Power Save function.

Attention This API shall be called before mesh is started.
Return
- ESP_OK
- ESP_ERR_WIFI_NOT_INIT
- ESP_ERR_MESH_NOT_ALLOWED

bool **esp_mesh_is_ps_enabled**(void)
Check whether the mesh Power Save function is enabled.

Return true/false

bool **esp_mesh_is_device_active**(void)
Check whether the device is in active state.

- If the device is not in active state, it will neither transmit nor receive frames.

Return true/false

esp_err_t **esp_mesh_set_active_duty_cycle**(int *dev_duty*, int *dev_duty_type*)

Set the device duty cycle and type.

- The range of dev_duty values is 1 to 100. The default value is 10.
- dev_duty = 100, the PS will be stopped.
- dev_duty is better to not less than 5.
- dev_duty_type could be MESH_PS_DEVICE_DUTY_REQUEST or MESH_PS_DEVICE_DUTY_DEMAND.
- If dev_duty_type is set to MESH_PS_DEVICE_DUTY_REQUEST, the device will use a nwk_duty provided by the network.
- If dev_duty_type is set to MESH_PS_DEVICE_DUTY_DEMAND, the device will use the specified dev_duty.

Attention This API can be called at any time after mesh is started.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] dev_duty`: device duty cycle
- `[in] dev_duty_type`: device PS duty cycle type, not accept MESH_PS_NETWORK_DUTY_MASTER

esp_err_t **esp_mesh_get_active_duty_cycle**(int **dev_duty*, int **dev_duty_type*)

Get device duty cycle and type.

Return

- ESP_OK

Parameters

- `[out] dev_duty`: device duty cycle
- `[out] dev_duty_type`: device PS duty cycle type

esp_err_t **esp_mesh_set_network_duty_cycle**(int *nwk_duty*, int *duration_mins*, int *applied_rule*)

Set the network duty cycle, duration and rule.

- The range of nwk_duty values is 1 to 100. The default value is 10.
- nwk_duty is the network duty cycle the entire network or the up-link path will use. A device that successfully sets the nwk_duty is known as a NWK-DUTY-MASTER.
- duration_mins specifies how long the specified nwk_duty will be used. Once duration_mins expires, the root will take over as the NWK-DUTY-MASTER. If an existing NWK-DUTY-MASTER leaves the network, the root will take over as the NWK-DUTY-MASTER again.
- duration_mins = (-1) represents nwk_duty will be used until a new NWK-DUTY-MASTER with a different nwk_duty appears.
- Only the root can set duration_mins to (-1).
- If applied_rule is set to MESH_PS_NETWORK_DUTY_APPLIED_ENTIRE, the nwk_duty will be used by the entire network.
- If applied_rule is set to MESH_PS_NETWORK_DUTY_APPLIED_UPLINK, the nwk_duty will only be used by the up-link path nodes.
- The root does not accept MESH_PS_NETWORK_DUTY_APPLIED_UPLINK.
- A nwk_duty with duration_mins(-1) set by the root is the default network duty cycle used by the entire network.

Attention This API can be called at any time after mesh is started.

- In self-organized network, if this API is called before mesh is started in all devices, (1)nwk_duty shall be set to the same value for all devices; (2)duration_mins shall be set to (-1); (3)applied_rule shall be set to MESH_PS_NETWORK_DUTY_APPLIED_ENTIRE; after the voted root appears, the root will become the NWK-DUTY-MASTER and broadcast the nwk_duty and its identity of NWK-DUTY-MASTER.
- If the root is specified (FIXED-ROOT), call this API in the root to provide a default nwk_duty for the entire network.
- After joins the network, any device can call this API to change the nwk_duty, duration_mins or applied_rule.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] nwk_duty`: network duty cycle
- `[in] duration_mins`: duration (unit: minutes)
- `[in] applied_rule`: only support MESH_PS_NETWORK_DUTY_APPLIED_ENTIRE

esp_err_t **esp_mesh_get_network_duty_cycle**(int **nwk_duty*, int **duration_mins*, int **dev_duty_type*, int **applied_rule*)

Get the network duty cycle, duration, type and rule.

Return

- ESP_OK

Parameters

- `[out] nwk_duty`: current network duty cycle
- `[out] duration_mins`: the duration of current nwk_duty
- `[out] dev_duty_type`: if it includes MESH_PS_DEVICE_DUTY_MASTER, this device is the current NWK-DUTY-MASTER.
- `[out] applied_rule`: MESH_PS_NETWORK_DUTY_APPLIED_ENTIRE

int **esp_mesh_get_running_active_duty_cycle**(void)

Get the running active duty cycle.

- The running active duty cycle of the root is 100.
- If duty type is set to MESH_PS_DEVICE_DUTY_REQUEST, the running active duty cycle is nwk_duty provided by the network.
- If duty type is set to MESH_PS_DEVICE_DUTY_DEMAND, the running active duty cycle is dev_duty specified by the users.
- In a mesh network, devices are typically working with a certain duty-cycle (transmitting, receiving and sleep) to reduce the power consumption. The running active duty cycle decides the amount of awake time within a beacon interval. At each start of beacon interval, all devices wake up, broadcast beacons, and transmit packets if they do have pending packets for their parents or for their children. Note that Low-duty-cycle means devices may not be active in most of the time, the latency of data transmission might be greater.

Return the running active duty cycle

esp_err_t **esp_mesh_ps_duty_signaling**(int *fwd_times*)

Duty signaling.

Return

- ESP_OK

Parameters

- `[in] fwd_times`: the times of forwarding duty signaling packets

Unions

union mesh_addr_t

#include <esp_mesh.h> Mesh address.

Public Members

uint8_t **addr**[6]

mac address

mip_t **mip**

mip address

union mesh_event_info_t

#include <esp_mesh.h> Mesh event information.

Public Members

mesh_event_channel_switch_t **channel_switch**
channel switch

mesh_event_child_connected_t **child_connected**
child connected

mesh_event_child_disconnected_t **child_disconnected**
child disconnected

mesh_event_routing_table_change_t **routing_table**
routing table change

mesh_event_connected_t **connected**
parent connected

mesh_event_disconnected_t **disconnected**
parent disconnected

mesh_event_no_parent_found_t **no_parent**
no parent found

mesh_event_layer_change_t **layer_change**
layer change

mesh_event_toDS_state_t **toDS_state**
toDS state, devices shall check this state firstly before trying to send packets to external IP network. This state indicates right now whether the root is capable of sending packets out. If not, devices had better to wait until this state changes to be MESH_TODS_REACHABLE.

mesh_event_vote_started_t **vote_started**
vote started

mesh_event_root_address_t **root_addr**
root address

mesh_event_root_switch_req_t **switch_req**
root switch request

mesh_event_root_conflict_t **root_conflict**
other powerful root

mesh_event_root_fixed_t **root_fixed**
fixed root

mesh_event_scan_done_t **scan_done**
scan done

mesh_event_network_state_t **network_state**
network state, such as whether current mesh network has a root.

mesh_event_find_network_t **find_network**
network found that can join

mesh_event_router_switch_t **router_switch**
new router information

mesh_event_ps_duty_t **ps_duty**
PS duty information

union mesh_rc_config_t
#include <esp_mesh.h> Vote address configuration.

Public Members

int **attempts**
max vote attempts before a new root is elected automatically by mesh network. (min:15, 15 by default)

mesh_addr_t **rc_addr**
a new root address specified by users for API esp_mesh_waive_root()

Structures

struct mip_t
IP address and port.

Public Members

ip4_addr_t **ip4**
IP address

uint16_t **port**
port

struct mesh_event_channel_switch_t
Channel switch information.

Public Members

uint8_t **channel**
new channel

struct mesh_event_connected_t
Parent connected information.

Public Members

wifi_event_sta_connected_t **connected**
parent information, same as Wi-Fi event SYSTEM_EVENT_STA_CONNECTED does

uint16_t **self_layer**
layer

uint8_t **duty**
parent duty

struct mesh_event_no_parent_found_t
No parent found information.

Public Members

int **scan_times**
scan times being through

struct mesh_event_layer_change_t
Layer change information.

Public Members

uint16_t **new_layer**
new layer

struct mesh_event_vote_started_t
vote started information

Public Members

int **reason**
vote reason, vote could be initiated by children or by the root itself

int **attempts**
max vote attempts before stopped

mesh_addr_t **rc_addr**
root address specified by users via API esp_mesh_waive_root()

struct mesh_event_find_network_t
find a mesh network that this device can join

Public Members

uint8_t **channel**
channel number of the new found network

uint8_t **router_bssid**[6]
router BSSID

struct mesh_event_root_switch_req_t
Root switch request information.

Public Members

int **reason**
root switch reason, generally root switch is initialized by users via API esp_mesh_waive_root()

mesh_addr_t **rc_addr**
the address of root switch requester

struct mesh_event_root_conflict_t
Other powerful root address.

Public Members

int8_t **rssi**
rssi with router

uint16_t **capacity**
the number of devices in current network

uint8_t **addr**[6]
other powerful root address

struct mesh_event_routing_table_change_t
Routing table change.

Public Members

uint16_t **rt_size_new**
the new value

uint16_t **rt_size_change**
the changed value

struct mesh_event_root_fixed_t
Root fixed.

Public Members

bool **is_fixed**
status

struct mesh_event_scan_done_t
Scan done event information.

Public Members

uint8_t **number**
the number of APs scanned

struct mesh_event_network_state_t
Network state information.

Public Members

bool **is_rootless**
whether current mesh network has a root

struct mesh_event_ps_duty_t
PS duty information.

Public Members

uint8_t **duty**
parent or child duty

mesh_event_child_connected_t **child_connected**
child info

struct mesh_opt_t
Mesh option.

Public Members

uint8_t **type**
option type

uint16_t **len**
option length

uint8_t ***val**
option value

struct mesh_data_t
Mesh data for esp_mesh_send() and esp_mesh_recv()

Public Members

uint8_t ***data**
data

uint16_t **size**
data size

mesh_proto_t **proto**
data protocol

mesh_tos_t **tos**
data type of service

struct mesh_router_t
Router configuration.

Public Members

uint8_t **ssid**[32]
SSID

uint8_t **ssid_len**
length of SSID

uint8_t **bssid**[6]
BSSID, if this value is specified, users should also specify "allow_router_switch".

uint8_t **password**[64]
password

bool **allow_router_switch**
if the BSSID is specified and this value is also set, when the router of this specified BSSID fails to be found after "fail" (mesh_attempts_t) times, the whole network is allowed to switch to another router with the same SSID. The new router might also be on a different channel. The default value is false. There is a risk that if the password is different between the new switched router and the previous one, the mesh network could be established but the root will never connect to the new switched router.

struct mesh_ap_cfg_t
Mesh softAP configuration.

Public Members

uint8_t **password**[64]
mesh softAP password

uint8_t **max_connection**
max number of stations allowed to connect in, default 6, max 10 = max_connection + non-mesh_max_connectionmax mesh connections

uint8_t **nonmesh_max_connection**
max non-mesh connections

struct mesh_cfg_t
Mesh initialization configuration.

Public Members

uint8_t **channel**
channel, the mesh network on

bool **allow_channel_switch**
if this value is set, when "fail" (mesh_attempts_t) times is reached, device will change to a full channel scan for a network that could join. The default value is false.

mesh_addr_t **mesh_id**
mesh network identification

mesh_router_t **router**
router configuration

mesh_ap_cfg_t **mesh_ap**
mesh softAP configuration

const mesh_crypto_funcs_t ***crypto_funcs**
crypto functions

struct mesh_vote_t
Vote.

Public Members

float **percentage**
vote percentage threshold for approval of being a root

bool **is_rc_specified**
if true, rc_addr shall be specified (Unimplemented). if false, attempts value shall be specified to make network start root election.

mesh_rc_config_t **config**
vote address configuration

struct mesh_tx_pending_t
The number of packets pending in the queue waiting to be sent by the mesh stack.

Public Members

int **to_parent**
to parent queue

int **to_parent_p2p**
to parent (P2P) queue

int **to_child**
to child queue

int **to_child_p2p**
to child (P2P) queue

int **mgmt**
management queue

int **broadcast**
broadcast and multicast queue

struct mesh_rx_pending_t
The number of packets available in the queue waiting to be received by applications.

Public Members

int **toDS**
to external DS

int **toSelf**
to self

Macros

MESH_ROOT_LAYER
root layer value

MESH_MTU
max transmit unit(in bytes)

MESH_MPS
max payload size(in bytes)

ESP_ERR_MESH_WIFI_NOT_START
Mesh error code definition.

Wi-Fi isn’t started

ESP_ERR_MESH_NOT_INIT
mesh isn’t initialized

ESP_ERR_MESH_NOT_CONFIG
mesh isn’t configured

ESP_ERR_MESH_NOT_START
mesh isn’t started

ESP_ERR_MESH_NOT_SUPPORT
not supported yet

ESP_ERR_MESH_NOT_ALLOWED
operation is not allowed

ESP_ERR_MESH_NO_MEMORY
out of memory

ESP_ERR_MESH_ARGUMENT
illegal argument

ESP_ERR_MESH_EXCEED_MTU
packet size exceeds MTU

ESP_ERR_MESH_TIMEOUT
timeout

ESP_ERR_MESH_DISCONNECTED
disconnected with parent on station interface

ESP_ERR_MESH_QUEUE_FAIL
queue fail

ESP_ERR_MESH_QUEUE_FULL
queue full

ESP_ERR_MESH_NO_PARENT_FOUND
no parent found to join the mesh network

ESP_ERR_MESH_NO_ROUTE_FOUND
no route found to forward the packet

ESP_ERR_MESH_OPTION_NULL
no option found

ESP_ERR_MESH_OPTION_UNKNOWN
unknown option

ESP_ERR_MESH_XON_NO_WINDOW
no window for software flow control on upstream

ESP_ERR_MESH_INTERFACE
low-level Wi-Fi interface error

ESP_ERR_MESH_DISCARD_DUPLICATE
discard the packet due to the duplicate sequence number

ESP_ERR_MESH_DISCARD
discard the packet

ESP_ERR_MESH_VOTING
vote in progress

ESP_ERR_MESH_XMIT
XMIT

ESP_ERR_MESH_QUEUE_READ
error in reading queue

ESP_ERR_MESH_PS
mesh PS is not specified as enable or disable

ESP_ERR_MESH_RECV_RELEASE
release esp_mesh_recv_toDS

MESH_DATA_ENC
Flags bitmap for esp_mesh_send() and esp_mesh_recv()

data encrypted (Unimplemented)

MESH_DATA_P2P
point-to-point delivery over the mesh network

MESH_DATA_FROMDS
receive from external IP network

MESH_DATA_TODS
identify this packet is target to external IP network

MESH_DATA_NONBLOCK
esp_mesh_send() non-block

MESH_DATA_DROP
in the situation of the root having been changed, identify this packet can be dropped by new root

MESH_DATA_GROUP
identify this packet is target to a group address

MESH_OPT_SEND_GROUP
Option definitions for esp_mesh_send() and esp_mesh_recv()

data transmission by group; used with esp_mesh_send() and shall have payload

MESH_OPT_RECV_DS_ADDR
return a remote IP address; used with esp_mesh_send() and esp_mesh_recv()

MESH_ASSOC_FLAG_VOTE_IN_PROGRESS
Flag of mesh networking IE.

vote in progress

MESH_ASSOC_FLAG_NETWORK_FREE
no root in current network

MESH_ASSOC_FLAG_ROOTS_FOUND
root conflict is found

MESH_ASSOC_FLAG_ROOT_FIXED
fixed root

MESH_PS_DEVICE_DUTY_REQUEST
Mesh PS (Power Save) duty cycle type.

requests to join a network PS without specifying a device duty cycle. After the device joins the network, a network duty cycle will be provided by the network

MESH_PS_DEVICE_DUTY_DEMAND
requests to join a network PS and specifies a demanded device duty cycle

MESH_PS_NETWORK_DUTY_MASTER
indicates the device is the NWK-DUTY-MASTER (network duty cycle master)

MESH_PS_NETWORK_DUTY_APPLIED_ENTIRE
Mesh PS (Power Save) duty cycle applied rule.

MESH_PS_NETWORK_DUTY_APPLIED_UPLINK

MESH_INIT_CONFIG_DEFAULT()

Type Definitions

typedef *mesh_addr_t* **mesh_event_root_address_t**
Root address.

typedef *wifi_event_sta_disconnected_t* **mesh_event_disconnected_t**
Parent disconnected information.

typedef *wifi_event_ap_staconnected_t* **mesh_event_child_connected_t**
Child connected information.

typedef *wifi_event_ap_stadisconnected_t* **mesh_event_child_disconnected_t**
Child disconnected information.

typedef *wifi_event_sta_connected_t* **mesh_event_router_switch_t**
New router information.

Enumerations

enum mesh_event_id_t
Enumerated list of mesh event id.

Values:

MESH_EVENT_STARTED
mesh is started

MESH_EVENT_STOPPED
mesh is stopped

MESH_EVENT_CHANNEL_SWITCH
channel switch

MESH_EVENT_CHILD_CONNECTED
a child is connected on softAP interface

MESH_EVENT_CHILD_DISCONNECTED
a child is disconnected on softAP interface

MESH_EVENT_ROUTING_TABLE_ADD
routing table is changed by adding newly joined children

MESH_EVENT_ROUTING_TABLE_REMOVE
routing table is changed by removing leave children

MESH_EVENT_PARENT_CONNECTED
parent is connected on station interface

MESH_EVENT_PARENT_DISCONNECTED
parent is disconnected on station interface

MESH_EVENT_NO_PARENT_FOUND
no parent found

MESH_EVENT_LAYER_CHANGE
layer changes over the mesh network

MESH_EVENT_TODS_STATE
state represents whether the root is able to access external IP network. This state is a manual event that needs to be triggered with esp_mesh_post_toDS_state().

MESH_EVENT_VOTE_STARTED
the process of voting a new root is started either by children or by the root

MESH_EVENT_VOTE_STOPPED
the process of voting a new root is stopped

MESH_EVENT_ROOT_ADDRESS
the root address is obtained. It is posted by mesh stack automatically.

MESH_EVENT_ROOT_SWITCH_REQ
root switch request sent from a new voted root candidate

MESH_EVENT_ROOT_SWITCH_ACK
root switch acknowledgment responds the above request sent from current root

MESH_EVENT_ROOT_ASKED_YIELD
the root is asked yield by a more powerful existing root. If self organized is disabled and this device is specified to be a root by users, users should set a new parent for this device. if self organized is enabled, this device will find a new parent by itself, users could ignore this event.

MESH_EVENT_ROOT_FIXED
when devices join a network, if the setting of Fixed Root for one device is different from that of its parent, the device will update the setting the same as its parent's. Fixed Root Setting of each device is variable as that setting changes of the root.

MESH_EVENT_SCAN_DONE
if self-organized networking is disabled, user can call esp_wifi_scan_start() to trigger this event, and add the corresponding scan done handler in this event.

MESH_EVENT_NETWORK_STATE
network state, such as whether current mesh network has a root.

MESH_EVENT_STOP_RECONNECTION
the root stops reconnecting to the router and non-root devices stop reconnecting to their parents.

MESH_EVENT_FIND_NETWORK
when the channel field in mesh configuration is set to zero, mesh stack will perform a full channel scan to find a mesh network that can join, and return the channel value after finding it.

MESH_EVENT_ROUTER_SWITCH
if users specify BSSID of the router in mesh configuration, when the root connects to another router with the same SSID, this event will be posted and the new router information is attached.

MESH_EVENT_PS_PARENT_DUTY
parent duty

MESH_EVENT_PS_CHILD_DUTY
child duty

MESH_EVENT_PS_DEVICE_DUTY
device duty

MESH_EVENT_MAX

enum mesh_type_t
Device type.

Values:

MESH_IDLE
hasn't joined the mesh network yet

MESH_ROOT
the only sink of the mesh network. Has the ability to access external IP network

MESH_NODE
intermediate device. Has the ability to forward packets over the mesh network

MESH_LEAF
has no forwarding ability

MESH_STA
connect to router with a standlone Wi-Fi station mode, no network expansion capability

enum mesh_proto_t
Protocol of transmitted application data.

Values:

MESH_PROTO_BIN
binary

MESH_PROTO_HTTP
HTTP protocol

MESH_PROTO_JSON
JSON format

MESH_PROTO_MQTT
MQTT protocol

MESH_PROTO_AP
IP network mesh communication of node’ s AP interface

MESH_PROTO_STA
IP network mesh communication of node’ s STA interface

enum mesh_tos_t
For reliable transmission, mesh stack provides three type of services.

Values:

MESH_TOS_P2P
provide P2P (point-to-point) retransmission on mesh stack by default

MESH_TOS_E2E
provide E2E (end-to-end) retransmission on mesh stack (Unimplemented)

MESH_TOS_DEF
no retransmission on mesh stack

enum mesh_vote_reason_t
Vote reason.

Values:

MESH_VOTE_REASON_ROOT_INITIATED = 1
vote is initiated by the root

MESH_VOTE_REASON_CHILD_INITIATED
vote is initiated by children

enum mesh_disconnect_reason_t
Mesh disconnect reason code.

Values:

MESH_REASON_CYCLIC = 100
cyclic is detected

MESH_REASON_PARENT_IDLE
parent is idle

MESH_REASON_LEAF
the connected device is changed to a leaf

MESH_REASON_DIFF_ID
in different mesh ID

MESH_REASON_ROOTS
root conflict is detected

MESH_REASON_PARENT_STOPPED
parent has stopped the mesh

MESH_REASON_SCAN_FAIL
scan fail

MESH_REASON_IE_UNKNOWN
unknown IE

MESH_REASON_WAIVE_ROOT
waive root

MESH_REASON_PARENT_WORSE
parent with very poor RSSI

MESH_REASON_EMPTY_PASSWORD
use an empty password to connect to an encrypted parent

MESH_REASON_PARENT_UNENCRYPTED
connect to an unencrypted parent/router

enum esp_mesh_topology_t
Mesh topology.

Values:

MESH_TOPO_TREE
tree topology

MESH_TOPO_CHAIN
chain topology

enum mesh_event_toDS_state_t
The reachability of the root to a DS (distribute system)

Values:

MESH_TODS_UNREACHABLE
the root isn’t able to access external IP network

MESH_TODS_REACHABLE
the root is able to access external IP network

本部分的 Wi-Fi API 示例代码存放在 ESP-IDF 示例项目的 wifi 目录下。

ESP-MESH 的示例代码存放在 ESP-IDF 示例项目的 mesh 目录下。

2.1.2 以太网

以太网

应用示例

- 以太网基本示例：ethernet/basic.
- 以太网 iperf 示例：ethernet/iperf.

以太网驱动程序模型

- esp_eth/include/esp_eth.h

以太网通用接口

- esp_eth/include/esp_eth_com.h

以太网 MAC 接口

- esp_eth/include/esp_eth_mac.h

以太网 PHY 接口

- esp_eth/include/esp_eth_phy.h

以太网 PHY 公共寄存器

- esp_eth/include/eth_phy_regs_struct.h

API 参考-驱动程序模型

Header File

- esp_eth/include/esp_eth.h

Functions

esp_err_t **esp_eth_driver_install**(**const** *esp_eth_config_t* *config, *esp_eth_handle_t* *out_hdl)

Install Ethernet driver.

Return

- ESP_OK: install esp_eth driver successfully
- ESP_ERR_INVALID_ARG: install esp_eth driver failed because of some invalid argument
- ESP_ERR_NO_MEM: install esp_eth driver failed because there's no memory for driver
- ESP_FAIL: install esp_eth driver failed because some other error occurred

Parameters

- `[in] config`: configuration of the Ethernet driver
- `[out] out_hdl`: handle of Ethernet driver

esp_err_t **esp_eth_driver_uninstall**(*esp_eth_handle_t* *hdl*)

Uninstall Ethernet driver.

Note It's not recommended to uninstall Ethernet driver unless it won't get used any more in application code. To uninstall Ethernet driver, you have to make sure, all references to the driver are released. Ethernet driver can only be uninstalled successfully when reference counter equals to one.

Return

- ESP_OK: uninstall esp_eth driver successfully
- ESP_ERR_INVALID_ARG: uninstall esp_eth driver failed because of some invalid argument
- ESP_ERR_INVALID_STATE: uninstall esp_eth driver failed because it has more than one reference
- ESP_FAIL: uninstall esp_eth driver failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver

esp_err_t **esp_eth_start**(*esp_eth_handle_t* *hdl*)

Start Ethernet driver **ONLY** in standalone mode (i.e. without TCP/IP stack)

Note This API will start driver state machine and internal software timer (for checking link status).

Return

- ESP_OK: start esp_eth driver successfully
- ESP_ERR_INVALID_ARG: start esp_eth driver failed because of some invalid argument
- ESP_ERR_INVALID_STATE: start esp_eth driver failed because driver has started already
- ESP_FAIL: start esp_eth driver failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver

esp_err_t **esp_eth_stop**(*esp_eth_handle_t hdl*)

Stop Ethernet driver.

Note This function does the oppsite operation of `esp_eth_start`.

Return

- ESP_OK: stop esp_eth driver successfully
- ESP_ERR_INVALID_ARG: stop esp_eth driver failed because of some invalid argument
- ESP_ERR_INVALID_STATE: stop esp_eth driver failed because driver has not started yet
- ESP_FAIL: stop esp_eth driver failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver

esp_err_t **esp_eth_update_input_path**(*esp_eth_handle_t hdl*, *esp_err_t* (**stack_input*)) *esp_eth_handle_t* hdl, uint8_t *buffer, uint32_t length, void *priv , void **priv*Update Ethernet data input path (i.e. specify where to pass the input buffer)

Note After install driver, Ethernet still don't know where to deliver the input buffer. In fact, this API registers a callback function which get invoked when Ethernet received new packets.

Return

- ESP_OK: update input path successfully
- ESP_ERR_INVALID_ARG: update input path failed because of some invalid argument
- ESP_FAIL: update input path failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver
- `[in] stack_input`: function pointer, which does the actual process on incoming packets
- `[in] priv`: private resource, which gets passed to `stack_input` callback without any modification

esp_err_t **esp_eth_transmit**(*esp_eth_handle_t hdl*, void **buf*, uint32_t *length*)

General Transmit.

Return

- ESP_OK: transmit frame buffer successfully
- ESP_ERR_INVALID_ARG: transmit frame buffer failed because of some invalid argument
- ESP_FAIL: transmit frame buffer failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver
- `[in] buf`: buffer of the packet to transfer
- `[in] length`: length of the buffer to transfer

esp_err_t **esp_eth_receive**(*esp_eth_handle_t hdl*, uint8_t **buf*, uint32_t **length*)

General Receive.

Note Before this function got invoked, the value of "length" should set by user, equals the size of buffer. After the function returned, the value of "length" means the real length of received data.

Return

- ESP_OK: receive frame buffer successfully
- ESP_ERR_INVALID_ARG: receive frame buffer failed because of some invalid argument
- ESP_ERR_INVALID_SIZE: input buffer size is not enough to hold the incoming data. in this case, value of returned "length" indicates the real size of incoming data.
- ESP_FAIL: receive frame buffer failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver
- `[out] buf`: buffer to preserve the received packet

- `[out] length`: length of the received packet

esp_err_t **esp_eth_ioctl**(*esp_eth_handle_t hdl*, *esp_eth_io_cmd_t cmd*, void **data*)

Misc IO function of Etherent driver.

Return

- ESP_OK: process io command successfully
- ESP_ERR_INVALID_ARG: process io command failed because of some invalid argument
- ESP_FAIL: process io command failed because some other error occurred

Parameters

- `[in] hdl`: handle of Ethernet driver
- `[in] cmd`: IO control command
- `[in] data`: specificed data for command

esp_err_t **esp_eth_increase_reference**(*esp_eth_handle_t hdl*)

Increase Ethernet driver reference.

Note Ethernet driver handle can be obtained by os timer, netif, etc. It's dangerous when thread A is using Ethernet but thread B uninstall the driver. Using reference counter can prevent such risk, but care should be taken, when you obtain Ethernet driver, this API must be invoked so that the driver won't be uninstalled during your using time.

Return

- ESP_OK: increase reference successfully
- ESP_ERR_INVALID_ARG: increase reference failed because of some invalid argument

Parameters

- `[in] hdl`: handle of Ethernet driver

esp_err_t **esp_eth_decrease_reference**(*esp_eth_handle_t hdl*)

Decrease Ethernet driver reference.

Return

- ESP_OK: increase reference successfully
- ESP_ERR_INVALID_ARG: increase reference failed because of some invalid argument

Parameters

- `[in] hdl`: handle of Ethernet driver

Structures

struct esp_eth_config_t

Configuration of Ethernet driver.

Public Members

esp_eth_mac_t ***mac**

Ethernet MAC object.

esp_eth_phy_t ***phy**

Ethernet PHY object.

uint32_t **check_link_period_ms**

Period time of checking Ethernet link status.

esp_err_t (***stack_input**)(*esp_eth_handle_t* eth_handle, uint8_t *buffer, uint32_t length, void *priv)

Input frame buffer to user's stack.

Return

- ESP_OK: input frame buffer to upper stack successfully
- ESP_FAIL: error occurred when inputting buffer to upper stack

Parameters

- `[in] eth_handle`: handle of Ethernet driver
- `[in] buffer`: frame buffer that will get input to upper stack
- `[in] length`: length of the frame buffer

esp_err_t (***on_lowlevel_init_done**) (*esp_eth_handle_t* eth_handle)
Callback function invoked when lowlevel initialization is finished.

Return
- ESP_OK: process extra lowlevel initialization successfully
- ESP_FAIL: error occurred when processing extra lowlevel initialization

Parameters
- `[in] eth_handle`: handle of Ethernet driver

esp_err_t (***on_lowlevel_deinit_done**) (*esp_eth_handle_t* eth_handle)
Callback function invoked when lowlevel deinitialization is finished.

Return
- ESP_OK: process extra lowlevel deinitialization successfully
- ESP_FAIL: error occurred when processing extra lowlevel deinitialization

Parameters
- `[in] eth_handle`: handle of Ethernet driver

Macros

ETH_DEFAULT_CONFIG(emac, ephy)
Default configuration for Ethernet driver.

Type Definitions

typedef void ***esp_eth_handle_t**
Handle of Ethernet driver.

API 参考–通用接口

Header File

- esp_eth/include/esp_eth_com.h

Functions

esp_err_t **esp_eth_detect_phy_addr**(*esp_eth_mediator_t* **eth*, uint32_t **detected_addr*)
Detect PHY address.

Return
- ESP_OK: detect phy address successfully
- ESP_ERR_INVALID_ARG: invalid parameter
- ESP_ERR_NOT_FOUND: can't detect any PHY device
- ESP_FAIL: detect phy address failed because some error occurred

Parameters
- `[in] eth`: mediator of Ethernet driver
- `[out] detected_addr`: a valid address after detection

Structures

struct esp_eth_mediator_s
Ethernet mediator.

Public Members

esp_err_t (***phy_reg_read**) (*esp_eth_mediator_t* *eth, uint32_t phy_addr, uint32_t phy_reg, uint32_t *reg_value)
Read PHY register.

Return

- ESP_OK: read PHY register successfully
- ESP_FAIL: read PHY register failed because some error occurred

Parameters

- `[in] eth`: mediator of Ethernet driver
- `[in] phy_addr`: PHY Chip address (0~31)
- `[in] phy_reg`: PHY register index code
- `[out] reg_value`: PHY register value

esp_err_t (***phy_reg_write**) (*esp_eth_mediator_t* *eth, uint32_t phy_addr, uint32_t phy_reg, uint32_t reg_value)

Write PHY register.

Return

- ESP_OK: write PHY register successfully
- ESP_FAIL: write PHY register failed because some error occurred

Parameters

- `[in] eth`: mediator of Ethernet driver
- `[in] phy_addr`: PHY Chip address (0~31)
- `[in] phy_reg`: PHY register index code
- `[in] reg_value`: PHY register value

esp_err_t (***stack_input**) (*esp_eth_mediator_t* *eth, uint8_t *buffer, uint32_t length)

Deliver packet to upper stack.

Return

- ESP_OK: deliver packet to upper stack successfully
- ESP_FAIL: deliver packet failed because some error occurred

Parameters

- `[in] eth`: mediator of Ethernet driver
- `[in] buffer`: packet buffer
- `[in] length`: length of the packet

esp_err_t (***on_state_changed**) (*esp_eth_mediator_t* *eth, *esp_eth_state_t* state, void *args)

Callback on Ethernet state changed.

Return

- ESP_OK: process the new state successfully
- ESP_FAIL: process the new state failed because some error occurred

Parameters

- `[in] eth`: mediator of Ethernet driver
- `[in] state`: new state
- `[in] args`: optional argument for the new state

Macros

ETH_MAX_PAYLOAD_LEN

Maximum Ethernet payload size.

ETH_MIN_PAYLOAD_LEN

Minimum Ethernet payload size.

ETH_HEADER_LEN

Ethernet frame header size: Dest addr(6 Bytes) + Src addr(6 Bytes) + length/type(2 Bytes)

ETH_CRC_LEN

Ethernet frame CRC length.

ETH_VLAN_TAG_LEN

Optional 802.1q VLAN Tag length.

ETH_JUMBO_FRAME_PAYLOAD_LEN

Jumbo frame payload size.

ETH_MAX_PACKET_SIZE

Maximum frame size (1522 Bytes)

ETH_MIN_PACKET_SIZE
Minimum frame size (64 Bytes)

Type Definitions

typedef struct *esp_eth_mediator_s* **esp_eth_mediator_t**
Ethernet mediator.

Enumerations

enum esp_eth_state_t
Ethernet driver state.

Values:

ETH_STATE_LLINIT
Lowlevel init done

ETH_STATE_DEINIT
Deinit done

ETH_STATE_LINK
Link status changed

ETH_STATE_SPEED
Speed updated

ETH_STATE_DUPLEX
Duplex updated

enum esp_eth_io_cmd_t
Command list for ioctl API.

Values:

ETH_CMD_G_MAC_ADDR
Get MAC address

ETH_CMD_S_MAC_ADDR
Set MAC address

ETH_CMD_G_PHY_ADDR
Get PHY address

ETH_CMD_S_PHY_ADDR
Set PHY address

ETH_CMD_G_SPEED
Get Speed

ETH_CMD_S_PROMISCUOUS
Set promiscuous mode

enum eth_link_t
Ethernet link status.

Values:

ETH_LINK_UP
Ethernet link is up

ETH_LINK_DOWN
Ethernet link is down

enum eth_speed_t
Ethernet speed.

Values:

ETH_SPEED_10M
Ethernet speed is 10Mbps

ETH_SPEED_100M
Ethernet speed is 100Mbps

enum eth_duplex_t
Ethernet duplex mode.

Values:

ETH_DUPLEX_HALF
Ethernet is in half duplex

ETH_DUPLEX_FULL
Ethernet is in full duplex

enum eth_event_t
Ethernet event declarations.

Values:

ETHERNET_EVENT_START
Ethernet driver start

ETHERNET_EVENT_STOP
Ethernet driver stop

ETHERNET_EVENT_CONNECTED
Ethernet got a valid link

ETHERNET_EVENT_DISCONNECTED
Ethernet lost a valid link

API 参考–MAC 接口

Header File

- esp_eth/include/esp_eth_mac.h

Structures

struct esp_eth_mac_s
Ethernet MAC.

Public Members

esp_err_t (***set_mediator**)(*esp_eth_mac_t* *mac, *esp_eth_mediator_t* *eth)
Set mediator for Ethernet MAC.

Return
- ESP_OK: set mediator for Ethernet MAC successfully
- ESP_ERR_INVALID_ARG: set mediator for Ethernet MAC failed because of invalid argument

Parameters
- `[in] mac`: Ethernet MAC instance
- `[in] eth`: Ethernet mediator

esp_err_t (***init**)(*esp_eth_mac_t* *mac)
Initialize Ethernet MAC.

Return
- ESP_OK: initialize Ethernet MAC successfully
- ESP_ERR_TIMEOUT: initialize Ethernet MAC failed because of timeout

- ESP_FAIL: initialize Ethernet MAC failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance

esp_err_t (***deinit**) (*esp_eth_mac_t* *mac)

Deinitialize Ethernet MAC.

Return

- ESP_OK: deinitialize Ethernet MAC successfully
- ESP_FAIL: deinitialize Ethernet MAC failed because some error occurred

Parameters

- `[in] mac`: Ethernet MAC instance

esp_err_t (***start**) (*esp_eth_mac_t* *mac)

Start Ethernet MAC.

Return

- ESP_OK: start Ethernet MAC successfully
- ESP_FAIL: start Ethernet MAC failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance

esp_err_t (***stop**) (*esp_eth_mac_t* *mac)

Stop Ethernet MAC.

Return

- ESP_OK: stop Ethernet MAC successfully
- ESP_FAIL: stop Ethernet MAC failed because some error occurred

Parameters

- `[in] mac`: Ethernet MAC instance

esp_err_t (***transmit**) (*esp_eth_mac_t* *mac, uint8_t *buf, uint32_t length)

Transmit packet from Ethernet MAC.

Return

- ESP_OK: transmit packet successfully
- ESP_ERR_INVALID_ARG: transmit packet failed because of invalid argument
- ESP_ERR_INVALID_STATE: transmit packet failed because of wrong state of MAC
- ESP_FAIL: transmit packet failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance
- `[in] buf`: packet buffer to transmit
- `[in] length`: length of packet

esp_err_t (***receive**) (*esp_eth_mac_t* *mac, uint8_t *buf, uint32_t *length)

Receive packet from Ethernet MAC.

Note Memory of buf is allocated in the Layer2, make sure it get free after process.

Note Before this function got invoked, the value of "length" should set by user, equals the size of buffer. After the function returned, the value of "length" means the real length of received data.

Return

- ESP_OK: receive packet successfully
- ESP_ERR_INVALID_ARG: receive packet failed because of invalid argument
- ESP_ERR_INVALID_SIZE: input buffer size is not enough to hold the incoming data. in this case, value of returned "length" indicates the real size of incoming data.
- ESP_FAIL: receive packet failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance
- `[out] buf`: packet buffer which will preserve the received frame
- `[out] length`: length of the received packet

esp_err_t (***read_phy_reg**) (*esp_eth_mac_t* *mac, uint32_t phy_addr, uint32_t phy_reg, uint32_t *reg_value)

Read PHY register.

Return

- ESP_OK: read PHY register successfully
- ESP_ERR_INVALID_ARG: read PHY register failed because of invalid argument
- ESP_ERR_INVALID_STATE: read PHY register failed because of wrong state of MAC
- ESP_FAIL: read PHY register failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance
- `[in] phy_addr`: PHY chip address (0~31)
- `[in] phy_reg`: PHY register index code
- `[out] reg_value`: PHY register value

esp_err_t (***write_phy_reg**) (*esp_eth_mac_t* *mac, uint32_t phy_addr, uint32_t phy_reg, uint32_t reg_value)

Write PHY register.

Return

- ESP_OK: write PHY register successfully
- ESP_ERR_INVALID_STATE: write PHY register failed because of wrong state of MAC
- ESP_ERR_TIMEOUT: write PHY register failed because of timeout
- ESP_FAIL: write PHY register failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance
- `[in] phy_addr`: PHY chip address (0~31)
- `[in] phy_reg`: PHY register index code
- `[in] reg_value`: PHY register value

esp_err_t (***set_addr**) (*esp_eth_mac_t* *mac, uint8_t *addr)

Set MAC address.

Return

- ESP_OK: set MAC address successfully
- ESP_ERR_INVALID_ARG: set MAC address failed because of invalid argument
- ESP_FAIL: set MAC address failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance
- `[in] addr`: MAC address

esp_err_t (***get_addr**) (*esp_eth_mac_t* *mac, uint8_t *addr)

Get MAC address.

Return

- ESP_OK: get MAC address successfully
- ESP_ERR_INVALID_ARG: get MAC address failed because of invalid argument
- ESP_FAIL: get MAC address failed because some other error occurred

Parameters

- `[in] mac`: Ethernet MAC instance
- `[out] addr`: MAC address

esp_err_t (***set_speed**) (*esp_eth_mac_t* *mac, *eth_speed_t* speed)

Set speed of MAC.

Return

- ESP_OK: set MAC speed successfully
- ESP_ERR_INVALID_ARG: set MAC speed failed because of invalid argument
- ESP_FAIL: set MAC speed failed because some other error occurred

Parameters

- `[in] ma:c`: Ethernet MAC instance
- `[in] speed`: MAC speed

esp_err_t (***set_duplex**) (*esp_eth_mac_t* *mac, *eth_duplex_t* duplex)

Set duplex mode of MAC.

Return
- ESP_OK: set MAC duplex mode successfully
- ESP_ERR_INVALID_ARG: set MAC duplex failed because of invalid argument
- ESP_FAIL: set MAC duplex failed because some other error occurred

Parameters
- `[in] mac`: Ethernet MAC instance
- `[in] duplex`: MAC duplex

esp_err_t (***set_link**) (*esp_eth_mac_t* *mac, *eth_link_t* link)
Set link status of MAC.

Return
- ESP_OK: set link status successfully
- ESP_ERR_INVALID_ARG: set link status failed because of invalid argument
- ESP_FAIL: set link status failed because some other error occurred

Parameters
- `[in] mac`: Ethernet MAC instance
- `[in] link`: Link status

esp_err_t (***set_promiscuous**) (*esp_eth_mac_t* *mac, bool enable)
Set promiscuous of MAC.

Return
- ESP_OK: set promiscuous mode successfully
- ESP_FAIL: set promiscuous mode failed because some error occurred

Parameters
- `[in] mac`: Ethernet MAC instance
- `[in] enable`: set true to enable promiscuous mode; set false to disable promiscuous mode

esp_err_t (***del**) (*esp_eth_mac_t* *mac)
Free memory of Ethernet MAC.

Return
- ESP_OK: free Ethernet MAC instance successfully
- ESP_FAIL: free Ethernet MAC instance failed because some error occurred

Parameters
- `[in] mac`: Ethernet MAC instance

struct eth_mac_config_t
Configuration of Ethernet MAC object.

Public Members

uint32_t **sw_reset_timeout_ms**
Software reset timeout value (Unit: ms)

uint32_t **rx_task_stack_size**
Stack size of the receive task

uint32_t **rx_task_prio**
Priority of the receive task

int **smi_mdc_gpio_num**
SMI MDC GPIO number

int **smi_mdio_gpio_num**
SMI MDIO GPIO number

uint32_t **flags**
Flags that specify extra capability for mac driver

Macros

ETH_MAC_FLAG_WORK_WITH_CACHE_DISABLE
MAC driver can work when cache is disabled
ETH_MAC_FLAG_PIN_TO_CORE
Pin MAC task to the CPU core where driver installation happened

ETH_MAC_DEFAULT_CONFIG()
Default configuration for Ethernet MAC object.

Type Definitions

typedef struct *esp_eth_mac_s* **esp_eth_mac_t**
Ethernet MAC.

API 参考–PHY 接口

Header File

- esp_eth/include/esp_eth_phy.h

Functions

esp_eth_phy_t ***esp_eth_phy_new_ip101**(**const** *eth_phy_config_t* **config*)
Create a PHY instance of IP101.

Return
- instance: create PHY instance successfully
- NULL: create PHY instance failed because some error occurred

Parameters
- `[in] config`: configuration of PHY

esp_eth_phy_t ***esp_eth_phy_new_rtl8201**(**const** *eth_phy_config_t* **config*)
Create a PHY instance of RTL8201.

Return
- instance: create PHY instance successfully
- NULL: create PHY instance failed because some error occurred

Parameters
- `[in] config`: configuration of PHY

esp_eth_phy_t ***esp_eth_phy_new_lan8720**(**const** *eth_phy_config_t* **config*)
Create a PHY instance of LAN8720.

Return
- instance: create PHY instance successfully
- NULL: create PHY instance failed because some error occurred

Parameters
- `[in] config`: configuration of PHY

esp_eth_phy_t ***esp_eth_phy_new_dp83848**(**const** *eth_phy_config_t* **config*)
Create a PHY instance of DP83848.

Return
- instance: create PHY instance successfully
- NULL: create PHY instance failed because some error occurred

Parameters
- `[in] config`: configuration of PHY

Structures

struct esp_eth_phy_s
Ethernet PHY.

Public Members

esp_err_t (***set_mediator**) (*esp_eth_phy_t* *phy, *esp_eth_mediator_t* *mediator)
Set mediator for PHY.

Return
- ESP_OK: set mediator for Ethernet PHY instance successfully
- ESP_ERR_INVALID_ARG: set mediator for Ethernet PHY instance failed because of some invalid arguments

Parameters
- `[in] phy`: Ethernet PHY instance
- `[in] mediator`: mediator of Ethernet driver

esp_err_t (***reset**) (*esp_eth_phy_t* *phy)
Software Reset Ethernet PHY.

Return
- ESP_OK: reset Ethernet PHY successfully
- ESP_FAIL: reset Ethernet PHY failed because some error occurred

Parameters
- `[in] phy`: Ethernet PHY instance

esp_err_t (***reset_hw**) (*esp_eth_phy_t* *phy)
Hardware Reset Ethernet PHY.

Note Hardware reset is mostly done by pull down and up PHY' s nRST pin

Return
- ESP_OK: reset Ethernet PHY successfully
- ESP_FAIL: reset Ethernet PHY failed because some error occurred

Parameters
- `[in] phy`: Ethernet PHY instance

esp_err_t (***init**) (*esp_eth_phy_t* *phy)
Initialize Ethernet PHY.

Return
- ESP_OK: initialize Ethernet PHY successfully
- ESP_FAIL: initialize Ethernet PHY failed because some error occurred

Parameters
- `[in] phy`: Ethernet PHY instance

esp_err_t (***deinit**) (*esp_eth_phy_t* *phy)
Deinitialize Ethernet PHY.

Return
- ESP_OK: deinitialize Ethernet PHY successfully
- ESP_FAIL: deinitialize Ethernet PHY failed because some error occurred

Parameters
- `[in] phyL`: Ethernet PHY instance

esp_err_t (***negotiate**) (*esp_eth_phy_t* *phy)
Start auto negotiation.

Return
- ESP_OK: restart auto negotiation successfully
- ESP_FAIL: restart auto negotiation failed because some error occurred

Parameters
- `[in] phy`: Ethernet PHY instance

esp_err_t (***get_link**) (*esp_eth_phy_t* *phy)
Get Ethernet PHY link status.

Return
- ESP_OK: get Ethernet PHY link status successfully

- ESP_FAIL: get Ethernet PHY link status failed because some error occurred

Parameters

- `[in] phy`: Ethernet PHY instance

esp_err_t (***pwrctl**) (*esp_eth_phy_t* *phy, bool enable)

Power control of Ethernet PHY.

Return

- ESP_OK: control Ethernet PHY power successfully
- ESP_FAIL: control Ethernet PHY power failed because some error occurred

Parameters

- `[in] phy`: Ethernet PHY instance
- `[in] enable`: set true to power on Ethernet PHY; ser false to power off Ethernet PHY

esp_err_t (***set_addr**) (*esp_eth_phy_t* *phy, uint32_t addr)

Set PHY chip address.

Return

- ESP_OK: set Ethernet PHY address successfully
- ESP_FAIL: set Ethernet PHY address failed because some error occurred

Parameters

- `[in] phy`: Ethernet PHY instance
- `[in] addr`: PHY chip address

esp_err_t (***get_addr**) (*esp_eth_phy_t* *phy, uint32_t *addr)

Get PHY chip address.

Return

- ESP_OK: get Ethernet PHY address successfully
- ESP_ERR_INVALID_ARG: get Ethernet PHY address failed because of invalid argument

Parameters

- `[in] phy`: Ethernet PHY instance
- `[out] addr`: PHY chip address

esp_err_t (***del**) (*esp_eth_phy_t* *phy)

Free memory of Ethernet PHY instance.

Return

- ESP_OK: free PHY instance successfully
- ESP_FAIL: free PHY instance failed because some error occurred

Parameters

- `[in] phy`: Ethernet PHY instance

struct eth_phy_config_t

Ethernet PHY configuration.

Public Members

int32_t **phy_addr**

PHY address, set -1 to enable PHY address detection at initialization stage

uint32_t **reset_timeout_ms**

Reset timeout value (Unit: ms)

uint32_t **autonego_timeout_ms**

Auto-negotiation timeout value (Unit: ms)

int **reset_gpio_num**

Reset GPIO number, -1 means no hardware reset

Macros

ESP_ETH_PHY_ADDR_AUTO

ETH_PHY_DEFAULT_CONFIG()
Default configuration for Ethernet PHY object.

Type Definitions

typedef struct *esp_eth_phy_s* **esp_eth_phy_t**
Ethernet PHY.

API 参考–esp_netif 相关使用

Header File

- esp_eth/include/esp_eth_netif_glue.h

Functions

void ***esp_eth_new_netif_glue**(*esp_eth_handle_t eth_hdl*)
Create a netif glue for Ethernet driver.

Note netif glue is used to attach io driver to TCP/IP netif
Return glue object, which inherits esp_netif_driver_base_t
Parameters
- `eth_hdl`: Ethernet driver handle

esp_err_t **esp_eth_del_netif_glue**(void **glue*)
Delete netif glue of Ethernet driver.

Return -ESP_OK: delete netif glue successfully
Parameters
- `glue`: netif glue

esp_err_t **esp_eth_set_default_handlers**(void **esp_netif*)
Register default IP layer handlers for Ethernet.

Note : Ethernet handle might not yet properly initialized when setting up these default handlers
Return
- ESP_ERR_INVALID_ARG: invalid parameter (esp_netif is NULL)
- ESP_OK: set default IP layer handlers successfully
- others: other failure occurred during register esp_event handler

Parameters
- `[in] esp_netif`: esp network interface handle created for Ethernet driver

esp_err_t **esp_eth_clear_default_handlers**(void **esp_netif*)
Unregister default IP layer handlers for Ethernet.

Return
- ESP_ERR_INVALID_ARG: invalid parameter (esp_netif is NULL)
- ESP_OK: clear default IP layer handlers successfully
- others: other failure occurred during unregister esp_event handler

Parameters
- `[in] esp_netif`: esp network interface handle created for Ethernet driver

本部分的以太网 API 示例代码存放在 ESP-IDF 示例项目的 ethernet 目录下。

2.1.3 IP 网络层协议

ESP-NETIF

The purpose of ESP-NETIF library is twofold:

- It provides an abstraction layer for the application on top of the TCP/IP stack. This will allow applications to choose between IP stacks in the future.

- The APIs it provides are thread safe, even if the underlying TCP/IP stack APIs are not.

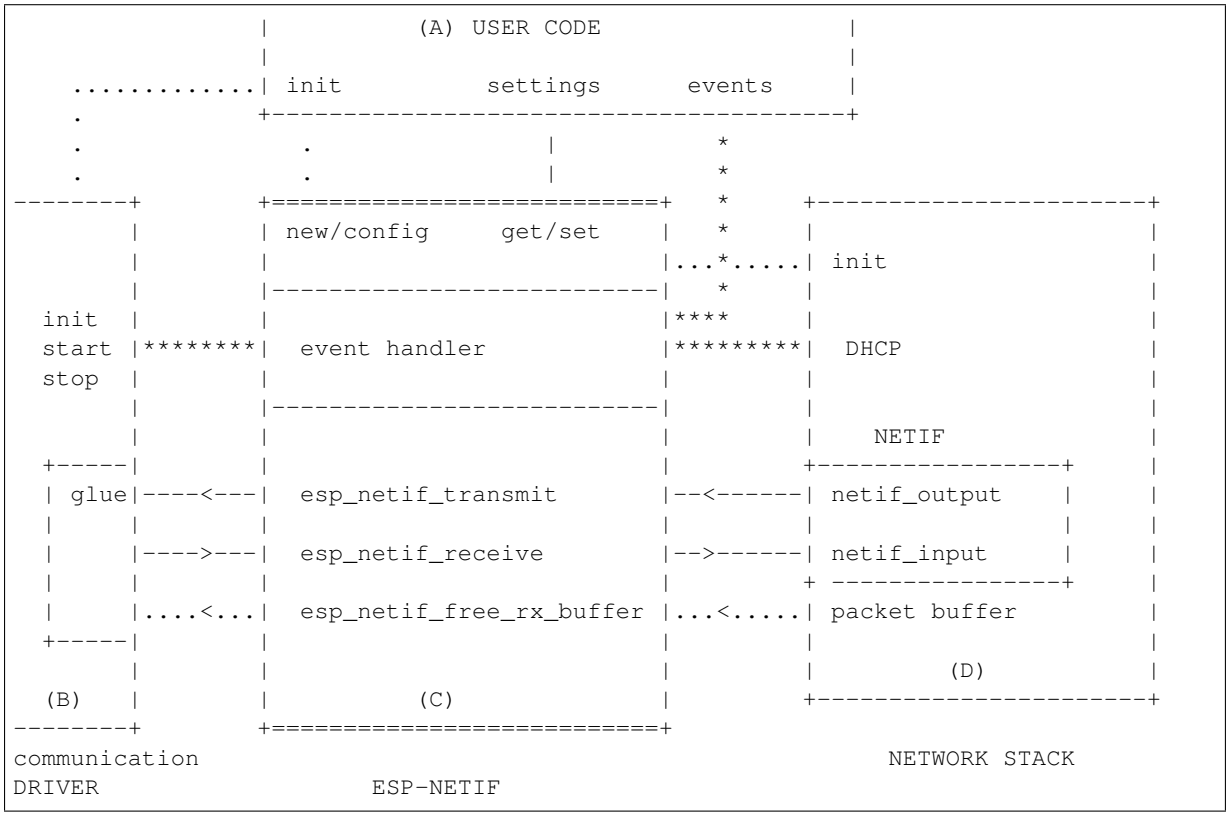
ESP-IDF currently implements ESP-NETIF for the lwIP TCP/IP stack only. However, the adapter itself is TCP/IP implementation agnostic and different implementations are possible.

Some ESP-NETIF API functions are intended to be called by application code, for example to get/set interface IP addresses, configure DHCP. Other functions are intended for internal ESP-IDF use by the network driver layer.

In many cases, applications do not need to call ESP-NETIF APIs directly as they are called from the default network event handlers.

ESP-NETIF component is a successor of the tcpip_adapter, former network interface abstraction, which has become deprecated since IDF v4.1. Please refer to the *TCP/IP 适配器迁移指南* section in case existing applications to be ported to use the esp-netif API instead.

ESP-NETIF architecture



Data and event flow in the diagram

- `. . . . . .` Initialization line from user code to ESP-NETIF and communication driver
- `--<---->--` Data packets going from communication media to TCP/IP stack and back
- `********` Events aggregated in ESP-NETIF propagates to driver, user code and network stack
- `|` User settings and runtime configuration

ESP-NETIF interaction

A) User code, boiler plate Overall application interaction with a specific IO driver for communication media and configured TCP/IP network stack is abstracted using ESP-NETIF APIs and outlined as below:

A) Initialization code

1) Initializes IO driver
2) Creates a new instance of ESP-NETIF and configure with

- ESP-NETIF specific options (flags, behaviour, name)
- Network stack options (netif init and input functions, not publicly available)
- IO driver specific options (transmit, free rx buffer functions, IO driver handle)

3) Attaches the IO driver handle to the ESP-NETIF instance created in the above steps
4) Configures event handlers

- use default handlers for common interfaces defined in IO drivers; or define a specific handlers for customised behaviour/new interfaces
- register handlers for app related events (such as IP lost/acquired)

B) Interaction with network interfaces using ESP-NETIF API

- Getting and setting TCP/IP related parameters (DHCP, IP, etc)
- Receiving IP events (connect/disconnect)
- Controlling application lifecycle (set interface up/down)

B) Communication driver, IO driver, media driver Communication driver plays these two important roles in relation with ESP-NETIF:

1) Event handlers: Define behaviour patterns of interaction with ESP-NETIF (for example: ethernet link-up -> turn netif on)
2) Glue IO layer: Adapts the input/output functions to use ESP-NETIF transmit, receive and free receive buffer

- Installs driver_transmit to appropriate ESP-NETIF object, so that outgoing packets from network stack are passed to the IO driver
- Calls `esp_netif_receive()` to pass incoming data to network stack

C) ESP-NETIF, former tcpip_adapter ESP-NETIF is an intermediary between an IO driver and a network stack, connecting packet data path between these two. As that it provides a set of interfaces for attaching a driver to ESP-NETIF object (runtime) and configuring a network stack (compile time). In addition to that a set of API is provided to control network interface lifecycle and its TCP/IP properties. As an overview, the ESP-NETIF public interface could be divided into these 6 groups:

1) Initialization APIs (to create and configure ESP-NETIF instance)
2) Input/Output API (for passing data between IO driver and network stack)
3) Event or Action API

- Used for network interface lifecycle management
- ESP-NETIF provides building blocks for designing event handlers

4) Setters and Getters for basic network interface properties
5) Network stack abstraction: enabling user interaction with TCP/IP stack

- Set interface up or down
- DHCP server and client API
- DNS API

6) Driver conversion utilities

D) Network stack Network stack has no public interaction with application code with regard to public interfaces and shall be fully abstracted by ESP-NETIF API.

ESP-NETIF programmer's manual Please refer to the example section for basic initialization of default interfaces:

- WiFi Station: wifi/getting_started/station/main/station_example_main.c
- WiFi Access Point: wifi/getting_started/softAP/main/softap_example_main.c
- Ethernet: ethernet/basic/main/ethernet_example_main.c

For more specific cases please consult this guide: *ESP-NETIF Custom I/O Driver*.

WiFi default initialization The initialization code as well as registering event handlers for default interfaces, such as softAP and station, are provided in two separate APIs to facilitate simple startup code for most applications:

- *esp_netif_create_default_wifi_ap()*
- *esp_netif_create_default_wifi_sta()*

Please note that these functions return the `esp_netif` handle, i.e. a pointer to a network interface object allocated and configured with default settings, which as a consequence, means that:

- The created object has to be destroyed if a network de-initialization is provided by an application.
- These *default* interfaces must not be created multiple times, unless the created handle is deleted using *esp_netif_destroy()*.
- When using Wifi in `AP+STA` mode, both these interfaces has to be created.

API Reference

Header File

- esp_netif/include/esp_netif.h

Functions

esp_err_t **esp_netif_init**(void)

Initialize the underlying TCP/IP stack.

Return
- ESP_OK on success
- ESP_FAIL if initializing failed

Note This function should be called exactly once from application code, when the application starts up.

esp_err_t **esp_netif_deinit**(void)

Deinitialize the esp-netif component (and the underlying TCP/IP stack)

Note: Deinitialization is not supported yet

Return
- ESP_ERR_INVALID_STATE if esp_netif not initialized
- ESP_ERR_NOT_SUPPORTED otherwise

esp_netif_t ***esp_netif_new**(**const** esp_netif_config_t **esp_netif_config*)

Creates an instance of new esp-netif object based on provided config.

Return
- pointer to esp-netif object on success
- NULL otherwise

Parameters
- `[in] esp_netif_config`: pointer esp-netif configuration

void **esp_netif_destroy**(esp_netif_t **esp_netif*)

Destroys the esp_netif object.

Parameters
- `[in] esp_netif`: pointer to the object to be deleted

esp_err_t **esp_netif_set_driver_config**(esp_netif_t **esp_netif*, **const** esp_netif_driver_ifconfig_t **driver_config*)

Configures driver related options of esp_netif object.

Return
- ESP_OK on success
- ESP_ERR_ESP_NETIF_INVALID_PARAMS if invalid parameters provided

Parameters
- `[inout] esp_netif`: pointer to the object to be configured
- `[in] driver_config`: pointer esp-netif io driver related configuration

esp_err_t **esp_netif_attach**(esp_netif_t **esp_netif*, esp_netif_iodriver_handle *driver_handle*)
Attaches esp_netif instance to the io driver handle.

Calling this function enables connecting specific esp_netif object with already initialized io driver to update esp_netif object with driver specific configuration (i.e. calls post_attach callback, which typically sets io driver callbacks to esp_netif instance and starts the driver)

Return
- ESP_OK on success
- ESP_ERR_ESP_NETIF_DRIVER_ATTACH_FAILED if driver' s pot_attach callback failed

Parameters
- `[inout] esp_netif`: pointer to esp_netif object to be attached
- `[in] driver_handle`: pointer to the driver handle

esp_err_t **esp_netif_receive**(esp_netif_t **esp_netif*, void **buffer*, size_t *len*, void **eb*)
Passes the raw packets from communication media to the appropriate TCP/IP stack.

This function is called from the configured (peripheral) driver layer. The data are then forwarded as frames to the TCP/IP stack.

Return
- ESP_OK

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[in] buffer`: Received data
- `[in] len`: Length of the data frame
- `[in] eb`: Pointer to internal buffer (used in Wi-Fi driver)

void **esp_netif_action_start**(void **esp_netif*, *esp_event_base_t* *base*, int32_t *event_id*, void **data*)
Default building block for network interface action upon IO driver start event Creates network interface, if AUTOUP enabled turns the interface on, if DHCPS enabled starts dhcp server.

Note This API can be directly used as event handler

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `base`:
- `event_id`:
- `data`:

void **esp_netif_action_stop**(void **esp_netif*, *esp_event_base_t* *base*, int32_t *event_id*, void **data*)
Default building block for network interface action upon IO driver stop event.

Note This API can be directly used as event handler

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `base`:
- `event_id`:
- `data`:

void **esp_netif_action_connected**(void **esp_netif*, *esp_event_base_t* *base*, int32_t *event_id*, void **data*)
Default building block for network interface action upon IO driver connected event.

Note This API can be directly used as event handler

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `base`:
- `event_id`:
- `data`:

void **esp_netif_action_disconnected**(void **esp_netif*, *esp_event_base_t* *base*, int32_t *event_id*, void **data*)
Default building block for network interface action upon IO driver disconnected event.

Note This API can be directly used as event handler

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `base`:
- `event_id`:
- `data`:

void **esp_netif_action_got_ip**(void **esp_netif*, *esp_event_base_t* *base*, int32_t *event_id*, void **data*)

Default building block for network interface action upon network got IP event.

Note This API can be directly used as event handler

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `base`:
- `event_id`:
- `data`:

esp_err_t **esp_netif_set_mac**(esp_netif_t **esp_netif*, uint8_t *mac*[])

Set the mac address for the interface instance.

Return
- ESP_OK - success
- ESP_ERR_ESP_NETIF_IF_NOT_READY - interface status error
- ESP_ERR_NOT_SUPPORTED - mac not supported on this interface

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[in] mac`: Desired mac address for the related network interface

esp_err_t **esp_netif_get_mac**(esp_netif_t **esp_netif*, uint8_t *mac*[])

Get the mac address for the interface instance.

Return
- ESP_OK - success
- ESP_ERR_ESP_NETIF_IF_NOT_READY - interface status error
- ESP_ERR_NOT_SUPPORTED - mac not supported on this interface

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] mac`: Resultant mac address for the related network interface

esp_err_t **esp_netif_set_hostname**(esp_netif_t **esp_netif*, **const** char **hostname*)

Set the hostname of an interface.

The configured hostname overrides the default configuration value CONFIG_LWIP_LOCAL_HOSTNAME. Please note that when the hostname is altered after interface started/connected the changes would only be reflected once the interface restarts/reconnects

Return
- ESP_OK - success
- ESP_ERR_ESP_NETIF_IF_NOT_READY - interface status error
- ESP_ERR_ESP_NETIF_INVALID_PARAMS - parameter error

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[in] hostname`: New hostname for the interface. Maximum length 32 bytes.

esp_err_t **esp_netif_get_hostname**(esp_netif_t **esp_netif*, **const** char ***hostname*)

Get interface hostname.

Return
- ESP_OK - success
- ESP_ERR_ESP_NETIF_IF_NOT_READY - interface status error
- ESP_ERR_ESP_NETIF_INVALID_PARAMS - parameter error

Parameters
- `[in] esp_netif`: Handle to esp-netif instance

- `[out] hostname`: Returns a pointer to the hostname. May be NULL if no hostname is set. If set non-NULL, pointer remains valid (and string may change if the hostname changes).

bool **esp_netif_is_netif_up**(esp_netif_t *esp_netif)

Test if supplied interface is up or down.

Return
- true - Interface is up
- false - Interface is down

Parameters
- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_get_ip_info**(esp_netif_t *esp_netif, esp_netif_ip_info_t *ip_info)

Get interface's IP address information.

If the interface is up, IP information is read directly from the TCP/IP stack. If the interface is down, IP information is read from a copy kept in the ESP-NETIF instance

Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] ip_info`: If successful, IP information will be returned in this argument.

esp_err_t **esp_netif_get_old_ip_info**(esp_netif_t *esp_netif, esp_netif_ip_info_t *ip_info)

Get interface's old IP information.

Returns an "old" IP address previously stored for the interface when the valid IP changed.

If the IP lost timer has expired (meaning the interface was down for longer than the configured interval) then the old IP information will be zero.

Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] ip_info`: If successful, IP information will be returned in this argument.

esp_err_t **esp_netif_set_ip_info**(esp_netif_t *esp_netif, **const** esp_netif_ip_info_t *ip_info)

Set interface's IP address information.

This function is mainly used to set a static IP on an interface.

If the interface is up, the new IP information is set directly in the TCP/IP stack.

The copy of IP information kept in the ESP-NETIF instance is also updated (this copy is returned if the IP is queried while the interface is still down.)

Note DHCP client/server must be stopped (if enabled for this interface) before setting new IP information.

Note Calling this interface for may generate a SYSTEM_EVENT_STA_GOT_IP or SYSTEM_EVENT_ETH_GOT_IP event.

Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_NOT_STOPPED If DHCP server or client is still running

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[in] ip_info`: IP information to set on the specified interface

esp_err_t **esp_netif_set_old_ip_info**(esp_netif_t *esp_netif, **const** esp_netif_ip_info_t *ip_info)

Set interface old IP information.

This function is called from the DHCP client (if enabled), before a new IP is set. It is also called from the default handlers for the SYSTEM_EVENT_STA_CONNECTED and SYSTEM_EVENT_ETH_CONNECTED events.

Calling this function stores the previously configured IP, which can be used to determine if the IP changes in the future.

If the interface is disconnected or down for too long, the "IP lost timer" will expire (after the configured interval) and set the old IP information to zero.

Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[in] ip_info`: Store the old IP information for the specified interface

int **esp_netif_get_netif_impl_index**(esp_netif_t *_esp_netif_)

Get net interface index from network stack implementation.

Note This index could be used in `setsockopt()` to bind socket with multicast interface
Return implementation specific index of interface represented with supplied esp_netif
Parameters
- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_get_netif_impl_name**(esp_netif_t *_esp_netif_, char *_name_)

Get net interface name from network stack implementation.

Note This name could be used in `setsockopt()` to bind socket with appropriate interface
Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] name`: Interface name as specified in underlying TCP/IP stack. Note that the actual name will be copied to the specified buffer, which must be allocated to hold maximum interface name size (6 characters for lwIP)

esp_err_t **esp_netif_dhcps_option**(esp_netif_t *_esp_netif_, esp_netif_dhcp_option_mode_t _opt_op_, esp_netif_dhcp_option_id_t _opt_id_, void *_opt_val_, uint32_t _opt_len_)

Set or Get DHCP server option.

Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STOPPED
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STARTED

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[in] opt_op`: ESP_NETIF_OP_SET to set an option, ESP_NETIF_OP_GET to get an option.
- `[in] opt_id`: Option index to get or set, must be one of the supported enum values.
- `[inout] opt_val`: Pointer to the option parameter.
- `[in] opt_len`: Length of the option parameter.

esp_err_t **esp_netif_dhcpc_option**(esp_netif_t *_esp_netif_, esp_netif_dhcp_option_mode_t _opt_op_, esp_netif_dhcp_option_id_t _opt_id_, void *_opt_val_, uint32_t _opt_len_)

Set or Get DHCP client option.

Return
- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STOPPED

- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STARTED

Parameters

- `[in] esp_netif`: Handle to esp-netif instance
- `[in] opt_op`: ESP_NETIF_OP_SET to set an option, ESP_NETIF_OP_GET to get an option.
- `[in] opt_id`: Option index to get or set, must be one of the supported enum values.
- `[inout] opt_val`: Pointer to the option parameter.
- `[in] opt_len`: Length of the option parameter.

esp_err_t **esp_netif_dhcpc_start**(esp_netif_t **esp_netif*)

Start DHCP client (only if enabled in interface object)

Note The default event handlers for the SYSTEM_EVENT_STA_CONNECTED and SYSTEM_EVENT_ETH_CONNECTED events call this function.

Return

- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STARTED
- ESP_ERR_ESP_NETIF_DHCPC_START_FAILED

Parameters

- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_dhcpc_stop**(esp_netif_t **esp_netif*)

Stop DHCP client (only if enabled in interface object)

Note Calling action_netif_stop() will also stop the DHCP Client if it is running.

Return

- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STOPPED
- ESP_ERR_ESP_NETIF_IF_NOT_READY

Parameters

- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_dhcpc_get_status**(esp_netif_t **esp_netif*, esp_netif_dhcp_status_t **status*)

Get DHCP client status.

Return

- ESP_OK

Parameters

- `[in] esp_netif`: Handle to esp-netif instance
- `[out] status`: If successful, the status of DHCP client will be returned in this argument.

esp_err_t **esp_netif_dhcps_get_status**(esp_netif_t **esp_netif*, esp_netif_dhcp_status_t **status*)

Get DHCP Server status.

Return

- ESP_OK

Parameters

- `[in] esp_netif`: Handle to esp-netif instance
- `[out] status`: If successful, the status of the DHCP server will be returned in this argument.

esp_err_t **esp_netif_dhcps_start**(esp_netif_t **esp_netif*)

Start DHCP server (only if enabled in interface object)

Return

- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STARTED

Parameters

- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_dhcps_stop**(esp_netif_t **esp_netif*)

Stop DHCP server (only if enabled in interface object)

Return

- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS
- ESP_ERR_ESP_NETIF_DHCP_ALREADY_STOPPED
- ESP_ERR_ESP_NETIF_IF_NOT_READY

Parameters

- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_set_dns_info**(esp_netif_t **esp_netif*, esp_netif_dns_type_t *type*, esp_netif_dns_info_t **dns*)

Set DNS Server information.

This function behaves differently if DHCP server or client is enabled

If DHCP client is enabled, main and backup DNS servers will be updated automatically from the DHCP lease if the relevant DHCP options are set. Fallback DNS Server is never updated from the DHCP lease and is designed to be set via this API. If DHCP client is disabled, all DNS server types can be set via this API only.

If DHCP server is enabled, the Main DNS Server setting is used by the DHCP server to provide a DNS Server option to DHCP clients (Wi-Fi stations).

- The default Main DNS server is typically the IP of the Wi-Fi AP interface itself.
- This function can override it by setting server type ESP_NETIF_DNS_MAIN.
- Other DNS Server types are not supported for the Wi-Fi AP interface.

Return

- ESP_OK on success
- ESP_ERR_ESP_NETIF_INVALID_PARAMS invalid params

Parameters

- `[in] esp_netif`: Handle to esp-netif instance
- `[in] type`: Type of DNS Server to set: ESP_NETIF_DNS_MAIN, ESP_NETIF_DNS_BACKUP, ESP_NETIF_DNS_FALLBACK
- `[in] dns`: DNS Server address to set

esp_err_t **esp_netif_get_dns_info**(esp_netif_t **esp_netif*, esp_netif_dns_type_t *type*, esp_netif_dns_info_t **dns*)

Get DNS Server information.

Return the currently configured DNS Server address for the specified interface and Server type.

This may be result of a previous call to *esp_netif_set_dns_info()*. If the interface's DHCP client is enabled, the Main or Backup DNS Server may be set by the current DHCP lease.

Return

- ESP_OK on success
- ESP_ERR_ESP_NETIF_INVALID_PARAMS invalid params

Parameters

- `[in] esp_netif`: Handle to esp-netif instance
- `[in] type`: Type of DNS Server to get: ESP_NETIF_DNS_MAIN, ESP_NETIF_DNS_BACKUP, ESP_NETIF_DNS_FALLBACK
- `[out] dns`: DNS Server result is written here on success

esp_err_t **esp_netif_create_ip6_linklocal**(esp_netif_t **esp_netif*)

Create interface link-local IPv6 address.

Cause the TCP/IP stack to create a link-local IPv6 address for the specified interface.

This function also registers a callback for the specified interface, so that if the link-local address becomes verified as the preferred address then a SYSTEM_EVENT_GOT_IP6 event will be sent.

Return

- ESP_OK
- ESP_ERR_ESP_NETIF_INVALID_PARAMS

Parameters

- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_get_ip6_linklocal** (esp_netif_t **esp_netif*, esp_ip6_addr_t **if_ip6*)
Get interface link-local IPv6 address.

If the specified interface is up and a preferred link-local IPv6 address has been created for the interface, return a copy of it.

Return
- ESP_OK
- ESP_FAIL If interface is down, does not have a link-local IPv6 address, or the link-local IPv6 address is not a preferred address.

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] if_ip6`: IPv6 information will be returned in this argument if successful.

esp_err_t **esp_netif_get_ip6_global** (esp_netif_t **esp_netif*, esp_ip6_addr_t **if_ip6*)
Get interface global IPv6 address.

If the specified interface is up and a preferred global IPv6 address has been created for the interface, return a copy of it.

Return
- ESP_OK
- ESP_FAIL If interface is down, does not have a global IPv6 address, or the global IPv6 address is not a preferred address.

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] if_ip6`: IPv6 information will be returned in this argument if successful.

int **esp_netif_get_all_ip6** (esp_netif_t **esp_netif*, esp_ip6_addr_t *if_ip6*[])
Get all IPv6 addresses of the specified interface.

Return number of returned IPv6 addresses

Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `[out] if_ip6`: Array of IPv6 addresses will be copied to the argument

void **esp_netif_set_ip4_addr** (esp_ip4_addr_t **addr*, uint8_t *a*, uint8_t *b*, uint8_t *c*, uint8_t *d*)
Sets IPv4 address to the specified octets.

Parameters
- `[out] addr`: IP address to be set
- `a`: the first octet (127 for IP 127.0.0.1)
- `b`:
- `c`:
- `d`:

char ***esp_ip4addr_ntoa** (**const** esp_ip4_addr_t **addr*, char **buf*, int *buflen*)
Converts numeric IP address into decimal dotted ASCII representation.

Return either pointer to buf which now holds the ASCII representation of addr or NULL if buf was too small

Parameters
- `addr`: ip address in network order to convert
- `buf`: target buffer where the string is stored
- `buflen`: length of buf

uint32_t **esp_ip4addr_aton** (**const** char **addr*)
Ascii internet address interpretation routine The value returned is in network order.

Return ip address in network order

Parameters
- `addr`: IP address in ascii representation (e.g. “127.0.0.1”)

esp_netif_iodriver_handle **esp_netif_get_io_driver** (esp_netif_t **esp_netif*)
Gets media driver handle for this esp-netif instance.

Return opaque pointer of related IO driver

Parameters
- `[in] esp_netif`: Handle to esp-netif instance

esp_netif_t ***esp_netif_get_handle_from_ifkey**(**const** char **if_key*)
Searches over a list of created objects to find an instance with supplied if key.

Return Handle to esp-netif instance
Parameters
- `if_key`: Textual description of network interface

esp_netif_flags_t **esp_netif_get_flags**(esp_netif_t **esp_netif*)
Returns configured flags for this interface.

Return Configuration flags
Parameters
- `[in] esp_netif`: Handle to esp-netif instance

const char ***esp_netif_get_ifkey**(esp_netif_t **esp_netif*)
Returns configured interface key for this esp-netif instance.

Return Textual description of related interface
Parameters
- `[in] esp_netif`: Handle to esp-netif instance

const char ***esp_netif_get_desc**(esp_netif_t **esp_netif*)
Returns configured interface type for this esp-netif instance.

Return Enumerated type of this interface, such as station, AP, ethernet
Parameters
- `[in] esp_netif`: Handle to esp-netif instance

int **esp_netif_get_route_prio**(esp_netif_t **esp_netif*)
Returns configured routing priority number.

Return Integer representing the instance' s route-prio, or -1 if invalid paramters
Parameters
- `[in] esp_netif`: Handle to esp-netif instance

int32_t **esp_netif_get_event_id**(esp_netif_t **esp_netif*, esp_netif_ip_event_type_t *event_type*)
Returns configured event for this esp-netif instance and supplied event type.

Return specific event id which is configured to be raised if the interface lost or acquired IP address -1 if supplied event_type is not known
Parameters
- `[in] esp_netif`: Handle to esp-netif instance
- `event_type`: (either get or lost IP)

esp_netif_t ***esp_netif_next**(esp_netif_t **esp_netif*)
Iterates over list of interfaces. Returns first netif if NULL given as parameter.

Return First netif from the list if supplied parameter is NULL, next one otherwise
Parameters
- `[in] esp_netif`: Handle to esp-netif instance

size_t **esp_netif_get_nr_of_ifs**(void)
Returns number of registered esp_netif objects.

Return Number of esp_netifs

void **esp_netif_netstack_buf_ref**(void **netstack_buf*)
increase the reference counter of net stack buffer

Parameters
- `[in] netstack_buf`: the net stack buffer

void **esp_netif_netstack_buf_free**(void **netstack_buf*)
free the netstack buffer

Parameters

- `[in] netstack_buf`: the net stack buffer

Macros

_ESP_NETIF_SUPPRESS_LEGACY_WARNING_

WiFi default API reference

Header File

- esp_wifi/include/esp_wifi_default.h

Functions

esp_err_t **esp_netif_attach_wifi_station**(esp_netif_t **esp_netif*)

Attaches wifi station interface to supplied netif.

Return

- ESP_OK on success
- ESP_FAIL if attach failed

Parameters

- `esp_netif`: instance to attach the wifi station to

esp_err_t **esp_netif_attach_wifi_ap**(esp_netif_t **esp_netif*)

Attaches wifi soft AP interface to supplied netif.

Return

- ESP_OK on success
- ESP_FAIL if attach failed

Parameters

- `esp_netif`: instance to attach the wifi AP to

esp_err_t **esp_wifi_set_default_wifi_sta_handlers**(void)

Sets default wifi event handlers for STA interface.

Return

- ESP_OK on success, error returned from esp_event_handler_register if failed

esp_err_t **esp_wifi_set_default_wifi_ap_handlers**(void)

Sets default wifi event handlers for STA interface.

Return

- ESP_OK on success, error returned from esp_event_handler_register if failed

esp_err_t **esp_wifi_clear_default_wifi_driver_and_handlers**(void **esp_netif*)

Clears default wifi event handlers for supplied network interface.

Return

- ESP_OK on success, error returned from esp_event_handler_register if failed

Parameters

- `esp_netif`: instance of corresponding if object

esp_netif_t ***esp_netif_create_default_wifi_ap**(void)

Creates default WIFI AP. In case of any init error this API aborts.

Return pointer to esp-netif instance

esp_netif_t ***esp_netif_create_default_wifi_sta**(void)

Creates default WIFI STA. In case of any init error this API aborts.

Return pointer to esp-netif instance

esp_netif_t ***esp_netif_create_wifi**(*wifi_interface_t* *wifi_if*, esp_netif_inherent_config_t **esp_netif_config*)

Creates esp_netif WiFi object based on the custom configuration.

Attention This API DOES NOT register default handlers!
Return pointer to esp-netif instance
Parameters
- `[in] wifi_if`: type of wifi interface
- `[in] esp_netif_config`: inherent esp-netif configuration pointer

esp_err_t **esp_netif_create_default_wifi_mesh_netifs**(esp_netif_t **p_netif_sta, esp_netif_t **p_netif_ap)

Creates default STA and AP network interfaces for esp-mesh.

Both netifs are almost identical to the default station and softAP, but with DHCP client and server disabled. Please note that the DHCP client is typically enabled only if the device is promoted to a root node.

Returns created interfaces which could be ignored setting parameters to NULL if an application code does not need to save the interface instances for further processing.

Return ESP_OK on success
Parameters
- `[out] p_netif_sta`: pointer where the resultant STA interface is saved (if non NULL)
- `[out] p_netif_ap`: pointer where the resultant AP interface is saved (if non NULL)

TCP/IP 适配器迁移指南

TCP/IP 适配器是在 IDF V4.1 之前使用的网络接口抽象组件。本文档概述了从 tcpip_adapter 移出至其后继者*ESP-NETIF* 的过程。

更新网络连接代码

网络软件栈初始化 只需将 `tcpip_adapter_init()` 替换为 `esp_netif_init()`。请注意，*ESP-NETIF* 初始化 API 可返回标准错误代码，还可以使用 `esp_netif_deinit()` 进行去初始化。

此外，还需将 `#include "tcpip_adapter.h"` 替换为 `#include "esp_netif.h"`。

创建网络接口 TCP/IP 适配器静态定义了三个接口：

- Wi-Fi Station
- Wi-Fi AP
- 以太网

网络接口的设计应严格参考*ESP-NETIF*，以使其能够连接到 TCP/IP 软件栈。例如，在 TCP/IP 软件栈和事件循环初始化完成后，Wi-Fi 的初始化代码必须显示调用 `esp_netif_create_default_wifi_sta();` 或 `esp_netif_create_default_wifi_ap();` 。请参阅这三个接口的初始化代码示例：

- Wi-Fi Station: wifi/getting_started/station/main/station_example_main.c
- Wi-Fi AP: wifi/getting_started/softAP/main/softap_example_main.c
- 以太网: ethernet/basic/main/ethernet_example_main.c

更换其他 tcpip_adapter API 所有 tcpip_adapter 函数都有对应的 esp-netif。具体请见 esp_netif 的内容：

- Setters/Getters
- DHCP
- DNS
- IP address

默认事件处理程序 事件处理程序已经从 tcpip_adapter 移动到相应的驱动程序代码。从应用程序的角度来看，这不会带来任何影响，所有事件仍以相同的方式处理。请注意，在与 IP 相关的事件处理程序中，应用程序代码通常以 esp-netif 结构体的形式接收 IP 地址（不是 LwIP 结构，但兼容二进制格式）。这是打印地址的首选方式：

```
ESP_LOGI(TAG, "got ip:" IPSTR "\n", IP2STR(&event->ip_info.ip));
```

而不是

```
ESP_LOGI(TAG, "got ip:%s\n", ip4addr_ntoa(&event->ip_info.ip));
```

由于 `ip4addr_ntoa()` 为 LwIP API，因此 esp-netif 还提供了替代函数 `esp_ip4addr_ntoa()`，但整体而言仍推荐上述方法。

IP 地址 推荐使用 esp-netif 定义的 IP 结构。请注意，在启用默认兼容性时，LwIP 结构体仍然可以工作。
* esp-netif IP address definitions

下一步 为了移植应用程序使其可以使用*ESP-NETIF* 还需完成的步骤包括：在组件配置中禁用 tcpip_adapter 兼容层。方法为：`ESP NETIF Adapter -> Enable backward compatible tcpip_adapter interface`，并检查工程是否编译成功。TCP/IP 适配器涉及大量依赖项，这一步可能有助于将应用程序与使用特定 TCP/IP 软件栈的 API 分离开来。

ESP-NETIF Custom I/O Driver

This section outlines implementing a new I/O driver with esp-netif connection capabilities. By convention the I/O driver has to register itself as an esp-netif driver and thus holds a dependency on esp-netif component and is responsible for providing data path functions, post-attach callback and in most cases also default event handlers to define network interface actions based on driver's lifecycle transitions.

Packet input/output As shown in the diagram, the following three API functions for the packet data path must be defined for connecting with esp-netif:

- *esp_netif_transmit()*
- *esp_netif_free_rx_buffer()*
- *esp_netif_receive()*

The first two functions for transmitting and freeing the rx buffer are provided as callbacks, i.e. they get called from esp-netif (and its underlying TCP/IP stack) and I/O driver provides their implementation.

The receiving function on the other hand gets called from the I/O driver, so that the driver's code simply calls *esp_netif_receive()* on a new data received event.

Post attach callback A final part of the network interface initialization consists of attaching the esp-netif instance to the I/O driver, by means of calling the following API:

```
esp_err_t esp_netif_attach(esp_netif_t *esp_netif, esp_netif_iodriver_handle␣
↪driver_handle);
```

It is assumed that the `esp_netif_iodriver_handle` is a pointer to driver's object, a struct derived from `struct esp_netif_driver_base_s`, so that the first member of I/O driver structure must be this base structure with pointers to

- post-attach function callback
- related esp-netif instance

As a consequence the I/O driver has to create an instance of the struct per below:

```
typedef struct my_netif_driver_s {
        esp_netif_driver_base_t base;           /*!< base structure reserved as␣
↪esp-netif driver */
        driver_impl             *h;             /*!< handle of driver␣
↪implementation */
    } my_netif_driver_t;
```

with actual values of `my_netif_driver_t::base.post_attach` and the actual drivers handle `my_netif_driver_t::h`. So when the *esp_netif_attach()* gets called from the initialization code, the post-attach callback from I/O driver's code gets executed to mutually register callbacks between esp-netif and I/O driver instances. Typically the driver is started as well in the post-attach callback. An example of a simple post-attach callback is outlined below:

```
static esp_err_t my_post_attach_start(esp_netif_t * esp_netif, void * args)
{
    my_netif_driver_t *driver = args;
    const esp_netif_driver_ifconfig_t driver_ifconfig = {
            .driver_free_rx_buffer = my_free_rx_buf,
            .transmit = my_transmit,
            .handle = driver->driver_impl
    };
    driver->base.netif = esp_netif;
    ESP_ERROR_CHECK(esp_netif_set_driver_config(esp_netif, &driver_ifconfig));
    my_driver_start(driver->driver_impl);
    return ESP_OK;
}
```

Default handlers I/O drivers also typically provide default definitions of lifecycle behaviour of related network interfaces based on state transitions of I/O drivers. For example *driver start -> network start*, etc. An example of such a default handler is provided below:

```
esp_err_t my_driver_netif_set_default_handlers(my_netif_driver_t *driver, esp_
↪netif_t * esp_netif)
{
    driver_set_event_handler(driver->driver_impl, esp_netif_action_start, MY_DRV_
↪EVENT_START, esp_netif);
    driver_set_event_handler(driver->driver_impl, esp_netif_action_stop, MY_DRV_
↪EVENT_STOP, esp_netif);
    return ESP_OK;
}
```

Network stack connection The packet data path functions for transmitting and freeing the rx buffer (defined in the I/O driver) are called from the esp-netif, specifically from its TCP/IP stack connecting layer. The following API reference outlines these network stack interaction with the esp-netif.

Header File

- esp_netif/include/esp_netif_net_stack.h

Functions

esp_netif_t ***esp_netif_get_handle_from_netif_impl**(void *dev*)

Returns esp-netif handle.

Return handle to related esp-netif instance

Parameters

- `[in] dev`: opaque ptr to network interface of specific TCP/IP stack

void ***esp_netif_get_netif_impl**(esp_netif_t *esp_netif*)

Returns network stack specific implementation handle (if supported)

Note that it is not supported to acquire PPP netif impl pointer and this function will return NULL for esp_netif instances configured to PPP mode

Return handle to related network stack netif handle

Parameters

- `[in] esp_netif`: Handle to esp-netif instance

esp_err_t **esp_netif_transmit**(esp_netif_t *esp_netif, void *data, size_t len)
Outputs packets from the TCP/IP stack to the media to be transmitted.

This function gets called from network stack to output packets to IO driver.

Return ESP_OK on success, an error passed from the I/O driver otherwise
Parameters
- [in] esp_netif: Handle to esp-netif instance
- [in] data: Data to be transmitted
- [in] len: Length of the data frame

esp_err_t **esp_netif_transmit_wrap**(esp_netif_t *esp_netif, void *data, size_t len, void *netstack_buf)
Outputs packets from the TCP/IP stack to the media to be transmitted.

This function gets called from network stack to output packets to IO driver.

Return ESP_OK on success, an error passed from the I/O driver otherwise
Parameters
- [in] esp_netif: Handle to esp-netif instance
- [in] data: Data to be transmitted
- [in] len: Length of the data frame
- [in] netstack_buf: net stack buffer

void **esp_netif_free_rx_buffer**(void *esp_netif, void *buffer)
Free the rx buffer allocated by the media driver.

This function gets called from network stack when the rx buffer to be freed in IO driver context, i.e. to deallocate a buffer owned by io driver (when data packets were passed to higher levels to avoid copying)

Parameters
- [in] esp_netif: Handle to esp-netif instance
- [in] buffer: Rx buffer pointer

TCP/IP 套接字 API 的示例代码存放在 ESP-IDF 示例项目的 protocols/sockets 目录下。

2.1.4 应用层协议

应用层网络协议（IP 网络层协议之上）的相关文档存放在应用层协议。

2.2 外设 API

2.2.1 Analog to Digital Converter

Overview

The ESP32-S2 integrates two 13-bit SAR (Successive Approximation Register) ADCs, supporting a total of 20 measurement channels (analog enabled pins).

The ADC driver API supports ADC1 (10 channels, attached to GPIOs 1 - 10), and ADC2 (10 channels, attached to GPIOs 11 - 20). However, the usage of ADC2 has some restrictions for the application:

1. Different from ADC1, the hardware arbiter function is added to ADC2, so when using the API of ADC2 to obtain the sampling voltage, you need to check whether the reading is successful.

Configuration and Reading ADC

Each ADC unit supports two work modes, ADC-RTC or ADC-DMA mode. ADC-RTC is controlled by the RTC controller and is suitable for low-frequency sampling operations. ADC-DMA is controlled by a digital controller and is suitable for high-frequency continuous sampling actions.

ADC-RTC mode The ADC should be configured before reading is taken.

- For ADC1, configure desired precision and attenuation by calling functions *adc1_config_width()* and *adc1_config_channel_atten()*.
- For ADC2, configure the attenuation by *adc2_config_channel_atten()*. The reading width of ADC2 is configured every time you take the reading.

Attenuation configuration is done per channel, see *adc1_channel_t* and *adc2_channel_t*, set as a parameter of above functions.

Then it is possible to read ADC conversion result with *adc1_get_raw()* and *adc2_get_raw()*. Reading width of ADC2 should be set as a parameter of *adc2_get_raw()* instead of in the configuration functions.

注解: Since the ADC2 is shared with the WIFI module, which has higher priority, reading operation of *adc2_get_raw()* will fail between *esp_wifi_start()* and *esp_wifi_stop()*. Use the return code to see whether the reading is successful.

This API provides convenient way to configure ADC1 for reading from *ULP*. To do so, call function *adc1_ulp_enable()* and then set precision and attenuation as discussed above.

There is another specific function *adc_vref_to_gpio()* used to route internal reference voltage to a GPIO pin. It comes handy to calibrate ADC reading and this is discussed in section *Minimizing Noise*.

Application Examples

Reading voltage on ADC1 channel 0 (GPIO 0):

```
#include <driver/adc.h>

...

    adc1_config_width(ADC_WIDTH_BIT_12);
    adc1_config_channel_atten(ADC1_CHANNEL_0,ADC_ATTEN_DB_0);
    int val = adc1_get_raw(ADC1_CHANNEL_0);
```

The input voltage in the above example is from 0 to 1.1 V (0 dB attenuation). The input range can be extended by setting a higher attenuation, see *adc_atten_t*. An example of using the ADC driver including calibration (discussed below) is available at esp-idf: peripherals/adc

Reading voltage on ADC2 channel 7 (GPIO 0):

```
#include <driver/adc.h>

...

    int read_raw;
    adc2_config_channel_atten( ADC2_CHANNEL_7, ADC_ATTEN_0db );

    esp_err_t r = adc2_get_raw( ADC2_CHANNEL_7, ADC_WIDTH_12Bit, &read_raw);
    if ( r == ESP_OK ) {
        printf("%d\n", read_raw );
    } else if ( r == ESP_ERR_TIMEOUT ) {
        printf("ADC2 used by Wi-Fi.\n");
    }
```

The reading may fail due to collision with Wi-Fi, should check it. An example using the ADC2 driver to read the output of DAC is available in esp-idf: peripherals/adc2

The value read in both these examples is 13 bits wide (range 0-8191).

Minimizing Noise

The ESP32-S2 ADC can be sensitive to noise leading to large discrepancies in ADC readings. To minimize noise, users may connect a 0.1uF capacitor to the ADC input pad in use. Multisampling may also be used to further mitigate the effects of noise.

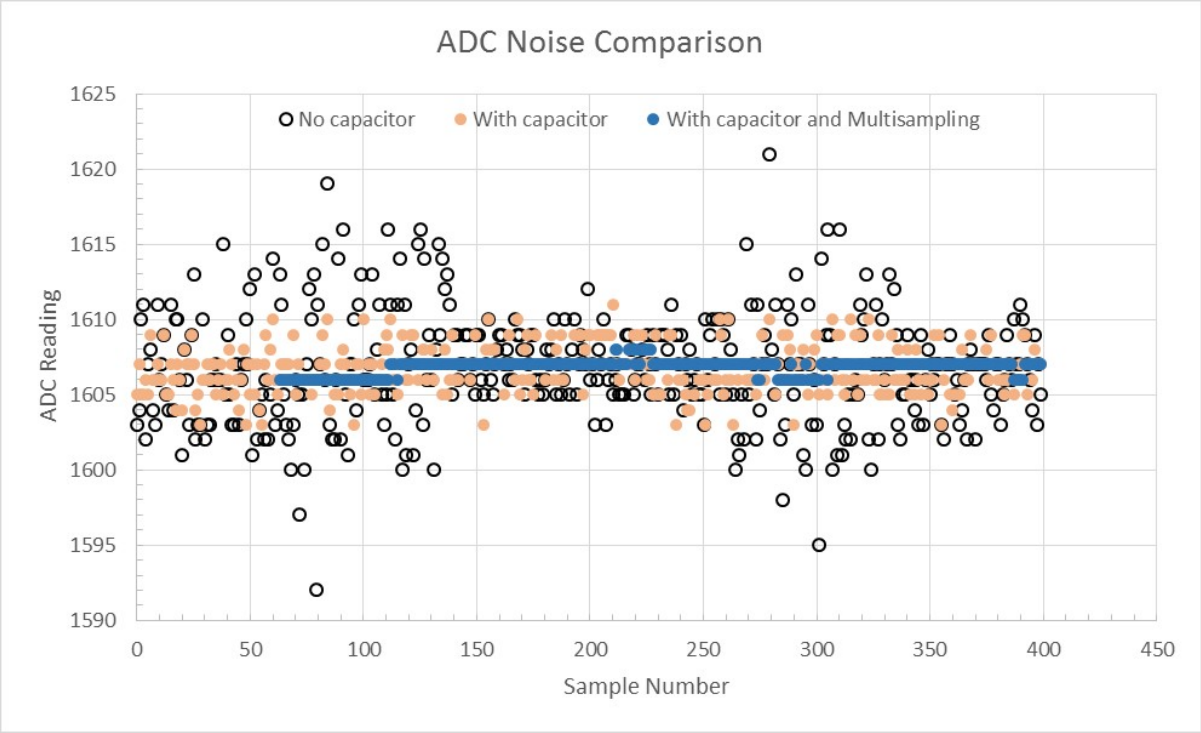


图 3: Graph illustrating noise mitigation using capacitor and multisampling of 64 samples.

ADC Calibration

The esp_adc_cal/include/esp_adc_cal.h API provides functions to correct for differences in measured voltages caused by variation of ADC reference voltages (Vref) between chips. Per design the ADC reference voltage is 1100 mV, however the true reference voltage can range from 1000 mV to 1200 mV amongst different ESP32-S2s.

Correcting ADC readings using this API involves characterizing one of the ADCs at a given attenuation to obtain a characteristics curve (ADC-Voltage curve) that takes into account the difference in ADC reference voltage. The characteristics curve is in the form of `y = coeff_a * x + coeff_b` and is used to convert ADC readings to voltages in mV. Calculation of the characteristics curve is based on calibration values which can be stored in eFuse or provided by the user.

Calibration Values Calibration values are used to generate characteristic curves that account for the variation of ADC reference voltage of a particular ESP32-S2 chip. There are currently three sources of calibration values on ESP32, and one source on ESP32-S2. The availability of these calibration values will depend on the type and production date of the ESP32-S2 chip/module.

- **eFuse Two Point** values calibrates the ADC output at two different voltages. This value is measured and burned into eFuse `BLOCK0` during factory calibration on newly manufactured ESP32-S2 chips and modules. If you would like to purchase chips or modules with calibration, double check with distributor or Espressif directly.

You can verify if **eFuse Two Point** is present by running the espefuse.py tool with `adc_info` parameter

```
$IDF_PATH/components/esptool_py/esptool/espefuse.py --port /dev/ttyUSB0 adc_info
```

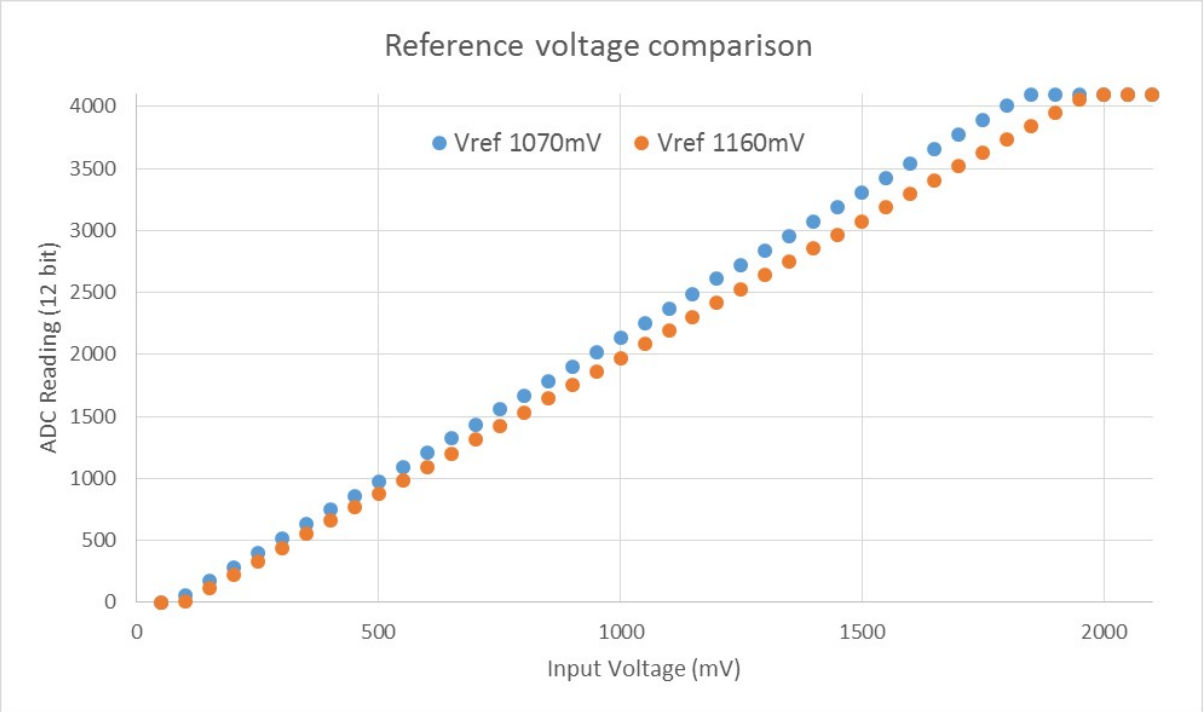


图 4: Graph illustrating effect of differing reference voltages on the ADC voltage curve.

Replace `/dev/ttyUSB0` with ESP32-S2 board’s port name.

Application Example For a full example see esp-idf: peripherals/adc

Characterizing an ADC at a particular attenuation:

```
#include "driver/adc.h"
#include "esp_adc_cal.h"

...

    //Characterize ADC at particular atten
    esp_adc_cal_characteristics_t *adc_chars = calloc(1, sizeof(esp_adc_cal_
↪characteristics_t));
    esp_adc_cal_value_t val_type = esp_adc_cal_characterize(unit, atten, ADC_WIDTH_
↪BIT_12, DEFAULT_VREF, adc_chars);
    //Check type of calibration value used to characterize ADC
    if (val_type == ESP_ADC_CAL_VAL_EFUSE_VREF) {
        printf("eFuse Vref");
    } else if (val_type == ESP_ADC_CAL_VAL_EFUSE_TP) {
        printf("Two Point");
    } else {
        printf("Default");
    }
```

Reading an ADC then converting the reading to a voltage:

```
#include "driver/adc.h"
#include "esp_adc_cal.h"

...
    uint32_t reading =  adc1_get_raw(ADC1_CHANNEL_5);
    uint32_t voltage = esp_adc_cal_raw_to_voltage(reading, adc_chars);
```

Routing ADC reference voltage to GPIO, so it can be manually measured (for **Default Vref**):

```
#include "driver/adc.h"

...

    esp_err_t status = adc_vref_to_gpio(ADC_UNIT_1, GPIO_NUM_25);
    if (status == ESP_OK) {
        printf("v_ref routed to GPIO\n");
    } else {
        printf("failed to route v_ref\n");
    }
```

GPIO Lookup Macros

There are macros available to specify the GPIO number of a ADC channel, or vice versa. e.g.

1. `ADC1_CHANNEL_0_GPIO_NUM` is the GPIO number of ADC1 channel 0 (36);
2. `ADC1_GPIO32_CHANNEL` is the ADC1 channel number of GPIO 32 (ADC1 channel 4).

API Reference

This reference covers three components:

- *ADC driver*
- *ADC Calibration*
- *GPIO Lookup Macros*

ADC driver

Header File

- driver/esp32s2/include/driver/adc.h

Functions

esp_err_t **adc_arbiter_config**(*adc_unit_t adc_unit*, *adc_arbiter_t *config*)

Config ADC module arbiter. The arbiter is to improve the use efficiency of ADC2. After the control right is robbed by the high priority, the low priority controller will read the invalid ADC2 data, and the validity of the data can be judged by the flag bit in the data.

Note Only ADC2 support arbiter.
Note Default priority: Wi-Fi > RTC > Digital;
Note In normal use, there is no need to call this interface to config arbiter.

Return

- ESP_OK Success
- ESP_ERR_NOT_SUPPORTED ADC unit not support arbiter.

Parameters

- `adc_unit`: ADC unit.
- `config`: Refer to *`adc_arbiter_t`*.

esp_err_t **adc_digi_start**(void)

Enable digital controller to trigger the measurement.

Return

- ESP_OK Success

esp_err_t **adc_digi_stop**(void)

Disable digital controller to trigger the measurement.

Return
- ESP_OK Success

esp_err_t **adc_digi_filter_reset**(*adc_digi_filter_idx_t idx*)
Reset adc digital controller filter.

Return
- ESP_OK Success

Parameters
- `idx`: Filter index.

esp_err_t **adc_digi_filter_set_config**(*adc_digi_filter_idx_t idx*, *adc_digi_filter_t *config*)
Set adc digital controller filter configuration.

Note For ESP32S2, Filter IDX0/IDX1 can only be used to filter all enabled channels of ADC1/ADC2 unit at the same time.

Return
- ESP_OK Success

Parameters
- `idx`: Filter index.
- `config`: See *adc_digi_filter_t*.

esp_err_t **adc_digi_filter_get_config**(*adc_digi_filter_idx_t idx*, *adc_digi_filter_t *config*)
Get adc digital controller filter configuration.

Note For ESP32S2, Filter IDX0/IDX1 can only be used to filter all enabled channels of ADC1/ADC2 unit at the same time.

Return
- ESP_OK Success

Parameters
- `idx`: Filter index.
- `config`: See *adc_digi_filter_t*.

esp_err_t **adc_digi_filter_enable**(*adc_digi_filter_idx_t idx*, bool *enable*)
Enable/disable adc digital controller filter. Filtering the ADC data to obtain smooth data at higher sampling rates.

Note For ESP32S2, Filter IDX0/IDX1 can only be used to filter all enabled channels of ADC1/ADC2 unit at the same time.

Return
- ESP_OK Success

Parameters
- `idx`: Filter index.
- `enable`: Enable/Disable filter.

esp_err_t **adc_digi_monitor_set_config**(*adc_digi_monitor_idx_t idx*, *adc_digi_monitor_t *config*)
Config monitor of adc digital controller.

Note For ESP32S2, The monitor will monitor all the enabled channel data of the each ADC unit at the same time.

Return
- ESP_OK Success

Parameters
- `idx`: Monitor index.
- `config`: See *adc_digi_monitor_t*.

esp_err_t **adc_digi_monitor_enable**(*adc_digi_monitor_idx_t idx*, bool *enable*)
Enable/disable monitor of adc digital controller.

Note For ESP32S2, The monitor will monitor all the enabled channel data of the each ADC unit at the same time.

Return
- ESP_OK Success

Parameters

- `idx`: Monitor index.
- `enable`: True or false enable monitor.

esp_err_t **adc_digi_intr_enable**(*adc_unit_t adc_unit*, *adc_digi_intr_t intr_mask*)

Enable interrupt of adc digital controller by bitmask.

Return

- ESP_OK Success

Parameters

- `adc_unit`: ADC unit.
- `intr_mask`: Interrupt bitmask. See `adc_digi_intr_t`.

esp_err_t **adc_digi_intr_disable**(*adc_unit_t adc_unit*, *adc_digi_intr_t intr_mask*)

Disable interrupt of adc digital controller by bitmask.

Return

- ESP_OK Success

Parameters

- `adc_unit`: ADC unit.
- `intr_mask`: Interrupt bitmask. See `adc_digi_intr_t`.

esp_err_t **adc_digi_intr_clear**(*adc_unit_t adc_unit*, *adc_digi_intr_t intr_mask*)

Clear interrupt of adc digital controller by bitmask.

Return

- ESP_OK Success

Parameters

- `adc_unit`: ADC unit.
- `intr_mask`: Interrupt bitmask. See `adc_digi_intr_t`.

uint32_t **adc_digi_intr_get_status**(*adc_unit_t adc_unit*)

Get interrupt status mask of adc digital controller.

Return

- intr Interrupt bitmask, See `adc_digi_intr_t`.

Parameters

- `adc_unit`: ADC unit.

esp_err_t **adc_digi_isr_register**(void (**fn*)) void *

, void **arg*, int *intr_alloc_flags*Register ADC interrupt handler, the handler is an ISR. The handler will be attached to the same CPU core that this function is running on.

Return

- ESP_OK Success
- ESP_ERR_NOT_FOUND Can not find the interrupt that matches the flags.
- ESP_ERR_INVALID_ARG Function pointer error.

Parameters

- `fn`: Interrupt handler function.
- `arg`: Parameter for handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.

esp_err_t **adc_digi_isr_deregister**(void)

Deregister ADC interrupt handler, the handler is an ISR.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG hander error.
- ESP_FAIL ISR not be registered.

esp_err_t **adc_set_i2s_data_source**(*adc_i2s_source_t src*)

Set I2S data source.

Parameters

- `src`: I2S DMA data source, I2S DMA can get data from digital signals or from ADC.

Return

- ESP_OK success

esp_err_t **adc_i2s_mode_init** (*adc_unit_t adc_unit*, *adc_channel_t channel*)

Initialize I2S ADC mode.

Parameters

- `adc_unit`: ADC unit index
- `channel`: ADC channel index

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Header File

- soc/include/hal/adc_types.h

Structures

struct adc_digi_pattern_table_t

ADC digital controller (DMA mode) conversion rules setting.

Public Members

uint8_t **atten** : 2

ADC sampling voltage attenuation configuration. Modification of attenuation affects the range of measurements. 0: measurement range 0 - 800mV, 1: measurement range 0 - 1100mV, 2: measurement range 0 - 1350mV, 3: measurement range 0 - 2600mV.

uint8_t **reserved** : 2

reserved0

uint8_t **channel** : 4

ADC channel index.

uint8_t **val**

Raw data value

struct adc_digi_output_data_t

ADC digital controller (DMA mode) output data format. Used to analyze the acquired ADC (DMA) data.

Note ESP32-S2: Member `channel` can be used to judge the validity of the ADC data, because the role of the arbiter may get invalid ADC data.

Public Members

uint16_t **data** : 12

ADC real output data info. Resolution: 12 bit.

ADC real output data info. Resolution: 11 bit.

uint16_t **channel** : 4

ADC channel index info. For ESP32-S2: If (channel < ADC_CHANNEL_MAX), The data is valid. If (channel > ADC_CHANNEL_MAX), The data is invalid.

struct *adc_digi_output_data_t*::**[anonymous]**::**[anonymous] type1**

When the configured output format is 12bit. `ADC_DIGI_FORMAT_12BIT`

uint16_t **unit** : 1

ADC unit index info. 0: ADC1; 1: ADC2.

struct *adc_digi_output_data_t*::**[anonymous]**::**[anonymous] type2**
When the configured output format is 11bit. `ADC_DIGI_FORMAT_11BIT`

uint16_t **val**
Raw data value

struct adc_digi_clk_t
ADC digital controller (DMA mode) clock system setting. Calculation formula: controller_clk = (`APLL` or `APB`) / (div_num + div_a / div_b + 1).

Note : The clocks of the DAC digital controller use the ADC digital controller clock divider.

Public Members

bool **use_apll**
true: use APLL clock; false: use APB clock.

uint32_t **div_num**
Division factor. Range: 0 ~ 255. Note: When a higher frequency clock is used (the division factor is less than 9), the ADC reading value will be slightly offset.

uint32_t **div_b**
Division factor. Range: 1 ~ 63.

uint32_t **div_a**
Division factor. Range: 0 ~ 63.

struct adc_digi_config_t
CONFIG_IDF_TARGET_ESP32.

ADC digital controller (DMA mode) configuration parameters.

Example setting: When using ADC1 channel0 to measure voltage, the sampling rate is required to be 1 kHz:

```
+---------------------+--------+--------+--------+
| sample rate         | 1 kHz  | 1 kHz  | 1 kHz  |
+---------------------+--------+--------+--------+
| conv_mode           | single |  both  |  alter |
| adc1_pattern_len    |    1   |    1   |    1   |
| dig_clk.use_apll    |    0   |    0   |    0   |
| dig_clk.div_num     |   99   |   99   |   99   |
| dig_clk.div_b       |    0   |    0   |    0   |
| dig_clk.div_a       |    0   |    0   |    0   |
| interval            |  400   |  400   |  200   |
+---------------------+--------+--------+--------+
| `trigger_meas_freq` | 1 kHz  | 1 kHz  | 2 kHz  |
+---------------------+--------+--------+--------+
```

Explanation of the relationship between `conv_limit_num`, `dma_eof_num` and the number of DMA outputs:

```
+---------------------+--------+--------+--------+
| conv_mode           | single |  both  |  alter |
+---------------------+--------+--------+--------+
| trigger meas times  |    1   |    1   |    1   |
+---------------------+--------+--------+--------+
| conv_limit_num      |   +1   |   +1   |   +1   |
| dma_eof_num         |   +1   |   +2   |   +1   |
| dma output (byte)   |   +2   |   +4   |   +2   |
+---------------------+--------+--------+--------+
```

Public Members

bool **conv_limit_en**
Enable the function of limiting ADC conversion times. If the number of ADC conversion trigger count is equal to the `limit_num`, the conversion is stopped.

uint32_t **conv_limit_num**
Set the upper limit of the number of ADC conversion triggers. Range: 1 ~ 255.

uint32_t **adc1_pattern_len**
Pattern table length for digital controller. Range: 0 ~ 16 (0: Don't change the pattern table setting). The pattern table that defines the conversion rules for each SAR ADC. Each table has 16 items, in which channel selection, resolution and attenuation are stored. When the conversion is started, the controller reads conversion rules from the pattern table one by one. For each controller the scan sequence has at most 16 different rules before repeating itself.

uint32_t **adc2_pattern_len**
Refer to `adc1_pattern_len`

adc_digi_pattern_table_t ***adc1_pattern**
Pointer to pattern table for digital controller. The table size defined by `adc1_pattern_len`.

adc_digi_pattern_table_t ***adc2_pattern**
Refer to `adc1_pattern`

adc_digi_convert_mode_t **conv_mode**
ADC conversion mode for digital controller. See `adc_digi_convert_mode_t`.

adc_digi_output_format_t **format**
ADC output data format for digital controller. See `adc_digi_output_format_t`.

uint32_t **interval**
The number of interval clock cycles for the digital controller to trigger the measurement. The unit is the divided clock. Range: 40 ~ 4095. Expression: `trigger_meas_freq` = `controller_clk` / 2 / interval. Refer to *`adc_digi_clk_t`*. Note: The sampling rate of each channel is also related to the conversion mode (See `adc_digi_convert_mode_t`) and pattern table settings.

adc_digi_clk_t **dig_clk**
ADC digital controller clock divider settings. Refer to *`adc_digi_clk_t`*. Note: The clocks of the DAC digital controller use the ADC digital controller clock divider.

uint32_t **dma_eof_num**
DMA eof num of adc digital controller. If the number of measurements reaches `dma_eof_num`, then `dma_in_suc_eof` signal is generated in DMA. Note: The converted data in the DMA in link buffer will be multiple of two bytes.

struct adc_arbiter_t
ADC arbiter work mode and priority setting.

Note ESP32-S2: Only ADC2 support arbiter.

Public Members

adc_arbiter_mode_t **mode**
Refer to `adc_arbiter_mode_t`. Note: only support ADC2.

uint8_t **rtc_pri**
RTC controller priority. Range: 0 ~ 2.

uint8_t **dig_pri**
Digital controller priority. Range: 0 ~ 2.

uint8_t **pwdet_pri**
Wi-Fi controller priority. Range: 0 ~ 2.

struct adc_digi_filter_t

ADC digital controller (DMA mode) filter configuration.

Note For ESP32-S2, The filter object of the ADC is fixed.
Note For ESP32-S2, The filter object is always all enabled channels.

Public Members

adc_unit_t **adc_unit**

Set adc unit number for filter. For ESP32-S2, Filter IDX0/IDX1 can only be used to filter all enabled channels of ADC1/ADC2 unit at the same time.

adc_channel_t **channel**

Set adc channel number for filter. For ESP32-S2, it's always `ADC_CHANNEL_MAX`

adc_digi_filter_mode_t **mode**

Set adc filter mode for filter. See `adc_digi_filter_mode_t`.

struct adc_digi_monitor_t

ADC digital controller (DMA mode) monitor configuration.

Note For ESP32-S2, The monitor object of the ADC is fixed.
Note For ESP32-S2, The monitor object is always all enabled channels.

Public Members

adc_unit_t **adc_unit**

Set adc unit number for monitor. For ESP32-S2, monitor IDX0/IDX1 can only be used to monitor all enabled channels of ADC1/ADC2 unit at the same time.

adc_channel_t **channel**

Set adc channel number for monitor. For ESP32-S2, it's always `ADC_CHANNEL_MAX`

adc_digi_monitor_mode_t **mode**

Set adc monitor mode. See `adc_digi_monitor_mode_t`.

uint32_t **threshold**

Set monitor threshold of adc digital controller.

Macros

ADC_ARBITER_CONFIG_DEFAULT()

ADC arbiter default configuration.

Note ESP32S2: Only ADC2 supports (needs) an arbiter.

Enumerations

enum adc_unit_t

ADC unit enumeration.

Note For ADC digital controller (DMA mode), ESP32 doesn't support `ADC_UNIT_2`, `ADC_UNIT_BOTH`, `ADC_UNIT_ALTER`.

Values:

ADC_UNIT_1 = 1

SAR ADC 1.

ADC_UNIT_2 = 2

SAR ADC 2.

ADC_UNIT_BOTH = 3

SAR ADC 1 and 2.

ADC_UNIT_ALTER = 7
SAR ADC 1 and 2 alternative mode.

ADC_UNIT_MAX

enum adc_channel_t
ADC channels handle. See `adc1_channel_t`, `adc2_channel_t`.

Note For ESP32 ADC1, don't use `ADC_CHANNEL_8`, `ADC_CHANNEL_9`. See `adc1_channel_t`.

Values:

ADC_CHANNEL_0 = 0
ADC channel

ADC_CHANNEL_1
ADC channel

ADC_CHANNEL_2
ADC channel

ADC_CHANNEL_3
ADC channel

ADC_CHANNEL_4
ADC channel

ADC_CHANNEL_5
ADC channel

ADC_CHANNEL_6
ADC channel

ADC_CHANNEL_7
ADC channel

ADC_CHANNEL_8
ADC channel

ADC_CHANNEL_9
ADC channel

ADC_CHANNEL_MAX

enum adc_atten_t
ADC attenuation parameter. Different parameters determine the range of the ADC. See `adc1_config_channel_atten`.

Values:

ADC_ATTEN_DB_0 = 0
No input attenumation, ADC can measure up to approx. 800 mV.

ADC_ATTEN_DB_2_5 = 1
The input voltage of ADC will be attenuated, extending the range of measurement to up to approx. 1100 mV.

ADC_ATTEN_DB_6 = 2
The input voltage of ADC will be attenuated, extending the range of measurement to up to approx. 1350 mV.

ADC_ATTEN_DB_11 = 3
The input voltage of ADC will be attenuated, extending the range of measurement to up to approx. 2600 mV.

ADC_ATTEN_MAX

enum adc_i2s_source_t
ESP32 ADC DMA source selection.

Values:

ADC_I2S_DATA_SRC_IO_SIG = 0
I2S data from GPIO matrix signal

ADC_I2S_DATA_SRC_ADC = 1
I2S data from ADC

ADC_I2S_DATA_SRC_MAX

enum adc_bits_width_t
ADC resolution setting option.

Note For ESP32-S2. Only 13 bit resolution is supported. For ESP32. 13 bit resolution is not supported.

Values:

ADC_WIDTH_BIT_9 = 0
ADC capture width is 9Bit. Only ESP32 is supported.

ADC_WIDTH_BIT_10 = 1
ADC capture width is 10Bit. Only ESP32 is supported.

ADC_WIDTH_BIT_11 = 2
ADC capture width is 11Bit. Only ESP32 is supported.

ADC_WIDTH_BIT_12 = 3
ADC capture width is 12Bit. Only ESP32 is supported.

ADC_WIDTH_BIT_13 = 4
ADC capture width is 13Bit. Only ESP32-S2 is supported.

ADC_WIDTH_MAX

enum adc_digi_convert_mode_t
ADC digital controller (DMA mode) work mode.

Note The conversion mode affects the sampling frequency: SINGLE_UNIT_1: When the measurement is triggered, only ADC1 is sampled once. SINGLE_UNIT_2: When the measurement is triggered, only ADC2 is sampled once. BOTH_UNIT : When the measurement is triggered, ADC1 and ADC2 are sampled at the same time. ALTER_UNIT : When the measurement is triggered, ADC1 or ADC2 samples alternately.

Values:

ADC_CONV_SINGLE_UNIT_1 = 1
SAR ADC 1.

ADC_CONV_SINGLE_UNIT_2 = 2
SAR ADC 2.

ADC_CONV_BOTH_UNIT = 3
SAR ADC 1 and 2.

ADC_CONV_ALTER_UNIT = 7
SAR ADC 1 and 2 alternative mode.

ADC_CONV_UNIT_MAX

enum adc_digi_output_format_t
ADC digital controller (DMA mode) output data format option.

Values:

ADC_DIGI_FORMAT_12BIT
ADC to DMA data format, [15:12]-channel, [11: 0]-12 bits ADC data (*adc_digi_output_data_t*). Note: For single convert mode.

ADC_DIGI_FORMAT_11BIT
ADC to DMA data format, [15]-adc unit, [14:11]-channel, [10: 0]-11 bits ADC data (*adc_digi_output_data_t*). Note: For multi or alter convert mode.

ADC_DIGI_FORMAT_MAX

enum adc_arbiter_mode_t

ADC arbiter work mode option.

Note ESP32-S2: Only ADC2 support arbiter.

Values:

ADC_ARB_MODE_SHIELD
Force shield arbiter, Select the highest priority controller to work.

ADC_ARB_MODE_FIX
Fixed priority switch controller mode.

ADC_ARB_MODE_LOOP
Loop priority switch controller mode. Each controller has the same priority, and the arbiter will switch to the next controller after the measurement is completed.

enum adc_digi_intr_t

ADC digital controller (DMA mode) interrupt type options.

Values:

ADC_DIGI_INTR_MASK_MONITOR = 0x1

ADC_DIGI_INTR_MASK_MEAS_DONE = 0x2

ADC_DIGI_INTR_MASK_ALL = 0x3

enum adc_digi_filter_idx_t

ADC digital controller (DMA mode) filter index options.

Note For ESP32-S2, The filter object of the ADC is fixed.

Values:

ADC_DIGI_FILTER_IDX0 = 0
The filter index 0. For ESP32-S2, It can only be used to filter all enabled channels of ADC1 unit at the same time.

ADC_DIGI_FILTER_IDX1
The filter index 1. For ESP32-S2, It can only be used to filter all enabled channels of ADC2 unit at the same time.

ADC_DIGI_FILTER_IDX_MAX

enum adc_digi_filter_mode_t

ADC digital controller (DMA mode) filter type options. Expression: filter_data = (k-1)/k * last_data + new_data / k.

Values:

ADC_DIGI_FILTER_IIR_2 = 0
The filter mode is first-order IIR filter. The coefficient is 2.

ADC_DIGI_FILTER_IIR_4
The filter mode is first-order IIR filter. The coefficient is 4.

ADC_DIGI_FILTER_IIR_8
The filter mode is first-order IIR filter. The coefficient is 8.

ADC_DIGI_FILTER_IIR_16
The filter mode is first-order IIR filter. The coefficient is 16.

ADC_DIGI_FILTER_IIR_64
The filter mode is first-order IIR filter. The coefficient is 64.

ADC_DIGI_FILTER_IIR_MAX

enum adc_digi_monitor_idx_t

ADC digital controller (DMA mode) monitor index options.

Note For ESP32-S2, The monitor object of the ADC is fixed.

Values:

ADC_DIGI_MONITOR_IDX0 = 0

The monitor index 0. For ESP32-S2, It can only be used to monitor all enabled channels of ADC1 unit at the same time.

ADC_DIGI_MONITOR_IDX1

The monitor index 1. For ESP32-S2, It can only be used to monitor all enabled channels of ADC2 unit at the same time.

ADC_DIGI_MONITOR_IDX_MAX

enum adc_digi_monitor_mode_t

Set monitor mode of adc digital controller. MONITOR_HIGH:If ADC_OUT > threshold, Generates monitor interrupt. MONITOR_LOW: If ADC_OUT < threshold, Generates monitor interrupt.

Values:

ADC_DIGI_MONITOR_HIGH = 0

If ADC_OUT > threshold, Generates monitor interrupt.

ADC_DIGI_MONITOR_LOW

If ADC_OUT < threshold, Generates monitor interrupt.

ADC_DIGI_MONITOR_MAX

Header File

- driver/include/driver/adc_common.h

Functions

void **adc_power_on**(void)

Enable ADC power.

void **adc_power_off**(void)

Power off SAR ADC.

void **adc_power_acquire**(void)

Increment the usage counter for ADC module. ADC will stay powered on while the counter is greater than 0. Call adc_power_release when done using the ADC.

void **adc_power_release**(void)

Decrement the usage counter for ADC module. ADC will stay powered on while the counter is greater than 0. Call this function when done using the ADC.

esp_err_t **adc_gpio_init**(*adc_unit_t adc_unit*, *adc_channel_t channel*)

Initialize ADC pad.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `adc_unit`: ADC unit index
- `channel`: ADC channel index

esp_err_t **adc1_pad_get_io_num**(*adc1_channel_t channel*, *gpio_num_t *gpio_num*)

Get the GPIO number of a specific ADC1 channel.

Return

- ESP_OK if success
- ESP_ERR_INVALID_ARG if channel not valid

Parameters

- `channel`: Channel to get the GPIO number
- `gpio_num`: output buffer to hold the GPIO number

esp_err_t **adc1_config_channel_atten**(*adc1_channel_t channel*, *adc_atten_t atten*)

Set the attenuation of a particular channel on ADC1, and configure its associated GPIO pin mux.

The default ADC voltage is for attenuation 0 dB and listed in the table below. By setting higher attenuation it is possible to read higher voltages.

Due to ADC characteristics, most accurate results are obtained within the "suggested range" shown in the following table.

SoC	attenuation (dB)	suggested range (mV)
ESP32	0	100 ~ 950
	2.5	100 ~ 1250
	6	150 ~ 1750
	11	150 ~ 2450
ESP32-S2	0	0 ~ 750
	2.5	0 ~ 1050
	6	0 ~ 1300
	11	0 ~ 2500

For maximum accuracy, use the ADC calibration APIs and measure voltages within these recommended ranges.

Note For any given channel, this function must be called before the first time `adc1_get_raw()` is called for that channel.

Note This function can be called multiple times to configure multiple ADC channels simultaneously. You may call `adc1_get_raw()` only after configuring a channel.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `channel`: ADC1 channel to configure
- `atten`: Attenuation level

esp_err_t **adc1_config_width**(*adc_bits_width_t width_bit*)

Configure ADC1 capture width, meanwhile enable output invert for ADC1. The configuration is for all channels of ADC1.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `width_bit`: Bit capture width for ADC1

int **adc1_get_raw**(*adc1_channel_t channel*)

Take an ADC1 reading from a single channel.

Note ESP32: When the power switch of SARADC1, SARADC2, HALL sensor and AMP sensor is turned on, the input of GPIO36 and GPIO39 will be pulled down for about 80ns. When enabling power

for any of these peripherals, ignore input from GPIO36 and GPIO39. Please refer to section 3.11 of 'ECO_and_Workarounds_for_Bugs_in_ESP32' for the description of this issue.

Note Call `adc1_config_width()` before the first time this function is called.

Note For any given channel, adc1_config_channel_atten(channel) must be called before the first time this function is called. Configuring a new channel does not prevent a previously configured channel from being read.

Return

- -1: Parameter error
- Other: ADC1 channel reading.

Parameters

- `channel`: ADC1 channel to read

esp_err_t **adc_set_data_inv**(*adc_unit_t adc_unit*, bool *inv_en*)

Set ADC data invert.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `adc_unit`: ADC unit index
- `inv_en`: whether enable data invert

esp_err_t **adc_set_clk_div**(uint8_t *clk_div*)

Set ADC source clock.

Return

- ESP_OK success

Parameters

- `clk_div`: ADC clock divider, ADC clock is divided from APB clock

esp_err_t **adc_set_data_width**(*adc_unit_t adc_unit*, *adc_bits_width_t width_bit*)

Configure ADC capture width.

Note ESP32-S2 only supports `ADC_WIDTH_BIT_13`.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `adc_unit`: ADC unit index
- `width_bit`: Bit capture width for ADC unit. ESP32-S2 only supports `ADC_WIDTH_BIT_13`.

void **adc1_ulp_enable**(void)

Configure ADC1 to be usable by the ULP.

This function reconfigures ADC1 to be controlled by the ULP. Effect of this function can be reverted using `adc1_get_raw()` function.

Note that adc1_config_channel_atten, `adc1_config_width()` functions need to be called to configure ADC1 channels, before ADC1 is used by the ULP.

esp_err_t **adc2_pad_get_io_num**(*adc2_channel_t channel*, *gpio_num_t *gpio_num*)

Get the GPIO number of a specific ADC2 channel.

Return

- ESP_OK if success
- ESP_ERR_INVALID_ARG if channel not valid

Parameters

- `channel`: Channel to get the GPIO number
- `gpio_num`: output buffer to hold the GPIO number

esp_err_t **adc2_config_channel_atten**(*adc2_channel_t channel*, *adc_atten_t atten*)

Configure the ADC2 channel, including setting attenuation.

The default ADC voltage is for attenuation 0 dB and listed in the table below. By setting higher attenuation it is possible to read higher voltages.

Due to ADC characteristics, most accurate results are obtained within the "suggested range" shown in the following table.

SoC	attenuation (dB)	suggested range (mV)
ESP32	0	100 ~ 950
	2.5	100 ~ 1250
	6	150 ~ 1750
	11	150 ~ 2450
ESP32-S2	0	0 ~ 750
	2.5	0 ~ 1050
	6	0 ~ 1300
	11	0 ~ 2500

For maximum accuracy, use the ADC calibration APIs and measure voltages within these recommended ranges.

Note This function also configures the input GPIO pin mux to connect it to the ADC2 channel. It must be called before calling `adc2_get_raw()` for this channel.

Note For any given channel, this function must be called before the first time `adc2_get_raw()` is called for that channel.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `channel`: ADC2 channel to configure
- `atten`: Attenuation level

esp_err_t **adc2_get_raw**(*adc2_channel_t channel*, *adc_bits_width_t width_bit*, int **raw_out*)

Take an ADC2 reading on a single channel.

Note ESP32: When the power switch of SARADC1, SARADC2, HALL sensor and AMP sensor is turned on, the input of GPIO36 and GPIO39 will be pulled down for about 80ns. When enabling power for any of these peripherals, ignore input from GPIO36 and GPIO39. Please refer to section 3.11 of 'ECO_and_Workarounds_for_Bugs_in_ESP32' for the description of this issue. As a workaround, call adc_power_acquire() in the app. This will result in higher power consumption (by ~1mA), but will remove the glitches on GPIO36 and GPIO39.

Note ESP32: For a given channel, `adc2_config_channel_atten()` must be called before the first time this function is called. If Wi-Fi is started via `esp_wifi_start()`, this function will always fail with `ESP_ERR_TIMEOUT`.

Note ESP32-S2: ADC2 support hardware arbiter. The arbiter is to improve the use efficiency of ADC2. After the control right is robbed by the high priority, the low priority controller will read the invalid ADC2 data. Default priority: Wi-Fi > RTC > Digital;

Return

- ESP_OK if success
- ESP_ERR_TIMEOUT ADC2 is being used by other controller and the request timed out.
- ESP_ERR_INVALID_STATE The controller status is invalid. Please try again.

Parameters

- `channel`: ADC2 channel to read
- `width_bit`: Bit capture width for ADC2. ESP32-S2 only supports `ADC_WIDTH_BIT_13`.
- `raw_out`: the variable to hold the output data.

esp_err_t **adc_vref_to_gpio**(*adc_unit_t adc_unit*, *gpio_num_t gpio*)

Output ADC1 or ADC2's reference voltage to adc2_channe_t's IO.

This function routes the internal reference voltage of ADCn to one of ADC2's channels. This reference voltage can then be manually measured for calibration purposes.

Note ESP32 only supports output of ADC2's internal reference voltage.

Return

- ESP_OK: v_ref successfully routed to selected GPIO
- ESP_ERR_INVALID_ARG: Unsupported GPIO

Parameters

- [in] adc_unit: ADC unit index
- [in] gpio: GPIO number (Only ADC2's channels IO are supported)

esp_err_t **adc2_vref_to_gpio**(*gpio_num_t gpio*)

Output ADC2 reference voltage to adc2_channe_t's IO.

This function routes the internal reference voltage of ADCn to one of ADC2's channels. This reference voltage can then be manually measured for calibration purposes.

Return

- ESP_OK: v_ref successfully routed to selected GPIO
- ESP_ERR_INVALID_ARG: Unsupported GPIO

Parameters

- [in] gpio: GPIO number (ADC2's channels are supported)

esp_err_t **adc_digi_init**(void)

ADC digital controller initialization.

Return

- ESP_OK Success

esp_err_t **adc_digi_deinit**(void)

ADC digital controller deinitialization.

Return

- ESP_OK Success

esp_err_t **adc_digi_controller_config**(**const** *adc_digi_config_t* **config*)

Setting the digital controller.

Return

- ESP_OK Success

Parameters

- config: Pointer to digital controller paramter. Refer to *adc_digi_config_t*.

Macros

ADC_ATTEN_0db

ADC rtc controller attenuation option.

Note This definitions are only for being back-compatible

ADC_ATTEN_2_5db

ADC_ATTEN_6db

ADC_ATTEN_11db

ADC_WIDTH_9Bit

ADC_WIDTH_10Bit

ADC_WIDTH_11Bit

ADC_WIDTH_12Bit

Enumerations

enum adc1_channel_t

Values:

ADC1_CHANNEL_0 = 0
ADC1 channel 0 is GPIO36 (ESP32), GPIO1 (ESP32-S2)

ADC1_CHANNEL_1
ADC1 channel 1 is GPIO37 (ESP32), GPIO2 (ESP32-S2)

ADC1_CHANNEL_2
ADC1 channel 2 is GPIO38 (ESP32), GPIO3 (ESP32-S2)

ADC1_CHANNEL_3
ADC1 channel 3 is GPIO39 (ESP32), GPIO4 (ESP32-S2)

ADC1_CHANNEL_4
ADC1 channel 4 is GPIO32 (ESP32), GPIO5 (ESP32-S2)

ADC1_CHANNEL_5
ADC1 channel 5 is GPIO33 (ESP32), GPIO6 (ESP32-S2)

ADC1_CHANNEL_6
ADC1 channel 6 is GPIO34 (ESP32), GPIO7 (ESP32-S2)

ADC1_CHANNEL_7
ADC1 channel 7 is GPIO35 (ESP32), GPIO8 (ESP32-S2)

ADC1_CHANNEL_8
ADC1 channel 6 is GPIO9 (ESP32-S2)

ADC1_CHANNEL_9
ADC1 channel 7 is GPIO10 (ESP32-S2)

ADC1_CHANNEL_MAX

enum adc2_channel_t

Values:

ADC2_CHANNEL_0 = 0
ADC2 channel 0 is GPIO4 (ESP32), GPIO11 (ESP32-S2)

ADC2_CHANNEL_1
ADC2 channel 1 is GPIO0 (ESP32), GPIO12 (ESP32-S2)

ADC2_CHANNEL_2
ADC2 channel 2 is GPIO2 (ESP32), GPIO13 (ESP32-S2)

ADC2_CHANNEL_3
ADC2 channel 3 is GPIO15 (ESP32), GPIO14 (ESP32-S2)

ADC2_CHANNEL_4
ADC2 channel 4 is GPIO13 (ESP32), GPIO15 (ESP32-S2)

ADC2_CHANNEL_5
ADC2 channel 5 is GPIO12 (ESP32), GPIO16 (ESP32-S2)

ADC2_CHANNEL_6
ADC2 channel 6 is GPIO14 (ESP32), GPIO17 (ESP32-S2)

ADC2_CHANNEL_7
ADC2 channel 7 is GPIO27 (ESP32), GPIO18 (ESP32-S2)

ADC2_CHANNEL_8
ADC2 channel 8 is GPIO25 (ESP32), GPIO19 (ESP32-S2)

ADC2_CHANNEL_9
ADC2 channel 9 is GPIO26 (ESP32), GPIO20 (ESP32-S2)

ADC2_CHANNEL_MAX

enum adc_i2s_encode_t

ADC digital controller encode option.

Values:

ADC_ENCODE_12BIT

ADC to DMA data format, , [15:12]-channel [11:0]-12 bits ADC data

ADC_ENCODE_11BIT

ADC to DMA data format, [15]-unit, [14:11]-channel [10:0]-11 bits ADC data

ADC_ENCODE_MAX

ADC Calibration

Header File

- esp_adc_cal/include/esp_adc_cal.h

Functions

esp_err_t **esp_adc_cal_check_efuse**(*esp_adc_cal_value_t value_type*)

Checks if ADC calibration values are burned into eFuse.

This function checks if ADC reference voltage or Two Point values have been burned to the eFuse of the current ESP32

Return

- ESP_OK: The calibration mode is supported in eFuse
- ESP_ERR_NOT_SUPPORTED: Error, eFuse values are not burned
- ESP_ERR_INVALID_ARG: Error, invalid argument (ESP_ADC_CAL_VAL_DEFAULT_VREF)

Parameters

- `value_type`: Type of calibration value (ESP_ADC_CAL_VAL_EFUSE_VREF or ESP_ADC_CAL_VAL_EFUSE_TP)

esp_adc_cal_value_t **esp_adc_cal_characterize**(*adc_unit_t adc_num*, *adc_atten_t atten*, *adc_bits_width_t bit_width*, uint32_t *default_vref*, *esp_adc_cal_characteristics_t* **chars*)

Characterize an ADC at a particular attenuation.

This function will characterize the ADC at a particular attenuation and generate the ADC-Voltage curve in the form of [y = coeff_a * x + coeff_b]. Characterization can be based on Two Point values, eFuse Vref, or default Vref and the calibration values will be prioritized in that order.

Note For ESP32, Two Point values and eFuse Vref calibration can be enabled/disabled using menuconfig. For ESP32s2, only Two Point values calibration and only ADC_WIDTH_BIT_13 is supported. The parameter default_vref is unused.

Return

- ESP_ADC_CAL_VAL_EFUSE_VREF: eFuse Vref used for characterization
- ESP_ADC_CAL_VAL_EFUSE_TP: Two Point value used for characterization (only in Linear Mode)
- ESP_ADC_CAL_VAL_DEFAULT_VREF: Default Vref used for characterization

Parameters

- `[in] adc_num`: ADC to characterize (ADC_UNIT_1 or ADC_UNIT_2)
- `[in] atten`: Attenuation to characterize
- `[in] bit_width`: Bit width configuration of ADC
- `[in] default_vref`: Default ADC reference voltage in mV (Only in ESP32, used if eFuse values is not available)
- `[out] chars`: Pointer to empty structure used to store ADC characteristics

uint32_t **esp_adc_cal_raw_to_voltage**(uint32_t *adc_reading*, **const** *esp_adc_cal_characteristics_t* **chars*)

Convert an ADC reading to voltage in mV.

This function converts an ADC reading to a voltage in mV based on the ADC' s characteristics.

Note Characteristics structure must be initialized before this function is called (call esp_adc_cal_characterize())

Return Voltage in mV

Parameters

- `[in] adc_reading`: ADC reading
- `[in] chars`: Pointer to initialized structure containing ADC characteristics

esp_err_t **esp_adc_cal_get_voltage** (*adc_channel_t channel*, **const** *esp_adc_cal_characteristics_t* **chars*, uint32_t **voltage*)

Reads an ADC and converts the reading to a voltage in mV.

This function reads an ADC then converts the raw reading to a voltage in mV based on the characteristics provided. The ADC that is read is also determined by the characteristics.

Note The Characteristics structure must be initialized before this function is called (call esp_adc_cal_characterize())

Return

- ESP_OK: ADC read and converted to mV
- ESP_ERR_TIMEOUT: Error, timed out attempting to read ADC
- ESP_ERR_INVALID_ARG: Error due to invalid arguments

Parameters

- `[in] channel`: ADC Channel to read
- `[in] chars`: Pointer to initialized ADC characteristics structure
- `[out] voltage`: Pointer to store converted voltage

Structures

struct esp_adc_cal_characteristics_t

Structure storing characteristics of an ADC.

Note Call esp_adc_cal_characterize() to initialize the structure

Public Members

adc_unit_t **adc_num**

ADC number

adc_atten_t **atten**

ADC attenuation

adc_bits_width_t **bit_width**

ADC bit width

uint32_t **coeff_a**

Gradient of ADC-Voltage curve

uint32_t **coeff_b**

Offset of ADC-Voltage curve

uint32_t **vref**

Vref used by lookup table

const uint32_t ***low_curve**

Pointer to low Vref curve of lookup table (NULL if unused)

const uint32_t ***high_curve**

Pointer to high Vref curve of lookup table (NULL if unused)

Enumerations

enum esp_adc_cal_value_t

Type of calibration value used in characterization.

Values:

ESP_ADC_CAL_VAL_EFUSE_VREF = 0

Characterization based on reference voltage stored in eFuse

ESP_ADC_CAL_VAL_EFUSE_TP = 1

Characterization based on Two Point values stored in eFuse

ESP_ADC_CAL_VAL_DEFAULT_VREF = 2

Characterization based on default reference voltage

ESP_ADC_CAL_VAL_MAX

GPIO Lookup Macros

Header File

- soc/soc/esp32s2/include/soc/adc_channel.h

Macros

ADC1_GPIO1_CHANNEL
ADC1_CHANNEL_0_GPIO_NUM

ADC1_GPIO2_CHANNEL

ADC1_CHANNEL_1_GPIO_NUM

ADC1_GPIO3_CHANNEL

ADC1_CHANNEL_2_GPIO_NUM

ADC1_GPIO4_CHANNEL

ADC1_CHANNEL_3_GPIO_NUM

ADC1_GPIO5_CHANNEL

ADC1_CHANNEL_4_GPIO_NUM

ADC1_GPIO6_CHANNEL

ADC1_CHANNEL_5_GPIO_NUM

ADC1_GPIO7_CHANNEL

ADC1_CHANNEL_6_GPIO_NUM

ADC1_GPIO8_CHANNEL

ADC1_CHANNEL_7_GPIO_NUM

ADC1_GPIO9_CHANNEL

ADC1_CHANNEL_8_GPIO_NUM

ADC1_GPIO10_CHANNEL

ADC1_CHANNEL_9_GPIO_NUM

ADC2_GPIO11_CHANNEL

ADC2_CHANNEL_0_GPIO_NUM

ADC2_GPIO12_CHANNEL

ADC2_CHANNEL_1_GPIO_NUM

ADC2_GPIO13_CHANNEL

ADC2_CHANNEL_2_GPIO_NUM

ADC2_GPIO14_CHANNEL

ADC2_CHANNEL_3_GPIO_NUM

ADC2_GPIO15_CHANNEL

ADC2_CHANNEL_4_GPIO_NUM

ADC2_GPIO16_CHANNEL

ADC2_CHANNEL_5_GPIO_NUM

ADC2_GPIO17_CHANNEL

ADC2_CHANNEL_6_GPIO_NUM

ADC2_GPIO18_CHANNEL

ADC2_CHANNEL_7_GPIO_NUM

ADC2_GPIO19_CHANNEL

ADC2_CHANNEL_8_GPIO_NUM

ADC2_GPIO20_CHANNEL

ADC2_CHANNEL_9_GPIO_NUM

2.2.2 Digital To Analog Converter

Overview

ESP32-S2 has two 8-bit DAC (digital to analog converter) channels, connected to GPIO17 (Channel 1) and GPIO18 (Channel 2).

The DAC driver allows these channels to be set to arbitrary voltages.

The DAC channels can also be driven with DMA-style written sample data, via the *I2S driver* when using the "built-in DAC mode".

For other analog output options, see the *Sigma-delta Modulation module* and the *LED Control module*. Both these modules produce high frequency PWM output, which can be hardware low-pass filtered in order to generate a lower frequency analog output.

Application Example

Setting DAC channel 1 (GPIO17) voltage to approx 0.78 of VDD_A voltage (VDD * 200 / 255). For VDD_A 3.3V, this is 2.59V:

```
#include <driver/dac.h>

...

    dac_output_enable(DAC_CHANNEL_1);
    dac_output_voltage(DAC_CHANNEL_1, 200);
```

API Reference

Header File

- driver/esp32s2/include/driver/dac.h

Functions

esp_err_t **dac_digi_init**(void)
DAC digital controller initialization.

Return
- ESP_OK success

esp_err_t **dac_digi_deinit**(void)
DAC digital controller deinitialization.

Return
- ESP_OK success

esp_err_t **dac_digi_controller_config**(**const** *dac_digi_config_t* *cfg)
Setting the DAC digital controller.

Return
- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `cfg`: Pointer to digital controller paramter. See *dac_digi_config_t*.

esp_err_t **dac_digi_start**(void)
DAC digital controller start output voltage.

Return
- ESP_OK success

esp_err_t **dac_digi_stop**(void)
DAC digital controller stop output voltage.

Return
- ESP_OK success

esp_err_t **dac_digi_fifo_reset**(void)
Reset DAC digital controller FIFO.

Return
- ESP_OK success

esp_err_t **dac_digi_reset**(void)
Reset DAC digital controller.

Return
- ESP_OK success

Header File

- driver/include/driver/dac_common.h

Functions

esp_err_t **dac_pad_get_io_num**(*dac_channel_t channel*, *gpio_num_t* *gpio_num)
Get the GPIO number of a specific DAC channel.

Return
- ESP_OK if success

Parameters
- `channel`: Channel to get the gpio number
- `gpio_num`: output buffer to hold the gpio number

esp_err_t **dac_output_voltage**(*dac_channel_t channel*, uint8_t *dac_value*)
Set DAC output voltage. DAC output is 8-bit. Maximum (255) corresponds to VDD3P3_RTC.

Note Need to configure DAC pad before calling this function. DAC channel 1 is attached to GPIO25, DAC channel 2 is attached to GPIO26

Return

- ESP_OK success

Parameters

- `channel`: DAC channel
- `dac_value`: DAC output value

esp_err_t **dac_output_enable**(*dac_channel_t channel*)

DAC pad output enable.

Note DAC channel 1 is attached to GPIO25, DAC channel 2 is attached to GPIO26 I2S left channel will be mapped to DAC channel 2 I2S right channel will be mapped to DAC channel 1

Parameters

- `channel`: DAC channel

esp_err_t **dac_output_disable**(*dac_channel_t channel*)

DAC pad output disable.

Note DAC channel 1 is attached to GPIO25, DAC channel 2 is attached to GPIO26

Return

- ESP_OK success

Parameters

- `channel`: DAC channel

esp_err_t **dac_cw_generator_enable**(void)

Enable cosine wave generator output.

Return

- ESP_OK success

esp_err_t **dac_cw_generator_disable**(void)

Disable cosine wave generator output.

Return

- ESP_OK success

esp_err_t **dac_cw_generator_config**(*dac_cw_config_t *cw*)

Config the cosine wave generator function in DAC module.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG The parameter is NULL.

Parameters

- `cw`: Configuration.

GPIO Lookup Macros Some useful macros can be used to specified the GPIO number of a DAC channel, or vice versa. e.g.

1. `DAC_CHANNEL_1_GPIO_NUM` is the GPIO number of channel 1 (GPIO17);
2. `DAC_GPIO18_CHANNEL` is the channel number of GPIO 26 (channel 2).

Header File

- soc/soc/esp32s2/include/soc/dac_channel.h

Macros

DAC_GPIO17_CHANNEL

DAC_CHANNEL_1_GPIO_NUM

DAC_GPIO18_CHANNEL

DAC_CHANNEL_2_GPIO_NUM

Header File

- soc/include/hal/dac_types.h

Structures

struct dac_cw_config_t

Config the cosine wave generator function in DAC module.

Public Members

dac_channel_t **en_ch**

Enable the cosine wave generator of DAC channel.

dac_cw_scale_t **scale**

Set the amplitude of the cosine wave generator output.

dac_cw_phase_t **phase**

Set the phase of the cosine wave generator output.

uint32_t **freq**

Set frequency of cosine wave generator output. Range: 130(130Hz) ~ 55000(100KHz).

int8_t **offset**

Set the voltage value of the DC component of the cosine wave generator output. Note: Unreasonable settings can cause waveform to be oversaturated. Range: -128 ~ 127.

struct dac_digi_config_t

DAC digital controller (DMA mode) configuration parameters.

Public Members

dac_digi_convert_mode_t **mode**

DAC digital controller (DMA mode) work mode. See `dac_digi_convert_mode_t`.

uint32_t **interval**

The number of interval clock cycles for the DAC digital controller to output voltage. The unit is the divided clock. Range: 1 ~ 4095. Expression: `dac_output_freq` = `controller_clk` / interval. Refer to *`adc_digi_clk_t`*. Note: The sampling rate of each channel is also related to the conversion mode (See `dac_digi_convert_mode_t`) and pattern table settings.

adc_digi_clk_t **dig_clk**

DAC digital controller clock divider settings. Refer to *`adc_digi_clk_t`*. Note: The clocks of the DAC digital controller use the ADC digital controller clock divider.

Enumerations

enum dac_channel_t

Values:

DAC_CHANNEL_1 = 0

DAC channel 1 is GPIO25(ESP32) / GPIO17(ESP32S2)

DAC_CHANNEL_2 = 1

DAC channel 2 is GPIO26(ESP32) / GPIO18(ESP32S2)

DAC_CHANNEL_MAX

enum dac_cw_scale_t

The multiple of the amplitude of the cosine wave generator. The max amplitude is VDD3P3_RTC.

Values:

DAC_CW_SCALE_1 = 0x0

1/1. Default.

DAC_CW_SCALE_2 = 0x1
1/2.

DAC_CW_SCALE_4 = 0x2
1/4.

DAC_CW_SCALE_8 = 0x3
1/8.

enum dac_cw_phase_t
Set the phase of the cosine wave generator output.

Values:

DAC_CW_PHASE_0 = 0x2
Phase shift +0°

DAC_CW_PHASE_180 = 0x3
Phase shift +180°

enum dac_digi_convert_mode_t
DAC digital controller (DMA mode) work mode.

Values:

DAC_CONV_NORMAL
The data in the DMA buffer is simultaneously output to the enable channel of the DAC.

DAC_CONV_ALTER
The data in the DMA buffer is alternately output to the enable channel of the DAC.

DAC_CONV_MAX

2.2.3 GPIO & RTC GPIO

Overview

The ESP32-S2 chip features 43 physical GPIO pads. Some GPIO pads cannot be used or do not have the corresponding pin on the chip package(refer to technical reference manual). Each pad can be used as a general purpose I/O or can be connected to an internal peripheral signal.

- Note that GPIO26-32 are usually used for SPI flash.
- GPIO46 is fixed to pull-down and is input only

There is also separate "RTC GPIO" support, which functions when GPIOs are routed to the "RTC" low-power and analog subsystem. These pin functions can be used when in deep sleep, when the *Ultra Low Power co-processor* is running, or when analog functions such as ADC/DAC/etc are in use.

Application Example

GPIO output and input interrupt example: peripherals/gpio.

API Reference - Normal GPIO

Header File

- driver/include/driver/gpio.h

Functions

esp_err_t **gpio_config**(**const** *gpio_config_t* **pGPIOConfig*)

GPIO common configuration.

Configure GPIO' s Mode,pull-up,PullDown,IntrType

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `pGPIOConfig`: Pointer to GPIO configure struct

esp_err_t **gpio_reset_pin**(*gpio_num_t gpio_num*)

Reset an gpio to default state (select gpio function, enable pullup and disable input and output).

Note This function also configures the IOMUX for this pin to the GPIO function, and disconnects any other peripheral output configured via GPIO Matrix.

Return Always return ESP_OK.

Parameters

- `gpio_num`: GPIO number.

esp_err_t **gpio_set_intr_type**(*gpio_num_t gpio_num*, *gpio_int_type_t intr_type*)

GPIO set interrupt trigger type.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number. If you want to set the trigger type of e.g. of GPIO16, gpio_num should be GPIO_NUM_16 (16);
- `intr_type`: Interrupt type, select from gpio_int_type_t

esp_err_t **gpio_intr_enable**(*gpio_num_t gpio_num*)

Enable GPIO module interrupt signal.

Note Please do not use the interrupt of GPIO36 and GPIO39 when using ADC or Wi-Fi and Bluetooth with sleep mode enabled. Please refer to the comments of `adc1_get_raw`. Please refer to section 3.11 of 'ECO_and_Workarounds_for_Bugs_in_ESP32' for the description of this issue. As a workaround, call adc_power_acquire() in the app. This will result in higher power consumption (by ~1mA), but will remove the glitches on GPIO36 and GPIO39.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number. If you want to enable an interrupt on e.g. GPIO16, gpio_num should be GPIO_NUM_16 (16);

esp_err_t **gpio_intr_disable**(*gpio_num_t gpio_num*)

Disable GPIO module interrupt signal.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number. If you want to disable the interrupt of e.g. GPIO16, gpio_num should be GPIO_NUM_16 (16);

esp_err_t **gpio_set_level**(*gpio_num_t gpio_num*, uint32_t *level*)

GPIO set output level.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO number error

Parameters

- `gpio_num`: GPIO number. If you want to set the output level of e.g. GPIO16, gpio_num should be GPIO_NUM_16 (16);
- `level`: Output level. 0: low ; 1: high

int **gpio_get_level** (*gpio_num_t gpio_num*)

GPIO get input level.

Warning If the pad is not configured for input (or input and output) the returned value is always 0.

Return

- 0 the GPIO input level is 0
- 1 the GPIO input level is 1

Parameters

- `gpio_num`: GPIO number. If you want to get the logic level of e.g. pin GPIO16, gpio_num should be GPIO_NUM_16 (16);

esp_err_t **gpio_set_direction** (*gpio_num_t gpio_num*, *gpio_mode_t mode*)

GPIO set direction.

Configure GPIO direction,such as output_only,input_only,output_and_input

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO error

Parameters

- `gpio_num`: Configure GPIO pins number, it should be GPIO number. If you want to set direction of e.g. GPIO16, gpio_num should be GPIO_NUM_16 (16);
- `mode`: GPIO direction

esp_err_t **gpio_set_pull_mode** (*gpio_num_t gpio_num*, *gpio_pull_mode_t pull*)

Configure GPIO pull-up/pull-down resistors.

Only pins that support both input & output have integrated pull-up and pull-down resistors. Input-only GPIOs 34-39 do not.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG : Parameter error

Parameters

- `gpio_num`: GPIO number. If you want to set pull up or down mode for e.g. GPIO16, gpio_num should be GPIO_NUM_16 (16);
- `pull`: GPIO pull up/down mode.

esp_err_t **gpio_wakeup_enable** (*gpio_num_t gpio_num*, *gpio_int_type_t intr_type*)

Enable GPIO wake-up function.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number.
- `intr_type`: GPIO wake-up type. Only GPIO_INTR_LOW_LEVEL or GPIO_INTR_HIGH_LEVEL can be used.

esp_err_t **gpio_wakeup_disable** (*gpio_num_t gpio_num*)

Disable GPIO wake-up function.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number

esp_err_t **gpio_isr_register** (void (**fn*)) void *

, void **arg*, int *intr_alloc_flags*, *gpio_isr_handle_t* **handle*Register GPIO interrupt handler, the handler is an ISR. The handler will be attached to the same CPU core that this function is running on.

This ISR function is called whenever any GPIO interrupt occurs. See the alternative gpio_install_isr_service() and gpio_isr_handler_add() API in order to have the driver support per-GPIO ISRs.

To disable or remove the ISR, pass the returned handle to the *interrupt allocation functions*.

Parameters

- `fn`: Interrupt handler function.
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `arg`: Parameter for handler function
- `handle`: Pointer to return handle. If non-NULL, a handle for the interrupt will be returned here.

Return

- ESP_OK Success ;
- ESP_ERR_INVALID_ARG GPIO error
- ESP_ERR_NOT_FOUND No free interrupt found with the specified flags

esp_err_t **gpio_pullup_en** (*gpio_num_t gpio_num*)

Enable pull-up on GPIO.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number

esp_err_t **gpio_pullup_dis** (*gpio_num_t gpio_num*)

Disable pull-up on GPIO.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number

esp_err_t **gpio_pulldown_en** (*gpio_num_t gpio_num*)

Enable pull-down on GPIO.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number

esp_err_t **gpio_pulldown_dis** (*gpio_num_t gpio_num*)

Disable pull-down on GPIO.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number

esp_err_t **gpio_install_isr_service** (int *intr_alloc_flags*)

Install the driver’ s GPIO ISR handler service, which allows per-pin GPIO interrupt handlers.

This function is incompatible with gpio_isr_register() - if that function is used, a single global ISR is registered for all GPIO interrupts. If this function is used, the ISR service provides a global GPIO ISR and individual pin handlers are registered via the gpio_isr_handler_add() function.

Return

- ESP_OK Success
- ESP_ERR_NO_MEM No memory to install this service
- ESP_ERR_INVALID_STATE ISR service already installed.
- ESP_ERR_NOT_FOUND No free interrupt found with the specified flags

- ESP_ERR_INVALID_ARG GPIO error

Parameters

- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.

void **gpio_uninstall_isr_service**(void)

Uninstall the driver's GPIO ISR service, freeing related resources.

esp_err_t **gpio_isr_handler_add**(*gpio_num_t gpio_num*, *gpio_isr_t isr_handler*, void **args*)

Add ISR handler for the corresponding GPIO pin.

Call this function after using gpio_install_isr_service() to install the driver's GPIO ISR handler service.

The pin ISR handlers no longer need to be declared with IRAM_ATTR, unless you pass the ESP_INTR_FLAG_IRAM flag when allocating the ISR in gpio_install_isr_service().

This ISR handler will be called from an ISR. So there is a stack size limit (configurable as "ISR stack size" in menuconfig). This limit is smaller compared to a global GPIO interrupt handler due to the additional level of indirection.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE Wrong state, the ISR service has not been initialized.
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number
- `isr_handler`: ISR handler function for the corresponding GPIO number.
- `args`: parameter for ISR handler.

esp_err_t **gpio_isr_handler_remove**(*gpio_num_t gpio_num*)

Remove ISR handler for the corresponding GPIO pin.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE Wrong state, the ISR service has not been initialized.
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number

esp_err_t **gpio_set_drive_capability**(*gpio_num_t gpio_num*, *gpio_drive_cap_t strength*)

Set GPIO pad drive capability.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number, only support output GPIOs
- `strength`: Drive capability of the pad

esp_err_t **gpio_get_drive_capability**(*gpio_num_t gpio_num*, *gpio_drive_cap_t *strength*)

Get GPIO pad drive capability.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `gpio_num`: GPIO number, only support output GPIOs
- `strength`: Pointer to accept drive capability of the pad

esp_err_t **gpio_hold_en**(*gpio_num_t gpio_num*)

Enable gpio pad hold function.

The gpio pad hold function works in both input and output modes, but must be output-capable gpios. If pad hold enabled: in output mode: the output level of the pad will be force locked and can not be changed. in input mode: the input value read will not change, regardless the changes of input signal.

The state of digital gpio cannot be held during Deep-sleep, and it will resume the hold function when the chip wakes up from Deep-sleep. If the digital gpio also needs to be held during Deep-sleep, `gpio_deep_sleep_hold_en` should also be called.

Power down or call gpio_hold_dis will disable this function.

Return
- ESP_OK Success
- ESP_ERR_NOT_SUPPORTED Not support pad hold function

Parameters
- `gpio_num`: GPIO number, only support output-capable GPIOs

esp_err_t **gpio_hold_dis** (*gpio_num_t gpio_num*)

Disable gpio pad hold function.

When the chip is woken up from Deep-sleep, the gpio will be set to the default mode, so, the gpio will output the default level if this function is called. If you don't want the level changes, the gpio should be configured to a known state before this function is called. e.g. If you hold gpio18 high during Deep-sleep, after the chip is woken up and `gpio_hold_dis` is called, gpio18 will output low level(because gpio18 is input mode by default). If you don't want this behavior, you should configure gpio18 as output mode and set it to hight level before calling `gpio_hold_dis`.

Return
- ESP_OK Success
- ESP_ERR_NOT_SUPPORTED Not support pad hold function

Parameters
- `gpio_num`: GPIO number, only support output-capable GPIOs

void **gpio_deep_sleep_hold_en** (void)

Enable all digital gpio pad hold function during Deep-sleep.

When the chip is in Deep-sleep mode, all digital gpio will hold the state before sleep, and when the chip is woken up, the status of digital gpio will not be held. Note that the pad hold feature only works when the chip is in Deep-sleep mode, when not in sleep mode, the digital gpio state can be changed even you have called this function.

Power down or call gpio_hold_dis will disable this function, otherwise, the digital gpio hold feature works as long as the chip enter Deep-sleep.

void **gpio_deep_sleep_hold_dis** (void)

Disable all digital gpio pad hold function during Deep-sleep.

void **gpio_iomux_in** (uint32_t *gpio_num*, uint32_t *signal_idx*)

Set pad input to a peripheral signal through the IOMUX.

Parameters
- `gpio_num`: GPIO number of the pad.
- `signal_idx`: Peripheral signal id to input. One of the `*_IN_IDX` signals in `soc/gpio_sig_map.h`.

void **gpio_iomux_out** (uint8_t *gpio_num*, int *func*, bool *oen_inv*)

Set peripheral output to an GPIO pad through the IOMUX.

Parameters
- `gpio_num`: gpio_num GPIO number of the pad.
- `func`: The function number of the peripheral pin to output pin. One of the `FUNC_X_*` of specified pin (X) in `soc/io_mux_reg.h`.
- `oen_inv`: True if the output enable needs to be inverted, otherwise False.

esp_err_t **gpio_force_hold_all** (void)

Force hold digital and rtc gpio pad.

Note GPIO force hold, whether the chip in sleep mode or wakeup mode.

esp_err_t **gpio_force_unhold_all** (void)

Force unhold digital and rtc gpio pad.

Note GPIO force unhold, whether the chip in sleep mode or wakeup mode.

Type Definitions

typedef *intr_handle_t* **gpio_isr_handle_t**

Header File

- soc/include/hal/gpio_types.h

Structures

struct gpio_config_t

Configuration parameters of GPIO pad for gpio_config function.

Public Members

uint64_t **pin_bit_mask**

GPIO pin: set with bit mask, each bit maps to a GPIO

gpio_mode_t **mode**

GPIO mode: set input/output mode

gpio_pullup_t **pull_up_en**

GPIO pull-up

gpio_pulldown_t **pull_down_en**

GPIO pull-down

gpio_int_type_t **intr_type**

GPIO interrupt type

Type Definitions

typedef void (***gpio_isr_t**)(void *)

Enumerations

enum gpio_port_t

Values:

GPIO_PORT_0 = 0

GPIO_PORT_MAX

enum gpio_num_t

Values:

GPIO_NUM_NC = -1

Use to signal not connected to S/W

GPIO_NUM_0 = 0

GPIO0, input and output

GPIO_NUM_1 = 1

GPIO1, input and output

GPIO_NUM_2 = 2

GPIO2, input and output

GPIO_NUM_3 = 3

GPIO3, input and output

GPIO_NUM_4 = 4

GPIO4, input and output

GPIO_NUM_5 = 5
GPIO5, input and output

GPIO_NUM_6 = 6
GPIO6, input and output

GPIO_NUM_7 = 7
GPIO7, input and output

GPIO_NUM_8 = 8
GPIO8, input and output

GPIO_NUM_9 = 9
GPIO9, input and output

GPIO_NUM_10 = 10
GPIO10, input and output

GPIO_NUM_11 = 11
GPIO11, input and output

GPIO_NUM_12 = 12
GPIO12, input and output

GPIO_NUM_13 = 13
GPIO13, input and output

GPIO_NUM_14 = 14
GPIO14, input and output

GPIO_NUM_15 = 15
GPIO15, input and output

GPIO_NUM_16 = 16
GPIO16, input and output

GPIO_NUM_17 = 17
GPIO17, input and output

GPIO_NUM_18 = 18
GPIO18, input and output

GPIO_NUM_19 = 19
GPIO19, input and output

GPIO_NUM_20 = 20
GPIO20, input and output

GPIO_NUM_21 = 21
GPIO21, input and output

GPIO_NUM_26 = 26
GPIO26, input and output

GPIO_NUM_27 = 27
GPIO27, input and output

GPIO_NUM_28 = 28
GPIO28, input and output

GPIO_NUM_29 = 29
GPIO29, input and output

GPIO_NUM_30 = 30
GPIO30, input and output

GPIO_NUM_31 = 31
GPIO31, input and output

GPIO_NUM_32 = 32
GPIO32, input and output

GPIO_NUM_33 = 33
GPIO33, input and output

GPIO_NUM_34 = 34
GPIO34, input mode only(ESP32) / input and output(ESP32-S2)

GPIO_NUM_35 = 35
GPIO35, input mode only(ESP32) / input and output(ESP32-S2)

GPIO_NUM_36 = 36
GPIO36, input mode only(ESP32) / input and output(ESP32-S2)

GPIO_NUM_37 = 37
GPIO37, input mode only(ESP32) / input and output(ESP32-S2)

GPIO_NUM_38 = 38
GPIO38, input mode only(ESP32) / input and output(ESP32-S2)

GPIO_NUM_39 = 39
GPIO39, input mode only(ESP32) / input and output(ESP32-S2)

GPIO_NUM_40 = 40
GPIO40, input and output

GPIO_NUM_41 = 41
GPIO41, input and output

GPIO_NUM_42 = 42
GPIO42, input and output

GPIO_NUM_43 = 43
GPIO43, input and output

GPIO_NUM_44 = 44
GPIO44, input and output

GPIO_NUM_45 = 45
GPIO45, input and output

GPIO_NUM_46 = 46
GPIO46, input mode only

GPIO_NUM_MAX

enum gpio_int_type_t
Values:

GPIO_INTR_DISABLE = 0
Disable GPIO interrupt

GPIO_INTR_POSEDGE = 1
GPIO interrupt type : rising edge

GPIO_INTR_NEGEDGE = 2
GPIO interrupt type : falling edge

GPIO_INTR_ANYEDGE = 3
GPIO interrupt type : both rising and falling edge

GPIO_INTR_LOW_LEVEL = 4
GPIO interrupt type : input low level trigger

GPIO_INTR_HIGH_LEVEL = 5
GPIO interrupt type : input high level trigger

GPIO_INTR_MAX

enum gpio_mode_t

Values:

GPIO_MODE_DISABLE = (0)
GPIO mode : disable input and output

GPIO_MODE_INPUT = (BIT0)
GPIO mode : input only

GPIO_MODE_OUTPUT = (BIT1)
GPIO mode : output only mode

GPIO_MODE_OUTPUT_OD = (((BIT1)) | ((BIT2)))
GPIO mode : output only with open-drain mode

GPIO_MODE_INPUT_OUTPUT_OD = (((BIT0)) | ((BIT1)) | ((BIT2)))
GPIO mode : output and input with open-drain mode

GPIO_MODE_INPUT_OUTPUT = (((BIT0)) | ((BIT1)))
GPIO mode : output and input mode

enum gpio_pullup_t

Values:

GPIO_PULLUP_DISABLE = 0x0
Disable GPIO pull-up resistor

GPIO_PULLUP_ENABLE = 0x1
Enable GPIO pull-up resistor

enum gpio_pulldown_t

Values:

GPIO_PULLDOWN_DISABLE = 0x0
Disable GPIO pull-down resistor

GPIO_PULLDOWN_ENABLE = 0x1
Enable GPIO pull-down resistor

enum gpio_pull_mode_t

Values:

GPIO_PULLUP_ONLY
Pad pull up

GPIO_PULLDOWN_ONLY
Pad pull down

GPIO_PULLUP_PULLDOWN
Pad pull up + pull down

GPIO_FLOATING
Pad floating

enum gpio_drive_cap_t

Values:

GPIO_DRIVE_CAP_0 = 0
Pad drive capability: weak

GPIO_DRIVE_CAP_1 = 1
Pad drive capability: stronger

GPIO_DRIVE_CAP_2 = 2
Pad drive capability: medium

GPIO_DRIVE_CAP_DEFAULT = 2
Pad drive capability: medium

GPIO_DRIVE_CAP_3 = 3
Pad drive capability: strongest

GPIO_DRIVE_CAP_MAX

API Reference - RTC GPIO

Header File

- driver/include/driver/rtc_io.h

Functions

static bool **rtc_gpio_is_valid_gpio** (*gpio_num_t gpio_num*)
Determine if the specified GPIO is a valid RTC GPIO.

Return true if GPIO is valid for RTC GPIO use. false otherwise.
Parameters
- gpio_num: GPIO number

static int **rtc_io_number_get** (*gpio_num_t gpio_num*)
Get RTC IO index number by gpio number.

Return >=0: Index of rtcio. -1 : The gpio is not rtcio.
Parameters
- gpio_num: GPIO number

esp_err_t **rtc_gpio_init** (*gpio_num_t gpio_num*)
Init a GPIO as RTC GPIO.

This function must be called when initializing a pad for an analog function.

Return
- ESP_OK success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- gpio_num: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_deinit** (*gpio_num_t gpio_num*)
Init a GPIO as digital GPIO.

Return
- ESP_OK success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- gpio_num: GPIO number (e.g. GPIO_NUM_12)

uint32_t **rtc_gpio_get_level** (*gpio_num_t gpio_num*)
Get the RTC IO input level.

Return
- 1 High level
- 0 Low level
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- gpio_num: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_set_level** (*gpio_num_t gpio_num*, uint32_t *level*)
Set the RTC IO output level.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- gpio_num: GPIO number (e.g. GPIO_NUM_12)

- `level`: output level

esp_err_t **rtc_gpio_set_direction** (*gpio_num_t gpio_num*, *rtc_gpio_mode_t mode*)

RTC GPIO set direction.

Configure RTC GPIO direction, such as output only, input only, output and input.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)
- `mode`: GPIO direction

esp_err_t **rtc_gpio_set_direction_in_sleep** (*gpio_num_t gpio_num*, *rtc_gpio_mode_t mode*)

RTC GPIO set direction in deep sleep mode or disable sleep status (default). In some application scenarios, IO needs to have another states during deep sleep.

NOTE: ESP32 support INPUT_ONLY mode. ESP32S2 support INPUT_ONLY, OUTPUT_ONLY, INPUT_OUTPUT mode.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)
- `mode`: GPIO direction

esp_err_t **rtc_gpio_pullup_en** (*gpio_num_t gpio_num*)

RTC GPIO pullup enable.

This function only works for RTC IOs. In general, call gpio_pullup_en, which will work both for normal GPIOs and RTC IOs.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_pulldown_en** (*gpio_num_t gpio_num*)

RTC GPIO pulldown enable.

This function only works for RTC IOs. In general, call gpio_pulldown_en, which will work both for normal GPIOs and RTC IOs.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_pullup_dis** (*gpio_num_t gpio_num*)

RTC GPIO pullup disable.

This function only works for RTC IOs. In general, call gpio_pullup_dis, which will work both for normal GPIOs and RTC IOs.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_pulldown_dis** (*gpio_num_t gpio_num*)

RTC GPIO pulldown disable.

This function only works for RTC IOs. In general, call gpio_pulldown_dis, which will work both for normal GPIOs and RTC IOs.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_hold_en** (*gpio_num_t gpio_num*)
Enable hold function on an RTC IO pad.

Enabling HOLD function will cause the pad to latch current values of input enable, output enable, output value, function, drive strength values. This function is useful when going into light or deep sleep mode to prevent the pin configuration from changing.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_hold_dis** (*gpio_num_t gpio_num*)
Disable hold function on an RTC IO pad.

Disabling hold function will allow the pad receive the values of input enable, output enable, output value, function, drive strength from RTC_IO peripheral.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12)

esp_err_t **rtc_gpio_isolate** (*gpio_num_t gpio_num*)
Helper function to disconnect internal circuits from an RTC IO This function disables input, output, pullup, pulldown, and enables hold feature for an RTC IO. Use this function if an RTC IO needs to be disconnected from internal circuits in deep sleep, to minimize leakage current.

In particular, for ESP32-WROVER module, call rtc_gpio_isolate(GPIO_NUM_12) before entering deep sleep, to reduce deep sleep current.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if GPIO is not an RTC IO

Parameters
- `gpio_num`: GPIO number (e.g. GPIO_NUM_12).

esp_err_t **rtc_gpio_force_hold_all** (void)
Enable force hold signal for all RTC IOs.

Each RTC pad has a "force hold" input signal from the RTC controller. If this signal is set, pad latches current values of input enable, function, output enable, and other signals which come from the RTC mux. Force hold signal is enabled before going into deep sleep for pins which are used for EXT1 wakeup.

esp_err_t **rtc_gpio_force_hold_dis_all** (void)
Disable force hold signal for all RTC IOs.

esp_err_t **rtc_gpio_set_drive_capability** (*gpio_num_t gpio_num*, *gpio_drive_cap_t strength*)
Set RTC GPIO pad drive capability.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `gpio_num`: GPIO number, only support output GPIOs

- `strength`: Drive capability of the pad

esp_err_t **rtc_gpio_get_drive_capability**(*gpio_num_t gpio_num*, *gpio_drive_cap_t *strength*)
Get RTC GPIO pad drive capability.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `gpio_num`: GPIO number, only support output GPIOs
- `strength`: Pointer to accept drive capability of the pad

esp_err_t **rtc_gpio_wakeup_enable**(*gpio_num_t gpio_num*, *gpio_int_type_t intr_type*)
Enable wakeup from sleep mode using specific GPIO.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if gpio_num is not an RTC IO, or intr_type is not one of GPIO_INTR_HIGH_LEVEL, GPIO_INTR_LOW_LEVEL.

Parameters
- `gpio_num`: GPIO number
- `intr_type`: Wakeup on high level (GPIO_INTR_HIGH_LEVEL) or low level (GPIO_INTR_LOW_LEVEL)

esp_err_t **rtc_gpio_wakeup_disable**(*gpio_num_t gpio_num*)
Disable wakeup from sleep mode using specific GPIO.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if gpio_num is not an RTC IO

Parameters
- `gpio_num`: GPIO number

Macros

RTC_GPIO_IS_VALID_GPIO(gpio_num)

Header File

- soc/include/hal/rtc_io_types.h

Enumerations

enum rtc_gpio_mode_t
RTCIO output/input mode type.

Values:

RTC_GPIO_MODE_INPUT_ONLY
Pad input

RTC_GPIO_MODE_OUTPUT_ONLY
Pad output

RTC_GPIO_MODE_INPUT_OUTPUT
Pad input + output

RTC_GPIO_MODE_DISABLED
Pad (output + input) disable

RTC_GPIO_MODE_OUTPUT_OD
Pad open-drain output

RTC_GPIO_MODE_INPUT_OUTPUT_OD
Pad input + open-drain output

2.2.4 HMAC

The HMAC (Hash-based Message Authentication Code) module provides hardware acceleration for SHA256-HMAC generation using a key burned into an eFuse block. HMACs work with pre-shared secret keys and provide authenticity and integrity to a message.

Look into the ESP32-S2 Technical Reference Manual (PDF) for more detailed information about the application workflow and the HMAC calculation process.

Generalized Application Scheme

Let there be two parties, A and B. They want to verify the authenticity and integrity of messages sent between each other. Before they can start sending messages, they need to exchange the secret key via a secure channel. To verify A's messages, B can do the following:

- A calculates the HMAC of the message it wants to send.
- A sends the message and the HMAC to B.
- B calculates HMAC of the received message itself.
- B checks wether the received and calculated HMACs match. If they do match, the message is authentic.

However, the HMAC itself isn't bound to this use case. It can also be used for challenge-response protocols supporting HMAC or as a key input for further security modules (see below), etc.

HMAC on the ESP32-S2

On the ESP32-S2, the HMAC module works with a secret key burnt into the eFuses. This eFuse key can be made completely inaccessible for any resources outside the cryptographic modules, thus avoiding key leakage.

Furthermore, the ESP32-S2 has three different application scenarios for its HMAC module:

1. HMAC is generated for software use
2. HMAC is used as a key for the Digital Signature (DS) module
3. HMAC is used for enabling the soft-disabled JTAG interface

The first mode is also called *Upstream* mode, while the last two modes are also called *Downstream* modes.

eFuse Keys for HMAC Six physical eFuse blocks can be used as keys for the HMAC module: block 4 up to block 9. The enum `hmac_key_id_t` in the API maps them to *HMAC_KEY0* ⋯*HMAC_KEY5*. Each key has a corresponding eFuse parameter *key purpose* determining for which of the three HMAC application scenarios (see below) the key may be used:

Key Purpose	Application Scenario
8	HMAC generated for software use
7	HMAC used as a key for the Digital Signature (DS) module
6	HMAC used for enabling the soft-disabled JTAG interface
5	HMAC both as a key for the DS module and for enabling JTAG

This is to prevent the usage of a key for a different function than originally intended.

To calculate an HMAC, the software has to provide the ID of the key block containing the secret key as well as the *key purpose* (see chapter *eFuse Controller* in the ESP32-S2 Technical Reference Manual). Before the HMAC key calculation, the HMAC module looks up the purpose of the provided key block. The calculation only proceeds if the provided key purpose matches the purpose stored in the eFuses of the key block provided by the ID.

HMAC Generation for Software Key Purpose value: 8

In this case, the HMAC is given out to the software (e.g. to authenticate a message).

The API to calculate the HMAC is *esp_hmac_calculate()*. Only the message, message length and the eFuse key block ID have to be provided to that function. The rest, like setting the key purpose, is done automatically.

HMAC for Digital Signature Key Purpose values: 7, 5

The HMAC can be used as a key derivation function to decrypt private key parameters which are used by the Digital Signature module. A standard message is used by the hardware in that case. The user only needs to provide the eFuse key block and purpose on the HMAC side (additional parameters are required for the Digital Signature component in that case). Neither the key nor the actual HMAC are ever exposed to outside the HMAC module and DS component. The calculation of the HMAC and its hand-over to the DS component happen internally.

For more details, check the chapter *Digital Signature* in the ESP32-S2 Technical Reference Manual.

HMAC for Enabling JTAG Key Purpose values: 6, 5

The third application is using the HMAC as a key to enable JTAG if it was soft-disabled before. This functionality is currently not implemented.

Application Outline

Following code is an outline of how to set an eFuse key and then use it to calculate an HMAC for software usage. We use *ets_efuse_write_key* to set physical key block 4 in the eFuse for the HMAC module together with its purpose. *ETS_EFUSE_KEY_PURPOSE_HMAC_UP* (8) means that this key can only be used for HMAC generation for software usage:

```
#include "esp32s2/rom/efuse.h"

const uint8_t key_data[32] = { ... };

int ets_status = ets_efuse_write_key(ETS_EFUSE_BLOCK_KEY4,
                    ETS_EFUSE_KEY_PURPOSE_HMAC_UP,
                    key_data, sizeof(key_data));

if (ets_status == ESP_OK) {
    // written key
} else {
    // writing key failed, maybe written already
}
```

Now we can use the saved key to calculate an HMAC for software usage.

```
#include "esp_hmac.h"

uint8_t hmac[32];

const char *message = "Hello, HMAC!";
const size_t msg_len = 12;

esp_err_t result = esp_hmac_calculate(HMAC_KEY4, message, msg_len, hmac);

if (result == ESP_OK) {
    // HMAC written to hmac now
} else {
    // failure calculating HMAC
}
```

API Reference

Header File

- esp32s2/include/esp_hmac.h

Functions

esp_err_t **esp_hmac_calculate**(*hmac_key_id_t key_id*, **const** void **message*, size_t *message_len*, uint8_t **hmac*)

Calculate the HMAC of a given message.

Calculate the HMAC `hmac` of a given message `message` with length `message_len`. SHA256 is used for the calculation (fixed on ESP32S2).

Note Uses the HMAC peripheral in "upstream" mode.

Return

- ESP_OK, if the calculation was successful,
- ESP_FAIL, if the hmac calculation failed

Parameters

- `key_id`: Determines which of the 6 key blocks in the efuses should be used for the HMAC calcuation. The corresponding purpose field of the key block in the efuse must be set to the HMAC upstream purpose value.
- `message`: the message for which to calculate the HMAC
- `message_len`: message length return ESP_ERR_INVALID_STATE if unsuccessful
- `[out] hmac`: the hmac result; the buffer behind the provided pointer must be 32 bytes long

Enumerations

enum hmac_key_id_t

The possible efuse keys for the HMAC peripheral

Values:

HMAC_KEY0 = 0

HMAC_KEY1

HMAC_KEY2

HMAC_KEY3

HMAC_KEY4

HMAC_KEY5

HMAC_KEY_MAX

2.2.5 Digital Signature

The Digital Signature (DS) module provides hardware acceleration of signing messages based on RSA. It uses pre-encrypted parameters to calculate a signature. The parameters are encrypted using HMAC as a key-derivation function. In turn, the HMAC uses eFuses as input key. The whole process happens in hardware so that neither the decryption key for the RSA parameters nor the input key for the HMAC key derivation function can be seen by the software while calculating the signature.

Look into the ESP32-S2 Technical Reference Manual (PDF) for more detailed information about the involved hardware during the signature calculation process and the used registers.

Private Key Parameters

The private key parameters for the RSA signature are stored in flash. To prevent unauthorized access, they are AES-encrypted. The HMAC module is used as a key-derivation function to calculate the AES encryption key for the private key parameters. In turn, the HMAC module uses a key from the eFuses key block which can be read-protected to prevent unauthorized access as well.

Upon signature calculation invocation, the software only specifies which eFuse key to use, the corresponding eFuse key purpose, the location of the encrypted RSA parameters and the message.

Key Generation

Both the HMAC key and the RSA private key have to be created and stored before the DS module can be used. This needs to be done in software on the ESP32-S2 or alternatively on a host. For this context, the IDF provides `esp_efuse_write_block()` to set the HMAC key and `esp_hmac_calculate()` to encrypt the private RSA key parameters.

Instructions on how to calculate and assemble the private key parameters are described in the ESP32-S2 Technical Reference Manual.

Signature Calculation with IDF

For thorough information about involved registers and the workflow, please have a look at the ESP32-S2 Technical Reference Manual.

Three parameters need to be prepared to calculate the digital signature:

1. the eFuse key block ID which is used as key for the HMAC,
2. the location of the encrypted private key parameters,
3. and the message to be signed.

Since the signature calculation takes some time, there are two possible API versions to use in IDF. The first one is `esp_ds_sign()` and simply blocks until the calculation is finished. If software needs to do something else during the calculation, `esp_ds_start_sign()` can be called, followed by periodic calls to `esp_ds_is_busy()` to check when the calculation has finished. Once the calculation has finished, `esp_ds_finish_sign()` can be called to get the resulting signature.

注解： Note that this is only the basic DS building block, the message length is fixed. To create signatures of arbitrary messages, the input is normally a hash of the actual message, padded up to the required length. An API to do this is planned in the future.

API Reference

Header File

- esp32s2/include/esp_ds.h

Functions

esp_err_t **esp_ds_sign**(**const** void **message*, **const** *esp_ds_data_t* **data*, *hmac_key_id_t key_id*, void **signature*)

Sign the message.

This function is a wrapper around `esp_ds_finish_sign()` and `esp_ds_start_sign()`, so do not use them in parallel. It blocks until the signing is finished and then returns the signature.

Note This function locks the HMAC, SHA, AES and RSA components during its entire execution time.

Return

- ESP_OK if successful, the signature was written to the parameter `signature`.
- ESP_ERR_INVALID_ARG if one of the parameters is NULL or data->rsa_length is too long or 0
- ESP_ERR_HW_CRYPTO_DS_HMAC_FAIL if there was an HMAC failure during retrieval of the decryption key
- ESP_ERR_NO_MEM if there hasn't been enough memory to allocate the context object
- ESP_ERR_HW_CRYPTO_DS_INVALID_KEY if there's a problem with passing the HMAC key to the DS component

- ESP_ERR_HW_CRYPTO_DS_INVALID_DIGEST if the message digest didn't match; the signature is invalid.
- ESP_ERR_HW_CRYPTO_DS_INVALID_PADDING if the message padding is incorrect, the signature can be read though since the message digest matches.

Parameters

- `message`: the message to be signed; its length is determined by data->rsa_length
- `data`: the encrypted signing key data (AES encrypted RSA key + IV)
- `key_id`: the HMAC key ID determining the HMAC key of the HMAC which will be used to decrypt the signing key data
- `signature`: the destination of the signature, should be (data->rsa_length + 1)*4 bytes long

esp_err_t **esp_ds_start_sign**(**const** void **message*, **const** *esp_ds_data_t* **data*, *hmac_key_id_t key_id*, *esp_ds_context_t* ***esp_ds_ctx*)

Start the signing process.

This function yields a context object which needs to be passed to `esp_ds_finish_sign()` to finish the signing process.

Note This function locks the HMAC, SHA, AES and RSA components, so the user has to ensure to call `esp_ds_finish_sign()` in a timely manner.

Return

- ESP_OK if successful, the ds operation was started now and has to be finished with `esp_ds_finish_sign()`
- ESP_ERR_INVALID_ARG if one of the parameters is NULL or data->rsa_length is too long or 0
- ESP_ERR_HW_CRYPTO_DS_HMAC_FAIL if there was an HMAC failure during retrieval of the decryption key
- ESP_ERR_NO_MEM if there hasn't been enough memory to allocate the context object
- ESP_ERR_HW_CRYPTO_DS_INVALID_KEY if there's a problem with passing the HMAC key to the DS component

Parameters

- `message`: the message to be signed; its length is determined by data->rsa_length
- `data`: the encrypted signing key data (AES encrypted RSA key + IV)
- `key_id`: the HMAC key ID determining the HMAC key of the HMAC which will be used to decrypt the signing key data
- `esp_ds_ctx`: the context object which is needed for finishing the signing process later

bool **esp_ds_is_busy**(void)

Return true if the DS peripheral is busy, otherwise false.

Note Only valid if `esp_ds_start_sign()` was called before.

esp_err_t **esp_ds_finish_sign**(void **signature*, *esp_ds_context_t* **esp_ds_ctx*)

Finish the signing process.

Return

- ESP_OK if successful, the ds operation has been finished and the result is written to signature.
- ESP_ERR_INVALID_ARG if one of the parameters is NULL
- ESP_ERR_HW_CRYPTO_DS_INVALID_DIGEST if the message digest didn't match; the signature is invalid.
- ESP_ERR_HW_CRYPTO_DS_INVALID_PADDING if the message padding is incorrect, the signature can be read though since the message digest matches.

Parameters

- `signature`: the destination of the signature, should be (data->rsa_length + 1)*4 bytes long
- `esp_ds_ctx`: the context object retreived by `esp_ds_start_sign()`

esp_err_t **esp_ds_encrypt_params**(*esp_ds_data_t* **data*, **const** void **iv*, **const** *esp_ds_p_data_t* **p_data*, **const** void **key*)

Encrypt the private key parameters.

Return

- ESP_OK if successful, the ds operation has been finished and the result is written to signature.
- ESP_ERR_INVALID_ARG if one of the parameters is NULL or p_data->rsa_length is too long

Parameters

- `data`: Output buffer to store encrypted data, suitable for later use generating signatures. The allocated memory must be in internal memory and word aligned since it's filled by DMA. Both is asserted at run time.
- `iv`: Pointer to 16 byte IV buffer, will be copied into 'data'. Should be randomly generated bytes each time.
- `p_data`: Pointer to input plaintext key data. The expectation is this data will be deleted after this process is done and 'data' is stored.
- `key`: Pointer to 32 bytes of key data. Type determined by key_type parameter. The expectation is the corresponding HMAC key will be stored to efuse and then permanently erased.

Structures

struct esp_digital_signature_data

Encrypted private key data. Recommended to store in flash in this format.

Note This struct has to match to one from the ROM code! This documentation is mostly taken from there.

Public Members

esp_digital_signature_length_t **rsa_length**

RSA LENGTH register parameters (number of words in RSA key & operands, minus one).

Max value 127 (for RSA 4096).

This value must match the length field encrypted and stored in 'c', or invalid results will be returned. (The DS peripheral will always use the value in 'c', not this value, so an attacker can't alter the DS peripheral results this way, it will just truncate or extend the message and the resulting signature in software.)

Note In IDF, the enum type length is the same as of type unsigned, so they can be used interchangably. See the ROM code for the original declaration of struct `ets_ds_data_t`.

uint8_t **iv**[**ESP_DS_IV_LEN**]

IV value used to encrypt 'c'

uint8_t **c**[**ESP_DS_C_LEN**]

Encrypted Digital Signature parameters. Result of AES-CBC encryption of plaintext values. Includes an encrypted message digest.

struct esp_ds_p_data_t

Plaintext parameters used by Digital Signature.

Not used for signing with DS peripheral, but can be encrypted in-device by calling esp_ds_encrypt_params()

Note This documentation is mostly taken from the ROM code.

Public Members

uint32_t **Y**[4096 / 32]

RSA exponent.

uint32_t **M**[4096 / 32]

RSA modulus.

uint32_t **Rb**[4096 / 32]

RSA r inverse operand.

uint32_t **M_prime**

RSA M prime operand.

esp_digital_signature_length_t **length**

RSA length.

Macros

ESP_ERR_HW_CRYPTO_DS_BASE
Starting number of HW cryptography module error codes

ESP_ERR_HW_CRYPTO_DS_HMAC_FAIL
HMAC peripheral problem

ESP_ERR_HW_CRYPTO_DS_INVALID_KEY
given HMAC key isn’ t correct, HMAC peripheral problem

ESP_ERR_HW_CRYPTO_DS_INVALID_DIGEST
message digest check failed, result is invalid

ESP_ERR_HW_CRYPTO_DS_INVALID_PADDING
padding check failed, but result is produced anyway and can be read

ESP_DS_IV_LEN

ESP_DS_C_LEN

Type Definitions

typedef struct esp_ds_context **esp_ds_context_t**
typedef struct *esp_digital_signature_data* **esp_ds_data_t**
Encrypted private key data. Recommended to store in flash in this format.

Note This struct has to match to one from the ROM code! This documentation is mostly taken from there.

Enumerations

enum esp_digital_signature_length_t
Values:

ESP_DS_RSA_1024 = (1024 / 32) - 1

ESP_DS_RSA_2048 = (2048 / 32) - 1

ESP_DS_RSA_3072 = (3072 / 32) - 1

ESP_DS_RSA_4096 = (4096 / 32) - 1

2.2.6 I2C Driver

Overview

I2C is a serial, synchronous, half-duplex communication protocol that allows co-existence of multiple masters and slaves on the same bus. The I2C bus consists of two lines: serial data line (SDA) and serial clock (SCL). Both lines require pull-up resistors.

With such advantages as simplicity and low manufacturing cost, I2C is mostly used for communication of low-speed peripheral devices over short distances (within one foot).

ESP32-S2 has two I2C controllers (also referred to as ports) which are responsible for handling communications on two I2C buses. Each I2C controller can operate as master or slave. As an example, one controller can act as a master and the other as a slave at the same time.

Driver Features

I2C driver governs communications of devices over the I2C bus. The driver supports the following features:

- Reading and writing bytes in Master mode
- Slave mode
- Reading and writing to registers which are in turn read/written by the master

Driver Usage

The following sections describe typical steps of configuring and operating the I2C driver:

1. *Configuration* - set the initialization parameters (master or slave mode, GPIO pins for SDA and SCL, clock speed, etc.)
2. *Install Driver*- activate the driver on one of the two I2C controllers as a master or slave
3. Depending on whether you configure the driver for a master or slave, choose the appropriate item
   a) *Communication as Master* - handle communications (master)
   b) *Communication as Slave* - respond to messages from the master (slave)
4. *Interrupt Handling* - configure and service I2C interrupts
5. *Customized Configuration* - adjust default I2C communication parameters (timings, bit order, etc.)
6. *Error Handling* - how to recognize and handle driver configuration and communication errors
7. *Delete Driver*- release resources used by the I2C driver when communication ends

Configuration To establish I2C communication, start by configuring the driver. This is done by setting the parameters of the structure `i2c_config_t`:

- Set I2C **mode of operation** - slave or master from `i2c_mode_t`
- Configure **communication pins**
  - Assign GPIO pins for SDA and SCL signals
  - Set whether to enable ESP32-S2’s internal pull-ups
- (Master only) Set I2C **clock speed**
- (Slave only) Configure the following
  - Whether to enable **10 bit address mode**
  - Define **slave address**

After that, initialize the configuration for a given I2C port. For this, call the function `i2c_param_config()` and pass to it the port number and the structure `i2c_config_t`.

At this stage, `i2c_param_config()` also sets a few other I2C configuration parameters to default values that are defined by the I2C specification. For more details on the values and how to modify them, see *Customized Configuration*.

Install Driver After the I2C driver is configured, install it by calling the function `i2c_driver_install()` with the following parameters:

- Port number, one of the two port numbers from `i2c_port_t`
- Master or slave, selected from `i2c_mode_t`
- (Slave only) Size of buffers to allocate for sending and receiving data. As I2C is a master-centric bus, data can only go from the slave to the master at the master’s request. Therefore, the slave will usually have a send buffer where the slave application writes data. The data remains in the send buffer to be read by the master at the master’s own discretion.
- Flags for allocating the interrupt (see ESP_INTR_FLAG_* values in esp32/include/esp_intr_alloc.h)

Communication as Master After installing the I2C driver, ESP32-S2 is ready to communicate with other I2C devices.

ESP32-S2’s I2C controller operating as master is responsible for establishing communication with I2C slave devices and sending commands to trigger a slave to action, for example, to take a measurement and send the readings back to the master.

For better process organization, the driver provides a container, called a “command link”, that should be populated with a sequence of commands and then passed to the I2C controller for execution.

Master Write The example below shows how to build a command link for an I2C master to send *n* bytes to a slave.

The following describes how a command link for a “master write” is set up and what comes inside:

1. Create a command link with `i2c_cmd_link_create()`.

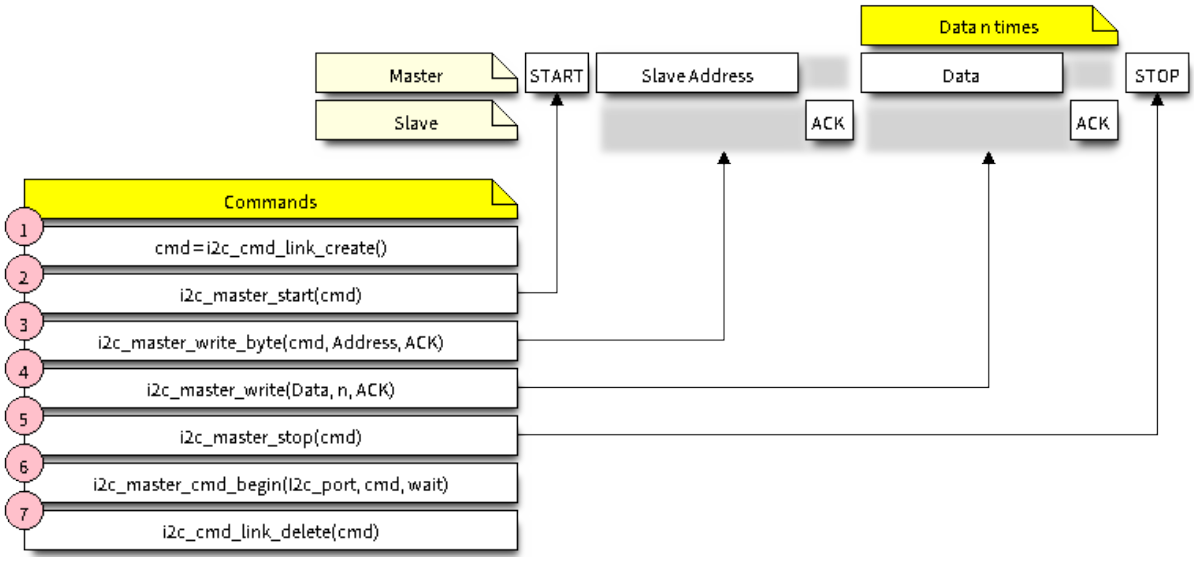


图 5: I2C command link - master write example

Then, populate it with the series of data to be sent to the slave:

a) **Start bit** - *i2c_master_start()*
b) **Slave address** - *i2c_master_write_byte()*. The single byte address is provided as an argument of this function call.
c) **Data** - One or more bytes as an argument of *i2c_master_write()*
d) **Stop bit** - *i2c_master_stop()*

Both functions *i2c_master_write_byte()* and *i2c_master_write()* have an additional argument specifying whether the master should ensure that it has received the ACK bit.

2. Trigger the execution of the command link by I2C controller by calling *i2c_master_cmd_begin()*. Once the execution is triggered, the command link cannot be modified.
3. After the commands are transmitted, release the resources used by the command link by calling *i2c_cmd_link_delete()*.

Master Read The example below shows how to build a command link for an I2C master to read *n* bytes from a slave.

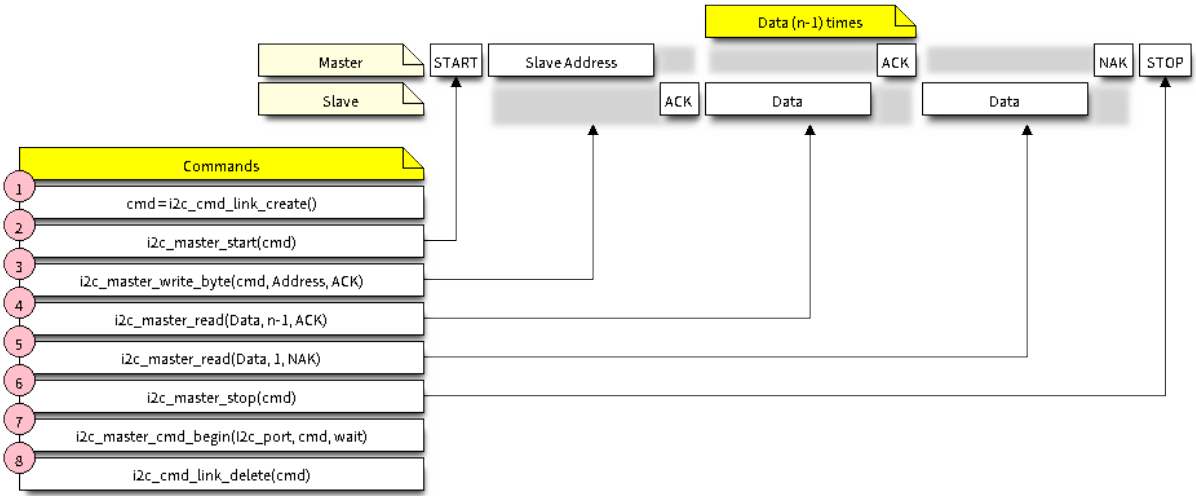


图 6: I2C command link - master read example

Compared to writing data, the command link is populated in Step 4 not with `i2c_master_write...` functions but with *i2c_master_read_byte()* and / or *i2c_master_read()*. Also, the last read in Step 5 is

configured so that the master does not provide the ACK bit.

Indicating Write or Read After sending a slave address (see Step 3 on both diagrams above), the master either writes or reads from the slave.

The information on what the master will actually do is hidden in the least significant bit of the slave's address.

For this reason, the command link sent by the master to write data to the slave contains the address `(ESP_SLAVE_ADDR << 1) | I2C_MASTER_WRITE` and looks as follows:

```
i2c_master_write_byte(cmd, (ESP_SLAVE_ADDR << 1) | I2C_MASTER_WRITE, ACK_EN);
```

Likewise, the command link to read from the slave looks as follows:

```
i2c_master_write_byte(cmd, (ESP_SLAVE_ADDR << 1) | I2C_MASTER_READ, ACK_EN);
```

Communication as Slave After installing the I2C driver, ESP32-S2 is ready to communicate with other I2C devices.

The API provides the following functions for slaves

- *i2c_slave_read_buffer()*
  Whenever the master writes data to the slave, the slave will automatically store it in the receive buffer. This allows the slave application to call the function *i2c_slave_read_buffer()* at its own discretion. This function also has a parameter to specify block time if no data is in the receive buffer. This will allow the slave application to wait with a specified timeout for data to arrive to the buffer.
- *i2c_slave_write_buffer()*
  The send buffer is used to store all the data that the slave wants to send to the master in FIFO order. The data stays there until the master requests for it. The function *i2c_slave_write_buffer()* has a parameter to specify block time if the send buffer is full. This will allow the slave application to wait with a specified timeout for the adequate amount of space to become available in the send buffer.

A code example showing how to use these functions can be found in peripherals/i2c.

Interrupt Handling During driver installation, an interrupt handler is installed by default. However, you can register your own interrupt handler instead of the default one by calling the function *i2c_isr_register()*. When implementing your own interrupt handler, refer to the ESP32-S2 Technical Reference Manual (PDF) for the description of interrupts triggered by the I2C controller.

To delete an interrupt handler, call *i2c_isr_free()*.

Customized Configuration As mentioned at the end of Section *Configuration*, when the function *i2c_param_config()* initializes the driver configuration for an I2C port, it also sets several I2C communication parameters to default values defined in the I2C specification. Some other related parameters are pre-configured in registers of the I2C controller.

All these parameters can be changed to user-defined values by calling dedicated functions given in the table below. Please note that the timing values are defined in APB clock cycles. The frequency of APB is specified in `I2C_APB_CLK_FREQ`.

表 1: Other Configurable I2C Communication Parameters

Parameters to Change	Function
High time and low time for SCL pulses	*i2c_set_period()*
SCL and SDA signal timing used during generation of **start** signals	*i2c_set_start_timing()*
SCL and SDA signal timing used during generation of **stop** signals	*i2c_set_stop_timing()*
Timing relationship between SCL and SDA signals when slave samples, as well as when master toggles	*i2c_set_data_timing()*
I2C timeout	*i2c_set_timeout()*
Choice between transmitting / receiving the LSB or MSB first, choose one of the modes defined in *i2c_trans_mode_t*	*i2c_set_data_mode()*

Each of the above functions has a *_get_* counterpart to check the currently set value. For example, to check the I2C timeout value, call *i2c_get_timeout()*.

To check the default parameter values which are set during the driver configuration process, please refer to the file driver/i2c.c and look for defines with the suffix `_DEFAULT`.

You can also select different pins for SDA and SCL signals and alter the configuration of pull-ups with the function *i2c_set_pin()*. If you want to modify already entered values, use the function *i2c_param_config()*.

注解： ESP32-S2’s internal pull-ups are in the range of tens of kOhm, which is, in most cases, insufficient for use as I2C pull-ups. Users are advised to use external pull-ups with values described in the I2C specification.

Error Handling The majority of I2C driver functions either return `ESP_OK` on successful completion or a specific error code on failure. It is a good practice to always check the returned values and implement error handling. The driver also prints out log messages that contain error details, e.g., when checking the validity of entered configuration. For details please refer to the file driver/i2c.c and look for defines with the suffix `_ERR_STR`.

Use dedicated interrupts to capture communication failures. For instance, if a slave stretches the clock for too long while preparing the data to send back to master, the interrupt `I2C_TIME_OUT_INT` will be triggered. For detailed information, see *Interrupt Handling*.

In case of a communication failure, you can reset the internal hardware buffers by calling the functions *i2c_reset_tx_fifo()* and *i2c_reset_rx_fifo()* for the send and receive buffers respectively.

Delete Driver When the I2C communication is established with the function *i2c_driver_install()* and is not required for some substantial amount of time, the driver may be deinitialized to release allocated resources by calling *i2c_driver_delete()*.

Before calling *i2c_driver_delete()* to remove i2c driver, please make sure that all threads have stopped using the driver in any way, because this function does not guarantee thread safety.

Application Example

I2C master and slave example: peripherals/i2c.

API Reference

Header File

- driver/include/driver/i2c.h

Functions

esp_err_t **i2c_driver_install**(*i2c_port_t i2c_num*, *i2c_mode_t mode*, size_t *slv_rx_buf_len*, size_t *slv_tx_buf_len*, int *intr_alloc_flags*)

I2C driver install.

Note Only slave mode will use this value, driver will ignore this value in master mode.

Note Only slave mode will use this value, driver will ignore this value in master mode.

Note In master mode, if the cache is likely to be disabled(such as write flash) and the slave is time-sensitive, `ESP_INTR_FLAG_IRAM` is suggested to be used. In this case, please use the memory allocated from internal RAM in i2c read and write function, because we can not access the psram(if psram is enabled) in interrupt handle function when cache is disabled.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_FAIL Driver install error

Parameters

- `i2c_num`: I2C port number
- `mode`: I2C mode(master or slave)
- `slv_rx_buf_len`: receiving buffer size for slave mode

Parameters

- `slv_tx_buf_len`: sending buffer size for slave mode

Parameters

- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.

esp_err_t **i2c_driver_delete**(*i2c_port_t i2c_num*)

I2C driver delete.

Note This function does not guarantee thread safety. Please make sure that no thread will continuously hold semaphores before calling the delete function.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number

esp_err_t **i2c_param_config**(*i2c_port_t i2c_num*, **const** *i2c_config_t* **i2c_conf*)

I2C parameter initialization.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `i2c_conf`: pointer to I2C parameter settings

esp_err_t **i2c_reset_tx_fifo**(*i2c_port_t i2c_num*)

reset I2C tx hardware fifo

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number

esp_err_t **i2c_reset_rx_fifo**(*i2c_port_t i2c_num*)

reset I2C rx fifo

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number

esp_err_t **i2c_isr_register**(*i2c_port_t i2c_num*, void (**fn*)) void *
, void **arg*, int *intr_alloc_flags*, *intr_handle_t* **handle*I2C isr handler register.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `fn`: isr handler function
- `arg`: parameter for isr handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `handle`: handle return from esp_intr_alloc.

esp_err_t **i2c_isr_free**(*intr_handle_t handle*)
to delete and free I2C isr.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `handle`: handle of isr.

esp_err_t **i2c_set_pin**(*i2c_port_t i2c_num*, int *sda_io_num*, int *scl_io_num*, bool *sda_pullup_en*, bool *scl_pullup_en*, *i2c_mode_t mode*)
Configure GPIO signal for I2C sck and sda.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `sda_io_num`: GPIO number for I2C sda signal
- `scl_io_num`: GPIO number for I2C scl signal
- `sda_pullup_en`: Whether to enable the internal pullup for sda pin
- `scl_pullup_en`: Whether to enable the internal pullup for scl pin
- `mode`: I2C mode

i2c_cmd_handle_t **i2c_cmd_link_create**(void)
Create and init I2C command link.

Note Before we build I2C command link, we need to call i2c_cmd_link_create() to create a command link. After we finish sending the commands, we need to call i2c_cmd_link_delete() to release and return the resources.

Return i2c command link handler

void **i2c_cmd_link_delete**(*i2c_cmd_handle_t cmd_handle*)
Free I2C command link.

Note Before we build I2C command link, we need to call i2c_cmd_link_create() to create a command link. After we finish sending the commands, we need to call i2c_cmd_link_delete() to release and return the resources.

Parameters
- `cmd_handle`: I2C command handle

esp_err_t **i2c_master_start**(*i2c_cmd_handle_t cmd_handle*)
Queue command for I2C master to generate a start signal.

Note Only call this function in I2C master mode Call i2c_master_cmd_begin() to send all queued commands

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `cmd_handle`: I2C cmd link

esp_err_t **i2c_master_write_byte**(*i2c_cmd_handle_t cmd_handle*, uint8_t *data*, bool *ack_en*)

Queue command for I2C master to write one byte to I2C bus.

Note Only call this function in I2C master mode Call i2c_master_cmd_begin() to send all queued commands

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `cmd_handle`: I2C cmd link
- `data`: I2C one byte command to write to bus
- `ack_en`: enable ack check for master

esp_err_t **i2c_master_write**(*i2c_cmd_handle_t cmd_handle*, uint8_t **data*, size_t *data_len*, bool *ack_en*)

Queue command for I2C master to write buffer to I2C bus.

Note Only call this function in I2C master mode Call i2c_master_cmd_begin() to send all queued commands

Note If the psram is enabled and intr_flag is `ESP_INTR_FLAG_IRAM`, please use the memory allocated from internal RAM.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `cmd_handle`: I2C cmd link
- `data`: data to send

Parameters

- `data_len`: data length
- `ack_en`: enable ack check for master

esp_err_t **i2c_master_read_byte**(*i2c_cmd_handle_t* *cmd_handle*, uint8_t **data*, *i2c_ack_type_t ack*)

Queue command for I2C master to read one byte from I2C bus.

Note Only call this function in I2C master mode Call i2c_master_cmd_begin() to send all queued commands

Note If the psram is enabled and intr_flag is `ESP_INTR_FLAG_IRAM`, please use the memory allocated from internal RAM.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `cmd_handle`: I2C cmd link
- `data`: pointer accept the data byte

Parameters

- `ack`: ack value for read command

esp_err_t **i2c_master_read**(*i2c_cmd_handle_t cmd_handle*, uint8_t **data*, size_t *data_len*, *i2c_ack_type_t ack*)

Queue command for I2C master to read data from I2C bus.

Note Only call this function in I2C master mode Call i2c_master_cmd_begin() to send all queued commands

Note If the psram is enabled and intr_flag is `ESP_INTR_FLAG_IRAM`, please use the memory allocated from internal RAM.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `cmd_handle`: I2C cmd link
- `data`: data buffer to accept the data from bus

Parameters

- `data_len`: read data length
- `ack`: ack value for read command

esp_err_t **i2c_master_stop**(*i2c_cmd_handle_t cmd_handle*)

Queue command for I2C master to generate a stop signal.

Note Only call this function in I2C master mode Call i2c_master_cmd_begin() to send all queued commands

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `cmd_handle`: I2C cmd link

esp_err_t **i2c_master_cmd_begin**(*i2c_port_t i2c_num*, *i2c_cmd_handle_t cmd_handle*, TickType_t *ticks_to_wait*)

I2C master send queued commands. This function will trigger sending all queued commands. The task will be blocked until all the commands have been sent out. The I2C APIs are not thread-safe, if you want to use one I2C port in different tasks, you need to take care of the multi-thread issue.

Note Only call this function in I2C master mode

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_FAIL Sending command error, slave doesn't ACK the transfer.
- ESP_ERR_INVALID_STATE I2C driver not installed or not in master mode.
- ESP_ERR_TIMEOUT Operation timeout because the bus is busy.

Parameters

- `i2c_num`: I2C port number
- `cmd_handle`: I2C command handler
- `ticks_to_wait`: maximum wait ticks.

int **i2c_slave_write_buffer**(*i2c_port_t i2c_num*, **const** uint8_t **data*, int *size*, TickType_t *ticks_to_wait*)

I2C slave write data to internal ringbuffer, when tx fifo empty, isr will fill the hardware fifo from the internal ringbuffer.

Note Only call this function in I2C slave mode

Return

- ESP_FAIL(-1) Parameter error
- Others(>=0) The number of data bytes that pushed to the I2C slave buffer.

Parameters

- `i2c_num`: I2C port number
- `data`: data pointer to write into internal buffer
- `size`: data size
- `ticks_to_wait`: Maximum waiting ticks

int **i2c_slave_read_buffer**(*i2c_port_t i2c_num*, uint8_t **data*, size_t *max_size*, TickType_t *ticks_to_wait*)

I2C slave read data from internal buffer. When I2C slave receive data, isr will copy received data from hardware rx fifo to internal ringbuffer. Then users can read from internal ringbuffer.

Note Only call this function in I2C slave mode

Return

- ESP_FAIL(-1) Parameter error
- Others(>=0) The number of data bytes that read from I2C slave buffer.

Parameters

- `i2c_num`: I2C port number
- `data`: data pointer to accept data from internal buffer
- `max_size`: Maximum data size to read
- `ticks_to_wait`: Maximum waiting ticks

esp_err_t **i2c_set_period**(*i2c_port_t i2c_num*, int *high_period*, int *low_period*)

set I2C master clock period

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `high_period`: clock cycle number during SCL is high level, high_period is a 14 bit value
- `low_period`: clock cycle number during SCL is low level, low_period is a 14 bit value

esp_err_t **i2c_get_period**(*i2c_port_t i2c_num*, int **high_period*, int **low_period*)

get I2C master clock period

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `high_period`: pointer to get clock cycle number during SCL is high level, will get a 14 bit value
- `low_period`: pointer to get clock cycle number during SCL is low level, will get a 14 bit value

esp_err_t **i2c_filter_enable**(*i2c_port_t i2c_num*, uint8_t *cyc_num*)

enable hardware filter on I2C bus Sometimes the I2C bus is disturbed by high frequency noise(about 20ns), or the rising edge of the SCL clock is very slow, these may cause the master state machine broken. enable hardware filter can filter out high frequency interference and make the master more stable.

Note Enable filter will slow the SCL clock.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `cyc_num`: the APB cycles need to be filtered(0<= cyc_num <=7). When the period of a pulse is less than cyc_num * APB_cycle, the I2C controller will ignore this pulse.

esp_err_t **i2c_filter_disable**(*i2c_port_t i2c_num*)

disable filter on I2C bus

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number

esp_err_t **i2c_set_start_timing**(*i2c_port_t i2c_num*, int *setup_time*, int *hold_time*)

set I2C master start signal timing

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `setup_time`: clock number between the falling-edge of SDA and rising-edge of SCL for start mark, it’ s a 10-bit value.
- `hold_time`: clock num between the falling-edge of SDA and falling-edge of SCL for start mark, it’ s a 10-bit value.

esp_err_t **i2c_get_start_timing**(*i2c_port_t i2c_num*, int **setup_time*, int **hold_time*)

get I2C master start signal timing

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `setup_time`: pointer to get setup time
- `hold_time`: pointer to get hold time

esp_err_t **i2c_set_stop_timing**(*i2c_port_t i2c_num*, int *setup_time*, int *hold_time*)
set I2C master stop signal timing

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `setup_time`: clock num between the rising-edge of SCL and the rising-edge of SDA, it's a 10-bit value.
- `hold_time`: clock number after the STOP bit's rising-edge, it's a 14-bit value.

esp_err_t **i2c_get_stop_timing**(*i2c_port_t i2c_num*, int **setup_time*, int **hold_time*)
get I2C master stop signal timing

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `setup_time`: pointer to get setup time.
- `hold_time`: pointer to get hold time.

esp_err_t **i2c_set_data_timing**(*i2c_port_t i2c_num*, int *sample_time*, int *hold_time*)
set I2C data signal timing

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `sample_time`: clock number I2C used to sample data on SDA after the rising-edge of SCL, it's a 10-bit value
- `hold_time`: clock number I2C used to hold the data after the falling-edge of SCL, it's a 10-bit value

esp_err_t **i2c_get_data_timing**(*i2c_port_t i2c_num*, int **sample_time*, int **hold_time*)
get I2C data signal timing

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `sample_time`: pointer to get sample time
- `hold_time`: pointer to get hold time

esp_err_t **i2c_set_timeout**(*i2c_port_t i2c_num*, int *timeout*)
set I2C timeout value

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2c_num`: I2C port number
- `timeout`: timeout value for I2C bus (unit: APB 80Mhz clock cycle)

esp_err_t **i2c_get_timeout**(*i2c_port_t i2c_num*, int **timeout*)
get I2C timeout value

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `timeout`: pointer to get timeout value

esp_err_t **i2c_set_data_mode**(*i2c_port_t i2c_num*, *i2c_trans_mode_t tx_trans_mode*, *i2c_trans_mode_t rx_trans_mode*)

set I2C data transfer mode

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `tx_trans_mode`: I2C sending data mode
- `rx_trans_mode`: I2C receving data mode

esp_err_t **i2c_get_data_mode**(*i2c_port_t i2c_num*, *i2c_trans_mode_t *tx_trans_mode*, *i2c_trans_mode_t *rx_trans_mode*)

get I2C data transfer mode

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2c_num`: I2C port number
- `tx_trans_mode`: pointer to get I2C sending data mode
- `rx_trans_mode`: pointer to get I2C receiving data mode

Macros

I2C_APB_CLK_FREQ

I2C source clock is APB clock, 80MHz

I2C_NUM_0

I2C port 0

I2C_NUM_1

I2C port 1

I2C_NUM_MAX

I2C port max

Type Definitions

typedef void ***i2c_cmd_handle_t**

I2C command handle

Header File

- soc/include/hal/i2c_types.h

Structures

struct i2c_config_t

I2C initialization parameters.

Public Members

i2c_mode_t **mode**

I2C mode

int **sda_io_num**

GPIO number for I2C sda signal

int **scl_io_num**
GPIO number for I2C scl signal

bool **sda_pullup_en**
Internal GPIO pull mode for I2C sda signal

bool **scl_pullup_en**
Internal GPIO pull mode for I2C scl signal

uint32_t **clk_speed**
I2C clock frequency for master mode, (no higher than 1MHz for now)

struct *i2c_config_t*::**[anonymous]**::**[anonymous] master**
I2C master config

uint8_t **addr_10bit_en**
I2C 10bit address mode enable for slave mode

uint16_t **slave_addr**
I2C address for slave mode

struct *i2c_config_t*::**[anonymous]**::**[anonymous] slave**
I2C slave config

Type Definitions

typedef int **i2c_port_t**
I2C port number, can be I2C_NUM_0 ~ (I2C_NUM_MAX-1).

Enumerations

enum i2c_mode_t
Values:

I2C_MODE_SLAVE = 0
I2C slave mode

I2C_MODE_MASTER
I2C master mode

I2C_MODE_MAX

enum i2c_rw_t
Values:

I2C_MASTER_WRITE = 0
I2C write data

I2C_MASTER_READ
I2C read data

enum i2c_opmode_t
Values:

I2C_CMD_RESTART = 0
I2C restart command

I2C_CMD_WRITE
I2C write command

I2C_CMD_READ
I2C read command

I2C_CMD_STOP
I2C stop command

I2C_CMD_END
I2C end command

enum i2c_trans_mode_t

Values:

I2C_DATA_MODE_MSB_FIRST = 0
I2C data msb first

I2C_DATA_MODE_LSB_FIRST = 1
I2C data lsb first

I2C_DATA_MODE_MAX

enum i2c_addr_mode_t

Values:

I2C_ADDR_BIT_7 = 0
I2C 7bit address for slave mode

I2C_ADDR_BIT_10
I2C 10bit address for slave mode

I2C_ADDR_BIT_MAX

enum i2c_ack_type_t

Values:

I2C_MASTER_ACK = 0x0
I2C ack for each byte read

I2C_MASTER_NACK = 0x1
I2C nack for each byte read

I2C_MASTER_LAST_NACK = 0x2
I2C nack for the last byte

I2C_MASTER_ACK_MAX

enum i2c_sclk_t

Values:

I2C_SCLK_REF_TICK
I2C source clock from REF_TICK

I2C_SCLK_APB
I2C source clock from APB

2.2.7 I2S

Overview

I2S (Inter-IC Sound) is a serial, synchronous communication protocol that is usually used for transmitting audio data between two digital audio devices.

ESP32-S2 contains one I2S peripheral. These peripherals can be configured to input and output sample data via the I2S driver.

An I2S bus consists of the following lines:

- Bit clock line
- Channel select line
- Serial data line

Each I2S controller has the following features that can be configured using the I2S driver:

- Operation as system master or slave
- Capable of acting as transmitter or receiver
- Dedicated DMA controller that allows for streaming sample data without requiring the CPU to copy each data sample

Each controller can operate in half-duplex communication mode. Thus, the two controllers can be combined to establish full-duplex communication.

I2S0 output can be routed directly to the digital-to-analog converter’ s (DAC) output channels (GPIO 25 & GPIO 26) to produce direct analog output without involving any external I2S codecs. I2S0 can also be used for transmitting PDM (Pulse-density modulation) signals.

The I2S peripherals also support LCD mode for communicating data over a parallel bus, as used by some LCD displays and camera modules. LCD mode has the following operational modes:

- LCD master transmitting mode
- Camera slave receiving mode
- ADC/DAC mode

注解： For high accuracy clock applications, use the APLL_CLK clock source, which has the frequency range of 16 ~ 128 MHz. You can enable the APLL_CLK clock source by setting *i2s_config_t::use_apll* to TRUE.

If *i2s_config_t::use_apll* = TRUE and *i2s_config_t::fixed_mclk* > 0, then the master clock output frequency for I2S will be equal to the value of *i2s_config_t::fixed_mclk*, which means that the mclk frequency is provided by the user, instead of being calculated by the driver.

The clock rate of the word select line, which is called audio left-right clock rate (LRCK) here, is always the divisor of the master clock output frequency and for which the following is always true: 0 < MCLK/LRCK/channels/bits_per_sample < 64.

Functional Overview

Installing the Driver Install the I2S driver by calling the function :cpp:func‘i2s_driver_install‘ and passing the following arguments:

- Port number
- The structure *i2s_config_t* with defined communication parameters
- Event queue size and handle

Configuration example:

```
static const int i2s_num = 0; // i2s port number

static const i2s_config_t i2s_config = {
    .mode = I2S_MODE_MASTER | I2S_MODE_TX,
    .sample_rate = 44100,
    .bits_per_sample = 16,
    .channel_format = I2S_CHANNEL_FMT_RIGHT_LEFT,
    .communication_format = I2S_COMM_FORMAT_STAND_I2S,
    .intr_alloc_flags = 0, // default interrupt priority
    .dma_buf_count = 8,
    .dma_buf_len = 64,
    .use_apll = false
};

i2s_driver_install(I2S_NUM, &i2s_config, 0, NULL);
```

Setting Communication Pins Once the driver is installed, configure physical GPIO pins to which signals will be routed. For this, call the function :cpp:func‘i2s_set_pin‘ and pass the following arguments to it:

- Port number
- The structure *i2s_pin_config_t* defining the GPIO pin numbers to which the driver should route the BCK, WS, DATA out, and DATA in signals. If you want to keep a currently allocated pin number for a specific signal, or if this signal is unused, then pass the macro *I2S_PIN_NO_CHANGE*. See the example below.

```
static const i2s_pin_config_t pin_config = {
    .bck_io_num = 26,
    .ws_io_num = 25,
    .data_out_num = 22,
    .data_in_num = I2S_PIN_NO_CHANGE
};

i2s_set_pin(i2s_num, &pin_config);
```

Running I2S Communication To perform a transmission:

- Prepare the data for sending
- Call the function *i2s_write()* and pass the data buffer address and data length to it

The function will write the data to the I2S DMA Tx buffer, and then the data will be transmitted automatically.

```
i2s_write(I2S_NUM, samples_data, ((bits+8)/16)*SAMPLE_PER_CYCLE*4, &i2s_bytes_
↪write, 100);
```

To retrieve received data, use the function *i2s_read()*. It will retrieve the data from the I2S DMA Rx buffer, once the data is received by the I2S controller.

You can temporarily stop the I2S driver by calling the function *i2s_stop()*, which will disable the I2S Tx/Rx units until the function *i2s_start()* is called. If the function :cpp:func‘i2s_driver_install‘ is used, the driver will start up automatically eliminating the need to call *i2s_start()*.

Deleting the Driver If the established communication is no longer required, the driver can be removed to free allocated resources by calling *i2s_driver_uninstall()*.

Application Example

A code example for the I2S driver can be found in the directory peripherals/i2s.

In addition, there are two short configuration examples for the I2S driver.

I2S configuration

```
#include "driver/i2s.h"
#include "freertos/queue.h"

static const int i2s_num = 0; // i2s port number

static const i2s_config_t i2s_config = {
    .mode = I2S_MODE_MASTER | I2S_MODE_TX,
    .sample_rate = 44100,
    .bits_per_sample = 16,
    .channel_format = I2S_CHANNEL_FMT_RIGHT_LEFT,
    .communication_format = I2S_COMM_FORMAT_STAND_I2S,
    .intr_alloc_flags = 0, // default interrupt priority
    .dma_buf_count = 8,
    .dma_buf_len = 64,
    .use_apll = false
};

static const i2s_pin_config_t pin_config = {
    .bck_io_num = 26,
    .ws_io_num = 25,
    .data_out_num = 22,
    .data_in_num = I2S_PIN_NO_CHANGE
```

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```
};

...

    i2s_driver_install(i2s_num, &i2s_config, 0, NULL);   //install and start i2s␣
↪driver

    i2s_set_pin(i2s_num, &pin_config);

    i2s_set_sample_rates(i2s_num, 22050); //set sample rates

    i2s_driver_uninstall(i2s_num); //stop & destroy i2s driver
```

Configuring I2S to use internal DAC for analog output

```
#include "driver/i2s.h"
#include "freertos/queue.h"

static const int i2s_num = 0; // i2s port number

static const i2s_config_t i2s_config = {
    .mode = I2S_MODE_MASTER | I2S_MODE_TX | I2S_MODE_DAC_BUILT_IN,
    .sample_rate = 44100,
    .bits_per_sample = 16, /* the DAC module will only take the 8bits from MSB */
    .channel_format = I2S_CHANNEL_FMT_RIGHT_LEFT,
    .intr_alloc_flags = 0, // default interrupt priority
    .dma_buf_count = 8,
    .dma_buf_len = 64,
    .use_apll = false
};

...

    i2s_driver_install(i2s_num, &i2s_config, 0, NULL);   //install and start i2s␣
↪driver

    i2s_set_pin(i2s_num, NULL); //for internal DAC, this will enable both of the␣
↪internal channels

    //You can call i2s_set_dac_mode to set built-in DAC output mode.
    //i2s_set_dac_mode(I2S_DAC_CHANNEL_BOTH_EN);

    i2s_set_sample_rates(i2s_num, 22050); //set sample rates

    i2s_driver_uninstall(i2s_num); //stop & destroy i2s driver
```

API Reference

Header File

- driver/include/driver/i2s.h

Functions

esp_err_t **i2s_set_pin** (*i2s_port_t i2s_num*, **const** *i2s_pin_config_t *pin*)

Set I2S pin number.

Inside the pin configuration structure, set I2S_PIN_NO_CHANGE for any pin where the current configuration should not be changed.

Note The I2S peripheral output signals can be connected to multiple GPIO pads. However, the I2S peripheral input signal can only be connected to one GPIO pad.

Parameters

- `i2s_num`: I2S_NUM_0 or I2S_NUM_1
- `pin`: I2S Pin structure, or NULL to set 2-channel 8-bit internal DAC pin configuration (GPIO25 & GPIO26)

Note if *pin is set as NULL, this function will initialize both of the built-in DAC channels by default. if you don't want this to happen and you want to initialize only one of the DAC channels, you can call i2s_set_dac_mode instead.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_FAIL IO error

esp_err_t **i2s_set_dac_mode**(*i2s_dac_mode_t dac_mode*)

Set I2S dac mode, I2S built-in DAC is disabled by default.

Note Built-in DAC functions are only supported on I2S0 for current ESP32 chip. If either of the built-in DAC channel are enabled, the other one can not be used as RTC DAC function at the same time.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `dac_mode`: DAC mode configurations - see i2s_dac_mode_t

esp_err_t **i2s_driver_install**(*i2s_port_t i2s_num*, **const** *i2s_config_t *i2s_config*, int *queue_size*, void **i2s_queue*)

Install and start I2S driver.

This function must be called before any I2S driver read/write operations.

Parameters

- `i2s_num`: I2S_NUM_0, I2S_NUM_1
- `i2s_config`: I2S configurations - see *i2s_config_t* struct
- `queue_size`: I2S event queue size/depth.
- `i2s_queue`: I2S event queue handle, if set NULL, driver will not use an event queue.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NO_MEM Out of memory

esp_err_t **i2s_driver_uninstall**(*i2s_port_t i2s_num*)

Uninstall I2S driver.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2s_num`: I2S_NUM_0, I2S_NUM_1

esp_err_t **i2s_write**(*i2s_port_t i2s_num*, **const** void **src*, size_t *size*, size_t **bytes_written*, TickType_t *ticks_to_wait*)

Write data to I2S DMA transmit buffer.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2s_num`: I2S_NUM_0, I2S_NUM_1
- `src`: Source address to write from
- `size`: Size of data in bytes

- `[out] bytes_written`: Number of bytes written, if timeout, the result will be less than the size passed in.
- `ticks_to_wait`: TX buffer wait timeout in RTOS ticks. If this many ticks pass without space becoming available in the DMA transmit buffer, then the function will return (note that if the data is written to the DMA buffer in pieces, the overall operation may still take longer than this timeout.) Pass portMAX_DELAY for no timeout.

esp_err_t **i2s_write_expand**(*i2s_port_t i2s_num*, **const** void **src*, size_t *size*, size_t *src_bits*, size_t *aim_bits*, size_t **bytes_written*, TickType_t *ticks_to_wait*)

Write data to I2S DMA transmit buffer while expanding the number of bits per sample. For example, expanding 16-bit PCM to 32-bit PCM.

Format of the data in source buffer is determined by the I2S configuration (see *i2s_config_t*).

Parameters

- `i2s_num`: I2S_NUM_0, I2S_NUM_1
- `src`: Source address to write from
- `size`: Size of data in bytes
- `src_bits`: Source audio bit
- `aim_bits`: Bit wanted, no more than 32, and must be greater than src_bits
- `[out] bytes_written`: Number of bytes written, if timeout, the result will be less than the size passed in.
- `ticks_to_wait`: TX buffer wait timeout in RTOS ticks. If this many ticks pass without space becoming available in the DMA transmit buffer, then the function will return (note that if the data is written to the DMA buffer in pieces, the overall operation may still take longer than this timeout.) Pass portMAX_DELAY for no timeout.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

esp_err_t **i2s_read**(*i2s_port_t i2s_num*, void **dest*, size_t *size*, size_t **bytes_read*, TickType_t *ticks_to_wait*)

Read data from I2S DMA receive buffer.

Note If the built-in ADC mode is enabled, we should call i2s_adc_enable and i2s_adc_disable around the whole reading process, to prevent the data getting corrupted.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `i2s_num`: I2S_NUM_0, I2S_NUM_1
- `dest`: Destination address to read into
- `size`: Size of data in bytes
- `[out] bytes_read`: Number of bytes read, if timeout, bytes read will be less than the size passed in.
- `ticks_to_wait`: RX buffer wait timeout in RTOS ticks. If this many ticks pass without bytes becoming available in the DMA receive buffer, then the function will return (note that if data is read from the DMA buffer in pieces, the overall operation may still take longer than this timeout.) Pass portMAX_DELAY for no timeout.

esp_err_t **i2s_set_sample_rates**(*i2s_port_t i2s_num*, uint32_t *rate*)

Set sample rate used for I2S RX and TX.

The bit clock rate is determined by the sample rate and *i2s_config_t* configuration parameters (number of channels, bits_per_sample).

```
bit_clock = rate * (number of channels) * bits_per_sample
```

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NO_MEM Out of memory

Parameters
- `i2s_num`: I2S_NUM_0, I2S_NUM_1
- `rate`: I2S sample rate (ex: 8000, 44100···)

esp_err_t **i2s_stop** (*i2s_port_t i2s_num*)
Stop I2S driver.

There is no need to call i2s_stop() before calling i2s_driver_uninstall().

Disables I2S TX/RX, until i2s_start() is called.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2s_num`: I2S_NUM_0, I2S_NUM_1

esp_err_t **i2s_start** (*i2s_port_t i2s_num*)
Start I2S driver.

It is not necessary to call this function after i2s_driver_install() (it is started automatically), however it is necessary to call it after i2s_stop().

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2s_num`: I2S_NUM_0, I2S_NUM_1

esp_err_t **i2s_zero_dma_buffer** (*i2s_port_t i2s_num*)
Zero the contents of the TX DMA buffer.

Pushes zero-byte samples into the TX DMA buffer, until it is full.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `i2s_num`: I2S_NUM_0, I2S_NUM_1

esp_err_t **i2s_set_clk** (*i2s_port_t i2s_num*, uint32_t *rate*, *i2s_bits_per_sample_t bits*, *i2s_channel_t ch*)
Set clock & bit width used for I2S RX and TX.

Similar to i2s_set_sample_rates(), but also sets bit width.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NO_MEM Out of memory

Parameters
- `i2s_num`: I2S_NUM_0, I2S_NUM_1
- `rate`: I2S sample rate (ex: 8000, 44100···)
- `bits`: I2S bit width (I2S_BITS_PER_SAMPLE_16BIT, I2S_BITS_PER_SAMPLE_24BIT, I2S_BITS_PER_SAMPLE_32BIT)
- `ch`: I2S channel, (I2S_CHANNEL_MONO, I2S_CHANNEL_STEREO)

float **i2s_get_clk** (*i2s_port_t i2s_num*)
get clock set on particular port number.

Return
- actual clock set by i2s driver

Parameters
- `i2s_num`: I2S_NUM_0, I2S_NUM_1

Macros

I2S_PIN_NO_CHANGE

Use in *i2s_pin_config_t* for pins which should not be changed

Type Definitions

typedef *intr_handle_t* **i2s_isr_handle_t**

Header File

- soc/include/hal/i2s_types.h

Structures

struct i2s_config_t

I2S configuration parameters for i2s_param_config function.

Public Members

i2s_mode_t **mode**

I2S work mode

int **sample_rate**

I2S sample rate

i2s_bits_per_sample_t **bits_per_sample**

I2S bits per sample

i2s_channel_fmt_t **channel_format**

I2S channel format

i2s_comm_format_t **communication_format**

I2S communication format

int **intr_alloc_flags**

Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info

int **dma_buf_count**

I2S DMA Buffer Count

int **dma_buf_len**

I2S DMA Buffer Length

bool **use_apll**

I2S using APLL as main I2S clock, enable it to get accurate clock

bool **tx_desc_auto_clear**

I2S auto clear tx descriptor if there is underflow condition (helps in avoiding noise in case of data unavailability)

int **fixed_mclk**

I2S using fixed MCLK output. If use_apll = true and fixed_mclk > 0, then the clock output for i2s is fixed and equal to the fixed_mclk value.

struct i2s_event_t

Event structure used in I2S event queue.

Public Members

i2s_event_type_t **type**

I2S event type

size_t **size**
I2S data size for I2S_DATA event

struct i2s_pin_config_t
I2S pin number for i2s_set_pin.

Public Members

int **bck_io_num**
BCK in out pin

int **ws_io_num**
WS in out pin

int **data_out_num**
DATA out pin

int **data_in_num**
DATA in pin

Enumerations

enum i2s_port_t
I2S port number, the max port number is (I2S_NUM_MAX -1).

Values:

I2S_NUM_0 = 0
I2S port 0

I2S_NUM_MAX
I2S port max

enum i2s_bits_per_sample_t
I2S bit width per sample.

Values:

I2S_BITS_PER_SAMPLE_8BIT = 8
I2S bits per sample: 8-bits

I2S_BITS_PER_SAMPLE_16BIT = 16
I2S bits per sample: 16-bits

I2S_BITS_PER_SAMPLE_24BIT = 24
I2S bits per sample: 24-bits

I2S_BITS_PER_SAMPLE_32BIT = 32
I2S bits per sample: 32-bits

enum i2s_channel_t
I2S channel.

Values:

I2S_CHANNEL_MONO = 1
I2S 1 channel (mono)

I2S_CHANNEL_STEREO = 2
I2S 2 channel (stereo)

enum i2s_comm_format_t
I2S communication standard format.

Values:

I2S_COMM_FORMAT_STAND_I2S = 0X01
I2S communication I2S Philips standard, data launch at second BCK

I2S_COMM_FORMAT_STAND_MSB = 0X03
I2S communication MSB alignment standard, data launch at first BCK

I2S_COMM_FORMAT_STAND_PCM_SHORT = 0x04
PCM Short standard

I2S_COMM_FORMAT_STAND_PCM_LONG = 0x0C
PCM Long standard

I2S_COMM_FORMAT_STAND_MAX
standard max

I2S_COMM_FORMAT_I2S = 0x01
I2S communication format I2S, correspond to `I2S_COMM_FORMAT_STAND_I2S`

I2S_COMM_FORMAT_I2S_MSB = 0x01
I2S format MSB, (I2S_COMM_FORMAT_I2S |I2S_COMM_FORMAT_I2S_MSB) correspond to `I2S_COMM_FORMAT_STAND_I2S`

I2S_COMM_FORMAT_I2S_LSB = 0x02
I2S format LSB, (I2S_COMM_FORMAT_I2S |I2S_COMM_FORMAT_I2S_LSB) correspond to `I2S_COMM_FORMAT_STAND_MSB`

I2S_COMM_FORMAT_PCM = 0x04
I2S communication format PCM, correspond to `I2S_COMM_FORMAT_STAND_PCM_SHORT`

I2S_COMM_FORMAT_PCM_SHORT = 0x04
PCM Short, (I2S_COMM_FORMAT_PCM | I2S_COMM_FORMAT_PCM_SHORT) correspond to `I2S_COMM_FORMAT_STAND_PCM_SHORT`

I2S_COMM_FORMAT_PCM_LONG = 0x08
PCM Long, (I2S_COMM_FORMAT_PCM | I2S_COMM_FORMAT_PCM_LONG) correspond to `I2S_COMM_FORMAT_STAND_PCM_LONG`

enum i2s_channel_fmt_t
I2S channel format type.

Values:

I2S_CHANNEL_FMT_RIGHT_LEFT = 0x00

I2S_CHANNEL_FMT_ALL_RIGHT

I2S_CHANNEL_FMT_ALL_LEFT

I2S_CHANNEL_FMT_ONLY_RIGHT

I2S_CHANNEL_FMT_ONLY_LEFT

enum i2s_mode_t
I2S Mode, defaut is I2S_MODE_MASTER | I2S_MODE_TX.

Note PDM and built-in DAC functions are only supported on I2S0 for current ESP32 chip.

Values:

I2S_MODE_MASTER = 1
Master mode

I2S_MODE_SLAVE = 2
Slave mode

I2S_MODE_TX = 4
TX mode

I2S_MODE_RX = 8
RX mode

enum i2s_clock_src_t
I2S source clock.

Values:

I2S_CLK_D2CLK = 0
Clock from PLL_D2_CLK(160M)

I2S_CLK_APLL
Clock from APLL

enum i2s_event_type_t
I2S event types.

Values:

I2S_EVENT_DMA_ERROR

I2S_EVENT_TX_DONE
I2S DMA finish sent 1 buffer

I2S_EVENT_RX_DONE
I2S DMA finish received 1 buffer

I2S_EVENT_MAX
I2S event max index

enum i2s_dac_mode_t
I2S DAC mode for i2s_set_dac_mode.

Note PDM and built-in DAC functions are only supported on I2S0 for current ESP32 chip.

Values:

I2S_DAC_CHANNEL_DISABLE = 0
Disable I2S built-in DAC signals

I2S_DAC_CHANNEL_RIGHT_EN = 1
Enable I2S built-in DAC right channel, maps to DAC channel 1 on GPIO25

I2S_DAC_CHANNEL_LEFT_EN = 2
Enable I2S built-in DAC left channel, maps to DAC channel 2 on GPIO26

I2S_DAC_CHANNEL_BOTH_EN = 0x3
Enable both of the I2S built-in DAC channels.

I2S_DAC_CHANNEL_MAX = 0x4
I2S built-in DAC mode max index

2.2.8 LED PWM 控制器

概述

LED PWM 控制器主要用于控制 LED，也可产生 PWM 信号用于其他设备的控制。该控制器有 8 路高速通道和 8 路低速通道，可以产生独立的波形来驱动 RGB LED 设备等。

LED PWM 控制器的高速通道和低速通道均支持硬件渐变功能，可在无需 CPU 干预的情况下自动改变 PWM 信号的占空比，也可由软件改变 PWM 信号的占空比，实现亮度和颜色渐变。此外，低速通道在 Sleep 模式下仍可运行。

功能概览

要让指定 LED PWM 控制器高速模式或低速模式 通道运行，需进行如下配置：

1. 配置定时器 指定 PWM 信号的频率和占空比分辨率。
2. 配置通道 绑定定时器和输出 PWM 信号的 GPIO。

3. *改变 PWM 信号* 输出 PWM 信号来驱动 LED。可通过软件控制或使用硬件渐变功能来改变 LED 的亮度。

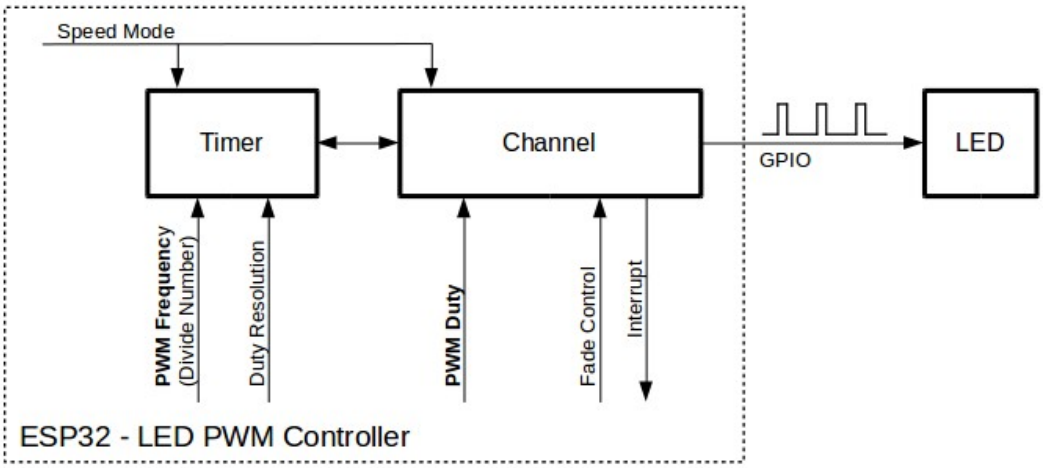


图 7: 配置 LED PWM 控制器的关键 API

配置定时器 要设置定时器，可调用函数*ledc_timer_config()*，并将包括如下配置参数的数据结构*ledc_timer_config_t* 传递给该函数：

- 定时器索引*ledc_timer_t*
- 速度模式*ledc_mode_t*
- PWM 信号频率
- PWM 占空比分辨率

频率和占空比分辨率相互关联。PWM 频率越高，占空比分辨率越低，反之则越高。使用该 API 用于除改变 LED 亮度以外的其他目的时，这一点很重要。更多信息详见*频率和占空比分辨率支持范围* 一节。

配置通道 定时器设置好后，请配置选定的通道（*ledc_channel_t* 之一）。配置通道需调用函数*ledc_channel_config()*。

通道的配置与定时器设置类似，需向通道配置函数传递包括通道配置参数的结构*ledc_channel_config_t* 。

此时，通道会按照*ledc_channel_config_t* 的配置，在绑定的 GPIO 上输出具有指定频率和占空比的 PWM 信号。在通道工作过程中，可以随时通过调用函数*ledc_stop()* 将其暂停。

改变 PWM 信号 通道开始运行、生成具有恒定占空比和频率的 PWM 信号之后，有几种方式可以改变该信号。驱动 LED 时，主要通过改变占空比来变化光线亮度。

以下两节介绍了如何使用软件和硬件改变占空比。如有需要，PWM 信号的频率也可更改，详见*改变 PWM 频率* 一节。

使用软件改变 PWM 占空比 调用函数*ledc_set_duty()* 可以设置新的占空比。之后，调用函数*ledc_update_duty()* 使新配置生效。要查看当前的占空比，可使用 _get_ 函数*ledc_get_duty()*。

另外一种设置占空比和其他通道参数的方式是调用*配置通道* 一节提到的函数*ledc_channel_config()*。

传递给函数的占空比数值范围取决于选定的 duty_resolution，应为 0 至 (2 ** duty_resolution) - 1。例如，如选定的占空比分辨率为 10，则占空比的数值范围为 0 至 1023。此时分辨率为 ~0.1%。

使用硬件渐变改变 PWM 占空比 LED PWM 控制器硬件可逐渐改变占空比的数值。要使用此功能，需用函数*ledc_fade_func_install()* 使能渐变，之后用下列可用渐变函数之一配置：

- *ledc_set_fade_with_time()*
- *ledc_set_fade_with_step()*
- *ledc_set_fade()*

最后用*ledc_fade_start()* 开启渐变。

如不需要渐变和渐变中断，可用函数*ledc_fade_func_uninstall()* 关闭。

改变 PWM 频率 LED PWM 控制器 API 有多种方式即时改变 PWM 频率：

- 通过调用函数*ledc_set_freq()* 设置频率。可用函数*ledc_get_freq()* 查看当前频率。
- 通过调用函数*ledc_bind_channel_timer()* 改变频率和占空比分辨率，将其他定时器与通道相连。
- 通过调用函数*ledc_channel_config()* 改变通道的定时器。

控制 PWM 的更多方式 要改变 PWM 设置，可使用低电平定时器的特定功能：

- *ledc_timer_set()*
- *ledc_timer_rst()*
- *ledc_timer_pause()*
- *ledc_timer_resume()*

前两个功能可通过函数*ledc_channel_config()* 在后台运行，在定时器配置后启动。

使用中断 配置 LED PWM 控制器通道时，可在*ledc_channel_config_t* 中选取参数*ledc_intr_type_t*，在渐变完成时触发中断。

要注册处理程序来处理中断，可调用函数*ledc_isr_register()*。

LED PWM 控制器高速和低速模式

LED PWM 控制器有 8 个可用定时器和 16 路通道，其中有一半专以高速模式运行，另一半则以低速模式运行。要选取高速或低速定时器或通道，可借助所调用函数中的参数*ledc_mode_t*。

高速模式的优点是可平稳地改变定时器设置。也就是说，高速模式下如定时器设置改变，此变更会自动应用于定时器的下一次溢出中断。而更新低速定时器时，设置变更应由软件显式触发。LED PWM 驱动器在后台更改设置，比如在调用函数*ledc_timer_config()* 或*ledc_timer_set()* 时。

更多关于速度模式的详细信息请参阅 ESP32 技术参考手册 (PDF)。注意，该手册中提到的支持 SLOW_CLOCK 暂不适用于 LED PWM 驱动器。

频率和占空比分辨率支持范围

LED PWM 控制器主要用于驱动 LED。该控制器 PWM 占空比设置的分辨率范围较广。比如，PWM 频率为 5 kHz 时，占空比分辨率最大可为 13 位。这意味着占空比可为 0 至 100% 之间的任意值，分辨率为 ~0.012%（2 ** 13 = 8192 LED 亮度的离散电平）。

LED PWM 控制器可用于生成频率较高的信号，足以为数码相机模组等其他设备计时。此时，最大频率可为 40 MHz，占空比分辨率为 1 位。也就是说，占空比固定为 50%，无法调整。

LED PWM 控制器 API 可在设定的频率和占空比分辨率超过 LED PWM 控制器硬件范围时报错。例如，试图将频率设置为 20 MHz、占空比分辨率设置为 3 位时，串行端口监视器上会报告如下错误：

```
E (196) ledc: requested frequency and duty resolution cannot be achieved, try␣
↪reducing freq_hz or duty_resolution. div_param=128
```

此时，占空比分辨率或频率必须降低。比如，将占空比分辨率设置为 2 会解决这一问题，让占空比设置为 25% 的倍数，即 25%、50% 或 75%。

如设置的频率和占空比分辨率低于所支持的最低值，LED PWM 驱动器也会反映并报告，如：

```
E (196) ledc: requested frequency and duty resolution cannot be achieved, try␣
↪increasing freq_hz or duty_resolution. div_param=128000000
```

占空比分辨率通常用 *ledc_timer_bit_t* 设置，范围是 10 至 15 位。如需较低的占空比分辨率（上至 10，下至 1），可直接输入相应数值。

应用实例

LED PWM 改变占空比和渐变控制的实例请参照 peripherals/ledc。

API 参考

Header File

- driver/include/driver/ledc.h

Functions

esp_err_t **ledc_channel_config**(**const** *ledc_channel_config_t* **ledc_conf*)

LEDC channel configuration Configure LEDC channel with the given channel/output gpio_num/interrupt/source timer/frequency(Hz)/LEDC duty resolution.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `ledc_conf`: Pointer of LEDC channel configure struct

esp_err_t **ledc_timer_config**(**const** *ledc_timer_config_t* **timer_conf*)

LEDC timer configuration Configure LEDC timer with the given source timer/frequency(Hz)/duty_resolution.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_FAIL Can not find a proper pre-divider number base on the given frequency and the current duty_resolution.

Parameters
- `timer_conf`: Pointer of LEDC timer configure struct

esp_err_t **ledc_update_duty**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*)

LEDC update channel parameters.

Note Call this function to activate the LEDC updated parameters. After ledc_set_duty, we need to call this function to update the settings. And the new LEDC parameters don't take effect until the next PWM cycle.

Note ledc_set_duty, ledc_set_duty_with_hpoint and ledc_update_duty are not thread-safe, do not call these functions to control one LEDC channel in different tasks at the same time. A thread-safe version of API is ledc_set_duty_and_update

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- speed_mode: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- channel: LEDC channel (0-7), select from ledc_channel_t

esp_err_t **ledc_set_pin** (int *gpio_num*, *ledc_mode_t* *speed_mode*, *ledc_channel_t* *ledc_channel*)

Set LEDC output gpio.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- gpio_num: The LEDC output gpio
- speed_mode: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- ledc_channel: LEDC channel (0-7), select from ledc_channel_t

esp_err_t **ledc_stop** (*ledc_mode_t* *speed_mode*, *ledc_channel_t* *channel*, uint32_t *idle_level*)

LEDC stop. Disable LEDC output, and set idle level.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- speed_mode: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- channel: LEDC channel (0-7), select from ledc_channel_t
- idle_level: Set output idle level after LEDC stops.

esp_err_t **ledc_set_freq** (*ledc_mode_t* *speed_mode*, *ledc_timer_t* *timer_num*, uint32_t *freq_hz*)

LEDC set channel frequency (Hz)

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_FAIL Can not find a proper pre-divider number base on the given frequency and the current duty_resolution.

Parameters
- speed_mode: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- timer_num: LEDC timer index (0-3), select from ledc_timer_t
- freq_hz: Set the LEDC frequency

uint32_t **ledc_get_freq** (*ledc_mode_t* *speed_mode*, *ledc_timer_t* *timer_num*)

LEDC get channel frequency (Hz)

Return
- 0 error
- Others Current LEDC frequency

Parameters
- speed_mode: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- timer_num: LEDC timer index (0-3), select from ledc_timer_t

esp_err_t **ledc_set_duty_with_hpoint** (*ledc_mode_t* *speed_mode*, *ledc_channel_t* *channel*, uint32_t *duty*, uint32_t *hpoint*)

LEDC set duty and hpoint value Only after calling ledc_update_duty will the duty update.

Note ledc_set_duty, ledc_set_duty_with_hpoint and ledc_update_duty are not thread-safe, do not call these functions to control one LEDC channel in different tasks at the same time. A thread-safe version of API is ledc_set_duty_and_update

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel (0-7), select from ledc_channel_t
- `duty`: Set the LEDC duty, the range of duty setting is [0, (2**duty_resolution)]
- `hpoint`: Set the LEDC hpoint value(max: 0xfffff)

int **ledc_get_hpoint**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*)

LEDC get hpoint value, the counter value when the output is set high level.

Return

- LEDC_ERR_VAL if parameter error
- Others Current hpoint value of LEDC channel

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel (0-7), select from ledc_channel_t

esp_err_t **ledc_set_duty**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *duty*)

LEDC set duty This function do not change the hpoint value of this channel. if needed, please call ledc_set_duty_with_hpoint. only after calling ledc_update_duty will the duty update.

Note ledc_set_duty, ledc_set_duty_with_hpoint and ledc_update_duty are not thread-safe, do not call these functions to control one LEDC channel in different tasks at the same time. A thread-safe version of API is ledc_set_duty_and_update.

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel (0-7), select from ledc_channel_t
- `duty`: Set the LEDC duty, the range of duty setting is [0, (2**duty_resolution)]

uint32_t **ledc_get_duty**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*)

LEDC get duty This function returns the duty at the present PWM cycle. You shouldn't expect the function to return the new duty in the same cycle of calling ledc_update_duty, because duty update doesn't take effect until the next cycle.

Return

- LEDC_ERR_DUTY if parameter error
- Others Current LEDC duty

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel (0-7), select from ledc_channel_t

esp_err_t **ledc_set_fade**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *duty*, *ledc_duty_direction_t fade_direction*, uint32_t *step_num*, uint32_t *duty_cycle_num*, uint32_t *duty_scale*)

LEDC set gradient Set LEDC gradient, After the function calls the ledc_update_duty function, the function can take effect.

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel (0-7), select from ledc_channel_t
- `duty`: Set the start of the gradient duty, the range of duty setting is [0, (2**duty_resolution)]
- `fade_direction`: Set the direction of the gradient
- `step_num`: Set the number of the gradient
- `duty_cycle_num`: Set how many LEDC tick each time the gradient lasts
- `duty_scale`: Set gradient change amplitude

esp_err_t **ledc_isr_register**(void (**fn*)) void *
, void **arg*, int *intr_alloc_flags*, *ledc_isr_handle_t* **handle*Register LEDC interrupt handler, the handler is an ISR. The handler will be attached to the same CPU core that this function is running on.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Function pointer error.

Parameters

- `fn`: Interrupt handler function.
- `arg`: User-supplied argument passed to the handler function.
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `handle`: Pointer to return handle. If non-NULL, a handle for the interrupt will be returned here.

esp_err_t **ledc_timer_set**(*ledc_mode_t speed_mode*, *ledc_timer_t timer_sel*, uint32_t *clock_divider*, uint32_t *duty_resolution*, *ledc_clk_src_t clk_src*)
Configure LEDC settings.

Return

- (-1) Parameter error
- Other Current LEDC duty

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `timer_sel`: Timer index (0-3), there are 4 timers in LEDC module
- `clock_divider`: Timer clock divide value, the timer clock is divided from the selected clock source
- `duty_resolution`: Resolution of duty setting in number of bits. The range of duty values is [0, (2**duty_resolution)]
- `clk_src`: Select LEDC source clock.

esp_err_t **ledc_timer_rst**(*ledc_mode_t speed_mode*, *ledc_timer_t timer_sel*)
Reset LEDC timer.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `timer_sel`: LEDC timer index (0-3), select from ledc_timer_t

esp_err_t **ledc_timer_pause**(*ledc_mode_t speed_mode*, *ledc_timer_t timer_sel*)
Pause LEDC timer counter.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `timer_sel`: LEDC timer index (0-3), select from ledc_timer_t

esp_err_t **ledc_timer_resume**(*ledc_mode_t speed_mode*, *ledc_timer_t timer_sel*)

Resume LEDC timer.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `timer_sel`: LEDC timer index (0-3), select from ledc_timer_t

esp_err_t **ledc_bind_channel_timer**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*, *ledc_timer_t timer_sel*)

Bind LEDC channel with the selected timer.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel index (0-7), select from ledc_channel_t
- `timer_sel`: LEDC timer index (0-3), select from ledc_timer_t

esp_err_t **ledc_set_fade_with_step**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *target_duty*, uint32_t *scale*, uint32_t *cycle_num*)

Set LEDC fade function.

Note Call ledc_fade_func_install() once before calling this function. Call ledc_fade_start() after this to start fading.

Note ledc_set_fade_with_step, ledc_set_fade_with_time and ledc_fade_start are not thread-safe, do not call these functions to control one LEDC channel in different tasks at the same time. A thread-safe version of API is ledc_set_fade_step_and_start

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success
- ESP_ERR_INVALID_STATE Fade function not installed.
- ESP_FAIL Fade function init error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode. ,
- `channel`: LEDC channel index (0-7), select from ledc_channel_t
- `target_duty`: Target duty of fading [0, (2**duty_resolution) - 1]
- `scale`: Controls the increase or decrease step scale.
- `cycle_num`: increase or decrease the duty every cycle_num cycles

esp_err_t **ledc_set_fade_with_time**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *target_duty*, int *max_fade_time_ms*)

Set LEDC fade function, with a limited time.

Note Call ledc_fade_func_install() once before calling this function. Call ledc_fade_start() after this to start fading.

Note ledc_set_fade_with_step, ledc_set_fade_with_time and ledc_fade_start are not thread-safe, do not call these functions to control one LEDC channel in different tasks at the same time. A thread-safe version of API is ledc_set_fade_step_and_start

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

- ESP_ERR_INVALID_STATE Fade function not installed.
- ESP_FAIL Fade function init error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode. ,
- `channel`: LEDC channel index (0-7), select from ledc_channel_t
- `target_duty`: Target duty of fading.(0 - (2 ** duty_resolution - 1)))
- `max_fade_time_ms`: The maximum time of the fading (ms).

esp_err_t **ledc_fade_func_install** (int *intr_alloc_flags*)

Install LEDC fade function. This function will occupy interrupt of LEDC module.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE Fade function already installed.

Parameters

- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.

void **ledc_fade_func_uninstall** (void)

Uninstall LEDC fade function.

esp_err_t **ledc_fade_start** (*ledc_mode_t speed_mode*, *ledc_channel_t channel*, *ledc_fade_mode_t fade_mode*)

Start LEDC fading.

Note Call ledc_fade_func_install() once before calling this function. Call this API right after ledc_set_fade_with_time or ledc_set_fade_with_step before to start fading.

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE Fade function not installed.
- ESP_ERR_INVALID_ARG Parameter error.

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel number
- `fade_mode`: Whether to block until fading done. See ledc_types.h ledc_fade_mode_t for more info. Note that this function will not return until fading to the target duty if LEDC_FADE_WAIT_DONE mode is selected.

esp_err_t **ledc_set_duty_and_update** (*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *duty*, uint32_t *hpoint*)

A thread-safe API to set duty for LEDC channel and return when duty updated.

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel (0-7), select from ledc_channel_t
- `duty`: Set the LEDC duty, the range of duty setting is [0, (2**duty_resolution)]
- `hpoint`: Set the LEDC hpoint value(max: 0xfffff)

esp_err_t **ledc_set_fade_time_and_start** (*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *target_duty*, uint32_t *max_fade_time_ms*, *ledc_fade_mode_t fade_mode*)

A thread-safe API to set and start LEDC fade function, with a limited time.

Note Call ledc_fade_func_install() once, before calling this function.

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success
- ESP_ERR_INVALID_STATE Fade function not installed.
- ESP_FAIL Fade function init error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel index (0-7), select from ledc_channel_t
- `target_duty`: Target duty of fading.(0 - (2 ** duty_resolution - 1)))
- `max_fade_time_ms`: The maximum time of the fading (ms).
- `fade_mode`: choose blocking or non-blocking mode

esp_err_t **ledc_set_fade_step_and_start**(*ledc_mode_t speed_mode*, *ledc_channel_t channel*, uint32_t *target_duty*, uint32_t *scale*, uint32_t *cycle_num*, *ledc_fade_mode_t fade_mode*)

A thread-safe API to set and start LEDC fade function.

Note Call ledc_fade_func_install() once before calling this function.

Note If a fade operation is running in progress on that channel, the driver would not allow it to be stopped. Other duty operations will have to wait until the fade operation has finished.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success
- ESP_ERR_INVALID_STATE Fade function not installed.
- ESP_FAIL Fade function init error

Parameters

- `speed_mode`: Select the LEDC channel group with specified speed mode. Note that not all targets support high speed mode.
- `channel`: LEDC channel index (0-7), select from ledc_channel_t
- `target_duty`: Target duty of fading [0, (2**duty_resolution) - 1]
- `scale`: Controls the increase or decrease step scale.
- `cycle_num`: increase or decrease the duty every cycle_num cycles
- `fade_mode`: choose blocking or non-blocking mode

Macros

LEDC_APB_CLK_HZ

LEDC_REF_CLK_HZ

LEDC_ERR_DUTY

LEDC_ERR_VAL

Type Definitions

typedef *intr_handle_t* **ledc_isr_handle_t**

Header File

- soc/include/hal/ledc_types.h

Structures

struct ledc_channel_config_t

Configuration parameters of LEDC channel for ledc_channel_config function.

Public Members

int **gpio_num**
the LEDC output gpio_num, if you want to use gpio16, gpio_num = 16

ledc_mode_t **speed_mode**
LEDC speed speed_mode, high-speed mode or low-speed mode

ledc_channel_t **channel**
LEDC channel (0 - 7)

ledc_intr_type_t **intr_type**
configure interrupt, Fade interrupt enable or Fade interrupt disable

ledc_timer_t **timer_sel**
Select the timer source of channel (0 - 3)

uint32_t **duty**
LEDC channel duty, the range of duty setting is [0, (2**duty_resolution)]

int **hpoint**
LEDC channel hpoint value, the max value is 0xfffff

struct ledc_timer_config_t
Configuration parameters of LEDC Timer timer for ledc_timer_config function.

Public Members

ledc_mode_t **speed_mode**
LEDC speed speed_mode, high-speed mode or low-speed mode

ledc_timer_bit_t **duty_resolution**
LEDC channel duty resolution

ledc_timer_bit_t **bit_num**
Deprecated in ESP-IDF 3.0. This is an alias to ‘duty_resolution’ for backward compatibility with ESP-IDF 2.1

ledc_timer_t **timer_num**
The timer source of channel (0 - 3)

uint32_t **freq_hz**
LEDC timer frequency (Hz)

ledc_clk_cfg_t **clk_cfg**
Configure LEDC source clock. For low speed channels and high speed channels, you can specify the source clock using LEDC_USE_REF_TICK, LEDC_USE_APB_CLK or LEDC_AUTO_CLK. For low speed channels, you can also specify the source clock using LEDC_USE_RTC8M_CLK, in this case, all low speed channel’ s source clock must be RTC8M_CLK

Enumerations

enum ledc_mode_t
Values:

LEDC_LOW_SPEED_MODE
LEDC low speed speed_mode

LEDC_SPEED_MODE_MAX
LEDC speed limit

enum ledc_intr_type_t
Values:

LEDC_INTR_DISABLE = 0
Disable LEDC interrupt

LEDC_INTR_FADE_END
Enable LEDC interrupt

LEDC_INTR_MAX

enum ledc_duty_direction_t
Values:

LEDC_DUTY_DIR_DECREASE = 0
LEDC duty decrease direction

LEDC_DUTY_DIR_INCREASE = 1
LEDC duty increase direction

LEDC_DUTY_DIR_MAX

enum ledc_slow_clk_sel_t
Values:

LEDC_SLOW_CLK_RTC8M = 0
LEDC low speed timer clock source is 8MHz RTC clock

LEDC_SLOW_CLK_APB
LEDC low speed timer clock source is 80MHz APB clock

LEDC_SLOW_CLK_XTAL
LEDC low speed timer clock source XTAL clock

enum ledc_clk_cfg_t
Values:

LEDC_AUTO_CLK = 0
The driver will automatically select the source clock(REF_TICK or APB) based on the giving resolution and duty parameter when init the timer

LEDC_USE_REF_TICK
LEDC timer select REF_TICK clock as source clock

LEDC_USE_APB_CLK
LEDC timer select APB clock as source clock

LEDC_USE_RTC8M_CLK
LEDC timer select RTC8M_CLK as source clock. Only for low speed channels and this parameter must be the same for all low speed channels

LEDC_USE_XTAL_CLK
LEDC timer select XTAL clock as source clock

enum ledc_clk_src_t
Values:

LEDC_REF_TICK = *LEDC_USE_REF_TICK*
LEDC timer clock divided from reference tick (1Mhz)

LEDC_APB_CLK = *LEDC_USE_APB_CLK*
LEDC timer clock divided from APB clock (80Mhz)

enum ledc_timer_t
Values:

LEDC_TIMER_0 = 0
LEDC timer 0

LEDC_TIMER_1
LEDC timer 1

LEDC_TIMER_2
LEDC timer 2

LEDC_TIMER_3
LEDC timer 3

LEDC_TIMER_MAX

enum ledc_channel_t
Values:

LEDC_CHANNEL_0 = 0
LEDC channel 0

LEDC_CHANNEL_1
LEDC channel 1

LEDC_CHANNEL_2
LEDC channel 2

LEDC_CHANNEL_3
LEDC channel 3

LEDC_CHANNEL_4
LEDC channel 4

LEDC_CHANNEL_5
LEDC channel 5

LEDC_CHANNEL_6
LEDC channel 6

LEDC_CHANNEL_7
LEDC channel 7

LEDC_CHANNEL_MAX

enum ledc_timer_bit_t
Values:

LEDC_TIMER_1_BIT = 1
LEDC PWM duty resolution of 1 bits

LEDC_TIMER_2_BIT
LEDC PWM duty resolution of 2 bits

LEDC_TIMER_3_BIT
LEDC PWM duty resolution of 3 bits

LEDC_TIMER_4_BIT
LEDC PWM duty resolution of 4 bits

LEDC_TIMER_5_BIT
LEDC PWM duty resolution of 5 bits

LEDC_TIMER_6_BIT
LEDC PWM duty resolution of 6 bits

LEDC_TIMER_7_BIT
LEDC PWM duty resolution of 7 bits

LEDC_TIMER_8_BIT
LEDC PWM duty resolution of 8 bits

LEDC_TIMER_9_BIT
LEDC PWM duty resolution of 9 bits

LEDC_TIMER_10_BIT
LEDC PWM duty resolution of 10 bits

LEDC_TIMER_11_BIT
LEDC PWM duty resolution of 11 bits

LEDC_TIMER_12_BIT
LEDC PWM duty resolution of 12 bits

LEDC_TIMER_13_BIT
LEDC PWM duty resolution of 13 bits

LEDC_TIMER_14_BIT
LEDC PWM duty resolution of 14 bits

LEDC_TIMER_15_BIT
LEDC PWM duty resolution of 15 bits

LEDC_TIMER_16_BIT
LEDC PWM duty resolution of 16 bits

LEDC_TIMER_17_BIT
LEDC PWM duty resolution of 17 bits

LEDC_TIMER_18_BIT
LEDC PWM duty resolution of 18 bits

LEDC_TIMER_19_BIT
LEDC PWM duty resolution of 19 bits

LEDC_TIMER_20_BIT
LEDC PWM duty resolution of 20 bits

LEDC_TIMER_BIT_MAX

enum ledc_fade_mode_t
Values:

LEDC_FADE_NO_WAIT = 0
LEDC fade function will return immediately

LEDC_FADE_WAIT_DONE
LEDC fade function will block until fading to the target duty

LEDC_FADE_MAX

2.2.9 Pulse Counter

Introduction

The PCNT (Pulse Counter) module is designed to count the number of rising and/or falling edges of an input signal. Each pulse counter unit has a 16-bit signed counter register and two channels that can be configured to either increment or decrement the counter. Each channel has a signal input that accepts signal edges to be detected, as well as a control input that can be used to enable or disable the signal input. The inputs have optional filters that can be used to discard unwanted glitches in the signal.

Functionality Overview

Description of functionality of this API has been broken down into four sections:

- *Configuration* - describes counter’ s configuration parameters and how to setup the counter.
- *Operating the Counter* - provides information on control functions to pause, measure and clear the counter.
- *Filtering Pulses* - describes options to filtering pulses and the counter control signals.
- *Using Interrupts* - presents how to trigger interrupts on specific states of the counter.

Configuration

The PCNT module has four independent counting "units" numbered from 0 to 3. In the API they are referred to using *pcnt_unit_t*. Each unit has two independent channels numbered as 0 and 1 and specified with *pcnt_channel_t*.

The configuration is provided separately per unit's channel using *pcnt_config_t* and covers:

- The unit and the channel number this configuration refers to.
- GPIO numbers of the pulse input and the pulse gate input.
- Two pairs of parameters: *pcnt_ctrl_mode_t* and *pcnt_count_mode_t* to define how the counter reacts depending on the the status of control signal and how counting is done positive / negative edge of the pulses.
- Two limit values (minimum / maximum) that are used to establish watchpoints and trigger interrupts when the pulse count is meeting particular limit.

Setting up of particular channel is then done by calling a function *pcnt_unit_config()* with above *pcnt_config_t* as the input parameter.

To disable the pulse or the control input pin in configuration, provide `PCNT_PIN_NOT_USED` instead of the GPIO number.

Operating the Counter

After doing setup with *pcnt_unit_config()*, the counter immediately starts to operate. The accumulated pulse count can be checked by calling *pcnt_get_counter_value()*.

There are couple of functions that allow to control the counter's operation: *pcnt_counter_pause()*, *pcnt_counter_resume()* and *pcnt_counter_clear()*

It is also possible to dynamically change the previously set up counter modes with *pcnt_unit_config()* by calling *pcnt_set_mode()*.

If desired, the pulse input pin and the control input pin may be changed "on the fly" using *pcnt_set_pin()*. To disable particular input provide as a function parameter `PCNT_PIN_NOT_USED` instead of the GPIO number.

注解: For the counter not to miss any pulses, the pulse duration should be longer than one APB_CLK cycle (12.5 ns). The pulses are sampled on the edges of the APB_CLK clock and may be missed, if fall between the edges. This applies to counter operation with or without a *filer*.

Filtering Pulses

The PCNT unit features filters on each of the pulse and control inputs, adding the option to ignore short glitches in the signals.

The length of ignored pulses is provided in APB_CLK clock cycles by calling *pcnt_set_filter_value()*. The current filter setting may be checked with *pcnt_get_filter_value()*. The APB_CLK clock is running at 80 MHz.

The filter is put into operation / suspended by calling *pcnt_filter_enable()* / *pcnt_filter_disable()*.

Using Interrupts

There are five counter state watch events, defined in *pcnt_evt_type_t*, that are able to trigger an interrupt. The event happens on the pulse counter reaching specific values:

- Minimum or maximum count values: `counter_l_lim` or `counter_h_lim` provided in *pcnt_config_t* as discussed in *Configuration*

- Threshold 0 or Threshold 1 values set using function *pcnt_set_event_value()*.
- Pulse count = 0

To register, enable or disable an interrupt to service the above events, call *pcnt_isr_register()*, *pcnt_intr_enable()*. and *pcnt_intr_disable()*. To enable or disable events on reaching threshold values, you will also need to call functions *pcnt_event_enable()* and *pcnt_event_disable()*.

In order to check what are the threshold values currently set, use function *pcnt_get_event_value()*.

Application Example

Pulse counter with control signal and event interrupt example: peripherals/pcnt.

API Reference

Header File

- driver/include/driver/pcnt.h

Functions

esp_err_t **pcnt_unit_config**(**const** *pcnt_config_t* **pcnt_config*)

Configure Pulse Counter unit.

Note This function will disable three events: PCNT_EVT_L_LIM, PCNT_EVT_H_LIM, PCNT_EVT_ZERO.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver already initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `pcnt_config`: Pointer of Pulse Counter unit configure parameter

esp_err_t **pcnt_get_counter_value**(*pcnt_unit_t pcnt_unit*, int16_t **count*)

Get pulse counter value.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `pcnt_unit`: Pulse Counter unit number
- `count`: Pointer to accept counter value

esp_err_t **pcnt_counter_pause**(*pcnt_unit_t pcnt_unit*)

Pause PCNT counter of PCNT unit.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `pcnt_unit`: PCNT unit number

esp_err_t **pcnt_counter_resume**(*pcnt_unit_t pcnt_unit*)

Resume counting for PCNT counter.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `pcnt_unit`: PCNT unit number, select from pcnt_unit_t

esp_err_t **pcnt_counter_clear**(*pcnt_unit_t pcnt_unit*)

Clear and reset PCNT counter value to zero.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `pcnt_unit`: PCNT unit number, select from pcnt_unit_t

esp_err_t **pcnt_intr_enable**(*pcnt_unit_t pcnt_unit*)

Enable PCNT interrupt for PCNT unit.

Note Each Pulse counter unit has five watch point events that share the same interrupt. Configure events with pcnt_event_enable() and pcnt_event_disable()

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `pcnt_unit`: PCNT unit number

esp_err_t **pcnt_intr_disable**(*pcnt_unit_t pcnt_unit*)

Disable PCNT interrupt for PCNT unit.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `pcnt_unit`: PCNT unit number

esp_err_t **pcnt_event_enable**(*pcnt_unit_t unit*, *pcnt_evt_type_t evt_type*)

Enable PCNT event of PCNT unit.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number
- `evt_type`: Watch point event type. All enabled events share the same interrupt (one interrupt per pulse counter unit).

esp_err_t **pcnt_event_disable**(*pcnt_unit_t unit*, *pcnt_evt_type_t evt_type*)

Disable PCNT event of PCNT unit.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number
- `evt_type`: Watch point event type. All enabled events share the same interrupt (one interrupt per pulse counter unit).

esp_err_t **pcnt_set_event_value**(*pcnt_unit_t unit*, *pcnt_evt_type_t evt_type*, int16_t *value*)

Set PCNT event value of PCNT unit.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized

- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `unit`: PCNT unit number
- `evt_type`: Watch point event type. All enabled events share the same interrupt (one interrupt per pulse counter unit).
- `value`: Counter value for PCNT event

esp_err_t **pcnt_get_event_value** (*pcnt_unit_t unit*, *pcnt_evt_type_t evt_type*, int16_t **value*)

Get PCNT event value of PCNT unit.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `unit`: PCNT unit number
- `evt_type`: Watch point event type. All enabled events share the same interrupt (one interrupt per pulse counter unit).
- `value`: Pointer to accept counter value for PCNT event

esp_err_t **pcnt_isr_unregister** (*pcnt_isr_handle_t handle*)

Unregister PCNT interrupt handler (registered by pcnt_isr_register), the handler is an ISR. The handler will be attached to the same CPU core that this function is running on. If the interrupt service is registered by pcnt_isr_service_install, please call pcnt_isr_service_uninstall instead.

Return

- ESP_OK Success
- ESP_ERR_NOT_FOUND Can not find the interrupt that matches the flags.
- ESP_ERR_INVALID_ARG Function pointer error.

Parameters

- `handle`: handle to unregister the ISR service.

esp_err_t **pcnt_isr_register** (void (***fn*)) void *

, void ***arg*, int *intr_alloc_flags*, *pcnt_isr_handle_t *handle*Register PCNT interrupt handler, the handler is an ISR. The handler will be attached to the same CPU core that this function is running on. Please do not use pcnt_isr_service_install if this function was called.

Return

- ESP_OK Success
- ESP_ERR_NOT_FOUND Can not find the interrupt that matches the flags.
- ESP_ERR_INVALID_ARG Function pointer error.

Parameters

- `fn`: Interrupt handler function.
- `arg`: Parameter for handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `handle`: Pointer to return handle. If non-NULL, a handle for the interrupt will be returned here. Calling pcnt_isr_unregister to unregister this ISR service if needed, but only if the handle is not NULL.

esp_err_t **pcnt_set_pin** (*pcnt_unit_t unit*, *pcnt_channel_t channel*, int *pulse_io*, int *ctrl_io*)

Configure PCNT pulse signal input pin and control input pin.

Note Set the signal input to PCNT_PIN_NOT_USED if unused.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `unit`: PCNT unit number
- `channel`: PCNT channel number
- `pulse_io`: Pulse signal input GPIO

- `ctrl_io`: Control signal input GPIO

esp_err_t **pcnt_filter_enable**(*pcnt_unit_t unit*)

Enable PCNT input filter.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number

esp_err_t **pcnt_filter_disable**(*pcnt_unit_t unit*)

Disable PCNT input filter.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number

esp_err_t **pcnt_set_filter_value**(*pcnt_unit_t unit*, uint16_t *filter_val*)

Set PCNT filter value.

Note filter_val is a 10-bit value, so the maximum filter_val should be limited to 1023.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number
- `filter_val`: PCNT signal filter value, counter in APB_CLK cycles. Any pulses lasting shorter than this will be ignored when the filter is enabled.

esp_err_t **pcnt_get_filter_value**(*pcnt_unit_t unit*, uint16_t **filter_val*)

Get PCNT filter value.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number
- `filter_val`: Pointer to accept PCNT filter value.

esp_err_t **pcnt_set_mode**(*pcnt_unit_t unit*, *pcnt_channel_t channel*, *pcnt_count_mode_t pos_mode*, *pcnt_count_mode_t neg_mode*, *pcnt_ctrl_mode_t hctrl_mode*, *pcnt_ctrl_mode_t lctrl_mode*)

Set PCNT counter mode.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number
- `channel`: PCNT channel number
- `pos_mode`: Counter mode when detecting positive edge
- `neg_mode`: Counter mode when detecting negative edge
- `hctrl_mode`: Counter mode when control signal is high level
- `lctrl_mode`: Counter mode when control signal is low level

esp_err_t **pcnt_isr_handler_add** (*pcnt_unit_t unit*, void (**isr_handler*)) void *
, void **args*Add ISR handler for specified unit.

Call this function after using pcnt_isr_service_install() to install the PCNT driver's ISR handler service.

The ISR handlers do not need to be declared with IRAM_ATTR, unless you pass the ESP_INTR_FLAG_IRAM flag when allocating the ISR in pcnt_isr_service_install().

This ISR handler will be called from an ISR. So there is a stack size limit (configurable as "ISR stack size" in menuconfig). This limit is smaller compared to a global PCNT interrupt handler due to the additional level of indirection.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number
- `isr_handler`: Interrupt handler function.
- `args`: Parameter for handler function

esp_err_t **pcnt_isr_service_install** (int *intr_alloc_flags*)
Install PCNT ISR service.

Note We can manage different interrupt service for each unit. This function will use the default ISR handle service, Calling pcnt_isr_service_uninstall to uninstall the default service if needed. Please do not use pcnt_isr_register if this function was called.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_NO_MEM No memory to install this service
- ESP_ERR_INVALID_STATE ISR service already installed

Parameters
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.

void **pcnt_isr_service_uninstall** (void)
Uninstall PCNT ISR service, freeing related resources.

esp_err_t **pcnt_isr_handler_remove** (*pcnt_unit_t unit*)
Delete ISR handler for specified unit.

Return
- ESP_OK Success
- ESP_ERR_INVALID_STATE pcnt driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `unit`: PCNT unit number

Type Definitions

typedef *intr_handle_t* **pcnt_isr_handle_t**

Header File

- soc/include/hal/pcnt_types.h

Structures

struct pcnt_config_t
Pulse Counter configuration for a single channel.

Public Members

int **pulse_gpio_num**
Pulse input GPIO number, if you want to use GPIO16, enter pulse_gpio_num = 16, a negative value will be ignored

int **ctrl_gpio_num**
Control signal input GPIO number, a negative value will be ignored

pcnt_ctrl_mode_t **lctrl_mode**
PCNT low control mode

pcnt_ctrl_mode_t **hctrl_mode**
PCNT high control mode

pcnt_count_mode_t **pos_mode**
PCNT positive edge count mode

pcnt_count_mode_t **neg_mode**
PCNT negative edge count mode

int16_t **counter_h_lim**
Maximum counter value

int16_t **counter_l_lim**
Minimum counter value

pcnt_unit_t **unit**
PCNT unit number

pcnt_channel_t **channel**
the PCNT channel

Macros

PCNT_PIN_NOT_USED
When selected for a pin, this pin will not be used

Enumerations

enum pcnt_port_t
PCNT port number, the max port number is (PCNT_PORT_MAX - 1).

Values:

PCNT_PORT_0 = 0
PCNT port 0

PCNT_PORT_MAX
PCNT port max

enum pcnt_unit_t
Selection of all available PCNT units.

Values:

PCNT_UNIT_0 = 0
PCNT unit 0

PCNT_UNIT_1 = 1
PCNT unit 1

PCNT_UNIT_2 = 2
PCNT unit 2

PCNT_UNIT_3 = 3
PCNT unit 3

PCNT_UNIT_MAX

enum pcnt_ctrl_mode_t

Selection of available modes that determine the counter’s action depending on the state of the control signal’s input GPIO.

Note Configuration covers two actions, one for high, and one for low level on the control input

Values:

PCNT_MODE_KEEP = 0
Control mode: won’t change counter mode

PCNT_MODE_REVERSE = 1
Control mode: invert counter mode(increase -> decrease, decrease -> increase)

PCNT_MODE_DISABLE = 2
Control mode: Inhibit counter(counter value will not change in this condition)

PCNT_MODE_MAX

enum pcnt_count_mode_t

Selection of available modes that determine the counter’s action on the edge of the pulse signal’s input GPIO.

Note Configuration covers two actions, one for positive, and one for negative edge on the pulse input

Values:

PCNT_COUNT_DIS = 0
Counter mode: Inhibit counter(counter value will not change in this condition)

PCNT_COUNT_INC = 1
Counter mode: Increase counter value

PCNT_COUNT_DEC = 2
Counter mode: Decrease counter value

PCNT_COUNT_MAX

enum pcnt_channel_t

Selection of channels available for a single PCNT unit.

Values:

PCNT_CHANNEL_0 = 0x00
PCNT channel 0

PCNT_CHANNEL_1 = 0x01
PCNT channel 1

PCNT_CHANNEL_MAX

enum pcnt_evt_type_t

Selection of counter’s events the may trigger an interrupt.

Values:

PCNT_EVT_THRES_1 = BIT(2)
PCNT watch point event: threshold1 value event

PCNT_EVT_THRES_0 = BIT(3)
PCNT watch point event: threshold0 value event

PCNT_EVT_L_LIM = BIT(4)
PCNT watch point event: Minimum counter value

PCNT_EVT_H_LIM = BIT(5)
PCNT watch point event: Maximum counter value

PCNT_EVT_ZERO = BIT(6)
PCNT watch point event: counter value zero event

PCNT_EVT_MAX

2.2.10 RMT

The RMT (Remote Control) module driver can be used to send and receive infrared remote control signals. Due to flexibility of RMT module, the driver can also be used to generate or receive many other types of signals.

The signal, which consists of a series of pulses, is generated by RMT's transmitter based on a list of values. The values define the pulse duration and a binary level, see below. The transmitter can also provide a carrier and modulate it with provided pulses.

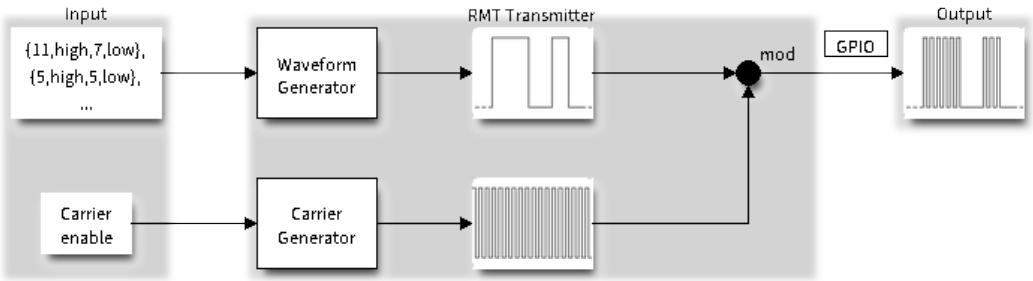


图 8: RMT Transmitter Overview

The reverse operation is performed by the receiver, where a series of pulses is decoded into a list of values containing the pulse duration and binary level. A filter may be applied to remove high frequency noise from the input signal.

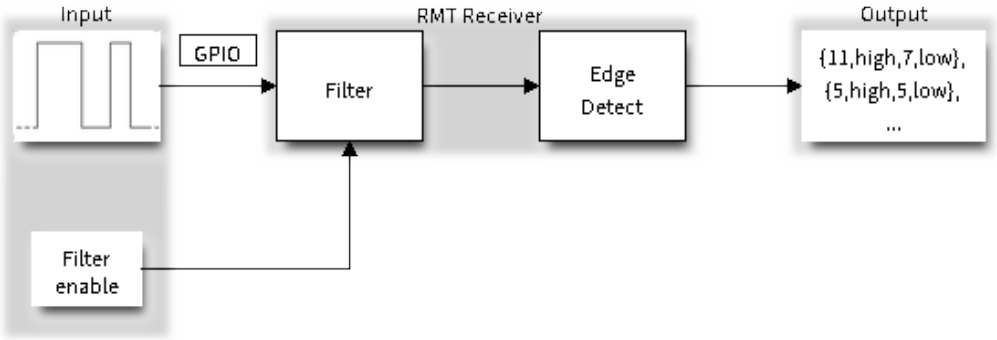


图 9: RMT Receiver Overview

There couple of typical steps to setup and operate the RMT and they are discussed in the following sections:

1. *Configure Driver*
2. *Transmit Data* or *Receive Data*
3. *Change Operation Parameters*
4. *Use Interrupts*

The RMT has four channels numbered from zero to three. Each channel is able to independently transmit or receive data. They are referred to using indexes defined in structure `rmt_channel_t`.

Configure Driver

There are several parameters that define how particular channel operates. Most of these parameters are configured by setting specific members of `rmt_config_t` structure. Some of the parameters are common to both transmit or receive mode, and some are mode specific. They are all discussed below.

Common Parameters

- The **channel** to be configured, select one from the *rmt_channel_t* enumerator.
- The RMT **operation mode** - whether this channel is used to transmit or receive data, selected by setting a **rmt_mode** members to one of the values from *rmt_mode_t*.
- What is the **pin number** to transmit or receive RMT signals, selected by setting **gpio_num**.
- How many **memory blocks** will be used by the channel, set with **mem_block_num**.
- Extra miscellaneous parameters for the channel can be set in the **flags**.
  - When **RMT_CHANNEL_FLAGS_ALWAYS_ON** is set, RMT channel will take REF_TICK as source clock. The benefit is, RMT channel can continue work even when APB clock is changing. See *power_management* for more information.
- A **clock divider**, that will determine the range of pulse length generated by the RMT transmitter or discriminated by the receiver. Selected by setting **clk_div** to a value within [1 .. 255] range. The RMT source clock is typically APB CLK, 80Mhz by default. But when **RMT_CHANNEL_FLAGS_ALWAYS_ON** is set in **flags**, RMT source clock is changed to REF_TICK.

注解: The period of a square wave after the clock divider is called a ‘tick’. The length of the pulses generated by the RMT transmitter or discriminated by the receiver is configured in number of ‘ticks’.

There are also couple of specific parameters that should be set up depending if selected channel is configured in *Transmit Mode* or *Receive Mode*:

Transmit Mode When configuring channel in transmit mode, set **tx_config** and the following members of *rmt_tx_config_t*:

- Transmit the currently configured data items in a loop - **loop_en**
- Enable the RMT carrier signal - **carrier_en**
- Frequency of the carrier in Hz - **carrier_freq_hz**
- Duty cycle of the carrier signal in percent (%) - **carrier_duty_percent**
- Level of the RMT output, when the carrier is applied - **carrier_level**
- Enable the RMT output if idle - **idle_output_en**
- Set the signal level on the RMT output if idle - **idle_level**
- Specify maximum number of transmissions in a loop - **loop_count**

Receive Mode In receive mode, set **rx_config** and the following members of *rmt_rx_config_t*:

- Enable a filter on the input of the RMT receiver - **filter_en**
- A threshold of the filter, set in the number of ticks - **filter_ticks_thresh**. Pulses shorter than this setting will be filtered out. Note, that the range of entered tick values is [0..255].
- A pulse length threshold that will turn the RMT receiver idle, set in number of ticks - **idle_threshold**. The receiver will ignore pulses longer than this setting.
- Enable the RMT carrier demodulation - **carrier_rm**
- Frequency of the carrier in Hz - **carrier_freq_hz**
- Duty cycle of the carrier signal in percent (%) - **carrier_duty_percent**
- Level of the RMT input, where the carrier is modulated to - **carrier_level**

Finalize Configuration Once the *rmt_config_t* structure is populated with parameters, it should be then invoked with *rmt_config()* to make the configuration effective.

The last configuration step is installation of the driver in memory by calling *rmt_driver_install()*. If `rx_buf_size` parameter of this function is > 0, then a ring buffer for incoming data will be allocated. A default ISR handler will be installed, see a note in *Use Interrupts*.

Now, depending on how the channel is configured, we are ready to either *Transmit Data* or *Receive Data*. This is described in next two sections.

Transmit Data

Before being able to transmit some RMT pulses, we need to define the pulse pattern. The minimum pattern recognized by the RMT controller, later called an ‘item’, is provided in a structure `rmt_item32_t`. Each item consists of two pairs of two values. The first value in a pair describes the signal duration in ticks and is 15 bits long, the second provides the signal level (high or low) and is contained in a single bit. A block of couple of items and the structure of an item is presented below.

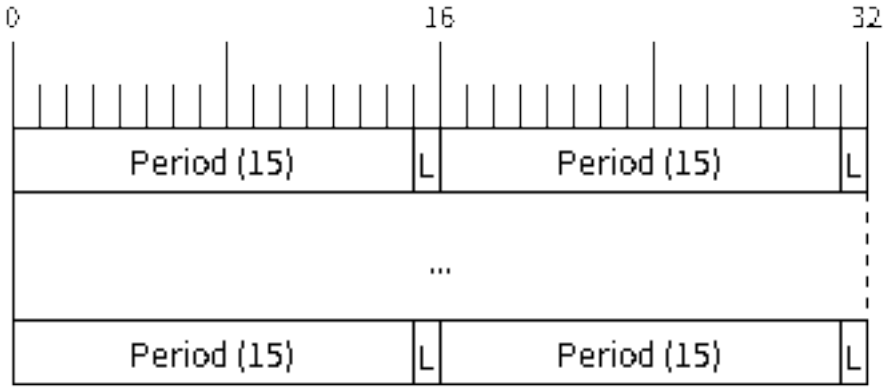


图 10: Structure of RMT items (L - signal level)

For a simple example how to define a block of items see peripherals/rmt/morse_code.

The items are provided to the RMT controller by calling function *`rmt_write_items()`*. This function also automatically triggers start of transmission. It may be called to wait for transmission completion or exit just after transmission start. In such case you can wait for the transmission end by calling *`rmt_wait_tx_done()`*. This function does not limit the number of data items to transmit. It is using an interrupt to successively copy the new data chunks to RMT’s internal memory as previously provided data are sent out.

Another way to provide data for transmission is by calling *`rmt_fill_tx_items()`*. In this case transmission is not started automatically. To control the transmission process use *`rmt_tx_start()`* and *`rmt_tx_stop()`*. The number of items to sent is restricted by the size of memory blocks allocated in the RMT controller’s internal memory, see *`rmt_set_mem_block_num()`*.

Receive Data

Before starting the receiver we need some storage for incoming items. The RMT controller has 256 x 32-bits of internal RAM shared between all four channels.

In typical scenarios it is not enough as an ultimate storage for all incoming (and outgoing) items. Therefore this API supports retrieval of incoming items on the fly to save them in a ring buffer of a size defined by the user. The size is provided when calling *`rmt_driver_install()`* discussed above. To get a handle to this buffer call *`rmt_get_ringbuf_handle()`*.

With the above steps complete we can start the receiver by calling *`rmt_rx_start()`* and then move to checking what’s inside the buffer. To do so, you can use common FreeRTOS functions that interact with the ring buffer. Please see an example how to do it in peripherals/rmt/ir_protocols.

To stop the receiver, call *`rmt_rx_stop()`*.

Change Operation Parameters

Previously described function *rmt_config()* provides a convenient way to set several configuration parameters in one shot. This is usually done on application start. Then, when the application is running, the API provides an alternate way to update individual parameters by calling dedicated functions. Each function refers to the specific RMT channel provided as the first input parameter. Most of the functions have *_get_* counterpart to read back the currently configured value.

Parameters Common to Transmit and Receive Mode

- Selection of a GPIO pin number on the input or output of the RMT - *rmt_set_pin()*
- Number of memory blocks allocated for the incoming or outgoing data - *rmt_set_mem_pd()*
- Setting of the clock divider - *rmt_set_clk_div()*
- Selection of the clock source, note that currently one clock source is supported, the APB clock which is 80Mhz - *rmt_set_source_clk()*

Transmit Mode Parameters

- Enable or disable the loop back mode for the transmitter - *rmt_set_tx_loop_mode()*
- Binary level on the output to apply the carrier - *rmt_set_tx_carrier()*, selected from *rmt_carrier_level_t*
- Determines the binary level on the output when transmitter is idle - *rmt_set_idle_level()*, selected from *rmt_idle_level_t*

Receive Mode Parameters

- The filter setting - *rmt_set_rx_filter()*
- The receiver threshold setting - *rmt_set_rx_idle_thresh()*
- Whether the transmitter or receiver is entitled to access RMT's memory - *rmt_set_memory_owner()*, selection is from *rmt_mem_owner_t*.

Use Interrupts

Registering of an interrupt handler for the RMT controller is done be calling *rmt_isr_register()*.

注解: When calling *rmt_driver_install()* to use the system RMT driver, a default ISR is being installed. In such a case you cannot register a generic ISR handler with *rmt_isr_register()*.

The RMT controller triggers interrupts on four specific events describes below. To enable interrupts on these events, the following functions are provided:

- The RMT receiver has finished receiving a signal - *rmt_set_rx_intr_en()*
- The RMT transmitter has finished transmitting the signal - *rmt_set_tx_intr_en()*
- The number of events the transmitter has sent matches a threshold value *rmt_set_tx_thr_intr_en()*
- Ownership to the RMT memory block has been violated - *rmt_set_err_intr_en()*

Setting or clearing an interrupt enable mask for specific channels and events may be also done by calling *rmt_set_intr_enable_mask()* or *rmt_clr_intr_enable_mask()*.

When servicing an interrupt within an ISR, the interrupt need to explicitly cleared. To do so, set specific bits described as `RMT.int_clr.val.chN_event_name` and defined as a `volatile struct` in soc/soc/esp32s2/include/soc/rmt_struct.h, where N is the RMT channel number [0, n] and the `event_name` is one of four events described above.

If you do not need an ISR anymore, you can deregister it by calling a function *rmt_isr_deregister()*.

Uninstall Driver

If the RMT driver has been installed with `rmt_driver_install()` for some specific period of time and then not required, the driver may be removed to free allocated resources by calling `rmt_driver_uninstall()`.

Application Examples

- A simple RMT TX example: peripherals/rmt/morse_code.
- Another RMT TX example, specific to drive a common RGB LED strip: peripherals/rmt/led_strip.
- NEC remote control TX and RX example: peripherals/rmt/ir_protocols.

API Reference

Header File

- driver/include/driver/rmt.h

Functions

esp_err_t **rmt_set_clk_div** (*rmt_channel_t channel*, uint8_t *div_cnt*)

Set RMT clock divider, channel clock is divided from source clock.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `div_cnt`: RMT counter clock divider

esp_err_t **rmt_get_clk_div** (*rmt_channel_t channel*, uint8_t **div_cnt*)

Get RMT clock divider, channel clock is divided from source clock.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `div_cnt`: pointer to accept RMT counter divider

esp_err_t **rmt_set_rx_idle_thresh** (*rmt_channel_t channel*, uint16_t *thresh*)

Set RMT RX idle threshold value.

In receive mode, when no edge is detected on the input signal for longer than idle_thres channel clock cycles, the receive process is finished.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `thresh`: RMT RX idle threshold

esp_err_t **rmt_get_rx_idle_thresh** (*rmt_channel_t channel*, uint16_t **thresh*)

Get RMT idle threshold value.

In receive mode, when no edge is detected on the input signal for longer than idle_thres channel clock cycles, the receive process is finished.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `thresh`: pointer to accept RMT RX idle threshold value

esp_err_t **rmt_set_mem_block_num** (*rmt_channel_t channel*, uint8_t *rmt_mem_num*)

Set RMT memory block number for RMT channel.

This function is used to configure the amount of memory blocks allocated to channel n The 8 channels share a 512x32-bit RAM block which can be read and written by the processor cores over the APB bus, as well as read by the transmitters and written by the receivers.

The RAM address range for channel n is start_addr_CHn to end_addr_CHn, which are defined by: Memory block start address is RMT_CHANNEL_MEM(n) (in soc/rmt_reg.h), that is, start_addr_chn = RMT base address + 0x800 + 64 ∗ 4 ∗ n, and end_addr_chn = RMT base address + 0x800 + 64 ∗ 4 ∗ n + 64 ∗ 4 ∗ RMT_MEM_SIZE_CHn mod 512 ∗ 4

Note If memory block number of one channel is set to a value greater than 1, this channel will occupy the memory block of the next channel. Channel 0 can use at most 8 blocks of memory, accordingly channel 7 can only use one memory block.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `rmt_mem_num`: RMT RX memory block number, one block has 64 * 32 bits.

esp_err_t **rmt_get_mem_block_num** (*rmt_channel_t channel*, uint8_t **rmt_mem_num*)

Get RMT memory block number.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `rmt_mem_num`: Pointer to accept RMT RX memory block number

esp_err_t **rmt_set_tx_carrier** (*rmt_channel_t channel*, bool *carrier_en*, uint16_t *high_level*, uint16_t *low_level*, *rmt_carrier_level_t carrier_level*)

Configure RMT carrier for TX signal.

Set different values for carrier_high and carrier_low to set different frequency of carrier. The unit of carrier_high/low is the source clock tick, not the divided channel counter clock.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `carrier_en`: Whether to enable output carrier.
- `high_level`: High level duration of carrier
- `low_level`: Low level duration of carrier.
- `carrier_level`: Configure the way carrier wave is modulated for channel.
  - 1’ b1:transmit on low output level
  - 1’ b0:transmit on high output level

esp_err_t **rmt_set_mem_pd** (*rmt_channel_t channel*, bool *pd_en*)

Set RMT memory in low power mode.

Reduce power consumed by memory. 1:memory is in low power state.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel

- `pd_en`: RMT memory low power enable.

esp_err_t **rmt_get_mem_pd**(*rmt_channel_t channel*, bool **pd_en*)

Get RMT memory low power mode.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `pd_en`: Pointer to accept RMT memory low power mode.

esp_err_t **rmt_tx_start**(*rmt_channel_t channel*, bool *tx_idx_rst*)

Set RMT start sending data from memory.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `tx_idx_rst`: Set true to reset memory index for TX. Otherwise, transmitter will continue sending from the last index in memory.

esp_err_t **rmt_tx_stop**(*rmt_channel_t channel*)

Set RMT stop sending.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel

esp_err_t **rmt_rx_start**(*rmt_channel_t channel*, bool *rx_idx_rst*)

Set RMT start receiving data.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `rx_idx_rst`: Set true to reset memory index for receiver. Otherwise, receiver will continue receiving data to the last index in memory.

esp_err_t **rmt_rx_stop**(*rmt_channel_t channel*)

Set RMT stop receiving data.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel

esp_err_t **rmt_memory_rw_rst**(*rmt_channel_t channel*)

Reset RMT TX/RX memory index.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel

esp_err_t **rmt_set_memory_owner**(*rmt_channel_t channel*, *rmt_mem_owner_t owner*)

Set RMT memory owner.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `owner`: To set when the transmitter or receiver can process the memory of channel.

esp_err_t **rmt_get_memory_owner**(*rmt_channel_t channel*, *rmt_mem_owner_t* **owner*)

Get RMT memory owner.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `owner`: Pointer to get memory owner.

esp_err_t **rmt_set_tx_loop_mode**(*rmt_channel_t channel*, bool *loop_en*)

Set RMT tx loop mode.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `loop_en`: Enable RMT transmitter loop sending mode. If set true, transmitter will continue sending from the first data to the last data in channel over and over again in a loop.

esp_err_t **rmt_get_tx_loop_mode**(*rmt_channel_t channel*, bool **loop_en*)

Get RMT tx loop mode.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `loop_en`: Pointer to accept RMT transmitter loop sending mode.

esp_err_t **rmt_set_rx_filter**(*rmt_channel_t channel*, bool *rx_filter_en*, uint8_t *thresh*)

Set RMT RX filter.

In receive mode, channel will ignore input pulse when the pulse width is smaller than threshold. Counted in source clock, not divided counter clock.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `rx_filter_en`: To enable RMT receiver filter.
- `thresh`: Threshold of pulse width for receiver.

esp_err_t **rmt_set_source_clk**(*rmt_channel_t channel*, *rmt_source_clk_t base_clk*)

Set RMT source clock.

RMT module has two clock sources:

1. APB clock which is 80Mhz
2. REF tick clock, which would be 1Mhz (not supported in this version).

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `base_clk`: To choose source clock for RMT module.

esp_err_t **rmt_get_source_clk**(*rmt_channel_t channel*, *rmt_source_clk_t *src_clk*)
Get RMT source clock.

RMT module has two clock sources:

1. APB clock which is 80Mhz
2. REF tick clock, which would be 1Mhz (not supported in this version).

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `src_clk`: Pointer to accept source clock for RMT module.

esp_err_t **rmt_set_idle_level**(*rmt_channel_t channel*, bool *idle_out_en*, *rmt_idle_level_t level*)
Set RMT idle output level for transmitter.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `idle_out_en`: To enable idle level output.
- `level`: To set the output signal’s level for channel in idle state.

esp_err_t **rmt_get_idle_level**(*rmt_channel_t channel*, bool **idle_out_en*, *rmt_idle_level_t *level*)
Get RMT idle output level for transmitter.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `idle_out_en`: Pointer to accept value of enable idle.
- `level`: Pointer to accept value of output signal’s level in idle state for specified channel.

esp_err_t **rmt_get_status**(*rmt_channel_t channel*, uint32_t **status*)
Get RMT status.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `status`: Pointer to accept channel status. Please refer to RMT_CHnSTATUS_REG(n=0~7) in `rmt_reg.h` for more details of each field.

void **rmt_set_intr_enable_mask**(uint32_t *mask*)
Set mask value to RMT interrupt enable register.

Parameters
- `mask`: Bit mask to set to the register

void **rmt_clr_intr_enable_mask**(uint32_t *mask*)
Clear mask value to RMT interrupt enable register.

Parameters
- `mask`: Bit mask to clear the register

esp_err_t **rmt_set_rx_intr_en**(*rmt_channel_t channel*, bool *en*)
Set RMT RX interrupt enable.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `en`: enable or disable RX interrupt.

esp_err_t **rmt_set_err_intr_en** (*rmt_channel_t channel*, bool *en*)
Set RMT RX error interrupt enable.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `en`: enable or disable RX err interrupt.

esp_err_t **rmt_set_tx_intr_en** (*rmt_channel_t channel*, bool *en*)
Set RMT TX interrupt enable.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `en`: enable or disable TX interrupt.

esp_err_t **rmt_set_tx_thr_intr_en** (*rmt_channel_t channel*, bool *en*, uint16_t *evt_thresh*)
Set RMT TX threshold event interrupt enable.

An interrupt will be triggered when the number of transmitted items reaches the threshold value

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `en`: enable or disable TX event interrupt.
- `evt_thresh`: RMT event interrupt threshold value

esp_err_t **rmt_set_pin** (*rmt_channel_t channel*, *rmt_mode_t mode*, *gpio_num_t gpio_num*)
Set RMT pin.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel
- `mode`: TX or RX mode for RMT
- `gpio_num`: GPIO number to transmit or receive the signal.

esp_err_t **rmt_config**(**const** *rmt_config_t *rmt_param*)
Configure RMT parameters.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `rmt_param`: RMT parameter struct

esp_err_t **rmt_isr_register** (void (**fn*)) void *
, void **arg*, int *intr_alloc_flags*, *rmt_isr_handle_t *handle*Register RMT interrupt handler, the handler is an ISR.

The handler will be attached to the same CPU core that this function is running on.

Note If you already called rmt_driver_install to use system RMT driver, please do not register ISR handler again.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Function pointer error.
- ESP_FAIL System driver installed, can not register ISR handler for RMT

Parameters

- `fn`: Interrupt handler function.
- `arg`: Parameter for the handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `handle`: If non-zero, a handle to later clean up the ISR gets stored here.

esp_err_t **rmt_isr_deregister**(*rmt_isr_handle_t handle*)

Deregister previously registered RMT interrupt handler.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Handle invalid

Parameters

- `handle`: Handle obtained from rmt_isr_register

esp_err_t **rmt_fill_tx_items**(*rmt_channel_t channel*, **const** rmt_item32_t **item*, uint16_t *item_num*, uint16_t *mem_offset*)

Fill memory data of channel with given RMT items.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `item`: Pointer of items.
- `item_num`: RMT sending items number.
- `mem_offset`: Index offset of memory.

esp_err_t **rmt_driver_install**(*rmt_channel_t channel*, size_t *rx_buf_size*, int *intr_alloc_flags*)

Initialize RMT driver.

Return

- ESP_ERR_INVALID_STATE Driver is already installed, call rmt_driver_uninstall first.
- ESP_ERR_NO_MEM Memory allocation failure
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `rx_buf_size`: Size of RMT RX ringbuffer. Can be 0 if the RX ringbuffer is not used.
- `intr_alloc_flags`: Flags for the RMT driver interrupt handler. Pass 0 for default flags. See esp_intr_alloc.h for details. If ESP_INTR_FLAG_IRAM is used, please do not use the memory allocated from psram when calling rmt_write_items.

esp_err_t **rmt_driver_uninstall**(*rmt_channel_t channel*)

Uninstall RMT driver.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel

esp_err_t **rmt_get_channel_status**(*rmt_channel_status_result_t *channel_status*)

Get the current status of eight channels.

Note Do not call this function if it is possible that `rmt_driver_uninstall` will be called at the same time.

Return

- ESP_ERR_INVALID_ARG Parameter is NULL
- ESP_OK Success

Parameters

- `[out] channel_status`: store the current status of each channel

esp_err_t **rmt_get_counter_clock**(*rmt_channel_t channel*, uint32_t **clock_hz*)

Get speed of channel’ s internal counter clock.

Return

- ESP_ERR_INVALID_ARG Parameter is NULL
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `[out] clock_hz`: counter clock speed, in hz

esp_err_t **rmt_write_items**(*rmt_channel_t channel*, **const** rmt_item32_t **rmt_item*, int *item_num*, bool *wait_tx_done*)

RMT send waveform from rmt_item array.

This API allows user to send waveform with any length.

Note This function will not copy data, instead, it will point to the original items, and send the waveform items. If wait_tx_done is set to true, this function will block and will not return until all items have been sent out. If wait_tx_done is set to false, this function will return immediately, and the driver interrupt will continue sending the items. We must make sure the item data will not be damaged when the driver is still sending items in driver interrupt.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `rmt_item`: head point of RMT items array. If ESP_INTR_FLAG_IRAM is used, please do not use the memory allocated from psram when calling rmt_write_items.
- `item_num`: RMT data item number.
- `wait_tx_done`:
  - If set 1, it will block the task and wait for sending done.
  - If set 0, it will not wait and return immediately.

esp_err_t **rmt_wait_tx_done**(*rmt_channel_t channel*, TickType_t *wait_time*)

Wait RMT TX finished.

Return

- ESP_OK RMT Tx done successfully
- ESP_ERR_TIMEOUT Exceeded the ‘wait_time’ given
- ESP_ERR_INVALID_ARG Parameter error
- ESP_FAIL Driver not installed

Parameters

- `channel`: RMT channel
- `wait_time`: Maximum time in ticks to wait for transmission to be complete. If set 0, return immediately with ESP_ERR_TIMEOUT if TX is busy (polling).

esp_err_t **rmt_get_ringbuf_handle**(*rmt_channel_t channel*, *RingbufHandle_t *buf_handle*)

Get ringbuffer from RMT.

Users can get the RMT RX ringbuffer handle, and process the RX data.

Return

- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters

- `channel`: RMT channel
- `buf_handle`: Pointer to buffer handle to accept RX ringbuffer handle.

esp_err_t **rmt_translator_init** (*rmt_channel_t channel*, *sample_to_rmt_t fn*)
Init rmt translator and register user callback. The callback will convert the raw data that needs to be sent to rmt format. If a channel is initialized more than once, tha user callback will be replaced by the later.

Return
- ESP_FAIL Init fail.
- ESP_OK Init success.

Parameters
- `channel`: RMT channel .
- `fn`: Point to the data conversion function.

esp_err_t **rmt_write_sample** (*rmt_channel_t channel*, **const** uint8_t **src*, size_t *src_size*, bool *wait_tx_done*)
Translate uint8_t type of data into rmt format and send it out. Requires rmt_translator_init to init the translator first.

Return
- ESP_FAIL Send fail
- ESP_OK Send success

Parameters
- `channel`: RMT channel .
- `src`: Pointer to the raw data.
- `src_size`: The size of the raw data.
- `wait_tx_done`: Set true to wait all data send done.

rmt_tx_end_callback_t **rmt_register_tx_end_callback** (*rmt_tx_end_fn_t function*, void **arg*)
Registers a callback that will be called when transmission ends.

Called by rmt_driver_isr_default in interrupt context.

Note Requires rmt_driver_install to install the default ISR handler.
Return the previous callback settings (members will be set to NULL if there was none)
Parameters
- `function`: Function to be called from the default interrupt handler or NULL.
- `arg`: Argument which will be provided to the callback when it is called.

esp_err_t **rmt_add_channel_to_group** (*rmt_channel_t channel*)
Add channel into a group (channels in the same group will transmit simultaneously)

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel

esp_err_t **rmt_remove_channel_from_group** (*rmt_channel_t channel*)
Remove channel out of a group.

Return
- ESP_ERR_INVALID_ARG Parameter error
- ESP_OK Success

Parameters
- `channel`: RMT channel

Structures

struct rmt_tx_config_t
Data struct of RMT TX configure parameters.

Public Members

uint32_t **carrier_freq_hz**
RMT carrier frequency

rmt_carrier_level_t **carrier_level**
Level of the RMT output, when the carrier is applied

rmt_idle_level_t **idle_level**
RMT idle level

uint8_t **carrier_duty_percent**
RMT carrier duty (%)

uint32_t **loop_count**
Maximum loop count

bool **carrier_en**
RMT carrier enable

bool **loop_en**
Enable sending RMT items in a loop

bool **idle_output_en**
RMT idle level output enable

struct rmt_rx_config_t
Data struct of RMT RX configure parameters.

Public Members

uint16_t **idle_threshold**
RMT RX idle threshold

uint8_t **filter_ticks_thresh**
RMT filter tick number

bool **filter_en**
RMT receiver filter enable

bool **rm_carrier**
RMT receiver remove carrier enable

uint32_t **carrier_freq_hz**
RMT carrier frequency

uint8_t **carrier_duty_percent**
RMT carrier duty (%)

rmt_carrier_level_t **carrier_level**
The level to remove the carrier

struct rmt_config_t
Data struct of RMT configure parameters.

Public Members

rmt_mode_t **rmt_mode**
RMT mode: transmitter or receiver

rmt_channel_t **channel**
RMT channel

gpio_num_t **gpio_num**
RMT GPIO number

uint8_t **clk_div**
RMT channel counter divider

uint8_t **mem_block_num**
RMT memory block number

uint32_t **flags**
RMT channel extra configurations, OR' d with RMT_CHANNEL_FLAGS_[*]

rmt_tx_config_t **tx_config**
RMT TX parameter

rmt_rx_config_t **rx_config**
RMT RX parameter

struct rmt_tx_end_callback_t
Structure encapsulating a RMT TX end callback.

Public Members

rmt_tx_end_fn_t **function**
Function which is called on RMT TX end

void ***arg**
Optional argument passed to function

Macros

RMT_CHANNEL_FLAGS_ALWAYS_ON
Channel can work when APB frequency is changing (RMT channel adopts REF_TICK as clock source)

RMT_MEM_ITEM_NUM
Define memory space of each RMT channel (in words = 4 bytes)

RMT_DEFAULT_CONFIG_TX(gpio, channel_id)
Default configuration for Tx channel.

RMT_DEFAULT_CONFIG_RX(gpio, channel_id)
Default configuration for RX channel.

Type Definitions

typedef *intr_handle_t* **rmt_isr_handle_t**
RMT interrupt handle.

typedef void (***rmt_tx_end_fn_t**)(*rmt_channel_t* channel, void *arg)
Type of RMT Tx End callback function.

typedef void (***sample_to_rmt_t**)(**const** void *src, rmt_item32_t *dest, size_t src_size, size_t wanted_num, size_t *translated_size, size_t *item_num)
User callback function to convert uint8_t type data to rmt format(rmt_item32_t).

This function may be called from an ISR, so, the code should be short and efficient.

Note In fact, item_num should be a multiple of translated_size, e.g. : When we convert each byte of uint8_t type data to rmt format data, the relation between item_num and translated_size should be `item_num = translated_size*8`.

Parameters

- `src`: Pointer to the buffer storing the raw data that needs to be converted to rmt format.
- `[out] dest`: Pointer to the buffer storing the rmt format data.
- `src_size`: The raw data size.
- `wanted_num`: The number of rmt format data that wanted to get.
- `[out] translated_size`: The size of the raw data that has been converted to rmt format, it should return 0 if no data is converted in user callback.
- `[out] item_num`: The number of the rmt format data that actually converted to, it can be less than wanted_num if there is not enough raw data, but cannot exceed wanted_num. it should return 0 if no data was converted.

Header File

- soc/include/hal/rmt_types.h

Structures

struct rmt_channel_status_result_t

Data struct of RMT channel status.

Public Members

rmt_channel_status_t **status**[**RMT_CHANNEL_MAX**]

Store the current status of each channel

Enumerations

enum rmt_channel_t

RMT channel ID.

Values:

RMT_CHANNEL_0

RMT channel number 0

RMT_CHANNEL_1

RMT channel number 1

RMT_CHANNEL_2

RMT channel number 2

RMT_CHANNEL_3

RMT channel number 3

RMT_CHANNEL_MAX

Number of RMT channels

enum rmt_mem_owner_t

RMT Internal Memory Owner.

Values:

RMT_MEM_OWNER_TX

RMT RX mode, RMT transmitter owns the memory block

RMT_MEM_OWNER_RX

RMT RX mode, RMT receiver owns the memory block

RMT_MEM_OWNER_MAX

enum rmt_source_clk_t

Clock Source of RMT Channel.

Values:

RMT_BASECLK_REF

RMT source clock is REF_TICK, 1MHz by default

RMT_BASECLK_APB

RMT source clock is APB CLK, 80Mhz by default

RMT_BASECLK_MAX

enum rmt_data_mode_t

RMT Data Mode.

Note We highly recommended to use MEM mode not FIFO mode since there will be some gotcha in FIFO mode.

Values:

RMT_DATA_MODE_FIFO

RMT_DATA_MODE_MEM

RMT_DATA_MODE_MAX

enum rmt_mode_t
RMT Channel Working Mode (TX or RX)

Values:

RMT_MODE_TX
RMT TX mode

RMT_MODE_RX
RMT RX mode

RMT_MODE_MAX

enum rmt_idle_level_t
RMT Idle Level.

Values:

RMT_IDLE_LEVEL_LOW
RMT TX idle level: low Level

RMT_IDLE_LEVEL_HIGH
RMT TX idle level: high Level

RMT_IDLE_LEVEL_MAX

enum rmt_carrier_level_t
RMT Carrier Level.

Values:

RMT_CARRIER_LEVEL_LOW
RMT carrier wave is modulated for low Level output

RMT_CARRIER_LEVEL_HIGH
RMT carrier wave is modulated for high Level output

RMT_CARRIER_LEVEL_MAX

enum rmt_channel_status_t
RMT Channel Status.

Values:

RMT_CHANNEL_UNINIT
RMT channel uninitialized

RMT_CHANNEL_IDLE
RMT channel status idle

RMT_CHANNEL_BUSY
RMT channel status busy

2.2.11 SD SPI Host Driver

Overview

The SD SPI host driver allows communicating with one or more SD cards by the SPI Master driver which makes use of the SPI host. Each card is accessed through an SD SPI device represented by an *sdspi_dev_handle_t* spi_handle returned when attaching the device to an SPI bus by calling *sdspi_host_init_device*. The bus should be already initialized before (by *spi_bus_initialize*).

With the help of *SPI Master driver* based on, the SPI bus can be shared among SD cards and other SPI devices. The SPI Master driver will handle exclusive access from different tasks.

The SD SPI driver uses software-controlled CS signal.

How to Use

Firstly, use the macro `SDSPI_DEVICE_CONFIG_DEFAULT` to initialize a structure `sdmmc_slot_config_t`, which is used to initialize an SD SPI device. This macro will also fill in the default pin mappings, which is same as the pin mappings of SDMMC host driver. Modify the host and pins of the structure to desired value. Then call *sdspi_host_init_device* to initialize the SD SPI device and attach to its bus.

Then use `SDSPI_HOST_DEFAULT` macro to initialize a `sdmmc_host_t` structure, which is used to store the state and configurations of upper layer (SD/SDIO/MMC driver). Modify the *slot* parameter of the structure to the SD SPI device spi_handle just returned from *sdspi_host_init_device*. Call *sdmmc_card_init* with the `sdmmc_host_t` to probe and initialize the SD card.

Now you can use SD/SDIO/MMC driver functions to access your card!

Other Details

Only the following driver's API functions are normally used by most applications:

- `sdspi_host_init()`
- `sdspi_host_init_device()`
- `sdspi_host_remove_device()`
- `sdspi_host_deinit()`

Other functions are mostly used by the protocol level SD/SDIO/MMC driver via function pointers in the `sdmmc_host_t` structure. For more details, see *the SD/SDIO/MMC Driver*.

注解: SD over SPI does not support speeds above `SDMMC_FREQ_DEFAULT` due to the limitations of the SPI driver.

API Reference

Header File

- driver/include/driver/sdspi_host.h

Functions

esp_err_t **sdspi_host_init**(void)

Initialize SD SPI driver.

Note This function is not thread safe

Return

- ESP_OK on success
- other error codes may be returned in future versions

esp_err_t **sdspi_host_init_device**(**const** *sdspi_device_config_t* *dev_config, *sdspi_dev_handle_t* *out_handle)

Attach and initialize an SD SPI device on the specific SPI bus.

Note This function is not thread safe

Note Initialize the SPI bus by `spi_bus_initialize()` before calling this function.

Note The SDIO over sdspi needs an extra interrupt line. Call `gpio_install_isr_service()` before this function.

Return

- ESP_OK on success

- ESP_ERR_INVALID_ARG if sdspi_host_init_device has invalid arguments
- ESP_ERR_NO_MEM if memory can not be allocated
- other errors from the underlying spi_master and gpio drivers

Parameters

- `dev_config`: pointer to device configuration structure
- `out_handle`: Output of the handle to the sdspi device.

esp_err_t **sdspi_host_remove_device**(*sdspi_dev_handle_t handle*)

Remove an SD SPI device.

Return Always ESP_OK

Parameters

- `handle`: Handle of the SD SPI device

esp_err_t **sdspi_host_do_transaction**(*sdspi_dev_handle_t handle*, *sdmmc_command_t *cmdinfo*)

Send command to the card and get response.

This function returns when command is sent and response is received, or data is transferred, or timeout occurs.

Note This function is not thread safe w.r.t. init/deinit functions, and bus width/clock speed configuration functions. Multiple tasks can call sdspi_host_do_transaction as long as other sdspi_host_* functions are not called.

Return

- ESP_OK on success
- ESP_ERR_TIMEOUT if response or data transfer has timed out
- ESP_ERR_INVALID_CRC if response or data transfer CRC check has failed
- ESP_ERR_INVALID_RESPONSE if the card has sent an invalid response

Parameters

- `handle`: Handle of the sdspi device
- `cmdinfo`: pointer to structure describing command and data to transfer

esp_err_t **sdspi_host_set_card_clk**(*sdspi_dev_handle_t host*, uint32_t *freq_khz*)

Set card clock frequency.

Currently only integer fractions of 40MHz clock can be used. For High Speed cards, 40MHz can be used. For Default Speed cards, 20MHz can be used.

Note This function is not thread safe

Return

- ESP_OK on success
- other error codes may be returned in the future

Parameters

- `host`: Handle of the sdspi device
- `freq_khz`: card clock frequency, in kHz

esp_err_t **sdspi_host_deinit**(void)

Release resources allocated using sdspi_host_init.

Note This function is not thread safe

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if sdspi_host_init function has not been called

esp_err_t **sdspi_host_io_int_enable**(*sdspi_dev_handle_t handle*)

Enable SDIO interrupt.

Return

- ESP_OK on success

Parameters

- `handle`: Handle of the sdspi device

esp_err_t **sdspi_host_io_int_wait**(*sdspi_dev_handle_t handle*, TickType_t *timeout_ticks*)

Wait for SDIO interrupt until timeout.

Return

- ESP_OK on success

Parameters

- `handle`: Handle of the sdspi device
- `timeout_ticks`: Ticks to wait before timeout.

esp_err_t **sdspi_host_init_slot**(int *slot*, **const** *sdspi_slot_config_t* **slot_config*)

Initialize SD SPI driver for the specific SPI controller.

Note This function is not thread safe

Note The SDIO over sdspi needs an extra interrupt line. Call `gpio_install_isr_service()` before this function.

Parameters

- `slot`: SPI controller to use (HSPI_HOST or VSPI_HOST)
- `slot_config`: pointer to slot configuration structure

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if sdspi_init_slot has invalid arguments
- ESP_ERR_NO_MEM if memory can not be allocated
- other errors from the underlying spi_master and gpio drivers

Structures

struct sdspi_device_config_t

Extra configuration for SD SPI device.

Public Members

spi_host_device_t **host_id**

SPI host to use, SPIx_HOST (see spi_types.h).

gpio_num_t **gpio_cs**

GPIO number of CS signal.

gpio_num_t **gpio_cd**

GPIO number of card detect signal.

gpio_num_t **gpio_wp**

GPIO number of write protect signal.

gpio_num_t **gpio_int**

GPIO number of interrupt line (input) for SDIO card.

struct sdspi_slot_config_t

Extra configuration for SPI host.

Public Members

gpio_num_t **gpio_cs**

GPIO number of CS signal.

gpio_num_t **gpio_cd**

GPIO number of card detect signal.

gpio_num_t **gpio_wp**

GPIO number of write protect signal.

gpio_num_t **gpio_int**

GPIO number of interrupt line (input) for SDIO card.

gpio_num_t **gpio_miso**

GPIO number of MISO signal.

gpio_num_t **gpio_mosi**
GPIO number of MOSI signal.

gpio_num_t **gpio_sck**
GPIO number of SCK signal.

int **dma_channel**
DMA channel to be used by SPI driver (1 or 2).

Macros

SDSPI_HOST_DEFAULT()
Default *sdmmc_host_t* structure initializer for SD over SPI driver.

Uses SPI mode and max frequency set to 20MHz

'slot' should be set to an sdspi device initialized by `sdspi_host_init_device()`.

SDSPI_SLOT_NO_CD
indicates that card detect line is not used

SDSPI_SLOT_NO_WP
indicates that write protect line is not used

SDSPI_SLOT_NO_INT
indicates that interrupt line is not used

SDSPI_DEVICE_CONFIG_DEFAULT()
Macro defining default configuration of SD SPI device.

SDSPI_SLOT_CONFIG_DEFAULT()
Macro defining default configuration of SPI host

Type Definitions

typedef int **sdspi_dev_handle_t**
Handle representing an SD SPI device.

2.2.12 Sigma-delta Modulation

Introduction

ESP32-S2 has a second-order sigma-delta modulation module. This driver configures the channels of the sigma-delta module.

Functionality Overview

There are eight independent sigma-delta modulation channels identified with *sigmadelta_channel_t*. Each channel is capable to output the binary, hardware generated signal with the sigma-delta modulation.

Selected channel should be set up by providing configuration parameters in *sigmadelta_config_t* and then applying this configuration with *sigmadelta_config()*.

Another option is to call individual functions, that will configure all required parameters one by one:

- **Prescaler** of the sigma-delta generator - *sigmadelta_set_prescale()*
- **Duty** of the output signal - *sigmadelta_set_duty()*
- **GPIO pin** to output modulated signal - *sigmadelta_set_pin()*

The range of the 'duty' input parameter of *sigmadelta_set_duty()* is from -128 to 127 (eight bit signed integer). If zero value is set, then the output signal's duty will be about 50%, see description of *sigmadelta_set_duty()*.

Application Example

Sigma-delta Modulation example: peripherals/sigmadelta.

API Reference

Header File

- driver/include/driver/sigmadelta.h

Functions

esp_err_t **sigmadelta_config**(**const** *sigmadelta_config_t* **config*)

Configure Sigma-delta channel.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE sigmadelta driver already initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `config`: Pointer of Sigma-delta channel configuration struct

esp_err_t **sigmadelta_set_duty**(*sigmadelta_channel_t channel*, int8_t *duty*)

Set Sigma-delta channel duty.

This function is used to set Sigma-delta channel duty, If you add a capacitor between the output pin and ground, the average output voltage will be Vdc = VDDIO / 256 * duty + VDDIO/2, where VDDIO is the power supply voltage.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE sigmadelta driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `channel`: Sigma-delta channel number
- `duty`: Sigma-delta duty of one channel, the value ranges from -128 to 127, recommended range is -90 ~ 90. The waveform is more like a random one in this range.

esp_err_t **sigmadelta_set_prescale**(*sigmadelta_channel_t channel*, uint8_t *prescale*)

Set Sigma-delta channel' s clock pre-scale value. The source clock is APP_CLK, 80MHz. The clock frequency of the sigma-delta channel is APP_CLK / pre_scale.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE sigmadelta driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `channel`: Sigma-delta channel number
- `prescale`: The divider of source clock, ranges from 0 to 255

esp_err_t **sigmadelta_set_pin**(*sigmadelta_channel_t channel*, *gpio_num_t gpio_num*)

Set Sigma-delta signal output pin.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE sigmadelta driver has not been initialized
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `channel`: Sigma-delta channel number
- `gpio_num`: GPIO number of output pin.

Header File

- soc/include/hal/sigmadelta_types.h

Structures

struct sigmadelta_config_t

Sigma-delta configure struct.

Public Members

sigmadelta_channel_t **channel**

Sigma-delta channel number

int8_t **sigmadelta_duty**

Sigma-delta duty, duty ranges from -128 to 127.

uint8_t **sigmadelta_prescale**

Sigma-delta prescale, prescale ranges from 0 to 255.

uint8_t **sigmadelta_gpio**

Sigma-delta output io number, refer to gpio.h for more details.

Type Definitions

typedef int **sigmadelta_port_t**

SIGMADELTA port number, the max port number is (SIGMADELTA_NUM_MAX -1).

typedef int **sigmadelta_channel_t**

Sigma-delta channel list.

2.2.13 SPI Master Driver

SPI Master driver is a program that controls ESP32-S2’s SPI peripherals while they function as masters.

Overview of ESP32-S2’s SPI peripherals

ESP32 integrates four SPI peripherals.

- SPI0 and SPI1 are used internally to access the ESP32’s attached flash memory and share an arbiter. Currently SPI Master driver hasn’t supported SPI1 bus.
- SPI2 and SPI3 are general purpose SPI controllers, sometimes referred to as HSPI and VSPI, respectively. They are open to users. SPI2 and SPI3 have independent signal buses with the same respective names. Each bus has three CS lines to drive up to three SPI slaves.

Terminology

The terms used in relation to the SPI master driver are given in the table below.

Term	Definition
Host	The SPI controller peripheral inside ESP32 that initiates SPI transmissions over the bus, and acts as an SPI Master. This may be the SPI2 or SPI3 peripheral. (The driver will also support the SPI1 peripheral in the future.)
Device	SPI slave device. An SPI bus may be connected to one or more Devices. Each Device shares the MOSI, MISO and SCLK signals but is only active on the bus when the Host asserts the Device's individual CS line.
Bus	A signal bus, common to all Devices connected to one Host. In general, a bus includes the following lines: MISO, MOSI, SCLK, one or more CS lines, and, optionally, QUADWP and QUADHD. So Devices are connected to the same lines, with the exception that each Device has its own CS line. Several Devices can also share one CS line if connected in the daisy-chain manner.
• **MISO**	Master In, Slave Out, a.k.a. Q. Data transmission from a Device to Host.
• **MOSI**	Master Out, Slave In, a.k.a. D. Data transmission from a Host to Device.
• **SCLK**	Serial Clock. Oscillating signal generated by a Host that keeps the transmission of data bits in sync.
• **CS**	Chip Select. Allows a Host to select individual Device(s) connected to the bus in order to send or receive data.
• **QUADWP**	Write Protect signal. Only used for 4-bit (qio/qout) transactions.
• **QUADHD**	Hold signal. Only used for 4-bit (qio/qout) transactions.
• **Assertion**	The action of activating a line. The opposite action of returning the line back to inactive (back to idle) is called *de-assertion*.
Transaction	One instance of a Host asserting a CS line, transferring data to and from a Device, and de-asserting the CS line. Transactions are atomic, which means they can never be interrupted by another transaction.
Launch edge	Edge of the clock at which the source register *launches* the signal onto the line.
Latch edge	Edge of the clock at which the destination register *latches in* the signal.

Driver Features

The SPI master driver governs communications of Hosts with Devices. The driver supports the following features:

- Multi-threaded environments
- Transparent handling of DMA transfers while reading and writing data
- Automatic time-division multiplexing of data coming from different Devices on the same signal bus, see *SPI Bus Lock*.

警告： The SPI master driver has the concept of multiple Devices connected to a single bus (sharing a single ESP32-S2 SPI peripheral). As long as each Device is accessed by only one task, the driver is thread safe. However, if multiple tasks try to access the same SPI Device, the driver is **not thread-safe**. In this case, it is recommended to either:

- Refactor your application so that each SPI peripheral is only accessed by a single task at a time.
- Add a mutex lock around the shared Device using *xSemaphoreCreateMutex*.

SPI Features

SPI Master

SPI Bus Lock To realize the multiplexing of different devices from different drivers (SPI Master, SPI Flash, etc.), an SPI bus lock is applied on each SPI bus. Drivers can attach their devices onto the bus with the arbitration of the lock.

Each bus lock are initialized with a BG (background) service registered, all devices request to do transactions on the bus should wait until the BG to be successfully disabled.

- For SPI1 bus, the BG is the cache, the bus lock will help to disable the cache before device operations starts, and enable it again after device releasing the lock. No devices on SPI1 is allowed using ISR (it's meaningless for the task to yield to other tasks when the cache is disabled).
  The SPI Master driver hasn't supported SPI1 bus. Only SPI Flash driver can attach to the bus.
- For other buses, the driver may register its ISR as the BG. The bus lock will block a device task when it requests for exclusive use of the bus, try to disable the ISR, and unblock the device task allowed to exclusively use the bus when the ISR is successfully disabled. When the task releases the lock, the lock will also try to resume the ISR if there are pending transactions to be done in the ISR.

SPI Transactions

An SPI bus transaction consists of five phases which can be found in the table below. Any of these phases can be skipped.

Phase	Description
Command	In this phase, a command (0-16 bit) is written to the bus by the Host.
Address	In this phase, an address (0-64 bit) is transmitted over the bus by the Host.
Write	Host sends data to a Device. This data follows the optional command and address phases and is indistinguishable from them at the electrical level.
Dummy	This phase is configurable and is used to meet the timing requirements.
Read	Device sends data to its Host.

The attributes of a transaction are determined by the bus configuration structure *spi_bus_config_t*, device configuration structure *spi_device_interface_config_t*, and transaction configuration structure *spi_transaction_t*.

An SPI Host can send full-duplex transactions, during which the read and write phases occur simultaneously. The total transaction length is determined by the sum of the following members:

- *spi_device_interface_config_t::command_bits*
- *spi_device_interface_config_t::address_bits*
- *spi_transaction_t::length*

While the member *spi_transaction_t::rxlength* only determines the length of data received into the buffer.

In half-duplex transactions, the read and write phases are not simultaneous (one direction at a time). The lengths of the write and read phases are determined by `length` and `rxlength` members of the struct *`spi_transaction_t`* respectively.

The command and address phases are optional, as not every SPI device requires a command and/or address. This is reflected in the Device's configuration: if `command_bits` and/or `address_bits` are set to zero, no command or address phase will occur.

The read and write phases can also be optional, as not every transaction requires both writing and reading data. If `rx_buffer` is NULL and `SPI_TRANS_USE_RXDATA` is not set, the read phase is skipped. If `tx_buffer` is NULL and `SPI_TRANS_USE_TXDATA` is not set, the write phase is skipped.

The driver supports two types of transactions: the interrupt transactions and polling transactions. The programmer can choose to use a different transaction type per Device. If your Device requires both transaction types, see *Notes on Sending Mixed Transactions to the Same Device*.

Interrupt Transactions Interrupt transactions will block the transaction routine until the transaction completes, thus allowing the CPU to run other tasks.

An application task can queue multiple transactions, and the driver will automatically handle them one-by-one in the interrupt service routine (ISR). It allows the task to switch to other procedures until all the transactions complete.

Polling Transactions Polling transactions do not use interrupts. The routine keeps polling the SPI Host's status bit until the transaction is finished.

All the tasks that use interrupt transactions can be blocked by the queue. At this point, they will need to wait for the ISR to run twice before the transaction is finished. Polling transactions save time otherwise spent on queue handling and context switching, which results in smaller transaction intervals. The disadvantage is that the CPU is busy while these transactions are in progress.

The *`spi_device_polling_end()`* routine needs an overhead of at least 1 us to unblock other tasks when the transaction is finished. It is strongly recommended to wrap a series of polling transactions using the functions *`spi_device_acquire_bus()`* and *`spi_device_release_bus()`* to avoid the overhead. For more information, see *Bus Acquiring*.

Command and Address Phases During the command and address phases, the members `cmd` and `addr` in the struct *`spi_transaction_t`* are sent to the bus, nothing is read at this time. The default lengths of the command and address phases are set in *`spi_device_interface_config_t`* by calling *`spi_bus_add_device()`*. If the flags `SPI_TRANS_VARIABLE_CMD` and `SPI_TRANS_VARIABLE_ADDR` in the member *`spi_transaction_t::flags`* are not set, the driver automatically sets the length of these phases to default values during Device initialization.

If the lengths of the command and address phases need to be variable, declare the struct *`spi_transaction_ext_t`*, set the flags `SPI_TRANS_VARIABLE_CMD` and/or `SPI_TRANS_VARIABLE_ADDR` in the member *`spi_transaction_ext_t::base`* and configure the rest of base as usual. Then the length of each phase will be equal to `command_bits` and `address_bits` set in the struct *`spi_transaction_ext_t`*.

Write and Read Phases Normally, the data that needs to be transferred to or from a Device will be read from or written to a chunk of memory indicated by the members `rx_buffer` and `tx_buffer` of the structure *`spi_transaction_t`*. If DMA is enabled for transfers, the buffers are required to be:

1. Allocated in DMA-capable internal memory. If *external PSRAM is enabled*, this means using `pvPortMallocCaps(size, MALLOC_CAP_DMA)`.
2. 32-bit aligned (staring from a 32-bit boundary and having a length of multiples of 4 bytes).

If these requirements are not satisfied, the transaction efficiency will be affected due to the allocation and copying of temporary buffers.

注解: Half-duplex transactions with both read and write phases are not supported when using DMA. For details and workarounds, see *Known Issues*.

Bus Acquiring Sometimes you might want to send SPI transactions exclusively and continuously so that it takes as little time as possible. For this, you can use bus acquiring, which helps to suspend transactions (both polling or interrupt) to other devices until the bus is released. To acquire and release a bus, use the functions `spi_device_acquire_bus()` and `spi_device_release_bus()`.

Driver Usage

- Initialize an SPI bus by calling the function `spi_bus_initialize()`. Make sure to set the correct I/O pins in the struct `spi_bus_config_t`. Set the signals that are not needed to `-1`.
- Register a Device connected to the bus with the driver by calling the function `spi_bus_add_device()`. Make sure to configure any timing requirements the device might need with the parameter `dev_config`. You should now have obtained the Device's handle which will be used when sending a transaction to it.
- To interact with the Device, fill one or more `spi_transaction_t` structs with any transaction parameters required. Then send the structs either using a polling transaction or an interrupt transaction:
  - ***Interrupt*** Either queue all transactions by calling the function `spi_device_queue_trans()` and, at a later time, query the result using the function `spi_device_get_trans_result()`, or handle all requests synchronously by feeding them into `spi_device_transmit()`.
  - ***Polling*** Call the function `spi_device_polling_transmit()` to send polling transactions. Alternatively, if you want to insert something in between, send the transactions by using `spi_device_polling_start()` and `spi_device_polling_end()`.
- (Optional) To perform back-to-back transactions with a Device, call the function `spi_device_acquire_bus()` before sending transactions and `spi_device_release_bus()` after the transactions have been sent.
- (Optional) To unload the driver for a certain Device, call `spi_bus_remove_device()` with the Device handle as an argument.
- (Optional) To remove the driver for a bus, make sure no more drivers are attached and call `spi_bus_free()`.

The example code for the SPI master driver can be found in the peripherals/spi_master directory of ESP-IDF examples.

Transactions with Data Not Exceeding 32 Bits When the transaction data size is equal to or less than 32 bits, it will be sub-optimal to allocate a buffer for the data. The data can be directly stored in the transaction struct instead. For transmitted data, it can be achieved by using the `tx_data` member and setting the `SPI_TRANS_USE_TXDATA` flag on the transmission. For received data, use `rx_data` and set `SPI_TRANS_USE_RXDATA`. In both cases, do not touch the `tx_buffer` or `rx_buffer` members, because they use the same memory locations as `tx_data` and `rx_data`.

Transactions with Integers Other Than `uint8_t` An SPI Host reads and writes data into memory byte by byte. By default, data is sent with the most significant bit (MSB) first, as LSB first used in rare cases. If a value less than 8 bits needs to be sent, the bits should be written into memory in the MSB first manner.

For example, if `0b00010` needs to be sent, it should be written into a `uint8_t` variable, and the length for reading should be set to 5 bits. The Device will still receive 8 bits with 3 additional "random" bits, so the reading must be performed correctly.

On top of that, ESP32-S2 is a little-endian chip, which means that the least significant byte of `uint16_t` and `uint32_t` variables is stored at the smallest address. Hence, if `uint16_t` is stored in memory, bits [7:0] are sent first, followed by bits [15:8].

For cases when the data to be transmitted has the size differing from `uint8_t` arrays, the following macros can be used to transform data to the format that can be sent by the SPI driver directly:

- *SPI_SWAP_DATA_TX* for data to be transmitted
- *SPI_SWAP_DATA_RX* for data received

Notes on Sending Mixed Transactions to the Same Device To reduce coding complexity, send only one type of transactions (interrupt or polling) to one Device. However, you still can send both interrupt and polling transactions alternately. The notes below explain how to do this.

The polling transactions should be initiated only after all the polling and interrupt transactions are finished.

Since an unfinished polling transaction blocks other transactions, please do not forget to call the function *spi_device_polling_end()* after *spi_device_polling_start()* to allow other transactions or to allow other Devices to use the bus. Remember that if there is no need to switch to other tasks during your polling transaction, you can initiate a transaction with *spi_device_polling_transmit()* so that it will be ended automatically.

In-flight polling transactions are disturbed by the ISR operation to accommodate interrupt transactions. Always make sure that all the interrupt transactions sent to the ISR are finished before you call *spi_device_polling_start()*. To do that, you can keep calling *spi_device_get_trans_result()* until all the transactions are returned.

To have better control of the calling sequence of functions, send mixed transactions to the same Device only within a single task.

Transfer Speed Considerations

There are three factors limiting the transfer speed:

- Transaction interval
- SPI clock frequency
- Cache miss of SPI functions, including callbacks

The main parameter that determines the transfer speed for large transactions is clock frequency. For multiple small transactions, the transfer speed is mostly determined by the length of transaction intervals.

Transaction Interval Transaction interval is the time that software requires to set up SPI peripheral registers and to copy data to FIFOs, or to set up DMA links.

Interrupt transactions allow appending extra overhead to accommodate the cost of FreeRTOS queues and the time needed for switching between tasks and the ISR.

For **interrupt transactions**, the CPU can switch to other tasks when a transaction is in progress. This saves the CPU time but increases the interval. See *Interrupt Transactions*. For **polling transactions**, it does not block the task but allows to do polling when the transaction is in progress. For more information, see *Polling Transactions*.

If DMA is enabled, setting up the linked list requires about 2 us per transaction. When a master is transferring data, it automatically reads the data from the linked list. If DMA is not enabled, the CPU has to write and read each byte from the FIFO by itself. Usually, this is faster than 2 us, but the transaction length is limited to 64 bytes for both write and read.

Typical transaction interval timings for one byte of data are given below.

	Typical Transaction Time (us)	
	Interrupt	Polling
DMA	24	8
No DMA	22	7

SPI Clock Frequency Transferring each byte takes eight times the clock period *8/fspi*. If the clock frequency is too high, the use of some functions might be limited. See *Timing Considerations*.

Cache Miss The default config puts only the ISR into the IRAM. Other SPI related functions, including the driver itself and the callback, might suffer from the cache miss and will need to wait until the code is read from the flash. Select *CONFIG_SPI_MASTER_IN_IRAM* to put the whole SPI driver into IRAM and put the entire callback(s) and its callee functions into IRAM to prevent cache miss.

For an interrupt transaction, the overall cost is *20+8n/Fspi[MHz]* [us] for n bytes transferred in one transaction. Hence, the transferring speed is: *n/(20+8n/Fspi)*. An example of transferring speed at 8 MHz clock speed is given in the following table.

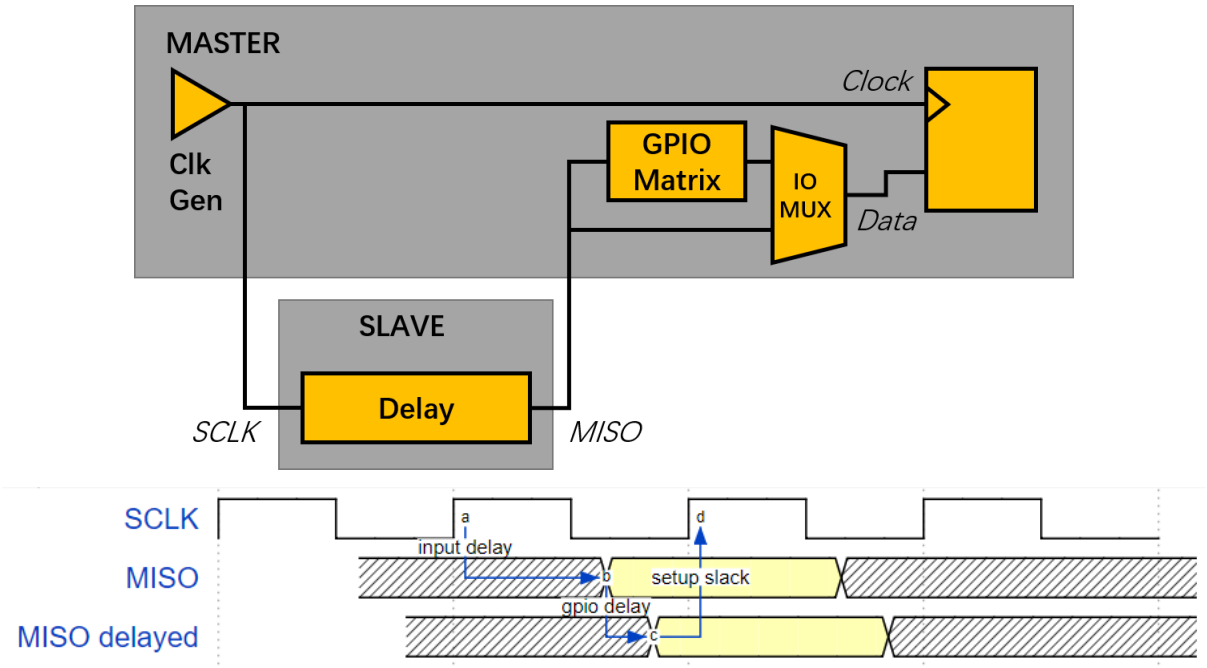
Frequency (MHz)	Transaction Interval (us)	Transaction Length (bytes)	Total Time (us)	Total Speed (KBps)
8	25	1	26	38.5
8	25	8	33	242.4
8	25	16	41	490.2
8	25	64	89	719.1
8	25	128	153	836.6

When a transaction length is short, the cost of transaction interval is high. If possible, try to squash several short transactions into one transaction to achieve a higher transfer speed.

Please note that the ISR is disabled during flash operation by default. To keep sending transactions during flash operations, enable *CONFIG_SPI_MASTER_ISR_IN_IRAM* and set `ESP_INTR_FLAG_IRAM` in the member *`spi_bus_config_t::intr_flags`*. In this case, all the transactions queued before starting flash operations will be handled by the ISR in parallel. Also note that the callback of each Device and their callee functions should be in IRAM, or your callback will crash due to cache miss. For more details, see *IRAM 安全中断处理程序*.

Timing Considerations

As shown in the figure below, there is a delay on the MISO line after the SCLK launch edge and before the signal is latched by the internal register. As a result, the MISO pin setup time is the limiting factor for the SPI clock speed. When the delay is too long, the setup slack is < 0, and the setup timing requirement is violated, which results in the failure to perform the reading correctly.



The maximum allowed frequency is dependent on:

- `input_delay_ns` - maximum data valid time on the MISO bus after a clock cycle on SCLK starts
- If the IO_MUX pin or the GPIO Matrix is used

When the GPIO matrix is used, the maximum allowed frequency is reduced to about 33~77% in comparison to the existing *input delay*. To retain a higher frequency, you have to use the IO_MUX pins or the *dummy bit workaround*. You can obtain the maximum reading frequency of the master by using the function `spi_get_freq_limit()`.

Dummy bit workaround: Dummy clocks, during which the Host does not read data, can be inserted before the read phase begins. The Device still sees the dummy clocks and sends out data, but the Host does not read until the read phase comes. This compensates for the lack of the MISO setup time required by the Host and allows the Host to do reading at a higher frequency.

In the ideal case, if the Device is so fast that the input delay is shorter than an APB clock cycle - 12.5 ns - the maximum frequency at which the Host can read (or read and write) in different conditions is as follows:

Frequency Limit (MHz)		Dummy Bits Used By Driver	Comments
GPIO matrix	IO_MUX pins		
26.6	80	No	
40	–	Yes	Half-duplex, no DMA allowed

If the Host only writes data, the *dummy bit workaround* and the frequency check can be disabled by setting the bit *SPI_DEVICE_NO_DUMMY* in the member `spi_device_interface_config_t::flags`. When disabled, the output frequency can be 80MHz, even if the GPIO matrix is used.

`spi_device_interface_config_t::flags`

The SPI master driver can work even if the `input_delay_ns` in the structure `spi_device_interface_config_t` is set to 0. However, setting an accurate value helps to:

- Calculate the frequency limit for full-duplex transactions
- Compensate the timing correctly with dummy bits for half-duplex transactions

You can approximate the maximum data valid time after the launch edge of SPI clocks by checking the statistics in the AC characteristics chapter of your Device's specification or measure the time on an oscilloscope or logic analyzer.

Please note that the actual PCB layout design and the excessive loads may increase the input delay. It means that non-optimal wiring and/or a load capacitor on the bus will most likely lead to the input delay values exceeding the values given in the Device specification or measured while the bus is floating.

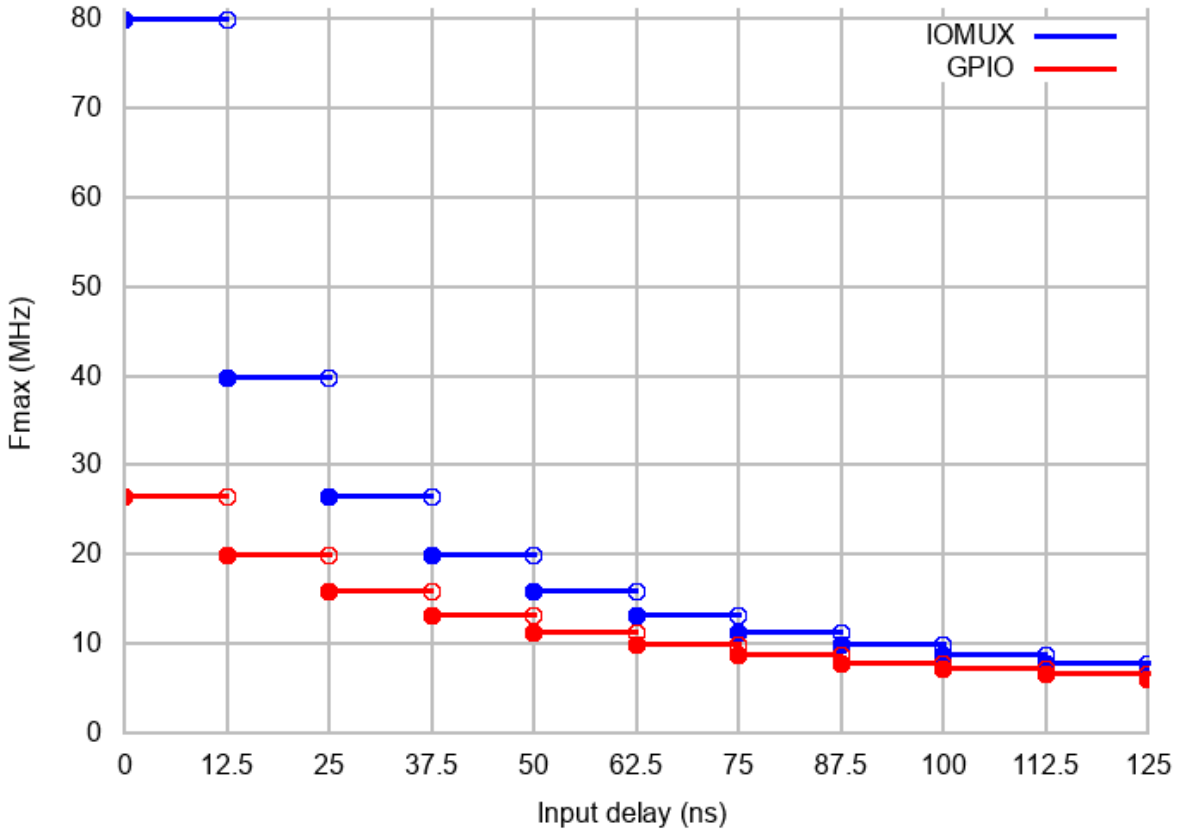
Some typical delay values are shown in the following table.

Device	Input delay (ns)
Ideal Device	0
ESP32 slave using IO_MUX*	50
ESP32 slave using GPIO_MUX*	75
ESP32's slave device is on a different physical chip.	

The MISO path delay (valid time) consists of a slave's *input delay* plus master's *GPIO matrix delay*. This delay determines the frequency limit above which full-duplex transfers will not work as well as the dummy bits used in the half-duplex transactions. The frequency limit is:

Freq limit [MHz] = 80 / (floor(MISO delay[ns]/12.5) + 1)

The figure below shows the relationship between frequency limit and input delay. Two extra APB clock cycle periods should be added to the MISO delay if the master uses the GPIO matrix.



Corresponding frequency limits for different Devices with different *input delay* times are shown in the table below.

Master	Input delay (ns)	MISO path delay (ns)	Freq. limit (MHz)
IO_MUX (0ns)	0	0	80
	50	50	16
	75	75	11.43
GPIO (25ns)	0	25	26.67
	50	75	11.43
	75	100	8.89

Known Issues

Application Example

The code example for displaying graphics on an ESP32-WROVER-KIT’ s 320x240 LCD screen can be found in the peripherals/spi_master directory of ESP-IDF examples.

API Reference - SPI Common

Header File

- soc/include/hal/spi_types.h

Enumerations

enum spi_host_device_t

Enum with the three SPI peripherals that are software-accessible in it.

Values:

SPI1_HOST = 0
SPI1.

SPI2_HOST = 1
SPI2.

SPI3_HOST = 2
SPI3.

Header File

- driver/include/driver/spi_common.h

Functions

esp_err_t **spi_bus_initialize**(*spi_host_device_t host_id*, **const** *spi_bus_config_t *bus_config*, int *dma_chan*)

Initialize a SPI bus.

Warning For now, only supports HSPI and VSPI.

Warning If a DMA channel is selected, any transmit and receive buffer used should be allocated in DMA-capable memory.

Warning The ISR of SPI is always executed on the core which calls this function. Never starve the ISR on this core or the SPI transactions will not be handled.

Return

- ESP_ERR_INVALID_ARG if configuration is invalid
- ESP_ERR_INVALID_STATE if host already is in use
- ESP_ERR_NO_MEM if out of memory
- ESP_OK on success

Parameters

- `host_id`: SPI peripheral that controls this bus
- `bus_config`: Pointer to a *spi_bus_config_t* struct specifying how the host should be initialized
- `dma_chan`: Either channel 1 or 2, or 0 in the case when no DMA is required. Selecting a DMA channel for a SPI bus allows transfers on the bus to have sizes only limited by the amount of internal memory. Selecting no DMA channel (by passing the value 0) limits the amount of bytes transfered to a maximum of 64. Set to 0 if only the SPI flash uses this bus.

esp_err_t **spi_bus_free**(*spi_host_device_t host_id*)

Free a SPI bus.

Warning In order for this to succeed, all devices have to be removed first.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_INVALID_STATE if bus hasn't been initialized before, or not all devices on the bus are freed
- ESP_OK on success

Parameters

- `host_id`: SPI peripheral to free

Structures

struct spi_bus_config_t

This is a configuration structure for a SPI bus.

You can use this structure to specify the GPIO pins of the bus. Normally, the driver will use the GPIO matrix to route the signals. An exception is made when all signals either can be routed through the IO_MUX or are -1. In that case, the IO_MUX is used, allowing for >40MHz speeds.

Note Be advised that the slave driver does not use the quadwp/quadhd lines and fields in *spi_bus_config_t* refering to these lines will be ignored and can thus safely be left uninitialized.

Public Members

int **mosi_io_num**
GPIO pin for Master Out Slave In (=spi_d) signal, or -1 if not used.

int **miso_io_num**
GPIO pin for Master In Slave Out (=spi_q) signal, or -1 if not used.

int **sclk_io_num**
GPIO pin for Spi CLocK signal, or -1 if not used.

int **quadwp_io_num**
GPIO pin for WP (Write Protect) signal which is used as D2 in 4-bit communication modes, or -1 if not used.

int **quadhd_io_num**
GPIO pin for HD (HolD) signal which is used as D3 in 4-bit communication modes, or -1 if not used.

int **max_transfer_sz**
Maximum transfer size, in bytes. Defaults to 4092 if 0 when DMA enabled, or to `SOC_SPI_MAXIMUM_BUFFER_SIZE` if DMA is disabled.

uint32_t **flags**
Abilities of bus to be checked by the driver. Or-ed value of `SPICOMMON_BUSFLAG_*` flags.

int **intr_flags**
Interrupt flag for the bus to set the priority, and IRAM attribute, see `esp_intr_alloc.h`. Note that the EDGE, INTRDISABLED attribute are ignored by the driver. Note that if ESP_INTR_FLAG_IRAM is set, ALL the callbacks of the driver, and their callee functions, should be put in the IRAM.

Macros

SPI_MAX_DMA_LEN

SPI_SWAP_DATA_TX(DATA, LEN)
Transform unsigned integer of length <= 32 bits to the format which can be sent by the SPI driver directly.

E.g. to send 9 bits of data, you can:

```
uint16_t data = SPI_SWAP_DATA_TX(0x145, 9);
```

Then points tx_buffer to `&data`.

Parameters

- `DATA`: Data to be sent, can be uint8_t, uint16_t or uint32_t.
- `LEN`: Length of data to be sent, since the SPI peripheral sends from the MSB, this helps to shift the data to the MSB.

SPI_SWAP_DATA_RX(DATA, LEN)
Transform received data of length <= 32 bits to the format of an unsigned integer.

E.g. to transform the data of 15 bits placed in a 4-byte array to integer:

```
uint16_t data = SPI_SWAP_DATA_RX(*(uint32_t*)t->rx_data, 15);
```

Parameters

- `DATA`: Data to be rearranged, can be uint8_t, uint16_t or uint32_t.
- `LEN`: Length of data received, since the SPI peripheral writes from the MSB, this helps to shift the data to the LSB.

SPICOMMON_BUSFLAG_SLAVE
Initialize I/O in slave mode.

SPICOMMON_BUSFLAG_MASTER
Initialize I/O in master mode.

SPICOMMON_BUSFLAG_IOMUX_PINS
Check using iomux pins. Or indicates the pins are configured through the IO mux rather than GPIO matrix.

SPICOMMON_BUSFLAG_SCLK
Check existing of SCLK pin. Or indicates CLK line initialized.

SPICOMMON_BUSFLAG_MISO
Check existing of MISO pin. Or indicates MISO line initialized.

SPICOMMON_BUSFLAG_MOSI
Check existing of MOSI pin. Or indicates CLK line initialized.

SPICOMMON_BUSFLAG_DUAL
Check MOSI and MISO pins can output. Or indicates bus able to work under DIO mode.

SPICOMMON_BUSFLAG_WPHD
Check existing of WP and HD pins. Or indicates WP & HD pins initialized.

SPICOMMON_BUSFLAG_QUAD
Check existing of MOSI/MISO/WP/HD pins as output. Or indicates bus able to work under QIO mode.

SPICOMMON_BUSFLAG_NATIVE_PINS

API Reference - SPI Master

Header File

- driver/include/driver/spi_master.h

Functions

esp_err_t **spi_bus_add_device**(*spi_host_device_t host_id*, **const** *spi_device_interface_config_t *dev_config*, *spi_device_handle_t *handle*)
Allocate a device on a SPI bus.

This initializes the internal structures for a device, plus allocates a CS pin on the indicated SPI master peripheral and routes it to the indicated GPIO. All SPI master devices have three CS pins and can thus control up to three devices.

Note While in general, speeds up to 80MHz on the dedicated SPI pins and 40MHz on GPIO-matrix-routed pins are supported, full-duplex transfers routed over the GPIO matrix only support speeds up to 26MHz.

Return
- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_NOT_FOUND if host doesn't have any free CS slots
- ESP_ERR_NO_MEM if out of memory
- ESP_OK on success

Parameters
- `host_id`: SPI peripheral to allocate device on
- `dev_config`: SPI interface protocol config for the device
- `handle`: Pointer to variable to hold the device handle

esp_err_t **spi_bus_remove_device**(*spi_device_handle_t handle*)
Remove a device from the SPI bus.

Return
- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_INVALID_STATE if device already is freed
- ESP_OK on success

Parameters
- `handle`: Device handle to free

esp_err_t **spi_device_queue_trans**(*spi_device_handle_t handle*, *spi_transaction_t *trans_desc*, TickType_t *ticks_to_wait*)
Queue a SPI transaction for interrupt transaction execution. Get the result by `spi_device_get_trans_result`.

Note Normally a device cannot start (queue) polling and interrupt transactions simultaneously.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_TIMEOUT if there was no room in the queue before ticks_to_wait expired
- ESP_ERR_NO_MEM if allocating DMA-capable temporary buffer failed
- ESP_ERR_INVALID_STATE if previous transactions are not finished
- ESP_OK on success

Parameters

- `handle`: Device handle obtained using spi_host_add_dev
- `trans_desc`: Description of transaction to execute
- `ticks_to_wait`: Ticks to wait until there's room in the queue; use portMAX_DELAY to never time out.

esp_err_t **spi_device_get_trans_result**(*spi_device_handle_t handle*, *spi_transaction_t* ***trans_desc*, TickType_t *ticks_to_wait*)

Get the result of a SPI transaction queued earlier by `spi_device_queue_trans`.

This routine will wait until a transaction to the given device succesfully completed. It will then return the description of the completed transaction so software can inspect the result and e.g. free the memory or re-use the buffers.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_TIMEOUT if there was no completed transaction before ticks_to_wait expired
- ESP_OK on success

Parameters

- `handle`: Device handle obtained using spi_host_add_dev
- `trans_desc`: Pointer to variable able to contain a pointer to the description of the transaction that is executed. The descriptor should not be modified until the descriptor is returned by spi_device_get_trans_result.
- `ticks_to_wait`: Ticks to wait until there's a returned item; use portMAX_DELAY to never time out.

esp_err_t **spi_device_transmit**(*spi_device_handle_t handle*, *spi_transaction_t *trans_desc*)

Send a SPI transaction, wait for it to complete, and return the result.

This function is the equivalent of calling spi_device_queue_trans() followed by spi_device_get_trans_result(). Do not use this when there is still a transaction separately queued (started) from spi_device_queue_trans() or polling_start/transmit that hasn't been finalized.

Note This function is not thread safe when multiple tasks access the same SPI device. Normally a device cannot start (queue) polling and interrupt transactions simutanuously.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_OK on success

Parameters

- `handle`: Device handle obtained using spi_host_add_dev
- `trans_desc`: Description of transaction to execute

esp_err_t **spi_device_polling_start**(*spi_device_handle_t handle*, *spi_transaction_t *trans_desc*, TickType_t *ticks_to_wait*)

Immediately start a polling transaction.

Note Normally a device cannot start (queue) polling and interrupt transactions simutanuously. Moreover, a device cannot start a new polling transaction if another polling transaction is not finished.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_TIMEOUT if the device cannot get control of the bus before `ticks_to_wait` expired
- ESP_ERR_NO_MEM if allocating DMA-capable temporary buffer failed
- ESP_ERR_INVALID_STATE if previous transactions are not finished
- ESP_OK on success

Parameters

- `handle`: Device handle obtained using spi_host_add_dev
- `trans_desc`: Description of transaction to execute
- `ticks_to_wait`: Ticks to wait until there's room in the queue; currently only portMAX_DELAY is supported.

esp_err_t **spi_device_polling_end**(*spi_device_handle_t handle*, TickType_t *ticks_to_wait*)

Poll until the polling transaction ends.

This routine will not return until the transaction to the given device has succesfully completed. The task is not blocked, but actively busy-spins for the transaction to be completed.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_TIMEOUT if the transaction cannot finish before ticks_to_wait expired
- ESP_OK on success

Parameters

- `handle`: Device handle obtained using spi_host_add_dev
- `ticks_to_wait`: Ticks to wait until there's a returned item; use portMAX_DELAY to never time out.

esp_err_t **spi_device_polling_transmit**(*spi_device_handle_t handle*, *spi_transaction_t *trans_desc*)

Send a polling transaction, wait for it to complete, and return the result.

This function is the equivalent of calling spi_device_polling_start() followed by spi_device_polling_end(). Do not use this when there is still a transaction that hasn't been finalized.

Note This function is not thread safe when multiple tasks access the same SPI device. Normally a device cannot start (queue) polling and interrupt transactions simutanuously.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_OK on success

Parameters

- `handle`: Device handle obtained using spi_host_add_dev
- `trans_desc`: Description of transaction to execute

esp_err_t **spi_device_acquire_bus**(*spi_device_handle_t device*, TickType_t *wait*)

Occupy the SPI bus for a device to do continuous transactions.

Transactions to all other devices will be put off until `spi_device_release_bus` is called.

Note The function will wait until all the existing transactions have been sent.

Return

- ESP_ERR_INVALID_ARG : `wait` is not set to portMAX_DELAY.
- ESP_OK : Success.

Parameters

- `device`: The device to occupy the bus.
- `wait`: Time to wait before the the bus is occupied by the device. Currently MUST set to portMAX_DELAY.

void **spi_device_release_bus**(*spi_device_handle_t dev*)

Release the SPI bus occupied by the device. All other devices can start sending transactions.

Parameters

- `dev`: The device to release the bus.

int **spi_cal_clock**(int *fapb*, int *hz*, int *duty_cycle*, uint32_t *reg_o*)

Calculate the working frequency that is most close to desired frequency, and also the register value.

Parameters

- `fapb`: The frequency of apb clock, should be `APB_CLK_FREQ`.
- `hz`: Desired working frequency
- `duty_cycle`: Duty cycle of the spi clock
- `reg_o`: Output of value to be set in clock register, or NULL if not needed.

Return Actual working frequency that most fit.

int **spi_get_actual_clock**(int *fapb*, int *hz*, int *duty_cycle*)

Calculate the working frequency that is most close to desired frequency.

Return Actual working frequency that most fit.

Parameters

- `fapb`: The frequency of apb clock, should be `APB_CLK_FREQ`.
- `hz`: Desired working frequency
- `duty_cycle`: Duty cycle of the spi clock

void **spi_get_timing**(bool *gpio_is_used*, int *input_delay_ns*, int *eff_clk*, int *dummy_o, int *cycles_remain_o)

Calculate the timing settings of specified frequency and settings.

Note If **dummy_o* is not zero, it means dummy bits should be applied in half duplex mode, and full duplex mode may not work.

Parameters

- `gpio_is_used`: True if using GPIO matrix, or False if iomux pins are used.
- `input_delay_ns`: Input delay from SCLK launch edge to MISO data valid.
- `eff_clk`: Effective clock frequency (in Hz) from spi_cal_clock.
- `dummy_o`: Address of dummy bits used output. Set to NULL if not needed.
- `cycles_remain_o`: Address of cycles remaining (after dummy bits are used) output.
  - -1 If too many cycles remaining, suggest to compensate half a clock.
  - 0 If no remaining cycles or dummy bits are not used.
  - positive value: cycles suggest to compensate.

int **spi_get_freq_limit**(bool *gpio_is_used*, int *input_delay_ns*)

Get the frequency limit of current configurations. SPI master working at this limit is OK, while above the limit, full duplex mode and DMA will not work, and dummy bits will be aplied in the half duplex mode.

Return Frequency limit of current configurations.

Parameters

- `gpio_is_used`: True if using GPIO matrix, or False if native pins are used.
- `input_delay_ns`: Input delay from SCLK launch edge to MISO data valid.

Structures

struct spi_device_interface_config_t

This is a configuration for a SPI slave device that is connected to one of the SPI buses.

Public Members

uint8_t **command_bits**

Default amount of bits in command phase (0-16), used when `SPI_TRANS_VARIABLE_CMD` is not used, otherwise ignored.

uint8_t **address_bits**

Default amount of bits in address phase (0-64), used when `SPI_TRANS_VARIABLE_ADDR` is not used, otherwise ignored.

uint8_t **dummy_bits**

Amount of dummy bits to insert between address and data phase.

uint8_t **mode**

SPI mode (0-3)

uint16_t **duty_cycle_pos**

Duty cycle of positive clock, in 1/256th increments (128 = 50%/50% duty). Setting this to 0 (=not setting it) is equivalent to setting this to 128.

uint16_t **cs_ena_pretrans**
: Amount of SPI bit-cycles the cs should be activated before the transmission (0-16). This only works on half-duplex transactions.

uint8_t **cs_ena_posttrans**
: Amount of SPI bit-cycles the cs should stay active after the transmission (0-16)

int **clock_speed_hz**
: Clock speed, divisors of 80MHz, in Hz. See `SPI_MASTER_FREQ_*`.

int **input_delay_ns**
: Maximum data valid time of slave. The time required between SCLK and MISO valid, including the possible clock delay from slave to master. The driver uses this value to give an extra delay before the MISO is ready on the line. Leave at 0 unless you know you need a delay. For better timing performance at high frequency (over 8MHz), it' s suggest to have the right value.

int **spics_io_num**
: CS GPIO pin for this device, or -1 if not used.

uint32_t **flags**
: Bitwise OR of SPI_DEVICE_* flags.

int **queue_size**
: Transaction queue size. This sets how many transactions can be 'in the air' (queued using spi_device_queue_trans but not yet finished using spi_device_get_trans_result) at the same time.

transaction_cb_t **pre_cb**
: Callback to be called before a transmission is started.

  This callback is called within interrupt context should be in IRAM for best performance, see "Transferring Speed" section in the SPI Master documentation for full details. If not, the callback may crash during flash operation when the driver is initialized with ESP_INTR_FLAG_IRAM.

transaction_cb_t **post_cb**
: Callback to be called after a transmission has completed.

  This callback is called within interrupt context should be in IRAM for best performance, see "Transferring Speed" section in the SPI Master documentation for full details. If not, the callback may crash during flash operation when the driver is initialized with ESP_INTR_FLAG_IRAM.

struct spi_transaction_t
: This structure describes one SPI transaction. The descriptor should not be modified until the transaction finishes.

Public Members

uint32_t **flags**
: Bitwise OR of SPI_TRANS_* flags.

uint16_t **cmd**
: Command data, of which the length is set in the `command_bits` of *spi_device_interface_config_t*.

  NOTE: this field, used to be "command" in ESP-IDF 2.1 and before, is re-written to be used in a new way in ESP-IDF 3.0.

  Example: write 0x0123 and command_bits=12 to send command 0x12, 0x3_ (in previous version, you may have to write 0x3_12).

uint64_t **addr**
: Address data, of which the length is set in the `address_bits` of *spi_device_interface_config_t*.

  NOTE: this field, used to be "address" in ESP-IDF 2.1 and before, is re-written to be used in a new way in ESP-IDF3.0.

  Example: write 0x123400 and address_bits=24 to send address of 0x12, 0x34, 0x00 (in previous version, you may have to write 0x12340000).

size_t **length**
Total data length, in bits.

size_t **rxlength**
Total data length received, should be not greater than `length` in full-duplex mode (0 defaults this to the value of `length`).

void ***user**
User-defined variable. Can be used to store eg transaction ID.

const void ***tx_buffer**
Pointer to transmit buffer, or NULL for no MOSI phase.

uint8_t **tx_data**[4]
If SPI_TRANS_USE_TXDATA is set, data set here is sent directly from this variable.

void ***rx_buffer**
Pointer to receive buffer, or NULL for no MISO phase. Written by 4 bytes-unit if DMA is used.

uint8_t **rx_data**[4]
If SPI_TRANS_USE_RXDATA is set, data is received directly to this variable.

struct spi_transaction_ext_t
This struct is for SPI transactions which may change their address and command length. Please do set the flags in base to `SPI_TRANS_VARIABLE_CMD_ADR` to use the bit length here.

Public Members

struct *spi_transaction_t* **base**
Transaction data, so that pointer to *spi_transaction_t* can be converted into *spi_transaction_ext_t*.

uint8_t **command_bits**
The command length in this transaction, in bits.

uint8_t **address_bits**
The address length in this transaction, in bits.

uint8_t **dummy_bits**
The dummy length in this transaction, in bits.

Macros

SPI_DEVICE_TXBIT_LSBFIRST
Transmit command/address/data LSB first instead of the default MSB first.

SPI master clock is divided by 80MHz apb clock. Below defines are example frequencies, and are accurate. Be free to specify a random frequency, it will be rounded to closest frequency (to macros below if above 8MHz). 8MHz

SPI_DEVICE_RXBIT_LSBFIRST
Receive data LSB first instead of the default MSB first.

SPI_DEVICE_BIT_LSBFIRST
Transmit and receive LSB first.

SPI_DEVICE_3WIRE
Use MOSI (=spid) for both sending and receiving data.

SPI_DEVICE_POSITIVE_CS
Make CS positive during a transaction instead of negative.

SPI_DEVICE_HALFDUPLEX
Transmit data before receiving it, instead of simultaneously.

SPI_DEVICE_CLK_AS_CS
Output clock on CS line if CS is active.

SPI_DEVICE_NO_DUMMY
There are timing issue when reading at high frequency (the frequency is related to whether iomux pins are used, valid time after slave sees the clock).

- In half-duplex mode, the driver automatically inserts dummy bits before reading phase to fix the timing issue. Set this flag to disable this feature.
- In full-duplex mode, however, the hardware cannot use dummy bits, so there is no way to prevent data being read from getting corrupted. Set this flag to confirm that you're going to work with output only, or read without dummy bits at your own risk.

SPI_DEVICE_DDRCLK

SPI_TRANS_MODE_DIO
Transmit/receive data in 2-bit mode.

SPI_TRANS_MODE_QIO
Transmit/receive data in 4-bit mode.

SPI_TRANS_USE_RXDATA
Receive into rx_data member of *spi_transaction_t* instead into memory at rx_buffer.

SPI_TRANS_USE_TXDATA
Transmit tx_data member of *spi_transaction_t* instead of data at tx_buffer. Do not set tx_buffer when using this.

SPI_TRANS_MODE_DIOQIO_ADDR
Also transmit address in mode selected by SPI_MODE_DIO/SPI_MODE_QIO.

SPI_TRANS_VARIABLE_CMD
Use the `command_bits` in *spi_transaction_ext_t* rather than default value in *spi_device_interface_config_t*.

SPI_TRANS_VARIABLE_ADDR
Use the `address_bits` in *spi_transaction_ext_t* rather than default value in *spi_device_interface_config_t*.

SPI_TRANS_VARIABLE_DUMMY
Use the `dummy_bits` in *spi_transaction_ext_t* rather than default value in *spi_device_interface_config_t*.

Type Definitions

typedef struct *spi_transaction_t* **spi_transaction_t**
typedef void (***transaction_cb_t**) (*spi_transaction_t* *trans)

typedef struct spi_device_t ***spi_device_handle_t**
Handle for a device on a SPI bus.

2.2.14 SPI Slave Driver

SPI Slave driver is a program that controls ESP32-S2's SPI peripherals while they function as slaves.

Overview of ESP32-S2's SPI peripherals

ESP32-S2 integrates two general purpose SPI controllers which can be used as slave nodes driven by an off-chip SPI master

- SPI2, sometimes referred to as HSPI
- SPI3, sometimes referred to as VSPI

SPI2 and SPI3 have independent signal buses with the same respective names.

Terminology

The terms used in relation to the SPI slave driver are given in the table below.

Term	Definition
Host	The SPI controller peripheral external to ESP32-S2 that initiates SPI transmissions over the bus, and acts as an SPI Master.
Device	SPI slave device, in this case the SPI2 and SPI3 controllers. Each Device shares the MOSI, MISO and SCLK signals but is only active on the bus when the Host asserts the Device' s individual CS line.
Bus	A signal bus, common to all Devices connected to one Host. In general, a bus includes the following lines: MISO, MOSI, SCLK, one or more CS lines, and, optionally, QUADWP and QUADHD. So Devices are connected to the same lines, with the exception that each Device has its own CS line. Several Devices can also share one CS line if connected in the daisy-chain manner.
• **MISO**	Master In, Slave Out, a.k.a. Q. Data transmission from a Device to Host.
• **MOSI**	Master Out, Slave in, a.k.a. D. Data transmission from a Host to Device.
• **SCLK**	Serial Clock. Oscillating signal generated by a Host that keeps the transmission of data bits in sync.
• **CS**	Chip Select. Allows a Host to select individual Device(s) connected to the bus in order to send or receive data.
• **QUADWP**	Write Protect signal. Only used for 4-bit (qio/qout) transactions.
• **QUADHD**	Hold signal. Only used for 4-bit (qio/qout) transactions.
• **Assertion**	The action of activating a line. The opposite action of returning the line back to inactive (back to idle) is called *de-assertion*.
Transaction	One instance of a Host asserting a CS line, transferring data to and from a Device, and de-asserting the CS line. Transactions are atomic, which means they can never be interrupted by another transaction.
Launch edge	Edge of the clock at which the source register *launches* the signal onto the line.
Latch edge	Edge of the clock at which the destination register *latches in* the signal.

Driver Features

The SPI slave driver allows using the SPI2 and/or SPI3 peripherals as full-duplex Devices. The driver can send/receive transactions up to 64 bytes in length, or utilize DMA to send/receive longer transactions. However, there are some *known issues* related to DMA.

SPI Transactions

A full-duplex SPI transaction begins when the Host asserts the CS line and starts sending out clock pulses on the SCLK line. Every clock pulse, a data bit is shifted from the Host to the Device on the MOSI line and back on the MISO line at the same time. At the end of the transaction, the Host de-asserts the CS line.

The attributes of a transaction are determined by the configuration structure for an SPI host acting as a slave device *spi_slave_interface_config_t*, and transaction configuration structure *spi_slave_transaction_t*.

As not every transaction requires both writing and reading data, you have a choice to configure the *spi_transaction_t* structure for TX only, RX only, or TX and RX transactions. If *spi_slave_transaction_t::rx_buffer* is set to NULL, the read phase will be skipped. If *spi_slave_transaction_t::tx_buffer* is set to NULL, the write phase will be skipped.

注解： A Host should not start a transaction before its Device is ready for receiving data. It is recommended to use another GPIO pin for a handshake signal to sync the Devices. For more details, see *Transaction Interval*.

Driver Usage

- Initialize an SPI peripheral as a Device by calling the function cpp:func:*spi_slave_initialize*. Make sure to set the correct I/O pins in the struct `bus_config`. Set the unused signals to `-1`.

If transactions will be longer than 32 bytes, allow a DMA channel by setting the parameter `dma_chan` to the host device. Otherwise, set `dma_chan` to `0`.

- Before initiating transactions, fill one or more *spi_slave_transaction_t* structs with the transaction parameters required. Either queue all transactions by calling the function *spi_slave_queue_trans()* and, at a later time, query the result by using the function *spi_slave_get_trans_result()*, or handle all requests individually by feeding them into *spi_slave_transmit()*. The latter two functions will be blocked until the Host has initiated and finished a transaction, causing the queued data to be sent and received.
- (Optional) To unload the SPI slave driver, call *spi_slave_free()*.

Transaction Data and Master/Slave Length Mismatches

Normally, the data that needs to be transferred to or from a Device is read or written to a chunk of memory indicated by the `rx_buffer` and `tx_buffer` members of the *spi_transaction_t* structure. The SPI driver can be configured to use DMA for transfers, in which case these buffers must be allocated in DMA-capable memory using `pvPortMallocCaps(size, MALLOC_CAP_DMA)`.

The amount of data that the driver can read or write to the buffers is limited by the member *spi_transaction_t::length*. However, this member does not define the actual length of an SPI transaction. A transaction's length is determined by a Host which drives the clock and CS lines. The actual length of the transmission can be read only after a transaction is finished from the member *spi_slave_transaction_t::trans_len*.

If the length of the transmission is greater than the buffer length, only the initial number of bits specified in the `length` member will be sent and received. In this case, `trans_len` is set to `length` instead of the actual transaction length. To meet the actual transaction length requirements, set `length` to a value greater than the maximum `trans_len` expected. If the transmission length is shorter than the buffer length, only the data equal to the length of the buffer will be transmitted.

Speed and Timing Considerations

Transaction Interval The ESP32-S2 SPI slave peripherals are designed as general purpose Devices controlled by a CPU. As opposed to dedicated slaves, CPU-based SPI Devices have a limited number of pre-defined registers. All

transactions must be handled by the CPU, which means that the transfers and responses are not real-time, and there might be noticeable latency.

As a solution, a Device's response rate can be doubled by using the functions `spi_slave_queue_trans()` and then `spi_slave_get_trans_result()` instead of using `spi_slave_transmit()`.

You can also configure a GPIO pin through which the Device will signal to the Host when it is ready for a new transaction. A code example of this can be found in peripherals/spi_slave.

SCLK Frequency Requirements The SPI slaves are designed to operate at up to 10 MHz. The data cannot be recognized or received correctly if the clock is too fast or does not have a 50% duty cycle.

On top of that, there are additional requirements for the data to meet the timing constraints:

- **Read (MOSI):** The Device can read data correctly only if the data is already set at the launch edge. Although it is usually the case for most masters.
- **Write (MISO):** The output delay of the MISO signal needs to be shorter than half of a clock cycle period so that the MISO line is stable before the next latch edge. Given that the clock is balanced, the output delay and frequency limitations in different cases are given below.

	Output delay of MISO (ns)	Freq. limit (MHz)
IO_MUX	43.75	<11.4
GPIO matrix	68.75	<7.2

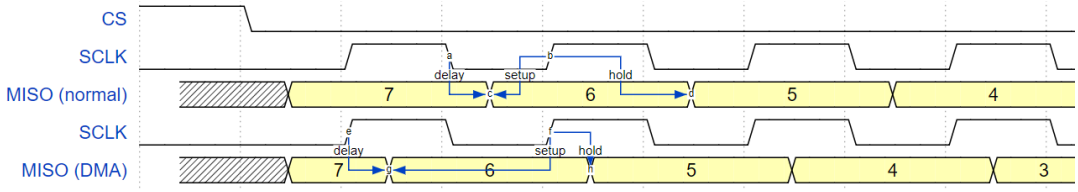
Note:

1. If the frequency is equal to the limitation, it can lead to random errors.
2. The clock uncertainty between Host and Device (12.5ns) is included.
3. The output delay is measured under ideal circumstances (no load). If the MISO pin is heavily loaded, the output delay will be longer, and the maximum allowed frequency will be lower.

Exception: The frequency is allowed to be higher if the master has more tolerance for the MISO setup time, e.g., latch data at the next edge than expected, or configurable latching time.

Restrictions and Known Issues

1. If DMA is enabled, the rx buffer should be word-aligned (starting from a 32-bit boundary and having a length of multiples of 4 bytes). Otherwise, DMA may write incorrectly or not in a boundary aligned manner. The driver reports an error if this condition is not satisfied.
   Also, a Host should write lengths that are multiples of 4 bytes. The data with inappropriate lengths will be discarded.
2. Furthermore, DMA requires SPI modes 1 and 3. For SPI modes 0 and 2, the MISO signal has to be launched half a clock cycle earlier to meet the timing. The new timing is as follows:



If DMA is enabled, a Device's launch edge is half of an SPI clock cycle ahead of the normal time, shifting to the Master's actual latch edge. In this case, if the GPIO matrix is bypassed, the hold time for data sampling is 68.75 ns and no longer a half of an SPI clock cycle. If the GPIO matrix is used, the hold time will increase to 93.75 ns. The Host should sample the data immediately at the latch edge or communicate in SPI modes 1 or 3. If your Host cannot meet these timing requirements, initialize your Device without DMA.

Application Example

The code example for Device/Host communication can be found in the peripherals/spi_slave directory of ESP-IDF examples.

API Reference

Header File

- driver/include/driver/spi_slave.h

Functions

esp_err_t **spi_slave_initialize**(*spi_host_device_t host*, **const** *spi_bus_config_t* **bus_config*, **const** *spi_slave_interface_config_t* **slave_config*, int *dma_chan*)

Initialize a SPI bus as a slave interface.

Warning For now, only supports HSPI and VSPI.

Warning If a DMA channel is selected, any transmit and receive buffer used should be allocated in DMA-capable memory.

Warning The ISR of SPI is always executed on the core which calls this function. Never starve the ISR on this core or the SPI transactions will not be handled.

Return

- ESP_ERR_INVALID_ARG if configuration is invalid
- ESP_ERR_INVALID_STATE if host already is in use
- ESP_ERR_NO_MEM if out of memory
- ESP_OK on success

Parameters

- `host`: SPI peripheral to use as a SPI slave interface
- `bus_config`: Pointer to a *spi_bus_config_t* struct specifying how the host should be initialized
- `slave_config`: Pointer to a *spi_slave_interface_config_t* struct specifying the details for the slave interface
- `dma_chan`: Either 1 or 2. A SPI bus used by this driver must have a DMA channel associated with it. The SPI hardware has two DMA channels to share. This parameter indicates which one to use.

esp_err_t **spi_slave_free**(*spi_host_device_t host*)

Free a SPI bus claimed as a SPI slave interface.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_ERR_INVALID_STATE if not all devices on the bus are freed
- ESP_OK on success

Parameters

- `host`: SPI peripheral to free

esp_err_t **spi_slave_queue_trans**(*spi_host_device_t host*, **const** *spi_slave_transaction_t* **trans_desc*, TickType_t *ticks_to_wait*)

Queue a SPI transaction for execution.

Queues a SPI transaction to be executed by this slave device. (The transaction queue size was specified when the slave device was initialised via spi_slave_initialize.) This function may block if the queue is full (depending on the ticks_to_wait parameter). No SPI operation is directly initiated by this function, the next queued transaction will happen when the master initiates a SPI transaction by pulling down CS and sending out clock signals.

This function hands over ownership of the buffers in `trans_desc` to the SPI slave driver; the application is not to access this memory until `spi_slave_queue_trans` is called to hand ownership back to the application.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_OK on success

Parameters

- `host`: SPI peripheral that is acting as a slave
- `trans_desc`: Description of transaction to execute. Not const because we may want to write status back into the transaction description.
- `ticks_to_wait`: Ticks to wait until there' s room in the queue; use portMAX_DELAY to never time out.

esp_err_t **spi_slave_get_trans_result**(*spi_host_device_t host*, *spi_slave_transaction_t* ***trans_desc*, TickType_t *ticks_to_wait*)

Get the result of a SPI transaction queued earlier.

This routine will wait until a transaction to the given device (queued earlier with spi_slave_queue_trans) has succesfully completed. It will then return the description of the completed transaction so software can inspect the result and e.g. free the memory or re-use the buffers.

It is mandatory to eventually use this function for any transaction queued by `spi_slave_queue_trans`.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_OK on success

Parameters

- `host`: SPI peripheral to that is acting as a slave
- `[out] trans_desc`: Pointer to variable able to contain a pointer to the description of the transaction that is executed
- `ticks_to_wait`: Ticks to wait until there's a returned item; use portMAX_DELAY to never time out.

esp_err_t **spi_slave_transmit**(*spi_host_device_t host*, *spi_slave_transaction_t* *trans_desc, TickType_t *ticks_to_wait*)

Do a SPI transaction.

Essentially does the same as spi_slave_queue_trans followed by spi_slave_get_trans_result. Do not use this when there is still a transaction queued that hasn't been finalized using spi_slave_get_trans_result.

Return

- ESP_ERR_INVALID_ARG if parameter is invalid
- ESP_OK on success

Parameters

- `host`: SPI peripheral to that is acting as a slave
- `trans_desc`: Pointer to variable able to contain a pointer to the description of the transaction that is executed. Not const because we may want to write status back into the transaction description.
- `ticks_to_wait`: Ticks to wait until there's a returned item; use portMAX_DELAY to never time out.

Structures

struct spi_slave_interface_config_t

This is a configuration for a SPI host acting as a slave device.

Public Members

int **spics_io_num**

CS GPIO pin for this device.

uint32_t **flags**

Bitwise OR of SPI_SLAVE_* flags.

int **queue_size**

Transaction queue size. This sets how many transactions can be 'in the air' (queued using spi_slave_queue_trans but not yet finished using spi_slave_get_trans_result) at the same time.

uint8_t **mode**

SPI mode (0-3)

slave_transaction_cb_t **post_setup_cb**

Callback called after the SPI registers are loaded with new data.

This callback is called within interrupt context should be in IRAM for best performance, see "Transferring Speed" section in the SPI Master documentation for full details. If not, the callback may crash during flash operation when the driver is initialized with ESP_INTR_FLAG_IRAM.

slave_transaction_cb_t **post_trans_cb**
Callback called after a transaction is done.

This callback is called within interrupt context should be in IRAM for best performance, see "Transferring Speed" section in the SPI Master documentation for full details. If not, the callback may crash during flash operation when the driver is initialized with ESP_INTR_FLAG_IRAM.

struct spi_slave_transaction_t
This structure describes one SPI transaction

Public Members

size_t **length**
Total data length, in bits.

size_t **trans_len**
Transaction data length, in bits.

const void ***tx_buffer**
Pointer to transmit buffer, or NULL for no MOSI phase.

void ***rx_buffer**
Pointer to receive buffer, or NULL for no MISO phase. When the DMA is anabled, must start at WORD boundary (`rx_buffer%4==0`), and has length of a multiple of 4 bytes.

void ***user**
User-defined variable. Can be used to store eg transaction ID.

Macros

SPI_SLAVE_TXBIT_LSBFIRST
Transmit command/address/data LSB first instead of the default MSB first.

SPI_SLAVE_RXBIT_LSBFIRST
Receive data LSB first instead of the default MSB first.

SPI_SLAVE_BIT_LSBFIRST
Transmit and receive LSB first.

Type Definitions

typedef struct *spi_slave_transaction_t* **spi_slave_transaction_t**

typedef void (***slave_transaction_cb_t**)(*spi_slave_transaction_t* *trans)

2.2.15 ESP32-S2 Temperature Sensor

Overview

The ESP32-S2 has a built-in temperature sensor. The temperature sensor module contains an 8-bit Sigma-Delta ADC and a temperature offset DAC.

The conversion relationship is the first columns of the table below. Among them, offset = 0 is the main measurement option, and other values are extended measurement options.

offset	measure range(Celsius)	measure error(Celsius)
-2	50 ~ 125	< 3
-1	20 ~ 100	< 2
0	-10 ~ 80	< 1
1	-30 ~ 50	< 2
2	-40 ~ 20	< 3

Application Example

Temperature sensor reading example: peripherals/temp_sensor_esp32s2.

API Reference - Normal Temp Sensor

Header File

- driver/esp32s2/include/driver/temp_sensor.h

Functions

esp_err_t **temp_sensor_set_config** (*temp_sensor_config_t tsens*)

Set parameter of temperature sensor.

Return
- ESP_OK Success

Parameters
- `tsens`:

esp_err_t **temp_sensor_get_config** (*temp_sensor_config_t *tsens*)

Get parameter of temperature sensor.

Return
- ESP_OK Success

Parameters
- `tsens`:

esp_err_t **temp_sensor_start** (void)

Start temperature sensor measure.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG

esp_err_t **temp_sensor_stop** (void)

Stop temperature sensor measure.

Return
- ESP_OK Success

esp_err_t **temp_sensor_read_raw** (uint32_t *tsens_out*)

Read temperature sensor raw data.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG `tsens_out` is NULL
- ESP_ERR_INVALID_STATE temperature sensor dont start

Parameters
- `tsens_out`: Pointer to raw data, Range: 0 ~ 255

esp_err_t **temp_sensor_read_celsius** (float *celsius*)

Read temperature sensor data that is converted to degrees Celsius.

Note Should not be called from interrupt.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG ARG is NULL.
- ESP_ERR_INVALID_STATE The ambient temperature is out of range.

Parameters
- `celsius`: The measure output value.

Structures

struct temp_sensor_config_t

Configuration for temperature sensor reading.

Public Members

temp_sensor_dac_offset_t **dac_offset**

The temperature measurement range is configured with a built-in temperature offset DAC.

uint8_t **clk_div**

Default: 6

Macros

TSENS_CONFIG_DEFAULT()

temperature sensor default setting.

Enumerations

enum temp_sensor_dac_offset_t

temperature sensor range option.

Values:

TSENS_DAC_L0 = 0

offset = -2, measure range: 50℃ ~ 125℃, error < 3℃.

TSENS_DAC_L1

offset = -1, measure range: 20℃ ~ 100℃, error < 2℃.

TSENS_DAC_L2

offset = 0, measure range:-10℃ ~ 80℃, error < 1℃.

TSENS_DAC_L3

offset = 1, measure range:-30℃ ~ 50℃, error < 2℃.

TSENS_DAC_L4

offset = 2, measure range:-40℃ ~ 20℃, error < 3℃.

TSENS_DAC_MAX

TSENS_DAC_DEFAULT = *TSENS_DAC_L2*

2.2.16 定时器

简介

ESP32 芯片提供两组硬件定时器，每组包含两个通用硬件定时器。所有定时器均为 64 位通用定时器，包括 16 位预分频器和 64 位自动重载向上/向下计数器。

功能概述

下文介绍了配置和操作定时器的常规步骤：

- 定时器初始化 - 启动定时器前应设置的参数，以及每个设置提供的具体功能。
- 定时器控制 - 如何读取定时器的值，如何暂停/启动定时器以及如何改变定时器的操作方式。
- 警报 - 如何设置和使用警报。
- 中断- 如何使能和使用中断。

定时器初始化 两个 ESP32 定时器组中，每组都有两个定时器，两组共有四个定时器供使用。ESP32 定时器组应使用*timer_group_t* 识别，而每组中的个体定时器则应使用*timer_idx_t* 识别。

首先调用*timer_init()* 函数，并将*timer_config_t* 结构体传递给此函数，用于定义定时器的工作方式，实现定时器初始化。特别注意以下定时器参数可设置为：

- **分频器**: 设置定时器中计数器计数的速度，divider 的设置将用作输入的 80 MHz APB_CLK 时钟的除数。
- **模式**: 设置计数器是递增还是递减。可通过从*timer_count_dir_t* 中选取一个值，后使用 counter_dir 来选择模式。
- **计数器使能**: 如果计数器已使能，则在调用*timer_init()* 后计数器将立即开始递增/递减。您可通过从*timer_start_t* 中选取一个值，后使用 counter_en 改变此行为。
- **报警使能**: 可使用 alarm_en 设置。
- **自动重载**: 设置计数器是否应该在定时器警报上使用 auto_reload 自动重载首个计数值，还是继续递增或递减。
- **中断类型**: 选择定时器警报上应触发的中断类型，请设置*timer_intr_mode_t* 中定义的值。

要获取定时器设置的当前值，请使用函数*timer_get_config()*。

定时器控制 定时器使能后便开始计数。要使能定时器，可首先设置 counter_en 为 true，然后调用函数*timer_init()*，或者直接调用函数*timer_start()*。您可通过调用函数*timer_set_counter_value()* 来指定定时器的首个计数值。要检查定时器的当前值，调用函数*timer_get_counter_value()* 或*timer_get_counter_time_sec()*。

可通过调用函数*timer_pause()* 随时暂停定时器。要再次启动它，调用函数*timer_start()*。

要重新配置定时器，可调用函数*timer_init()*，该函数详细介绍见定时器初始化。

除此之外，还可通过使用专有函数更改个别设置来重新配置定时器：

设置	专有函数	描述
分频器	*timer_set_divider()*	更改计数频率。为避免发生不可预测情况，更改分频器时应暂停定时器。如果定时器正在运行，则使用*timer_set_divider()* 将其暂停并更改设置，然后重启定时器。
模式	*timer_set_counter_mode()*	设置计数器应递增还是递减
自动重载	*timer_set_auto_reload()*	设置是否应在定时器警报上重载首个计数值

警报 要设置警报，先调用函数*timer_set_alarm_value()*，然后使用*timer_set_alarm()* 使能警报。当调用函数*timer_init()* 时，也可以在定时器初始化阶段使能警报。

警报已使能且定时器达到警报值后，根据配置，可能会出现以下两种行为：

- 如果先前已配置，此时将触发中断。有关如何配置中断，请参见中断。
- 如 auto_reload 已使能，定时器的计数器将重新加载，从先前配置好的值开始再次计数。应使用函数*timer_set_counter_value()* 预先设置该值。

注解:

- 如果已设置警报值且定时器已超过该值，则将立即触发警报。
- 一旦触发后，警报将自动关闭，需要重新使能以再次触发。

要检查某特定的警报值，调用函数*timer_get_alarm_value()*。

中断 可通过调用函数*timer_isr_register()* 为特定定时器组和定时器注册中断处理程序。

调用*timer_group_intr_enable()* 使能定时器组的中断程序，调用*timer_enable_intr()* 使能某定时器的中断程序。调用*timer_group_intr_disable()* 关闭定时器组的中断程序，调用*timer_disable_intr()* 关闭某定时器的中断程序。

在中断服务程序（ISR）中处理中断时，需要明确地清除中断状态位。为此，请设置定义在 soc/soc/esp32/include/soc/timer_group_struct.h 中的 `TIMERGN.int_clr_timers.tM` 结构。该结构中 `N` 是定时器组别编号 [0, 1]，`M` 是定时器编号 [0, 1]。例如，要清除定时器组别 0 中定时器 1 的中断状态位，请调用以下命令:

```
TIMERG0.int_clr_timers.t1 = 1
```

有关如何使用中断，请参阅应用示例。

应用示例

64 位硬件定时器示例：peripherals/timer_group。

API 参考

Header File

- driver/include/driver/timer.h

Functions

esp_err_t **timer_get_counter_value**(*timer_group_t group_num*, *timer_idx_t timer_num*, uint64_t **timer_val*)

Read the counter value of hardware timer.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `timer_val`: Pointer to accept timer counter value.

esp_err_t **timer_get_counter_time_sec**(*timer_group_t group_num*, *timer_idx_t timer_num*, double **time*)

Read the counter value of hardware timer, in unit of a given scale.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `time`: Pointer, type of double*, to accept timer counter value, in seconds.

esp_err_t **timer_set_counter_value**(*timer_group_t group_num*, *timer_idx_t timer_num*, uint64_t *load_val*)

Set counter value to hardware timer.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `load_val`: Counter value to write to the hardware timer.

esp_err_t **timer_start** (*timer_group_t group_num*, *timer_idx_t timer_num*)

Start the counter of hardware timer.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]

esp_err_t **timer_pause** (*timer_group_t group_num*, *timer_idx_t timer_num*)

Pause the counter of hardware timer.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]

esp_err_t **timer_set_counter_mode** (*timer_group_t group_num*, *timer_idx_t timer_num*, *timer_count_dir_t counter_dir*)

Set counting mode for hardware timer.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `counter_dir`: Counting direction of timer, count-up or count-down

esp_err_t **timer_set_auto_reload** (*timer_group_t group_num*, *timer_idx_t timer_num*, *timer_autoreload_t reload*)

Enable or disable counter reload function when alarm event occurs.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `reload`: Counter reload mode.

esp_err_t **timer_set_divider** (*timer_group_t group_num*, *timer_idx_t timer_num*, uint32_t *divider*)

Set hardware timer source clock divider. Timer groups clock are divider from APB clock.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `divider`: Timer clock divider value. The divider’ s range is from from 2 to 65536.

esp_err_t **timer_set_alarm_value** (*timer_group_t group_num*, *timer_idx_t timer_num*, uint64_t *alarm_value*)

Set timer alarm value.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group, 0 for TIMERG0 or 1 for TIMERG1

- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `alarm_value`: A 64-bit value to set the alarm value.

esp_err_t **timer_get_alarm_value**(*timer_group_t group_num*, *timer_idx_t timer_num*, uint64_t **alarm_value*)

Get timer alarm value.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `alarm_value`: Pointer of A 64-bit value to accept the alarm value.

esp_err_t **timer_set_alarm**(*timer_group_t group_num*, *timer_idx_t timer_num*, *timer_alarm_t alarm_en*)

Enable or disable generation of timer alarm events.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `alarm_en`: To enable or disable timer alarm function.

esp_err_t **timer_isr_callback_add**(*timer_group_t group_num*, *timer_idx_t timer_num*, *timer_isr_t isr_handler*, void **arg*, int *intr_alloc_flags*)

Add ISR handle callback for the corresponding timer.

The callback should return a bool value to determine whether need to do YIELD at the end of the ISR.

Note This ISR handler will be called from an ISR. This ISR handler do not need to handle interrupt status, and should be kept short. If you want to realize some specific applications or write the whole ISR, you can call timer_isr_register(⋯) to register ISR.

Parameters
- `group_num`: Timer group number
- `timer_num`: Timer index of timer group
- `isr_handler`: Interrupt handler function, it is a callback function.
- `arg`: Parameter for handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.

If the intr_alloc_flags value ESP_INTR_FLAG_IRAM is set, the handler function must be declared with IRAM_ATTR attribute and can only call functions in IRAM or ROM. It cannot call other timer APIs.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

esp_err_t **timer_isr_callback_remove**(*timer_group_t group_num*, *timer_idx_t timer_num*)

Remove ISR handle callback for the corresponding timer.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group number
- `timer_num`: Timer index of timer group

esp_err_t **timer_isr_register**(*timer_group_t group_num*, *timer_idx_t timer_num*, void (**fn*)) void *
, void **arg*, int *intr_alloc_flags*, *timer_isr_handle_t *handle*Register Timer interrupt handler, the handler is an ISR. The handler will be attached to the same CPU core that this function is running on.

If the intr_alloc_flags value ESP_INTR_FLAG_IRAM is set, the handler function must be declared with IRAM_ATTR attribute and can only call functions in IRAM or ROM. It cannot call other timer APIs. Use direct register access to configure timers from inside the ISR in this case.

Note If use this function to reigster ISR, you need to write the whole ISR. In the interrupt handler, you need to call timer_spinlock_take(..) before your handling, and call timer_spinlock_give(···) after your handling.

Parameters

- `group_num`: Timer group number
- `timer_num`: Timer index of timer group
- `fn`: Interrupt handler function.
- `arg`: Parameter for handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `handle`: Pointer to return handle. If non-NULL, a handle for the interrupt will be returned here.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

esp_err_t **timer_init**(*timer_group_t group_num*, *timer_idx_t timer_num*, **const** *timer_config_t *config*)

Initializes and configure the timer.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `config`: Pointer to timer initialization parameters.

esp_err_t **timer_deinit**(*timer_group_t group_num*, *timer_idx_t timer_num*)

Deinitializes the timer.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]

esp_err_t **timer_get_config**(*timer_group_t group_num*, *timer_idx_t timer_num*, *timer_config_t *config*)

Get timer configure value.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index, 0 for hw_timer[0] & 1 for hw_timer[1]
- `config`: Pointer of struct to accept timer parameters.

esp_err_t **timer_group_intr_enable**(*timer_group_t group_num*, *timer_intr_t intr_mask*)

Enable timer group interrupt, by enable mask.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `intr_mask`: Timer interrupt enable mask.
  - TIMER_INTR_T0: t0 interrupt
  - TIMER_INTR_T1: t1 interrupt

- TIMER_INTR_WDT: watchdog interrupt

esp_err_t **timer_group_intr_disable** (*timer_group_t group_num*, *timer_intr_t intr_mask*)
Disable timer group interrupt, by disable mask.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `intr_mask`: Timer interrupt disable mask.
  - TIMER_INTR_T0: t0 interrupt
  - TIMER_INTR_T1: t1 interrupt
  - TIMER_INTR_WDT: watchdog interrupt

esp_err_t **timer_enable_intr** (*timer_group_t group_num*, *timer_idx_t timer_num*)
Enable timer interrupt.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.

esp_err_t **timer_disable_intr** (*timer_group_t group_num*, *timer_idx_t timer_num*)
Disable timer interrupt.

Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.

void **timer_group_intr_clr_in_isr** (*timer_group_t group_num*, *timer_idx_t timer_num*)
Clear timer interrupt status, just used in ISR.

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.

void **timer_group_clr_intr_status_in_isr** (*timer_group_t group_num*, *timer_idx_t timer_num*)
Clear timer interrupt status, just used in ISR.

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.

void **timer_group_enable_alarm_in_isr** (*timer_group_t group_num*, *timer_idx_t timer_num*)
Enable alarm interrupt, just used in ISR.

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.

uint64_t **timer_group_get_counter_value_in_isr** (*timer_group_t group_num*, *timer_idx_t timer_num*)
Get the current counter value, just used in ISR.

Return
- Counter value

Parameters
- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.

void **timer_group_set_alarm_value_in_isr**(*timer_group_t group_num*, *timer_idx_t timer_num*, uint64_t *alarm_val*)

Set the alarm threshold for the timer, just used in ISR.

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.
- `alarm_val`: Alarm threshold.

void **timer_group_set_counter_enable_in_isr**(*timer_group_t group_num*, *timer_idx_t timer_num*, *timer_start_t counter_en*)

Enable/disable a counter, just used in ISR.

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index.
- `counter_en`: Enable/disable.

timer_intr_t **timer_group_intr_get_in_isr**(*timer_group_t group_num*)

Get the masked interrupt status, just used in ISR.

Return

- Interrupt status

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1

uint32_t **timer_group_get_intr_status_in_isr**(*timer_group_t group_num*)

Get interrupt status, just used in ISR.

Return

- Interrupt status

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1

void **timer_group_clr_intr_sta_in_isr**(*timer_group_t group_num*, *timer_intr_t intr_mask*)

Clear the masked interrupt status, just used in ISR.

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `intr_mask`: Masked interrupt.

bool **timer_group_get_auto_reload_in_isr**(*timer_group_t group_num*, *timer_idx_t timer_num*)

Get auto reload enable status, just used in ISR.

Return

- True Auto reload enabled
- False Auto reload disabled

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1
- `timer_num`: Timer index

esp_err_t **timer_spinlock_take**(*timer_group_t group_num*)

Take timer spinlock to enter critical protect.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1

esp_err_t **timer_spinlock_give**(*timer_group_t group_num*)

Give timer spinlock to exit critical protect.

Return

- ESP_OK Success

- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `group_num`: Timer group number, 0 for TIMERG0 or 1 for TIMERG1

Macros

`TIMER_BASE_CLK`

Frequency of the clock on the input of the timer groups

Type Definitions

`typedef` bool (***`timer_isr_t`**) (void *)

Interrupt handle callback function. User need to retrun a bool value in callback.

Return

- True Do task yield at the end of ISR
- False Not do task yield at the end of ISR

Note If you called FreeRTOS functions in callback, you need to return true or false based on the retrun value of argument `pxHigherPriorityTaskWoken`. For example, `xQueueSendFromISR` is called in callback, if the return value `pxHigherPriorityTaskWoken` of any FreeRTOS calls is pdTRUE, return true; otherwise return false.

`typedef` *intr_handle_t* **`timer_isr_handle_t`**

Interrupt handle, used in order to free the isr after use. Aliases to an int handle for now.

Header File

- soc/include/hal/timer_types.h

Structures

`struct timer_config_t`

Data structure with timer's configuration settings.

Public Members

timer_alarm_t **`alarm_en`**

Timer alarm enable

timer_start_t **`counter_en`**

Counter enable

timer_intr_mode_t **`intr_type`**

Interrupt mode

timer_count_dir_t **`counter_dir`**

Counter direction

timer_autoreload_t **`auto_reload`**

Timer auto-reload

uint32_t **`divider`**

Counter clock divider. The divider's range is from from 2 to 65536.

timer_src_clk_t **`clk_src`**

Use XTAL as source clock.

Enumerations

`enum timer_group_t`

Selects a Timer-Group out of 2 available groups.

Values:

TIMER_GROUP_0 = 0
Hw timer group 0

TIMER_GROUP_1 = 1
Hw timer group 1

TIMER_GROUP_MAX

enum timer_idx_t
Select a hardware timer from timer groups.

Values:

TIMER_0 = 0
Select timer0 of GROUPx

TIMER_1 = 1
Select timer1 of GROUPx

TIMER_MAX

enum timer_count_dir_t
Decides the direction of counter.

Values:

TIMER_COUNT_DOWN = 0
Descending Count from cnt.high|cnt.low

TIMER_COUNT_UP = 1
Ascending Count from Zero

TIMER_COUNT_MAX

enum timer_start_t
Decides whether timer is on or paused.

Values:

TIMER_PAUSE = 0
Pause timer counter

TIMER_START = 1
Start timer counter

enum timer_intr_t
Interrupt types of the timer.

Values:

TIMER_INTR_T0 = BIT(0)
interrupt of timer 0

TIMER_INTR_T1 = BIT(1)
interrupt of timer 1

TIMER_INTR_WDT = BIT(2)
interrupt of watchdog

TIMER_INTR_NONE = 0

enum timer_alarm_t
Decides whether to enable alarm mode.

Values:

TIMER_ALARM_DIS = 0
Disable timer alarm

TIMER_ALARM_EN = 1
Enable timer alarm

TIMER_ALARM_MAX

enum timer_intr_mode_t
Select interrupt type if running in alarm mode.

Values:

TIMER_INTR_LEVEL = 0
Interrupt mode: level mode

TIMER_INTR_MAX

enum timer_autoreload_t
Select if Alarm needs to be loaded by software or automatically reload by hardware.

Values:

TIMER_AUTORELOAD_DIS = 0
Disable auto-reload: hardware will not load counter value after an alarm event

TIMER_AUTORELOAD_EN = 1
Enable auto-reload: hardware will load counter value after an alarm event

TIMER_AUTORELOAD_MAX

enum timer_src_clk_t
Select timer source clock.

Values:

TIMER_SRC_CLK_APB = 0
Select APB as the source clock

TIMER_SRC_CLK_XTAL = 1
Select XTAL as the source clock

2.2.17 触摸传感器

概述

触摸传感器系统由保护覆盖层、触摸电极、绝缘基板和走线组成，保护覆盖层位于最上层，绝缘基板上设有电极及走线。用户触摸覆盖层将产生电容变化，根据电容变化判断此次触摸是否为有效触摸行为。

ESP32 可支持最多 10 个电容式触摸板/GPIO，触摸板可以以矩阵或滑条等方式组合使用，从而覆盖更大触感区域及更多触感点。触摸传感由有限状态机 (FSM) 硬件控制，由软件或专用硬件计时器发起。

如需了解触摸传感器设计、操作及其控制寄存器等相关信息，请参考《ESP32-S2 技术参考手册》(PDF)，您也可以在《ESP32 技术参考手册》中查看这一子系统是如何运行的。

请参考 触摸传感器应用方案简介，查看触摸传感器设计详情和固件开发指南。如果不想亲自在多种配置环境下测试触摸传感器，请查看 ESP32 触摸功能开发套件。

功能介绍

下面将 API 分解成几个函数组进行介绍，帮助您快速了解以下功能：

- 初始化触摸板驱动程序
- 配置触摸板 GPIO 管脚
- 触摸状态测量
- 调整测量参数（优化测量）
- 过滤触摸测量
- 触摸监测方式
- 设置中断信号监测触碰动作
- 中断触发

请前往*API 参考* 章节，查看某一函数的具体描述。*应用示例* 章节则介绍了此 API 的具体实现。

初始化触摸板驱动程序 使用触摸板之前，需要先调用`touch_pad_init()` 函数初始化触摸板驱动程序。此函数设置了*API 参考* 项下的 *Macros* 中列出的几项 `..._DEFAULT` 驱动程序参数，同时删除之前设置过的触摸板信息（如有），并禁用中断。

如果不再需要该驱动程序，可以调用`touch_pad_deinit()` 释放已初始化的驱动程序。

配置触摸板 GPIO 管脚 调用`touch_pad_config()` 使能某一 GPIO 的触感功能。

使用`touch_pad_set_fsm_mode()` 选择触摸板测量（由 FSM 操作）是由硬件计时器自动启动，还是由软件自动启动。如果选择软件模式，请使用`touch_pad_sw_start()` 启动 FSM。

触摸状态测量 借助以下两个函数从传感器读取原始数据或过滤后的数据：

- `touch_pad_read()`
- `touch_pad_read_filtered()`

这两个函数也可以用于检查触碰和释放触摸板时传感器读数变化范围，来评估触摸板设计，然后根据这些信息设定触摸阈值。

注解： 使用 `touch_pad_read_filtered()` 之前，需要先调用过滤触摸测量 中特定的滤波器函数初始化并配置该滤波器。

请参考应用示例 peripherals/touch_pad_read，查看如何使用这两个读值函数。

优化测量 触摸传感器设有数个可配置参数，以适应触摸板设计特点。例如，如果需要感知较细微的电容变化，则可以缩小触摸板充放电的参考电压范围。您可以使用`touch_pad_set_voltage()` 函数设置电压参考低值和参考高值。

优化测量除了可以识别细微的电容变化之外，还可以降低应用程序功耗，但可能会增加测量噪声干扰。如果得到的动态读数范围结果比较理想，则可以调用`touch_pad_set_meas_time()` 函数来减少测量时间，从而进一步降低功耗。

可用的测量参数及相应的‘set’函数总结如下：

- 触摸板充放电参数：
  - 电压门限：`touch_pad_set_voltage()`
  - 速率（斜率）`touch_pad_set_cnt_mode()`
- 测量时间：`touch_pad_set_meas_time()`

电压门限（参考低值/参考高值）、速率（斜率）与测量时间的关系如下图所示：

上图中的 *Output* 代表触摸传感器读值，即一个测量周期内测得的脉冲计数值。

所有函数均成对出现，用于设定某一特定参数，并获取当前参数值。例如：`touch_pad_set_voltage()` 和`touch_pad_get_voltage()`。

过滤触摸测量 如果测量中存在噪声，可以使用提供的 API 函数对测量进行过滤。使用滤波器之前，请先调用 `touch_pad_filter_start()` 启动该滤波器。

滤波器类型为 IIR（无限脉冲响应滤波器），您可以调用 `touch_pad_set_filter_period()` 配置此类滤波器的采样周期。

如需停止滤波器，请调用 `touch_pad_filter_stop()` 函数。如果不再使用该滤波器，请调用 `touch_pad_filter_delete()` 删除此滤波器。

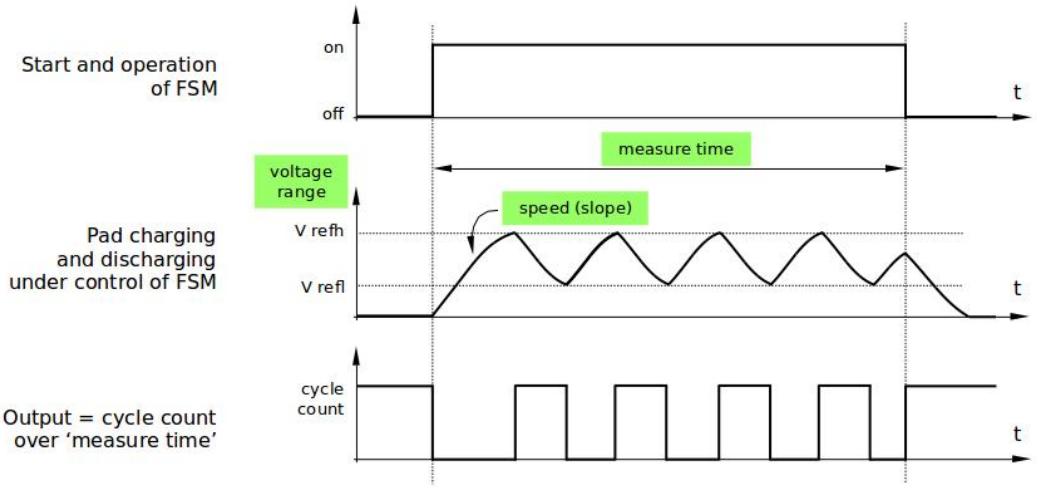


图 11: 触摸板 - 测量参数之间的关系

触摸监测 触摸监测基于用户配置的阈值和 FSM 执行的原始测量，并由 ESP32 硬件实现。你可以调用 *touch_pad_get_status()* 查看被触碰的触摸板，或调用 *touch_pad_clear_status()* 清除触摸状态信息。

您也可以将硬件触摸监测连接至中断，详细介绍见下一章节。

如果测量中存在噪声，且电容变化幅度较小，硬件触摸监测结果可能就不太理想。如需解决这一问题，不建议使用硬件监测或中断信号，建议您在自己的应用程序中采用测量过滤，并执行触摸监测。请参考 peripherals/touch_pad_interrupt，查看以上两种触摸监测的实现方式。

中断触发 在对触摸监测启用中断之前，请先设置一个触摸监测阈值。然后使用触摸状态测量 中所述的函数读取并显示触摸和释放触摸板时测得的结果。如果测量中存在噪声且相对电容变化较小，请使用滤波器。您也可以根据应用程序和环境条件，测试温度和电源电压变化对测量值的影响。

确定监测阈值后就可以在初始化时调用 *touch_pad_config()* 设置此阈值，或在运行时调用 *touch_pad_set_thresh()* 设置此阈值。

下一步就是设置如何触发中断。您可以设置在阈值以下或以上触发中断，具体触发模式由函数 `touch_pad_set_trigger_mode()` 设置。

最后您可以使用以下函数配置和管理中断调用：

- *touch_pad_isr_register()* / *touch_pad_isr_deregister()*
- *touch_pad_intr_enable()* / *touch_pad_intr_disable()*

中断配置完成后，您可以调用 *touch_pad_get_status()* 查看中断信号来自哪个触摸板，也可以调用 *touch_pad_clear_status()* 清除触摸板状态信息。

注解: 触摸监测中的中断信号基于原始/未经过滤的测量值（对比用户设置的阈值），并在硬件中实现。启用软件滤波 API 并不会影响这一过程，见过滤触摸测量。

从睡眠模式唤醒 如果使用触摸板中断将芯片从睡眠模式唤醒，您可以选择配置一些触摸板，例如 SET1 或 SET1 和 SET2，触摸这些触摸板将触发中断并唤醒芯片。请调用 `touch_pad_set_trigger_source()` 实现上述操作。

您可以使用以下函数管理‘SET’中触摸板所需的位模式配置：

- `touch_pad_set_group_mask()` / `touch_pad_get_group_mask()`
- `touch_pad_clear_group_mask()`

应用示例

- 触摸传感器读值示例：peripherals/touch_pad_read
- 触摸传感器中断示例：peripherals/touch_pad_interrupt

API 参考

Header File

- driver/esp32s2/include/driver/touch_sensor.h

Functions

esp_err_t **touch_pad_fsm_start**(void)

Set touch sensor FSM start.

Note Start FSM after the touch sensor FSM mode is set.
Note Call this function will reset benchmark of all touch channels.
Return
- ESP_OK on success

esp_err_t **touch_pad_fsm_stop**(void)

Stop touch sensor FSM.

Return
- ESP_OK on success

esp_err_t **touch_pad_sw_start**(void)

Trigger a touch sensor measurement, only support in SW mode of FSM.

Return
- ESP_OK on success

esp_err_t **touch_pad_set_meas_time**(uint16_t *sleep_cycle*, uint16_t *meas_times*)

Set touch sensor times of charge and discharge and sleep time. Excessive total time will slow down the touch response. Too small measurement time will not be sampled enough, resulting in inaccurate measurements.

Note The greater the duty cycle of the measurement time, the more system power is consumed.
Return
- ESP_OK on success

Parameters
- `sleep_cycle`: The touch sensor will sleep after each measurement. sleep_cycle decide the interval between each measurement. t_sleep = sleep_cycle / (RTC_SLOW_CLK frequency). The approximate frequency value of RTC_SLOW_CLK can be obtained using rtc_clk_slow_freq_get_hz function.
- `meas_times`: The times of charge and discharge in each measure process of touch channels. The timer frequency is 8Mhz. Range: 0 ~ 0xffff. Recommended typical value: Modify this value to make the measurement time around 1ms.

esp_err_t **touch_pad_get_meas_time**(uint16_t **sleep_cycle*, uint16_t **meas_times*)

Get touch sensor times of charge and discharge and sleep time.

Return
- ESP_OK on success

Parameters
- `sleep_cycle`: Pointer to accept sleep cycle number
- `meas_times`: Pointer to accept measurement times count.

esp_err_t **touch_pad_set_idle_channel_connect** (*touch_pad_conn_type_t type*)

Set connection type of touch channel in idle status. When a channel is in measurement mode, other initialized channels are in idle mode. The touch channel is generally adjacent to the trace, so the connection state of the idle channel affects the stability and sensitivity of the test channel. The CONN_HIGHZ(high resistance) setting increases the sensitivity of touch channels. The CONN_GND(grounding) setting increases the stability of touch channels.

Return
- ESP_OK on success

Parameters
- type: Select idle channel connect to high resistance state or ground.

esp_err_t **touch_pad_get_idle_channel_connect** (*touch_pad_conn_type_t *type*)

Set connection type of touch channel in idle status. When a channel is in measurement mode, other initialized channels are in idle mode. The touch channel is generally adjacent to the trace, so the connection state of the idle channel affects the stability and sensitivity of the test channel. The CONN_HIGHZ(high resistance) setting increases the sensitivity of touch channels. The CONN_GND(grounding) setting increases the stability of touch channels.

Return
- ESP_OK on success

Parameters
- type: Pointer to connection type.

esp_err_t **touch_pad_set_thresh** (*touch_pad_t touch_num*, uint32_t *threshold*)

Set the trigger threshold of touch sensor. The threshold determines the sensitivity of the touch sensor. The threshold is the original value of the trigger state minus the benchmark value.

Note If set "TOUCH_PAD_THRESHOLD_MAX", the touch is never be triggered.

Return
- ESP_OK on success

Parameters
- touch_num: touch pad index
- threshold: threshold of touch sensor. Should be less than the max change value of touch.

esp_err_t **touch_pad_get_thresh** (*touch_pad_t touch_num*, uint32_t *threshold*)

Get touch sensor trigger threshold.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if argument is wrong

Parameters
- touch_num: touch pad index
- threshold: pointer to accept threshold

esp_err_t **touch_pad_set_channel_mask** (uint16_t *enable_mask*)

Register touch channel into touch sensor scan group. The working mode of the touch sensor is cyclically scanned. This function will set the scan bits according to the given bitmask.

Note If set this mask, the FSM timer should be stop firsty.

Note The touch sensor that in scan map, should be deinit GPIO function firstly by touch_pad_io_init.

Return
- ESP_OK on success

Parameters
- enable_mask: bitmask of touch sensor scan group. e.g. TOUCH_PAD_NUM14 -> BIT(14)

esp_err_t **touch_pad_get_channel_mask** (uint16_t *enable_mask*)

Get the touch sensor scan group bit mask.

Return
- ESP_OK on success

Parameters
- enable_mask: Pointer to bitmask of touch sensor scan group. e.g. TOUCH_PAD_NUM14 -> BIT(14)

esp_err_t **touch_pad_clear_channel_mask** (uint16_t *enable_mask*)
Clear touch channel from touch sensor scan group. The working mode of the touch sensor is cyclically scanned. This function will clear the scan bits according to the given bitmask.

Note If clear all mask, the FSM timer should be stop firsty.
Return
- ESP_OK on success

Parameters
- `enable_mask`: bitmask of touch sensor scan group. e.g. TOUCH_PAD_NUM14 -> BIT(14)

esp_err_t **touch_pad_config** (*touch_pad_t touch_num*)
Configure parameter for each touch channel.

Note Touch num 0 is denoise channel, please use `touch_pad_denoise_enable` to set denoise function
Return
- ESP_OK Success
- ESP_ERR_INVALID_ARG if argument wrong
- ESP_FAIL if touch pad not initialized

Parameters
- `touch_num`: touch pad index

esp_err_t **touch_pad_reset** (void)
Reset the FSM of touch module.

Note Call this function after `touch_pad_fsm_stop`.
Return
- ESP_OK Success

touch_pad_t **touch_pad_get_current_meas_channel** (void)
Get the current measure channel.

Note Should be called when touch sensor measurement is in cyclic scan mode.
Return
- touch channel number

uint32_t **touch_pad_read_intr_status_mask** (void)
Get the touch sensor interrupt status mask.

Return
- touch interrupt bit

esp_err_t **touch_pad_intr_enable** (*touch_pad_intr_mask_t int_mask*)
Enable touch sensor interrupt by bitmask.

Return
- ESP_OK on success

Parameters
- `int_mask`: Pad mask to enable interrupts

esp_err_t **touch_pad_intr_disable** (*touch_pad_intr_mask_t int_mask*)
Disable touch sensor interrupt by bitmask.

Return
- ESP_OK on success

Parameters
- `int_mask`: Pad mask to disable interrupts

esp_err_t **touch_pad_intr_clear** (*touch_pad_intr_mask_t int_mask*)
Clear touch sensor interrupt by bitmask.

Return
- ESP_OK on success

Parameters
- `int_mask`: Pad mask to clear interrupts

esp_err_t **touch_pad_isr_register** (*intr_handler_t fn*, void **arg*, *touch_pad_intr_mask_t intr_mask*)

Register touch-pad ISR. The handler will be attached to the same CPU core that this function is running on.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Arguments error
- ESP_ERR_NO_MEM No memory

Parameters

- `fn`: Pointer to ISR handler
- `arg`: Parameter for ISR
- `intr_mask`: Enable touch sensor interrupt handler by bitmask.

esp_err_t **touch_pad_timeout_set** (bool *enable*, uint32_t *threshold*)

Enable/disable the timeout check and set timeout threshold for all touch sensor channels measurements. If enable: When the touch reading of a touch channel exceeds the measurement threshold, a timeout interrupt will be generated. If disable: the FSM does not check if the channel under measurement times out.

Note The threshold compared with touch readings.

Note In order to avoid abnormal short circuit of some touch channels. This function should be turned on. Ensure the normal operation of other touch channels.

Return

- ESP_OK Success

Parameters

- `enable`: true(default): Enable the timeout check; false: Disable the timeout check.
- `threshold`: For all channels, the maximum value that will not be exceeded during normal operation.

esp_err_t **touch_pad_timeout_resume** (void)

Call this interface after timeout to make the touch channel resume normal work. Point on the next channel to measure. If this API is not called, the touch FSM will stop the measurement after timeout interrupt.

Note Call this API after finishes the exception handling by user.

Return

- ESP_OK Success

esp_err_t **touch_pad_read_raw_data** (*touch_pad_t touch_num*, uint32_t **raw_data*)

get raw data of touch sensor.

Note After the initialization is complete, the "raw_data" is max value. You need to wait for a measurement cycle before you can read the correct touch value.

Return

- ESP_OK Success
- ESP_FAIL Touch channel 0 haven't this parameter.

Parameters

- `touch_num`: touch pad index
- `raw_data`: pointer to accept touch sensor value

esp_err_t **touch_pad_read_benchmark** (*touch_pad_t touch_num*, uint32_t **benchmark*)

get benchmark of touch sensor.

Note After initialization, the benchmark value is the maximum during the first measurement period.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Touch channel 0 haven't this parameter.

Parameters

- `touch_num`: touch pad index
- `benchmark`: pointer to accept touch sensor benchmark value

esp_err_t **touch_pad_filter_read_smooth** (*touch_pad_t touch_num*, uint32_t **smooth*)

Get smoothed data that obtained by filtering the raw data.

Parameters

- `touch_num`: touch pad index
- `smooth`: pointer to smoothed data

esp_err_t **touch_pad_reset_benchmark** (*touch_pad_t touch_num*)
Force reset benchmark to raw data of touch sensor.

Return
- ESP_OK Success

Parameters
- `touch_num`: touch pad index
  - TOUCH_PAD_MAX Reset basaline of all channels

esp_err_t **touch_pad_filter_set_config** (*touch_filter_config_t *filter_info*)
set parameter of touch sensor filter and detection algorithm. For more details on the detection algorithm, please refer to the application documentation.

Return
- ESP_OK Success

Parameters
- `filter_info`: select filter type and threshold of detection algorithm

esp_err_t **touch_pad_filter_get_config** (*touch_filter_config_t *filter_info*)
get parameter of touch sensor filter and detection algorithm. For more details on the detection algorithm, please refer to the application documentation.

Return
- ESP_OK Success

Parameters
- `filter_info`: select filter type and threshold of detection algorithm

esp_err_t **touch_pad_filter_enable** (void)
enable touch sensor filter for detection algorithm. For more details on the detection algorithm, please refer to the application documentation.

Return
- ESP_OK Success

esp_err_t **touch_pad_filter_disable** (void)
disable touch sensor filter for detection algorithm. For more details on the detection algorithm, please refer to the application documentation.

Return
- ESP_OK Success

esp_err_t **touch_pad_denoise_set_config** (*touch_pad_denoise_t *denoise*)
set parameter of denoise pad (TOUCH_PAD_NUM0). T0 is an internal channel that does not have a corresponding external GPIO. T0 will work simultaneously with the measured channel Tn. Finally, the actual measured value of Tn is the value after subtracting lower bits of T0. The noise reduction function filters out interference introduced simultaneously on all channels, such as noise introduced by power supplies and external EMI.

Return
- ESP_OK Success

Parameters
- `denoise`: parameter of denoise

esp_err_t **touch_pad_denoise_get_config** (*touch_pad_denoise_t *denoise*)
get parameter of denoise pad (TOUCH_PAD_NUM0).

Return
- ESP_OK Success

Parameters
- `denoise`: Pointer to parameter of denoise

esp_err_t **touch_pad_denoise_enable** (void)
enable denoise function. T0 is an internal channel that does not have a corresponding external GPIO. T0 will work simultaneously with the measured channel Tn. Finally, the actual measured value of Tn is the value after

subtracting lower bits of T0. The noise reduction function filters out interference introduced simultaneously on all channels, such as noise introduced by power supplies and external EMI.

Return
- ESP_OK Success

esp_err_t **touch_pad_denoise_disable**(void)
disable denoise function.

Return
- ESP_OK Success

esp_err_t **touch_pad_denoise_read_data**(uint32_t **data*)
Get denoise measure value (TOUCH_PAD_NUM0).

Return
- ESP_OK Success

Parameters
- `data`: Pointer to receive denoise value

esp_err_t **touch_pad_waterproof_set_config**(*touch_pad_waterproof_t *waterproof*)
set parameter of waterproof function.

The waterproof function includes a shielded channel (TOUCH_PAD_NUM14) and a guard channel. Guard pad is used to detect the large area of water covering the touch panel. Shield pad is used to shield the influence of water droplets covering the touch panel. It is generally designed as a grid and is placed around the touch buttons.

Return
- ESP_OK Success

Parameters
- `waterproof`: parameter of waterproof

esp_err_t **touch_pad_waterproof_get_config**(*touch_pad_waterproof_t *waterproof*)
get parameter of waterproof function.

Return
- ESP_OK Success

Parameters
- `waterproof`: parameter of waterproof

esp_err_t **touch_pad_waterproof_enable**(void)
Enable parameter of waterproof function. Should be called after function `touch_pad_waterproof_set_config`.

Return
- ESP_OK Success

esp_err_t **touch_pad_waterproof_disable**(void)
Disable parameter of waterproof function.

Return
- ESP_OK Success

esp_err_t **touch_pad_proximity_enable**(*touch_pad_t touch_num*, bool *enabled*)
Enable/disable proximity function of touch channels. The proximity sensor measurement is the accumulation of touch channel measurements.

Note Supports up to three touch channels configured as proximity sensors.

Return
- ESP_OK: Configured correctly.
- ESP_ERR_INVALID_ARG: Touch channel number error.
- ESP_ERR_NOT_SUPPORTED: Don't support configured.

Parameters
- `touch_num`: touch pad index
- `enabled`: true: enable the proximity function; false: disable the proximity function

esp_err_t **touch_pad_proximity_set_count** (*touch_pad_t touch_num*, uint32_t *count*)

Set measure count of proximity channel. The proximity sensor measurement is the accumulation of touch channel measurements.

Note All proximity channels use the same `count` value. So please pass the parameter `TOUCH_PAD_MAX`.

Return

- ESP_OK: Configured correctly.
- ESP_ERR_INVALID_ARG: Touch channel number error.

Parameters

- `touch_num`: Touch pad index. In this version, pass the parameter `TOUCH_PAD_MAX`.
- `count`: The cumulative times of measurements for proximity pad. Range: 0 ~ 255.

esp_err_t **touch_pad_proximity_get_count** (*touch_pad_t touch_num*, uint32_t **count*)

Get measure count of proximity channel. The proximity sensor measurement is the accumulation of touch channel measurements.

Note All proximity channels use the same `count` value. So please pass the parameter `TOUCH_PAD_MAX`.

Return

- ESP_OK: Configured correctly.
- ESP_ERR_INVALID_ARG: Touch channel number error.

Parameters

- `touch_num`: Touch pad index. In this version, pass the parameter `TOUCH_PAD_MAX`.
- `count`: The cumulative times of measurements for proximity pad. Range: 0 ~ 255.

esp_err_t **touch_pad_proximity_get_data** (*touch_pad_t touch_num*, uint32_t **measure_out*)

Get the accumulated measurement of the proximity sensor. The proximity sensor measurement is the accumulation of touch channel measurements.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Touch num is not proximity

Parameters

- `touch_num`: touch pad index
- `measure_out`: If the accumulation process does not end, the `measure_out` is the process value.

esp_err_t **touch_pad_sleep_channel_get_info** (*touch_pad_sleep_channel_t *slp_config*)

Get parameter of touch sensor sleep channel. The touch sensor can works in sleep mode to wake up sleep.

Note After the sleep channel is configured, Please use special functions for sleep channel. e.g. The user should uses `touch_pad_sleep_channel_read_data` instead of `touch_pad_read_raw_data` to obtain the sleep channel reading.

Return

- ESP_OK Success

Parameters

- `slp_config`: touch sleep pad config.

esp_err_t **touch_pad_sleep_channel_enable** (*touch_pad_t pad_num*, bool *enable*)

Enable/Disable sleep channel function for touch sensor. The touch sensor can works in sleep mode to wake up sleep.

Note ESP32S2 only support one sleep channel.

Note After the sleep channel is configured, Please use special functions for sleep channel. e.g. The user should uses `touch_pad_sleep_channel_read_data` instead of `touch_pad_read_raw_data` to obtain the sleep channel reading.

Return

- ESP_OK Success

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `enable`: true: enable sleep pad for touch sensor; false: disable sleep pad for touch sensor;

esp_err_t **touch_pad_sleep_channel_enable_proximity**(*touch_pad_t pad_num*, bool *enable*)

Enable/Disable proximity function for sleep channel. The touch sensor can works in sleep mode to wake up sleep.

Note ESP32S2 only support one sleep channel.

Return

- ESP_OK Success

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `enable`: true: enable proximity for sleep channel; false: disable proximity for sleep channel;

esp_err_t **touch_pad_sleep_set_threshold**(*touch_pad_t pad_num*, uint32_t *touch_thres*)

Set the trigger threshold of touch sensor in deep sleep. The threshold determines the sensitivity of the touch sensor.

Note In general, the touch threshold during sleep can use the threshold parameter parameters before sleep.

Return

- ESP_OK Success

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `touch_thres`: touch sleep pad threshold

esp_err_t **touch_pad_sleep_get_threshold**(*touch_pad_t pad_num*, uint32_t **touch_thres*)

Get the trigger threshold of touch sensor in deep sleep. The threshold determines the sensitivity of the touch sensor.

Note In general, the touch threshold during sleep can use the threshold parameter parameters before sleep.

Return

- ESP_OK Success

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `touch_thres`: touch sleep pad threshold

esp_err_t **touch_pad_sleep_channel_read_benchmark**(*touch_pad_t pad_num*, uint32_t **benchmark*)

Read benchmark of touch sensor sleep channel.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG parameter is NULL

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `benchmark`: pointer to accept touch sensor benchmark value

esp_err_t **touch_pad_sleep_channel_read_smooth**(*touch_pad_t pad_num*, uint32_t **smooth_data*)

Read smoothed data of touch sensor sleep channel. Smoothed data is filtered from the raw data.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG parameter is NULL

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `smooth_data`: pointer to accept touch sensor smoothed data

esp_err_t **touch_pad_sleep_channel_read_data**(*touch_pad_t pad_num*, uint32_t **raw_data*)

Read raw data of touch sensor sleep channel.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG parameter is NULL

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `raw_data`: pointer to accept touch sensor raw data

esp_err_t **touch_pad_sleep_channel_reset_benchmark**(void)

Reset benchmark of touch sensor sleep channel.

Return

- ESP_OK Success

esp_err_t **touch_pad_sleep_channel_read_proximity_cnt**(*touch_pad_t pad_num*, uint32_t **proximity_cnt*)

Read proximity count of touch sensor sleep channel.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG parameter is NULL

Parameters

- `pad_num`: Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode.
- `proximity_cnt`: pointer to accept touch sensor proximity count value

Header File

- driver/include/driver/touch_sensor_common.h

Functions

esp_err_t **touch_pad_init**(void)

Initialize touch module.

Note If default parameter don't match the usage scenario, it can be changed after this function.

Return

- ESP_OK Success
- ESP_ERR_NO_MEM Touch pad init error
- ESP_ERR_NOT_SUPPORTED Touch pad is providing current to external XTAL

esp_err_t **touch_pad_deinit**(void)

Un-install touch pad driver.

Note After this function is called, other touch functions are prohibited from being called.

Return

- ESP_OK Success
- ESP_FAIL Touch pad driver not initialized

esp_err_t **touch_pad_io_init**(*touch_pad_t touch_num*)

Initialize touch pad GPIO.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if argument is wrong

Parameters

- `touch_num`: touch pad index

esp_err_t **touch_pad_set_voltage**(*touch_high_volt_t refh*, *touch_low_volt_t refl*, *touch_volt_atten_t atten*)

Set touch sensor high voltage threshold of chanrge. The touch sensor measures the channel capacitance value by charging and discharging the channel. So the high threshold should be less than the supply voltage.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if argument is wrong

Parameters

- `refh`: the value of DREFH
- `refl`: the value of DREFL
- `atten`: the attenuation on DREFH

esp_err_t **touch_pad_get_voltage**(*touch_high_volt_t *refh*, *touch_low_volt_t *refl*, *touch_volt_atten_t *atten*)

Get touch sensor reference voltage,.

Return

- ESP_OK on success

Parameters

- `refh`: pointer to accept DREFH value
- `refl`: pointer to accept DREFL value
- `atten`: pointer to accept the attenuation on DREFH

esp_err_t **touch_pad_set_cnt_mode**(*touch_pad_t touch_num*, *touch_cnt_slope_t slope*, *touch_tie_opt_t opt*)

Set touch sensor charge/discharge speed for each pad. If the slope is 0, the counter would always be zero. If the slope is 1, the charging and discharging would be slow, accordingly. If the slope is set 7, which is the maximum value, the charging and discharging would be fast.

Note The higher the charge and discharge current, the greater the immunity of the touch channel, but it will increase the system power consumption.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if argument is wrong

Parameters

- `touch_num`: touch pad index
- `slope`: touch pad charge/discharge speed
- `opt`: the initial voltage

esp_err_t **touch_pad_get_cnt_mode**(*touch_pad_t touch_num*, *touch_cnt_slope_t *slope*, *touch_tie_opt_t *opt*)

Get touch sensor charge/discharge speed for each pad.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if argument is wrong

Parameters

- `touch_num`: touch pad index
- `slope`: pointer to accept touch pad charge/discharge slope
- `opt`: pointer to accept the initial voltage

esp_err_t **touch_pad_isr_deregister**(void (*fn)) void *
, void *arg)Deregister the handler previously registered using touch_pad_isr_handler_register.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if a handler matching both fn and arg isn't registered

Parameters

- `fn`: handler function to call (as passed to touch_pad_isr_handler_register)
- `arg`: argument of the handler (as passed to touch_pad_isr_handler_register)

esp_err_t **touch_pad_get_wakeup_status**(*touch_pad_t *pad_num*)

Get the touch pad which caused wakeup from deep sleep.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG parameter is NULL

Parameters

- `pad_num`: pointer to touch pad which caused wakeup

esp_err_t **touch_pad_set_fsm_mode**(*touch_fsm_mode_t mode*)
Set touch sensor FSM mode, the test action can be triggered by the timer, as well as by the software.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if argument is wrong

Parameters
- `mode`: FSM mode

esp_err_t **touch_pad_get_fsm_mode**(*touch_fsm_mode_t *mode*)
Get touch sensor FSM mode.

Return
- ESP_OK on success

Parameters
- `mode`: pointer to accept FSM mode

esp_err_t **touch_pad_clear_status**(void)
To clear the touch sensor channel active status.

Note The FSM automatically updates the touch sensor status. It is generally not necessary to call this API to clear the status.

Return
- ESP_OK on success

uint32_t **touch_pad_get_status**(void)
Get the touch sensor channel active status mask. The bit position represents the channel number. The 0/1 status of the bit represents the trigger status.

Return
- The touch sensor status. e.g. Touch1 trigger status is `status_mask & (BIT1)`.

bool **touch_pad_meas_is_done**(void)
Check touch sensor measurement status.

Return
- True measurement is under way
- False measurement done

GPIO 宏查找表 您可以使用宏定义某一触摸板通道的 GPIO，或定义某一 GPIO 的通道。例如：

1. `TOUCH_PAD_NUM5_GPIO_NUM` 定义了通道 5 的 GPIO（即 GPIO 12）；
2. `TOUCH_PAD_GPIO4_CHANNEL` 定义了 GPIO 4 的通道（即通道 0）。

Header File

- soc/soc/esp32s2/include/soc/touch_sensor_channel.h

Macros

TOUCH_PAD_GPIO1_CHANNEL
TOUCH_PAD_NUM1_GPIO_NUM

TOUCH_PAD_GPIO2_CHANNEL

TOUCH_PAD_NUM2_GPIO_NUM

TOUCH_PAD_GPIO3_CHANNEL

TOUCH_PAD_NUM3_GPIO_NUM

TOUCH_PAD_GPIO4_CHANNEL

TOUCH_PAD_NUM4_GPIO_NUM

TOUCH_PAD_GPIO5_CHANNEL

TOUCH_PAD_NUM5_GPIO_NUM

TOUCH_PAD_GPIO6_CHANNEL

TOUCH_PAD_NUM6_GPIO_NUM

TOUCH_PAD_GPIO7_CHANNEL

TOUCH_PAD_NUM7_GPIO_NUM

TOUCH_PAD_GPIO8_CHANNEL

TOUCH_PAD_NUM8_GPIO_NUM

TOUCH_PAD_GPIO9_CHANNEL

TOUCH_PAD_NUM9_GPIO_NUM

TOUCH_PAD_GPIO10_CHANNEL

TOUCH_PAD_NUM10_GPIO_NUM

TOUCH_PAD_GPIO11_CHANNEL

TOUCH_PAD_NUM11_GPIO_NUM

TOUCH_PAD_GPIO12_CHANNEL

TOUCH_PAD_NUM12_GPIO_NUM

TOUCH_PAD_GPIO13_CHANNEL

TOUCH_PAD_NUM13_GPIO_NUM

TOUCH_PAD_GPIO14_CHANNEL

TOUCH_PAD_NUM14_GPIO_NUM

Header File

- soc/include/hal/touch_sensor_types.h

Structures

struct touch_pad_denoise

Touch sensor denoise configuration

Public Members

touch_pad_denoise_grade_t **grade**

Select denoise range of denoise channel. Determined by measuring the noise amplitude of the denoise channel.

touch_pad_denoise_cap_t **cap_level**

Select internal reference capacitance of denoise channel. Ensure that the denoise readings are closest to the readings of the channel being measured. Use `touch_pad_denoise_read_data` to get the reading of denoise channel. The equivalent capacitance of the shielded channel can be calculated from the reading of denoise channel.

struct touch_pad_waterproof

Touch sensor waterproof configuration

Public Members

touch_pad_t **guard_ring_pad**

Waterproof. Select touch channel use for guard pad. Guard pad is used to detect the large area of water covering the touch panel.

touch_pad_shield_driver_t **shield_driver**
Waterproof. Shield channel drive capability configuration. Shield pad is used to shield the influence of water droplets covering the touch panel. When the waterproof function is enabled, Touch14 is set as shield channel by default. The larger the parasitic capacitance on the shielding channel, the higher the drive capability needs to be set. The equivalent capacitance of the shield channel can be estimated through the reading value of the denoise channel(Touch0).

struct touch_filter_config
Touch sensor filter configuration

Public Members

touch_filter_mode_t **mode**
Set filter mode. The input of the filter is the raw value of touch reading, and the output of the filter is involved in the judgment of the touch state.

uint32_t **debounce_cnt**
Set debounce count, such as `n`. If the measured values continue to exceed the threshold for `n+1` times, the touch sensor state changes. Range: 0 ~ 7

uint32_t **noise_thr**
Noise threshold coefficient. Higher = More noise resistance. The actual noise should be less than (noise coefficient * touch threshold). Range: 0 ~ 3. The coefficient is 0: 4/8; 1: 3/8; 2: 2/8; 3: 1;

uint32_t **jitter_step**
Set jitter filter step size. Range: 0 ~ 15

touch_smooth_mode_t **smh_lvl**
Level of filter applied on the original data against large noise interference.

struct touch_pad_sleep_channel_t
Touch sensor channel sleep configuration

Public Members

touch_pad_t **touch_num**
Set touch channel number for sleep pad. Only one touch sensor channel is supported in deep sleep mode. If clear the sleep channel, point this pad to `TOUCH_PAD_NUM0`

bool **en_proximity**
enable proximity function for sleep pad

Macros

TOUCH_PAD_SLOPE_DEFAULT
TOUCH_PAD_TIE_OPT_DEFAULT

TOUCH_PAD_BIT_MASK_MAX

TOUCH_PAD_HIGH_VOLTAGE_THRESHOLD

TOUCH_PAD_LOW_VOLTAGE_THRESHOLD

TOUCH_PAD_ATTEN_VOLTAGE_THRESHOLD

TOUCH_PAD_IDLE_CH_CONNECT_DEFAULT

TOUCH_PAD_THRESHOLD_MAX
If set touch threshold max value, The touch sensor can't be in touched status

TOUCH_PAD_SLEEP_CYCLE_DEFAULT
Excessive total time will slow down the touch response. Too small measurement time will not be sampled enough, resulting in inaccurate measurements.

Note The greater the duty cycle of the measurement time, the more system power is consumed.The number of sleep cycle in each measure process of touch channels. The timer frequency is RTC_SLOW_CLK (can be 150k or 32k depending on the options). Range: 0 ~ 0xffff

TOUCH_PAD_MEASURE_CYCLE_DEFAULT
The times of charge and discharge in each measure process of touch channels. The timer frequency is 8Mhz. Recommended typical value: Modify this value to make the measurement time around 1ms. Range: 0 ~ 0xffff

TOUCH_PAD_INTR_MASK_ALL
All touch interrupt type enable.

TOUCH_PROXIMITY_MEAS_NUM_MAX
Touch sensor proximity detection configuration

TOUCH_DEBOUNCE_CNT_MAX

TOUCH_NOISE_THR_MAX

TOUCH_JITTER_STEP_MAX

Type Definitions

typedef struct *touch_pad_denoise* **touch_pad_denoise_t**
Touch sensor denoise configuration

typedef struct *touch_pad_waterproof* **touch_pad_waterproof_t**
Touch sensor waterproof configuration

typedef struct *touch_filter_config* **touch_filter_config_t**
Touch sensor filter configuration

Enumerations

enum touch_pad_t
Touch pad channel

Values:

TOUCH_PAD_NUM0 = 0
Touch pad channel 0 is GPIO4(ESP32)

TOUCH_PAD_NUM1
Touch pad channel 1 is GPIO0(ESP32) / GPIO1(ESP32-S2)

TOUCH_PAD_NUM2
Touch pad channel 2 is GPIO2(ESP32) / GPIO2(ESP32-S2)

TOUCH_PAD_NUM3
Touch pad channel 3 is GPIO15(ESP32) / GPIO3(ESP32-S2)

TOUCH_PAD_NUM4
Touch pad channel 4 is GPIO13(ESP32) / GPIO4(ESP32-S2)

TOUCH_PAD_NUM5
Touch pad channel 5 is GPIO12(ESP32) / GPIO5(ESP32-S2)

TOUCH_PAD_NUM6
Touch pad channel 6 is GPIO14(ESP32) / GPIO6(ESP32-S2)

TOUCH_PAD_NUM7
Touch pad channel 7 is GPIO27(ESP32) / GPIO7(ESP32-S2)

TOUCH_PAD_NUM8
Touch pad channel 8 is GPIO33(ESP32) / GPIO8(ESP32-S2)

TOUCH_PAD_NUM9
Touch pad channel 9 is GPIO32(ESP32) / GPIO9(ESP32-S2)

TOUCH_PAD_NUM10
Touch channel 10 is GPIO10(ESP32-S2)

TOUCH_PAD_NUM11
Touch channel 11 is GPIO11(ESP32-S2)

TOUCH_PAD_NUM12
Touch channel 12 is GPIO12(ESP32-S2)

TOUCH_PAD_NUM13
Touch channel 13 is GPIO13(ESP32-S2)

TOUCH_PAD_NUM14
Touch channel 14 is GPIO14(ESP32-S2)

TOUCH_PAD_MAX

enum touch_high_volt_t
Touch sensor high reference voltage

Values:

TOUCH_HVOLT_KEEP = -1
Touch sensor high reference voltage, no change

TOUCH_HVOLT_2V4 = 0
Touch sensor high reference voltage, 2.4V

TOUCH_HVOLT_2V5
Touch sensor high reference voltage, 2.5V

TOUCH_HVOLT_2V6
Touch sensor high reference voltage, 2.6V

TOUCH_HVOLT_2V7
Touch sensor high reference voltage, 2.7V

TOUCH_HVOLT_MAX

enum touch_low_volt_t
Touch sensor low reference voltage

Values:

TOUCH_LVOLT_KEEP = -1
Touch sensor low reference voltage, no change

TOUCH_LVOLT_0V5 = 0
Touch sensor low reference voltage, 0.5V

TOUCH_LVOLT_0V6
Touch sensor low reference voltage, 0.6V

TOUCH_LVOLT_0V7
Touch sensor low reference voltage, 0.7V

TOUCH_LVOLT_0V8
Touch sensor low reference voltage, 0.8V

TOUCH_LVOLT_MAX

enum touch_volt_atten_t
Touch sensor high reference voltage attenuation

Values:

TOUCH_HVOLT_ATTEN_KEEP = -1
Touch sensor high reference voltage attenuation, no change

TOUCH_HVOLT_ATTEN_1V5 = 0
Touch sensor high reference voltage attenuation, 1.5V attenuation

TOUCH_HVOLT_ATTEN_1V
Touch sensor high reference voltage attenuation, 1.0V attenuation

TOUCH_HVOLT_ATTEN_0V5
Touch sensor high reference voltage attenuation, 0.5V attenuation

TOUCH_HVOLT_ATTEN_0V
Touch sensor high reference voltage attenuation, 0V attenuation

TOUCH_HVOLT_ATTEN_MAX

enum touch_cnt_slope_t
Touch sensor charge/discharge speed

Values:

TOUCH_PAD_SLOPE_0 = 0
Touch sensor charge / discharge speed, always zero

TOUCH_PAD_SLOPE_1 = 1
Touch sensor charge / discharge speed, slowest

TOUCH_PAD_SLOPE_2 = 2
Touch sensor charge / discharge speed

TOUCH_PAD_SLOPE_3 = 3
Touch sensor charge / discharge speed

TOUCH_PAD_SLOPE_4 = 4
Touch sensor charge / discharge speed

TOUCH_PAD_SLOPE_5 = 5
Touch sensor charge / discharge speed

TOUCH_PAD_SLOPE_6 = 6
Touch sensor charge / discharge speed

TOUCH_PAD_SLOPE_7 = 7
Touch sensor charge / discharge speed, fast

TOUCH_PAD_SLOPE_MAX

enum touch_tie_opt_t
Touch sensor initial charge level

Values:

TOUCH_PAD_TIE_OPT_LOW = 0
Initial level of charging voltage, low level

TOUCH_PAD_TIE_OPT_HIGH = 1
Initial level of charging voltage, high level

TOUCH_PAD_TIE_OPT_MAX

enum touch_fsm_mode_t
Touch sensor FSM mode

Values:

TOUCH_FSM_MODE_TIMER = 0
To start touch FSM by timer

TOUCH_FSM_MODE_SW
To start touch FSM by software trigger

TOUCH_FSM_MODE_MAX

enum touch_trigger_mode_t
Values:

TOUCH_TRIGGER_BELOW = 0
Touch interrupt will happen if counter value is less than threshold.

TOUCH_TRIGGER_ABOVE = 1
Touch interrupt will happen if counter value is larger than threshold.

TOUCH_TRIGGER_MAX

enum touch_trigger_src_t
Values:

TOUCH_TRIGGER_SOURCE_BOTH = 0
wakeup interrupt is generated if both SET1 and SET2 are "touched"

TOUCH_TRIGGER_SOURCE_SET1 = 1
wakeup interrupt is generated if SET1 is "touched"

TOUCH_TRIGGER_SOURCE_MAX

enum touch_pad_intr_mask_t
Values:

TOUCH_PAD_INTR_MASK_DONE = BIT(0)
Measurement done for one of the enabled channels.

TOUCH_PAD_INTR_MASK_ACTIVE = BIT(1)
Active for one of the enabled channels.

TOUCH_PAD_INTR_MASK_INACTIVE = BIT(2)
Inactive for one of the enabled channels.

TOUCH_PAD_INTR_MASK_SCAN_DONE = BIT(3)
Measurement done for all the enabled channels.

TOUCH_PAD_INTR_MASK_TIMEOUT = BIT(4)
Timeout for one of the enabled channels.

enum touch_pad_denoise_grade_t
Values:

TOUCH_PAD_DENOISE_BIT12 = 0
Denoise range is 12bit

TOUCH_PAD_DENOISE_BIT10 = 1
Denoise range is 10bit

TOUCH_PAD_DENOISE_BIT8 = 2
Denoise range is 8bit

TOUCH_PAD_DENOISE_BIT4 = 3
Denoise range is 4bit

TOUCH_PAD_DENOISE_MAX

enum touch_pad_denoise_cap_t
Values:

TOUCH_PAD_DENOISE_CAP_L0 = 0
Denoise channel internal reference capacitance is 5pf

TOUCH_PAD_DENOISE_CAP_L1 = 1
Denoise channel internal reference capacitance is 6.4pf

TOUCH_PAD_DENOISE_CAP_L2 = 2
Denoise channel internal reference capacitance is 7.8pf

TOUCH_PAD_DENOISE_CAP_L3 = 3
Denoise channel internal reference capacitance is 9.2pf

TOUCH_PAD_DENOISE_CAP_L4 = 4
Denoise channel internal reference capacitance is 10.6pf

TOUCH_PAD_DENOISE_CAP_L5 = 5
Denoise channel internal reference capacitance is 12.0pf

TOUCH_PAD_DENOISE_CAP_L6 = 6
Denoise channel internal reference capacitance is 13.4pf

TOUCH_PAD_DENOISE_CAP_L7 = 7
Denoise channel internal reference capacitance is 14.8pf

TOUCH_PAD_DENOISE_CAP_MAX = 8

enum touch_pad_shield_driver_t
Touch sensor shield channel drive capability level

Values:

TOUCH_PAD_SHIELD_DRV_L0 = 0
The max equivalent capacitance in shield channel is 40pf

TOUCH_PAD_SHIELD_DRV_L1
The max equivalent capacitance in shield channel is 80pf

TOUCH_PAD_SHIELD_DRV_L2
The max equivalent capacitance in shield channel is 120pf

TOUCH_PAD_SHIELD_DRV_L3
The max equivalent capacitance in shield channel is 160pf

TOUCH_PAD_SHIELD_DRV_L4
The max equivalent capacitance in shield channel is 200pf

TOUCH_PAD_SHIELD_DRV_L5
The max equivalent capacitance in shield channel is 240pf

TOUCH_PAD_SHIELD_DRV_L6
The max equivalent capacitance in shield channel is 280pf

TOUCH_PAD_SHIELD_DRV_L7
The max equivalent capacitance in shield channel is 320pf

TOUCH_PAD_SHIELD_DRV_MAX

enum touch_pad_conn_type_t
Touch channel idle state configuration

Values:

TOUCH_PAD_CONN_HIGHZ = 0
Idle status of touch channel is high resistance state

TOUCH_PAD_CONN_GND = 1
Idle status of touch channel is ground connection

TOUCH_PAD_CONN_MAX

enum touch_filter_mode_t
Touch channel IIR filter coefficient configuration.

Note On ESP32S2. There is an error in the IIR calculation. The magnitude of the error is twice the filter coefficient. So please select a smaller filter coefficient on the basis of meeting the filtering requirements. Recommended filter coefficient selection `IIR_16`.

Values:

TOUCH_PAD_FILTER_IIR_4 = 0
The filter mode is first-order IIR filter. The coefficient is 4.

TOUCH_PAD_FILTER_IIR_8
The filter mode is first-order IIR filter. The coefficient is 8.

TOUCH_PAD_FILTER_IIR_16
The filter mode is first-order IIR filter. The coefficient is 16 (Typical value).

TOUCH_PAD_FILTER_IIR_32
The filter mode is first-order IIR filter. The coefficient is 32.

TOUCH_PAD_FILTER_IIR_64
The filter mode is first-order IIR filter. The coefficient is 64.

TOUCH_PAD_FILTER_IIR_128
The filter mode is first-order IIR filter. The coefficient is 128.

TOUCH_PAD_FILTER_IIR_256
The filter mode is first-order IIR filter. The coefficient is 256.

TOUCH_PAD_FILTER_JITTER
The filter mode is jitter filter

TOUCH_PAD_FILTER_MAX

enum touch_smooth_mode_t
Level of filter applied on the original data against large noise interference.

Note On ESP32S2. There is an error in the IIR calculation. The magnitude of the error is twice the filter coefficient. So please select a smaller filter coefficient on the basis of meeting the filtering requirements. Recommended filter coefficient selection `IIR_2`.

Values:

TOUCH_PAD_SMOOTH_OFF = 0
No filtering of raw data.

TOUCH_PAD_SMOOTH_IIR_2 = 1
Filter the raw data. The coefficient is 2 (Typical value).

TOUCH_PAD_SMOOTH_IIR_4 = 2
Filter the raw data. The coefficient is 4.

TOUCH_PAD_SMOOTH_IIR_8 = 3
Filter the raw data. The coefficient is 8.

TOUCH_PAD_SMOOTH_MAX

2.2.18 TWAI

Overview

The Two-Wire Automotive Interface (TWAI) is a real-time serial communication protocol suited for automotive and industrial applications. It is compatible with ISO11898-1 Classical frames, thus can support Standard Frame Format (11-bit ID) and Extended Frame Format (29-bit ID). The ESP32-S2' s peripherals contains a TWAI controller that can be configured to communicate on a TWAI bus via an external transceiver.

警告: The TWAI controller is not compatible with ISO11898-1 FD Format frames, and will interpret such frames as errors.

This programming guide is split into the following sections:

1. *TWAI Protocol Summary*
2. *Signals Lines and Transceiver*
3. *Driver Configuration*
4. *Driver Operation*

5. *Examples*

TWAI Protocol Summary

The TWAI is a multi-master, multi-cast, asynchronous, serial communication protocol. TWAI also supports error detection and signalling, and inbuilt message prioritization.

Multi-master: Any node on the bus can initiate the transfer of a message.

Multi-cast: When a node transmits a message, all nodes on the bus will receive the message (i.e., broadcast) thus ensuring data consistency across all nodes. However, some nodes can selectively choose which messages to accept via the use of acceptance filtering (multi-cast).

Asynchronous: The bus does not contain a clock signal. All nodes on the bus operate at the same bit rate and synchronize using the edges of the bits transmitted on the bus.

Error Detection and Signalling: Every node will constantly monitor the bus. When any node detects an error, it will signal the detection by transmitting an error frame. Other nodes will receive the error frame and transmit their own error frames in response. This will result in an error detection being propagated to all nodes on the bus.

Message Priorities: Messages contain an ID field. If two or more nodes attempt to transmit simultaneously, the node transmitting the message with the lower ID value will win arbitration of the bus. All other nodes will become receivers ensuring that there is at most one transmitter at any time.

TWAI Messages TWAI Messages are split into Data Frames and Remote Frames. Data Frames are used to deliver a data payload to other nodes, whereas a Remote Frame is used to request a Data Frame from other nodes (other nodes can optionally respond with a Data Frame). Data and Remote Frames have two frame formats known as **Extended Frame** and **Standard Frame** which contain a 29-bit ID and an 11-bit ID respectively. A TWAI message consists of the following fields:

- 29-bit or 11-bit ID: Determines the priority of the message (lower value has higher priority).
- Data Length Code (DLC) between 0 to 8: Indicates the size (in bytes) of the data payload for a Data Frame, or the amount of data to request for a Remote Frame.
- Up to 8 bytes of data for a Data Frame (should match DLC).

Error States and Counters The TWAI protocol implements a feature known as "fault confinement" where a persistently erroneous node will eventually eliminate itself form the bus. This is implemented by requiring every node to maintain two internal error counters known as the **Transmit Error Counter (TEC)** and the **Receive Error Counter (REC)**. The two error counters are incremented and decremented according to a set of rules (where the counters increase on an error, and decrease on a successful message transmission/reception). The values of the counters are used to determine a node's **error state**, namely **Error Active**, **Error Passive**, and **Bus-Off**.

Error Active: A node is Error Active when **both TEC and REC are less than 128** and indicates that the node is operating normally. Error Active nodes are allowed to participate in bus communications, and will actively signal the detection of any errors by automatically transmitting an **Active Error Flag** over the bus.

Error Passive: A node is Error Passive when **either the TEC or REC becomes greater than or equal to 128**. Error Passive nodes are still able to take part in bus communications, but will instead transmit a **Passive Error Flag** upon detection of an error.

Bus-Off: A node becomes Bus-Off when the **TEC becomes greater than or equal to 256**. A Bus-Off node is unable influence the bus in any manner (essentially disconnected from the bus) thus eliminating itself from the bus. A node will remain in the Bus-Off state until it undergoes bus-off recovery.

Signals Lines and Transceiver

The TWAI controller does not contain a integrated transceiver. Therefore, to connect the TWAI controller to a TWAI bus, **an external transceiver is required**. The type of external transceiver used should depend on the application's physical layer specification (e.g. using SN65HVD23x transceivers for ISO 11898-2 compatibility).

The TWAI controller’s interface consists of 4 signal lines known as **TX, RX, BUS-OFF, and CLKOUT**. These four signal lines can be routed through the GPIO Matrix to the ESP32-S2’s GPIO pads.

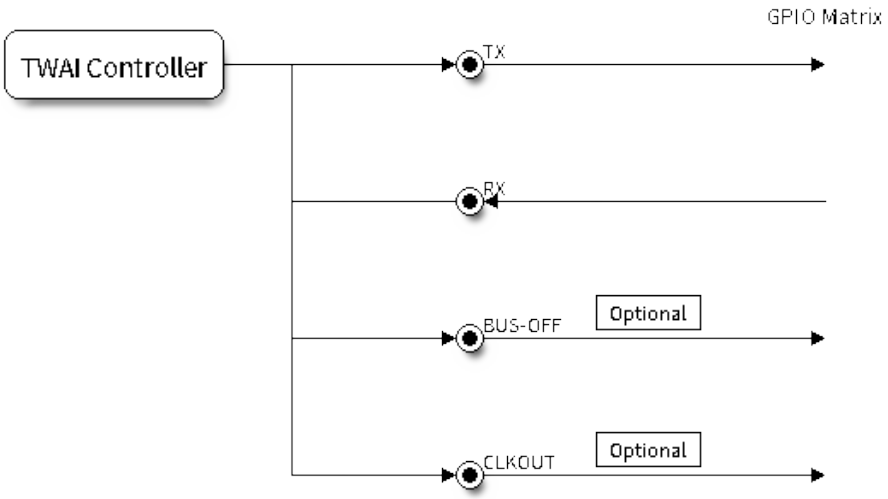


图 12: Signal lines of the TWAI controller

TX and RX: The TX and RX signal lines are required to interface with an external transceiver. Both signal lines represent/interpret a dominant bit as a low logic level (0V), and a recessive bit as a high logic level (3.3V).

BUS-OFF: The BUS-OFF signal line is **optional** and is set to a low logic level (0V) whenever the TWAI controller reaches a bus-off state. The BUS-OFF signal line is set to a high logic level (3.3V) otherwise.

CLKOUT: The CLKOUT signal line is **optional** and outputs a prescaled version of the controller’s source clock (APB Clock).

注解: An external transceiver **must internally loopback the TX to RX** such that a change in logic level to the TX signal line can be observed on the RX line. Failing to do so will cause the TWAI controller to interpret differences in logic levels between the two signal lines as a loss in arbitration or a bit error.

Driver Configuration

This section covers how to configure the TWAI driver.

Operating Modes The TWAI driver supports the following modes of operations:

Normal Mode: The normal operating mode allows the TWAI controller to take part in bus activities such as transmitting and receiving messages/error frames. Acknowledgement from another node is required when transmitting a message.

No Ack Mode: The No Acknowledgement mode is similar to normal mode, however acknowledgements are not required for a message transmission to be considered successful. This mode is useful when self testing the TWAI controller (loopback of transmissions).

Listen Only Mode: This mode will prevent the TWAI controller from influencing the bus. Therefore, transmission of messages/acknowledgement/error frames will be disabled. However the TWAI controller will still be able to receive messages but will not acknowledge the message. This mode is suited for bus monitor applications.

Alerts The TWAI driver contains an alert feature that is used to notify the application layer of certain TWAI controller or TWAI bus events. Alerts are selectively enabled when the TWAI driver is installed, but can be reconfigured during runtime by calling *twai_reconfigure_alerts()*. The application can then wait for any enabled alerts to occur by calling *twai_read_alerts()*. The TWAI driver supports the following alerts:

表 2: TWAI Driver Alerts

Alert Flag	Description
TWAI_ALERT_TX_IDLE	No more messages queued for transmission
TWAI_ALERT_TX_SUCCESS	The previous transmission was successful
TWAI_ALERT_BELOW_ERR_WARN	Both error counters have dropped below error warning limit
TWAI_ALERT_ERR_ACTIVE	TWAI controller has become error active
TWAI_ALERT_RECOVERY_IN_PROGRESS	TWAI controller is undergoing bus recovery
TWAI_ALERT_BUS_RECOVERED	TWAI controller has successfully completed bus recovery
TWAI_ALERT_ARB_LOST	The previous transmission lost arbitration
TWAI_ALERT_ABOVE_ERR_WARN	One of the error counters have exceeded the error warning limit
TWAI_ALERT_BUS_ERROR	A (Bit, Stuff, CRC, Form, ACK) error has occurred on the bus
TWAI_ALERT_TX_FAILED	The previous transmission has failed
TWAI_ALERT_RX_QUEUE_FULL	The RX queue is full causing a received frame to be lost
TWAI_ALERT_ERR_PASS	TWAI controller has become error passive
TWAI_ALERT_BUS_OFF	Bus-off condition occurred. TWAI controller can no longer influence bus

注解: The TWAI controller's **error warning limit** is used to preemptively warn the application of bus errors before the error passive state is reached. By default, the TWAI driver sets the **error warning limit** to **96**. The TWAI_ALERT_ABOVE_ERR_WARN is raised when the TEC or REC becomes larger then or equal to the error warning limit. The TWAI_ALERT_BELOW_ERR_WARN is raised when both TEC and REC return back to values below **96**.

注解: When enabling alerts, the TWAI_ALERT_AND_LOG flag can be used to cause the TWAI driver to log any raised alerts to UART. However, alert logging is disabled and TWAI_ALERT_AND_LOG if the *CONFIG_TWAI_ISR_IN_IRAM* option is enabled (see *Placing ISR into IRAM*).

注解: The TWAI_ALERT_ALL and TWAI_ALERT_NONE macros can also be used to enable/disable all alerts during configuration/reconfiguration.

Bit Timing The operating bit rate of the TWAI driver is configured using the *twai_timing_config_t* structure. The period of each bit is made up of multiple **time quanta**, and the period of a **time quanta** is determined by a prescaled version of the TWAI controller's source clock. A single bit contains the following segments in the following order:

1. The **Synchronization Segment** consists of a single time quanta
2. **Timing Segment 1** consists of 1 to 16 time quanta before sample point
3. **Timing Segment 2** consists of 1 to 8 time quanta after sample point

The **Baudrate Prescaler** is used to determine the period of each time quanta by dividing the TWAI controller's source clock (80 MHz APB clock). On the ESP32-S2, the brp can be **any even number from 2 to 32768**.

The sample point of a bit is located on the intersection of Timing Segment 1 and 2. Enabling **Triple Sampling** will cause 3 time quanta to be sampled per bit instead of 1 (extra samples are located at the tail end of Timing Segment 1).

The **Synchronization Jump Width** is used to determine the maximum number of time quanta a single bit time can be lengthened/shortened for synchronization purposes. sjw can **range from 1 to 4**.

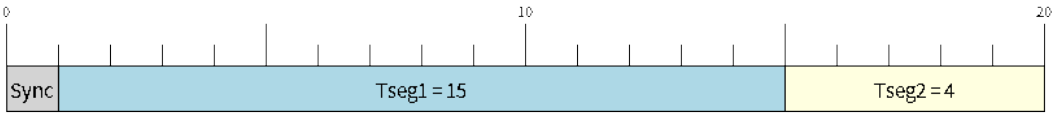


图 13: Bit timing configuration for 500kbit/s given BRP = 8

注解: Multiple combinations of `brp`, `tseg_1`, `tseg_2`, and `sjw` can achieve the same bit rate. Users should tune these values to the physical characteristics of their bus by taking into account factors such as **propagation delay, node information processing time, and phase errors**.

Bit timing **macro initializers** are also available for commonly used bit rates. The following macro initializers are provided by the TWAI driver.

- `TWAI_TIMING_CONFIG_1MBITS()`
- `TWAI_TIMING_CONFIG_800KBITS()`
- `TWAI_TIMING_CONFIG_500KBITS()`
- `TWAI_TIMING_CONFIG_250KBITS()`
- `TWAI_TIMING_CONFIG_125KBITS()`
- `TWAI_TIMING_CONFIG_100KBITS()`
- `TWAI_TIMING_CONFIG_50KBITS()`
- `TWAI_TIMING_CONFIG_25KBITS()`
- `TWAI_TIMING_CONFIG_20KBITS()`
- `TWAI_TIMING_CONFIG_16KBITS()`
- `TWAI_TIMING_CONFIG_12_5KBITS()`
- `TWAI_TIMING_CONFIG_10KBITS()`
- `TWAI_TIMING_CONFIG_5KBITS()`
- `TWAI_TIMING_CONFIG_1KBITS()`

Acceptance Filter The TWAI controller contains a hardware acceptance filter which can be used to filter messages of a particular ID. A node that filters out a message **will not receive the message, but will still acknowledge it**. Acceptance filters can make a node more efficient by filtering out messages sent over the bus that are irrelevant to the node. The acceptance filter is configured using two 32-bit values within *`twai_filter_config_t`* known as the **acceptance code** and the **acceptance mask**.

The **acceptance code** specifies the bit sequence which a message's ID, RTR, and data bytes must match in order for the message to be received by the TWAI controller. The **acceptance mask** is a bit sequence specifying which bits of the acceptance code can be ignored. This allows for a messages of different IDs to be accepted by a single acceptance code.

The acceptance filter can be used under **Single or Dual Filter Mode**. Single Filter Mode will use the acceptance code and mask to define a single filter. This allows for the first two data bytes of a standard frame to be filtered, or the entirety of an extended frame's 29-bit ID. The following diagram illustrates how the 32-bit acceptance code and mask will be interpreted under Single Filter Mode (Note: The yellow and blue fields represent standard and extended frame formats respectively).

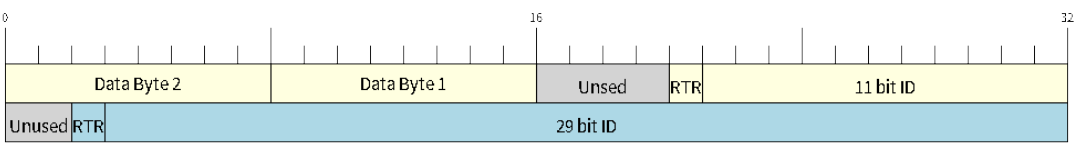


图 14: Bit layout of single filter mode (Right side MSBit)

Dual Filter Mode will use the acceptance code and mask to define two separate filters allowing for increased flexibility of ID's to accept, but does not allow for all 29-bits of an extended ID to be filtered. The following diagram illustrates how the 32-bit acceptance code and mask will be interpreted under **Dual Filter Mode** (Note: The yellow and blue fields represent standard and extended frame formats respectively).

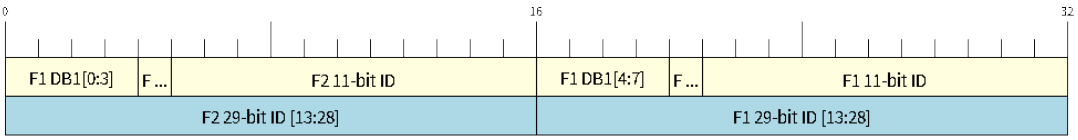


图 15: Bit layout of dual filter mode (Right side MSBit)

Disabling TX Queue The TX queue can be disabled during configuration by setting the `tx_queue_len` member of *`twai_general_config_t`* to `0`. This will allow applications that do not require message transmission to save a small amount of memory when using the TWAI driver.

Placing ISR into IRAM The TWAI driver's ISR (Interrupt Service Routine) can be placed into IRAM so that the ISR can still run whilst the cache is disabled. Placing the ISR into IRAM may be necessary to maintain the TWAI driver's functionality during lengthy cache disabling operations (such as SPI Flash writes, OTA updates etc). Whilst the cache is disabled, the ISR will continue to:

- Read received messages from the RX buffer and place them into the driver's RX queue.
- Load messages pending transmission from the driver's TX queue and write them into the TX buffer.

To place the TWAI driver's ISR, users must do the following:

- Enable the *CONFIG_TWAI_ISR_IN_IRAM* option using `idf.py menuconfig`.
- When calling *`twai_driver_install()`*, the *intr_flags* member of *`twai_general_config_t`* should set the *`ESP_INTR_FLAG_IRAM`* set.

注解: When the *CONFIG_TWAI_ISR_IN_IRAM* option is enabled, the TWAI driver will no longer log any alerts (i.e., the `TWAI_ALERT_AND_LOG` flag will not have any effect).

Driver Operation

The TWAI driver is designed with distinct states and strict rules regarding the functions or conditions that trigger a state transition. The following diagram illustrates the various states and their transitions.

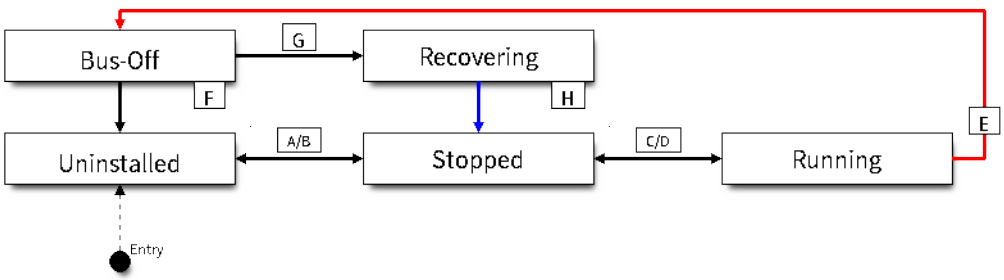


图 16: State transition diagram of the TWAI driver (see table below)

Label	Transition	Action/Condition
A	Uninstalled -> Stopped	*twai_driver_install()*
B	Stopped -> Uninstalled	*twai_driver_uninstall()*
C	Stopped -> Running	*twai_start()*
D	Running -> Stopped	*twai_stop()*
E	Running -> Bus-Off	Transmit Error Counter >= 256
F	Bus-Off -> Uninstalled	*twai_driver_uninstall()*
G	Bus-Off -> Recovering	*twai_initiate_recovery()*
H	Recovering -> Stopped	128 occurrences of 11 consecutive recessive bits.

Driver States **Uninstalled**: In the uninstalled state, no memory is allocated for the driver and the TWAI controller is powered OFF.

Stopped: In this state, the TWAI controller is powered ON and the TWAI driver has been installed. However the TWAI controller will be unable to take part in any bus activities such as transmitting, receiving, or acknowledging messages.

Running: In the running state, the TWAI controller is able to take part in bus activities. Therefore messages can be transmitted/received/acknowledged. Furthermore the TWAI controller will be able to transmit error frames upon detection of errors on the bus.

Bus-Off: The bus-off state is automatically entered when the TWAI controller's Transmit Error Counter becomes greater than or equal to 256. The bus-off state indicates the occurrence of severe errors on the bus or in the TWAI controller. Whilst in the bus-off state, the TWAI controller will be unable to take part in any bus activities. To exit the bus-off state, the TWAI controller must undergo the bus recovery process.

Recovering: The recovering state is entered when the TWAI controller undergoes bus recovery. The TWAI controller/TWAI driver will remain in the recovering state until the 128 occurrences of 11 consecutive recessive bits is observed on the bus.

Message Fields and Flags The TWAI driver distinguishes different types of messages by using the various bit field members of the *twai_message_t* structure. These bit field members determine whether a message is in standard or extended format, a remote frame, and the type of transmission to use when transmitting such a message.

These bit field members can also be toggled using the the *flags* member of *twai_message_t* and the following message flags:

Message Flag	Description
TWAI_MSG_FLAG_EXTD	Message is in Extended Frame Format (29bit ID)
TWAI_MSG_FLAG_RTR	Message is a Remote Frame (Remote Transmission Request)
TWAI_MSG_FLAG_SS	Transmit message using Single Shot Transmission (Message will not be retransmitted upon error or loss of arbitration). Unused for received message.
TWAI_MSG_FLAG_SELF	Transmit message using Self Reception Request (Transmitted message will also received by the same node). Unused for received message.
TWAI_MSG_FLAG_DLC_NON_COMP	Message's Data length code is larger than 8. This will break compliance with TWAI
TWAI_MSG_FLAG_NONE	Clears all bit fields. Equivalent to a Standard Frame Format (11bit ID) Data Frame.

Examples

Configuration & Installation The following code snippet demonstrates how to configure, install, and start the TWAI driver via the use of the various configuration structures, macro initializers, the *twai_driver_install()* function, and the *twai_start()* function.

```
#include "driver/gpio.h"
#include "driver/twai.h"

void app_main()
{
    //Initialize configuration structures using macro initializers
    twai_general_config_t g_config = TWAI_GENERAL_CONFIG_DEFAULT(GPIO_NUM_21, GPIO_
↪NUM_22, TWAI_MODE_NORMAL);
    twai_timing_config_t t_config = TWAI_TIMING_CONFIG_500KBITS();
    twai_filter_config_t f_config = TWAI_FILTER_CONFIG_ACCEPT_ALL();

    //Install TWAI driver
    if (twai_driver_install(&g_config, &t_config, &f_config) == ESP_OK) {
        printf("Driver installed\n");
    } else {
        printf("Failed to install driver\n");
        return;
    }

    //Start TWAI driver
    if (twai_start() == ESP_OK) {
        printf("Driver started\n");
    } else {
        printf("Failed to start driver\n");
        return;
    }

    ...

}
```

The usage of macro initializers is not mandatory and each of the configuration structures can be manually.

Message Transmission The following code snippet demonstrates how to transmit a message via the usage of the *twai_message_t* type and *twai_transmit()* function.

```
#include "driver/twai.h"

...

//Configure message to transmit
twai_message_t message;
message.identifier = 0xAAAA;
message.extd = 1;
message.data_length_code = 4;
for (int i = 0; i < 4; i++) {
    message.data[i] = 0;
}

//Queue message for transmission
if (twai_transmit(&message, pdMS_TO_TICKS(1000)) == ESP_OK) {
    printf("Message queued for transmission\n");
} else {
    printf("Failed to queue message for transmission\n");
}
```

Message Reception The following code snippet demonstrates how to receive a message via the usage of the *twai_message_t* type and *twai_receive()* function.

```
#include "driver/twai.h"

...

//Wait for message to be received
twai_message_t message;
if (twai_receive(&message, pdMS_TO_TICKS(10000)) == ESP_OK) {
    printf("Message received\n");
} else {
    printf("Failed to receive message\n");
    return;
}

//Process received message
if (message.extd) {
    printf("Message is in Extended Format\n");
} else {
    printf("Message is in Standard Format\n");
}
printf("ID is %d\n", message.identifier);
if (!(message.rtr)) {
    for (int i = 0; i < message.data_length_code; i++) {
        printf("Data byte %d = %d\n", i, message.data[i]);
    }
}
```

Reconfiguring and Reading Alerts The following code snippet demonstrates how to reconfigure and read TWAI driver alerts via the use of the *twai_reconfigure_alerts()* and *twai_read_alerts()* functions.

```
#include "driver/twai.h"

...

//Reconfigure alerts to detect Error Passive and Bus-Off error states
uint32_t alerts_to_enable = TWAI_ALERT_ERR_PASS | TWAI_ALERT_BUS_OFF;
if (twai_reconfigure_alerts(alerts_to_enable, NULL) == ESP_OK) {
    printf("Alerts reconfigured\n");
} else {
    printf("Failed to reconfigure alerts");
}

//Block indefinitely until an alert occurs
uint32_t alerts_triggered;
twai_read_alerts(&alerts_triggered, portMAX_DELAY);
```

Stop and Uninstall The following code demonstrates how to stop and uninstall the TWAI driver via the use of the *twai_stop()* and *twai_driver_uninstall()* functions.

```
#include "driver/twai.h"

...

//Stop the TWAI driver
if (twai_stop() == ESP_OK) {
    printf("Driver stopped\n");
} else {
    printf("Failed to stop driver\n");
    return;
}
```

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```
//Uninstall the TWAI driver
if (twai_driver_uninstall() == ESP_OK) {
    printf("Driver uninstalled\n");
} else {
    printf("Failed to uninstall driver\n");
    return;
}
```

Multiple ID Filter Configuration The acceptance mask in *twai_filter_config_t* can be configured such that two or more IDs will be accepted for a single filter. For a particular filter to accept multiple IDs, the conflicting bit positions amongst the IDs must be set in the acceptance mask. The acceptance code can be set to any one of the IDs.

The following example shows how the calculate the acceptance mask given multiple IDs:

```
ID1 =  11'b101 1010 0000
ID2 =  11'b101 1010 0001
ID3 =  11'b101 1010 0100
ID4 =  11'b101 1010 1000
//Acceptance Mask
MASK = 11'b000 0000 1101
```

Application Examples **Network Example:** The TWAI Network example demonstrates communication between two ESP32-S2s using the TWAI driver API. One TWAI node acts as a network master that initiates and ceases the transfer of a data from another node acting as a network slave. The example can be found via peripherals/twai/twai_network.

Alert and Recovery Example: This example demonstrates how to use the TWAI driver's alert and bus-off recovery API. The example purposely introduces errors on the bus to put the TWAI controller into the Bus-Off state. An alert is used to detect the Bus-Off state and trigger the bus recovery process. The example can be found via peripherals/twai/twai_alert_and_recovery.

Self Test Example: This example uses the No Acknowledge Mode and Self Reception Request to cause the TWAI controller to send and simultaneously receive a series of messages. This example can be used to verify if the connections between the TWAI controller and the external transceiver are working correctly. The example can be found via peripherals/twai/twai_self_test.

API Reference

Header File

- soc/include/hal/twai_types.h

Structures

struct twai_message_t

Structure to store a TWAI message.

Note The flags member is deprecated

Public Members

uint32_t **extd** : 1

Extended Frame Format (29bit ID)

uint32_t **rtr** : 1

Message is a Remote Frame

uint32_t **ss** : 1
Transmit as a Single Shot Transmission. Unused for received.

uint32_t **self** : 1
Transmit as a Self Reception Request. Unused for received.

uint32_t **dlc_non_comp** : 1
Message's Data length code is larger than 8. This will break compliance with ISO 11898-1

uint32_t **reserved** : 27
Reserved bits

uint32_t **flags**
Deprecated: Alternate way to set bits using message flags

uint32_t **identifier**
11 or 29 bit identifier

uint8_t **data_length_code**
Data length code

uint8_t **data**[**TWAI_FRAME_MAX_DLC**]
Data bytes (not relevant in RTR frame)

struct twai_timing_config_t
Structure for bit timing configuration of the TWAI driver.

Note Macro initializers are available for this structure

Public Members

uint32_t **brp**
Baudrate prescaler (i.e., APB clock divider). Any even number from 2 to 128 for ESP32, 2 to 32768 for ESP32S2. For ESP32 Rev 2 or later, multiples of 4 from 132 to 256 are also supported

uint8_t **tseg_1**
Timing segment 1 (Number of time quanta, between 1 to 16)

uint8_t **tseg_2**
Timing segment 2 (Number of time quanta, 1 to 8)

uint8_t **sjw**
Synchronization Jump Width (Max time quanta jump for synchronize from 1 to 4)

bool **triple_sampling**
Enables triple sampling when the TWAI controller samples a bit

struct twai_filter_config_t
Structure for acceptance filter configuration of the TWAI driver (see documentation)

Note Macro initializers are available for this structure

Public Members

uint32_t **acceptance_code**
32-bit acceptance code

uint32_t **acceptance_mask**
32-bit acceptance mask

bool **single_filter**
Use Single Filter Mode (see documentation)

Macros

TWAI_EXTD_ID_MASK
TWAI Constants.

Bit mask for 29 bit Extended Frame Format ID

TWAI_STD_ID_MASK
Bit mask for 11 bit Standard Frame Format ID

TWAI_FRAME_MAX_DLC
Max data bytes allowed in TWAI

TWAI_FRAME_EXTD_ID_LEN_BYTES
EFF ID requires 4 bytes (29bit)

TWAI_FRAME_STD_ID_LEN_BYTES
SFF ID requires 2 bytes (11bit)

TWAI_ERR_PASS_THRESH
Error counter threshold for error passive

Enumerations

enum twai_mode_t
TWAI Controller operating modes.

Values:

TWAI_MODE_NORMAL
Normal operating mode where TWAI controller can send/receive/acknowledge messages

TWAI_MODE_NO_ACK
Transmission does not require acknowledgment. Use this mode for self testing

TWAI_MODE_LISTEN_ONLY
The TWAI controller will not influence the bus (No transmissions or acknowledgments) but can receive messages

Header File

- driver/include/driver/twai.h

Functions

esp_err_t **twai_driver_install**(**const** *twai_general_config_t* *g_config, **const** *twai_timing_config_t* *t_config, **const** *twai_filter_config_t* *f_config)
Install TWAI driver.

This function installs the TWAI driver using three configuration structures. The required memory is allocated and the TWAI driver is placed in the stopped state after running this function.

Note Macro initializers are available for the configuration structures (see documentation)
Note To reinstall the TWAI driver, call twai_driver_uninstall() first
Return
- ESP_OK: Successfully installed TWAI driver
- ESP_ERR_INVALID_ARG: Arguments are invalid
- ESP_ERR_NO_MEM: Insufficient memory
- ESP_ERR_INVALID_STATE: Driver is already installed

Parameters
- `[in] g_config`: General configuration structure
- `[in] t_config`: Timing configuration structure
- `[in] f_config`: Filter configuration structure

esp_err_t **twai_driver_uninstall**(void)

Uninstall the TWAI driver.

This function uninstalls the TWAI driver, freeing the memory utilized by the driver. This function can only be called when the driver is in the stopped state or the bus-off state.

Warning The application must ensure that no tasks are blocked on TX/RX queues or alerts when this function is called.

Return

- ESP_OK: Successfully uninstalled TWAI driver
- ESP_ERR_INVALID_STATE: Driver is not in stopped/bus-off state, or is not installed

esp_err_t **twai_start**(void)

Start the TWAI driver.

This function starts the TWAI driver, putting the TWAI driver into the running state. This allows the TWAI driver to participate in TWAI bus activities such as transmitting/receiving messages. The TX and RX queue are reset in this function, clearing any messages that are unread or pending transmission. This function can only be called when the TWAI driver is in the stopped state.

Return

- ESP_OK: TWAI driver is now running
- ESP_ERR_INVALID_STATE: Driver is not in stopped state, or is not installed

esp_err_t **twai_stop**(void)

Stop the TWAI driver.

This function stops the TWAI driver, preventing any further message from being transmitted or received until twai_start() is called. Any messages in the TX queue are cleared. Any messages in the RX queue should be read by the application after this function is called. This function can only be called when the TWAI driver is in the running state.

Warning A message currently being transmitted/received on the TWAI bus will be ceased immediately. This may lead to other TWAI nodes interpreting the unfinished message as an error.

Return

- ESP_OK: TWAI driver is now Stopped
- ESP_ERR_INVALID_STATE: Driver is not in running state, or is not installed

esp_err_t **twai_transmit**(**const** *twai_message_t* **message*, TickType_t *ticks_to_wait*)

Transmit a TWAI message.

This function queues a TWAI message for transmission. Transmission will start immediately if no other messages are queued for transmission. If the TX queue is full, this function will block until more space becomes available or until it times out. If the TX queue is disabled (TX queue length = 0 in configuration), this function will return immediately if another message is undergoing transmission. This function can only be called when the TWAI driver is in the running state and cannot be called under Listen Only Mode.

Note This function does not guarantee that the transmission is successful. The TX_SUCCESS/TX_FAILED alert can be enabled to alert the application upon the success/failure of a transmission.

Note The TX_IDLE alert can be used to alert the application when no other messages are awaiting transmission.

Return

- ESP_OK: Transmission successfully queued/initiated
- ESP_ERR_INVALID_ARG: Arguments are invalid
- ESP_ERR_TIMEOUT: Timed out waiting for space on TX queue
- ESP_FAIL: TX queue is disabled and another message is currently transmitting
- ESP_ERR_INVALID_STATE: TWAI driver is not in running state, or is not installed
- ESP_ERR_NOT_SUPPORTED: Listen Only Mode does not support transmissions

Parameters

- `[in] message`: Message to transmit
- `[in] ticks_to_wait`: Number of FreeRTOS ticks to block on the TX queue

esp_err_t **twai_receive**(*twai_message_t* **message*, TickType_t *ticks_to_wait*)

Receive a TWAI message.

This function receives a message from the RX queue. The flags field of the message structure will indicate the type of message received. This function will block if there are no messages in the RX queue

Warning The flags field of the received message should be checked to determine if the received message contains any data bytes.

Return

- ESP_OK: Message successfully received from RX queue
- ESP_ERR_TIMEOUT: Timed out waiting for message
- ESP_ERR_INVALID_ARG: Arguments are invalid
- ESP_ERR_INVALID_STATE: TWAI driver is not installed

Parameters

- `[out] message`: Received message
- `[in] ticks_to_wait`: Number of FreeRTOS ticks to block on RX queue

esp_err_t **twai_read_alerts** (uint32_t **alerts*, TickType_t *ticks_to_wait*)

Read TWAI driver alerts.

This function will read the alerts raised by the TWAI driver. If no alert has been issued when this function is called, this function will block until an alert occurs or until it timeouts.

Note Multiple alerts can be raised simultaneously. The application should check for all alerts that have been enabled.

Return

- ESP_OK: Alerts read
- ESP_ERR_TIMEOUT: Timed out waiting for alerts
- ESP_ERR_INVALID_ARG: Arguments are invalid
- ESP_ERR_INVALID_STATE: TWAI driver is not installed

Parameters

- `[out] alerts`: Bit field of raised alerts (see documentation for alert flags)
- `[in] ticks_to_wait`: Number of FreeRTOS ticks to block for alert

esp_err_t **twai_reconfigure_alerts** (uint32_t *alerts_enabled*, uint32_t **current_alerts*)

Reconfigure which alerts are enabled.

This function reconfigures which alerts are enabled. If there are alerts which have not been read whilst reconfiguring, this function can read those alerts.

Return

- ESP_OK: Alerts reconfigured
- ESP_ERR_INVALID_STATE: TWAI driver is not installed

Parameters

- `[in] alerts_enabled`: Bit field of alerts to enable (see documentation for alert flags)
- `[out] current_alerts`: Bit field of currently raised alerts. Set to NULL if unused

esp_err_t **twai_initiate_recovery** (void)

Start the bus recovery process.

This function initiates the bus recovery process when the TWAI driver is in the bus-off state. Once initiated, the TWAI driver will enter the recovering state and wait for 128 occurrences of the bus-free signal on the TWAI bus before returning to the stopped state. This function will reset the TX queue, clearing any messages pending transmission.

Note The BUS_RECOVERED alert can be enabled to alert the application when the bus recovery process completes.

Return

- ESP_OK: Bus recovery started
- ESP_ERR_INVALID_STATE: TWAI driver is not in the bus-off state, or is not installed

esp_err_t **twai_get_status_info** (*twai_status_info_t* **status_info*)

Get current status information of the TWAI driver.

Return

- ESP_OK: Status information retrieved
- ESP_ERR_INVALID_ARG: Arguments are invalid

- ESP_ERR_INVALID_STATE: TWAI driver is not installed

Parameters

- `[out] status_info`: Status information

esp_err_t **twai_clear_transmit_queue**(void)

Clear the transmit queue.

This function will clear the transmit queue of all messages.

Note The transmit queue is automatically cleared when twai_stop() or twai_initiate_recovery() is called.

Return

- ESP_OK: Transmit queue cleared
- ESP_ERR_INVALID_STATE: TWAI driver is not installed or TX queue is disabled

esp_err_t **twai_clear_receive_queue**(void)

Clear the receive queue.

This function will clear the receive queue of all messages.

Note The receive queue is automatically cleared when twai_start() is called.

Return

- ESP_OK: Transmit queue cleared
- ESP_ERR_INVALID_STATE: TWAI driver is not installed

Structures

struct twai_general_config_t

Structure for general configuration of the TWAI driver.

Note Macro initializers are available for this structure

Public Members

twai_mode_t **mode**

Mode of TWAI controller

gpio_num_t **tx_io**

Transmit GPIO number

gpio_num_t **rx_io**

Receive GPIO number

gpio_num_t **clkout_io**

CLKOUT GPIO number (optional, set to -1 if unused)

gpio_num_t **bus_off_io**

Bus off indicator GPIO number (optional, set to -1 if unused)

uint32_t **tx_queue_len**

Number of messages TX queue can hold (set to 0 to disable TX Queue)

uint32_t **rx_queue_len**

Number of messages RX queue can hold

uint32_t **alerts_enabled**

Bit field of alerts to enable (see documentation)

uint32_t **clkout_divider**

CLKOUT divider. Can be 1 or any even number from 2 to 14 (optional, set to 0 if unused)

int **intr_flags**

Interrupt flags to set the priority of the driver’ s ISR. Note that to use the ESP_INTR_FLAG_IRAM, the CONFIG_TWAI_ISR_IN_IRAM option should be enabled first.

struct twai_status_info_t

Structure to store status information of TWAI driver.

Public Members

twai_state_t **state**
Current state of TWAI controller (Stopped/Running/Bus-Off/Recovery)

uint32_t **msgs_to_tx**
Number of messages queued for transmission or awaiting transmission completion

uint32_t **msgs_to_rx**
Number of messages in RX queue waiting to be read

uint32_t **tx_error_counter**
Current value of Transmit Error Counter

uint32_t **rx_error_counter**
Current value of Receive Error Counter

uint32_t **tx_failed_count**
Number of messages that failed transmissions

uint32_t **rx_missed_count**
Number of messages that were lost due to a full RX queue (or errata workaround if enabled)

uint32_t **rx_overrun_count**
Number of messages that were lost due to a RX FIFO overrun

uint32_t **arb_lost_count**
Number of instances arbitration was lost

uint32_t **bus_error_count**
Number of instances a bus error has occurred

Macros

TWAI_IO_UNUSED
Marks GPIO as unused in TWAI configuration

Enumerations

enum twai_state_t
TWAI driver states.

Values:

TWAI_STATE_STOPPED
Stopped state. The TWAI controller will not participate in any TWAI bus activities

TWAI_STATE_RUNNING
Running state. The TWAI controller can transmit and receive messages

TWAI_STATE_BUS_OFF
Bus-off state. The TWAI controller cannot participate in bus activities until it has recovered

TWAI_STATE_RECOVERING
Recovering state. The TWAI controller is undergoing bus recovery

2.2.19 UART

Overview

A Universal Asynchronous Receiver/Transmitter (UART) is a hardware feature that handles communication (i.e., timing requirements and data framing) using widely-adapted asynchronous serial communication interfaces, such as RS232, RS422, RS485. A UART provides a widely adopted and cheap method to realize full-duplex or half-duplex data exchange among different devices.

The ESP32-S2 chip has two UART controllers (UART0 and UART1) that feature an identical set of registers for ease of programming and flexibility.

Each UART controller is independently configurable with parameters such as baud rate, data bit length, bit ordering, number of stop bits, parity bit etc. All the controllers are compatible with UART-enabled devices from various manufacturers and can also support Infrared Data Association protocols (IrDA).

Functional Overview

The following overview describes how to establish communication between an ESP32-S2 and other UART devices using the functions and data types of the UART driver. The overview reflects a typical programming workflow and is broken down into the sections provided below:

1. *Setting Communication Parameters* - Setting baud rate, data bits, stop bits, etc.
2. *Setting Communication Pins* - Assigning pins for connection to a device.
3. *Driver Installation* - Allocating ESP32-S2's resources for the UART driver.
4. *Running UART Communication* - Sending / receiving data
5. *Using Interrupts* - Triggering interrupts on specific communication events
6. *Deleting a Driver* - Freeing allocated resources if a UART communication is no longer required

Steps 1 to 3 comprise the configuration stage. Step 4 is where the UART starts operating. Steps 5 and 6 are optional.

The UART driver's functions identify each of the UART controllers using `uart_port_t`. This identification is needed for all the following function calls.

Setting Communication Parameters UART communication parameters can be configured all in a single step or individually in multiple steps.

Single Step Call the function `uart_param_config()` and pass to it a `uart_config_t` structure. The `uart_config_t` structure should contain all the required parameters. See the example below.

```
const int uart_num = UART_NUM_1;
uart_config_t uart_config = {
    .baud_rate = 115200,
    .data_bits = UART_DATA_8_BITS,
    .parity = UART_PARITY_DISABLE,
    .stop_bits = UART_STOP_BITS_1,
    .flow_ctrl = UART_HW_FLOWCTRL_CTS_RTS,
    .rx_flow_ctrl_thresh = 122,
};
// Configure UART parameters
ESP_ERROR_CHECK(uart_param_config(uart_num, &uart_config));
```

Multiple Steps Configure specific parameters individually by calling a dedicated function from the table given below. These functions are also useful if re-configuring a single parameter.

表 3: Functions for Configuring specific parameters individually

Parameter to Configure	Function
Baud rate	`uart_set_baudrate()`
Number of transmitted bits	`uart_set_word_length()` selected out of `uart_word_length_t`
Parity control	`uart_set_parity()` selected out of `uart_parity_t`
Number of stop bits	`uart_set_stop_bits()` selected out of `uart_stop_bits_t`
Hardware flow control mode	`uart_set_hw_flow_ctrl()` selected out of `uart_hw_flowcontrol_t`
Communication mode	`uart_set_mode()` selected out of `uart_mode_t`

Each of the above functions has a _get_ counterpart to check the currently set value. For example, to check the current baud rate value, call *uart_get_baudrate()*.

Setting Communication Pins After setting communication parameters, configure the physical GPIO pins to which the other UART device will be connected. For this, call the function *uart_set_pin()* and specify the GPIO pin numbers to which the driver should route the Tx, Rx, RTS, and CTS signals. If you want to keep a currently allocated pin number for a specific signal, pass the macro *UART_PIN_NO_CHANGE*.

The same macro should be specified for pins that will not be used.

```
// Set UART pins(TX: IO17 (UART1 default), RX: IO18 (UART1 default), RTS: IO19,␣
↪CTS: IO20)
ESP_ERROR_CHECK(uart_set_pin(UART_NUM_1, UART_PIN_NO_CHANGE, UART_PIN_NO_CHANGE,␣
↪19, 20));
```

Driver Installation Once the communication pins are set, install the driver by calling *uart_driver_install()* and specify the following parameters:

- Size of Tx ring buffer
- Size of Rx ring buffer
- Event queue handle and size
- Flags to allocate an interrupt

The function will allocate the required internal resources for the UART driver.

```
// Setup UART buffered IO with event queue
const int uart_buffer_size = (1024 * 2);
QueueHandle_t uart_queue;
// Install UART driver using an event queue here
ESP_ERROR_CHECK(uart_driver_install(UART_NUM_1, uart_buffer_size, \
                                        uart_buffer_size, 10, &uart_queue, 0));
```

Once this step is complete, you can connect the external UART device and check the communication.

Running UART Communication Serial communication is controlled by each UART controller's finite state machine (FSM).

The process of sending data involves the following steps:

1. Write data into Tx FIFO buffer
2. FSM serializes the data
3. FSM sends the data out

The process of receiving data is similar, but the steps are reversed:

1. FSM processes an incoming serial stream and parallelizes it
2. FSM writes the data into Rx FIFO buffer
3. Read the data from Rx FIFO buffer

Therefore, an application will be limited to writing and reading data from a respective buffer using *uart_write_bytes()* and *uart_read_bytes()* respectively, and the FSM will do the rest.

Transmitting After preparing the data for transmission, call the function *uart_write_bytes()* and pass the data buffer's address and data length to it. The function will copy the data to the Tx ring buffer (either immediately or after enough space is available), and then exit. When there is free space in the Tx FIFO buffer, an interrupt service routine (ISR) moves the data from the Tx ring buffer to the Tx FIFO buffer in the background. The code below demonstrates the use of this function.

```
// Write data to UART.
char* test_str = "This is a test string.\n";
uart_write_bytes(uart_num, (const char*)test_str, strlen(test_str));
```

The function *uart_write_bytes_with_break()* is similar to *uart_write_bytes()* but adds a serial break signal at the end of the transmission. A 'serial break signal' means holding the Tx line low for a period longer than one data frame.

```
// Write data to UART, end with a break signal.
uart_write_bytes_with_break(uart_num, "test break\n",strlen("test break\n"), 100);
```

Another function for writing data to the Tx FIFO buffer is *uart_tx_chars()*. Unlike *uart_write_bytes()*, this function will not block until space is available. Instead, it will write all data which can immediately fit into the hardware Tx FIFO, and then return the number of bytes that were written.

There is a 'companion' function *uart_wait_tx_done()* that monitors the status of the Tx FIFO buffer and returns once it is empty.

```
// Wait for packet to be sent
const int uart_num = UART_NUM_1;
ESP_ERROR_CHECK(uart_wait_tx_done(uart_num, 100)); // wait timeout is 100 RTOS␣
↪ticks (TickType_t)
```

Receiving Once the data is received by the UART and saved in the Rx FIFO buffer, it needs to be retrieved using the function *uart_read_bytes()*. Before reading data, you can check the number of bytes available in the Rx FIFO buffer by calling *uart_get_buffered_data_len()*. An example of using these functions is given below.

```
// Read data from UART.
const int uart_num = UART_NUM_1;
uint8_t data[128];
int length = 0;
ESP_ERROR_CHECK(uart_get_buffered_data_len(uart_num, (size_t*)&length));
length = uart_read_bytes(uart_num, data, length, 100);
```

If the data in the Rx FIFO buffer is no longer needed, you can clear the buffer by calling *uart_flush()*.

Software Flow Control If the hardware flow control is disabled, you can manually set the RTS and DTR signal levels by using the functions *uart_set_rts()* and *uart_set_dtr()* respectively.

Communication Mode Selection The UART controller supports a number of communication modes. A mode can be selected using the function *uart_set_mode()*. Once a specific mode is selected, the UART driver will handle the behavior of a connected UART device accordingly. As an example, it can control the RS485 driver chip using the RTS line to allow half-duplex RS485 communication.

```
// Setup UART in rs485 half duplex mode
ESP_ERROR_CHECK(uart_set_mode(uart_num, UART_MODE_RS485_HALF_DUPLEX));
```

Using Interrupts There are many interrupts that can be generated following specific UART states or detected errors. The full list of available interrupts is provided ESP32-S2 Technical Reference Manual (PDF).. You can enable or disable specific interrupts by calling *uart_enable_intr_mask()* or *uart_disable_intr_mask()* respectively. The mask of all interrupts is available as `UART_INTR_MASK`.

By default, the *uart_driver_install()* function installs the driver's internal interrupt handler to manage the Tx and Rx ring buffers and provides high-level API functions like events (see below). It is also possible to register a lower level interrupt handler instead using *uart_isr_register()*, and to free it again using *uart_isr_free()*. Some UART driver functions which use the Tx and Rx ring buffers, events, etc. will not

automatically work in this case - it is necessary to handle the interrupts directly in the ISR. Inside the custom handler implementation, clear the interrupt status bits using *uart_clear_intr_status()*.

The API provides a convenient way to handle specific interrupts discussed in this document by wrapping them into dedicated functions:

- **Event detection**: There are several events defined in *uart_event_type_t* that may be reported to a user application using the FreeRTOS queue functionality. You can enable this functionality when calling *uart_driver_install()* described in *Driver Installation*. An example of using Event detection can be found in peripherals/uart/uart_events.
- **FIFO space threshold or transmission timeout reached**: The Tx and Rx FIFO buffers can trigger an interrupt when they are filled with a specific number of characters, or on a timeout of sending or receiving data. To use these interrupts, do the following:
  - Configure respective threshold values of the buffer length and timeout by entering them in the structure *uart_intr_config_t* and calling *uart_intr_config()*
  - Enable the interrupts using the functions *uart_enable_tx_intr()* and *uart_enable_rx_intr()*
  - Disable these interrupts using the corresponding functions *uart_disable_tx_intr()* or *uart_disable_rx_intr()*
- **Pattern detection**: An interrupt triggered on detecting a 'pattern' of the same character being received/sent repeatedly for a number of times. This functionality is demonstrated in the example peripherals/uart/uart_events. It can be used, e.g., to detect a command string followed by a specific number of identical characters (the 'pattern') added at the end of the command string. The following functions are available:
  - Configure and enable this interrupt using `uart_enable_pattern_det_intr()`
  - Disable the interrupt using *uart_disable_pattern_det_intr()*

Macros The API also defines several macros. For example, `UART_FIFO_LEN` defines the length of hardware FIFO buffers; `UART_BITRATE_MAX` gives the maximum baud rate supported by the UART controllers, etc.

Deleting a Driver If the communication established with *uart_driver_install()* is no longer required, the driver can be removed to free allocated resources by calling *uart_driver_delete()*.

Overview of RS485 specific communication options

注解: The following section will use `[UART_REGISTER_NAME].[UART_FIELD_BIT]` to refer to UART register fields/bits. To find more information on a specific option bit, open the Register Summary section of the SoC Technical Reference Manual. Use the register name to navigate to the register description and then find the field/bit.

- `UART_RS485_CONF_REG.UART_RS485_EN`: setting this bit enables RS485 communication mode support.
- `UART_RS485_CONF_REG.UART_RS485TX_RX_EN`: if this bit is set, the transmitter's output signal loops back to the receiver's input signal.
- `UART_RS485_CONF_REG.UART_RS485RXBY_TX_EN`: if this bit is set, the transmitter will still be sending data if the receiver is busy (remove collisions automatically by hardware).

The ESP32-S2's RS485 UART hardware can detect signal collisions during transmission of a datagram and generate the interrupt `UART_RS485_CLASH_INT` if this interrupt is enabled. The term collision means that a transmitted datagram is not equal to the one received on the other end. Data collisions are usually associated with the presence of other active devices on the bus or might occur due to bus errors.

The collision detection feature allows handling collisions when their interrupts are activated and triggered. The interrupts `UART_RS485_FRM_ERR_INT` and `UART_RS485_PARITY_ERR_INT` can be used with the collision detection feature to control frame errors and parity bit errors accordingly in RS485 mode. This functionality is supported in the UART driver and can be used by selecting the *UART_MODE_RS485_APP_CTRL* mode (see the function *uart_set_mode()*).

The collision detection feature can work with circuit A and circuit C (see Section *Interface Connection Options*). In the case of using circuit A or B, the RTS pin connected to the DE pin of the bus driver should be controlled by the user application. Use the function `uart_get_collision_flag()` to check if the collision detection flag has been raised.

The ESP32-S2 UART controllers themselves do not support half-duplex communication as they cannot provide automatic control of the RTS pin connected to the ~RE/DE input of RS485 bus driver. However, half-duplex communication can be achieved via software control of the RTS pin by the UART driver. This can be enabled by selecting the `UART_MODE_RS485_HALF_DUPLEX` mode when calling `uart_set_mode()`.

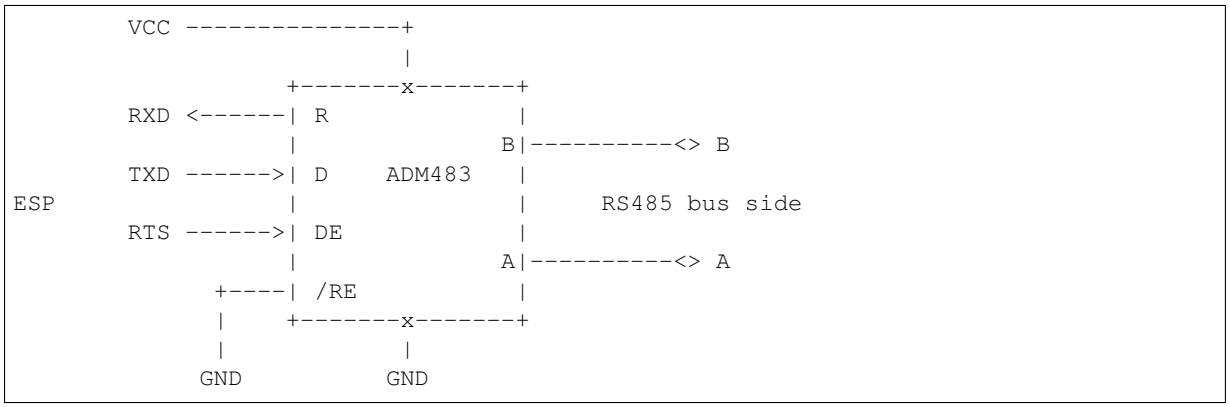
Once the host starts writing data to the Tx FIFO buffer, the UART driver automatically asserts the RTS pin (logic 1); once the last bit of the data has been transmitted, the driver de-asserts the RTS pin (logic 0). To use this mode, the software would have to disable the hardware flow control function. This mode works with all the used circuits shown below.

Interface Connection Options This section provides example schematics to demonstrate the basic aspects of ESP32-S2’s RS485 interface connection.

注解:

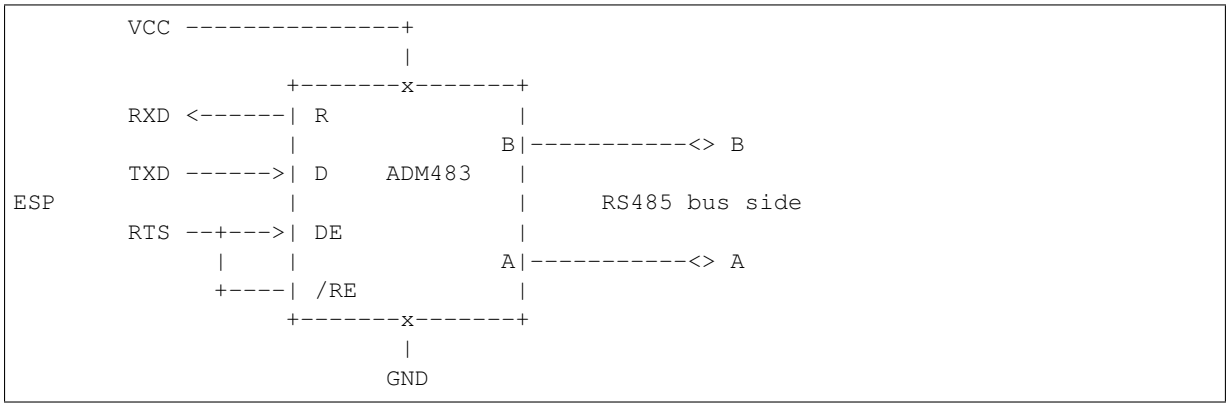
- The schematics below do **not** necessarily contain **all required elements**.
- The **analog devices** ADM483 & ADM2483 are examples of common RS485 transceivers and **can be replaced** with other similar transceivers.

Circuit A: Collision Detection Circuit



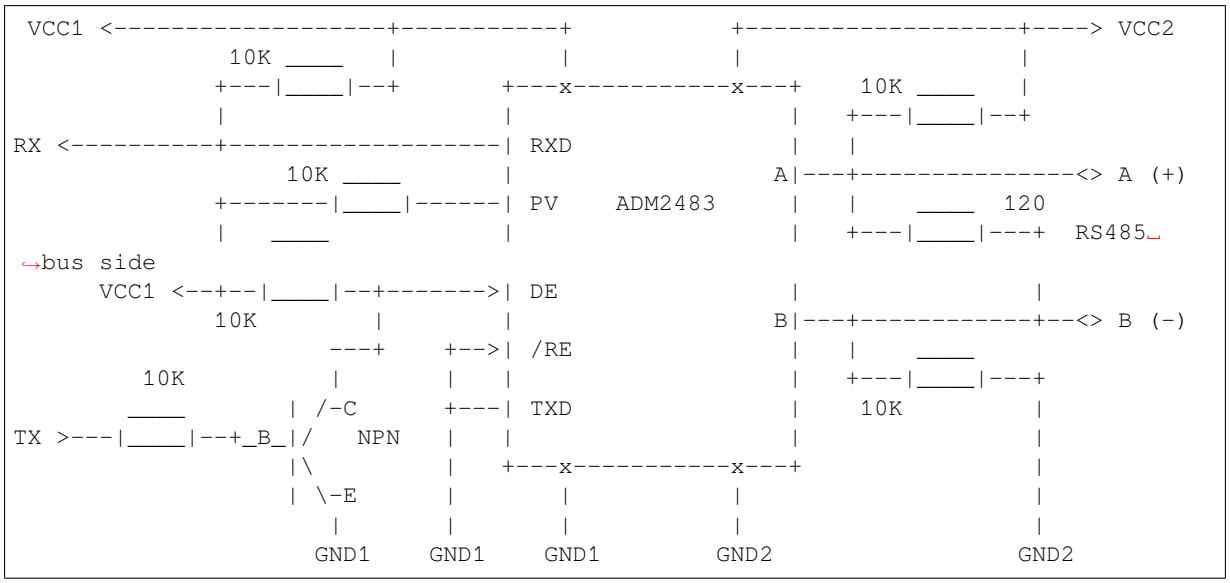
This circuit is preferable because it allows for collision detection and is quite simple at the same time. The receiver in the line driver is constantly enabled, which allows the UART to monitor the RS485 bus. Echo suppression is performed by the UART peripheral when the bit `UART_RS485_CONF_REG.UART_RS485TX_RX_EN` is enabled.

Circuit B: Manual Switching Transmitter/Receiver Without Collision Detection



This circuit does not allow for collision detection. It suppresses the null bytes that the hardware receives when the bit `UART_RS485_CONF_REG.UART_RS485TX_RX_EN` is set. The bit `UART_RS485_CONF_REG.UART_RS485RXBY_TX_EN` is not applicable in this case.

Circuit C: Auto Switching Transmitter/Receiver



This galvanically isolated circuit does not require RTS pin control by a software application or driver because it controls the transceiver direction automatically. However, it requires suppressing null bytes during transmission by setting `UART_RS485_CONF_REG.UART_RS485RXBY_TX_EN` to 1 and `UART_RS485_CONF_REG.UART_RS485TX_RX_EN` to 0. This setup can work in any RS485 UART mode or even in *`UART_MODE_UART`*.

Application Examples

The table below describes the code examples available in the directory peripherals/uart/.

Code Example	Description
peripherals/uart/uart_echo	Configuring UART settings, installing the UART driver, and reading/writing over the UART1 interface.
peripherals/uart/uart_events	Reporting various communication events, using pattern detection interrupts.
peripherals/uart/uart_async_rxtxtasks	Transmitting and receiving data in two separate FreeRTOS tasks over the same UART.
peripherals/uart/uart_select	Using synchronous I/O multiplexing for UART file descriptors.
peripherals/uart/uart_echo_rs485	Setting up UART driver to communicate over RS485 interface in half-duplex mode. This example is similar to peripherals/uart/uart_echo but allows communication through an RS485 interface chip connected to ESP32-S2 pins.
peripherals/uart/nmea0183_parser	Obtaining GPS information by parsing NMEA0183 statements received from GPS via the UART peripheral.

API Reference

Header File

- driver/include/driver/uart.h

Functions

esp_err_t **uart_driver_install**(*uart_port_t uart_num*, int *rx_buffer_size*, int *tx_buffer_size*, int *queue_size*, *QueueHandle_t* **uart_queue*, int *intr_alloc_flags*)

Install UART driver and set the UART to the default configuration.

UART ISR handler will be attached to the same CPU core that this function is running on.

Note Rx_buffer_size should be greater than UART_FIFO_LEN. Tx_buffer_size should be either zero or greater than UART_FIFO_LEN.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `rx_buffer_size`: UART RX ring buffer size.
- `tx_buffer_size`: UART TX ring buffer size. If set to zero, driver will not use TX buffer, TX function will block task until all data have been sent out.
- `queue_size`: UART event queue size/depth.
- `uart_queue`: UART event queue handle (out param). On success, a new queue handle is written here to provide access to UART events. If set to NULL, driver will not use an event queue.
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info. Do not set ESP_INTR_FLAG_IRAM here (the driver's ISR handler is not located in IRAM)

esp_err_t **uart_driver_delete**(*uart_port_t uart_num*)

Uninstall UART driver.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

bool **uart_is_driver_installed**(*uart_port_t uart_num*)

Checks whether the driver is installed or not.

Return

- true driver is installed
- false driver is not installed

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_set_word_length**(*uart_port_t uart_num*, *uart_word_length_t data_bit*)

Set UART data bits.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `data_bit`: UART data bits

esp_err_t **uart_get_word_length**(*uart_port_t uart_num*, *uart_word_length_t* **data_bit*)

Get the UART data bit configuration.

Return

- ESP_FAIL Parameter error
- ESP_OK Success, result will be put in (*data_bit)

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `data_bit`: Pointer to accept value of UART data bits.

esp_err_t **uart_set_stop_bits**(*uart_port_t uart_num*, *uart_stop_bits_t stop_bits*)

Set UART stop bits.

Return
- ESP_OK Success
- ESP_FAIL Fail

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `stop_bits`: UART stop bits

esp_err_t **uart_get_stop_bits** (*uart_port_t uart_num*, *uart_stop_bits_t *stop_bits*)
Get the UART stop bit configuration.

Return
- ESP_FAIL Parameter error
- ESP_OK Success, result will be put in (*stop_bit)

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `stop_bits`: Pointer to accept value of UART stop bits.

esp_err_t **uart_set_parity** (*uart_port_t uart_num*, *uart_parity_t parity_mode*)
Set UART parity mode.

Return
- ESP_FAIL Parameter error
- ESP_OK Success

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `parity_mode`: the enum of uart parity configuration

esp_err_t **uart_get_parity** (*uart_port_t uart_num*, *uart_parity_t *parity_mode*)
Get the UART parity mode configuration.

Return
- ESP_FAIL Parameter error
- ESP_OK Success, result will be put in (*parity_mode)

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `parity_mode`: Pointer to accept value of UART parity mode.

esp_err_t **uart_set_baudrate** (*uart_port_t uart_num*, uint32_t *baudrate*)
Set UART baud rate.

Return
- ESP_FAIL Parameter error
- ESP_OK Success

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `baudrate`: UART baud rate.

esp_err_t **uart_get_baudrate** (*uart_port_t uart_num*, uint32_t *baudrate*)
Get the UART baud rate configuration.

Return
- ESP_FAIL Parameter error
- ESP_OK Success, result will be put in (*baudrate)

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `baudrate`: Pointer to accept value of UART baud rate

esp_err_t **uart_set_line_inverse** (*uart_port_t uart_num*, uint32_t *inverse_mask*)
Set UART line inverse mode.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `inverse_mask`: Choose the wires that need to be inverted. Using the ORred mask of `uart_signal_inv_t`

esp_err_t **`uart_set_hw_flow_ctrl`** (*uart_port_t uart_num*, *uart_hw_flowcontrol_t flow_ctrl*, uint8_t *rx_thresh*)

Set hardware flow control.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `flow_ctrl`: Hardware flow control mode
- `rx_thresh`: Threshold of Hardware RX flow control (0 ~ UART_FIFO_LEN). Only when UART_HW_FLOWCTRL_RTS is set, will the rx_thresh value be set.

esp_err_t **`uart_set_sw_flow_ctrl`** (*uart_port_t uart_num*, bool *enable*, uint8_t *rx_thresh_xon*, uint8_t *rx_thresh_xoff*)

Set software flow control.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters
- `uart_num`: UART_NUM_0, UART_NUM_1 or UART_NUM_2
- `enable`: switch on or off
- `rx_thresh_xon`: low water mark
- `rx_thresh_xoff`: high water mark

esp_err_t **`uart_get_hw_flow_ctrl`** (*uart_port_t uart_num*, *uart_hw_flowcontrol_t *flow_ctrl*)

Get the UART hardware flow control configuration.

Return
- ESP_FAIL Parameter error
- ESP_OK Success, result will be put in (*flow_ctrl)

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `flow_ctrl`: Option for different flow control mode.

esp_err_t **`uart_clear_intr_status`** (*uart_port_t uart_num*, uint32_t *clr_mask*)

Clear UART interrupt status.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `clr_mask`: Bit mask of the interrupt status to be cleared.

esp_err_t **`uart_enable_intr_mask`** (*uart_port_t uart_num*, uint32_t *enable_mask*)

Set UART interrupt enable.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `enable_mask`: Bit mask of the enable bits.

esp_err_t **`uart_disable_intr_mask`** (*uart_port_t uart_num*, uint32_t *disable_mask*)

Clear UART interrupt enable bits.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `disable_mask`: Bit mask of the disable bits.

esp_err_t **uart_enable_rx_intr**(*uart_port_t uart_num*)

Enable UART RX interrupt (RX_FULL & RX_TIMEOUT INTERRUPT)

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_disable_rx_intr**(*uart_port_t uart_num*)

Disable UART RX interrupt (RX_FULL & RX_TIMEOUT INTERRUPT)

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_disable_tx_intr**(*uart_port_t uart_num*)

Disable UART TX interrupt (TX_FULL & TX_TIMEOUT INTERRUPT)

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number

esp_err_t **uart_enable_tx_intr**(*uart_port_t uart_num*, int *enable*, int *thresh*)

Enable UART TX interrupt (TX_FULL & TX_TIMEOUT INTERRUPT)

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `enable`: 1: enable; 0: disable
- `thresh`: Threshold of TX interrupt, 0 ~ UART_FIFO_LEN

esp_err_t **uart_isr_register**(*uart_port_t uart_num*, void (**fn*)) void *

, void **arg*, int *intr_alloc_flags*, *uart_isr_handle_t* **handle*Register UART interrupt handler (ISR).

Note UART ISR handler will be attached to the same CPU core that this function is running on.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `fn`: Interrupt handler function.
- `arg`: parameter for handler function
- `intr_alloc_flags`: Flags used to allocate the interrupt. One or multiple (ORred) ESP_INTR_FLAG_* values. See esp_intr_alloc.h for more info.
- `handle`: Pointer to return handle. If non-NULL, a handle for the interrupt will be returned here.

esp_err_t **uart_isr_free**(*uart_port_t uart_num*)

Free UART interrupt handler registered by uart_isr_register. Must be called on the same core as uart_isr_register was called.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_set_pin**(*uart_port_t uart_num*, int *tx_io_num*, int *rx_io_num*, int *rts_io_num*, int *cts_io_num*)

Set UART pin number.

Note Internal signal can be output to multiple GPIO pads. Only one GPIO pad can connect with input signal.

Note Instead of GPIO number a macro 'UART_PIN_NO_CHANGE' may be provided to keep the currently allocated pin.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `tx_io_num`: UART TX pin GPIO number.
- `rx_io_num`: UART RX pin GPIO number.
- `rts_io_num`: UART RTS pin GPIO number.
- `cts_io_num`: UART CTS pin GPIO number.

esp_err_t **uart_set_rts**(*uart_port_t uart_num*, int *level*)

Manually set the UART RTS pin level.

Note UART must be configured with hardware flow control disabled.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `level`: 1: RTS output low (active); 0: RTS output high (block)

esp_err_t **uart_set_dtr**(*uart_port_t uart_num*, int *level*)

Manually set the UART DTR pin level.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `level`: 1: DTR output low; 0: DTR output high

esp_err_t **uart_set_tx_idle_num**(*uart_port_t uart_num*, uint16_t *idle_num*)

Set UART idle interval after tx FIFO is empty.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `idle_num`: idle interval after tx FIFO is empty(unit: the time it takes to send one bit under current baudrate)

esp_err_t **uart_param_config**(*uart_port_t uart_num*, **const** *uart_config_t* **uart_config*)

Set UART configuration parameters.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `uart_config`: UART parameter settings

esp_err_t **uart_intr_config**(*uart_port_t uart_num*, **const** *uart_intr_config_t* **intr_conf*)
Configure UART interrupts.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `intr_conf`: UART interrupt settings

esp_err_t **uart_wait_tx_done**(*uart_port_t uart_num*, TickType_t *ticks_to_wait*)
Wait until UART TX FIFO is empty.

Return
- ESP_OK Success
- ESP_FAIL Parameter error
- ESP_ERR_TIMEOUT Timeout

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `ticks_to_wait`: Timeout, count in RTOS ticks

int **uart_tx_chars**(*uart_port_t uart_num*, **const** char **buffer*, uint32_t *len*)
Send data to the UART port from a given buffer and length.

This function will not wait for enough space in TX FIFO. It will just fill the available TX FIFO and return when the FIFO is full.

Note This function should only be used when UART TX buffer is not enabled.

Return
- (-1) Parameter error
- OTHERS (>=0) The number of bytes pushed to the TX FIFO

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `buffer`: data buffer address
- `len`: data length to send

int **uart_write_bytes**(*uart_port_t uart_num*, **const** char **src*, size_t *size*)
Send data to the UART port from a given buffer and length,.

If the UART driver's parameter 'tx_buffer_size' is set to zero: This function will not return until all the data have been sent out, or at least pushed into TX FIFO.

Otherwise, if the 'tx_buffer_size' > 0, this function will return after copying all the data to tx ring buffer, UART ISR will then move data from the ring buffer to TX FIFO gradually.

Return
- (-1) Parameter error
- OTHERS (>=0) The number of bytes pushed to the TX FIFO

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `src`: data buffer address
- `size`: data length to send

int **uart_write_bytes_with_break**(*uart_port_t uart_num*, **const** char **src*, size_t *size*, int *brk_len*)
Send data to the UART port from a given buffer and length,.

If the UART driver's parameter 'tx_buffer_size' is set to zero: This function will not return until all the data and the break signal have been sent out. After all data is sent out, send a break signal.

Otherwise, if the 'tx_buffer_size' > 0, this function will return after copying all the data to tx ring buffer, UART ISR will then move data from the ring buffer to TX FIFO gradually. After all data sent out, send a break signal.

Return

- (-1) Parameter error
- OTHERS (>=0) The number of bytes pushed to the TX FIFO

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `src`: data buffer address
- `size`: data length to send
- `brk_len`: break signal duration(unit: the time it takes to send one bit at current baudrate)

int **uart_read_bytes**(*uart_port_t uart_num*, uint8_t **buf*, uint32_t *length*, TickType_t *ticks_to_wait*)

UART read bytes from UART buffer.

Return

- (-1) Error
- OTHERS (>=0) The number of bytes read from UART FIFO

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `buf`: pointer to the buffer.
- `length`: data length
- `ticks_to_wait`: sTimeout, count in RTOS ticks

esp_err_t **uart_flush**(*uart_port_t uart_num*)

Alias of uart_flush_input. UART ring buffer flush. This will discard all data in the UART RX buffer.

Note Instead of waiting the data sent out, this function will clear UART rx buffer. In order to send all the data in tx FIFO, we can use uart_wait_tx_done function.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_flush_input**(*uart_port_t uart_num*)

Clear input buffer, discard all the data is in the ring-buffer.

Note In order to send all the data in tx FIFO, we can use uart_wait_tx_done function.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_get_buffered_data_len**(*uart_port_t uart_num*, size_t **size*)

UART get RX ring buffer cached data length.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).
- `size`: Pointer of size_t to accept cached data length

esp_err_t **uart_disable_pattern_det_intr**(*uart_port_t uart_num*)

UART disable pattern detect function. Designed for applications like ‘AT commands’. When the hardware detects a series of one same character, the interrupt will be triggered.

Return

- ESP_OK Success
- ESP_FAIL Parameter error

Parameters

- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_enable_pattern_det_baud_intr**(*uart_port_t uart_num*, char *pattern_chr*, uint8_t *chr_num*, int *chr_tout*, int *post_idle*, int *pre_idle*)

UART enable pattern detect function. Designed for applications like ‘AT commands’. When the hardware detect a series of one same character, the interrupt will be triggered.

Return
- ESP_OK Success
- ESP_FAIL Parameter error

Parameters
- `uart_num`: UART port number.
- `pattern_chr`: character of the pattern.
- `chr_num`: number of the character, 8bit value.
- `chr_tout`: timeout of the interval between each pattern characters, 16bit value, unit is the baud-rate cycle you configured. When the duration is more than this value, it will not take this data as at_cmd char.
- `post_idle`: idle time after the last pattern character, 16bit value, unit is the baud-rate cycle you configured. When the duration is less than this value, it will not take the previous data as the last at_cmd char
- `pre_idle`: idle time before the first pattern character, 16bit value, unit is the baud-rate cycle you configured. When the duration is less than this value, it will not take this data as the first at_cmd char.

int **uart_pattern_pop_pos**(*uart_port_t uart_num*)

Return the nearest detected pattern position in buffer. The positions of the detected pattern are saved in a queue, this function will dequeue the first pattern position and move the pointer to next pattern position.

The following APIs will modify the pattern position info: uart_flush_input, uart_read_bytes, uart_driver_delete, uart_pop_pattern_pos It is the application’s responsibility to ensure atomic access to the pattern queue and the rx data buffer when using pattern detect feature.

Note If the RX buffer is full and flow control is not enabled, the detected pattern may not be found in the rx buffer due to overflow.

Return
- (-1) No pattern found for current index or parameter error
- others the pattern position in rx buffer.

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

int **uart_pattern_get_pos**(*uart_port_t uart_num*)

Return the nearest detected pattern position in buffer. The positions of the detected pattern are saved in a queue, This function do nothing to the queue.

The following APIs will modify the pattern position info: uart_flush_input, uart_read_bytes, uart_driver_delete, uart_pop_pattern_pos It is the application’s responsibility to ensure atomic access to the pattern queue and the rx data buffer when using pattern detect feature.

Note If the RX buffer is full and flow control is not enabled, the detected pattern may not be found in the rx buffer due to overflow.

Return
- (-1) No pattern found for current index or parameter error
- others the pattern position in rx buffer.

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

esp_err_t **uart_pattern_queue_reset**(*uart_port_t uart_num*, int *queue_length*)

Allocate a new memory with the given length to save record the detected pattern position in rx buffer.

Return
- ESP_ERR_NO_MEM No enough memory
- ESP_ERR_INVALID_STATE Driver not installed
- ESP_FAIL Parameter error
- ESP_OK Success

Parameters
- `uart_num`: UART port number, the max port number is (UART_NUM_MAX -1).

- `queue_length`: Max queue length for the detected pattern. If the queue length is not large enough, some pattern positions might be lost. Set this value to the maximum number of patterns that could be saved in data buffer at the same time.

esp_err_t **uart_set_mode** (*uart_port_t uart_num*, *uart_mode_t mode*)

UART set communication mode.

Note This function must be executed after uart_driver_install(), when the driver object is initialized.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `uart_num`: Uart number to configure, the max port number is (UART_NUM_MAX -1).
- `mode`: UART UART mode to set

esp_err_t **uart_set_rx_full_threshold** (*uart_port_t uart_num*, int *threshold*)

Set uart threshold value for RX fifo full.

Note If application is using higher baudrate and it is observed that bytes in hardware RX fifo are overwritten then this threshold can be reduced

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_INVALID_STATE Driver is not installed

Parameters

- `uart_num`: UART_NUM_0, UART_NUM_1 or UART_NUM_2
- `threshold`: Threshold value above which RX fifo full interrupt is generated

esp_err_t **uart_set_tx_empty_threshold** (*uart_port_t uart_num*, int *threshold*)

Set uart threshold values for TX fifo empty.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_INVALID_STATE Driver is not installed

Parameters

- `uart_num`: UART_NUM_0, UART_NUM_1 or UART_NUM_2
- `threshold`: Threshold value below which TX fifo empty interrupt is generated

esp_err_t **uart_set_rx_timeout** (*uart_port_t uart_num*, **const** uint8_t *tout_thresh*)

UART set threshold timeout for TOUT feature.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_INVALID_STATE Driver is not installed

Parameters

- `uart_num`: Uart number to configure, the max port number is (UART_NUM_MAX -1).
- `tout_thresh`: This parameter defines timeout threshold in uart symbol periods. The maximum value of threshold is 126. tout_thresh = 1, defines TOUT interrupt timeout equal to transmission time of one symbol (~11 bit) on current baudrate. If the time is expired the UART_RXFIFO_TOUT_INT interrupt is triggered. If tout_thresh == 0, the TOUT feature is disabled.

esp_err_t **uart_get_collision_flag** (*uart_port_t uart_num*, bool **collision_flag*)

Returns collision detection flag for RS485 mode Function returns the collision detection flag into variable pointed by collision_flag. *collision_flag = true, if collision detected else it is equal to false. This function should be executed when actual transmission is completed (after uart_write_bytes()).

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Parameter error

Parameters

- `uart_num`: Uart number to configure the max port number is (UART_NUM_MAX -1).

- `collision_flag`: Pointer to variable of type bool to return collision flag.

esp_err_t **uart_set_wakeup_threshold**(*uart_port_t uart_num*, int *wakeup_threshold*)

Set the number of RX pin signal edges for light sleep wakeup.

UART can be used to wake up the system from light sleep. This feature works by counting the number of positive edges on RX pin and comparing the count to the threshold. When the count exceeds the threshold, system is woken up from light sleep. This function allows setting the threshold value.

Stop bit and parity bits (if enabled) also contribute to the number of edges. For example, letter 'a' with ASCII code 97 is encoded as 0100001101 on the wire (with 8n1 configuration), start and stop bits included. This sequence has 3 positive edges (transitions from 0 to 1). Therefore, to wake up the system when 'a' is sent, set wakeup_threshold=3.

The character that triggers wakeup is not received by UART (i.e. it can not be obtained from UART FIFO). Depending on the baud rate, a few characters after that will also not be received. Note that when the chip enters and exits light sleep mode, APB frequency will be changing. To make sure that UART has correct baud rate all the time, select REF_TICK as UART clock source, by setting use_ref_tick field in *uart_config_t* to true.

Note in ESP32, the wakeup signal can only be input via IO_MUX (i.e. GPIO3 should be configured as function_1 to wake up UART0, GPIO9 should be configured as function_5 to wake up UART1), UART2 does not support light sleep wakeup feature.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if uart_num is incorrect or wakeup_threshold is outside of [3, 0x3ff] range.

Parameters

- `uart_num`: UART number, the max port number is (UART_NUM_MAX -1).
- `wakeup_threshold`: number of RX edges for light sleep wakeup, value is 3 .. 0x3ff.

esp_err_t **uart_get_wakeup_threshold**(*uart_port_t uart_num*, int **out_wakeup_threshold*)

Get the number of RX pin signal edges for light sleep wakeup.

See description of uart_set_wakeup_threshold for the explanation of UART wakeup feature.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if out_wakeup_threshold is NULL

Parameters

- `uart_num`: UART number, the max port number is (UART_NUM_MAX -1).
- `[out] out_wakeup_threshold`: output, set to the current value of wakeup threshold for the given UART.

esp_err_t **uart_wait_tx_idle_polling**(*uart_port_t uart_num*)

Wait until UART tx memory empty and the last char send ok (polling mode).

- **Return**
  - ESP_OK on success
  - ESP_ERR_INVALID_ARG Parameter error
  - ESP_FAIL Driver not installed

Parameters

- `uart_num`: UART number

esp_err_t **uart_set_loop_back**(*uart_port_t uart_num*, bool *loop_back_en*)

Configure TX signal loop back to RX module, just for the test usage.

- **Return**
  - ESP_OK on success
  - ESP_ERR_INVALID_ARG Parameter error
  - ESP_FAIL Driver not installed

Parameters

- `uart_num`: UART number
- `loop_back_en`: Set ture to enable the loop back function, else set it false.

void **uart_set_always_rx_timeout** (*uart_port_t uart_num*, bool *always_rx_timeout_en*)

Configure behavior of UART RX timeout interrupt.

When always_rx_timeout is true, timeout interrupt is triggered even if FIFO is full. This function can cause extra timeout interrupts triggered only to send the timeout event. Call this function only if you want to ensure timeout interrupt will always happen after a byte stream.

Parameters

- `uart_num`: UART number
- `always_rx_timeout_en`: Set to false enable the default behavior of timeout interrupt, set it to true to always trigger timeout interrupt.

Structures

struct uart_intr_config_t

UART interrupt configuration parameters for uart_intr_config function.

Public Members

uint32_t **intr_enable_mask**

UART interrupt enable mask, choose from UART_XXXX_INT_ENA_M under UART_INT_ENA_REG(i), connect with bit-or operator

uint8_t **rx_timeout_thresh**

UART timeout interrupt threshold (unit: time of sending one byte)

uint8_t **txfifo_empty_intr_thresh**

UART TX empty interrupt threshold.

uint8_t **rxfifo_full_thresh**

UART RX full interrupt threshold.

struct uart_event_t

Event structure used in UART event queue.

Public Members

uart_event_type_t **type**

UART event type

size_t **size**

UART data size for UART_DATA event

bool **timeout_flag**

UART data read timeout flag for UART_DATA event (no new data received during configured RX TOUT) If the event is caused by FIFO-full interrupt, then there will be no event with the timeout flag before the next byte coming.

Macros

UART_NUM_0

UART port 0

UART_NUM_1

UART port 1

UART_NUM_MAX

UART port max

UART_PIN_NO_CHANGE

Constant for uart_set_pin function which indicates that UART pin should not be changed

Type Definitions

typedef *intr_handle_t* **uart_isr_handle_t**

Enumerations

enum uart_event_type_t

UART event types used in the ring buffer.

Values:

UART_DATA
UART data event

UART_BREAK
UART break event

UART_BUFFER_FULL
UART RX buffer full event

UART_FIFO_OVF
UART FIFO overflow event

UART_FRAME_ERR
UART RX frame error event

UART_PARITY_ERR
UART RX parity event

UART_DATA_BREAK
UART TX data and break event

UART_PATTERN_DET
UART pattern detected

UART_EVENT_MAX
UART event max index

Header File

- soc/include/hal/uart_types.h

Structures

struct uart_at_cmd_t

UART AT cmd char configuration parameters Note that this function may different on different chip. Please refer to the TRM at confirguration.

Public Members

uint8_t **cmd_char**
UART AT cmd char

uint8_t **char_num**
AT cmd char repeat number

uint32_t **gap_tout**
gap time(in baud-rate) between AT cmd char

uint32_t **pre_idle**
the idle time(in baud-rate) between the non AT char and first AT char

uint32_t **post_idle**
the idle time(in baud-rate) between the last AT char and the none AT char

struct uart_sw_flowctrl_t

UART software flow control configuration parameters.

Public Members

uint8_t **xon_char**
Xon flow control char

uint8_t **xoff_char**
Xoff flow control char

uint8_t **xon_thrd**
If the software flow control is enabled and the data amount in rxfifo is less than xon_thrd, an xon_char will be sent

uint8_t **xoff_thrd**
If the software flow control is enabled and the data amount in rxfifo is more than xoff_thrd, an xoff_char will be sent

struct uart_config_t
UART configuration parameters for uart_param_config function.

Public Members

int **baud_rate**
UART baud rate

uart_word_length_t **data_bits**
UART byte size

uart_parity_t **parity**
UART parity mode

uart_stop_bits_t **stop_bits**
UART stop bits

uart_hw_flowcontrol_t **flow_ctrl**
UART HW flow control mode (cts/rts)

uint8_t **rx_flow_ctrl_thresh**
UART HW RTS threshold

uart_sclk_t **source_clk**
UART source clock selection

bool **use_ref_tick**
Deprecated method to select ref tick clock source, set source_clk field instead

Type Definitions

typedef int **uart_port_t**
UART port number, can be UART_NUM_0 ~ (UART_NUM_MAX -1).

Enumerations

enum uart_mode_t
UART mode selection.

Values:

UART_MODE_UART = 0x00
mode: regular UART mode

UART_MODE_RS485_HALF_DUPLEX = 0x01
mode: half duplex RS485 UART mode control by RTS pin

UART_MODE_IRDA = 0x02
mode: IRDA UART mode

UART_MODE_RS485_COLLISION_DETECT = 0x03
mode: RS485 collision detection UART mode (used for test purposes)

UART_MODE_RS485_APP_CTRL = 0x04
mode: application control RS485 UART mode (used for test purposes)

enum uart_word_length_t
UART word length constants.

Values:

UART_DATA_5_BITS = 0x0
word length: 5bits

UART_DATA_6_BITS = 0x1
word length: 6bits

UART_DATA_7_BITS = 0x2
word length: 7bits

UART_DATA_8_BITS = 0x3
word length: 8bits

UART_DATA_BITS_MAX = 0x4

enum uart_stop_bits_t
UART stop bits number.

Values:

UART_STOP_BITS_1 = 0x1
stop bit: 1bit

UART_STOP_BITS_1_5 = 0x2
stop bit: 1.5bits

UART_STOP_BITS_2 = 0x3
stop bit: 2bits

UART_STOP_BITS_MAX = 0x4

enum uart_parity_t
UART parity constants.

Values:

UART_PARITY_DISABLE = 0x0
Disable UART parity

UART_PARITY_EVEN = 0x2
Enable UART even parity

UART_PARITY_ODD = 0x3
Enable UART odd parity

enum uart_hw_flowcontrol_t
UART hardware flow control modes.

Values:

UART_HW_FLOWCTRL_DISABLE = 0x0
disable hardware flow control

UART_HW_FLOWCTRL_RTS = 0x1
enable RX hardware flow control (rts)

UART_HW_FLOWCTRL_CTS = 0x2
enable TX hardware flow control (cts)

UART_HW_FLOWCTRL_CTS_RTS = 0x3
enable hardware flow control

UART_HW_FLOWCTRL_MAX = 0x4

enum uart_signal_inv_t
UART signal bit map.

Values:

UART_SIGNAL_INV_DISABLE = 0
Disable UART signal inverse

UART_SIGNAL_IRDA_TX_INV = (0x1 « 0)
inverse the UART irda_tx signal

UART_SIGNAL_IRDA_RX_INV = (0x1 « 1)
inverse the UART irda_rx signal

UART_SIGNAL_RXD_INV = (0x1 « 2)
inverse the UART rxd signal

UART_SIGNAL_CTS_INV = (0x1 « 3)
inverse the UART cts signal

UART_SIGNAL_DSR_INV = (0x1 « 4)
inverse the UART dsr signal

UART_SIGNAL_TXD_INV = (0x1 « 5)
inverse the UART txd signal

UART_SIGNAL_RTS_INV = (0x1 « 6)
inverse the UART rts signal

UART_SIGNAL_DTR_INV = (0x1 « 7)
inverse the UART dtr signal

enum uart_sclk_t
UART source clock.

Values:

UART_SCLK_APB = 0x0
UART source clock from APB

UART_SCLK_REF_TICK = 0x01
UART source clock from REF_TICK

GPIO Lookup Macros The UART peripherals have dedicated IO_MUX pins to which they are connected directly. However, signals can also be routed to other pins using the less direct GPIO matrix. To use direct routes, you need to know which pin is a dedicated IO_MUX pin for a UART channel. GPIO Lookup Macros simplify the process of finding and assigning IO_MUX pins. You choose a macro based on either the IO_MUX pin number, or a required UART channel name, and the macro will return the matching counterpart for you. See some examples below.

注解： These macros are useful if you need very high UART baud rates (over 40 MHz), which means you will have to use IO_MUX pins only. In other cases, these macros can be ignored, and you can use the GPIO Matrix as it allows you to configure any GPIO pin for any UART function.

---

1. `UART_NUM_2_TXD_DIRECT_GPIO_NUM` returns the IO_MUX pin number of UART channel 2 TXD pin (pin 17)
2. *`UART_GPIO19_DIRECT_CHANNEL`* returns the UART number of GPIO 19 when connected to the UART peripheral via IO_MUX (this is UART_NUM_0)
3. `UART_CTS_GPIO19_DIRECT_CHANNEL` returns the UART number of GPIO 19 when used as the UART CTS pin via IO_MUX (this is UART_NUM_0). Similar to the above macro but specifies the pin function which is also part of the IO_MUX assignment.

Header File

- soc/soc/esp32s2/include/soc/uart_channel.h

Macros

UART_GPIO43_DIRECT_CHANNEL
UART_NUM_0_TXD_DIRECT_GPIO_NUM

UART_GPIO44_DIRECT_CHANNEL

UART_NUM_0_RXD_DIRECT_GPIO_NUM

UART_GPIO16_DIRECT_CHANNEL

UART_NUM_0_CTS_DIRECT_GPIO_NUM

UART_GPIO15_DIRECT_CHANNEL

UART_NUM_0_RTS_DIRECT_GPIO_NUM

UART_TXD_GPIO43_DIRECT_CHANNEL

UART_RXD_GPIO44_DIRECT_CHANNEL

UART_CTS_GPIO16_DIRECT_CHANNEL

UART_RTS_GPIO15_DIRECT_CHANNEL

UART_GPIO17_DIRECT_CHANNEL

UART_NUM_1_TXD_DIRECT_GPIO_NUM

UART_GPIO18_DIRECT_CHANNEL

UART_NUM_1_RXD_DIRECT_GPIO_NUM

UART_GPIO20_DIRECT_CHANNEL

UART_NUM_1_CTS_DIRECT_GPIO_NUM

UART_GPIO19_DIRECT_CHANNEL

UART_NUM_1_RTS_DIRECT_GPIO_NUM

UART_TXD_GPIO17_DIRECT_CHANNEL

UART_RXD_GPIO18_DIRECT_CHANNEL

UART_CTS_GPIO20_DIRECT_CHANNEL

UART_RTS_GPIO19_DIRECT_CHANNEL

本部分的 API 示例代码存放在 ESP-IDF 示例项目的 peripherals 目录下。

2.3 应用层协议

2.3.1 mDNS 服务

概述

mDNS 是一种组播 UDP 服务，用来提供本地网络服务和主机发现。

绝大多数的操作系统默认都会安装 mDNS 服务，或者提供单独的安装包。`Mac OS` 默认会安装名为 `Bonjour` 的服务（该服务基于 mDNS），此外 Apple 还发布了适用于 Windows 系统的安装程序，可以在官方支持 找到。在 `Linux` 上，mDNS 服务由 avahi 提供，通常也会被默认安装。

mDNS 属性

- hostname：设备会去响应的主机名，如果没有设置，会根据设备的网络接口名定义 hostname 。例如，my-esp32s2 会被解析为 my-esp32s2.local。
- default_instance：默认实例名（即易记的设备名），例如 Jhon's ESP32-S2 Thing。如果没有设置，将会使用 hostname。

以下为 STA 接口启动 mDNS 服务并设置 hostname 和 default_instance 的示例方法：

```
void start_mdns_service()
{
    // 初始化 mDNS 服务
    esp_err_t err = mdns_init();
    if (err) {
        printf("MDNS Init failed: %d\n", err);
        return;
    }

    // 设置 hostname
    mdns_hostname_set("my-esp32s2");
    // 设置默认实例
    mdns_instance_name_set("Jhon's ESP32-S2 Thing");
}
```

mDNS 服务 mDNS 可以广播设备能够提供的网络服务的相关信息，每个服务会由以下属性构成。

- instance_name：实例名（即易记的服务名），例如 Jhon's ESP32-S2 Web Server。如果没有定义，会使用 default_instance。
- service_type：（必需）服务类型，以下划线为前缀，这里 列出了常见的类型。
- proto：（必需）服务运行所依赖的协议，以下划线为前缀，例如 _tcp 或者 _udp。
- port：（必需）服务运行所用的端口号。
- txt：形如 {var, val} 的字符串数组，用于定义服务的属性。

添加一些服务和不同属性的示例方法：

```
void add_mdns_services()
{
    // 添加服务
    mdns_service_add(NULL, "_http", "_tcp", 80, NULL, 0);
    mdns_service_add(NULL, "_arduino", "_tcp", 3232, NULL, 0);
    mdns_service_add(NULL, "_myservice", "_udp", 1234, NULL, 0);

    // 注意：必须先添加服务，然后才能设置其属性
    // web 服务器使用自定义的实例名
    mdns_service_instance_name_set("_http", "_tcp", "Jhon's ESP32-S2 Web Server");

    mdns_txt_item_t serviceTxtData[3] = {
        {"board","esp32s2"},
        {"u","user"},
        {"p","password"}
    };
    // 设置服务的文本数据（会释放并替换当前数据）
    mdns_service_txt_set("_http", "_tcp", serviceTxtData, 3);

    // 修改服务端口号
    mdns_service_port_set("_myservice", "_udp", 4321);
}
```

mDNS 查询 mDNS 提供查询服务和解析主机 IP/IPv6 地址的方法。

服务查询的结果会作为 mdns_result_t 类型对象的链表返回。

解析主机 IP 地址的示例方法:

```
void resolve_mdns_host(const char * host_name)
{
    printf("Query A: %s.local", host_name);

    struct ip4_addr addr;
    addr.addr = 0;

    esp_err_t err = mdns_query_a(host_name, 2000,  &addr);
    if(err){
        if(err == ESP_ERR_NOT_FOUND){
            printf("Host was not found!");
            return;
        }
        printf("Query Failed");
        return;
    }

    printf(IPSTR, IP2STR(&addr));
}
```

解析本地服务的示例方法:

```
static const char * if_str[] = {"STA", "AP", "ETH", "MAX"};
static const char * ip_protocol_str[] = {"V4", "V6", "MAX"};

void mdns_print_results(mdns_result_t * results){
    mdns_result_t * r = results;
    mdns_ip_addr_t * a = NULL;
    int i = 1, t;
    while(r){
        printf("%d: Interface: %s, Type: %s\n", i++, if_str[r->tcpip_if], ip_
↪protocol_str[r->ip_protocol]);
        if(r->instance_name){
            printf("  PTR : %s\n", r->instance_name);
        }
        if(r->hostname){
            printf("  SRV : %s.local:%u\n", r->hostname, r->port);
        }
        if(r->txt_count){
            printf("  TXT : [%u] ", r->txt_count);
            for(t=0; t<r->txt_count; t++){
                printf("%s=%s; ", r->txt[t].key, r->txt[t].value);
            }
            printf("\n");
        }
        a = r->addr;
        while(a){
            if(a->addr.type == IPADDR_TYPE_V6){
                printf("  AAAA: " IPV6STR "\n", IPV62STR(a->addr.u_addr.ip6));
            } else {
                printf("  A   : " IPSTR "\n", IP2STR(&(a->addr.u_addr.ip4)));
            }
            a = a->next;
        }
        r = r->next;
    }

}

void find_mdns_service(const char * service_name, const char * proto)
{
```

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```
    ESP_LOGI(TAG, "Query PTR: %s.%s.local", service_name, proto);

    mdns_result_t * results = NULL;
    esp_err_t err = mdns_query_ptr(service_name, proto, 3000, 20,  &results);
    if(err){
        ESP_LOGE(TAG, "Query Failed");
        return;
    }
    if(!results){
        ESP_LOGW(TAG, "No results found!");
        return;
    }

    mdns_print_results(results);
    mdns_query_results_free(results);
}
```

使用上述方法的示例:

```
void my_app_some_method(){
    // 搜索 esp32s2-mdns.local
    resolve_mdns_host("esp32s2-mdns");

    // 搜索 HTTP 服务器
    find_mdns_service("_http", "_tcp");
    // 或者搜索文件服务器
    find_mdns_service("_smb", "_tcp"); // Windows 系统的共享服务
    find_mdns_service("_afpovertcp", "_tcp"); // Apple AFP 文件共享服务
    find_mdns_service("_nfs", "_tcp"); // NFS 服务器
    find_mdns_service("_ftp", "_tcp"); // FTP 服务器
    // 或者网络打印机
    find_mdns_service("_printer", "_tcp");
    find_mdns_service("_ipp", "_tcp");
}
```

应用示例

有关 mDNS 服务器和查询器的应用示例请参考 protocols/mdns。

API 参考

Header File

- mdns/include/mdns.h

Functions

esp_err_t **mdns_init** (void)

Initialize mDNS on given interface.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE when failed to register event handler
- ESP_ERR_NO_MEM on memory error
- ESP_FAIL when failed to start mdns task

void **mdns_free** (void)

Stop and free mDNS server.

esp_err_t **mdns_hostname_set**(**const** char **hostname*)

Set the hostname for mDNS server required if you want to advertise services.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NO_MEM memory error

Parameters

- `hostname`: Hostname to set

esp_err_t **mdns_instance_name_set**(**const** char **instance_name*)

Set the default instance name for mDNS server.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NO_MEM memory error

Parameters

- `instance_name`: Instance name to set

esp_err_t **mdns_service_add**(**const** char **instance_name*, **const** char **service_type*, **const** char **proto*, uint16_t *port*, *mdns_txt_item_t* *txt*[], size_t *num_items*)

Add service to mDNS server.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NO_MEM memory error
- ESP_FAIL failed to add service

Parameters

- `instance_name`: instance name to set. If NULL, global instance name or hostname will be used
- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)
- `port`: service port
- `txt`: string array of TXT data (eg. {{ “var” ,” val” },{ “other” ,” 2” }})
- `num_items`: number of items in TXT data

esp_err_t **mdns_service_remove**(**const** char **service_type*, **const** char **proto*)

Remove service from mDNS server.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NOT_FOUND Service not found
- ESP_ERR_NO_MEM memory error

Parameters

- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)

esp_err_t **mdns_service_instance_name_set**(**const** char **service_type*, **const** char **proto*, **const** char **instance_name*)

Set instance name for service.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NOT_FOUND Service not found
- ESP_ERR_NO_MEM memory error

Parameters

- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)
- `instance_name`: instance name to set

esp_err_t **mdns_service_port_set**(**const** char **service_type*, **const** char **proto*, uint16_t *port*)

Set service port.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NOT_FOUND Service not found
- ESP_ERR_NO_MEM memory error

Parameters

- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)
- `port`: service port

esp_err_t **mdns_service_txt_set**(**const** char **service_type*, **const** char **proto*, *mdns_txt_item_t* *txt*[], uint8_t *num_items*)

Replace all TXT items for service.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NOT_FOUND Service not found
- ESP_ERR_NO_MEM memory error

Parameters

- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)
- `txt`: array of TXT data (eg. {{ “var” ,” val” },{ “other” ,” 2” }})
- `num_items`: number of items in TXT data

esp_err_t **mdns_service_txt_item_set**(**const** char **service_type*, **const** char **proto*, **const** char **key*, **const** char **value*)

Set/Add TXT item for service TXT record.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NOT_FOUND Service not found
- ESP_ERR_NO_MEM memory error

Parameters

- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)
- `key`: the key that you want to add/update
- `value`: the new value of the key

esp_err_t **mdns_service_txt_item_remove**(**const** char **service_type*, **const** char **proto*, **const** char **key*)

Remove TXT item for service TXT record.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error
- ESP_ERR_NOT_FOUND Service not found
- ESP_ERR_NO_MEM memory error

Parameters

- `service_type`: service type (_http, _ftp, etc)
- `proto`: service protocol (_tcp, _udp)
- `key`: the key that you want to remove

esp_err_t **mdns_service_remove_all**(void)

Remove and free all services from mDNS server.

Return

- ESP_OK success
- ESP_ERR_INVALID_ARG Parameter error

esp_err_t **mdns_query**(**const** char **name*, **const** char **service_type*, **const** char **proto*, uint16_t *type*, uint32_t *timeout*, size_t *max_results*, *mdns_result_t* ***results*)

Query mDNS for host or service All following query methods are derived from this one.

Return

- ESP_OK success
- ESP_ERR_INVALID_STATE mDNS is not running
- ESP_ERR_NO_MEM memory error
- ESP_ERR_INVALID_ARG timeout was not given

Parameters

- `name`: service instance or host name (NULL for PTR queries)
- `service_type`: service type (_http, _arduino, _ftp etc.) (NULL for host queries)
- `proto`: service protocol (_tcp, _udp, etc.) (NULL for host queries)
- `type`: type of query (MDNS_TYPE_*)
- `timeout`: time in milliseconds to wait for answers.
- `max_results`: maximum results to be collected
- `results`: pointer to the results of the query results must be freed using mdns_query_results_free below

void **mdns_query_results_free**(*mdns_result_t* **results*)

Free query results.

Parameters

- `results`: linked list of results to be freed

esp_err_t **mdns_query_ptr**(**const** char **service_type*, **const** char **proto*, uint32_t *timeout*, size_t *max_results*, *mdns_result_t* ***results*)

Query mDNS for service.

Return

- ESP_OK success
- ESP_ERR_INVALID_STATE mDNS is not running
- ESP_ERR_NO_MEM memory error
- ESP_ERR_INVALID_ARG parameter error

Parameters

- `service_type`: service type (_http, _arduino, _ftp etc.)
- `proto`: service protocol (_tcp, _udp, etc.)
- `timeout`: time in milliseconds to wait for answer.
- `max_results`: maximum results to be collected
- `results`: pointer to the results of the query

esp_err_t **mdns_query_srv**(**const** char **instance_name*, **const** char **service_type*, **const** char **proto*, uint32_t *timeout*, *mdns_result_t* ***result*)

Query mDNS for SRV record.

Return

- ESP_OK success
- ESP_ERR_INVALID_STATE mDNS is not running
- ESP_ERR_NO_MEM memory error
- ESP_ERR_INVALID_ARG parameter error

Parameters

- `instance_name`: service instance name
- `service_type`: service type (_http, _arduino, _ftp etc.)
- `proto`: service protocol (_tcp, _udp, etc.)
- `timeout`: time in milliseconds to wait for answer.
- `result`: pointer to the result of the query

esp_err_t **mdns_query_txt**(**const** char **instance_name*, **const** char **service_type*, **const** char **proto*, uint32_t *timeout*, *mdns_result_t* ***result*)

Query mDNS for TXT record.

Return

- ESP_OK success

- ESP_ERR_INVALID_STATE mDNS is not running
- ESP_ERR_NO_MEM memory error
- ESP_ERR_INVALID_ARG parameter error

Parameters

- `instance_name`: service instance name
- `service_type`: service type (_http, _arduino, _ftp etc.)
- `proto`: service protocol (_tcp, _udp, etc.)
- `timeout`: time in milliseconds to wait for answer.
- `result`: pointer to the result of the query

esp_err_t **mdns_query_a**(**const** char **host_name*, uint32_t *timeout*, esp_ip4_addr_t **addr*)

Query mDNS for A record.

Return

- ESP_OK success
- ESP_ERR_INVALID_STATE mDNS is not running
- ESP_ERR_NO_MEM memory error
- ESP_ERR_INVALID_ARG parameter error

Parameters

- `host_name`: host name to look for
- `timeout`: time in milliseconds to wait for answer.
- `addr`: pointer to the resulting IP4 address

esp_err_t **mdns_query_aaaa**(**const** char **host_name*, uint32_t *timeout*, esp_ip6_addr_t **addr*)

Query mDNS for A record.

Return

- ESP_OK success
- ESP_ERR_INVALID_STATE mDNS is not running
- ESP_ERR_NO_MEM memory error
- ESP_ERR_INVALID_ARG parameter error

Parameters

- `host_name`: host name to look for
- `timeout`: time in milliseconds to wait for answer. If 0, max_results needs to be defined
- `addr`: pointer to the resulting IP6 address

esp_err_t **mdns_handle_system_event**(void **ctx*, *system_event_t* **event*)

System event handler This method controls the service state on all active interfaces and applications are required to call it from the system event handler for normal operation of mDNS service.

Parameters

- `ctx`: The system event context
- `event`: The system event

Structures

struct mdns_txt_item_t

mDNS basic text item structure Used in mdns_service_add()

Public Members

const char ***key**

item key name

const char ***value**

item value string

struct mdns_ip_addr_s

mDNS query linked list IP item

Public Members

esp_ip_addr_t **addr**
IP address

struct *mdns_ip_addr_s* ***next**
next IP, or NULL for the last IP in the list

struct mdns_result_s
mDNS query result structure

Public Members

struct *mdns_result_s* ***next**
next result, or NULL for the last result in the list

mdns_if_t **tcpip_if**
interface index

mdns_ip_protocol_t **ip_protocol**
ip_protocol type of the interface (v4/v6)

char ***instance_name**
instance name

char ***hostname**
hostname

uint16_t **port**
service port

mdns_txt_item_t ***txt**
txt record

size_t **txt_count**
number of txt items

mdns_ip_addr_t ***addr**
linked list of IP addresses found

Macros

MDNS_TYPE_A
MDNS_TYPE_PTR
MDNS_TYPE_TXT
MDNS_TYPE_AAAA
MDNS_TYPE_SRV
MDNS_TYPE_OPT
MDNS_TYPE_NSEC
MDNS_TYPE_ANY

Type Definitions

typedef struct *mdns_ip_addr_s* **mdns_ip_addr_t**
mDNS query linked list IP item

typedef enum *mdns_if_internal* **mdns_if_t**

typedef struct *mdns_result_s* **mdns_result_t**
mDNS query result structure

Enumerations

enum mdns_ip_protocol_t

mDNS enum to specify the ip_protocol type

Values:

MDNS_IP_PROTOCOL_V4

MDNS_IP_PROTOCOL_V6

MDNS_IP_PROTOCOL_MAX

enum mdns_if_internal

Values:

MDNS_IF_STA = 0

MDNS_IF_AP = 1

MDNS_IF_ETH = 2

MDNS_IF_MAX

2.3.2 ESP-TLS

Overview

The ESP-TLS component provides a simplified API interface for accessing the commonly used TLS functionality. It supports common scenarios like CA certification validation, SNI, ALPN negotiation, non-blocking connection among others. All the configuration can be specified in the `esp_tls_cfg_t` data structure. Once done, TLS communication can be conducted using the following APIs:

- *`esp_tls_conn_new()`*: for opening a new TLS connection.
- *`esp_tls_conn_read()`*: for reading from the connection.
- *`esp_tls_conn_write()`*: for writing into the connection.
- *`esp_tls_conn_delete()`*: for freeing up the connection.

Any application layer protocol like HTTP1, HTTP2 etc can be executed on top of this layer.

Application Example

Simple HTTPS example that uses ESP-TLS to establish a secure socket connection: protocols/https_request.

Tree structure for ESP-TLS component

```
├── esp_tls.c
├── esp_tls.h
├── esp_tls_mbedtls.c
├── esp_tls_wolfssl.c
└── private_include
    ├── esp_tls_mbedtls.h
    └── esp_tls_wolfssl.h
```

The ESP-TLS component has a file esp-tls/esp_tls.h which contain the public API headers for the component. Internally ESP-TLS component uses one of the two SSL/TLS Libraries between mbedtls and wolfssl for its operation. API specific to mbedtls are present in esp-tls/private_include/esp_tls_mbedtls.h and API specific to wolfssl are present in esp-tls/private_include/esp_tls_wolfssl.h.

Underlying SSL/TLS Library Options

The ESP-TLS component has an option to use mbedtls or wolfssl as their underlying SSL/TLS library. By default only mbedtls is available and is used, wolfssl SSL/TLS library is available publicly at https://github.com/espressif/esp-wolfssl. The repository provides wolfssl component in binary format, it also provides few examples which are useful for understanding the API. Please refer the repository README.md for information on licensing and other options. Please see below option for using wolfssl in your project.

注解： *As the library options are internal to ESP-TLS, switching the libries will not change ESP-TLS specific code for a project.*

How to use wolfssl with ESP-IDF

There are two ways to use wolfssl in your project

1) Directly add wolfssl as a component in your project with following three commands.:

```
(First change directory (cd) to your project directory)
mkdir components
cd components
git clone https://github.com/espressif/esp-wolfssl.git
```

2) Add wolfssl as an extra component in your project.

- Download wolfssl with:

```
git clone https://github.com/espressif/esp-wolfssl.git
```

- Include esp-wolfssl in ESP-IDF with setting EXTRA_COMPONENT_DIRS in CMakeLists.txt/Makefile of your project as done in wolfssl/examples. For reference see Optional Project variables in *build-system*.

After above steps, you will have option to choose wolfssl as underlying SSL/TLS library in configuration menu of your project as follows:

```
idf.py/make menuconfig -> ESP-TLS -> choose SSL/TLS Library -> mbedtls/wolfssl
```

Comparison between mbedtls and wolfssl

The following table shows a typical comparison between wolfssl and mbedtls when protocols/https_request example *(which has server authentication)* was run with both SSL/TLS libraries and with all respective configurations set to default. *(mbedtls IN_CONTENT length and OUT_CONTENT length were set to 16384 bytes and 4096 bytes respectively)*

Property	Wolfssl	Mbedtls
Total Heap Consumed	~19 Kb	~37 Kb
Task Stack Used	~2.2 Kb	~3.6 Kb
Bin size	~858 Kb	~736 Kb

注解： *These values are subject to change with change in configuration options and version of respective libraries.*

ATECC608A (Secure Element) with ESP-TLS

ESP-TLS provides support for using ATECC608A cryptoauth chip with ESP32-WROOM-32SE. Use of ATECC608A is supported only when ESP-TLS is used with mbedTLS as its underlying SSL/TLS stack. ESP-TLS uses mbedtls as its underlying TLS/SSL stack by default unless changed manually.

注解： ATECC608A chip on ESP32-WROOM-32SE must be already configured and provisioned, for details refer esp_cryptoauth_utility

To enable the secure element support, and use it in you project for TLS connection, you will have to follow below steps

1) Add esp-cryptoauthlib in your project, for details please refer esp-cryptoauthlib with ESP_IDF
2) Enable following menuconfig option:

```
menuconfig->Component config->ESP-TLS->Use Secure Element (ATECC608A) with ESP-TLS
```

3) Select type of ATECC608A chip with following option:

```
menuconfig->Component config->esp-cryptoauthlib->Choose Type of ATECC608A chip
```

to know more about different types of ATECC608A chips and how to obtain type of ATECC608A connected to your ESP module please visit ATECC608A chip type

4) Enable use of ATECC608A in ESP-TLS by providing following config option in *esp_tls_cfg_t*

```
esp_tls_cfg_t cfg = {
    /* other configurations options */
    .use_secure_element = true,
};
```

API Reference

Header File

- esp-tls/esp_tls.h

Functions

esp_tls_t ***esp_tls_init**(void)

Create TLS connection.

This function allocates and initializes esp-tls structure handle.

Return tls Pointer to esp-tls as esp-tls handle if successfully initialized, NULL if allocation error

esp_tls_t ***esp_tls_conn_new**(**const** char **hostname*, int *hostlen*, int *port*, **const** *esp_tls_cfg_t* **cfg*)

Create a new blocking TLS/SSL connection.

This function establishes a TLS/SSL connection with the specified host in blocking manner.

Note: This API is present for backward compatibility reasons. Alternative function with the same functionality is `esp_tls_conn_new_sync` (and its asynchronous version `esp_tls_conn_new_async`)

Return pointer to esp_tls_t, or NULL if connection couldn't be opened.

Parameters

- `[in] hostname`: Hostname of the host.
- `[in] hostlen`: Length of hostname.
- `[in] port`: Port number of the host.
- `[in] cfg`: TLS configuration as esp_tls_cfg_t. If you wish to open non-TLS connection, keep this NULL. For TLS connection, a pass pointer to esp_tls_cfg_t. At a minimum, this structure should be zero-initialized.

int **esp_tls_conn_new_sync**(**const** char **hostname*, int *hostlen*, int *port*, **const** *esp_tls_cfg_t* **cfg*, *esp_tls_t* **tls*)

Create a new blocking TLS/SSL connection.

This function establishes a TLS/SSL connection with the specified host in blocking manner.

Return

- -1 If connection establishment fails.
- 1 If connection establishment is successful.
- 0 If connection state is in progress.

Parameters

- `[in] hostname`: Hostname of the host.
- `[in] hostlen`: Length of hostname.
- `[in] port`: Port number of the host.
- `[in] cfg`: TLS configuration as esp_tls_cfg_t. If you wish to open non-TLS connection, keep this NULL. For TLS connection, a pass pointer to esp_tls_cfg_t. At a minimum, this structure should be zero-initialized.
- `[in] tls`: Pointer to esp-tls as esp-tls handle.

esp_tls_t ***esp_tls_conn_http_new**(**const** char **url*, **const** *esp_tls_cfg_t* **cfg*)

Create a new blocking TLS/SSL connection with a given "HTTP" url.

The behaviour is same as esp_tls_conn_new() API. However this API accepts host's url.

Return pointer to esp_tls_t, or NULL if connection couldn't be opened.

Parameters

- `[in] url`: url of host.
- `[in] cfg`: TLS configuration as esp_tls_cfg_t. If you wish to open non-TLS connection, keep this NULL. For TLS connection, a pass pointer to 'esp_tls_cfg_t'. At a minimum, this structure should be zero-initialized.

int **esp_tls_conn_new_async**(**const** char **hostname*, int *hostlen*, int *port*, **const** *esp_tls_cfg_t* **cfg*, *esp_tls_t* **tls*)

Create a new non-blocking TLS/SSL connection.

This function initiates a non-blocking TLS/SSL connection with the specified host, but due to its non-blocking nature, it doesn't wait for the connection to get established.

Return

- -1 If connection establishment fails.
- 0 If connection establishment is in progress.
- 1 If connection establishment is successful.

Parameters

- `[in] hostname`: Hostname of the host.
- `[in] hostlen`: Length of hostname.
- `[in] port`: Port number of the host.
- `[in] cfg`: TLS configuration as esp_tls_cfg_t. `non_block` member of this structure should be set to be true.
- `[in] tls`: pointer to esp-tls as esp-tls handle.

int **esp_tls_conn_http_new_async**(**const** char **url*, **const** *esp_tls_cfg_t* **cfg*, *esp_tls_t* **tls*)

Create a new non-blocking TLS/SSL connection with a given "HTTP" url.

The behaviour is same as esp_tls_conn_new() API. However this API accepts host's url.

Return

- -1 If connection establishment fails.
- 0 If connection establishment is in progress.
- 1 If connection establishment is successful.

Parameters

- `[in] url`: url of host.
- `[in] cfg`: TLS configuration as esp_tls_cfg_t.
- `[in] tls`: pointer to esp-tls as esp-tls handle.

static ssize_t **esp_tls_conn_write**(*esp_tls_t* **tls*, **const** void **data*, size_t *datalen*)

Write from buffer 'data' into specified tls connection.

Return

- >=0 if write operation was successful, the return value is the number of bytes actually written to the TLS/SSL connection.

- <0 if write operation was not successful, because either an error occured or an action must be taken by the calling process.
- ESP_TLS_ERR_SSL_WANT_READ/ ESP_TLS_ERR_SSL_WANT_WRITE. if the handshake is incomplete and waiting for data to be available for reading. In this case this functions needs to be called again when the underlying transport is ready for operation.

Parameters

- `[in] tls`: pointer to esp-tls as esp-tls handle.
- `[in] data`: Buffer from which data will be written.
- `[in] datalen`: Length of data buffer.

static ssize_t **esp_tls_conn_read**(*esp_tls_t* *tls, void *data, size_t datalen)

Read from specified tls connection into the buffer ‘data’.

Return

- >0 if read operation was successful, the return value is the number of bytes actually read from the TLS/SSL connection.
- 0 if read operation was not successful. The underlying connection was closed.
- <0 if read operation was not successful, because either an error occured or an action must be taken by the calling process.

Parameters

- `[in] tls`: pointer to esp-tls as esp-tls handle.
- `[in] data`: Buffer to hold read data.
- `[in] datalen`: Length of data buffer.

void **esp_tls_conn_delete**(*esp_tls_t* *tls)

Compatible version of esp_tls_conn_destroy() to close the TLS/SSL connection.

Note This API will be removed in IDFv5.0

Parameters

- `[in] tls`: pointer to esp-tls as esp-tls handle.

int **esp_tls_conn_destroy**(*esp_tls_t* *tls)

Close the TLS/SSL connection and free any allocated resources.

This function should be called to close each tls connection opened with esp_tls_conn_new() or esp_tls_conn_http_new() APIs.

Return - 0 on success

- -1 if socket error or an invalid argument

Parameters

- `[in] tls`: pointer to esp-tls as esp-tls handle.

ssize_t **esp_tls_get_bytes_avail**(*esp_tls_t* *tls)

Return the number of application data bytes remaining to be read from the current record.

This API is a wrapper over mbedtls’s mbedtls_ssl_get_bytes_avail() API.

Return

- -1 in case of invalid arg
- bytes available in the application data record read buffer

Parameters

- `[in] tls`: pointer to esp-tls as esp-tls handle.

esp_err_t **esp_tls_get_conn_sockfd**(*esp_tls_t* *tls, int *sockfd)

Returns the connection socket file descriptor from *esp_tls* session.

Return - ESP_OK on success and value of sockfd will be updated with socket file descriptor for connection

- ESP_ERR_INVALID_ARG if (tls == NULL || sockfd == NULL)

Parameters

- `[in] tls`: handle to *esp_tls* context
- `[out] sockfd`: int pointer to sockfd value.

esp_err_t **esp_tls_init_global_ca_store**(void)

Create a global CA store, initially empty.

This function should be called if the application wants to use the same CA store for multiple connections. This function initialises the global CA store which can be then set by calling esp_tls_set_global_ca_store(). To be effective, this function must be called before any call to esp_tls_set_global_ca_store().

Return

- ESP_OK if creating global CA store was successful.
- ESP_ERR_NO_MEM if an error occured when allocating the mbedTLS resources.

esp_err_t **esp_tls_set_global_ca_store**(**const** unsigned char **cacert_pem_buf*, **const** unsigned int *cacert_pem_bytes*)

Set the global CA store with the buffer provided in pem format.

This function should be called if the application wants to set the global CA store for multiple connections i.e. to add the certificates in the provided buffer to the certificate chain. This function implicitly calls esp_tls_init_global_ca_store() if it has not already been called. The application must call this function before calling esp_tls_conn_new().

Return

- ESP_OK if adding certificates was successful.
- Other if an error occured or an action must be taken by the calling process.

Parameters

- `[in] cacert_pem_buf`: Buffer which has certificates in pem format. This buffer is used for creating a global CA store, which can be used by other tls connections.
- `[in] cacert_pem_bytes`: Length of the buffer.

void **esp_tls_free_global_ca_store**(void)

Free the global CA store currently being used.

The memory being used by the global CA store to store all the parsed certificates is freed up. The application can call this API if it no longer needs the global CA store.

esp_err_t **esp_tls_get_and_clear_last_error**(*esp_tls_error_handle_t h*, int **esp_tls_code*, int **esp_tls_flags*)

Returns last error in *esp_tls* with detailed mbedtls related error codes. The error information is cleared internally upon return.

Return

- ESP_ERR_INVALID_STATE if invalid parameters
- ESP_OK (0) if no error occurred
- specific error code (based on ESP_ERR_ESP_TLS_BASE) otherwise

Parameters

- `[in] h`: esp-tls error handle.
- `[out] esp_tls_code`: last error code returned from mbedtls api (set to zero if none) This pointer could be NULL if caller does not care about esp_tls_code
- `[out] esp_tls_flags`: last certification verification flags (set to zero if none) This pointer could be NULL if caller does not care about esp_tls_code

mbedtls_x509_crt ***esp_tls_get_global_ca_store**(void)

Get the pointer to the global CA store currently being used.

The application must first call esp_tls_set_global_ca_store(). Then the same CA store could be used by the application for APIs other than *esp_tls*.

Note Modifying the pointer might cause a failure in verifying the certificates.

Return

- Pointer to the global CA store currently being used if successful.
- NULL if there is no global CA store set.

Structures

struct esp_tls_last_error

Error structure containing relevant errors in case tls error occurred.

Public Members

esp_err_t **last_error**
error code (based on ESP_ERR_ESP_TLS_BASE) of the last occurred error

int **esp_tls_error_code**
esp_tls error code from last *esp_tls* failed api

int **esp_tls_flags**
last certification verification flags

struct psk_key_hint
ESP-TLS preshared key and hint structure.

Public Members

const uint8_t ***key**
key in PSK authentication mode in binary format

const size_t **key_size**
length of the key

const char ***hint**
hint in PSK authentication mode in string format

struct tls_keep_alive_cfg
Keep alive parameters structure.

Public Members

bool **keep_alive_enable**
Enable keep-alive timeout

int **keep_alive_idle**
Keep-alive idle time (second)

int **keep_alive_interval**
Keep-alive interval time (second)

int **keep_alive_count**
Keep-alive packet retry send count

struct esp_tls_cfg
ESP-TLS configuration parameters.

Note Note about format of certificates:
- This structure includes certificates of a Certificate Authority, of client or server as well as private keys, which may be of PEM or DER format. In case of PEM format, the buffer must be NULL terminated (with NULL character included in certificate size).
- Certificate Authority’ s certificate may be a chain of certificates in case of PEM format, but could be only one certificate in case of DER format
- Variables names of certificates and private key buffers and sizes are defined as unions providing backward compatibility for legacy *_pem_buf and *_pem_bytes names which suggested only PEM format was supported. It is encouraged to use generic names such as cacert_buf and cacert_bytes.

Public Members

const char ****alpn_protos**
Application protocols required for HTTP2. If HTTP2/ALPN support is required, a list of protocols that should be negotiated. The format is length followed by protocol name. For the most common cases the following is ok: const char **alpn_protos = { “h2”, NULL };

- where ‘h2’ is the protocol name

const unsigned char ***cacert_buf**
Certificate Authority’s certificate in a buffer. Format may be PEM or DER, depending on mbedtls-support This buffer should be NULL terminated in case of PEM

const unsigned char ***cacert_pem_buf**
CA certificate buffer legacy name

unsigned int **cacert_bytes**
Size of Certificate Authority certificate pointed to by cacert_buf (including NULL-terminator in case of PEM format)

unsigned int **cacert_pem_bytes**
Size of Certificate Authority certificate legacy name

const unsigned char ***clientcert_buf**
Client certificate in a buffer Format may be PEM or DER, depending on mbedtls-support This buffer should be NULL terminated in case of PEM

const unsigned char ***clientcert_pem_buf**
Client certificate legacy name

unsigned int **clientcert_bytes**
Size of client certificate pointed to by clientcert_pem_buf (including NULL-terminator in case of PEM format)

unsigned int **clientcert_pem_bytes**
Size of client certificate legacy name

const unsigned char ***clientkey_buf**
Client key in a buffer Format may be PEM or DER, depending on mbedtls-support This buffer should be NULL terminated in case of PEM

const unsigned char ***clientkey_pem_buf**
Client key legacy name

unsigned int **clientkey_bytes**
Size of client key pointed to by clientkey_pem_buf (including NULL-terminator in case of PEM format)

unsigned int **clientkey_pem_bytes**
Size of client key legacy name

const unsigned char ***clientkey_password**
Client key decryption password string

unsigned int **clientkey_password_len**
String length of the password pointed to by clientkey_password

bool **non_block**
Configure non-blocking mode. If set to true the underneath socket will be configured in non blocking mode after tls session is established

bool **use_secure_element**
Enable this option to use secure element or atecc608a chip (Integrated with ESP32-WROOM-32SE)

int **timeout_ms**
Network timeout in milliseconds

bool **use_global_ca_store**
Use a global ca_store for all the connections in which this bool is set.

const char ***common_name**
If non-NULL, server certificate CN must match this name. If NULL, server certificate CN must match hostname.

bool **skip_common_name**
Skip any validation of server certificate CN field

const *psk_hint_key_t* ***psk_hint_key**
Pointer to PSK hint and key. if not NULL (and certificates are NULL) then PSK authentication is enabled with configured setup. Important note: the pointer must be valid for connection

esp_err_t (***crt_bundle_attach**) (void *conf)
Function pointer to esp_crt_bundle_attach. Enables the use of certification bundle for server verification, must be enabled in menuconfig

tls_keep_alive_cfg_t ***keep_alive_cfg**
Enable TCP keep-alive timeout for SSL connection

struct esp_tls
ESP-TLS Connection Handle.

Public Members

mbedtls_ssl_context **ssl**
TLS/SSL context

mbedtls_entropy_context **entropy**
mbedTLS entropy context structure

mbedtls_ctr_drbg_context **ctr_drbg**
mbedTLS ctr drbg context structure. CTR_DRBG is deterministic random bit generation based on AES-256

mbedtls_ssl_config **conf**
TLS/SSL configuration to be shared between mbedtls_ssl_context structures

mbedtls_net_context **server_fd**
mbedTLS wrapper type for sockets

mbedtls_x509_crt **cacert**
Container for the X.509 CA certificate

mbedtls_x509_crt ***cacert_ptr**
Pointer to the cacert being used.

mbedtls_x509_crt **clientcert**
Container for the X.509 client certificate

mbedtls_pk_context **clientkey**
Container for the private key of the client certificate

int **sockfd**
Underlying socket file descriptor.

ssize_t (***read**) (**struct** *esp_tls* *tls, char *data, size_t datalen)
Callback function for reading data from TLS/SSL connection.

ssize_t (***write**) (**struct** *esp_tls* *tls, **const** char *data, size_t datalen)
Callback function for writing data to TLS/SSL connection.

esp_tls_conn_state_t **conn_state**
ESP-TLS Connection state

fd_set **rset**
read file descriptors

fd_set **wset**
write file descriptors

bool **is_tls**
indicates connection type (TLS or NON-TLS)

esp_tls_role_t **role**
esp-tls role

- ESP_TLS_CLIENT
- ESP_TLS_SERVER

esp_tls_error_handle_t **error_handle**
handle to error descriptor

Macros

ESP_ERR_ESP_TLS_BASE
Starting number of ESP-TLS error codes

ESP_ERR_ESP_TLS_CANNOT_RESOLVE_HOSTNAME
Error if hostname couldn't be resolved upon tls connection

ESP_ERR_ESP_TLS_CANNOT_CREATE_SOCKET
Failed to create socket

ESP_ERR_ESP_TLS_UNSUPPORTED_PROTOCOL_FAMILY
Unsupported protocol family

ESP_ERR_ESP_TLS_FAILED_CONNECT_TO_HOST
Failed to connect to host

ESP_ERR_ESP_TLS_SOCKET_SETOPT_FAILED
failed to set socket option

ESP_ERR_MBEDTLS_CERT_PARTLY_OK
mbedtls parse certificates was partly successful

ESP_ERR_MBEDTLS_CTR_DRBG_SEED_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_SET_HOSTNAME_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_CONFIG_DEFAULTS_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_CONF_ALPN_PROTOCOLS_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_X509_CRT_PARSE_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_CONF_OWN_CERT_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_SETUP_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_WRITE_FAILED
mbedtls api returned error

ESP_ERR_MBEDTLS_PK_PARSE_KEY_FAILED
mbedtls api returned failed

ESP_ERR_MBEDTLS_SSL_HANDSHAKE_FAILED
mbedtls api returned failed

ESP_ERR_MBEDTLS_SSL_CONF_PSK_FAILED
mbedtls api returned failed

ESP_ERR_ESP_TLS_CONNECTION_TIMEOUT
new connection in esp_tls_low_level_conn connection timeouted

ESP_ERR_WOLFSSL_SSL_SET_HOSTNAME_FAILED
wolfSSL api returned error

ESP_ERR_WOLFSSL_SSL_CONF_ALPN_PROTOCOLS_FAILED
wolfSSL api returned error

ESP_ERR_WOLFSSL_CERT_VERIFY_SETUP_FAILED
wolfSSL api returned error

ESP_ERR_WOLFSSL_KEY_VERIFY_SETUP_FAILED
wolfSSL api returned error

ESP_ERR_WOLFSSL_SSL_HANDSHAKE_FAILED
wolfSSL api returned failed

ESP_ERR_WOLFSSL_CTX_SETUP_FAILED
wolfSSL api returned failed

ESP_ERR_WOLFSSL_SSL_SETUP_FAILED
wolfSSL api returned failed

ESP_ERR_WOLFSSL_SSL_WRITE_FAILED
wolfSSL api returned failed

ESP_ERR_ESP_TLS_SE_FAILED

ESP_TLS_ERR_SSL_WANT_READ

ESP_TLS_ERR_SSL_WANT_WRITE

ESP_TLS_ERR_SSL_TIMEOUT

Type Definitions

typedef struct *esp_tls_last_error* ***esp_tls_error_handle_t**
typedef struct *esp_tls_last_error* **esp_tls_last_error_t**
Error structure containing relevant errors in case tls error occurred.

typedef enum *esp_tls_conn_state* **esp_tls_conn_state_t**
ESP-TLS Connection State.

typedef enum *esp_tls_role* **esp_tls_role_t**

typedef struct *psk_key_hint* **psk_hint_key_t**
ESP-TLS preshared key and hint structure.

typedef struct *tls_keep_alive_cfg* **tls_keep_alive_cfg_t**
Keep alive parameters structure.

typedef struct *esp_tls_cfg* **esp_tls_cfg_t**
ESP-TLS configuration parameters.

Note Note about format of certificates:

- This structure includes certificates of a Certificate Authority, of client or server as well as private keys, which may be of PEM or DER format. In case of PEM format, the buffer must be NULL terminated (with NULL character included in certificate size).
- Certificate Authority’ s certificate may be a chain of certificates in case of PEM format, but could be only one certificate in case of DER format
- Variables names of certificates and private key buffers and sizes are defined as unions providing backward compatibility for legacy *_pem_buf and *_pem_bytes names which suggested only PEM format was supported. It is encouraged to use generic names such as cacert_buf and cacert_bytes.

typedef struct *esp_tls* **esp_tls_t**
ESP-TLS Connection Handle.

Enumerations

enum esp_tls_conn_state
ESP-TLS Connection State.

Values:

ESP_TLS_INIT = 0

ESP_TLS_CONNECTING

ESP_TLS_HANDSHAKE

ESP_TLS_FAIL

ESP_TLS_DONE

enum esp_tls_role

Values:

ESP_TLS_CLIENT = 0

ESP_TLS_SERVER

2.3.3 ESP HTTP Client

Overview

`esp_http_client` provides an API for making HTTP/S requests from ESP-IDF programs. The steps to use this API for an HTTP request are:

- *`esp_http_client_init()`*: To use the HTTP client, the first thing we must do is create an `esp_http_client` by pass into this function with the *`esp_http_client_config_t`* configurations. Which configuration values we do not define, the library will use default.
- *`esp_http_client_perform()`*: The `esp_http_client` argument created from the init function is needed. This function performs all operations of the esp_http_client, from opening the connection, sending data, downloading data and closing the connection if necessary. All related events will be invoked in the event_handle (defined by *`esp_http_client_config_t`*). This function performs its job and blocks the current task until it's done
- *`esp_http_client_cleanup()`*: After completing our **esp_http_client's** task, this is the last function to be called. It will close the connection (if any) and free up all the memory allocated to the HTTP client

Application Example

```
esp_err_t _http_event_handle(esp_http_client_event_t *evt)
{
    switch(evt->event_id) {
        case HTTP_EVENT_ERROR:
            ESP_LOGI(TAG, "HTTP_EVENT_ERROR");
            break;
        case HTTP_EVENT_ON_CONNECTED:
            ESP_LOGI(TAG, "HTTP_EVENT_ON_CONNECTED");
            break;
        case HTTP_EVENT_HEADER_SENT:
            ESP_LOGI(TAG, "HTTP_EVENT_HEADER_SENT");
            break;
        case HTTP_EVENT_ON_HEADER:
            ESP_LOGI(TAG, "HTTP_EVENT_ON_HEADER");
            printf("%.*s", evt->data_len, (char*)evt->data);
            break;
        case HTTP_EVENT_ON_DATA:
            ESP_LOGI(TAG, "HTTP_EVENT_ON_DATA, len=%d", evt->data_len);
            if (!esp_http_client_is_chunked_response(evt->client)) {
                printf("%.*s", evt->data_len, (char*)evt->data);
            }

            break;
        case HTTP_EVENT_ON_FINISH:
            ESP_LOGI(TAG, "HTTP_EVENT_ON_FINISH");
```

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```
            break;
        case HTTP_EVENT_DISCONNECTED:
            ESP_LOGI(TAG, "HTTP_EVENT_DISCONNECTED");
            break;
    }
    return ESP_OK;
}

esp_http_client_config_t config = {
   .url = "http://httpbin.org/redirect/2",
   .event_handler = _http_event_handle,
};
esp_http_client_handle_t client = esp_http_client_init(&config);
esp_err_t err = esp_http_client_perform(client);

if (err == ESP_OK) {
   ESP_LOGI(TAG, "Status = %d, content_length = %d",
            esp_http_client_get_status_code(client),
            esp_http_client_get_content_length(client));
}
esp_http_client_cleanup(client);
```

Persistent Connections

Persistent connections means that the HTTP client can re-use the same connection for several transfers. If the server does not request to close the connection with the `Connection: close` header, the new transfer with sample ip address, port, and protocol.

To allow the HTTP client to take full advantage of persistent connections, you should do as many of your file transfers as possible using the same handle.

Persistent Connections example

```
esp_err_t err;
esp_http_client_config_t config = {
    .url = "http://httpbin.org/get",
};
esp_http_client_handle_t client = esp_http_client_init(&config);
// first request
err = esp_http_client_perform(client);

// second request
esp_http_client_set_url(client, "http://httpbin.org/anything")
esp_http_client_set_method(client, HTTP_METHOD_DELETE);
esp_http_client_set_header(client, "HeaderKey", "HeaderValue");
err = esp_http_client_perform(client);

esp_http_client_cleanup(client);
```

HTTPS

The HTTP client supports SSL connections using **mbedtls**, with the **url** configuration starting with `https` scheme (or `transport_type = HTTP_TRANSPORT_OVER_SSL`). HTTPS support can be configured via *CONFIG_ESP_HTTP_CLIENT_ENABLE_HTTPS* (enabled by default).

注解： By providing information using HTTPS, the library will use the SSL transport type to connect to the server. If you want to verify server, then need to provide additional certificate in PEM format, and provide to `cert_pem`

in esp_http_client_config_t

HTTPS example

```
static void https()
{
    esp_http_client_config_t config = {
        .url = "https://www.howsmyssl.com",
        .cert_pem = howsmyssl_com_root_cert_pem_start,
    };
    esp_http_client_handle_t client = esp_http_client_init(&config);
    esp_err_t err = esp_http_client_perform(client);

    if (err == ESP_OK) {
        ESP_LOGI(TAG, "Status = %d, content_length = %d",
                esp_http_client_get_status_code(client),
                esp_http_client_get_content_length(client));
    }
    esp_http_client_cleanup(client);
}
```

HTTP Stream

Some applications need to open the connection and control the reading of the data in an active manner. the HTTP client supports some functions to make this easier, of course, once you use these functions you should not use the *esp_http_client_perform()* function with that handle, and *esp_http_client_init()* alway to called first to get the handle. Perform that functions in the order below:

- *esp_http_client_init()*: to create and handle
- esp_http_client_set_* or esp_http_client_delete_*: to modify the http connection information (optional)
- *esp_http_client_open()*: Open the http connection with write_len parameter, write_len=0 if we only need read
- *esp_http_client_write()*: Upload data, max length equal to write_len of *esp_http_client_open()* function. We may not need to call it if write_len=0
- *esp_http_client_fetch_headers()*: After sending the headers and write data (if any) to the server, this function will read the HTTP Server response headers. Calling this function will return the content-length from the Server, and we can call *esp_http_client_get_status_code()* for the HTTP status of the connection.
- *esp_http_client_read()*: Now, we can read the HTTP stream by this function.
- *esp_http_client_close()*: We should the connection after finish
- *esp_http_client_cleanup()*: And release the resources

Perform HTTP request as Stream reader Check the example function http_perform_as_stream_reader at protocols/esp_http_client.

HTTP Authentication

The HTTP client supports both **Basic** and **Digest** Authentication. By providing usernames and passwords in url or in the username, password of config entry. And with auth_type = HTTP_AUTH_TYPE_BASIC, the HTTP client takes only 1 perform to pass the authentication process. If auth_type = HTTP_AUTH_TYPE_NONE, but there are username and password in the configuration, the HTTP client takes 2 performs. The first time it connects to the server and receives the UNAUTHORIZED header. Based on this information, it will know which authentication method to choose, and perform it on the second.

Config authentication example with URI

```
esp_http_client_config_t config = {
    .url = "http://user:passwd@httpbin.org/basic-auth/user/passwd",
    .auth_type = HTTP_AUTH_TYPE_BASIC,
};
```

Config authentication example with username, password entry

```
esp_http_client_config_t config = {
    .url = "http://httpbin.org/basic-auth/user/passwd",
    .username = "user",
    .password = "passwd",
    .auth_type = HTTP_AUTH_TYPE_BASIC,
};
```

HTTP Client example: protocols/esp_http_client.

API Reference

Header File

- esp_http_client/include/esp_http_client.h

Functions

esp_http_client_handle_t **esp_http_client_init**(**const** *esp_http_client_config_t* **config*)

Start a HTTP session This function must be the first function to call, and it returns a esp_http_client_handle_t that you must use as input to other functions in the interface. This call MUST have a corresponding call to esp_http_client_cleanup when the operation is complete.

Return

- `esp_http_client_handle_t`
- NULL if any errors

Parameters

- `[in] config`: The configurations, see `http_client_config_t`

esp_err_t **esp_http_client_perform**(*esp_http_client_handle_t client*)

Invoke this function after `esp_http_client_init` and all the options calls are made, and will perform the transfer as described in the options. It must be called with the same esp_http_client_handle_t as input as the esp_http_client_init call returned. esp_http_client_perform performs the entire request in either blocking or non-blocking manner. By default, the API performs request in a blocking manner and returns when done, or if it failed, and in non-blocking manner, it returns if EAGAIN/EWOULDBLOCK or EINPROGRESS is encountered, or if it failed. And in case of non-blocking request, the user may call this API multiple times unless request & response is complete or there is a failure. To enable non-blocking esp_http_client_perform(), `is_async` member of *esp_http_client_config_t* must be set while making a call to esp_http_client_init() API. You can do any amount of calls to esp_http_client_perform while using the same esp_http_client_handle_t. The underlying connection may be kept open if the server allows it. If you intend to transfer more than one file, you are even encouraged to do so. esp_http_client will then attempt to re-use the same connection for the following transfers, thus making the operations faster, less CPU intense and using less network resources. Just note that you will have to use `esp_http_client_set_**` between the invokes to set options for the following esp_http_client_perform.

Note You must never call this function simultaneously from two places using the same client handle. Let the function return first before invoking it another time. If you want parallel transfers, you must use several esp_http_client_handle_t. This function include `esp_http_client_open` -> `esp_http_client_write` -> `esp_http_client_fetch_headers` -> `esp_http_client_read` (and option) `esp_http_client_close`.

Return

- ESP_OK on successful
- ESP_FAIL on error

Parameters

- `client`: The esp_http_client handle

esp_err_t **esp_http_client_set_url**(*esp_http_client_handle_t client*, **const** char **url*)

Set URL for client, when performing this behavior, the options in the URL will replace the old ones.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle
- `[in] url`: The url

esp_err_t **esp_http_client_set_post_field**(*esp_http_client_handle_t client*, **const** char **data*, int *len*)

Set post data, this function must be called before `esp_http_client_perform`. Note: The data parameter passed to this function is a pointer and this function will not copy the data.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle
- `[in] data`: post data pointer
- `[in] len`: post length

int **esp_http_client_get_post_field**(*esp_http_client_handle_t client*, char ***data*)

Get current post field information.

Return Size of post data

Parameters

- `[in] client`: The client
- `[out] data`: Point to post data pointer

esp_err_t **esp_http_client_set_header**(*esp_http_client_handle_t client*, **const** char **key*, **const** char **value*)

Set http request header, this function must be called after esp_http_client_init and before any perform function.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle
- `[in] key`: The header key
- `[in] value`: The header value

esp_err_t **esp_http_client_get_header**(*esp_http_client_handle_t client*, **const** char **key*, char ***value*)

Get http request header. The value parameter will be set to NULL if there is no header which is same as the key specified, otherwise the address of header value will be assigned to value parameter. This function must be called after `esp_http_client_init`.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle
- `[in] key`: The header key
- `[out] value`: The header value

esp_err_t **esp_http_client_get_username**(*esp_http_client_handle_t client*, char ***value*)

Get http request username. The address of username buffer will be assigned to value parameter. This function must be called after `esp_http_client_init`.

Return

- ESP_OK
- ESP_ERR_INVALID_ARG

Parameters

- `[in] client`: The esp_http_client handle
- `[out] value`: The username value

esp_err_t **esp_http_client_set_username**(*esp_http_client_handle_t client*, **const** char **username*)

Set http request username. The value of username parameter will be assigned to username buffer. If the username parameter is NULL then username buffer will be freed.

Return

- ESP_OK
- ESP_ERR_INVALID_ARG

Parameters

- `[in] client`: The esp_http_client handle
- `[in] username`: The username value

esp_err_t **esp_http_client_get_password**(*esp_http_client_handle_t client*, char ***value*)

Get http request password. The address of password buffer will be assigned to value parameter. This function must be called after `esp_http_client_init`.

Return

- ESP_OK
- ESP_ERR_INVALID_ARG

Parameters

- `[in] client`: The esp_http_client handle
- `[out] value`: The password value

esp_err_t **esp_http_client_set_password**(*esp_http_client_handle_t client*, char **password*)

Set http request password. The value of password parameter will be assigned to password buffer. If the password parameter is NULL then password buffer will be freed.

Return

- ESP_OK
- ESP_ERR_INVALID_ARG

Parameters

- `[in] client`: The esp_http_client handle
- `[in] password`: The password value

esp_err_t **esp_http_client_set_authtype**(*esp_http_client_handle_t client*, *esp_http_client_auth_type_t auth_type*)

Set http request auth_type.

Return

- ESP_OK
- ESP_ERR_INVALID_ARG

Parameters

- `[in] client`: The esp_http_client handle
- `[in] auth_type`: The esp_http_client auth type

esp_err_t **esp_http_client_set_method**(*esp_http_client_handle_t client*, *esp_http_client_method_t method*)

Set http request method.

Return

- ESP_OK
- ESP_ERR_INVALID_ARG

Parameters

- `[in] client`: The esp_http_client handle
- `[in] method`: The method

esp_err_t **esp_http_client_delete_header**(*esp_http_client_handle_t client*, **const** char **key*)

Delete http request header.

Return
- ESP_OK
- ESP_FAIL

Parameters
- `[in] client`: The esp_http_client handle
- `[in] key`: The key

esp_err_t **esp_http_client_open**(*esp_http_client_handle_t client*, int *write_len*)
This function will be open the connection, write all header strings and return.

Return
- ESP_OK
- ESP_FAIL

Parameters
- `[in] client`: The esp_http_client handle
- `[in] write_len`: HTTP Content length need to write to the server

int **esp_http_client_write**(*esp_http_client_handle_t client*, **const** char ***buffer*, int *len*)
This function will write data to the HTTP connection previously opened by esp_http_client_open()

Return
- (-1) if any errors
- Length of data written

Parameters
- `[in] client`: The esp_http_client handle
- `buffer`: The buffer
- `[in] len`: This value must not be larger than the write_len parameter provided to esp_http_client_open()

int **esp_http_client_fetch_headers**(*esp_http_client_handle_t client*)
This function need to call after esp_http_client_open, it will read from http stream, process all receive headers.

Return
- (0) if stream doesn’ t contain content-length header, or chunked encoding (checked by `esp_http_client_is_chunked` response)
- (-1: ESP_FAIL) if any errors
- Download data length defined by content-length header

Parameters
- `[in] client`: The esp_http_client handle

bool **esp_http_client_is_chunked_response**(*esp_http_client_handle_t client*)
Check response data is chunked.

Return true or false

Parameters
- `[in] client`: The esp_http_client handle

int **esp_http_client_read**(*esp_http_client_handle_t client*, char ***buffer*, int *len*)
Read data from http stream.

Return
- (-1) if any errors
- Length of data was read

Parameters
- `[in] client`: The esp_http_client handle
- `buffer`: The buffer
- `[in] len`: The length

int **esp_http_client_get_status_code**(*esp_http_client_handle_t client*)
Get http response status code, the valid value if this function invoke after `esp_http_client_perform`

Return Status code

Parameters
- `[in] client`: The esp_http_client handle

int **esp_http_client_get_content_length**(*esp_http_client_handle_t client*)

Get http response content length (from header Content-Length) the valid value if this function invoke after `esp_http_client_perform`

Return

- (-1) Chunked transfer
- Content-Length value as bytes

Parameters

- `[in] client`: The esp_http_client handle

esp_err_t **esp_http_client_close**(*esp_http_client_handle_t client*)

Close http connection, still kept all http request resources.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle

esp_err_t **esp_http_client_cleanup**(*esp_http_client_handle_t client*)

This function must be the last function to call for an session. It is the opposite of the esp_http_client_init function and must be called with the same handle as input that a esp_http_client_init call returned. This might close all connections this handle has used and possibly has kept open until now. Don't call this function if you intend to transfer more files, re-using handles is a key to good performance with esp_http_client.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle

esp_http_client_transport_t **esp_http_client_get_transport_type**(*esp_http_client_handle_t client*)

Get transport type.

Return

- HTTP_TRANSPORT_UNKNOWN
- HTTP_TRANSPORT_OVER_TCP
- HTTP_TRANSPORT_OVER_SSL

Parameters

- `[in] client`: The esp_http_client handle

esp_err_t **esp_http_client_set_redirection**(*esp_http_client_handle_t client*)

Set redirection URL. When received the 30x code from the server, the client stores the redirect URL provided by the server. This function will set the current URL to redirect to enable client to execute the redirection request.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle

void **esp_http_client_add_auth**(*esp_http_client_handle_t client*)

On receiving HTTP Status code 401, this API can be invoked to add authorization information.

Note There is a possibility of receiving body message with redirection status codes, thus make sure to flush off body data after calling this API.

Parameters

- `[in] client`: The esp_http_client handle

bool **esp_http_client_is_complete_data_received**(*esp_http_client_handle_t client*)

Checks if entire data in the response has been read without any error.

Return

- true
- false

Parameters

- `[in] client`: The esp_http_client handle

int **esp_http_client_read_response**(*esp_http_client_handle_t client*, char ***buffer*, int *len*)

Helper API to read larger data chunks This is a helper API which internally calls `esp_http_client_read` multiple times till the end of data is reached or till the buffer gets full.

Return

- Length of data was read

Parameters

- `[in] client`: The esp_http_client handle
- `buffer`: The buffer
- `[in] len`: The buffer length

esp_err_t **esp_http_client_get_url**(*esp_http_client_handle_t client*, char ***url*, **const** int *len*)

Get URL from client.

Return

- ESP_OK
- ESP_FAIL

Parameters

- `[in] client`: The esp_http_client handle
- `[inout] url`: The buffer to store URL
- `[in] len`: The buffer length

Structures

struct esp_http_client_event

HTTP Client events data.

Public Members

esp_http_client_event_id_t **event_id**

event_id, to know the cause of the event

esp_http_client_handle_t **client**

esp_http_client_handle_t context

void ***data**

data of the event

int **data_len**

data length of data

void ***user_data**

user_data context, from *esp_http_client_config_t* user_data

char ***header_key**

For HTTP_EVENT_ON_HEADER event_id, it's store current http header key

char ***header_value**

For HTTP_EVENT_ON_HEADER event_id, it's store current http header value

struct esp_http_client_config_t

HTTP configuration.

Public Members

const char ***url**

HTTP URL, the information on the URL is most important, it overrides the other fields below, if any

const char ***host**
Domain or IP as string

int **port**
Port to connect, default depend on esp_http_client_transport_t (80 or 443)

const char ***username**
Using for Http authentication

const char ***password**
Using for Http authentication

esp_http_client_auth_type_t **auth_type**
Http authentication type, see `esp_http_client_auth_type_t`

const char ***path**
HTTP Path, if not set, default is `/`

const char ***query**
HTTP query

const char ***cert_pem**
SSL server certification, PEM format as string, if the client requires to verify server

const char ***client_cert_pem**
SSL client certification, PEM format as string, if the server requires to verify client

const char ***client_key_pem**
SSL client key, PEM format as string, if the server requires to verify client

const char ***user_agent**
The User Agent string to send with HTTP requests

esp_http_client_method_t **method**
HTTP Method

int **timeout_ms**
Network timeout in milliseconds

bool **disable_auto_redirect**
Disable HTTP automatic redirects

int **max_redirection_count**
Max redirection number, using default value if zero

http_event_handle_cb **event_handler**
HTTP Event Handle

esp_http_client_transport_t **transport_type**
HTTP transport type, see `esp_http_client_transport_t`

int **buffer_size**
HTTP receive buffer size

int **buffer_size_tx**
HTTP transmit buffer size

void ***user_data**
HTTP user_data context

bool **is_async**
Set asynchronous mode, only supported with HTTPS for now

bool **use_global_ca_store**
Use a global ca_store for all the connections in which this bool is set.

bool **skip_cert_common_name_check**
Skip any validation of server certificate CN field

bool **keep_alive_enable**
Enable keep-alive timeout

int **keep_alive_idle**
Keep-alive idle time. Default is 5 (second)

int **keep_alive_interval**
Keep-alive interval time. Default is 5 (second)

int **keep_alive_count**
Keep-alive packet retry send count. Default is 3 counts

Macros

DEFAULT_HTTP_BUF_SIZE

ESP_ERR_HTTP_BASE
Starting number of HTTP error codes

ESP_ERR_HTTP_MAX_REDIRECT
The error exceeds the number of HTTP redirects

ESP_ERR_HTTP_CONNECT
Error open the HTTP connection

ESP_ERR_HTTP_WRITE_DATA
Error write HTTP data

ESP_ERR_HTTP_FETCH_HEADER
Error read HTTP header from server

ESP_ERR_HTTP_INVALID_TRANSPORT
There are no transport support for the input scheme

ESP_ERR_HTTP_CONNECTING
HTTP connection hasn’t been established yet

ESP_ERR_HTTP_EAGAIN
Mapping of errno EAGAIN to esp_err_t

Type Definitions

typedef struct esp_http_client ***esp_http_client_handle_t**

typedef struct *esp_http_client_event* ***esp_http_client_event_handle_t**

typedef struct *esp_http_client_event* **esp_http_client_event_t**
HTTP Client events data.

typedef *esp_err_t* (***http_event_handle_cb**) (*esp_http_client_event_t* *evt)

Enumerations

enum esp_http_client_event_id_t
HTTP Client events id.

Values:

HTTP_EVENT_ERROR = 0
This event occurs when there are any errors during execution

HTTP_EVENT_ON_CONNECTED
Once the HTTP has been connected to the server, no data exchange has been performed

HTTP_EVENT_HEADERS_SENT
After sending all the headers to the server

HTTP_EVENT_HEADER_SENT = *HTTP_EVENT_HEADERS_SENT*
This header has been kept for backward compatability and will be deprecated in future versions esp-idf

HTTP_EVENT_ON_HEADER
Occurs when receiving each header sent from the server

HTTP_EVENT_ON_DATA
Occurs when receiving data from the server, possibly multiple portions of the packet

HTTP_EVENT_ON_FINISH
Occurs when finish a HTTP session

HTTP_EVENT_DISCONNECTED
The connection has been disconnected

enum esp_http_client_transport_t
HTTP Client transport.

Values:

HTTP_TRANSPORT_UNKNOWN = 0x0
Unknown

HTTP_TRANSPORT_OVER_TCP
Transport over tcp

HTTP_TRANSPORT_OVER_SSL
Transport over ssl

enum esp_http_client_method_t
HTTP method.

Values:

HTTP_METHOD_GET = 0
HTTP GET Method

HTTP_METHOD_POST
HTTP POST Method

HTTP_METHOD_PUT
HTTP PUT Method

HTTP_METHOD_PATCH
HTTP PATCH Method

HTTP_METHOD_DELETE
HTTP DELETE Method

HTTP_METHOD_HEAD
HTTP HEAD Method

HTTP_METHOD_NOTIFY
HTTP NOTIFY Method

HTTP_METHOD_SUBSCRIBE
HTTP SUBSCRIBE Method

HTTP_METHOD_UNSUBSCRIBE
HTTP UNSUBSCRIBE Method

HTTP_METHOD_OPTIONS
HTTP OPTIONS Method

HTTP_METHOD_MAX

enum esp_http_client_auth_type_t
HTTP Authentication type.

Values:

HTTP_AUTH_TYPE_NONE = 0
No authention

HTTP_AUTH_TYPE_BASIC
HTTP Basic authentication

HTTP_AUTH_TYPE_DIGEST
HTTP Disgest authentication

enum HttpStatus_Code
Enum for the HTTP status codes.

Values:

HttpStatus_Ok = 200

HttpStatus_MultipleChoices = 300

HttpStatus_MovedPermanently = 301

HttpStatus_Found = 302

HttpStatus_TemporaryRedirect = 307

HttpStatus_Unauthorized = 401

2.3.4 ESP WebSocket Client

Overview

The ESP WebSocket client is an implementation of WebSocket protocol client for ESP32

Features

- supports WebSocket over TCP, SSL with mbedtls
- Easy to setup with URI
- Multiple instances (Multiple clients in one application)

Configuration

URI

- Supports `ws`, `wss` schemes
- WebSocket samples:
  - `ws://websocket.org`: WebSocket over TCP, default port 80
  - `wss://websocket.org`: WebSocket over SSL, default port 443
- Minimal configurations:

```
const esp_websocket_client_config_t ws_cfg = {
    .uri = "ws://websocket.org",
};
```

- If there are any options related to the URI in `esp_websocket_client_config_t`, the option defined by the URI will be overridden. Sample:

```
const esp_websocket_client_config_t ws_cfg = {
    .uri = "ws://websocket.org:123",
    .port = 4567,
};
//WebSocket client will connect to websocket.org using port 4567
```

SSL

- Get certificate from server, example: `websocket.org openssl s_client -showcerts -connect websocket.org:443 </dev/null 2>/dev/null|openssl x509 -outform PEM >websocket_org.pem`
- Configuration:

```
const esp_websocket_client_config_t ws_cfg = {
    .uri = "wss://websocket.org",
    .cert_pem = (const char *)websocket_org_pem_start,
};
```

For more options on `esp_websocket_client_config_t`, please refer to API reference below

Application Example

Simple WebSocket example that uses esp_websocket_client to establish a websocket connection and send/receive data with the websocket.org Server: protocols/websocket.

API Reference

Header File

- esp_websocket_client/include/esp_websocket_client.h

Functions

esp_websocket_client_handle_t **esp_websocket_client_init**(**const** *esp_websocket_client_config_t* **config*)

Start a Websocket session This function must be the first function to call, and it returns a esp_websocket_client_handle_t that you must use as input to other functions in the interface. This call MUST have a corresponding call to esp_websocket_client_destroy when the operation is complete.

Return

- `esp_websocket_client_handle_t`
- NULL if any errors

Parameters

- `[in] config`: The configuration

esp_err_t **esp_websocket_client_set_uri**(*esp_websocket_client_handle_t client*, **const** char **uri*)

Set URL for client, when performing this behavior, the options in the URL will replace the old ones Must stop the WebSocket client before set URI if the client has been connected.

Return esp_err_t

Parameters

- `[in] client`: The client
- `[in] uri`: The uri

esp_err_t **esp_websocket_client_start**(*esp_websocket_client_handle_t client*)

Open the WebSocket connection.

Return esp_err_t

Parameters

- `[in] client`: The client

esp_err_t **esp_websocket_client_stop**(*esp_websocket_client_handle_t client*)

Close the WebSocket connection.

Notes:

- Cannot be called from the websocket event handler

Return esp_err_t

Parameters

- `[in] client`: The client

esp_err_t **esp_websocket_client_destroy** (*esp_websocket_client_handle_t client*)

Destroy the WebSocket connection and free all resources. This function must be the last function to call for an session. It is the opposite of the esp_websocket_client_init function and must be called with the same handle as input that a esp_websocket_client_init call returned. This might close all connections this handle has used.

Notes:

- Cannot be called from the websocket event handler

Return esp_err_t

Parameters

- `[in] client`: The client

int **esp_websocket_client_send** (*esp_websocket_client_handle_t client*, **const** char **data*, int *len*, TickType_t *timeout*)

Generic write data to the WebSocket connection; defaults to binary send.

Return

- Number of data was sent
- (-1) if any errors

Parameters

- `[in] client`: The client
- `[in] data`: The data
- `[in] len`: The length
- `[in] timeout`: Write data timeout in RTOS ticks

int **esp_websocket_client_send_bin** (*esp_websocket_client_handle_t client*, **const** char **data*, int *len*, TickType_t *timeout*)

Write binary data to the WebSocket connection (data send with WS OPCODE=02, i.e. binary)

Return

- Number of data was sent
- (-1) if any errors

Parameters

- `[in] client`: The client
- `[in] data`: The data
- `[in] len`: The length
- `[in] timeout`: Write data timeout in RTOS ticks

int **esp_websocket_client_send_text** (*esp_websocket_client_handle_t client*, **const** char **data*, int *len*, TickType_t *timeout*)

Write textual data to the WebSocket connection (data send with WS OPCODE=01, i.e. text)

Return

- Number of data was sent
- (-1) if any errors

Parameters

- `[in] client`: The client
- `[in] data`: The data
- `[in] len`: The length
- `[in] timeout`: Write data timeout in RTOS ticks

bool **esp_websocket_client_is_connected** (*esp_websocket_client_handle_t client*)

Check the WebSocket connection status.

Return

- true
- false

Parameters

- `[in] client`: The client handle

esp_err_t **esp_websocket_register_events**(*esp_websocket_client_handle_t client*, *esp_websocket_event_id_t event*, *esp_event_handler_t event_handler*, void **event_handler_arg*)

Register the Websocket Events.

Return esp_err_t

Parameters

- `client`: The client handle
- `event`: The event id
- `event_handler`: The callback function
- `event_handler_arg`: User context

Structures

struct esp_websocket_event_data_t

Websocket event data.

Public Members

const char ***data_ptr**

Data pointer

int **data_len**

Data length

uint8_t **op_code**

Received opcode

esp_websocket_client_handle_t **client**

esp_websocket_client_handle_t context

void ***user_context**

user_data context, from *esp_websocket_client_config_t* user_data

int **payload_len**

Total payload length, payloads exceeding buffer will be posted through multiple events

int **payload_offset**

Actual offset for the data associated with this event

struct esp_websocket_client_config_t

Websocket client setup configuration.

Public Members

const char ***uri**

Websocket URI, the information on the URI can be overrides the other fields below, if any

const char ***host**

Domain or IP as string

int **port**

Port to connect, default depend on esp_websocket_transport_t (80 or 443)

const char ***username**

Using for Http authentication - Not supported for now

const char ***password**

Using for Http authentication - Not supported for now

const char ***path**

HTTP Path, if not set, default is /

bool **disable_auto_reconnect**

Disable the automatic reconnect function when disconnected

void ***user_context**
HTTP user data context

int **task_prio**
Websocket task priority

int **task_stack**
Websocket task stack

int **buffer_size**
Websocket buffer size

const char ***cert_pem**
SSL Certification, PEM format as string, if the client requires to verify server

esp_websocket_transport_t **transport**
Websocket transport type, see 'esp_websocket_transport_t

char ***subprotocol**
Websocket subprotocol

char ***user_agent**
Websocket user-agent

char ***headers**
Websocket additional headers

bool **keep_alive_enable**
Enable keep-alive timeout

int **keep_alive_idle**
Keep-alive idle time. Default is 5 (second)

int **keep_alive_interval**
Keep-alive interval time. Default is 5 (second)

int **keep_alive_count**
Keep-alive packet retry send count. Default is 3 counts

Type Definitions

typedef struct esp_websocket_client ***esp_websocket_client_handle_t**

Enumerations

enum esp_websocket_event_id_t
Websocket Client events id.

Values:

WEBSOCKET_EVENT_ANY = -1

WEBSOCKET_EVENT_ERROR = 0
This event occurs when there are any errors during execution

WEBSOCKET_EVENT_CONNECTED
Once the Websocket has been connected to the server, no data exchange has been performed

WEBSOCKET_EVENT_DISCONNECTED
The connection has been disconnected

WEBSOCKET_EVENT_DATA
When receiving data from the server, possibly multiple portions of the packet

WEBSOCKET_EVENT_MAX

enum esp_websocket_transport_t
Websocket Client transport.

Values:

WEBSOCKET_TRANSPORT_UNKNOWN = 0x0
Transport unknown

WEBSOCKET_TRANSPORT_OVER_TCP
Transport over tcp

WEBSOCKET_TRANSPORT_OVER_SSL
Transport over ssl

2.3.5 HTTP 服务器

概述

HTTP Server 组件提供了在 ESP32 上运行轻量级 Web 服务器的功能，下面介绍使用 HTTP Server 组件 API 的详细步骤：

- *httpd_start()*：创建 HTTP 服务器的实例，根据具体的配置为其分配内存和资源，并返回该服务器实例的句柄。服务器使用了两个套接字，一个用来监听 HTTP 流量（TCP 类型），另一个用来处理控制信号（UDP 类型），它们在服务器的任务循环中轮流使用。通过向 httpd_start() 传递 httpd_config_t 结构体，可以在创建服务器实例时配置任务的优先级和堆栈的大小。TCP 流量被解析为 HTTP 请求，根据请求的 URI 来调用用户注册的处理程序，在处理程序中需要发送回 HTTP 响应数据包。
- *httpd_stop()*：根据传入的句柄停止服务器，并释放相关联的内存和资源。这是一个阻塞函数，首先给服务器任务发送停止信号，然后等待其终止。期间服务器任务会关闭所有已打开的连接，删除已注册的 URI 处理程序，并将所有会话的上下文数据重置为空。
- *httpd_register_uri_handler()*：通过传入 `httpd_uri_t` 结构体类型的对象来注册 URI 处理程序。该结构体包含如下成员：`uri` 名字，`method` 类型（比如 `HTTPD_GET/HTTPD_POST/HTTPD_PUT` 等等），`esp_err_t *handler (httpd_req_t *req)` 类型的函数指针，指向用户上下文数据的 `user_ctx` 指针。

应用示例

```
/* URI 处理函数，在客户端发起 GET /uri 请求时被调用 */
esp_err_t get_handler(httpd_req_t *req)
{
    /* 发送回简单的响应数据包 */
    const char[] resp = "URI GET Response";
    httpd_resp_send(req, resp, strlen(resp));
    return ESP_OK;
}

/* URI 处理函数，在客户端发起 POST /uri 请求时被调用 */
esp_err_t post_handler(httpd_req_t *req)
{
    /* 定义 HTTP POST 请求数据的目标缓存区
     * httpd_req_recv() 只接收 char* 数据，但也可以是任意二进制数据（需要类型转换）
     * 对于字符串数据，null 终止符会被省略，content_len 会给出字符串的长度 */
    char content[100];

    /* 如果内容长度大于缓冲区则截断 */
    size_t recv_size = MIN(req->content_len, sizeof(content));

    int ret = httpd_req_recv(req, content, recv_size);
    if (ret <= 0) {  /* 返回 0 表示连接已关闭 */
        /* 检查是否超时 */
        if (ret == HTTPD_SOCK_ERR_TIMEOUT) {
            /* 如果是超时，可以调用 httpd_req_recv() 重试
             * 简单起见，这里我们直接响应 HTTP 408（请求超时）错误给客户端 */
            httpd_resp_send_408(req);
```

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```
        }
        /* 如果发生了错误，返回 ESP_FAIL 可以确保底层套接字被关闭 */
        return ESP_FAIL;
    }

    /* 发送简单的响应数据包 */
    const char[] resp = "URI POST Response";
    httpd_resp_send(req, resp, strlen(resp));
    return ESP_OK;
}

/* GET /uri 的 URI 处理结构 */
httpd_uri_t uri_get = {
    .uri      = "/uri",
    .method   = HTTP_GET,
    .handler  = get_handler,
    .user_ctx = NULL
};

/* POST /uri 的 URI 处理结构 */
httpd_uri_t uri_post = {
    .uri      = "/uri",
    .method   = HTTP_POST,
    .handler  = post_handler,
    .user_ctx = NULL
};

/* 启动 Web 服务器的函数 */
httpd_handle_t start_webserver(void)
{
    /* 生成默认的配置参数 */
    httpd_config_t config = HTTPD_DEFAULT_CONFIG();

    /* 置空 esp_http_server 的实例句柄 */
    httpd_handle_t server = NULL;

    /* 启动 httpd server */
    if (httpd_start(&server, &config) == ESP_OK) {
        /* 注册 URI 处理程序 */
        httpd_register_uri_handler(server, &uri_get);
        httpd_register_uri_handler(server, &uri_post);
    }
    /* 如果服务器启动失败，返回的句柄是 NULL */
    return server;
}

/* 停止 Web 服务器的函数 */
void stop_webserver(httpd_handle_t server)
{
    if (server) {
        /* 停止 httpd server */
        httpd_stop(server);
    }
}
```

简单 HTTP 服务器示例 请查看位于 protocols/http_server/simple 的 HTTP 服务器示例，该示例演示了如何处理任意内容长度的数据，读取请求头和 URL 查询参数，设置响应头。

HTTP 长连接

HTTP 服务器具有长连接的功能，允许重复使用同一个连接（会话）进行多次传输，同时保持会话的上下文数据。上下文数据可由处理程序动态分配，在这种情况下需要提前指定好自定义的回调函数，以便在连接/会话被关闭时释放这部分内存资源。

长连接示例

```
/* 自定义函数，用来释放上下文数据 */
void free_ctx_func(void *ctx)
{
    /* 也可以是 free 以外的代码逻辑 */
    free(ctx);
}

esp_err_t adder_post_handler(httpd_req_t *req)
{
    /* 如果会话上下文还不存在则新建一个 */
    if (! req->sess_ctx) {
        req->sess_ctx = malloc(sizeof(ANY_DATA_TYPE));   /*!< 指向上下文数据 */
        req->free_ctx = free_ctx_func;                   /*!< 释放上下文数据的函数 */
    }

    /* 访问上下文数据 */
    ANY_DATA_TYPE *ctx_data = (ANY_DATA_TYPE *)req->sess_ctx;

    /* 响应 */
    ................
    ................
    ................

    return ESP_OK;
}
```

详情请参考位于 protocols/http_server/persistent_sockets 的示例代码。

API 参考

Header File

- esp_http_server/include/esp_http_server.h

Functions

esp_err_t **httpd_register_uri_handler**(*httpd_handle_t handle*, **const** *httpd_uri_t *uri_handler*)

Registers a URI handler.

Example usage:

```
esp_err_t my_uri_handler(httpd_req_t* req)
{
    // Recv , Process and Send
    ....
    ....
    ....

    // Fail condition
    if (....) {
        // Return fail to close session //
        return ESP_FAIL;
    }
```

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```
    // On success
    return ESP_OK;
}

// URI handler structure
httpd_uri_t my_uri {
    .uri      = "/my_uri/path/xyz",
    .method   = HTTPD_GET,
    .handler  = my_uri_handler,
    .user_ctx = NULL
};

// Register handler
if (httpd_register_uri_handler(server_handle, &my_uri) != ESP_OK) {
   // If failed to register handler
   ....
}
```

Note URI handlers can be registered in real time as long as the server handle is valid.

Return

- ESP_OK : On successfully registering the handler
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_HANDLERS_FULL : If no slots left for new handler
- ESP_ERR_HTTPD_HANDLER_EXISTS : If handler with same URI and method is already registered

Parameters

- `[in] handle`: handle to HTTPD server instance
- `[in] uri_handler`: pointer to handler that needs to be registered

esp_err_t **httpd_unregister_uri_handler**(*httpd_handle_t handle*, **const** char **uri*, *httpd_method_t method*)

Unregister a URI handler.

Return

- ESP_OK : On successfully deregistering the handler
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_NOT_FOUND : Handler with specified URI and method not found

Parameters

- `[in] handle`: handle to HTTPD server instance
- `[in] uri`: URI string
- `[in] method`: HTTP method

esp_err_t **httpd_unregister_uri**(*httpd_handle_t handle*, **const** char **uri*)

Unregister all URI handlers with the specified uri string.

Return

- ESP_OK : On successfully deregistering all such handlers
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_NOT_FOUND : No handler registered with specified uri string

Parameters

- `[in] handle`: handle to HTTPD server instance
- `[in] uri`: uri string specifying all handlers that need to be deregisterd

esp_err_t **httpd_sess_set_recv_override**(*httpd_handle_t hd*, int *sockfd*, *httpd_recv_func_t recv_func*)

Override web server’s receive function (by session FD)

This function overrides the web server’s receive function. This same function is used to read HTTP request packets.

Note This API is supposed to be called either from the context of

- an http session APIs where sockfd is a valid parameter
- a URI handler where sockfd is obtained using httpd_req_to_sockfd()

Return

- ESP_OK : On successfully registering override
- ESP_ERR_INVALID_ARG : Null arguments

Parameters

- `[in] hd`: HTTPD instance handle
- `[in] sockfd`: Session socket FD
- `[in] recv_func`: The receive function to be set for this session

esp_err_t **httpd_sess_set_send_override** (*httpd_handle_t hd*, int *sockfd*, *httpd_send_func_t send_func*)

Override web server's send function (by session FD)

This function overrides the web server's send function. This same function is used to send out any response to any HTTP request.

Note This API is supposed to be called either from the context of

- an http session APIs where sockfd is a valid parameter
- a URI handler where sockfd is obtained using httpd_req_to_sockfd()

Return

- ESP_OK : On successfully registering override
- ESP_ERR_INVALID_ARG : Null arguments

Parameters

- `[in] hd`: HTTPD instance handle
- `[in] sockfd`: Session socket FD
- `[in] send_func`: The send function to be set for this session

esp_err_t **httpd_sess_set_pending_override** (*httpd_handle_t hd*, int *sockfd*, *httpd_pending_func_t pending_func*)

Override web server's pending function (by session FD)

This function overrides the web server's pending function. This function is used to test for pending bytes in a socket.

Note This API is supposed to be called either from the context of

- an http session APIs where sockfd is a valid parameter
- a URI handler where sockfd is obtained using httpd_req_to_sockfd()

Return

- ESP_OK : On successfully registering override
- ESP_ERR_INVALID_ARG : Null arguments

Parameters

- `[in] hd`: HTTPD instance handle
- `[in] sockfd`: Session socket FD
- `[in] pending_func`: The receive function to be set for this session

int **httpd_req_to_sockfd** (*httpd_req_t* *r)

Get the Socket Descriptor from the HTTP request.

This API will return the socket descriptor of the session for which URI handler was executed on reception of HTTP request. This is useful when user wants to call functions that require session socket fd, from within a URI handler, ie. : httpd_sess_get_ctx(), httpd_sess_trigger_close(), httpd_sess_update_lru_counter().

Note This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.

Return

- Socket descriptor : The socket descriptor for this request
- -1 : Invalid/NULL request pointer

Parameters

- `[in] r`: The request whose socket descriptor should be found

int **httpd_req_recv** (*httpd_req_t* *r, char *buf, size_t *buf_len*)

API to read content data from the HTTP request.

This API will read HTTP content data from the HTTP request into provided buffer. Use content_len provided in httpd_req_t structure to know the length of data to be fetched. If content_len is too large for the buffer then user may have to make multiple calls to this function, each time fetching 'buf_len' number of bytes, while the pointer to content data is incremented internally by the same number.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- If an error is returned, the URI handler must further return an error. This will ensure that the erroneous socket is closed and cleaned up by the web server.
- Presently Chunked Encoding is not supported

Return

- Bytes : Number of bytes read into the buffer successfully
- 0 : Buffer length parameter is zero / connection closed by peer
- HTTPD_SOCK_ERR_INVALID : Invalid arguments
- HTTPD_SOCK_ERR_TIMEOUT : Timeout/interrupted while calling socket recv()
- HTTPD_SOCK_ERR_FAIL : Unrecoverable error while calling socket recv()

Parameters

- `[in] r`: The request being responded to
- `[in] buf`: Pointer to a buffer that the data will be read into
- `[in] buf_len`: Length of the buffer

size_t **httpd_req_get_hdr_value_len**(*httpd_req_t* *r, **const** char *field)

Search for a field in request headers and return the string length of it's value.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once httpd_resp_send() API is called all request headers are purged, so request headers need be copied into separate buffers if they are required later.

Return

- Length : If field is found in the request URL
- Zero : Field not found / Invalid request / Null arguments

Parameters

- `[in] r`: The request being responded to
- `[in] field`: The header field to be searched in the request

esp_err_t **httpd_req_get_hdr_value_str**(*httpd_req_t* *r, **const** char *field, char *val, size_t *val_size*)

Get the value string of a field from the request headers.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once httpd_resp_send() API is called all request headers are purged, so request headers need be copied into separate buffers if they are required later.
- If output size is greater than input, then the value is truncated, accompanied by truncation error as return value.
- Use httpd_req_get_hdr_value_len() to know the right buffer length

Return

- ESP_OK : Field found in the request header and value string copied
- ESP_ERR_NOT_FOUND : Key not found
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_INVALID_REQ : Invalid HTTP request pointer
- ESP_ERR_HTTPD_RESULT_TRUNC : Value string truncated

Parameters

- `[in] r`: The request being responded to
- `[in] field`: The field to be searched in the header
- `[out] val`: Pointer to the buffer into which the value will be copied if the field is found
- `[in] val_size`: Size of the user buffer "val"

size_t **httpd_req_get_url_query_len**(*httpd_req_t* *r)

Get Query string length from the request URL.

Note This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid

Return

- Length : Query is found in the request URL
- Zero : Query not found / Null arguments / Invalid request

Parameters

- `[in] r`: The request being responded to

esp_err_t **httpd_req_get_url_query_str**(*httpd_req_t* *r, char *buf, size_t buf_len)

Get Query string from the request URL.

Note

- Presently, the user can fetch the full URL query string, but decoding will have to be performed by the user. Request headers can be read using httpd_req_get_hdr_value_str() to know the 'Content-Type' (eg. Content-Type: application/x-www-form-urlencoded) and then the appropriate decoding algorithm needs to be applied.
- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid
- If output size is greater than input, then the value is truncated, accompanied by truncation error as return value
- Prior to calling this function, one can use httpd_req_get_url_query_len() to know the query string length beforehand and hence allocate the buffer of right size (usually query string length + 1 for null termination) for storing the query string

Return

- ESP_OK : Query is found in the request URL and copied to buffer
- ESP_ERR_NOT_FOUND : Query not found
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_INVALID_REQ : Invalid HTTP request pointer
- ESP_ERR_HTTPD_RESULT_TRUNC : Query string truncated

Parameters

- `[in] r`: The request being responded to
- `[out] buf`: Pointer to the buffer into which the query string will be copied (if found)
- `[in] buf_len`: Length of output buffer

esp_err_t **httpd_query_key_value**(**const** char *qry, **const** char *key, char *val, size_t val_size)

Helper function to get a URL query tag from a query string of the type param1=val1&param2=val2.

Note

- The components of URL query string (keys and values) are not URLdecoded. The user must check for 'Content-Type' field in the request headers and then depending upon the specified encoding (URLencoded or otherwise) apply the appropriate decoding algorithm.
- If actual value size is greater than val_size, then the value is truncated, accompanied by truncation error as return value.

Return

- ESP_OK : Key is found in the URL query string and copied to buffer
- ESP_ERR_NOT_FOUND : Key not found
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_RESULT_TRUNC : Value string truncated

Parameters

- `[in] qry`: Pointer to query string
- `[in] key`: The key to be searched in the query string
- `[out] val`: Pointer to the buffer into which the value will be copied if the key is found
- `[in] val_size`: Size of the user buffer "val"

bool **httpd_uri_match_wildcard**(**const** char *uri_template, **const** char *uri_to_match, size_t match_upto)

Test if a URI matches the given wildcard template.

Template may end with "?" to make the previous character optional (typically a slash), "*" for a wildcard

match, and “?*” to make the previous character optional, and if present, allow anything to follow.

Example:

- * matches everything
- /foo/? matches /foo and /foo/
- /foo/* (sans the backslash) matches /foo/ and /foo/bar, but not /foo or /fo
- /foo/?* or /foo/*? (sans the backslash) matches /foo/, /foo/bar, and also /foo, but not /foox or /fo

The special characters “?” and “*” anywhere else in the template will be taken literally.

Return true if a match was found

Parameters

- `[in] uri_template`: URI template (pattern)
- `[in] uri_to_match`: URI to be matched
- `[in] match_upto`: how many characters of the URI buffer to test (there may be trailing query string etc.)

esp_err_t **httpd_resp_send** (*httpd_req_t* *r, **const** char *buf, ssize_t buf_len)

API to send a complete HTTP response.

This API will send the data as an HTTP response to the request. This assumes that you have the entire response ready in a single buffer. If you wish to send response in incremental chunks use httpd_resp_send_chunk() instead.

If no status code and content-type were set, by default this will send 200 OK status code and content type as text/html. You may call the following functions before this API to configure the response headers : httpd_resp_set_status() - for setting the HTTP status string, httpd_resp_set_type() - for setting the Content Type, httpd_resp_set_hdr() - for appending any additional field value entries in the response header

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once this API is called, the request has been responded to.
- No additional data can then be sent for the request.
- Once this API is called, all request headers are purged, so request headers need be copied into separate buffers if they are required later.

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null request pointer
- ESP_ERR_HTTPD_RESP_HDR : Essential headers are too large for internal buffer
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request

Parameters

- `[in] r`: The request being responded to
- `[in] buf`: Buffer from where the content is to be fetched
- `[in] buf_len`: Length of the buffer, HTTPD_RESP_USE_STRLEN to use strlen()

esp_err_t **httpd_resp_send_chunk** (*httpd_req_t* *r, **const** char *buf, ssize_t buf_len)

API to send one HTTP chunk.

This API will send the data as an HTTP response to the request. This API will use chunked-encoding and send the response in the form of chunks. If you have the entire response contained in a single buffer, please use httpd_resp_send() instead.

If no status code and content-type were set, by default this will send 200 OK status code and content type as text/html. You may call the following functions before this API to configure the response headers httpd_resp_set_status() - for setting the HTTP status string, httpd_resp_set_type() - for setting the Content Type, httpd_resp_set_hdr() - for appending any additional field value entries in the response header

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- When you are finished sending all your chunks, you must call this function with buf_len as 0.

- Once this API is called, all request headers are purged, so request headers need be copied into separate buffers if they are required later.

Return

- ESP_OK : On successfully sending the response packet chunk
- ESP_ERR_INVALID_ARG : Null request pointer
- ESP_ERR_HTTPD_RESP_HDR : Essential headers are too large for internal buffer
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to
- `[in] buf`: Pointer to a buffer that stores the data
- `[in] buf_len`: Length of the buffer, HTTPD_RESP_USE_STRLEN to use strlen()

static *esp_err_t* **httpd_resp_sendstr** (*httpd_req_t* *r, **const** char *str)

API to send a complete string as HTTP response.

This API simply calls http_resp_send with buffer length set to string length assuming the buffer contains a null terminated string

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null request pointer
- ESP_ERR_HTTPD_RESP_HDR : Essential headers are too large for internal buffer
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request

Parameters

- `[in] r`: The request being responded to
- `[in] str`: String to be sent as response body

static *esp_err_t* **httpd_resp_sendstr_chunk** (*httpd_req_t* *r, **const** char *str)

API to send a string as an HTTP response chunk.

This API simply calls http_resp_send_chunk with buffer length set to string length assuming the buffer contains a null terminated string

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null request pointer
- ESP_ERR_HTTPD_RESP_HDR : Essential headers are too large for internal buffer
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request

Parameters

- `[in] r`: The request being responded to
- `[in] str`: String to be sent as response body (NULL to finish response packet)

esp_err_t **httpd_resp_set_status** (*httpd_req_t* *r, **const** char *status)

API to set the HTTP status code.

This API sets the status of the HTTP response to the value specified. By default, the '200 OK' response is sent as the response.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- This API only sets the status to this value. The status isn't sent out until any of the send APIs is executed.
- Make sure that the lifetime of the status string is valid till send function is called.

Return

- ESP_OK : On success
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to

- `[in] status`: The HTTP status code of this response

esp_err_t **httpd_resp_set_type** (*httpd_req_t* *r, **const** char *type)

API to set the HTTP content type.

This API sets the 'Content Type' field of the response. The default content type is 'text/html'.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- This API only sets the content type to this value. The type isn't sent out until any of the send APIs is executed.
- Make sure that the lifetime of the type string is valid till send function is called.

Return

- ESP_OK : On success
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to
- `[in] type`: The Content Type of the response

esp_err_t **httpd_resp_set_hdr** (*httpd_req_t* *r, **const** char *field, **const** char *value)

API to append any additional headers.

This API sets any additional header fields that need to be sent in the response.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- The header isn't sent out until any of the send APIs is executed.
- The maximum allowed number of additional headers is limited to value of max_resp_headers in config structure.
- Make sure that the lifetime of the field value strings are valid till send function is called.

Return

- ESP_OK : On successfully appending new header
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_RESP_HDR : Total additional headers exceed max allowed
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to
- `[in] field`: The field name of the HTTP header
- `[in] value`: The value of this HTTP header

esp_err_t **httpd_resp_send_err** (*httpd_req_t* *req, *httpd_err_code_t* error, **const** char *msg)

For sending out error code in response to HTTP request.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once this API is called, all request headers are purged, so request headers need be copied into separate buffers if they are required later.
- If you wish to send additional data in the body of the response, please use the lower-level functions directly.

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] req`: Pointer to the HTTP request for which the response needs to be sent
- `[in] error`: Error type to send
- `[in] msg`: Error message string (pass NULL for default message)

static *esp_err_t* **httpd_resp_send_404** (*httpd_req_t* *r)

Helper function for HTTP 404.

Send HTTP 404 message. If you wish to send additional data in the body of the response, please use the lower-level functions directly.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once this API is called, all request headers are purged, so request headers need be copied into separate buffers if they are required later.

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to

static *esp_err_t* **httpd_resp_send_408** (*httpd_req_t* *r)

Helper function for HTTP 408.

Send HTTP 408 message. If you wish to send additional data in the body of the response, please use the lower-level functions directly.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once this API is called, all request headers are purged, so request headers need be copied into separate buffers if they are required later.

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to

static *esp_err_t* **httpd_resp_send_500** (*httpd_req_t* *r)

Helper function for HTTP 500.

Send HTTP 500 message. If you wish to send additional data in the body of the response, please use the lower-level functions directly.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Once this API is called, all request headers are purged, so request headers need be copied into separate buffers if they are required later.

Return

- ESP_OK : On successfully sending the response packet
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_HTTPD_RESP_SEND : Error in raw send
- ESP_ERR_HTTPD_INVALID_REQ : Invalid request pointer

Parameters

- `[in] r`: The request being responded to

int **httpd_send** (*httpd_req_t* *r, **const** char *buf, size_t buf_len)

Raw HTTP send.

Call this API if you wish to construct your custom response packet. When using this, all essential header, eg. HTTP version, Status Code, Content Type and Length, Encoding, etc. will have to be constructed manually,

and HTTP delimeters (CRLF) will need to be placed correctly for separating sub-sections of the HTTP response packet.

If the send override function is set, this API will end up calling that function eventually to send data out.

Note

- This API is supposed to be called only from the context of a URI handler where httpd_req_t* request pointer is valid.
- Unless the response has the correct HTTP structure (which the user must now ensure) it is not guaranteed that it will be recognized by the client. For most cases, you wouldn't have to call this API, but you would rather use either of : httpd_resp_send(), httpd_resp_send_chunk()

Return

- Bytes : Number of bytes that were sent successfully
- HTTPD_SOCK_ERR_INVALID : Invalid arguments
- HTTPD_SOCK_ERR_TIMEOUT : Timeout/interrupted while calling socket send()
- HTTPD_SOCK_ERR_FAIL : Unrecoverable error while calling socket send()

Parameters

- `[in] r`: The request being responded to
- `[in] buf`: Buffer from where the fully constructed packet is to be read
- `[in] buf_len`: Length of the buffer

esp_err_t **httpd_register_err_handler**(*httpd_handle_t handle*, *httpd_err_code_t error*, *httpd_err_handler_func_t handler_fn*)

Function for registering HTTP error handlers.

This function maps a handler function to any supported error code given by `httpd_err_code_t`. See prototype `httpd_err_handler_func_t` above for details.

Return

- ESP_OK : handler registered successfully
- ESP_ERR_INVALID_ARG : invalid error code or server handle

Parameters

- `[in] handle`: HTTP server handle
- `[in] error`: Error type
- `[in] handler_fn`: User implemented handler function (Pass NULL to unset any previously set handler)

esp_err_t **httpd_start**(*httpd_handle_t* **handle*, **const** *httpd_config_t* **config*)

Starts the web server.

Create an instance of HTTP server and allocate memory/resources for it depending upon the specified configuration.

Example usage:

```
//Function for starting the webserver
httpd_handle_t start_webserver(void)
{
     // Generate default configuration
     httpd_config_t config = HTTPD_DEFAULT_CONFIG();

     // Empty handle to http_server
     httpd_handle_t server = NULL;

     // Start the httpd server
     if (httpd_start(&server, &config) == ESP_OK) {
         // Register URI handlers
         httpd_register_uri_handler(server, &uri_get);
         httpd_register_uri_handler(server, &uri_post);
     }
     // If server failed to start, handle will be NULL
     return server;
}
```

Return

- ESP_OK : Instance created successfully
- ESP_ERR_INVALID_ARG : Null argument(s)
- ESP_ERR_HTTPD_ALLOC_MEM : Failed to allocate memory for instance
- ESP_ERR_HTTPD_TASK : Failed to launch server task

Parameters

- `[in] config`: Configuration for new instance of the server
- `[out] handle`: Handle to newly created instance of the server. NULL on error

esp_err_t **httpd_stop**(*httpd_handle_t handle*)

Stops the web server.

Deallocates memory/resources used by an HTTP server instance and deletes it. Once deleted the handle can no longer be used for accessing the instance.

Example usage:

```
// Function for stopping the webserver
void stop_webserver(httpd_handle_t server)
{
    // Ensure handle is non NULL
    if (server != NULL) {
        // Stop the httpd server
        httpd_stop(server);
    }
}
```

Return

- ESP_OK : Server stopped successfully
- ESP_ERR_INVALID_ARG : Handle argument is Null

Parameters

- `[in] handle`: Handle to server returned by httpd_start

esp_err_t **httpd_queue_work**(*httpd_handle_t handle*, *httpd_work_fn_t work*, void **arg*)

Queue execution of a function in HTTPD's context.

This API queues a work function for asynchronous execution

Note Some protocols require that the web server generate some asynchronous data and send it to the persistently opened connection. This facility is for use by such protocols.

Return

- ESP_OK : On successfully queueing the work
- ESP_FAIL : Failure in ctrl socket
- ESP_ERR_INVALID_ARG : Null arguments

Parameters

- `[in] handle`: Handle to server returned by httpd_start
- `[in] work`: Pointer to the function to be executed in the HTTPD's context
- `[in] arg`: Pointer to the arguments that should be passed to this function

void ***httpd_sess_get_ctx**(*httpd_handle_t handle*, int *sockfd*)

Get session context from socket descriptor.

Typically if a session context is created, it is available to URI handlers through the httpd_req_t structure. But, there are cases where the web server's send/receive functions may require the context (for example, for accessing keying information etc). Since the send/receive function only have the socket descriptor at their disposal, this API provides them with a way to retrieve the session context.

Return

- void* : Pointer to the context associated with this session
- NULL : Empty context / Invalid handle / Invalid socket fd

Parameters

- `[in] handle`: Handle to server returned by httpd_start
- `[in] sockfd`: The socket descriptor for which the context should be extracted.

void **httpd_sess_set_ctx**(*httpd_handle_t handle*, int *sockfd*, void **ctx*, *httpd_free_ctx_fn_t free_fn*)
Set session context by socket descriptor.

Parameters
- `[in] handle`: Handle to server returned by httpd_start
- `[in] sockfd`: The socket descriptor for which the context should be extracted.
- `[in] ctx`: Context object to assign to the session
- `[in] free_fn`: Function that should be called to free the context

void ****httpd_sess_get_transport_ctx**(*httpd_handle_t handle*, int *sockfd*)
Get session 'transport' context by socket descriptor.

This context is used by the send/receive functions, for example to manage SSL context.

See httpd_sess_get_ctx()

Return
- void* : Pointer to the transport context associated with this session
- NULL : Empty context / Invalid handle / Invalid socket fd

Parameters
- `[in] handle`: Handle to server returned by httpd_start
- `[in] sockfd`: The socket descriptor for which the context should be extracted.

void **httpd_sess_set_transport_ctx**(*httpd_handle_t handle*, int *sockfd*, void **ctx*, *httpd_free_ctx_fn_t free_fn*)
Set session 'transport' context by socket descriptor.

See httpd_sess_set_ctx()

Parameters
- `[in] handle`: Handle to server returned by httpd_start
- `[in] sockfd`: The socket descriptor for which the context should be extracted.
- `[in] ctx`: Transport context object to assign to the session
- `[in] free_fn`: Function that should be called to free the transport context

void ****httpd_get_global_user_ctx**(*httpd_handle_t handle*)
Get HTTPD global user context (it was set in the server config struct)

Return global user context

Parameters
- `[in] handle`: Handle to server returned by httpd_start

void ****httpd_get_global_transport_ctx**(*httpd_handle_t handle*)
Get HTTPD global transport context (it was set in the server config struct)

Return global transport context

Parameters
- `[in] handle`: Handle to server returned by httpd_start

esp_err_t **httpd_sess_trigger_close**(*httpd_handle_t handle*, int *sockfd*)
Trigger an httpd session close externally.

Note Calling this API is only required in special circumstances wherein some application requires to close an httpd client session asynchronously.

Return
- ESP_OK : On successfully initiating closure
- ESP_FAIL : Failure to queue work
- ESP_ERR_NOT_FOUND : Socket fd not found
- ESP_ERR_INVALID_ARG : Null arguments

Parameters
- `[in] handle`: Handle to server returned by httpd_start
- `[in] sockfd`: The socket descriptor of the session to be closed

esp_err_t **httpd_sess_update_lru_counter**(*httpd_handle_t handle*, int *sockfd*)
Update LRU counter for a given socket.

LRU Counters are internally associated with each session to monitor how recently a session exchanged traffic. When LRU purge is enabled, if a client is requesting for connection but maximum number of sockets/sessions is reached, then the session having the earliest LRU counter is closed automatically.

Updating the LRU counter manually prevents the socket from being purged due to the Least Recently Used (LRU) logic, even though it might not have received traffic for some time. This is useful when all open sockets/session are frequently exchanging traffic but the user specifically wants one of the sessions to be kept open, irrespective of when it last exchanged a packet.

Note Calling this API is only necessary if the LRU Purge Enable option is enabled.

Return

- ESP_OK : Socket found and LRU counter updated
- ESP_ERR_NOT_FOUND : Socket not found
- ESP_ERR_INVALID_ARG : Null arguments

Parameters

- `[in] handle`: Handle to server returned by httpd_start
- `[in] sockfd`: The socket descriptor of the session for which LRU counter is to be updated

Structures

`struct httpd_config`

HTTP Server Configuration Structure.

Note Use HTTPD_DEFAULT_CONFIG() to initialize the configuration to a default value and then modify only those fields that are specifically determined by the use case.

Public Members

unsigned **`task_priority`**
Priority of FreeRTOS task which runs the server

size_t **`stack_size`**
The maximum stack size allowed for the server task

BaseType_t **`core_id`**
The core the HTTP server task will run on

uint16_t **`server_port`**
TCP Port number for receiving and transmitting HTTP traffic

uint16_t **`ctrl_port`**
UDP Port number for asynchronously exchanging control signals between various components of the server

uint16_t **`max_open_sockets`**
Max number of sockets/clients connected at any time

uint16_t **`max_uri_handlers`**
Maximum allowed uri handlers

uint16_t **`max_resp_headers`**
Maximum allowed additional headers in HTTP response

uint16_t **`backlog_conn`**
Number of backlog connections

bool **`lru_purge_enable`**
Purge "Least Recently Used" connection

uint16_t **`recv_wait_timeout`**
Timeout for recv function (in seconds)

uint16_t **`send_wait_timeout`**
Timeout for send function (in seconds)

void ***global_user_ctx**
Global user context.

This field can be used to store arbitrary user data within the server context. The value can be retrieved using the server handle, available e.g. in the httpd_req_t struct.

When shutting down, the server frees up the user context by calling free() on the global_user_ctx field. If you wish to use a custom function for freeing the global user context, please specify that here.

httpd_free_ctx_fn_t **global_user_ctx_free_fn**
Free function for global user context

void ***global_transport_ctx**
Global transport context.

Similar to global_user_ctx, but used for session encoding or encryption (e.g. to hold the SSL context). It will be freed using free(), unless global_transport_ctx_free_fn is specified.

httpd_free_ctx_fn_t **global_transport_ctx_free_fn**
Free function for global transport context

httpd_open_func_t **open_fn**
Custom session opening callback.

Called on a new session socket just after accept(), but before reading any data.

This is an opportunity to set up e.g. SSL encryption using global_transport_ctx and the send/recv/pending session overrides.

If a context needs to be maintained between these functions, store it in the session using httpd_sess_set_transport_ctx() and retrieve it later with httpd_sess_get_transport_ctx()

Returning a value other than ESP_OK will immediately close the new socket.

httpd_close_func_t **close_fn**
Custom session closing callback.

Called when a session is deleted, before freeing user and transport contexts and before closing the socket. This is a place for custom de-init code common to all sockets.

Set the user or transport context to NULL if it was freed here, so the server does not try to free it again.

This function is run for all terminated sessions, including sessions where the socket was closed by the network stack - that is, the file descriptor may not be valid anymore.

httpd_uri_match_func_t **uri_match_fn**
URI matcher function.

Called when searching for a matching URI: 1) whose request handler is to be executed right after an HTTP request is successfully parsed 2) in order to prevent duplication while registering a new URI handler using `httpd_register_uri_handler()`

Available options are: 1) NULL : Internally do basic matching using `strncmp()` 2) `httpd_uri_match_wildcard()` : URI wildcard matcher

Users can implement their own matching functions (See description of the `httpd_uri_match_func_t` function prototype)

struct httpd_req
HTTP Request Data Structure.

Public Members

httpd_handle_t **handle**
Handle to server instance

int **method**
The type of HTTP request, -1 if unsupported method

const char **uri**[**HTTPD_MAX_URI_LEN** + 1]
The URI of this request (1 byte extra for null termination)

size_t **content_len**
Length of the request body

void ***aux**
Internally used members

void ***user_ctx**
User context pointer passed during URI registration.

void ***sess_ctx**
Session Context Pointer

A session context. Contexts are maintained across 'sessions' for a given open TCP connection. One session could have multiple request responses. The web server will ensure that the context persists across all these request and responses.

By default, this is NULL. URI Handlers can set this to any meaningful value.

If the underlying socket gets closed, and this pointer is non-NULL, the web server will free up the context by calling free(), unless free_ctx function is set.

httpd_free_ctx_fn_t **free_ctx**
Pointer to free context hook

Function to free session context

If the web server's socket closes, it frees up the session context by calling free() on the sess_ctx member. If you wish to use a custom function for freeing the session context, please specify that here.

bool **ignore_sess_ctx_changes**
Flag indicating if Session Context changes should be ignored

By default, if you change the sess_ctx in some URI handler, the http server will internally free the earlier context (if non NULL), after the URI handler returns. If you want to manage the allocation/reallocation/freeing of sess_ctx yourself, set this flag to true, so that the server will not perform any checks on it. The context will be cleared by the server (by calling free_ctx or free()) only if the socket gets closed.

struct httpd_uri
Structure for URI handler.

Public Members

const char ***uri**
The URI to handle

httpd_method_t **method**
Method supported by the URI

esp_err_t (***handler**) (*httpd_req_t* *r)
Handler to call for supported request method. This must return ESP_OK, or else the underlying socket will be closed.

void ***user_ctx**
Pointer to user context data which will be available to handler

Macros

HTTPD_MAX_REQ_HDR_LEN
HTTPD_MAX_URI_LEN

HTTPD_SOCK_ERR_FAIL

HTTPD_SOCK_ERR_INVALID

HTTPD_SOCK_ERR_TIMEOUT

HTTPD_200
HTTP Response 200

HTTPD_204
HTTP Response 204

HTTPD_207
HTTP Response 207

HTTPD_400
HTTP Response 400

HTTPD_404
HTTP Response 404

HTTPD_408
HTTP Response 408

HTTPD_500
HTTP Response 500

HTTPD_TYPE_JSON
HTTP Content type JSON

HTTPD_TYPE_TEXT
HTTP Content type text/HTML

HTTPD_TYPE_OCTET
HTTP Content type octext-stream

HTTPD_DEFAULT_CONFIG()

ESP_ERR_HTTPD_BASE
Starting number of HTTPD error codes

ESP_ERR_HTTPD_HANDLERS_FULL
All slots for registering URI handlers have been consumed

ESP_ERR_HTTPD_HANDLER_EXISTS
URI handler with same method and target URI already registered

ESP_ERR_HTTPD_INVALID_REQ
Invalid request pointer

ESP_ERR_HTTPD_RESULT_TRUNC
Result string truncated

ESP_ERR_HTTPD_RESP_HDR
Response header field larger than supported

ESP_ERR_HTTPD_RESP_SEND
Error occured while sending response packet

ESP_ERR_HTTPD_ALLOC_MEM
Failed to dynamically allocate memory for resource

ESP_ERR_HTTPD_TASK
Failed to launch server task/thread

HTTPD_RESP_USE_STRLEN

Type Definitions

typedef struct *httpd_req* **httpd_req_t**
HTTP Request Data Structure.

typedef struct *httpd_uri* **httpd_uri_t**
Structure for URI handler.

typedef int (***httpd_send_func_t**) (*httpd_handle_t* hd, int sockfd, **const** char *buf, size_t buf_len, int flags)

Prototype for HTTPDs low-level send function.

Note User specified send function must handle errors internally, depending upon the set value of errno, and return specific HTTPD_SOCK_ERR_ codes, which will eventually be conveyed as return value of httpd_send() function

Return

- Bytes : The number of bytes sent successfully
- HTTPD_SOCK_ERR_INVALID : Invalid arguments
- HTTPD_SOCK_ERR_TIMEOUT : Timeout/interrupted while calling socket send()
- HTTPD_SOCK_ERR_FAIL : Unrecoverable error while calling socket send()

Parameters

- `[in] hd`: server instance
- `[in] sockfd`: session socket file descriptor
- `[in] buf`: buffer with bytes to send
- `[in] buf_len`: data size
- `[in] flags`: flags for the send() function

typedef int (***httpd_recv_func_t**) (*httpd_handle_t* hd, int sockfd, char *buf, size_t buf_len, int flags)

Prototype for HTTPDs low-level recv function.

Note User specified recv function must handle errors internally, depending upon the set value of errno, and return specific HTTPD_SOCK_ERR_ codes, which will eventually be conveyed as return value of httpd_req_recv() function

Return

- Bytes : The number of bytes received successfully
- 0 : Buffer length parameter is zero / connection closed by peer
- HTTPD_SOCK_ERR_INVALID : Invalid arguments
- HTTPD_SOCK_ERR_TIMEOUT : Timeout/interrupted while calling socket recv()
- HTTPD_SOCK_ERR_FAIL : Unrecoverable error while calling socket recv()

Parameters

- `[in] hd`: server instance
- `[in] sockfd`: session socket file descriptor
- `[in] buf`: buffer with bytes to send
- `[in] buf_len`: data size
- `[in] flags`: flags for the send() function

typedef int (***httpd_pending_func_t**) (*httpd_handle_t* hd, int sockfd)

Prototype for HTTPDs low-level "get pending bytes" function.

Note User specified pending function must handle errors internally, depending upon the set value of errno, and return specific HTTPD_SOCK_ERR_ codes, which will be handled accordingly in the server task.

Return

- Bytes : The number of bytes waiting to be received
- HTTPD_SOCK_ERR_INVALID : Invalid arguments
- HTTPD_SOCK_ERR_TIMEOUT : Timeout/interrupted while calling socket pending()
- HTTPD_SOCK_ERR_FAIL : Unrecoverable error while calling socket pending()

Parameters

- `[in] hd`: server instance
- `[in] sockfd`: session socket file descriptor

typedef *esp_err_t* (***httpd_err_handler_func_t**) (*httpd_req_t* *req, *httpd_err_code_t* error)

Function prototype for HTTP error handling.

This function is executed upon HTTP errors generated during internal processing of an HTTP request. This is used to override the default behavior on error, which is to send HTTP error response and close the underlying socket.

Note

- If implemented, the server will not automatically send out HTTP error response codes, therefore,

httpd_resp_send_err() must be invoked inside this function if user wishes to generate HTTP error responses.
- When invoked, the validity of `uri`, `method`, `content_len` and `user_ctx` fields of the httpd_req_t parameter is not guaranteed as the HTTP request may be partially received/parsed.
- The function must return ESP_OK if underlying socket needs to be kept open. Any other value will ensure that the socket is closed. The return value is ignored when error is of type `HTTPD_500_INTERNAL_SERVER_ERROR` and the socket closed anyway.

Return
- ESP_OK : error handled successful
- ESP_FAIL : failure indicates that the underlying socket needs to be closed

Parameters
- `[in] req`: HTTP request for which the error needs to be handled
- `[in] error`: Error type

typedef void ***httpd_handle_t**
HTTP Server Instance Handle.

Every instance of the server will have a unique handle.

typedef enum http_method **httpd_method_t**
HTTP Method Type wrapper over "enum http_method" available in "http_parser" library.

typedef void (***httpd_free_ctx_fn_t**)(void *ctx)
Prototype for freeing context data (if any)

Parameters
- `[in] ctx`: object to free

typedef *esp_err_t* (***httpd_open_func_t**)(*httpd_handle_t* hd, int sockfd)
Function prototype for opening a session.

Called immediately after the socket was opened to set up the send/recv functions and other parameters of the socket.

Return
- ESP_OK : On success
- Any value other than ESP_OK will signal the server to close the socket immediately

Parameters
- `[in] hd`: server instance
- `[in] sockfd`: session socket file descriptor

typedef void (***httpd_close_func_t**)(*httpd_handle_t* hd, int sockfd)
Function prototype for closing a session.

Note It's possible that the socket descriptor is invalid at this point, the function is called for all terminated sessions. Ensure proper handling of return codes.

Parameters
- `[in] hd`: server instance
- `[in] sockfd`: session socket file descriptor

typedef bool (***httpd_uri_match_func_t**)(**const** char *reference_uri, **const** char *uri_to_match, size_t match_upto)
Function prototype for URI matching.

Return true on match

Parameters
- `[in] reference_uri`: URI/template with respect to which the other URI is matched
- `[in] uri_to_match`: URI/template being matched to the reference URI/template
- `[in] match_upto`: For specifying the actual length of `uri_to_match` up to which the matching algorithm is to be applied (The maximum value is `strlen(uri_to_match)`, independent of the length of `reference_uri`)

typedef struct *httpd_config* **httpd_config_t**
HTTP Server Configuration Structure.

Note Use HTTPD_DEFAULT_CONFIG() to initialize the configuration to a default value and then modify only those fields that are specifically determined by the use case.

typedef void (***httpd_work_fn_t**) (void *arg)

Prototype of the HTTPD work function Please refer to httpd_queue_work() for more details.

Parameters

- `[in] arg`: The arguments for this work function

Enumerations

enum httpd_err_code_t

Error codes sent as HTTP response in case of errors encountered during processing of an HTTP request.

Values:

HTTPD_500_INTERNAL_SERVER_ERROR = 0

HTTPD_501_METHOD_NOT_IMPLEMENTED

HTTPD_505_VERSION_NOT_SUPPORTED

HTTPD_400_BAD_REQUEST

HTTPD_404_NOT_FOUND

HTTPD_405_METHOD_NOT_ALLOWED

HTTPD_408_REQ_TIMEOUT

HTTPD_411_LENGTH_REQUIRED

HTTPD_414_URI_TOO_LONG

HTTPD_431_REQ_HDR_FIELDS_TOO_LARGE

HTTPD_ERR_CODE_MAX

2.3.6 HTTPS server

Overview

This component is built on top of *esp_http_server*. The HTTPS server takes advantage of hooks and function overrides in the regular HTTP server to provide encryption using OpenSSL.

All documentation for *esp_http_server* applies also to a server you create this way.

Used APIs

The following API of *esp_http_server* should not be used with *esp_https_server*, as they are used internally to handle secure sessions and to maintain internal state:

- "send", "receive" and "pending" function overrides - secure socket handling
  - *httpd_sess_set_send_override()*
  - *httpd_sess_set_recv_override()*
  - *httpd_sess_set_pending_override()*
- "transport context" - both global and session
  - *httpd_sess_get_transport_ctx()* - returns SSL used for the session
  - *httpd_sess_set_transport_ctx()*
  - *httpd_get_global_transport_ctx()* - returns the shared SSL context
  - `httpd_config_t.global_transport_ctx`
  - `httpd_config_t.global_transport_ctx_free_fn`
  - `httpd_config_t.open_fn` - used to set up secure sockets

Everything else can be used without limitations.

Usage

Please see the example protocols/https_server to learn how to set up a secure server.

Basically all you need is to generate a certificate, embed it in the firmware, and provide its pointers and lengths to the start function via the init struct.

The server can be started with or without SSL by changing a flag in the init struct - `httpd_ssl_config.transport_mode`. This could be used e.g. for testing or in trusted environments where you prefer speed over security.

Performance

The initial session setup can take about two seconds, or more with slower clock speeds or more verbose logging. Subsequent requests through the open secure socket are much faster (down to under 100 ms).

API Reference

Header File

- esp_https_server/include/esp_https_server.h

Functions

esp_err_t **httpd_ssl_start** (*httpd_handle_t *handle*, *httpd_ssl_config_t *config*)

Create a SSL capable HTTP server (secure mode may be disabled in config)

Return success

Parameters

- `[inout] config`: - server config, must not be const. Does not have to stay valid after calling this function.
- `[out] handle`: - storage for the server handle, must be a valid pointer

void **httpd_ssl_stop** (*httpd_handle_t handle*)

Stop the server. Blocks until the server is shut down.

Parameters

- `[in] handle`:

Structures

struct httpd_ssl_config

HTTPS server config struct

Please use HTTPD_SSL_CONFIG_DEFAULT() to initialize it.

Public Members

httpd_config_t **httpd**

Underlying HTTPD server config

Parameters like task stack size and priority can be adjusted here.

const uint8_t ***cacert_pem**

CA certificate (here it is treated as server cert) Todo: Fix this change in release/v5.0 as it would be a breaking change i.e. Rename the nomenclature of variables holding different certs in https_server component as well as example 1)The cacert variable should hold the CA which is used to authenticate clients (should inherit current role of client_verify_cert_pem var) 2)There should be another variable servercert which whould hold servers own certificate (should inherit current role of cacert var)

size_t **cacert_len**

CA certificate byte length

const uint8_t ***client_verify_cert_pem**
Client verify authority certificate (CA used to sign clients, or client cert itself

size_t **client_verify_cert_len**
Client verify authority cert len

const uint8_t ***prvtkey_pem**
Private key

size_t **prvtkey_len**
Private key byte length

httpd_ssl_transport_mode_t **transport_mode**
Transport Mode (default secure)

uint16_t **port_secure**
Port used when transport mode is secure (default 443)

uint16_t **port_insecure**
Port used when transport mode is insecure (default 80)

Macros

HTTPD_SSL_CONFIG_DEFAULT()
Default config struct init

(http_server default config had to be copied for customization)

Notes:

- port is set when starting the server, according to 'transport_mode'
- one socket uses ~ 40kB RAM with SSL, we reduce the default socket count to 4
- SSL sockets are usually long-lived, closing LRU prevents pool exhaustion DOS
- Stack size may need adjustments depending on the user application

Type Definitions

typedef struct *httpd_ssl_config* **httpd_ssl_config_t**

Enumerations

enum httpd_ssl_transport_mode_t
Values:

HTTPD_SSL_TRANSPORT_SECURE

HTTPD_SSL_TRANSPORT_INSECURE

2.3.7 ICMP Echo

Overview

ICMP (Internet Control Message Protocol) is used for diagnostic or control purposes or generated in response to errors in IP operations. The common network util `ping` is implemented based on the ICMP packets with the type field value of 0, also called `Echo Reply`.

During a ping session, the source host firstly sends out an ICMP echo request packet and wait for an ICMP echo reply with specific times. In this way, it also measures the round-trip time for the messages. After receiving a valid ICMP echo reply, the source host will generate statistics about the IP link layer (e.g. packet loss, elapsed time, etc).

It is common that IoT device needs to check whether a remote server is alive or not. The device should show the warnings to users when it got offline. It can be achieved by creating a ping session and sending/parsing ICMP echo packets periodically.

To make this internal procedure much easier for users, ESP-IDF provides some out-of-box APIs.

Create a new ping session To create a ping session, you need to fill in the `esp_ping_config_t` configuration structure firstly, specifying target IP address, interval times, and etc. Optionally, you can also register some callback functions with the *esp_ping_callbacks_t*‘ structure.

Example method to create a new ping session and register callbacks:

```
static void test_on_ping_success(esp_ping_handle_t hdl, void *args)
{
    // optionally, get callback arguments
    // const char* str = (const char*) args;
    // printf("%s\r\n", str); // "foo"
    uint8_t ttl;
    uint16_t seqno;
    uint32_t elapsed_time, recv_len;
    ip_addr_t target_addr;
    esp_ping_get_profile(hdl, ESP_PING_PROF_SEQNO, &seqno, sizeof(seqno));
    esp_ping_get_profile(hdl, ESP_PING_PROF_TTL, &ttl, sizeof(ttl));
    esp_ping_get_profile(hdl, ESP_PING_PROF_IPADDR, &target_addr, sizeof(target_
↪addr));
    esp_ping_get_profile(hdl, ESP_PING_PROF_SIZE, &recv_len, sizeof(recv_len));
    esp_ping_get_profile(hdl, ESP_PING_PROF_TIMEGAP, &elapsed_time, sizeof(elapsed_
↪time));
    printf("%d bytes from %s icmp_seq=%d ttl=%d time=%d ms\n",
           recv_len, inet_ntoa(target_addr.u_addr.ip4), seqno, ttl, elapsed_time);
}

static void test_on_ping_timeout(esp_ping_handle_t hdl, void *args)
{
    uint16_t seqno;
    ip_addr_t target_addr;
    esp_ping_get_profile(hdl, ESP_PING_PROF_SEQNO, &seqno, sizeof(seqno));
    esp_ping_get_profile(hdl, ESP_PING_PROF_IPADDR, &target_addr, sizeof(target_
↪addr));
    printf("From %s icmp_seq=%d timeout\n", inet_ntoa(target_addr.u_addr.ip4),␣
↪seqno);
}

static void test_on_ping_end(esp_ping_handle_t hdl, void *args)
{
    uint32_t transmitted;
    uint32_t received;
    uint32_t total_time_ms;

    esp_ping_get_profile(hdl, ESP_PING_PROF_REQUEST, &transmitted,␣
↪sizeof(transmitted));
    esp_ping_get_profile(hdl, ESP_PING_PROF_REPLY, &received, sizeof(received));
    esp_ping_get_profile(hdl, ESP_PING_PROF_DURATION, &total_time_ms, sizeof(total_
↪time_ms));
    printf("%d packets transmitted, %d received, time %dms\n", transmitted,␣
↪received, total_time_ms);
}

void initialize_ping()
{
    /* convert URL to IP address */
    ip_addr_t target_addr;
    struct addrinfo hint;
    struct addrinfo *res = NULL;
    memset(&hint, 0, sizeof(hint));
    memset(&target_addr, 0, sizeof(target_addr));
    getaddrinfo("www.espressif.com", NULL, &hint, &res);
    struct in_addr addr4 = ((struct sockaddr_in *) (res->ai_addr))->sin_addr;
    inet_addr_to_ip4addr(ip_2_ip4(&target_addr), &addr4);
```

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```
    freeaddrinfo(res);

    esp_ping_config_t ping_config = ESP_PING_DEFAULT_CONFIG();
    ping_config.target_addr = target_addr;          // target IP address
    ping_config.count = ESP_PING_COUNT_INFINITE;    // ping in infinite mode, esp_
↪ping_stop can stop it

    /* set callback functions */
    esp_ping_callbacks_t cbs;
    cbs.on_ping_success = test_on_ping_success;
    cbs.on_ping_timeout = test_on_ping_timeout;
    cbs.on_ping_end = test_on_ping_end;
    cbs.cb_args = "foo";  // arguments that will feed to all callback functions,␣
↪can be NULL
    cbs.cb_args = eth_event_group;

    esp_ping_handle_t ping;
    esp_ping_new_session(&ping_config, &cbs, &ping);
}
```

Start and Stop ping session You can start and stop ping session with the handle returned by `esp_ping_new_session`. Note that, the ping session won't start automatically after creation. If the ping session is stopped, and restart again, the sequence number in ICMP packets will recount from zero again.

Delete a ping session If a ping session won't be used any more, you can delete it with `esp_ping_delete_session`. Please make sure the ping session is in stop state (i.e. you have called `esp_ping_stop` before or the ping session has finished all the procedures) when you call this function.

Get runtime statistics As the example code above, you can call `esp_ping_get_profile` to get different runtime statistics of ping session in the callback function.

Application Example

ICMP echo example: protocols/icmp_echo

API Reference

Header File

- lwip/include/apps/ping/ping_sock.h

Functions

esp_err_t **esp_ping_new_session**(**const** *esp_ping_config_t* *config, **const** *esp_ping_callbacks_t* *cbs, *esp_ping_handle_t* *hdl_out)

Create a ping session.

Return

- ESP_ERR_INVALID_ARG: invalid parameters (e.g. configuration is null, etc)
- ESP_ERR_NO_MEM: out of memory
- ESP_FAIL: other internal error (e.g. socket error)
- ESP_OK: create ping session successfully, user can take the ping handle to do follow-on jobs

Parameters

- `config`: ping configuration
- `cbs`: a bunch of callback functions invoked by internal ping task
- `hdl_out`: handle of ping session

esp_err_t **esp_ping_delete_session**(*esp_ping_handle_t hdl*)
Delete a ping session.

Return
- ESP_ERR_INVALID_ARG: invalid parameters (e.g. ping handle is null, etc)
- ESP_OK: delete ping session successfully

Parameters
- `hdl`: handle of ping session

esp_err_t **esp_ping_start**(*esp_ping_handle_t hdl*)
Start the ping session.

Return
- ESP_ERR_INVALID_ARG: invalid parameters (e.g. ping handle is null, etc)
- ESP_OK: start ping session successfully

Parameters
- `hdl`: handle of ping session

esp_err_t **esp_ping_stop**(*esp_ping_handle_t hdl*)
Stop the ping session.

Return
- ESP_ERR_INVALID_ARG: invalid parameters (e.g. ping handle is null, etc)
- ESP_OK: stop ping session successfully

Parameters
- `hdl`: handle of ping session

esp_err_t **esp_ping_get_profile**(*esp_ping_handle_t hdl*, *esp_ping_profile_t profile*, void **data*, uint32_t *size*)
Get runtime profile of ping session.

Return
- ESP_ERR_INVALID_ARG: invalid parameters (e.g. ping handle is null, etc)
- ESP_ERR_INVALID_SIZE: the actual profile data size doesn't match the "size" parameter
- ESP_OK: get profile successfully

Parameters
- `hdl`: handle of ping session
- `profile`: type of profile
- `data`: profile data
- `size`: profile data size

Structures

struct esp_ping_callbacks_t
Type of "ping" callback functions.

Public Members

void ***cb_args**
arguments for callback functions

void (***on_ping_success**)(*esp_ping_handle_t* hdl, void *args)
Invoked by internal ping thread when received ICMP echo reply packet.

void (***on_ping_timeout**)(*esp_ping_handle_t* hdl, void *args)
Invoked by internal ping thread when receive ICMP echo reply packet timeout.

void (***on_ping_end**)(*esp_ping_handle_t* hdl, void *args)
Invoked by internal ping thread when a ping session is finished.

struct esp_ping_config_t
Type of "ping" configuration.

Public Members

uint32_t **count**
A "ping" session contains count procedures

uint32_t **interval_ms**
Milliseconds between each ping procedure

uint32_t **timeout_ms**
Timeout value (in milliseconds) of each ping procedure

uint32_t **data_size**
Size of the data next to ICMP packet header

uint8_t **tos**
Type of Service, a field specified in the IP header

ip_addr_t **target_addr**
Target IP address, either IPv4 or IPv6

uint32_t **task_stack_size**
Stack size of internal ping task

uint32_t **task_prio**
Priority of internal ping task

uint32_t **interface**
Netif index, interface=0 means NETIF_NO_INDEX

Macros

ESP_PING_DEFAULT_CONFIG()
Default ping configuration.

ESP_PING_COUNT_INFINITE
Set ping count to zero will ping target infinitely

Type Definitions

typedef void ***esp_ping_handle_t**
Type of "ping" session handle.

Enumerations

enum esp_ping_profile_t
Profile of ping session.

Values:

ESP_PING_PROF_SEQNO
Sequence number of a ping procedure

ESP_PING_PROF_TTL
Time to live of a ping procedure

ESP_PING_PROF_REQUEST
Number of request packets sent out

ESP_PING_PROF_REPLY
Number of reply packets received

ESP_PING_PROF_IPADDR
IP address of replied target

ESP_PING_PROF_SIZE
Size of received packet

ESP_PING_PROF_TIMEGAP
Elapsed time between request and reply packet

ESP_PING_PROF_DURATION
Elapsed time of the whole ping session

2.3.8 ASIO port

Overview

Asio is a cross-platform C++ library, see https://think-async.com. It provides a consistent asynchronous model using a modern C++ approach.

ASIO documentation Please refer to the original asio documentation at https://think-async.com/Asio/Documentation. Asio also comes with a number of examples which could be find under Documentation/Examples on that web site.

Supported features ESP platform port currently supports only network asynchronous socket operations; does not support serial port and ssl. Internal asio settings for ESP include

- EXCEPTIONS are enabled in ASIO if enabled in menuconfig
- TYPEID is enabled in ASIO if enabled in menuconfig

Application Example

ESP examples are based on standard asio protocols/asio:

- protocols/asio/udp_echo_server
- protocols/asio/tcp_echo_server
- protocols/asio/chat_client
- protocols/asio/chat_server

Please refer to the specific example README.md for details

2.3.9 ESP-MQTT

Overview

ESP-MQTT is an implementation of MQTT protocol client (MQTT is a lightweight publish/subscribe messaging protocol).

Features

- Supports MQTT over TCP, SSL with mbedtls, MQTT over Websocket, MQTT over Websocket Secure.
- Easy to setup with URI
- Multiple instances (Multiple clients in one application)
- Support subscribing, publishing, authentication, last will messages, keep alive pings and all 3 QoS levels (it should be a fully functional client).

Application Example

- protocols/mqtt/tcp: MQTT over tcp, default port 1883
- protocols/mqtt/ssl: MQTT over tcp, default port 8883
- protocols/mqtt/ssl_psk: MQTT over tcp using pre-shared keys for authentication, default port 8883

- protocols/mqtt/ws: MQTT over Websocket, default port 80
- protocols/mqtt/wss: MQTT over Websocket Secure, default port 443

Configuration

URI

- Curently support `mqtt`, `mqtts`, `ws`, `wss` schemes
- MQTT over TCP samples:
  - `mqtt://mqtt.eclipse.org`: MQTT over TCP, default port 1883:
  - `mqtt://mqtt.eclipse.org:1884` MQTT over TCP, port 1884:
  - `mqtt://username:password@mqtt.eclipse.org:1884` MQTT over TCP, port 1884, with username and password
- MQTT over SSL samples:
  - `mqtts://mqtt.eclipse.org`: MQTT over SSL, port 8883
  - `mqtts://mqtt.eclipse.org:8884`: MQTT over SSL, port 8884
- MQTT over Websocket samples:
  - `ws://mqtt.eclipse.org:80/mqtt`
- MQTT over Websocket Secure samples:
  - `wss://mqtt.eclipse.org:443/mqtt`
- Minimal configurations:

```
const esp_mqtt_client_config_t mqtt_cfg = {
    .uri = "mqtt://mqtt.eclipse.org",
    // .user_context = (void *)your_context
};
esp_mqtt_client_handle_t client = esp_mqtt_client_init(&mqtt_cfg);
esp_mqtt_client_register_event(client, ESP_EVENT_ANY_ID, mqtt_event_handler,␣
↪client);
esp_mqtt_client_start(client);
```

- Note: By default mqtt client uses event loop library to post related mqtt events (connected, subscribed, published, etc.)

SSL

- Get certificate from server, example: `mqtt.eclipse.org openssl s_client -showcerts -connect mqtt.eclipse.org:8883 </dev/null 2>/dev/null|openssl x509 -outform PEM >mqtt_eclipse_org.pem`
- Check the sample application: `examples/mqtt_ssl`
- Configuration:

```
const esp_mqtt_client_config_t mqtt_cfg = {
    .uri = "mqtts://mqtt.eclipse.org:8883",
    .event_handle = mqtt_event_handler,
    .cert_pem = (const char *)mqtt_eclipse_org_pem_start,
};
```

If the certificate is not null-terminated then `cert_len` should also be set. Other SSL related configuration parameters are:

- `use_global_ca_store`: use the global certificate store to verify server certificate, see `esp-tls.h` for more information
- `client_cert_pem`: pointer to certificate data in PEM or DER format for SSL mutual authentication, default is NULL, not required if mutual authentication is not needed.
- `client_cert_len`: length of the buffer pointed to by client_cert_pem. May be 0 for null-terminated pem.
- `client_key_pem`: pointer to private key data in PEM or DER format for SSL mutual authentication, default is NULL, not required if mutual authentication is not needed.
- `client_key_len`: length of the buffer pointed to by client_key_pem. May be 0 for null-terminated pem.

- psk_hint_key: pointer to PSK struct defined in esp_tls.h to enable PSK authentication (as alternative to certificate verification). If not NULL and server/client certificates are NULL, PSK is enabled
- alpn_protos: NULL-terminated list of protocols to be used for ALPN.

Last Will and Testament MQTT allows for a last will and testament (LWT) message to notify other clients when a client ungracefully disconnects. This is configured by the following fields in the esp_mqtt_client_config_t-struct.

- lwt_topic: pointer to the LWT message topic
- lwt_msg: pointer to the LWT message
- lwt_msg_len: length of the LWT message, required if lwt_msg is not null-terminated
- lwt_qos: quality of service for the LWT message
- lwt_retain: specifies the retain flag of the LWT message

Other Configuration Parameters

- disable_clean_session: determines the clean session flag for the connect message, defaults to a clean session
- keepalive: determines how many seconds the client will wait for a ping response before disconnecting, default is 120 seconds.
- disable_auto_reconnect: enable to stop the client from reconnecting to server after errors or disconnects
- user_context: custom context that will be passed to the event handler
- task_prio: MQTT task priority, defaults to 5
- task_stack: MQTT task stack size, defaults to 6144 bytes, setting this will override setting from menuconfig
- buffer_size: size of MQTT send/receive buffer, default is 1024 bytes
- username: pointer to the username used for connecting to the broker
- password: pointer to the password used for connecting to the broker
- client_id: pointer to the client id, defaults to ESP32_%CHIPID% where %CHIPID% are the last 3 bytes of MAC address in hex format
- host: MQTT broker domain (ipv4 as string), setting the uri will override this
- port: MQTT broker port, specifying the port in the uri will override this
- transport: sets the transport protocol, setting the uri will override this
- refresh_connection_after_ms: refresh connection after this value (in milliseconds)
- event_handle: handle for MQTT events as a callback in legacy mode
- event_loop_handle: handle for MQTT event loop library

For more options on esp_mqtt_client_config_t, please refer to API reference below

Change settings in Project Configuration Menu The settings for MQTT can be found using idf.py menuconfig, under Component config -> ESP-MQTT Configuration

The following settings are available:

- *CONFIG_MQTT_PROTOCOL_311*: Enables 3.1.1 version of MQTT protocol
- *CONFIG_MQTT_TRANSPORT_SSL*, *CONFIG_MQTT_TRANSPORT_WEBSOCKET*: Enables specific MQTT transport layer, such as SSL, WEBSOCKET, WEBSOCKET_SECURE
- *CONFIG_MQTT_CUSTOM_OUTBOX*: Disables default implementation of mqtt_outbox, so a specific implementaion can be supplied

Events

The following events may be posted by the MQTT client:

- MQTT_EVENT_BEFORE_CONNECT: The client is initialized and about to start connecting to the broker.
- MQTT_EVENT_CONNECTED: The client has successfully established a connection to the broker. The client is now ready to send and receive data.

- MQTT_EVENT_DISCONNECTED: The client has aborted the connection due to being unable to read or write data, e.g. because the server is unavailable.
- MQTT_EVENT_SUBSCRIBED: The broker has acknowledged the client's subscribe request. The event data will contain the message ID of the subscribe message.
- MQTT_EVENT_UNSUBSCRIBED: The broker has acknowledged the client's unsubscribe request. The event data will contain the message ID of the unsubscribe message.
- MQTT_EVENT_PUBLISHED: The broker has acknowledged the client's publish message. This will only be posted for Quality of Service level 1 and 2, as level 0 does not use acknowledgements. The event data will contain the message ID of the publish message.
- MQTT_EVENT_DATA: The client has received a publish message. The event data contains: message ID, name of the topic it was published to, received data and its length. For data that exceeds the internal buffer multiple *MQTT_EVENT_DATA* will be posted and *current_data_offset* and *total_data_len* from event data updated to keep track of the fragmented message.
- MQTT_EVENT_ERROR: The client has encountered an error. *esp_mqtt_error_type_t* from *error_handle* in the event data can be used to further determine the type of the error. The type of error will determine which parts of the *error_handle* struct is filled.

API Reference

Header File

- mqtt/esp-mqtt/include/mqtt_client.h

Functions

esp_mqtt_client_handle_t **esp_mqtt_client_init**(**const** *esp_mqtt_client_config_t* **config*)

Creates mqtt client handle based on the configuration.

Return mqtt_client_handle if successfully created, NULL on error

Parameters

- config: mqtt configuration structure

esp_err_t **esp_mqtt_client_set_uri**(*esp_mqtt_client_handle_t client*, **const** char **uri*)

Sets mqtt connection URI. This API is usually used to overrides the URI configured in esp_mqtt_client_init.

Return ESP_FAIL if URI parse error, ESP_OK on success

Parameters

- client: mqtt client handle
- uri:

esp_err_t **esp_mqtt_client_start**(*esp_mqtt_client_handle_t client*)

Starts mqtt client with already created client handle.

Return ESP_OK on success ESP_ERR_INVALID_ARG on wrong initialization ESP_FAIL on other error

Parameters

- client: mqtt client handle

esp_err_t **esp_mqtt_client_reconnect**(*esp_mqtt_client_handle_t client*)

This api is typically used to force reconnection upon a specific event.

Return ESP_OK on success ESP_FAIL if client is in invalid state

Parameters

- client: mqtt client handle

esp_err_t **esp_mqtt_client_disconnect**(*esp_mqtt_client_handle_t client*)

This api is typically used to force disconnection from the broker.

Return ESP_OK on success

Parameters

- client: mqtt client handle

esp_err_t **esp_mqtt_client_stop**(*esp_mqtt_client_handle_t client*)

Stops mqtt client tasks.

- Notes:
- Cannot be called from the mqtt event handler

Return ESP_OK on success ESP_FAIL if client is in invalid state

Parameters

- `client`: mqtt client handle

int **esp_mqtt_client_subscribe**(*esp_mqtt_client_handle_t client*, **const** char **topic*, int *qos*)

Subscribe the client to defined topic with defined qos.

Notes:

- Client must be connected to send subscribe message
- This API is could be executed from a user task or from a mqtt event callback i.e. internal mqtt task (API is protected by internal mutex, so it might block if a longer data receive operation is in progress.

Return message_id of the subscribe message on success -1 on failure

Parameters

- `client`: mqtt client handle
- `topic`:
- `qos`:

int **esp_mqtt_client_unsubscribe**(*esp_mqtt_client_handle_t client*, **const** char **topic*)

Unsubscribe the client from defined topic.

Notes:

- Client must be connected to send unsubscribe message
- It is thread safe, please refer to `esp_mqtt_client_subscribe` for details

Return message_id of the subscribe message on success -1 on failure

Parameters

- `client`: mqtt client handle
- `topic`:

int **esp_mqtt_client_publish**(*esp_mqtt_client_handle_t client*, **const** char **topic*, **const** char **data*, int *len*, int *qos*, int *retain*)

Client to send a publish message to the broker.

Notes:

- This API might block for several seconds, either due to network timeout (10s) or if publishing payloads longer than internal buffer (due to message fragmentation)
- Client doesn't have to be connected for this API to work, enqueueing the messages with qos>1 (returning -1 for all the qos=0 messages if disconnected). If MQTT_SKIP_PUBLISH_IF_DISCONNECTED is enabled, this API will not attempt to publish when the client is not connected and will always return -1.
- It is thread safe, please refer to `esp_mqtt_client_subscribe` for details

Return message_id of the publish message (for QoS 0 message_id will always be zero) on success. -1 on failure.

Parameters

- `client`: mqtt client handle
- `topic`: topic string
- `data`: payload string (set to NULL, sending empty payload message)
- `len`: data length, if set to 0, length is calculated from payload string
- `qos`: qos of publish message
- `retain`: retain flag

int **esp_mqtt_client_enqueue**(*esp_mqtt_client_handle_t client*, **const** char **topic*, **const** char **data*, int *len*, int *qos*, int *retain*, bool *store*)

Enqueue a message to the outbox, to be sent later. Typically used for messages with qos>0, but could be also used for qos=0 messages if store=true.

This API generates and stores the publish message into the internal outbox and the actual sending to the network is performed in the mqtt-task context (in contrast to the esp_mqtt_client_publish() which sends the pub-

lish message immediately in the user task’ s context). Thus, it could be used as a non blocking version of esp_mqtt_client_publish().

Return message_id if queued successfully, -1 otherwise
Parameters
- `client`: mqtt client handle
- `topic`: topic string
- `data`: payload string (set to NULL, sending empty payload message)
- `len`: data length, if set to 0, length is calculated from payload string
- `qos`: qos of publish message
- `retain`: retain flag
- `store`: if true, all messages are enqueued; otherwise only qos1 and qos 2 are enqueued

esp_err_t **esp_mqtt_client_destroy**(*esp_mqtt_client_handle_t client*)
Destroys the client handle.

Notes:

- Cannot be called from the mqtt event handler

Return ESP_OK
Parameters
- `client`: mqtt client handle

esp_err_t **esp_mqtt_set_config**(*esp_mqtt_client_handle_t client*, **const** *esp_mqtt_client_config_t *config*)
Set configuration structure, typically used when updating the config (i.e. on “before_connect” event.

Return ESP_ERR_NO_MEM if failed to allocate ESP_OK on success
Parameters
- `client`: mqtt client handle
- `config`: mqtt configuration structure

esp_err_t **esp_mqtt_client_register_event**(*esp_mqtt_client_handle_t client*, *esp_mqtt_event_id_t event*, *esp_event_handler_t event_handler*, void **event_handler_arg*)
Registers mqtt event.

Return ESP_ERR_NO_MEM if failed to allocate ESP_OK on success
Parameters
- `client`: mqtt client handle
- `event`: event type
- `event_handler`: handler callback
- `event_handler_arg`: handlers context

int **esp_mqtt_client_get_outbox_size**(*esp_mqtt_client_handle_t client*)
Get outbox size.

Return outbox size
Parameters
- `client`: mqtt client handle

Structures

struct esp_mqtt_error_codes
MQTT error code structure to be passed as a contextual information into ERROR event.

Important: This structure extends `esp_tls_last_error` error structure and is backward compatible with it (so might be down-casted and treated as `esp_tls_last_error` error, but recommended to update applications if used this way previously)

Use this structure directly checking error_type first and then appropriate error code depending on the source of the error:

error_type	related member variables	note
MQTT_ERROR_TYPE_TCP_TRANSPORT	esp_tls_last_esp_err, esp_tls_stack_err, esp_tls_cert_verify_flags, sock_errno	Error reported from

tcp_transport/esp-tls | | MQTT_ERROR_TYPE_CONNECTION_REFUSED | connect_return_code | Internal error reported from MQTT broker on connection |

Public Members

esp_err_t **esp_tls_last_esp_err**
last esp_err code reported from esp-tls component

int **esp_tls_stack_err**
tls specific error code reported from underlying tls stack

int **esp_tls_cert_verify_flags**
tls flags reported from underlying tls stack during certificate verification

esp_mqtt_error_type_t **error_type**
error type referring to the source of the error

esp_mqtt_connect_return_code_t **connect_return_code**
connection refused error code reported from MQTT broker on connection

int **esp_transport_sock_errno**
errno from the underlying socket

struct esp_mqtt_event_t
MQTT event configuration structure

Public Members

esp_mqtt_event_id_t **event_id**
MQTT event type

esp_mqtt_client_handle_t **client**
MQTT client handle for this event

void ***user_context**
User context passed from MQTT client config

char ***data**
Data associated with this event

int **data_len**
Length of the data for this event

int **total_data_len**
Total length of the data (longer data are supplied with multiple events)

int **current_data_offset**
Actual offset for the data associated with this event

char ***topic**
Topic associated with this event

int **topic_len**
Length of the topic for this event associated with this event

int **msg_id**
MQTT messaged id of message

int **session_present**
MQTT session_present flag for connection event

esp_mqtt_error_codes_t ***error_handle**
esp-mqtt error handle including esp-tls errors as well as internal mqtt errors

struct esp_mqtt_client_config_t
MQTT client configuration structure

Public Members

mqtt_event_callback_t **event_handle**
handle for MQTT events as a callback in legacy mode

esp_event_loop_handle_t **event_loop_handle**
handle for MQTT event loop library

const char ***host**
MQTT server domain (ipv4 as string)

const char ***uri**
Complete MQTT broker URI

uint32_t **port**
MQTT server port

const char ***client_id**
default client id is `ESP32_CHIPID%` where CHIPID% are last 3 bytes of MAC address in hex format

const char ***username**
MQTT username

const char ***password**
MQTT password

const char ***lwt_topic**
LWT (Last Will and Testament) message topic (NULL by default)

const char ***lwt_msg**
LWT message (NULL by default)

int **lwt_qos**
LWT message qos

int **lwt_retain**
LWT retained message flag

int **lwt_msg_len**
LWT message length

int **disable_clean_session**
mqtt clean session, default clean_session is true

int **keepalive**
mqtt keepalive, default is 120 seconds

bool **disable_auto_reconnect**
this mqtt client will reconnect to server (when errors/disconnect). Set disable_auto_reconnect=true to disable

void ***user_context**
pass user context to this option, then can receive that context in `event->user_context`

int **task_prio**
MQTT task priority, default is 5, can be changed in `make menuconfig`

int **task_stack**
MQTT task stack size, default is 6144 bytes, can be changed in `make menuconfig`

int **buffer_size**
size of MQTT send/receive buffer, default is 1024 (only receive buffer size if `out_buffer_size` defined)

const char ***cert_pem**
Pointer to certificate data in PEM or DER format for server verify (with SSL), default is NULL, not required to verify the server. PEM-format must have a terminating NULL-character. DER-format requires the length to be passed in cert_len.

size_t **cert_len**
Length of the buffer pointed to by cert_pem. May be 0 for null-terminated pem

const char ***client_cert_pem**
Pointer to certificate data in PEM or DER format for SSL mutual authentication, default is NULL, not required if mutual authentication is not needed. If it is not NULL, also `client_key_pem` has to be provided. PEM-format must have a terminating NULL-character. DER-format requires the length to be passed in client_cert_len.

size_t **client_cert_len**
Length of the buffer pointed to by client_cert_pem. May be 0 for null-terminated pem

const char ***client_key_pem**
Pointer to private key data in PEM or DER format for SSL mutual authentication, default is NULL, not required if mutual authentication is not needed. If it is not NULL, also `client_cert_pem` has to be provided. PEM-format must have a terminating NULL-character. DER-format requires the length to be passed in client_key_len

size_t **client_key_len**
Length of the buffer pointed to by client_key_pem. May be 0 for null-terminated pem

esp_mqtt_transport_t **transport**
overrides URI transport

int **refresh_connection_after_ms**
Refresh connection after this value (in milliseconds)

const struct *psk_key_hint* ***psk_hint_key**
Pointer to PSK struct defined in esp_tls.h to enable PSK authentication (as alternative to certificate verification). If not NULL and server/client certificates are NULL, PSK is enabled

bool **use_global_ca_store**
Use a global ca_store for all the connections in which this bool is set.

int **reconnect_timeout_ms**
Reconnect to the broker after this value in miliseconds if auto reconnect is not disabled (defaults to 10s)

const char ****alpn_protos**
NULL-terminated list of supported application protocols to be used for ALPN

const char ***clientkey_password**
Client key decryption password string

int **clientkey_password_len**
String length of the password pointed to by clientkey_password

esp_mqtt_protocol_ver_t **protocol_ver**
MQTT protocol version used for connection, defaults to value from menuconfig

int **out_buffer_size**
size of MQTT output buffer. If not defined, both output and input buffers have the same size defined as `buffer_size`

bool **skip_cert_common_name_check**
Skip any validation of server certificate CN field, this reduces the security of TLS and makes the mqtt client susceptible to MITM attacks

bool **use_secure_element**
enable secure element for enabling SSL connection

void ***ds_data**
carrier of handle for digital signature parameters

int **network_timeout_ms**
Abort network operation if it is not completed after this value, in milliseconds (defaults to 10s)

bool **disable_keepalive**
Set disable_keepalive=true to turn off keep-alive mechanism, false by default (keepalive is active by

default). Note: setting the config value `keepalive` to `0` doesn't disable keepalive feature, but uses a default keepalive period

Macros

MQTT_ERROR_TYPE_ESP_TLS

MQTT_ERROR_TYPE_TCP_TRANSPORT error type hold all sorts of transport layer errors, including ESP-TLS error, but in the past only the errors from MQTT_ERROR_TYPE_ESP_TLS layer were reported, so the ESP-TLS error type is re-defined here for backward compatibility

Type Definitions

typedef struct esp_mqtt_client ***esp_mqtt_client_handle_t**

typedef struct *esp_mqtt_error_codes* **esp_mqtt_error_codes_t**

MQTT error code structure to be passed as a contextual information into ERROR event.

Important: This structure extends *esp_tls_last_error* error structure and is backward compatible with it (so might be down-casted and treated as *esp_tls_last_error* error, but recommended to update applications if used this way previously)

Use this structure directly checking error_type first and then appropriate error code depending on the source of the error:

| error_type | related member variables | note | | MQTT_ERROR_TYPE_TCP_TRANSPORT | esp_tls_last_esp_err, esp_tls_stack_err, esp_tls_cert_verify_flags, sock_errno | Error reported from tcp_transport/esp-tls | | MQTT_ERROR_TYPE_CONNECTION_REFUSED | connect_return_code | Internal error reported from MQTT broker on connection |

typedef *esp_mqtt_event_t* ***esp_mqtt_event_handle_t**

typedef *esp_err_t* (***mqtt_event_callback_t**) (*esp_mqtt_event_handle_t* event)

Enumerations

enum esp_mqtt_event_id_t

MQTT event types.

User event handler receives context data in *esp_mqtt_event_t* structure with

- `user_context` - user data from *esp_mqtt_client_config_t*
- `client` - mqtt client handle
- various other data depending on event type

Values:

MQTT_EVENT_ANY = -1

MQTT_EVENT_ERROR = 0

on error event, additional context: connection return code, error handle from *esp_tls* (if supported)

MQTT_EVENT_CONNECTED

connected event, additional context: session_present flag

MQTT_EVENT_DISCONNECTED

disconnected event

MQTT_EVENT_SUBSCRIBED

subscribed event, additional context: msg_id

MQTT_EVENT_UNSUBSCRIBED

unsubscribed event

MQTT_EVENT_PUBLISHED

published event, additional context: msg_id

MQTT_EVENT_DATA

data event, additional context:

- msg_id message id
- topic pointer to the received topic
- topic_len length of the topic
- data pointer to the received data
- data_len length of the data for this event
- current_data_offset offset of the current data for this event
- total_data_len total length of the data received Note: Multiple MQTT_EVENT_DATA could be fired for one message, if it is longer than internal buffer. In that case only first event contains topic pointer and length, other contain data only with current data length and current data offset updating.

MQTT_EVENT_BEFORE_CONNECT
The event occurs before connecting

MQTT_EVENT_DELETED
Notification on delete of one message from the internal outbox, if the message couldn't have been sent and acknowledged before expiring defined in OUTBOX_EXPIRED_TIMEOUT_MS. (events are not posted upon deletion of successfully acknowledged messages)

- This event id is posted only if MQTT_REPORT_DELETED_MESSAGES==1
- Additional context: msg_id (id of the deleted message).

enum esp_mqtt_connect_return_code_t
MQTT connection error codes propagated via ERROR event

Values:

MQTT_CONNECTION_ACCEPTED = 0
Connection accepted

MQTT_CONNECTION_REFUSE_PROTOCOL
MQTT connection refused reason: Wrong protocol

MQTT_CONNECTION_REFUSE_ID_REJECTED
MQTT connection refused reason: ID rejected

MQTT_CONNECTION_REFUSE_SERVER_UNAVAILABLE
MQTT connection refused reason: Server unavailable

MQTT_CONNECTION_REFUSE_BAD_USERNAME
MQTT connection refused reason: Wrong user

MQTT_CONNECTION_REFUSE_NOT_AUTHORIZED
MQTT connection refused reason: Wrong username or password

enum esp_mqtt_error_type_t
MQTT connection error codes propagated via ERROR event

Values:

MQTT_ERROR_TYPE_NONE = 0

MQTT_ERROR_TYPE_TCP_TRANSPORT

MQTT_ERROR_TYPE_CONNECTION_REFUSED

enum esp_mqtt_transport_t
Values:

MQTT_TRANSPORT_UNKNOWN = 0x0

MQTT_TRANSPORT_OVER_TCP
MQTT over TCP, using scheme: `mqtt`

MQTT_TRANSPORT_OVER_SSL
MQTT over SSL, using scheme: `mqtts`

MQTT_TRANSPORT_OVER_WS
MQTT over Websocket, using scheme:: `ws`

MQTT_TRANSPORT_OVER_WSS
MQTT over Websocket Secure, using scheme: `wss`

enum esp_mqtt_protocol_ver_t
MQTT protocol version used for connection

Values:

MQTT_PROTOCOL_UNDEFINED = 0

MQTT_PROTOCOL_V_3_1

MQTT_PROTOCOL_V_3_1_1

2.3.10 ESP-Modbus

Overview

The Modbus serial communication protocol is de facto standard protocol widely used to connect industrial electronic devices. Modbus allows communication among many devices connected to the same network, for example, a system that measures temperature and humidity and communicates the results to a computer. The Modbus protocol uses several types of data: Holding Registers, Input Registers, Coils (single bit output), Discrete Inputs. Versions of the Modbus protocol exist for serial port and for Ethernet and other protocols that support the Internet protocol suite. There are many variants of Modbus protocols, some of them are:

- `Modbus RTU` —This is used in serial communication and makes use of a compact, binary representation of the data for protocol communication. The RTU format follows the commands/data with a cyclic redundancy check checksum as an error check mechanism to ensure the reliability of data. Modbus RTU is the most common implementation available for Modbus. A Modbus RTU message must be transmitted continuously without inter-character hesitations. Modbus messages are framed (separated) by idle (silent) periods. The RS-485 interface communication is usually used for this type.
- `Modbus ASCII` —This is used in serial communication and makes use of ASCII characters for protocol communication. The ASCII format uses a longitudinal redundancy check checksum. Modbus ASCII messages are framed by leading colon (":") and trailing newline (CR/LF).
- `Modbus TCP/IP or Modbus TCP` —This is a Modbus variant used for communications over TCP/IP networks, connecting over port 502. It does not require a checksum calculation, as lower layers already provide checksum protection.

Modbus common interface API overview The API functions below provide common functionality to setup Modbus stack for slave and master implementation accordingly. ISP-IDF supports Modbus serial slave and master protocol stacks and provides modbus_controller interface API to interact with user application.

esp_err_t **mbc_slave_init**(mb_port_type_t *port_type*, void **handler*)
Initialize Modbus controller and stack.

Return
- ESP_OK Success
- ESP_ERR_NO_MEM Parameter error

Parameters
- `[out] handler`: handler(pointer) to master data structure
- `[in] port_type`: type of stack

esp_err_t **mbc_master_init**(mb_port_type_t *port_type*, void **handler*)
Initialize Modbus controller and stack.

Return
- ESP_OK Success
- ESP_ERR_NO_MEM Parameter error

Parameters
- `[out] handler`: handler(pointer) to master data structure
- `[in] port_type`: the type of port

The function initializes the Modbus controller interface and its active context (tasks, RTOS objects and other resources).

esp_err_t **mbc_slave_setup**(void *comm_info)

Set Modbus communication parameters for the controller.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Incorrect parameter data

Parameters

- `comm_info`: Communication parameters structure.

esp_err_t **mbc_master_setup**(void *comm_info)

Set Modbus communication parameters for the controller.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Incorrect parameter data

Parameters

- `comm_info`: Communication parameters structure.

The function is used to setup communication parameters of the Modbus stack. See the Modbus controller API documentation for more information.

mbc_slave_set_descriptor(): Initialization of slave descriptor.

mbc_master_set_descriptor(): Initialization of master descriptor.

The Modbus stack uses parameter description tables (descriptors) for communication. These are different for master and slave implementation of stack and should be assigned by the API call before start of communication.

esp_err_t **mbc_slave_start**(void)

Start Modbus communication stack.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Modbus stack start error

esp_err_t **mbc_master_start**(void)

Start Modbus communication stack.

Return

- ESP_OK Success
- ESP_ERR_INVALID_ARG Modbus stack start error

Modbus controller start function. Starts stack and interface and allows communication.

esp_err_t **mbc_slave_destroy**(void)

Destroy Modbus controller and stack.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE Parameter error

esp_err_t **mbc_master_destroy**(void)

Destroy Modbus controller and stack.

Return

- ESP_OK Success
- ESP_ERR_INVALID_STATE Parameter error

This function stops Modbus communication stack and destroys controller interface.

There are some configurable parameters of modbus_controller interface and Modbus stack that can be configured using KConfig values in "Modbus configuration" menu. The most important option in KConfig menu is "Selection of Modbus stack support mode" that allows to select master or slave stack for implementation. See the examples for more information about how to use these API functions.

Modbus serial slave interface API overview The slave stack requires the user defined structures which represent Modbus parameters accessed by stack. These structures should be prepared by user and be assigned to the modbus_controller interface using *mbc_slave_set_descriptor()* API call before start of communication. The interface API functions below are used for Modbus slave application:

esp_err_t **mbc_slave_set_descriptor**(mb_register_area_descriptor_t *descr_data*)

Set Modbus area descriptor.

Return

- ESP_OK: The appropriate descriptor is set
- ESP_ERR_INVALID_ARG: The argument is incorrect

Parameters

- `descr_data`: Modbus registers area descriptor structure

The function initializes Modbus communication descriptors for each type of Modbus register area (Holding Registers, Input Registers, Coils (single bit output), Discrete Inputs). Once areas are initialized and the *mbc_slave_start()* API is called the Modbus stack can access the data in user data structures by request from master. See the `mb_register_area_descriptor_t` and example for more information.

mb_event_group_t **mbc_slave_check_event**(mb_event_group_t *group*)

Wait for specific event on parameter change.

Return

- mb_event_group_t event bits triggered

Parameters

- `group`: Group event bit mask to wait for change

The blocking call to function waits for event specified in the input parameter as event mask. Once master access the parameter and event mask matches the parameter the application task will be unblocked and function will return ESP_OK. See the `mb_event_group_t` for more information about Modbus event masks.

esp_err_t **mbc_slave_get_param_info**(mb_param_info_t **reg_info*, uint32_t *timeout*)

Get parameter information.

Return

- ESP_OK Success
- ESP_ERR_TIMEOUT Can not get data from parameter queue or queue overflow

Parameters

- `[out] reg_info`: parameter info structure
- `timeout`: Timeout in milliseconds to read information from parameter queue

The function gets information about accessed parameters from modbus controller event queue. The KConfig 'CONFIG_FMB_CONTROLLER_NOTIFY_QUEUE_SIZE' key can be used to configure the notification queue size. The timeout parameter allows to specify timeout for waiting notification. The `mb_param_info_t` structure contain information about accessed parameter.

Modbus serial master interface API overview The Modbus master implementation requires parameter description table be defined before start of stack. This table describes characteristics (physical parameters like temperature, humidity, etc.) and links them to Modbus registers in specific slave device in the Modbus segment. The table has to be assigned to the modbus_controller interface using *mbc_master_set_descriptor()* API call before start of communication.

Below are the interface API functions that are used to setup and use Modbus master stack from user application and can be executed in next order:

esp_err_t **mbc_master_set_descriptor**(**const** mb_parameter_descriptor_t **descriptor*, **const** uint16_t *num_elements*)

Assign parameter description table for Modbus controller interface.

Return

- esp_err_t ESP_OK - set descriptor successfully
- esp_err_t ESP_ERR_INVALID_ARG - invalid argument in function call

Parameters

- `[in] descriptor`: pointer to parameter description table

- `num_elements`: number of elements in the table

Assigns parameter description table for Modbus controller interface. The table has to be prepared by user according to particular

esp_err_t **mbc_master_send_request**(mb_param_request_t *_request_, void *_data_ptr_)

Send data request as defined in parameter request, waits response from slave and returns status of command execution. This function provides standard way for read/write access to Modbus devices in the network.

Return

- esp_err_t ESP_OK - request was successful
- esp_err_t ESP_ERR_INVALID_ARG - invalid argument of function
- esp_err_t ESP_ERR_INVALID_RESPONSE - an invalid response from slave
- esp_err_t ESP_ERR_TIMEOUT - operation timeout or no response from slave
- esp_err_t ESP_ERR_NOT_SUPPORTED - the request command is not supported by slave
- esp_err_t ESP_FAIL - slave returned an exception or other failure

Parameters

- `[in] request`: pointer to request structure of type mb_param_request_t
- `[in] data_ptr`: pointer to data buffer to send or received data (dependent of command field in request)

This function sends data request as defined in parameter request, waits response from corresponded slave and returns status of command execution. This function provides a standard way for read/write access to Modbus devices in the network.

esp_err_t **mbc_master_get_cid_info**(uint16_t *cid*, **const** mb_parameter_descriptor_t **_param_info_)

Get information about supported characteristic defined as cid. Uses parameter description table to get this information. The function will check if characteristic defined as a cid parameter is supported and returns its description in param_info. Returns ESP_ERR_NOT_FOUND if characteristic is not supported.

Return

- esp_err_t ESP_OK - request was successful and buffer contains the supported characteristic name
- esp_err_t ESP_ERR_INVALID_ARG - invalid argument of function
- esp_err_t ESP_ERR_NOT_FOUND - the characteristic (cid) not found
- esp_err_t ESP_FAIL - unknown error during lookup table processing

Parameters

- `[in] cid`: characteristic id
- `param_info`: pointer to pointer of characteristic data.

The function gets information about supported characteristic defined as cid. It will check if characteristic is supported and returns its description.

esp_err_t **mbc_master_get_parameter**(uint16_t *cid*, char *_name_, uint8_t *_value_, uint8_t *_type_)

Read parameter from modbus slave device whose name is defined by name and has cid. The additional data for request is taken from parameter description (lookup) table.

Return

- esp_err_t ESP_OK - request was successful and value buffer contains representation of actual parameter data from slave
- esp_err_t ESP_ERR_INVALID_ARG - invalid argument of function or parameter descriptor
- esp_err_t ESP_ERR_INVALID_RESPONSE - an invalid response from slave
- esp_err_t ESP_ERR_INVALID_STATE - invalid state during data processing or allocation failure
- esp_err_t ESP_ERR_TIMEOUT - operation timed out and no response from slave
- esp_err_t ESP_ERR_NOT_SUPPORTED - the request command is not supported by slave
- esp_err_t ESP_ERR_NOT_FOUND - the parameter is not found in the parameter description table
- esp_err_t ESP_FAIL - slave returned an exception or other failure

Parameters

- `[in] cid`: id of the characteristic for parameter
- `[in] name`: pointer into string name (key) of parameter (null terminated)
- `[out] value`: pointer to data buffer of parameter
- `[out] type`: parameter type associated with the name returned from parameter description table.

The function reads data of characteristic defined in parameters from Modbus slave device and returns its data. The additional data for request is taken from parameter description table.

esp_err_t **mbc_master_set_parameter**(uint16_t *cid*, char **name*, uint8_t **value*, uint8_t **type*)

Set characteristic's value defined as a name and cid parameter. The additional data for cid parameter request is taken from master parameter lookup table.

Return

- esp_err_t ESP_OK - request was successful and value was saved in the slave device registers
- esp_err_t ESP_ERR_INVALID_ARG - invalid argument of function or parameter descriptor
- esp_err_t ESP_ERR_INVALID_RESPONSE - an invalid response from slave during processing of parameter
- esp_err_t ESP_ERR_INVALID_STATE - invalid state during data processing or allocation failure
- esp_err_t ESP_ERR_TIMEOUT - operation timed out and no response from slave
- esp_err_t ESP_ERR_NOT_SUPPORTED - the request command is not supported by slave
- esp_err_t ESP_FAIL - slave returned an exception or other failure

Parameters

- `[in] cid`: id of the characteristic for parameter
- `[in] name`: pointer into string name (key) of parameter (null terminated)
- `[out] value`: pointer to data buffer of parameter (actual representation of json value field in binary form)
- `[out] type`: pointer to parameter type associated with the name returned from parameter lookup table.

The function writes characteristic's value defined as a name and cid parameter in corresponded slave device. The additional data for parameter request is taken from master parameter description table.

Application Example

The examples below use the FreeModbus library port for serial slave and master implementation accordingly. The selection of stack is performed through KConfig menu "Selection of Modbus stack support mode" and related configuration keys.

protocols/modbus/serial/mb_slave

protocols/modbus/serial/mb_master

Please refer to the specific example README.md for details.

2.3.11 ESP Local Control

Overview

ESP Local Control (**esp_local_ctrl**) component in ESP-IDF provides capability to control an ESP device over Wi-Fi + HTTPS or BLE. It provides access to application defined **properties** that are available for reading / writing via a set of configurable handlers.

Initialization of the **esp_local_ctrl** service over BLE transport is performed as follows:

```
esp_local_ctrl_config_t config = {
    .transport = ESP_LOCAL_CTRL_TRANSPORT_BLE,
    .transport_config = {
        .ble = & (protocomm_ble_config_t) {
            .device_name  = SERVICE_NAME,
            .service_uuid = {
                /* LSB <---------------------------------------
                * ---------------------------------------> MSB */
                0x21, 0xd5, 0x3b, 0x8d, 0xbd, 0x75, 0x68, 0x8a,
                0xb4, 0x42, 0xeb, 0x31, 0x4a, 0x1e, 0x98, 0x3d
            }
```

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```
        }
    },
    .proto_sec = {
        .version = PROTOCOM_SEC0,
        .custom_handle = NULL,
        .pop = NULL,
    },
    .handlers = {
        /* User defined handler functions */
        .get_prop_values = get_property_values,
        .set_prop_values = set_property_values,
        .usr_ctx         = NULL,
        .usr_ctx_free_fn = NULL
    },
    /* Maximum number of properties that may be set */
    .max_properties = 10
};

/* Start esp_local_ctrl service */
ESP_ERROR_CHECK(esp_local_ctrl_start(&config));
```

Similarly for HTTPS transport:

```
/* Set the configuration */
httpd_ssl_config_t https_conf = HTTPD_SSL_CONFIG_DEFAULT();

/* Load server certificate */
extern const unsigned char cacert_pem_start[] asm("_binary_cacert_pem_
↪start");
extern const unsigned char cacert_pem_end[]   asm("_binary_cacert_pem_end
↪");
https_conf.cacert_pem = cacert_pem_start;
https_conf.cacert_len = cacert_pem_end - cacert_pem_start;

/* Load server private key */
extern const unsigned char prvtkey_pem_start[] asm("_binary_prvtkey_pem_
↪start");
extern const unsigned char prvtkey_pem_end[]   asm("_binary_prvtkey_pem_
↪end");
https_conf.prvtkey_pem = prvtkey_pem_start;
https_conf.prvtkey_len = prvtkey_pem_end - prvtkey_pem_start;

esp_local_ctrl_config_t config = {
    .transport = ESP_LOCAL_CTRL_TRANSPORT_HTTPD,
    .transport_config = {
        .httpd = &https_conf
    },
    .proto_sec = {
        .version = PROTOCOM_SEC0,
        .custom_handle = NULL,
        .pop = NULL,
    },
    .handlers = {
        /* User defined handler functions */
        .get_prop_values = get_property_values,
        .set_prop_values = set_property_values,
        .usr_ctx         = NULL,
        .usr_ctx_free_fn = NULL
    },
    /* Maximum number of properties that may be set */
    .max_properties = 10
```

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```
};

/* Start esp_local_ctrl service */
ESP_ERROR_CHECK(esp_local_ctrl_start(&config));
```

You may set security for transport in ESP local control using following options:

1. *PROTOCOM_SEC1*: specifies that end to end encryption is used.
2. *PROTOCOM_SEC0*: specifies that data will be exchanged as a plain text.
3. *PROTOCOM_SEC_CUSTOM*: you can define your own security requirement. Please note that you will also have to provide *custom_handle* of type *protocomm_security_t* * in this context.

Creating a property

Now that we know how to start the **esp_local_ctrl** service, let's add a property to it. Each property must have a unique *name* (string), a *type* (e.g. enum), *flags* (bit fields) and *size*.

The *size* is to be kept 0, if we want our property value to be of variable length (e.g. if its a string or bytestream). For fixed length property value data-types, like int, float, etc., setting the *size* field to the right value, helps **esp_local_ctrl** to perform internal checks on arguments received with write requests.

The interpretation of *type* and *flags* fields is totally upto the application, hence they may be used as enumerations, bit-fields, or even simple integers. One way is to use *type* values to classify properties, while *flags* to specify characteristics of a property.

Here is an example property which is to function as a timestamp. It is assumed that the application defines *TYPE_TIMESTAMP* and *READONLY*, which are used for setting the *type* and *flags* fields here.

```
/* Create a timestamp property */
esp_local_ctrl_prop_t timestamp = {
    .name        = "timestamp",
    .type        = TYPE_TIMESTAMP,
    .size        = sizeof(int32_t),
    .flags       = READONLY,
    .ctx         = func_get_time,
    .ctx_free_fn = NULL
};

/* Now register the property */
esp_local_ctrl_add_property(&timestamp);
```

Also notice that there is a ctx field, which is set to point to some custom *func_get_time()*. This can be used inside the property get / set handlers to retrieve timestamp.

Here is an example of *get_prop_values()* handler, which is used for retrieving the timestamp.

```
static esp_err_t get_property_values(size_t props_count,
                                     const esp_local_ctrl_prop_t *props,
                                     esp_local_ctrl_prop_val_t *prop_
↪values,
                                     void *usr_ctx)
{
    for (uint32_t i = 0; i < props_count; i++) {
        ESP_LOGI(TAG, "Reading %s", props[i].name);
        if (props[i].type == TYPE_TIMESTAMP) {
            /* Obtain the timer function from ctx */
            int32_t (*func_get_time)(void) = props[i].ctx;

            /* Use static variable for saving the value.
             * This is essential because the value has to be
             * valid even after this function returns.
```

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```
            * Alternative is to use dynamic allocation
            * and set the free_fn field */
            static int32_t ts = func_get_time();
            prop_values[i].data = &ts;
        }
    }
    return ESP_OK;
}
```

Here is an example of *set_prop_values()* handler. Notice how we restrict from writing to read-only properties.

```
static esp_err_t set_property_values(size_t props_count,
                                     const esp_local_ctrl_prop_t *props,
                                     const esp_local_ctrl_prop_val_t␣
↪*prop_values,
                                     void *usr_ctx)
{
    for (uint32_t i = 0; i < props_count; i++) {
        if (props[i].flags & READONLY) {
            ESP_LOGE(TAG, "Cannot write to read-only property %s",␣
↪props[i].name);
            return ESP_ERR_INVALID_ARG;
        } else {
            ESP_LOGI(TAG, "Setting %s", props[i].name);

            /* For keeping it simple, lets only log the incoming data */
            ESP_LOG_BUFFER_HEX_LEVEL(TAG, prop_values[i].data,
                                     prop_values[i].size, ESP_LOG_INFO);
        }
    }
    return ESP_OK;
}
```

For complete example see protocols/esp_local_ctrl

Client Side Implementation

The client side implementation will have establish a protocomm session with the device first, over the supported mode of transport, and then send and receive protobuf messages understood by the **esp_local_ctrl** service. The service will translate these messages into requests and then call the appropriate handlers (set / get). Then, the generated response for each handler is again packed into a protobuf message and transmitted back to the client.

See below the various protobuf messages understood by the **esp_local_ctrl** service:

1. *get_prop_count* : This should simply return the total number of properties supported by the service
2. *get_prop_values* : This accepts an array of indices and should return the information (name, type, flags) and values of the properties corresponding to those indices
3. *set_prop_values* : This accepts an array of indices and an array of new values, which are used for setting the values of the properties corresponding to the indices

Note that indices may or may not be the same for a property, across multiple sessions. Therefore, the client must only use the names of the properties to uniquely identify them. So, every time a new session is established, the client should first call *get_prop_count* and then *get_prop_values*, hence form an index to name mapping for all properties. Now when calling *set_prop_values* for a set of properties, it must first convert the names to indexes, using the created mapping. As emphasized earlier, the client must refresh the index to name mapping every time a new session is established with the same device.

The various protocomm endpoints provided by **esp_local_ctrl** are listed below:

表 4: Endpoints provided by ESP Local Control

Endpoint Name (BLE + GATT Server)	URI (HTTPS Server + mDNS)	Description
esp_local_ctrl/version	https://<mdns-hostname>.local/esp_local_ctrl/version	Endpoint used for retrieving version string
esp_local_ctrl/control	https://<mdns-hostname>.local/esp_local_ctrl/control	Endpoint used for sending / receiving control messages

API Reference

Header File

- esp_local_ctrl/include/esp_local_ctrl.h

Functions

const *esp_local_ctrl_transport_t* ***esp_local_ctrl_get_transport_ble**(void)

Function for obtaining BLE transport mode.

const *esp_local_ctrl_transport_t* ***esp_local_ctrl_get_transport_httpd**(void)

Function for obtaining HTTPD transport mode.

esp_err_t **esp_local_ctrl_start**(**const** *esp_local_ctrl_config_t* **config*)

Start local control service.

Return

- ESP_OK : Success
- ESP_FAIL : Failure

Parameters

- `[in] config`: Pointer to configuration structure

esp_err_t **esp_local_ctrl_stop**(void)

Stop local control service.

esp_err_t **esp_local_ctrl_add_property**(**const** *esp_local_ctrl_prop_t* **prop*)

Add a new property.

This adds a new property and allocates internal resources for it. The total number of properties that could be added is limited by configuration option `max_properties`

Return

- ESP_OK : Success
- ESP_FAIL : Failure

Parameters

- `[in] prop`: Property description structure

esp_err_t **esp_local_ctrl_remove_property**(**const** char **name*)

Remove a property.

This finds a property by name, and releases the internal resources which are associated with it.

Return

- ESP_OK : Success
- ESP_ERR_NOT_FOUND : Failure

Parameters

- `[in] name`: Name of the property to remove

const *esp_local_ctrl_prop_t* ***esp_local_ctrl_get_property**(**const** char **name*)

Get property description structure by name.

This API may be used to get a property' s context structure `esp_local_ctrl_prop_t` when its name is known

Return
- Pointer to property
- NULL if not found

Parameters
- `[in] name`: Name of the property to find

esp_err_t **esp_local_ctrl_set_handler**(**const** char **ep_name*, *protocomm_req_handler_t handler*, void **user_ctx*)

Register protocomm handler for a custom endpoint.

This API can be called by the application to register a protocomm handler for an endpoint after the local control service has started.

Note In case of BLE transport the names and uuids of all custom endpoints must be provided beforehand as a part of the `protocomm_ble_config_t` structure set in `esp_local_ctrl_config_t`, and passed to `esp_local_ctrl_start()`.

Return
- ESP_OK : Success
- ESP_FAIL : Failure

Parameters
- `[in] ep_name`: Name of the endpoint
- `[in] handler`: Endpoint handler function
- `[in] user_ctx`: User data

Unions

union esp_local_ctrl_transport_config_t

#include <esp_local_ctrl.h> Transport mode (BLE / HTTPD) configuration.

Public Members

esp_local_ctrl_transport_config_ble_t ***ble**

This is same as `protocomm_ble_config_t`. See `protocomm_ble.h` for available configuration parameters.

esp_local_ctrl_transport_config_httpd_t ***httpd**

This is same as `httpd_ssl_config_t`. See `esp_https_server.h` for available configuration parameters.

Structures

struct esp_local_ctrl_prop

Property description data structure, which is to be populated and passed to the `esp_local_ctrl_add_property()` function.

Once a property is added, its structure is available for read-only access inside `get_prop_values()` and `set_prop_values()` handlers.

Public Members

char ***name**

Unique name of property

uint32_t **type**

Type of property. This may be set to application defined enums

size_t **size**

Size of the property value, which:

- if zero, the property can have values of variable size
- if non-zero, the property can have values of fixed size only, therefore, checks are performed internally by esp_local_ctrl when setting the value of such a property

uint32_t **flags**
Flags set for this property. This could be a bit field. A flag may indicate property behavior, e.g. read-only / constant

void ***ctx**
Pointer to some context data relevant for this property. This will be available for use inside the `get_prop_values` and `set_prop_values` handlers as a part of this property structure. When set, this is valid throughout the lifetime of a property, till either the property is removed or the esp_local_ctrl service is stopped.

void (***ctx_free_fn**) (void *ctx)
Function used by esp_local_ctrl to internally free the property context when `esp_local_ctrl_remove_property()` or `esp_local_ctrl_stop()` is called.

struct esp_local_ctrl_prop_val
Property value data structure. This gets passed to the `get_prop_values()` and `set_prop_values()` handlers for the purpose of retrieving or setting the present value of a property.

Public Members

void ***data**
Pointer to memory holding property value

size_t **size**
Size of property value

void (***free_fn**) (void *data)
This may be set by the application in `get_prop_values()` handler to tell `esp_local_ctrl` to call this function on the data pointer above, for freeing its resources after sending the `get_prop_values` response.

struct esp_local_ctrl_handlers
Handlers for receiving and responding to local control commands for getting and setting properties.

Public Members

esp_err_t (***get_prop_values**) (size_t props_count, **const** *esp_local_ctrl_prop_t* props[], *esp_local_ctrl_prop_val_t* prop_values[], void *usr_ctx)
Handler function to be implemented for retrieving current values of properties.

Note If any of the properties have fixed sizes, the size field of corresponding element in `prop_values` need to be set

Return Returning different error codes will convey the corresponding protocol level errors to the client :

- ESP_OK : Success
- ESP_ERR_INVALID_ARG : InvalidArgument
- ESP_ERR_INVALID_STATE : InvalidProto
- All other error codes : InternalError

Parameters

- `[in] props_count`: Total elements in the props array
- `[in] props`: Array of properties, the current values for which have been requested by the client
- `[out] prop_values`: Array of empty property values, the elements of which need to be populated with the current values of those properties specified by props argument
- `[in] usr_ctx`: This provides value of the `usr_ctx` field of `esp_local_ctrl_handlers_t` structure

esp_err_t (***set_prop_values**) (size_t props_count, **const** *esp_local_ctrl_prop_t* props[], **const** *esp_local_ctrl_prop_val_t* prop_values[], void *usr_ctx)

Handler function to be implemented for changing values of properties.

Note If any of the properties have variable sizes, the size field of the corresponding element in `prop_values` must be checked explicitly before making any assumptions on the size.

Return Returning different error codes will convey the corresponding protocol level errors to the client :

- ESP_OK : Success
- ESP_ERR_INVALID_ARG : InvalidArgument
- ESP_ERR_INVALID_STATE : InvalidProto
- All other error codes : InternalError

Parameters

- `[in] props_count`: Total elements in the props array
- `[in] props`: Array of properties, the values for which the client requests to change
- `[in] prop_values`: Array of property values, the elements of which need to be used for updating those properties specified by props argument
- `[in] usr_ctx`: This provides value of the `usr_ctx` field of `esp_local_ctrl_handlers_t` structure

void ***usr_ctx**

Context pointer to be passed to above handler functions upon invocation. This is different from the property level context, as this is valid throughout the lifetime of the `esp_local_ctrl` service, and freed only when the service is stopped.

void (***usr_ctx_free_fn**) (void *usr_ctx)

Pointer to function which will be internally invoked on `usr_ctx` for freeing the context resources when `esp_local_ctrl_stop()` is called.

struct esp_local_ctrl_proto_sec_cfg

Protocom security configs

Public Members

esp_local_ctrl_proto_sec_t **version**

This sets protocom security version, sec0/sec1 or custom If custom, user must provide handle via `proto_sec_custom_handle` below

void ***custom_handle**

Custom security handle if security is set custom via `proto_sec` above This handle must follow `protocomm_security_t` signature

void ***pop**

Proof of possession to be used for local control. Could be NULL.

struct esp_local_ctrl_config

Configuration structure to pass to `esp_local_ctrl_start()`

Public Members

const *esp_local_ctrl_transport_t* ***transport**

Transport layer over which service will be provided

esp_local_ctrl_transport_config_t **transport_config**

Transport layer over which service will be provided

esp_local_ctrl_proto_sec_cfg_t **proto_sec**

Security version and POP

esp_local_ctrl_handlers_t **handlers**

Register handlers for responding to get/set requests on properties

size_t **max_properties**
This limits the number of properties that are available at a time

Macros

ESP_LOCAL_CTRL_TRANSPORT_BLE
ESP_LOCAL_CTRL_TRANSPORT_HTTPD

Type Definitions

typedef struct *esp_local_ctrl_prop* **esp_local_ctrl_prop_t**
Property description data structure, which is to be populated and passed to the `esp_local_ctrl_add_property()` function.

Once a property is added, its structure is available for read-only access inside `get_prop_values()` and `set_prop_values()` handlers.

typedef struct *esp_local_ctrl_prop_val* **esp_local_ctrl_prop_val_t**
Property value data structure. This gets passed to the `get_prop_values()` and `set_prop_values()` handlers for the purpose of retrieving or setting the present value of a property.

typedef struct *esp_local_ctrl_handlers* **esp_local_ctrl_handlers_t**
Handlers for receiving and responding to local control commands for getting and setting properties.

typedef struct esp_local_ctrl_transport **esp_local_ctrl_transport_t**
Transport mode (BLE / HTTPD) over which the service will be provided.

This is forward declaration of a private structure, implemented internally by `esp_local_ctrl`.

typedef struct *protocomm_ble_config* **esp_local_ctrl_transport_config_ble_t**
Configuration for transport mode BLE.

This is a forward declaration for `protocomm_ble_config_t`. To use this, application must set CONFIG_BT_BLUEDROID_ENABLED and include `protocomm_ble.h`.

typedef struct *httpd_ssl_config* **esp_local_ctrl_transport_config_httpd_t**
Configuration for transport mode HTTPD.

This is a forward declaration for `httpd_ssl_config_t`. To use this, application must set CONFIG_ESP_HTTPS_SERVER_ENABLE and include `esp_https_server.h`

typedef enum *esp_local_ctrl_proto_sec* **esp_local_ctrl_proto_sec_t**
Security types for esp_local_control.

typedef struct *esp_local_ctrl_proto_sec_cfg* **esp_local_ctrl_proto_sec_cfg_t**
Protocom security configs

typedef struct *esp_local_ctrl_config* **esp_local_ctrl_config_t**
Configuration structure to pass to `esp_local_ctrl_start()`

Enumerations

enum esp_local_ctrl_proto_sec
Security types for esp_local_control.

Values:

PROTOCOM_SEC0 = 0

PROTOCOM_SEC1

PROTOCOM_SEC_CUSTOM

2.3.12 ESP x509 Certificate Bundle

Overview

The ESP x509 Certificate Bundle API provides an easy way to include a bundle of custom x509 root certificates for TLS server verification.

注解： The bundle is currently not available when using WolfSSL.

The bundle comes with the complete list of root certificates from Mozilla's NSS root certificate store. Using the gen_crt_bundle.py python utility the certificates' subject name and public key are stored in a file and embedded in the ESP32-S2 binary.

When generating the bundle you may choose between:

- The full root certificate bundle from Mozilla, containing more than 130 certificates. The current bundle was updated Wed Jan 23 04:12:09 2019 GMT.
- A pre-selected filter list of the name of the most commonly used root certificates, reducing the amount of certificates to around 35 while still having around 90 % coverage according to market share statistics.

In addition it is possible to specify a path to a certificate file or a directory containing certificates which then will be added to the generated bundle.

注解： Trusting all root certificates means the list will have to be updated if any of the certificates are retracted. This includes removing them from *cacrt_all.pem*.

Configuration

Most configuration is done through menuconfig. Make and CMake will generate the bundle according to the configuration and embed it.

- *CONFIG_MBEDTLS_CERTIFICATE_BUNDLE*: automatically build and attach the bundle.
- *CONFIG_MBEDTLS_DEFAULT_CERTIFICATE_BUNDLE*: decide which certificates to include from the complete root list.
- *CONFIG_MBEDTLS_CUSTOM_CERTIFICATE_BUNDLE_PATH*: specify the path of any additional certificates to embed in the bundle.

To enable the bundle when using ESP-TLS simply pass the function pointer to the bundle attach function:

```
esp_tls_cfg_t cfg = {
     .crt_bundle_attach = esp_crt_bundle_attach,
};
```

This is done to avoid embedding the certificate bundle unless activated by the user.

If using mbedTLS directly then the bundle may be activated by directly calling the attach function during the setup process:

```
mbedtls_ssl_config conf;
mbedtls_ssl_config_init(&conf);

esp_crt_bundle_attach(&conf);
```

Generating the List of Root Certificates

The list of root certificates comes from Mozilla's NSS root certificate store, which can be found here The list can be downloaded and created by running the script `mk-ca-bundle.pl` that is distributed as a part of curl. Another

alternative would be to download the finished list directly from the curl website: CA certificates extracted from Mozilla

The common certificates bundle were made by selecting the authorities with a market share of more than 1 % from w3tech's SSL Survey. These authorities were then used to pick the names of the certificates for the filter list, *cmn_crt_authorities.csv*, from this list provided by Mozilla.

Updating the Certificate Bundle

The bundle is embedded into the app and can be updated along with the app by an OTA update. If you want to include a more up-to-date bundle than the bundle currently included in IDF, then the certificate list can be downloaded from Mozilla as described in *Updating the Certificate Bundle*.

Application Example

Simple HTTPS example that uses ESP-TLS to establish a secure socket connection using the certificate bundle with two custom certificates added for verification: protocols/https_x509_bundle.

HTTPS example that uses ESP-TLS and the default bundle: protocols/https_request.

HTTPS example that uses mbedTLS and the default bundle: protocols/https_mbedtls.

API Reference

Header File

- mbedtls/esp_crt_bundle/include/esp_crt_bundle.h

Functions

esp_err_t **esp_crt_bundle_attach** (void *conf)

Attach and enable use of a bundle for certificate verification.

Attach and enable use of a bundle for certificate verification through a verification callback. If no specific bundle has been set through esp_crt_bundle_set() it will default to the bundle defined in menuconfig and embedded in the binary.

Return

- ESP_OK if adding certificates was successful.
- Other if an error occured or an action must be taken by the calling process.

Parameters

- `[in] conf`: The config struct for the SSL connection.

void **esp_crt_bundle_detach** (mbedtls_ssl_config *conf)

Disable and dealloc the certification bundle.

Removes the certificate verification callback and deallocates used resources

Parameters

- `[in] conf`: The config struct for the SSL connection.

void **esp_crt_bundle_set** (**const** uint8_t *x509_bundle)

Set the default certificate bundle used for verification.

Overrides the default certificate bundle. In most use cases the bundle should be set through menuconfig. The bundle needs to be sorted by subject name since binary search is used to find certificates.

Parameters

- `[in] x509_bundle`: A pointer to the certificate bundle.

此 API 部分的示例代码在 ESP-IDF 示例工程的 protocols 目录下提供。

2.3.13 IP 网络层协议

IP 网络层协议（应用层协议之下）的文档位于连网 *API* 。

2.4 配网 API

2.4.1 Unified Provisioning

Overview

Unified provisioning support in the ESP-IDF provides an extensible mechanism to the developers to configure the device with the Wi-Fi credentials and/or other custom configuration using various transports and different security schemes. Depending on the use-case it provides a complete and ready solution for Wi-Fi network provisioning along with example iOS and Android applications. Or developers can extend the device-side and phone-app side implementations to accommodate their requirements for sending additional configuration data. Following are the important features of this implementation.

1. *Extensible Protocol:* The protocol is completely flexible and it offers the ability for the developers to send custom configuration in the provisioning process. The data representation too is left to the application to decide.
2. *Transport Flexibility:* The protocol can work on Wi-Fi (SoftAP + HTTP server) or on BLE as a transport protocol. The framework provides an ability to add support for any other transport easily as long as command-response behaviour can be supported on the transport.
3. *Security Scheme Flexibility:* It's understood that each use-case may require different security scheme to secure the data that is exchanged in the provisioning process. Some applications may work with SoftAP that's WPA2 protected or BLE with "just-works" security. Or the applications may consider the transport to be insecure and may want application level security. The unified provisioning framework allows application to choose the security as deemed suitable.
4. *Compact Data Representation:* The protocol uses Google Protobufs as a data representation for session setup and Wi-Fi provisioning. They provide a compact data representation and ability to parse the data in multiple programming languages in native format. Please note that this data representation is not forced on application specific data and the developers may choose the representation of their choice.

Typical Provisioning Process

Deciding on Transport

Unified provisioning subsystem supports Wi-Fi (SoftAP+HTTP server) and BLE (GATT based) transport schemes. Following points need to be considered while selecting the best possible transport for provisioning.

1. BLE based transport has an advantage that in the provisioning process, the BLE communication channel stays intact between the device and the client. That provides reliable provisioning feedback.
2. BLE based provisioning implementation makes the user-experience better from the phone apps as on Android and iOS both, the phone app can discover and connect to the device without requiring user to go out of the phone app
3. BLE transport however consumes ~110KB memory at runtime. If the product does not use the BLE or BT functionality after provisioning is done, almost all the memory can be reclaimed back and can be added into the heap.
4. SoftAP based transport is highly interoperable; however as the same radio is shared between SoftAP and Station interface, the transport is not reliable in the phase when the Wi-Fi connection to external AP is attempted. Also, the client may roam back to different network when the SoftAP changes the channel at the time of Station connection.
5. SoftAP transport does not require much additional memory for the Wi-Fi use-cases
6. SoftAP based provisioning requires the phone app user to go to "System Settings" to connect to Wi-Fi network hosted by the device in case of iOS. The discovery (scanning) as well as connection API is not available for the iOS applications.

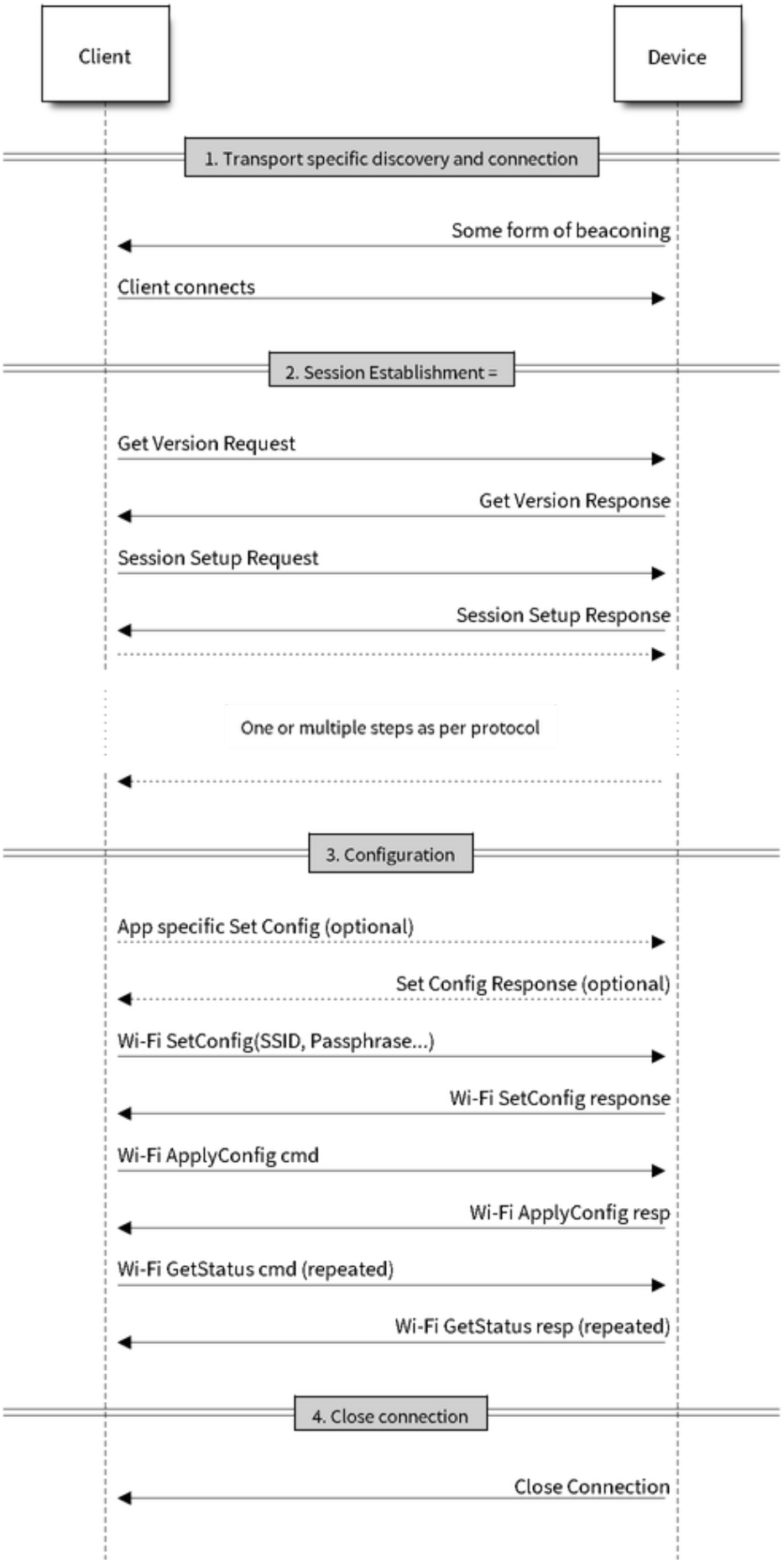

图 17: Typical Provisioning Process

Deciding on Security

Depending on the transport and other constraints the security scheme needs to be selected by the application developers. Following considerations need to be given from the provisioning security perspective: 1. The configuration data sent from the client to the device and the response has to be secured. 2. The client should authenticate the device it is connected to. 3. The device manufacturer may choose proof-of-possession - a unique per device secret to be entered on the provisioning client as a security measure to make sure that the user can provisions the device in the possession.

There are two levels of security schemes. The developer may select one or combination depending on requirements.

1. *Transport Security:* SoftAP provisioning may choose WPA2 protected security with unique per-device passphrase. Per-device unique passphrase can also act as a proof-of-possession. For BLE, "just-works" security can be used as a transport level security after understanding the level of security it provides.
2. *Application Security:* The unified provisioning subsystem provides application level security (*security1*) that provides data protection and authentication (through proof-of-possession) if the application does not use the transport level security or if the transport level security is not sufficient for the use-case.

Device Discovery

The advertisement and device discovery is left to the application and depending on the protocol chosen, the phone apps and device firmware application can choose appropriate method to advertise and discovery.

For the SoftAP+HTTP transport, typically the SSID (network name) of the AP hosted by the device can be used for discovery.

For the BLE transport device name or primary service included in the advertisement or combination of both can be used for discovery.

Architecture

The below diagram shows architecture of unified provisioning.

It relies on the base layer called *Protocol Communication* (Protocol Communication) which provides a framework for security schemes and transport mechanisms. Wi-Fi Provisioning layer uses Protocomm to provide simple callbacks to the application for setting the configuration and getting the Wi-Fi status. The application has control over implementation of these callbacks. In addition application can directly use protocomm to register custom handlers.

Application creates a protocomm instance which is mapped to a specific transport and specific security scheme. Each transport in the protocomm has a concept of an "end-point" which corresponds to logical channel for communication for specific type of information. For example security handshake happens on a different endpoint than the Wi-Fi configuration endpoint. Each end-point is identified using a string and depending on the transport internal representation of the end-point changes. In case of SoftAP+HTTP transport the end-point corresponds to URI whereas in case of BLE the end-point corresponds to GATT characteristic with specific UUID. Developers can create custom end-points and implement handler for the data that is received or sent over the same end-point.

Security Schemes

At present unified provisioning supports two security schemes: 1. Security0 - No security (No encryption) 2. Security1 - Curve25519 based key exchange, shared key derivation and AES256-CTR mode encryption of the data. It supports two modes :

a. Authorized - Proof of Possession (PoP) string used to authorize session and derive shared key
b. No Auth (Null PoP) - Shared key derived through key exchange only

Security1 scheme details are shown in the below sequence diagram

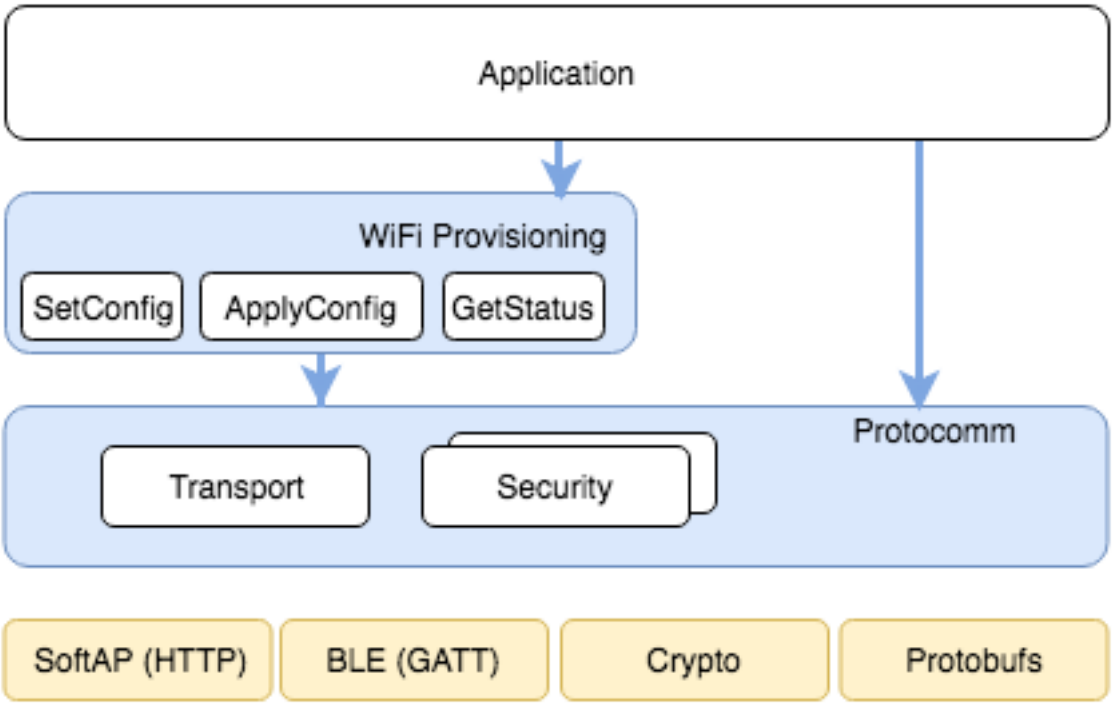

图 18: Unified Provisioning Architecture

Sample Code

Please refer to *Protocol Communication* and *Wi-Fi Provisioning* for API guides and code snippets on example usage.

Application implementation can be found as an example under provisioning.

Provisioning Tools

Provisioning applications are available for various platforms, along with source code:

- **Android:**
  - BLE Provisioning app on Play Store.
  - SoftAP Provisioning app on Play Store.
  - Source code on GitHub: esp-idf-provisioning-android.
- **iOS:**
  - BLE Provisioning app on app store.
  - SoftAP Provisioning app on app Store.
  - Source code on GitHub: esp-idf-provisioning-ios.
- Linux/MacOS/Windows : tools/esp_prov (a python based command line tool for provisioning)

The phone applications offer simple UI and thus more user centric, while the command line application is useful as a debugging tool for developers.

2.4.2 Protocol Communication

Overview

Protocol Communication (protocomm) component manages secure sessions and provides framework for multiple transports. The application can also use protocomm layer directly to have application specific extensions for the provisioning (or non-provisioning) use cases.

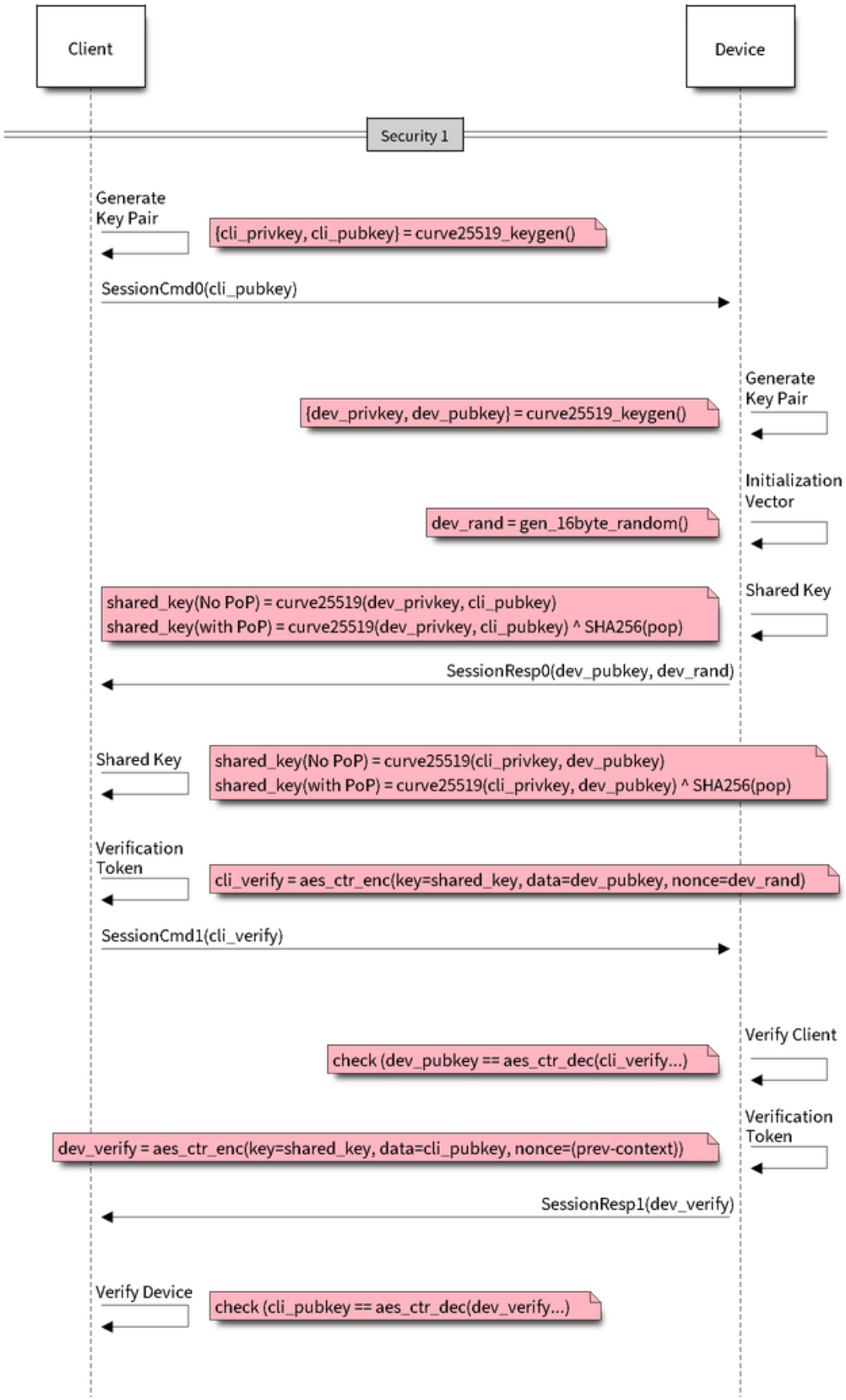

图 19: Security1

Following features are available for provisioning :

- **Communication security at application level -**
  - protocomm_security0 (no security)
  - protocomm_security1 (curve25519 key exchange + AES-CTR encryption)
- Proof-of-possession (support with protocomm_security1 only)

Protocomm internally uses protobuf (protocol buffers) for secure session establishment. Though users can implement their own security (even without using protobuf). One can even use protocomm without any security layer.

Protocomm provides framework for various transports - WiFi (SoftAP+HTTPD), BLE, console - in which case the handler invocation is automatically taken care of on the device side (see Transport Examples below for code snippets).

Note that the client still needs to establish session (only for protocomm_security1) by performing the two way handshake. See *Unified Provisioning* for more details about the secure handshake logic.

Transport Example (SoftAP + HTTP) with Security 1

For complete example see provisioning/legacy/softap_prov

```
/* Endpoint handler to be registered with protocomm.
 * This simply echoes back the received data. */
esp_err_t echo_req_handler (uint32_t session_id,
                            const uint8_t *inbuf, ssize_t inlen,
                            uint8_t **outbuf, ssize_t *outlen,
                            void *priv_data)
{
    /* Session ID may be used for persistence */
    printf("Session ID : %d", session_id);

    /* Echo back the received data */
    *outlen = inlen;            /* Output data length updated */
    *outbuf = malloc(inlen);    /* This will be deallocated outside */
    memcpy(*outbuf, inbuf, inlen);

    /* Private data that was passed at the time of endpoint creation */
    uint32_t *priv = (uint32_t *) priv_data;
    if (priv) {
        printf("Private data : %d", *priv);
    }

    return ESP_OK;
}

/* Example function for launching a protocomm instance over HTTP */
protocomm_t *start_pc(const char *pop_string)
{
    protocomm_t *pc = protocomm_new();


    /* Config for protocomm_httpd_start() */
    protocomm_httpd_config_t pc_config = {
        .data = {
        .config = PROTOCOMM_HTTPD_DEFAULT_CONFIG()
        }
    };

    /* Start protocomm server on top of HTTP */
    protocomm_httpd_start(pc, &pc_config);

    /* Create Proof of Possession object from pop_string. It must be valid
     * throughout the scope of protocomm endpoint. This need not be␣
↪static,
```

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```
     * ie. could be dynamically allocated and freed at the time of␣
↪endpoint
     * removal */
    const static protocomm_security_pop_t pop_obj = {
        .data = (const uint8_t *) strdup(pop_string),
        .len = strlen(pop_string)
    };

    /* Set security for communication at application level. Just like for
     * request handlers, setting security creates an endpoint and␣
↪registers
     * the handler provided by protocomm_security1. One can similarly use
     * protocomm_security0. Only one type of security can be set for a
     * protocomm instance at a time. */
    protocomm_set_security(pc, "security_endpoint", &protocomm_security1,␣
↪&pop_obj);

    /* Private data passed to the endpoint must be valid throughout the␣
↪scope
     * of protocomm endpoint. This need not be static, ie. could be␣
↪dynamically
     * allocated and freed at the time of endpoint removal */
    static uint32_t priv_data = 1234;

    /* Add a new endpoint for the protocomm instance, identified by a␣
↪unique name
     * and register a handler function along with private data to be␣
↪passed at the
     * time of handler execution. Multiple endpoints can be added as long␣
↪as they
     * are identified by unique names */
    protocomm_add_endpoint(pc, "echo_req_endpoint",
                           echo_req_handler, (void *) &priv_data);
    return pc;
}

/* Example function for stopping a protocomm instance */
void stop_pc(protocomm_t *pc)
{
    /* Remove endpoint identified by it's unique name */
    protocomm_remove_endpoint(pc, "echo_req_endpoint");

    /* Remove security endpoint identified by it's name */
    protocomm_unset_security(pc, "security_endpoint");

    /* Stop HTTP server */
    protocomm_httpd_stop(pc);

    /* Delete (deallocate) the protocomm instance */
    protocomm_delete(pc);
}
```

Transport Example (BLE) with Security 0

For complete example see provisioning/legacy/ble_prov

```
/* Example function for launching a secure protocomm instance over BLE */
protocomm_t *start_pc()
{
    protocomm_t *pc = protocomm_new();
```

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```
    /* Endpoint UUIDs */
    protocomm_ble_name_uuid_t nu_lookup_table[] = {
        {"security_endpoint", 0xFF51},
        {"echo_req_endpoint", 0xFF52}
    };

    /* Config for protocomm_ble_start() */
    protocomm_ble_config_t config = {
        .service_uuid = {
            /* LSB <---------------------------------------
            * ---------------------------------------> MSB */
            0xfb, 0x34, 0x9b, 0x5f, 0x80, 0x00, 0x00, 0x80,
            0x00, 0x10, 0x00, 0x00, 0xFF, 0xFF, 0x00, 0x00,
        },
        .nu_lookup_count = sizeof(nu_lookup_table)/sizeof(nu_lookup_
↪table[0]),
        .nu_lookup = nu_lookup_table
    };

    /* Start protocomm layer on top of BLE */
    protocomm_ble_start(pc, &config);

    /* For protocomm_security0, Proof of Possession is not used, and can␣
↪be kept NULL */
    protocomm_set_security(pc, "security_endpoint", &protocomm_security0,␣
↪NULL);
    protocomm_add_endpoint(pc, "echo_req_endpoint", echo_req_handler,␣
↪NULL);
    return pc;
}

/* Example function for stopping a protocomm instance */
void stop_pc(protocomm_t *pc)
{
    protocomm_remove_endpoint(pc, "echo_req_endpoint");
    protocomm_unset_security(pc, "security_endpoint");

    /* Stop BLE protocomm service */
    protocomm_ble_stop(pc);

    protocomm_delete(pc);
}
```

API Reference

Header File

- protocomm/include/common/protocomm.h

Functions

protocomm_t ***protocomm_new**(void)

Create a new protocomm instance.

This API will return a new dynamically allocated protocomm instance with all elements of the protocomm_t structure initialized to NULL.

Return

- protocomm_t* : On success
- NULL : No memory for allocating new instance

void **protocomm_delete**(*protocomm_t* *pc)

Delete a protocomm instance.

This API will deallocate a protocomm instance that was created using `protocomm_new()`.

Parameters

- `[in] pc`: Pointer to the protocomm instance to be deleted

esp_err_t **protocomm_add_endpoint**(*protocomm_t* *pc, **const** char *ep_name, *protocomm_req_handler_t* h, void *priv_data)

Add endpoint request handler for a protocomm instance.

This API will bind an endpoint handler function to the specified endpoint name, along with any private data that needs to be pass to the handler at the time of call.

Note

- An endpoint must be bound to a valid protocomm instance, created using `protocomm_new()`.
- This function internally calls the registered `add_endpoint()` function of the selected transport which is a member of the protocomm_t instance structure.

Return

- ESP_OK : Success
- ESP_FAIL : Error adding endpoint / Endpoint with this name already exists
- ESP_ERR_NO_MEM : Error allocating endpoint resource
- ESP_ERR_INVALID_ARG : Null instance/name/handler arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string
- `[in] h`: Endpoint handler function
- `[in] priv_data`: Pointer to private data to be passed as a parameter to the handler function on call. Pass NULL if not needed.

esp_err_t **protocomm_remove_endpoint**(*protocomm_t* *pc, **const** char *ep_name)

Remove endpoint request handler for a protocomm instance.

This API will remove a registered endpoint handler identified by an endpoint name.

Note

- This function internally calls the registered `remove_endpoint()` function which is a member of the protocomm_t instance structure.

Return

- ESP_OK : Success
- ESP_ERR_NOT_FOUND : Endpoint with specified name doesn’ t exist
- ESP_ERR_INVALID_ARG : Null instance/name arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string

esp_err_t **protocomm_open_session**(*protocomm_t* *pc, uint32_t *session_id*)

Allocates internal resources for new transport session.

Note

- An endpoint must be bound to a valid protocomm instance, created using `protocomm_new()`.

Return

- ESP_OK : Request handled successfully
- ESP_ERR_NO_MEM : Error allocating internal resource
- ESP_ERR_INVALID_ARG : Null instance/name arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] session_id`: Unique ID for a communication session

esp_err_t **protocomm_close_session**(*protocomm_t* *pc, uint32_t *session_id*)

Frees internal resources used by a transport session.

Note

- An endpoint must be bound to a valid protocomm instance, created using `protocomm_new()`.

Return

- ESP_OK : Request handled successfully
- ESP_ERR_INVALID_ARG : Null instance/name arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] session_id`: Unique ID for a communication session

esp_err_t **protocomm_req_handle**(*protocomm_t* *pc, **const** char *ep_name, uint32_t *session_id*, **const** uint8_t *inbuf, ssize_t *inlen*, uint8_t **outbuf, ssize_t *outlen*)

Calls the registered handler of an endpoint session for processing incoming data and generating the response.

Note

- An endpoint must be bound to a valid protocomm instance, created using `protocomm_new()`.
- Resulting output buffer must be deallocated by the caller.

Return

- ESP_OK : Request handled successfully
- ESP_FAIL : Internal error in execution of registered handler
- ESP_ERR_NO_MEM : Error allocating internal resource
- ESP_ERR_NOT_FOUND : Endpoint with specified name doesn't exist
- ESP_ERR_INVALID_ARG : Null instance/name arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string
- `[in] session_id`: Unique ID for a communication session
- `[in] inbuf`: Input buffer contains input request data which is to be processed by the registered handler
- `[in] inlen`: Length of the input buffer
- `[out] outbuf`: Pointer to internally allocated output buffer, where the resulting response data output from the registered handler is to be stored
- `[out] outlen`: Buffer length of the allocated output buffer

esp_err_t **protocomm_set_security**(*protocomm_t* *pc, **const** char *ep_name, **const** *protocomm_security_t* *sec, **const** *protocomm_security_pop_t* *pop)

Add endpoint security for a protocomm instance.

This API will bind a security session establisher to the specified endpoint name, along with any proof of possession that may be required for authenticating a session client.

Note

- An endpoint must be bound to a valid protocomm instance, created using `protocomm_new()`.
- The choice of security can be any `protocomm_security_t` instance. Choices `protocomm_security0` and `protocomm_security1` are readily available.

Return

- ESP_OK : Success
- ESP_FAIL : Error adding endpoint / Endpoint with this name already exists
- ESP_ERR_INVALID_STATE : Security endpoint already set
- ESP_ERR_NO_MEM : Error allocating endpoint resource
- ESP_ERR_INVALID_ARG : Null instance/name/handler arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string
- `[in] sec`: Pointer to endpoint security instance
- `[in] pop`: Pointer to proof of possession for authenticating a client

esp_err_t **protocomm_unset_security**(*protocomm_t* *pc, **const** char *ep_name)

Remove endpoint security for a protocomm instance.

This API will remove a registered security endpoint identified by an endpoint name.

Return

- ESP_OK : Success
- ESP_ERR_NOT_FOUND : Endpoint with specified name doesn't exist
- ESP_ERR_INVALID_ARG : Null instance/name arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string

esp_err_t **protocomm_set_version** (*protocomm_t* *pc, **const** char *ep_name, **const** char *version)

Set endpoint for version verification.

This API can be used for setting an application specific protocol version which can be verified by clients through the endpoint.

Note

- An endpoint must be bound to a valid protocomm instance, created using `protocomm_new()`.

Return

- ESP_OK : Success
- ESP_FAIL : Error adding endpoint / Endpoint with this name already exists
- ESP_ERR_INVALID_STATE : Version endpoint already set
- ESP_ERR_NO_MEM : Error allocating endpoint resource
- ESP_ERR_INVALID_ARG : Null instance/name/handler arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string
- `[in] version`: Version identifier(name) string

esp_err_t **protocomm_unset_version** (*protocomm_t* *pc, **const** char *ep_name)

Remove version verification endpoint from a protocomm instance.

This API will remove a registered version endpoint identified by an endpoint name.

Return

- ESP_OK : Success
- ESP_ERR_NOT_FOUND : Endpoint with specified name doesn't exist
- ESP_ERR_INVALID_ARG : Null instance/name arguments

Parameters

- `[in] pc`: Pointer to the protocomm instance
- `[in] ep_name`: Endpoint identifier(name) string

Type Definitions

typedef *esp_err_t* (***protocomm_req_handler_t**) (uint32_t session_id, **const** uint8_t *inbuf, ssize_t inlen, uint8_t **outbuf, ssize_t *outlen, void *priv_data)

Function prototype for protocomm endpoint handler.

typedef struct protocomm **protocomm_t**

This structure corresponds to a unique instance of protocomm returned when the API `protocomm_new()` is called. The remaining Protocomm APIs require this object as the first parameter.

Note Structure of the protocomm object is kept private

Header File

- protocomm/include/security/protocomm_security.h

Structures

struct protocomm_security_pop

Proof Of Possession for authenticating a secure session.

Public Members

const uint8_t ***data**
Pointer to buffer containing the proof of possession data

uint16_t **len**
Length (in bytes) of the proof of possession data

struct protocomm_security
Protocomm security object structure.

The member functions are used for implementing secure protocomm sessions.

Note This structure should not have any dynamic members to allow re-entrancy

Public Members

int **ver**
Unique version number of security implementation

esp_err_t (***init**) (*protocomm_security_handle_t* *handle)
Function for initializing/allocating security infrastructure

esp_err_t (***cleanup**) (*protocomm_security_handle_t* handle)
Function for deallocating security infrastructure

esp_err_t (***new_transport_session**) (*protocomm_security_handle_t* handle, uint32_t session_id)
Starts new secure transport session with specified ID

esp_err_t (***close_transport_session**) (*protocomm_security_handle_t* handle, uint32_t session_id)
Closes a secure transport session with specified ID

esp_err_t (***security_req_handler**) (*protocomm_security_handle_t* handle, **const** *protocomm_security_pop_t* *pop, uint32_t session_id, **const** uint8_t *inbuf, ssize_t inlen, uint8_t **outbuf, ssize_t *outlen, void *priv_data)
Handler function for authenticating connection request and establishing secure session

esp_err_t (***encrypt**) (*protocomm_security_handle_t* handle, uint32_t session_id, **const** uint8_t *inbuf, ssize_t inlen, uint8_t *outbuf, ssize_t *outlen)
Function which implements the encryption algorithm

esp_err_t (***decrypt**) (*protocomm_security_handle_t* handle, uint32_t session_id, **const** uint8_t *inbuf, ssize_t inlen, uint8_t *outbuf, ssize_t *outlen)
Function which implements the decryption algorithm

Type Definitions

typedef struct *protocomm_security_pop* **protocomm_security_pop_t**
Proof Of Possession for authenticating a secure session.

typedef void ***protocomm_security_handle_t**

typedef struct *protocomm_security* **protocomm_security_t**
Protocomm security object structure.

The member functions are used for implementing secure protocomm sessions.

Note This structure should not have any dynamic members to allow re-entrancy

Header File

- protocomm/include/security/protocomm_security0.h

Header File

- protocomm/include/security/protocomm_security1.h

Header File

- protocomm/include/transports/protocomm_httpd.h

Functions

esp_err_t **protocomm_httpd_start** (*protocomm_t *pc*, **const** *protocomm_httpd_config_t *config*)

Start HTTPD protocomm transport.

This API internally creates a framework to allow endpoint registration and security configuration for the protocomm.

Note This is a singleton. ie. Protocomm can have multiple instances, but only one instance can be bound to an HTTP transport layer.

Return

- ESP_OK : Success
- ESP_ERR_INVALID_ARG : Null arguments
- ESP_ERR_NOT_SUPPORTED : Transport layer bound to another protocomm instance
- ESP_ERR_INVALID_STATE : Transport layer already bound to this protocomm instance
- ESP_ERR_NO_MEM : Memory allocation for server resource failed
- ESP_ERR_HTTPD_* : HTTP server error on start

Parameters

- `[in] pc`: Protocomm instance pointer obtained from protocomm_new()
- `[in] config`: Pointer to config structure for initializing HTTP server

esp_err_t **protocomm_httpd_stop** (*protocomm_t *pc*)

Stop HTTPD protocomm transport.

This API cleans up the HTTPD transport protocomm and frees all the handlers registered with the protocomm.

Return

- ESP_OK : Success
- ESP_ERR_INVALID_ARG : Null / incorrect protocomm instance pointer

Parameters

- `[in] pc`: Same protocomm instance that was passed to protocomm_httpd_start()

Unions

union protocomm_httpd_config_data_t

#include <protocomm_httpd.h> Protocomm HTTPD Configuration Data

Public Members

void ***handle**

HTTP Server Handle, if ext_handle_provided is set to true

protocomm_http_server_config_t **config**

HTTP Server Configuration, if a server is not already active

Structures

struct protocomm_http_server_config_t

Config parameters for protocomm HTTP server.

Public Members

uint16_t **port**
Port on which the HTTP server will listen

size_t **stack_size**
Stack size of server task, adjusted depending upon stack usage of endpoint handler

unsigned **task_priority**
Priority of server task

struct protocomm_httpd_config_t
Config parameters for protocomm HTTP server.

Public Members

bool **ext_handle_provided**
Flag to indicate of an external HTTP Server Handle has been provided. In such as case, protocomm will use the same HTTP Server and not start a new one internally.

protocomm_httpd_config_data_t **data**
Protocomm HTTPD Configuration Data

Macros

PROTOCOMM_HTTPD_DEFAULT_CONFIG()

Header File

- protocomm/include/transports/protocomm_ble.h

Functions

esp_err_t **protocomm_ble_start**(*protocomm_t* *pc, **const** *protocomm_ble_config_t* *config)
Start Bluetooth Low Energy based transport layer for provisioning.

Initialize and start required BLE service for provisioning. This includes the initialization for characteristics/service for BLE.

Return
- ESP_OK : Success
- ESP_FAIL : Simple BLE start error
- ESP_ERR_NO_MEM : Error allocating memory for internal resources
- ESP_ERR_INVALID_STATE : Error in ble config
- ESP_ERR_INVALID_ARG : Null arguments

Parameters
- `[in] pc`: Protocomm instance pointer obtained from protocomm_new()
- `[in] config`: Pointer to config structure for initializing BLE

esp_err_t **protocomm_ble_stop**(*protocomm_t* *pc)
Stop Bluetooth Low Energy based transport layer for provisioning.

Stops service/task responsible for BLE based interactions for provisioning

Note You might want to optionally reclaim memory from Bluetooth. Refer to the documentation of `esp_bt_mem_release` in that case.

Return
- ESP_OK : Success
- ESP_FAIL : Simple BLE stop error
- ESP_ERR_INVALID_ARG : Null / incorrect protocomm instance

Parameters
- `[in] pc`: Same protocomm instance that was passed to protocomm_ble_start()

Structures

struct name_uuid

This structure maps handler required by protocomm layer to UUIDs which are used to uniquely identify BLE characteristics from a smartphone or a similar client device.

Public Members

const char ***name**

Name of the handler, which is passed to protocomm layer

uint16_t **uuid**

UUID to be assigned to the BLE characteristic which is mapped to the handler

struct protocomm_ble_config

Config parameters for protocomm BLE service.

Public Members

char **device_name**[**MAX_BLE_DEVNAME_LEN**]

BLE device name being broadcast at the time of provisioning

uint8_t **service_uuid**[**BLE_UUID128_VAL_LENGTH**]

128 bit UUID of the provisioning service

uint8_t ***manufacturer_data**

BLE device manufacturer data pointer in advertisement

ssize_t **manufacturer_data_len**

BLE device manufacturer data length in advertisement

ssize_t **nu_lookup_count**

Number of entries in the Name-UUID lookup table

protocomm_ble_name_uuid_t ***nu_lookup**

Pointer to the Name-UUID lookup table

Macros

MAX_BLE_DEVNAME_LEN

BLE device name cannot be larger than this value 31 bytes (max scan response size) - 1 byte (length) - 1 byte (type) = 29 bytes

BLE_UUID128_VAL_LENGTH

MAX_BLE_MANUFACTURER_DATA_LEN

Theoretically, the limit for max manufacturer length remains same as BLE device name i.e. 31 bytes (max scan response size) - 1 byte (length) - 1 byte (type) = 29 bytes However, manufacturer data goes along with BLE device name in scan response. So, it is important to understand the actual length should be smaller than (29 - (BLE device name length) - 2).

Type Definitions

typedef struct *name_uuid* **protocomm_ble_name_uuid_t**

This structure maps handler required by protocomm layer to UUIDs which are used to uniquely identify BLE characteristics from a smartphone or a similar client device.

typedef struct *protocomm_ble_config* **protocomm_ble_config_t**

Config parameters for protocomm BLE service.

2.4.3 Wi-Fi Provisioning

Overview

This component provides APIs that control Wi-Fi provisioning service for receiving and configuring Wi-Fi credentials over SoftAP or BLE transport via secure *Protocol Communication (protocomm)* sessions. The set of `wifi_prov_mgr_` APIs help in quickly implementing a provisioning service having necessary features with minimal amount of code and sufficient flexibility.

Initialization *`wifi_prov_mgr_init()`* is called to configure and initialize the provisioning manager and thus this must be called prior to invoking any other `wifi_prov_mgr_` APIs. Note that the manager relies on other components of IDF, namely NVS, TCP/IP, Event Loop and Wi-Fi (and optionally mDNS), hence these must be initialized beforehand. The manager can be de-initialized at any moment by making a call to *`wifi_prov_mgr_deinit()`*.

```
wifi_prov_mgr_config_t config = {
    .scheme = wifi_prov_scheme_ble,
    .scheme_event_handler = WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BTDM
};

ESP_ERR_CHECK( wifi_prov_mgr_init(config) );
```

The configuration structure `wifi_prov_mgr_config_t` has a few fields to specify the behavior desired of the manager :

- *scheme* : This is used to specify the provisioning scheme. Each scheme corresponds to one of the modes of transport supported by protocomm. Hence, we have three options :
  - `wifi_prov_scheme_ble` : BLE transport and GATT Server for handling provisioning commands
  - `wifi_prov_scheme_softap` : Wi-Fi SoftAP transport and HTTP Server for handling provisioning commands
  - `wifi_prov_scheme_console` : Serial transport and console for handling provisioning commands
- *scheme_event_handler* : An event handler defined along with scheme. Choosing appropriate scheme specific event handler allows the manager to take care of certain matters automatically. Presently this is not used for either SoftAP or Console based provisioning, but is very convenient for BLE. To understand how, we must recall that Bluetooth requires quite some amount of memory to function and once provisioning is finished, the main application may want to reclaim back this memory (or part of it, if it needs to use either BLE or classic BT). Also, upon every future reboot of a provisioned device, this reclamation of memory needs to be performed again. To reduce this complication in using `wifi_prov_scheme_ble`, the scheme specific handlers have been defined, and depending upon the chosen handler, the BLE / classic BT / BTDM memory will be freed automatically when the provisioning manager is de-initialized. The available options are:
  - `WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BTDM` - Free both classic BT and BLE (BTDM) memory. Used when main application doesn't require Bluetooth at all.
  - `WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BLE` - Free only BLE memory. Used when main application requires classic BT.
  - `WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BT` - Free only classic BT. Used when main application requires BLE. In this case freeing happens right when the manager is initialized.
  - `WIFI_PROV_EVENT_HANDLER_NONE` - Don't use any scheme specific handler. Used when provisioning scheme is not BLE (i.e. SoftAP or Console), or when main application wants to handle the memory reclaiming on its own, or needs both BLE and classic BT to function.
- *app_event_handler* (Deprecated) : It is now recommended to catch `WIFI_PROV_EVENT``s that are emitted to the default event loop handler. See definition of ``wifi_prov_cb_event_t` for the list of events that are generated by the provisioning service. Here is an excerpt showing some of the provisioning events:

```
static void event_handler(void* arg, esp_event_base_t event_base,
                          int event_id, void* event_data)
```

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```
{
    if (event_base == WIFI_PROV_EVENT) {
        switch (event_id) {
            case WIFI_PROV_START:
                ESP_LOGI(TAG, "Provisioning started");
                break;
            case WIFI_PROV_CRED_RECV: {
                wifi_sta_config_t *wifi_sta_cfg = (wifi_sta_config_t␣
↪*)event_data;
                ESP_LOGI(TAG, "Received Wi-Fi credentials"
                         "\n\tSSID     : %s\n\tPassword : %s",
                         (const char *) wifi_sta_cfg->ssid,
                         (const char *) wifi_sta_cfg->password);
                break;
            }
            case WIFI_PROV_CRED_FAIL: {
                wifi_prov_sta_fail_reason_t *reason = (wifi_prov_sta_fail_
↪reason_t *)event_data;
                ESP_LOGE(TAG, "Provisioning failed!\n\tReason : %s"
                         "\n\tPlease reset to factory and retry␣
↪provisioning",
                         (*reason == WIFI_PROV_STA_AUTH_ERROR) ?
                         "Wi-Fi station authentication failed" : "Wi-Fi␣
↪access-point not found");
                break;
            }
            case WIFI_PROV_CRED_SUCCESS:
                ESP_LOGI(TAG, "Provisioning successful");
                break;
            case WIFI_PROV_END:
                /* De-initialize manager once provisioning is finished */
                wifi_prov_mgr_deinit();
                break;
            default:
                break;
        }
    }
}
```

The manager can be de-initialized at any moment by making a call to *wifi_prov_mgr_deinit()*.

Check Provisioning State Whether device is provisioned or not can be checked at runtime by calling *wifi_prov_mgr_is_provisioned()*. This internally checks if the Wi-Fi credentials are stored in NVS.

Note that presently manager does not have its own NVS namespace for storage of Wi-Fi credentials, instead it relies on the `esp_wifi_` APIs to set and get the credentials stored in NVS from the default location.

If provisioning state needs to be reset, any of the following approaches may be taken :

- the associated part of NVS partition has to be erased manually
- main application must implement some logic to call `esp_wifi_` APIs for erasing the credentials at runtime
- main application must implement some logic to force start the provisioning irrespective of the provisioning state

```
bool provisioned = false;
ESP_ERR_CHECK( wifi_prov_mgr_is_provisioned(&provisioned) );
```

Start Provisioning Service At the time of starting provisioning we need to specify a service name and the corresponding key. These translate to :

- Wi-Fi SoftAP SSID and passphrase, respectively, when scheme is `wifi_prov_scheme_softap`
- BLE Device name (service key is ignored) when scheme is `wifi_prov_scheme_ble`

Also, since internally the manager uses *protocomm*, we have the option of choosing one of the security features provided by it :

- Security 1 is secure communication which consists of a prior handshake involving X25519 key exchange along with authentication using a proof of possession (*pop*), followed by AES-CTR for encryption/decryption of subsequent messages
- Security 0 is simply plain text communication. In this case the *pop* is simply ignored

See *Provisioning* for details about the security features.

```
const char *service_name = "my_device";
const char *service_key  = "password";

wifi_prov_security_t security = WIFI_PROV_SECURITY_1;
const char *pop = "abcd1234";

ESP_ERR_CHECK( wifi_prov_mgr_start_provisioning(security, pop, service_
↪name, service_key) );
```

The provisioning service will automatically finish only if it receives valid Wi-Fi AP credentials followed by successfully connection of device to the AP (IP obtained). Regardless of that, the provisioning service can be stopped at any moment by making a call to *`wifi_prov_mgr_stop_provisioning()`*.

注解: If the device fails to connect with the provided credentials, it won't accept new credentials anymore, but the provisioning service will keep on running (only to convey failure to the client), until the device is restarted. Upon restart the provisioning state will turn out to be true this time (as credentials will be found in NVS), but device will again fail to connect with those same credentials (unless an AP with the matching credentials somehow does become available). This situation can be fixed by resetting the credentials in NVS or force starting the provisioning service. This has been explained above in *Check Provisioning State*.

Waiting For Completion Typically, the main application will wait for the provisioning to finish, then de-initialize the manager to free up resources and finally start executing its own logic.

There are two ways for making this possible. The simpler way is to use a blocking call to *`wifi_prov_mgr_wait()`*.

```
// Start provisioning service
ESP_ERR_CHECK( wifi_prov_mgr_start_provisioning(security, pop, service_
↪name, service_key) );

// Wait for service to complete
wifi_prov_mgr_wait();

// Finally de-initialize the manager
wifi_prov_mgr_deinit();
```

The other way is to use the default event loop handler to catch `WIFI_PROV_EVENT``s and call :cpp:func:`wifi_prov_mgr_deinit()` when event ID is ``WIFI_PROV_END:`

```
static void event_handler(void* arg, esp_event_base_t event_base,
                          int event_id, void* event_data)
{
    if (event_base == WIFI_PROV_EVENT && event_id == WIFI_PROV_END) {
        /* De-initialize manager once provisioning is finished */
        wifi_prov_mgr_deinit();
    }
}
```

User Side Implementation When the service is started, the device to be provisioned is identified by the advertised service name which, depending upon the selected transport, is either the BLE device name or the SoftAP SSID.

When using SoftAP transport, for allowing service discovery, mDNS must be initialized before starting provisioning. In this case the hostname set by the main application is used, and the service type is internally set to *_esp_wifi_prov*.

When using BLE transport, a custom 128 bit UUID should be set using `wifi_prov_scheme_ble_set_service_uuid()`. This UUID will be included in the BLE advertisement and will correspond to the primary GATT service that provides provisioning endpoints as GATT characteristics. Each GATT characteristic will be formed using the primary service UUID as base, with different auto assigned 12th and 13th bytes (assume counting starts from 0th byte). Since, an endpoint characteristic UUID is auto assigned, it shouldn't be used to identify the endpoint. Instead, client side applications should identify the endpoints by reading the User Characteristic Description (0x2901) descriptor for each characteristic, which contains the endpoint name of the characteristic. For example, if the service UUID is set to *55cc035e-fb27-4f80-be02-3c60828b7451*, each endpoint characteristic will be assigned a UUID like *55cc____-fb27-4f80-be02-3c60828b7451*, with unique values at the 12th and 13th bytes.

Once connected to the device, the provisioning related protocomm endpoints can be identified as follows :

表 5: Endpoints provided by Provisioning Service

Endpoint Name (BLE + GATT Server)	URI (SoftAP + HTTP Server + mDNS)	Description
prov-session	http://<mdns-hostname>.local/prov-session	Security endpoint used for session establishment
prov-scan	http://wifi-prov.local/prov-scan	Endpoint used for starting Wi-Fi scan and receiving scan results
prov-config	http://<mdns-hostname>.local/prov-config	Endpoint used for configuring Wi-Fi credentials on device
proto-ver	http://<mdns-hostname>.local/proto-ver	Endpoint for retrieving version info

Immediately after connecting, the client application may fetch the version / capabilities information from the *proto-ver* endpoint. All communications to this endpoint are un-encrypted, hence necessary information (that may be relevant for deciding compatibility) can be retrieved before establishing a secure session. The response is in JSON format and looks like : `prov: { ver: v1.1, cap: [no_pop] }, my_app: { ver: 1.345, cap: [cloud, local_ctrl] },....` Here label *prov* provides provisioning service version (*ver*) and capabilities (*cap*). For now, only *no_pop* capability is supported, which indicates that the service doesn't require proof of possession for authentication. Any application related version / capabilities will be given by other labels (like *my_app* in this example). These additional fields are set using `wifi_prov_mgr_set_app_info()`.

User side applications need to implement the signature handshaking required for establishing and authenticating secure protocomm sessions as per the security scheme configured for use (this is not needed when manager is configured to use protocomm security 0).

See Unified Provisioning for more details about the secure handshake and encryption used. Applications must use the *.proto* files found under protocomm/proto, which define the Protobuf message structures supported by *prov-session* endpoint.

Once a session is established, Wi-Fi credentials are configured using the following set of *wifi_config* commands, serialized as Protobuf messages (the corresponding *.proto* files can be found under wifi_provisioning/proto) :

- *get_status* - For querying the Wi-Fi connection status. The device will respond with a status which will be one of connecting / connected / disconnected. If status is disconnected, a disconnection reason will also be included in the status response.
- *set_config* - For setting the Wi-Fi connection credentials
- *apply_config* - For applying the credentials saved during *set_config* and start the Wi-Fi station

After session establishment, client can also request Wi-Fi scan results from the device. The results returned is a list

of AP SSIDs, sorted in descending order of signal strength. This allows client applications to display APs nearby to the device at the time of provisioning, and users can select one of the SSIDs and provide the password which is then sent using the *wifi_config* commands described above. The *wifi_scan* endpoint supports the following protobuf commands :

- *scan_start* - For starting Wi-Fi scan with various options :
  - *blocking* (input) - If true, the command returns only when the scanning is finished
  - *passive* (input) - If true scan is started in passive mode (this may be slower) instead of active mode
  - *group_channels* (input) - This specifies whether to scan all channels in one go (when zero) or perform scanning of channels in groups, with 120ms delay between scanning of consecutive groups, and the value of this parameter sets the number of channels in each group. This is useful when transport mode is SoftAP, where scanning all channels in one go may not give the Wi-Fi driver enough time to send out beacons, and hence may cause disconnection with any connected stations. When scanning in groups, the manager will wait for atleast 120ms after completing scan on a group of channels, and thus allow the driver to send out the beacons. For example, given that the total number of Wi-Fi channels is 14, then setting group_channels to 4, will create 5 groups, with each group having 3 channels, except the last one which will have 14 % 3 = 2 channels. So, when scan is started, the first 3 channels will be scanned, followed by a 120ms delay, and then the next 3 channels, and so on, until all the 14 channels have been scanned. One may need to adjust this parameter as having only few channels in a group may slow down the overall scan time, while having too many may again cause disconnection. Usually a value of 4 should work for most cases. Note that for any other mode of transport, e.g. BLE, this can be safely set to 0, and hence achieve the fastest overall scanning time.
  - *period_ms* (input) - Scan parameter specifying how long to wait on each channel
- *scan_status* - Gives the status of scanning process :
  - *scan_finished* (output) - When scan has finished this returns true
  - *result_count* (output) - This gives the total number of results obtained till now. If scan is yet happening this number will keep on updating
- *scan_result* - For fetching scan results. This can be called even if scan is still on going
  - *start_index* (input) - Starting index from where to fetch the entries from the results list
  - *count* (input) - Number of entries to fetch from the starting index
  - *entries* (output) - List of entries returned. Each entry consists of *ssid*, *channel* and *rssi* information

Additional Endpoints In case users want to have some additional protocomm endpoints customized to their requirements, this is done in two steps. First is creation of an endpoint with a specific name, and the second step is the registration of a handler for this endpoint. See *protocomm* for the function signature of an endpoint handler. A custom endpoint must be created after initialization and before starting the provisioning service. Whereas, the protocomm handler is registered for this endpoint only after starting the provisioning service.

```
wifi_prov_mgr_init(config);
wifi_prov_mgr_endpoint_create("custom-endpoint");
wifi_prov_mgr_start_provisioning(security, pop, service_name, service_
↪key);
wifi_prov_mgr_endpoint_register("custom-endpoint", custom_ep_handler,␣
↪custom_ep_data);
```

When the provisioning service stops, the endpoint is unregistered automatically.

One can also choose to call *wifi_prov_mgr_endpoint_unregister()* to manually deactivate an endpoint at runtime. This can also be used to deactivate the internal endpoints used by the provisioning service.

When / How To Stop Provisioning Service? The default behavior is that once the device successfully connects using the Wi-Fi credentials set by the *apply_config* command, the provisioning service will be stopped (and BLE / SoftAP turned off) automatically after responding to the next *get_status* command. If *get_status* command is not received by the device, the service will be stopped after a 30s timeout.

On the other hand, if device was not able to connect using the provided Wi-Fi credentials, due to incorrect SSID / passphrase, the service will keep running, and *get_status* will keep responding with disconnected status and reason for disconnection. Any further attempts to provide another set of Wi-Fi credentials, will be rejected. These credentials will be preserved, unless the provisioning service is force started, or NVS erased.

If this default behavior is not desired, it can be disabled by calling `wifi_prov_mgr_disable_auto_stop()`. Now the provisioning service will only be stopped after an explicit call to `wifi_prov_mgr_stop_provisioning()`, which returns immediately after scheduling a task for stopping the service. The service stops after a certain delay and WIFI_PROV_END event gets emitted. This delay is specified by the argument to `wifi_prov_mgr_disable_auto_stop()`.

The customized behavior is useful for applications which want the provisioning service to be stopped some time after the Wi-Fi connection is successfully established. For example, if the application requires the device to connect to some cloud service and obtain another set of credentials, and exchange this credentials over a custom protocomm endpoint, then after successfully doing so stop the provisioning service by calling `wifi_prov_mgr_stop_provisioning()` inside the protocomm handler itself. The right amount of delay ensures that the transport resources are freed only after the response from the protocomm handler reaches the client side application.

Application Examples

For complete example implementation see provisioning/wifi_prov_mgr

Provisioning Tools

Provisioning applications are available for various platforms, along with source code:

- **Android:**
  - BLE Provisioning app on Play Store.
  - SoftAP Provisioning app on Play Store.
  - Source code on GitHub: esp-idf-provisioning-android.
- **iOS:**
  - BLE Provisioning app on app store.
  - SoftAP Provisioning app on app Store.
  - Source code on GitHub: esp-idf-provisioning-ios.
- Linux/MacOS/Windows : tools/esp_prov (a python based command line tool for provisioning)

The phone applications offer simple UI and thus more user centric, while the command line application is useful as a debugging tool for developers.

API Reference

Header File

- wifi_provisioning/include/wifi_provisioning/manager.h

Functions

esp_err_t **wifi_prov_mgr_init** (*wifi_prov_mgr_config_t config*)

Initialize provisioning manager instance.

Configures the manager and allocates internal resources

Configuration specifies the provisioning scheme (transport) and event handlers

Event WIFI_PROV_INIT is emitted right after initialization is complete

Return

- ESP_OK : Success
- ESP_FAIL : Fail

Parameters

- `[in] config`: Configuration structure

void **wifi_prov_mgr_deinit**(void)

Stop provisioning (if running) and release resource used by the manager.

Event WIFI_PROV_DEINIT is emitted right after de-initialization is finished

If provisioning service is still active when this API is called, it first stops the service, hence emitting WIFI_PROV_END, and then performs the de-initialization

esp_err_t **wifi_prov_mgr_is_provisioned**(bool **provisioned*)

Checks if device is provisioned.

This checks if Wi-Fi credentials are present on the NVS

The Wi-Fi credentials are assumed to be kept in the same NVS namespace as used by esp_wifi component

If one were to call esp_wifi_set_config() directly instead of going through the provisioning process, this function will still yield true (i.e. device will be found to be provisioned)

Note Calling wifi_prov_mgr_start_provisioning() automatically resets the provision state, irrespective of what the state was prior to making the call.

Return

- ESP_OK : Retrieved provision state successfully
- ESP_FAIL : Wi-Fi not initialized
- ESP_ERR_INVALID_ARG : Null argument supplied
- ESP_ERR_INVALID_STATE : Manager not initialized

Parameters

- `[out] provisioned`: True if provisioned, else false

esp_err_t **wifi_prov_mgr_start_provisioning**(*wifi_prov_security_t* *security*, **const** char **pop*, **const** char **service_name*, **const** char **service_key*)

Start provisioning service.

This starts the provisioning service according to the scheme configured at the time of initialization. For scheme :

- wifi_prov_scheme_ble : This starts protocomm_ble, which internally initializes BLE transport and starts GATT server for handling provisioning requests
- wifi_prov_scheme_softap : This activates SoftAP mode of Wi-Fi and starts protocomm_httpd, which internally starts an HTTP server for handling provisioning requests (If mDNS is active it also starts advertising service with type _esp_wifi_prov._tcp)

Event WIFI_PROV_START is emitted right after provisioning starts without failure

Note This API will start provisioning service even if device is found to be already provisioned, i.e. wifi_prov_mgr_is_provisioned() yields true

Return

- ESP_OK : Provisioning started successfully
- ESP_FAIL : Failed to start provisioning service
- ESP_ERR_INVALID_STATE : Provisioning manager not initialized or already started

Parameters

- `[in] security`: Specify which protocomm security scheme to use :
  - WIFI_PROV_SECURITY_0 : For no security
  - WIFI_PROV_SECURITY_1 : x25519 secure handshake for session establishment followed by AES-CTR encryption of provisioning messages
- `[in] pop`: Pointer to proof of possession string (NULL if not needed). This is relevant only for protocomm security 1, in which case it is used for authenticating secure session
- `[in] service_name`: Unique name of the service. This translates to:
  - Wi-Fi SSID when provisioning mode is softAP
  - Device name when provisioning mode is BLE
- `[in] service_key`: Key required by client to access the service (NULL if not needed). This translates to:
  - Wi-Fi password when provisioning mode is softAP
  - ignored when provisioning mode is BLE

void **wifi_prov_mgr_stop_provisioning**(void)

Stop provisioning service.

If provisioning service is active, this API will initiate a process to stop the service and return. Once the service actually stops, the event WIFI_PROV_END will be emitted.

If wifi_prov_mgr_deinit() is called without calling this API first, it will automatically stop the provisioning service and emit the WIFI_PROV_END, followed by WIFI_PROV_DEINIT, before returning.

This API will generally be used along with wifi_prov_mgr_disable_auto_stop() in the scenario when the main application has registered its own endpoints, and wishes that the provisioning service is stopped only when some protocomm command from the client side application is received.

Calling this API inside an endpoint handler, with sufficient cleanup_delay, will allow the response / acknowledgment to be sent successfully before the underlying protocomm service is stopped.

Cleaup_delay is set when calling wifi_prov_mgr_disable_auto_stop(). If not specified, it defaults to 1000ms.

For straightforward cases, using this API is usually not necessary as provisioning is stopped automatically once WIFI_PROV_CRED_SUCCESS is emitted. Stopping is delayed (maximum 30 seconds) thus allowing the client side application to query for Wi-Fi state, i.e. after receiving the first query and sending `Wi-Fi state connected` response the service is stopped immediately.

void **wifi_prov_mgr_wait**(void)

Wait for provisioning service to finish.

Calling this API will block until provisioning service is stopped i.e. till event WIFI_PROV_END is emitted.

This will not block if provisioning is not started or not initialized.

esp_err_t **wifi_prov_mgr_disable_auto_stop**(uint32_t *cleanup_delay*)

Disable auto stopping of provisioning service upon completion.

By default, once provisioning is complete, the provisioning service is automatically stopped, and all endpoints (along with those registered by main application) are deactivated.

This API is useful in the case when main application wishes to close provisioning service only after it receives some protocomm command from the client side app. For example, after connecting to Wi-Fi, the device may want to connect to the cloud, and only once that is successfully, the device is said to be fully configured. But, then it is upto the main application to explicitly call wifi_prov_mgr_stop_provisioning() later when the device is fully configured and the provisioning service is no longer required.

Note This must be called before executing wifi_prov_mgr_start_provisioning()

Return

- ESP_OK : Success
- ESP_ERR_INVALID_STATE : Manager not initialized or provisioning service already started

Parameters

- `[in] cleanup_delay`: Sets the delay after which the actual cleanup of transport related resources is done after a call to wifi_prov_mgr_stop_provisioning() returns. Minimum allowed value is 100ms. If not specified, this will default to 1000ms.

esp_err_t **wifi_prov_mgr_set_app_info**(**const** char **label*, **const** char **version*, **const** char ***capabilities*, size_t *total_capabilities*)

Set application version and capabilities in the JSON data returned by proto-ver endpoint.

This function can be called multiple times, to specify information about the various application specific services running on the device, identified by unique labels.

The provisioning service itself registers an entry in the JSON data, by the label "prov", containing only provisioning service version and capabilities. Application services should use a label other than "prov" so as not to overwrite this.

Note This must be called before executing wifi_prov_mgr_start_provisioning()

Return

- ESP_OK : Success
- ESP_ERR_INVALID_STATE : Manager not initialized or provisioning service already started
- ESP_ERR_NO_MEM : Failed to allocate memory for version string

- ESP_ERR_INVALID_ARG : Null argument

Parameters

- `[in] label`: String indicating the application name.
- `[in] version`: String indicating the application version. There is no constraint on format.
- `[in] capabilities`: Array of strings with capabilities. These could be used by the client side app to know the application registered endpoint capabilities
- `[in] total_capabilities`: Size of capabilities array

esp_err_t **wifi_prov_mgr_endpoint_create**(**const** char *_ep_name_)

Create an additional endpoint and allocate internal resources for it.

This API is to be called by the application if it wants to create an additional endpoint. All additional endpoints will be assigned UUIDs starting from 0xFF54 and so on in the order of execution.

protocomm handler for the created endpoint is to be registered later using wifi_prov_mgr_endpoint_register() after provisioning has started.

Note This API can only be called BEFORE provisioning is started

Note Additional endpoints can be used for configuring client provided parameters other than Wi-Fi credentials, that are necessary for the main application and hence must be set prior to starting the application

Note After session establishment, the additional endpoints must be targeted first by the client side application before sending Wi-Fi configuration, because once Wi-Fi configuration finishes the provisioning service is stopped and hence all endpoints are unregistered

Return

- ESP_OK : Success
- ESP_FAIL : Failure

Parameters

- `[in] ep_name`: unique name of the endpoint

esp_err_t **wifi_prov_mgr_endpoint_register**(**const** char *_ep_name_, *protocomm_req_handler_t* _handler_, void *_user_ctx_)

Register a handler for the previously created endpoint.

This API can be called by the application to register a protocomm handler to any endpoint that was created using wifi_prov_mgr_endpoint_create().

Note This API can only be called AFTER provisioning has started

Note Additional endpoints can be used for configuring client provided parameters other than Wi-Fi credentials, that are necessary for the main application and hence must be set prior to starting the application

Note After session establishment, the additional endpoints must be targeted first by the client side application before sending Wi-Fi configuration, because once Wi-Fi configuration finishes the provisioning service is stopped and hence all endpoints are unregistered

Return

- ESP_OK : Success
- ESP_FAIL : Failure

Parameters

- `[in] ep_name`: Name of the endpoint
- `[in] handler`: Endpoint handler function
- `[in] user_ctx`: User data

void **wifi_prov_mgr_endpoint_unregister**(**const** char *_ep_name_)

Unregister the handler for an endpoint.

This API can be called if the application wants to selectively unregister the handler of an endpoint while the provisioning is still in progress.

All the endpoint handlers are unregistered automatically when the provisioning stops.

Parameters

- `[in] ep_name`: Name of the endpoint

esp_err_t **wifi_prov_mgr_event_handler**(void *_ctx_, *system_event_t* *_event_)

Event handler for provisioning manager.

This is called from the main event handler and controls the provisioning manager’ s internal state machine depending on incoming Wi-Fi events

Note : This function is DEPRECATED, because events are now handled internally using the event loop library, esp_event. Calling this will do nothing and simply return ESP_OK.

Return

- ESP_OK : Event handled successfully

Parameters

- `[in] ctx`: Event context data
- `[in] event`: Event info

esp_err_t **wifi_prov_mgr_get_wifi_state**(*wifi_prov_sta_state_t *state*)

Get state of Wi-Fi Station during provisioning.

Return

- ESP_OK : Successfully retrieved Wi-Fi state
- ESP_FAIL : Provisioning app not running

Parameters

- `[out] state`: Pointer to wifi_prov_sta_state_t variable to be filled

esp_err_t **wifi_prov_mgr_get_wifi_disconnect_reason**(*wifi_prov_sta_fail_reason_t *reason*)

Get reason code in case of Wi-Fi station disconnection during provisioning.

Return

- ESP_OK : Successfully retrieved Wi-Fi disconnect reason
- ESP_FAIL : Provisioning app not running

Parameters

- `[out] reason`: Pointer to wifi_prov_sta_fail_reason_t variable to be filled

esp_err_t **wifi_prov_mgr_configure_sta**(*wifi_config_t *wifi_cfg*)

Runs Wi-Fi as Station with the supplied configuration.

Configures the Wi-Fi station mode to connect to the AP with SSID and password specified in config structure and sets Wi-Fi to run as station.

This is automatically called by provisioning service upon receiving new credentials.

If credentials are to be supplied to the manager via a different mode other than through protocomm, then this API needs to be called.

Event WIFI_PROV_CRED_RECV is emitted after credentials have been applied and Wi-Fi station started

Return

- ESP_OK : Wi-Fi configured and started successfully
- ESP_FAIL : Failed to set configuration

Parameters

- `[in] wifi_cfg`: Pointer to Wi-Fi configuration structure

Structures

struct wifi_prov_event_handler_t

Event handler that is used by the manager while provisioning service is active.

Public Members

wifi_prov_cb_func_t **event_cb**

Callback function to be executed on provisioning events

void ***user_data**

User context data to pass as parameter to callback function

struct wifi_prov_scheme

Structure for specifying the provisioning scheme to be followed by the manager.

Note Ready to use schemes are available:

- wifi_prov_scheme_ble : for provisioning over BLE transport + GATT server
- wifi_prov_scheme_softap : for provisioning over SoftAP transport + HTTP server
- wifi_prov_scheme_console : for provisioning over Serial UART transport + Console (for debugging)

Public Members

esp_err_t (***prov_start**) (*protocomm_t* *pc, void *config)
Function which is to be called by the manager when it is to start the provisioning service associated with a protocomm instance and a scheme specific configuration

esp_err_t (***prov_stop**) (*protocomm_t* *pc)
Function which is to be called by the manager to stop the provisioning service previously associated with a protocomm instance

void *(***new_config**) (void)
Function which is to be called by the manager to generate a new configuration for the provisioning service, that is to be passed to *prov_start()*

void (***delete_config**) (void *config)
Function which is to be called by the manager to delete a configuration generated using *new_config()*

esp_err_t (***set_config_service**) (void *config, **const** char *service_name, **const** char *service_key)
Function which is to be called by the manager to set the service name and key values in the configuration structure

esp_err_t (***set_config_endpoint**) (void *config, **const** char *endpoint_name, uint16_t uuid)
Function which is to be called by the manager to set a protocomm endpoint with an identifying name and UUID in the configuration structure

wifi_mode_t **wifi_mode**
Sets mode of operation of Wi-Fi during provisioning This is set to :

- WIFI_MODE_APSTA for SoftAP transport
- WIFI_MODE_STA for BLE transport

struct wifi_prov_mgr_config_t
Structure for specifying the manager configuration.

Public Members

wifi_prov_scheme_t **scheme**
Provisioning scheme to use. Following schemes are already available:

- wifi_prov_scheme_ble : for provisioning over BLE transport + GATT server
- wifi_prov_scheme_softap : for provisioning over SoftAP transport + HTTP server + mDNS (optional)
- wifi_prov_scheme_console : for provisioning over Serial UART transport + Console (for debugging)

wifi_prov_event_handler_t **scheme_event_handler**
Event handler required by the scheme for incorporating scheme specific behavior while provisioning manager is running. Various options may be provided by the scheme for setting this field. Use WIFI_PROV_EVENT_HANDLER_NONE when not used. When using scheme wifi_prov_scheme_ble, the following options are available:

- WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BTDM
- WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BLE
- WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BT

wifi_prov_event_handler_t **app_event_handler**
Event handler that can be set for the purpose of incorporating application specific behavior. Use WIFI_PROV_EVENT_HANDLER_NONE when not used.

Macros

WIFI_PROV_EVENT_HANDLER_NONE

Event handler can be set to none if not used.

Type Definitions

typedef void (***wifi_prov_cb_func_t**)(void *user_data, *wifi_prov_cb_event_t* event, void *event_data)

typedef struct *wifi_prov_scheme* **wifi_prov_scheme_t**

Structure for specifying the provisioning scheme to be followed by the manager.

Note Ready to use schemes are available:

- wifi_prov_scheme_ble : for provisioning over BLE transport + GATT server
- wifi_prov_scheme_softap : for provisioning over SoftAP transport + HTTP server
- wifi_prov_scheme_console : for provisioning over Serial UART transport + Console (for debugging)

typedef enum *wifi_prov_security* **wifi_prov_security_t**

Security modes supported by the Provisioning Manager.

These are same as the security modes provided by protocomm

Enumerations

enum wifi_prov_cb_event_t

Events generated by manager.

These events are generated in order of declaration and, for the stretch of time between initialization and de-initialization of the manager, each event is signaled only once

Values:

WIFI_PROV_INIT

Emitted when the manager is initialized

WIFI_PROV_START

Indicates that provisioning has started

WIFI_PROV_CRED_RECV

Emitted when Wi-Fi AP credentials are received via `protocomm` endpoint `wifi_config`. The event data in this case is a pointer to the corresponding *`wifi_sta_config_t`* structure

WIFI_PROV_CRED_FAIL

Emitted when device fails to connect to the AP of which the credentials were received earlier on event `WIFI_PROV_CRED_RECV`. The event data in this case is a pointer to the disconnection reason code with type `wifi_prov_sta_fail_reason_t`

WIFI_PROV_CRED_SUCCESS

Emitted when device successfully connects to the AP of which the credentials were received earlier on event `WIFI_PROV_CRED_RECV`

WIFI_PROV_END

Signals that provisioning service has stopped

WIFI_PROV_DEINIT

Signals that manager has been de-initialized

enum wifi_prov_security

Security modes supported by the Provisioning Manager.

These are same as the security modes provided by protocomm

Values:

WIFI_PROV_SECURITY_0 = 0

No security (plain-text communication)

WIFI_PROV_SECURITY_1

This secure communication mode consists of X25519 key exchange

- proof of possession (pop) based authentication
- AES-CTR encryption

Header File

- wifi_provisioning/include/wifi_provisioning/scheme_ble.h

Functions

void **wifi_prov_scheme_ble_event_cb_free_btdm**(void **user_data*, *wifi_prov_cb_event_t event*, void **event_data*)

void **wifi_prov_scheme_ble_event_cb_free_ble**(void **user_data*, *wifi_prov_cb_event_t event*, void **event_data*)

void **wifi_prov_scheme_ble_event_cb_free_bt**(void **user_data*, *wifi_prov_cb_event_t event*, void **event_data*)

esp_err_t **wifi_prov_scheme_ble_set_service_uuid**(uint8_t **uuid128*)

Set the 128 bit GATT service UUID used for provisioning.

This API is used to override the default 128 bit provisioning service UUID, which is 0000ffff-0000-1000-8000-00805f9b34fb.

This must be called before starting provisioning, i.e. before making a call to wifi_prov_mgr_start_provisioning(), otherwise the default UUID will be used.

Note The data being pointed to by the argument must be valid atleast till provisioning is started. Upon start, the manager will store an internal copy of this UUID, and this data can be freed or invalidated afterwords.

Return

- ESP_OK : Success
- ESP_ERR_INVALID_ARG : Null argument

Parameters

- `[in] uuid128`: A custom 128 bit UUID

esp_err_t **wifi_prov_scheme_ble_set_mfg_data**(uint8_t **mfg_data*, ssize_t *mfg_data_len*)

Set manufacturer specific data in scan response.

This must be called before starting provisioning, i.e. before making a call to wifi_prov_mgr_start_provisioning().

Note It is important to understand that length of custom manufacturer data should be within limits. The manufacturer data goes into scan response along with BLE device name. By default, BLE device name length is of 11 Bytes, however it can vary as per application use case. So, one has to honour the scan response data size limits i.e. (mfg_data_len + 2) < 31 - (device_name_length + 2). If the mfg_data length exceeds this limit, the length will be truncated.

Return

- ESP_OK : Success
- ESP_ERR_INVALID_ARG : Null argument

Parameters

- `[in] mfg_data`: Custom manufacturer data
- `[in] mfg_data_len`: Manufacturer data length

Macros

WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BTDM

WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BLE

WIFI_PROV_SCHEME_BLE_EVENT_HANDLER_FREE_BT

Header File

- wifi_provisioning/include/wifi_provisioning/scheme_softap.h

Functions

void **wifi_prov_scheme_softap_set_httpd_handle**(void *_handle_)

Provide HTTPD Server handle externally.

Useful in cases wherein applications need the webserver for some different operations, and do not want the wifi provisioning component to start/stop a new instance.

Note This API should be called before wifi_prov_mgr_start_provisioning()

Parameters

- `[in] handle`: Handle to HTTPD server instance

Header File

- wifi_provisioning/include/wifi_provisioning/scheme_console.h

Header File

- wifi_provisioning/include/wifi_provisioning/wifi_config.h

Functions

esp_err_t **wifi_prov_config_data_handler**(uint32_t *session_id*, **const** uint8_t *_inbuf_, ssize_t *inlen*, uint8_t **_outbuf_, ssize_t *_outlen_, void *_priv_data_)

Handler for receiving and responding to requests from master.

This is to be registered as the `wifi_config` endpoint handler (protocomm `protocomm_req_handler_t`) using `protocomm_add_endpoint()`

Structures

struct wifi_prov_sta_conn_info_t

WiFi STA connected status information.

Public Members

char **ip_addr**[**IP4ADDR_STRLEN_MAX**]

IP Address received by station

char **bssid**[6]

BSSID of the AP to which connection was estalished

char **ssid**[33]

SSID of the to which connection was estalished

uint8_t **channel**

Channel of the AP

uint8_t **auth_mode**

Authorization mode of the AP

struct wifi_prov_config_get_data_t

WiFi status data to be sent in response to `get_status` request from master.

Public Members

wifi_prov_sta_state_t **wifi_state**

WiFi state of the station

wifi_prov_sta_fail_reason_t **fail_reason**

Reason for disconnection (valid only when `wifi_state` is `WIFI_STATION_DISCONNECTED`)

wifi_prov_sta_conn_info_t **conn_info**
Connection information (valid only when `wifi_state` is `WIFI_STATION_CONNECTED`)

struct wifi_prov_config_set_data_t
WiFi config data received by slave during `set_config` request from master.

Public Members

char **ssid**[33]
SSID of the AP to which the slave is to be connected

char **password**[64]
Password of the AP

char **bssid**[6]
BSSID of the AP

uint8_t **channel**
Channel of the AP

struct wifi_prov_config_handlers
Internal handlers for receiving and responding to protocomm requests from master.

This is to be passed as priv_data for protocomm request handler (refer to `wifi_prov_config_data_handler()`) when calling `protocomm_add_endpoint()`.

Public Members

esp_err_t (***get_status_handler**)(*wifi_prov_config_get_data_t* *resp_data, *wifi_prov_ctx_t* **ctx)
Handler function called when connection status of the slave (in WiFi station mode) is requested

esp_err_t (***set_config_handler**)(**const** *wifi_prov_config_set_data_t* *req_data, *wifi_prov_ctx_t* **ctx)
Handler function called when WiFi connection configuration (eg. AP SSID, password, etc.) of the slave (in WiFi station mode) is to be set to user provided values

esp_err_t (***apply_config_handler**)(*wifi_prov_ctx_t* **ctx)
Handler function for applying the configuration that was set in `set_config_handler`. After applying the station may get connected to the AP or may fail to connect. The slave must be ready to convey the updated connection status information when `get_status_handler` is invoked again by the master.

wifi_prov_ctx_t ***ctx**
Context pointer to be passed to above handler functions upon invocation

Type Definitions

typedef struct wifi_prov_ctx **wifi_prov_ctx_t**
Type of context data passed to each get/set/apply handler function set in *`wifi_prov_config_handlers`* structure.

This is passed as an opaque pointer, thereby allowing it be defined later in application code as per requirements.

typedef struct *wifi_prov_config_handlers* **wifi_prov_config_handlers_t**
Internal handlers for receiving and responding to protocomm requests from master.

This is to be passed as priv_data for protocomm request handler (refer to `wifi_prov_config_data_handler()`) when calling `protocomm_add_endpoint()`.

Enumerations

enum wifi_prov_sta_state_t
WiFi STA status for conveying back to the provisioning master.

Values:

WIFI_PROV_STA_CONNECTING

WIFI_PROV_STA_CONNECTED

WIFI_PROV_STA_DISCONNECTED

enum wifi_prov_sta_fail_reason_t

WiFi STA connection fail reason.

Values:

WIFI_PROV_STA_AUTH_ERROR

WIFI_PROV_STA_AP_NOT_FOUND

本部分的 API 示例代码存放在 ESP-IDF 示例项目的 provisioning 目录下。

2.5 存储 API

2.5.1 SPI Flash API

概述

SPI Flash 组件提供外部 flash 数据读取、写入、擦除和内存映射相关的 API 函数，同时也提供了更高层级的，面向分区的 API 函数（定义在分区表 中）。

与 ESP-IDF V4.0 之前的 API 不同，这一版 API 功能并不局限于主 SPI Flash 芯片（即运行程序的 SPI Flash 芯片）。使用不同的芯片指针，您可以通过 SPI0/1 或 HSPI/VSPI 总线访问外部 flash。

注解： ESP-IDF V4.0 之后的 flash API 不再是原子的。因此，如果 flash 操作地址有重叠，且写操作与读操作同时执行，读操作可能会返回一部分写入之前的数据，返回一部分写入之后的数据。

Kconfig 选项*CONFIG_SPI_FLASH_USE_LEGACY_IMPL* 可将 `spi_flash_*` 函数切换至 ESP-IDF V4.0 之前的实现。但是，如果同时使用新旧 API，代码量可能会增多。

即便未启用*CONFIG_SPI_FLASH_USE_LEGACY_IMPL*，加密读取和加密写入操作也均使用旧实现。因此，仅有主 flash 芯片支持加密操作，其他不同片选（经 SPI1 访问的 flash 芯片）则不支持加密操作。

初始化 Flash 设备

在使用 `esp_flash_*` API 之前，您需要在 SPI 总线上初始化芯片。

1. 调用*`spi_bus_initialize()`* 初始化 SPI 总线，此函数将初始化总线上设备间共享的资源，如 I/O、DMA 及中断等。
2. 调用*`spi_bus_add_flash_device()`* 将 flash 设备连接到总线上。然后分配内存，填充 `esp_flash_t` 结构体，同时初始化 CS I/O。
3. 调用*`esp_flash_init()`* 与芯片进行通信。后续操作会依据芯片类型不同而有差异。

注解： 目前，多个 flash 芯片可连接到同一总线。但尚不支持在同一个 SPI 总线上使用 `esp_flash_*` 和 `spi_device_*` 设备。

SPI Flash 访问 API

如下所示为处理 flash 中数据的函数集：

- *`esp_flash_read()`*：将数据从 flash 读取到 RAM；
- *`esp_flash_write()`*：将数据从 RAM 写入到 flash；
- *`esp_flash_erase_region()`*：擦除 flash 中指定区域的数据；

- *esp_flash_erase_chip()*：擦除整个 flash；
- esp_flash_get_chip_size()：返回 menuconfig 中设置的 flash 芯片容量（以字节为单位）。

一般来说，请尽量避免对主 SPI flash 芯片直接使用原始 SPI flash 函数，如需对主 SPI flash 芯片进行操作，请使用分区专用函数。

SPI Flash 容量

SPI flash 容量存储于引导程序映像头部（烧录偏移量为 0x1000）的一个字段。

默认情况下，引导程序写入 flash 时，esptool.py 将引导程序写入 flash 时，会自动检测 SPI flash 容量，同时使用正确容量更新引导程序的头部。您也可以在工程配置中设置 CONFIG_ESPTOOLPY_FLASHSIZE，生成固定的 flash 容量。

如需在运行时覆盖已配置的 flash 容量，请配置 g_rom_flashchip 结构中的 chip_size。esp_flash_* 函数使用此容量（于软件和 ROM 中）进行边界检查。

SPI1 Flash 并发约束

由于 SPI1 flash 也被用于执行固件（通过指令 cache 或数据 cache ），因此在执行读取、写入及擦除操作时，必须禁用这些 cache。这意味着在执行 flash 写操作时，两个 CPU 必须从 IRAM 运行代码，且只能从 DRAM 中读取数据。

如果您使用本文档中 API 函数，上述限制将自动生效且透明（无需您额外关注），但这些限制可能会影响系统中的其他任务的性能。

除 SPI0/1 以外的 SPI 总线上的其它 flash 芯片则不受这种限制。

请参阅应用程序内存分布，查看 IRAM、DRAM 和 flash cache 的区别。

为避免意外读取 flash cache，一个 CPU 在启动 flash 写入或擦除操作时，另一个 CPU 将阻塞，并且在 flash 操作完成前，两个 CPU 上的所有的非 IRAM 安全的中断都会被禁用。

IRAM 安全中断处理程序 如果您需要在 flash 操作期间运行中断处理程序（比如低延迟操作），请在注册中断处理程序 时设置 ESP_INTR_FLAG_IRAM。

请确保中断处理程序访问的所有数据和函数（包括其调用的数据和函数）都存储在 IRAM 或 DRAM 中。

为函数添加 IRAM_ATTR 属性:

```
#include "esp_attr.h"

void IRAM_ATTR gpio_isr_handler(void* arg)
{
    // ...
}
```

为常量添加 DRAM_ATTR 和 DRAM_STR 属性:

```
void IRAM_ATTR gpio_isr_handler(void* arg)
{
   const static DRAM_ATTR uint8_t INDEX_DATA[] = { 45, 33, 12, 0 };
   const static char *MSG = DRAM_STR("I am a string stored in RAM");
}
```

辨别哪些数据应标记为 DRAM_ATTR 可能会比较困难，除非明确标记为 DRAM_ATTR，否则编译器依然可能将某些变量或表达式当做常量（即便没有 const 标记），并将其放入 flash。

如果函数或符号未被正确放入 IRAM/DRAM 中，当中断处理程序在 flash 操作期间从 flash cache 中读取数据，则会产生非法指令异常（这是因为代码未被正确放入 IRAM）或读取垃圾数据（这是因为常数未被正确放入 DRAM），而导致崩溃。

分区表 API

ESP-IDF 工程使用分区表保存 SPI flash 各区信息，包括引导程序、各种应用程序二进制文件、数据及文件系统等。请参考分区表，查看详细信息。

该组件在 esp_partition.h 中声明了一些 API 函数，用以枚举在分区表中找到的分区，并对这些分区执行操作：

- *esp_partition_find()*：在分区表中查找特定类型的条目，返回一个不透明迭代器；
- *esp_partition_get()*：返回一个结构，描述给定迭代器的分区；
- *esp_partition_next()*：将迭代器移至下一个找到的分区；
- *esp_partition_iterator_release()*：释放 esp_partition_find 中返回的迭代器；
- *esp_partition_find_first()*：返回一个结构，描述 esp_partition_find 中找到的第一个分区；
- *esp_partition_read()*、*esp_partition_write()* 和 *esp_partition_erase_range()* 在分区边界内执行，等同于 spi_flash_read()、spi_flash_write() 和 spi_flash_erase_range()。

注解： 请在应用程序代码中使用上述 esp_partition_* API 函数，而非低层级的 spi_flash_* API 函数。分区表 API 函数根据存储在分区表中的数据，进行边界检查并计算在 flash 中的正确偏移量。

SPI Flash 加密

您可以对 SPI flash 内容进行加密，并在硬件层对其进行透明解密。

请参阅*Flash 加密*，查看详细信息。

内存映射 API

ESP32 内存硬件可以将 flash 部分区域映射到指令地址空间和数据地址空间，此映射仅用于读操作。不能通过写入 flash 映射的存储区域来改变 flash 中内容。

Flash 以 64 KB 页为单位进行地址映射。内存映射硬件最多可将 4 MB flash 映射到数据地址空间，将 16 MB flash 映射到指令地址空间。请参考《ESP32 技术参考手册》查看内存映射硬件的详细信息。

请注意，有些 64 KB 页还用于将应用程序映射到内存中，因此实际可用的 64 KB 页会更少一些。

Flash 加密 启用时，使用内存映射区域从 flash 读取数据是解密 flash 的唯一方法，解密需在硬件层进行。

内存映射 API 在 esp_spi_flash.h 和 esp_partition.h 中声明：

- spi_flash_mmap()：将 flash 物理地址区域映射到 CPU 指令空间或数据空间；
- spi_flash_munmap()：取消上述区域的映射；
- *esp_partition_mmap()*：将分区的一部分映射至 CPU 指令空间或数据空间；

spi_flash_mmap() 和*esp_partition_mmap()* 的区别如下：

- spi_flash_mmap()：需要给定一个 64 KB 对齐的物理地址；
- *esp_partition_mmap()*：给定分区内任意偏移量即可，此函数根据需要将返回的指针调整至指向映射内存。

内存映射在 64 KB 块中进行，如果分区已传递给 esp_partition_mmap，则可读取分区外数据。

实现

esp_flash_t 结构包含芯片数据和该 API 的三个重要部分：

1. 主机驱动，为访问芯片提供硬件支持；
2. 芯片驱动，为不同芯片提供兼容性服务；

3. OS 函数，在不同阶段（一级或二级 Boot 或者应用程序阶段）为部分 OS 函数提供支持（如一些锁、延迟）。

主机驱动 主机驱动依赖 `soc/include/hal` 文件夹下 `spi_flash_host_drv.h` 定义的 `spi_flash_host_driver_t` 接口。该接口提供了一些与芯片通信常用的函数。

在 SPI HAL 文件中，有些函数是基于现有的 ESP32 memory-spi 来实现的。但是，由于 ESP32 速度限制，HAL 层无法提供某些读命令的高速实现（所以这些命令根本没有在 HAL 的文件中被实现）。`memspi_host_driver.h` 和 `.c` 文件使用 HAL 提供的 `common_command` 函数实现上述读命令的高速版本，并将所有它实现的及 HAL 函数封装为 `spi_flash_host_driver_t` 供更上层调用。

您也可以实现自己的主机驱动，甚至只通过简单的 GPIO。只要实现了 `spi_flash_host_driver_t` 中所有函数，不管底层硬件是什么，esp_flash API 都可以访问 flash。

芯片驱动 芯片驱动在 `spi_flash_chip_driver.h` 中进行定义，并将主机驱动提供的基本函数进行封装以供 API 层使用。

有些操作需在执行前先发送命令，或在执行后读取状态，因此有些芯片需要不同的命令或值以及通信方式。

`generic chip` 芯片代表了常见的 flash 芯片，其他芯片驱动可以在通用芯片的基础上进行开发。

芯片驱动依赖主机驱动。

OS 函数 OS 函数层提供访问锁和延迟的方法。

该锁定用于解决 SPI Flash 芯片访问和其他函数之间的冲突。例如，经 SPI0/1 访问 flash 芯片时，应当禁用 cache（平时用于取代码和 PSRAM 数据）。另一种情况是，一些没有 CS 线或者 CS 线受软件控制的设备（如通过 SPI 接口的 SD 卡控制）需要在一段时间内独占总线。

延时则用于某些长时操作，需要主机处于等待状态或执行轮询。

顶层 API 将芯片驱动和 OS 函数封装成一个完整的组件，并提供参数检查。

另请参考

- 分区表
- *OTA API* 提供了高层 API 用于更新存储在 flash 中的 app 固件。
- *NVS API* 提供了结构化 API 用于存储 SPI flash 中的碎片数据。

实现细节

必须确保操作期间，两个 CPU 均未从 flash 运行代码，实现细节如下：

- 单核模式下，SDK 在执行 flash 操作前将禁用中断或调度算法。
- 双核模式下，实现细节更为复杂，SDK 需确保两个 CPU 均未运行 flash 代码。

如果有 SPI flash API 在 CPU A（PRO 或 APP）上调用，它使用 `esp_ipc_call` API 在 CPU B 上运行 `spi_flash_op_block_func` 函数。`esp_ipc_call` API 在 CPU B 上唤醒一个高优先级任务，即运行 `spi_flash_op_block_func` 函数。运行该函数将禁用 CPU B 上的 cache，并使用 `s_flash_op_can_start` 旗帜来标志 cache 已禁用。然后，CPU A 上的任务也会禁用 cache 并继续执行 flash 操作。

执行 flash 操作时，CPU A 和 CPU B 仍然可以执行中断操作。默认中断代码均存储于 RAM 中，如果新添加了中断分配 API，则应添加一个标志位以请求在 flash 操作期间禁用该新分配的中断。

Flash 操作完成后，CPU A 上的函数将设置另一标志位，即 `s_flash_op_complete`，用以通知 CPU B 上的任务可以重新启用 cache 并释放 CPU。接着，CPU A 上的函数也重新启用 cache，并将控制权返还给调用者。

另外，所有 API 函数均受互斥量 `s_flash_op_mutex` 保护。

在单核环境中（启用*CONFIG_FREERTOS_UNICORE*），您需要禁用上述两个 cache 以防发生 CPU 间通信。

SPI Flash API 参考

Header File

- spi_flash/include/esp_flash_spi_init.h

Functions

esp_err_t **spi_bus_add_flash_device**(*esp_flash_t* ***out_chip*, **const** *esp_flash_spi_device_config_t* **config*)

Add a SPI Flash device onto the SPI bus.

The bus should be already initialized by `spi_bus_initialization`.

Return

- ESP_ERR_INVALID_ARG: out_chip is NULL, or some field in the config is invalid.
- ESP_ERR_NO_MEM: failed to allocate memory for the chip structures.
- ESP_OK: success.

Parameters

- `out_chip`: Pointer to hold the initialized chip.
- `config`: Configuration of the chips to initialize.

esp_err_t **spi_bus_remove_flash_device**(*esp_flash_t* **chip*)

Remove a SPI Flash device from the SPI bus.

Return

- ESP_ERR_INVALID_ARG: The chip is invalid.
- ESP_OK: success.

Parameters

- `chip`: The flash device to remove.

Structures

struct esp_flash_spi_device_config_t

Configurations for the SPI Flash to init.

Public Members

spi_host_device_t **host_id**

Bus to use.

int **cs_io_num**

GPIO pin to output the CS signal.

esp_flash_io_mode_t **io_mode**

IO mode to read from the Flash.

esp_flash_speed_t **speed**

Speed of the Flash clock.

int **input_delay_ns**

Input delay of the data pins, in ns. Set to 0 if unknown.

int **cs_id**

CS line ID, ignored when not `host_id` is not SPI1_HOST, or `CONFIG_SPI_FLASH_SHARE_SPI1_BUS` is enabled. In this case, the CS line used is automatically assigned by the SPI bus lock.

Header File

- spi_flash/include/esp_flash.h

Functions

esp_err_t **esp_flash_init**(*esp_flash_t *chip*)

Initialise SPI flash chip interface.

This function must be called before any other API functions are called for this chip.

Note Only the `host` and `read_mode` fields of the chip structure must be initialised before this function is called. Other fields may be auto-detected if left set to zero or NULL.

Note If the chip->drv pointer is NULL, chip chip_drv will be auto-detected based on its manufacturer & product IDs. See `esp_flash_registered_flash_drivers` pointer for details of this process.

Return ESP_OK on success, or a flash error code if initialisation fails.

Parameters

- `chip`: Pointer to SPI flash chip to use. If NULL, esp_flash_default_chip is substituted.

bool **esp_flash_chip_driver_initialized**(**const** *esp_flash_t *chip*)

Check if appropriate chip driver is set.

Return true if set, otherwise false.

Parameters

- `chip`: Pointer to SPI flash chip to use. If NULL, esp_flash_default_chip is substituted.

esp_err_t **esp_flash_read_id**(*esp_flash_t *chip*, uint32_t *out_id*)

Read flash ID via the common "RDID" SPI flash command.

ID is a 24-bit value. Lower 16 bits of 'id' are the chip ID, upper 8 bits are the manufacturer ID.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `[out] out_id`: Pointer to receive ID value.

Return ESP_OK on success, or a flash error code if operation failed.

esp_err_t **esp_flash_get_size**(*esp_flash_t *chip*, uint32_t *out_size*)

Detect flash size based on flash ID.

Note Most flash chips use a common format for flash ID, where the lower 4 bits specify the size as a power of 2. If the manufacturer doesn't follow this convention, the size may be incorrectly detected.

Return ESP_OK on success, or a flash error code if operation failed.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `[out] out_size`: Detected size in bytes.

esp_err_t **esp_flash_erase_chip**(*esp_flash_t *chip*)

Erase flash chip contents.

Return ESP_OK on success, or a flash error code if operation failed.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()

esp_err_t **esp_flash_erase_region**(*esp_flash_t *chip*, uint32_t *start*, uint32_t *len*)

Erase a region of the flash chip.

Sector size is specifyed in chip->drv->sector_size field (typically 4096 bytes.) ESP_ERR_INVALID_ARG will be returned if the start & length are not a multiple of this size.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `start`: Address to start erasing flash. Must be sector aligned.
- `len`: Length of region to erase. Must also be sector aligned.

Erase is performed using block (multi-sector) erases where possible (block size is specified in chip->drv->block_erase_size field, typically 65536 bytes). Remaining sectors are erased using individual sector erase commands.

Return ESP_OK on success, or a flash error code if operation failed.

esp_err_t **esp_flash_get_chip_write_protect**(*esp_flash_t *chip*, bool *write_protected*)

Read if the entire chip is write protected.

Note A correct result for this flag depends on the SPI flash chip model and chip_drv in use (via the 'chip->drv' field).

Return ESP_OK on success, or a flash error code if operation failed.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `[out] write_protected`: Pointer to boolean, set to the value of the write protect flag.

esp_err_t **esp_flash_set_chip_write_protect**(*esp_flash_t* *chip, bool *write_protect*)

Set write protection for the SPI flash chip.

Some SPI flash chips may require a power cycle before write protect status can be cleared. Otherwise, write protection can be removed via a follow-up call to this function.

Note Correct behaviour of this function depends on the SPI flash chip model and chip_drv in use (via the 'chip->drv' field).

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `write_protect`: Boolean value for the write protect flag

Return ESP_OK on success, or a flash error code if operation failed.

esp_err_t **esp_flash_get_protectable_regions**(**const** *esp_flash_t* *chip, **const** *esp_flash_region_t* **out_regions, uint32_t *out_num_regions)

Read the list of individually protectable regions of this SPI flash chip.

Note Correct behaviour of this function depends on the SPI flash chip model and chip_drv in use (via the 'chip->drv' field).

Return ESP_OK on success, or a flash error code if operation failed.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `[out] out_regions`: Pointer to receive a pointer to the array of protectable regions of the chip.
- `[out] out_num_regions`: Pointer to an integer receiving the count of protectable regions in the array returned in 'regions'.

esp_err_t **esp_flash_get_protected_region**(*esp_flash_t* *chip, **const** *esp_flash_region_t* *region, bool *out_protected)

Detect if a region of the SPI flash chip is protected.

Note It is possible for this result to be false and write operations to still fail, if protection is enabled for the entire chip.

Note Correct behaviour of this function depends on the SPI flash chip model and chip_drv in use (via the 'chip->drv' field).

Return ESP_OK on success, or a flash error code if operation failed.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `region`: Pointer to a struct describing a protected region. This must match one of the regions returned from esp_flash_get_protectable_regions(···).
- `[out] out_protected`: Pointer to a flag which is set based on the protected status for this region.

esp_err_t **esp_flash_set_protected_region**(*esp_flash_t* *chip, **const** *esp_flash_region_t* *region, bool *protect*)

Update the protected status for a region of the SPI flash chip.

Note It is possible for the region protection flag to be cleared and write operations to still fail, if protection is enabled for the entire chip.

Note Correct behaviour of this function depends on the SPI flash chip model and chip_drv in use (via the 'chip->drv' field).

Return ESP_OK on success, or a flash error code if operation failed.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()

- `region`: Pointer to a struct describing a protected region. This must match one of the regions returned from esp_flash_get_protectable_regions(…).
- `protect`: Write protection flag to set.

esp_err_t **esp_flash_read**(*esp_flash_t* *chip, void *buffer, uint32_t address, uint32_t length)

Read data from the SPI flash chip.

There are no alignment constraints on buffer, address or length.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `buffer`: Pointer to a buffer where the data will be read. To get better performance, this should be in the DRAM and word aligned.
- `address`: Address on flash to read from. Must be less than chip->size field.
- `length`: Length (in bytes) of data to read.

Note If on-chip flash encryption is used, this function returns raw (ie encrypted) data. Use the flash cache to transparently decrypt data.

Return

- ESP_OK: success
- ESP_ERR_NO_MEM: Buffer is in external PSRAM which cannot be concurrently accessed, and a temporary internal buffer could not be allocated.
- or a flash error code if operation failed.

esp_err_t **esp_flash_write**(*esp_flash_t* *chip, **const** void *buffer, uint32_t address, uint32_t length)

Write data to the SPI flash chip.

There are no alignment constraints on buffer, address or length.

Parameters

- `chip`: Pointer to identify flash chip. Must have been successfully initialised via esp_flash_init()
- `address`: Address on flash to write to. Must be previously erased (SPI NOR flash can only write bits 1->0).
- `buffer`: Pointer to a buffer with the data to write. To get better performance, this should be in the DRAM and word aligned.
- `length`: Length (in bytes) of data to write.

Return ESP_OK on success, or a flash error code if operation failed.

esp_err_t **esp_flash_write_encrypted**(*esp_flash_t* *chip, uint32_t address, **const** void *buffer, uint32_t length)

Encrypted and write data to the SPI flash chip using on-chip hardware flash encryption.

Note Both address & length must be 16 byte aligned, as this is the encryption block size

Return

- ESP_OK: on success
- ESP_ERR_NOT_SUPPORTED: encrypted write not supported for this chip.
- ESP_ERR_INVALID_ARG: Either the address, buffer or length is invalid.
- or other flash error code from spi_flash_write_encrypted().

Parameters

- `chip`: Pointer to identify flash chip. Must be NULL (the main flash chip). For other chips, encrypted write is not supported.
- `address`: Address on flash to write to. 16 byte aligned. Must be previously erased (SPI NOR flash can only write bits 1->0).
- `buffer`: Pointer to a buffer with the data to write.
- `length`: Length (in bytes) of data to write. 16 byte aligned.

esp_err_t **esp_flash_read_encrypted**(*esp_flash_t* *chip, uint32_t address, void *out_buffer, uint32_t length)

Read and decrypt data from the SPI flash chip using on-chip hardware flash encryption.

Return

- ESP_OK: on success
- ESP_ERR_NOT_SUPPORTED: encrypted read not supported for this chip.

- or other flash error code from spi_flash_read_encrypted().

Parameters

- `chip`: Pointer to identify flash chip. Must be NULL (the main flash chip). For other chips, encrypted read is not supported.
- `address`: Address on flash to read from.
- `out_buffer`: Pointer to a buffer for the data to read to.
- `length`: Length (in bytes) of data to read.

static bool **esp_flash_is_quad_mode**(**const** *esp_flash_t* *chip)

Returns true if chip is configured for Quad I/O or Quad Fast Read.

Return true if flash works in quad mode, otherwise false

Parameters

- `chip`: Pointer to SPI flash chip to use. If NULL, esp_flash_default_chip is substituted.

Structures

struct esp_flash_region_t

Structure for describing a region of flash.

Public Members

uint32_t **offset**

Start address of this region.

uint32_t **size**

Size of the region.

struct esp_flash_os_functions_t

OS-level integration hooks for accessing flash chips inside a running OS

Public Members

esp_err_t (***start**)(void *arg)

Called before commencing any flash operation. Does not need to be recursive (ie is called at most once for each call to 'end').

esp_err_t (***end**)(void *arg)

Called after completing any flash operation.

esp_err_t (***region_protected**)(void *arg, size_t start_addr, size_t size)

Called before any erase/write operations to check whether the region is limited by the OS

esp_err_t (***delay_us**)(void *arg, unsigned us)

Delay for at least 'us' microseconds. Called in between 'start' and 'end'.

esp_err_t (***yield**)(void *arg)

Yield to other tasks. Called during erase operations.

struct esp_flash_t

Structure to describe a SPI flash chip connected to the system.

Structure must be initialized before use (passed to esp_flash_init()).

Public Members

spi_flash_host_driver_t ***host**

Pointer to hardware-specific "host_driver" structure. Must be initialized before used.

const *spi_flash_chip_t* ***chip_drv**

Pointer to chip-model-specific "adapter" structure. If NULL, will be detected during initialisation.

const *esp_flash_os_functions_t* ***os_func**
Pointer to os-specific hook structure. Call `esp_flash_init_os_functions()` to setup this field, after the host is properly initialized.

void ***os_func_data**
Pointer to argument for os-specific hooks. Left NULL and will be initialized with `os_func`.

esp_flash_io_mode_t **read_mode**
Configured SPI flash read mode. Set before `esp_flash_init` is called.

uint32_t **size**
Size of SPI flash in bytes. If 0, size will be detected during initialisation.

uint32_t **chip_id**
Detected chip id.

Type Definitions

typedef struct *spi_flash_chip_t* **spi_flash_chip_t**
typedef struct *esp_flash_t* **esp_flash_t**

Header File

- soc/include/hal/spi_flash_types.h

Structures

struct spi_flash_trans_t
Definition of a common transaction. Also holds the return value.

Public Members

uint8_t **command**
Command to send, always 8bits.

uint8_t **mosi_len**
Output data length, in bytes.

uint8_t **miso_len**
Input data length, in bytes.

uint8_t **address_bitlen**
Length of address in bits, set to 0 if command does not need an address.

uint32_t **address**
Address to perform operation on.

const uint8_t ***mosi_data**
Output data to salve.

uint8_t ***miso_data**
[out] Input data from slave, little endian

struct spi_flash_host_driver_t
Host driver configuration and context structure.

Public Members

void ***driver_data**
Configuration and static data used by the specific host driver. The type is determined by the host driver.

esp_err_t (***dev_config**) (*spi_flash_host_driver_t* *driver)
Configure the device-related register before transactions. This saves some time to re-configure those registers when we send continuously

esp_err_t (***common_command**) (*spi_flash_host_driver_t* *driver, *spi_flash_trans_t* *t)
Send an user-defined spi transaction to the device.

esp_err_t (***read_id**) (*spi_flash_host_driver_t* *driver, uint32_t *id)
Read flash ID.

void (***erase_chip**) (*spi_flash_host_driver_t* *driver)
Erase whole flash chip.

void (***erase_sector**) (*spi_flash_host_driver_t* *driver, uint32_t start_address)
Erase a specific sector by its start address.

void (***erase_block**) (*spi_flash_host_driver_t* *driver, uint32_t start_address)
Erase a specific block by its start address.

esp_err_t (***read_status**) (*spi_flash_host_driver_t* *driver, uint8_t *out_sr)
Read the status of the flash chip.

esp_err_t (***set_write_protect**) (*spi_flash_host_driver_t* *driver, bool wp)
Disable write protection.

void (***program_page**) (*spi_flash_host_driver_t* *driver, **const** void *buffer, uint32_t address, uint32_t length)
Program a page of the flash. Check `max_write_bytes` for the maximum allowed writing length.

bool (***supports_direct_write**) (*spi_flash_host_driver_t* *driver, **const** void *p)
Check whether need to allocate new buffer to write

bool (***supports_direct_read**) (*spi_flash_host_driver_t* *driver, **const** void *p)
Check whether need to allocate new buffer to read

int **max_write_bytes**
maximum length of program_page

esp_err_t (***read**) (*spi_flash_host_driver_t* *driver, void *buffer, uint32_t address, uint32_t read_len)
Read data from the flash. Check `max_read_bytes` for the maximum allowed reading length.

int **max_read_bytes**
maximum length of read

bool (***host_idle**) (*spi_flash_host_driver_t* *driver)
Check whether the host is idle to perform new operations.

esp_err_t (***configure_host_io_mode**) (*spi_flash_host_driver_t* *driver, uint32_t command, uint32_t addr_bitlen, int dummy_bitlen_base, *esp_flash_io_mode_t* io_mode)
Configure the host to work at different read mode. Responsible to compensate the timing and set IO mode.

void (***poll_cmd_done**) (*spi_flash_host_driver_t* *driver)
Internal use, poll the HW until the last operation is done.

esp_err_t (***flush_cache**) (*spi_flash_host_driver_t* *driver, uint32_t addr, uint32_t size)
For some host (SPI1), they are shared with a cache. When the data is modified, the cache needs to be flushed. Left NULL if not supported.

Macros

ESP_FLASH_SPEED_MIN
Lowest speed supported by the driver, currently 5 MHz.

SPI_FLASH_READ_MODE_MIN
Slowest io mode supported by ESP32, currently SlowRd.

Type Definitions

typedef struct *spi_flash_host_driver_t* **spi_flash_host_driver_t**

Enumerations

enum esp_flash_speed_t

SPI flash clock speed values, always refer to them by the enum rather than the actual value (more speed may be appended into the list).

A strategy to select the maximum allowed speed is to enumerate from the `ESP_FLSH_SPEED_MAX-1` or highest frequency supported by your flash, and decrease the speed until the probing success.

Values:

ESP_FLASH_5MHZ = 0
The flash runs under 5MHz.

ESP_FLASH_10MHZ
The flash runs under 10MHz.

ESP_FLASH_20MHZ
The flash runs under 20MHz.

ESP_FLASH_26MHZ
The flash runs under 26MHz.

ESP_FLASH_40MHZ
The flash runs under 40MHz.

ESP_FLASH_80MHZ
The flash runs under 80MHz.

ESP_FLASH_SPEED_MAX
The maximum frequency supported by the host is `ESP_FLASH_SPEED_MAX-1`.

enum esp_flash_io_mode_t

Mode used for reading from SPI flash.

Values:

SPI_FLASH_SLOWRD = 0
Data read using single I/O, some limits on speed.

SPI_FLASH_FASTRD
Data read using single I/O, no limit on speed.

SPI_FLASH_DOUT
Data read using dual I/O.

SPI_FLASH_DIO
Both address & data transferred using dual I/O.

SPI_FLASH_QOUT
Data read using quad I/O.

SPI_FLASH_QIO
Both address & data transferred using quad I/O.

SPI_FLASH_READ_MODE_MAX
The fastest io mode supported by the host is `ESP_FLASH_READ_MODE_MAX-1`.

分区表 API 参考

Header File

- spi_flash/include/esp_partition.h

Functions

esp_partition_iterator_t **esp_partition_find**(*esp_partition_type_t type*, *esp_partition_subtype_t subtype*, **const** char **label*)

Find partition based on one or more parameters.

Return iterator which can be used to enumerate all the partitions found, or NULL if no partitions were found. Iterator obtained through this function has to be released using esp_partition_iterator_release when not used any more.

Parameters

- `type`: Partition type, one of esp_partition_type_t values or an 8-bit unsigned integer
- `subtype`: Partition subtype, one of esp_partition_subtype_t values or an 8-bit unsigned integer. To find all partitions of given type, use ESP_PARTITION_SUBTYPE_ANY.
- `label`: (optional) Partition label. Set this value if looking for partition with a specific name. Pass NULL otherwise.

const *esp_partition_t* ***esp_partition_find_first**(*esp_partition_type_t* *type*, *esp_partition_subtype_t* *subtype*, **const** char **label*)

Find first partition based on one or more parameters.

Return pointer to *esp_partition_t* structure, or NULL if no partition is found. This pointer is valid for the lifetime of the application.

Parameters

- `type`: Partition type, one of esp_partition_type_t values or an 8-bit unsigned integer
- `subtype`: Partition subtype, one of esp_partition_subtype_t values or an 8-bit unsigned integer To find all partitions of given type, use ESP_PARTITION_SUBTYPE_ANY.
- `label`: (optional) Partition label. Set this value if looking for partition with a specific name. Pass NULL otherwise.

const *esp_partition_t* ***esp_partition_get**(*esp_partition_iterator_t iterator*)

Get *esp_partition_t* structure for given partition.

Return pointer to *esp_partition_t* structure. This pointer is valid for the lifetime of the application.

Parameters

- `iterator`: Iterator obtained using esp_partition_find. Must be non-NULL.

esp_partition_iterator_t **esp_partition_next**(*esp_partition_iterator_t iterator*)

Move partition iterator to the next partition found.

Any copies of the iterator will be invalid after this call.

Return NULL if no partition was found, valid esp_partition_iterator_t otherwise.

Parameters

- `iterator`: Iterator obtained using esp_partition_find. Must be non-NULL.

void **esp_partition_iterator_release**(*esp_partition_iterator_t iterator*)

Release partition iterator.

Parameters

- `iterator`: Iterator obtained using esp_partition_find. Must be non-NULL.

const *esp_partition_t* ***esp_partition_verify**(**const** *esp_partition_t* **partition*)

Verify partition data.

Given a pointer to partition data, verify this partition exists in the partition table (all fields match.)

This function is also useful to take partition data which may be in a RAM buffer and convert it to a pointer to the permanent partition data stored in flash.

Pointers returned from this function can be compared directly to the address of any pointer returned from *esp_partition_get()*, as a test for equality.

Return

- If partition not found, returns NULL.
- If found, returns a pointer to the *esp_partition_t* structure in flash. This pointer is always valid for the lifetime of the application.

Parameters

- `partition`: Pointer to partition data to verify. Must be non-NULL. All fields of this structure must match the partition table entry in flash for this function to return a successful match.

esp_err_t **esp_partition_read**(**const** *esp_partition_t* **partition*, size_t *src_offset*, void **dst*, size_t *size*)

Read data from the partition.

Return ESP_OK, if data was read successfully; ESP_ERR_INVALID_ARG, if src_offset exceeds partition size; ESP_ERR_INVALID_SIZE, if read would go out of bounds of the partition; or one of error codes from lower-level flash driver.

Parameters

- `partition`: Pointer to partition structure obtained using esp_partition_find_first or esp_partition_get. Must be non-NULL.
- `dst`: Pointer to the buffer where data should be stored. Pointer must be non-NULL and buffer must be at least 'size' bytes long.
- `src_offset`: Address of the data to be read, relative to the beginning of the partition.
- `size`: Size of data to be read, in bytes.

esp_err_t **esp_partition_write**(**const** *esp_partition_t* **partition*, size_t *dst_offset*, **const** void **src*, size_t *size*)

Write data to the partition.

Before writing data to flash, corresponding region of flash needs to be erased. This can be done using esp_partition_erase_range function.

Partitions marked with an encryption flag will automatically be written via the spi_flash_write_encrypted() function. If writing to an encrypted partition, all write offsets and lengths must be multiples of 16 bytes. See the spi_flash_write_encrypted() function for more details. Unencrypted partitions do not have this restriction.

Note Prior to writing to flash memory, make sure it has been erased with esp_partition_erase_range call.

Return ESP_OK, if data was written successfully; ESP_ERR_INVALID_ARG, if dst_offset exceeds partition size; ESP_ERR_INVALID_SIZE, if write would go out of bounds of the partition; or one of error codes from lower-level flash driver.

Parameters

- `partition`: Pointer to partition structure obtained using esp_partition_find_first or esp_partition_get. Must be non-NULL.
- `dst_offset`: Address where the data should be written, relative to the beginning of the partition.
- `src`: Pointer to the source buffer. Pointer must be non-NULL and buffer must be at least 'size' bytes long.
- `size`: Size of data to be written, in bytes.

esp_err_t **esp_partition_erase_range**(**const** *esp_partition_t* **partition*, size_t *offset*, size_t *size*)

Erase part of the partition.

Return ESP_OK, if the range was erased successfully; ESP_ERR_INVALID_ARG, if iterator or dst are NULL; ESP_ERR_INVALID_SIZE, if erase would go out of bounds of the partition; or one of error codes from lower-level flash driver.

Parameters

- `partition`: Pointer to partition structure obtained using esp_partition_find_first or esp_partition_get. Must be non-NULL.
- `offset`: Offset from the beginning of partition where erase operation should start. Must be aligned to 4 kilobytes.
- `size`: Size of the range which should be erased, in bytes. Must be divisible by 4 kilobytes.

esp_err_t **esp_partition_mmap**(**const** *esp_partition_t* **partition*, size_t *offset*, size_t *size*, spi_flash_mmap_memory_t *memory*, **const** void ***out_ptr*, spi_flash_mmap_handle_t **out_handle*)

Configure MMU to map partition into data memory.

Unlike spi_flash_mmap function, which requires a 64kB aligned base address, this function doesn't impose such a requirement. If offset results in a flash address which is not aligned to 64kB boundary, address will be rounded to the lower 64kB boundary, so that mapped region includes requested range. Pointer returned via out_ptr argument will be adjusted to point to the requested offset (not necessarily to the beginning of mmap-ed region).

To release mapped memory, pass handle returned via out_handle argument to spi_flash_munmap function.

Return ESP_OK, if successful

Parameters

- `partition`: Pointer to partition structure obtained using esp_partition_find_first or esp_partition_get. Must be non-NULL.
- `offset`: Offset from the beginning of partition where mapping should start.
- `size`: Size of the area to be mapped.
- `memory`: Memory space where the region should be mapped
- `out_ptr`: Output, pointer to the mapped memory region
- `out_handle`: Output, handle which should be used for spi_flash_munmap call

esp_err_t **esp_partition_get_sha256**(**const** *esp_partition_t* **partition*, uint8_t **sha_256*)

Get SHA-256 digest for required partition.

For apps with SHA-256 appended to the app image, the result is the appended SHA-256 value for the app image content. The hash is verified before returning, if app content is invalid then the function returns ESP_ERR_IMAGE_INVALID. For apps without SHA-256 appended to the image, the result is the SHA-256 of all bytes in the app image. For other partition types, the result is the SHA-256 of the entire partition.

Return

- ESP_OK: In case of successful operation.
- ESP_ERR_INVALID_ARG: The size was 0 or the sha_256 was NULL.
- ESP_ERR_NO_MEM: Cannot allocate memory for sha256 operation.
- ESP_ERR_IMAGE_INVALID: App partition doesn't contain a valid app image.
- ESP_FAIL: An allocation error occurred.

Parameters

- `[in] partition`: Pointer to info for partition containing app or data. (fields: address, size and type, are required to be filled).
- `[out] sha_256`: Returned SHA-256 digest for a given partition.

bool **esp_partition_check_identity**(**const** *esp_partition_t* **partition_1*, **const** *esp_partition_t* **partition_2*)

Check for the identity of two partitions by SHA-256 digest.

Return

- True: In case of the two firmware is equal.
- False: Otherwise

Parameters

- `[in] partition_1`: Pointer to info for partition 1 containing app or data. (fields: address, size and type, are required to be filled).
- `[in] partition_2`: Pointer to info for partition 2 containing app or data. (fields: address, size and type, are required to be filled).

esp_err_t **esp_partition_register_external**(*esp_flash_t* **flash_chip*, size_t *offset*, size_t *size*, **const** char **label*, *esp_partition_type_t* *type*, *esp_partition_subtype_t* *subtype*, **const** *esp_partition_t* ***out_partition*)

Register a partition on an external flash chip.

This API allows designating certain areas of external flash chips (identified by the *esp_flash_t* structure) as partitions. This allows using them with components which access SPI flash through the esp_partition API.

Return

- ESP_OK on success
- ESP_ERR_NOT_SUPPORTED if CONFIG_CONFIG_SPI_FLASH_USE_LEGACY_IMPL is enabled
- ESP_ERR_NO_MEM if memory allocation has failed
- ESP_ERR_INVALID_ARG if the new partition overlaps another partition on the same flash chip
- ESP_ERR_INVALID_SIZE if the partition doesn't fit into the flash chip size

Parameters

- `flash_chip`: Pointer to the structure identifying the flash chip
- `offset`: Address in bytes, where the partition starts
- `size`: Size of the partition in bytes

- `label`: Partition name
- `type`: One of the partition types (ESP_PARTITION_TYPE_*), or an integer. Note that applications can not be booted from external flash chips, so using ESP_PARTITION_TYPE_APP is not supported.
- `subtype`: One of the partition subtypes (ESP_PARTITION_SUBTYPE_*), or an integer.
- `[out] out_partition`: Output, if non-NULL, receives the pointer to the resulting *esp_partition_t* structure

esp_err_t **esp_partition_deregister_external**(**const** *esp_partition_t* **partition*)

Deregister the partition previously registered using esp_partition_register_external.

Return

- ESP_OK on success
- ESP_ERR_NOT_FOUND if the partition pointer is not found
- ESP_ERR_INVALID_ARG if the partition comes from the partition table
- ESP_ERR_INVALID_ARG if the partition was not registered using esp_partition_register_external function.

Parameters

- `partition`: pointer to the partition structure obtained from esp_partition_register_external,

Structures

struct esp_partition_t

partition information structure

This is not the format in flash, that format is esp_partition_info_t.

However, this is the format used by this API.

Public Members

esp_flash_t ***flash_chip**

SPI flash chip on which the partition resides

esp_partition_type_t **type**

partition type (app/data)

esp_partition_subtype_t **subtype**

partition subtype

uint32_t **address**

starting address of the partition in flash

uint32_t **size**

size of the partition, in bytes

char **label**[17]

partition label, zero-terminated ASCII string

bool **encrypted**

flag is set to true if partition is encrypted

Macros

ESP_PARTITION_SUBTYPE_OTA(i)

Convenience macro to get esp_partition_subtype_t value for the i-th OTA partition.

Type Definitions

typedef struct esp_partition_iterator_opaque_ ***esp_partition_iterator_t**

Opaque partition iterator type.

Enumerations

enum esp_partition_type_t

Partition type.

Note Partition types with integer value 0x00-0x3F are reserved for partition types defined by ESP-IDF. Any other integer value 0x40-0xFE can be used by individual applications, without restriction.

Values:

ESP_PARTITION_TYPE_APP = 0x00
Application partition type.

ESP_PARTITION_TYPE_DATA = 0x01
Data partition type.

enum esp_partition_subtype_t

Partition subtype.

Application-defined partition types (0x40-0xFE) can set any numeric subtype value.

Note These ESP-IDF-defined partition subtypes apply to partitions of type ESP_PARTITION_TYPE_APP and ESP_PARTITION_TYPE_DATA.

Values:

ESP_PARTITION_SUBTYPE_APP_FACTORY = 0x00
Factory application partition.

ESP_PARTITION_SUBTYPE_APP_OTA_MIN = 0x10
Base for OTA partition subtypes.

ESP_PARTITION_SUBTYPE_APP_OTA_0 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 0
OTA partition 0.

ESP_PARTITION_SUBTYPE_APP_OTA_1 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 1
OTA partition 1.

ESP_PARTITION_SUBTYPE_APP_OTA_2 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 2
OTA partition 2.

ESP_PARTITION_SUBTYPE_APP_OTA_3 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 3
OTA partition 3.

ESP_PARTITION_SUBTYPE_APP_OTA_4 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 4
OTA partition 4.

ESP_PARTITION_SUBTYPE_APP_OTA_5 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 5
OTA partition 5.

ESP_PARTITION_SUBTYPE_APP_OTA_6 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 6
OTA partition 6.

ESP_PARTITION_SUBTYPE_APP_OTA_7 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 7
OTA partition 7.

ESP_PARTITION_SUBTYPE_APP_OTA_8 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 8
OTA partition 8.

ESP_PARTITION_SUBTYPE_APP_OTA_9 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 9
OTA partition 9.

ESP_PARTITION_SUBTYPE_APP_OTA_10 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 10
OTA partition 10.

ESP_PARTITION_SUBTYPE_APP_OTA_11 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 11
OTA partition 11.

ESP_PARTITION_SUBTYPE_APP_OTA_12 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 12
OTA partition 12.

ESP_PARTITION_SUBTYPE_APP_OTA_13 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 13
OTA partition 13.

ESP_PARTITION_SUBTYPE_APP_OTA_14 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 14
OTA partition 14.

ESP_PARTITION_SUBTYPE_APP_OTA_15 = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 15
OTA partition 15.

ESP_PARTITION_SUBTYPE_APP_OTA_MAX = *ESP_PARTITION_SUBTYPE_APP_OTA_MIN* + 16
Max subtype of OTA partition.

ESP_PARTITION_SUBTYPE_APP_TEST = 0x20
Test application partition.

ESP_PARTITION_SUBTYPE_DATA_OTA = 0x00
OTA selection partition.

ESP_PARTITION_SUBTYPE_DATA_PHY = 0x01
PHY init data partition.

ESP_PARTITION_SUBTYPE_DATA_NVS = 0x02
NVS partition.

ESP_PARTITION_SUBTYPE_DATA_COREDUMP = 0x03
COREDUMP partition.

ESP_PARTITION_SUBTYPE_DATA_NVS_KEYS = 0x04
Partition for NVS keys.

ESP_PARTITION_SUBTYPE_DATA_EFUSE_EM = 0x05
Partition for emulate eFuse bits.

ESP_PARTITION_SUBTYPE_DATA_UNDEFINED = 0x06
Undefined (or unspecified) data partition.

ESP_PARTITION_SUBTYPE_DATA_ESPHTTPD = 0x80
ESPHTTPD partition.

ESP_PARTITION_SUBTYPE_DATA_FAT = 0x81
FAT partition.

ESP_PARTITION_SUBTYPE_DATA_SPIFFS = 0x82
SPIFFS partition.

ESP_PARTITION_SUBTYPE_ANY = 0xff
Used to search for partitions with any subtype.

Flash 加密 API 参考

Header File

- bootloader_support/include/esp_flash_encrypt.h

Functions

static bool **esp_flash_encryption_enabled**(void)
Is flash encryption currently enabled in hardware?

Flash encryption is enabled if the FLASH_CRYPT_CNT efuse has an odd number of bits set.

Return true if flash encryption is enabled.

esp_err_t **esp_flash_encrypt_check_and_update**(void)

esp_err_t **esp_flash_encrypt_region**(uint32_t *src_addr*, size_t *data_length*)
Encrypt-in-place a block of flash sectors.

Note This function resets RTC_WDT between operations with sectors.

Return ESP_OK if all operations succeeded, ESP_ERR_FLASH_OP_FAIL if SPI flash fails, ESP_ERR_FLASH_OP_TIMEOUT if flash times out.

Parameters

- `src_addr`: Source offset in flash. Should be multiple of 4096 bytes.
- `data_length`: Length of data to encrypt in bytes. Will be rounded up to next multiple of 4096 bytes.

void **esp_flash_write_protect_crypt_cnt**(void)

Write protect FLASH_CRYPT_CNT.

Intended to be called as a part of boot process if flash encryption is enabled but secure boot is not used. This should protect against serial re-flashing of an unauthorised code in absence of secure boot.

Note On ESP32 V3 only, write protecting FLASH_CRYPT_CNT will also prevent disabling UART Download Mode. If both are wanted, call esp_efuse_disable_rom_download_mode() before calling this function.

esp_flash_enc_mode_t **esp_get_flash_encryption_mode**(void)

Return the flash encryption mode.

The API is called during boot process but can also be called by application to check the current flash encryption mode of ESP32

Return

void **esp_flash_encryption_init_checks**(void)

Check the flash encryption mode during startup.

Verifies the flash encryption config during startup:

Note This function is called automatically during app startup, it doesn't need to be called from the app.

- Correct any insecure flash encryption settings if hardware Secure Boot is enabled.
- Log warnings if the efuse config doesn't match the project config in any way

Enumerations

enum esp_flash_enc_mode_t

Values:

ESP_FLASH_ENC_MODE_DISABLED

ESP_FLASH_ENC_MODE_DEVELOPMENT

ESP_FLASH_ENC_MODE_RELEASE

2.5.2 SD/SDIO/MMC 驱动程序

概述

SD/SDIO/MMC 驱动是一种基于 SDMMC 和 SD SPI 主机驱动的协议级驱动程序，目前已支持 SD 存储器、SDIO 卡和 eMMC 芯片。

SDMMC 主机驱动和 SD SPI 主机驱动（driver/include/driver/sdmmc_host.h）为以下功能提供 API：

- 发送命令至从设备
- 接收和发送数据
- 处理总线错误

初始化函数及配置函数：

- 如需初始化和配置 SD SPI 主机，请参阅*SD SPI 主机 API*

本文档中所述的 SDMMC 协议层仅处理 SD 协议相关事项，例如卡初始化和数据传输命令。

协议层通过`sdmmc_host_t`结构体和主机协同工作，该结构体包含指向主机各类函数的指针。

应用示例

ESP-IDF storage/sd_card 目录下提供了 SDMMC 驱动与 FatFs 库组合使用的示例，演示了先初始化卡，然后使用 POSIX 和 C 库 API 向卡读写数据。请参考示例目录下 README.md 文件，查看更多详细信息。

协议层 API

协议层具备 *sdmmc_host_t* 结构体，此结构体描述了 SD/MMC 主机驱动，列出了其功能，并提供指向驱动程序函数的指针。协议层将卡信息储存于 *sdmmc_card_t* 结构体中。向 SD/MMC 主机发送命令时，协议层调用时需要一个 *sdmmc_command_t* 结构体来描述命令、参数、预期返回值和需传输的数据（如有）。

用于 SD 存储卡的 API

1. 初始化主机，请调用主机驱动函数，例如 `sdmmc_host_init()` 和 `sdmmc_host_init_slot()`；
2. 初始化卡，请调用 *sdmmc_card_init()*，并将参数 `host`（即主机驱动信息）和参数 `card`（指向 *sdmmc_card_t* 结构体的指针）传递给此函数。函数运行结束后，将会向 *sdmmc_card_t* 结构体填充该卡的信息；
3. 读取或写入卡的扇区，请分别调用 *sdmmc_read_sectors()* 和 *sdmmc_write_sectors()*，并将参数 `card`（指向卡信息结构的指针）传递给函数；
4. 如果不再使用该卡，请调用主机驱动函数，例如 `sdmmc_host_deinit()`，以禁用主机外设，并释放驱动程序分配的资源。

用于 eMMC 芯片的 API 从协议层的角度而言，eMMC 存储芯片与 SD 存储卡相同。尽管 eMMC 是芯片，不具备卡的外形，但由于协议相似 (*sdmmc_card_t*, *sdmmc_card_init*)，用于 SD 卡的一些概念同样适用于 eMMC 芯片。注意，eMMC 芯片不可通过 SPI 使用，因此它与 SD API 主机驱动不兼容。

如需初始化 eMMC 内存并执行读/写操作，请参照上一章节 SD 卡操作步骤。

用于 SDIO 卡的 API SDIO 卡初始化和检测过程与 SD 存储卡相同，唯一的区别是 SDIO 模式下数据传输命令不同。

在卡初始化和卡检测（通过运行 *sdmmc_card_init()*）期间，驱动仅配置 SDIO 卡如下寄存器：

1. I/O 中止 (0x06) 寄存器：在该寄存器中设置 RES 位可重置卡的 I/O 部分；
2. 总线接口控制 (0x07) 寄存器：如果主机和插槽配置中启用 4 线模式，则驱动程序会尝试在该寄存器中设置总线宽度字段。如果字段设置成功，则从机支持 4 线模式，主机也切换至 4 线模式；
3. 高速（0x13）寄存器：如果主机配置中启用高速模式，则会在该寄存器中设置 SHS 位。

注意，驱动程序不会在 (1) I/O 使能寄存器和 Int 使能寄存器，及 (2) I/O 块大小中，设置任何位。应用程序可通过调用 *sdmmc_io_write_byte()* 来设置相关位。

如需设置卡配置或传输数据，请根据您的具体情况选择下表中的函数：

操作	读函数	写函数
使用 IO_RW_DIRECT (CMD52) 读写单个字节。	*sdmmc_io_read_byte()*	*sdmmc_io_write_byte()*
使用 IO_RW_EXTENDED (CMD53) 的字节模式读写多个字节。	*sdmmc_io_read_bytes()*	*sdmmc_io_write_bytes()*
块模式下，使用 IO_RW_EXTENDED (CMD53) 读写数据块。	*sdmmc_io_read_blocks()*	*sdmmc_io_write_blocks()*

使用 *sdmmc_io_enable_int()* 函数，应用程序可启用 SDIO 中断。

在单线模式下使用 SDIO 时，还需要连接 D1 线来启用 SDIO 中断。

如果您需要应用程序保持等待直至发生 SDIO 中断，请使用 *sdmmc_io_wait_int()* 函数。

复合卡（存储 + SDIO） 该驱动程序不支持 SDIO 复合卡，复合卡会被视为 SDIO 卡。

线程安全 多数应用程序仅需在一个任务中使用协议层。因此，协议层在 `sdmmc_card_t` 结构体或在访问 SDMMC 或 SD SPI 主机驱动程序时不使用任何类型的锁。这种锁通常在较高层级实现，例如文件系统驱动程序。

API 参考

Header File

- sdmmc/include/sdmmc_cmd.h

Functions

esp_err_t **sdmmc_card_init**(**const** *sdmmc_host_t* **host*, *sdmmc_card_t* **out_card*)

Probe and initialize SD/MMC card using given host

Note Only SD cards (SDSC and SDHC/SDXC) are supported now. Support for MMC/eMMC cards will be added later.

Return

- ESP_OK on success
- One of the error codes from SDMMC host controller

Parameters

- `host`: pointer to structure defining host controller
- `out_card`: pointer to structure which will receive information about the card when the function completes

void **sdmmc_card_print_info**(FILE **stream*, **const** *sdmmc_card_t* **card*)

Print information about the card to a stream.

Parameters

- `stream`: stream obtained using fopen or fdopen
- `card`: card information structure initialized using sdmmc_card_init

esp_err_t **sdmmc_write_sectors**(*sdmmc_card_t* **card*, **const** void **src*, size_t *start_sector*, size_t *sector_count*)

Write given number of sectors to SD/MMC card

Return

- ESP_OK on success
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `src`: pointer to data buffer to read data from; data size must be equal to sector_count * card->csd.sector_size
- `start_sector`: sector where to start writing
- `sector_count`: number of sectors to write

esp_err_t **sdmmc_read_sectors**(*sdmmc_card_t* **card*, void **dst*, size_t *start_sector*, size_t *sector_count*)

Write given number of sectors to SD/MMC card

Return

- ESP_OK on success
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `dst`: pointer to data buffer to write into; buffer size must be at least sector_count * card->csd.sector_size
- `start_sector`: sector where to start reading
- `sector_count`: number of sectors to read

esp_err_t **sdmmc_io_read_byte**(*sdmmc_card_t* *card, uint32_t *function*, uint32_t *reg*, uint8_t *out_byte)

Read one byte from an SDIO card using IO_RW_DIRECT (CMD52)

Return

- ESP_OK on success
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `function`: IO function number
- `reg`: byte address within IO function
- `[out] out_byte`: output, receives the value read from the card

esp_err_t **sdmmc_io_write_byte**(*sdmmc_card_t* *card, uint32_t *function*, uint32_t *reg*, uint8_t *in_byte*, uint8_t *out_byte)

Write one byte to an SDIO card using IO_RW_DIRECT (CMD52)

Return

- ESP_OK on success
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `function`: IO function number
- `reg`: byte address within IO function
- `in_byte`: value to be written
- `[out] out_byte`: if not NULL, receives new byte value read from the card (read-after-write).

esp_err_t **sdmmc_io_read_bytes**(*sdmmc_card_t* *card, uint32_t *function*, uint32_t *addr*, void *dst, size_t *size*)

Read multiple bytes from an SDIO card using IO_RW_EXTENDED (CMD53)

This function performs read operation using CMD53 in byte mode. For block mode, see sdmmc_io_read_blocks.

Return

- ESP_OK on success
- ESP_ERR_INVALID_SIZE if size exceeds 512 bytes
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `function`: IO function number
- `addr`: byte address within IO function where reading starts
- `dst`: buffer which receives the data read from card
- `size`: number of bytes to read

esp_err_t **sdmmc_io_write_bytes**(*sdmmc_card_t* *card, uint32_t *function*, uint32_t *addr*, **const** void *src, size_t *size*)

Write multiple bytes to an SDIO card using IO_RW_EXTENDED (CMD53)

This function performs write operation using CMD53 in byte mode. For block mode, see sdmmc_io_write_blocks.

Return

- ESP_OK on success
- ESP_ERR_INVALID_SIZE if size exceeds 512 bytes
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `function`: IO function number
- `addr`: byte address within IO function where writing starts
- `src`: data to be written
- `size`: number of bytes to write

esp_err_t **sdmmc_io_read_blocks**(*sdmmc_card_t* **card*, uint32_t *function*, uint32_t *addr*, void **dst*, size_t *size*)

Read blocks of data from an SDIO card using IO_RW_EXTENDED (CMD53)

This function performs read operation using CMD53 in block mode. For byte mode, see sdmmc_io_read_bytes.

Return

- ESP_OK on success
- ESP_ERR_INVALID_SIZE if size is not divisible by 512 bytes
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `function`: IO function number
- `addr`: byte address within IO function where writing starts
- `dst`: buffer which receives the data read from card
- `size`: number of bytes to read, must be divisible by the card block size.

esp_err_t **sdmmc_io_write_blocks**(*sdmmc_card_t* **card*, uint32_t *function*, uint32_t *addr*, **const** void **src*, size_t *size*)

Write blocks of data to an SDIO card using IO_RW_EXTENDED (CMD53)

This function performs write operation using CMD53 in block mode. For byte mode, see sdmmc_io_write_bytes.

Return

- ESP_OK on success
- ESP_ERR_INVALID_SIZE if size is not divisible by 512 bytes
- One of the error codes from SDMMC host controller

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `function`: IO function number
- `addr`: byte address within IO function where writing starts
- `src`: data to be written
- `size`: number of bytes to read, must be divisible by the card block size.

esp_err_t **sdmmc_io_enable_int**(*sdmmc_card_t* **card*)

Enable SDIO interrupt in the SDMMC host

Return

- ESP_OK on success
- ESP_ERR_NOT_SUPPORTED if the host controller does not support IO interrupts

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init

esp_err_t **sdmmc_io_wait_int**(*sdmmc_card_t* **card*, TickType_t *timeout_ticks*)

Block until an SDIO interrupt is received

Slave uses D1 line to signal interrupt condition to the host. This function can be used to wait for the interrupt.

Return

- ESP_OK if the interrupt is received
- ESP_ERR_NOT_SUPPORTED if the host controller does not support IO interrupts
- ESP_ERR_TIMEOUT if the interrupt does not happen in timeout_ticks

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `timeout_ticks`: time to wait for the interrupt, in RTOS ticks

esp_err_t **sdmmc_io_get_cis_data**(*sdmmc_card_t* **card*, uint8_t **out_buffer*, size_t *buffer_size*, size_t **inout_cis_size*)

Get the data of CIS region of a SDIO card.

You may provide a buffer not sufficient to store all the CIS data. In this case, this functions store as much data into your buffer as possible. Also, this function will try to get and return the size required for you.

Return

- ESP_OK: on success
- ESP_ERR_INVALID_RESPONSE: if the card does not (correctly) support CIS.
- ESP_ERR_INVALID_SIZE: CIS_CODE_END found, but buffer_size is less than required size, which is stored in the inout_cis_size then.
- ESP_ERR_NOT_FOUND: if the CIS_CODE_END not found. Increase input value of inout_cis_size or set it to 0, if you still want to search for the end; output value of inout_cis_size is invalid in this case.
- and other error code return from sdmmc_io_read_bytes

Parameters

- `card`: pointer to card information structure previously initialized using sdmmc_card_init
- `out_buffer`: Output buffer of the CIS data
- `buffer_size`: Size of the buffer.
- `inout_cis_size`: Mandatory, pointer to a size, input and output.
  - input: Limitation of maximum searching range, should be 0 or larger than buffer_size. The function searches for CIS_CODE_END until this range. Set to 0 to search infinitely.
  - output: The size required to store all the CIS data, if CIS_CODE_END is found.

esp_err_t **sdmmc_io_print_cis_info**(uint8_t **buffer*, size_t *buffer_size*, FILE **fp*)

Parse and print the CIS information of a SDIO card.

Note Not all the CIS codes and all kinds of tuples are supported. If you see some unresolved code, you can add the parsing of these code in sdmmc_io.c and contribute to the IDF through the Github repository.

```
        using sdmmc_card_init
```

Return

- ESP_OK: on success
- ESP_ERR_NOT_SUPPORTED: if the value from the card is not supported to be parsed.
- ESP_ERR_INVALID_SIZE: if the CIS size fields are not correct.

Parameters

- `buffer`: Buffer to parse
- `buffer_size`: Size of the buffer.
- `fp`: File pointer to print to, set to NULL to print to stdout.

Header File

- driver/include/driver/sdmmc_types.h

Structures

struct sdmmc_csd_t

Decoded values from SD card Card Specific Data register

Public Members

int **csd_ver**

CSD structure format

int **mmc_ver**

MMC version (for CID format)

int **capacity**

total number of sectors

int **sector_size**

sector size in bytes

int **read_block_len**

block length for reads

int **card_command_class**
Card Command Class for SD

int **tr_speed**
Max transfer speed

struct sdmmc_cid_t
Decoded values from SD card Card IDentification register

Public Members

int **mfg_id**
manufacturer identification number

int **oem_id**
OEM/product identification number

char **name**[8]
product name (MMC v1 has the longest)

int **revision**
product revision

int **serial**
product serial number

int **date**
manufacturing date

struct sdmmc_scr_t
Decoded values from SD Configuration Register

Public Members

int **sd_spec**
SD Physical layer specification version, reported by card

int **bus_width**
bus widths supported by card: BIT(0) —1-bit bus, BIT(2) —4-bit bus

struct sdmmc_ext_csd_t
Decoded values of Extended Card Specific Data

Public Members

uint8_t **power_class**
Power class used by the card

struct sdmmc_switch_func_rsp_t
SD SWITCH_FUNC response buffer

Public Members

uint32_t **data**[512 / 8 / sizeof(uint32_t)]
response data

struct sdmmc_command_t
SD/MMC command information

Public Members

uint32_t **opcode**
SD or MMC command index

uint32_t **arg**
SD/MMC command argument

sdmmc_response_t **response**
response buffer

void ***data**
buffer to send or read into

size_t **datalen**
length of data buffer

size_t **blklen**
block length

int **flags**
see below

esp_err_t **error**
error returned from transfer

int **timeout_ms**
response timeout, in milliseconds

struct sdmmc_host_t
SD/MMC Host description

This structure defines properties of SD/MMC host and functions of SD/MMC host which can be used by upper layers.

Public Members

uint32_t **flags**
flags defining host properties

int **slot**
slot number, to be passed to host functions

int **max_freq_khz**
max frequency supported by the host

float **io_voltage**
I/O voltage used by the controller (voltage switching is not supported)

esp_err_t (***init**)(void)
Host function to initialize the driver

esp_err_t (***set_bus_width**)(int slot, size_t width)
host function to set bus width

size_t (***get_bus_width**)(int slot)
host function to get bus width

esp_err_t (***set_bus_ddr_mode**)(int slot, bool ddr_enable)
host function to set DDR mode

esp_err_t (***set_card_clk**)(int slot, uint32_t freq_khz)
host function to set card clock frequency

esp_err_t (***do_transaction**)(int slot, *sdmmc_command_t* *cmdinfo)
host function to do a transaction

esp_err_t (***deinit**) (void)
host function to deinitialize the driver

esp_err_t (***deinit_p**) (int slot)
host function to deinitialize the driver, called with the `slot`

esp_err_t (***io_int_enable**) (int slot)
Host function to enable SDIO interrupt line

esp_err_t (***io_int_wait**) (int slot, TickType_t timeout_ticks)
Host function to wait for SDIO interrupt line to be active

int **command_timeout_ms**
timeout, in milliseconds, of a single command. Set to 0 to use the default value.

struct sdmmc_card_t
SD/MMC card information structure

Public Members

sdmmc_host_t **host**
Host with which the card is associated

uint32_t **ocr**
OCR (Operation Conditions Register) value

sdmmc_cid_t **cid**
decoded CID (Card IDentification) register value

sdmmc_response_t **raw_cid**
raw CID of MMC card to be decoded after the CSD is fetched in the data transfer mode

sdmmc_csd_t **csd**
decoded CSD (Card-Specific Data) register value

sdmmc_scr_t **scr**
decoded SCR (SD card Configuration Register) value

sdmmc_ext_csd_t **ext_csd**
decoded EXT_CSD (Extended Card Specific Data) register value

uint16_t **rca**
RCA (Relative Card Address)

uint16_t **max_freq_khz**
Maximum frequency, in kHz, supported by the card

uint32_t **is_mem** : 1
Bit indicates if the card is a memory card

uint32_t **is_sdio** : 1
Bit indicates if the card is an IO card

uint32_t **is_mmc** : 1
Bit indicates if the card is MMC

uint32_t **num_io_functions** : 3
If is_sdio is 1, contains the number of IO functions on the card

uint32_t **log_bus_width** : 2
log2(bus width supported by card)

uint32_t **is_ddr** : 1
Card supports DDR mode

uint32_t **reserved** : 23
Reserved for future expansion

Macros

SDMMC_HOST_FLAG_1BIT
host supports 1-line SD and MMC protocol

SDMMC_HOST_FLAG_4BIT
host supports 4-line SD and MMC protocol

SDMMC_HOST_FLAG_8BIT
host supports 8-line MMC protocol

SDMMC_HOST_FLAG_SPI
host supports SPI protocol

SDMMC_HOST_FLAG_DDR
host supports DDR mode for SD/MMC

SDMMC_HOST_FLAG_DEINIT_ARG
host `deinit` function called with the slot argument

SDMMC_FREQ_DEFAULT
SD/MMC Default speed (limited by clock divider)

SDMMC_FREQ_HIGHSPEED
SD High speed (limited by clock divider)

SDMMC_FREQ_PROBING
SD/MMC probing speed

SDMMC_FREQ_52M
MMC 52MHz speed

SDMMC_FREQ_26M
MMC 26MHz speed

Type Definitions

typedef uint32_t **sdmmc_response_t**[4]
SD/MMC command response buffer

2.5.3 非易失性存储库

简介

非易失性存储 (NVS) 库主要用于在 flash 中存储键值格式的数据。本文档将详细介绍 NVS 常用的一些概念。

底层存储 NVS 通过调用 `spi_flash_{read|write|erase}` API 对主 flash 的部分空间进行读、写、擦除操作，包括 `data` 类型和 `nvs` 子类型的所有分区。应用程序可调用 `nvs_open` API 选择使用带有 `nvs` 标签的分区，也可以通过调用 `nvs_open_from_part` API 选择使用指定名称的任意分区。

NVS 库后续版本可能会增加其他存储器后端，实现将数据保存至其他 flash 芯片（SPI 或 I2C 接口）、RTC 或 FRAM 中。

注解： 如果 NVS 分区被截断（例如，更改分区表布局时），则应擦除分区内容。可以使用 ESP-IDF 构建系统中的 `idf.py erase_flash` 命令擦除 flash 上的所有内容。

注解： NVS 最适合存储一些较小的数据，而非字符串或二进制大对象 (BLOB) 等较大的数据。如需存储较大的 BLOB 或者字符串，请考虑使用基于磨损均衡库的 FAT 文件系统。

键值对 NVS 的操作对象为键值对，其中键是 ASCII 字符串，当前支持最大键长为 15 个字符，值可以为以下几种类型：

- 整数型：`uint8_t`、`int8_t`、`uint16_t`、`int16_t`、`uint32_t`、`int32_t`、`uint64_t` 和 `int64_t`；
- 以 `\0` 结尾的字符串；
- 可变长度的二进制数据 (BLOB)

注解： 字符串值当前上限为 4000 字节，其中包括空终止符。BLOB 值上限为 508,000 字节或分区大小减去 4000 字节的 97.6%，以较低值为准。

后续可能会增加对 `float` 和 `double` 等其他类型数据的支持。

键必须唯一。为现有的键写入新的值可能产生如下结果：

- 如果新旧值数据类型相同，则更新值；
- 如果新旧值数据类型不同，则返回错误。

读取值时也会执行数据类型检查。如果读取操作的数据类型与该值的数据类型不匹配，则返回错误。

命名空间 为了减少不同组件之间键名的潜在冲突，NVS 将每个键值对分配给一个命名空间。命名空间的命名规则遵循键名的命名规则，即最多可占 15 个字符。命名空间的名称在调用 `nvs_open` 或 `nvs_open_from_part` 中指定，调用后将返回一个不透明句柄，用于后续调用 `nvs_get_*`、`nvs_set_*` 和 `nvs_commit` 函数。这样，一个句柄关联一个命名空间，键名便不会与其他命名空间中相同键名冲突。请注意，不同 NVS 分区中具有相同名称的命名空间将被视为不同的命名空间。

安全性、篡改性及鲁棒性 NVS 与 ESP32 flash 加密系统不直接兼容。但如果 NVS 加密与 ESP32 flash 加密一起使用时，数据仍可以加密形式存储。更多详情请参阅*NVS 加密*。

如果未启用 NVS 加密，任何对 flash 芯片有物理访问权限的人都可以修改、擦除或添加键值对。NVS 加密启用后，如果不知道相应的 NVS 加密密钥，则无法修改或添加键值对并将其识别为有效键值。但是，针对擦除操作没有相应的防篡改功能。

当 flash 处于不一致状态时，NVS 库会尝试恢复。在任何时间点关闭设备电源，然后重新打开电源，不会导致数据丢失；但如果关闭设备电源时正在写入新的键值对，这一键值对可能会丢失。该库还应当能对 flash 中的任意数据进行正确初始化。

内部实现

键值对日志 NVS 按顺序存储键值对，新的键值对添加在最后。因此，如需更新某一键值对，实际是在日志最后增加一对新的键值对，同时将旧的键值对标记为已擦除。

页面和条目 NVS 库在其操作中主要使用两个实体：页面和条目。页面是一个逻辑结构，用于存储部分的整体日志。逻辑页面对应 flash 的一个物理扇区，正在使用中的页面具有与之相关联的序列号。序列号赋予了页面顺序，较高的序列号对应较晚创建的页面。页面有以下几种状态：

空或未初始化 页面对应的 flash 扇区为空白状态（所有字节均为 `0xff`）。此时，页面未存储任何数据且没有关联的序列号。

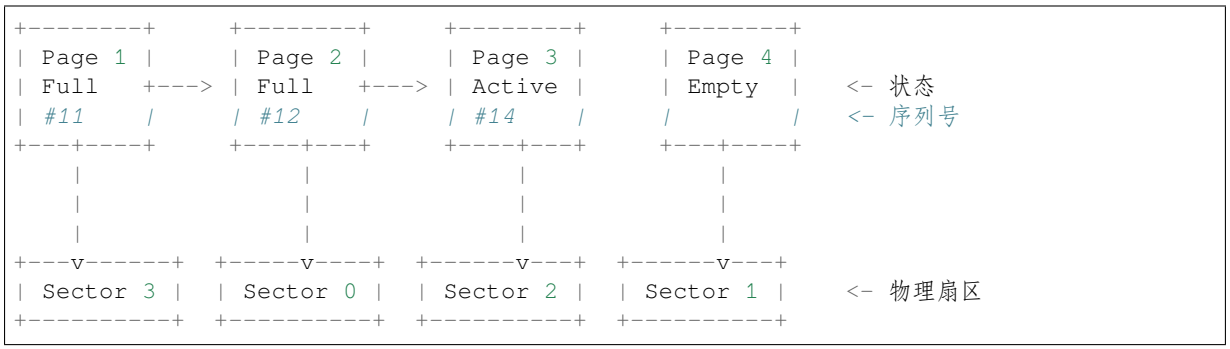
活跃状态 此时 flash 已完成初始化，页头部写入 flash，页面已具备有效序列号。页面中存在一些空条目，可写入数据。任意时刻，至多有一个页面处于活跃状态。

写满状态 Flash 已写满键值对，状态不再改变。用户无法向写满状态下的页面写入新键值对，但仍可将一些键值对标记为已擦除。

擦除状态 未擦除的键值对将移至其他页面，以便擦除当前页面。这一状态仅为暂时性状态，即 API 调用返回时，页面应脱离这一状态。如果设备突然断电，下次开机时，设备将继续把未擦除的键值对移至其他页面，并继续擦除当前页面。

损坏状态 页头部包含无效数据，无法进一步解析该页面中的数据，因此之前写入该页面的所有条目均无法访问。相应的 flash 扇区并不会被立即擦除，而是与其他处于未初始化状态的扇区一起等待后续使用。这一状态可能对调试有用。

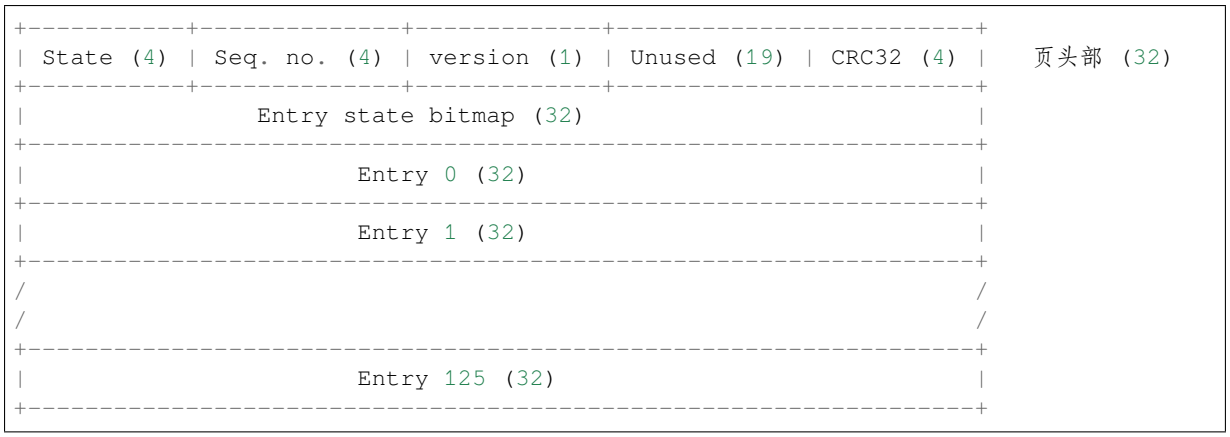
Flash 扇区映射至逻辑页面并没有特定的顺序，NVS 库会检查存储在 flash 扇区的页面序列号，并根据序列号组织页面。



页面结构 当前，我们假设 flash 扇区大小为 4096 字节，并且 ESP32 flash 加密硬件在 32 字节块上运行。未来有可能引入一些编译时可配置项（可通过 menuconfig 进行配置），以适配具有不同扇区大小的 flash 芯片。但目前尚不清楚 SPI flash 驱动和 SPI flash cache 之类的系统组件是否支持其他扇区大小。

页面由头部、条目状态位图和条目三部分组成。为了实现与 ESP32 flash 加密功能兼容，条目大小设置为 32 字节。如果键值为整数型，条目则保存一个键值对；如果键值为字符串或 BLOB 类型，则条目仅保存一个键值对的部分内容（更多信息详见条目结构描述）。

页面结构如下图所示，括号内数字表示该部分的大小（以字节为单位）：



头部和条目状态位图写入 flash 时不加密。如果启用了 ESP32 flash 加密功能，则条目写入 flash 时将会加密。

通过将 0 写入某些位可以定义页面状态值，表示状态改变。因此，如果需要变更页面状态，并不一定要擦除页面，除非要将其变更为擦除状态。

头部中的 `version` 字段反映了所用的 NVS 格式版本。为实现向后兼容，版本升级从 0xff 开始依次递减（例如，version-1 为 0xff，version-2 为 0xfe 等）。

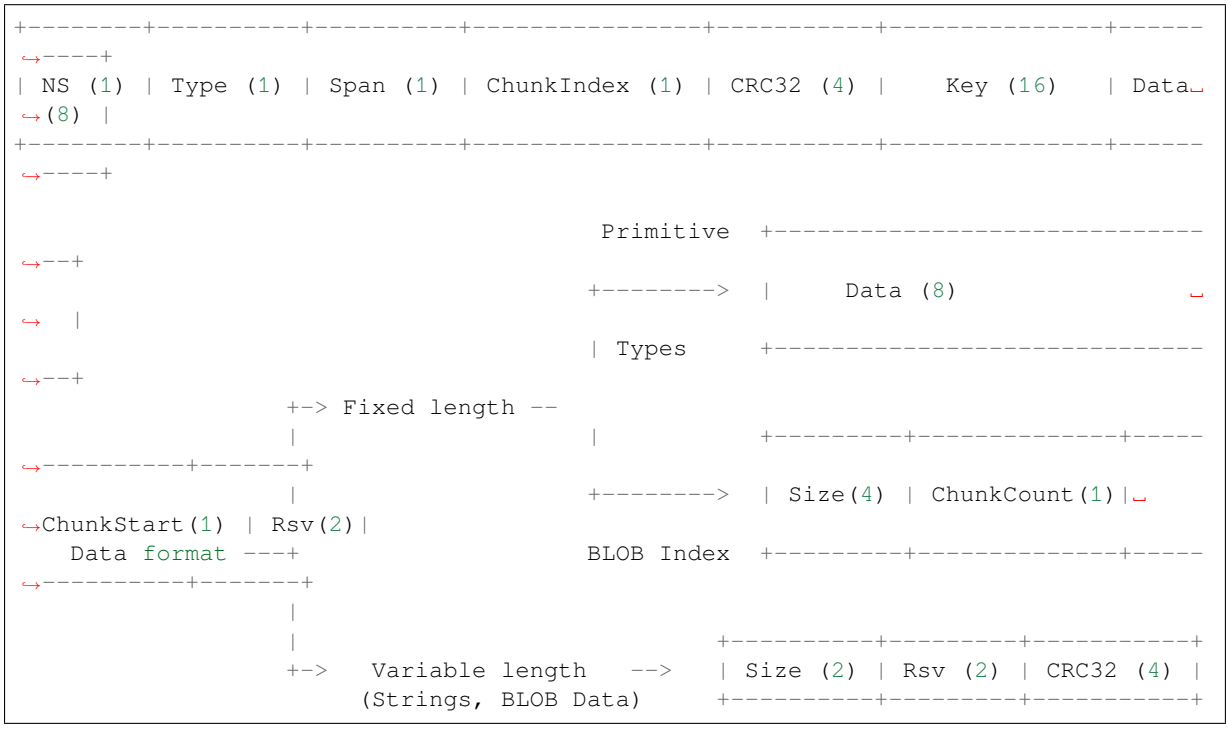
头部中 CRC32 值是由不包含状态值的条目计算所得（4 到 28 字节）。当前未使用的条目用 `0xff` 字节填充。

条目结构和条目状态位图详细信息见下文描述。

条目和条目状态位图 每个条目可处于以下三种状态之一，每个状态在条目状态位图中用两位表示。位图中的最后四位 (256 - 2 * 126) 未使用。

空 (2’ b11) 条目还未写入任何内容，处于未初始化状态（全部字节为 `0xff`）。
写入（2’ b10） 一个键值对（或跨多个条目的键值对的部分内容）已写入条目中。
擦除（2’ b00） 条目中的键值对已丢弃，条目内容不再解析。

条目结构 如果键值类型为基础类型，即 1 - 8 个字节长度的整数型，条目将保存一个键值对；如果键值类型为字符串或 BLOB 类型，条目将保存整个键值对的部分内容。另外，如果键值为字符串类型且跨多个条目，则键值所跨的所有条目均保存在同一页面。BLOB 则可以切分为多个块，实现跨多个页面。BLOB 索引是一个附加的固定长度元数据条目，用于追踪 BLOB 块。目前条目仍支持早期 BLOB 格式（可读取可修改），但这些 BLOB 一经修改，即以新格式储存至条目。



条目结构中各个字段含义如下：

命名空间 (NS, NameSpace) 该条目的命名空间索引，详细信息见命名空间实现章节。
类型 (Type) 一个字节表示的值的数据类型，可能的类型见 `nvs_types.h` 中 `ItemType` 枚举。
跨度 (Span) 该键值对所用的条目数量。如果键值为整数型，条目数量即为 1。如果键值为字符串或 BLOB，则条目数量取决于值的长度。
块索引 (ChunkIndex) 用于存储 BLOB 类型数据块的索引。如果键值为其他数据类型，则此处索引应写入 `0xff`。
CRC32 对条目下所有字节进行校验，所得的校验和（CRC32 字段不计算在内）。
键 (Key) 即以零结尾的 ASCII 字符串，字符串最长为 15 字节，不包含最后一个字节的 NULL (`\0`) 终止符。
数据 (Data) 如果键值类型为整数型，则数据字段仅包含键值。如果键值小于八个字节，使用 `0xff` 填充未使用的部分（右侧）。
如果键值类型为 BLOB 索引条目，则该字段的八个字节将保存以下数据块信息：
- **块大小** 整个 BLOB 数据的大小（以字节为单位）。该字段仅用于 BLOB 索引类型条目。
- **ChunkCount** 存储过程中 BLOB 分成的数据块数量。该字段仅用于 BLOB 索引类型条目。
- **ChunkStart** BLOB 第一个数据块的块索引，后续数据块索引依次递增，步长为 1。该字段仅用于 BLOB 索引类型条目。

如果键值类型为字符串或 BLOB 数据块，数据字段的这八个字节将保存该键值的一些附加信息，如下所示：
- **数据大小** 实际数据的大小（以字节为单位）。如果键值类型为字符串，此字段也应将零终止符包含在内。此字段仅用于字符串和 BLOB 类型条目。
- **CRC32** 数据所有字节的校验和，该字段仅用于字符串和 BLOB 类型条目。

可变长度值（字符串和 BLOB）写入后续条目，每个条目 32 字节。第一个条目的 span 字段将指明使用了多少条目。

命名空间 如上所述，每个键值对属于一个命名空间。命名空间标识符（字符串）也作为键值对的键，存储在索引为 0 的命名空间中。与这些键对应的值就是这些命名空间的索引。

```
+-----------------------------------------+
| NS=0 Type=uint8_t Key="wifi" Value=1    |   Entry describing namespace "wifi"
+-----------------------------------------+
| NS=1 Type=uint32_t Key="channel" Value=6|   Key "channel" in namespace "wifi"
+-----------------------------------------+
| NS=0 Type=uint8_t Key="pwm" Value=2     |   Entry describing namespace "pwm"
+-----------------------------------------+
| NS=2 Type=uint16_t Key="channel" Value=20|  Key "channel" in namespace "pwm"
+-----------------------------------------+
```

条目哈希列表 为了减少对 flash 执行的读操作次数，Page 类对象均设有一个列表，包含一对数据：条目索引和条目哈希值。该列表可大大提高检索速度，而无需迭代所有条目并逐个从 flash 中读取。`Page::findItem` 首先从哈希列表中检索条目哈希值，如果条目存在，则在页面内给出条目索引。由于哈希冲突，在哈希列表中检索条目哈希值可能会得到不同的条目，对 flash 中条目再次迭代可解决这一冲突。

哈希列表中每个节点均包含一个 24 位哈希值和 8 位条目索引。哈希值根据条目命名空间、键名和块索引由 CRC32 计算所得，计算结果保留 24 位。为减少将 32 位条目存储在链表中的开销，链表采用了数组的双向链表。每个数组占用 128 个字节，包含 29 个条目、两个链表指针和一个 32 位计数字段。因此，每页额外需要的 RAM 最少为 128 字节，最多为 640 字节。

NVS 加密

NVS 分区内存储的数据可使用 AES-XTS 进行加密，类似于 IEEE P1619 磁盘加密标准中提到的加密方式。为了实现加密，每个条目被均视为一个扇区，并将条目相对地址（相对于分区开头）传递给加密算法，用作扇区号。NVS 加密所需的密钥存储于其他分区，并进行了*flash 加密*。因此，在使用 NVS 加密前应先启用*flash 加密*。

NVS 密钥分区 应用程序如果想使用 NVS 加密，则需要编译进一个类型为 `data`，子类型为 `key` 的密钥分区。该分区应标记为已加密，且最小为 4096 字节，具体结构见下表。如需了解更多详细信息，请参考分区表。

```
+-----------+--------------+-------------+----+
|               XTS encryption key(32)         |
+----------------------------------------------+
|               XTS tweak key (32)             |
+----------------------------------------------+
|                   CRC32(4)                   |
+----------------------------------------------+
```

使用 NVS 分区生成程序生成上述分区表，并烧录至设备。由于分区已标记为已加密，而且启用了*flash 加密*，引导程序在首次启动时将使用 flash 加密对密钥分区进行加密。您也可以在设备启动后调用 `nvs_flash.h` 提供的 `nvs_flash_generate_keys` API 生成加密密钥，然后再将密钥以加密形式写入密钥分区。

应用程序可以使用不同的密钥对不同的 NVS 分区进行加密，这样就会需要多个加密密钥分区。应用程序应为加解密操作提供正确的密钥或密钥分区。

加密读取/写入 `nvs_get_*` 和 `nvs_set_*` 等 NVS API 函数同样可以对 NVS 加密分区执行读写操作。但用于初始化 NVS 非加密分区和加密分区的 API 则有所不同：初始化 NVS 非加密分区可以使用 `nvs_flash_init` 和 `nvs_flash_init_partition`，但初始化 NVS 加密分区则需调用 `nvs_flash_secure_init` 和 `nvs_flash_secure_init_partition`。上述 API 函数所需的 `nvs_sec_cfg_t` 结构可使用 `nvs_flash_generate_keys` 或者 `nvs_flash_read_security_cfg` 进行填充。

应用程序如需在加密状态下执行 NVS 读写操作，应遵循以下步骤：

1. 使用 `esp_partition_find*` API 查找密钥分区和 NVS 数据分区；

2. 使用 `nvs_flash_read_security_cfg` 或 `nvs_flash_generate_keys` API 填充 `nvs_sec_cfg_t` 结构；
3. 使用 `nvs_flash_secure_init` 或 `nvs_flash_secure_init_partition` API 初始化 NVS flash 分区；
4. 使用 `nvs_open` 或 `nvs_open_from_part` API 打开命名空间；
5. 使用 `nvs_get_*` 或 `nvs_set_*` API 执行 NVS 读取/写入操作；
6. 使用 `nvs_flash_deinit` API 释放已初始化的 NVS 分区。

NVS 迭代器 迭代器允许根据指定的分区名称、命名空间和数据类型轮询 NVS 中存储的键值对。

您可以使用以下函数，执行相关操作：

- `nvs_entry_find`：返回一个不透明句柄，用于后续调用 `nvs_entry_next` 和 `nvs_entry_info` 函数；
- `nvs_entry_next`：返回指向下一个键值对的迭代器；
- `nvs_entry_info`：返回每个键值对的信息。

如果未找到符合标准的键值对，`nvs_entry_find` 和 `nvs_entry_next` 将返回 NULL，此时不必释放迭代器。若不再需要迭代器，可使用 `nvs_release_iterator` 释放迭代器。

NVS 分区生成程序

NVS 分区生成程序帮助生成 NVS 分区二进制文件，可使用烧录程序将二进制文件单独烧录至特定分区。烧录至分区上的键值对由 CSV 文件提供，详情请参考*NVS 分区生成程序*。

应用示例

ESP-IDF storage 目录下提供了两个代码示例：

storage/nvs_rw_value

演示如何读取及写入 NVS 单个整数值。

此示例中的值表示 ESP32 模组重启次数。NVS 中数据不会因为模组重启而丢失，因此只有将这一值存储于 NVS 中，才能起到重启次数计数器的作用。

该示例也演示了如何检测读取/写入操作是否成功，以及某个特定值是否在 NVS 中尚未初始化。诊断程序以纯文本形式提供，帮助您追踪程序流程，及时发现问题。

storage/nvs_rw_blob

演示如何读取及写入 NVS 单个整数值和 Blob（二进制大对象），并在 NVS 中存储这一数值，即便 ESP32 模组重启也不会消失。

- value - 记录 ESP32 模组软重启次数和硬重启次数。
- blob - 内含记录模组运行次数的表格。此表格将被从 NVS 读取至动态分配的 RAM 上。每次手动软重启后，表格内运行次数即增加一次，新加的运行次数被写入 NVS。下拉 GPIO0 即可手动软重启。

该示例也演示了如何执行诊断程序以检测读取/写入操作是否成功。

API 参考

Header File

- nvs_flash/include/nvs_flash.h

Functions

esp_err_t **nvs_flash_init**(void)

Initialize the default NVS partition.

This API initialises the default NVS partition. The default NVS partition is the one that is labeled "nvs" in the partition table.

Return

- ESP_OK if storage was successfully initialized.
- ESP_ERR_NVS_NO_FREE_PAGES if the NVS storage contains no empty pages (which may happen if NVS partition was truncated)
- ESP_ERR_NOT_FOUND if no partition with label "nvs" is found in the partition table
- one of the error codes from the underlying flash storage driver

esp_err_t **nvs_flash_init_partition**(**const** char ***partition_label*)

Initialize NVS flash storage for the specified partition.

Return

- ESP_OK if storage was successfully initialized.
- ESP_ERR_NVS_NO_FREE_PAGES if the NVS storage contains no empty pages (which may happen if NVS partition was truncated)
- ESP_ERR_NOT_FOUND if specified partition is not found in the partition table
- ESP_ERR_NOT_SUPPORTED if the partition with name partition_label is not in internal flash
- one of the error codes from the underlying flash storage driver

Parameters

- `[in] partition_label`: Label of the partition. Must be no longer than 16 characters.

esp_err_t **nvs_flash_init_partition_ptr**(**const** *esp_partition_t* **partition*)

Initialize NVS flash storage for the partition specified by partition pointer.

Return

- ESP_OK if storage was successfully initialized
- ESP_ERR_NVS_NO_FREE_PAGES if the NVS storage contains no empty pages (which may happen if NVS partition was truncated)
- ESP_ERR_INVALID_ARG in case partition is NULL
- ESP_ERR_NOT_SUPPORTED if the partition is not in internal flash
- one of the error codes from the underlying flash storage driver

Parameters

- `[in] partition`: pointer to a partition obtained by the ESP partition API.

esp_err_t **nvs_flash_deinit**(void)

Deinitialize NVS storage for the default NVS partition.

Default NVS partition is the partition with "nvs" label in the partition table.

Return

- ESP_OK on success (storage was deinitialized)
- ESP_ERR_NVS_NOT_INITIALIZED if the storage was not initialized prior to this call

esp_err_t **nvs_flash_deinit_partition**(**const** char ***partition_label*)

Deinitialize NVS storage for the given NVS partition.

Return

- ESP_OK on success
- ESP_ERR_NVS_NOT_INITIALIZED if the storage for given partition was not initialized prior to this call

Parameters

- `[in] partition_label`: Label of the partition

esp_err_t **nvs_flash_erase**(void)

Erase the default NVS partition.

Erases all contents of the default NVS partition (one with label "nvs").

Note If the partition is initialized, this function first de-initializes it. Afterwards, the partition has to be initialized again to be used.

Return

- ESP_OK on success
- ESP_ERR_NOT_FOUND if there is no NVS partition labeled “nvs” in the partition table
- different error in case de-initialization fails (shouldn’t happen)

esp_err_t **nvs_flash_erase_partition**(**const** char **part_name*)

Erase specified NVS partition.

Erase all content of a specified NVS partition

Note If the partition is initialized, this function first de-initializes it. Afterwards, the partition has to be initialized again to be used.

Return

- ESP_OK on success
- ESP_ERR_NOT_FOUND if there is no NVS partition with the specified name in the partition table
- ESP_ERR_NOT_SUPPORTED if the partition with part_name is not in internal flash
- different error in case de-initialization fails (shouldn’t happen)

Parameters

- `[in] part_name`: Name (label) of the partition which should be erased

esp_err_t **nvs_flash_erase_partition_ptr**(**const** *esp_partition_t* **partition*)

Erase custom partition.

Erase all content of specified custom partition.

Note If the partition is initialized, this function first de-initializes it. Afterwards, the partition has to be initialized again to be used.

Return

- ESP_OK on success
- ESP_ERR_NOT_FOUND if there is no partition with the specified parameters in the partition table
- ESP_ERR_INVALID_ARG in case partition is NULL
- ESP_ERR_NOT_SUPPORTED if the partition is not in internal flash
- one of the error codes from the underlying flash storage driver

Parameters

- `[in] partition`: pointer to a partition obtained by the ESP partition API.

esp_err_t **nvs_flash_secure_init**(*nvs_sec_cfg_t* **cfg*)

Initialize the default NVS partition.

This API initialises the default NVS partition. The default NVS partition is the one that is labeled “nvs” in the partition table.

Return

- ESP_OK if storage was successfully initialized.
- ESP_ERR_NVS_NO_FREE_PAGES if the NVS storage contains no empty pages (which may happen if NVS partition was truncated)
- ESP_ERR_NOT_FOUND if no partition with label “nvs” is found in the partition table
- one of the error codes from the underlying flash storage driver

Parameters

- `[in] cfg`: Security configuration (keys) to be used for NVS encryption/decryption. If cfg is NULL, no encryption is used.

esp_err_t **nvs_flash_secure_init_partition**(**const** char **partition_label*, *nvs_sec_cfg_t* **cfg*)

Initialize NVS flash storage for the specified partition.

Return

- ESP_OK if storage was successfully initialized.
- ESP_ERR_NVS_NO_FREE_PAGES if the NVS storage contains no empty pages (which may happen if NVS partition was truncated)
- ESP_ERR_NOT_FOUND if specified partition is not found in the partition table
- ESP_ERR_NOT_SUPPORTED if the partition is not in internal flash
- one of the error codes from the underlying flash storage driver

Parameters

- `[in] partition_label`: Label of the partition. Note that internally a reference to passed value is kept and it should be accessible for future operations
- `[in] cfg`: Security configuration (keys) to be used for NVS encryption/decryption. If cfg is null, no encryption/decryption is used.

esp_err_t **nvs_flash_generate_keys**(**const** *esp_partition_t* *partition, *nvs_sec_cfg_t* *cfg)

Generate and store NVS keys in the provided esp partition.

Return

- ESP_OK, if cfg was read successfully;
- ESP_ERR_NOT_SUPPORTED if the partition is not in internal flash
- or error codes from esp_partition_write/erase APIs.

Parameters

- `[in] partition`: Pointer to partition structure obtained using esp_partition_find_first or esp_partition_get. Must be non-NULL.
- `[out] cfg`: Pointer to nvs security configuration structure. Pointer must be non-NULL. Generated keys will be populated in this structure.

esp_err_t **nvs_flash_read_security_cfg**(**const** *esp_partition_t* *partition, *nvs_sec_cfg_t* *cfg)

Read NVS security configuration from a partition.

Note Provided parition is assumed to be marked ‘encrypted’.

Return

- ESP_OK, if cfg was read successfully;
- ESP_ERR_NVS_KEYS_NOT_INITIALIZED, if the partition is not yet written with keys.
- ESP_ERR_NVS_CORRUPT_KEY_PART, if the partition containing keys is found to be corrupt
- ESP_ERR_NOT_SUPPORTED if the partition is not in internal flash
- or error codes from esp_partition_read API.

Parameters

- `[in] partition`: Pointer to partition structure obtained using esp_partition_find_first or esp_partition_get. Must be non-NULL.
- `[out] cfg`: Pointer to nvs security configuration structure. Pointer must be non-NULL.

Structures

struct nvs_sec_cfg_t

Key for encryption and decryption.

Public Members

uint8_t **eky**[**NVS_KEY_SIZE**]

XTS encryption and decryption key

uint8_t **tky**[**NVS_KEY_SIZE**]

XTS tweak key

Macros

NVS_KEY_SIZE

Header File

- nvs_flash/include/nvs.h

Functions

esp_err_t **nvs_set_i8**(*nvs_handle_t* *handle*, **const** char *key, int8_t *value*)

set int8_t value for given key

Set value for the key, given its name. Note that the actual storage will not be updated until `nvs_commit` is called.

Return

- ESP_OK if value was set successfully
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_READ_ONLY if storage handle was opened as read only
- ESP_ERR_NVS_INVALID_NAME if key name doesn't satisfy constraints
- ESP_ERR_NVS_NOT_ENOUGH_SPACE if there is not enough space in the underlying storage to save the value
- ESP_ERR_NVS_REMOVE_FAILED if the value wasn't updated because flash write operation has failed. The value was written however, and update will be finished after re-initialization of nvs, provided that flash operation doesn't fail again.

Parameters

- `[in] handle`: Handle obtained from nvs_open function. Handles that were opened read only cannot be used.
- `[in] key`: Key name. Maximal length is determined by the underlying implementation, but is guaranteed to be at least 15 characters. Shouldn't be empty.
- `[in] value`: The value to set.

esp_err_t **nvs_set_u8** (*nvs_handle_t handle*, **const** char ***key*, uint8_t *value*)

set uint8_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_i16** (*nvs_handle_t handle*, **const** char ***key*, int16_t *value*)

set int16_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_u16** (*nvs_handle_t handle*, **const** char ***key*, uint16_t *value*)

set uint16_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_i32** (*nvs_handle_t handle*, **const** char ***key*, int32_t *value*)

set int32_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_u32** (*nvs_handle_t handle*, **const** char ***key*, uint32_t *value*)

set uint32_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_i64** (*nvs_handle_t handle*, **const** char ***key*, int64_t *value*)

set int64_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_u64** (*nvs_handle_t handle*, **const** char ***key*, uint64_t *value*)

set uint64_t value for given key

This function is the same as `nvs_set_i8` except for the data type.

esp_err_t **nvs_set_str** (*nvs_handle_t handle*, **const** char ***key*, **const** char ***value*)

set string for given key

Set value for the key, given its name. Note that the actual storage will not be updated until `nvs_commit` is called.

Return

- ESP_OK if value was set successfully
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_READ_ONLY if storage handle was opened as read only
- ESP_ERR_NVS_INVALID_NAME if key name doesn't satisfy constraints
- ESP_ERR_NVS_NOT_ENOUGH_SPACE if there is not enough space in the underlying storage to save the value

- ESP_ERR_NVS_REMOVE_FAILED if the value wasn't updated because flash write operation has failed. The value was written however, and update will be finished after re-initialization of nvs, provided that flash operation doesn't fail again.
- ESP_ERR_NVS_VALUE_TOO_LONG if the string value is too long

Parameters

- `[in] handle`: Handle obtained from nvs_open function. Handles that were opened read only cannot be used.
- `[in] key`: Key name. Maximal length is (NVS_KEY_NAME_MAX_SIZE-1) characters. Shouldn't be empty.
- `[in] value`: The value to set. For strings, the maximum length (including null character) is 4000 bytes, if there is one complete page free for writing. This decreases, however, if the free space is fragmented.

esp_err_t **nvs_get_i8** (*nvs_handle_t handle*, **const** char ***key*, int8_t ***out_value*)

get int8_t value for given key

These functions retrieve value for the key, given its name. If `key` does not exist, or the requested variable type doesn't match the type which was used when setting a value, an error is returned.

In case of any error, out_value is not modified.

`out_value` has to be a pointer to an already allocated variable of the given type.

```
// Example of using nvs_get_i32:
int32_t max_buffer_size = 4096; // default value
esp_err_t err = nvs_get_i32(my_handle, "max_buffer_size", &max_buffer_size);
assert(err == ESP_OK || err == ESP_ERR_NVS_NOT_FOUND);
// if ESP_ERR_NVS_NOT_FOUND was returned, max_buffer_size will still
// have its default value.
```

Return

- ESP_OK if the value was retrieved successfully
- ESP_ERR_NVS_NOT_FOUND if the requested key doesn't exist
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_INVALID_NAME if key name doesn't satisfy constraints
- ESP_ERR_NVS_INVALID_LENGTH if length is not sufficient to store data

Parameters

- `[in] handle`: Handle obtained from nvs_open function.
- `[in] key`: Key name. Maximal length is determined by the underlying implementation, but is guaranteed to be at least 15 characters. Shouldn't be empty.
- `out_value`: Pointer to the output value. May be NULL for nvs_get_str and nvs_get_blob, in this case required length will be returned in length argument.

esp_err_t **nvs_get_u8** (*nvs_handle_t handle*, **const** char ***key*, uint8_t ***out_value*)

get uint8_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_i16** (*nvs_handle_t handle*, **const** char ***key*, int16_t ***out_value*)

get int16_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_u16** (*nvs_handle_t handle*, **const** char ***key*, uint16_t ***out_value*)

get uint16_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_i32** (*nvs_handle_t handle*, **const** char ***key*, int32_t ***out_value*)

get int32_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_u32** (*nvs_handle_t handle*, **const** char ***key*, uint32_t ***out_value*)

get uint32_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_i64** (*nvs_handle_t handle*, **const** char ***key*, int64_t ***out_value*)
get int64_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_u64** (*nvs_handle_t handle*, **const** char ***key*, uint64_t ***out_value*)
get uint64_t value for given key

This function is the same as `nvs_get_i8` except for the data type.

esp_err_t **nvs_get_str** (*nvs_handle_t handle*, **const** char ***key*, char ***out_value*, size_t ***length*)
get string value for given key

These functions retrieve the data of an entry, given its key. If key does not exist, or the requested variable type doesn't match the type which was used when setting a value, an error is returned.

In case of any error, out_value is not modified.

All functions expect out_value to be a pointer to an already allocated variable of the given type.

nvs_get_str and nvs_get_blob functions support WinAPI-style length queries. To get the size necessary to store the value, call nvs_get_str or nvs_get_blob with zero out_value and non-zero pointer to length. Variable pointed to by length argument will be set to the required length. For nvs_get_str, this length includes the zero terminator. When calling nvs_get_str and nvs_get_blob with non-zero out_value, length has to be non-zero and has to point to the length available in out_value. It is suggested that nvs_get/set_str is used for zero-terminated C strings, and nvs_get/set_blob used for arbitrary data structures.

```
// Example (without error checking) of using nvs_get_str to get a string into␣
↪dynamic array:
size_t required_size;
nvs_get_str(my_handle, "server_name", NULL, &required_size);
char* server_name = malloc(required_size);
nvs_get_str(my_handle, "server_name", server_name, &required_size);

// Example (without error checking) of using nvs_get_blob to get a binary data
into a static array:
uint8_t mac_addr[6];
size_t size = sizeof(mac_addr);
nvs_get_blob(my_handle, "dst_mac_addr", mac_addr, &size);
```

Return

- ESP_OK if the value was retrieved successfully
- ESP_ERR_NVS_NOT_FOUND if the requested key doesn't exist
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_INVALID_NAME if key name doesn't satisfy constraints
- ESP_ERR_NVS_INVALID_LENGTH if `length` is not sufficient to store data

Parameters

- `[in] handle`: Handle obtained from nvs_open function.
- `[in] key`: Key name. Maximal length is (NVS_KEY_NAME_MAX_SIZE-1) characters. Shouldn't be empty.
- `[out] out_value`: Pointer to the output value. May be NULL for nvs_get_str and nvs_get_blob, in this case required length will be returned in length argument.
- `[inout] length`: A non-zero pointer to the variable holding the length of out_value. In case out_value a zero, will be set to the length required to hold the value. In case out_value is not zero, will be set to the actual length of the value written. For nvs_get_str this includes zero terminator.

esp_err_t **nvs_get_blob** (*nvs_handle_t handle*, **const** char ***key*, void ***out_value*, size_t ***length*)
get blob value for given key

This function behaves the same as `nvs_get_str`, except for the data type.

esp_err_t **nvs_open** (**const** char ***name*, *nvs_open_mode_t open_mode*, *nvs_handle_t* ***out_handle*)
Open non-volatile storage with a given namespace from the default NVS partition.

Multiple internal ESP-IDF and third party application modules can store their key-value pairs in the NVS module. In order to reduce possible conflicts on key names, each module can use its own namespace. The default NVS partition is the one that is labelled "nvs" in the partition table.

Return

- ESP_OK if storage handle was opened successfully
- ESP_ERR_NVS_NOT_INITIALIZED if the storage driver is not initialized
- ESP_ERR_NVS_PART_NOT_FOUND if the partition with label "nvs" is not found
- ESP_ERR_NVS_NOT_FOUND id namespace doesn't exist yet and mode is NVS_READONLY
- ESP_ERR_NVS_INVALID_NAME if namespace name doesn't satisfy constraints
- other error codes from the underlying storage driver

Parameters

- `[in] name`: Namespace name. Maximal length is determined by the underlying implementation, but is guaranteed to be at least 15 characters. Shouldn't be empty.
- `[in] open_mode`: NVS_READWRITE or NVS_READONLY. If NVS_READONLY, will open a handle for reading only. All write requests will be rejected for this handle.
- `[out] out_handle`: If successful (return code is zero), handle will be returned in this argument.

esp_err_t **nvs_open_from_partition**(**const** char **part_name*, **const** char **name*, *nvs_open_mode_t open_mode*, *nvs_handle_t *out_handle*)

Open non-volatile storage with a given namespace from specified partition.

The behaviour is same as nvs_open() API. However this API can operate on a specified NVS partition instead of default NVS partition. Note that the specified partition must be registered with NVS using nvs_flash_init_partition() API.

Return

- ESP_OK if storage handle was opened successfully
- ESP_ERR_NVS_NOT_INITIALIZED if the storage driver is not initialized
- ESP_ERR_NVS_PART_NOT_FOUND if the partition with specified name is not found
- ESP_ERR_NVS_NOT_FOUND id namespace doesn't exist yet and mode is NVS_READONLY
- ESP_ERR_NVS_INVALID_NAME if namespace name doesn't satisfy constraints
- other error codes from the underlying storage driver

Parameters

- `[in] part_name`: Label (name) of the partition of interest for object read/write/erase
- `[in] name`: Namespace name. Maximal length is determined by the underlying implementation, but is guaranteed to be at least 15 characters. Shouldn't be empty.
- `[in] open_mode`: NVS_READWRITE or NVS_READONLY. If NVS_READONLY, will open a handle for reading only. All write requests will be rejected for this handle.
- `[out] out_handle`: If successful (return code is zero), handle will be returned in this argument.

esp_err_t **nvs_set_blob**(*nvs_handle_t handle*, **const** char **key*, **const** void **value*, size_t *length*)

set variable length binary value for given key

This family of functions set value for the key, given its name. Note that actual storage will not be updated until nvs_commit function is called.

Return

- ESP_OK if value was set successfully
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_READ_ONLY if storage handle was opened as read only
- ESP_ERR_NVS_INVALID_NAME if key name doesn't satisfy constraints
- ESP_ERR_NVS_NOT_ENOUGH_SPACE if there is not enough space in the underlying storage to save the value
- ESP_ERR_NVS_REMOVE_FAILED if the value wasn't updated because flash write operation has failed. The value was written however, and update will be finished after re-initialization of nvs, provided that flash operation doesn't fail again.
- ESP_ERR_NVS_VALUE_TOO_LONG if the value is too long

Parameters

- `[in] handle`: Handle obtained from nvs_open function. Handles that were opened read only cannot be used.
- `[in] key`: Key name. Maximal length is 15 characters. Shouldn't be empty.

- `[in] value`: The value to set.
- `[in] length`: length of binary value to set, in bytes; Maximum length is 508000 bytes or (97.6% of the partition size - 4000) bytes whichever is lower.

esp_err_t **nvs_erase_key**(*nvs_handle_t handle*, **const** char **key*)

Erase key-value pair with given key name.

Note that actual storage may not be updated until nvs_commit function is called.

Return

- ESP_OK if erase operation was successful
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_READ_ONLY if handle was opened as read only
- ESP_ERR_NVS_NOT_FOUND if the requested key doesn't exist
- other error codes from the underlying storage driver

Parameters

- `[in] handle`: Storage handle obtained with nvs_open. Handles that were opened read only cannot be used.
- `[in] key`: Key name. Maximal length is determined by the underlying implementation, but is guaranteed to be at least 15 characters. Shouldn't be empty.

esp_err_t **nvs_erase_all**(*nvs_handle_t handle*)

Erase all key-value pairs in a namespace.

Note that actual storage may not be updated until nvs_commit function is called.

Return

- ESP_OK if erase operation was successful
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- ESP_ERR_NVS_READ_ONLY if handle was opened as read only
- other error codes from the underlying storage driver

Parameters

- `[in] handle`: Storage handle obtained with nvs_open. Handles that were opened read only cannot be used.

esp_err_t **nvs_commit**(*nvs_handle_t handle*)

Write any pending changes to non-volatile storage.

After setting any values, nvs_commit() must be called to ensure changes are written to non-volatile storage. Individual implementations may write to storage at other times, but this is not guaranteed.

Return

- ESP_OK if the changes have been written successfully
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
- other error codes from the underlying storage driver

Parameters

- `[in] handle`: Storage handle obtained with nvs_open. Handles that were opened read only cannot be used.

void **nvs_close**(*nvs_handle_t handle*)

Close the storage handle and free any allocated resources.

This function should be called for each handle opened with nvs_open once the handle is not in use any more. Closing the handle may not automatically write the changes to nonvolatile storage. This has to be done explicitly using nvs_commit function. Once this function is called on a handle, the handle should no longer be used.

Parameters

- `[in] handle`: Storage handle to close

esp_err_t **nvs_get_stats**(**const** char **part_name*, *nvs_stats_t *nvs_stats*)

Fill structure *nvs_stats_t*. It provides info about used memory the partition.

This function calculates to runtime the number of used entries, free entries, total entries, and amount namespace in partition.

```
// Example of nvs_get_stats() to get the number of used entries and free␣
↪entries:
nvs_stats_t nvs_stats;
nvs_get_stats(NULL, &nvs_stats);
printf("Count: UsedEntries = (%d), FreeEntries = (%d), AllEntries = (%d)\n",
       nvs_stats.used_entries, nvs_stats.free_entries, nvs_stats.total_
↪entries);
```

Return

- ESP_OK if the changes have been written successfully. Return param nvs_stats will be filled.
- ESP_ERR_NVS_PART_NOT_FOUND if the partition with label "name" is not found. Return param nvs_stats will be filled 0.
- ESP_ERR_NVS_NOT_INITIALIZED if the storage driver is not initialized. Return param nvs_stats will be filled 0.
- ESP_ERR_INVALID_ARG if nvs_stats equal to NULL.
- ESP_ERR_INVALID_STATE if there is page with the status of INVALID. Return param nvs_stats will be filled not with correct values because not all pages will be counted. Counting will be interrupted at the first INVALID page.

Parameters

- `[in] part_name`: Partition name NVS in the partition table. If pass a NULL than will use NVS_DEFAULT_PART_NAME ("nvs").
- `[out] nvs_stats`: Returns filled structure nvs_states_t. It provides info about used memory the partition.

esp_err_t **nvs_get_used_entry_count** (*nvs_handle_t handle*, size_t **used_entries*)

Calculate all entries in a namespace.

An entry represents the smallest storage unit in NVS. Strings and blobs may occupy more than one entry. Note that to find out the total number of entries occupied by the namespace, add one to the returned value used_entries (if err is equal to ESP_OK). Because the name space entry takes one entry.

```
// Example of nvs_get_used_entry_count() to get amount of all key-value pairs␣
↪in one namespace:
nvs_handle_t handle;
nvs_open("namespace1", NVS_READWRITE, &handle);
...
size_t used_entries;
size_t total_entries_namespace;
if(nvs_get_used_entry_count(handle, &used_entries) == ESP_OK){
    // the total number of entries occupied by the namespace
    total_entries_namespace = used_entries + 1;
}
```

Return

- ESP_OK if the changes have been written successfully. Return param used_entries will be filled valid value.
- ESP_ERR_NVS_NOT_INITIALIZED if the storage driver is not initialized. Return param used_entries will be filled 0.
- ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL. Return param used_entries will be filled 0.
- ESP_ERR_INVALID_ARG if used_entries equal to NULL.
- Other error codes from the underlying storage driver. Return param used_entries will be filled 0.

Parameters

- `[in] handle`: Handle obtained from nvs_open function.
- `[out] used_entries`: Returns amount of used entries from a namespace.

nvs_iterator_t **nvs_entry_find**(**const** char **part_name*, **const** char **namespace_name*, *nvs_type_t type*)

Create an iterator to enumerate NVS entries based on one or more parameters.

```
// Example of listing all the key-value pairs of any type under specified␣
↪partition and namespace
nvs_iterator_t it = nvs_entry_find(partition, namespace, NVS_TYPE_ANY);
while (it != NULL) {
        nvs_entry_info_t info;
        nvs_entry_info(it, &info);
        it = nvs_entry_next(it);
        printf("key '%s', type '%d' \n", info.key, info.type);
};
// Note: no need to release iterator obtained from nvs_entry_find function when
//       nvs_entry_find or nvs_entry_next function return NULL, indicating no␣
↪other
//       element for specified criteria was found.
}
```

Return Iterator used to enumerate all the entries found, or NULL if no entry satisfying criteria was found. Iterator obtained through this function has to be released using nvs_release_iterator when not used any more.

Parameters

- `[in] part_name`: Partition name
- `[in] namespace_name`: Set this value if looking for entries with a specific namespace. Pass NULL otherwise.
- `[in] type`: One of nvs_type_t values.

nvs_iterator_t **nvs_entry_next** (*nvs_iterator_t iterator*)

Returns next item matching the iterator criteria, NULL if no such item exists.

Note that any copies of the iterator will be invalid after this call.

Return NULL if no entry was found, valid nvs_iterator_t otherwise.

Parameters

- `[in] iterator`: Iterator obtained from nvs_entry_find function. Must be non-NULL.

void **nvs_entry_info** (*nvs_iterator_t iterator*, *nvs_entry_info_t *out_info*)

Fills *nvs_entry_info_t* structure with information about entry pointed to by the iterator.

Parameters

- `[in] iterator`: Iterator obtained from nvs_entry_find or nvs_entry_next function. Must be non-NULL.
- `[out] out_info`: Structure to which entry information is copied.

void **nvs_release_iterator** (*nvs_iterator_t iterator*)

Release iterator.

Parameters

- `[in] iterator`: Release iterator obtained from nvs_entry_find function. NULL argument is allowed.

Structures

struct nvs_entry_info_t

information about entry obtained from nvs_entry_info function

Public Members

char **namespace_name**[16]

Namespace to which key-value belong

char **key**[**NVS_KEY_NAME_MAX_SIZE**]

Key of stored key-value pair

nvs_type_t **type**

Type of stored key-value pair

struct nvs_stats_t

Note Info about storage space NVS.

Public Members

size_t **used_entries**
Amount of used entries.

size_t **free_entries**
Amount of free entries.

size_t **total_entries**
Amount all available entries.

size_t **namespace_count**
Amount name space.

Macros

ESP_ERR_NVS_BASE
Starting number of error codes

ESP_ERR_NVS_NOT_INITIALIZED
The storage driver is not initialized

ESP_ERR_NVS_NOT_FOUND
Id namespace doesn't exist yet and mode is NVS_READONLY

ESP_ERR_NVS_TYPE_MISMATCH
The type of set or get operation doesn't match the type of value stored in NVS

ESP_ERR_NVS_READ_ONLY
Storage handle was opened as read only

ESP_ERR_NVS_NOT_ENOUGH_SPACE
There is not enough space in the underlying storage to save the value

ESP_ERR_NVS_INVALID_NAME
Namespace name doesn't satisfy constraints

ESP_ERR_NVS_INVALID_HANDLE
Handle has been closed or is NULL

ESP_ERR_NVS_REMOVE_FAILED
The value wasn't updated because flash write operation has failed. The value was written however, and update will be finished after re-initialization of nvs, provided that flash operation doesn't fail again.

ESP_ERR_NVS_KEY_TOO_LONG
Key name is too long

ESP_ERR_NVS_PAGE_FULL
Internal error; never returned by nvs API functions

ESP_ERR_NVS_INVALID_STATE
NVS is in an inconsistent state due to a previous error. Call nvs_flash_init and nvs_open again, then retry.

ESP_ERR_NVS_INVALID_LENGTH
String or blob length is not sufficient to store data

ESP_ERR_NVS_NO_FREE_PAGES
NVS partition doesn't contain any empty pages. This may happen if NVS partition was truncated. Erase the whole partition and call nvs_flash_init again.

ESP_ERR_NVS_VALUE_TOO_LONG
String or blob length is longer than supported by the implementation

ESP_ERR_NVS_PART_NOT_FOUND
Partition with specified name is not found in the partition table

ESP_ERR_NVS_NEW_VERSION_FOUND
NVS partition contains data in new format and cannot be recognized by this version of code

ESP_ERR_NVS_XTS_ENCR_FAILED
XTS encryption failed while writing NVS entry

ESP_ERR_NVS_XTS_DECR_FAILED
XTS decryption failed while reading NVS entry

ESP_ERR_NVS_XTS_CFG_FAILED
XTS configuration setting failed

ESP_ERR_NVS_XTS_CFG_NOT_FOUND
XTS configuration not found

ESP_ERR_NVS_ENCR_NOT_SUPPORTED
NVS encryption is not supported in this version

ESP_ERR_NVS_KEYS_NOT_INITIALIZED
NVS key partition is uninitialized

ESP_ERR_NVS_CORRUPT_KEY_PART
NVS key partition is corrupt

ESP_ERR_NVS_CONTENT_DIFFERS
Internal error; never returned by nvs API functions. NVS key is different in comparison

NVS_DEFAULT_PART_NAME
Default partition name of the NVS partition in the partition table

NVS_PART_NAME_MAX_SIZE
maximum length of partition name (excluding null terminator)

NVS_KEY_NAME_MAX_SIZE
Maximal length of NVS key name (including null terminator)

Type Definitions

typedef uint32_t **nvs_handle_t**
Opaque pointer type representing non-volatile storage handle

typedef *nvs_handle_t* **nvs_handle**

typedef *nvs_open_mode_t* **nvs_open_mode**

typedef struct nvs_opaque_iterator_t ***nvs_iterator_t**
Opaque pointer type representing iterator to nvs entries

Enumerations

enum nvs_open_mode_t
Mode of opening the non-volatile storage.

Values:

NVS_READONLY
Read only

NVS_READWRITE
Read and write

enum nvs_type_t
Types of variables.

Values:

NVS_TYPE_U8 = 0x01
Type uint8_t

NVS_TYPE_I8 = 0x11
Type int8_t

NVS_TYPE_U16 = 0x02
Type uint16_t

NVS_TYPE_I16 = 0x12
Type int16_t

NVS_TYPE_U32 = 0x04
Type uint32_t

NVS_TYPE_I32 = 0x14
Type int32_t

NVS_TYPE_U64 = 0x08
Type uint64_t

NVS_TYPE_I64 = 0x18
Type int64_t

NVS_TYPE_STR = 0x21
Type string

NVS_TYPE_BLOB = 0x42
Type blob

NVS_TYPE_ANY = 0xff
Must be last

2.5.4 NVS 分区生成程序

介绍

NVS 分区生成程序 (nvs_flash/nvs_partition_generator/nvs_partition_gen.py) 根据 CSV 文件中的键值对生成二进制文件。该二进制文件与非易失性存储器 *(NVS)* 中定义的 NVS 结构兼容。NVS 分区生成程序适合用于生成二进制数据（Blob），其中包括设备生产时可从外部烧录的 ODM/OEM 数据。这也使得生产制造商在使用同一个固件的基础上，通过自定义参数，如序列号等，为每个设备生成不同配置。

准备工作

在加密模式下使用该程序，需安装下列软件包：

- cryptography package

根目录下的 *requirements.txt* 包含必需 python 包，请预先安装。

CSV 文件格式

.csv 文件每行需包含四个参数，以逗号隔开。具体参数描述见下表：

序号	参数	描述	说明
1	Key	主键，应用程序可通过查询此键来获取数据。	
2	Type	支持 `file`、`data` 和 `namespace`。	
3	Encoding	支持 `u8`、`i8`、`u16`、`u32`、`i32`、`string`、`hex2bin`、`base64` 和 `binary`。决定二进制 `bin` 文件中 value 被编码成的类型。`string` 和 `binary` 编码的区别在于，`string` 数据以 NULL 字符结尾，`binary` 数据则不是。	`file` 类型当前仅支持 `hex2bin`、`base64`、`string` 和 `binary` 编码。
4	Value	Data value	`namespace` 字段的 `encoding` 和 `value` 应为空。`namespace` 的 `encoding` 和 `value` 为固定值，不可设置。这些单元格中的所有值都会被忽视。

注解： CSV 文件的第一行应为列标题，不可设置。

此类 CSV 文件的 Dump 示例如下:

```
key,type,encoding,value      <-- 列标题
namespace_name,namespace,,   <-- 第一个条目为 "namespace"
key1,data,u8,1
key2,file,string,/path/to/file
```

注解:

请确保:

- 逗号‘,’前后无空格;
- CSV 文件每行末尾无空格。

NVS 条目和命名空间 (namespace)

如 CSV 文件中出现命名空间条目，后续条目均会被视为该命名空间的一部分，直至找到下一个命名空间条目。找到新命名空间条目后，后续所有条目都会被视为新命名空间的一部分。

注解： CSV 文件中第一个条目应始终为 `namespace`。

支持多页 Blob

默认情况下，二进制 Blob 可跨多页，格式参考条目结构 章节。如需使用旧版格式，可在程序中禁用该功能。

支持加密

NVS 分区生成程序还可使用 AES-XTS 加密生成二进制加密文件。更多信息详见*NVS 加密*。

支持解密

如果 NVS 二进制文件采用了 AES-XTS 加密，该程序还可对此类文件进行解密，更多信息详见*NVS 加密*。

运行程序

使用方法:

```
python nvs_partition_gen.py [-h] {generate,generate-key,encrypt,decrypt} ...
```

可选参数:

序号	参数	描述
1	-h, –help	显示帮助信息并退出

命令:

```
运行 nvs_partition_gen.py {command} -h 查看更多帮助信息
```

序号	参数	描述
1	generate	生成 NVS 分区
2	generate-key	生成加密密钥
3	encrypt	加密 NVS 分区
4	decrypt	解密 NVS 分区

生成 NVS 分区（默认模式）

使用方法:

```
python nvs_partition_gen.py generate [-h] [--version {1,2}] [--outdir OUTDIR]
                                     input output size
```

位置参数:

参数	描述
input	待解析的 CSV 文件路径
output	NVS 二进制文件的输出路径
size	NVS 分区大小（以字节为单位，且为 4096 的整数倍）

可选参数:

参数	描述
-h, –help	显示帮助信息并退出
–version {1,2}	• 设置多页 Blob 版本。 • 版本 1：禁用多页 Blob； • 版本 2：启用多页 Blob； • 默认版本：版本 2。
–outdir OUTDIR	输出目录，用于存储创建的文件。（默认当前目录）

运行如下命令创建 NVS 分区，该程序同时会提供 CSV 示例文件:

```
python nvs_partition_gen.py generate sample_singlepage_blob.csv sample.bin 0x3000
```

仅生成加密密钥

使用方法:

```
python nvs_partition_gen.py generate-key [-h] [--keyfile KEYFILE]
                                         [--outdir OUTDIR]
```

可选参数:

参数	描述
-h, –help	显示帮助信息并退出
–keyfile KEYFILE	加密密钥文件的输出路径
–outdir OUTDIR	输出目录，用于存储创建的文件。（默认当前目录）

运行以下命令仅生成加密密钥:

```
python nvs_partition_gen.py generate-key
```

生成 NVS 加密分区　使用方法:

```
python nvs_partition_gen.py encrypt [-h] [--version {1,2}] [--keygen]
                                    [--keyfile KEYFILE] [--inputkey INPUTKEY]
                                    [--outdir OUTDIR]
                                    input output size
```

位置参数:

参数	描述
input	待解析 CSV 文件的路径
output	NVS 二进制文件的输出路径
size	NVS 分区大小（以字节为单位，且为 4096 的整数倍）

可选参数:

参数	描述
-h, –help	显示帮助信息并退出
–version {1,2}	• 设置多页 Blob 版本。 • 版本 1：禁用多页 Blob； • 版本 2：启用多页 Blob； • 默认版本：版本 2。
–keygen	生成 NVS 分区加密密钥
–keyfile KEYFILE	密钥文件的输出路径
–inputkey INPUTKEY	内含 NVS 分区加密密钥的文件
–outdir OUTDIR	输出目录，用于存储创建的文件（默认当前目录）

运行以下命令加密 NVS 分区，该程序同时会提供一个 CSV 示例文件。

- 通过 NVS 分区生成程序生成加密密钥来加密:

```
python nvs_partition_gen.py encrypt sample_singlepage_blob.csv sample_encr.bin␣
↪0x3000 --keygen
```

注解: 创建的加密密钥格式为 `<outdir>/keys/keys-<timestamp>.bin`。

- 通过 NVS 分区生成程序生成加密密钥，并将密钥存储于自定义的文件中:

```
python nvs_partition_gen.py encrypt sample_singlepage_blob.csv sample_encr.bin␣
↪0x3000 --keygen --keyfile sample_keys.bin
```

注解： 创建的加密密钥格式为 <outdir>/keys/keys-<timestamp>.bin。

注解： 加密密钥存储于新建文件的 keys/ 目录下，与 NVS 密钥分区结构兼容。更多信息请参考*NVS 密钥分区*。

- 将加密密钥用作二进制输入文件来进行加密:

```
python nvs_partition_gen.py encrypt sample_singlepage_blob.csv sample_encr.bin␣
↪0x3000 --inputkey sample_keys.bin
```

解密 NVS 分区　使用方法:

```
python nvs_partition_gen.py decrypt [-h] [--outdir OUTDIR] input key output
```

位置参数:

参数	描述
input	待解析的 NVS 加密分区文件路径
key	含有解密密钥的文件路径
output	已解密的二进制文件输出路径

可选参数:

参数	描述
-h, –help	显示帮助信息并退出
–outdir OUTDIR	输出目录，用于存储创建的文件（默认当前目录）

运行以下命令解密已加密的 NVS 分区:

```
python nvs_partition_gen.py decrypt sample_encr.bin sample_keys.bin sample_decr.bin
```

您可以自定义格式版本号:

- 版本 1：禁用多页 Blob
- 版本 2：启用多页 Blob

版本 1：禁用多页 Blob　如需禁用多页 Blob，请按照如下命令将版本参数设置为 1，以此格式运行分区生成程序。该程序同时会提供一个 CSV 示例文件:

```
python nvs_partition_gen.py generate sample_singlepage_blob.csv sample.bin 0x3000 -
↪-version 1
```

版本 2：启用多页 Blob　如需启用多页 Blob，请按照如下命令将版本参数设置为 2，以此格式运行分区生成程序。该程序同时会提供一个 CSV 示例文件:

```
python nvs_partition_gen.py generate sample_multipage_blob.csv sample.bin 0x4000 --
↪version 2
```

注解： NVS 分区最小为 0x3000 字节。

注解： 将二进制文件烧录至设备时，请确保与应用的 sdkconfig 设置一致。

说明

- 分区生成程序不会对重复键进行检查，而将数据同时写入这两个重复键中。请注意不要使用同名的键；
- 新页面创建后，前一页的空白处不会再写入数据。CSV 文件中的字段须按次序排列以优化内存；
- 暂不支持 64 位数据类型。

2.5.5 虚拟文件系统组件

概述

虚拟文件系统 (VFS) 组件可为一些驱动提供一个统一接口。有了该接口，用户可像操作普通文件一样操作虚拟文件。这类驱动程序可以是 FAT、SPIFFS 等真实文件系统，也可以是有文件类接口的设备驱动程序。

VFS 组件支持 C 库函数（如 fopen 和 fprintf 等）与文件系统 (FS) 驱动程序协同工作。在高层级，每个 FS 驱动程序均与某些路径前缀相关联。当一个 C 库函数需要打开文件时，VFS 组件将搜索与该文件所在文件路径相关联的 FS 驱动程序，并将调用传递给该驱动程序。针对该文件的读取、写入等其他操作的调用也将传递给这个驱动程序。

例如，您可以使用 `/fat` 前缀注册 FAT 文件系统驱动，之后即可调用 `fopen("/fat/file.txt", "w")`。之后，VFS 将调用 FAT 驱动的 `open` 函数，并将参数 `/file.txt` 和合适的打开模式传递给 `open` 函数；后续对返回的 `FILE*` 数据流调用 C 库函数也同样会传递给 FAT 驱动。

注册 FS 驱动程序

如需注册 FS 驱动程序，首先要定义一个*esp_vfs_t* 结构体实例，并用指向 FS API 的函数指针填充它。

```
esp_vfs_t myfs = {
    .flags = ESP_VFS_FLAG_DEFAULT,
    .write = &myfs_write,
    .open = &myfs_open,
    .fstat = &myfs_fstat,
    .close = &myfs_close,
    .read = &myfs_read,
};

ESP_ERROR_CHECK(esp_vfs_register("/data", &myfs, NULL));
```

在上述代码中需要用到 `read`、`write` 或 `read_p`、`write_p`，具体使用哪组函数由 FS 驱动程序 API 的声明方式决定。

示例 1：声明 API 函数时不带额外的上下文指针参数，即 FS 驱动程序为单例模式，此时使用 `write`

```
ssize_t myfs_write(int fd, const void * data, size_t size);

// In definition of esp_vfs_t:
    .flags = ESP_VFS_FLAG_DEFAULT,
    .write = &myfs_write,
// ... other members initialized
```

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```
// When registering FS, context pointer (third argument) is NULL:
ESP_ERROR_CHECK(esp_vfs_register("/data", &myfs, NULL));
```

示例 2：声明 API 函数时需要一个额外的上下文指针作为参数，即可支持多个 FS 驱动程序实例，此时使用 `write_p`

```
ssize_t myfs_write(myfs_t* fs, int fd, const void * data, size_t size);

// In definition of esp_vfs_t:
    .flags = ESP_VFS_FLAG_CONTEXT_PTR,
    .write_p = &myfs_write,
// ... other members initialized

// When registering FS, pass the FS context pointer into the third argument
// (hypothetical myfs_mount function is used for illustrative purposes)
myfs_t* myfs_inst1 = myfs_mount(partition1->offset, partition1->size);
ESP_ERROR_CHECK(esp_vfs_register("/data1", &myfs, myfs_inst1));

// Can register another instance:
myfs_t* myfs_inst2 = myfs_mount(partition2->offset, partition2->size);
ESP_ERROR_CHECK(esp_vfs_register("/data2", &myfs, myfs_inst2));
```

同步输入/输出多路复用 如需通过 `select()` 使用同步输入/输出多路复用，首先需要把 `start_select()` 和 `end_select()` 注册到 VFS，如下所示：

```
// In definition of esp_vfs_t:
    .start_select = &uart_start_select,
    .end_select = &uart_end_select,
// ... other members initialized
```

调用 `start_select()` 设置环境，用以检测某一 VFS 文件描述符的读取/写入/错误条件。调用 `end_select()` 终止、析构或释放 `start_select()` 设置的资源。请在 vfs/vfs_uart.c 中查看 UART 外设参考实现、*`esp_vfs_dev_uart_register()`*、`uart_start_select()` 和 `uart_end_select()` 函数。

请参考以下示例，查看如何使用 VFS 文件描述符调用 `select()`：

- peripherals/uart/uart_select
- system/select

如果 `select()` 用于套接字文件描述符，您可以启用 `CONFIG_LWIP_USE_ONLY_LWIP_SELECT` 选项来减少代码量，提高性能。

路径

已注册的 FS 驱动程序均有一个路径前缀与之关联，此路径前缀即为分区的挂载点。

如果挂载点中嵌套了其他挂载点，则在打开文件时使用具有最长匹配路径前缀的挂载点。例如，假设以下文件系统已在 VFS 中注册：

- 在 /data 下注册 FS 驱动程序 1
- 在 /data/static 下注册 FS 驱动程序 2

那么：

- 打开 `/data/log.txt` 会调用驱动程序 FS 1；
- 打开 `/data/static/index.html` 需调用 FS 驱动程序 2；
- 即便 FS 驱动程序 2 中没有 `/index.html`，也不会在 FS 驱动程序 1 中查找 `/static/index.html`。

挂载点名称必须以路径分隔符 (/) 开头，且分隔符后至少包含一个字符。但在以下情况中，VFS 同样支持空的挂载点名称：1. 应用程序需要提供一个”最后方案“下使用的文件系统；2. 应用程序需要同时覆盖 VFS 功能。如果没有与路径匹配的前缀，就会使用到这种文件系统。

VFS 不会对路径中的点 (.) 进行特殊处理，也不会将 .. 视为对父目录的引用。在上述示例中，使用 /data/static/../log.txt 路径不会调用 FS 驱动程序 1 打开 /log.txt。特定的 FS 驱动程序（如 FATFS）可能以不同的方式处理文件名中的点。

执行打开文件操作时，FS 驱动程序仅得到文件的相对路径（挂载点前缀已经被去除）：

1. 以 /data 为路径前缀注册 myfs 驱动；
2. 应用程序调用 fopen("/data/config.json", ...)；
3. VFS 调用 myfs_open("/config.json", ...)；
4. myfs 驱动打开 /config.json 文件。

VFS 对文件路径长度没有限制，但文件系统路径前缀受 ESP_VFS_PATH_MAX 限制，即路径前缀上限为 ESP_VFS_PATH_MAX。各个文件系统驱动则可能会对自己的文件名长度设置一些限制。

文件描述符

文件描述符是一组很小的正整数，从 0 到 FD_SETSIZE - 1，FD_SETSIZE 在 newlib sys/types.h 中定义。最大文件描述符由 CONFIG_LWIP_MAX_SOCKETS 定义，且为套接字保留。VFS 中包含一个名为 s_fd_table 的查找表，用于将全局文件描述符映射至 s_vfs 数组中注册的 VFS 驱动索引。

标准 IO 流 (stdin, stdout, stderr)

如果 menuconfig 中 UART for console output 选项没有设置为 None，则 stdin、stdout 和 stderr 将默认从 UART 读取或写入。UART0 或 UART1 可用作标准 IO。默认情况下，UART0 使用 115200 波特率，TX 管脚为 GPIO1，RX 管脚为 GPIO3。您可以在 menuconfig 中更改上述参数。

对 stdout 或 stderr 执行写入操作将会向 UART 发送 FIFO 发送字符，对 stdin 执行读取操作则会从 UART 接收 FIFO 中取出字符。

默认情况下，VFS 使用简单的函数对 UART 进行读写操作。在所有数据放进 UART FIFO 之前，写操作将处于 busy-wait 状态，读操处于非阻塞状态，仅返回 FIFO 中已有数据。由于读操作为非阻塞，高层级 C 库函数调用（如 fscanf("%d\n", &var);）可能获取不到所需结果。

如果应用程序使用 UART 驱动，则可以调用 esp_vfs_dev_uart_use_driver 函数来指导 VFS 使用驱动中断、读写阻塞功能等。您也可以调用 esp_vfs_dev_uart_use_nonblocking 来恢复非阻塞函数。

VFS 还为输入和输出提供换行符转换功能（可选）。多数应用程序在程序内部发送或接收以 LF（‘ n’’）结尾的行，但不同的终端程序可能需要不同的换行符，比如 CR 或 CRLF。应用程序可以通过 menuconfig 或者调用 esp_vfs_dev_uart_port_set_rx_line_endings 和 esp_vfs_dev_uart_port_set_tx_line_endings 为输入输出配置换行符。

标准流和 FreeRTOS 任务 stdin、stdout 和 stderr 的 FILE 对象在所有 FreeRTOS 任务之间共享，指向这些对象的指针分别存储在每个任务的 struct _reent 中。

预处理器把如下代码：

```
fprintf(stderr, "42\n");
```

解释为：

```
fprintf(__getreent()->_stderr, "42\n");
```

其中 __getreent() 函数将为每个任务返回一个指向 struct _reent 的指针。每个任务的 TCB 均拥有一个 struct _reent 结构体，任务初始化后，struct _reent 结构体中的 _stdin、_stdout 和 _stderr 将会被赋予 _GLOBAL_REENT 中 _stdin、_stdout 和 _stderr 的值，_GLOBAL_REENT 即为 FreeRTOS 启动之前所用结构体。

这样设计带来的结果是：

- 允许重定向给定任务的 stdin、stdout 和 stderr，而不影响其他任务，例如通过 stdin = fopen("/dev/uart/1", "r")；
- 但使用 fclose 关闭默认 stdin、stdout 或 stderr 将同时关闭相应的 FILE 流对象，因此会影响其他任务；
- 如需更改新任务的默认 stdin、stdout 和 stderr 流，请在创建新任务之前修改 _GLOBAL_REENT->_stdin (_stdout、_stderr)。

应用示例

指南 （未完成）

API 参考

Header File

- vfs/include/esp_vfs.h

Functions

ssize_t **esp_vfs_write**(**struct** _reent *r, int *fd*, **const** void **data*, size_t *size*)

These functions are to be used in newlib syscall table. They will be called by newlib when it needs to use any of the syscalls.

off_t **esp_vfs_lseek**(**struct** _reent *r, int *fd*, off_t *size*, int *mode*)

ssize_t **esp_vfs_read**(**struct** _reent *r, int *fd*, void **dst*, size_t *size*)

int **esp_vfs_open**(**struct** _reent *r, **const** char **path*, int *flags*, int *mode*)

int **esp_vfs_close**(**struct** _reent *r, int *fd*)

int **esp_vfs_fstat**(**struct** _reent *r, int *fd*, **struct** stat **st*)

int **esp_vfs_stat**(**struct** _reent *r, **const** char **path*, **struct** stat **st*)

int **esp_vfs_link**(**struct** _reent *r, **const** char **n1*, **const** char **n2*)

int **esp_vfs_unlink**(**struct** _reent *r, **const** char **path*)

int **esp_vfs_rename**(**struct** _reent *r, **const** char **src*, **const** char **dst*)

int **esp_vfs_utime**(**const** char **path*, **const struct** utimbuf **times*)

esp_err_t **esp_vfs_register**(**const** char **base_path*, **const** *esp_vfs_t* **vfs*, void **ctx*)

Register a virtual filesystem for given path prefix.

Return ESP_OK if successful, ESP_ERR_NO_MEM if too many VFSes are registered.

Parameters

- base_path: file path prefix associated with the filesystem. Must be a zero-terminated C string, up to ESP_VFS_PATH_MAX characters long, and at least 2 characters long. Name must start with a “/” and must not end with “/”. For example, “/data” or “/dev/spi” are valid. These VFSes would then be called to handle file paths such as “/data/myfile.txt” or “/dev/spi/0”.
- vfs: Pointer to *esp_vfs_t*, a structure which maps syscalls to the filesystem driver functions. VFS component doesn’t assume ownership of this pointer.
- ctx: If vfs->flags has ESP_VFS_FLAG_CONTEXT_PTR set, a pointer which should be passed to VFS functions. Otherwise, NULL.

esp_err_t **esp_vfs_register_fd_range**(**const** *esp_vfs_t* **vfs*, void **ctx*, int *min_fd*, int *max_fd*)

Special case function for registering a VFS that uses a method other than open() to open new file descriptors from the interval <min_fd; max_fd).

This is a special-purpose function intended for registering LWIP sockets to VFS.

Return ESP_OK if successful, ESP_ERR_NO_MEM if too many VFSes are registered, ESP_ERR_INVALID_ARG if the file descriptor boundaries are incorrect.

Parameters

- `vfs`: Pointer to *esp_vfs_t*. Meaning is the same as for esp_vfs_register().
- `ctx`: Pointer to context structure. Meaning is the same as for esp_vfs_register().
- `min_fd`: The smallest file descriptor this VFS will use.
- `max_fd`: Upper boundary for file descriptors this VFS will use (the biggest file descriptor plus one).

esp_err_t **esp_vfs_register_with_id**(**const** *esp_vfs_t* **vfs*, void **ctx*, *esp_vfs_id_t* **vfs_id*)

Special case function for registering a VFS that uses a method other than open() to open new file descriptors. In comparison with esp_vfs_register_fd_range, this function doesn't pre-registers an interval of file descriptors. File descriptors can be registered later, by using esp_vfs_register_fd.

Return ESP_OK if successful, ESP_ERR_NO_MEM if too many VFSes are registered, ESP_ERR_INVALID_ARG if the file descriptor boundaries are incorrect.

Parameters

- `vfs`: Pointer to *esp_vfs_t*. Meaning is the same as for esp_vfs_register().
- `ctx`: Pointer to context structure. Meaning is the same as for esp_vfs_register().
- `vfs_id`: Here will be written the VFS ID which can be passed to esp_vfs_register_fd for registering file descriptors.

esp_err_t **esp_vfs_unregister**(**const** char **base_path*)

Unregister a virtual filesystem for given path prefix

Return ESP_OK if successful, ESP_ERR_INVALID_STATE if VFS for given prefix hasn't been registered

Parameters

- `base_path`: file prefix previously used in esp_vfs_register call

esp_err_t **esp_vfs_register_fd**(*esp_vfs_id_t vfs_id*, int **fd*)

Special function for registering another file descriptor for a VFS registered by esp_vfs_register_with_id.

Return ESP_OK if the registration is successful, ESP_ERR_NO_MEM if too many file descriptors are registered, ESP_ERR_INVALID_ARG if the arguments are incorrect.

Parameters

- `vfs_id`: VFS identificator returned by esp_vfs_register_with_id.
- `fd`: The registered file descriptor will be written to this address.

esp_err_t **esp_vfs_unregister_fd**(*esp_vfs_id_t vfs_id*, int *fd*)

Special function for unregistering a file descriptor belonging to a VFS registered by esp_vfs_register_with_id.

Return ESP_OK if the registration is successful, ESP_ERR_INVALID_ARG if the arguments are incorrect.

Parameters

- `vfs_id`: VFS identificator returned by esp_vfs_register_with_id.
- `fd`: File descriptor which should be unregistered.

int **esp_vfs_select**(int *nfds*, fd_set **readfds*, fd_set **writefds*, fd_set **errorfds*, **struct** timeval **timeout*)

Synchronous I/O multiplexing which implements the functionality of POSIX select() for VFS.

Return The number of descriptors set in the descriptor sets, or -1 when an error (specified by errno) have occurred.

Parameters

- `nfds`: Specifies the range of descriptors which should be checked. The first nfds descriptors will be checked in each set.
- `readfds`: If not NULL, then points to a descriptor set that on input specifies which descriptors should be checked for being ready to read, and on output indicates which descriptors are ready to read.
- `writefds`: If not NULL, then points to a descriptor set that on input specifies which descriptors should be checked for being ready to write, and on output indicates which descriptors are ready to write.
- `errorfds`: If not NULL, then points to a descriptor set that on input specifies which descriptors should be checked for error conditions, and on output indicates which descriptors have error conditions.

- `timeout`: If not NULL, then points to timeval structure which specifies the time period after which the functions should time-out and return. If it is NULL, then the function will not time-out. Note that the timeout period is rounded up to the system tick and incremented by one.

void **esp_vfs_select_triggered**(*esp_vfs_select_sem_t sem*)
Notification from a VFS driver about a read/write/error condition.

This function is called when the VFS driver detects a read/write/error condition as it was requested by the previous call to start_select.

Parameters
- `sem`: semaphore structure which was passed to the driver by the start_select call

void **esp_vfs_select_triggered_isr**(*esp_vfs_select_sem_t sem*, BaseType_t **woken*)
Notification from a VFS driver about a read/write/error condition (ISR version)

This function is called when the VFS driver detects a read/write/error condition as it was requested by the previous call to start_select.

Parameters
- `sem`: semaphore structure which was passed to the driver by the start_select call
- `woken`: is set to pdTRUE if the function wakes up a task with higher priority

ssize_t **esp_vfs_pread**(int *fd*, void **dst*, size_t *size*, off_t *offset*)
Implements the VFS layer of POSIX pread()

Return A positive return value indicates the number of bytes read. -1 is return on failure and errno is set accordingly.

Parameters
- `fd`: File descriptor used for read
- `dst`: Pointer to the buffer where the output will be written
- `size`: Number of bytes to be read
- `offset`: Starting offset of the read

ssize_t **esp_vfs_pwrite**(int *fd*, **const** void **src*, size_t *size*, off_t *offset*)
Implements the VFS layer of POSIX pwrite()

Return A positive return value indicates the number of bytes written. -1 is return on failure and errno is set accordingly.

Parameters
- `fd`: File descriptor used for write
- `src`: Pointer to the buffer from where the output will be read
- `size`: Number of bytes to write
- `offset`: Starting offset of the write

Structures

struct esp_vfs_select_sem_t
VFS semaphore type for select()

Public Members

bool **is_sem_local**
type of "sem" is SemaphoreHandle_t when true, defined by socket driver otherwise

void ***sem**
semaphore instance

struct esp_vfs_t
VFS definition structure.

This structure should be filled with pointers to corresponding FS driver functions.

VFS component will translate all FDs so that the filesystem implementation sees them starting at zero. The caller sees a global FD which is prefixed with an pre-filesystem-implementation.

Some FS implementations expect some state (e.g. pointer to some structure) to be passed in as a first argument. For these implementations, populate the members of this structure which have _p suffix, set flags member to ESP_VFS_FLAG_CONTEXT_PTR and provide the context pointer to esp_vfs_register function. If the implementation doesn't use this extra argument, populate the members without _p suffix and set flags member to ESP_VFS_FLAG_DEFAULT.

If the FS driver doesn't provide some of the functions, set corresponding members to NULL.

Public Members

int **flags**
ESP_VFS_FLAG_CONTEXT_PTR or ESP_VFS_FLAG_DEFAULT

ssize_t (***write_p**) (void *p, int fd, **const** void *data, size_t size)
Write with context pointer

ssize_t (***write**) (int fd, **const** void *data, size_t size)
Write without context pointer

off_t (***lseek_p**) (void *p, int fd, off_t size, int mode)
Seek with context pointer

off_t (***lseek**) (int fd, off_t size, int mode)
Seek without context pointer

ssize_t (***read_p**) (void *ctx, int fd, void *dst, size_t size)
Read with context pointer

ssize_t (***read**) (int fd, void *dst, size_t size)
Read without context pointer

ssize_t (***pread_p**) (void *ctx, int fd, void *dst, size_t size, off_t offset)
pread with context pointer

ssize_t (***pread**) (int fd, void *dst, size_t size, off_t offset)
pread without context pointer

ssize_t (***pwrite_p**) (void *ctx, int fd, **const** void *src, size_t size, off_t offset)
pwrite with context pointer

ssize_t (***pwrite**) (int fd, **const** void *src, size_t size, off_t offset)
pwrite without context pointer

int (***open_p**) (void *ctx, **const** char *path, int flags, int mode)
open with context pointer

int (***open**) (**const** char *path, int flags, int mode)
open without context pointer

int (***close_p**) (void *ctx, int fd)
close with context pointer

int (***close**) (int fd)
close without context pointer

int (***fstat_p**) (void *ctx, int fd, **struct** *stat* *st)
fstat with context pointer

int (***fstat**) (int fd, **struct** *stat* *st)
fstat without context pointer

int (***stat_p**) (void *ctx, **const** char *path, **struct** *stat* *st)
stat with context pointer

int (***stat**) (**const** char *path, **struct** *stat* *st)
stat without context pointer

int (***link_p**) (void *ctx, **const** char *n1, **const** char *n2)
link with context pointer

int (***link**) (**const** char *n1, **const** char *n2)
link without context pointer

int (***unlink_p**) (void *ctx, **const** char *path)
unlink with context pointer

int (***unlink**) (**const** char *path)
unlink without context pointer

int (***rename_p**) (void *ctx, **const** char *src, **const** char *dst)
rename with context pointer

int (***rename**) (**const** char *src, **const** char *dst)
rename without context pointer

DIR *(***opendir_p**) (void *ctx, **const** char *name)
opendir with context pointer

DIR *(***opendir**) (**const** char *name)
opendir without context pointer

struct dirent *(***readdir_p**) (void *ctx, DIR *pdir)
readdir with context pointer

struct dirent *(***readdir**) (DIR *pdir)
readdir without context pointer

int (***readdir_r_p**) (void *ctx, DIR *pdir, **struct** dirent *entry, **struct** dirent **out_dirent)
readdir_r with context pointer

int (***readdir_r**) (DIR *pdir, **struct** dirent *entry, **struct** dirent **out_dirent)
readdir_r without context pointer

long (***telldir_p**) (void *ctx, DIR *pdir)
telldir with context pointer

long (***telldir**) (DIR *pdir)
telldir without context pointer

void (***seekdir_p**) (void *ctx, DIR *pdir, long offset)
seekdir with context pointer

void (***seekdir**) (DIR *pdir, long offset)
seekdir without context pointer

int (***closedir_p**) (void *ctx, DIR *pdir)
closedir with context pointer

int (***closedir**) (DIR *pdir)
closedir without context pointer

int (***mkdir_p**) (void *ctx, **const** char *name, mode_t mode)
mkdir with context pointer

int (***mkdir**) (**const** char *name, mode_t mode)
mkdir without context pointer

int (***rmdir_p**) (void *ctx, **const** char *name)
rmdir with context pointer

int (***rmdir**) (**const** char *name)
rmdir without context pointer

int (***fcntl_p**) (void *ctx, int fd, int cmd, int arg)
fcntl with context pointer

int (***fcntl**) (int fd, int cmd, int arg)
fcntl without context pointer

int (***ioctl_p**) (void *ctx, int fd, int cmd, va_list args)
ioctl with context pointer

int (***ioctl**) (int fd, int cmd, va_list args)
ioctl without context pointer

int (***fsync_p**) (void *ctx, int fd)
fsync with context pointer

int (***fsync**) (int fd)
fsync without context pointer

int (***access_p**) (void *ctx, **const** char *path, int amode)
access with context pointer

int (***access**) (**const** char *path, int amode)
access without context pointer

int (***truncate_p**) (void *ctx, **const** char *path, off_t length)
truncate with context pointer

int (***truncate**) (**const** char *path, off_t length)
truncate without context pointer

int (***utime_p**) (void *ctx, **const** char *path, **const struct** utimbuf *times)
utime with context pointer

int (***utime**) (**const** char *path, **const struct** utimbuf *times)
utime without context pointer

int (***tcsetattr_p**) (void *ctx, int fd, int optional_actions, **const struct** termios *p)
tcsetattr with context pointer

int (***tcsetattr**) (int fd, int optional_actions, **const struct** termios *p)
tcsetattr without context pointer

int (***tcgetattr_p**) (void *ctx, int fd, **struct** termios *p)
tcgetattr with context pointer

int (***tcgetattr**) (int fd, **struct** termios *p)
tcgetattr without context pointer

int (***tcdrain_p**) (void *ctx, int fd)
tcdrain with context pointer

int (***tcdrain**) (int fd)
tcdrain without context pointer

int (***tcflush_p**) (void *ctx, int fd, int select)
tcflush with context pointer

int (***tcflush**) (int fd, int select)
tcflush without context pointer

int (***tcflow_p**) (void *ctx, int fd, int action)
tcflow with context pointer

int (***tcflow**) (int fd, int action)
tcflow without context pointer

pid_t (***tcgetsid_p**) (void *ctx, int fd)
tcgetsid with context pointer

pid_t (***tcgetsid**) (int fd)
tcgetsid without context pointer

int (***tcsendbreak_p**) (void *ctx, int fd, int duration)
tcsendbreak with context pointer

int (***tcsendbreak**) (int fd, int duration)
tcsendbreak without context pointer

esp_err_t (***start_select**) (int nfds, fd_set *readfds, fd_set *writefds, fd_set *exceptfds, *esp_vfs_select_sem_t* sem, void **end_select_args)
start_select is called for setting up synchronous I/O multiplexing of the desired file descriptors in the given VFS

int (***socket_select**) (int nfds, fd_set *readfds, fd_set *writefds, fd_set *errorfds, **struct** timeval *timeout)
socket select function for socket FDs with the functionality of POSIX select(); this should be set only for the socket VFS

void (***stop_socket_select**) (void *sem)
called by VFS to interrupt the socket_select call when select is activated from a non-socket VFS driver; set only for the socket driver

void (***stop_socket_select_isr**) (void *sem, BaseType_t *woken)
stop_socket_select which can be called from ISR; set only for the socket driver

void *(***get_socket_select_semaphore**) (void)
end_select is called to stop the I/O multiplexing and deinitialize the environment created by start_select for the given VFS

esp_err_t (***end_select**) (void *end_select_args)
get_socket_select_semaphore returns semaphore allocated in the socket driver; set only for the socket driver

Macros

MAX_FDS
Maximum number of (global) file descriptors.

ESP_VFS_PATH_MAX
Maximum length of path prefix (not including zero terminator)

ESP_VFS_FLAG_DEFAULT
Default value of flags member in *esp_vfs_t* structure.

ESP_VFS_FLAG_CONTEXT_PTR
Flag which indicates that FS needs extra context pointer in syscalls.

Type Definitions

typedef int **esp_vfs_id_t**

Header File

- vfs/include/esp_vfs_dev.h

Functions

void **esp_vfs_dev_uart_register** (void)
add /dev/uart virtual filesystem driver

This function is called from startup code to enable serial output

void **esp_vfs_dev_uart_set_rx_line_endings** (*esp_line_endings_t mode*)
Set the line endings expected to be received on UART.

This specifies the conversion between line endings received on UART and newlines (‘’ , LF) passed into stdin:

- ESP_LINE_ENDINGS_CRLF: convert CRLF to LF

- ESP_LINE_ENDINGS_CR: convert CR to LF
- ESP_LINE_ENDINGS_LF: no modification

Note this function is not thread safe w.r.t. reading from UART
Parameters
- `mode`: line endings expected on UART

void **esp_vfs_dev_uart_set_tx_line_endings**(*esp_line_endings_t mode*)
Set the line endings to sent to UART.

This specifies the conversion between newlines (‘’ , LF) on stdout and line endings sent over UART:

- ESP_LINE_ENDINGS_CRLF: convert LF to CRLF
- ESP_LINE_ENDINGS_CR: convert LF to CR
- ESP_LINE_ENDINGS_LF: no modification

Note this function is not thread safe w.r.t. writing to UART
Parameters
- `mode`: line endings to send to UART

int **esp_vfs_dev_uart_port_set_rx_line_endings**(int *uart_num*, *esp_line_endings_t mode*)
Set the line endings expected to be received on specified UART.

This specifies the conversion between line endings received on UART and newlines (‘’ , LF) passed into stdin:

- ESP_LINE_ENDINGS_CRLF: convert CRLF to LF
- ESP_LINE_ENDINGS_CR: convert CR to LF
- ESP_LINE_ENDINGS_LF: no modification

Note this function is not thread safe w.r.t. reading from UART
Return 0 if successed, or -1 when an error (specified by errno) have occurred.
Parameters
- `uart_num`: the UART number
- `mode`: line endings to send to UART

int **esp_vfs_dev_uart_port_set_tx_line_endings**(int *uart_num*, *esp_line_endings_t mode*)
Set the line endings to sent to specified UART.

This specifies the conversion between newlines (‘’ , LF) on stdout and line endings sent over UART:

- ESP_LINE_ENDINGS_CRLF: convert LF to CRLF
- ESP_LINE_ENDINGS_CR: convert LF to CR
- ESP_LINE_ENDINGS_LF: no modification

Note this function is not thread safe w.r.t. writing to UART
Return 0 if successed, or -1 when an error (specified by errno) have occurred.
Parameters
- `uart_num`: the UART number
- `mode`: line endings to send to UART

void **esp_vfs_dev_uart_use_nonblocking**(int *uart_num*)
set VFS to use simple functions for reading and writing UART Read is non-blocking, write is busy waiting until TX FIFO has enough space. These functions are used by default.

Parameters
- `uart_num`: UART peripheral number

void **esp_vfs_dev_uart_use_driver**(int *uart_num*)
set VFS to use UART driver for reading and writing

Note application must configure UART driver before calling these functions With these functions, read and write are blocking and interrupt-driven.
Parameters
- `uart_num`: UART peripheral number

Enumerations

enum esp_line_endings_t

Line ending settings.

Values:

ESP_LINE_ENDINGS_CRLF

CR + LF.

ESP_LINE_ENDINGS_CR

CR.

ESP_LINE_ENDINGS_LF

LF.

2.5.6 FAT 文件系统

ESP-IDF 使用 FatFs 库来实现 FAT 文件系统。FatFs 库位于 `fatfs` 组件中，您可以直接使用，也可以借助 C 标准库和 POSIX API 通过 VFS（虚拟文件系统）使用 FatFs 库的大多数功能。

此外，我们对 FatFs 库进行了扩展，新增了支持可插拔磁盘 I/O 调度层，从而允许在运行时将 FatFs 驱动映射到物理磁盘。

FatFs 与 VFS 配合使用

fatfs/vfs/esp_vfs_fat.h 头文件定义了连接 FatFs 和 VFS 的函数。

函数*esp_vfs_fat_register()* 分配一个 `FATFS` 结构，并在 VFS 中注册特定路径前缀。如果文件路径以此前缀开头，则对此文件的后续操作将转至 FatFs API。函数*esp_vfs_fat_unregister_path()* 删除在 VFS 中的注册，并释放 `FATFS` 结构。

多数应用程序在使用 `esp_vfs_fat_` 函数时，采用如下步骤：

1. 调用*esp_vfs_fat_register()*，指定：
   - 挂载文件系统的路径前缀（例如，`"/sdcard"` 或 `"/spiflash"`）
   - FatFs 驱动编号
   - 一个用于接收指向 `FATFS` 结构指针的变量
2. 调用*ff_diskio_register()* 为上述步骤中的驱动编号注册磁盘 I/O 驱动；
3. 调用 FatFs 函数 `f_mount`，或 `f_fdisk`，`f_mkfs`，并使用与传递到*esp_vfs_fat_register()* 相同的驱动编号挂载文件系统。请参考 FatFs 文档，查看更多信息；
4. 调用 C 标准库和 POSIX API 对路径中带有步骤 1 中所述前缀的文件（例如，`"/sdcard/hello.txt"`）执行打开、读取、写入、擦除、复制等操作。
5. 您可以选择直接调用 FatFs 库函数，但需要使用没有 VFS 前缀的路径（例如，`"/hello.txt"`）；
6. 关闭所有打开的文件；
7. 调用 `f_mount` 并使用 NULL `FATFS*` 参数为与上述编号相同的驱动卸载文件系统；
8. 调用 FatFs 函数*ff_diskio_register()* 并使用 NULL `ff_diskio_impl_t*` 参数和相同的驱动编号来释放注册的磁盘 I/O 驱动。
9. 调用*esp_vfs_fat_unregister_path()* 并使用文件系统挂载的路径将 FatFs 从 NVS 中移除，并释放步骤 1 中分配的 FatFs 结构。

`esp_vfs_fat_sdmmc_mount` 和 `esp_vfs_fat_sdmmc_unmount` 这两个便捷函数对上述步骤进行了封装，并加入对 SD 卡初始化的处理，非常便捷。我们将在下一章节详细介绍这两个函数。

esp_err_t **esp_vfs_fat_register**(**const** char **base_path*, **const** char **fat_drive*, size_t *max_files*, FATFS ***out_fs*)

Register FATFS with VFS component.

This function registers given FAT drive in VFS, at the specified base path. If only one drive is used, fat_drive argument can be an empty string. Refer to FATFS library documentation on how to specify FAT drive. This function also allocates FATFS structure which should be used for f_mount call.

Note This function doesn’t mount the drive into FATFS, it just connects POSIX and C standard library IO function with FATFS. You need to mount desired drive into FATFS separately.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if esp_vfs_fat_register was already called
- ESP_ERR_NO_MEM if not enough memory or too many VFSes already registered

Parameters

- `base_path`: path prefix where FATFS should be registered
- `fat_drive`: FATFS drive specification; if only one drive is used, can be an empty string
- `max_files`: maximum number of files which can be open at the same time
- `[out] out_fs`: pointer to FATFS structure which can be used for FATFS f_mount call is returned via this argument.

esp_err_t **esp_vfs_fat_unregister_path**(**const** char **base_path*)

Un-register FATFS from VFS.

Note FATFS structure returned by esp_vfs_fat_register is destroyed after this call. Make sure to call f_mount function to unmount it before calling esp_vfs_fat_unregister_ctx. Difference between this function and the one above is that this one will release the correct drive, while the one above will release the last registered one

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if FATFS is not registered in VFS

Parameters

- `base_path`: path prefix where FATFS is registered. This is the same used when esp_vfs_fat_register was called

FatFs 与 VFS 和 SD 卡配合使用

fatfs/vfs/esp_vfs_fat.h 头文件定义了两个便捷函数*esp_vfs_fat_sdmmc_mount()*和*esp_vfs_fat_sdmmc_unmount()*。这两个函数分别执行上一章节的步骤 1-3 和步骤 7-9，并初始化 SD 卡，但仅提供有限的错误处理功能。我们鼓励开发人员查看源代码并将更多高级功能集成到产品应用中。

esp_vfs_fat_sdmmc_unmount() 函数用于卸载文件系统并释放从*esp_vfs_fat_sdmmc_mount()* 函数获取的资源。

esp_err_t **esp_vfs_fat_sdmmc_mount**(**const** char **base_path*, **const** *sdmmc_host_t* **host_config*, **const** void **slot_config*, **const** *esp_vfs_fat_mount_config_t* **mount_config*, *sdmmc_card_t* ***out_card*)

Convenience function to get FAT filesystem on SD card registered in VFS.

This is an all-in-one function which does the following:

- initializes SDMMC driver or SPI driver with configuration in host_config
- initializes SD card with configuration in slot_config
- mounts FAT partition on SD card using FATFS library, with configuration in mount_config
- registers FATFS library with VFS, with prefix given by base_prefix variable

This function is intended to make example code more compact. For real world applications, developers should implement the logic of probing SD card, locating and mounting partition, and registering FATFS in VFS, with proper error checking and handling of exceptional conditions.

Note Use this API to mount a card through SDSPI is deprecated. Please call `esp_vfs_fat_sdspi_mount()` instead for that case.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if esp_vfs_fat_sdmmc_mount was already called
- ESP_ERR_NO_MEM if memory can not be allocated
- ESP_FAIL if partition can not be mounted
- other error codes from SDMMC or SPI drivers, SDMMC protocol, or FATFS drivers

Parameters

- `base_path`: path where partition should be registered (e.g. “/sdcard”)

- `host_config`: Pointer to structure describing SDMMC host. When using SDMMC peripheral, this structure can be initialized using SDMMC_HOST_DEFAULT() macro. When using SPI peripheral, this structure can be initialized using SDSPI_HOST_DEFAULT() macro.
- `slot_config`: Pointer to structure with slot configuration. For SDMMC peripheral, pass a pointer to sdmmc_slot_config_t structure initialized using SDMMC_SLOT_CONFIG_DEFAULT. (Deprecated) For SPI peripheral, pass a pointer to *sdspi_slot_config_t* structure initialized using SDSPI_SLOT_CONFIG_DEFAULT().
- `mount_config`: pointer to structure with extra parameters for mounting FATFS
- `[out] out_card`: if not NULL, pointer to the card information structure will be returned via this argument

struct esp_vfs_fat_mount_config_t

Configuration arguments for esp_vfs_fat_sdmmc_mount and esp_vfs_fat_spiflash_mount functions.

Public Members

bool **format_if_mount_failed**

If FAT partition can not be mounted, and this parameter is true, create partition table and format the filesystem.

int **max_files**

Max number of open files.

size_t **allocation_unit_size**

If format_if_mount_failed is set, and mount fails, format the card with given allocation unit size. Must be a power of 2, between sector size and 128 * sector size. For SD cards, sector size is always 512 bytes. For wear_levelling, sector size is determined by CONFIG_WL_SECTOR_SIZE option.

Using larger allocation unit size will result in higher read/write performance and higher overhead when storing small files.

Setting this field to 0 will result in allocation unit set to the sector size.

esp_err_t **esp_vfs_fat_sdmmc_unmount**(void)

Unmount FAT filesystem and release resources acquired using esp_vfs_fat_sdmmc_mount.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if esp_vfs_fat_sdmmc_mount hasn't been called

FatFs 与 VFS 配合使用（只读模式下）

fatfs/vfs/esp_vfs_fat.h 头文件也定义了两个便捷函数*esp_vfs_fat_rawflash_mount()*和*esp_vfs_fat_rawflash_unmount()*。上述两个函数分别对 FAT 只读分区执行步骤 1-3 和步骤 7-9。有些数据分区仅在工厂时写入一次，之后在整个硬件生命周期内都不会再有任何改动。利用上述两个函数处理这种数据分区非常方便。

esp_err_t **esp_vfs_fat_rawflash_mount**(**const** char **base_path*, **const** char **partition_label*, **const** *esp_vfs_fat_mount_config_t* **mount_config*)

Convenience function to initialize read-only FAT filesystem and register it in VFS.

This is an all-in-one function which does the following:

- finds the partition with defined partition_label. Partition label should be configured in the partition table.
- mounts FAT partition using FATFS library
- registers FATFS library with VFS, with prefix given by base_prefix variable

Note Wear levelling is not used when FAT is mounted in read-only mode using this function.

Return

- ESP_OK on success
- ESP_ERR_NOT_FOUND if the partition table does not contain FATFS partition with given label

- ESP_ERR_INVALID_STATE if esp_vfs_fat_rawflash_mount was already called for the same partition
- ESP_ERR_NO_MEM if memory can not be allocated
- ESP_FAIL if partition can not be mounted
- other error codes from SPI flash driver, or FATFS drivers

Parameters

- `base_path`: path where FATFS partition should be mounted (e.g. "/spiflash")
- `partition_label`: label of the partition which should be used
- `mount_config`: pointer to structure with extra parameters for mounting FATFS

esp_err_t **esp_vfs_fat_rawflash_unmount**(**const** char **base_path*, **const** char **partition_label*)

Unmount FAT filesystem and release resources acquired using esp_vfs_fat_rawflash_mount.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if esp_vfs_fat_spiflash_mount hasn't been called

Parameters

- `base_path`: path where partition should be registered (e.g. "/spiflash")
- `partition_label`: label of partition to be unmounted

FatFs 磁盘 I/O 层

我们对 FatFs API 函数进行了扩展，实现了运行期间注册磁盘 I/O 驱动。

上述 API 为 SD/MMC 卡提供了磁盘 I/O 函数实现方式，可使用 *`ff_diskio_register_sdmmc()`* 注册指定的 FatFs 驱动编号。

void **ff_diskio_register**(BYTE *pdrv*, **const** *ff_diskio_impl_t* **discio_impl*)

Register or unregister diskio driver for given drive number.

When FATFS library calls one of disk_xxx functions for driver number pdrv, corresponding function in discio_impl for given pdrv will be called.

Parameters

- `pdrv`: drive number
- `discio_impl`: pointer to *ff_diskio_impl_t* structure with diskio functions or NULL to unregister and free previously registered drive

struct ff_diskio_impl_t

Structure of pointers to disk IO driver functions.

See FatFs documentation for details about these functions

Public Members

DSTATUS (***init**)(unsigned char pdrv)

disk initialization function

DSTATUS (***status**)(unsigned char pdrv)

disk status check function

DRESULT (***read**)(unsigned char pdrv, unsigned char *buff, uint32_t sector, unsigned count)

sector read function

DRESULT (***write**)(unsigned char pdrv, **const** unsigned char *buff, uint32_t sector, unsigned count)

sector write function

DRESULT (***ioctl**)(unsigned char pdrv, unsigned char cmd, void *buff)

function to get info about disk and do some misc operations

void **ff_diskio_register_sdmmc**(unsigned char *pdrv*, *sdmmc_card_t* **card*)
Register SD/MMC diskio driver

Parameters
- `pdrv`: drive number
- `card`: pointer to *sdmmc_card_t* structure describing a card; card should be initialized before calling f_mount.

esp_err_t **ff_diskio_register_wl_partition**(unsigned char *pdrv*, *wl_handle_t flash_handle*)
Register spi flash partition

Parameters
- `pdrv`: drive number
- `flash_handle`: handle of the wear levelling partition.

esp_err_t **ff_diskio_register_raw_partition**(unsigned char *pdrv*, **const** *esp_partition_t* **part_handle*)
Register spi flash partition

Parameters
- `pdrv`: drive number
- `part_handle`: pointer to raw flash partition.

2.5.7 磨损均衡 API

概述

ESP32 所使用的 flash，特别是 SPI flash 多数具备扇区结构，且每个扇区仅允许有限次数的擦除/修改操作。为了避免过度使用某一扇区，乐鑫提供了磨损均衡组件，无需用户介入即可帮助用户均衡各个扇区之间的磨损。

磨损均衡组件包含了通过分区组件对外部 SPI flash 进行数据读取、写入、擦除和存储器映射相关的 API 函数。磨损均衡组件还具有软件上更高级别的 API 函数，与*FAT 文件系统* 协同工作。

磨损均衡组件与 FAT 文件系统组件共用 FAT 文件系统的扇区，扇区大小为 4096 字节，是标准 flash 扇区的大小。在这种模式下，磨损均衡组件性能达到最佳，但需要在 RAM 中占用更多内存。

为了节省内存，磨损均衡组件还提供了另外两种模式，均使用 512 字节大小的扇区：

- **性能模式**：先将数据保存在 RAM 中，擦除扇区，然后将数据存储回 flash。如果设备在扇区擦写过程中突然断电，则整个扇区（4096 字节）数据将全部丢失。
- **安全模式**：数据先保存在 flash 中空余扇区，擦除扇区后，数据即存储回去。如果设备断电，上电后可立即恢复数据。

设备默认设置如下：

- 定义扇区大小为 512 字节
- 默认使用性能模式

您可以使用配置菜单更改设置。

磨损均衡组件不会将数据缓存在 RAM 中。写入和擦除函数直接修改 flash，函数返回后，flash 即完成修改。

磨损均衡访问 API

处理 flash 数据常用的 API 如下所示：

- `wl_mount` - 为指定分区挂载并初始化磨损均衡模块
- `wl_unmount` - 卸载分区并释放磨损均衡模块
- `wl_erase_range` - 擦除 flash 中指定的地址范围
- `wl_write` - 将数据写入分区
- `wl_read` - 从分区读取数据

- `wl_size` - 返回可用内存的大小（以字节为单位）
- `wl_sector_size` - 返回一个扇区的大小

请尽量避免直接使用原始磨损均衡函数，建议您使用文件系统特定的函数。

内存大小

内存大小是根据分区参数在磨损均衡模块中计算所得，由于模块使用 flash 部分扇区存储内部数据，因此计算所得内存大小有少许偏差。

另请参阅

- *FAT 文件系统*
- 分区表

应用示例

storage/wear_levelling 中提供了一款磨损均衡驱动与 FatFs 库结合使用的示例。该示例初始化磨损均衡驱动，挂载 FAT 文件系统分区，并使用 POSIX（可移植操作系统接口）和 C 库 API 从中写入和读取数据。如需了解更多信息，请参考 storage/wear_levelling/README.md。

高级 API 参考

头文件

- fatfs/vfs/esp_vfs_fat.h

函数

esp_err_t **esp_vfs_fat_spiflash_mount**(**const** char **base_path*, **const** char **partition_label*, **const** *esp_vfs_fat_mount_config_t* **mount_config*, *wl_handle_t* **wl_handle*)

Convenience function to initialize FAT filesystem in SPI flash and register it in VFS.

This is an all-in-one function which does the following:

- finds the partition with defined partition_label. Partition label should be configured in the partition table.
- initializes flash wear levelling library on top of the given partition
- mounts FAT partition using FATFS library on top of flash wear levelling library
- registers FATFS library with VFS, with prefix given by base_prefix variable

This function is intended to make example code more compact.

Return

- ESP_OK on success
- ESP_ERR_NOT_FOUND if the partition table does not contain FATFS partition with given label
- ESP_ERR_INVALID_STATE if esp_vfs_fat_spiflash_mount was already called
- ESP_ERR_NO_MEM if memory can not be allocated
- ESP_FAIL if partition can not be mounted
- other error codes from wear levelling library, SPI flash driver, or FATFS drivers

Parameters

- `base_path`: path where FATFS partition should be mounted (e.g. “/spiflash”)
- `partition_label`: label of the partition which should be used
- `mount_config`: pointer to structure with extra parameters for mounting FATFS
- `[out] wl_handle`: wear levelling driver handle

struct esp_vfs_fat_mount_config_t

Configuration arguments for esp_vfs_fat_sdmmc_mount and esp_vfs_fat_spiflash_mount functions.

Public Members

bool **format_if_mount_failed**
If FAT partition can not be mounted, and this parameter is true, create partition table and format the filesystem.

int **max_files**
Max number of open files.

size_t **allocation_unit_size**
If format_if_mount_failed is set, and mount fails, format the card with given allocation unit size. Must be a power of 2, between sector size and 128 * sector size. For SD cards, sector size is always 512 bytes. For wear_levelling, sector size is determined by CONFIG_WL_SECTOR_SIZE option.

Using larger allocation unit size will result in higher read/write performance and higher overhead when storing small files.

Setting this field to 0 will result in allocation unit set to the sector size.

esp_err_t **esp_vfs_fat_spiflash_unmount**(**const** char **base_path*, *wl_handle_t wl_handle*)
Unmount FAT filesystem and release resources acquired using esp_vfs_fat_spiflash_mount.

Return
- ESP_OK on success
- ESP_ERR_INVALID_STATE if esp_vfs_fat_spiflash_mount hasn't been called

Parameters
- `base_path`: path where partition should be registered (e.g. "/spiflash")
- `wl_handle`: wear levelling driver handle returned by esp_vfs_fat_spiflash_mount

中层 API 参考

Header File

- wear_levelling/include/wear_levelling.h

Functions

esp_err_t **wl_mount**(**const** *esp_partition_t *partition*, *wl_handle_t *out_handle*)
Mount WL for defined partition.

Return
- ESP_OK, if the allocation was successfully;
- ESP_ERR_INVALID_ARG, if WL allocation was unsuccessful;
- ESP_ERR_NO_MEM, if there was no memory to allocate WL components;

Parameters
- `partition`: that will be used for access
- `out_handle`: handle of the WL instance

esp_err_t **wl_unmount**(*wl_handle_t handle*)
Unmount WL for defined partition.

Return
- ESP_OK, if the operation completed successfully;
- or one of error codes from lower-level flash driver.

Parameters
- `handle`: WL partition handle

esp_err_t **wl_erase_range**(*wl_handle_t handle*, size_t *start_addr*, size_t *size*)
Erase part of the WL storage.

Return
- ESP_OK, if the range was erased successfully;
- ESP_ERR_INVALID_ARG, if iterator or dst are NULL;
- ESP_ERR_INVALID_SIZE, if erase would go out of bounds of the partition;

- or one of error codes from lower-level flash driver.

Parameters

- `handle`: WL handle that are related to the partition
- `start_addr`: Address where erase operation should start. Must be aligned to the result of function wl_sector_size(⋯).
- `size`: Size of the range which should be erased, in bytes. Must be divisible by result of function wl_sector_size(⋯)..

esp_err_t **wl_write**(*wl_handle_t handle*, size_t *dest_addr*, **const** void **src*, size_t *size*)

Write data to the WL storage.

Before writing data to flash, corresponding region of flash needs to be erased. This can be done using wl_erase_range function.

Note Prior to writing to WL storage, make sure it has been erased with wl_erase_range call.

Return

- ESP_OK, if data was written successfully;
- ESP_ERR_INVALID_ARG, if dst_offset exceeds partition size;
- ESP_ERR_INVALID_SIZE, if write would go out of bounds of the partition;
- or one of error codes from lower-level flash driver.

Parameters

- `handle`: WL handle that are related to the partition
- `dest_addr`: Address where the data should be written, relative to the beginning of the partition.
- `src`: Pointer to the source buffer. Pointer must be non-NULL and buffer must be at least ‘size’ bytes long.
- `size`: Size of data to be written, in bytes.

esp_err_t **wl_read**(*wl_handle_t handle*, size_t *src_addr*, void **dest*, size_t *size*)

Read data from the WL storage.

Return

- ESP_OK, if data was read successfully;
- ESP_ERR_INVALID_ARG, if src_offset exceeds partition size;
- ESP_ERR_INVALID_SIZE, if read would go out of bounds of the partition;
- or one of error codes from lower-level flash driver.

Parameters

- `handle`: WL module instance that was initialized before
- `dest`: Pointer to the buffer where data should be stored. Pointer must be non-NULL and buffer must be at least ‘size’ bytes long.
- `src_addr`: Address of the data to be read, relative to the beginning of the partition.
- `size`: Size of data to be read, in bytes.

size_t **wl_size**(*wl_handle_t handle*)

Get size of the WL storage.

Return usable size, in bytes

Parameters

- `handle`: WL module handle that was initialized before

size_t **wl_sector_size**(*wl_handle_t handle*)

Get sector size of the WL instance.

Return sector size, in bytes

Parameters

- `handle`: WL module handle that was initialized before

Macros

WL_INVALID_HANDLE

Type Definitions

typedef int32_t **wl_handle_t**
wear levelling handle

2.5.8 SPIFFS 文件系统

概述

SPIFFS 是一个用于 SPI NOR flash 设备的嵌入式文件系统，支持磨损均衡、文件系统一致性检查等功能。

说明

- 目前，SPIFFS 尚不支持目录，但可以生成扁平结构。如果 SPIFFS 挂载在 `/spiffs` 下，在 `/spiffs/tmp/myfile.txt` 路径下创建一个文件则会在 SPIFFS 中生成一个名为 `/tmp/myfile.txt` 的文件，而不是在 `/spiffs/tmp` 下生成名为 `myfile.txt` 的文件；
- SPIFFS 并非实时栈，每次写操作耗时不等；
- 目前，SPIFFS 尚不支持检测或处理已损坏的块。

工具

spiffsgen.py spiffsgen.py:

```
python spiffsgen.py <image_size> <base_dir> <output_file>
```

参数（必选）说明如下：

- **image_size**：分区大小，用于烧录生成的 SPIFFS 映像；
- **base_dir**：创建 SPIFFS 映像的目录；
- **output_file**：SPIFFS 映像输出文件。

其他参数（可选）也参与控制映像的生成，您可以运行以下帮助命令，查看这些参数的具体信息:

```
python spiffsgen.py --help
```

上述可选参数对应 SPIFFS 构建配置选项。若想顺利生成可用的映像，请确保使用的参数或配置与构建 SPIFFS 时所用的参数或配置相同。运行帮助命令将显示参数所对应的 SPIFFS 构建配置。如未指定参数，将使用帮助信息中的默认值。

映像生成后，您可以使用 `esptool.py` 或 `parttool.py` 烧录映像。

您可以在命令行或脚本中手动单独调用 `spiffsgen.py`，也可以直接从构建系统调用 `spiffs_create_partition_image` 来使用 `spiffsgen.py`。

在 Make 构建系统中运行:

```
SPIFFS_IMAGE_FLASH_IN_PROJECT := ...
SPIFFS_IMAGE_DEPENDS := ...
$(eval $(call spiffs_create_partition_image,<partition>,<base_dir>))
```

在 CMake 构建系统中运行:

```
spiffs_create_partition_image(<partition> <base_dir> [FLASH_IN_PROJECT] [DEPENDS␣
↪dep dep dep...])
```

在构建系统中使用 `spiffsgen.py` 更为方便，构建配置自动传递给 `spiffsgen.py` 工具，确保生成的映像可用于构建。比如，单独调用 `spiffsgen.py` 时需要用到 *image_size* 参数，但在构建系统中调用 `spiffs_create_partition_image` 时，仅需要 *partition* 参数，映像大小将直接从工程分区表中获取。

Make 构建系统和 CMake 构建系统结构有所不同，请注意以下几点：

- 在 Make 构建系统中使用 `spiffs_create_partition_image`，需从工程 Makefile 中调用；

- 在 CMake 构建系统中使用 `spiffs_create_partition_image`，需从组件 CMakeLists.txt 文件调用。

您也可以指定 `FLASH_IN_PROJECT`，然后使用 `idf.py flash` 或 `make flash` 将映像与应用程序二进制文件、分区表等一起自动烧录至设备，例如：

在 Make 构建系统中运行：

```
SPIFFS_IMAGE_FLASH_IN_PROJECT := 1
$(eval $(call spiffs_create_partition_image,<partition>,<base_dir>))
```

在 CMake 构建系统中运行：

```
spiffs_create_partition_image(my_spiffs_partition my_folder FLASH_IN_PROJECT)
```

不指定 FLASH_IN_PROJECT/SPIFFS_IMAGE_FLASH_IN_PROJECT 也可以生成映像，但须使用 `esptool.py`、`parttool.py` 或自定义构建系统目标手动烧录。

有时基本目录中的内容是在构建时生成的，您可以使用 DEPENDS/SPIFFS_IMAGE_DEPENDS 指定目标，因此可以在生成映像之前执行此目标。

在 Make 构建系统中运行：

```
dep:
    ...

SPIFFS_IMAGE_DEPENDS := dep
$(eval $(call spiffs_create_partition_image,<partition>,<base_dir>))
```

在 CMake 构建系统中运行：

```
add_custom_target(dep COMMAND ...)

spiffs_create_partition_image(my_spiffs_partition my_folder DEPENDS dep)
```

请参考 storage/spiffsgen，查看示例。

mkspiffs 您也可以使用 mkspiffs 工具创建 SPIFFS 分区映像。与 `spiffsgen.py` 相似，mkspiffs 也可以用于从指定文件夹中生成映像，然后使用 `esptool.py` 烧录映像。

该工具需要获取以下参数：

- **Block Size**：4096（SPI flash 标准）
- **Page Size**：256（SPI flash 标准）
- **Image Size**：分区大小（以字节为单位，可从分区表中获取）
- **Partition Offset**：分区起始地址（可从分区表内获取）

运行以下命令，将文件夹打包成 1 MB 大小的映像：

```
mkspiffs -c [src_folder] -b 4096 -p 256 -s 0x100000 spiffs.bin
```

运行以下命令，将映像烧录到 ESP32（偏移量：0x110000）：

```
python esptool.py --chip esp32s2 --port [port] --baud [baud] write_flash -z␣
↪0x110000 spiffs.bin
```

选择合适的 SPIFFS 工具 上面介绍的两款 SPIFFS 工具功能相似，需根据实际情况，选择合适的一款。

以下情况优先选用 `spiffsgen.py` 工具：

1. 仅需在构建时简单生成 SPIFFS 映像，请选择使用 `spiffsgen.py`，`spiffsgen.py` 可以直接在构建系统中使用函数或命令生成 SPIFFS 映像。

2. 主机没有可用的 C/C++ 编译器时，可以选择使用 `spiffsgen.py` 工具，`spiffsgen.py` 不需要编译。

以下情况优先选用 `mkspiffs` 工具：

1. 如果您除了需要生成映像外，还需要拆包 SPIFFS 映像，请选择使用 `mkspiffs` 工具。`spiffsgen.py` 目前尚不支持此功能。
2. 如果您当前环境中 Python 解释器不可用，但主机编译器可用，或者有预编译的 `mkspiffs` 二进制文件，此时请选择使用 `mkspiffs` 工具。但是，`mkspiffs` 没有集成到构建系统，用户必须自己完成以下工作：在构建期间编译 `mkspiffs` （如果未使用预编译的二进制文件），为输出文件创建构建规则或目标，将适当的参数传递给工具等。

另请参阅

- 分区表

应用示例

storage/spiffs 目录下提供了 SPIFFS 应用示例。该示例初始化并挂载了一个 SPIFFS 分区，然后使用 POSIX 和 C 库 API 写入和读取数据。请参考 `example` 目录下的 README.md 文件，查看详细信息。

高级 API 参考

Header File

- spiffs/include/esp_spiffs.h

Functions

esp_err_t **esp_vfs_spiffs_register**(**const** *esp_vfs_spiffs_conf_t* **conf*)

Register and mount SPIFFS to VFS with given path prefix.

Return

- ESP_OK if success
- ESP_ERR_NO_MEM if objects could not be allocated
- ESP_ERR_INVALID_STATE if already mounted or partition is encrypted
- ESP_ERR_NOT_FOUND if partition for SPIFFS was not found
- ESP_FAIL if mount or format fails

Parameters

- `conf`: Pointer to *esp_vfs_spiffs_conf_t* configuration structure

esp_err_t **esp_vfs_spiffs_unregister**(**const** char **partition_label*)

Unregister and unmount SPIFFS from VFS

Return

- ESP_OK if successful
- ESP_ERR_INVALID_STATE already unregistered

Parameters

- `partition_label`: Same label as passed to esp_vfs_spiffs_register.

bool **esp_spiffs_mounted**(**const** char **partition_label*)

Check if SPIFFS is mounted

Return

- true if mounted
- false if not mounted

Parameters

- `partition_label`: Optional, label of the partition to check. If not specified, first partition with subtype=spiffs is used.

esp_err_t **esp_spiffs_format**(**const** char **partition_label*)
Format the SPIFFS partition

Return
- ESP_OK if successful
- ESP_FAIL on error

Parameters
- `partition_label`: Same label as passed to esp_vfs_spiffs_register.

esp_err_t **esp_spiffs_info**(**const** char **partition_label*, size_t **total_bytes*, size_t **used_bytes*)
Get information for SPIFFS

Return
- ESP_OK if success
- ESP_ERR_INVALID_STATE if not mounted

Parameters
- `partition_label`: Same label as passed to esp_vfs_spiffs_register
- `[out] total_bytes`: Size of the file system
- `[out] used_bytes`: Current used bytes in the file system

Structures

struct esp_vfs_spiffs_conf_t
Configuration structure for esp_vfs_spiffs_register.

Public Members

const char ***base_path**
File path prefix associated with the filesystem.

const char ***partition_label**
Optional, label of SPIFFS partition to use. If set to NULL, first partition with subtype=spiffs will be used.

size_t **max_files**
Maximum files that could be open at the same time.

bool **format_if_mount_failed**
If true, it will format the file system if it fails to mount.

2.5.9 量产程序

介绍

这一程序主要用于量产时为每一设备创建工厂 NVS（非易失性存储器）分区映像。NVS 分区映像由 CSV（逗号分隔值）文件生成，文件中包含了用户提供的配置项及配置值。

注意，该程序仅创建用于量产的二进制映像，您需要使用以下工具将映像烧录到设备上：

- esptool.py
- Flash 下载工具（仅适用于 Windows）
- 直接烧录程序

准备工作

该程序需要用到分区公用程序。

- **操作系统要求：**
  - Linux、MacOS 或 Windows（标准版）
- **安装依赖包：**

– Python：https://www.python.org/downloads/。

注解：

使用该程序之前，请确保：

- Python 路径已添加到 PATH 环境变量中；
- 已经安装 *requirement.txt* 中的软件包，*requirement.txt* 在 esp-idf 根目录下。

具体流程

CSV 配置文件

CSV 配置文件中包含设备待烧录的配置信息，定义了待烧录的配置项。例如定义 `firmware_key` (`key`) 的 `type` 为 `data`，`encoding` 为 `hex2bin`。

配置文件中数据格式如下（*REPEAT* 标签可选）：

```
name1,namespace,     <-- 第一行为 "namespace" 条目
key1,type1,encoding1
key2,type2,encoding2,REPEAT
name2,namespace,
key3,type3,encoding3
key4,type4,encoding4
```

注解： 文件第一行应始终为 `namespace` 条目。

每行应包含三个参数：`key`、`type` 和 `encoding`，并以逗号分隔。如果有 `REPEAT` 标签，则主 CSV 文件中所有设备此键值均相同。

有关各个参数的详细说明，请参阅 *NVS* 分区生成程序的 *README* 文件。

CSV 配置文件示例如下:

```
app,namespace,
firmware_key,data,hex2bin
serial_no,data,string,REPEAT <-- "serial_no" 被标记为 "REPEAT"
device_no,data,i32
```

注解：

请确保：

- 逗号‘,’前后无空格；
- CSV 文件每行末尾无空格。

主 CSV 文件

主 CSV 文件中包含设备待烧录的详细信息，文件中每行均对应一个设备实体。主 CSV 文件中的 `key` 应首先在 CSV 配置文件中定义。

主 CSV 文件的数据格式如下:

```
key1,key2,key3,.....
value1,value2,value3,.... <-- 对应一个设备实体
value4,value5,value6,.... <-- 对应一个设备实体
value7,value8,value9,.... <-- 对应一个设备实体
```

注解: 文件中键 (`key`) 名应始终置于文件首行。从配置文件中获取的键，在此文件中的排列顺序应与其在配置文件中的排列顺序相同。主 CSV 文件同时可以包含其它列（键），这些列将被视为元数据，而不会编译进最终二进制文件。

每行应包含相应键的键值 (`value`)，并用逗号隔开。如果某键带有 `REPEAT` 标签，则仅需在第二行（即第一个条目）输入对应的值，后面其他行为空。

参数描述如下:

`value` Data value

`value` 是与键对应的键值。

主 CSV 文件示例如下:

```
id,firmware_key,serial_no,device_no
1,1a2b3c4d5e6faabb,A1,101 <-- 对应一个设备实体（在 CSV 配置文件中标记为 `REPEAT` 的键，除第一个条目外，其他均为空）
2,1a2b3c4d5e6fccdd,,102   <-- 对应一个设备实体
3,1a2b3c4d5e6feeff,,103   <-- 对应一个设备实体
```

注解: 如果出现 *REPEAT* 标签，则会在相同目录下生成一个新的主 CSV 文件用作主输入文件，并在每行为带有 *REPEAT* 标签的键插入键值。

量产程序还会创建中间 CSV 文件，NVS 分区程序将使用此 CSV 文件作为输入，然后生成二进制文件。

中间 CSV 文件的格式如下:

```
key,type,encoding,value
key,namespace, ,
key1,type1,encoding1,value1
key2,type2,encoding2,value2
```

此步骤将为每一设备生成一个中间 CSV 文件。

运行量产程序

使用方法:

```
python mfg_gen.py [-h] {generate,generate-key} ...
```

可选参数:

序号	参数	描述
1	-h, –help	显示帮助信息并退出

命令:

运行 mfg_gen.py {command} -h 查看更多帮助信息

序号	参数	描述
1	generate	生成 NVS 分区
2	generate-key	生成加密密钥

为每个设备生成工厂映像（默认）

使用方法:

```
python mfg_gen.py generate [-h] [--fileid FILEID] [--version {1,2}] [--keygen]
                           [--keyfile KEYFILE] [--inputkey INPUTKEY]
                           [--outdir OUTDIR]
                           conf values prefix size
```

位置参数:

参数	描述
conf	待解析的 CSV 配置文件路径
values	待解析的主 CSV 文件路径
prefix	每个输出文件名前缀的唯一名称
size	NVS 分区大小（以字节为单位，且为 4096 的整数倍）

可选参数:

参数	描述
-h, –help	显示帮助信息并退出
–fileid FILEID	每个文件名后缀的唯一文件标识符（主 CSV 文件中的任意键），默认为数值 1、2、3…
–version {1,2}	• 设置多页 Blob 版本。 • 版本 1 - 禁用多页 Blob； • 版本 2 - 启用多页 Blob； • 默认版本：版本 2
–keygen	生成 NVS 分区加密密钥
–inputkey INPUTKEY	内含 NVS 分区加密密钥的文件
–outdir OUTDIR	输出目录，用于存储创建的文件（默认当前目录）

请运行以下命令为每个设备生成工厂映像，量产程序同时提供了一个 CSV 示例文件:

```
python mfg_gen.py generate samples/sample_config.csv samples/sample_values_
↪singlepage_blob.csv Sample 0x3000
```

主 CSV 文件应在 `file` 类型下设置一个相对路径，指向运行该程序的当前目录。

为每个设备生成工厂加密映像

运行以下命令为每一设备生成工厂加密映像，量产程序同时提供了一个 CSV 示例文件。

- 通过量产程序生成加密密钥来进行加密:

```
python mfg_gen.py generate samples/sample_config.csv samples/sample_values_
↪singlepage_blob.csv Sample 0x3000 --keygen
```

注解: 创建的加密密钥格式为 `<outdir>/keys/keys-<prefix>-<fileid>.bin`。

注解： 加密密钥存储于新建文件的 keys/ 目录下，与 NVS 密钥分区结构兼容。更多信息请参考*NVS 密钥分区*。

- 提供加密密钥用作二进制输入文件来进行加密:

```
python mfg_gen.py generate samples/sample_config.csv samples/sample_values_
↪singlepage_blob.csv Sample 0x3000 --inputkey keys/sample_keys.bin
```

仅生成加密密钥

使用方法:

```
python mfg_gen.py generate-key [-h] [--keyfile KEYFILE] [--outdir OUTDIR]
```

可选参数：

参数	描述
-h, –help	显示帮助信息并退出
–keyfile KEYFILE	加密密钥文件的输出路径
–outdir OUTDIR	输出目录，用于存储创建的文件（默认当前目录）

运行以下命令仅生成加密密钥:

```
python mfg_gen.py generate-key
```

注解： 创建的加密密钥格式为 <outdir>/keys/keys-<timestamp>.bin。时间戳格式为：%m-%d_%H-%M。

注解： 如需自定义目标文件名，请使用–keyfile 参数。

生成的加密密钥二进制文件还可以用于为每个设备的工厂映像加密。

fileid 参数的默认值为 1、2、3…，与主 CSV 文件中的行一一对应，内含设备配置值。

运行量产程序时，将在指定的 outdir 目录下创建以下文件夹：

- bin/ 存储生成的二进制文件
- csv/ 存储生成的中间 CSV 文件
- keys/ 存储加密密钥（创建工厂加密映像时会用到）

此部分 API 代码示例详见 ESP-IDF 项下 storage 目录。

2.6 System API

2.6.1 App Image Format

An application image consists of the following structures:

1. The *esp_image_header_t* structure describes the mode of SPI flash and the count of memory segments.
2. The *esp_image_segment_header_t* structure describes each segment, its length, and its location in ESP32-S2' s memory, followed by the data with a length of data_len. The data offset for each segment in the image is calculated in the following way:

- offset for 0 Segment = sizeof(*esp_image_header_t*) + sizeof(*esp_image_segment_header_t*).
- offset for 1 Segment = offset for 0 Segment + length of 0 Segment + sizeof(*esp_image_segment_header_t*).

- offset for 2 Segment = offset for 1 Segment + length of 1 Segment + sizeof(*esp_image_segment_header_t*).
- …

The count of each segment is defined in the `segment_count` field that is stored in *esp_image_header_t*. The count cannot be more than `ESP_IMAGE_MAX_SEGMENTS`.

To get the list of your image segments, please run the following command:

```
esptool.py --chip esp32s2 image_info build/app.bin
```

```
esptool.py v2.3.1
Image version: 1
Entry point: 40080ea4
13 segments
Segment 1: len 0x13ce0 load 0x3f400020 file_offs 0x00000018 SOC_DROM
Segment 2: len 0x00000 load 0x3ff80000 file_offs 0x00013d00 SOC_RTC_DRAM
Segment 3: len 0x00000 load 0x3ff80000 file_offs 0x00013d08 SOC_RTC_DRAM
Segment 4: len 0x028e0 load 0x3ffb0000 file_offs 0x00013d10 DRAM
Segment 5: len 0x00000 load 0x3ffb28e0 file_offs 0x000165f8 DRAM
Segment 6: len 0x00400 load 0x40080000 file_offs 0x00016600 SOC_IRAM
Segment 7: len 0x09600 load 0x40080400 file_offs 0x00016a08 SOC_IRAM
Segment 8: len 0x62e4c load 0x400d0018 file_offs 0x00020010 SOC_IROM
Segment 9: len 0x06cec load 0x40089a00 file_offs 0x00082e64 SOC_IROM
Segment 10: len 0x00000 load 0x400c0000 file_offs 0x00089b58 SOC_RTC_IRAM
Segment 11: len 0x00004 load 0x50000000 file_offs 0x00089b60 SOC_RTC_DATA
Segment 12: len 0x00000 load 0x50000004 file_offs 0x00089b6c SOC_RTC_DATA
Segment 13: len 0x00000 load 0x50000004 file_offs 0x00089b74 SOC_RTC_DATA
Checksum: e8 (valid)Validation Hash:␣
↪407089ca0eae2bbf83b4120979d3354b1c938a49cb7a0c997f240474ef2ec76b (valid)
```

You can also see the information on segments in the IDF logs while your application is booting:

```
I (443) esp_image: segment 0: paddr=0x00020020 vaddr=0x3f400020 size=0x13ce0 (␣
↪81120) map
I (489) esp_image: segment 1: paddr=0x00033d08 vaddr=0x3ff80000 size=0x00000 ( 0)␣
↪load
I (530) esp_image: segment 2: paddr=0x00033d10 vaddr=0x3ff80000 size=0x00000 ( 0)␣
↪load
I (571) esp_image: segment 3: paddr=0x00033d18 vaddr=0x3ffb0000 size=0x028e0 (␣
↪10464) load
I (612) esp_image: segment 4: paddr=0x00036600 vaddr=0x3ffb28e0 size=0x00000 ( 0)␣
↪load
I (654) esp_image: segment 5: paddr=0x00036608 vaddr=0x40080000 size=0x00400 (␣
↪1024) load
I (695) esp_image: segment 6: paddr=0x00036a10 vaddr=0x40080400 size=0x09600 (␣
↪38400) load
I (737) esp_image: segment 7: paddr=0x00040018 vaddr=0x400d0018 size=0x62e4c␣
↪(405068) map
I (847) esp_image: segment 8: paddr=0x000a2e6c vaddr=0x40089a00 size=0x06cec (␣
↪27884) load
I (888) esp_image: segment 9: paddr=0x000a9b60 vaddr=0x400c0000 size=0x00000 ( 0)␣
↪load
I (929) esp_image: segment 10: paddr=0x000a9b68 vaddr=0x50000000 size=0x00004 ( 4)␣
↪load
I (971) esp_image: segment 11: paddr=0x000a9b74 vaddr=0x50000004 size=0x00000 ( 0)␣
↪load
I (1012) esp_image: segment 12: paddr=0x000a9b7c vaddr=0x50000004 size=0x00000 (␣
↪0) load
```

For more details on the type of memory segments and their address ranges, see the ESP32-S2 Technical Reference Manual, Section 1.3.2 *Embedded Memory*.

3. The image has a single checksum byte after the last segment. This byte is written on a sixteen byte padded

boundary, so the application image might need padding.

4. If the `hash_appended` field from *esp_image_header_t* is set then a SHA256 checksum will be appended. The value of SHA256 is calculated on the range from first byte and up to this field. The length of this field is 32 bytes.
5. If the options *CONFIG_SECURE_SIGNED_APPS_SCHEME* is set to ECDSA then the application image will have additional 68 bytes for an ECDSA signature, which includes:

- version word (4 bytes),
- signature data (64 bytes).

Application Description

The `DROM` segment starts with the *esp_app_desc_t* structure which carries specific fields describing the application:

- `secure_version` - see *Anti-rollback*.
- `version` - see *App version*. *
- `project_name` is filled from `PROJECT_NAME`. *
- `time` and `date` - compile time and date.
- `idf_ver` - version of ESP-IDF. *
- `app_elf_sha256` - contains sha256 for the elf application file.

* - The maximum length is 32 characters, including null-termination character. For example, if the length of `PROJECT_NAME` exceeds 32 characters, the excess characters will be disregarded.

This structure is useful for identification of images uploaded OTA because it has a fixed offset = sizeof(*esp_image_header_t*) + sizeof(*esp_image_segment_header_t*). As soon as a device receives the first fragment containing this structure, it has all the information to determine whether the update should be continued or not.

Adding a Custom Structure to an Application

Customer also has the opportunity to have similar structure with a fixed offset relative to the beginning of the image. The following pattern can be used to add a custom structure to your image:

```
const __attribute__((section(".rodata_custom_desc"))) esp_custom_app_desc_t custom_
↪app_desc = { ... }
```

Offset for custom structure is sizeof(*esp_image_header_t*) + sizeof(*esp_image_segment_header_t*) + sizeof(*esp_app_desc_t*).

To guarantee that the custom structure is located in the image even if it is not used, you need to add:

- For Make: add `COMPONENT_ADD_LDFLAGS += -u custom_app_desc` into `component.mk`
- For Cmake: add `target_link_libraries(${COMPONENT_TARGET} "-u custom_app_desc")` into `CMakeLists.txt`

API Reference

Header File

- bootloader_support/include/esp_app_format.h

Structures

`struct esp_image_header_t`

Main header of binary image.

Public Members

uint8_t **magic**
Magic word ESP_IMAGE_HEADER_MAGIC

uint8_t **segment_count**
Count of memory segments

uint8_t **spi_mode**
flash read mode (esp_image_spi_mode_t as uint8_t)

uint8_t **spi_speed** : 4
flash frequency (esp_image_spi_freq_t as uint8_t)

uint8_t **spi_size** : 4
flash chip size (esp_image_flash_size_t as uint8_t)

uint32_t **entry_addr**
Entry address

uint8_t **wp_pin**
WP pin when SPI pins set via efuse (read by ROM bootloader, the IDF bootloader uses software to configure the WP pin and sets this field to 0xEE=disabled)

uint8_t **spi_pin_drv**[3]
Drive settings for the SPI flash pins (read by ROM bootloader)

esp_chip_id_t **chip_id**
Chip identification number

uint8_t **min_chip_rev**
Minimum chip revision supported by image

uint8_t **reserved**[8]
Reserved bytes in additional header space, currently unused

uint8_t **hash_appended**
If 1, a SHA256 digest "simple hash" (of the entire image) is appended after the checksum. Included in image length. This digest is separate to secure boot and only used for detecting corruption. For secure boot signed images, the signature is appended after this (and the simple hash is included in the signed data).

struct esp_image_segment_header_t
Header of binary image segment.

Public Members

uint32_t **load_addr**
Address of segment

uint32_t **data_len**
Length of data

struct esp_app_desc_t
Description about application.

Public Members

uint32_t **magic_word**
Magic word ESP_APP_DESC_MAGIC_WORD

uint32_t **secure_version**
Secure version

uint32_t **reserv1**[2]
reserv1

char **version**[32]
Application version

char **project_name**[32]
Project name

char **time**[16]
Compile time

char **date**[16]
Compile date

char **idf_ver**[32]
Version IDF

uint8_t **app_elf_sha256**[32]
sha256 of elf file

uint32_t **reserv2**[20]
reserv2

Macros

ESP_IMAGE_HEADER_MAGIC
The magic word for the *esp_image_header_t* structure.

ESP_IMAGE_MAX_SEGMENTS
Max count of segments in the image.

ESP_APP_DESC_MAGIC_WORD
The magic word for the esp_app_desc structure that is in DROM.

Enumerations

enum esp_chip_id_t
ESP chip ID.

Values:

ESP_CHIP_ID_ESP32 = 0x0000
chip ID: ESP32

ESP_CHIP_ID_ESP32S2 = 0x0002
chip ID: ESP32S2

ESP_CHIP_ID_INVALID = 0xFFFF
Invalid chip ID (we defined it to make sure the esp_chip_id_t is 2 bytes size)

enum esp_image_spi_mode_t
SPI flash mode, used in *esp_image_header_t*.

Values:

ESP_IMAGE_SPI_MODE_QIO
SPI mode QIO

ESP_IMAGE_SPI_MODE_QOUT
SPI mode QOUT

ESP_IMAGE_SPI_MODE_DIO
SPI mode DIO

ESP_IMAGE_SPI_MODE_DOUT
SPI mode DOUT

ESP_IMAGE_SPI_MODE_FAST_READ
SPI mode FAST_READ

ESP_IMAGE_SPI_MODE_SLOW_READ
SPI mode SLOW_READ

enum esp_image_spi_freq_t
SPI flash clock frequency.

Values:

ESP_IMAGE_SPI_SPEED_40M
SPI clock frequency 40 MHz

ESP_IMAGE_SPI_SPEED_26M
SPI clock frequency 26 MHz

ESP_IMAGE_SPI_SPEED_20M
SPI clock frequency 20 MHz

ESP_IMAGE_SPI_SPEED_80M = 0xF
SPI clock frequency 80 MHz

enum esp_image_flash_size_t
Supported SPI flash sizes.

Values:

ESP_IMAGE_FLASH_SIZE_1MB = 0
SPI flash size 1 MB

ESP_IMAGE_FLASH_SIZE_2MB
SPI flash size 2 MB

ESP_IMAGE_FLASH_SIZE_4MB
SPI flash size 4 MB

ESP_IMAGE_FLASH_SIZE_8MB
SPI flash size 8 MB

ESP_IMAGE_FLASH_SIZE_16MB
SPI flash size 16 MB

ESP_IMAGE_FLASH_SIZE_MAX
SPI flash size MAX

2.6.2 Application Level Tracing

Overview

IDF provides useful feature for program behaviour analysis: application level tracing. It is implemented in the corresponding library and can be enabled via menuconfig. This feature allows to transfer arbitrary data between host and ESP32-S2 via JTAG interface with small overhead on program execution. Developers can use this library to send application specific state of execution to the host and receive commands or other type of info in the opposite direction at runtime. The main use cases of this library are:

1. Collecting application specific data, see 特定应用程序的跟踪
2. Lightweight logging to the host, see 记录日志到主机
3. System behaviour analysis, see 基于 *SEGGER SystemView* 的系统行为分析

API Reference

Header File

- app_trace/include/esp_app_trace.h

Functions

esp_err_t **esp_apptrace_init**(void)
Initializes application tracing module.

Note Should be called before any esp_apptrace_xxx call.
Return ESP_OK on success, otherwise see esp_err_t

void **esp_apptrace_down_buffer_config**(uint8_t **buf*, uint32_t *size*)
Configures down buffer.

Note Needs to be called before initiating any data transfer using esp_apptrace_buffer_get and esp_apptrace_write. This function does not protect internal data by lock.
Parameters
- `buf`: Address of buffer to use for down channel (host to target) data.
- `size`: Size of the buffer.

uint8_t ***esp_apptrace_buffer_get**(*esp_apptrace_dest_t* *dest*, uint32_t *size*, uint32_t *tmo*)
Allocates buffer for trace data. After data in buffer are ready to be sent off esp_apptrace_buffer_put must be called to indicate it.

Return non-NULL on success, otherwise NULL.
Parameters
- `dest`: Indicates HW interface to send data.
- `size`: Size of data to write to trace buffer.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

esp_err_t **esp_apptrace_buffer_put**(*esp_apptrace_dest_t* *dest*, uint8_t **ptr*, uint32_t *tmo*)
Indicates that the data in buffer are ready to be sent off. This function is a counterpart of and must be preceeded by esp_apptrace_buffer_get.

Return ESP_OK on success, otherwise see esp_err_t
Parameters
- `dest`: Indicates HW interface to send data. Should be identical to the same parameter in call to esp_apptrace_buffer_get.
- `ptr`: Address of trace buffer to release. Should be the value returned by call to esp_apptrace_buffer_get.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

esp_err_t **esp_apptrace_write**(*esp_apptrace_dest_t* *dest*, **const** void **data*, uint32_t *size*, uint32_t *tmo*)
Writes data to trace buffer.

Return ESP_OK on success, otherwise see esp_err_t
Parameters
- `dest`: Indicates HW interface to send data.
- `data`: Address of data to write to trace buffer.
- `size`: Size of data to write to trace buffer.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

int **esp_apptrace_vprintf_to**(*esp_apptrace_dest_t* *dest*, uint32_t *tmo*, **const** char **fmt*, va_list *ap*)
vprintf-like function to sent log messages to host via specified HW interface.

Return Number of bytes written.
Parameters
- `dest`: Indicates HW interface to send data.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.
- `fmt`: Address of format string.
- `ap`: List of arguments.

int **esp_apptrace_vprintf**(**const** char **fmt*, va_list *ap*)
vprintf-like function to sent log messages to host.

Return Number of bytes written.
Parameters
- `fmt`: Address of format string.

- `ap`: List of arguments.

esp_err_t **esp_apptrace_flush** (*esp_apptrace_dest_t dest*, uint32_t *tmo*)
Flushes remaining data in trace buffer to host.

Return ESP_OK on success, otherwise see esp_err_t
Parameters
- `dest`: Indicates HW interface to flush data on.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

esp_err_t **esp_apptrace_flush_nolock** (*esp_apptrace_dest_t dest*, uint32_t *min_sz*, uint32_t *tmo*)
Flushes remaining data in trace buffer to host without locking internal data. This is special version of esp_apptrace_flush which should be called from panic handler.

Return ESP_OK on success, otherwise see esp_err_t
Parameters
- `dest`: Indicates HW interface to flush data on.
- `min_sz`: Threshold for flushing data. If current filling level is above this value, data will be flushed. TRAX destinations only.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

esp_err_t **esp_apptrace_read** (*esp_apptrace_dest_t dest*, void **data*, uint32_t **size*, uint32_t *tmo*)
Reads host data from trace buffer.

Return ESP_OK on success, otherwise see esp_err_t
Parameters
- `dest`: Indicates HW interface to read the data on.
- `data`: Address of buffer to put data from trace buffer.
- `size`: Pointer to store size of read data. Before call to this function pointed memory must hold requested size of data
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

uint8_t ***esp_apptrace_down_buffer_get** (*esp_apptrace_dest_t dest*, uint32_t **size*, uint32_t *tmo*)
Retrieves incoming data buffer if any. After data in buffer are processed esp_apptrace_down_buffer_put must be called to indicate it.

Return non-NULL on success, otherwise NULL.
Parameters
- `dest`: Indicates HW interface to receive data.
- `size`: Address to store size of available data in down buffer. Must be initialized with requested value.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

esp_err_t **esp_apptrace_down_buffer_put** (*esp_apptrace_dest_t dest*, uint8_t **ptr*, uint32_t *tmo*)
Indicates that the data in down buffer are processed. This function is a counterpart of and must be preceeded by esp_apptrace_down_buffer_get.

Return ESP_OK on success, otherwise see esp_err_t
Parameters
- `dest`: Indicates HW interface to receive data. Should be identical to the same parameter in call to esp_apptrace_down_buffer_get.
- `ptr`: Address of trace buffer to release. Should be the value returned by call to esp_apptrace_down_buffer_get.
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinitely.

bool **esp_apptrace_host_is_connected** (*esp_apptrace_dest_t dest*)
Checks whether host is connected.

Return true if host is connected, otherwise false
Parameters
- `dest`: Indicates HW interface to use.

void ***esp_apptrace_fopen** (*esp_apptrace_dest_t dest*, **const** char **path*, **const** char **mode*)
Opens file on host. This function has the same semantic as ‘fopen’ except for the first argument.

Return non zero file handle on success, otherwise 0

Parameters

- `dest`: Indicates HW interface to use.
- `path`: Path to file.
- `mode`: Mode string. See fopen for details.

int **esp_apptrace_fclose**(*esp_apptrace_dest_t dest*, void **stream*)

Closes file on host. This function has the same semantic as 'fclose' except for the first argument.

Return Zero on success, otherwise non-zero. See fclose for details.

Parameters

- `dest`: Indicates HW interface to use.
- `stream`: File handle returned by esp_apptrace_fopen.

size_t **esp_apptrace_fwrite**(*esp_apptrace_dest_t dest*, **const** void **ptr*, size_t *size*, size_t *nmemb*, void **stream*)

Writes to file on host. This function has the same semantic as 'fwrite' except for the first argument.

Return Number of written items. See fwrite for details.

Parameters

- `dest`: Indicates HW interface to use.
- `ptr`: Address of data to write.
- `size`: Size of an item.
- `nmemb`: Number of items to write.
- `stream`: File handle returned by esp_apptrace_fopen.

size_t **esp_apptrace_fread**(*esp_apptrace_dest_t dest*, void **ptr*, size_t *size*, size_t *nmemb*, void **stream*)

Read file on host. This function has the same semantic as 'fread' except for the first argument.

Return Number of read items. See fread for details.

Parameters

- `dest`: Indicates HW interface to use.
- `ptr`: Address to store read data.
- `size`: Size of an item.
- `nmemb`: Number of items to read.
- `stream`: File handle returned by esp_apptrace_fopen.

int **esp_apptrace_fseek**(*esp_apptrace_dest_t dest*, void **stream*, long *offset*, int *whence*)

Set position indicator in file on host. This function has the same semantic as 'fseek' except for the first argument.

Return Zero on success, otherwise non-zero. See fseek for details.

Parameters

- `dest`: Indicates HW interface to use.
- `stream`: File handle returned by esp_apptrace_fopen.
- `offset`: Offset. See fseek for details.
- `whence`: Position in file. See fseek for details.

int **esp_apptrace_ftell**(*esp_apptrace_dest_t dest*, void **stream*)

Get current position indicator for file on host. This function has the same semantic as 'ftell' except for the first argument.

Return Current position in file. See ftell for details.

Parameters

- `dest`: Indicates HW interface to use.
- `stream`: File handle returned by esp_apptrace_fopen.

int **esp_apptrace_fstop**(*esp_apptrace_dest_t dest*)

Indicates to the host that all file operations are completed. This function should be called after all file operations are finished and indicate to the host that it can perform cleanup operations (close open files etc.).

Return ESP_OK on success, otherwise see esp_err_t

Parameters

- `dest`: Indicates HW interface to use.

void **esp_gcov_dump**(void)
Triggers gcov info dump. This function waits for the host to connect to target before dumping data.

Enumerations

enum esp_apptrace_dest_t
Application trace data destinations bits.

Values:

ESP_APPTRACE_DEST_TRAX = 0x1
JTAG destination.

ESP_APPTRACE_DEST_UART0 = 0x2
UART destination.

Header File

- app_trace/include/esp_sysview_trace.h

Functions

static *esp_err_t* **esp_sysview_flush**(uint32_t *tmo*)
Flushes remaining data in SystemView trace buffer to host.

Return ESP_OK.
Parameters
- `tmo`: Timeout for operation (in us). Use ESP_APPTRACE_TMO_INFINITE to wait indefinetly.

int **esp_sysview_vprintf**(**const** char ***format*, va_list *args*)
vprintf-like function to sent log messages to the host.

Return Number of bytes written.
Parameters
- `format`: Address of format string.
- `args`: List of arguments.

esp_err_t **esp_sysview_heap_trace_start**(uint32_t *tmo*)
Starts SystemView heap tracing.

Return ESP_OK on success, ESP_ERR_TIMEOUT if operation has been timed out.
Parameters
- `tmo`: Timeout (in us) to wait for the host to be connected. Use -1 to wait forever.

esp_err_t **esp_sysview_heap_trace_stop**(void)
Stops SystemView heap tracing.

Return ESP_OK.

void **esp_sysview_heap_trace_alloc**(void **addr*, uint32_t *size*, **const** void **callers*)
Sends heap allocation event to the host.

Parameters
- `addr`: Address of allocated block.
- `size`: Size of allocated block.
- `callers`: Pointer to array with callstack addresses. Array size must be CONFIG_HEAP_TRACING_STACK_DEPTH.

void **esp_sysview_heap_trace_free**(void **addr*, **const** void **callers*)
Sends heap de-allocation event to the host.

Parameters
- `addr`: Address of de-allocated block.
- `callers`: Pointer to array with callstack addresses. Array size must be CONFIG_HEAP_TRACING_STACK_DEPTH.

2.6.3 控制台终端

ESP-IDF 提供了 `console` 组件，它包含了开发基于串口的交互式控制终端所需要的所有模块，主要支持以下功能：

- 行编辑，由 linenoise 库具体实现，它支持处理退格键和方向键，支持回看命令的历史记录，支持命令的自动补全和参数提示。
- 将命令行拆分为参数列表。
- 参数解析，由 argtable3 库具体实现，它支持解析 GNU 样式的命令行参数。
- 用于注册和调度命令的函数。
- 帮助创建 REPL (Read-Evaluate-Print-Loop) 环境的函数。

注解： 这些功能模块可以一起使用也可以独立使用，例如仅使用行编辑和命令注册的功能，然后使用 `getopt` 函数或者自定义的函数来实现参数解析，而不是直接使用 argtable3 库。同样地，还可以使用更简单的命令输入方法（比如 `fgets` 函数）和其他用于命令分割和参数解析的方法。

行编辑

行编辑功能允许用户通过按键输入来编辑命令，使用退格键删除符号，使用左/右键在命令中移动光标，使用上/下键导航到之前输入的命令，使用制表键（“Tab”）来自动补全命令。

注解： 此功能依赖于终端应用程序对 ANSI 转移符的支持，显示原始 UART 数据的串口监视器不能与行编辑库一同使用。如果运行 system/console 示例程序的时候观察到的输出结果是 `[6n` 或者类似的转义字符而不是命令行提示符 `esp>` ``时，就表明当前的串口监视器不支持 ANSI 转移字符。已知可用的串口监视程序有 GNU screen，minicom 和 idf_monitor.py（可以通过在项目目录下执行 ``idf.py monitor` 来调用）。

前往这里可以查看 linenoise 库提供的所有函数的描述。

配置 Linenoise 库不需要显式地初始化，但是在调用行编辑函数之前，可能需要对某些配置的默认值稍作修改。

`linenoiseClearScreen()`

使用转移字符清除终端屏幕，并将光标定位在左上角。

`linenoiseSetMultiLine()`

在单行和多行编辑模式之间进行切换。单行模式下，如果命令的长度超过终端的宽度，会在行内滚动命令文本以显示文本的结尾，在这种情况下，文本的开头部分会被隐藏。单行模式在每次按下按键时发送给屏幕刷新的数据比较少，与多行模式相比更不容易发生故障。另一方面，在单行模式下编辑命令和复制命令将变得更加困难。默认情况下开启的是单行模式。

主循环 `linenoise()`

在大多数情况下，控制台应用程序都会具有相同的工作形式——在某个循环中不断读取输入的内容，然后解析再处理。`linenoise()` 是专门用来获取用户按键输入的函数，当回车键被按下后会便返回完整的一行内容。因此可以用它来完成前面循环中的“读取”任务。

`linenoiseFree()`

必须调用此函数才能释放从 `linenoise()` 函数获取的命令行缓冲。

提示和补全 `linenoiseSetCompletionCallback()`

当用户按下制表键时，linenoise 会调用 **补全回调函数**，该回调函数会检查当前已经输入的内容，然后调用 `linenoiseAddCompletion()` 函数来提供所有可能的补全后的命令列表。启用补全功能，需要事先调用 `linenoiseSetCompletionCallback()` 函数来注册补全回调函数。

`console` 组件提供了一个现成的函数来为注册的命令提供补全功能*esp_console_get_completion()*（见下文）。

`linenoiseAddCompletion()`

补全回调函数会通过调用此函数来通知 linenoise 库当前键入命令所有可能的补全结果。

`linenoiseSetHintsCallback()`

每当用户的输入改变时，linenoise 就会调用此回调函数，检查到目前为止输入的命令行内容，然后提供带有提示信息的字符串（例如命令参数列表），然后会在同一行上用不同的颜色显示出该文本。

`linenoiseSetFreeHintsCallback()`

如果 **提示回调函数**返回的提示字符串是动态分配的或者需要以其它方式回收，就需要使用 `linenoiseSetFreeHintsCallback()` 注册具体的清理函数。

历史记录 `linenoiseHistorySetMaxLen()`

该函数设置要保留在内存中的最近输入的命令的数量。用户通过使用向上/向下箭头来导航历史记录。

`linenoiseHistoryAdd()`

Linenoise 不会自动向历史记录中添加命令，应用程序需要调用此函数来将命令字符串添加到历史记录中。

`linenoiseHistorySave()`

该函数将命令的历史记录从 RAM 中保存为文本文件，例如保存到 SD 卡或者 Flash 的文件系统中。

`linenoiseHistoryLoad()`

与 `linenoiseHistorySave` 相对应，从文件中加载历史记录。

`linenoiseHistoryFree()`

释放用于存储命令历史记录的内存。当使用完 linenoise 库后需要调用此函数。

将命令行拆分成参数列表

`console` 组件提供*esp_console_split_argv()* 函数来将命令行字符串拆分为参数列表。该函数会返回参数的数量（`argc`）和一个指针数组，该指针数组可以作为 `argv` 参数传递给任何接受 `argc`，`argv` 格式参数的函数。

根据以下规则来将命令行拆分成参数列表：

- 参数由空格分隔
- 如果参数本身需要使用空格，可以使用 \（反斜杠）对它们进行转义
- 其它能被识别的转义字符有 \\（显示反斜杠本身）和 \"（显示双引号）
- 可以使用双引号来引用参数，引号只可能出现在参数的开头和结尾。参数中的引号必须如上所述进行转移。参数周围的引号会被*esp_console_split_argv()* 函数删除

示例：

- `abc def 1 20 .3` → [`abc`, `def`, `1`, `20`, `.3`]
- `abc "123 456" def` → [`abc`, `123 456`, `def`]
- `` `a\ b\\c\" `` → [`a b\c"`]

参数解析

对于参数解析，console 组件使用 argtable3 库。有关 argtable3 的介绍请查看 教程 或者 Github 仓库中的 示例代码 。

命令的注册与调度

console 组件包含了一些工具函数，用来注册命令，将用户输入的命令和已经注册的命令进行匹配，使用命令行输入的参数调用命令。

应用程序首先调用*esp_console_init()* 来初始化命令注册模块，然后调用*esp_console_cmd_register()* 函数注册命令处理程序。

对于每个命令，应用程序需要提供以下信息（需要以 esp_console_cmd_t 结构体的形式给出）：

- 命令名字（不含空格的字符串）
- 帮助文档，解释该命令的用途
- 可选的提示文本，列出命令的参数。如果应用程序使用 Argtable3 库来解析参数，则可以通过提供指向 argtable 参数定义结构体的指针来自动生成提示文本
- 命令处理函数

命令注册模块还提供了其它函数：

esp_console_run()

> 该函数接受命令行字符串，使用*esp_console_split_argv()* 函数将其拆分为 argc/argv 形式的参数列表，在已经注册的组件列表中查找命令，如果找到，则执行其对应的处理程序。

esp_console_register_help_command()

> 将 help 命令添加到已注册命令列表中，此命令将会以列表的方式打印所有注册的命令及其参数和帮助文本。

esp_console_get_completion()

> 与 linenoise 库中的 linenoiseSetCompletionCallback() 一同使用的回调函数，根据已经注册的命令列表为 linenoise 提供补全功能。

esp_console_get_hint()

> 与 linenoise 库中 linenoiseSetHintsCallback() 一同使用的回调函数，为 linenoise 提供已经注册的命令的参数提示功能。

初始化 REPL 环境

console 组建还提供了一些 API 来帮助创建一个基本的 REPL 环境。

在一个典型的 console 应用中，你只需要调用*esp_console_new_repl_uart()*，它会为你初始化好构建在 UART 基础上的 REPL 环境，其中包括安装 UART 驱动，基本的 console 配置，创建一个新的线程来执行 REPL 任务，注册一些基本的命令（比如 *help* 命令）。

完了之后你可以使用*esp_console_cmd_register()* 来注册其它命令。REPL 环境在初始化后需要再调用*esp_console_start_repl()* 函数才能开始运行。

应用程序示例

system/console 目录下提供了 console 组件的示例应用程序，展示了具体的使用方法。该示例介绍了如何初始化 UART 和 VFS 的功能，设置 linenoise 库，从 UART 中读取命令并加以处理，然后将历史命令存储到 Flash 中。更多信息，请参阅示例代码目录中的 README.md 文件。

此外，ESP-IDF 还提供了众多基于 *console* 组件的示例程序，它们可以辅助应用程序的开发。例如，peripherals/i2c/i2c_tools，wifi/iperf 等等。

API 参考

Header File

- console/esp_console.h

Functions

esp_err_t **esp_console_init**(**const** *esp_console_config_t* **config*)

initialize console module

Note Call this once before using other console module features

Return

- ESP_OK on success
- ESP_ERR_NO_MEM if out of memory
- ESP_ERR_INVALID_STATE if already initialized
- ESP_ERR_INVALID_ARG if the configuration is invalid

Parameters

- `config`: console configuration

esp_err_t **esp_console_deinit**(void)

de-initialize console module

Note Call this once when done using console module functions

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if not initialized yet

esp_err_t **esp_console_cmd_register**(**const** *esp_console_cmd_t* **cmd*)

Register console command.

Return

- ESP_OK on success
- ESP_ERR_NO_MEM if out of memory
- ESP_ERR_INVALID_ARG if command description includes invalid arguments

Parameters

- `cmd`: pointer to the command description; can point to a temporary value

esp_err_t **esp_console_run**(**const** char **cmdline*, int **cmd_ret*)

Run command line.

Return

- ESP_OK, if command was run
- ESP_ERR_INVALID_ARG, if the command line is empty, or only contained whitespace
- ESP_ERR_NOT_FOUND, if command with given name wasn't registered
- ESP_ERR_INVALID_STATE, if esp_console_init wasn't called

Parameters

- `cmdline`: command line (command name followed by a number of arguments)
- `[out] cmd_ret`: return code from the command (set if command was run)

size_t **esp_console_split_argv**(char **line*, char ***argv*, size_t *argv_size*)

Split command line into arguments in place.

```
* - This function finds whitespace-separated arguments in the given input line.
*
*     'abc def 1 20 .3' -> [ 'abc', 'def', '1', '20', '.3' ]
*
* - Argument which include spaces may be surrounded with quotes. In this case
*   spaces are preserved and quotes are stripped.
*
*     'abc "123 456" def' -> [ 'abc', '123 456', 'def' ]
*
* - Escape sequences may be used to produce backslash, double quote, and space:
*
```

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```
*      'a\ b\\c\"' -> [ 'a b\c"' ]
*
```

Note Pointers to at most argv_size - 1 arguments are returned in argv array. The pointer after the last one (i.e. argv[argc]) is set to NULL.

Return number of arguments found (argc)

Parameters

- `line`: pointer to buffer to parse; it is modified in place
- `argv`: array where the pointers to arguments are written
- `argv_size`: number of elements in argv_array (max. number of arguments)

void **esp_console_get_completion**(**const** char *_buf_, _linenoiseCompletions_ *_lc_)

Callback which provides command completion for linenoise library.

When using linenoise for line editing, command completion support can be enabled like this:

linenoiseSetCompletionCallback(&esp_console_get_completion);

Parameters

- `buf`: the string typed by the user
- `lc`: linenoiseCompletions to be filled in

const char ***esp_console_get_hint**(**const** char *_buf_, int *_color_, int *_bold_)

Callback which provides command hints for linenoise library.

When using linenoise for line editing, hints support can be enabled as follows:

linenoiseSetHintsCallback((linenoiseHintsCallback*) &esp_console_get_hint);

The extra cast is needed because linenoiseHintsCallback is defined as returning a char* instead of const char*.

Return string containing the hint text. This string is persistent and should not be freed (i.e. linenoiseSetFreeHintsCallback should not be used).

Parameters

- `buf`: line typed by the user
- `[out] color`: ANSI color code to be used when displaying the hint
- `[out] bold`: set to 1 if hint has to be displayed in bold

_esp_err_t_ **esp_console_register_help_command**(void)

Register a ‘help’ command.

Default ‘help’ command prints the list of registered commands along with hints and help strings.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE, if esp_console_init wasn’t called

_esp_err_t_ **esp_console_new_repl_uart**(**const** _esp_console_dev_uart_config_t_ *_dev_config_, **const** _esp_console_repl_config_t_ *_repl_config_, _esp_console_repl_t_ **_ret_repl_)

Establish a console REPL environment over UART driver.

Note This is a all-in-one function to establish the environment needed for REPL, includes:

- Install the UART driver on the console UART (8n1, 115200, REF_TICK clock source)
- Configures the stdin/stdout to go through the UART driver
- Initializes linenoise
- Spawn new thread to run REPL in the background

Attention This function is meant to be used in the examples to make the code more compact. Applications which use console functionality should be based on the underlying linenoise and esp_console functions.

Return

- ESP_OK on success
- ESP_FAIL Parameter error

Parameters

- `[in] dev_config`: UART device configuration

- `[in] repl_config`: REPL configuration
- `[out] ret_repl`: return REPL handle after initialization succeed, return NULL otherwise

esp_err_t **esp_console_start_repl** (*esp_console_repl_t* **repl*)

Start REPL environment.

Note Once the REPL got started, it won't be stopped until user call repl->del(repl) to destory the REPL environment.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE, if repl has started already

Parameters

- `[in] repl`: REPL handle returned from esp_console_new_repl_xxx

Structures

struct esp_console_config_t

Parameters for console initialization.

Public Members

size_t **max_cmdline_length**

length of command line buffer, in bytes

size_t **max_cmdline_args**

maximum number of command line arguments to parse

int **hint_color**

ASCII color code of hint text.

int **hint_bold**

Set to 1 to print hint text in bold.

struct esp_console_repl_config_t

Parameters for console REPL (Read Eval Print Loop)

Public Members

uint32_t **max_history_len**

maximum length for the history

const char ***history_save_path**

file path used to save history commands, set to NULL won't save to file system

uint32_t **task_stack_size**

repl task stack size

uint32_t **task_priority**

repl task priority

const char ***prompt**

prompt (NULL represents default: "esp> ")

struct esp_console_dev_uart_config_t

Parameters for console device: UART.

Public Members

int **channel**

UART channel number (count from zero)

int **baud_rate**

Comunication baud rate.

int **tx_gpio_num**
GPIO number for TX path, -1 means using default one.

int **rx_gpio_num**
GPIO number for RX path, -1 means using default one.

struct esp_console_cmd_t
Console command description.

Public Members

const char ***command**
Command name. Must not be NULL, must not contain spaces. The pointer must be valid until the call to esp_console_deinit.

const char ***help**
Help text for the command, shown by help command. If set, the pointer must be valid until the call to esp_console_deinit. If not set, the command will not be listed in 'help' output.

const char ***hint**
Hint text, usually lists possible arguments. If set to NULL, and 'argtable' field is non-NULL, hint will be generated automatically

esp_console_cmd_func_t **func**
Pointer to a function which implements the command.

void ***argtable**
Array or structure of pointers to arg_xxx structures, may be NULL. Used to generate hint text if 'hint' is set to NULL. Array/structure which this field points to must end with an arg_end. Only used for the duration of esp_console_cmd_register call.

struct esp_console_repl_s
Console REPL base structure.

Public Members

esp_err_t (***del**)(*esp_console_repl_t* *repl)
Delete console REPL environment.

Return
- ESP_OK on success
- ESP_FAIL on errors

Parameters
- `[in] repl`: REPL handle returned from esp_console_new_repl_xxx

Macros

ESP_CONSOLE_CONFIG_DEFAULT()
Default console configuration value.

ESP_CONSOLE_REPL_CONFIG_DEFAULT()
Default console repl configuration value.

ESP_CONSOLE_DEV_UART_CONFIG_DEFAULT()

Type Definitions

typedef struct *linenoiseCompletions* **linenoiseCompletions**

typedef int (***esp_console_cmd_func_t**)(int argc, char **argv)
Console command main function.

Return console command return code, 0 indicates "success"
Parameters

- `argc`: number of arguments
- `argv`: array with argc entries, each pointing to a zero-terminated string argument

typedef struct *esp_console_repl_s* **esp_console_repl_t**
Type defined for console REPL.

2.6.4 eFuse Manager

Introduction

The eFuse Manager library is designed to structure access to eFuse bits and make using these easy. This library operates eFuse bits by a structure name which is assigned in eFuse table. This sections introduces some concepts used by eFuse Manager.

Hardware description

The ESP32-S2 has a number of eFuses which can store system and user parameters. Each eFuse is a one-bit field which can be programmed to 1 after which it cannot be reverted back to 0. Some of system parameters are using these eFuse bits directly by hardware modules and have special place (for example EFUSE_BLK0).

For more details see ESP32-S2 Technical Reference Manual. Some eFuse bits are available for user applications.

ESP32-S2 has 11 eFuse blocks each of the size of 256 bits (not all bits are available):

Each block is divided into 8 32-bits registers.

eFuse Manager component

The component has API functions for reading and writing fields. Access to the fields is carried out through the structures that describe the location of the eFuse bits in the blocks. The component provides the ability to form fields of any length and from any number of individual bits. The description of the fields is made in a CSV file in a table form. To generate from a tabular form (CSV file) in the C-source uses the tool *efuse_table_gen.py*. The tool checks the CSV file for uniqueness of field names and bit intersection, in case of using a *custom* file from the user's project directory, the utility will check with the *common* CSV file.

CSV files:

- common (*esp_efuse_table.csv*) - contains eFuse fields which are used inside the IDF. C-source generation should be done manually when changing this file (run command `idf.py efuse_common_table`). Note that changes in this file can lead to incorrect operation.
- custom - (optional and can be enabled by `CONFIG_EFUSE_CUSTOM_TABLE`) contains eFuse fields that are used by the user in their application. C-source generation should be done manually when changing this file and running `idf.py efuse_custom_table`.

Description CSV file

The CSV file contains a description of the eFuse fields. In the simple case, one field has one line of description. Table header:

```
# field_name,  efuse_block(EFUSE_BLK0..EFUSE_BLK3), bit_start(0..255),  bit_
↪count(1..256),              comment
```

Individual params in CSV file the following meanings:

field_name Name of field. The prefix *ESP_EFUSE_* will be added to the name, and this field name will be available in the code. This name will be used to access the fields. The name must be unique for all fields. If the line has an empty name, then this line is combined with the previous field. This allows you to set an arbitrary order of bits in the field, and expand the field as well (see `MAC_FACTORY` field in the common table).

efuse_block Block number. It determines where the eFuse bits will be placed for this field. Available EFUSE_BLK0..EFUSE_BLK3.

bit_start Start bit number (0..255). The bit_start field can be omitted. In this case, it will be set to bit_start + bit_count from the previous record, if it has the same efuse_block. Otherwise (if efuse_block is different, or this is the first entry), an error will be generated.

bit_count The number of bits to use in this field (1..-). This parameter can not be omitted. This field also may be `MAX_BLK_LEN` in this case, the field length will have the maximum block length, taking into account the coding scheme (applicable for `ESP_EFUSE_SECURE_BOOT_KEY` and `ESP_EFUSE_ENCRYPT_FLASH_KEY` fields). The value `MAX_BLK_LEN` depends on `CONFIG_EFUSE_MAX_BLK_LEN`, will be replaced with "None" - 256, "3/4" - 192, "REPEAT" - 128.

comment This param is using for comment field, it also move to C-header file. The comment field can be omitted.

If a non-sequential bit order is required to describe a field, then the field description in the following lines should be continued without specifying a name, this will indicate that it belongs to one field. For example two fields MAC_FACTORY and MAC_FACTORY_CRC:

```
# Factory MAC address #
#######################
MAC_FACTORY,            EFUSE_BLK0,     72,     8,    Factory MAC addr [0]
,                       EFUSE_BLK0,     64,     8,    Factory MAC addr [1]
,                       EFUSE_BLK0,     56,     8,    Factory MAC addr [2]
,                       EFUSE_BLK0,     48,     8,    Factory MAC addr [3]
,                       EFUSE_BLK0,     40,     8,    Factory MAC addr [4]
,                       EFUSE_BLK0,     32,     8,    Factory MAC addr [5]
MAC_FACTORY_CRC,        EFUSE_BLK0,     80,     8,    CRC8 for factory MAC address
```

This field will available in code as ESP_EFUSE_MAC_FACTORY and ESP_EFUSE_MAC_FACTORY_CRC.

efuse_table_gen.py tool

The tool is designed to generate C-source files from CSV file and validate fields. First of all, the check is carried out on the uniqueness of the names and overlaps of the field bits. If an additional *custom* file is used, it will be checked with the existing *common* file (esp_efuse_table.csv). In case of errors, a message will be displayed and the string that caused the error. C-source files contain structures of type *esp_efuse_desc_t*.

To generate a *common* files, use the following command `idf.py efuse_common_table` or:

```
cd $IDF_PATH/components/efuse/
./efuse_table_gen.py esp32s2/esp_efuse_table.csv
```

After generation in the folder *esp32s2* create:

- *esp_efuse_table.c* file.
- In *include* folder *esp_efuse_table.c* file.

To generate a *custom* files, use the following command `idf.py efuse_custom_table` or:

```
cd $IDF_PATH/components/efuse/
./efuse_table_gen.py esp32s2/esp_efuse_table.csv PROJECT_PATH/main/esp_efuse_
↪custom_table.csv
```

After generation in the folder PROJECT_PATH/main create:

- *esp_efuse_custom_table.c* file.
- In *include* folder *esp_efuse_custom_table.c* file.

To use the generated fields, you need to include two files:

```
#include "esp_efuse.h"
#include "esp_efuse_table.h" or "esp_efuse_custom_table.h"
```

Support coding scheme

eFuse have three coding schemes:

- None (value 0).
- 3/4 (value 1).
- Repeat (value 2).

The coding scheme affects only EFUSE_BLK1, EFUSE_BLK2 and EFUSE_BLK3 blocks. EUSE_BLK0 block always has a coding scheme None. Coding changes the number of bits that can be written into a block, the block length is constant 256, some of these bits are used for encoding and are not used.

When using a coding scheme, the length of the payload that can be written is limited (for more details 20.3.1.3 System Parameter coding_scheme):

- None 256 bits.
- 3/4 192 bits.
- Repeat 128 bits.

You can find out the coding scheme of your chip:

- run a espefuse.py -p COM4 summary command.
- from esptool utility logs (during flashing).
- calling the function in the code *esp_efuse_get_coding_scheme()* for the EFUSE_BLK3 block.

eFuse tables must always comply with the coding scheme in the chip. There is an EFUSE_CODE_SCHEME_SELECTOR option to select the coding type for tables in a Kconfig. When generating source files, if your tables do not follow the coding scheme, an error message will be displayed. Adjust the length or offset fields. If your program was compiled with None encoding and 3/4 is used in the chip, then the ESP_ERR_CODING error may occur when calling the eFuse API (the field is outside the block boundaries). If the field matches the new block boundaries, then the API will work without errors.

Also, 3/4 coding scheme imposes restrictions on writing bits belonging to one coding unit. The whole block with a length of 256 bits is divided into 4 coding units, and in each coding unit there are 6 bytes of useful data and 2 service bytes. These 2 service bytes contain the checksum of the previous 6 data bytes.

It turns out that only one field can be written into one coding unit. Repeated rewriting in one coding unit is prohibited. But if the record was made in advance or through a *esp_efuse_write_block()* function, then reading the fields belonging to one coding unit is possible.

In case 3/4 coding scheme, the writing process is divided into the coding units and we can not use the usual mode of writing some fields. We can prepare all the data for writing and burn it in one time. You can also use this mode for None coding scheme but it is not necessary. It is important for 3/4 coding scheme. To write some fields in one time need to use the batch writing mode. Firstly set this mode through *esp_efuse_batch_write_begin()* function then write some fields as usual use the esp_efuse_write_... functions. At the end to burn they, need to call the *esp_efuse_batch_write_commit()* function. It burns prepared data to the efuse blocks and disable the batch recording mode. The batch writing mode blocks esp_efuse_read_... operations.

After changing the coding scheme, run efuse_common_table and efuse_custom_table commands to check the tables of the new coding scheme.

eFuse API

Access to the fields is via a pointer to the description structure. API functions have some basic operation:

- *esp_efuse_read_field_blob()* - returns an array of read eFuse bits.
- *esp_efuse_read_field_cnt()* - returns the number of bits programmed as "1".
- *esp_efuse_write_field_blob()* - writes an array.
- *esp_efuse_write_field_cnt()* - writes a required count of bits as "1".
- *esp_efuse_get_field_size()* - returns the number of bits by the field name.
- *esp_efuse_read_reg()* - returns value of eFuse register.
- *esp_efuse_write_reg()* - writes value to eFuse register.

- *esp_efuse_get_coding_scheme()* - returns eFuse coding scheme for blocks.
- *esp_efuse_read_block()* - reads key to eFuse block starting at the offset and the required size.
- *esp_efuse_write_block()* - writes key to eFuse block starting at the offset and the required size.
- *esp_efuse_batch_write_begin()* - set the batch mode of writing fields.
- *esp_efuse_batch_write_commit()* - writes all prepared data for batch writing mode and reset the batch writing mode.
- *esp_efuse_batch_write_cancel()* - reset the batch writing mode and prepared data.

For frequently used fields, special functions are made, like this *esp_efuse_get_chip_ver()*, *esp_efuse_get_pkg_ver()*.

How add a new field

1. Find a free bits for field. Show *esp_efuse_table.csv* file or run `idf.py show_efuse_table` or the next command:

```
$ ./efuse_table_gen.py esp32s2/esp_efuse_table.csv --info
eFuse coding scheme: NONE
#       field_name                      efuse_block     bit_start       bit_count
1       WR_DIS_FLASH_CRYPT_CNT          EFUSE_BLK0          2               1
2       WR_DIS_BLK1                     EFUSE_BLK0          7               1
3       WR_DIS_BLK2                     EFUSE_BLK0          8               1
4       WR_DIS_BLK3                     EFUSE_BLK0          9               1
5       RD_DIS_BLK1                     EFUSE_BLK0          16              1
6       RD_DIS_BLK2                     EFUSE_BLK0          17              1
7       RD_DIS_BLK3                     EFUSE_BLK0          18              1
8       FLASH_CRYPT_CNT                 EFUSE_BLK0          20              7
9       MAC_FACTORY                     EFUSE_BLK0          32              8
10      MAC_FACTORY                     EFUSE_BLK0          40              8
11      MAC_FACTORY                     EFUSE_BLK0          48              8
12      MAC_FACTORY                     EFUSE_BLK0          56              8
13      MAC_FACTORY                     EFUSE_BLK0          64              8
14      MAC_FACTORY                     EFUSE_BLK0          72              8
15      MAC_FACTORY_CRC                 EFUSE_BLK0          80              8
16      CHIP_VER_DIS_APP_CPU            EFUSE_BLK0          96              1
17      CHIP_VER_DIS_BT                 EFUSE_BLK0          97              1
18      CHIP_VER_PKG                    EFUSE_BLK0          105             3
19      CHIP_CPU_FREQ_LOW               EFUSE_BLK0          108             1
20      CHIP_CPU_FREQ_RATED             EFUSE_BLK0          109             1
21      CHIP_VER_REV1                   EFUSE_BLK0          111             1
22      ADC_VREF_AND_SDIO_DREF          EFUSE_BLK0          136             6
23      XPD_SDIO_REG                    EFUSE_BLK0          142             1
24      SDIO_TIEH                       EFUSE_BLK0          143             1
25      SDIO_FORCE                      EFUSE_BLK0          144             1
26      ENCRYPT_CONFIG                  EFUSE_BLK0          188             4
27      CONSOLE_DEBUG_DISABLE           EFUSE_BLK0          194             1
28      ABS_DONE_0                      EFUSE_BLK0          196             1
29      DISABLE_JTAG                    EFUSE_BLK0          198             1
30      DISABLE_DL_ENCRYPT              EFUSE_BLK0          199             1
31      DISABLE_DL_DECRYPT              EFUSE_BLK0          200             1
32      DISABLE_DL_CACHE                EFUSE_BLK0          201             1
33      ENCRYPT_FLASH_KEY               EFUSE_BLK1          0               256
34      SECURE_BOOT_KEY                 EFUSE_BLK2          0               256
35      MAC_CUSTOM_CRC                  EFUSE_BLK3          0               8
36      MAC_CUSTOM                      EFUSE_BLK3          8               48
37      ADC1_TP_LOW                     EFUSE_BLK3          96              7
38      ADC1_TP_HIGH                    EFUSE_BLK3          103             9
39      ADC2_TP_LOW                     EFUSE_BLK3          112             7
40      ADC2_TP_HIGH                    EFUSE_BLK3          119             9
41      SECURE_VERSION                  EFUSE_BLK3          128             32
42      MAC_CUSTOM_VER                  EFUSE_BLK3          184             8
```

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```
Used bits in eFuse table:
EFUSE_BLK0
[2 2] [7 9] [16 18] [20 27] [32 87] [96 97] [105 109] [111 111] [136 144] [188␣
↪191] [194 194] [196 196] [198 201]

EFUSE_BLK1
[0 255]

EFUSE_BLK2
[0 255]

EFUSE_BLK3
[0 55] [96 159] [184 191]

Note: Not printed ranges are free for using. (bits in EFUSE_BLK0 are reserved for␣
↪Espressif)

Parsing eFuse CSV input file $IDF_PATH/components/efuse/esp32s2/esp_efuse_table.
↪csv ...
Verifying eFuse table...
```

The number of bits not included in square brackets is free (bits in EFUSE_BLK0 are reserved for Espressif). All fields are checked for overlapping.

2. Fill a line for field: field_name, efuse_block, bit_start, bit_count, comment.
3. Run a `show_efuse_table` command to check eFuse table. To generate source files run `efuse_common_table` or `efuse_custom_table` command.

Debug eFuse & Unit tests

Virtual eFuses The Kconfig option `CONFIG_EFUSE_VIRTUAL` will virtualize eFuse values inside the eFuse Manager, so writes are emulated and no eFuse values are permanently changed. This can be useful for debugging app and unit tests.

espefuse.py esptool includes a useful tool for reading/writing ESP32-S2 eFuse bits - espefuse.py.

```
espefuse.py -p COM4 summary

espefuse.py v2.3.1
Connecting........_
Security fuses:
FLASH_CRYPT_CNT         Flash encryption mode counter                    = 0 R/W␣
↪(0x0)
FLASH_CRYPT_CONFIG      Flash encryption config (key tweak bits)         = 0 R/W␣
↪(0x0)
CONSOLE_DEBUG_DISABLE   Disable ROM BASIC interpreter fallback           = 1 R/W␣
↪(0x1)
ABS_DONE_0              secure boot enabled for bootloader               = 0 R/W␣
↪(0x0)
ABS_DONE_1              secure boot abstract 1 locked                    = 0 R/W␣
↪(0x0)
JTAG_DISABLE            Disable JTAG                                     = 0 R/W␣
↪(0x0)
DISABLE_DL_ENCRYPT      Disable flash encryption in UART bootloader      = 0 R/W␣
↪(0x0)
DISABLE_DL_DECRYPT      Disable flash decryption in UART bootloader      = 0 R/W␣
↪(0x0)
DISABLE_DL_CACHE        Disable flash cache in UART bootloader           = 0 R/W␣
↪(0x0)
```

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```
BLK1                    Flash encryption key
  = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00␣
↪00 00 00 00 00 00 R/W
BLK2                    Secure boot key
  = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00␣
↪00 00 00 00 00 00 R/W
BLK3                    Variable Block 3
  = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 fa 87 02 91 00 00 00 00 00␣
↪00 00 00 00 00 00 R/W

Efuse fuses:
WR_DIS                  Efuse write disable mask                          = 0 R/W␣
↪(0x0)
RD_DIS                  Efuse read disablemask                            = 0 R/W␣
↪(0x0)
CODING_SCHEME           Efuse variable block length scheme                = 1 R/W␣
↪(0x1) (3/4)
KEY_STATUS              Usage of efuse block 3 (reserved)                 = 0 R/W␣
↪(0x0)

Config fuses:
XPD_SDIO_FORCE          Ignore MTDI pin (GPIO12) for VDD_SDIO on reset    = 0 R/W␣
↪(0x0)
XPD_SDIO_REG            If XPD_SDIO_FORCE, enable VDD_SDIO reg on reset   = 0 R/W␣
↪(0x0)
XPD_SDIO_TIEH           If XPD_SDIO_FORCE & XPD_SDIO_REG, 1=3.3V 0=1.8V   = 0 R/W␣
↪(0x0)
SPI_PAD_CONFIG_CLK      Override SD_CLK pad (GPIO6/SPICLK)                = 0 R/W␣
↪(0x0)
SPI_PAD_CONFIG_Q        Override SD_DATA_0 pad (GPIO7/SPIQ)               = 0 R/W␣
↪(0x0)
SPI_PAD_CONFIG_D        Override SD_DATA_1 pad (GPIO8/SPID)               = 0 R/W␣
↪(0x0)
SPI_PAD_CONFIG_HD       Override SD_DATA_2 pad (GPIO9/SPIHD)              = 0 R/W␣
↪(0x0)
SPI_PAD_CONFIG_CS0      Override SD_CMD pad (GPIO11/SPICS0)               = 0 R/W␣
↪(0x0)
DISABLE_SDIO_HOST       Disable SDIO host                                 = 0 R/W␣
↪(0x0)

Identity fuses:
MAC                     MAC Address
  = 84:0d:8e:18:8e:44 (CRC ad OK) R/W
CHIP_VER_REV1           Silicon Revision 1                                = 1 R/W␣
↪(0x1)
CHIP_VERSION            Reserved for future chip versions                 = 2 R/W␣
↪(0x2)
CHIP_PACKAGE            Chip package identifier                           = 0 R/W␣
↪(0x0)

Calibration fuses:
BLK3_PART_RESERVE       BLOCK3 partially served for ADC calibration data  = 1 R/W␣
↪(0x1)
ADC_VREF                Voltage reference calibration                     = 1114 R/
↪W (0x2)
ADC1_TP_LOW             ADC1 150mV reading                                = 346 R/W␣
↪(0x11)
ADC1_TP_HIGH            ADC1 850mV reading                                = 3285 R/
↪W (0x5)
ADC2_TP_LOW             ADC2 150mV reading                                = 449 R/W␣
↪(0x7)
```

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```
ADC2_TP_HIGH           ADC2 850mV reading                                    = 3362 R/
↪W (0x1f5)

Flash voltage (VDD_SDIO) determined by GPIO12 on reset (High for 1.8V, Low/NC for␣
↪3.3V).
```

To get a dump for all eFuse registers.

```
espefuse.py -p COM4 dump

$ espefuse.py -p COM4 dump
espefuse.py v2.3.1
Connecting........__
EFUSE block 0:
00000000 c403bb68 0082240a 00000000 00000035 00000000 00000000
EFUSE block 1:
00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000
EFUSE block 2:
00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000
EFUSE block 3:
00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000
```

Header File

- efuse/include/esp_efuse.h

Functions

esp_err_t **esp_efuse_read_field_blob**(**const** *esp_efuse_desc_t* **field*[], void **dst*, size_t *dst_size_bits*)

Reads bits from EFUSE field and writes it into an array.

The number of read bits will be limited to the minimum value from the description of the bits in "field" structure or "dst_size_bits" required size. Use "esp_efuse_get_field_size()" function to determine the length of the field.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.

Parameters

- `[in] field`: A pointer to the structure describing the fields of efuse.
- `[out] dst`: A pointer to array that will contain the result of reading.
- `[in] dst_size_bits`: The number of bits required to read. If the requested number of bits is greater than the field, the number will be limited to the field size.

bool **esp_efuse_read_field_bit**(**const** *esp_efuse_desc_t* **field*[])

Read a single bit eFuse field as a boolean value.

Note The value must exist and must be a single bit wide. If there is any possibility of an error in the provided arguments, call esp_efuse_read_field_blob() and check the returned value instead.

Note If assertions are enabled and the parameter is invalid, execution will abort

Return

- true: The field parameter is valid and the bit is set.
- false: The bit is not set, or the parameter is invalid and assertions are disabled.

Parameters

- `[in] field`: A pointer to the structure describing the fields of efuse.

esp_err_t **esp_efuse_read_field_cnt**(**const** *esp_efuse_desc_t* **field*[], size_t **out_cnt*)

Reads bits from EFUSE field and returns number of bits programmed as "1".

If the bits are set not sequentially, they will still be counted.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.

Parameters

- `[in] field`: A pointer to the structure describing the fields of efuse.
- `[out] out_cnt`: A pointer that will contain the number of programmed as "1" bits.

esp_err_t **esp_efuse_write_field_blob**(**const** *esp_efuse_desc_t* ***field*[], **const** void ***src*, size_t *src_size_bits*)

Writes array to EFUSE field.

The number of write bits will be limited to the minimum value from the description of the bits in "field" structure or "src_size_bits" required size. Use "esp_efuse_get_field_size()" function to determine the length of the field. After the function is completed, the writing registers are cleared.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.
- ESP_ERR_EFUSE_REPEATED_PROG: Error repeated programming of programmed bits is strictly forbidden.
- ESP_ERR_CODING: Error range of data does not match the coding scheme.

Parameters

- `[in] field`: A pointer to the structure describing the fields of efuse.
- `[in] src`: A pointer to array that contains the data for writing.
- `[in] src_size_bits`: The number of bits required to write.

esp_err_t **esp_efuse_write_field_cnt**(**const** *esp_efuse_desc_t* ***field*[], size_t *cnt*)

Writes a required count of bits as "1" to EFUSE field.

If there are no free bits in the field to set the required number of bits to "1", ESP_ERR_EFUSE_CNT_IS_FULL error is returned, the field will not be partially recorded. After the function is completed, the writing registers are cleared.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.
- ESP_ERR_EFUSE_CNT_IS_FULL: Not all requested cnt bits is set.

Parameters

- `[in] field`: A pointer to the structure describing the fields of efuse.
- `[in] cnt`: Required number of programmed as "1" bits.

esp_err_t **esp_efuse_write_field_bit**(**const** *esp_efuse_desc_t* ***field*[])

Write a single bit eFuse field to 1.

For use with eFuse fields that are a single bit. This function will write the bit to value 1 if it is not already set, or does nothing if the bit is already set.

This is equivalent to calling esp_efuse_write_field_cnt() with the cnt parameter equal to 1, except that it will return ESP_OK if the field is already set to 1.

Return

- ESP_OK: The operation was successfully completed, or the bit was already set to value 1.
- ESP_ERR_INVALID_ARG: Error in the passed arugments, including if the efuse field is not 1 bit wide.

Parameters

- `[in] field`: Pointer to the structure describing the efuse field.

esp_err_t **esp_efuse_set_write_protect**(esp_efuse_block_t *blk*)

Sets a write protection for the whole block.

After that, it is impossible to write to this block. The write protection does not apply to block 0.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.

- ESP_ERR_EFUSE_CNT_IS_FULL: Not all requested cnt bits is set.
- ESP_ERR_NOT_SUPPORTED: The block does not support this command.

Parameters

- `[in] blk`: Block number of eFuse. (EFUSE_BLK1, EFUSE_BLK2 and EFUSE_BLK3)

esp_err_t **esp_efuse_set_read_protect**(esp_efuse_block_t *blk*)

Sets a read protection for the whole block.

After that, it is impossible to read from this block. The read protection does not apply to block 0.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.
- ESP_ERR_EFUSE_CNT_IS_FULL: Not all requested cnt bits is set.
- ESP_ERR_NOT_SUPPORTED: The block does not support this command.

Parameters

- `[in] blk`: Block number of eFuse. (EFUSE_BLK1, EFUSE_BLK2 and EFUSE_BLK3)

int **esp_efuse_get_field_size**(**const** *esp_efuse_desc_t* **field*[])

Returns the number of bits used by field.

Return Returns the number of bits used by field.

Parameters

- `[in] field`: A pointer to the structure describing the fields of efuse.

uint32_t **esp_efuse_read_reg**(esp_efuse_block_t *blk*, unsigned int *num_reg*)

Returns value of efuse register.

This is a thread-safe implementation. Example: EFUSE_BLK2_RDATA3_REG where (blk=2, num_reg=3)

Return Value of register

Parameters

- `[in] blk`: Block number of eFuse.
- `[in] num_reg`: The register number in the block.

esp_err_t **esp_efuse_write_reg**(esp_efuse_block_t *blk*, unsigned int *num_reg*, uint32_t *val*)

Write value to efuse register.

Apply a coding scheme if necessary. This is a thread-safe implementation. Example: EFUSE_BLK3_WDATA0_REG where (blk=3, num_reg=0)

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_EFUSE_REPEATED_PROG: Error repeated programming of programmed bits is strictly forbidden.

Parameters

- `[in] blk`: Block number of eFuse.
- `[in] num_reg`: The register number in the block.
- `[in] val`: Value to write.

esp_efuse_coding_scheme_t **esp_efuse_get_coding_scheme**(esp_efuse_block_t *blk*)

Return efuse coding scheme for blocks.

Note: The coding scheme is applicable only to 1, 2 and 3 blocks. For 0 block, the coding scheme is always `NONE`.

Return Return efuse coding scheme for blocks

Parameters

- `[in] blk`: Block number of eFuse.

esp_err_t **esp_efuse_read_block**(esp_efuse_block_t *blk*, void **dst_key*, size_t *offset_in_bits*, size_t *size_bits*)

Read key to efuse block starting at the offset and the required size.

Return

- ESP_OK: The operation was successfully completed.

- ESP_ERR_INVALID_ARG: Error in the passed arguments.
- ESP_ERR_CODING: Error range of data does not match the coding scheme.

Parameters

- `[in] blk`: Block number of eFuse.
- `[in] dst_key`: A pointer to array that will contain the result of reading.
- `[in] offset_in_bits`: Start bit in block.
- `[in] size_bits`: The number of bits required to read.

esp_err_t **esp_efuse_write_block**(esp_efuse_block_t *blk*, **const** void **src_key*, size_t *offset_in_bits*, size_t *size_bits*)

Write key to efuse block starting at the offset and the required size.

Return

- ESP_OK: The operation was successfully completed.
- ESP_ERR_INVALID_ARG: Error in the passed arguments.
- ESP_ERR_CODING: Error range of data does not match the coding scheme.
- ESP_ERR_EFUSE_REPEATED_PROG: Error repeated programming of programmed bits

Parameters

- `[in] blk`: Block number of eFuse.
- `[in] src_key`: A pointer to array that contains the key for writing.
- `[in] offset_in_bits`: Start bit in block.
- `[in] size_bits`: The number of bits required to write.

uint8_t **esp_efuse_get_chip_ver**(void)

Returns chip version from efuse.

Return chip version

uint32_t **esp_efuse_get_pkg_ver**(void)

Returns chip package from efuse.

Return chip package

void **esp_efuse_burn_new_values**(void)

void **esp_efuse_reset**(void)

esp_err_t **esp_efuse_disable_rom_download_mode**(void)

esp_err_t **esp_efuse_enable_rom_secure_download_mode**(void)

void **esp_efuse_write_random_key**(uint32_t *blk_wdata0_reg*)

uint32_t **esp_efuse_read_secure_version**(void)

bool **esp_efuse_check_secure_version**(uint32_t *secure_version*)

esp_err_t **esp_efuse_update_secure_version**(uint32_t *secure_version*)

void **esp_efuse_init**(uint32_t *offset*, uint32_t *size*)

esp_err_t **esp_efuse_batch_write_begin**(void)

esp_err_t **esp_efuse_batch_write_cancel**(void)

esp_err_t **esp_efuse_batch_write_commit**(void)

Structures

struct esp_efuse_desc_s

Structure eFuse field.

Public Members

esp_efuse_block_t **efuse_block** : 8

Block of eFuse

uint8_t **bit_start**
Start bit [0..255]

uint16_t **bit_count**
Length of bit field [1..-]

Macros

ESP_ERR_EFUSE
Base error code for efuse api.

ESP_OK_EFUSE_CNT
OK the required number of bits is set.

ESP_ERR_EFUSE_CNT_IS_FULL
Error field is full.

ESP_ERR_EFUSE_REPEATED_PROG
Error repeated programming of programmed bits is strictly forbidden.

ESP_ERR_CODING
Error while a encoding operation.

Type Definitions

typedef struct *esp_efuse_desc_s* **esp_efuse_desc_t**
Type definition for an eFuse field.

2.6.5 Error Codes and Helper Functions

This section lists definitions of common ESP-IDF error codes and several helper functions related to error handling.

For general information about error codes in ESP-IDF, see *Error Handling*.

For the full list of error codes defined in ESP-IDF, see *Error Code Reference*.

API Reference

Header File

- esp_common/include/esp_err.h

Functions

const char ***esp_err_to_name**(*esp_err_t code*)
Returns string for esp_err_t error codes.

This function finds the error code in a pre-generated lookup-table and returns its string representation.

The function is generated by the Python script tools/gen_esp_err_to_name.py which should be run each time an esp_err_t error is modified, created or removed from the IDF project.

Return string error message
Parameters
- `code`: esp_err_t error code

const char ***esp_err_to_name_r**(*esp_err_t code*, char **buf*, size_t *buflen*)
Returns string for esp_err_t and system error codes.

This function finds the error code in a pre-generated lookup-table of esp_err_t errors and returns its string representation. If the error code is not found then it is attempted to be found among system errors.

The function is generated by the Python script tools/gen_esp_err_to_name.py which should be run each time an esp_err_t error is modified, created or removed from the IDF project.

Return buf containing the string error message

Parameters

- `code`: esp_err_t error code
- `[out] buf`: buffer where the error message should be written
- `buflen`: Size of buffer buf. At most buflen bytes are written into the buf buffer (including the terminating null byte).

Macros

ESP_OK
esp_err_t value indicating success (no error)

ESP_FAIL
Generic esp_err_t code indicating failure

ESP_ERR_NO_MEM
Out of memory

ESP_ERR_INVALID_ARG
Invalid argument

ESP_ERR_INVALID_STATE
Invalid state

ESP_ERR_INVALID_SIZE
Invalid size

ESP_ERR_NOT_FOUND
Requested resource not found

ESP_ERR_NOT_SUPPORTED
Operation or feature not supported

ESP_ERR_TIMEOUT
Operation timed out

ESP_ERR_INVALID_RESPONSE
Received response was invalid

ESP_ERR_INVALID_CRC
CRC or checksum was invalid

ESP_ERR_INVALID_VERSION
Version was invalid

ESP_ERR_INVALID_MAC
MAC address was invalid

ESP_ERR_WIFI_BASE
Starting number of WiFi error codes

ESP_ERR_MESH_BASE
Starting number of MESH error codes

ESP_ERR_FLASH_BASE
Starting number of flash error codes

ESP_ERROR_CHECK(x)
Macro which can be used to check the error code, and terminate the program in case the code is not ESP_OK. Prints the error code, error location, and the failed statement to serial output.

Disabled if assertions are disabled.

ESP_ERROR_CHECK_WITHOUT_ABORT(x)
Macro which can be used to check the error code. Prints the error code, error location, and the failed statement to serial output. In comparison with ESP_ERROR_CHECK(), this prints the same error message but isn't terminating the program.

Type Definitions

typedef int32_t **esp_err_t**

2.6.6 ESP HTTPS OTA

Overview

`esp_https_ota` provides simplified APIs to perform firmware upgrades over HTTPS. It's an abstraction layer over existing OTA APIs.

Application Example

```
esp_err_t do_firmware_upgrade()
{
    esp_http_client_config_t config = {
        .url = CONFIG_FIRMWARE_UPGRADE_URL,
        .cert_pem = (char *)server_cert_pem_start,
    };
    esp_err_t ret = esp_https_ota(&config);
    if (ret == ESP_OK) {
        esp_restart();
    } else {
        return ESP_FAIL;
    }
    return ESP_OK;
}
```

API Reference

Header File

- esp_https_ota/include/esp_https_ota.h

Functions

esp_err_t **esp_https_ota**(**const** *esp_http_client_config_t* **config*)

HTTPS OTA Firmware upgrade.

This function allocates HTTPS OTA Firmware upgrade context, establishes HTTPS connection, reads image data from HTTP stream and writes it to OTA partition and finishes HTTPS OTA Firmware upgrade operation. This API supports URL redirection, but if CA cert of URLs differ then it should be appended to `cert_pem` member of `config`.

Note This API handles the entire OTA operation, so if this API is being used then no other APIs from `esp_https_ota` component should be called. If more information and control is needed during the HTTPS OTA process, then one can use `esp_https_ota_begin` and subsequent APIs. If this API returns successfully, esp_restart() must be called to boot from the new firmware image.

Return

- ESP_OK: OTA data updated, next reboot will use specified partition.
- ESP_FAIL: For generic failure.
- ESP_ERR_INVALID_ARG: Invalid argument
- ESP_ERR_OTA_VALIDATE_FAILED: Invalid app image
- ESP_ERR_NO_MEM: Cannot allocate memory for OTA operation.
- ESP_ERR_FLASH_OP_TIMEOUT or ESP_ERR_FLASH_OP_FAIL: Flash write failed.
- For other return codes, refer OTA documentation in esp-idf's app_update component.

Parameters

- `[in] config`: pointer to *esp_http_client_config_t* structure.

esp_err_t **esp_https_ota_begin**(*esp_https_ota_config_t *ota_config*, *esp_https_ota_handle_t *handle*)

Start HTTPS OTA Firmware upgrade.

This function initializes ESP HTTPS OTA context and establishes HTTPS connection. This function must be invoked first. If this function returns successfully, then `esp_https_ota_perform` should be called to continue with the OTA process and there should be a call to `esp_https_ota_finish` on completion of OTA operation or on failure in subsequent operations. This API supports URL redirection, but if CA cert of URLs differ then it should be appended to `cert_pem` member of `http_config`, which is a part of `ota_config`. In case of error, this API explicitly sets `handle` to NULL.

Note This API is blocking, so setting `is_async` member of `http_config` structure will result in an error.

Return

- ESP_OK: HTTPS OTA Firmware upgrade context initialised and HTTPS connection established
- ESP_FAIL: For generic failure.
- ESP_ERR_INVALID_ARG: Invalid argument (missing/incorrect config, certificate, etc.)
- For other return codes, refer documentation in app_update component and esp_http_client component in esp-idf.

Parameters

- `[in] ota_config`: pointer to *esp_https_ota_config_t* structure
- `[out] handle`: pointer to an allocated data of type `esp_https_ota_handle_t` which will be initialised in this function

esp_err_t **esp_https_ota_perform**(*esp_https_ota_handle_t https_ota_handle*)

Read image data from HTTP stream and write it to OTA partition.

This function reads image data from HTTP stream and writes it to OTA partition. This function must be called only if esp_https_ota_begin() returns successfully. This function must be called in a loop since it returns after every HTTP read operation thus giving you the flexibility to stop OTA operation midway.

Return

- ESP_ERR_HTTPS_OTA_IN_PROGRESS: OTA update is in progress, call this API again to continue.
- ESP_OK: OTA update was successful
- ESP_FAIL: OTA update failed
- ESP_ERR_INVALID_ARG: Invalid argument
- ESP_ERR_OTA_VALIDATE_FAILED: Invalid app image
- ESP_ERR_NO_MEM: Cannot allocate memory for OTA operation.
- ESP_ERR_FLASH_OP_TIMEOUT or ESP_ERR_FLASH_OP_FAIL: Flash write failed.
- For other return codes, refer OTA documentation in esp-idf's app_update component.

Parameters

- `[in] https_ota_handle`: pointer to esp_https_ota_handle_t structure

bool **esp_https_ota_is_complete_data_received**(*esp_https_ota_handle_t https_ota_handle*)

Checks if complete data was received or not.

Note This API can be called just before esp_https_ota_finish() to validate if the complete image was indeed received.

Return

- false
- true

Parameters

- `[in] https_ota_handle`: pointer to esp_https_ota_handle_t structure

esp_err_t **esp_https_ota_finish**(*esp_https_ota_handle_t https_ota_handle*)

Clean-up HTTPS OTA Firmware upgrade and close HTTPS connection.

This function closes the HTTP connection and frees the ESP HTTPS OTA context. This function switches the boot partition to the OTA partition containing the new firmware image.

Note If this API returns successfully, esp_restart() must be called to boot from the new firmware image

Return

- ESP_OK: Clean-up successful

- ESP_ERR_INVALID_STATE
- ESP_ERR_INVALID_ARG: Invalid argument
- ESP_ERR_OTA_VALIDATE_FAILED: Invalid app image

Parameters

- `[in] https_ota_handle`: pointer to esp_https_ota_handle_t structure

esp_err_t **esp_https_ota_get_img_desc**(*esp_https_ota_handle_t https_ota_handle*, *esp_app_desc_t *new_app_info*)

Reads app description from image header. The app description provides information like the "Firmware version" of the image.

Note This API can be called only after esp_https_ota_begin() and before esp_https_ota_perform(). Calling this API is not mandatory.

Return

- ESP_ERR_INVALID_ARG: Invalid arguments
- ESP_FAIL: Failed to read image descriptor
- ESP_OK: Successfully read image descriptor

Parameters

- `[in] https_ota_handle`: pointer to esp_https_ota_handle_t structure
- `[out] new_app_info`: pointer to an allocated *esp_app_desc_t* structure

int **esp_https_ota_get_image_len_read**(*esp_https_ota_handle_t https_ota_handle*)

This function returns OTA image data read so far.

Note This API should be called only if `esp_https_ota_perform()` has been called atleast once or if `esp_https_ota_get_img_desc` has been called before.

Return

- -1 On failure
- total bytes read so far

Parameters

- `[in] https_ota_handle`: pointer to esp_https_ota_handle_t structure

Structures

struct esp_https_ota_config_t

ESP HTTPS OTA configuration.

Public Members

const *esp_http_client_config_t* ***http_config**

ESP HTTP client configuration

Macros

ESP_ERR_HTTPS_OTA_BASE

ESP_ERR_HTTPS_OTA_IN_PROGRESS

Type Definitions

typedef void ***esp_https_ota_handle_t**

2.6.7 ESP-pthread

Overview

This module offers Espressif specific extensions to the pthread library that can be used to influence the behaviour of pthreads

- Stack size of the pthreads
- Priority of the created pthreads

- Inheriting this configuration across threads
- Thread name
- Core affinity / core pinning.

Example to tune the stack size of the pthread:

```
void * thread_func(void * p)
{
    printf("In thread_func\n");
    return NULL;
}

void app_main(void)
{
    pthread_t t1;

    esp_pthread_cfg_t cfg = esp_create_default_pthread_config();
    cfg.stack_size = (4 * 1024);
    esp_pthread_set_cfg(&cfg);

    pthread_create(&t1, NULL, thread_func);
}
```

The API can also be used for inheriting the settings across threads. For example:

```
void * my_thread2(void * p)
{
    /* This thread will inherit the stack size of 4K */
    printf("In my_thread2\n");

    return NULL;
}

void * my_thread1(void * p)
{
    printf("In my_thread1\n");
    pthread_t t2;
    pthread_create(&t2, NULL, my_thread2);

    return NULL;
}

void app_main(void)
{
    pthread_t t1;

    esp_pthread_cfg_t cfg = esp_create_default_pthread_config();
    cfg.stack_size = (4 * 1024);
    cfg.inherit_cfg = true;
    esp_pthread_set_cfg(&cfg);

    pthread_create(&t1, NULL, my_thread1);
}
```

API Reference

Header File

- pthread/include/esp_pthread.h

Functions

esp_pthread_cfg_t **esp_pthread_get_default_config**(void)
Creates a default pthread configuration based on the values set via menuconfig.

Return A default configuration structure.

esp_err_t **esp_pthread_set_cfg**(**const** *esp_pthread_cfg_t* *cfg)
Configure parameters for creating pthread.

This API allows you to configure how the subsequent pthread_create() call will behave. This call can be used to setup configuration parameters like stack size, priority, configuration inheritance etc.

If the 'inherit' flag in the configuration structure is enabled, then the same configuration is also inherited in the thread subtree.

Note Passing non-NULL attributes to pthread_create() will override the stack_size parameter set using this API

Return
- ESP_OK if configuration was successfully set
- ESP_ERR_NO_MEM if out of memory
- ESP_ERR_INVALID_ARG if stack_size is less than PTHREAD_STACK_MIN

Parameters
- `cfg`: The pthread config parameters

esp_err_t **esp_pthread_get_cfg**(*esp_pthread_cfg_t* *p)
Get current pthread creation configuration.

This will retrieve the current configuration that will be used for creating threads.

Return
- ESP_OK if the configuration was available
- ESP_ERR_NOT_FOUND if a configuration wasn't previously set

Parameters
- `p`: Pointer to the pthread config structure that will be updated with the currently configured parameters

Structures

struct esp_pthread_cfg_t
pthread configuration structure that influences pthread creation

Public Members

size_t **stack_size**
The stack size of the pthread.

size_t **prio**
The thread's priority.

bool **inherit_cfg**
Inherit this configuration further.

const char ***thread_name**
The thread name.

int **pin_to_core**
The core id to pin the thread to. Has the same value range as xCoreId argument of xTaskCreatePinnedToCore.

Macros

PTHREAD_STACK_MIN

2.6.8 Event Loop Library

Overview

The event loop library allows components to declare events to which other components can register handlers –code which will execute when those events occur. This allows loosely coupled components to attach desired behavior to changes in state of other components without application involvement. For instance, a high level connection handling library may subscribe to events produced by the wifi subsystem directly and act on those events. This also simplifies event processing by serializing and deferring code execution to another context.

Using `esp_event` APIs

There are two objects of concern for users of this library: events and event loops.

Events are occurrences of note. For example, for WiFi, a successful connection to the access point may be an event. Events are referenced using a two part identifier which are discussed more *here*. Event loops are the vehicle by which events get posted by event sources and handled by event handler functions. These two appear prominently in the event loop library APIs.

Using this library roughly entails the following flow:

1. A user defines a function that should run when an event is posted to a loop. This function is referred to as the event handler. It should have the same signature as `esp_event_handler_t`.
2. An event loop is created using `esp_event_loop_create()`, which outputs a handle to the loop of type `esp_event_loop_handle_t`. Event loops created using this API are referred to as user event loops. There is, however, a special type of event loop called the default event loop which are discussed *here*.
3. Components register event handlers to the loop using `esp_event_handler_register_with()`. Handlers can be registered with multiple loops, more on that *here*.
4. Event sources post an event to the loop using `esp_event_post_to()`.
5. Components wanting to remove their handlers from being called can do so by unregistering from the loop using `esp_event_handler_unregister_with()`.
6. Event loops which are no longer needed can be deleted using `esp_event_loop_delete()`.

In code, the flow above may look like as follows:

```
// 1. Define the event handler
void run_on_event(void* handler_arg, esp_event_base_t base, int32_t id, void*␣
↪event_data)
{
    // Event handler logic
}

void app_main()
{
    // 2. A configuration structure of type esp_event_loop_args_t is needed to␣
↪specify the properties of the loop to be
    // created. A handle of type esp_event_loop_handle_t is obtained, which is␣
↪needed by the other APIs to reference the loop
    // to perform their operations on.
    esp_event_loop_args_t loop_args = {
        .queue_size = ...,
        .task_name = ...
        .task_priority = ...,
        .task_stack_size = ...,
        .task_core_id = ...
    };

    esp_event_loop_handle_t loop_handle;

    esp_event_loop_create(&loop_args, &loop_handle);
```

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```
    // 3. Register event handler defined in (1). MY_EVENT_BASE and MY_EVENT_ID␣
↪specifies a hypothetical
    // event that handler run_on_event should execute on when it gets posted to␣
↪the loop.
    esp_event_handler_register_with(loop_handle, MY_EVENT_BASE, MY_EVENT_ID, run_
↪on_event, ...);

    ...

    // 4. Post events to the loop. This queues the event on the event loop. At␣
↪some point in time
    // the event loop executes the event handler registered to the posted event,␣
↪in this case run_on_event.
    // For simplicity sake this example calls esp_event_post_to from app_main, but␣
↪posting can be done from
    // any other tasks (which is the more interesting use case).
    esp_event_post_to(loop_handle, MY_EVENT_BASE, MY_EVENT_ID, ...);

    ...

    // 5. Unregistering an unneeded handler
    esp_event_handler_unregister_with(loop_handle, MY_EVENT_BASE, MY_EVENT_ID, run_
↪on_event);

    ...

    // 6. Deleting an unneeded event loop
    esp_event_loop_delete(loop_handle);
}
```

Declaring and defining events

As mentioned previously, events consists of two-part identifers: the event base and the event ID. The event base identifies an independent group of events; the event ID identifies the event within that group. Think of the event base and event ID as a person's last name and first name, respectively. A last name identifies a family, and the first name identifies a person within that family.

The event loop library provides macros to declare and define the event base easily.

Event base declaration:

```
ESP_EVENT_DECLARE_BASE(EVENT_BASE)
```

Event base definition:

```
ESP_EVENT_DEFINE_BASE(EVENT_BASE)
```

注解: In IDF, the base identifiers for system events are uppercase and are postfixed with `_EVENT`. For example, the base for wifi events is declared and defined as `WIFI_EVENT`, the ethernet event base `ETHERNET_EVENT`, and so on. The purpose is to have event bases look like constants (although they are global variables considering the defintions of macros `ESP_EVENT_DECLARE_BASE` and `ESP_EVENT_DEFINE_BASE`).

For event ID's, declaring them as enumerations is recommended. Once again, for visibility, these are typically placed in public header files.

Event ID:

```
enum {
    EVENT_ID_1,
    EVENT_ID_2,
    EVENT_ID_3,
    ...
}
```

Default Event Loop

The default event loop is a special type of loop used for system events (WiFi events, for example). The handle for this loop is hidden from the user. The creation, deletion, handler registration/unregistration and posting of events is done through a variant of the APIs for user event loops. The table below enumerates those variants, and the user event loops equivalent.

User Event Loops	Default Event Loops
esp_event_loop_create()	*esp_event_loop_create_default()*
esp_event_loop_delete()	*esp_event_loop_delete_default()*
esp_event_handler_register_with()	*esp_event_handler_register()*
esp_event_handler_unregister_with()	*esp_event_handler_unregister()*
esp_event_post_to()	*esp_event_post()*

If you compare the signatures for both, they are mostly similar except the for the lack of loop handle specification for the default event loop APIs.

Other than the API difference and the special designation to which system events are posted to, there is no difference to how default event loops and user event loops behave. It is even possible for users to post their own events to the default event loop, should the user opt to not create their own loops to save memory.

Notes on Handler Registration

It is possible to register a single handler to multiple events individually, i.e. using multiple calls to *esp_event_handler_register_with()*. For those multiple calls, the specific event base and event ID can be specified with which the handler should execute.

However, in some cases it is desirable for a handler to execute on (1) all events that get posted to a loop or (2) all events of a particular base identifier. This is possible using the special event base identifier `ESP_EVENT_ANY_BASE` and special event ID `ESP_EVENT_ANY_ID`. These special identifiers may be passed as the event base and event ID arguments for *esp_event_handler_register_with()*.

Therefore, the valid arguments to *esp_event_handler_register_with()* are:

1. <event base>, <event ID> - handler executes when the event with base <event base> and event ID <event ID> gets posted to the loop
2. <event base>, ESP_EVENT_ANY_ID - handler executes when any event with base <event base> gets posted to the loop
3. ESP_EVENT_ANY_BASE, ESP_EVENT_ANY_ID - handler executes when any event gets posted to the loop

As an example, suppose the following handler registrations were performed:

```
esp_event_handler_register_with(loop_handle, MY_EVENT_BASE, MY_EVENT_ID, run_on_
↪event_1, ...);
esp_event_handler_register_with(loop_handle, MY_EVENT_BASE, ESP_EVENT_ANY_ID, run_
↪on_event_2, ...);
esp_event_handler_register_with(loop_handle, ESP_EVENT_ANY_BASE, ESP_EVENT_ANY_ID,␣
↪run_on_event_3, ...);
```

If the hypothetical event `MY_EVENT_BASE`, `MY_EVENT_ID` is posted, all three handlers `run_on_event_1`, `run_on_event_2`, and `run_on_event_3` would execute.

If the hypothetical event `MY_EVENT_BASE`, `MY_OTHER_EVENT_ID` is posted, only `run_on_event_2` and `run_on_event_3` would execute.

If the hypothetical event `MY_OTHER_EVENT_BASE`, `MY_OTHER_EVENT_ID` is posted, only `run_on_event_3` would execute.

Handler Registration and Handler Dispatch Order The general rule is that for handlers that match a certain posted event during dispatch, those which are registered first also gets executed first. The user can then control which handlers get executed first by registering them before other handlers, provided that all registrations are performed using a single task. If the user plans to take advantage of this behavior, caution must be exercised if there are multiple tasks registering handlers. While the 'first registered, first executed' behavior still holds true, the task which gets executed first will also get their handlers registered first. Handlers registered one after the other by a single task will still be dispatched in the order relative to each other, but if that task gets pre-empted in between registration by another task which also registers handlers; then during dispatch those handlers will also get executed in between.

Event loop profiling

A configuration option *CONFIG_ESP_EVENT_LOOP_PROFILING* can be enabled in order to activate statistics collection for all event loops created. The function `esp_event_dump()` can be used to output the collected statistics to a file stream. More details on the information included in the dump can be found in the `esp_event_dump()` API Reference.

Application Example

Examples on using the `esp_event` library can be found in system/esp_event. The examples cover event declaration, loop creation, handler registration and unregistration and event posting.

Other examples which also adopt esp_event library:

- NMEA Parser , which will decode the statements received from GPS.

API Reference

Header File

- esp_event/include/esp_event.h

Functions

esp_err_t **esp_event_loop_create**(**const** *esp_event_loop_args_t* **event_loop_args*, *esp_event_loop_handle_t* **event_loop*)

Create a new event loop.

Return

- ESP_OK: Success
- ESP_ERR_INVALID_ARG: event_loop_args or event_loop was NULL
- ESP_ERR_NO_MEM: Cannot allocate memory for event loops list
- ESP_FAIL: Failed to create task loop
- Others: Fail

Parameters

- `[in] event_loop_args`: configuration structure for the event loop to create
- `[out] event_loop`: handle to the created event loop

esp_err_t **esp_event_loop_delete**(*esp_event_loop_handle_t event_loop*)

Delete an existing event loop.

Return

- ESP_OK: Success
- Others: Fail

Parameters

- `[in] event_loop`: event loop to delete, must not be NULL

esp_err_t **esp_event_loop_create_default**(void)

Create default event loop.

Return

- ESP_OK: Success
- ESP_ERR_NO_MEM: Cannot allocate memory for event loops list
- ESP_FAIL: Failed to create task loop
- Others: Fail

esp_err_t **esp_event_loop_delete_default**(void)

Delete the default event loop.

Return

- ESP_OK: Success
- Others: Fail

esp_err_t **esp_event_loop_run**(*esp_event_loop_handle_t event_loop*, TickType_t *ticks_to_run*)

Dispatch events posted to an event loop.

This function is used to dispatch events posted to a loop with no dedicated task, i.e task name was set to NULL in event_loop_args argument during loop creation. This function includes an argument to limit the amount of time it runs, returning control to the caller when that time expires (or some time afterwards). There is no guarantee that a call to this function will exit at exactly the time of expiry. There is also no guarantee that events have been dispatched during the call, as the function might have spent all of the alloted time waiting on the event queue. Once an event has been unqueued, however, it is guaranteed to be dispatched. This guarantee contributes to not being able to exit exactly at time of expiry as (1) blocking on internal mutexes is necessary for dispatching the unqueued event, and (2) during dispatch of the unqueued event there is no way to control the time occupied by handler code execution. The guaranteed time of exit is therefore the alloted time + amount of time required to dispatch the last unqueued event.

In cases where waiting on the queue times out, ESP_OK is returned and not ESP_ERR_TIMEOUT, since it is normal behavior.

Note encountering an unknown event that has been posted to the loop will only generate a warning, not an error.

Return

- ESP_OK: Success
- Others: Fail

Parameters

- `[in] event_loop`: event loop to dispatch posted events from, must not be NULL
- `[in] ticks_to_run`: number of ticks to run the loop

esp_err_t **esp_event_handler_register**(*esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_t event_handler*, void **event_handler_arg*)

Register an event handler to the system event loop (legacy).

This function can be used to register a handler for either: (1) specific events, (2) all events of a certain event base, or (3) all events known by the system event loop.

Note This function is obsolete and will be deprecated soon, please use esp_event_handler_instance_register() instead.

- specific events: specify exact event_base and event_id
- all events of a certain base: specify exact event_base and use ESP_EVENT_ANY_ID as the event_id
- all events known by the loop: use ESP_EVENT_ANY_BASE for event_base and ESP_EVENT_ANY_ID as the event_id

Registering multiple handlers to events is possible. Registering a single handler to multiple events is also possible. However, registering the same handler to the same event multiple times would cause the previous

registrations to be overwritten.

Note the event loop library does not maintain a copy of event_handler_arg, therefore the user should ensure that event_handler_arg still points to a valid location by the time the handler gets called

Return

- ESP_OK: Success
- ESP_ERR_NO_MEM: Cannot allocate memory for the handler
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- `[in] event_base`: the base id of the event to register the handler for
- `[in] event_id`: the id of the event to register the handler for
- `[in] event_handler`: the handler function which gets called when the event is dispatched
- `[in] event_handler_arg`: data, aside from event data, that is passed to the handler when it is called

esp_err_t **esp_event_handler_register_with**(*esp_event_loop_handle_t event_loop*, *esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_t event_handler*, void **event_handler_arg*)

Register an event handler to a specific loop (legacy).

This function behaves in the same manner as esp_event_handler_register, except the additional specification of the event loop to register the handler to.

Note This function is obsolete and will be deprecated soon, please use esp_event_handler_instance_register_with() instead.

Note the event loop library does not maintain a copy of event_handler_arg, therefore the user should ensure that event_handler_arg still points to a valid location by the time the handler gets called

Return

- ESP_OK: Success
- ESP_ERR_NO_MEM: Cannot allocate memory for the handler
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- `[in] event_loop`: the event loop to register this handler function to, must not be NULL
- `[in] event_base`: the base id of the event to register the handler for
- `[in] event_id`: the id of the event to register the handler for
- `[in] event_handler`: the handler function which gets called when the event is dispatched
- `[in] event_handler_arg`: data, aside from event data, that is passed to the handler when it is called

esp_err_t **esp_event_handler_instance_register_with**(*esp_event_loop_handle_t event_loop*, *esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_t event_handler*, void **event_handler_arg*, *esp_event_handler_instance_t* **instance*)

Register an instance of event handler to a specific loop.

This function can be used to register a handler for either: (1) specific events, (2) all events of a certain event base, or (3) all events known by the system event loop.

- specific events: specify exact event_base and event_id
- all events of a certain base: specify exact event_base and use ESP_EVENT_ANY_ID as the event_id
- all events known by the loop: use ESP_EVENT_ANY_BASE for event_base and ESP_EVENT_ANY_ID as the event_id

Besides the error, the function returns an instance object as output parameter to identify each registration. This is necessary to remove (unregister) the registration before the event loop is deleted.

Registering multiple handlers to events, registering a single handler to multiple events as well as registering the same handler to the same event multiple times is possible. Each registration yields a distinct instance object which identifies it over the registration lifetime.

Note the event loop library does not maintain a copy of event_handler_arg, therefore the user should ensure that event_handler_arg still points to a valid location by the time the handler gets called

Return

- ESP_OK: Success
- ESP_ERR_NO_MEM: Cannot allocate memory for the handler
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id or instance is NULL
- Others: Fail

Parameters

- `[in] event_loop`: the event loop to register this handler function to, must not be NULL
- `[in] event_base`: the base id of the event to register the handler for
- `[in] event_id`: the id of the event to register the handler for
- `[in] event_handler`: the handler function which gets called when the event is dispatched
- `[in] event_handler_arg`: data, aside from event data, that is passed to the handler when it is called
- `[out] instance`: An event handler instance object related to the registered event handler and data, can be NULL. This needs to be kept if the specific callback instance should be unregistered before deleting the whole event loop. Registering the same event handler multiple times is possible and yields distinct instance objects. The data can be the same for all registrations. If no unregistration is needed but the handler should be deleted when the event loop is deleted, instance can be NULL.

esp_err_t **esp_event_handler_instance_register**(*esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_t event_handler*, void **event_handler_arg*, *esp_event_handler_instance_t *instance*)

Register an instance of event handler to the default loop.

This function does the same as esp_event_handler_instance_register_with, except that it registers the handler to the default event loop.

Note the event loop library does not maintain a copy of event_handler_arg, therefore the user should ensure that event_handler_arg still points to a valid location by the time the handler gets called

Return

- ESP_OK: Success
- ESP_ERR_NO_MEM: Cannot allocate memory for the handler
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id or instance is NULL
- Others: Fail

Parameters

- `[in] event_base`: the base id of the event to register the handler for
- `[in] event_id`: the id of the event to register the handler for
- `[in] event_handler`: the handler function which gets called when the event is dispatched
- `[in] event_handler_arg`: data, aside from event data, that is passed to the handler when it is called
- `[out] instance`: An event handler instance object related to the registered event handler and data, can be NULL. This needs to be kept if the specific callback instance should be unregistered before deleting the whole event loop. Registering the same event handler multiple times is possible and yields distinct instance objects. The data can be the same for all registrations. If no unregistration is needed but the handler should be deleted when the event loop is deleted, instance can be NULL.

esp_err_t **esp_event_handler_unregister**(*esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_t event_handler*)

Unregister a handler with the system event loop (legacy).

Unregisters a handler so it will no longer be called during dispatch. Handlers can be unregistered for any combination of event_base and event_id which were previously registered. To unregister a handler, the event_base and event_id arguments must match exactly the arguments passed to esp_event_handler_register() when that handler was registered. Passing ESP_EVENT_ANY_BASE and/or ESP_EVENT_ANY_ID will only unregister handlers that were registered with the same wildcard arguments.

Note This function is obsolete and will be deprecated soon, please use esp_event_handler_instance_unregister() instead.

Note When using ESP_EVENT_ANY_ID, handlers registered to specific event IDs using the same base will not be unregistered. When using ESP_EVENT_ANY_BASE, events registered to specific bases will also not be unregistered. This avoids accidental unregistration of handlers registered by other users or components.

Return ESP_OK success

Return ESP_ERR_INVALID_ARG invalid combination of event base and event id

Return others fail

Parameters

- `[in] event_base`: the base of the event with which to unregister the handler
- `[in] event_id`: the id of the event with which to unregister the handler
- `[in] event_handler`: the handler to unregister

esp_err_t **esp_event_handler_unregister_with**(*esp_event_loop_handle_t event_loop*, *esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_t event_handler*)

Unregister a handler from a specific event loop (legacy).

This function behaves in the same manner as esp_event_handler_unregister, except the additional specification of the event loop to unregister the handler with.

Note This function is obsolete and will be deprecated soon, please use esp_event_handler_instance_unregister_with() instead.

Return

- ESP_OK: Success
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- `[in] event_loop`: the event loop with which to unregister this handler function, must not be NULL
- `[in] event_base`: the base of the event with which to unregister the handler
- `[in] event_id`: the id of the event with which to unregister the handler
- `[in] event_handler`: the handler to unregister

esp_err_t **esp_event_handler_instance_unregister_with**(*esp_event_loop_handle_t event_loop*, *esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_instance_t instance*)

Unregister a handler instance from a specific event loop.

Unregisters a handler instance so it will no longer be called during dispatch. Handler instances can be unregistered for any combination of event_base and event_id which were previously registered. To unregister a handler instance, the event_base and event_id arguments must match exactly the arguments passed to esp_event_handler_instance_register() when that handler instance was registered. Passing ESP_EVENT_ANY_BASE and/or ESP_EVENT_ANY_ID will only unregister handler instances that were registered with the same wildcard arguments.

Note When using ESP_EVENT_ANY_ID, handlers registered to specific event IDs using the same base will not be unregistered. When using ESP_EVENT_ANY_BASE, events registered to specific bases will also not be unregistered. This avoids accidental unregistration of handlers registered by other users or components.

Return

- ESP_OK: Success
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- `[in] event_loop`: the event loop with which to unregister this handler function, must not be NULL

- [in] event_base: the base of the event with which to unregister the handler
- [in] event_id: the id of the event with which to unregister the handler
- [in] instance: the instance object of the registration to be unregistered

esp_err_t **esp_event_handler_instance_unregister**(*esp_event_base_t event_base*, int32_t *event_id*, *esp_event_handler_instance_t instance*)

Unregister a handler from the system event loop.

This function does the same as esp_event_handler_instance_unregister_with, except that it unregisters the handler instance from the default event loop.

Return

- ESP_OK: Success
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- [in] event_base: the base of the event with which to unregister the handler
- [in] event_id: the id of the event with which to unregister the handler
- [in] instance: the instance object of the registration to be unregistered

esp_err_t **esp_event_post**(*esp_event_base_t event_base*, int32_t *event_id*, void **event_data*, size_t *event_data_size*, TickType_t *ticks_to_wait*)

Posts an event to the system default event loop. The event loop library keeps a copy of event_data and manages the copy's lifetime automatically (allocation + deletion); this ensures that the data the handler recieves is always valid.

Return

- ESP_OK: Success
- ESP_ERR_TIMEOUT: Time to wait for event queue to unblock expired, queue full when posting from ISR
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- [in] event_base: the event base that identifies the event
- [in] event_id: the event id that identifies the event
- [in] event_data: the data, specific to the event occurence, that gets passed to the handler
- [in] event_data_size: the size of the event data
- [in] ticks_to_wait: number of ticks to block on a full event queue

esp_err_t **esp_event_post_to**(*esp_event_loop_handle_t event_loop*, *esp_event_base_t event_base*, int32_t *event_id*, void **event_data*, size_t *event_data_size*, TickType_t *ticks_to_wait*)

Posts an event to the specified event loop. The event loop library keeps a copy of event_data and manages the copy's lifetime automatically (allocation + deletion); this ensures that the data the handler recieves is always valid.

This function behaves in the same manner as esp_event_post_to, except the additional specification of the event loop to post the event to.

Return

- ESP_OK: Success
- ESP_ERR_TIMEOUT: Time to wait for event queue to unblock expired, queue full when posting from ISR
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id
- Others: Fail

Parameters

- [in] event_loop: the event loop to post to, must not be NULL
- [in] event_base: the event base that identifies the event
- [in] event_id: the event id that identifies the event
- [in] event_data: the data, specific to the event occurence, that gets passed to the handler
- [in] event_data_size: the size of the event data
- [in] ticks_to_wait: number of ticks to block on a full event queue

esp_err_t **esp_event_isr_post**(*esp_event_base_t event_base*, int32_t *event_id*, void **event_data*, size_t *event_data_size*, BaseType_t **task_unblocked*)

Special variant of esp_event_post for posting events from interrupt handlers.

Note this function is only available when CONFIG_ESP_EVENT_POST_FROM_ISR is enabled

Note when this function is called from an interrupt handler placed in IRAM, this function should be placed in IRAM as well by enabling CONFIG_ESP_EVENT_POST_FROM_IRAM_ISR

Return

- ESP_OK: Success
- ESP_FAIL: Event queue for the default event loop full
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id, data size of more than 4 bytes
- Others: Fail

Parameters

- `[in] event_base`: the event base that identifies the event
- `[in] event_id`: the event id that identifies the event
- `[in] event_data`: the data, specific to the event occurence, that gets passed to the handler
- `[in] event_data_size`: the size of the event data; max is 4 bytes
- `[out] task_unblocked`: an optional parameter (can be NULL) which indicates that an event task with higher priority than currently running task has been unblocked by the posted event; a context switch should be requested before the interrupt is existed.

esp_err_t **esp_event_isr_post_to**(*esp_event_loop_handle_t event_loop*, *esp_event_base_t event_base*, int32_t *event_id*, void **event_data*, size_t *event_data_size*, BaseType_t **task_unblocked*)

Special variant of esp_event_post_to for posting events from interrupt handlers.

Note this function is only available when CONFIG_ESP_EVENT_POST_FROM_ISR is enabled

Note when this function is called from an interrupt handler placed in IRAM, this function should be placed in IRAM as well by enabling CONFIG_ESP_EVENT_POST_FROM_IRAM_ISR

Return

- ESP_OK: Success
- ESP_FAIL: Event queue for the loop full
- ESP_ERR_INVALID_ARG: Invalid combination of event base and event id, data size of more than 4 bytes
- Others: Fail

Parameters

- `[in] event_loop`: the event loop to post to, must not be NULL
- `[in] event_base`: the event base that identifies the event
- `[in] event_id`: the event id that identifies the event
- `[in] event_data`: the data, specific to the event occurence, that gets passed to the handler
- `[in] event_data_size`: the size of the event data
- `[out] task_unblocked`: an optional parameter (can be NULL) which indicates that an event task with higher priority than currently running task has been unblocked by the posted event; a context switch should be requested before the interrupt is existed.

esp_err_t **esp_event_dump**(FILE **file*)

Dumps statistics of all event loops.

Dumps event loop info in the format:

```
      event loop
          handler
          handler
          ...
      event loop
          handler
          handler
          ...

  where:
```

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```
 event loop
     format: address,name rx:total_recieved dr:total_dropped
     where:
         address - memory address of the event loop
         name - name of the event loop, 'none' if no dedicated task
         total_recieved - number of successfully posted events
         total_dropped - number of events unsuccessfully posted due to queue␣
↪being full

 handler
     format: address ev:base,id inv:total_invoked run:total_runtime
     where:
         address - address of the handler function
         base,id - the event specified by event base and id this handler␣
↪executes
         total_invoked - number of times this handler has been invoked
         total_runtime - total amount of time used for invoking this handler
```

Note this function is a noop when CONFIG_ESP_EVENT_LOOP_PROFILING is disabled

Return

- ESP_OK: Success
- ESP_ERR_NO_MEM: Cannot allocate memory for event loops list
- Others: Fail

Parameters

- `[in] file`: the file stream to output to

Structures

struct esp_event_loop_args_t

Configuration for creating event loops.

Public Members

int32_t **queue_size**

size of the event loop queue

const char ***task_name**

name of the event loop task; if NULL, a dedicated task is not created for event loop

UBaseType_t **task_priority**

priority of the event loop task, ignored if task name is NULL

uint32_t **task_stack_size**

stack size of the event loop task, ignored if task name is NULL

BaseType_t **task_core_id**

core to which the event loop task is pinned to, ignored if task name is NULL

Header File

- esp_event/include/esp_event_base.h

Macros

ESP_EVENT_DECLARE_BASE(id)

ESP_EVENT_DEFINE_BASE(id)

ESP_EVENT_ANY_BASE

register handler for any event base

ESP_EVENT_ANY_ID
register handler for any event id

Type Definitions

typedef const char ***esp_event_base_t**
unique pointer to a subsystem that exposes events

typedef void ***esp_event_loop_handle_t**
a number that identifies an event with respect to a base

typedef void (***esp_event_handler_t**)(void *event_handler_arg, *esp_event_base_t* event_base, int32_t event_id, void *event_data)
function called when an event is posted to the queue

typedef void ***esp_event_handler_instance_t**
context identifying an instance of a registered event handler

Related Documents

Legacy event loop

API Reference

Header File

- esp_event/include/esp_event_legacy.h

Functions

esp_err_t **esp_event_send**(*system_event_t *event*)
Send a event to event task.

Other task/modules, such as the tcpip_adapter, can call this API to send an event to event task

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

Return ESP_OK : succeed
Return others : fail
Parameters
- `event`: Event to send

esp_err_t **esp_event_send_internal**(*esp_event_base_t event_base*, int32_t *event_id*, void *event_data*, size_t *event_data_size*, TickType_t *ticks_to_wait*)
Send a event to event task.

Other task/modules, such as the tcpip_adapter, can call this API to send an event to event task

Note This API is used by WiFi Driver only.

Return ESP_OK : succeed
Return others : fail
Parameters
- `[in] event_base`: the event base that identifies the event
- `[in] event_id`: the event id that identifies the event
- `[in] event_data`: the data, specific to the event occurence, that gets passed to the handler
- `[in] event_data_size`: the size of the event data
- `[in] ticks_to_wait`: number of ticks to block on a full event queue

esp_err_t **esp_event_process_default**(*system_event_t *event*)
Default event handler for system events.

This function performs default handling of system events. When using esp_event_loop APIs, it is called automatically before invoking the user-provided callback function.

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

Applications which implement a custom event loop must call this function as part of event processing.

Return ESP_OK if an event was handled successfully

Parameters

- `event`: pointer to event to be handled

void **esp_event_set_default_eth_handlers**(void)

Install default event handlers for Ethernet interface.

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

void **esp_event_set_default_wifi_handlers**(void)

Install default event handlers for Wi-Fi interfaces (station and AP)

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

esp_err_t **esp_event_loop_init**(*system_event_cb_t cb*, void **ctx*)

Initialize event loop.

Create the event handler and task

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

Return

- ESP_OK: succeed
- others: fail

Parameters

- `cb`: application specified event callback, it can be modified by call esp_event_set_cb
- `ctx`: reserved for user

system_event_cb_t **esp_event_loop_set_cb**(*system_event_cb_t cb*, void **ctx*)

Set application specified event callback function.

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

Attention 1. If cb is NULL, means application don't need to handle If cb is not NULL, it will be call when an event is received, after the default event callback is completed

Return old callback

Parameters

- `cb`: application callback function
- `ctx`: argument to be passed to callback

Unions

union system_event_info_t

#include <esp_event_legacy.h> Union of all possible system_event argument structures

Public Members

system_event_sta_connected_t **connected**

ESP32 station connected to AP

system_event_sta_disconnected_t **disconnected**

ESP32 station disconnected to AP

system_event_sta_scan_done_t **scan_done**

ESP32 station scan (APs) done

system_event_sta_authmode_change_t **auth_change**

the auth mode of AP ESP32 station connected to changed

system_event_sta_got_ip_t **got_ip**

ESP32 station got IP, first time got IP or when IP is changed

system_event_sta_wps_er_pin_t **sta_er_pin**
ESP32 station WPS enrollee mode PIN code received

system_event_sta_wps_fail_reason_t **sta_er_fail_reason**
ESP32 station WPS enrollee mode failed reason code received

system_event_sta_wps_er_success_t **sta_er_success**
ESP32 station WPS enrollee success

system_event_ap_staconnected_t **sta_connected**
a station connected to ESP32 soft-AP

system_event_ap_stadisconnected_t **sta_disconnected**
a station disconnected to ESP32 soft-AP

system_event_ap_probe_req_rx_t **ap_probereqrecved**
ESP32 soft-AP receive probe request packet

system_event_ap_staipassigned_t **ap_staipassigned**
ESP32 soft-AP assign an IP to the station

system_event_got_ip6_t **got_ip6**
ESP32 station or ap or ethernet ipv6 addr state change to preferred

Structures

struct system_event_t
Event, as a tagged enum

Public Members

system_event_id_t **event_id**
event ID

system_event_info_t **event_info**
event information

Macros

SYSTEM_EVENT_AP_STA_GOT_IP6

Type Definitions

typedef *wifi_event_sta_wps_fail_reason_t* **system_event_sta_wps_fail_reason_t**
Argument structure of SYSTEM_EVENT_STA_WPS_ER_FAILED event

typedef *wifi_event_sta_scan_done_t* **system_event_sta_scan_done_t**
Argument structure of SYSTEM_EVENT_SCAN_DONE event

typedef *wifi_event_sta_connected_t* **system_event_sta_connected_t**
Argument structure of SYSTEM_EVENT_STA_CONNECTED event

typedef *wifi_event_sta_disconnected_t* **system_event_sta_disconnected_t**
Argument structure of SYSTEM_EVENT_STA_DISCONNECTED event

typedef *wifi_event_sta_authmode_change_t* **system_event_sta_authmode_change_t**
Argument structure of SYSTEM_EVENT_STA_AUTHMODE_CHANGE event

typedef *wifi_event_sta_wps_er_pin_t* **system_event_sta_wps_er_pin_t**
Argument structure of SYSTEM_EVENT_STA_WPS_ER_PIN event

typedef *wifi_event_sta_wps_er_success_t* **system_event_sta_wps_er_success_t**
Argument structure of SYSTEM_EVENT_STA_WPS_ER_PIN event

typedef *wifi_event_ap_staconnected_t* **system_event_ap_staconnected_t**
Argument structure of event

typedef *wifi_event_ap_stadisconnected_t* **system_event_ap_stadisconnected_t**
Argument structure of event

typedef *wifi_event_ap_probe_req_rx_t* **system_event_ap_probe_req_rx_t**
Argument structure of event

typedef ip_event_ap_staipassigned_t **system_event_ap_staipassigned_t**
Argument structure of event

typedef ip_event_got_ip_t **system_event_sta_got_ip_t**
Argument structure of event

typedef ip_event_got_ip6_t **system_event_got_ip6_t**
Argument structure of event

typedef *esp_err_t* (***system_event_handler_t**) (*esp_event_base_t* event_base, int32_t event_id, void *event_data, size_t event_data_size, TickType_t ticks_to_wait)
Event handler function type

typedef *esp_err_t* (***system_event_cb_t**) (void *ctx, *system_event_t* *event)
Application specified event callback function.

Note This API is part of the legacy event system. New code should use event library API in esp_event.h

Return
- ESP_OK: succeed
- others: fail

Parameters
- `ctx`: reserved for user
- `event`: event type defined in this file

Enumerations

enum system_event_id_t
System event types enumeration

Values:

SYSTEM_EVENT_WIFI_READY = 0
ESP32 WiFi ready

SYSTEM_EVENT_SCAN_DONE
ESP32 finish scanning AP

SYSTEM_EVENT_STA_START
ESP32 station start

SYSTEM_EVENT_STA_STOP
ESP32 station stop

SYSTEM_EVENT_STA_CONNECTED
ESP32 station connected to AP

SYSTEM_EVENT_STA_DISCONNECTED
ESP32 station disconnected from AP

SYSTEM_EVENT_STA_AUTHMODE_CHANGE
the auth mode of AP connected by ESP32 station changed

SYSTEM_EVENT_STA_GOT_IP
ESP32 station got IP from connected AP

SYSTEM_EVENT_STA_LOST_IP
ESP32 station lost IP and the IP is reset to 0

SYSTEM_EVENT_STA_WPS_ER_SUCCESS
ESP32 station wps succeeds in enrollee mode

SYSTEM_EVENT_STA_WPS_ER_FAILED
ESP32 station wps fails in enrollee mode

SYSTEM_EVENT_STA_WPS_ER_TIMEOUT
ESP32 station wps timeout in enrollee mode

SYSTEM_EVENT_STA_WPS_ER_PIN
ESP32 station wps pin code in enrollee mode

SYSTEM_EVENT_STA_WPS_ER_PBC_OVERLAP
ESP32 station wps overlap in enrollee mode

SYSTEM_EVENT_AP_START
ESP32 soft-AP start

SYSTEM_EVENT_AP_STOP
ESP32 soft-AP stop

SYSTEM_EVENT_AP_STACONNECTED
a station connected to ESP32 soft-AP

SYSTEM_EVENT_AP_STADISCONNECTED
a station disconnected from ESP32 soft-AP

SYSTEM_EVENT_AP_STAIPASSIGNED
ESP32 soft-AP assign an IP to a connected station

SYSTEM_EVENT_AP_PROBEREQRECVED
Receive probe request packet in soft-AP interface

SYSTEM_EVENT_STA_BEACON_TIMEOUT
ESP32 station beacon timeout

SYSTEM_EVENT_GOT_IP6
ESP32 station or ap or ethernet interface v6IP addr is preferred

SYSTEM_EVENT_ETH_START
ESP32 ethernet start

SYSTEM_EVENT_ETH_STOP
ESP32 ethernet stop

SYSTEM_EVENT_ETH_CONNECTED
ESP32 ethernet phy link up

SYSTEM_EVENT_ETH_DISCONNECTED
ESP32 ethernet phy link down

SYSTEM_EVENT_ETH_GOT_IP
ESP32 ethernet got IP from connected AP

SYSTEM_EVENT_MAX
Number of members in this enum

2.6.9 FreeRTOS

Overview

This section contains documentation of FreeRTOS types, functions, and macros. It is automatically generated from FreeRTOS header files.

注解： ESP-IDF FreeRTOS is based on the Xtensa port of FreeRTOS v8.2.0, however some functions of FreeRTOS v9.0.0 have been backported. See the *Backported Features* for more information.

For more information about FreeRTOS features specific to ESP-IDF, see *ESP-IDF FreeRTOS SMP Changes* and *ESP-IDF FreeRTOS Additions*.

Task API

Header File

- freertos/include/freertos/task.h

Functions

BaseType_t **xTaskCreatePinnedToCore**(TaskFunction_t *pvTaskCode*, **const** char ***const** *pcName*, **const** uint32_t *usStackDepth*, void ***const** *pvParameters*, UBaseType_t *uxPriority*, *TaskHandle_t* ***const** *pvCreatedTask*, **const** BaseType_t *xCoreID*)

Create a new task with a specified affinity.

This function is similar to xTaskCreate, but allows setting task affinity in SMP system.

Return pdPASS if the task was successfully created and added to a ready list, otherwise an error code defined in the file projdefs.h

Parameters

- `pvTaskCode`: Pointer to the task entry function. Tasks must be implemented to never return (i.e. continuous loop).
- `pcName`: A descriptive name for the task. This is mainly used to facilitate debugging. Max length defined by configMAX_TASK_NAME_LEN - default is 16.
- `usStackDepth`: The size of the task stack specified as the number of bytes. Note that this differs from vanilla FreeRTOS.
- `pvParameters`: Pointer that will be used as the parameter for the task being created.
- `uxPriority`: The priority at which the task should run. Systems that include MPU support can optionally create tasks in a privileged (system) mode by setting bit portPRIVILEGE_BIT of the priority parameter. For example, to create a privileged task at priority 2 the uxPriority parameter should be set to (2 | portPRIVILEGE_BIT).
- `pvCreatedTask`: Used to pass back a handle by which the created task can be referenced.
- `xCoreID`: If the value is tskNO_AFFINITY, the created task is not pinned to any CPU, and the scheduler can run it on any core available. Values 0 or 1 indicate the index number of the CPU which the task should be pinned to. Specifying values larger than (portNUM_PROCESSORS - 1) will cause the function to fail.

static BaseType_t **xTaskCreate**(TaskFunction_t *pvTaskCode*, **const** char ***const** *pcName*, **const** uint32_t *usStackDepth*, void ***const** *pvParameters*, UBaseType_t *uxPriority*, *TaskHandle_t* ***const** *pvCreatedTask*)

Create a new task and add it to the list of tasks that are ready to run.

Internally, within the FreeRTOS implementation, tasks use two blocks of memory. The first block is used to hold the task's data structures. The second block is used by the task as its stack. If a task is created using xTaskCreate() then both blocks of memory are automatically dynamically allocated inside the xTaskCreate() function. (see http://www.freertos.org/a00111.html). If a task is created using xTaskCreateStatic() then the application writer must provide the required memory. xTaskCreateStatic() therefore allows a task to be created without using any dynamic memory allocation.

See xTaskCreateStatic() for a version that does not use any dynamic memory allocation.

xTaskCreate() can only be used to create a task that has unrestricted access to the entire microcontroller memory map. Systems that include MPU support can alternatively create an MPU constrained task using xTaskCreateRestricted().

Example usage:

```
// Task to be created.
void vTaskCode( void * pvParameters )
{
```

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```
 for( ;; )
 {
     // Task code goes here.
 }
}

// Function that creates a task.
void vOtherFunction( void )
{
static uint8_t ucParameterToPass;
TaskHandle_t xHandle = NULL;

 // Create the task, storing the handle.  Note that the passed parameter␣
↪ucParameterToPass
 // must exist for the lifetime of the task, so in this case is declared␣
↪static.  If it was just an
 // an automatic stack variable it might no longer exist, or at least have␣
↪been corrupted, by the time
 // the new task attempts to access it.
 xTaskCreate( vTaskCode, "NAME", STACK_SIZE, &ucParameterToPass, tskIDLE_
↪PRIORITY, &xHandle );
    configASSERT( xHandle );

 // Use the handle to delete the task.
    if( xHandle != NULL )
    {
     vTaskDelete( xHandle );
    }
}
```

Return pdPASS if the task was successfully created and added to a ready list, otherwise an error code defined in the file projdefs.h

Note If program uses thread local variables (ones specified with "__thread" keyword) then storage for them will be allocated on the task's stack.

Parameters

- `pvTaskCode`: Pointer to the task entry function. Tasks must be implemented to never return (i.e. continuous loop).
- `pcName`: A descriptive name for the task. This is mainly used to facilitate debugging. Max length defined by configMAX_TASK_NAME_LEN - default is 16.
- `usStackDepth`: The size of the task stack specified as the number of bytes. Note that this differs from vanilla FreeRTOS.
- `pvParameters`: Pointer that will be used as the parameter for the task being created.
- `uxPriority`: The priority at which the task should run. Systems that include MPU support can optionally create tasks in a privileged (system) mode by setting bit portPRIVILEGE_BIT of the priority parameter. For example, to create a privileged task at priority 2 the uxPriority parameter should be set to (2 | portPRIVILEGE_BIT).
- `pvCreatedTask`: Used to pass back a handle by which the created task can be referenced.

TaskHandle_t **xTaskCreateStaticPinnedToCore**(TaskFunction_t *pvTaskCode*, **const** char ***const** *pcName*, **const** uint32_t *ulStackDepth*, void ***const** *pvParameters*, UBaseType_t *uxPriority*, StackType_t ***const** *pxStackBuffer*, StaticTask_t ***const** *pxTaskBuffer*, **const** BaseType_t *xCoreID*)

Create a new task with a specified affinity.

This function is similar to xTaskCreateStatic, but allows specifying task affinity in an SMP system.

Return If neither pxStackBuffer or pxTaskBuffer are NULL, then the task will be created and a task handle will be returned by which the created task can be referenced. If either pxStackBuffer or pxTaskBuffer are NULL then the task will not be created and NULL is returned.

Parameters

- `pvTaskCode`: Pointer to the task entry function. Tasks must be implemented to never return (i.e. continuous loop).
- `pcName`: A descriptive name for the task. This is mainly used to facilitate debugging. The maximum length of the string is defined by configMAX_TASK_NAME_LEN in FreeRTOSConfig.h.
- `ulStackDepth`: The size of the task stack specified as the number of bytes. Note that this differs from vanilla FreeRTOS.
- `pvParameters`: Pointer that will be used as the parameter for the task being created.
- `uxPriority`: The priority at which the task will run.
- `pxStackBuffer`: Must point to a StackType_t array that has at least ulStackDepth indexes - the array will then be used as the task's stack, removing the need for the stack to be allocated dynamically.
- `pxTaskBuffer`: Must point to a variable of type StaticTask_t, which will then be used to hold the task's data structures, removing the need for the memory to be allocated dynamically.
- `xCoreID`: If the value is tskNO_AFFINITY, the created task is not pinned to any CPU, and the scheduler can run it on any core available. Values 0 or 1 indicate the index number of the CPU which the task should be pinned to. Specifying values larger than (portNUM_PROCESSORS - 1) will cause the function to fail.

static *TaskHandle_t* **xTaskCreateStatic**(TaskFunction_t *pvTaskCode*, **const** char ***const** *pcName*, **const** uint32_t *ulStackDepth*, void ***const** *pvParameters*, UBaseType_t *uxPriority*, StackType_t ***const** *pxStackBuffer*, StaticTask_t ***const** *pxTaskBuffer*)

Create a new task and add it to the list of tasks that are ready to run.

Internally, within the FreeRTOS implementation, tasks use two blocks of memory. The first block is used to hold the task's data structures. The second block is used by the task as its stack. If a task is created using xTaskCreate() then both blocks of memory are automatically dynamically allocated inside the xTaskCreate() function. (see http://www.freertos.org/a00111.html). If a task is created using xTaskCreateStatic() then the application writer must provide the required memory. xTaskCreateStatic() therefore allows a task to be created without using any dynamic memory allocation.

Example usage:

```
// Dimensions the buffer that the task being created will use as its stack.
// NOTE:  This is the number of bytes the stack will hold, not the number of
// words as found in vanilla FreeRTOS.
#define STACK_SIZE 200

// Structure that will hold the TCB of the task being created.
StaticTask_t xTaskBuffer;

// Buffer that the task being created will use as its stack.  Note this is
// an array of StackType_t variables.  The size of StackType_t is dependent on
// the RTOS port.
StackType_t xStack[ STACK_SIZE ];

// Function that implements the task being created.
void vTaskCode( void * pvParameters )
{
    // The parameter value is expected to be 1 as 1 is passed in the
    // pvParameters value in the call to xTaskCreateStatic().
    configASSERT( ( uint32_t ) pvParameters == 1UL );

    for( ;; )
    {
        // Task code goes here.
    }
}
```

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```
// Function that creates a task.
void vOtherFunction( void )
{
    TaskHandle_t xHandle = NULL;

    // Create the task without using any dynamic memory allocation.
    xHandle = xTaskCreateStatic(
                  vTaskCode,       // Function that implements the task.
                  "NAME",          // Text name for the task.
                  STACK_SIZE,      // Stack size in bytes, not words.
                  ( void * ) 1,    // Parameter passed into the task.
                  tskIDLE_PRIORITY,// Priority at which the task is created.
                  xStack,          // Array to use as the task's stack.
                  &xTaskBuffer );  // Variable to hold the task's data␣
↪structure.

    // puxStackBuffer and pxTaskBuffer were not NULL, so the task will have
    // been created, and xHandle will be the task's handle.  Use the handle
    // to suspend the task.
    vTaskSuspend( xHandle );
}
```

Return If neither pxStackBuffer or pxTaskBuffer are NULL, then the task will be created and a task handle will be returned by which the created task can be referenced. If either pxStackBuffer or pxTaskBuffer are NULL then the task will not be created and NULL is returned.

Note If program uses thread local variables (ones specified with "__thread" keyword) then storage for them will be allocated on the task's stack.

Parameters

- `pvTaskCode`: Pointer to the task entry function. Tasks must be implemented to never return (i.e. continuous loop).
- `pcName`: A descriptive name for the task. This is mainly used to facilitate debugging. The maximum length of the string is defined by configMAX_TASK_NAME_LEN in FreeRTOSConfig.h.
- `ulStackDepth`: The size of the task stack specified as the number of bytes. Note that this differs from vanilla FreeRTOS.
- `pvParameters`: Pointer that will be used as the parameter for the task being created.
- `uxPriority`: The priority at which the task will run.
- `pxStackBuffer`: Must point to a StackType_t array that has at least ulStackDepth indexes - the array will then be used as the task's stack, removing the need for the stack to be allocated dynamically.
- `pxTaskBuffer`: Must point to a variable of type StaticTask_t, which will then be used to hold the task's data structures, removing the need for the memory to be allocated dynamically.

void **vTaskDelete**(*TaskHandle_t xTaskToDelete*)

Remove a task from the RTOS real time kernel's management.

The task being deleted will be removed from all ready, blocked, suspended and event lists.

INCLUDE_vTaskDelete must be defined as 1 for this function to be available. See the configuration section for more information.

See the demo application file death.c for sample code that utilises vTaskDelete ().

Note The idle task is responsible for freeing the kernel allocated memory from tasks that have been deleted. It is therefore important that the idle task is not starved of microcontroller processing time if your application makes any calls to vTaskDelete (). Memory allocated by the task code is not automatically freed, and should be freed before the task is deleted.

Example usage:

```
void vOtherFunction( void )
{
```

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```
TaskHandle_t xHandle;

    // Create the task, storing the handle.
    xTaskCreate( vTaskCode, "NAME", STACK_SIZE, NULL, tskIDLE_PRIORITY, &
↪xHandle );

    // Use the handle to delete the task.
    vTaskDelete( xHandle );
}
```

Parameters

- `xTaskToDelete`: The handle of the task to be deleted. Passing NULL will cause the calling task to be deleted.

void **vTaskDelay**(**const** TickType_t *xTicksToDelay*)

Delay a task for a given number of ticks.

The actual time that the task remains blocked depends on the tick rate. The constant portTICK_PERIOD_MS can be used to calculate real time from the tick rate - with the resolution of one tick period.

INCLUDE_vTaskDelay must be defined as 1 for this function to be available. See the configuration section for more information.

vTaskDelay() specifies a time at which the task wishes to unblock relative to the time at which vTaskDelay() is called. For example, specifying a block period of 100 ticks will cause the task to unblock 100 ticks after vTaskDelay() is called. vTaskDelay() does not therefore provide a good method of controlling the frequency of a periodic task as the path taken through the code, as well as other task and interrupt activity, will effect the frequency at which vTaskDelay() gets called and therefore the time at which the task next executes. See vTaskDelayUntil() for an alternative API function designed to facilitate fixed frequency execution. It does this by specifying an absolute time (rather than a relative time) at which the calling task should unblock.

Example usage:

```
void vTaskFunction( void * pvParameters )
{
// Block for 500ms.
const TickType_t xDelay = 500 / portTICK_PERIOD_MS;

    for( ;; )
    {
        // Simply toggle the LED every 500ms, blocking between each toggle.
        vToggleLED();
        vTaskDelay( xDelay );
    }
}
```

Parameters

- `xTicksToDelay`: The amount of time, in tick periods, that the calling task should block.

void **vTaskDelayUntil**(TickType_t *__const__ *pxPreviousWakeTime*, **const** TickType_t *xTimeIncrement*)

Delay a task until a specified time.

INCLUDE_vTaskDelayUntil must be defined as 1 for this function to be available. See the configuration section for more information.

This function can be used by periodic tasks to ensure a constant execution frequency.

This function differs from vTaskDelay () in one important aspect: vTaskDelay () will cause a task to block for the specified number of ticks from the time vTaskDelay () is called. It is therefore difficult to use vTaskDelay () by itself to generate a fixed execution frequency as the time between a task starting to execute and that task calling vTaskDelay () may not be fixed [the task may take a different path though the code between calls, or may get interrupted or preempted a different number of times each time it executes].

Whereas vTaskDelay () specifies a wake time relative to the time at which the function is called, vTaskDelayUntil () specifies the absolute (exact) time at which it wishes to unblock.

The constant portTICK_PERIOD_MS can be used to calculate real time from the tick rate - with the resolution of one tick period.

Example usage:

```
// Perform an action every 10 ticks.
void vTaskFunction( void * pvParameters )
{
TickType_t xLastWakeTime;
const TickType_t xFrequency = 10;

    // Initialise the xLastWakeTime variable with the current time.
    xLastWakeTime = xTaskGetTickCount ();
    for( ;; )
    {
        // Wait for the next cycle.
        vTaskDelayUntil( &xLastWakeTime, xFrequency );

        // Perform action here.
    }
}
```

Parameters

- `pxPreviousWakeTime`: Pointer to a variable that holds the time at which the task was last unblocked. The variable must be initialised with the current time prior to its first use (see the example below). Following this the variable is automatically updated within vTaskDelayUntil ().
- `xTimeIncrement`: The cycle time period. The task will be unblocked at time *pxPreviousWakeTime + xTimeIncrement. Calling vTaskDelayUntil with the same xTimeIncrement parameter value will cause the task to execute with a fixed interface period.

UBaseType_t **uxTaskPriorityGet** (*TaskHandle_t xTask*)

Obtain the priority of any task.

INCLUDE_uxTaskPriorityGet must be defined as 1 for this function to be available. See the configuration section for more information.

Example usage:

```
 void vAFunction( void )
 {
 TaskHandle_t xHandle;

  // Create a task, storing the handle.
  xTaskCreate( vTaskCode, "NAME", STACK_SIZE, NULL, tskIDLE_PRIORITY, &xHandle␣
↪);

  // ...

  // Use the handle to obtain the priority of the created task.
  // It was created with tskIDLE_PRIORITY, but may have changed
  // it itself.
  if( uxTaskPriorityGet( xHandle ) != tskIDLE_PRIORITY )
  {
      // The task has changed it's priority.
  }

  // ...

  // Is our priority higher than the created task?
  if( uxTaskPriorityGet( xHandle ) < uxTaskPriorityGet( NULL ) )
```

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```
  {
      // Our priority (obtained using NULL handle) is higher.
  }
}
```

Return The priority of xTask.

Parameters

- `xTask`: Handle of the task to be queried. Passing a NULL handle results in the priority of the calling task being returned.

UBaseType_t **uxTaskPriorityGetFromISR** (*TaskHandle_t xTask*)

A version of uxTaskPriorityGet() that can be used from an ISR.

Return The priority of xTask.

Parameters

- `xTask`: Handle of the task to be queried. Passing a NULL handle results in the priority of the calling task being returned.

eTaskState **eTaskGetState** (*TaskHandle_t xTask*)

Obtain the state of any task.

States are encoded by the eTaskState enumerated type.

INCLUDE_eTaskGetState must be defined as 1 for this function to be available. See the configuration section for more information.

Return The state of xTask at the time the function was called. Note the state of the task might change between the function being called, and the functions return value being tested by the calling task.

Parameters

- `xTask`: Handle of the task to be queried.

void **vTaskPrioritySet** (*TaskHandle_t xTask*, UBaseType_t *uxNewPriority*)

Set the priority of any task.

INCLUDE_vTaskPrioritySet must be defined as 1 for this function to be available. See the configuration section for more information.

A context switch will occur before the function returns if the priority being set is higher than the currently executing task.

Example usage:

```
void vAFunction( void )
{
TaskHandle_t xHandle;

 // Create a task, storing the handle.
 xTaskCreate( vTaskCode, "NAME", STACK_SIZE, NULL, tskIDLE_PRIORITY, &xHandle␣
↪);

 // ...

 // Use the handle to raise the priority of the created task.
 vTaskPrioritySet( xHandle, tskIDLE_PRIORITY + 1 );

 // ...

 // Use a NULL handle to raise our priority to the same value.
 vTaskPrioritySet( NULL, tskIDLE_PRIORITY + 1 );
}
```

Parameters

- `xTask`: Handle to the task for which the priority is being set. Passing a NULL handle results in the priority of the calling task being set.

- `uxNewPriority`: The priority to which the task will be set.

void **vTaskSuspend** (*TaskHandle_t xTaskToSuspend*)

Suspend a task.

INCLUDE_vTaskSuspend must be defined as 1 for this function to be available. See the configuration section for more information.

When suspended, a task will never get any microcontroller processing time, no matter what its priority.

Calls to vTaskSuspend are not accumulative - i.e. calling vTaskSuspend () twice on the same task still only requires one call to vTaskResume () to ready the suspended task.

Example usage:

```
void vAFunction( void )
{
TaskHandle_t xHandle;

 // Create a task, storing the handle.
 xTaskCreate( vTaskCode, "NAME", STACK_SIZE, NULL, tskIDLE_PRIORITY, &xHandle␣
↪);

 // ...

 // Use the handle to suspend the created task.
 vTaskSuspend( xHandle );

 // ...

 // The created task will not run during this period, unless
 // another task calls vTaskResume( xHandle ).

 //...


 // Suspend ourselves.
 vTaskSuspend( NULL );

 // We cannot get here unless another task calls vTaskResume
 // with our handle as the parameter.
}
```

Parameters

- `xTaskToSuspend`: Handle to the task being suspended. Passing a NULL handle will cause the calling task to be suspended.

void **vTaskResume** (*TaskHandle_t xTaskToResume*)

Resumes a suspended task.

INCLUDE_vTaskSuspend must be defined as 1 for this function to be available. See the configuration section for more information.

A task that has been suspended by one or more calls to vTaskSuspend () will be made available for running again by a single call to vTaskResume ().

Example usage:

```
void vAFunction( void )
{
TaskHandle_t xHandle;

 // Create a task, storing the handle.
 xTaskCreate( vTaskCode, "NAME", STACK_SIZE, NULL, tskIDLE_PRIORITY, &xHandle␣
↪);
```

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```
 // ...

 // Use the handle to suspend the created task.
 vTaskSuspend( xHandle );

 // ...

 // The created task will not run during this period, unless
 // another task calls vTaskResume( xHandle ).

 //...


 // Resume the suspended task ourselves.
 vTaskResume( xHandle );

 // The created task will once again get microcontroller processing
 // time in accordance with its priority within the system.
}
```

Parameters

- `xTaskToResume`: Handle to the task being readied.

BaseType_t **`xTaskResumeFromISR`** (*TaskHandle_t xTaskToResume*)

An implementation of vTaskResume() that can be called from within an ISR.

INCLUDE_xTaskResumeFromISR must be defined as 1 for this function to be available. See the configuration section for more information.

A task that has been suspended by one or more calls to vTaskSuspend () will be made available for running again by a single call to xTaskResumeFromISR ().

xTaskResumeFromISR() should not be used to synchronise a task with an interrupt if there is a chance that the interrupt could arrive prior to the task being suspended - as this can lead to interrupts being missed. Use of a semaphore as a synchronisation mechanism would avoid this eventuality.

Return pdTRUE if resuming the task should result in a context switch, otherwise pdFALSE. This is used by the ISR to determine if a context switch may be required following the ISR.

Parameters

- `xTaskToResume`: Handle to the task being readied.

void **`vTaskSuspendAll`** (void)

Suspends the scheduler without disabling interrupts.

Context switches will not occur while the scheduler is suspended.

After calling vTaskSuspendAll () the calling task will continue to execute without risk of being swapped out until a call to xTaskResumeAll () has been made.

API functions that have the potential to cause a context switch (for example, vTaskDelayUntil(), xQueueSend(), etc.) must not be called while the scheduler is suspended.

Example usage:

```
void vTask1( void * pvParameters )
{
 for( ;; )
 {
     // Task code goes here.

     // ...

     // At some point the task wants to perform a long operation during
```

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```
    // which it does not want to get swapped out.  It cannot use
    // taskENTER_CRITICAL ()/taskEXIT_CRITICAL () as the length of the
    // operation may cause interrupts to be missed - including the
    // ticks.

    // Prevent the real time kernel swapping out the task.
    vTaskSuspendAll ();

    // Perform the operation here.  There is no need to use critical
    // sections as we have all the microcontroller processing time.
    // During this time interrupts will still operate and the kernel
    // tick count will be maintained.

    // ...

    // The operation is complete.  Restart the kernel.
    xTaskResumeAll ();
 }
}
```

BaseType_t **xTaskResumeAll** (void)

Resumes scheduler activity after it was suspended by a call to vTaskSuspendAll().

xTaskResumeAll() only resumes the scheduler. It does not unsuspend tasks that were previously suspended by a call to vTaskSuspend().

Example usage:

```
void vTask1( void * pvParameters )
{
 for( ;; )
 {
    // Task code goes here.

    // ...

    // At some point the task wants to perform a long operation during
    // which it does not want to get swapped out.  It cannot use
    // taskENTER_CRITICAL ()/taskEXIT_CRITICAL () as the length of the
    // operation may cause interrupts to be missed - including the
    // ticks.

    // Prevent the real time kernel swapping out the task.
    vTaskSuspendAll ();

    // Perform the operation here.  There is no need to use critical
    // sections as we have all the microcontroller processing time.
    // During this time interrupts will still operate and the real
    // time kernel tick count will be maintained.

    // ...

    // The operation is complete.  Restart the kernel.  We want to force
    // a context switch - but there is no point if resuming the scheduler
    // caused a context switch already.
    if( !xTaskResumeAll () )
    {
         taskYIELD ();
    }
 }
}
```

Return If resuming the scheduler caused a context switch then pdTRUE is returned, otherwise pdFALSE is returned.

TickType_t **xTaskGetTickCount** (void)
Get tick count

Return The count of ticks since vTaskStartScheduler was called.

TickType_t **xTaskGetTickCountFromISR** (void)
Get tick count from ISR

This is a version of xTaskGetTickCount() that is safe to be called from an ISR - provided that TickType_t is the natural word size of the microcontroller being used or interrupt nesting is either not supported or not being used.

Return The count of ticks since vTaskStartScheduler was called.

UBaseType_t **uxTaskGetNumberOfTasks** (void)
Get current number of tasks

Return The number of tasks that the real time kernel is currently managing. This includes all ready, blocked and suspended tasks. A task that has been deleted but not yet freed by the idle task will also be included in the count.

char ***pcTaskGetTaskName** (*TaskHandle_t xTaskToQuery*)
Get task name

Return The text (human readable) name of the task referenced by the handle xTaskToQuery. A task can query its own name by either passing in its own handle, or by setting xTaskToQuery to NULL. INCLUDE_pcTaskGetTaskName must be set to 1 in FreeRTOSConfig.h for pcTaskGetTaskName() to be available.

UBaseType_t **uxTaskGetStackHighWaterMark** (*TaskHandle_t xTask*)
Returns the high water mark of the stack associated with xTask.

INCLUDE_uxTaskGetStackHighWaterMark must be set to 1 in FreeRTOSConfig.h for this function to be available.

High water mark is the minimum free stack space there has been (in bytes rather than words as found in vanilla FreeRTOS) since the task started. The smaller the returned number the closer the task has come to overflowing its stack.

Return The smallest amount of free stack space there has been (in bytes rather than words as found in vanilla FreeRTOS) since the task referenced by xTask was created.

Parameters

- `xTask`: Handle of the task associated with the stack to be checked. Set xTask to NULL to check the stack of the calling task.

uint8_t ***pxTaskGetStackStart** (*TaskHandle_t xTask*)
Returns the start of the stack associated with xTask.

INCLUDE_pxTaskGetStackStart must be set to 1 in FreeRTOSConfig.h for this function to be available.

Returns the highest stack memory address on architectures where the stack grows down from high memory, and the lowest memory address on architectures where the stack grows up from low memory.

Return A pointer to the start of the stack.

Parameters

- `xTask`: Handle of the task associated with the stack returned. Set xTask to NULL to return the stack of the calling task.

void **vTaskSetApplicationTaskTag** (*TaskHandle_t xTask*, *TaskHookFunction_t pxHookFunction*)
Sets pxHookFunction to be the task hook function used by the task xTask.

Parameters

- `xTask`: Handle of the task to set the hook function for Passing xTask as NULL has the effect of setting the calling tasks hook function.
- `pxHookFunction`: Pointer to the hook function.

TaskHookFunction_t **xTaskGetApplicationTaskTag** (*TaskHandle_t xTask*)

Get the hook function assigned to given task.

Return The pxHookFunction value assigned to the task xTask.

Parameters

- `xTask`: Handle of the task to get the hook function for Passing xTask as NULL has the effect of getting the calling tasks hook function.

void **vTaskSetThreadLocalStoragePointer** (*TaskHandle_t xTaskToSet*, BaseType_t *xIndex*, void **pvValue*)

Set local storage pointer specific to the given task.

Each task contains an array of pointers that is dimensioned by the configNUM_THREAD_LOCAL_STORAGE_POINTERS setting in FreeRTOSConfig.h. The kernel does not use the pointers itself, so the application writer can use the pointers for any purpose they wish.

Parameters

- `xTaskToSet`: Task to set thread local storage pointer for
- `xIndex`: The index of the pointer to set, from 0 to configNUM_THREAD_LOCAL_STORAGE_POINTERS - 1.
- `pvValue`: Pointer value to set.

void ***pvTaskGetThreadLocalStoragePointer** (*TaskHandle_t xTaskToQuery*, BaseType_t *xIndex*)

Get local storage pointer specific to the given task.

Each task contains an array of pointers that is dimensioned by the configNUM_THREAD_LOCAL_STORAGE_POINTERS setting in FreeRTOSConfig.h. The kernel does not use the pointers itself, so the application writer can use the pointers for any purpose they wish.

Return Pointer value

Parameters

- `xTaskToQuery`: Task to get thread local storage pointer for
- `xIndex`: The index of the pointer to get, from 0 to configNUM_THREAD_LOCAL_STORAGE_POINTERS - 1.

void **vTaskSetThreadLocalStoragePointerAndDelCallback** (*TaskHandle_t xTaskToSet*, BaseType_t *xIndex*, void **pvValue*, *TlsDeleteCallbackFunction_t pvDelCallback*)

Set local storage pointer and deletion callback.

Each task contains an array of pointers that is dimensioned by the configNUM_THREAD_LOCAL_STORAGE_POINTERS setting in FreeRTOSConfig.h. The kernel does not use the pointers itself, so the application writer can use the pointers for any purpose they wish.

Local storage pointers set for a task can reference dynamically allocated resources. This function is similar to vTaskSetThreadLocalStoragePointer, but provides a way to release these resources when the task gets deleted. For each pointer, a callback function can be set. This function will be called when task is deleted, with the local storage pointer index and value as arguments.

Parameters

- `xTaskToSet`: Task to set thread local storage pointer for
- `xIndex`: The index of the pointer to set, from 0 to configNUM_THREAD_LOCAL_STORAGE_POINTERS - 1.
- `pvValue`: Pointer value to set.
- `pvDelCallback`: Function to call to dispose of the local storage pointer when the task is deleted.

BaseType_t **xTaskCallApplicationTaskHook** (*TaskHandle_t xTask*, void **pvParameter*)

Calls the hook function associated with xTask. Passing xTask as NULL has the effect of calling the Running tasks (the calling task) hook function.

Parameters

- `xTask`: Handle of the task to call the hook for.

- `pvParameter`: Parameter passed to the hook function for the task to interpret as it wants. The return value is the value returned by the task hook function registered by the user.

TaskHandle_t **xTaskGetIdleTaskHandle**(void)

Get the handle of idle task for the current CPU.

xTaskGetIdleTaskHandle() is only available if INCLUDE_xTaskGetIdleTaskHandle is set to 1 in FreeRTOSConfig.h.

Return The handle of the idle task. It is not valid to call xTaskGetIdleTaskHandle() before the scheduler has been started.

TaskHandle_t **xTaskGetIdleTaskHandleForCPU**(UBaseType_t *cpuid*)

Get the handle of idle task for the given CPU.

xTaskGetIdleTaskHandleForCPU() is only available if INCLUDE_xTaskGetIdleTaskHandle is set to 1 in FreeRTOSConfig.h.

Return Idle task handle of a given cpu. It is not valid to call xTaskGetIdleTaskHandleForCPU() before the scheduler has been started.

Parameters

- `cpuid`: The CPU to get the handle for

UBaseType_t **uxTaskGetSystemState**(*TaskStatus_t* ***const** *pxTaskStatusArray*, **const** UBaseType_t *uxArraySize*, uint32_t ***const** *pulTotalRunTime*)

Get the state of tasks in the system.

configUSE_TRACE_FACILITY must be defined as 1 in FreeRTOSConfig.h for uxTaskGetSystemState() to be available.

uxTaskGetSystemState() populates an TaskStatus_t structure for each task in the system. TaskStatus_t structures contain, among other things, members for the task handle, task name, task priority, task state, and total amount of run time consumed by the task. See the TaskStatus_t structure definition in this file for the full member list.

Example usage:

```
// This example demonstrates how a human readable table of run time stats
// information is generated from raw data provided by uxTaskGetSystemState().
// The human readable table is written to pcWriteBuffer
void vTaskGetRunTimeStats( char *pcWriteBuffer )
{
TaskStatus_t *pxTaskStatusArray;
volatile UBaseType_t uxArraySize, x;
uint32_t ulTotalRunTime, ulStatsAsPercentage;

 // Make sure the write buffer does not contain a string.
 *pcWriteBuffer = 0x00;

 // Take a snapshot of the number of tasks in case it changes while this
 // function is executing.
 uxArraySize = uxTaskGetNumberOfTasks();

 // Allocate a TaskStatus_t structure for each task.  An array could be
 // allocated statically at compile time.
 pxTaskStatusArray = pvPortMalloc( uxArraySize * sizeof( TaskStatus_t ) );

 if( pxTaskStatusArray != NULL )
 {
     // Generate raw status information about each task.
     uxArraySize = uxTaskGetSystemState( pxTaskStatusArray, uxArraySize, &
↪ulTotalRunTime );

     // For percentage calculations.
     ulTotalRunTime /= 100UL;
```

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```
    // Avoid divide by zero errors.
    if( ulTotalRunTime > 0 )
    {
        // For each populated position in the pxTaskStatusArray array,
        // format the raw data as human readable ASCII data
        for( x = 0; x < uxArraySize; x++ )
        {
            // What percentage of the total run time has the task used?
            // This will always be rounded down to the nearest integer.
            // ulTotalRunTimeDiv100 has already been divided by 100.
            ulStatsAsPercentage = pxTaskStatusArray[ x ].ulRunTimeCounter /␣
↪ulTotalRunTime;

            if( ulStatsAsPercentage > 0UL )
            {
                sprintf( pcWriteBuffer, "%s\t\t%lu\t\t%lu%%\r\n",␣
↪pxTaskStatusArray[ x ].pcTaskName, pxTaskStatusArray[ x ].ulRunTimeCounter,␣
↪ulStatsAsPercentage );
            }
            else
            {
                // If the percentage is zero here then the task has
                // consumed less than 1% of the total run time.
                sprintf( pcWriteBuffer, "%s\t\t%lu\t\t<1%%\r\n",␣
↪pxTaskStatusArray[ x ].pcTaskName, pxTaskStatusArray[ x ].ulRunTimeCounter );
            }

            pcWriteBuffer += strlen( ( char * ) pcWriteBuffer );
        }
    }

    // The array is no longer needed, free the memory it consumes.
    vPortFree( pxTaskStatusArray );
 }
}
```

Note This function is intended for debugging use only as its use results in the scheduler remaining suspended for an extended period.

Return The number of TaskStatus_t structures that were populated by uxTaskGetSystemState(). This should equal the number returned by the uxTaskGetNumberOfTasks() API function, but will be zero if the value passed in the uxArraySize parameter was too small.

Parameters

- `pxTaskStatusArray`: A pointer to an array of TaskStatus_t structures. The array must contain at least one TaskStatus_t structure for each task that is under the control of the RTOS. The number of tasks under the control of the RTOS can be determined using the uxTaskGetNumberOfTasks() API function.
- `uxArraySize`: The size of the array pointed to by the pxTaskStatusArray parameter. The size is specified as the number of indexes in the array, or the number of TaskStatus_t structures contained in the array, not by the number of bytes in the array.
- `pulTotalRunTime`: If configGENERATE_RUN_TIME_STATS is set to 1 in FreeRTOSConfig.h then *pulTotalRunTime is set by uxTaskGetSystemState() to the total run time (as defined by the run time stats clock, see http://www.freertos.org/rtos-run-time-stats.html) since the target booted. pulTotalRunTime can be set to NULL to omit the total run time information.

void **vTaskList**(char *_pcWriteBuffer_)

List all the current tasks.

configUSE_TRACE_FACILITY and configUSE_STATS_FORMATTING_FUNCTIONS must both be defined as 1 for this function to be available. See the configuration section of the FreeRTOS.org website for more information.

Lists all the current tasks, along with their current state and stack usage high water mark.

Note This function will disable interrupts for its duration. It is not intended for normal application runtime use but as a debug aid.

Tasks are reported as blocked ('B'), ready ('R'), deleted ('D') or suspended ('S').

vTaskList() calls uxTaskGetSystemState(), then formats part of the uxTaskGetSystemState() output into a human readable table that displays task names, states and stack usage.

Note This function is provided for convenience only, and is used by many of the demo applications. Do not consider it to be part of the scheduler.

vTaskList() has a dependency on the sprintf() C library function that might bloat the code size, use a lot of stack, and provide different results on different platforms. An alternative, tiny, third party, and limited functionality implementation of sprintf() is provided in many of the FreeRTOS/Demo sub-directories in a file called printf-stdarg.c (note printf-stdarg.c does not provide a full snprintf() implementation!).

It is recommended that production systems call uxTaskGetSystemState() directly to get access to raw stats data, rather than indirectly through a call to vTaskList().

Parameters

- `pcWriteBuffer`: A buffer into which the above mentioned details will be written, in ASCII form. This buffer is assumed to be large enough to contain the generated report. Approximately 40 bytes per task should be sufficient.

void **`vTaskGetRunTimeStats`**(char **pcWriteBuffer*)

Get the state of running tasks as a string

configGENERATE_RUN_TIME_STATS and configUSE_STATS_FORMATTING_FUNCTIONS must both be defined as 1 for this function to be available. The application must also then provide definitions for portCONFIGURE_TIMER_FOR_RUN_TIME_STATS() and portGET_RUN_TIME_COUNTER_VALUE() to configure a peripheral timer/counter and return the timers current count value respectively. The counter should be at least 10 times the frequency of the tick count.

Setting configGENERATE_RUN_TIME_STATS to 1 will result in a total accumulated execution time being stored for each task. The resolution of the accumulated time value depends on the frequency of the timer configured by the portCONFIGURE_TIMER_FOR_RUN_TIME_STATS() macro. Calling vTaskGetRunTimeStats() writes the total execution time of each task into a buffer, both as an absolute count value and as a percentage of the total system execution time.

Note This function will disable interrupts for its duration. It is not intended for normal application runtime use but as a debug aid.

vTaskGetRunTimeStats() calls uxTaskGetSystemState(), then formats part of the uxTaskGetSystemState() output into a human readable table that displays the amount of time each task has spent in the Running state in both absolute and percentage terms.

Note This function is provided for convenience only, and is used by many of the demo applications. Do not consider it to be part of the scheduler.

vTaskGetRunTimeStats() has a dependency on the sprintf() C library function that might bloat the code size, use a lot of stack, and provide different results on different platforms. An alternative, tiny, third party, and limited functionality implementation of sprintf() is provided in many of the FreeRTOS/Demo sub-directories in a file called printf-stdarg.c (note printf-stdarg.c does not provide a full snprintf() implementation!).

It is recommended that production systems call uxTaskGetSystemState() directly to get access to raw stats data, rather than indirectly through a call to vTaskGetRunTimeStats().

Parameters

- `pcWriteBuffer`: A buffer into which the execution times will be written, in ASCII form. This buffer is assumed to be large enough to contain the generated report. Approximately 40 bytes per task should be sufficient.

BaseType_t **`xTaskNotify`**(*TaskHandle_t xTaskToNotify*, uint32_t *ulValue*, *eNotifyAction eAction*)

Send task notification.

configUSE_TASK_NOTIFICATIONS must be undefined or defined as 1 for this function to be available.

When configUSE_TASK_NOTIFICATIONS is set to one each task has its own private "notification value", which is a 32-bit unsigned integer (uint32_t).

Events can be sent to a task using an intermediary object. Examples of such objects are queues, semaphores, mutexes and event groups. Task notifications are a method of sending an event directly to a task without the need for such an intermediary object.

A notification sent to a task can optionally perform an action, such as update, overwrite or increment the task's notification value. In that way task notifications can be used to send data to a task, or be used as light weight and fast binary or counting semaphores.

A notification sent to a task will remain pending until it is cleared by the task calling xTaskNotifyWait() or ulTaskNotifyTake(). If the task was already in the Blocked state to wait for a notification when the notification arrives then the task will automatically be removed from the Blocked state (unblocked) and the notification cleared.

A task can use xTaskNotifyWait() to [optionally] block to wait for a notification to be pending, or ulTaskNotifyTake() to [optionally] block to wait for its notification value to have a non-zero value. The task does not consume any CPU time while it is in the Blocked state.

See http://www.FreeRTOS.org/RTOS-task-notifications.html for details.

Return Dependent on the value of eAction. See the description of the eAction parameter.

Parameters

- `xTaskToNotify`: The handle of the task being notified. The handle to a task can be returned from the xTaskCreate() API function used to create the task, and the handle of the currently running task can be obtained by calling xTaskGetCurrentTaskHandle().
- `ulValue`: Data that can be sent with the notification. How the data is used depends on the value of the eAction parameter.
- `eAction`: Specifies how the notification updates the task's notification value, if at all. Valid values for eAction are as follows:
  - eSetBits: The task's notification value is bitwise ORed with ulValue. xTaskNofify() always returns pdPASS in this case.
  - eIncrement: The task's notification value is incremented. ulValue is not used and xTaskNotify() always returns pdPASS in this case.
  - eSetValueWithOverwrite: The task's notification value is set to the value of ulValue, even if the task being notified had not yet processed the previous notification (the task already had a notification pending). xTaskNotify() always returns pdPASS in this case.
  - eSetValueWithoutOverwrite: If the task being notified did not already have a notification pending then the task's notification value is set to ulValue and xTaskNotify() will return pdPASS. If the task being notified already had a notification pending then no action is performed and pdFAIL is returned.
  - eNoAction: The task receives a notification without its notification value being updated. ulValue is not used and xTaskNotify() always returns pdPASS in this case.

BaseType_t **xTaskNotifyFromISR**(*TaskHandle_t xTaskToNotify*, uint32_t *ulValue*, *eNotifyAction eAction*, BaseType_t **pxHigherPriorityTaskWoken*)

Send task notification from an ISR.

configUSE_TASK_NOTIFICATIONS must be undefined or defined as 1 for this function to be available.

When configUSE_TASK_NOTIFICATIONS is set to one each task has its own private "notification value", which is a 32-bit unsigned integer (uint32_t).

A version of xTaskNotify() that can be used from an interrupt service routine (ISR).

Events can be sent to a task using an intermediary object. Examples of such objects are queues, semaphores, mutexes and event groups. Task notifications are a method of sending an event directly to a task without the need for such an intermediary object.

A notification sent to a task can optionally perform an action, such as update, overwrite or increment the task's notification value. In that way task notifications can be used to send data to a task, or be used as light weight and fast binary or counting semaphores.

A notification sent to a task will remain pending until it is cleared by the task calling xTaskNotifyWait() or ulTaskNotifyTake(). If the task was already in the Blocked state to wait for a notification when the notification arrives then the task will automatically be removed from the Blocked state (unblocked) and the notification cleared.

A task can use xTaskNotifyWait() to [optionally] block to wait for a notification to be pending, or ulTaskNotifyTake() to [optionally] block to wait for its notification value to have a non-zero value. The task does not consume any CPU time while it is in the Blocked state.

See http://www.FreeRTOS.org/RTOS-task-notifications.html for details.

Return Dependent on the value of eAction. See the description of the eAction parameter.
Parameters

- `xTaskToNotify`: The handle of the task being notified. The handle to a task can be returned from the xTaskCreate() API function used to create the task, and the handle of the currently running task can be obtained by calling xTaskGetCurrentTaskHandle().
- `ulValue`: Data that can be sent with the notification. How the data is used depends on the value of the eAction parameter.
- `eAction`: Specifies how the notification updates the task ' s notification value, if at all. Valid values for eAction are as follows:
  - eSetBits: The task ' s notification value is bitwise ORed with ulValue. xTaskNofify() always returns pdPASS in this case.
  - eIncrement: The task ' s notification value is incremented. ulValue is not used and xTaskNotify() always returns pdPASS in this case.
  - eSetValueWithOverwrite: The task ' s notification value is set to the value of ulValue, even if the task being notified had not yet processed the previous notification (the task already had a notification pending). xTaskNotify() always returns pdPASS in this case.
  - eSetValueWithoutOverwrite: If the task being notified did not already have a notification pending then the task ' s notification value is set to ulValue and xTaskNotify() will return pdPASS. If the task being notified already had a notification pending then no action is performed and pdFAIL is returned.
  - eNoAction: The task receives a notification without its notification value being updated. ulValue is not used and xTaskNotify() always returns pdPASS in this case.
- `pxHigherPriorityTaskWoken`: xTaskNotifyFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending the notification caused the task to which the notification was sent to leave the Blocked state, and the unblocked task has a priority higher than the currently running task. If xTaskNotifyFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited. How a context switch is requested from an ISR is dependent on the port - see the documentation page for the port in use.

BaseType_t **`xTaskNotifyWait`**(uint32_t *ulBitsToClearOnEntry*, uint32_t *ulBitsToClearOnExit*, uint32_t **pulNotificationValue*, TickType_t *xTicksToWait*)

Wait for task notification

configUSE_TASK_NOTIFICATIONS must be undefined or defined as 1 for this function to be available.

When configUSE_TASK_NOTIFICATIONS is set to one each task has its own private "notification value", which is a 32-bit unsigned integer (uint32_t).

Events can be sent to a task using an intermediary object. Examples of such objects are queues, semaphores, mutexes and event groups. Task notifications are a method of sending an event directly to a task without the need for such an intermediary object.

A notification sent to a task can optionally perform an action, such as update, overwrite or increment the task ' s notification value. In that way task notifications can be used to send data to a task, or be used as light weight and fast binary or counting semaphores.

A notification sent to a task will remain pending until it is cleared by the task calling xTaskNotifyWait() or ulTaskNotifyTake(). If the task was already in the Blocked state to wait for a notification when the notification arrives then the task will automatically be removed from the Blocked state (unblocked) and the notification cleared.

A task can use xTaskNotifyWait() to [optionally] block to wait for a notification to be pending, or ulTaskNo-

tifyTake() to [optionally] block to wait for its notification value to have a non-zero value. The task does not consume any CPU time while it is in the Blocked state.

See http://www.FreeRTOS.org/RTOS-task-notifications.html for details.

Return If a notification was received (including notifications that were already pending when xTaskNotifyWait was called) then pdPASS is returned. Otherwise pdFAIL is returned.

Parameters

- `ulBitsToClearOnEntry`: Bits that are set in ulBitsToClearOnEntry value will be cleared in the calling task's notification value before the task checks to see if any notifications are pending, and optionally blocks if no notifications are pending. Setting ulBitsToClearOnEntry to ULONG_MAX (if limits.h is included) or 0xffffffffUL (if limits.h is not included) will have the effect of resetting the task's notification value to 0. Setting ulBitsToClearOnEntry to 0 will leave the task's notification value unchanged.
- `ulBitsToClearOnExit`: If a notification is pending or received before the calling task exits the xTaskNotifyWait() function then the task's notification value (see the xTaskNotify() API function) is passed out using the pulNotificationValue parameter. Then any bits that are set in ulBitsToClearOnExit will be cleared in the task's notification value (note *pulNotificationValue is set before any bits are cleared). Setting ulBitsToClearOnExit to ULONG_MAX (if limits.h is included) or 0xffffffffUL (if limits.h is not included) will have the effect of resetting the task's notification value to 0 before the function exits. Setting ulBitsToClearOnExit to 0 will leave the task's notification value unchanged when the function exits (in which case the value passed out in pulNotificationValue will match the task's notification value).
- `pulNotificationValue`: Used to pass the task's notification value out of the function. Note the value passed out will not be effected by the clearing of any bits caused by ulBitsToClearOnExit being non-zero.
- `xTicksToWait`: The maximum amount of time that the task should wait in the Blocked state for a notification to be received, should a notification not already be pending when xTaskNotifyWait() was called. The task will not consume any processing time while it is in the Blocked state. This is specified in kernel ticks, the macro pdMS_TO_TICSK(value_in_ms) can be used to convert a time specified in milliseconds to a time specified in ticks.

void **vTaskNotifyGiveFromISR**(*TaskHandle_t xTaskToNotify*, BaseType_t **pxHigherPriorityTaskWoken*)

Simplified macro for sending task notification from ISR.

configUSE_TASK_NOTIFICATIONS must be undefined or defined as 1 for this macro to be available.

When configUSE_TASK_NOTIFICATIONS is set to one each task has its own private "notification value", which is a 32-bit unsigned integer (uint32_t).

A version of xTaskNotifyGive() that can be called from an interrupt service routine (ISR).

Events can be sent to a task using an intermediary object. Examples of such objects are queues, semaphores, mutexes and event groups. Task notifications are a method of sending an event directly to a task without the need for such an intermediary object.

A notification sent to a task can optionally perform an action, such as update, overwrite or increment the task's notification value. In that way task notifications can be used to send data to a task, or be used as light weight and fast binary or counting semaphores.

vTaskNotifyGiveFromISR() is intended for use when task notifications are used as light weight and faster binary or counting semaphore equivalents. Actual FreeRTOS semaphores are given from an ISR using the xSemaphoreGiveFromISR() API function, the equivalent action that instead uses a task notification is vTaskNotifyGiveFromISR().

When task notifications are being used as a binary or counting semaphore equivalent then the task being notified should wait for the notification using the ulTaskNotificationTake() API function rather than the xTaskNotifyWait() API function.

See http://www.FreeRTOS.org/RTOS-task-notifications.html for more details.

Parameters

- `xTaskToNotify`: The handle of the task being notified. The handle to a task can be returned from the xTaskCreate() API function used to create the task, and the handle of the currently running task can be obtained by calling xTaskGetCurrentTaskHandle().
- `pxHigherPriorityTaskWoken`: vTaskNotifyGiveFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending the notification caused the task to which the notification was sent to leave the Blocked state, and the unblocked task has a priority higher than the currently running task. If vTaskNotifyGiveFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited. How a context switch is requested from an ISR is dependent on the port - see the documentation page for the port in use.

uint32_t **ulTaskNotifyTake**(BaseType_t *xClearCountOnExit*, TickType_t *xTicksToWait*)

Simplified macro for receiving task notification.

configUSE_TASK_NOTIFICATIONS must be undefined or defined as 1 for this function to be available.

When configUSE_TASK_NOTIFICATIONS is set to one each task has its own private "notification value", which is a 32-bit unsigned integer (uint32_t).

Events can be sent to a task using an intermediary object. Examples of such objects are queues, semaphores, mutexes and event groups. Task notifications are a method of sending an event directly to a task without the need for such an intermediary object.

A notification sent to a task can optionally perform an action, such as update, overwrite or increment the task's notification value. In that way task notifications can be used to send data to a task, or be used as light weight and fast binary or counting semaphores.

ulTaskNotifyTake() is intended for use when a task notification is used as a faster and lighter weight binary or counting semaphore alternative. Actual FreeRTOS semaphores are taken using the xSemaphoreTake() API function, the equivalent action that instead uses a task notification is ulTaskNotifyTake().

When a task is using its notification value as a binary or counting semaphore other tasks should send notifications to it using the xTaskNotifyGive() macro, or xTaskNotify() function with the eAction parameter set to eIncrement.

ulTaskNotifyTake() can either clear the task's notification value to zero on exit, in which case the notification value acts like a binary semaphore, or decrement the task's notification value on exit, in which case the notification value acts like a counting semaphore.

A task can use ulTaskNotifyTake() to [optionally] block to wait for a the task's notification value to be non-zero. The task does not consume any CPU time while it is in the Blocked state.

Where as xTaskNotifyWait() will return when a notification is pending, ulTaskNotifyTake() will return when the task's notification value is not zero.

See http://www.FreeRTOS.org/RTOS-task-notifications.html for details.

Return The task's notification count before it is either cleared to zero or decremented (see the xClearCountOnExit parameter).

Parameters

- `xClearCountOnExit`: if xClearCountOnExit is pdFALSE then the task's notification value is decremented when the function exits. In this way the notification value acts like a counting semaphore. If xClearCountOnExit is not pdFALSE then the task's notification value is cleared to zero when the function exits. In this way the notification value acts like a binary semaphore.
- `xTicksToWait`: The maximum amount of time that the task should wait in the Blocked state for the task's notification value to be greater than zero, should the count not already be greater than zero when ulTaskNotifyTake() was called. The task will not consume any processing time while it is in the Blocked state. This is specified in kernel ticks, the macro pdMS_TO_TICSK(value_in_ms) can be used to convert a time specified in milliseconds to a time specified in ticks.

Structures

struct xTASK_STATUS

Used with the uxTaskGetSystemState() function to return the state of each task in the system.

Public Members

TaskHandle_t **xHandle**
The handle of the task to which the rest of the information in the structure relates.

const char ***pcTaskName**
A pointer to the task’ s name. This value will be invalid if the task was deleted since the structure was populated!

UBaseType_t **xTaskNumber**
A number unique to the task.

eTaskState **eCurrentState**
The state in which the task existed when the structure was populated.

UBaseType_t **uxCurrentPriority**
The priority at which the task was running (may be inherited) when the structure was populated.

UBaseType_t **uxBasePriority**
The priority to which the task will return if the task’ s current priority has been inherited to avoid unbounded priority inversion when obtaining a mutex. Only valid if configUSE_MUTEXES is defined as 1 in FreeRTOSConfig.h.

uint32_t **ulRunTimeCounter**
The total run time allocated to the task so far, as defined by the run time stats clock. See http://www.freertos.org/rtos-run-time-stats.html. Only valid when configGENERATE_RUN_TIME_STATS is defined as 1 in FreeRTOSConfig.h.

StackType_t ***pxStackBase**
Points to the lowest address of the task’ s stack area.

uint32_t **usStackHighWaterMark**
The minimum amount of stack space that has remained for the task since the task was created. The closer this value is to zero the closer the task has come to overflowing its stack.

BaseType_t **xCoreID**
Core this task is pinned to (0, 1, or -1 for tskNO_AFFINITY). This field is present if CONFIG_FREERTOS_VTASKLIST_INCLUDE_COREID is set.

struct xTASK_SNAPSHOT
Used with the uxTaskGetSnapshotAll() function to save memory snapshot of each task in the system. We need this struct because TCB_t is defined (hidden) in tasks.c.

Public Members

void ***pxTCB**
Address of task control block.

StackType_t ***pxTopOfStack**
Points to the location of the last item placed on the tasks stack.

StackType_t ***pxEndOfStack**
Points to the end of the stack. pxTopOfStack < pxEndOfStack, stack grows hi2lo pxTopOfStack > pxEndOfStack, stack grows lo2hi

Macros

tskKERNEL_VERSION_NUMBER
tskKERNEL_VERSION_MAJOR

tskKERNEL_VERSION_MINOR

tskKERNEL_VERSION_BUILD

tskNO_AFFINITY
Argument of xTaskCreatePinnedToCore indicating that task has no affinity.

tskIDLE_PRIORITY
Defines the priority used by the idle task. This must not be modified.

taskYIELD()
task. h

Macro for forcing a context switch.

taskENTER_CRITICAL(mux)
task. h

Macro to mark the start of a critical code region. Preemptive context switches cannot occur when in a critical region.

Note This may alter the stack (depending on the portable implementation) so must be used with care!

taskENTER_CRITICAL_ISR(mux)

taskEXIT_CRITICAL(mux)
task. h

Macro to mark the end of a critical code region. Preemptive context switches cannot occur when in a critical region.

Note This may alter the stack (depending on the portable implementation) so must be used with care!

taskEXIT_CRITICAL_ISR(mux)

taskDISABLE_INTERRUPTS()
task. h

Macro to disable all maskable interrupts.

taskENABLE_INTERRUPTS()
task. h

Macro to enable microcontroller interrupts.

taskSCHEDULER_SUSPENDED

taskSCHEDULER_NOT_STARTED

taskSCHEDULER_RUNNING

xTaskNotifyGive(xTaskToNotify)
Simplified macro for sending task notification.

configUSE_TASK_NOTIFICATIONS must be undefined or defined as 1 for this macro to be available.

When configUSE_TASK_NOTIFICATIONS is set to one each task has its own private "notification value", which is a 32-bit unsigned integer (uint32_t).

Events can be sent to a task using an intermediary object. Examples of such objects are queues, semaphores, mutexes and event groups. Task notifications are a method of sending an event directly to a task without the need for such an intermediary object.

A notification sent to a task can optionally perform an action, such as update, overwrite or increment the task's notification value. In that way task notifications can be used to send data to a task, or be used as light weight and fast binary or counting semaphores.

xTaskNotifyGive() is a helper macro intended for use when task notifications are used as light weight and faster binary or counting semaphore equivalents. Actual FreeRTOS semaphores are given using the xSemaphoreGive() API function, the equivalent action that instead uses a task notification is xTaskNotifyGive().

When task notifications are being used as a binary or counting semaphore equivalent then the task being notified should wait for the notification using the ulTaskNotificationTake() API function rather than the xTaskNotifyWait() API function.

See http://www.FreeRTOS.org/RTOS-task-notifications.html for more details.

Return xTaskNotifyGive() is a macro that calls xTaskNotify() with the eAction parameter set to eIncrement - so pdPASS is always returned.

Parameters

- `xTaskToNotify`: The handle of the task being notified. The handle to a task can be returned from the xTaskCreate() API function used to create the task, and the handle of the currently running task can be obtained by calling xTaskGetCurrentTaskHandle().

Type Definitions

typedef void ***TaskHandle_t**

task. h

Type by which tasks are referenced. For example, a call to xTaskCreate returns (via a pointer parameter) an TaskHandle_t variable that can then be used as a parameter to vTaskDelete to delete the task.

typedef BaseType_t (***TaskHookFunction_t**) (void *)

Defines the prototype to which the application task hook function must conform.

typedef struct *xTASK_STATUS* **TaskStatus_t**

Used with the uxTaskGetSystemState() function to return the state of each task in the system.

typedef struct *xTASK_SNAPSHOT* **TaskSnapshot_t**

Used with the uxTaskGetSnapshotAll() function to save memory snapshot of each task in the system. We need this struct because TCB_t is defined (hidden) in tasks.c.

typedef void (***TlsDeleteCallbackFunction_t**) (int, void *)

Prototype of local storage pointer deletion callback.

Enumerations

enum eTaskState

Task states returned by eTaskGetState.

Values:

eRunning = 0

A task is querying the state of itself, so must be running.

eReady

The task being queried is in a read or pending ready list.

eBlocked

The task being queried is in the Blocked state.

eSuspended

The task being queried is in the Suspended state, or is in the Blocked state with an infinite time out.

eDeleted

The task being queried has been deleted, but its TCB has not yet been freed.

enum eNotifyAction

Actions that can be performed when vTaskNotify() is called.

Values:

eNoAction = 0

Notify the task without updating its notify value.

eSetBits

Set bits in the task' s notification value.

eIncrement

Increment the task' s notification value.

eSetValueWithOverwrite

Set the task' s notification value to a specific value even if the previous value has not yet been read by the task.

eSetValueWithoutOverwrite
Set the task's notification value if the previous value has been read by the task.

enum eSleepModeStatus
Possible return values for eTaskConfirmSleepModeStatus().

Values:

eAbortSleep = 0
A task has been made ready or a context switch pended since portSUPPORESS_TICKS_AND_SLEEP() was called - abort entering a sleep mode.

eStandardSleep
Enter a sleep mode that will not last any longer than the expected idle time.

eNoTasksWaitingTimeout
No tasks are waiting for a timeout so it is safe to enter a sleep mode that can only be exited by an external interrupt.

Queue API

Header File

- freertos/include/freertos/queue.h

Functions

BaseType_t **xQueueGenericSendFromISR**(*QueueHandle_t xQueue*, **const** void *__const__ *pvItemToQueue*, BaseType_t *__const__ *pxHigherPriorityTaskWoken*, **const** BaseType_t *xCopyPosition*)

It is preferred that the macros xQueueSendFromISR(), xQueueSendToFrontFromISR() and xQueueSendToBackFromISR() be used in place of calling this function directly. xQueueGiveFromISR() is an equivalent for use by semaphores that don't actually copy any data.

Post an item on a queue. It is safe to use this function from within an interrupt service routine.

Items are queued by copy not reference so it is preferable to only queue small items, especially when called from an ISR. In most cases it would be preferable to store a pointer to the item being queued.

Example usage for buffered IO (where the ISR can obtain more than one value per call):

```
void vBufferISR( void )
{
char cIn;
BaseType_t xHigherPriorityTaskWokenByPost;

   // We have not woken a task at the start of the ISR.
   xHigherPriorityTaskWokenByPost = pdFALSE;

   // Loop until the buffer is empty.
   do
   {
        // Obtain a byte from the buffer.
        cIn = portINPUT_BYTE( RX_REGISTER_ADDRESS );

        // Post each byte.
        xQueueGenericSendFromISR( xRxQueue, &cIn, &
↪xHigherPriorityTaskWokenByPost, queueSEND_TO_BACK );

   } while( portINPUT_BYTE( BUFFER_COUNT ) );

   // Now the buffer is empty we can switch context if necessary.  Note that␣
↪the
   // name of the yield function required is port specific.
```

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```
    if( xHigherPriorityTaskWokenByPost )
    {
        taskYIELD_YIELD_FROM_ISR();
    }
}
```

Return pdTRUE if the data was successfully sent to the queue, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `[out] pxHigherPriorityTaskWoken`: xQueueGenericSendFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending to the queue caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xQueueGenericSendFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.
- `xCopyPosition`: Can take the value queueSEND_TO_BACK to place the item at the back of the queue, or queueSEND_TO_FRONT to place the item at the front of the queue (for high priority messages).

BaseType_t **xQueueGiveFromISR**(*QueueHandle_t xQueue*, BaseType_t ***const** *pxHigherPriorityTaskWoken*)

BaseType_t **xQueueIsQueueEmptyFromISR**(**const** *QueueHandle_t xQueue*)

Utilities to query queues that are safe to use from an ISR. These utilities should be used only from witin an ISR, or within a critical section.

BaseType_t **xQueueIsQueueFullFromISR**(**const** *QueueHandle_t xQueue*)

UBaseType_t **uxQueueMessagesWaitingFromISR**(**const** *QueueHandle_t xQueue*)

BaseType_t **xQueueGenericSend**(*QueueHandle_t xQueue*, **const** void ***const** *pvItemToQueue*, TickType_t *xTicksToWait*, **const** BaseType_t *xCopyPosition*)

It is preferred that the macros xQueueSend(), xQueueSendToFront() and xQueueSendToBack() are used in place of calling this function directly.

Post an item on a queue. The item is queued by copy, not by reference. This function must not be called from an interrupt service routine. See xQueueSendFromISR () for an alternative which may be used in an ISR.

Example usage:

```
struct AMessage
{
char ucMessageID;
char ucData[ 20 ];
} xMessage;

uint32_t ulVar = 10UL;

void vATask( void *pvParameters )
{
QueueHandle_t xQueue1, xQueue2;
struct AMessage *pxMessage;

// Create a queue capable of containing 10 uint32_t values.
xQueue1 = xQueueCreate( 10, sizeof( uint32_t ) );

// Create a queue capable of containing 10 pointers to AMessage structures.
// These should be passed by pointer as they contain a lot of data.
xQueue2 = xQueueCreate( 10, sizeof( struct AMessage * ) );

// ...
```

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```
if( xQueue1 != 0 )
{
    // Send an uint32_t.  Wait for 10 ticks for space to become
    // available if necessary.
    if( xQueueGenericSend( xQueue1, ( void * ) &ulVar, ( TickType_t ) 10,␣
↪queueSEND_TO_BACK ) != pdPASS )
    {
        // Failed to post the message, even after 10 ticks.
    }
}

if( xQueue2 != 0 )
{
    // Send a pointer to a struct AMessage object.  Don't block if the
    // queue is already full.
    pxMessage = & xMessage;
    xQueueGenericSend( xQueue2, ( void * ) &pxMessage, ( TickType_t ) 0,␣
↪queueSEND_TO_BACK );
}

// ... Rest of task code.
}
```

Return pdTRUE if the item was successfully posted, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `xTicksToWait`: The maximum amount of time the task should block waiting for space to become available on the queue, should it already be full. The call will return immediately if this is set to 0 and the queue is full. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required.
- `xCopyPosition`: Can take the value queueSEND_TO_BACK to place the item at the back of the queue, or queueSEND_TO_FRONT to place the item at the front of the queue (for high priority messages).

BaseType_t **xQueuePeekFromISR**(*QueueHandle_t xQueue*, void *const *pvBuffer*)

A version of xQueuePeek() that can be called from an interrupt service routine (ISR).

Receive an item from a queue without removing the item from the queue. The item is received by copy so a buffer of adequate size must be provided. The number of bytes copied into the buffer was defined when the queue was created.

Successfully received items remain on the queue so will be returned again by the next call, or a call to xQueueReceive().

Return pdTRUE if an item was successfully received from the queue, otherwise pdFALSE.

Parameters

- `xQueue`: The handle to the queue from which the item is to be received.
- `pvBuffer`: Pointer to the buffer into which the received item will be copied.

BaseType_t **xQueueGenericReceive**(*QueueHandle_t xQueue*, void *const *pvBuffer*, TickType_t *xTicksToWait*, const BaseType_t *xJustPeek*)

It is preferred that the macro xQueueReceive() be used rather than calling this function directly.

Receive an item from a queue. The item is received by copy so a buffer of adequate size must be provided. The number of bytes copied into the buffer was defined when the queue was created.

This function must not be used in an interrupt service routine. See xQueueReceiveFromISR for an alternative that can.

Example usage:

```
struct AMessage
{
    char ucMessageID;
    char ucData[ 20 ];
} xMessage;

QueueHandle_t xQueue;

// Task to create a queue and post a value.
void vATask( void *pvParameters )
{
struct AMessage *pxMessage;

    // Create a queue capable of containing 10 pointers to AMessage structures.
    // These should be passed by pointer as they contain a lot of data.
    xQueue = xQueueCreate( 10, sizeof( struct AMessage * ) );
    if( xQueue == 0 )
    {
        // Failed to create the queue.
    }

    // ...

    // Send a pointer to a struct AMessage object.  Don't block if the
    // queue is already full.
    pxMessage = & xMessage;
    xQueueSend( xQueue, ( void * ) &pxMessage, ( TickType_t ) 0 );

    // ... Rest of task code.
}

// Task to receive from the queue.
void vADifferentTask( void *pvParameters )
{
struct AMessage *pxRxedMessage;

    if( xQueue != 0 )
    {
        // Receive a message on the created queue.  Block for 10 ticks if a
        // message is not immediately available.
        if( xQueueGenericReceive( xQueue, &( pxRxedMessage ), ( TickType_t ) 10
↪) )
        {
            // pcRxedMessage now points to the struct AMessage variable posted
            // by vATask.
        }
    }

    // ... Rest of task code.
}
```

Return pdTRUE if an item was successfully received from the queue, otherwise pdFALSE.

Parameters

- `xQueue`: The handle to the queue from which the item is to be received.
- `pvBuffer`: Pointer to the buffer into which the received item will be copied.
- `xTicksToWait`: The maximum amount of time the task should block waiting for an item to receive should the queue be empty at the time of the call. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required. xQueueGenericReceive() will return immediately if the queue is empty and xTicksToWait is 0.
- `xJustPeek`: When set to true, the item received from the queue is not actually removed from the queue - meaning a subsequent call to xQueueReceive() will return the same item. When set to false, the item being received from the queue is also removed from the queue.

UBaseType_t **uxQueueMessagesWaiting**(**const** *QueueHandle_t xQueue*)

Return the number of messages stored in a queue.

Return The number of messages available in the queue.

Parameters

- `xQueue`: A handle to the queue being queried.

UBaseType_t **uxQueueSpacesAvailable**(**const** *QueueHandle_t xQueue*)

Return the number of free spaces available in a queue. This is equal to the number of items that can be sent to the queue before the queue becomes full if no items are removed.

Return The number of spaces available in the queue.

Parameters

- `xQueue`: A handle to the queue being queried.

void **vQueueDelete**(*QueueHandle_t xQueue*)

Delete a queue - freeing all the memory allocated for storing of items placed on the queue.

Parameters

- `xQueue`: A handle to the queue to be deleted.

BaseType_t **xQueueReceiveFromISR**(*QueueHandle_t xQueue*, void ***const** *pvBuffer*, BaseType_t ***const** *pxHigherPriorityTaskWoken*)

Receive an item from a queue. It is safe to use this function from within an interrupt service routine.

Example usage:

```
QueueHandle_t xQueue;

// Function to create a queue and post some values.
void vAFunction( void *pvParameters )
{
char cValueToPost;
const TickType_t xTicksToWait = ( TickType_t )0xff;

   // Create a queue capable of containing 10 characters.
   xQueue = xQueueCreate( 10, sizeof( char ) );
   if( xQueue == 0 )
   {
       // Failed to create the queue.
   }

   // ...

   // Post some characters that will be used within an ISR.  If the queue
   // is full then this task will block for xTicksToWait ticks.
   cValueToPost = 'a';
   xQueueSend( xQueue, ( void * ) &cValueToPost, xTicksToWait );
   cValueToPost = 'b';
   xQueueSend( xQueue, ( void * ) &cValueToPost, xTicksToWait );

   // ... keep posting characters ... this task may block when the queue
   // becomes full.

   cValueToPost = 'c';
   xQueueSend( xQueue, ( void * ) &cValueToPost, xTicksToWait );
}

// ISR that outputs all the characters received on the queue.
void vISR_Routine( void )
{
BaseType_t xTaskWokenByReceive = pdFALSE;
char cRxedChar;
```

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```
   while( xQueueReceiveFromISR( xQueue, ( void * ) &cRxedChar, &
↪xTaskWokenByReceive) )
   {
       // A character was received.  Output the character now.
       vOutputCharacter( cRxedChar );

       // If removing the character from the queue woke the task that was
       // posting onto the queue cTaskWokenByReceive will have been set to
       // pdTRUE.  No matter how many times this loop iterates only one
       // task will be woken.
   }

   if( cTaskWokenByPost != ( char ) pdFALSE;
   {
       taskYIELD ();
   }
}
```

Return pdTRUE if an item was successfully received from the queue, otherwise pdFALSE.

Parameters

- `xQueue`: The handle to the queue from which the item is to be received.
- `pvBuffer`: Pointer to the buffer into which the received item will be copied.
- `[out] pxHigherPriorityTaskWoken`: A task may be blocked waiting for space to become available on the queue. If xQueueReceiveFromISR causes such a task to unblock *pxTaskWoken will get set to pdTRUE, otherwise *pxTaskWoken will remain unchanged.

void **vQueueAddToRegistry** (*QueueHandle_t xQueue*, **const** char **pcName*)

The registry is provided as a means for kernel aware debuggers to locate queues, semaphores and mutexes. Call vQueueAddToRegistry() add a queue, semaphore or mutex handle to the registry if you want the handle to be available to a kernel aware debugger. If you are not using a kernel aware debugger then this function can be ignored.

configQUEUE_REGISTRY_SIZE defines the maximum number of handles the registry can hold. configQUEUE_REGISTRY_SIZE must be greater than 0 within FreeRTOSConfig.h for the registry to be available. Its value does not effect the number of queues, semaphores and mutexes that can be created - just the number that the registry can hold.

Parameters

- `xQueue`: The handle of the queue being added to the registry. This is the handle returned by a call to xQueueCreate(). Semaphore and mutex handles can also be passed in here.
- `pcName`: The name to be associated with the handle. This is the name that the kernel aware debugger will display. The queue registry only stores a pointer to the string - so the string must be persistent (global or preferably in ROM/Flash), not on the stack.

void **vQueueUnregisterQueue** (*QueueHandle_t xQueue*)

The registry is provided as a means for kernel aware debuggers to locate queues, semaphores and mutexes. Call vQueueAddToRegistry() add a queue, semaphore or mutex handle to the registry if you want the handle to be available to a kernel aware debugger, and vQueueUnregisterQueue() to remove the queue, semaphore or mutex from the register. If you are not using a kernel aware debugger then this function can be ignored.

Parameters

- `xQueue`: The handle of the queue being removed from the registry.

const char ***pcQueueGetName** (*QueueHandle_t xQueue*)

The queue registry is provided as a means for kernel aware debuggers to locate queues, semaphores and mutexes. Call pcQueueGetName() to look up and return the name of a queue in the queue registry from the queue's handle.

Note This function has been back ported from FreeRTOS v9.0.0

Return If the queue is in the registry then a pointer to the name of the queue is returned. If the queue is not in the registry then NULL is returned.

Parameters

- `xQueue`: The handle of the queue the name of which will be returned.

QueueHandle_t **xQueueGenericCreate**(**const** UBaseType_t *uxQueueLength*, **const** UBaseType_t *uxItemSize*, **const** uint8_t *ucQueueType*)

Generic version of the function used to creaet a queue using dynamic memory allocation. This is called by other functions and macros that create other RTOS objects that use the queue structure as their base.

QueueHandle_t **xQueueGenericCreateStatic**(**const** UBaseType_t *uxQueueLength*, **const** UBaseType_t *uxItemSize*, uint8_t **pucQueueStorage*, StaticQueue_t **pxStaticQueue*, **const** uint8_t *ucQueueType*)

Generic version of the function used to creaet a queue using dynamic memory allocation. This is called by other functions and macros that create other RTOS objects that use the queue structure as their base.

QueueSetHandle_t **xQueueCreateSet**(**const** UBaseType_t *uxEventQueueLength*)

Queue sets provide a mechanism to allow a task to block (pend) on a read operation from multiple queues or semaphores simultaneously.

See FreeRTOS/Source/Demo/Common/Minimal/QueueSet.c for an example using this function.

A queue set must be explicitly created using a call to xQueueCreateSet() before it can be used. Once created, standard FreeRTOS queues and semaphores can be added to the set using calls to xQueueAddToSet(). xQueueSelectFromSet() is then used to determine which, if any, of the queues or semaphores contained in the set is in a state where a queue read or semaphore take operation would be successful.

Note 1: See the documentation on http://www.FreeRTOS.org/RTOS-queue-sets.html for reasons why queue sets are very rarely needed in practice as there are simpler methods of blocking on multiple objects.

Note 2: Blocking on a queue set that contains a mutex will not cause the mutex holder to inherit the priority of the blocked task.

Note 3: An additional 4 bytes of RAM is required for each space in a every queue added to a queue set. Therefore counting semaphores that have a high maximum count value should not be added to a queue set.

Note 4: A receive (in the case of a queue) or take (in the case of a semaphore) operation must not be performed on a member of a queue set unless a call to xQueueSelectFromSet() has first returned a handle to that set member.

Return If the queue set is created successfully then a handle to the created queue set is returned. Otherwise NULL is returned.

Parameters

- `uxEventQueueLength`: Queue sets store events that occur on the queues and semaphores contained in the set. uxEventQueueLength specifies the maximum number of events that can be queued at once. To be absolutely certain that events are not lost uxEventQueueLength should be set to the total sum of the length of the queues added to the set, where binary semaphores and mutexes have a length of 1, and counting semaphores have a length set by their maximum count value. Examples:
  - If a queue set is to hold a queue of length 5, another queue of length 12, and a binary semaphore, then uxEventQueueLength should be set to (5 + 12 + 1), or 18.
  - If a queue set is to hold three binary semaphores then uxEventQueueLength should be set to (1 + 1 + 1), or 3.
  - If a queue set is to hold a counting semaphore that has a maximum count of 5, and a counting semaphore that has a maximum count of 3, then uxEventQueueLength should be set to (5 + 3), or 8.

BaseType_t **xQueueAddToSet**(*QueueSetMemberHandle_t* *xQueueOrSemaphore*, *QueueSetHandle_t* *xQueueSet*)

Adds a queue or semaphore to a queue set that was previously created by a call to xQueueCreateSet().

See FreeRTOS/Source/Demo/Common/Minimal/QueueSet.c for an example using this function.

Note 1: A receive (in the case of a queue) or take (in the case of a semaphore) operation must not be performed on a member of a queue set unless a call to xQueueSelectFromSet() has first returned a handle to that set member.

Return If the queue or semaphore was successfully added to the queue set then pdPASS is returned. If the queue could not be successfully added to the queue set because it is already a member of a different queue set then pdFAIL is returned.

Parameters

- `xQueueOrSemaphore`: The handle of the queue or semaphore being added to the queue set (cast to an QueueSetMemberHandle_t type).
- `xQueueSet`: The handle of the queue set to which the queue or semaphore is being added.

BaseType_t **xQueueRemoveFromSet**(*QueueSetMemberHandle_t xQueueOrSemaphore*, *QueueSetHandle_t xQueueSet*)

Removes a queue or semaphore from a queue set. A queue or semaphore can only be removed from a set if the queue or semaphore is empty.

See FreeRTOS/Source/Demo/Common/Minimal/QueueSet.c for an example using this function.

Return If the queue or semaphore was successfully removed from the queue set then pdPASS is returned. If the queue was not in the queue set, or the queue (or semaphore) was not empty, then pdFAIL is returned.

Parameters

- `xQueueOrSemaphore`: The handle of the queue or semaphore being removed from the queue set (cast to an QueueSetMemberHandle_t type).
- `xQueueSet`: The handle of the queue set in which the queue or semaphore is included.

QueueSetMemberHandle_t **xQueueSelectFromSet**(*QueueSetHandle_t xQueueSet*, **const** TickType_t *xTicksToWait*)

xQueueSelectFromSet() selects from the members of a queue set a queue or semaphore that either contains data (in the case of a queue) or is available to take (in the case of a semaphore). xQueueSelectFromSet() effectively allows a task to block (pend) on a read operation on all the queues and semaphores in a queue set simultaneously.

See FreeRTOS/Source/Demo/Common/Minimal/QueueSet.c for an example using this function.

Note 1: See the documentation on http://www.FreeRTOS.org/RTOS-queue-sets.html for reasons why queue sets are very rarely needed in practice as there are simpler methods of blocking on multiple objects.

Note 2: Blocking on a queue set that contains a mutex will not cause the mutex holder to inherit the priority of the blocked task.

Note 3: A receive (in the case of a queue) or take (in the case of a semaphore) operation must not be performed on a member of a queue set unless a call to xQueueSelectFromSet() has first returned a handle to that set member.

Return xQueueSelectFromSet() will return the handle of a queue (cast to a QueueSetMemberHandle_t type) contained in the queue set that contains data, or the handle of a semaphore (cast to a QueueSetMemberHandle_t type) contained in the queue set that is available, or NULL if no such queue or semaphore exists before before the specified block time expires.

Parameters

- `xQueueSet`: The queue set on which the task will (potentially) block.
- `xTicksToWait`: The maximum time, in ticks, that the calling task will remain in the Blocked state (with other tasks executing) to wait for a member of the queue set to be ready for a successful queue read or semaphore take operation.

QueueSetMemberHandle_t **xQueueSelectFromSetFromISR**(*QueueSetHandle_t xQueueSet*)

A version of xQueueSelectFromSet() that can be used from an ISR.

Macros

xQueueCreate(uxQueueLength, uxItemSize)

Creates a new queue instance. This allocates the storage required by the new queue and returns a handle for the queue.

Example usage:

```
struct AMessage
{
```

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```
char ucMessageID;
char ucData[ 20 ];
};

void vATask( void *pvParameters )
{
QueueHandle_t xQueue1, xQueue2;

// Create a queue capable of containing 10 uint32_t values.
xQueue1 = xQueueCreate( 10, sizeof( uint32_t ) );
if( xQueue1 == 0 )
{
    // Queue was not created and must not be used.
}

// Create a queue capable of containing 10 pointers to AMessage structures.
// These should be passed by pointer as they contain a lot of data.
xQueue2 = xQueueCreate( 10, sizeof( struct AMessage * ) );
if( xQueue2 == 0 )
{
    // Queue was not created and must not be used.
}

// ... Rest of task code.
}
```

Return If the queue is successfully create then a handle to the newly created queue is returned. If the queue cannot be created then 0 is returned.

Parameters

- `uxQueueLength`: The maximum number of items that the queue can contain.
- `uxItemSize`: The number of bytes each item in the queue will require. Items are queued by copy, not by reference, so this is the number of bytes that will be copied for each posted item. Each item on the queue must be the same size.

`xQueueCreateStatic`(uxQueueLength, uxItemSize, pucQueueStorage, pxQueueBuffer)

Creates a new queue instance, and returns a handle by which the new queue can be referenced.

Internally, within the FreeRTOS implementation, queues use two blocks of memory. The first block is used to hold the queue's data structures. The second block is used to hold items placed into the queue. If a queue is created using xQueueCreate() then both blocks of memory are automatically dynamically allocated inside the xQueueCreate() function. (see http://www.freertos.org/a00111.html). If a queue is created using xQueueCreateStatic() then the application writer must provide the memory that will get used by the queue. xQueueCreateStatic() therefore allows a queue to be created without using any dynamic memory allocation.

http://www.FreeRTOS.org/Embedded-RTOS-Queues.html

Example usage:

```
struct AMessage
{
 char ucMessageID;
 char ucData[ 20 ];
};

#define QUEUE_LENGTH 10
#define ITEM_SIZE sizeof( uint32_t )

// xQueueBuffer will hold the queue structure.
StaticQueue_t xQueueBuffer;

// ucQueueStorage will hold the items posted to the queue.  Must be at least
// [(queue length) * ( queue item size)] bytes long.
```

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```
uint8_t ucQueueStorage[ QUEUE_LENGTH * ITEM_SIZE ];

void vATask( void *pvParameters )
{
 QueueHandle_t xQueue1;

 // Create a queue capable of containing 10 uint32_t values.
 xQueue1 = xQueueCreate( QUEUE_LENGTH, // The number of items the queue can␣
↪hold.
                         ITEM_SIZE     // The size of each item in the queue
                         &( ucQueueStorage[ 0 ] ), // The buffer that will␣
↪hold the items in the queue.
                         &xQueueBuffer ); // The buffer that will hold the␣
↪queue structure.

 // The queue is guaranteed to be created successfully as no dynamic memory
 // allocation is used.  Therefore xQueue1 is now a handle to a valid queue.

 // ... Rest of task code.
}
```

Return If the queue is created then a handle to the created queue is returned. If pxQueueBuffer is NULL then NULL is returned.

Parameters

- `uxQueueLength`: The maximum number of items that the queue can contain.
- `uxItemSize`: The number of bytes each item in the queue will require. Items are queued by copy, not by reference, so this is the number of bytes that will be copied for each posted item. Each item on the queue must be the same size.
- `pucQueueStorage`: If uxItemSize is not zero then pucQueueStorageBuffer must point to a uint8_t array that is at least large enough to hold the maximum number of items that can be in the queue at any one time - which is (uxQueueLength * uxItemsSize) bytes. If uxItemSize is zero then pucQueueStorageBuffer can be NULL.
- `pxQueueBuffer`: Must point to a variable of type StaticQueue_t, which will be used to hold the queue's data structure.

`xQueueSendToFront` (xQueue, pvItemToQueue, xTicksToWait)

This is a macro that calls xQueueGenericSend().

Post an item to the front of a queue. The item is queued by copy, not by reference. This function must not be called from an interrupt service routine. See xQueueSendFromISR () for an alternative which may be used in an ISR.

Example usage:

```
struct AMessage
{
char ucMessageID;
char ucData[ 20 ];
} xMessage;

uint32_t ulVar = 10UL;

void vATask( void *pvParameters )
{
QueueHandle_t xQueue1, xQueue2;
struct AMessage *pxMessage;

// Create a queue capable of containing 10 uint32_t values.
xQueue1 = xQueueCreate( 10, sizeof( uint32_t ) );

// Create a queue capable of containing 10 pointers to AMessage structures.
```

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```
// These should be passed by pointer as they contain a lot of data.
xQueue2 = xQueueCreate( 10, sizeof( struct AMessage * ) );

// ...

if( xQueue1 != 0 )
{
    // Send an uint32_t.  Wait for 10 ticks for space to become
    // available if necessary.
    if( xQueueSendToFront( xQueue1, ( void * ) &ulVar, ( TickType_t ) 10 ) !=␣
↪pdPASS )
    {
        // Failed to post the message, even after 10 ticks.
    }
}

if( xQueue2 != 0 )
{
    // Send a pointer to a struct AMessage object.  Don't block if the
    // queue is already full.
    pxMessage = & xMessage;
    xQueueSendToFront( xQueue2, ( void * ) &pxMessage, ( TickType_t ) 0 );
}

// ... Rest of task code.
}
```

Return pdTRUE if the item was successfully posted, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `xTicksToWait`: The maximum amount of time the task should block waiting for space to become available on the queue, should it already be full. The call will return immediately if this is set to 0 and the queue is full. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required.

xQueueSendToBack (xQueue, pvItemToQueue, xTicksToWait)

This is a macro that calls xQueueGenericSend().

Post an item to the back of a queue. The item is queued by copy, not by reference. This function must not be called from an interrupt service routine. See xQueueSendFromISR () for an alternative which may be used in an ISR.

Example usage:

```
struct AMessage
{
char ucMessageID;
char ucData[ 20 ];
} xMessage;

uint32_t ulVar = 10UL;

void vATask( void *pvParameters )
{
QueueHandle_t xQueue1, xQueue2;
struct AMessage *pxMessage;

// Create a queue capable of containing 10 uint32_t values.
```

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```
xQueue1 = xQueueCreate( 10, sizeof( uint32_t ) );

// Create a queue capable of containing 10 pointers to AMessage structures.
// These should be passed by pointer as they contain a lot of data.
xQueue2 = xQueueCreate( 10, sizeof( struct AMessage * ) );

// ...

if( xQueue1 != 0 )
{
    // Send an uint32_t.  Wait for 10 ticks for space to become
    // available if necessary.
    if( xQueueSendToBack( xQueue1, ( void * ) &ulVar, ( TickType_t ) 10 ) !=␣
↪pdPASS )
    {
        // Failed to post the message, even after 10 ticks.
    }
}

if( xQueue2 != 0 )
{
    // Send a pointer to a struct AMessage object.  Don't block if the
    // queue is already full.
    pxMessage = & xMessage;
    xQueueSendToBack( xQueue2, ( void * ) &pxMessage, ( TickType_t ) 0 );
}

// ... Rest of task code.
}
```

Return pdTRUE if the item was successfully posted, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `xTicksToWait`: The maximum amount of time the task should block waiting for space to become available on the queue, should it already be full. The call will return immediately if this is set to 0 and the queue is full. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required.

xQueueSend(xQueue, pvItemToQueue, xTicksToWait)

This is a macro that calls xQueueGenericSend(). It is included for backward compatibility with versions of FreeRTOS.org that did not include the xQueueSendToFront() and xQueueSendToBack() macros. It is equivalent to xQueueSendToBack().

Post an item on a queue. The item is queued by copy, not by reference. This function must not be called from an interrupt service routine. See xQueueSendFromISR () for an alternative which may be used in an ISR.

Example usage:

```
struct AMessage
{
char ucMessageID;
char ucData[ 20 ];
} xMessage;

uint32_t ulVar = 10UL;

void vATask( void *pvParameters )
{
```

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```
QueueHandle_t xQueue1, xQueue2;
struct AMessage *pxMessage;

// Create a queue capable of containing 10 uint32_t values.
xQueue1 = xQueueCreate( 10, sizeof( uint32_t ) );

// Create a queue capable of containing 10 pointers to AMessage structures.
// These should be passed by pointer as they contain a lot of data.
xQueue2 = xQueueCreate( 10, sizeof( struct AMessage * ) );

// ...

if( xQueue1 != 0 )
{
    // Send an uint32_t.  Wait for 10 ticks for space to become
    // available if necessary.
    if( xQueueSend( xQueue1, ( void * ) &ulVar, ( TickType_t ) 10 ) != pdPASS )
    {
        // Failed to post the message, even after 10 ticks.
    }
}

if( xQueue2 != 0 )
{
    // Send a pointer to a struct AMessage object.  Don't block if the
    // queue is already full.
    pxMessage = & xMessage;
    xQueueSend( xQueue2, ( void * ) &pxMessage, ( TickType_t ) 0 );
}

// ... Rest of task code.
}
```

Return pdTRUE if the item was successfully posted, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `xTicksToWait`: The maximum amount of time the task should block waiting for space to become available on the queue, should it already be full. The call will return immediately if this is set to 0 and the queue is full. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required.

`xQueueOverwrite`(xQueue, pvItemToQueue)

Only for use with queues that have a length of one - so the queue is either empty or full.

Post an item on a queue. If the queue is already full then overwrite the value held in the queue. The item is queued by copy, not by reference.

This function must not be called from an interrupt service routine. See xQueueOverwriteFromISR () for an alternative which may be used in an ISR.

Example usage:

```
void vFunction( void *pvParameters )
{
QueueHandle_t xQueue;
uint32_t ulVarToSend, ulValReceived;

// Create a queue to hold one uint32_t value.  It is strongly
```

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```
// recommended *not* to use xQueueOverwrite() on queues that can
// contain more than one value, and doing so will trigger an assertion
// if configASSERT() is defined.
xQueue = xQueueCreate( 1, sizeof( uint32_t ) );

// Write the value 10 to the queue using xQueueOverwrite().
ulVarToSend = 10;
xQueueOverwrite( xQueue, &ulVarToSend );

// Peeking the queue should now return 10, but leave the value 10 in
// the queue.  A block time of zero is used as it is known that the
// queue holds a value.
ulValReceived = 0;
xQueuePeek( xQueue, &ulValReceived, 0 );

if( ulValReceived != 10 )
{
    // Error unless the item was removed by a different task.
}

// The queue is still full.  Use xQueueOverwrite() to overwrite the
// value held in the queue with 100.
ulVarToSend = 100;
xQueueOverwrite( xQueue, &ulVarToSend );

// This time read from the queue, leaving the queue empty once more.
// A block time of 0 is used again.
xQueueReceive( xQueue, &ulValReceived, 0 );

// The value read should be the last value written, even though the
// queue was already full when the value was written.
if( ulValReceived != 100 )
{
    // Error!
}

// ...
}
```

Return xQueueOverwrite() is a macro that calls xQueueGenericSend(), and therefore has the same return values as xQueueSendToFront(). However, pdPASS is the only value that can be returned because xQueueOverwrite() will write to the queue even when the queue is already full.

Parameters

- `xQueue`: The handle of the queue to which the data is being sent.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.

`xQueuePeek`(xQueue, pvBuffer, xTicksToWait)

This is a macro that calls the xQueueGenericReceive() function.

Receive an item from a queue without removing the item from the queue. The item is received by copy so a buffer of adequate size must be provided. The number of bytes copied into the buffer was defined when the queue was created.

Successfully received items remain on the queue so will be returned again by the next call, or a call to xQueueReceive().

This macro must not be used in an interrupt service routine. See xQueuePeekFromISR() for an alternative that can be called from an interrupt service routine.

Example usage:

```
struct AMessage
{
char ucMessageID;
char ucData[ 20 ];
} xMessage;

QueueHandle_t xQueue;

// Task to create a queue and post a value.
void vATask( void *pvParameters )
{
struct AMessage *pxMessage;

// Create a queue capable of containing 10 pointers to AMessage structures.
// These should be passed by pointer as they contain a lot of data.
xQueue = xQueueCreate( 10, sizeof( struct AMessage * ) );
if( xQueue == 0 )
{
    // Failed to create the queue.
}

// ...

// Send a pointer to a struct AMessage object.  Don't block if the
// queue is already full.
pxMessage = & xMessage;
xQueueSend( xQueue, ( void * ) &pxMessage, ( TickType_t ) 0 );

// ... Rest of task code.
}

// Task to peek the data from the queue.
void vADifferentTask( void *pvParameters )
{
struct AMessage *pxRxedMessage;

if( xQueue != 0 )
{
    // Peek a message on the created queue.  Block for 10 ticks if a
    // message is not immediately available.
    if( xQueuePeek( xQueue, &( pxRxedMessage ), ( TickType_t ) 10 ) )
    {
        // pcRxedMessage now points to the struct AMessage variable posted
        // by vATask, but the item still remains on the queue.
    }
}

// ... Rest of task code.
}
```

Return pdTRUE if an item was successfully received from the queue, otherwise pdFALSE.

Parameters

- `xQueue`: The handle to the queue from which the item is to be received.
- `pvBuffer`: Pointer to the buffer into which the received item will be copied.
- `xTicksToWait`: The maximum amount of time the task should block waiting for an item to receive should the queue be empty at the time of the call. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required. xQueuePeek() will return immediately if xTicksToWait is 0 and the queue is empty.

xQueueReceive (xQueue, pvBuffer, xTicksToWait)

queue. h

This is a macro that calls the xQueueGenericReceive() function.

Receive an item from a queue. The item is received by copy so a buffer of adequate size must be provided. The number of bytes copied into the buffer was defined when the queue was created.

Successfully received items are removed from the queue.

This function must not be used in an interrupt service routine. See xQueueReceiveFromISR for an alternative that can.

Example usage:

```
struct AMessage
{
 char ucMessageID;
 char ucData[ 20 ];
} xMessage;

QueueHandle_t xQueue;

// Task to create a queue and post a value.
void vATask( void *pvParameters )
{
struct AMessage *pxMessage;

 // Create a queue capable of containing 10 pointers to AMessage structures.
 // These should be passed by pointer as they contain a lot of data.
 xQueue = xQueueCreate( 10, sizeof( struct AMessage * ) );
 if( xQueue == 0 )
 {
     // Failed to create the queue.
 }

 // ...

 // Send a pointer to a struct AMessage object.  Don't block if the
 // queue is already full.
 pxMessage = & xMessage;
 xQueueSend( xQueue, ( void * ) &pxMessage, ( TickType_t ) 0 );

 // ... Rest of task code.
}

// Task to receive from the queue.
void vADifferentTask( void *pvParameters )
{
struct AMessage *pxRxedMessage;

 if( xQueue != 0 )
 {
     // Receive a message on the created queue.  Block for 10 ticks if a
     // message is not immediately available.
     if( xQueueReceive( xQueue, &( pxRxedMessage ), ( TickType_t ) 10 ) )
     {
         // pcRxedMessage now points to the struct AMessage variable posted
         // by vATask.
     }
 }

 // ... Rest of task code.
}
```

Return pdTRUE if an item was successfully received from the queue, otherwise pdFALSE.
Parameters

- `xQueue`: The handle to the queue from which the item is to be received.
- `pvBuffer`: Pointer to the buffer into which the received item will be copied.

- `xTicksToWait`: The maximum amount of time the task should block waiting for an item to receive should the queue be empty at the time of the call. xQueueReceive() will return immediately if xTicksToWait is zero and the queue is empty. The time is defined in tick periods so the constant portTICK_PERIOD_MS should be used to convert to real time if this is required.

xQueueSendToFrontFromISR(xQueue, pvItemToQueue, pxHigherPriorityTaskWoken)

This is a macro that calls xQueueGenericSendFromISR().

Post an item to the front of a queue. It is safe to use this macro from within an interrupt service routine.

Items are queued by copy not reference so it is preferable to only queue small items, especially when called from an ISR. In most cases it would be preferable to store a pointer to the item being queued.

Example usage for buffered IO (where the ISR can obtain more than one value per call):

```
void vBufferISR( void )
{
char cIn;
BaseType_t xHigherPrioritTaskWoken;

   // We have not woken a task at the start of the ISR.
   xHigherPriorityTaskWoken = pdFALSE;

   // Loop until the buffer is empty.
   do
   {
       // Obtain a byte from the buffer.
       cIn = portINPUT_BYTE( RX_REGISTER_ADDRESS );

       // Post the byte.
       xQueueSendToFrontFromISR( xRxQueue, &cIn, &xHigherPriorityTaskWoken );

   } while( portINPUT_BYTE( BUFFER_COUNT ) );

   // Now the buffer is empty we can switch context if necessary.
   if( xHigherPriorityTaskWoken )
   {
       portYIELD_FROM_ISR ();
   }
}
```

Return pdTRUE if the data was successfully sent to the queue, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `[out] pxHigherPriorityTaskWoken`: xQueueSendToFrontFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending to the queue caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xQueueSendToFromFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.

xQueueSendToBackFromISR(xQueue, pvItemToQueue, pxHigherPriorityTaskWoken)

This is a macro that calls xQueueGenericSendFromISR().

Post an item to the back of a queue. It is safe to use this macro from within an interrupt service routine.

Items are queued by copy not reference so it is preferable to only queue small items, especially when called from an ISR. In most cases it would be preferable to store a pointer to the item being queued.

Example usage for buffered IO (where the ISR can obtain more than one value per call):

```
void vBufferISR( void )
{
```

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```
char cIn;
BaseType_t xHigherPriorityTaskWoken;

   // We have not woken a task at the start of the ISR.
   xHigherPriorityTaskWoken = pdFALSE;

   // Loop until the buffer is empty.
   do
   {
       // Obtain a byte from the buffer.
       cIn = portINPUT_BYTE( RX_REGISTER_ADDRESS );

       // Post the byte.
       xQueueSendToBackFromISR( xRxQueue, &cIn, &xHigherPriorityTaskWoken );

   } while( portINPUT_BYTE( BUFFER_COUNT ) );

   // Now the buffer is empty we can switch context if necessary.
   if( xHigherPriorityTaskWoken )
   {
       portYIELD_FROM_ISR ();
   }
}
```

Return pdTRUE if the data was successfully sent to the queue, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `[out] pxHigherPriorityTaskWoken`: xQueueSendToBackFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending to the queue caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xQueueSendToBackFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.

xQueueOverwriteFromISR(xQueue, pvItemToQueue, pxHigherPriorityTaskWoken)

A version of xQueueOverwrite() that can be used in an interrupt service routine (ISR).

Only for use with queues that can hold a single item - so the queue is either empty or full.

Post an item on a queue. If the queue is already full then overwrite the value held in the queue. The item is queued by copy, not by reference.

Example usage:

```
 QueueHandle_t xQueue;

 void vFunction( void *pvParameters )
 {
    // Create a queue to hold one uint32_t value.  It is strongly
    // recommended *not* to use xQueueOverwriteFromISR() on queues that can
    // contain more than one value, and doing so will trigger an assertion
    // if configASSERT() is defined.
    xQueue = xQueueCreate( 1, sizeof( uint32_t ) );
}

void vAnInterruptHandler( void )
{
// xHigherPriorityTaskWoken must be set to pdFALSE before it is used.
BaseType_t xHigherPriorityTaskWoken = pdFALSE;
uint32_t ulVarToSend, ulValReceived;
```

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```
    // Write the value 10 to the queue using xQueueOverwriteFromISR().
    ulVarToSend = 10;
    xQueueOverwriteFromISR( xQueue, &ulVarToSend, &xHigherPriorityTaskWoken );

    // The queue is full, but calling xQueueOverwriteFromISR() again will still
    // pass because the value held in the queue will be overwritten with the
    // new value.
    ulVarToSend = 100;
    xQueueOverwriteFromISR( xQueue, &ulVarToSend, &xHigherPriorityTaskWoken );

    // Reading from the queue will now return 100.

    // ...

    if( xHigherPritytaskWoken == pdTRUE )
    {
        // Writing to the queue caused a task to unblock and the unblocked task
        // has a priority higher than or equal to the priority of the currently
        // executing task (the task this interrupt interrupted).  Perform a␣
↪context
        // switch so this interrupt returns directly to the unblocked task.
        portYIELD_FROM_ISR(); // or portEND_SWITCHING_ISR() depending on the␣
↪port.
    }
}
```

Return xQueueOverwriteFromISR() is a macro that calls xQueueGenericSendFromISR(), and therefore has the same return values as xQueueSendToFrontFromISR(). However, pdPASS is the only value that can be returned because xQueueOverwriteFromISR() will write to the queue even when the queue is already full.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `[out] pxHigherPriorityTaskWoken`: xQueueOverwriteFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending to the queue caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xQueueOverwriteFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.

xQueueSendFromISR (xQueue, pvItemToQueue, pxHigherPriorityTaskWoken)

This is a macro that calls xQueueGenericSendFromISR(). It is included for backward compatibility with versions of FreeRTOS.org that did not include the xQueueSendToBackFromISR() and xQueueSendToFrontFromISR() macros.

Post an item to the back of a queue. It is safe to use this function from within an interrupt service routine.

Items are queued by copy not reference so it is preferable to only queue small items, especially when called from an ISR. In most cases it would be preferable to store a pointer to the item being queued.

Example usage for buffered IO (where the ISR can obtain more than one value per call):

```
void vBufferISR( void )
{
char cIn;
BaseType_t xHigherPriorityTaskWoken;

   // We have not woken a task at the start of the ISR.
   xHigherPriorityTaskWoken = pdFALSE;
```

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```
    // Loop until the buffer is empty.
    do
    {
        // Obtain a byte from the buffer.
        cIn = portINPUT_BYTE( RX_REGISTER_ADDRESS );

        // Post the byte.
        xQueueSendFromISR( xRxQueue, &cIn, &xHigherPriorityTaskWoken );

    } while( portINPUT_BYTE( BUFFER_COUNT ) );

    // Now the buffer is empty we can switch context if necessary.
    if( xHigherPriorityTaskWoken )
    {
        // Actual macro used here is port specific.
        portYIELD_FROM_ISR ();
    }
}
```

Return pdTRUE if the data was successfully sent to the queue, otherwise errQUEUE_FULL.

Parameters

- `xQueue`: The handle to the queue on which the item is to be posted.
- `pvItemToQueue`: A pointer to the item that is to be placed on the queue. The size of the items the queue will hold was defined when the queue was created, so this many bytes will be copied from pvItemToQueue into the queue storage area.
- `[out] pxHigherPriorityTaskWoken`: xQueueSendFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if sending to the queue caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xQueueSendFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.

xQueueReset(xQueue)

Reset a queue back to its original empty state. pdPASS is returned if the queue is successfully reset. pdFAIL is returned if the queue could not be reset because there are tasks blocked on the queue waiting to either receive from the queue or send to the queue.

Return always returns pdPASS

Parameters

- `xQueue`: The queue to reset

Type Definitions

typedef void ***QueueHandle_t**

Type by which queues are referenced. For example, a call to xQueueCreate() returns an QueueHandle_t variable that can then be used as a parameter to xQueueSend(), xQueueReceive(), etc.

typedef void ***QueueSetHandle_t**

Type by which queue sets are referenced. For example, a call to xQueueCreateSet() returns an xQueueSet variable that can then be used as a parameter to xQueueSelectFromSet(), xQueueAddToSet(), etc.

typedef void ***QueueSetMemberHandle_t**

Queue sets can contain both queues and semaphores, so the QueueSetMemberHandle_t is defined as a type to be used where a parameter or return value can be either an QueueHandle_t or an SemaphoreHandle_t.

Semaphore API

Header File

- freertos/include/freertos/semphr.h

Macros

semBINARY_SEMAPHORE_QUEUE_LENGTH
semSEMAPHORE_QUEUE_ITEM_LENGTH

semGIVE_BLOCK_TIME

xSemaphoreCreateBinary()
Creates a new binary semaphore instance, and returns a handle by which the new semaphore can be referenced.

In many usage scenarios it is faster and more memory efficient to use a direct to task notification in place of a binary semaphore! http://www.freertos.org/RTOS-task-notifications.html

Internally, within the FreeRTOS implementation, binary semaphores use a block of memory, in which the semaphore structure is stored. If a binary semaphore is created using xSemaphoreCreateBinary() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateBinary() function. (see http://www.freertos.org/a00111.html). If a binary semaphore is created using xSemaphoreCreateBinaryStatic() then the application writer must provide the memory. xSemaphoreCreateBinaryStatic() therefore allows a binary semaphore to be created without using any dynamic memory allocation.

The old vSemaphoreCreateBinary() macro is now deprecated in favour of this xSemaphoreCreateBinary() function. Note that binary semaphores created using the vSemaphoreCreateBinary() macro are created in a state such that the first call to 'take' the semaphore would pass, whereas binary semaphores created using xSemaphoreCreateBinary() are created in a state such that the the semaphore must first be 'given' before it can be 'taken'.

Function that creates a semaphore by using the existing queue mechanism. The queue length is 1 as this is a binary semaphore. The data size is 0 as nothing is actually stored - all that is important is whether the queue is empty or full (the binary semaphore is available or not).

This type of semaphore can be used for pure synchronisation between tasks or between an interrupt and a task. The semaphore need not be given back once obtained, so one task/interrupt can continuously 'give' the semaphore while another continuously 'takes' the semaphore. For this reason this type of semaphore does not use a priority inheritance mechanism. For an alternative that does use priority inheritance see xSemaphoreCreateMutex().

Example usage:

```
SemaphoreHandle_t xSemaphore = NULL;

void vATask( void * pvParameters )
{
   // Semaphore cannot be used before a call to vSemaphoreCreateBinary ().
   // This is a macro so pass the variable in directly.
   xSemaphore = xSemaphoreCreateBinary();

   if( xSemaphore != NULL )
   {
       // The semaphore was created successfully.
       // The semaphore can now be used.
   }
}
```

Return Handle to the created semaphore.

xSemaphoreCreateBinaryStatic(pxStaticSemaphore)
Creates a new binary semaphore instance, and returns a handle by which the new semaphore can be referenced.

NOTE: In many usage scenarios it is faster and more memory efficient to use a direct to task notification in place of a binary semaphore! http://www.freertos.org/RTOS-task-notifications.html

Internally, within the FreeRTOS implementation, binary semaphores use a block of memory, in which the semaphore structure is stored. If a binary semaphore is created using xSemaphoreCreateBinary() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateBinary() function. (see http://www.freertos.org/a00111.html). If a binary semaphore is created using xSemaphoreCreateBinaryStatic() then the application writer must provide the memory. xSemaphoreCreateBinaryStatic() therefore allows a binary semaphore to be created without using any dynamic memory allocation.

This type of semaphore can be used for pure synchronisation between tasks or between an interrupt and a task. The semaphore need not be given back once obtained, so one task/interrupt can continuously 'give' the semaphore while another continuously 'takes' the semaphore. For this reason this type of semaphore does not use a priority inheritance mechanism. For an alternative that does use priority inheritance see xSemaphoreCreateMutex().

Example usage:

```
SemaphoreHandle_t xSemaphore = NULL;
StaticSemaphore_t xSemaphoreBuffer;

void vATask( void * pvParameters )
{
   // Semaphore cannot be used before a call to xSemaphoreCreateBinary() or
   // xSemaphoreCreateBinaryStatic().
   // The semaphore's data structures will be placed in the xSemaphoreBuffer
   // variable, the address of which is passed into the function.  The
   // function's parameter is not NULL, so the function will not attempt any
   // dynamic memory allocation, and therefore the function will not return
   // return NULL.
   xSemaphore = xSemaphoreCreateBinaryStatic( &xSemaphoreBuffer );

   // Rest of task code goes here.
}
```

Return If the semaphore is created then a handle to the created semaphore is returned. If pxSemaphoreBuffer is NULL then NULL is returned.

Parameters

- `pxStaticSemaphore`: Must point to a variable of type StaticSemaphore_t, which will then be used to hold the semaphore's data structure, removing the need for the memory to be allocated dynamically.

`xSemaphoreTake`(xSemaphore, xBlockTime)

Macro to obtain a semaphore. The semaphore must have previously been created with a call to vSemaphoreCreateBinary(), xSemaphoreCreateMutex() or xSemaphoreCreateCounting().

Example usage:

```
SemaphoreHandle_t xSemaphore = NULL;

// A task that creates a semaphore.
void vATask( void * pvParameters )
{
   // Create the semaphore to guard a shared resource.
   vSemaphoreCreateBinary( xSemaphore );
}

// A task that uses the semaphore.
void vAnotherTask( void * pvParameters )
{
   // ... Do other things.

   if( xSemaphore != NULL )
   {
       // See if we can obtain the semaphore.  If the semaphore is not␣
↪available
       // wait 10 ticks to see if it becomes free.
       if( xSemaphoreTake( xSemaphore, ( TickType_t ) 10 ) == pdTRUE )
       {
           // We were able to obtain the semaphore and can now access the
           // shared resource.

           // ...
```

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```
            // We have finished accessing the shared resource.  Release the
            // semaphore.
            xSemaphoreGive( xSemaphore );
        }
        else
        {
            // We could not obtain the semaphore and can therefore not access
            // the shared resource safely.
        }
    }
}
```

Return pdTRUE if the semaphore was obtained. pdFALSE if xBlockTime expired without the semaphore becoming available.

Parameters

- `xSemaphore`: A handle to the semaphore being taken - obtained when the semaphore was created.
- `xBlockTime`: The time in ticks to wait for the semaphore to become available. The macro portTICK_PERIOD_MS can be used to convert this to a real time. A block time of zero can be used to poll the semaphore. A block time of portMAX_DELAY can be used to block indefinitely (provided INCLUDE_vTaskSuspend is set to 1 in FreeRTOSConfig.h).

xSemaphoreTakeRecursive(xMutex, xBlockTime)

Macro to recursively obtain, or ‘take’, a mutex type semaphore. The mutex must have previously been created using a call to xSemaphoreCreateRecursiveMutex();

configUSE_RECURSIVE_MUTEXES must be set to 1 in FreeRTOSConfig.h for this macro to be available.

This macro must not be used on mutexes created using xSemaphoreCreateMutex().

A mutex used recursively can be ‘taken’ repeatedly by the owner. The mutex doesn’t become available again until the owner has called xSemaphoreGiveRecursive() for each successful ‘take’ request. For example, if a task successfully ‘takes’ the same mutex 5 times then the mutex will not be available to any other task until it has also ‘given’ the mutex back exactly five times.

Example usage:

```
SemaphoreHandle_t xMutex = NULL;

// A task that creates a mutex.
void vATask( void * pvParameters )
{
   // Create the mutex to guard a shared resource.
   xMutex = xSemaphoreCreateRecursiveMutex();
}

// A task that uses the mutex.
void vAnotherTask( void * pvParameters )
{
   // ... Do other things.

   if( xMutex != NULL )
   {
       // See if we can obtain the mutex.  If the mutex is not available
       // wait 10 ticks to see if it becomes free.
       if( xSemaphoreTakeRecursive( xSemaphore, ( TickType_t ) 10 ) == pdTRUE )
       {
           // We were able to obtain the mutex and can now access the
           // shared resource.

           // ...
           // For some reason due to the nature of the code further calls to
```

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```
            // xSemaphoreTakeRecursive() are made on the same mutex.  In real
            // code these would not be just sequential calls as this would make
            // no sense.  Instead the calls are likely to be buried inside
            // a more complex call structure.
            xSemaphoreTakeRecursive( xMutex, ( TickType_t ) 10 );
            xSemaphoreTakeRecursive( xMutex, ( TickType_t ) 10 );

            // The mutex has now been 'taken' three times, so will not be
            // available to another task until it has also been given back
            // three times.  Again it is unlikely that real code would have
            // these calls sequentially, but instead buried in a more complex
            // call structure.  This is just for illustrative purposes.
            xSemaphoreGiveRecursive( xMutex );
            xSemaphoreGiveRecursive( xMutex );
            xSemaphoreGiveRecursive( xMutex );

            // Now the mutex can be taken by other tasks.
        }
        else
        {
            // We could not obtain the mutex and can therefore not access
            // the shared resource safely.
        }
    }
}
```

Return pdTRUE if the semaphore was obtained. pdFALSE if xBlockTime expired without the semaphore becoming available.

Parameters

- `xMutex`: A handle to the mutex being obtained. This is the handle returned by xSemaphoreCreateRecursiveMutex();
- `xBlockTime`: The time in ticks to wait for the semaphore to become available. The macro portTICK_PERIOD_MS can be used to convert this to a real time. A block time of zero can be used to poll the semaphore. If the task already owns the semaphore then xSemaphoreTakeRecursive() will return immediately no matter what the value of xBlockTime.

xSemaphoreGive(xSemaphore)

Macro to release a semaphore. The semaphore must have previously been created with a call to vSemaphoreCreateBinary(), xSemaphoreCreateMutex() or xSemaphoreCreateCounting(). and obtained using sSemaphoreTake().

This macro must not be used from an ISR. See xSemaphoreGiveFromISR () for an alternative which can be used from an ISR.

This macro must also not be used on semaphores created using xSemaphoreCreateRecursiveMutex().

Example usage:

```
SemaphoreHandle_t xSemaphore = NULL;

void vATask( void * pvParameters )
{
   // Create the semaphore to guard a shared resource.
   vSemaphoreCreateBinary( xSemaphore );

   if( xSemaphore != NULL )
   {
       if( xSemaphoreGive( xSemaphore ) != pdTRUE )
       {
           // We would expect this call to fail because we cannot give
           // a semaphore without first "taking" it!
```

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```
        }

        // Obtain the semaphore - don't block if the semaphore is not
        // immediately available.
        if( xSemaphoreTake( xSemaphore, ( TickType_t ) 0 ) )
        {
            // We now have the semaphore and can access the shared resource.

            // ...

            // We have finished accessing the shared resource so can free the
            // semaphore.
            if( xSemaphoreGive( xSemaphore ) != pdTRUE )
            {
                // We would not expect this call to fail because we must have
                // obtained the semaphore to get here.
            }
        }
    }
}
```

Return pdTRUE if the semaphore was released. pdFALSE if an error occurred. Semaphores are implemented using queues. An error can occur if there is no space on the queue to post a message - indicating that the semaphore was not first obtained correctly.

Parameters

- `xSemaphore`: A handle to the semaphore being released. This is the handle returned when the semaphore was created.

xSemaphoreGiveRecursive(xMutex)

Macro to recursively release, or ‘give’, a mutex type semaphore. The mutex must have previously been created using a call to xSemaphoreCreateRecursiveMutex();

configUSE_RECURSIVE_MUTEXES must be set to 1 in FreeRTOSConfig.h for this macro to be available.

This macro must not be used on mutexes created using xSemaphoreCreateMutex().

A mutex used recursively can be ‘taken’ repeatedly by the owner. The mutex doesn’t become available again until the owner has called xSemaphoreGiveRecursive() for each successful ‘take’ request. For example, if a task successfully ‘takes’ the same mutex 5 times then the mutex will not be available to any other task until it has also ‘given’ the mutex back exactly five times.

Example usage:

```
SemaphoreHandle_t xMutex = NULL;

// A task that creates a mutex.
void vATask( void * pvParameters )
{
   // Create the mutex to guard a shared resource.
   xMutex = xSemaphoreCreateRecursiveMutex();
}

// A task that uses the mutex.
void vAnotherTask( void * pvParameters )
{
   // ... Do other things.

   if( xMutex != NULL )
   {
       // See if we can obtain the mutex.  If the mutex is not available
       // wait 10 ticks to see if it becomes free.
       if( xSemaphoreTakeRecursive( xMutex, ( TickType_t ) 10 ) == pdTRUE )
```

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```
        {
            // We were able to obtain the mutex and can now access the
            // shared resource.

            // ...
            // For some reason due to the nature of the code further calls to
            // xSemaphoreTakeRecursive() are made on the same mutex.  In real
            // code these would not be just sequential calls as this would make
            // no sense.  Instead the calls are likely to be buried inside
            // a more complex call structure.
            xSemaphoreTakeRecursive( xMutex, ( TickType_t ) 10 );
            xSemaphoreTakeRecursive( xMutex, ( TickType_t ) 10 );

            // The mutex has now been 'taken' three times, so will not be
            // available to another task until it has also been given back
            // three times.  Again it is unlikely that real code would have
            // these calls sequentially, it would be more likely that the calls
            // to xSemaphoreGiveRecursive() would be called as a call stack
            // unwound.  This is just for demonstrative purposes.
            xSemaphoreGiveRecursive( xMutex );
            xSemaphoreGiveRecursive( xMutex );
            xSemaphoreGiveRecursive( xMutex );

            // Now the mutex can be taken by other tasks.
        }
        else
        {
            // We could not obtain the mutex and can therefore not access
            // the shared resource safely.
        }
    }
}
```

Return pdTRUE if the semaphore was given.

Parameters

- `xMutex`: A handle to the mutex being released, or 'given'. This is the handle returned by xSemaphoreCreateMutex();

xSemaphoreGiveFromISR (xSemaphore, pxHigherPriorityTaskWoken)

Macro to release a semaphore. The semaphore must have previously been created with a call to vSemaphoreCreateBinary() or xSemaphoreCreateCounting().

Mutex type semaphores (those created using a call to xSemaphoreCreateMutex()) must not be used with this macro.

This macro can be used from an ISR.

Example usage:

```
\#define LONG_TIME 0xffff
\#define TICKS_TO_WAIT 10
SemaphoreHandle_t xSemaphore = NULL;

// Repetitive task.
void vATask( void * pvParameters )
{
   for( ;; )
   {
       // We want this task to run every 10 ticks of a timer.  The semaphore
       // was created before this task was started.

       // Block waiting for the semaphore to become available.
```

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```
        if( xSemaphoreTake( xSemaphore, LONG_TIME ) == pdTRUE )
        {
            // It is time to execute.

            // ...

            // We have finished our task.  Return to the top of the loop where
            // we will block on the semaphore until it is time to execute
            // again.  Note when using the semaphore for synchronisation with an
            // ISR in this manner there is no need to 'give' the semaphore back.
        }
    }
}

// Timer ISR
void vTimerISR( void * pvParameters )
{
static uint8_t ucLocalTickCount = 0;
static BaseType_t xHigherPriorityTaskWoken;

    // A timer tick has occurred.

    // ... Do other time functions.

    // Is it time for vATask () to run?
    xHigherPriorityTaskWoken = pdFALSE;
    ucLocalTickCount++;
    if( ucLocalTickCount >= TICKS_TO_WAIT )
    {
        // Unblock the task by releasing the semaphore.
        xSemaphoreGiveFromISR( xSemaphore, &xHigherPriorityTaskWoken );

        // Reset the count so we release the semaphore again in 10 ticks time.
        ucLocalTickCount = 0;
    }

    if( xHigherPriorityTaskWoken != pdFALSE )
    {
        // We can force a context switch here.  Context switching from an
        // ISR uses port specific syntax.  Check the demo task for your port
        // to find the syntax required.
    }
}
```

Return pdTRUE if the semaphore was successfully given, otherwise errQUEUE_FULL.

Parameters

- `xSemaphore`: A handle to the semaphore being released. This is the handle returned when the semaphore was created.
- `[out] pxHigherPriorityTaskWoken`: xSemaphoreGiveFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if giving the semaphore caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xSemaphoreGiveFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.

xSemaphoreTakeFromISR (xSemaphore, pxHigherPriorityTaskWoken)

Macro to take a semaphore from an ISR. The semaphore must have previously been created with a call to vSemaphoreCreateBinary() or xSemaphoreCreateCounting().

Mutex type semaphores (those created using a call to xSemaphoreCreateMutex()) must not be used with this macro.

This macro can be used from an ISR, however taking a semaphore from an ISR is not a common operation. It is likely to only be useful when taking a counting semaphore when an interrupt is obtaining an object from a

resource pool (when the semaphore count indicates the number of resources available).

Return pdTRUE if the semaphore was successfully taken, otherwise pdFALSE
Parameters

- `xSemaphore`: A handle to the semaphore being taken. This is the handle returned when the semaphore was created.
- `[out] pxHigherPriorityTaskWoken`: xSemaphoreTakeFromISR() will set *pxHigherPriorityTaskWoken to pdTRUE if taking the semaphore caused a task to unblock, and the unblocked task has a priority higher than the currently running task. If xSemaphoreTakeFromISR() sets this value to pdTRUE then a context switch should be requested before the interrupt is exited.

xSemaphoreCreateMutex()

Macro that implements a mutex semaphore by using the existing queue mechanism.

Internally, within the FreeRTOS implementation, mutex semaphores use a block of memory, in which the mutex structure is stored. If a mutex is created using xSemaphoreCreateMutex() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateMutex() function. (see http://www.freertos.org/a00111.html). If a mutex is created using xSemaphoreCreateMutexStatic() then the application writer must provided the memory. xSemaphoreCreateMutexStatic() therefore allows a mutex to be created without using any dynamic memory allocation.

Mutexes created using this function can be accessed using the xSemaphoreTake() and xSemaphoreGive() macros. The xSemaphoreTakeRecursive() and xSemaphoreGiveRecursive() macros must not be used.

This type of semaphore uses a priority inheritance mechanism so a task ‘taking’ a semaphore MUST ALWAYS ‘give’ the semaphore back once the semaphore it is no longer required.

Mutex type semaphores cannot be used from within interrupt service routines.

See vSemaphoreCreateBinary() for an alternative implementation that can be used for pure synchronisation (where one task or interrupt always ‘gives’ the semaphore and another always ‘takes’ the semaphore) and from within interrupt service routines.

Example usage:

```
SemaphoreHandle_t xSemaphore;

void vATask( void * pvParameters )
{
   // Semaphore cannot be used before a call to xSemaphoreCreateMutex().
   // This is a macro so pass the variable in directly.
   xSemaphore = xSemaphoreCreateMutex();

   if( xSemaphore != NULL )
   {
       // The semaphore was created successfully.
       // The semaphore can now be used.
   }
}
```

Return If the mutex was successfully created then a handle to the created semaphore is returned. If there was not enough heap to allocate the mutex data structures then NULL is returned.

xSemaphoreCreateMutexStatic(pxMutexBuffer)

Creates a new mutex type semaphore instance, and returns a handle by which the new mutex can be referenced.

Internally, within the FreeRTOS implementation, mutex semaphores use a block of memory, in which the mutex structure is stored. If a mutex is created using xSemaphoreCreateMutex() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateMutex() function. (see http://www.freertos.org/a00111.html). If a mutex is created using xSemaphoreCreateMutexStatic() then the application writer must provided the memory. xSemaphoreCreateMutexStatic() therefore allows a mutex to be created without using any dynamic memory allocation.

Mutexes created using this function can be accessed using the xSemaphoreTake() and xSemaphoreGive() macros. The xSemaphoreTakeRecursive() and xSemaphoreGiveRecursive() macros must not be used.

This type of semaphore uses a priority inheritance mechanism so a task ‘taking’ a semaphore MUST ALWAYS ‘give’ the semaphore back once the semaphore it is no longer required.

Mutex type semaphores cannot be used from within interrupt service routines.

See xSemaphoreCreateBinary() for an alternative implementation that can be used for pure synchronisation (where one task or interrupt always ‘gives’ the semaphore and another always ‘takes’ the semaphore) and from within interrupt service routines.

Example usage:

```
SemaphoreHandle_t xSemaphore;
StaticSemaphore_t xMutexBuffer;

void vATask( void * pvParameters )
{
   // A mutex cannot be used before it has been created.  xMutexBuffer is
   // into xSemaphoreCreateMutexStatic() so no dynamic memory allocation is
   // attempted.
   xSemaphore = xSemaphoreCreateMutexStatic( &xMutexBuffer );

   // As no dynamic memory allocation was performed, xSemaphore cannot be NULL,
   // so there is no need to check it.
}
```

Return If the mutex was successfully created then a handle to the created mutex is returned. If pxMutexBuffer was NULL then NULL is returned.

Parameters

- `pxMutexBuffer`: Must point to a variable of type StaticSemaphore_t, which will be used to hold the mutex’s data structure, removing the need for the memory to be allocated dynamically.

`xSemaphoreCreateRecursiveMutex`()

Creates a new recursive mutex type semaphore instance, and returns a handle by which the new recursive mutex can be referenced.

Internally, within the FreeRTOS implementation, recursive mutexs use a block of memory, in which the mutex structure is stored. If a recursive mutex is created using xSemaphoreCreateRecursiveMutex() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateRecursiveMutex() function. (see http://www.freertos.org/a00111.html). If a recursive mutex is created using xSemaphoreCreateRecursiveMutexStatic() then the application writer must provide the memory that will get used by the mutex. xSemaphoreCreateRecursiveMutexStatic() therefore allows a recursive mutex to be created without using any dynamic memory allocation.

Mutexes created using this macro can be accessed using the xSemaphoreTakeRecursive() and xSemaphoreGiveRecursive() macros. The xSemaphoreTake() and xSemaphoreGive() macros must not be used.

A mutex used recursively can be ‘taken’ repeatedly by the owner. The mutex doesn’t become available again until the owner has called xSemaphoreGiveRecursive() for each successful ‘take’ request. For example, if a task successfully ‘takes’ the same mutex 5 times then the mutex will not be available to any other task until it has also ‘given’ the mutex back exactly five times.

This type of semaphore uses a priority inheritance mechanism so a task ‘taking’ a semaphore MUST ALWAYS ‘give’ the semaphore back once the semaphore it is no longer required.

Mutex type semaphores cannot be used from within interrupt service routines.

See vSemaphoreCreateBinary() for an alternative implementation that can be used for pure synchronisation (where one task or interrupt always ‘gives’ the semaphore and another always ‘takes’ the semaphore) and from within interrupt service routines.

Example usage:

```
SemaphoreHandle_t xSemaphore;
```

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```
void vATask( void * pvParameters )
{
   // Semaphore cannot be used before a call to xSemaphoreCreateMutex().
   // This is a macro so pass the variable in directly.
   xSemaphore = xSemaphoreCreateRecursiveMutex();

   if( xSemaphore != NULL )
   {
       // The semaphore was created successfully.
       // The semaphore can now be used.
   }
}
```

Return xSemaphore Handle to the created mutex semaphore. Should be of type SemaphoreHandle_t.

xSemaphoreCreateRecursiveMutexStatic(pxStaticSemaphore)

Creates a new recursive mutex type semaphore instance, and returns a handle by which the new recursive mutex can be referenced.

Internally, within the FreeRTOS implementation, recursive mutexs use a block of memory, in which the mutex structure is stored. If a recursive mutex is created using xSemaphoreCreateRecursiveMutex() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateRecursiveMutex() function. (see http://www.freertos.org/a00111.html). If a recursive mutex is created using xSemaphoreCreateRecursiveMutexStatic() then the application writer must provide the memory that will get used by the mutex. xSemaphoreCreateRecursiveMutexStatic() therefore allows a recursive mutex to be created without using any dynamic memory allocation.

Mutexes created using this macro can be accessed using the xSemaphoreTakeRecursive() and xSemaphoreGiveRecursive() macros. The xSemaphoreTake() and xSemaphoreGive() macros must not be used.

A mutex used recursively can be ‘taken’ repeatedly by the owner. The mutex doesn’t become available again until the owner has called xSemaphoreGiveRecursive() for each successful ‘take’ request. For example, if a task successfully ‘takes’ the same mutex 5 times then the mutex will not be available to any other task until it has also ‘given’ the mutex back exactly five times.

This type of semaphore uses a priority inheritance mechanism so a task ‘taking’ a semaphore MUST ALWAYS ‘give’ the semaphore back once the semaphore it is no longer required.

Mutex type semaphores cannot be used from within interrupt service routines.

See xSemaphoreCreateBinary() for an alternative implementation that can be used for pure synchronisation (where one task or interrupt always ‘gives’ the semaphore and another always ‘takes’ the semaphore) and from within interrupt service routines.

Example usage:

```
SemaphoreHandle_t xSemaphore;
StaticSemaphore_t xMutexBuffer;

void vATask( void * pvParameters )
{
   // A recursive semaphore cannot be used before it is created.  Here a
   // recursive mutex is created using xSemaphoreCreateRecursiveMutexStatic().
   // The address of xMutexBuffer is passed into the function, and will hold
   // the mutexes data structures - so no dynamic memory allocation will be
   // attempted.
   xSemaphore = xSemaphoreCreateRecursiveMutexStatic( &xMutexBuffer );

   // As no dynamic memory allocation was performed, xSemaphore cannot be NULL,
   // so there is no need to check it.
}
```

Return If the recursive mutex was successfully created then a handle to the created recursive mutex is returned. If pxMutexBuffer was NULL then NULL is returned.

Parameters

- `pxStaticSemaphore`: Must point to a variable of type StaticSemaphore_t, which will then be used to hold the recursive mutex's data structure, removing the need for the memory to be allocated dynamically.

`xSemaphoreCreateCounting`(uxMaxCount, uxInitialCount)

Creates a new counting semaphore instance, and returns a handle by which the new counting semaphore can be referenced.

In many usage scenarios it is faster and more memory efficient to use a direct to task notification in place of a counting semaphore! http://www.freertos.org/RTOS-task-notifications.html

Internally, within the FreeRTOS implementation, counting semaphores use a block of memory, in which the counting semaphore structure is stored. If a counting semaphore is created using xSemaphoreCreateCounting() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateCounting() function. (see http://www.freertos.org/a00111.html). If a counting semaphore is created using xSemaphoreCreateCountingStatic() then the application writer can instead optionally provide the memory that will get used by the counting semaphore. xSemaphoreCreateCountingStatic() therefore allows a counting semaphore to be created without using any dynamic memory allocation.

Counting semaphores are typically used for two things:

1) Counting events.

In this usage scenario an event handler will 'give' a semaphore each time an event occurs (incrementing the semaphore count value), and a handler task will 'take' a semaphore each time it processes an event (decrementing the semaphore count value). The count value is therefore the difference between the number of events that have occurred and the number that have been processed. In this case it is desirable for the initial count value to be zero.

2) Resource management.

In this usage scenario the count value indicates the number of resources available. To obtain control of a resource a task must first obtain a semaphore - decrementing the semaphore count value. When the count value reaches zero there are no free resources. When a task finishes with the resource it 'gives' the semaphore back - incrementing the semaphore count value. In this case it is desirable for the initial count value to be equal to the maximum count value, indicating that all resources are free.

Example usage:

```
SemaphoreHandle_t xSemaphore;

void vATask( void * pvParameters )
{
SemaphoreHandle_t xSemaphore = NULL;

   // Semaphore cannot be used before a call to xSemaphoreCreateCounting().
   // The max value to which the semaphore can count should be 10, and the
   // initial value assigned to the count should be 0.
   xSemaphore = xSemaphoreCreateCounting( 10, 0 );

   if( xSemaphore != NULL )
   {
        // The semaphore was created successfully.
        // The semaphore can now be used.
   }
}
```

Return Handle to the created semaphore. Null if the semaphore could not be created.

Parameters

- `uxMaxCount`: The maximum count value that can be reached. When the semaphore reaches this value it can no longer be 'given'.

- uxInitialCount: The count value assigned to the semaphore when it is created.

xSemaphoreCreateCountingStatic(uxMaxCount, uxInitialCount, pxSemaphoreBuffer)

Creates a new counting semaphore instance, and returns a handle by which the new counting semaphore can be referenced.

In many usage scenarios it is faster and more memory efficient to use a direct to task notification in place of a counting semaphore! http://www.freertos.org/RTOS-task-notifications.html

Internally, within the FreeRTOS implementation, counting semaphores use a block of memory, in which the counting semaphore structure is stored. If a counting semaphore is created using xSemaphoreCreateCounting() then the required memory is automatically dynamically allocated inside the xSemaphoreCreateCounting() function. (see http://www.freertos.org/a00111.html). If a counting semaphore is created using xSemaphoreCreateCountingStatic() then the application writer must provide the memory. xSemaphoreCreateCountingStatic() therefore allows a counting semaphore to be created without using any dynamic memory allocation.

Counting semaphores are typically used for two things:

1) Counting events.

In this usage scenario an event handler will ‘give’ a semaphore each time an event occurs (incrementing the semaphore count value), and a handler task will ‘take’ a semaphore each time it processes an event (decrementing the semaphore count value). The count value is therefore the difference between the number of events that have occurred and the number that have been processed. In this case it is desirable for the initial count value to be zero.

2) Resource management.

In this usage scenario the count value indicates the number of resources available. To obtain control of a resource a task must first obtain a semaphore - decrementing the semaphore count value. When the count value reaches zero there are no free resources. When a task finishes with the resource it ‘gives’ the semaphore back - incrementing the semaphore count value. In this case it is desirable for the initial count value to be equal to the maximum count value, indicating that all resources are free.

Example usage:

```
SemaphoreHandle_t xSemaphore;
StaticSemaphore_t xSemaphoreBuffer;

void vATask( void * pvParameters )
{
SemaphoreHandle_t xSemaphore = NULL;

   // Counting semaphore cannot be used before they have been created.  Create
   // a counting semaphore using xSemaphoreCreateCountingStatic().  The max
   // value to which the semaphore can count is 10, and the initial value
   // assigned to the count will be 0.  The address of xSemaphoreBuffer is
   // passed in and will be used to hold the semaphore structure, so no dynamic
   // memory allocation will be used.
   xSemaphore = xSemaphoreCreateCounting( 10, 0, &xSemaphoreBuffer );

   // No memory allocation was attempted so xSemaphore cannot be NULL, so there
   // is no need to check its value.
}
```

Return If the counting semaphore was successfully created then a handle to the created counting semaphore is returned. If pxSemaphoreBuffer was NULL then NULL is returned.

Parameters

- uxMaxCount: The maximum count value that can be reached. When the semaphore reaches this value it can no longer be ‘given’.
- uxInitialCount: The count value assigned to the semaphore when it is created.
- pxSemaphoreBuffer: Must point to a variable of type StaticSemaphore_t, which will then be used to hold the semaphore’s data structure, removing the need for the memory to be allocated

dynamically.

vSemaphoreDelete(xSemaphore)

Delete a semaphore. This function must be used with care. For example, do not delete a mutex type semaphore if the mutex is held by a task.

Parameters

- `xSemaphore`: A handle to the semaphore to be deleted.

xSemaphoreGetMutexHolder(xSemaphore)

If xMutex is indeed a mutex type semaphore, return the current mutex holder. If xMutex is not a mutex type semaphore, or the mutex is available (not held by a task), return NULL.

Note: This is a good way of determining if the calling task is the mutex holder, but not a good way of determining the identity of the mutex holder as the holder may change between the function exiting and the returned value being tested.

uxSemaphoreGetCount(xSemaphore)

If the semaphore is a counting semaphore then uxSemaphoreGetCount() returns its current count value. If the semaphore is a binary semaphore then uxSemaphoreGetCount() returns 1 if the semaphore is available, and 0 if the semaphore is not available.

Type Definitions

typedef *QueueHandle_t* **SemaphoreHandle_t**

Timer API

Header File

- freertos/include/freertos/timers.h

Functions

TimerHandle_t **xTimerCreate**(**const** char ***const** *pcTimerName*, **const** TickType_t *xTimerPeriodInTicks*, **const** UBaseType_t *uxAutoReload*, void ***const** *pvTimerID*, *TimerCallbackFunction_t pxCallbackFunction*)

Creates a new software timer instance, and returns a handle by which the created software timer can be referenced.

Internally, within the FreeRTOS implementation, software timers use a block of memory, in which the timer data structure is stored. If a software timer is created using xTimerCreate() then the required memory is automatically dynamically allocated inside the xTimerCreate() function. (see http://www.freertos.org/a00111.html). If a software timer is created using xTimerCreateStatic() then the application writer must provide the memory that will get used by the software timer. xTimerCreateStatic() therefore allows a software timer to be created without using any dynamic memory allocation.

Timers are created in the dormant state. The xTimerStart(), xTimerReset(), xTimerStartFromISR(), xTimerResetFromISR(), xTimerChangePeriod() and xTimerChangePeriodFromISR() API functions can all be used to transition a timer into the active state.

Example usage:

```
#define NUM_TIMERS 5

// An array to hold handles to the created timers.
TimerHandle_t xTimers[ NUM_TIMERS ];

// An array to hold a count of the number of times each timer expires.
int32_t lExpireCounters[ NUM_TIMERS ] = { 0 };

// Define a callback function that will be used by multiple timer instances.
// The callback function does nothing but count the number of times the
```

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```
// associated timer expires, and stop the timer once the timer has expired
// 10 times.
void vTimerCallback( TimerHandle_t pxTimer )
{
int32_t lArrayIndex;
const int32_t xMaxExpiryCountBeforeStopping = 10;

       // Optionally do something if the pxTimer parameter is NULL.
       configASSERT( pxTimer );

    // Which timer expired?
    lArrayIndex = ( int32_t ) pvTimerGetTimerID( pxTimer );

    // Increment the number of times that pxTimer has expired.
    lExpireCounters[ lArrayIndex ] += 1;

    // If the timer has expired 10 times then stop it from running.
    if( lExpireCounters[ lArrayIndex ] == xMaxExpiryCountBeforeStopping )
    {
        // Do not use a block time if calling a timer API function from a
        // timer callback function, as doing so could cause a deadlock!
        xTimerStop( pxTimer, 0 );
    }
}

void main( void )
{
int32_t x;

    // Create then start some timers.  Starting the timers before the scheduler
    // has been started means the timers will start running immediately that
    // the scheduler starts.
    for( x = 0; x < NUM_TIMERS; x++ )
    {
        xTimers[ x ] = xTimerCreate(    "Timer",       // Just a text name,␣
↪not used by the kernel.
                                        ( 100 * x ),   // The timer period in␣
↪ticks.
                                        pdTRUE,        // The timers will auto-
↪reload themselves when they expire.
                                        ( void * ) x,  // Assign each timer a␣
↪unique id equal to its array index.
                                        vTimerCallback // Each timer calls the␣
↪same callback when it expires.
                                    );

        if( xTimers[ x ] == NULL )
        {
            // The timer was not created.
        }
        else
        {
            // Start the timer.  No block time is specified, and even if one␣
↪was
            // it would be ignored because the scheduler has not yet been
            // started.
            if( xTimerStart( xTimers[ x ], 0 ) != pdPASS )
            {
                // The timer could not be set into the Active state.
            }
        }
```

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```
    }

    // ...
    // Create tasks here.
    // ...

    // Starting the scheduler will start the timers running as they have␣
↪already
    // been set into the active state.
    vTaskStartScheduler();

    // Should not reach here.
    for( ;; );
}
```

Return If the timer is successfully created then a handle to the newly created timer is returned. If the timer cannot be created (because either there is insufficient FreeRTOS heap remaining to allocate the timer structures, or the timer period was set to 0) then NULL is returned.

Parameters

- `pcTimerName`: A text name that is assigned to the timer. This is done purely to assist debugging. The kernel itself only ever references a timer by its handle, and never by its name.
- `xTimerPeriodInTicks`: The timer period. The time is defined in tick periods so the constant portTICK_PERIOD_MS can be used to convert a time that has been specified in milliseconds. For example, if the timer must expire after 100 ticks, then xTimerPeriodInTicks should be set to 100. Alternatively, if the timer must expire after 500ms, then xPeriod can be set to (500 / portTICK_PERIOD_MS) provided configTICK_RATE_HZ is less than or equal to 1000.
- `uxAutoReload`: If uxAutoReload is set to pdTRUE then the timer will expire repeatedly with a frequency set by the xTimerPeriodInTicks parameter. If uxAutoReload is set to pdFALSE then the timer will be a one-shot timer and enter the dormant state after it expires.
- `pvTimerID`: An identifier that is assigned to the timer being created. Typically this would be used in the timer callback function to identify which timer expired when the same callback function is assigned to more than one timer.
- `pxCallbackFunction`: The function to call when the timer expires. Callback functions must have the prototype defined by TimerCallbackFunction_t, which is "void vCallbackFunction(TimerHandle_t xTimer);" .

TimerHandle_t **xTimerCreateStatic**(**const** char ***const** *pcTimerName*, **const** TickType_t *xTimerPeriodInTicks*, **const** UBaseType_t *uxAutoReload*, void ***const** *pvTimerID*, *TimerCallbackFunction_t* *pxCallbackFunction*, StaticTimer_t **pxTimerBuffer*)

Creates a new software timer instance, and returns a handle by which the created software timer can be referenced.

Internally, within the FreeRTOS implementation, software timers use a block of memory, in which the timer data structure is stored. If a software timer is created using xTimerCreate() then the required memory is automatically dynamically allocated inside the xTimerCreate() function. (see http://www.freertos.org/a00111.html). If a software timer is created using xTimerCreateStatic() then the application writer must provide the memory that will get used by the software timer. xTimerCreateStatic() therefore allows a software timer to be created without using any dynamic memory allocation.

Timers are created in the dormant state. The xTimerStart(), xTimerReset(), xTimerStartFromISR(), xTimerResetFromISR(), xTimerChangePeriod() and xTimerChangePeriodFromISR() API functions can all be used to transition a timer into the active state.

Example usage:

```
// The buffer used to hold the software timer's data structure.
static StaticTimer_t xTimerBuffer;

// A variable that will be incremented by the software timer's callback
```

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```
// function.
UBaseType_t uxVariableToIncrement = 0;

// A software timer callback function that increments a variable passed to
// it when the software timer was created.  After the 5th increment the
// callback function stops the software timer.
static void prvTimerCallback( TimerHandle_t xExpiredTimer )
{
UBaseType_t *puxVariableToIncrement;
BaseType_t xReturned;

    // Obtain the address of the variable to increment from the timer ID.
    puxVariableToIncrement = ( UBaseType_t * ) pvTimerGetTimerID(␣
↪xExpiredTimer );

    // Increment the variable to show the timer callback has executed.
    ( *puxVariableToIncrement )++;

    // If this callback has executed the required number of times, stop the
    // timer.
    if( *puxVariableToIncrement == 5 )
    {
        // This is called from a timer callback so must not block.
        xTimerStop( xExpiredTimer, staticDONT_BLOCK );
    }
}


void main( void )
{
    // Create the software time.  xTimerCreateStatic() has an extra parameter
    // than the normal xTimerCreate() API function.  The parameter is a pointer
    // to the StaticTimer_t structure that will hold the software timer
    // structure.  If the parameter is passed as NULL then the structure will␣
↪be
    // allocated dynamically, just as if xTimerCreate() had been called.
    xTimer = xTimerCreateStatic( "T1",             // Text name for the task. ␣
↪Helps debugging only.  Not used by FreeRTOS.
                                 xTimerPeriod,     // The period of the timer␣
↪in ticks.
                                 pdTRUE,           // This is an auto-reload␣
↪timer.
                                 ( void * ) &uxVariableToIncrement,    // A␣
↪variable incremented by the software timer's callback function
                                 prvTimerCallback, // The function to execute␣
↪when the timer expires.
                                 &xTimerBuffer );  // The buffer that will␣
↪hold the software timer structure.

    // The scheduler has not started yet so a block time is not used.
    xReturned = xTimerStart( xTimer, 0 );

    // ...
    // Create tasks here.
    // ...

    // Starting the scheduler will start the timers running as they have␣
↪already
    // been set into the active state.
    vTaskStartScheduler();
```

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```
    // Should not reach here.
    for( ;; );
}
```

Return If the timer is created then a handle to the created timer is returned. If pxTimerBuffer was NULL then NULL is returned.

Parameters

- `pcTimerName`: A text name that is assigned to the timer. This is done purely to assist debugging. The kernel itself only ever references a timer by its handle, and never by its name.
- `xTimerPeriodInTicks`: The timer period. The time is defined in tick periods so the constant portTICK_PERIOD_MS can be used to convert a time that has been specified in milliseconds. For example, if the timer must expire after 100 ticks, then xTimerPeriodInTicks should be set to 100. Alternatively, if the timer must expire after 500ms, then xPeriod can be set to (500 / portTICK_PERIOD_MS) provided configTICK_RATE_HZ is less than or equal to 1000.
- `uxAutoReload`: If uxAutoReload is set to pdTRUE then the timer will expire repeatedly with a frequency set by the xTimerPeriodInTicks parameter. If uxAutoReload is set to pdFALSE then the timer will be a one-shot timer and enter the dormant state after it expires.
- `pvTimerID`: An identifier that is assigned to the timer being created. Typically this would be used in the timer callback function to identify which timer expired when the same callback function is assigned to more than one timer.
- `pxCallbackFunction`: The function to call when the timer expires. Callback functions must have the prototype defined by TimerCallbackFunction_t, which is "void vCallbackFunction(TimerHandle_t xTimer);" .
- `pxTimerBuffer`: Must point to a variable of type StaticTimer_t, which will be then be used to hold the software timer's data structures, removing the need for the memory to be allocated dynamically.

void ***pvTimerGetTimerID**(*TimerHandle_t xTimer*)

Returns the ID assigned to the timer.

IDs are assigned to timers using the pvTimerID parameter of the call to xTimerCreated() that was used to create the timer.

If the same callback function is assigned to multiple timers then the timer ID can be used within the callback function to identify which timer actually expired.

Example usage:

Return The ID assigned to the timer being queried.

Parameters

- `xTimer`: The timer being queried.

See the xTimerCreate() API function example usage scenario.

void **vTimerSetTimerID**(*TimerHandle_t xTimer*, void **pvNewID*)

Sets the ID assigned to the timer.

IDs are assigned to timers using the pvTimerID parameter of the call to xTimerCreated() that was used to create the timer.

If the same callback function is assigned to multiple timers then the timer ID can be used as time specific (timer local) storage.

Example usage:

Parameters

- `xTimer`: The timer being updated.
- `pvNewID`: The ID to assign to the timer.

See the xTimerCreate() API function example usage scenario.

BaseType_t **xTimerIsTimerActive**(*TimerHandle_t xTimer*)

Queries a timer to see if it is active or dormant.

A timer will be dormant if:

```
1) It has been created but not started, or

2) It is an expired one-shot timer that has not been restarted.
```

Timers are created in the dormant state. The xTimerStart(), xTimerReset(), xTimerStartFromISR(), xTimerResetFromISR(), xTimerChangePeriod() and xTimerChangePeriodFromISR() API functions can all be used to transition a timer into the active state.

Example usage:

```
// This function assumes xTimer has already been created.
void vAFunction( TimerHandle_t xTimer )
{
    if( xTimerIsTimerActive( xTimer ) != pdFALSE ) // or more simply and␣
↪equivalently "if( xTimerIsTimerActive( xTimer ) )"
    {
        // xTimer is active, do something.
    }
    else
    {
        // xTimer is not active, do something else.
    }
}
```

Return pdFALSE will be returned if the timer is dormant. A value other than pdFALSE will be returned if the timer is active.

Parameters

- `xTimer`: The timer being queried.

TaskHandle_t **xTimerGetTimerDaemonTaskHandle**(void)

xTimerGetTimerDaemonTaskHandle() is only available if INCLUDE_xTimerGetTimerDaemonTaskHandle is set to 1 in FreeRTOSConfig.h.

Simply returns the handle of the timer service/daemon task. It it not valid to call xTimerGetTimerDaemonTaskHandle() before the scheduler has been started.

TickType_t **xTimerGetPeriod**(*TimerHandle_t xTimer*)

Returns the period of a timer.

Return The period of the timer in ticks.

Parameters

- `xTimer`: The handle of the timer being queried.

TickType_t **xTimerGetExpiryTime**(*TimerHandle_t xTimer*)

Returns the time in ticks at which the timer will expire. If this is less than the current tick count then the expiry time has overflowed from the current time.

Return If the timer is running then the time in ticks at which the timer will next expire is returned. If the timer is not running then the return value is undefined.

Parameters

- `xTimer`: The handle of the timer being queried.

BaseType_t **xTimerPendFunctionCallFromISR**(*PendedFunction_t xFunctionToPend*, void **pvParameter1*, uint32_t *ulParameter2*, BaseType_t **pxHigherPriorityTaskWoken*)

Used from application interrupt service routines to defer the execution of a function to the RTOS daemon task (the timer service task, hence this function is implemented in timers.c and is prefixed with 'Timer').

Ideally an interrupt service routine (ISR) is kept as short as possible, but sometimes an ISR either has a lot of processing to do, or needs to perform processing that is not deterministic. In these cases xTimerPendFunctionCallFromISR() can be used to defer processing of a function to the RTOS daemon task.

A mechanism is provided that allows the interrupt to return directly to the task that will subsequently execute the pended callback function. This allows the callback function to execute contiguously in time with the interrupt - just as if the callback had executed in the interrupt itself.

Example usage:

```
// The callback function that will execute in the context of the daemon task.
// Note callback functions must all use this same prototype.
void vProcessInterface( void *pvParameter1, uint32_t ulParameter2 )
{
   BaseType_t xInterfaceToService;

   // The interface that requires servicing is passed in the second
    // parameter.  The first parameter is not used in this case.
   xInterfaceToService = ( BaseType_t ) ulParameter2;

   // ...Perform the processing here...
}

// An ISR that receives data packets from multiple interfaces
void vAnISR( void )
{
   BaseType_t xInterfaceToService, xHigherPriorityTaskWoken;

   // Query the hardware to determine which interface needs processing.
   xInterfaceToService = prvCheckInterfaces();

    // The actual processing is to be deferred to a task.  Request the
    // vProcessInterface() callback function is executed, passing in the
   // number of the interface that needs processing.  The interface to
   // service is passed in the second parameter.  The first parameter is
   // not used in this case.
   xHigherPriorityTaskWoken = pdFALSE;
   xTimerPendFunctionCallFromISR( vProcessInterface, NULL, ( uint32_t )␣
↪xInterfaceToService, &xHigherPriorityTaskWoken );

   // If xHigherPriorityTaskWoken is now set to pdTRUE then a context
   // switch should be requested.  The macro used is port specific and will
   // be either portYIELD_FROM_ISR() or portEND_SWITCHING_ISR() - refer to
   // the documentation page for the port being used.
   portYIELD_FROM_ISR( xHigherPriorityTaskWoken );

}
```

Return pdPASS is returned if the message was successfully sent to the timer daemon task, otherwise pdFALSE is returned.

Parameters

- `xFunctionToPend`: The function to execute from the timer service/ daemon task. The function must conform to the PendedFunction_t prototype.
- `pvParameter1`: The value of the callback function's first parameter. The parameter has a void * type to allow it to be used to pass any type. For example, unsigned longs can be cast to a void *, or the void * can be used to point to a structure.
- `ulParameter2`: The value of the callback function's second parameter.
- `pxHigherPriorityTaskWoken`: As mentioned above, calling this function will result in a message being sent to the timer daemon task. If the priority of the timer daemon task (which is set using configTIMER_TASK_PRIORITY in FreeRTOSConfig.h) is higher than the priority of the currently running task (the task the interrupt interrupted) then *pxHigherPriorityTaskWoken will be set to pdTRUE within xTimerPendFunctionCallFromISR(), indicating that a context switch should be requested before the interrupt exits. For that reason *pxHigherPriorityTaskWoken must be initialised to pdFALSE. See the example code below.

BaseType_t **xTimerPendFunctionCall**(*PendedFunction_t xFunctionToPend*, void **pvParameter1*, uint32_t *ulParameter2*, TickType_t *xTicksToWait*)

Used to defer the execution of a function to the RTOS daemon task (the timer service task, hence this function is implemented in timers.c and is prefixed with 'Timer').

Return pdPASS is returned if the message was successfully sent to the timer daemon task, otherwise pdFALSE is returned.

Parameters

- `xFunctionToPend`: The function to execute from the timer service/ daemon task. The function must conform to the PendedFunction_t prototype.
- `pvParameter1`: The value of the callback function's first parameter. The parameter has a void * type to allow it to be used to pass any type. For example, unsigned longs can be cast to a void *, or the void * can be used to point to a structure.
- `ulParameter2`: The value of the callback function's second parameter.
- `xTicksToWait`: Calling this function will result in a message being sent to the timer daemon task on a queue. xTicksToWait is the amount of time the calling task should remain in the Blocked state (so not using any processing time) for space to become available on the timer queue if the queue is found to be full.

const char ***pcTimerGetTimerName** (*TimerHandle_t xTimer*)

Returns the name that was assigned to a timer when the timer was created.

Return The name assigned to the timer specified by the xTimer parameter.

Parameters

- `xTimer`: The handle of the timer being queried.

Macros

tmrCOMMAND_EXECUTE_CALLBACK_FROM_ISR

tmrCOMMAND_EXECUTE_CALLBACK

tmrCOMMAND_START_DONT_TRACE

tmrCOMMAND_START

tmrCOMMAND_RESET

tmrCOMMAND_STOP

tmrCOMMAND_CHANGE_PERIOD

tmrCOMMAND_DELETE

tmrFIRST_FROM_ISR_COMMAND

tmrCOMMAND_START_FROM_ISR

tmrCOMMAND_RESET_FROM_ISR

tmrCOMMAND_STOP_FROM_ISR

tmrCOMMAND_CHANGE_PERIOD_FROM_ISR

xTimerStart (xTimer, xTicksToWait)

Timer functionality is provided by a timer service/daemon task. Many of the public FreeRTOS timer API functions send commands to the timer service task through a queue called the timer command queue. The timer command queue is private to the kernel itself and is not directly accessible to application code. The length of the timer command queue is set by the configTIMER_QUEUE_LENGTH configuration constant.

xTimerStart() starts a timer that was previously created using the xTimerCreate() API function. If the timer had already been started and was already in the active state, then xTimerStart() has equivalent functionality to the xTimerReset() API function.

Starting a timer ensures the timer is in the active state. If the timer is not stopped, deleted, or reset in the mean time, the callback function associated with the timer will get called 'n' ticks after xTimerStart() was called, where 'n' is the timers defined period.

It is valid to call xTimerStart() before the scheduler has been started, but when this is done the timer will not actually start until the scheduler is started, and the timers expiry time will be relative to when the scheduler is started, not relative to when xTimerStart() was called.

The configUSE_TIMERS configuration constant must be set to 1 for xTimerStart() to be available.

Example usage:

Return pdFAIL will be returned if the start command could not be sent to the timer command queue even after xTicksToWait ticks had passed. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system, although the timers expiry time is relative to when xTimerStart() is actually called. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer being started/restarted.
- `xTicksToWait`: Specifies the time, in ticks, that the calling task should be held in the Blocked state to wait for the start command to be successfully sent to the timer command queue, should the queue already be full when xTimerStart() was called. xTicksToWait is ignored if xTimerStart() is called before the scheduler is started.

See the xTimerCreate() API function example usage scenario.

`xTimerStop`(xTimer, xTicksToWait)

Timer functionality is provided by a timer service/daemon task. Many of the public FreeRTOS timer API functions send commands to the timer service task through a queue called the timer command queue. The timer command queue is private to the kernel itself and is not directly accessible to application code. The length of the timer command queue is set by the configTIMER_QUEUE_LENGTH configuration constant.

xTimerStop() stops a timer that was previously started using either of the The xTimerStart(), xTimerReset(), xTimerStartFromISR(), xTimerResetFromISR(), xTimerChangePeriod() or xTimerChangePeriodFromISR() API functions.

Stopping a timer ensures the timer is not in the active state.

The configUSE_TIMERS configuration constant must be set to 1 for xTimerStop() to be available.

Example usage:

Return pdFAIL will be returned if the stop command could not be sent to the timer command queue even after xTicksToWait ticks had passed. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer being stopped.
- `xTicksToWait`: Specifies the time, in ticks, that the calling task should be held in the Blocked state to wait for the stop command to be successfully sent to the timer command queue, should the queue already be full when xTimerStop() was called. xTicksToWait is ignored if xTimerStop() is called before the scheduler is started.

See the xTimerCreate() API function example usage scenario.

`xTimerChangePeriod`(xTimer, xNewPeriod, xTicksToWait)

Timer functionality is provided by a timer service/daemon task. Many of the public FreeRTOS timer API functions send commands to the timer service task through a queue called the timer command queue. The timer command queue is private to the kernel itself and is not directly accessible to application code. The length of the timer command queue is set by the configTIMER_QUEUE_LENGTH configuration constant.

xTimerChangePeriod() changes the period of a timer that was previously created using the xTimerCreate() API function.

xTimerChangePeriod() can be called to change the period of an active or dormant state timer.

The configUSE_TIMERS configuration constant must be set to 1 for xTimerChangePeriod() to be available.

Example usage:

```
// This function assumes xTimer has already been created.  If the timer
// referenced by xTimer is already active when it is called, then the timer
// is deleted.  If the timer referenced by xTimer is not active when it is
// called, then the period of the timer is set to 500ms and the timer is
// started.
void vAFunction( TimerHandle_t xTimer )
{
    if( xTimerIsTimerActive( xTimer ) != pdFALSE ) // or more simply and␣
↪equivalently "if( xTimerIsTimerActive( xTimer ) )"
    {
        // xTimer is already active - delete it.
        xTimerDelete( xTimer );
    }
    else
    {
        // xTimer is not active, change its period to 500ms.  This will also
        // cause the timer to start.  Block for a maximum of 100 ticks if the
        // change period command cannot immediately be sent to the timer
        // command queue.
        if( xTimerChangePeriod( xTimer, 500 / portTICK_PERIOD_MS, 100 ) ==␣
↪pdPASS )
        {
            // The command was successfully sent.
        }
        else
        {
            // The command could not be sent, even after waiting for 100 ticks
            // to pass.  Take appropriate action here.
        }
    }
}
```

Return pdFAIL will be returned if the change period command could not be sent to the timer command queue even after xTicksToWait ticks had passed. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer that is having its period changed.
- `xNewPeriod`: The new period for xTimer. Timer periods are specified in tick periods, so the constant portTICK_PERIOD_MS can be used to convert a time that has been specified in milliseconds. For example, if the timer must expire after 100 ticks, then xNewPeriod should be set to 100. Alternatively, if the timer must expire after 500ms, then xNewPeriod can be set to (500 / portTICK_PERIOD_MS) provided configTICK_RATE_HZ is less than or equal to 1000.
- `xTicksToWait`: Specifies the time, in ticks, that the calling task should be held in the Blocked state to wait for the change period command to be successfully sent to the timer command queue, should the queue already be full when xTimerChangePeriod() was called. xTicksToWait is ignored if xTimerChangePeriod() is called before the scheduler is started.

`xTimerDelete` (xTimer, xTicksToWait)

Timer functionality is provided by a timer service/daemon task. Many of the public FreeRTOS timer API functions send commands to the timer service task through a queue called the timer command queue. The timer command queue is private to the kernel itself and is not directly accessible to application code. The length of the timer command queue is set by the configTIMER_QUEUE_LENGTH configuration constant.

xTimerDelete() deletes a timer that was previously created using the xTimerCreate() API function.

The configUSE_TIMERS configuration constant must be set to 1 for xTimerDelete() to be available.

Example usage:

Return pdFAIL will be returned if the delete command could not be sent to the timer command queue even after xTicksToWait ticks had passed. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer being deleted.
- `xTicksToWait`: Specifies the time, in ticks, that the calling task should be held in the Blocked state to wait for the delete command to be successfully sent to the timer command queue, should the queue already be full when xTimerDelete() was called. xTicksToWait is ignored if xTimerDelete() is called before the scheduler is started.

See the xTimerChangePeriod() API function example usage scenario.

`xTimerReset`(xTimer, xTicksToWait)

Timer functionality is provided by a timer service/daemon task. Many of the public FreeRTOS timer API functions send commands to the timer service task through a queue called the timer command queue. The timer command queue is private to the kernel itself and is not directly accessible to application code. The length of the timer command queue is set by the configTIMER_QUEUE_LENGTH configuration constant.

xTimerReset() re-starts a timer that was previously created using the xTimerCreate() API function. If the timer had already been started and was already in the active state, then xTimerReset() will cause the timer to re-evaluate its expiry time so that it is relative to when xTimerReset() was called. If the timer was in the dormant state then xTimerReset() has equivalent functionality to the xTimerStart() API function.

Resetting a timer ensures the timer is in the active state. If the timer is not stopped, deleted, or reset in the mean time, the callback function associated with the timer will get called 'n' ticks after xTimerReset() was called, where 'n' is the timers defined period.

It is valid to call xTimerReset() before the scheduler has been started, but when this is done the timer will not actually start until the scheduler is started, and the timers expiry time will be relative to when the scheduler is started, not relative to when xTimerReset() was called.

The configUSE_TIMERS configuration constant must be set to 1 for xTimerReset() to be available.

Example usage:

```
// When a key is pressed, an LCD back-light is switched on.  If 5 seconds pass
// without a key being pressed, then the LCD back-light is switched off.  In
// this case, the timer is a one-shot timer.

TimerHandle_t xBacklightTimer = NULL;

// The callback function assigned to the one-shot timer.  In this case the
// parameter is not used.
void vBacklightTimerCallback( TimerHandle_t pxTimer )
{
    // The timer expired, therefore 5 seconds must have passed since a key
    // was pressed.  Switch off the LCD back-light.
    vSetBacklightState( BACKLIGHT_OFF );
}

// The key press event handler.
void vKeyPressEventHandler( char cKey )
{
    // Ensure the LCD back-light is on, then reset the timer that is
    // responsible for turning the back-light off after 5 seconds of
    // key inactivity.  Wait 10 ticks for the command to be successfully sent
    // if it cannot be sent immediately.
    vSetBacklightState( BACKLIGHT_ON );
    if( xTimerReset( xBacklightTimer, 100 ) != pdPASS )
    {
        // The reset command was not executed successfully.  Take appropriate
```

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```
        // action here.
    }

    // Perform the rest of the key processing here.
}

void main( void )
{
int32_t x;

    // Create then start the one-shot timer that is responsible for turning
    // the back-light off if no keys are pressed within a 5 second period.
    xBacklightTimer = xTimerCreate( "BacklightTimer",           // Just a text␣
↪name, not used by the kernel.
                                    ( 5000 / portTICK_PERIOD_MS), // The timer␣
↪period in ticks.
                                    pdFALSE,                    // The timer␣
↪is a one-shot timer.
                                    0,                          // The id is␣
↪not used by the callback so can take any value.
                                    vBacklightTimerCallback     // The␣
↪callback function that switches the LCD back-light off.
                                  );

    if( xBacklightTimer == NULL )
    {
        // The timer was not created.
    }
    else
    {
        // Start the timer.  No block time is specified, and even if one was
        // it would be ignored because the scheduler has not yet been
        // started.
        if( xTimerStart( xBacklightTimer, 0 ) != pdPASS )
        {
            // The timer could not be set into the Active state.
        }
    }

    // ...
    // Create tasks here.
    // ...

    // Starting the scheduler will start the timer running as it has already
    // been set into the active state.
    xTaskStartScheduler();

    // Should not reach here.
    for( ;; );
}
```

Return pdFAIL will be returned if the reset command could not be sent to the timer command queue even after xTicksToWait ticks had passed. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system, although the timers expiry time is relative to when xTimerStart() is actually called. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer being reset/started/restarted.
- `xTicksToWait`: Specifies the time, in ticks, that the calling task should be held in the Blocked state to wait for the reset command to be successfully sent to the timer command queue, should the

queue already be full when xTimerReset() was called. xTicksToWait is ignored if xTimerReset() is called before the scheduler is started.

xTimerStartFromISR(xTimer, pxHigherPriorityTaskWoken)

A version of xTimerStart() that can be called from an interrupt service routine.

Example usage:

```
// This scenario assumes xBacklightTimer has already been created.  When a
// key is pressed, an LCD back-light is switched on.  If 5 seconds pass
// without a key being pressed, then the LCD back-light is switched off.  In
// this case, the timer is a one-shot timer, and unlike the example given for
// the xTimerReset() function, the key press event handler is an interrupt
// service routine.

// The callback function assigned to the one-shot timer.  In this case the
// parameter is not used.
void vBacklightTimerCallback( TimerHandle_t pxTimer )
{
    // The timer expired, therefore 5 seconds must have passed since a key
    // was pressed.  Switch off the LCD back-light.
    vSetBacklightState( BACKLIGHT_OFF );
}

// The key press interrupt service routine.
void vKeyPressEventInterruptHandler( void )
{
BaseType_t xHigherPriorityTaskWoken = pdFALSE;

    // Ensure the LCD back-light is on, then restart the timer that is
    // responsible for turning the back-light off after 5 seconds of
    // key inactivity.  This is an interrupt service routine so can only
    // call FreeRTOS API functions that end in "FromISR".
    vSetBacklightState( BACKLIGHT_ON );

    // xTimerStartFromISR() or xTimerResetFromISR() could be called here
    // as both cause the timer to re-calculate its expiry time.
    // xHigherPriorityTaskWoken was initialised to pdFALSE when it was
    // declared (in this function).
    if( xTimerStartFromISR( xBacklightTimer, &xHigherPriorityTaskWoken ) !=␣
↪pdPASS )
    {
        // The start command was not executed successfully.  Take appropriate
        // action here.
    }

    // Perform the rest of the key processing here.

    // If xHigherPriorityTaskWoken equals pdTRUE, then a context switch
    // should be performed.  The syntax required to perform a context switch
    // from inside an ISR varies from port to port, and from compiler to
    // compiler.  Inspect the demos for the port you are using to find the
    // actual syntax required.
    if( xHigherPriorityTaskWoken != pdFALSE )
    {
        // Call the interrupt safe yield function here (actual function
        // depends on the FreeRTOS port being used).
    }
}
```

Return pdFAIL will be returned if the start command could not be sent to the timer command queue. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system, although the timers expiry time is relative to when xTimerStartFromISR() is actually called.

The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer being started/restarted.
- `pxHigherPriorityTaskWoken`: The timer service/daemon task spends most of its time in the Blocked state, waiting for messages to arrive on the timer command queue. Calling xTimerStartFromISR() writes a message to the timer command queue, so has the potential to transition the timer service/daemon task out of the Blocked state. If calling xTimerStartFromISR() causes the timer service/daemon task to leave the Blocked state, and the timer service/ daemon task has a priority equal to or greater than the currently executing task (the task that was interrupted), then *pxHigherPriorityTaskWoken will get set to pdTRUE internally within the xTimerStartFromISR() function. If xTimerStartFromISR() sets this value to pdTRUE then a context switch should be performed before the interrupt exits.

`xTimerStopFromISR`(xTimer, pxHigherPriorityTaskWoken)

A version of xTimerStop() that can be called from an interrupt service routine.

Example usage:

```
// This scenario assumes xTimer has already been created and started.  When
// an interrupt occurs, the timer should be simply stopped.

// The interrupt service routine that stops the timer.
void vAnExampleInterruptServiceRoutine( void )
{
BaseType_t xHigherPriorityTaskWoken = pdFALSE;

    // The interrupt has occurred - simply stop the timer.
    // xHigherPriorityTaskWoken was set to pdFALSE where it was defined
    // (within this function).  As this is an interrupt service routine, only
    // FreeRTOS API functions that end in "FromISR" can be used.
    if( xTimerStopFromISR( xTimer, &xHigherPriorityTaskWoken ) != pdPASS )
    {
        // The stop command was not executed successfully.  Take appropriate
        // action here.
    }

    // If xHigherPriorityTaskWoken equals pdTRUE, then a context switch
    // should be performed.  The syntax required to perform a context switch
    // from inside an ISR varies from port to port, and from compiler to
    // compiler.  Inspect the demos for the port you are using to find the
    // actual syntax required.
    if( xHigherPriorityTaskWoken != pdFALSE )
    {
        // Call the interrupt safe yield function here (actual function
        // depends on the FreeRTOS port being used).
    }
}
```

Return pdFAIL will be returned if the stop command could not be sent to the timer command queue. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer being stopped.
- `pxHigherPriorityTaskWoken`: The timer service/daemon task spends most of its time in the Blocked state, waiting for messages to arrive on the timer command queue. Calling xTimerStopFromISR() writes a message to the timer command queue, so has the potential to transition the timer service/daemon task out of the Blocked state. If calling xTimerStopFromISR() causes the timer service/daemon task to leave the Blocked state, and the timer service/ daemon task has a priority equal to or greater than the currently executing task (the task that was interrupted), then

*pxHigherPriorityTaskWoken will get set to pdTRUE internally within the xTimerStopFromISR() function. If xTimerStopFromISR() sets this value to pdTRUE then a context switch should be performed before the interrupt exits.

xTimerChangePeriodFromISR(xTimer, xNewPeriod, pxHigherPriorityTaskWoken)

A version of xTimerChangePeriod() that can be called from an interrupt service routine.

Example usage:

```
// This scenario assumes xTimer has already been created and started.  When
// an interrupt occurs, the period of xTimer should be changed to 500ms.

// The interrupt service routine that changes the period of xTimer.
void vAnExampleInterruptServiceRoutine( void )
{
BaseType_t xHigherPriorityTaskWoken = pdFALSE;

    // The interrupt has occurred - change the period of xTimer to 500ms.
    // xHigherPriorityTaskWoken was set to pdFALSE where it was defined
    // (within this function).  As this is an interrupt service routine, only
    // FreeRTOS API functions that end in "FromISR" can be used.
    if( xTimerChangePeriodFromISR( xTimer, &xHigherPriorityTaskWoken ) !=␣
↪pdPASS )
    {
        // The command to change the timers period was not executed
        // successfully.  Take appropriate action here.
    }

    // If xHigherPriorityTaskWoken equals pdTRUE, then a context switch
    // should be performed.  The syntax required to perform a context switch
    // from inside an ISR varies from port to port, and from compiler to
    // compiler.  Inspect the demos for the port you are using to find the
    // actual syntax required.
    if( xHigherPriorityTaskWoken != pdFALSE )
    {
        // Call the interrupt safe yield function here (actual function
        // depends on the FreeRTOS port being used).
    }
}
```

Return pdFAIL will be returned if the command to change the timers period could not be sent to the timer command queue. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer that is having its period changed.
- `xNewPeriod`: The new period for xTimer. Timer periods are specified in tick periods, so the constant portTICK_PERIOD_MS can be used to convert a time that has been specified in milliseconds. For example, if the timer must expire after 100 ticks, then xNewPeriod should be set to 100. Alternatively, if the timer must expire after 500ms, then xNewPeriod can be set to (500 / portTICK_PERIOD_MS) provided configTICK_RATE_HZ is less than or equal to 1000.
- `pxHigherPriorityTaskWoken`: The timer service/daemon task spends most of its time in the Blocked state, waiting for messages to arrive on the timer command queue. Calling xTimerChangePeriodFromISR() writes a message to the timer command queue, so has the potential to transition the timer service/ daemon task out of the Blocked state. If calling xTimerChangePeriodFromISR() causes the timer service/daemon task to leave the Blocked state, and the timer service/daemon task has a priority equal to or greater than the currently executing task (the task that was interrupted), then *pxHigherPriorityTaskWoken will get set to pdTRUE internally within the xTimerChangePeriodFromISR() function. If xTimerChangePeriodFromISR() sets this value to pdTRUE then a context switch should be performed before the interrupt exits.

xTimerResetFromISR (xTimer, pxHigherPriorityTaskWoken)

A version of xTimerReset() that can be called from an interrupt service routine.

Example usage:

```
// This scenario assumes xBacklightTimer has already been created.  When a
// key is pressed, an LCD back-light is switched on.  If 5 seconds pass
// without a key being pressed, then the LCD back-light is switched off.  In
// this case, the timer is a one-shot timer, and unlike the example given for
// the xTimerReset() function, the key press event handler is an interrupt
// service routine.

// The callback function assigned to the one-shot timer.  In this case the
// parameter is not used.
void vBacklightTimerCallback( TimerHandle_t pxTimer )
{
    // The timer expired, therefore 5 seconds must have passed since a key
    // was pressed.  Switch off the LCD back-light.
    vSetBacklightState( BACKLIGHT_OFF );
}

// The key press interrupt service routine.
void vKeyPressEventInterruptHandler( void )
{
BaseType_t xHigherPriorityTaskWoken = pdFALSE;

    // Ensure the LCD back-light is on, then reset the timer that is
    // responsible for turning the back-light off after 5 seconds of
    // key inactivity.  This is an interrupt service routine so can only
    // call FreeRTOS API functions that end in "FromISR".
    vSetBacklightState( BACKLIGHT_ON );

    // xTimerStartFromISR() or xTimerResetFromISR() could be called here
    // as both cause the timer to re-calculate its expiry time.
    // xHigherPriorityTaskWoken was initialised to pdFALSE when it was
    // declared (in this function).
    if( xTimerResetFromISR( xBacklightTimer, &xHigherPriorityTaskWoken ) !=␣
↪pdPASS )
    {
        // The reset command was not executed successfully.  Take appropriate
        // action here.
    }

    // Perform the rest of the key processing here.

    // If xHigherPriorityTaskWoken equals pdTRUE, then a context switch
    // should be performed.  The syntax required to perform a context switch
    // from inside an ISR varies from port to port, and from compiler to
    // compiler.  Inspect the demos for the port you are using to find the
    // actual syntax required.
    if( xHigherPriorityTaskWoken != pdFALSE )
    {
        // Call the interrupt safe yield function here (actual function
        // depends on the FreeRTOS port being used).
    }
}
```

Return pdFAIL will be returned if the reset command could not be sent to the timer command queue. pdPASS will be returned if the command was successfully sent to the timer command queue. When the command is actually processed will depend on the priority of the timer service/daemon task relative to other tasks in the system, although the timers expiry time is relative to when xTimerResetFromISR() is actually called. The timer service/daemon task priority is set by the configTIMER_TASK_PRIORITY configuration constant.

Parameters

- `xTimer`: The handle of the timer that is to be started, reset, or restarted.
- `pxHigherPriorityTaskWoken`: The timer service/daemon task spends most of its time in the Blocked state, waiting for messages to arrive on the timer command queue. Calling xTimerResetFromISR() writes a message to the timer command queue, so has the potential to transition the timer service/daemon task out of the Blocked state. If calling xTimerResetFromISR() causes the timer service/daemon task to leave the Blocked state, and the timer service/ daemon task has a priority equal to or greater than the currently executing task (the task that was interrupted), then *pxHigherPriorityTaskWoken will get set to pdTRUE internally within the xTimerResetFromISR() function. If xTimerResetFromISR() sets this value to pdTRUE then a context switch should be performed before the interrupt exits.

Type Definitions

typedef void ***TimerHandle_t**

Type by which software timers are referenced. For example, a call to xTimerCreate() returns an TimerHandle_t variable that can then be used to reference the subject timer in calls to other software timer API functions (for example, xTimerStart(), xTimerReset(), etc.).

typedef void (***TimerCallbackFunction_t**) (*TimerHandle_t* xTimer)

Defines the prototype to which timer callback functions must conform.

typedef void (***PendedFunction_t**) (void *, uint32_t)

Defines the prototype to which functions used with the xTimerPendFunctionCallFromISR() function must conform.

Event Group API

Header File

- freertos/include/freertos/event_groups.h

Functions

EventGroupHandle_t **xEventGroupCreate** (void)

Create a new event group.

Internally, within the FreeRTOS implementation, event groups use a [small] block of memory, in which the event group's structure is stored. If an event groups is created using xEventGroupCreate() then the required memory is automatically dynamically allocated inside the xEventGroupCreate() function. (see http://www.freertos.org/a00111.html). If an event group is created using xEventGropuCreateStatic() then the application writer must instead provide the memory that will get used by the event group. xEventGroupCreateStatic() therefore allows an event group to be created without using any dynamic memory allocation.

Although event groups are not related to ticks, for internal implementation reasons the number of bits available for use in an event group is dependent on the configUSE_16_BIT_TICKS setting in FreeRTOSConfig.h. If configUSE_16_BIT_TICKS is 1 then each event group contains 8 usable bits (bit 0 to bit 7). If configUSE_16_BIT_TICKS is set to 0 then each event group has 24 usable bits (bit 0 to bit 23). The EventBits_t type is used to store event bits within an event group.

Example usage:

```
// Declare a variable to hold the created event group.
EventGroupHandle_t xCreatedEventGroup;

// Attempt to create the event group.
xCreatedEventGroup = xEventGroupCreate();

// Was the event group created successfully?
if( xCreatedEventGroup == NULL )
{
    // The event group was not created because there was insufficient
```

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```
    // FreeRTOS heap available.
}
else
{
    // The event group was created.
}
```

Return If the event group was created then a handle to the event group is returned. If there was insufficient FreeRTOS heap available to create the event group then NULL is returned. See http://www.freertos.org/a00111.html

EventGroupHandle_t **xEventGroupCreateStatic**(StaticEventGroup_t **pxEventGroupBuffer*)

Create a new event group.

Internally, within the FreeRTOS implementation, event groups use a [small] block of memory, in which the event group's structure is stored. If an event groups is created using xEventGropuCreate() then the required memory is automatically dynamically allocated inside the xEventGroupCreate() function. (see http://www.freertos.org/a00111.html). If an event group is created using xEventGropuCreateStatic() then the application writer must instead provide the memory that will get used by the event group. xEventGroupCreateStatic() therefore allows an event group to be created without using any dynamic memory allocation.

Although event groups are not related to ticks, for internal implementation reasons the number of bits available for use in an event group is dependent on the configUSE_16_BIT_TICKS setting in FreeRTOSConfig.h. If configUSE_16_BIT_TICKS is 1 then each event group contains 8 usable bits (bit 0 to bit 7). If configUSE_16_BIT_TICKS is set to 0 then each event group has 24 usable bits (bit 0 to bit 23). The EventBits_t type is used to store event bits within an event group.

Example usage:

```
// StaticEventGroup_t is a publicly accessible structure that has the same
// size and alignment requirements as the real event group structure.  It is
// provided as a mechanism for applications to know the size of the event
// group (which is dependent on the architecture and configuration file
// settings) without breaking the strict data hiding policy by exposing the
// real event group internals.  This StaticEventGroup_t variable is passed
// into the xSemaphoreCreateEventGroupStatic() function and is used to store
// the event group's data structures
StaticEventGroup_t xEventGroupBuffer;

// Create the event group without dynamically allocating any memory.
xEventGroup = xEventGroupCreateStatic( &xEventGroupBuffer );
```

Return If the event group was created then a handle to the event group is returned. If pxEventGroupBuffer was NULL then NULL is returned.

Parameters

- `pxEventGroupBuffer`: pxEventGroupBuffer must point to a variable of type StaticEventGroup_t, which will be then be used to hold the event group's data structures, removing the need for the memory to be allocated dynamically.

EventBits_t **xEventGroupWaitBits**(*EventGroupHandle_t* *xEventGroup*, **const** *EventBits_t* *uxBitsToWaitFor*, **const** BaseType_t *xClearOnExit*, **const** BaseType_t *xWaitForAllBits*, TickType_t *xTicksToWait*)

[Potentially] block to wait for one or more bits to be set within a previously created event group.

This function cannot be called from an interrupt.

Example usage:

```
#define BIT_0    ( 1 << 0 )
#define BIT_4    ( 1 << 4 )

void aFunction( EventGroupHandle_t xEventGroup )
```

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```
{
EventBits_t uxBits;
const TickType_t xTicksToWait = 100 / portTICK_PERIOD_MS;

    // Wait a maximum of 100ms for either bit 0 or bit 4 to be set within
    // the event group.  Clear the bits before exiting.
    uxBits = xEventGroupWaitBits(
                xEventGroup,    // The event group being tested.
                BIT_0 | BIT_4,  // The bits within the event group to wait␣
↪for.
                pdTRUE,         // BIT_0 and BIT_4 should be cleared before␣
↪returning.
                pdFALSE,        // Don't wait for both bits, either bit will␣
↪do.
                xTicksToWait ); // Wait a maximum of 100ms for either bit to␣
↪be set.

    if( ( uxBits & ( BIT_0 | BIT_4 ) ) == ( BIT_0 | BIT_4 ) )
    {
        // xEventGroupWaitBits() returned because both bits were set.
    }
    else if( ( uxBits & BIT_0 ) != 0 )
    {
        // xEventGroupWaitBits() returned because just BIT_0 was set.
    }
    else if( ( uxBits & BIT_4 ) != 0 )
    {
        // xEventGroupWaitBits() returned because just BIT_4 was set.
    }
    else
    {
        // xEventGroupWaitBits() returned because xTicksToWait ticks passed
        // without either BIT_0 or BIT_4 becoming set.
    }
}
```

{c}

Return The value of the event group at the time either the bits being waited for became set, or the block time expired. Test the return value to know which bits were set. If xEventGroupWaitBits() returned because its timeout expired then not all the bits being waited for will be set. If xEventGroupWaitBits() returned because the bits it was waiting for were set then the returned value is the event group value before any bits were automatically cleared in the case that xClearOnExit parameter was set to pdTRUE.

Parameters

- `xEventGroup`: The event group in which the bits are being tested. The event group must have previously been created using a call to xEventGroupCreate().
- `uxBitsToWaitFor`: A bitwise value that indicates the bit or bits to test inside the event group. For example, to wait for bit 0 and/or bit 2 set uxBitsToWaitFor to 0x05. To wait for bits 0 and/or bit 1 and/or bit 2 set uxBitsToWaitFor to 0x07. Etc.
- `xClearOnExit`: If xClearOnExit is set to pdTRUE then any bits within uxBitsToWaitFor that are set within the event group will be cleared before xEventGroupWaitBits() returns if the wait condition was met (if the function returns for a reason other than a timeout). If xClearOnExit is set to pdFALSE then the bits set in the event group are not altered when the call to xEventGroupWaitBits() returns.
- `xWaitForAllBits`: If xWaitForAllBits is set to pdTRUE then xEventGroupWaitBits() will return when either all the bits in uxBitsToWaitFor are set or the specified block time expires. If xWaitForAllBits is set to pdFALSE then xEventGroupWaitBits() will return when any one of the bits set in uxBitsToWaitFor is set or the specified block time expires. The block time is specified by the xTicksToWait parameter.
- `xTicksToWait`: The maximum amount of time (specified in ‘ticks’) to wait for one/all (de-

pending on the xWaitForAllBits value) of the bits specified by uxBitsToWaitFor to become set.

EventBits_t **xEventGroupClearBits** (*EventGroupHandle_t xEventGroup*, **const** *EventBits_t uxBitsToClear*)

Clear bits within an event group. This function cannot be called from an interrupt.

Example usage:

```
#define BIT_0    ( 1 << 0 )
#define BIT_4    ( 1 << 4 )

void aFunction( EventGroupHandle_t xEventGroup )
{
EventBits_t uxBits;

    // Clear bit 0 and bit 4 in xEventGroup.
    uxBits = xEventGroupClearBits(
                            xEventGroup,    // The event group being updated.
                            BIT_0 | BIT_4 );// The bits being cleared.

    if( ( uxBits & ( BIT_0 | BIT_4 ) ) == ( BIT_0 | BIT_4 ) )
    {
        // Both bit 0 and bit 4 were set before xEventGroupClearBits() was
        // called.  Both will now be clear (not set).
    }
    else if( ( uxBits & BIT_0 ) != 0 )
    {
        // Bit 0 was set before xEventGroupClearBits() was called.  It will
        // now be clear.
    }
    else if( ( uxBits & BIT_4 ) != 0 )
    {
        // Bit 4 was set before xEventGroupClearBits() was called.  It will
        // now be clear.
    }
    else
    {
        // Neither bit 0 nor bit 4 were set in the first place.
    }
}
```

Return The value of the event group before the specified bits were cleared.

Parameters

- `xEventGroup`: The event group in which the bits are to be cleared.
- `uxBitsToClear`: A bitwise value that indicates the bit or bits to clear in the event group. For example, to clear bit 3 only, set uxBitsToClear to 0x08. To clear bit 3 and bit 0 set uxBitsToClear to 0x09.

EventBits_t **xEventGroupSetBits** (*EventGroupHandle_t xEventGroup*, **const** *EventBits_t uxBitsToSet*)

Set bits within an event group. This function cannot be called from an interrupt. xEventGroupSetBitsFromISR() is a version that can be called from an interrupt.

Setting bits in an event group will automatically unblock tasks that are blocked waiting for the bits.

Example usage:

```
#define BIT_0    ( 1 << 0 )
#define BIT_4    ( 1 << 4 )

void aFunction( EventGroupHandle_t xEventGroup )
{
EventBits_t uxBits;

    // Set bit 0 and bit 4 in xEventGroup.
```

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```
    uxBits = xEventGroupSetBits(
                        xEventGroup,    // The event group being updated.
                        BIT_0 | BIT_4 );// The bits being set.

    if( ( uxBits & ( BIT_0 | BIT_4 ) ) == ( BIT_0 | BIT_4 ) )
    {
        // Both bit 0 and bit 4 remained set when the function returned.
    }
    else if( ( uxBits & BIT_0 ) != 0 )
    {
        // Bit 0 remained set when the function returned, but bit 4 was
        // cleared.  It might be that bit 4 was cleared automatically as a
        // task that was waiting for bit 4 was removed from the Blocked
        // state.
    }
    else if( ( uxBits & BIT_4 ) != 0 )
    {
        // Bit 4 remained set when the function returned, but bit 0 was
        // cleared.  It might be that bit 0 was cleared automatically as a
        // task that was waiting for bit 0 was removed from the Blocked
        // state.
    }
    else
    {
        // Neither bit 0 nor bit 4 remained set.  It might be that a task
        // was waiting for both of the bits to be set, and the bits were
        // cleared as the task left the Blocked state.
    }
}
```

{c}

Return The value of the event group at the time the call to xEventGroupSetBits() returns. There are two reasons why the returned value might have the bits specified by the uxBitsToSet parameter cleared. First, if setting a bit results in a task that was waiting for the bit leaving the blocked state then it is possible the bit will be cleared automatically (see the xClearBitOnExit parameter of xEventGroupWaitBits()). Second, any unblocked (or otherwise Ready state) task that has a priority above that of the task that called xEventGroupSetBits() will execute and may change the event group value before the call to xEventGroupSetBits() returns.

Parameters

- `xEventGroup`: The event group in which the bits are to be set.
- `uxBitsToSet`: A bitwise value that indicates the bit or bits to set. For example, to set bit 3 only, set uxBitsToSet to 0x08. To set bit 3 and bit 0 set uxBitsToSet to 0x09.

EventBits_t **xEventGroupSync**(*EventGroupHandle_t* *xEventGroup*, **const** *EventBits_t* *uxBitsToSet*, **const** *EventBits_t* *uxBitsToWaitFor*, TickType_t *xTicksToWait*)

Atomically set bits within an event group, then wait for a combination of bits to be set within the same event group. This functionality is typically used to synchronise multiple tasks, where each task has to wait for the other tasks to reach a synchronisation point before proceeding.

This function cannot be used from an interrupt.

The function will return before its block time expires if the bits specified by the uxBitsToWait parameter are set, or become set within that time. In this case all the bits specified by uxBitsToWait will be automatically cleared before the function returns.

Example usage:

```
// Bits used by the three tasks.
#define TASK_0_BIT      ( 1 << 0 )
#define TASK_1_BIT      ( 1 << 1 )
#define TASK_2_BIT      ( 1 << 2 )
```

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```
#define ALL_SYNC_BITS ( TASK_0_BIT | TASK_1_BIT | TASK_2_BIT )

// Use an event group to synchronise three tasks.  It is assumed this event
// group has already been created elsewhere.
EventGroupHandle_t xEventBits;

void vTask0( void *pvParameters )
{
EventBits_t uxReturn;
TickType_t xTicksToWait = 100 / portTICK_PERIOD_MS;

    for( ;; )
    {
       // Perform task functionality here.

       // Set bit 0 in the event flag to note this task has reached the
       // sync point.  The other two tasks will set the other two bits defined
       // by ALL_SYNC_BITS.  All three tasks have reached the synchronisation
       // point when all the ALL_SYNC_BITS are set.  Wait a maximum of 100ms
       // for this to happen.
       uxReturn = xEventGroupSync( xEventBits, TASK_0_BIT, ALL_SYNC_BITS,␣
↪xTicksToWait );

       if( ( uxReturn & ALL_SYNC_BITS ) == ALL_SYNC_BITS )
       {
           // All three tasks reached the synchronisation point before the call
           // to xEventGroupSync() timed out.
       }
   }
}

void vTask1( void *pvParameters )
{
    for( ;; )
    {
       // Perform task functionality here.

       // Set bit 1 in the event flag to note this task has reached the
       // synchronisation point.  The other two tasks will set the other two
       // bits defined by ALL_SYNC_BITS.  All three tasks have reached the
       // synchronisation point when all the ALL_SYNC_BITS are set.  Wait
       // indefinitely for this to happen.
       xEventGroupSync( xEventBits, TASK_1_BIT, ALL_SYNC_BITS, portMAX_DELAY );

       // xEventGroupSync() was called with an indefinite block time, so
       // this task will only reach here if the syncrhonisation was made by all
       // three tasks, so there is no need to test the return value.
    }
}

void vTask2( void *pvParameters )
{
    for( ;; )
    {
       // Perform task functionality here.

       // Set bit 2 in the event flag to note this task has reached the
       // synchronisation point.  The other two tasks will set the other two
       // bits defined by ALL_SYNC_BITS.  All three tasks have reached the
       // synchronisation point when all the ALL_SYNC_BITS are set.  Wait
```

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```
        // indefinitely for this to happen.
        xEventGroupSync( xEventBits, TASK_2_BIT, ALL_SYNC_BITS, portMAX_DELAY );

        // xEventGroupSync() was called with an indefinite block time, so
        // this task will only reach here if the syncrhonisation was made by all
        // three tasks, so there is no need to test the return value.
    }
}
```

Return The value of the event group at the time either the bits being waited for became set, or the block time expired. Test the return value to know which bits were set. If xEventGroupSync() returned because its timeout expired then not all the bits being waited for will be set. If xEventGroupSync() returned because all the bits it was waiting for were set then the returned value is the event group value before any bits were automatically cleared.

Parameters

- `xEventGroup`: The event group in which the bits are being tested. The event group must have previously been created using a call to xEventGroupCreate().
- `uxBitsToSet`: The bits to set in the event group before determining if, and possibly waiting for, all the bits specified by the uxBitsToWait parameter are set.
- `uxBitsToWaitFor`: A bitwise value that indicates the bit or bits to test inside the event group. For example, to wait for bit 0 and bit 2 set uxBitsToWaitFor to 0x05. To wait for bits 0 and bit 1 and bit 2 set uxBitsToWaitFor to 0x07. Etc.
- `xTicksToWait`: The maximum amount of time (specified in 'ticks') to wait for all of the bits specified by uxBitsToWaitFor to become set.

EventBits_t **xEventGroupGetBitsFromISR**(*EventGroupHandle_t xEventGroup*)

A version of xEventGroupGetBits() that can be called from an ISR.

Return The event group bits at the time xEventGroupGetBitsFromISR() was called.

Parameters

- `xEventGroup`: The event group being queried.

void **vEventGroupDelete**(*EventGroupHandle_t xEventGroup*)

Delete an event group that was previously created by a call to xEventGroupCreate(). Tasks that are blocked on the event group will be unblocked and obtain 0 as the event group's value.

Parameters

- `xEventGroup`: The event group being deleted.

Macros

xEventGroupClearBitsFromISR(xEventGroup, uxBitsToClear)

A version of xEventGroupClearBits() that can be called from an interrupt.

Setting bits in an event group is not a deterministic operation because there are an unknown number of tasks that may be waiting for the bit or bits being set. FreeRTOS does not allow nondeterministic operations to be performed while interrupts are disabled, so protects event groups that are accessed from tasks by suspending the scheduler rather than disabling interrupts. As a result event groups cannot be accessed directly from an interrupt service routine. Therefore xEventGroupClearBitsFromISR() sends a message to the timer task to have the clear operation performed in the context of the timer task.

Example usage:

```
#define BIT_0    ( 1 << 0 )
#define BIT_4    ( 1 << 4 )

// An event group which it is assumed has already been created by a call to
// xEventGroupCreate().
EventGroupHandle_t xEventGroup;

void anInterruptHandler( void )
```

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```
{
      // Clear bit 0 and bit 4 in xEventGroup.
      xResult = xEventGroupClearBitsFromISR(
                          xEventGroup,     // The event group being updated.
                          BIT_0 | BIT_4 ); // The bits being set.

      if( xResult == pdPASS )
      {
          // The message was posted successfully.
      }
}
```

Return If the request to execute the function was posted successfully then pdPASS is returned, otherwise pdFALSE is returned. pdFALSE will be returned if the timer service queue was full.

Parameters

- `xEventGroup`: The event group in which the bits are to be cleared.
- `uxBitsToClear`: A bitwise value that indicates the bit or bits to clear. For example, to clear bit 3 only, set uxBitsToClear to 0x08. To clear bit 3 and bit 0 set uxBitsToClear to 0x09.

xEventGroupSetBitsFromISR(xEventGroup, uxBitsToSet, pxHigherPriorityTaskWoken)

A version of xEventGroupSetBits() that can be called from an interrupt.

Setting bits in an event group is not a deterministic operation because there are an unknown number of tasks that may be waiting for the bit or bits being set. FreeRTOS does not allow nondeterministic operations to be performed in interrupts or from critical sections. Therefore xEventGroupSetBitFromISR() sends a message to the timer task to have the set operation performed in the context of the timer task - where a scheduler lock is used in place of a critical section.

Example usage:

```
#define BIT_0    ( 1 << 0 )
#define BIT_4    ( 1 << 4 )

// An event group which it is assumed has already been created by a call to
// xEventGroupCreate().
EventGroupHandle_t xEventGroup;

void anInterruptHandler( void )
{
BaseType_t xHigherPriorityTaskWoken, xResult;

      // xHigherPriorityTaskWoken must be initialised to pdFALSE.
      xHigherPriorityTaskWoken = pdFALSE;

      // Set bit 0 and bit 4 in xEventGroup.
      xResult = xEventGroupSetBitsFromISR(
                          xEventGroup,    // The event group being updated.
                          BIT_0 | BIT_4   // The bits being set.
                          &xHigherPriorityTaskWoken );

      // Was the message posted successfully?
      if( xResult == pdPASS )
      {
          // If xHigherPriorityTaskWoken is now set to pdTRUE then a context
          // switch should be requested.  The macro used is port specific and
          // will be either portYIELD_FROM_ISR() or portEND_SWITCHING_ISR() -
          // refer to the documentation page for the port being used.
          portYIELD_FROM_ISR( xHigherPriorityTaskWoken );
      }
}
```

Return If the request to execute the function was posted successfully then pdPASS is returned, otherwise

pdFALSE is returned. pdFALSE will be returned if the timer service queue was full.

Parameters

- `xEventGroup`: The event group in which the bits are to be set.
- `uxBitsToSet`: A bitwise value that indicates the bit or bits to set. For example, to set bit 3 only, set uxBitsToSet to 0x08. To set bit 3 and bit 0 set uxBitsToSet to 0x09.
- `pxHigherPriorityTaskWoken`: As mentioned above, calling this function will result in a message being sent to the timer daemon task. If the priority of the timer daemon task is higher than the priority of the currently running task (the task the interrupt interrupted) then *pxHigherPriorityTaskWoken will be set to pdTRUE by xEventGroupSetBitsFromISR(), indicating that a context switch should be requested before the interrupt exits. For that reason *pxHigherPriorityTaskWoken must be initialised to pdFALSE. See the example code below.

xEventGroupGetBits(xEventGroup)

Returns the current value of the bits in an event group. This function cannot be used from an interrupt.

Return The event group bits at the time xEventGroupGetBits() was called.

Parameters

- `xEventGroup`: The event group being queried.

Type Definitions

typedef void ***EventGroupHandle_t**

An event group is a collection of bits to which an application can assign a meaning. For example, an application may create an event group to convey the status of various CAN bus related events in which bit 0 might mean "A CAN

message has been received and is ready for processing" , bit 1 might mean "The application has queued a message that is ready for sending onto the CAN network" , and bit 2 might mean "It is time to send a SYNC message onto the CAN network" etc. A task can then test the bit values to see which events are active, and optionally enter the Blocked state to wait for a specified bit or a group of specified bits to be active. To continue the CAN bus example, a CAN controlling task can enter the Blocked state (and therefore not consume any processing time) until either bit 0, bit 1 or bit 2 are active, at which time the bit that was actually active would inform the task which action it had to take (process a received message, send a message, or send a SYNC).

The event groups implementation contains intelligence to avoid race conditions that would otherwise occur were an application to use a simple variable for the same purpose. This is particularly important with respect to when a bit within an event group is to be cleared, and when bits have to be set and then tested atomically - as is the case where event groups are used to create a synchronisation point between multiple tasks (a 'rendezvous'). event_groups.h

Type by which event groups are referenced. For example, a call to xEventGroupCreate() returns an EventGroupHandle_t variable that can then be used as a parameter to other event group functions.

typedef TickType_t **EventBits_t**

2.6.10 FreeRTOS Additions

Overview

ESP-IDF FreeRTOS is based on the Xtensa port of FreeRTOS v8.2.0 with significant modifications for SMP compatibility (see *ESP-IDF FreeRTOS SMP Changes*). However various features specific to ESP-IDF FreeRTOS have been added. The features are as follows:

Ring Buffers: Ring buffers were added to provide a form of buffer that could accept entries of arbitrary lengths.

Hooks: ESP-IDF FreeRTOS hooks provides support for registering extra Idle and Tick hooks at run time. Moreover, the hooks can be asymmetric amongst both CPUs.

Ring Buffers

The ESP-IDF FreeRTOS ring buffer is a strictly FIFO buffer that supports arbitrarily sized items. Ring buffers are a more memory efficient alternative to FreeRTOS queues in situations where the size of items is variable. The capacity of a ring buffer is not measured by the number of items it can store, but rather by the amount of memory used for storing items. You may apply for a piece of memory on the ring buffer to send an item, or just use the API to copy your data and send (according to the send API you call). For efficiency reasons, **items are always retrieved from the ring buffer by reference**. As a result, all retrieved items *must also be returned* in order for them to be removed from the ring buffer completely. The ring buffers are split into the three following types:

No-Split buffers will guarantee that an item is stored in contiguous memory and will not attempt to split an item under any circumstances. Use no-split buffers when items must occupy contiguous memory. *Only this buffer type allows you getting the data item address and writting to the item by yourself.*

Allow-Split buffers will allow an item to be split when wrapping around if doing so will allow the item to be stored. Allow-split buffers are more memory efficient than no-split buffers but can return an item in two parts when retrieving.

Byte buffers do not store data as separate items. All data is stored as a sequence of bytes, and any number of bytes and be sent or retrieved each time. Use byte buffers when separate items do not need to be maintained (e.g. a byte stream).

注解: No-split/allow-split buffers will always store items at 32-bit aligned addresses. Therefore when retrieving an item, the item pointer is guaranteed to be 32-bit aligned. This is useful especially when you need to send some data to the DMA.

注解: Each item stored in no-split/allow-split buffers will **require an additional 8 bytes for a header**. Item sizes will also be rounded up to a 32-bit aligned size (multiple of 4 bytes), however the true item size is recorded within the header. The sizes of no-split/allow-split buffers will also be rounded up when created.

Usage The following example demonstrates the usage of *xRingbufferCreate()* and *xRingbufferSend()* to create a ring buffer then send an item to it.

```
#include "freertos/ringbuf.h"
static char tx_item[] = "test_item";

...

    //Create ring buffer
    RingbufHandle_t buf_handle;
    buf_handle = xRingbufferCreate(1028, RINGBUF_TYPE_NOSPLIT);
    if (buf_handle == NULL) {
        printf("Failed to create ring buffer\n");
    }

    //Send an item
    UBaseType_t res =  xRingbufferSend(buf_handle, tx_item, sizeof(tx_item), pdMS_
↪TO_TICKS(1000));
    if (res != pdTRUE) {
        printf("Failed to send item\n");
    }
```

The following example demonstrates the usage of *xRingbufferSendAcquire()* and *xRingbufferSendComplete()* instead of *xRingbufferSend()* to apply for the memory on the ring buffer (of type *RINGBUF_TYPE_NOSPLIT*) and then send an item to it. This way adds one more step, but allows getting the address of the memory to write to, and writing to the memory yourself.

```
#include "freertos/ringbuf.h"
#include "soc/lldesc.h"

typedef struct {
    lldesc_t dma_desc;
    uint8_t buf[1];
} dma_item_t;

#define DMA_ITEM_SIZE(N) (sizeof(lldesc_t)+(((N)+3)&(~3)))

...

    //Retrieve space for DMA descriptor and corresponding data buffer
    //This has to be done with SendAcquire, or the address may be different when␣
↪copy
    dma_item_t item;
    UBaseType_t res =  xRingbufferSendAcquire(buf_handle,
                        &item, DMA_ITEM_SIZE(buffer_size), pdMS_TO_TICKS(1000));
    if (res != pdTRUE) {
        printf("Failed to acquire memory for item\n");
    }
    item->dma_desc = (lldesc_t) {
        .size = buffer_size,
        .length = buffer_size,
        .eof = 0,
        .owner = 1,
        .buf = &item->buf,
    };
    //Actually send to the ring buffer for consumer to use
    res = xRingbufferSendComplete(buf_handle, &item);
    if (res != pdTRUE) {
        printf("Failed to send item\n");
    }
```

The following example demonstrates retrieving and returning an item from a **no-split ring buffer** using *xRingbufferReceive()* and *vRingbufferReturnItem()*

```
...

    //Receive an item from no-split ring buffer
    size_t item_size;
    char *item = (char *)xRingbufferReceive(buf_handle, &item_size, pdMS_TO_
↪TICKS(1000));

    //Check received item
    if (item != NULL) {
        //Print item
        for (int i = 0; i < item_size; i++) {
            printf("%c", item[i]);
        }
        printf("\n");
        //Return Item
        vRingbufferReturnItem(buf_handle, (void *)item);
    } else {
        //Failed to receive item
        printf("Failed to receive item\n");
    }
```

The following example demonstrates retrieving and returning an item from an **allow-split ring buffer** using *xRingbufferReceiveSplit()* and *vRingbufferReturnItem()*

```
...

    //Receive an item from allow-split ring buffer
    size_t item_size1, item_size2;
    char *item1, *item2;
    BaseType_t ret = xRingbufferReceiveSplit(buf_handle, (void **)&item1, (void
↪**)&item2, &item_size1, &item_size2, pdMS_TO_TICKS(1000));

    //Check received item
    if (ret == pdTRUE && item1 != NULL) {
        for (int i = 0; i < item_size1; i++) {
            printf("%c", item1[i]);
        }
        vRingbufferReturnItem(buf_handle, (void *)item1);
        //Check if item was split
        if (item2 != NULL) {
            for (int i = 0; i < item_size2; i++) {
                printf("%c", item2[i]);
            }
            vRingbufferReturnItem(buf_handle, (void *)item2);
        }
        printf("\n");
    } else {
        //Failed to receive item
        printf("Failed to receive item\n");
    }
```

The following example demonstrates retrieving and returning an item from a **byte buffer** using *xRingbufferReceiveUpTo()* and *vRingbufferReturnItem()*

```
...

    //Receive data from byte buffer
    size_t item_size;
    char *item = (char *)xRingbufferReceiveUpTo(buf_handle, &item_size, pdMS_TO_
↪TICKS(1000), sizeof(tx_item));

    //Check received data
    if (item != NULL) {
        //Print item
        for (int i = 0; i < item_size; i++) {
            printf("%c", item[i]);
        }
        printf("\n");
        //Return Item
        vRingbufferReturnItem(buf_handle, (void *)item);
    } else {
        //Failed to receive item
        printf("Failed to receive item\n");
    }
```

For ISR safe versions of the functions used above, call *xRingbufferSendFromISR()*, *xRingbufferReceiveFromISR()*, *xRingbufferReceiveSplitFromISR()*, *xRingbufferReceiveUpToFromISR()*, and *vRingbufferReturnItemFromISR()*

Sending to Ring Buffer The following diagrams illustrate the differences between no-split/allow-split buffers and byte buffers with regards to sending items/data. The diagrams assume that three items of sizes **18, 3, and 27 bytes** are sent respectively to a **buffer of 128 bytes**.

For no-split/allow-split buffers, a header of 8 bytes precedes every data item. Furthermore, the space occupied by each item is **rounded up to the nearest 32-bit aligned size** in order to maintain overall 32-bit alignment. However

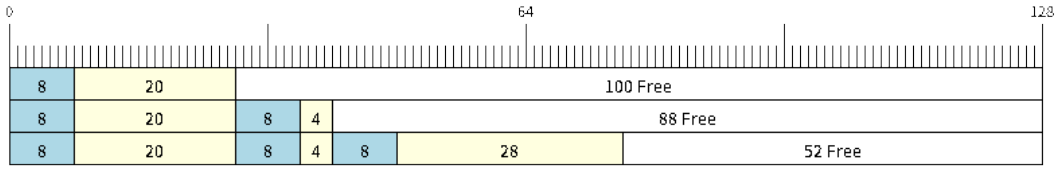

图 20: Sending items to no-split/allow-split ring buffers

the true size of the item is recorded inside the header which will be returned when the item is retrieved.

Referring to the diagram above, the 18, 3, and 27 byte items are **rounded up to 20, 4, and 28 bytes** respectively. An 8 byte header is then added in front of each item.

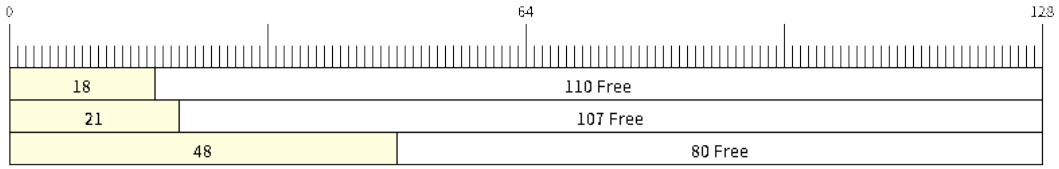

图 21: Sending items to byte buffers

Byte buffers treat data as a sequence of bytes and does not incur any overhead (no headers). As a result, all data sent to a byte buffer is merged into a single item.

Referring to the diagram above, the 18, 3, and 27 byte items are sequentially written to the byte buffer and **merged into a single item of 48 bytes**.

Using SendAcquire and SendComplete Items in no-split buffers are acquired (by SendAcquire) in strict FIFO order and must be sent to the buffer by SendComplete for the data to be accessible by the consumer. Multiple items can be sent or acquired without calling SendComplete, and the items do not necessarily need to be completed in the order they were acquired. However the receiving of data items must occur in FIFO order, therefore not calling SendComplete the earliest acquired item will prevent the subsequent items from being received.

The following diagrams illustrate what will happen when SendAcquire/SendComplete don't happen in the same order. At the beginning, there is already an data item of 16 bytes sent to the ring buffer. Then SendAcquire is called to acquire space of 20, 8, 24 bytes on the ring buffer.

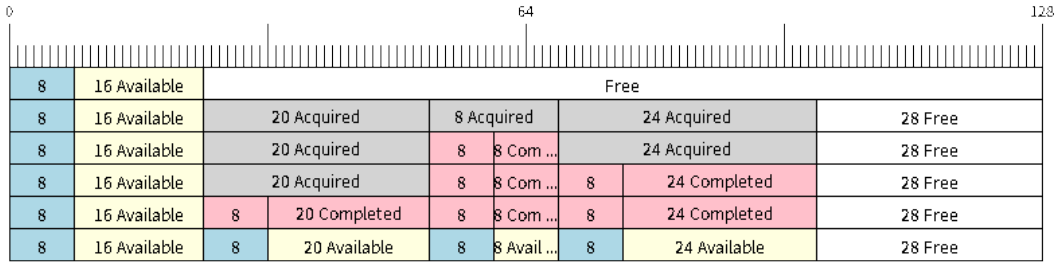

图 22: SendAcquire/SendComplete items in no-split ring buffers

After that, we fill (use) the buffers, and send them to the ring buffer by SendComplete in the order of 8, 24, 20. When

8 bytes and 24 bytes data are sent, the consumer still can only get the 16 bytes data item. Due to the usage if 20 bytes item is not complete, it’ s not available, nor the following data items.

When the 20 bytes item is finally completed, all the 3 data items can be received now, in the order of 20, 8, 24 bytes, right after the 16 bytes item existing in the buffer at the beginning.

Allow-split/byte buffers do not allow using SendAcquire/SendComplete since acquired buffers are required to be complete (not wrapped).

Wrap around The following diagrams illustrate the differences between no-split, allow-split, and byte buffers when a sent item requires a wrap around. The diagrams assumes a buffer of **128 bytes** with **56 bytes of free space that wraps around** and a sent item of **28 bytes**.

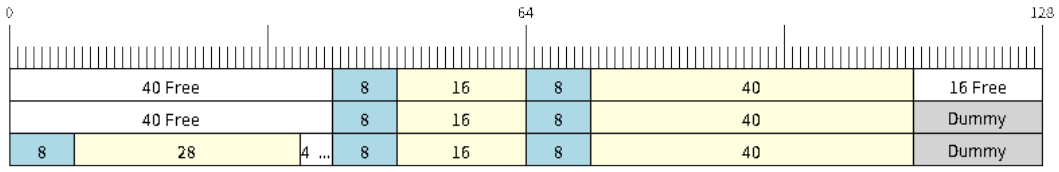

图 23: Wrap around in no-split buffers

No-split buffers will **only store an item in continuous free space and will not split an item under any circumstances**. When the free space at the tail of the buffer is insufficient to completely store the item and its header, the free space at the tail will be **marked as dummy data**. The buffer will then wrap around and store the item in the free space at the head of the buffer.

Referring to the diagram above, the 16 bytes of free space at the tail of the buffer is insufficient to store the 28 byte item. Therefore the 16 bytes is marked as dummy data and the item is written to the free space at the head of the buffer instead.

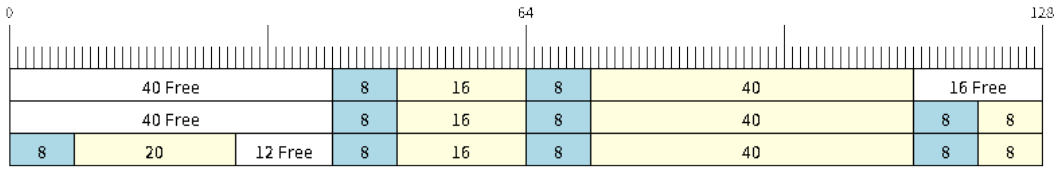

图 24: Wrap around in allow-split buffers

Allow-split buffers will attempt to **split the item into two parts** when the free space at the tail of the buffer is insufficient to store the item data and its header. Both parts of the split item will have their own headers (therefore incurring an extra 8 bytes of overhead).

Referring to the diagram above, the 16 bytes of free space at the tail of the buffer is insufficient to store the 28 byte item. Therefore the item is split into two parts (8 and 20 bytes) and written as two parts to the buffer.

注解: Allow-split buffers treats the both parts of the split item as two separate items, therefore call `xRingbufferReceiveSplit()` instead of `xRingbufferReceive()` to receive both parts of a split item in a thread safe manner.

Byte buffers will **store as much data as possible into the free space at the tail of buffer**. The remaining data will then be stored in the free space at the head of the buffer. No overhead is incurred when wrapping around in byte buffers.

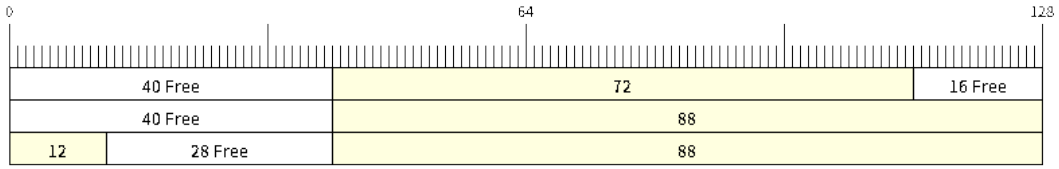

图 25: Wrap around in byte buffers

Referring to the diagram above, the 16 bytes of free space at the tail of the buffer is insufficient to completely store the 28 bytes of data. Therefore the 16 bytes of free space is filled with data, and the remaining 12 bytes are written to the free space at the head of the buffer. The buffer now contains data in two separate continuous parts, and each part continuous will be treated as a separate item by the byte buffer.

Retrieving/Returning The following diagrams illustrates the differences between no-split/allow-split and byte buffers in retrieving and returning data.

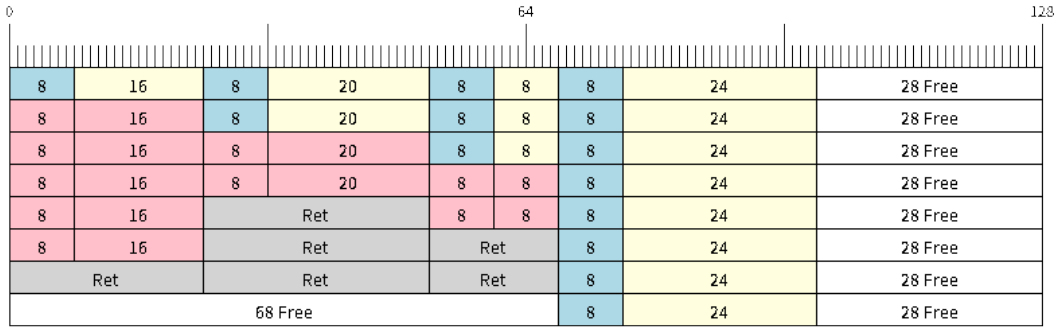

图 26: Retrieving/Returning items in no-split/allow-split ring buffers

Items in no-split/allow-split buffers are **retrieved in strict FIFO order** and **must be returned** for the occupied space to be freed. Multiple items can be retrieved before returning, and the items do not necessarily need to be returned in the order they were retrieved. However the freeing of space must occur in FIFO order, therefore not returning the earliest retrieved item will prevent the space of subsequent items from being freed.

Referring to the diagram above, the **16, 20, and 8 byte items are retrieved in FIFO order**. However the items are not returned in they were retrieved (20, 8, 16). As such, the space is not freed until the first item (16 byte) is returned.

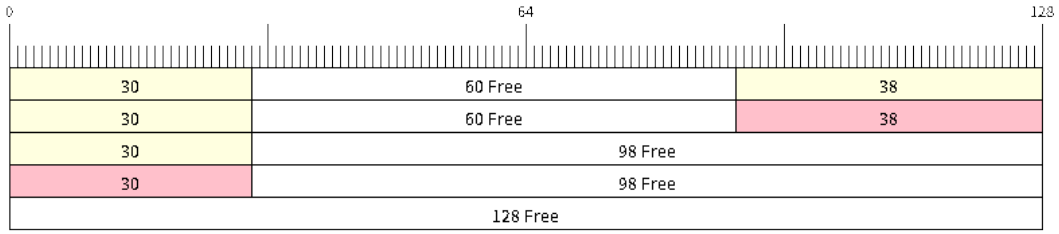

图 27: Retrieving/Returning data in byte buffers

Byte buffers **do not allow multiple retrievals before returning** (every retrieval must be followed by a return

before another retrieval is permitted). When using *xRingbufferReceive()* or *xRingbufferReceiveFromISR()*, all continuous stored data will be retrieved. *xRingbufferReceiveUpTo()* or *xRingbufferReceiveUpToFromISR()* can be used to restrict the maximum number of bytes retrieved. Since every retrieval must be followed by a return, the space will be freed as soon as the data is returned.

Referring to the diagram above, the 38 bytes of continuous stored data at the tail of the buffer is retrieved, returned, and freed. The next call to *xRingbufferReceive()* or *xRingbufferReceiveFromISR()* then wraps around and does the same to the 30 bytes of continuous stored data at the head of the buffer.

Ring Buffers with Queue Sets Ring buffers can be added to FreeRTOS queue sets using *xRingbufferAddToQueueSetRead()* such that every time a ring buffer receives an item or data, the queue set is notified. Once added to a queue set, every attempt to retrieve an item from a ring buffer should be preceded by a call to *xQueueSelectFromSet()*. To check whether the selected queue set member is the ring buffer, call *xRingbufferCanRead()*.

The following example demonstrates queue set usage with ring buffers.

```
#include "freertos/queue.h"
#include "freertos/ringbuf.h"

...

    //Create ring buffer and queue set
    RingbufHandle_t buf_handle = xRingbufferCreate(1028, RINGBUF_TYPE_NOSPLIT);
    QueueSetHandle_t queue_set = xQueueCreateSet(3);

    //Add ring buffer to queue set
    if (xRingbufferAddToQueueSetRead(buf_handle, queue_set) != pdTRUE) {
        printf("Failed to add to queue set\n");
    }

...

    //Block on queue set
    xQueueSetMemberHandle member = xQueueSelectFromSet(queue_set, pdMS_TO_
↪TICKS(1000));

    //Check if member is ring buffer
    if (member != NULL && xRingbufferCanRead(buf_handle, member) == pdTRUE) {
        //Member is ring buffer, receive item from ring buffer
        size_t item_size;
        char *item = (char *)xRingbufferReceive(buf_handle, &item_size, 0);

        //Handle item
        ...

    } else {
        ...
    }
```

Ring Buffers with Static Allocation The *xRingbufferCreateStatic()* can be used to create ring buffers with specific memory requirements (such as a ring buffer being allocated in external RAM). All blocks of memory used by a ring buffer must be manually allocated beforehand then passed to the *xRingbufferCreateStatic()* to be initialized as a ring buffer. These blocks include the following:

- The ring buffer's data structure of type *StaticRingbuffer_t*
- The ring buffer's storage area of size `xBufferSize`. Note that `xBufferSize` must be 32-bit aligned for no-split/allow-split buffers.

The manner in which these blocks are allocated will depend on the users requirements (e.g. all blocks being statically declared, or dynamically allocated with specific capabilities such as external RAM).

注解: The *CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION* option must be enabled in *menuconfig* for statically allocated ring buffers to be available.

注解: When deleting a ring buffer created via *xRingbufferCreateStatic()*, the function *vRingbufferDelete()* will not free any of the memory blocks. This must be done manually by the user after *vRingbufferDelete()* is called.

The code snippet below demonstrates a ring buffer being allocated entirely in external RAM.

```
#include "freertos/ringbuf.h"
#include "freertos/semphr.h"
#include "esp_heap_caps.h"

#define BUFFER_SIZE     400      //32-bit aligned size
#define BUFFER_TYPE     RINGBUF_TYPE_NOSPLIT
...

//Allocate ring buffer data structure and storage area into external RAM
StaticRingbuffer_t *buffer_struct = (StaticRingbuffer_t *)heap_caps_
↪malloc(sizeof(StaticRingbuffer_t), MALLOC_CAP_SPIRAM);
uint8_t *buffer_storage = (uint8_t *)heap_caps_malloc(sizeof(uint8_t)*BUFFER_SIZE,␣
↪MALLOC_CAP_SPIRAM);

//Create a ring buffer with manually allocated memory
RingbufHandle_t handle = xRingbufferCreateStatic(BUFFER_SIZE, BUFFER_TYPE, buffer_
↪storage, buffer_struct);

...

//Delete the ring buffer after used
vRingbufferDelete(handle);

//Manually free all blocks of memory
free(buffer_struct);
free(buffer_storage);
```

Ring Buffer API Reference

注解: Ideally, ring buffers can be used with multiple tasks in an SMP fashion where the **highest priority task will always be serviced first.** However due to the usage of binary semaphores in the ring buffer's underlying implementation, priority inversion may occur under very specific circumstances.

The ring buffer governs sending by a binary semaphore which is given whenever space is freed on the ring buffer. The highest priority task waiting to send will repeatedly take the semaphore until sufficient free space becomes available or until it times out. Ideally this should prevent any lower priority tasks from being serviced as the semaphore should always be given to the highest priority task.

However in between iterations of acquiring the semaphore, there is a **gap in the critical section** which may permit another task (on the other core or with an even higher priority) to free some space on the ring buffer and as a result give the semaphore. Therefore the semaphore will be given before the highest priority task can re-acquire the semaphore. This will result in the **semaphore being acquired by the second highest priority task** waiting to send, hence causing priority inversion.

This side effect will not affect ring buffer performance drastically given if the number of tasks using the ring buffer simultaneously is low, and the ring buffer is not operating near maximum capacity.

Header File

- esp_ringbuf/include/freertos/ringbuf.h

Functions

RingbufHandle_t **xRingbufferCreate**(size_t *xBufferSize*, *RingbufferType_t* *xBufferType*)

Create a ring buffer.

Note xBufferSize of no-split/allow-split buffers will be rounded up to the nearest 32-bit aligned size.

Return A handle to the created ring buffer, or NULL in case of error.

Parameters

- `[in] xBufferSize`: Size of the buffer in bytes. Note that items require space for overhead in no-split/allow-split buffers
- `[in] xBufferType`: Type of ring buffer, see documentation.

RingbufHandle_t **xRingbufferCreateNoSplit**(size_t *xItemSize*, size_t *xItemNum*)

Create a ring buffer of type RINGBUF_TYPE_NOSPLIT for a fixed item_size.

This API is similar to xRingbufferCreate(), but it will internally allocate additional space for the headers.

Return A RingbufHandle_t handle to the created ring buffer, or NULL in case of error.

Parameters

- `[in] xItemSize`: Size of each item to be put into the ring buffer
- `[in] xItemNum`: Maximum number of items the buffer needs to hold simultaneously

RingbufHandle_t **xRingbufferCreateStatic**(size_t *xBufferSize*, *RingbufferType_t* *xBufferType*, uint8_t **pucRingbufferStorage*, *StaticRingbuffer_t* **pxStaticRingbuffer*)

Create a ring buffer but manually provide the required memory.

Note The CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION option must be enabled for this to be available

Note xBufferSize of no-split/allow-split buffers MUST be 32-bit aligned.

Return A handle to the created ring buffer

Parameters

- `[in] xBufferSize`: Size of the buffer in bytes.
- `[in] xBufferType`: Type of ring buffer, see documentation
- `[in] pucRingbufferStorage`: Pointer to the ring buffer's storage area. Storage area must of the same size as specified by xBufferSize
- `[in] pxStaticRingbuffer`: Pointed to a struct of type StaticRingbuffer_t which will be used to hold the ring buffer's data structure

BaseType_t **xRingbufferSend**(*RingbufHandle_t* *xRingbuffer*, **const** void **pvItem*, size_t *xItemSize*, TickType_t *xTicksToWait*)

Insert an item into the ring buffer.

Attempt to insert an item into the ring buffer. This function will block until enough free space is available or until it times out.

Note For no-split/allow-split ring buffers, the actual size of memory that the item will occupy will be rounded up to the nearest 32-bit aligned size. This is done to ensure all items are always stored in 32-bit aligned fashion.

Return

- pdTRUE if succeeded
- pdFALSE on time-out or when the data is larger than the maximum permissible size of the buffer

Parameters

- `[in] xRingbuffer`: Ring buffer to insert the item into
- `[in] pvItem`: Pointer to data to insert. NULL is allowed if xItemSize is 0.
- `[in] xItemSize`: Size of data to insert.
- `[in] xTicksToWait`: Ticks to wait for room in the ring buffer.

BaseType_t **xRingbufferSendFromISR**(*RingbufHandle_t* *xRingbuffer*, **const** void **pvItem*, size_t *xItemSize*, BaseType_t **pxHigherPriorityTaskWoken*)

Insert an item into the ring buffer in an ISR.

Attempt to insert an item into the ring buffer from an ISR. This function will return immediately if there is insufficient free space in the buffer.

Note For no-split/allow-split ring buffers, the actual size of memory that the item will occupy will be rounded up to the nearest 32-bit aligned size. This is done to ensure all items are always stored in 32-bit aligned fashion.

Return

- pdTRUE if succeeded
- pdFALSE when the ring buffer does not have space.

Parameters

- `[in] xRingbuffer`: Ring buffer to insert the item into
- `[in] pvItem`: Pointer to data to insert. NULL is allowed if xItemSize is 0.
- `[in] xItemSize`: Size of data to insert.
- `[out] pxHigherPriorityTaskWoken`: Value pointed to will be set to pdTRUE if the function woke up a higher priority task.

BaseType_t **`xRingbufferSendAcquire`**(*RingbufHandle_t xRingbuffer*, void **ppvItem*, size_t *xItemSize*, TickType_t *xTicksToWait*)

Acquire memory from the ring buffer to be written to by an external source and to be sent later.

Attempt to allocate buffer for an item to be sent into the ring buffer. This function will block until enough free space is available or until it timesout.

The item, as well as the following items `SendAcquire` or `Send` after it, will not be able to be read from the ring buffer until this item is actually sent into the ring buffer.

Note Only applicable for no-split ring buffers now, the actual size of memory that the item will occupy will be rounded up to the nearest 32-bit aligned size. This is done to ensure all items are always stored in 32-bit aligned fashion.

Return

- pdTRUE if succeeded
- pdFALSE on time-out or when the data is larger than the maximum permissible size of the buffer

Parameters

- `[in] xRingbuffer`: Ring buffer to allocate the memory
- `[out] ppvItem`: Double pointer to memory acquired (set to NULL if no memory were retrieved)
- `[in] xItemSize`: Size of item to acquire.
- `[in] xTicksToWait`: Ticks to wait for room in the ring buffer.

BaseType_t **`xRingbufferSendComplete`**(*RingbufHandle_t xRingbuffer*, void **pvItem*)

Actually send an item into the ring buffer allocated before by `xRingbufferSendAcquire`.

Note Only applicable for no-split ring buffers. Only call for items allocated by `xRingbufferSendAcquire`.

Return

- pdTRUE if succeeded
- pdFALSE if fail for some reason.

Parameters

- `[in] xRingbuffer`: Ring buffer to insert the item into
- `[in] pvItem`: Pointer to item in allocated memory to insert.

void ***`xRingbufferReceive`**(*RingbufHandle_t xRingbuffer*, size_t **pxItemSize*, TickType_t *xTicksToWait*)

Retrieve an item from the ring buffer.

Attempt to retrieve an item from the ring buffer. This function will block until an item is available or until it times out.

Note A call to vRingbufferReturnItem() is required after this to free the item retrieved.

Return

- Pointer to the retrieved item on success; *pxItemSize filled with the length of the item.
- NULL on timeout, *pxItemSize is untouched in that case.

Parameters

- `[in] xRingbuffer`: Ring buffer to retrieve the item from
- `[out] pxItemSize`: Pointer to a variable to which the size of the retrieved item will be written.
- `[in] xTicksToWait`: Ticks to wait for items in the ring buffer.

void ***xRingbufferReceiveFromISR**(*RingbufHandle_t xRingbuffer*, size_t **pxItemSize*)

Retrieve an item from the ring buffer in an ISR.

Attempt to retrieve an item from the ring buffer. This function returns immediately if there are no items available for retrieval

Note A call to vRingbufferReturnItemFromISR() is required after this to free the item retrieved.
Note Byte buffers do not allow multiple retrievals before returning an item
Return

- Pointer to the retrieved item on success; *pxItemSize filled with the length of the item.
- NULL when the ring buffer is empty, *pxItemSize is untouched in that case.

Parameters

- `[in] xRingbuffer`: Ring buffer to retrieve the item from
- `[out] pxItemSize`: Pointer to a variable to which the size of the retrieved item will be written.

BaseType_t **xRingbufferReceiveSplit**(*RingbufHandle_t xRingbuffer*, void **ppvHeadItem*, void **ppvTailItem*, size_t **pxHeadItemSize*, size_t **pxTailItemSize*, TickType_t *xTicksToWait*)

Retrieve a split item from an allow-split ring buffer.

Attempt to retrieve a split item from an allow-split ring buffer. If the item is not split, only a single item is retried. If the item is split, both parts will be retrieved. This function will block until an item is available or until it times out.

Note Call(s) to vRingbufferReturnItem() is required after this to free up the item(s) retrieved.
Note This function should only be called on allow-split buffers
Return

- pdTRUE if an item (split or unsplit) was retrieved
- pdFALSE when no item was retrieved

Parameters

- `[in] xRingbuffer`: Ring buffer to retrieve the item from
- `[out] ppvHeadItem`: Double pointer to first part (set to NULL if no items were retrieved)
- `[out] ppvTailItem`: Double pointer to second part (set to NULL if item is not split)
- `[out] pxHeadItemSize`: Pointer to size of first part (unmodified if no items were retrieved)
- `[out] pxTailItemSize`: Pointer to size of second part (unmodified if item is not split)
- `[in] xTicksToWait`: Ticks to wait for items in the ring buffer.

BaseType_t **xRingbufferReceiveSplitFromISR**(*RingbufHandle_t xRingbuffer*, void **ppvHeadItem*, void **ppvTailItem*, size_t **pxHeadItemSize*, size_t **pxTailItemSize*)

Retrieve a split item from an allow-split ring buffer in an ISR.

Attempt to retrieve a split item from an allow-split ring buffer. If the item is not split, only a single item is retried. If the item is split, both parts will be retrieved. This function returns immediately if there are no items available for retrieval

Note Calls to vRingbufferReturnItemFromISR() is required after this to free up the item(s) retrieved.
Note This function should only be called on allow-split buffers
Return

- pdTRUE if an item (split or unsplit) was retrieved
- pdFALSE when no item was retrieved

Parameters

- `[in] xRingbuffer`: Ring buffer to retrieve the item from
- `[out] ppvHeadItem`: Double pointer to first part (set to NULL if no items were retrieved)
- `[out] ppvTailItem`: Double pointer to second part (set to NULL if item is not split)
- `[out] pxHeadItemSize`: Pointer to size of first part (unmodified if no items were retrieved)
- `[out] pxTailItemSize`: Pointer to size of second part (unmodified if item is not split)

void ***xRingbufferReceiveUpTo**(*RingbufHandle_t xRingbuffer*, size_t *pxItemSize, TickType_t *xTicksToWait*, size_t *xMaxSize*)

Retrieve bytes from a byte buffer, specifying the maximum amount of bytes to retrieve.

Attempt to retrieve data from a byte buffer whilst specifying a maximum number of bytes to retrieve. This function will block until there is data available for retrieval or until it times out.

Note A call to vRingbufferReturnItem() is required after this to free up the data retrieved.
Note This function should only be called on byte buffers
Note Byte buffers do not allow multiple retrievals before returning an item
Return

- Pointer to the retrieved item on success; *pxItemSize filled with the length of the item.
- NULL on timeout, *pxItemSize is untouched in that case.

Parameters

- `[in] xRingbuffer`: Ring buffer to retrieve the item from
- `[out] pxItemSize`: Pointer to a variable to which the size of the retrieved item will be written.
- `[in] xTicksToWait`: Ticks to wait for items in the ring buffer.
- `[in] xMaxSize`: Maximum number of bytes to return.

void ***xRingbufferReceiveUpToFromISR**(*RingbufHandle_t xRingbuffer*, size_t *pxItemSize, size_t *xMaxSize*)

Retrieve bytes from a byte buffer, specifying the maximum amount of bytes to retrieve. Call this from an ISR.

Attempt to retrieve bytes from a byte buffer whilst specifying a maximum number of bytes to retrieve. This function will return immediately if there is no data available for retrieval.

Note A call to vRingbufferReturnItemFromISR() is required after this to free up the data received.
Note This function should only be called on byte buffers
Note Byte buffers do not allow multiple retrievals before returning an item
Return

- Pointer to the retrieved item on success; *pxItemSize filled with the length of the item.
- NULL when the ring buffer is empty, *pxItemSize is untouched in that case.

Parameters

- `[in] xRingbuffer`: Ring buffer to retrieve the item from
- `[out] pxItemSize`: Pointer to a variable to which the size of the retrieved item will be written.
- `[in] xMaxSize`: Maximum number of bytes to return.

void **vRingbufferReturnItem**(*RingbufHandle_t xRingbuffer*, void *pvItem)

Return a previously-retrieved item to the ring buffer.

Note If a split item is retrieved, both parts should be returned by calling this function twice
Parameters

- `[in] xRingbuffer`: Ring buffer the item was retrieved from
- `[in] pvItem`: Item that was received earlier

void **vRingbufferReturnItemFromISR**(*RingbufHandle_t xRingbuffer*, void *pvItem, BaseType_t *pxHigherPriorityTaskWoken*)

Return a previously-retrieved item to the ring buffer from an ISR.

Note If a split item is retrieved, both parts should be returned by calling this function twice
Parameters

- `[in] xRingbuffer`: Ring buffer the item was retrieved from
- `[in] pvItem`: Item that was received earlier
- `[out] pxHigherPriorityTaskWoken`: Value pointed to will be set to pdTRUE if the function woke up a higher priority task.

void **vRingbufferDelete**(*RingbufHandle_t xRingbuffer*)

Delete a ring buffer.

Note This function will not deallocate any memory if the ring buffer was created using xRingbufferCreateStatic(). Deallocation must be done manually be the user.
Parameters

- `[in] xRingbuffer`: Ring buffer to delete

size_t **xRingbufferGetMaxItemSize**(*RingbufHandle_t xRingbuffer*)

Get maximum size of an item that can be placed in the ring buffer.

This function returns the maximum size an item can have if it was placed in an empty ring buffer.

Note The max item size for a no-split buffer is limited to ((buffer_size/2)-header_size). This limit is imposed so that an item of max item size can always be sent to the an empty no-split buffer regardless of the internal positions of the buffer's read/write/free pointers.

Return Maximum size, in bytes, of an item that can be placed in a ring buffer.

Parameters

- `[in] xRingbuffer`: Ring buffer to query

size_t **xRingbufferGetCurFreeSize**(*RingbufHandle_t xRingbuffer*)

Get current free size available for an item/data in the buffer.

This gives the real time free space available for an item/data in the ring buffer. This represents the maximum size an item/data can have if it was currently sent to the ring buffer.

Warning This API is not thread safe. So, if multiple threads are accessing the same ring buffer, it is the application's responsibility to ensure atomic access to this API and the subsequent Send

Note An empty no-split buffer has a max current free size for an item that is limited to ((buffer_size/2)-header_size). See API reference for xRingbufferGetMaxItemSize().

Return Current free size, in bytes, available for an entry

Parameters

- `[in] xRingbuffer`: Ring buffer to query

BaseType_t **xRingbufferAddToQueueSetRead**(*RingbufHandle_t xRingbuffer*, *QueueSetHandle_t xQueueSet*)

Add the ring buffer's read semaphore to a queue set.

The ring buffer's read semaphore indicates that data has been written to the ring buffer. This function adds the ring buffer's read semaphore to a queue set.

Return

- pdTRUE on success, pdFALSE otherwise

Parameters

- `[in] xRingbuffer`: Ring buffer to add to the queue set
- `[in] xQueueSet`: Queue set to add the ring buffer's read semaphore to

BaseType_t **xRingbufferCanRead**(*RingbufHandle_t xRingbuffer*, *QueueSetMemberHandle_t xMember*)

Check if the selected queue set member is the ring buffer's read semaphore.

This API checks if queue set member returned from xQueueSelectFromSet() is the read semaphore of this ring buffer. If so, this indicates the ring buffer has items waiting to be retrieved.

Return

- pdTRUE when semaphore belongs to ring buffer
- pdFALSE otherwise.

Parameters

- `[in] xRingbuffer`: Ring buffer which should be checked
- `[in] xMember`: Member returned from xQueueSelectFromSet

BaseType_t **xRingbufferRemoveFromQueueSetRead**(*RingbufHandle_t xRingbuffer*, *QueueSetHandle_t xQueueSet*)

Remove the ring buffer's read semaphore from a queue set.

This specifically removes a ring buffer's read semaphore from a queue set. The read semaphore is used to indicate when data has been written to the ring buffer

Return

- pdTRUE on success
- pdFALSE otherwise

Parameters

- `[in] xRingbuffer`: Ring buffer to remove from the queue set
- `[in] xQueueSet`: Queue set to remove the ring buffer's read semaphore from

void **vRingbufferGetInfo**(*RingbufHandle_t xRingbuffer*, UBaseType_t **uxFree*, UBaseType_t **uxRead*, UBaseType_t **uxWrite*, UBaseType_t **uxAcquire*, UBaseType_t **uxItemsWaiting*)

Get information about ring buffer status.

Get information of the a ring buffer's current status such as free/read/write pointer positions, and number of items waiting to be retrieved. Arguments can be set to NULL if they are not required.

Parameters

- `[in] xRingbuffer`: Ring buffer to remove from the queue set
- `[out] uxFree`: Pointer use to store free pointer position
- `[out] uxRead`: Pointer use to store read pointer position
- `[out] uxWrite`: Pointer use to store write pointer position
- `[out] uxAcquire`: Pointer use to store acquire pointer position
- `[out] uxItemsWaiting`: Pointer use to store number of items (bytes for byte buffer) waiting to be retrieved

void **xRingbufferPrintInfo**(*RingbufHandle_t xRingbuffer*)

Debugging function to print the internal pointers in the ring buffer.

Parameters

- `xRingbuffer`: Ring buffer to show

Structures

struct xSTATIC_RINGBUFFER

Struct that is equivalent in size to the ring buffer's data structure.

The contents of this struct are not meant to be used directly. This structure is meant to be used when creating a statically allocated ring buffer where this struct is of the exact size required to store a ring buffer's control data structure.

Note The CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION option must be enabled for this structure to be available.

Type Definitions

typedef void ***RingbufHandle_t**

Type by which ring buffers are referenced. For example, a call to xRingbufferCreate() returns a RingbufHandle_t variable that can then be used as a parameter to xRingbufferSend(), xRingbufferReceive(), etc.

typedef struct *xSTATIC_RINGBUFFER* **StaticRingbuffer_t**

Struct that is equivalent in size to the ring buffer's data structure.

The contents of this struct are not meant to be used directly. This structure is meant to be used when creating a statically allocated ring buffer where this struct is of the exact size required to store a ring buffer's control data structure.

Note The CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION option must be enabled for this structure to be available.

Enumerations

enum RingbufferType_t

Values:

RINGBUF_TYPE_NOSPLIT = 0

No-split buffers will only store an item in contiguous memory and will never split an item. Each item requires an 8 byte overhead for a header and will always internally occupy a 32-bit aligned size of space.

RINGBUF_TYPE_ALLOWSPLIT

Allow-split buffers will split an item into two parts if necessary in order to store it. Each item requires an 8 byte overhead for a header, splitting incurs an extra header. Each item will always internally occupy a 32-bit aligned size of space.

RINGBUF_TYPE_BYTEBUF
Byte buffers store data as a sequence of bytes and do not maintain separate items, therefore byte buffers have no overhead. All data is stored as a sequence of byte and any number of bytes can be sent or retrieved each time.

RINGBUF_TYPE_MAX

Hooks

FreeRTOS consists of Idle Hooks and Tick Hooks which allow for application specific functionality to be added to the Idle Task and Tick Interrupt. ESP-IDF provides its own Idle and Tick Hook API in addition to the hooks provided by Vanilla FreeRTOS. ESP-IDF hooks have the added benefit of being run time configurable and asymmetrical.

Vanilla FreeRTOS Hooks Idle and Tick Hooks in vanilla FreeRTOS are implemented by the user defining the functions `vApplicationIdleHook()` and `vApplicationTickHook()` respectively somewhere in the application. Vanilla FreeRTOS will run the user defined Idle Hook and Tick Hook on every iteration of the Idle Task and Tick Interrupt respectively.

Vanilla FreeRTOS hooks are referred to as **Legacy Hooks** in ESP-IDF FreeRTOS. To enable legacy hooks, *CONFIG_FREERTOS_LEGACY_HOOKS* should be enabled in *project configuration menu*.

Due to vanilla FreeRTOS being designed for single core, `vApplicationIdleHook()` and `vApplicationTickHook()` can only be defined once. However, the ESP32 is dual core in nature, therefore same Idle Hook and Tick Hook are used for both cores (in other words, the hooks are symmetrical for both cores).

ESP-IDF Idle and Tick Hooks Due to the the dual core nature of the ESP32, it may be necessary for some applications to have separate hooks for each core. Furthermore, it may be necessary for the Idle Tasks or Tick Interrupts to execute multiple hooks that are configurable at run time. Therefore the ESP-IDF provides it's own hooks API in addition to the legacy hooks provided by Vanilla FreeRTOS.

The ESP-IDF tick/idle hooks are registered at run time, and each tick/idle hook must be registered to a specific CPU. When the idle task runs/tick Interrupt occurs on a particular CPU, the CPU will run each of its registered idle/tick hooks in turn.

Hooks API Reference

Header File

- esp_common/include/esp_freertos_hooks.h

Functions

esp_err_t **esp_register_freertos_idle_hook_for_cpu**(*esp_freertos_idle_cb_t* *new_idle_cb*, UBaseType_t *cpuid*)

Register a callback to be called from the specified core's idle hook. The callback should return true if it should be called by the idle hook once per interrupt (or FreeRTOS tick), and return false if it should be called repeatedly as fast as possible by the idle hook.

Warning Idle callbacks MUST NOT, UNDER ANY CIRCUMSTANCES, CALL A FUNCTION THAT MIGHT BLOCK.

Return
- ESP_OK: Callback registered to the specified core's idle hook
- ESP_ERR_NO_MEM: No more space on the specified core's idle hook to register callback
- ESP_ERR_INVALID_ARG: cpuid is invalid

Parameters
- `[in] new_idle_cb`: Callback to be called
- `[in] cpuid`: id of the core

esp_err_t **esp_register_freertos_idle_hook**(*esp_freertos_idle_cb_t new_idle_cb*)

Register a callback to the idle hook of the core that calls this function. The callback should return true if it should be called by the idle hook once per interrupt (or FreeRTOS tick), and return false if it should be called repeatedly as fast as possible by the idle hook.

Warning Idle callbacks MUST NOT, UNDER ANY CIRCUMSTANCES, CALL A FUNCTION THAT MIGHT BLOCK.

Return

- ESP_OK: Callback registered to the calling core' s idle hook
- ESP_ERR_NO_MEM: No more space on the calling core' s idle hook to register callback

Parameters

- `[in] new_idle_cb`: Callback to be called

esp_err_t **esp_register_freertos_tick_hook_for_cpu**(*esp_freertos_tick_cb_t new_tick_cb*, UBaseType_t *cpuid*)

Register a callback to be called from the specified core' s tick hook.

Return

- ESP_OK: Callback registered to specified core' s tick hook
- ESP_ERR_NO_MEM: No more space on the specified core' s tick hook to register the callback
- ESP_ERR_INVALID_ARG: cpuid is invalid

Parameters

- `[in] new_tick_cb`: Callback to be called
- `[in] cpuid`: id of the core

esp_err_t **esp_register_freertos_tick_hook**(*esp_freertos_tick_cb_t new_tick_cb*)

Register a callback to be called from the calling core' s tick hook.

Return

- ESP_OK: Callback registered to the calling core' s tick hook
- ESP_ERR_NO_MEM: No more space on the calling core' s tick hook to register the callback

Parameters

- `[in] new_tick_cb`: Callback to be called

void **esp_deregister_freertos_idle_hook_for_cpu**(*esp_freertos_idle_cb_t old_idle_cb*, UBaseType_t *cpuid*)

Unregister an idle callback from the idle hook of the specified core.

Parameters

- `[in] old_idle_cb`: Callback to be unregistered
- `[in] cpuid`: id of the core

void **esp_deregister_freertos_idle_hook**(*esp_freertos_idle_cb_t old_idle_cb*)

Unregister an idle callback. If the idle callback is registered to the idle hooks of both cores, the idle hook will be unregistered from both cores.

Parameters

- `[in] old_idle_cb`: Callback to be unregistered

void **esp_deregister_freertos_tick_hook_for_cpu**(*esp_freertos_tick_cb_t old_tick_cb*, UBaseType_t *cpuid*)

Unregister a tick callback from the tick hook of the specified core.

Parameters

- `[in] old_tick_cb`: Callback to be unregistered
- `[in] cpuid`: id of the core

void **esp_deregister_freertos_tick_hook**(*esp_freertos_tick_cb_t old_tick_cb*)

Unregister a tick callback. If the tick callback is registered to the tick hooks of both cores, the tick hook will be unregistered from both cores.

Parameters

- `[in] old_tick_cb`: Callback to be unregistered

Type Definitions

typedef bool (***esp_freertos_idle_cb_t**) (void)
typedef void (***esp_freertos_tick_cb_t**) (void)

2.6.11 Heap Memory Allocation

Stack and Heap

ESP-IDF applications use the common computer architecture patterns of *stack* (dynamic memory allocated by program control flow) and *heap* (dynamic memory allocated by function calls), as well as statically allocated memory (allocated at compile time).

Because ESP-IDF is a multi-threaded RTOS environment, each RTOS task has its own stack. By default, each of these stacks is allocated from the heap when the task is created. (See `xTaskCreateStatic()` for the alternative where stacks are statically allocated.)

Because ESP32-S2 uses multiple types of RAM, it also contains multiple heaps with different capabilities. A capabilities-based memory allocator allows apps to make heap allocations for different purposes.

For most purposes, the standard libc `malloc()` and `free()` functions can be used for heap allocation without any special consideration.

However, in order to fully make use of all of the memory types and their characteristics, ESP-IDF also has a capabilities-based heap memory allocator. If you want to have memory with certain properties (for example, *DMA-Capable Memory* or executable-memory), you can create an OR-mask of the required capabilities and pass that to `heap_caps_malloc()`.

Memory Capabilities

The ESP32-S2 contains multiple types of RAM:

- DRAM (Data RAM) is memory used to hold data. This is the most common kind of memory accessed as heap.
- IRAM (Instruction RAM) usually holds executable data only. If accessed as generic memory, all accesses must be *32-bit aligned*.
- D/IRAM is RAM which can be used as either Instruction or Data RAM.

For more details on these internal memory types, see *应用程序的内存布局*.

It's also possible to connect external SPI RAM to the ESP32-S2 - *external RAM* can be integrated into the ESP32-S2's memory map using the flash cache, and accessed similarly to DRAM.

DRAM uses capability `MALLOC_CAP_8BIT` (accessible in single byte reads and writes). When calling `malloc()`, the ESP-IDF `malloc()` implementation internally calls `heap_caps_malloc(size, MALLOC_CAP_8BIT)` in order to allocate DRAM that is byte-addressable. To test the free DRAM heap size at runtime, call cpp:func:*heap_caps_get_free_size(MALLOC_CAP_8BIT)*.

Because malloc uses the capabilities-based allocation system, memory allocated using `heap_caps_malloc()` can be freed by calling the standard `free()` function.

Available Heap

DRAM At startup, the DRAM heap contains all data memory which is not statically allocated by the app. Reducing statically allocated buffers will increase the amount of available free heap.

To find the amount of statically allocated memory, use the *idf.py size* command.

注解: Due to a technical limitation, the maximum statically allocated DRAM usage is 160KB. The remaining 160KB (for a total of 320KB of DRAM) can only be allocated at runtime as heap.

注解: At runtime, the available heap DRAM may be less than calculated at compile time, because at startup some memory is allocated from the heap before the FreeRTOS scheduler is started (including memory for the stacks of initial FreeRTOS tasks).

IRAM At startup, the IRAM heap contains all instruction memory which is not used by the app executable code.

The *idf.py size* command can be used to find the amount of IRAM used by the app.

D/IRAM Some memory in the ESP32-S2 is available as either DRAM or IRAM. If memory is allocated from a D/IRAM region, the free heap size for both types of memory will decrease.

Heap Sizes At startup, all ESP-IDF apps log a summary of all heap addresses (and sizes) at level Info:

```
I (252) heap_init: Initializing. RAM available for dynamic allocation:
I (259) heap_init: At 3FFAE6E0 len 00001920 (6 KiB): DRAM
I (265) heap_init: At 3FFB2EC8 len 0002D138 (180 KiB): DRAM
I (272) heap_init: At 3FFE0440 len 00003AE0 (14 KiB): D/IRAM
I (278) heap_init: At 3FFE4350 len 0001BCB0 (111 KiB): D/IRAM
I (284) heap_init: At 4008944C len 00016BB4 (90 KiB): IRAM
```

Finding available heap See *Heap Information*.

Special Capabilities

DMA-Capable Memory Use the `MALLOC_CAP_DMA` flag to allocate memory which is suitable for use with hardware DMA engines (for example SPI and I2S). This capability flag excludes any external PSRAM.

32-Bit Accessible Memory If a certain memory structure is only addressed in 32-bit units, for example an array of ints or pointers, it can be useful to allocate it with the `MALLOC_CAP_32BIT` flag. This also allows the allocator to give out IRAM memory; something which it can't do for a normal malloc() call. This can help to use all the available memory in the ESP32-S2.

Memory allocated with `MALLOC_CAP_32BIT` can *only* be accessed via 32-bit reads and writes, any other type of access will generate a fatal LoadStoreError exception.

External SPI Memory When *external RAM* is enabled, external SPI RAM under 4MiB in size can be allocated using standard `malloc` calls, or via `heap_caps_malloc(MALLOC_CAP_SPIRAM)`, depending on configuration. See 配置片外 *RAM* for more details.

API Reference - Heap Allocation

Header File

- heap/include/esp_heap_caps.h

Functions

esp_err_t **heap_caps_register_failed_alloc_callback**(*esp_alloc_failed_hook_t callback*)

registers a callback function to be invoked if a memory allocation operation fails

Return ESP_OK if callback was registered.

Parameters

- `callback`: caller defined callback to be invoked

void ***heap_caps_malloc**(size_t *size*, uint32_t *caps*)

Allocate a chunk of memory which has the given capabilities.

Equivalent semantics to libc malloc(), for capability-aware memory.

In IDF, `malloc(p)` is equivalent to `heap_caps_malloc(p, MALLOC_CAP_8BIT)`.

Return A pointer to the memory allocated on success, NULL on failure

Parameters

- `size`: Size, in bytes, of the amount of memory to allocate
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory to be returned

void **heap_caps_free**(void **ptr*)

Free memory previously allocated via heap_caps_malloc() or heap_caps_realloc().

Equivalent semantics to libc free(), for capability-aware memory.

In IDF, `free(p)` is equivalent to `heap_caps_free(p)`.

Parameters

- `ptr`: Pointer to memory previously returned from heap_caps_malloc() or heap_caps_realloc(). Can be NULL.

void ***heap_caps_realloc**(void **ptr*, size_t *size*, int *caps*)

Reallocate memory previously allocated via heap_caps_malloc() or heap_caps_realloc().

Equivalent semantics to libc realloc(), for capability-aware memory.

In IDF, `realloc(p, s)` is equivalent to `heap_caps_realloc(p, s, MALLOC_CAP_8BIT)`.

‘caps’ parameter can be different to the capabilities that any original ‘ptr’ was allocated with. In this way, realloc can be used to “move” a buffer if necessary to ensure it meets a new set of capabilities.

Return Pointer to a new buffer of size ‘size’ with capabilities ‘caps’, or NULL if allocation failed.

Parameters

- `ptr`: Pointer to previously allocated memory, or NULL for a new allocation.
- `size`: Size of the new buffer requested, or 0 to free the buffer.
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory desired for the new allocation.

void ***heap_caps_aligned_alloc**(size_t *alignment*, size_t *size*, int *caps*)

Allocate a aligned chunk of memory which has the given capabilities.

Equivalent semantics to libc aligned_alloc(), for capability-aware memory.

Return A pointer to the memory allocated on success, NULL on failure

Note Any memory allocated with heaps_caps_aligned_alloc() MUST be freed with heap_caps_aligned_free() and CANNOT be passed to free()

Parameters

- `alignment`: How the pointer received needs to be aligned must be a power of two
- `size`: Size, in bytes, of the amount of memory to allocate
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory to be returned

void ***heap_caps_aligned_calloc**(size_t *alignment*, size_t *n*, size_t *size*, uint32_t *caps*)

Allocate a aligned chunk of memory which has the given capabilities. The initialized value in the memory is set to zero.

Return A pointer to the memory allocated on success, NULL on failure

Note Any memory allocated with heap_caps_aligned_calloc() MUST be freed with heap_caps_aligned_free() and CANNOT be passed to free()

Parameters

- `alignment`: How the pointer received needs to be aligned must be a power of two
- `n`: Number of continuing chunks of memory to allocate
- `size`: Size, in bytes, of a chunk of memory to allocate
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory to be returned

void **heap_caps_aligned_free**(void *ptr)

Used to deallocate memory previously allocated with heap_caps_aligned_alloc.

Note This function is aimed to deallocate only memory allocated with heap_caps_aligned_alloc, memory allocated with heap_caps_malloc MUST not be passed to this function

Parameters

- `ptr`: Pointer to the memory allocated

void ***heap_caps_calloc**(size_t *n*, size_t *size*, uint32_t *caps*)

Allocate a chunk of memory which has the given capabilities. The initialized value in the memory is set to zero.

Equivalent semantics to libc calloc(), for capability-aware memory.

In IDF, `calloc(p)` is equivalent to `heap_caps_calloc(p, MALLOC_CAP_8BIT)`.

Return A pointer to the memory allocated on success, NULL on failure

Parameters

- `n`: Number of continuing chunks of memory to allocate
- `size`: Size, in bytes, of a chunk of memory to allocate
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory to be returned

size_t **heap_caps_get_total_size**(uint32_t *caps*)

Get the total size of all the regions that have the given capabilities.

This function takes all regions capable of having the given capabilities allocated in them and adds up the total space they have.

Return total size in bytes

Parameters

- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

size_t **heap_caps_get_free_size**(uint32_t *caps*)

Get the total free size of all the regions that have the given capabilities.

This function takes all regions capable of having the given capabilities allocated in them and adds up the free space they have.

Note that because of heap fragmentation it is probably not possible to allocate a single block of memory of this size. Use heap_caps_get_largest_free_block() for this purpose.

Return Amount of free bytes in the regions

Parameters

- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

size_t **heap_caps_get_minimum_free_size**(uint32_t *caps*)

Get the total minimum free memory of all regions with the given capabilities.

This adds all the low water marks of the regions capable of delivering the memory with the given capabilities.

Note the result may be less than the global all-time minimum available heap of this kind, as "low water marks" are tracked per-region. Individual regions' heaps may have reached their "low water marks" at different points in time. However this result still gives a "worst case" indication for all-time minimum free heap.

Return Amount of free bytes in the regions

Parameters

- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

size_t **heap_caps_get_largest_free_block**(uint32_t *caps*)

Get the largest free block of memory able to be allocated with the given capabilities.

Returns the largest value of `s` for which `heap_caps_malloc(s, caps)` will succeed.

Return Size of largest free block in bytes.

Parameters

- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

void **heap_caps_get_info**(*multi_heap_info_t* *info, uint32_t *caps*)
Get heap info for all regions with the given capabilities.

Calls multi_heap_info() on all heaps which share the given capabilities. The information returned is an aggregate across all matching heaps. The meanings of fields are the same as defined for *multi_heap_info_t*, except that `minimum_free_bytes` has the same caveats described in heap_caps_get_minimum_free_size().

Parameters
- `info`: Pointer to a structure which will be filled with relevant heap metadata.
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

void **heap_caps_print_heap_info**(uint32_t *caps*)
Print a summary of all memory with the given capabilities.

Calls multi_heap_info on all heaps which share the given capabilities, and prints a two-line summary for each, then a total summary.

Parameters
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

bool **heap_caps_check_integrity_all**(bool *print_errors*)
Check integrity of all heap memory in the system.

Calls multi_heap_check on all heaps. Optionally print errors if heaps are corrupt.

Calling this function is equivalent to calling heap_caps_check_integrity with the caps argument set to MALLOC_CAP_INVALID.

Return True if all heaps are valid, False if at least one heap is corrupt.
Parameters
- `print_errors`: Print specific errors if heap corruption is found.

bool **heap_caps_check_integrity**(uint32_t *caps*, bool *print_errors*)
Check integrity of all heaps with the given capabilities.

Calls multi_heap_check on all heaps which share the given capabilities. Optionally print errors if the heaps are corrupt.

See also heap_caps_check_integrity_all to check all heap memory in the system and heap_caps_check_integrity_addr to check memory around a single address.

Return True if all heaps are valid, False if at least one heap is corrupt.
Parameters
- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory
- `print_errors`: Print specific errors if heap corruption is found.

bool **heap_caps_check_integrity_addr**(intptr_t *addr*, bool *print_errors*)
Check integrity of heap memory around a given address.

This function can be used to check the integrity of a single region of heap memory, which contains the given address.

This can be useful if debugging heap integrity for corruption at a known address, as it has a lower overhead than checking all heap regions. Note that if the corrupt address moves around between runs (due to timing or other factors) then this approach won't work and you should call heap_caps_check_integrity or heap_caps_check_integrity_all instead.

Note The entire heap region around the address is checked, not only the adjacent heap blocks.
Return True if the heap containing the specified address is valid, False if at least one heap is corrupt or the address doesn't belong to a heap region.
Parameters
- `addr`: Address in memory. Check for corruption in region containing this address.
- `print_errors`: Print specific errors if heap corruption is found.

void **heap_caps_malloc_extmem_enable**(size_t *limit*)
Enable malloc() in external memory and set limit below which malloc() attempts are placed in internal memory.

When external memory is in use, the allocation strategy is to initially try to satisfy smaller allocation requests with internal memory and larger requests with external memory. This sets the limit between the two, as well as generally enabling allocation in external memory.

Parameters

- `limit`: Limit, in bytes.

void ***heap_caps_malloc_prefer**(size_t *size*, size_t *num*, ...)

Allocate a chunk of memory as preference in decreasing order.

Attention The variable parameters are bitwise OR of MALLOC_CAP_* flags indicating the type of memory. This API prefers to allocate memory with the first parameter. If failed, allocate memory with the next parameter. It will try in this order until allocating a chunk of memory successfully or fail to allocate memories with any of the parameters.

Return A pointer to the memory allocated on success, NULL on failure

Parameters

- `size`: Size, in bytes, of the amount of memory to allocate
- `num`: Number of variable paramters

void ***heap_caps_realloc_prefer**(void **ptr*, size_t *size*, size_t *num*, ...)

Allocate a chunk of memory as preference in decreasing order.

Return Pointer to a new buffer of size ‘size’, or NULL if allocation failed.

Parameters

- `ptr`: Pointer to previously allocated memory, or NULL for a new allocation.
- `size`: Size of the new buffer requested, or 0 to free the buffer.
- `num`: Number of variable paramters

void ***heap_caps_calloc_prefer**(size_t *n*, size_t *size*, size_t *num*, ...)

Allocate a chunk of memory as preference in decreasing order.

Return A pointer to the memory allocated on success, NULL on failure

Parameters

- `n`: Number of continuing chunks of memory to allocate
- `size`: Size, in bytes, of a chunk of memory to allocate
- `num`: Number of variable paramters

void **heap_caps_dump**(uint32_t *caps*)

Dump the full structure of all heaps with matching capabilities.

Prints a large amount of output to serial (because of locking limitations, the output bypasses stdout/stderr). For each (variable sized) block in each matching heap, the following output is printed on a single line:

- Block address (the data buffer returned by malloc is 4 bytes after this if heap debugging is set to Basic, or 8 bytes otherwise).
- Data size (the data size may be larger than the size requested by malloc, either due to heap fragmentation or because of heap debugging level).
- Address of next block in the heap.
- If the block is free, the address of the next free block is also printed.

Parameters

- `caps`: Bitwise OR of MALLOC_CAP_* flags indicating the type of memory

void **heap_caps_dump_all**(void)

Dump the full structure of all heaps.

Covers all registered heaps. Prints a large amount of output to serial.

Output is the same as for heap_caps_dump.

size_t **heap_caps_get_allocated_size**(void **ptr*)

Return the size that a particular pointer was allocated with.

Note The app will crash with an assertion failure if the pointer is not valid.

Return Size of the memory allocated at this block.

Parameters

- `ptr`: Pointer to currently allocated heap memory. Must be a pointer value previously returned by heap_caps_malloc,malloc,calloc, etc. and not yet freed.

Macros

MALLOC_CAP_EXEC

Flags to indicate the capabilities of the various memory systems.

Memory must be able to run executable code

MALLOC_CAP_32BIT

Memory must allow for aligned 32-bit data accesses.

MALLOC_CAP_8BIT

Memory must allow for 8/16/…-bit data accesses.

MALLOC_CAP_DMA

Memory must be able to accessed by DMA.

MALLOC_CAP_PID2

Memory must be mapped to PID2 memory space (PIDs are not currently used)

MALLOC_CAP_PID3

Memory must be mapped to PID3 memory space (PIDs are not currently used)

MALLOC_CAP_PID4

Memory must be mapped to PID4 memory space (PIDs are not currently used)

MALLOC_CAP_PID5

Memory must be mapped to PID5 memory space (PIDs are not currently used)

MALLOC_CAP_PID6

Memory must be mapped to PID6 memory space (PIDs are not currently used)

MALLOC_CAP_PID7

Memory must be mapped to PID7 memory space (PIDs are not currently used)

MALLOC_CAP_SPIRAM

Memory must be in SPI RAM.

MALLOC_CAP_INTERNAL

Memory must be internal; specifically it should not disappear when flash/spiram cache is switched off.

MALLOC_CAP_DEFAULT

Memory can be returned in a non-capability-specific memory allocation (e.g. malloc(), calloc()) call.

MALLOC_CAP_IRAM_8BIT

Memory must be in IRAM and allow unaligned access.

MALLOC_CAP_INVALID

Memory can't be used / list end marker.

Type Definitions

typedef void (***esp_alloc_failed_hook_t**)(size_t size, uint32_t caps, **const** char *function_name)

callback called when a allocation operation fails, if registered

Parameters

- `size`: in bytes of failed allocation
- `caps`: capabillites requested of failed allocation
- `function_name`: function which generated the failure

Thread Safety Heap functions are thread safe, meaning they can be called from different tasks simultaneously without any limitations.

It is technically possible to call `malloc`, `free`, and related functions from interrupt handler (ISR) context. However this is not recommended, as heap function calls may delay other interrupts. It is strongly recommended to refactor applications so that any buffers used by an ISR are pre-allocated outside of the ISR. Support for calling heap functions from ISRs may be removed in a future update.

Heap Tracing & Debugging

The following features are documented on the *Heap Memory Debugging* page:

- *Heap Information* (free space, etc.)
- *Heap Corruption Detection*
- *Heap Tracing* (memory leak detection, monitoring, etc.)

API Reference - Initialisation

Header File

- heap/include/esp_heap_caps_init.h

Functions

void **heap_caps_init**(void)

Initialize the capability-aware heap allocator.

This is called once in the IDF startup code. Do not call it at other times.

void **heap_caps_enable_nonos_stack_heaps**(void)

Enable heap(s) in memory regions where the startup stacks are located.

On startup, the pro/app CPUs have a certain memory region they use as stack, so we cannot do allocations in the regions these stack frames are. When FreeRTOS is completely started, they do not use that memory anymore and heap(s) there can be enabled.

esp_err_t **heap_caps_add_region**(intptr_t *start*, intptr_t *end*)

Add a region of memory to the collection of heaps at runtime.

Most memory regions are defined in soc_memory_layout.c for the SoC, and are registered via heap_caps_init(). Some regions can't be used immediately and are later enabled via heap_caps_enable_nonos_stack_heaps().

Call this function to add a region of memory to the heap at some later time.

This function does not consider any of the "reserved" regions or other data in soc_memory_layout, caller needs to consider this themselves.

All memory within the region specified by start & end parameters must be otherwise unused.

The capabilities of the newly registered memory will be determined by the start address, as looked up in the regions specified in soc_memory_layout.c.

Use heap_caps_add_region_with_caps() to register a region with custom capabilities.

Return ESP_OK on success, ESP_ERR_INVALID_ARG if a parameter is invalid, ESP_ERR_NOT_FOUND if the specified start address doesn't reside in a known region, or any error returned by heap_caps_add_region_with_caps().

Parameters

- `start`: Start address of new region.
- `end`: End address of new region.

esp_err_t **heap_caps_add_region_with_caps**(**const** uint32_t *caps*[], intptr_t *start*, intptr_t *end*)

Add a region of memory to the collection of heaps at runtime, with custom capabilities.

Similar to heap_caps_add_region(), only custom memory capabilities are specified by the caller.

Return

- ESP_OK on success

- ESP_ERR_INVALID_ARG if a parameter is invalid
- ESP_ERR_NO_MEM if no memory to register new heap.
- ESP_ERR_INVALID_SIZE if the memory region is too small to fit a heap
- ESP_FAIL if region overlaps the start and/or end of an existing region

Parameters

- `caps`: Ordered array of capability masks for the new region, in order of priority. Must have length SOC_MEMORY_TYPE_NO_PRIOS. Does not need to remain valid after the call returns.
- `start`: Start address of new region.
- `end`: End address of new region.

Implementation Notes

Knowledge about the regions of memory in the chip comes from the "soc" component, which contains memory layout information for the chip, and the different capabilities of each region. Each region's capabilities are prioritised, so that (for example) dedicated DRAM and IRAM regions will be used for allocations ahead of the more versatile D/IRAM regions.

Each contiguous region of memory contains its own memory heap. The heaps are created using the multi_heap functionality. multi_heap allows any contiguous region of memory to be used as a heap.

The heap capabilities allocator uses knowledge of the memory regions to initialize each individual heap. Allocation functions in the heap capabilities API will find the most appropriate heap for the allocation (based on desired capabilities, available space, and preferences for each region's use) and then calling *`multi_heap_malloc()`* or `multi_heap_calloc()` for the heap situated in that particular region.

Calling `free()` involves finding the particular heap corresponding to the freed address, and then calling *`multi_heap_free()`* on that particular multi_heap instance.

API Reference - Multi Heap API

(Note: The multi heap API is used internally by the heap capabilities allocator. Most IDF programs will never need to call this API directly.)

Header File

- heap/include/multi_heap.h

Functions

void ***multi_heap_aligned_alloc**(*multi_heap_handle_t heap*, size_t *size*, size_t *alignment*)

allocate a chunk of memory with specific alignment

Return pointer to the memory allocated, NULL on failure

Parameters

- `heap`: Handle to a registered heap.
- `size`: size in bytes of memory chunk
- `alignment`: how the memory must be aligned

void ***multi_heap_malloc**(*multi_heap_handle_t heap*, size_t *size*)

malloc() a buffer in a given heap

Semantics are the same as standard malloc(), only the returned buffer will be allocated in the specified heap.

Return Pointer to new memory, or NULL if allocation fails.

Parameters

- `heap`: Handle to a registered heap.
- `size`: Size of desired buffer.

void **multi_heap_aligned_free**(*multi_heap_handle_t heap*, void **p*)

free() a buffer aligned in a given heap.

Parameters

- heap: Handle to a registered heap.
- p: NULL, or a pointer previously returned from multi_heap_aligned_alloc() for the same heap.

void **multi_heap_free**(*multi_heap_handle_t heap*, void *p)

free() a buffer in a given heap.

Semantics are the same as standard free(), only the argument 'p' must be NULL or have been allocated in the specified heap.

Parameters

- heap: Handle to a registered heap.
- p: NULL, or a pointer previously returned from multi_heap_malloc() or multi_heap_realloc() for the same heap.

void ***multi_heap_realloc**(*multi_heap_handle_t heap*, void *p, size_t *size*)

realloc() a buffer in a given heap.

Semantics are the same as standard realloc(), only the argument 'p' must be NULL or have been allocated in the specified heap.

Return New buffer of 'size' containing contents of 'p', or NULL if reallocation failed.

Parameters

- heap: Handle to a registered heap.
- p: NULL, or a pointer previously returned from multi_heap_malloc() or multi_heap_realloc() for the same heap.
- size: Desired new size for buffer.

size_t **multi_heap_get_allocated_size**(*multi_heap_handle_t heap*, void *p)

Return the size that a particular pointer was allocated with.

Return Size of the memory allocated at this block. May be more than the original size argument, due to padding and minimum block sizes.

Parameters

- heap: Handle to a registered heap.
- p: Pointer, must have been previously returned from multi_heap_malloc() or multi_heap_realloc() for the same heap.

multi_heap_handle_t **multi_heap_register**(void *start, size_t *size*)

Register a new heap for use.

This function initialises a heap at the specified address, and returns a handle for future heap operations.

There is no equivalent function for deregistering a heap - if all blocks in the heap are free, you can immediately start using the memory for other purposes.

Return Handle of a new heap ready for use, or NULL if the heap region was too small to be initialised.

Parameters

- start: Start address of the memory to use for a new heap.
- size: Size (in bytes) of the new heap.

void **multi_heap_set_lock**(*multi_heap_handle_t heap*, void *lock)

Associate a private lock pointer with a heap.

The lock argument is supplied to the MULTI_HEAP_LOCK() and MULTI_HEAP_UNLOCK() macros, defined in multi_heap_platform.h.

The lock in question must be recursive.

When the heap is first registered, the associated lock is NULL.

Parameters

- heap: Handle to a registered heap.
- lock: Optional pointer to a locking structure to associate with this heap.

void **multi_heap_dump**(*multi_heap_handle_t heap*)

Dump heap information to stdout.

For debugging purposes, this function dumps information about every block in the heap to stdout.

Parameters
- `heap`: Handle to a registered heap.

bool **multi_heap_check**(*multi_heap_handle_t heap*, bool *print_errors*)
Check heap integrity.

Walks the heap and checks all heap data structures are valid. If any errors are detected, an error-specific message can be optionally printed to stderr. Print behaviour can be overriden at compile time by defining MULTI_CHECK_FAIL_PRINTF in multi_heap_platform.h.

Return true if heap is valid, false otherwise.
Parameters
- `heap`: Handle to a registered heap.
- `print_errors`: If true, errors will be printed to stderr.

size_t **multi_heap_free_size**(*multi_heap_handle_t heap*)
Return free heap size.

Returns the number of bytes available in the heap.

Equivalent to the total_free_bytes member returned by multi_heap_get_heap_info().

Note that the heap may be fragmented, so the actual maximum size for a single malloc() may be lower. To know this size, see the largest_free_block member returned by multi_heap_get_heap_info().

Return Number of free bytes.
Parameters
- `heap`: Handle to a registered heap.

size_t **multi_heap_minimum_free_size**(*multi_heap_handle_t heap*)
Return the lifetime minimum free heap size.

Equivalent to the minimum_free_bytes member returned by multi_heap_get_info().

Returns the lifetime "low water mark" of possible values returned from multi_free_heap_size(), for the specified heap.

Return Number of free bytes.
Parameters
- `heap`: Handle to a registered heap.

void **multi_heap_get_info**(*multi_heap_handle_t heap*, *multi_heap_info_t *info*)
Return metadata about a given heap.

Fills a *multi_heap_info_t* structure with information about the specified heap.

Parameters
- `heap`: Handle to a registered heap.
- `info`: Pointer to a structure to fill with heap metadata.

Structures

struct multi_heap_info_t
Structure to access heap metadata via multi_heap_get_info.

Public Members

size_t **total_free_bytes**
Total free bytes in the heap. Equivalent to multi_free_heap_size().

size_t **total_allocated_bytes**
Total bytes allocated to data in the heap.

size_t **largest_free_block**
Size of largest free block in the heap. This is the largest malloc-able size.

size_t **minimum_free_bytes**
Lifetime minimum free heap size. Equivalent to multi_minimum_free_heap_size().

size_t **allocated_blocks**
Number of (variable size) blocks allocated in the heap.

size_t **free_blocks**
Number of (variable size) free blocks in the heap.

size_t **total_blocks**
Total number of (variable size) blocks in the heap.

Type Definitions

typedef struct multi_heap_info ***multi_heap_handle_t**
Opaque handle to a registered heap.

2.6.12 Heap Memory Debugging

Overview

ESP-IDF integrates tools for requesting *heap information*, *detecting heap corruption*, and *tracing memory leaks*. These can help track down memory-related bugs.

For general information about the heap memory allocator, see the *Heap Memory Allocation* page.

Heap Information

To obtain information about the state of the heap:

- `xPortGetFreeHeapSize()` is a FreeRTOS function which returns the number of free bytes in the (data memory) heap. This is equivalent to calling `heap_caps_get_free_size(MALLOC_CAP_8BIT)`.
- *`heap_caps_get_free_size()`* can also be used to return the current free memory for different memory capabilities.
- *`heap_caps_get_largest_free_block()`* can be used to return the largest free block in the heap. This is the largest single allocation which is currently possible. Tracking this value and comparing to total free heap allows you to detect heap fragmentation.
- `xPortGetMinimumEverFreeHeapSize()` and the related *`heap_caps_get_minimum_free_size()`* can be used to track the heap "low water mark" since boot.
- *`heap_caps_get_info()`* returns a *`multi_heap_info_t`* structure which contains the information from the above functions, plus some additional heap-specific data (number of allocations, etc.).
- *`heap_caps_print_heap_info()`* prints a summary to stdout of the information returned by *`heap_caps_get_info()`*.
- *`heap_caps_dump()`* and *`heap_caps_dump_all()`* will output detailed information about the structure of each block in the heap. Note that this can be large amount of output.

Heap Corruption Detection

Heap corruption detection allows you to detect various types of heap memory errors:

- Out of bounds writes & buffer overflow.
- Writes to freed memory.
- Reads from freed or uninitialized memory,

Assertions The heap implementation (`multi_heap.c`, etc.) includes a lot of assertions which will fail if the heap memory is corrupted. To detect heap corruption most effectively, ensure that assertions are enabled in the project configuration menu under `Compiler options` -> *CONFIG_COMPILER_OPTIMIZATION_ASSERTION_LEVEL*.

If a heap integrity assertion fails, a line will be printed like `CORRUPT HEAP: multi_heap.c:225 detected at 0x3ffbb71c`. The memory address which is printed is the address of the heap structure which has corrupt content.

It's also possible to manually check heap integrity by calling *heap_caps_check_integrity_all()* or related functions. This function checks all of requested heap memory for integrity, and can be used even if assertions are disabled. If the integrity check prints an error, it will also contain the address(es) of corrupt heap structures.

Memory Allocation Failed Hook Users can use *heap_caps_register_failed_alloc_callback()* to register a callback that will be invoked every time a allocation operation fails.

Additionaly user can enable a generation of a system abort if allocation operation fails by following the steps below: - In the project configuration menu, navigate to `Component config -> Heap Memory Debugging` and select `Abort if memory allocation fails` option (see *CONFIG_HEAP_ABORT_WHEN_ALLOCATION_FAILS*).

The example below show how to register a allocation failure callback:

```
#include "esp_heap_caps.h"

void heap_caps_alloc_failed_hook(size_t requested_size, uint32_t caps, const char␣
↪*function_name)
{
  printf("%s was called but failed to allocate %d bytes with 0x%X capabilities. \n
↪",function_name, requested_size, caps);
}

void app_main()
{
    ...
    esp_err_t error = heap_caps_register_failed_alloc_callback(heap_caps_alloc_
↪failed_hook);
    ...
    void *ptr = heap_caps_malloc(allocation_size, MALLOC_CAP_DEFAULT);
    ...
}
```

Finding Heap Corruption Memory corruption can be one of the hardest classes of bugs to find and fix, as one area of memory can be corrupted from a totally different place. Some tips:

- A crash with a `CORRUPT HEAP:` message will usually include a stack trace, but this stack trace is rarely useful. The crash is the symptom of memory corruption when the system realises the heap is corrupt, but usually the corruption happened elsewhere and earlier in time.
- Increasing the Heap memory debugging *Configuration* level to "Light impact" or "Comprehensive" can give you a more accurate message with the first corrupt memory address.
- Adding regular calls to *heap_caps_check_integrity_all()* or *heap_caps_check_integrity_addr()* in your code will help you pin down the exact time that the corruption happened. You can move these checks around to "close in on" the section of code that corrupted the heap.
- Based on the memory address which is being corrupted, you can use *JTAG debugging* to set a watchpoint on this address and have the CPU halt when it is written to.
- If you don't have JTAG, but you do know roughly when the corruption happens, then you can set a watchpoint in software just beforehand via `esp_set_watchpoint()`. A fatal exception will occur when the watchpoint triggers. For example `esp_set_watchpoint(0, (void *)addr, 4, ESP_WATCHPOINT_STORE`. Note that watchpoints are per-CPU and are set on the current running CPU only, so if you don't know which CPU is corrupting memory then you will need to call this function on both CPUs.
- For buffer overflows, *heap tracing* in `HEAP_TRACE_ALL` mode lets you see which callers are allocating which addresses from the heap. See *Heap Tracing To Find Heap Corruption* for more details. If you can find the

function which allocates memory with an address immediately before the address which is corrupted, this will probably be the function which overflows the buffer.
- Calling *heap_caps_dump()* or *heap_caps_dump_all()* can give an indication of what heap blocks are surrounding the corrupted region and may have overflowed/underflowed/etc.

Configuration Temporarily increasing the heap corruption detection level can give more detailed information about heap corruption errors.

In the project configuration menu, under `Component config` there is a menu `Heap memory debugging`. The setting *CONFIG_HEAP_CORRUPTION_DETECTION* can be set to one of three levels:

Basic (no poisoning) This is the default level. No special heap corruption features are enabled, but provided assertions are enabled (the default configuration) then a heap corruption error will be printed if any of the heap's internal data structures appear overwritten or corrupted. This usually indicates a buffer overrun or out of bounds write.

If assertions are enabled, an assertion will also trigger if a double-free occurs (the same memory is freed twice).

Calling *heap_caps_check_integrity()* in Basic mode will check the integrity of all heap structures, and print errors if any appear to be corrupted.

Light Impact At this level, heap memory is additionally "poisoned" with head and tail "canary bytes" before and after each block which is allocated. If an application writes outside the bounds of allocated buffers, the canary bytes will be corrupted and the integrity check will fail.

The head canary word is 0xABBA1234 (3412BAAB in byte order), and the tail canary word is 0xBAAD5678 (7856ADBA in byte order).

"Basic" heap corruption checks can also detect most out of bounds writes, but this setting is more precise as even a single byte overrun can be detected. With Basic heap checks, the number of overrun bytes before a failure is detected will depend on the properties of the heap.

Enabling "Light Impact" checking increases memory usage, each individual allocation will use 9 to 12 additional bytes of memory (depending on alignment).

Each time `free()` is called in Light Impact mode, the head and tail canary bytes of the buffer being freed are checked against the expected values.

When *heap_caps_check_integrity()* is called, all allocated blocks of heap memory have their canary bytes checked against the expected values.

In both cases, the check is that the first 4 bytes of an allocated block (before the buffer returned to the user) should be the word 0xABBA1234. Then the last 4 bytes of the allocated block (after the buffer returned to the user) should be the word 0xBAAD5678.

Different values usually indicate buffer underrun or overrun, respectively.

Comprehensive This level incorporates the "light impact" detection features plus additional checks for uninitialised-access and use-after-free bugs. In this mode, all freshly allocated memory is filled with the pattern 0xCE, and all freed memory is filled with the pattern 0xFE.

Enabling "Comprehensive" detection has a substantial runtime performance impact (as all memory needs to be set to the allocation patterns each time a malloc/free completes, and the memory also needs to be checked each time.) However it allows easier detection of memory corruption bugs which are much more subtle to find otherwise. It is recommended to only enable this mode when debugging, not in production.

Crashes in Comprehensive Mode If an application crashes reading/writing an address related to 0xCECECECE in Comprehensive mode, this indicates it has read uninitialized memory. The application should be changed to either use calloc() (which zeroes memory), or initialize the memory before using it. The value 0xCECECECE may also be

seen in stack-allocated automatic variables, because in IDF most task stacks are originally allocated from the heap and in C stack memory is uninitialized by default.

If an application crashes and the exception register dump indicates that some addresses or values were 0xFEFEFEFE, this indicates it is reading heap memory after it has been freed (a "use after free bug".) The application should be changed to not access heap memory after it has been freed.

If a call to malloc() or realloc() causes a crash because it expected to find the pattern 0xFEFEFEFE in free memory and a different pattern was found, then this indicates the app has a use-after-free bug where it is writing to memory which has already been freed.

Manual Heap Checks in Comprehensive Mode Calls to *heap_caps_check_integrity()* may print errors relating to 0xFEFEFEFE, 0xABBA1234 or 0xBAAD5678. In each case the checker is expecting to find a given pattern, and will error out if this is not found:

- For free heap blocks, the checker expects to find all bytes set to 0xFE. Any other values indicate a use-after-free bug where free memory has been incorrectly overwritten.
- For allocated heap blocks, the behaviour is the same as for *Light Impact* mode. The canary bytes 0xABBA1234 and 0xBAAD5678 are checked at the head and tail of each allocated buffer, and any variation indicates a buffer overrun/underrun.

Heap Task Tracking

Heap Task Tracking can be used to get per task info for heap memory allocation. Application has to specify the heap capabilities for which the heap allocation is to be tracked.

Example code is provided in system/heap_task_tracking

Heap Tracing

Heap Tracing allows tracing of code which allocates/frees memory. Two tracing modes are supported:

- Standalone. In this mode trace data are kept on-board, so the size of gathered information is limited by the buffer assigned for that purposes. Analysis is done by the on-board code. There are a couple of APIs available for accessing and dumping collected info.
- Host-based. This mode does not have the limitation of the standalone mode, because trace data are sent to the host over JTAG connection using app_trace library. Later on they can be analysed using special tools.

Heap tracing can perform two functions:

- Leak checking: find memory which is allocated and never freed.
- Heap use analysis: show all functions that are allocating/freeing memory while the trace is running.

How To Diagnose Memory Leaks If you suspect a memory leak, the first step is to figure out which part of the program is leaking memory. Use the `xPortGetFreeHeapSize()`, *heap_caps_get_free_size()*, or *related functions* to track memory use over the life of the application. Try to narrow the leak down to a single function or sequence of functions where free memory always decreases and never recovers.

Standalone Mode Once you've identified the code which you think is leaking:

- In the project configuration menu, navigate to `Component settings` -> `Heap Memory Debugging` -> `Heap tracing` and select `Standalone` option (see *CONFIG_HEAP_TRACING_DEST*).
- Call the function *heap_trace_init_standalone()* early in the program, to register a buffer which can be used to record the memory trace.
- Call the function *heap_trace_start()* to begin recording all mallocs/frees in the system. Call this immediately before the piece of code which you suspect is leaking memory.
- Call the function *heap_trace_stop()* to stop the trace once the suspect piece of code has finished executing.

- Call the function *heap_trace_dump()* to dump the results of the heap trace.

An example:

```
#include "esp_heap_trace.h"

#define NUM_RECORDS 100
static heap_trace_record_t trace_record[NUM_RECORDS]; // This buffer must be in␣
↪internal RAM

...

void app_main()
{
    ...
    ESP_ERROR_CHECK( heap_trace_init_standalone(trace_record, NUM_RECORDS) );
    ...
}

void some_function()
{
    ESP_ERROR_CHECK( heap_trace_start(HEAP_TRACE_LEAKS) );

    do_something_you_suspect_is_leaking();

    ESP_ERROR_CHECK( heap_trace_stop() );
    heap_trace_dump();
    ...
}
```

The output from the heap trace will look something like this:

```
2 allocations trace (100 entry buffer)
32 bytes (@ 0x3ffaf214) allocated CPU 0 ccount 0x2e9b7384 caller␣
↪0x400d276d:0x400d27c1
0x400d276d: leak_some_memory at /path/to/idf/examples/get-started/blink/main/./
↪blink.c:27

0x400d27c1: blink_task at /path/to/idf/examples/get-started/blink/main/./blink.c:52

8 bytes (@ 0x3ffaf804) allocated CPU 0 ccount 0x2e9b79c0 caller␣
↪0x400d2776:0x400d27c1
0x400d2776: leak_some_memory at /path/to/idf/examples/get-started/blink/main/./
↪blink.c:29

0x400d27c1: blink_task at /path/to/idf/examples/get-started/blink/main/./blink.c:52

40 bytes 'leaked' in trace (2 allocations)
total allocations 2 total frees 0
```

(Above example output is using *IDF Monitor* to automatically decode PC addresses to their source files & line number.)

The first line indicates how many allocation entries are in the buffer, compared to its total size.

In `HEAP_TRACE_LEAKS` mode, for each traced memory allocation which has not already been freed a line is printed with:

- `XX bytes` is number of bytes allocated
- `@ 0x...` is the heap address returned from malloc/calloc.
- `CPU x` is the CPU (0 or 1) running when the allocation was made.
- `ccount 0x...` is the CCOUNT (CPU cycle count) register value when the allocation was mode. Is different for CPU 0 vs CPU 1.
- `caller 0x...` gives the call stack of the call to malloc()/free(), as a list of PC addresses. These can be decoded to source files and line numbers, as shown above.

The depth of the call stack recorded for each trace entry can be configured in the project configuration menu, under `Heap Memory Debugging` -> `Enable heap tracing` -> `Heap tracing stack depth`. Up to 10 stack frames can be recorded for each allocation (the default is 2). Each additional stack frame increases the memory usage of each `heap_trace_record_t` record by eight bytes.

Finally, the total number of 'leaked' bytes (bytes allocated but not freed while trace was running) is printed, and the total number of allocations this represents.

A warning will be printed if the trace buffer was not large enough to hold all the allocations which happened. If you see this warning, consider either shortening the tracing period or increasing the number of records in the trace buffer.

Host-Based Mode Once you've identified the code which you think is leaking:

- In the project configuration menu, navigate to `Component settings` -> `Heap Memory Debugging` -> *CONFIG_HEAP_TRACING_DEST* and select `Host-Based`.
- In the project configuration menu, navigate to `Component settings` -> `Application Level Tracing` -> *CONFIG_APPTRACE_DESTINATION* and select `Trace memory`.
- In the project configuration menu, navigate to `Component settings` -> `Application Level Tracing` -> `FreeRTOS SystemView Tracing` and enable *CONFIG_SYSVIEW_ENABLE*.
- Call the function *heap_trace_init_tohost()* early in the program, to initialize JTAG heap tracing module.
- Call the function *heap_trace_start()* to begin recording all mallocs/frees in the system. Call this immediately before the piece of code which you suspect is leaking memory. In host-based mode argument to this function is ignored and heap tracing module behaves like `HEAP_TRACE_ALL` was passed: all allocations and deallocations are sent to the host.
- Call the function *heap_trace_stop()* to stop the trace once the suspect piece of code has finished executing.

An example:

```
#include "esp_heap_trace.h"

...

void app_main()
{
    ...
    ESP_ERROR_CHECK( heap_trace_init_tohost() );
    ...
}

void some_function()
{
    ESP_ERROR_CHECK( heap_trace_start(HEAP_TRACE_LEAKS) );

    do_something_you_suspect_is_leaking();

    ESP_ERROR_CHECK( heap_trace_stop() );
    ...
}
```

To gather and analyse heap trace do the following on the host:

1. Build the program and download it to the target as described in *Getting Started Guide*.
2. Run OpenOCD (see *JTAG Debugging*).

注解: In order to use this feature you need OpenOCD version *v0.10.0-esp32-20181105* or later.

---

3. You can use GDB to start and/or stop tracing automatically. To do this you need to prepare special `gdbinit` file:

```
target remote :3333

mon reset halt
flushregs

tb heap_trace_start
commands
mon esp32 sysview start file:///tmp/heap.svdat
c
end

tb heap_trace_stop
commands
mon esp32 sysview stop
end

c
```

Using this file GDB will connect to the target, reset it, and start tracing when program hits breakpoint at *heap_trace_start()*. Trace data will be saved to `/tmp/heap_log.svdat`. Tracing will be stopped when program hits breakpoint at *heap_trace_stop()*.

4. Run GDB using the following command `xtensa-esp32s2-elf-gdb -x gdbinit </path/to/program/elf>`
5. Quit GDB when program stops at *heap_trace_stop()*. Trace data are saved in `/tmp/heap.svdat`
6. Run processing script `$IDF_PATH/tools/esp_app_trace/sysviewtrace_proc.py -p -b </path/to/program/elf> /tmp/heap_log.svdat`

The output from the heap trace will look something like this:

```
Parse trace from '/tmp/heap.svdat'...
Stop parsing trace. (Timeout 0.000000 sec while reading 1 bytes!)
Process events from '['/tmp/heap.svdat']'...
[0.002244575] HEAP: Allocated 1 bytes @ 0x3ffaffd8 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.002258425] HEAP: Allocated 2 bytes @ 0x3ffaffe0 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:48
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.002563725] HEAP: Freed bytes @ 0x3ffaffe0 from task "free" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:31 (discriminator 9)
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.002782950] HEAP: Freed bytes @ 0x3ffb40b8 from task "main" on core 0 by:
/home/user/projects/esp/esp-idf/components/freertos/tasks.c:4590
/home/user/projects/esp/esp-idf/components/freertos/tasks.c:4590

[0.002798700] HEAP: Freed bytes @ 0x3ffb50bc from task "main" on core 0 by:
/home/user/projects/esp/esp-idf/components/freertos/tasks.c:4590
/home/user/projects/esp/esp-idf/components/freertos/tasks.c:4590

[0.102436025] HEAP: Allocated 2 bytes @ 0x3ffaffe0 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.102449800] HEAP: Allocated 4 bytes @ 0x3ffaffe8 from task "alloc" on core 0 by:
```

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```
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:48
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.102666150] HEAP: Freed bytes @ 0x3ffaffe8 from task "free" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:31 (discriminator 9)
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.202436200] HEAP: Allocated 3 bytes @ 0x3ffaffe8 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.202451725] HEAP: Allocated 6 bytes @ 0x3ffafff0 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:48
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.202667075] HEAP: Freed bytes @ 0x3ffafff0 from task "free" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:31 (discriminator 9)
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.302436000] HEAP: Allocated 4 bytes @ 0x3ffafff0 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.302451475] HEAP: Allocated 8 bytes @ 0x3ffb40b8 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:48
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.302667500] HEAP: Freed bytes @ 0x3ffb40b8 from task "free" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:31 (discriminator 9)
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

Processing completed.
Processed 1019 events
=============== HEAP TRACE REPORT ===============
Processed 14 heap events.
[0.002244575] HEAP: Allocated 1 bytes @ 0x3ffaffd8 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.102436025] HEAP: Allocated 2 bytes @ 0x3ffaffe0 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.202436200] HEAP: Allocated 3 bytes @ 0x3ffaffe8 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

[0.302436000] HEAP: Allocated 4 bytes @ 0x3ffafff0 from task "alloc" on core 0 by:
/home/user/projects/esp/esp-idf/examples/system/sysview_tracing_heap_log/main/
↪sysview_heap_log.c:47
```

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```
/home/user/projects/esp/esp-idf/components/freertos/port.c:355 (discriminator 1)

Found 10 leaked bytes in 4 blocks.
```

Heap Tracing To Find Heap Corruption Heap tracing can also be used to help track down heap corruption. When a region in heap is corrupted, it may be from some other part of the program which allocated memory at a nearby address.

If you have some idea at what time the corruption occurred, enabling heap tracing in `HEAP_TRACE_ALL` mode allows you to record all of the functions which allocated memory, and the addresses of the allocations.

Using heap tracing in this way is very similar to memory leak detection as described above. For memory which is allocated and not freed, the output is the same. However, records will also be shown for memory which has been freed.

Performance Impact Enabling heap tracing in menuconfig increases the code size of your program, and has a very small negative impact on performance of heap allocation/free operations even when heap tracing is not running.

When heap tracing is running, heap allocation/free operations are substantially slower than when heap tracing is stopped. Increasing the depth of stack frames recorded for each allocation (see above) will also increase this performance impact.

False-Positive Memory Leaks Not everything printed by *`heap_trace_dump()`* is necessarily a memory leak. Among things which may show up here, but are not memory leaks:

- Any memory which is allocated after *`heap_trace_start()`* but then freed after *`heap_trace_stop()`* will appear in the leak dump.
- Allocations may be made by other tasks in the system. Depending on the timing of these tasks, it's quite possible this memory is freed after *`heap_trace_stop()`* is called.
- The first time a task uses stdio - for example, when it calls `printf()` - a lock (RTOS mutex semaphore) is allocated by the libc. This allocation lasts until the task is deleted.
- Certain uses of `printf()`, such as printing floating point numbers, will allocate some memory from the heap on demand. These allocations last until the task is deleted.
- The Bluetooth, WiFi, and TCP/IP libraries will allocate heap memory buffers to handle incoming or outgoing data. These memory buffers are usually short lived, but some may be shown in the heap leak trace if the data was received/transmitted by the lower levels of the network while the leak trace was running.
- TCP connections will continue to use some memory after they are closed, because of the `TIME_WAIT` state. After the `TIME_WAIT` period has completed, this memory will be freed.

One way to differentiate between "real" and "false positive" memory leaks is to call the suspect code multiple times while tracing is running, and look for patterns (multiple matching allocations) in the heap trace output.

API Reference - Heap Tracing

Header File

- heap/include/esp_heap_trace.h

Functions

esp_err_t **heap_trace_init_standalone** (*heap_trace_record_t* **record_buffer*, size_t *num_records*)

Initialise heap tracing in standalone mode.

This function must be called before any other heap tracing functions.

To disable heap tracing and allow the buffer to be freed, stop tracing and then call heap_trace_init_standalone(NULL, 0);

Return

- ESP_ERR_NOT_SUPPORTED Project was compiled without heap tracing enabled in menuconfig.
- ESP_ERR_INVALID_STATE Heap tracing is currently in progress.
- ESP_OK Heap tracing initialised successfully.

Parameters

- `record_buffer`: Provide a buffer to use for heap trace data. Must remain valid any time heap tracing is enabled, meaning it must be allocated from internal memory not in PSRAM.
- `num_records`: Size of the heap trace buffer, as number of record structures.

esp_err_t **heap_trace_init_tohost** (void)

Initialise heap tracing in host-based mode.

This function must be called before any other heap tracing functions.

Return

- ESP_ERR_INVALID_STATE Heap tracing is currently in progress.
- ESP_OK Heap tracing initialised successfully.

esp_err_t **heap_trace_start** (*heap_trace_mode_t mode*)

Start heap tracing. All heap allocations & frees will be traced, until heap_trace_stop() is called.

Note heap_trace_init_standalone() must be called to provide a valid buffer, before this function is called.

Note Calling this function while heap tracing is running will reset the heap trace state and continue tracing.

Return

- ESP_ERR_NOT_SUPPORTED Project was compiled without heap tracing enabled in menuconfig.
- ESP_ERR_INVALID_STATE A non-zero-length buffer has not been set via heap_trace_init_standalone().
- ESP_OK Tracing is started.

Parameters

- `mode`: Mode for tracing.
  - HEAP_TRACE_ALL means all heap allocations and frees are traced.
  - HEAP_TRACE_LEAKS means only suspected memory leaks are traced. (When memory is freed, the record is removed from the trace buffer.)

esp_err_t **heap_trace_stop** (void)

Stop heap tracing.

Return

- ESP_ERR_NOT_SUPPORTED Project was compiled without heap tracing enabled in menuconfig.
- ESP_ERR_INVALID_STATE Heap tracing was not in progress.
- ESP_OK Heap tracing stopped..

esp_err_t **heap_trace_resume** (void)

Resume heap tracing which was previously stopped.

Unlike heap_trace_start(), this function does not clear the buffer of any pre-existing trace records.

The heap trace mode is the same as when heap_trace_start() was last called (or HEAP_TRACE_ALL if heap_trace_start() was never called).

Return

- ESP_ERR_NOT_SUPPORTED Project was compiled without heap tracing enabled in menuconfig.
- ESP_ERR_INVALID_STATE Heap tracing was already started.
- ESP_OK Heap tracing resumed.

size_t **heap_trace_get_count** (void)

Return number of records in the heap trace buffer.

It is safe to call this function while heap tracing is running.

esp_err_t **heap_trace_get** (size_t *index*, *heap_trace_record_t* **record*)

Return a raw record from the heap trace buffer.

Note It is safe to call this function while heap tracing is running, however in HEAP_TRACE_LEAK mode record indexing may skip entries unless heap tracing is stopped first.

Return

- ESP_ERR_NOT_SUPPORTED Project was compiled without heap tracing enabled in menuconfig.
- ESP_ERR_INVALID_STATE Heap tracing was not initialised.
- ESP_ERR_INVALID_ARG Index is out of bounds for current heap trace record count.
- ESP_OK Record returned successfully.

Parameters

- `index`: Index (zero-based) of the record to return.
- `[out] record`: Record where the heap trace record will be copied.

void **heap_trace_dump**(void)

Dump heap trace record data to stdout.

Note It is safe to call this function while heap tracing is running, however in HEAP_TRACE_LEAK mode the dump may skip entries unless heap tracing is stopped first.

Structures

struct heap_trace_record_t

Trace record data type. Stores information about an allocated region of memory.

Public Members

uint32_t **ccount**

CCOUNT of the CPU when the allocation was made. LSB (bit value 1) is the CPU number (0 or 1).

void ***address**

Address which was allocated.

size_t **size**

Size of the allocation.

void ***alloced_by**[**CONFIG_HEAP_TRACING_STACK_DEPTH**]

Call stack of the caller which allocated the memory.

void ***freed_by**[**CONFIG_HEAP_TRACING_STACK_DEPTH**]

Call stack of the caller which freed the memory (all zero if not freed.)

Macros

CONFIG_HEAP_TRACING_STACK_DEPTH

Enumerations

enum heap_trace_mode_t

Values:

HEAP_TRACE_ALL

HEAP_TRACE_LEAKS

2.6.13 High Resolution Timer

Overview

Although FreeRTOS provides software timers, these timers have a few limitations:

- Maximum resolution is equal to RTOS tick period
- Timer callbacks are dispatched from a low-priority task

Hardware timers are free from both of the limitations, but often they are less convenient to use. For example, application components may need timer events to fire at certain times in the future, but the hardware timer only contains one "compare" value used for interrupt generation. This means that some facility needs to be built on top

of the hardware timer to manage the list of pending events can dispatch the callbacks for these events as corresponding hardware interrupts happen.

esp_timer set of APIs provides one-shot and periodic timers, microsecond time resolution, and 64-bit range.

Internally, esp_timer uses a 64-bit hardware timer *CONFIG_ESP_TIMER_IMPL*:

- LAC timer (ESP32)
- (legacy) FRC2 timer (ESP32)
- SYSTIMER for (ESP32-S2)

Timer callbacks are dispatched from a high-priority esp_timer task. Because all the callbacks are dispatched from the same task, it is recommended to only do the minimal possible amount of work from the callback itself, posting an event to a lower priority task using a queue instead.

If other tasks with priority higher than esp_timer are running, callback dispatching will be delayed until esp_timer task has a chance to run. For example, this will happen if a SPI Flash operation is in progress.

Creating and starting a timer, and dispatching the callback takes some time. Therefore there is a lower limit to the timeout value of one-shot esp_timer. If *esp_timer_start_once()* is called with a timeout value less than 20us, the callback will be dispatched only after approximately 20us.

Periodic esp_timer also imposes a 50us restriction on the minimal timer period. Periodic software timers with period of less than 50us are not practical since they would consume most of the CPU time. Consider using dedicated hardware peripherals or DMA features if you find that a timer with small period is required.

Using esp_timer APIs

Single timer is represented by *esp_timer_handle_t* type. Timer has a callback function associated with it. This callback function is called from the esp_timer task each time the timer elapses.

- To create a timer, call *esp_timer_create()*.
- To delete the timer when it is no longer needed, call *esp_timer_delete()*.

The timer can be started in one-shot mode or in periodic mode.

- To start the timer in one-shot mode, call *esp_timer_start_once()*, passing the time interval after which the callback should be called. When the callback gets called, the timer is considered to be stopped.
- To start the timer in periodic mode, call *esp_timer_start_periodic()*, passing the period with which the callback should be called. The timer keeps running until *esp_timer_stop()* is called.

Note that the timer must not be running when *esp_timer_start_once()* or *esp_timer_start_periodic()* is called. To restart a running timer, call *esp_timer_stop()* first, then call one of the start functions.

Obtaining Current Time

esp_timer also provides a convenience function to obtain the time passed since start-up, with microsecond precision: *esp_timer_get_time()*. This function returns the number of microseconds since esp_timer was initialized, which usually happens shortly before app_main function is called.

Unlike *gettimeofday* function, values returned by *esp_timer_get_time()*:

- Start from zero after the chip wakes up from deep sleep
- Do not have timezone or DST adjustments applied

Application Example

The following example illustrates usage of esp_timer APIs: system/esp_timer.

API Reference

Header File

- esp_timer/include/esp_timer.h

Functions

esp_err_t **esp_timer_init**(void)

Initialize esp_timer library.

Note This function is called from startup code. Applications do not need to call this function before using other esp_timer APIs.

Return

- ESP_OK on success
- ESP_ERR_NO_MEM if allocation has failed
- ESP_ERR_INVALID_STATE if already initialized
- other errors from interrupt allocator

esp_err_t **esp_timer_deinit**(void)

De-initialize esp_timer library.

Note Normally this function should not be called from applications

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if not yet initialized

esp_err_t **esp_timer_create**(**const** *esp_timer_create_args_t* **create_args*, *esp_timer_handle_t* **out_handle*)

Create an esp_timer instance.

Note When done using the timer, delete it with esp_timer_delete function.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if some of the create_args are not valid
- ESP_ERR_INVALID_STATE if esp_timer library is not initialized yet
- ESP_ERR_NO_MEM if memory allocation fails

Parameters

- `create_args`: Pointer to a structure with timer creation arguments. Not saved by the library, can be allocated on the stack.
- `[out] out_handle`: Output, pointer to esp_timer_handle_t variable which will hold the created timer handle.

esp_err_t **esp_timer_start_once**(*esp_timer_handle_t timer*, uint64_t *timeout_us*)

Start one-shot timer.

Timer should not be running when this function is called.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the handle is invalid
- ESP_ERR_INVALID_STATE if the timer is already running

Parameters

- `timer`: timer handle created using esp_timer_create
- `timeout_us`: timer timeout, in microseconds relative to the current moment

esp_err_t **esp_timer_start_periodic**(*esp_timer_handle_t timer*, uint64_t *period*)

Start a periodic timer.

Timer should not be running when this function is called. This function will start the timer which will trigger every ‘period’ microseconds.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the handle is invalid

- ESP_ERR_INVALID_STATE if the timer is already running

Parameters

- `timer`: timer handle created using esp_timer_create
- `period`: timer period, in microseconds

esp_err_t **esp_timer_stop**(*esp_timer_handle_t timer*)

Stop the timer.

This function stops the timer previously started using esp_timer_start_once or esp_timer_start_periodic.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if the timer is not running

Parameters

- `timer`: timer handle created using esp_timer_create

esp_err_t **esp_timer_delete**(*esp_timer_handle_t timer*)

Delete an esp_timer instance.

The timer must be stopped before deleting. A one-shot timer which has expired does not need to be stopped.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if the timer is not running

Parameters

- `timer`: timer handle allocated using esp_timer_create

int64_t **esp_timer_get_time**(void)

Get time in microseconds since boot.

Return number of microseconds since esp_timer_init was called (this normally happens early during application startup).

int64_t **esp_timer_get_next_alarm**(void)

Get the timestamp when the next timeout is expected to occur.

Return Timestamp of the nearest timer event, in microseconds. The timebase is the same as for the values returned by esp_timer_get_time.

esp_err_t **esp_timer_dump**(FILE **stream*)

Dump the list of timers to a stream.

If CONFIG_ESP_TIMER_PROFILING option is enabled, this prints the list of all the existing timers. Otherwise, only the list active timers is printed.

The format is:

name period alarm times_armed times_triggered total_callback_run_time

where:

name —timer name (if CONFIG_ESP_TIMER_PROFILING is defined), or timer pointer period —period of timer, in microseconds, or 0 for one-shot timer alarm - time of the next alarm, in microseconds since boot, or 0 if the timer is not started

The following fields are printed if CONFIG_ESP_TIMER_PROFILING is defined:

times_armed —number of times the timer was armed via esp_timer_start_X times_triggered - number of times the callback was called total_callback_run_time - total time taken by callback to execute, across all calls

Return

- ESP_OK on success
- ESP_ERR_NO_MEM if can not allocate temporary buffer for the output

Parameters

- `stream`: stream (such as stdout) to dump the information to

Structures

struct esp_timer_create_args_t

Timer configuration passed to esp_timer_create.

Public Members

esp_timer_cb_t **callback**

Function to call when timer expires.

void ***arg**

Argument to pass to the callback.

esp_timer_dispatch_t **dispatch_method**

Call the callback from task or from ISR.

const char ***name**

Timer name, used in esp_timer_dump function.

Type Definitions

typedef struct esp_timer ***esp_timer_handle_t**

Opaque type representing a single esp_timer.

typedef void (***esp_timer_cb_t**) (void *arg)

Timer callback function type.

Parameters

- `arg`: pointer to opaque user-specific data

Enumerations

enum esp_timer_dispatch_t

Method for dispatching timer callback.

Values:

ESP_TIMER_TASK

Callback is called from timer task.

2.6.14 Call function with external stack

Overview

A given function can be executed with a user allocated stack space which is independent of current task stack, this mechanism can be used to save stack space wasted by tasks which call a common function with intensive stack usage such as *printf*. The given function can be called inside the shared stack space which is a callback function deferred by calling *esp_execute_shared_stack_function()*, passing that function as parameter

Usage

esp_execute_shared_stack_function() takes four arguments, a mutex object allocated by the caller, which is used to protect if the same function shares its allocated stack, a pointer to the top of stack used to that fuction, the size in bytes of stack and, a pointer to a user function where the shared stack space will reside, after calling the function, the user defined function will be deferred as a callback where functions can be called using the user allocated space without taking space from current task stack.

The usage may looks like the code below:

```
void external_stack_function(void)
{
    printf("Executing this printf from external stack! \n");
}

//Let's suppose we wanting to call printf using a separated stack space
//allowing app to reduce its stack size.
void app_main()
{
    //Allocate a stack buffer, from heap or as a static form:
    portSTACK_TYPE *shared_stack = malloc(8192 * sizeof(portSTACK_TYPE));
    assert(shared_stack != NULL);

    //Allocate a mutex to protect its usage:
    SemaphoreHandle_t printf_lock = xSemaphoreCreateMutex();
    assert(printf_lock != NULL);

    //Call the desired function using the macro helper:
    esp_execute_shared_stack_function(printf_lock,
                                      shared_stack,
                                      8192,
                                      external_stack_function);

    vSemaphoreDelete(printf_lock);
    free(shared_stack);
}
```

API Reference

Header File

- esp_common/include/esp_expression_with_stack.h

Functions

void **esp_execute_shared_stack_function**(*SemaphoreHandle_t* *lock*, void **stack*, size_t *stack_size*, *shared_stack_function* *function*)

Calls user defined shared stack space function.

Note if either lock, stack or stack size is invalid, the expression will be called using the current stack.

Parameters

- `lock`: Mutex object to protect in case of shared stack
- `stack`: Pointer to user alocated stack
- `stack_size`: Size of current stack in bytes
- `function`: pointer to the shared stack function to be executed

Macros

ESP_EXECUTE_EXPRESSION_WITH_STACK(lock, stack, stack_size, expression)

Type Definitions

typedef void (***shared_stack_function**)(void)

2.6.15 Interrupt allocation

Overview

The ESP32-S2 has one core, with 32 interrupts. Each interrupt has a certain priority level, most (but not all) interrupts are connected to the interrupt mux. Because there are more interrupt sources than interrupts, sometimes it makes

sense to share an interrupt in multiple drivers. The esp_intr_alloc abstraction exists to hide all these implementation details.

A driver can allocate an interrupt for a certain peripheral by calling esp_intr_alloc (or esp_intr_alloc_sintrstatus). It can use the flags passed to this function to set the type of interrupt allocated, specifying a specific level or trigger method. The interrupt allocation code will then find an applicable interrupt, use the interrupt mux to hook it up to the peripheral, and install the given interrupt handler and ISR to it.

This code has two different types of interrupts it handles differently: Shared interrupts and non-shared interrupts. The simplest of the two are non-shared interrupts: a separate interrupt is allocated per esp_intr_alloc call and this interrupt is solely used for the peripheral attached to it, with only one ISR that will get called. Shared interrupts can have multiple peripherals triggering it, with multiple ISRs being called when one of the peripherals attached signals an interrupt. Thus, ISRs that are intended for shared interrupts should check the interrupt status of the peripheral they service in order to see if any action is required.

Non-shared interrupts can be either level- or edge-triggered. Shared interrupts can only be level interrupts (because of the chance of missed interrupts when edge interrupts are used.) (The logic behind this: DevA and DevB share an int. DevB signals an int. Int line goes high. ISR handler calls code for DevA -> does nothing. ISR handler calls code for DevB, but while doing that, DevA signals an int. ISR DevB is done, clears int for DevB, exits interrupt code. Now an interrupt for DevA is still pending, but because the int line never went low (DevA kept it high even when the int for DevB was cleared) the interrupt is never serviced.)

Multicore issues

Peripherals that can generate interrupts can be divided in two types:

- External peripherals, within the ESP32-S2 but outside the Xtensa cores themselves. Most ESP32-S2 peripherals are of this type.
- Internal peripherals, part of the Xtensa CPU cores themselves.

Interrupt handling differs slightly between these two types of peripherals.

Internal peripheral interrupts Each Xtensa CPU core has its own set of six internal peripherals:

- Three timer comparators
- A performance monitor
- Two software interrupts.

Internal interrupt sources are defined in esp_intr_alloc.h as `ETS_INTERNAL_*_INTR_SOURCE`.

These peripherals can only be configured from the core they are associated with. When generating an interrupt, the interrupt they generate is hard-wired to their associated core; it’ s not possible to have e.g. an internal timer comparator of one core generate an interrupt on another core. That is why these sources can only be managed using a task running on that specific core. Internal interrupt sources are still allocatable using esp_intr_alloc as normal, but they cannot be shared and will always have a fixed interrupt level (namely, the one associated in hardware with the peripheral).

External Peripheral Interrupts The remaining interrupt sources are from external peripherals. These are defined in soc/soc.h as `ETS_*_INTR_SOURCE`.

Non-internal interrupt slots in both CPU cores are wired to an interrupt multiplexer, which can be used to route any external interrupt source to any of these interrupt slots.

- Allocating an external interrupt will always allocate it on the core that does the allocation.
- Freeing an external interrupt must always happen on the same core it was allocated on.
- Disabling and enabling external interrupts from another core is allowed.
- Multiple external interrupt sources can share an interrupt slot by passing `ESP_INTR_FLAG_SHARED` as a flag to esp_intr_alloc().

Care should be taken when calling esp_intr_alloc() from a task which is not pinned to a core. During task switching, these tasks can migrate between cores. Therefore it is impossible to tell which CPU the interrupt is allocated on,

which makes it difficult to free the interrupt handle and may also cause debugging difficulties. It is advised to use xTaskCreatePinnedToCore() with a specific CoreID argument to create tasks that will allocate interrupts. In the case of internal interrupt sources, this is required.

IRAM-Safe Interrupt Handlers

The `ESP_INTR_FLAG_IRAM` flag registers an interrupt handler that always runs from IRAM (and reads all its data from DRAM), and therefore does not need to be disabled during flash erase and write operations.

This is useful for interrupts which need a guaranteed minimum execution latency, as flash write and erase operations can be slow (erases can take tens or hundreds of milliseconds to complete).

It can also be useful to keep an interrupt handler in IRAM if it is called very frequently, to avoid flash cache misses.

Refer to the *SPI flash API documentation* for more details.

Multiple Handlers Sharing A Source

Several handlers can be assigned to a same source, given that all handlers are allocated using the `ESP_INTR_FLAG_SHARED` flag. They' ll be all allocated to the interrupt, which the source is attached to, and called sequentially when the source is active. The handlers can be disabled and freed individually. The source is attached to the interrupt (enabled), if one or more handlers are enabled, otherwise detached. A handler will never be called when disabled, while **its source may still be triggered** if any one of its handler enabled.

Sources attached to non-shared interrupt do not support this feature.

Though the framework support this feature, you have to use it *very carefully*. There usually exist 2 ways to stop a interrupt from being triggered: *disable the source* or *mask peripheral interrupt status*. IDF only handles the enabling and disabling of the source itself, leaving status and mask bits to be handled by users. **Status bits should always be masked before the handler responsible for it is disabled, or the status should be handled in other enabled interrupt properly**. You may leave some status bits unhandled if you just disable one of all the handlers without masking the status bits, which causes the interrupt to trigger infinitely resulting in a system crash.

API Reference

Header File

- esp32s2/include/esp_intr_alloc.h

Functions

esp_err_t **esp_intr_mark_shared**(int *intno*, int *cpu*, bool *is_in_iram*)

Mark an interrupt as a shared interrupt.

This will mark a certain interrupt on the specified CPU as an interrupt that can be used to hook shared interrupt handlers to.

Return ESP_ERR_INVALID_ARG if cpu or intno is invalid ESP_OK otherwise

Parameters

- `intno`: The number of the interrupt (0-31)
- `cpu`: CPU on which the interrupt should be marked as shared (0 or 1)
- `is_in_iram`: Shared interrupt is for handlers that reside in IRAM and the int can be left enabled while the flash cache is disabled.

esp_err_t **esp_intr_reserve**(int *intno*, int *cpu*)

Reserve an interrupt to be used outside of this framework.

This will mark a certain interrupt on the specified CPU as reserved, not to be allocated for any reason.

Return ESP_ERR_INVALID_ARG if cpu or intno is invalid ESP_OK otherwise

Parameters

- `intno`: The number of the interrupt (0-31)

- `cpu`: CPU on which the interrupt should be marked as shared (0 or 1)

esp_err_t **esp_intr_alloc**(int *source*, int *flags*, *intr_handler_t* *handler*, void **arg*, *intr_handle_t* **ret_handle*)

Allocate an interrupt with the given parameters.

This finds an interrupt that matches the restrictions as given in the flags parameter, maps the given interrupt source to it and hooks up the given interrupt handler (with optional argument) as well. If needed, it can return a handle for the interrupt as well.

The interrupt will always be allocated on the core that runs this function.

If ESP_INTR_FLAG_IRAM flag is used, and handler address is not in IRAM or RTC_FAST_MEM, then ESP_ERR_INVALID_ARG is returned.

Return ESP_ERR_INVALID_ARG if the combination of arguments is invalid. ESP_ERR_NOT_FOUND No free interrupt found with the specified flags ESP_OK otherwise

Parameters

- `source`: The interrupt source. One of the ETS_*_INTR_SOURCE interrupt mux sources, as defined in soc/soc.h, or one of the internal ETS_INTERNAL_*_INTR_SOURCE sources as defined in this header.
- `flags`: An ORred mask of the ESP_INTR_FLAG_* defines. These restrict the choice of interrupts that this routine can choose from. If this value is 0, it will default to allocating a non-shared interrupt of level 1, 2 or 3. If this is ESP_INTR_FLAG_SHARED, it will allocate a shared interrupt of level 1. Setting ESP_INTR_FLAG_INTRDISABLED will return from this function with the interrupt disabled.
- `handler`: The interrupt handler. Must be NULL when an interrupt of level >3 is requested, because these types of interrupts aren't C-callable.
- `arg`: Optional argument for passed to the interrupt handler
- `ret_handle`: Pointer to an intr_handle_t to store a handle that can later be used to request details or free the interrupt. Can be NULL if no handle is required.

esp_err_t **esp_intr_alloc_intrstatus**(int *source*, int *flags*, uint32_t *intrstatusreg*, uint32_t *intrstatusmask*, *intr_handler_t* *handler*, void **arg*, *intr_handle_t* **ret_handle*)

Allocate an interrupt with the given parameters.

This essentially does the same as esp_intr_alloc, but allows specifying a register and mask combo. For shared interrupts, the handler is only called if a read from the specified register, ANDed with the mask, returns non-zero. By passing an interrupt status register address and a fitting mask, this can be used to accelerate interrupt handling in the case a shared interrupt is triggered; by checking the interrupt statuses first, the code can decide which ISRs can be skipped

Return ESP_ERR_INVALID_ARG if the combination of arguments is invalid. ESP_ERR_NOT_FOUND No free interrupt found with the specified flags ESP_OK otherwise

Parameters

- `source`: The interrupt source. One of the ETS_*_INTR_SOURCE interrupt mux sources, as defined in soc/soc.h, or one of the internal ETS_INTERNAL_*_INTR_SOURCE sources as defined in this header.
- `flags`: An ORred mask of the ESP_INTR_FLAG_* defines. These restrict the choice of interrupts that this routine can choose from. If this value is 0, it will default to allocating a non-shared interrupt of level 1, 2 or 3. If this is ESP_INTR_FLAG_SHARED, it will allocate a shared interrupt of level 1. Setting ESP_INTR_FLAG_INTRDISABLED will return from this function with the interrupt disabled.
- `intrstatusreg`: The address of an interrupt status register
- `intrstatusmask`: A mask. If a read of address intrstatusreg has any of the bits that are 1 in the mask set, the ISR will be called. If not, it will be skipped.
- `handler`: The interrupt handler. Must be NULL when an interrupt of level >3 is requested, because these types of interrupts aren't C-callable.
- `arg`: Optional argument for passed to the interrupt handler
- `ret_handle`: Pointer to an intr_handle_t to store a handle that can later be used to request details or free the interrupt. Can be NULL if no handle is required.

esp_err_t **esp_intr_free** (*intr_handle_t handle*)
Disable and free an interrupt.

Use an interrupt handle to disable the interrupt and release the resources associated with it.

Note When the handler shares its source with other handlers, the interrupt status bits it's responsible for should be managed properly before freeing it. see `esp_intr_disable` for more details.
Return ESP_ERR_INVALID_ARG if handle is invalid, or esp_intr_free runs on another core than where the interrupt is allocated on. ESP_OK otherwise
Parameters
- `handle`: The handle, as obtained by esp_intr_alloc or esp_intr_alloc_intrstatus

int **esp_intr_get_cpu** (*intr_handle_t handle*)
Get CPU number an interrupt is tied to.

Return The core number where the interrupt is allocated
Parameters
- `handle`: The handle, as obtained by esp_intr_alloc or esp_intr_alloc_intrstatus

int **esp_intr_get_intno** (*intr_handle_t handle*)
Get the allocated interrupt for a certain handle.

Return The interrupt number
Parameters
- `handle`: The handle, as obtained by esp_intr_alloc or esp_intr_alloc_intrstatus

esp_err_t **esp_intr_disable** (*intr_handle_t handle*)
Disable the interrupt associated with the handle.

Note
1. For local interrupts (ESP_INTERNAL_* sources), this function has to be called on the CPU the interrupt is allocated on. Other interrupts have no such restriction.
2. When several handlers sharing a same interrupt source, interrupt status bits, which are handled in the handler to be disabled, should be masked before the disabling, or handled in other enabled interrupts properly. Miss of interrupt status handling will cause infinite interrupt calls and finally system crash.

Return ESP_ERR_INVALID_ARG if the combination of arguments is invalid. ESP_OK otherwise
Parameters
- `handle`: The handle, as obtained by esp_intr_alloc or esp_intr_alloc_intrstatus

esp_err_t **esp_intr_enable** (*intr_handle_t handle*)
Enable the interrupt associated with the handle.

Note For local interrupts (ESP_INTERNAL_* sources), this function has to be called on the CPU the interrupt is allocated on. Other interrupts have no such restriction.
Return ESP_ERR_INVALID_ARG if the combination of arguments is invalid. ESP_OK otherwise
Parameters
- `handle`: The handle, as obtained by esp_intr_alloc or esp_intr_alloc_intrstatus

esp_err_t **esp_intr_set_in_iram** (*intr_handle_t handle*, bool *is_in_iram*)
Set the "in IRAM" status of the handler.

Note Does not work on shared interrupts.
Return ESP_ERR_INVALID_ARG if the combination of arguments is invalid. ESP_OK otherwise
Parameters
- `handle`: The handle, as obtained by esp_intr_alloc or esp_intr_alloc_intrstatus
- `is_in_iram`: Whether the handler associated with this handle resides in IRAM. Handlers residing in IRAM can be called when cache is disabled.

void **esp_intr_noniram_disable** (void)
Disable interrupts that aren't specifically marked as running from IRAM.

void **esp_intr_noniram_enable** (void)
Re-enable interrupts disabled by esp_intr_noniram_disable.

Macros

ESP_INTR_FLAG_LEVEL1
Interrupt allocation flags.

These flags can be used to specify which interrupt qualities the code calling esp_intr_alloc* needs.Accept a Level 1 interrupt vector (lowest priority)

ESP_INTR_FLAG_LEVEL2
Accept a Level 2 interrupt vector.

ESP_INTR_FLAG_LEVEL3
Accept a Level 3 interrupt vector.

ESP_INTR_FLAG_LEVEL4
Accept a Level 4 interrupt vector.

ESP_INTR_FLAG_LEVEL5
Accept a Level 5 interrupt vector.

ESP_INTR_FLAG_LEVEL6
Accept a Level 6 interrupt vector.

ESP_INTR_FLAG_NMI
Accept a Level 7 interrupt vector (highest priority)

ESP_INTR_FLAG_SHARED
Interrupt can be shared between ISRs.

ESP_INTR_FLAG_EDGE
Edge-triggered interrupt.

ESP_INTR_FLAG_IRAM
ISR can be called if cache is disabled.

ESP_INTR_FLAG_INTRDISABLED
Return with this interrupt disabled.

ESP_INTR_FLAG_LOWMED
Low and medium prio interrupts. These can be handled in C.

ESP_INTR_FLAG_HIGH
High level interrupts. Need to be handled in assembly.

ESP_INTR_FLAG_LEVELMASK
Mask for all level flags.

ETS_INTERNAL_TIMER0_INTR_SOURCE
Xtensa timer 0 interrupt source.

The esp_intr_alloc* functions can allocate an int for all ETS_*_INTR_SOURCE interrupt sources that are routed through the interrupt mux. Apart from these sources, each core also has some internal sources that do not pass through the interrupt mux. To allocate an interrupt for these sources, pass these pseudo-sources to the functions.

ETS_INTERNAL_TIMER1_INTR_SOURCE
Xtensa timer 1 interrupt source.

ETS_INTERNAL_TIMER2_INTR_SOURCE
Xtensa timer 2 interrupt source.

ETS_INTERNAL_SW0_INTR_SOURCE
Software int source 1.

ETS_INTERNAL_SW1_INTR_SOURCE
Software int source 2.

ETS_INTERNAL_PROFILING_INTR_SOURCE
Int source for profiling.

ETS_INTERNAL_INTR_SOURCE_OFF
Provides SystemView with positive IRQ IDs, otherwise scheduler events are not shown properly

ESP_INTR_ENABLE(inum)
Enable interrupt by interrupt number

ESP_INTR_DISABLE(inum)
Disable interrupt by interrupt number

Type Definitions

typedef void (***intr_handler_t**)(void *arg)
Function prototype for interrupt handler function

typedef struct *intr_handle_data_t* **intr_handle_data_t**
Interrupt handler associated data structure

typedef *intr_handle_data_t* ***intr_handle_t**
Handle to an interrupt handler

2.6.16 Logging library

Overview

The logging library provides two ways for setting log verbosity:

- **At compile time**: in menuconfig, set the verbosity level using the option `CONFIG_LOG_DEFAULT_LEVEL`. All logging statements for verbosity levels higher than `CONFIG_LOG_DEFAULT_LEVEL` will be removed by the preprocessor.
- **At runtime**: all logs for verbosity levels lower than `CONFIG_LOG_DEFAULT_LEVEL` are enabled by default. The function *`esp_log_level_set()`* can be used to set a logging level on a per module basis. Modules are identified by their tags, which are human-readable ASCII zero-terminated strings.

There are the following verbosity levels:

- Error (lowest)
- Warning
- Info
- Debug
- Verbose (highest)

注解: The function *`esp_log_level_set()`* cannot set logging levels higher than specified by `CONFIG_LOG_DEFAULT_LEVEL`. To increase log level for a specific file at compile time, use the macro *LOG_LOCAL_LEVEL* (see the details below).

How to use this library

In each C file that uses logging functionality, define the TAG variable as shown below:

```
static const char* TAG = "MyModule";
```

Then use one of logging macros to produce output, e.g:

```
ESP_LOGW(TAG, "Baud rate error %.1f%%. Requested: %d baud, actual: %d baud", error_
↪* 100, baud_req, baud_real);
```

Several macros are available for different verbosity levels:

- `ESP_LOGE` - error (lowest)
- `ESP_LOGW` - warning

- ESP_LOGI - info
- ESP_LOGD - debug
- ESP_LOGV - verbose (highest)

Additionally, there are ESP_EARLY_LOGx versions for each of these macros, e.g., *ESP_EARLY_LOGE*. These versions have to be used explicitly in the early startup code only, before heap allocator and syscalls have been initialized. Normal ESP_LOGx macros can also be used while compiling the bootloader, but they will fall back to the same implementation as ESP_EARLY_LOGx macros.

To override default verbosity level at file or component scope, define the LOG_LOCAL_LEVEL macro.

At file scope, define it before including esp_log.h, e.g.:

```
#define LOG_LOCAL_LEVEL ESP_LOG_VERBOSE
#include "esp_log.h"
```

At component scope, define it in the component makefile:

```
CFLAGS += -D LOG_LOCAL_LEVEL=ESP_LOG_DEBUG
```

To configure logging output per module at runtime, add calls to the function *esp_log_level_set()* as follows:

```
esp_log_level_set("*", ESP_LOG_ERROR);        // set all components to ERROR level
esp_log_level_set("wifi", ESP_LOG_WARN);      // enable WARN logs from WiFi stack
esp_log_level_set("dhcpc", ESP_LOG_INFO);     // enable INFO logs from DHCP client
```

Logging to Host via JTAG By default, the logging library uses the vprintf-like function to write formatted output to the dedicated UART. By calling a simple API, all log output may be routed to JTAG instead, making logging several times faster. For details, please refer to Section 记录日志到主机.

Application Example

The logging library is commonly used by most esp-idf components and examples. For demonstration of log functionality, check ESP-IDF's examples directory. The most revelant examples that deal with logging are the following:

- system/ota
- storage/sd_card
- protocols/https_request

API Reference

Header File

- log/include/esp_log.h

Functions

void **esp_log_level_set**(**const** char *tag, *esp_log_level_t level*)

Set log level for given tag.

If logging for given component has already been enabled, changes previous setting.

Note that this function can not raise log level above the level set using CONFIG_LOG_DEFAULT_LEVEL setting in menuconfig.

To raise log level above the default one for a given file, define LOG_LOCAL_LEVEL to one of the ESP_LOG_* values, before including esp_log.h in this file.

Parameters

- tag: Tag of the log entries to enable. Must be a non-NULL zero terminated string. Value "*" resets log level for all tags to the given value.

- `level`: Selects log level to enable. Only logs at this and lower verbosity levels will be shown.

vprintf_like_t **esp_log_set_vprintf** (*vprintf_like_t func*)
Set function used to output log entries.

By default, log output goes to UART0. This function can be used to redirect log output to some other destination, such as file or network. Returns the original log handler, which may be necessary to return output to the previous destination.

Return func old Function used for output.
Parameters
- `func`: new Function used for output. Must have same signature as vprintf.

uint32_t **esp_log_timestamp** (void)
Function which returns timestamp to be used in log output.

This function is used in expansion of ESP_LOGx macros. In the 2nd stage bootloader, and at early application startup stage this function uses CPU cycle counter as time source. Later when FreeRTOS scheduler start running, it switches to FreeRTOS tick count.

For now, we ignore millisecond counter overflow.

Return timestamp, in milliseconds

char ***esp_log_system_timestamp** (void)
Function which returns system timestamp to be used in log output.

This function is used in expansion of ESP_LOGx macros to print the system time as "HH:MM:SS.sss". The system time is initialized to 0 on startup, this can be set to the correct time with an SNTP sync, or manually with standard POSIX time functions.

Currently this will not get used in logging from binary blobs (i.e WiFi & Bluetooth libraries), these will still print the RTOS tick time.

Return timestamp, in "HH:MM:SS.sss"

uint32_t **esp_log_early_timestamp** (void)
Function which returns timestamp to be used in log output.

This function uses HW cycle counter and does not depend on OS, so it can be safely used after application crash.

Return timestamp, in milliseconds

void **esp_log_write** (*esp_log_level_t level*, **const** char **tag*, **const** char **format*, ...)
Write message into the log.

This function is not intended to be used directly. Instead, use one of ESP_LOGE, ESP_LOGW, ESP_LOGI, ESP_LOGD, ESP_LOGV macros.

This function or these macros should not be used from an interrupt.

void **esp_log_writev** (*esp_log_level_t level*, **const** char **tag*, **const** char **format*, va_list *args*)
Write message into the log, va_list variant.

This function is provided to ease integration toward other logging framework, so that esp_log can be used as a log sink.

See esp_log_write()

Macros

ESP_LOG_BUFFER_HEX_LEVEL (tag, buffer, buff_len, level)
Log a buffer of hex bytes at specified level, separated into 16 bytes each line.

Parameters
- `tag`: description tag
- `buffer`: Pointer to the buffer array
- `buff_len`: length of buffer in bytes

- `level`: level of the log

ESP_LOG_BUFFER_CHAR_LEVEL(tag, buffer, buff_len, level)

Log a buffer of characters at specified level, separated into 16 bytes each line. Buffer should contain only printable characters.

Parameters

- `tag`: description tag
- `buffer`: Pointer to the buffer array
- `buff_len`: length of buffer in bytes
- `level`: level of the log

ESP_LOG_BUFFER_HEXDUMP(tag, buffer, buff_len, level)

Dump a buffer to the log at specified level.

The dump log shows just like the one below:

```
 W (195) log_example: 0x3ffb4280   45 53 50 33 32 20 69 73  20 67 72 65 61 74␣
↪2c 20  |ESP32 is great, |
 W (195) log_example: 0x3ffb4290   77 6f 72 6b 69 6e 67 20  61 6c 6f 6e 67 20␣
↪77 69  |working along wi|
 W (205) log_example: 0x3ffb42a0   74 68 20 74 68 65 20 49  44 46 2e 00       ␣
↪       |th the IDF..|
```

It is highly recommend to use terminals with over 102 text width.

Parameters

- `tag`: description tag
- `buffer`: Pointer to the buffer array
- `buff_len`: length of buffer in bytes
- `level`: level of the log

ESP_LOG_BUFFER_HEX(tag, buffer, buff_len)

Log a buffer of hex bytes at Info level.

See `esp_log_buffer_hex_level`

Parameters

- `tag`: description tag
- `buffer`: Pointer to the buffer array
- `buff_len`: length of buffer in bytes

ESP_LOG_BUFFER_CHAR(tag, buffer, buff_len)

Log a buffer of characters at Info level. Buffer should contain only printable characters.

See `esp_log_buffer_char_level`

Parameters

- `tag`: description tag
- `buffer`: Pointer to the buffer array
- `buff_len`: length of buffer in bytes

ESP_EARLY_LOGE(tag, format, ...)

macro to output logs in startup code, before heap allocator and syscalls have been initialized. log at `ESP_LOG_ERROR` level.

See `printf`,`ESP_LOGE`

ESP_EARLY_LOGW(tag, format, ...)

macro to output logs in startup code at `ESP_LOG_WARN` level.

See `ESP_EARLY_LOGE`,`ESP_LOGE`, `printf`

ESP_EARLY_LOGI(tag, format, ...)

macro to output logs in startup code at `ESP_LOG_INFO` level.

See `ESP_EARLY_LOGE`,`ESP_LOGE`, `printf`

ESP_EARLY_LOGD (tag, format, ...)
macro to output logs in startup code at ESP_LOG_DEBUG level.

See ESP_EARLY_LOGE,ESP_LOGE, printf

ESP_EARLY_LOGV (tag, format, ...)
macro to output logs in startup code at ESP_LOG_VERBOSE level.

See ESP_EARLY_LOGE,ESP_LOGE, printf

ESP_LOG_EARLY_IMPL (tag, format, log_level, log_tag_letter, ...)

ESP_LOGE (tag, format, ...)

ESP_LOGW (tag, format, ...)

ESP_LOGI (tag, format, ...)

ESP_LOGD (tag, format, ...)

ESP_LOGV (tag, format, ...)

ESP_LOG_LEVEL (level, tag, format, ...)
runtime macro to output logs at a specified level.

See printf
Parameters
- tag: tag of the log, which can be used to change the log level by esp_log_level_set at runtime.
- level: level of the output log.
- format: format of the output log. see printf
- ...: variables to be replaced into the log. see printf

ESP_LOG_LEVEL_LOCAL (level, tag, format, ...)
runtime macro to output logs at a specified level. Also check the level with LOG_LOCAL_LEVEL.

See printf, ESP_LOG_LEVEL

ESP_DRAM_LOGE (tag, format, ...)
Macro to output logs when the cache is disabled. log at ESP_LOG_ERROR level.

Similar to ESP_EARLY_LOGE, the log level cannot be changed by esp_log_level_set.

Usage: ESP_DRAM_LOGE(DRAM_STR("my_tag"), "format", orESP_DRAM_LOGE(TAG, "format" , ⋯)', where TAG is a char* that points to a str in the DRAM.

Note Placing log strings in DRAM reduces available DRAM, so only use when absolutely essential.
See ets_printf,ESP_LOGE

ESP_DRAM_LOGW (tag, format, ...)
macro to output logs when the cache is disabled at ESP_LOG_WARN level.

See ESP_DRAM_LOGW,ESP_LOGW, ets_printf

ESP_DRAM_LOGI (tag, format, ...)
macro to output logs when the cache is disabled at ESP_LOG_INFO level.

See ESP_DRAM_LOGI,ESP_LOGI, ets_printf

ESP_DRAM_LOGD (tag, format, ...)
macro to output logs when the cache is disabled at ESP_LOG_DEBUG level.

See ESP_DRAM_LOGD,ESP_LOGD, ets_printf

ESP_DRAM_LOGV (tag, format, ...)
macro to output logs when the cache is disabled at ESP_LOG_VERBOSE level.

See ESP_DRAM_LOGV,ESP_LOGV, ets_printf

Type Definitions

typedef int (***vprintf_like_t**)(**const** char *, va_list)

Enumerations

enum esp_log_level_t

Log level.

Values:

ESP_LOG_NONE
No log output

ESP_LOG_ERROR
Critical errors, software module can not recover on its own

ESP_LOG_WARN
Error conditions from which recovery measures have been taken

ESP_LOG_INFO
Information messages which describe normal flow of events

ESP_LOG_DEBUG
Extra information which is not necessary for normal use (values, pointers, sizes, etc).

ESP_LOG_VERBOSE
Bigger chunks of debugging information, or frequent messages which can potentially flood the output.

2.6.17 Miscellaneous System APIs

Software reset

To perform software reset of the chip, *esp_restart()* function is provided. When the function is called, execution of the program will stop, both CPUs will be reset, application will be loaded by the bootloader and started again.

Additionally, *esp_register_shutdown_handler()* function is provided to register a routine which needs to be called prior to restart (when done by *esp_restart()*). This is similar to the functionality of `atexit` POSIX function.

Reset reason

ESP-IDF application can be started or restarted due to a variety of reasons. To get the last reset reason, call *esp_reset_reason()* function. See description of *esp_reset_reason_t* for the list of possible reset reasons.

Heap memory

Two heap memory related functions are provided:

- *esp_get_free_heap_size()* returns the current size of free heap memory
- *esp_get_minimum_free_heap_size()* returns the minimum size of free heap memory that was available during program execution.

Note that ESP-IDF supports multiple heaps with different capabilities. Functions mentioned in this section return the size of heap memory which can be allocated using `malloc` family of functions. For further information about heap memory see *Heap Memory Allocation*.

Random number generation

ESP32-S2 contains a hardware random number generator, values from it can be obtained using `esp_random()`.

When Wi-Fi or Bluetooth are enabled, numbers returned by hardware random number generator (RNG) can be considered true random numbers. Without Wi-Fi or Bluetooth enabled, hardware RNG is a pseudo-random number generator. At startup, ESP-IDF bootloader seeds the hardware RNG with entropy, but care must be taken when reading random values between the start of `app_main` and initialization of Wi-Fi or Bluetooth drivers.

MAC Address

These APIs allow querying and customizing MAC addresses for different network interfaces that supported (e.g. Wi-Fi, Bluetooth, Ethernet).

In ESP-IDF these addresses are calculated from *Base MAC address*. Base MAC address can be initialized with factory-programmed value from internal eFuse, or with a user-defined value. In addition to setting the base MAC address, applications can specify the way in which MAC addresses are allocated to devices. See *Number of universally administered MAC address* section for more details.

Interface	MAC address (2 universally administered)	MAC address (1 universally administered)
Wi-Fi Station	base_mac	base_mac
Wi-Fi SoftAP	base_mac, +1 to the last octet	base_mac, first octet randomized

Base MAC address To fetch MAC address for a specific interface (e.g. Wi-Fi, Bluetooth, Ethernet), you can simply use `esp_read_mac()` function.

By default, this function takes the eFuse value burned at a pre-defined block (e.g. BLK0 for ESP32, BLK1 for ESP32-S2) as the base MAC address. Per-interface MAC addresses will be calculated according to the table above.

Applications who want to customize base MAC address (not the one provided by Espressif) should call `esp_base_mac_addr_set()` before `esp_read_mac()`. The customized MAC address can be stored in any supported storage device (e.g. Flash, NVS, etc).

Note that, calls to `esp_base_mac_addr_set()` should take place before the initialization of network stack, for example, early in `app_main`.

Custom MAC address in eFuse To facilitate the usage of custom MAC addresses, ESP-IDF provides `esp_efuse_mac_get_custom()` function, which loads MAC address from internal pre-defined eFuse block (e.g. BLK3 for ESP32). This function assumes that custom MAC address is stored in the following format:

Field	# of bits	Range of bits	Notes
Version	8	191:184	0: invalid, others —valid
Reserved	128	183:56	
MAC address	48	55:8	
MAC address CRC	8	7:0	CRC-8-CCITT, polynomial 0x07

Once MAC address has been obtained using `esp_efuse_mac_get_custom()`, call `esp_base_mac_addr_set()` to set this MAC address as base MAC address.

Number of universally administered MAC address Several MAC addresses (universally administered by IEEE) are uniquely assigned to the networking interfaces (Wi-Fi/BT/Ethernet). The final octet of each universally administered MAC address increases by one. Only the first one of them (which is called base MAC address) is stored in eFuse or external storage, the others are generated from it. Here, ‘generate’ means adding 0, 1, 2 and 3 (respectively) to the final octet of the base MAC address.

If the universally administered MAC addresses are not enough for all of the networking interfaces, locally administered MAC addresses which are derived from universally administered MAC addresses are assigned to the rest of networking interfaces.

See this article for the definition of local and universally administered MAC addresses.

Chip version

esp_chip_info() function fills *esp_chip_info_t* structure with information about the chip. This includes the chip revision, number of CPU cores, and a bit mask of features enabled in the chip.

SDK version

esp_get_idf_version() returns a string describing the IDF version which was used to compile the application. This is the same value as the one available through `IDF_VER` variable of the build system. The version string generally has the format of `git describe` output.

To get the version at build time, additional version macros are provided. They can be used to enable or disable parts of the program depending on IDF version.

- *ESP_IDF_VERSION_MAJOR*, *ESP_IDF_VERSION_MINOR*, *ESP_IDF_VERSION_PATCH* are defined to integers representing major, minor, and patch version.
- *ESP_IDF_VERSION_VAL* and *ESP_IDF_VERSION* can be used when implementing version checks:

```
#include "esp_idf_version.h"

#if ESP_IDF_VERSION >= ESP_IDF_VERSION_VAL(4, 0, 0)
    // enable functionality present in IDF v4.0
#endif
```

App version

Application version is stored in *esp_app_desc_t* structure. It is located in DROM sector and has a fixed offset from the beginning of the binary file. The structure is located after *esp_image_header_t* and *esp_image_segment_header_t* structures. The field version has string type and max length 32 chars.

To set version in your project manually you need to set `PROJECT_VER` variable in your project CMakeLists.txt/Makefile:

- In application CMakeLists.txt put `set(PROJECT_VER "0.1.0.1")` before including `project.cmake`.

(For legacy GNU Make build system: in application Makefile put `PROJECT_VER = "0.1.0.1"` before including `project.mk`.)

If *CONFIG_APP_PROJECT_VER_FROM_CONFIG* option is set, the value of *CONFIG_APP_PROJECT_VER* will be used. Otherwise if `PROJECT_VER` variable is not set in the project then it will be retrieved from either `$(PROJECT_PATH)/version.txt` file (if present) else using git command `git describe`. If neither is available then `PROJECT_VER` will be set to "1". Application can make use of this by calling *esp_ota_get_app_description()* or *esp_ota_get_partition_description()* functions.

API Reference

Header File

- esp_system/include/esp_system.h

Functions

esp_err_t **esp_register_shutdown_handler**(*shutdown_handler_t handle*)

Register shutdown handler.

This function allows you to register a handler that gets invoked before the application is restarted using esp_restart function.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if the handler has already been registered
- ESP_ERR_NO_MEM if no more shutdown handler slots are available

Parameters

- `handle`: function to execute on restart

esp_err_t **esp_unregister_shutdown_handler**(*shutdown_handler_t handle*)

Unregister shutdown handler.

This function allows you to unregister a handler which was previously registered using esp_register_shutdown_handler function.

- ESP_OK on success
- ESP_ERR_INVALID_STATE if the given handler hasn't been registered before

void **esp_restart**(void)

Restart PRO and APP CPUs.

This function can be called both from PRO and APP CPUs. After successful restart, CPU reset reason will be SW_CPU_RESET. Peripherals (except for WiFi, BT, UART0, SPI1, and legacy timers) are not reset. This function does not return.

esp_reset_reason_t **esp_reset_reason**(void)

Get reason of last reset.

Return See description of esp_reset_reason_t for explanation of each value.

uint32_t **esp_get_free_heap_size**(void)

Get the size of available heap.

Note that the returned value may be larger than the maximum contiguous block which can be allocated.

Return Available heap size, in bytes.

uint32_t **esp_get_free_internal_heap_size**(void)

Get the size of available internal heap.

Note that the returned value may be larger than the maximum contiguous block which can be allocated.

Return Available internal heap size, in bytes.

uint32_t **esp_get_minimum_free_heap_size**(void)

Get the minimum heap that has ever been available.

Return Minimum free heap ever available

uint32_t **esp_random**(void)

Get one random 32-bit word from hardware RNG.

The hardware RNG is fully functional whenever an RF subsystem is running (ie Bluetooth or WiFi is enabled). For random values, call this function after WiFi or Bluetooth are started.

If the RF subsystem is not used by the program, the function bootloader_random_enable() can be called to enable an entropy source. bootloader_random_disable() must be called before RF subsystem or I2S peripheral are used. See these functions' documentation for more details.

Any time the app is running without an RF subsystem (or bootloader_random) enabled, RNG hardware should be considered a PRNG. A very small amount of entropy is available due to pre-seeding while the IDF bootloader is running, but this should not be relied upon for any use.

Return Random value between 0 and UINT32_MAX

void **esp_fill_random**(void *buf, size_t len)

Fill a buffer with random bytes from hardware RNG.

Note This function has the same restrictions regarding available entropy as esp_random()

Parameters

- `buf`: Pointer to buffer to fill with random numbers.
- `len`: Length of buffer in bytes

esp_err_t **esp_base_mac_addr_set**(**const** uint8_t *mac)

Set base MAC address with the MAC address which is stored in BLK3 of EFUSE or external storage e.g. flash and EEPROM.

Base MAC address is used to generate the MAC addresses used by the networking interfaces. If using base MAC address stored in BLK3 of EFUSE or external storage, call this API to set base MAC address with the MAC address which is stored in BLK3 of EFUSE or external storage before initializing WiFi/BT/Ethernet.

Note Base MAC must be a unicast MAC (least significant bit of first byte must be zero).

Note If not using a valid OUI, set the "locally administered" bit (bit value 0x02 in the first byte) to avoid collisions.

Return ESP_OK on success ESP_ERR_INVALID_ARG If mac is NULL or is not a unicast MAC

Parameters

- `mac`: base MAC address, length: 6 bytes.

esp_err_t **esp_base_mac_addr_get**(uint8_t *mac)

Return base MAC address which is set using esp_base_mac_addr_set.

Return ESP_OK on success ESP_ERR_INVALID_MAC base MAC address has not been set

Parameters

- `mac`: base MAC address, length: 6 bytes.

esp_err_t **esp_efuse_mac_get_custom**(uint8_t *mac)

Return base MAC address which was previously written to BLK3 of EFUSE.

Base MAC address is used to generate the MAC addresses used by the networking interfaces. This API returns the custom base MAC address which was previously written to BLK3 of EFUSE. Writing this EFUSE allows setting of a different (non-Espressif) base MAC address. It is also possible to store a custom base MAC address elsewhere, see esp_base_mac_addr_set() for details.

Return ESP_OK on success ESP_ERR_INVALID_VERSION An invalid MAC version field was read from BLK3 of EFUSE ESP_ERR_INVALID_CRC An invalid MAC CRC was read from BLK3 of EFUSE

Parameters

- `mac`: base MAC address, length: 6 bytes.

esp_err_t **esp_efuse_mac_get_default**(uint8_t *mac)

Return base MAC address which is factory-programmed by Espressif in BLK0 of EFUSE.

Return ESP_OK on success

Parameters

- `mac`: base MAC address, length: 6 bytes.

esp_err_t **esp_read_mac**(uint8_t *mac, esp_mac_type_t type)

Read base MAC address and set MAC address of the interface.

This function first get base MAC address using esp_base_mac_addr_get or reads base MAC address from BLK0 of EFUSE. Then set the MAC address of the interface including wifi station, wifi softap, bluetooth and ethernet.

Return ESP_OK on success

Parameters

- `mac`: MAC address of the interface, length: 6 bytes.
- `type`: type of MAC address, 0:wifi station, 1:wifi softap, 2:bluetooth, 3:ethernet.

esp_err_t **esp_derive_local_mac**(uint8_t *local_mac, **const** uint8_t *universal_mac)

Derive local MAC address from universal MAC address.

This function derives a local MAC address from an universal MAC address. A `definition of local vs universal MAC address can be found on Wikipedia <>`. In ESP32, universal MAC address is generated from base MAC address in EFUSE or other external storage. Local MAC address is derived from the universal MAC address.

Return ESP_OK on success

Parameters

- `local_mac`: Derived local MAC address, length: 6 bytes.
- `universal_mac`: Source universal MAC address, length: 6 bytes.

void **esp_system_abort**(**const** char *_details_)

Trigger a software abort.

Parameters

- `details`: Details that will be displayed during panic handling.

void **esp_chip_info**(*esp_chip_info_t* *_out_info_)

Fill an *esp_chip_info_t* structure with information about the chip.

Parameters

- `[out] out_info`: structure to be filled

Structures

struct esp_chip_info_t

The structure represents information about the chip.

Public Members

esp_chip_model_t **model**

chip model, one of esp_chip_model_t

uint32_t **features**

bit mask of CHIP_FEATURE_x feature flags

uint8_t **cores**

number of CPU cores

uint8_t **revision**

chip revision number

Macros

CHIP_FEATURE_EMB_FLASH

Chip has embedded flash memory.

CHIP_FEATURE_WIFI_BGN

Chip has 2.4GHz WiFi.

CHIP_FEATURE_BLE

Chip has Bluetooth LE.

CHIP_FEATURE_BT

Chip has Bluetooth Classic.

Type Definitions

typedef void (***shutdown_handler_t**)(void)

Shutdown handler type

Enumerations

enum esp_mac_type_t

Values:

ESP_MAC_WIFI_STA

ESP_MAC_WIFI_SOFTAP

ESP_MAC_BT

ESP_MAC_ETH

enum esp_reset_reason_t

Reset reasons.

Values:

ESP_RST_UNKNOWN
Reset reason can not be determined.

ESP_RST_POWERON
Reset due to power-on event.

ESP_RST_EXT
Reset by external pin (not applicable for ESP32)

ESP_RST_SW
Software reset via esp_restart.

ESP_RST_PANIC
Software reset due to exception/panic.

ESP_RST_INT_WDT
Reset (software or hardware) due to interrupt watchdog.

ESP_RST_TASK_WDT
Reset due to task watchdog.

ESP_RST_WDT
Reset due to other watchdogs.

ESP_RST_DEEPSLEEP
Reset after exiting deep sleep mode.

ESP_RST_BROWNOUT
Brownout reset (software or hardware)

ESP_RST_SDIO
Reset over SDIO.

enum esp_chip_model_t

Chip models.

Values:

CHIP_ESP32 = 1
ESP32.

CHIP_ESP32S2 = 2
ESP32-S2.

Header File

- esp_common/include/esp_idf_version.h

Functions

const char ***esp_get_idf_version**(void)

Return full IDF version string, same as 'git describe' output.

Note If you are printing the ESP-IDF version in a log file or other information, this function provides more information than using the numerical version macros. For example, numerical version macros don't differentiate between development, pre-release and release versions, but the output of this function does.

Return constant string from IDF_VER

Macros

ESP_IDF_VERSION_MAJOR
Major version number (X.x.x)

ESP_IDF_VERSION_MINOR
Minor version number (x.X.x)

ESP_IDF_VERSION_PATCH
Patch version number (x.x.X)

ESP_IDF_VERSION_VAL(major, minor, patch)
Macro to convert IDF version number into an integer

To be used in comparisons, such as ESP_IDF_VERSION >= ESP_IDF_VERSION_VAL(4, 0, 0)

ESP_IDF_VERSION
Current IDF version, as an integer

To be used in comparisons, such as ESP_IDF_VERSION >= ESP_IDF_VERSION_VAL(4, 0, 0)

2.6.18 空中升级 (OTA)

OTA 流程概览

OTA 升级机制可以让设备在固件正常运行时根据接收数据更新（如通过 Wi-Fi 或蓝牙）。

要运行 OTA 机制，需配置设备的:doc:‘ 分区表 <../../api-guides/partition-tables>‘，该分区表至少包括两个 OTA 应用程序分区（即 *ota_0* 和 *ota_1*）和一个 OTA 数据分区。

OTA 功能启动后，向当前未用于启动的 OTA 应用分区写入新的应用固件镜像。镜像验证后，OTA 数据分区更新，指定在下一次启动时使用该镜像。

OTA 数据分区

所有使用 OTA 功能项目，其:doc:‘ 分区表 <../../api-guides/partition-tables>‘必须包含一个 OTA 数据分区（类型为 `data`，子类型为 `ota`）。

工厂启动设置下，OTA 数据分区中应没有数据（所有字节擦写成 0xFF）。如果分区表中有工厂应用程序，ESP-IDF 软件启动加载器会启动工厂应用程序。如果分区表中没有工厂应用程序，则启动第一个可用的 OTA 分区（通常是 `ota_0`）。

第一次 OTA 升级后，OTA 数据分区更新，指定下一次启动哪个 OTA 应用程序分区。

OTA 数据分区是两个 0x2000 字节大小的 flash 扇区，防止写入时电源故障引发问题。两个扇区单独擦除、写入匹配数据，若存在不一致，则用计数器字段判定哪个扇区为最新数据。

应用程序回滚

应用程序回滚的主要目的是确保设备更新后正常运转。该功能可使设备在更新新版本后出现严重错误时，回滚到之前正常运行的应用版本。回滚使能，OTA 升级，应用更新至新版本，之后可能有以下三种情况：

- 应用程序运行正常*esp_ota_mark_app_valid_cancel_rollback()* 将应用程序状态标记为 `ESP_OTA_IMG_VALID`，启动无限制。
- 应用程序出现严重错误，无法继续工作，必须回滚到此前的版本，*esp_ota_mark_app_invalid_rollback_and_reboot()* 将正在运行的版本标记为 `ESP_OTA_IMG_INVALID` 然后复位。启动加载器不会选取此版本，而是此前正常运行的版本。
- 如果*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 使能，则无需调用函数便可复位，回滚至之前的应用版本。

注解：应用程序的状态不是写到程序的二进制镜像，而是写到 otadata 分区。该分区有一个 ota_seq 计数器，该计数器是 OTA 应用分区的指针，指向下次启动时选取应用所在的分区 (ota_0, ota_1, ···)。

应用程序 OTA 状态 状态控制了选取启动应用程序的过程：

状态	启动加载器选取启动应用程序的限制
ESP_OTA_IMG_VALID	没有限制，可以选取。
ESP_OTA_IMG_UNDEFINED	没有限制，可以选取。
ESP_OTA_IMG_INVALID	不会选取。
ESP_OTA_IMG_ABORTED	不会选取。
ESP_OTA_IMG_NEW	如使能*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE*，则仅会选取一次。在启动加载器中，状态立即变为 ESP_OTA_IMG_PENDING_VERIFY。
ESP_OTA_IMG_PENDING_VERIFY	如使能*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE*，则不会选取，状态变为 “ESP_OTA_IMG_ABORTED“。

如果*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 没有使能（默认情况），则*esp_ota_mark_app_valid_cancel_rollback()* 和*esp_ota_mark_app_invalid_rollback_and_reboot* 为可选功能，ESP_OTA_IMG_NEW 和 ESP_OTA_IMG_PENDING_VERIFY 不会使用。

Kconfig 中的*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 可以帮助用户追踪新版应用程序的第一次启动。应用程序需调用*esp_ota_mark_app_valid_cancel_rollback()* 函数确认可以运行，否则将会在重启时回滚至旧版本。该功能可让用户在启动阶段控制应用程序的可操作性。新版应用程序仅有一次机会尝试是否能成功启动。

回滚过程 *CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 使能时，回滚过程如下：

- 新版应用程序下载成功，*esp_ota_set_boot_partition()* 函数将分区设为可启动，状态设为 ESP_OTA_IMG_NEW。该状态表示应用程序为新版本，第一次启动需要监测。
- 重新启动*esp_restart()*。
- 启动加载器检查新版应用程序，若状态设置为 ESP_OTA_IMG_PENDING_VERIFY，则写入 ESP_OTA_IMG_ABORTED。
- 启动加载器选取新版应用程序启动，应用程序状态不设置为 ESP_OTA_IMG_INVALID 或 ESP_OTA_IMG_ABORTED。
- 启动加载器检查所选取的新版应用程序，若状态设置为 ESP_OTA_IMG_NEW，则写入 ESP_OTA_IMG_PENDING_VERIFY。该状态表示，需确认应用程序的可操作性，如不确认，发生重启，则状态会重写为 ESP_OTA_IMG_ABORTED （见上文），该应用程序不可再启动，将回滚至上一版本。
- 新版应用程序启动，应进行自测。
- 若通过自测，则必须调用函数*esp_ota_mark_app_valid_cancel_rollback()*，因为新版应用程序在等待确认其可操作性 (ESP_OTA_IMG_PENDING_VERIFY 状态)。
- 若未通过自测，则调用函数*esp_ota_mark_app_invalid_rollback_and_reboot()*，回滚至之前的版本，同时无效的新版本设置为 ESP_OTA_IMG_INVALID。
- 如果新版应用程序可操作性没有确认，则状态一直为 ESP_OTA_IMG_PENDING_VERIFY。下一次启动时，状态变更为 ESP_OTA_IMG_ABORTED，阻止其再次启动，之后回滚到之前的版本。

意外复位 如果在新版应用第一次启动时发生断电或意外崩溃，则会回滚至之前正常运行的版本。

建议：尽快完成自测，防止因断电回滚。

只有 OTA 分区可以回滚。工厂分区不会回滚。

启动无效/中止的应用程序 用户可以启动此前设置为 ESP_OTA_IMG_INVALID 或 ESP_OTA_IMG_ABORTED 的应用程序：

- 获取最后一个无效应用分区*esp_ota_get_last_invalid_partition()*。
- 将获取的分区传递给*esp_ota_set_boot_partition()*，更新 otadata。
- 重启*esp_restart()*。启动加载器会启动指定应用程序。

要确定是否在应用程序启动时进行自测，可以调用*esp_ota_get_state_partition()* 函数。如果结果为 ESP_OTA_IMG_PENDING_VERIFY，则需要自测，后续确认应用程序的可操作性。

如何设置状态 下文简单描述了如何设置应用程序状态：

- ESP_OTA_IMG_VALID 由函数*esp_ota_mark_app_valid_cancel_rollback()* 设置。
- 如果*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 没有使能，ESP_OTA_IMG_UNDEFINED 由函数*esp_ota_set_boot_partition()* 设置。
- 如果*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 没有使能，ESP_OTA_IMG_NEW 由函数*esp_ota_set_boot_partition()* 设置。
- ESP_OTA_IMG_INVALID 由函数*esp_ota_mark_app_invalid_rollback_and_reboot()* 设置。
- 如果应用程序的可操作性无法确认，发生重启（*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 使能），则设置 ESP_OTA_IMG_ABORTED。
- 如果*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 使能，选取的应用程序状态为 ESP_OTA_IMG_NEW，则在启动加载器中设置 ESP_OTA_IMG_PENDING_VERIFY。

防回滚

防回滚机制可以防止回滚到安全版本号低于芯片 eFuse 中烧录程序的应用程序版本。

设置*CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK*，启动防回滚机制。在启动加载器中选取可启动的应用程序，会额外检查芯片和应用程序镜像的安全版本号。可启动固件中的应用安全版本号必须等于或高于芯片中的应用安全版本号。

CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK 和*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE* 一起使用。此时，只有安全版本号等于或高于芯片中的应用安全版本号时才会回滚。

典型的防回滚机制

- 新发布的固件解决了此前版本的安全问题。
- 开发者在确保固件可以运行之后，增加安全版本号，发布固件。
- 下载新版应用程序。
- 运行函数*esp_ota_set_boot_partition()*，将新版应用程序设为可启动。如果新版应用程序的安全版本号低于芯片中的应用安全版本号，新版应用程序会被擦除，无法更新到新固件。
- 重新启动。
- 在启动加载器中选取安全版本号等于或高于芯片中应用安全版本号的应用程序。如果 otadata 处于初始阶段，通过串行通道加载了安全版本号高于芯片中应用安全版本号的固件，则启动加载器中 eFuse 的安全版本号会立即更新。
- 新版应用程序启动，之后进行可操作性检测，如果通过检测，则调用函数*esp_ota_mark_app_valid_cancel_rollback()*，将应用程序标记为 ESP_OTA_IMG_VALID，更新芯片中应用程序的安全版本号。注意，如果调用函数*esp_ota_mark_app_invalid_rollback_and_reboot()*，可能会因为设备中没有可启动的应用程序而回滚失败，返回 ESP_ERR_OTA_ROLLBACK_FAILED 错误，应用程序状态一直为 ESP_OTA_IMG_PENDING_VERIFY。
- 如果运行的应用程序处于 ESP_OTA_IMG_VALID 状态，则可再次更新。

建议：

如果想避免因服务器应用程序的安全版本号低于运行的应用程序，造成不必要的下载和擦除，必须从镜像的第一个包中获取 new_app_info.secure_version，和 eFuse 的安全版本号比较。如果 esp_efuse_check_secure_version(new_app_info.secure_version) 函数为真，则下载继续，反之则中断。

```
....
bool image_header_was_checked = false;
while (1) {
    int data_read = esp_http_client_read(client, ota_write_data, BUFFSIZE);
    ...
```

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```
    if (data_read > 0) {
        if (image_header_was_checked == false) {
            esp_app_desc_t new_app_info;
            if (data_read > sizeof(esp_image_header_t) + sizeof(esp_image_segment_
↪header_t) + sizeof(esp_app_desc_t)) {
                // check current version with downloading
                if (esp_efuse_check_secure_version(new_app_info.secure_version) ==␣
↪false) {
                    ESP_LOGE(TAG, "This a new app can not be downloaded due to a␣
↪secure version is lower than stored in efuse.");
                    http_cleanup(client);
                    task_fatal_error();
                }

                image_header_was_checked = true;

                esp_ota_begin(update_partition, OTA_SIZE_UNKNOWN, &update_handle);
            }
        }
        esp_ota_write( update_handle, (const void *)ota_write_data, data_read);
    }
}
...
```

限制：

- secure_version 字段最多有 32 位。也就是说，防回滚最多可以做 32 次。用户可以使用 *CONFIG_BOOTLOADER_APP_SEC_VER_SIZE_EFUSE_FIELD* 减少该 eFuse 字段的长度。
- 防回滚仅在 eFuse 编码机制设置为 NONE 时生效。
- 分区表不应有工厂分区，应仅有两个应用程序分区。

security_version:

- 存储在应用程序镜像中的 esp_app_desc 里。版本号用 *CONFIG_BOOTLOADER_APP_SECURE_VERSION* 设置。

OTA 工具 (otatool.py)

app_update 组件中有 otatool.py 工具，用于在目标设备上完成下列 OTA 分区相关操作：

- 读取 otadata 分区 (read_otadata)
- 擦除 otadata 分区，将设备复位至工厂应用程序 (erase_otadata)
- 切换 OTA 分区 (switch_ota_partition)
- 擦除 OTA 分区 (erase_ota_partition)
- 写入 OTA 分区 (write_ota_partition)
- 读取 OTA 分区 (read_ota_partition)

用户若想通过编程方式完成相关操作，可从另一个 Python 脚本导入并使用分区工具，或者从 Shell 脚本调用分区工具。前者可使用工具的 Python API，后者可使用命令行界面。

Python API 首先，确保已导入 *otatool* 模块。

```
import sys
import os

idf_path = os.environ["IDF_PATH"]  # 从环境中获取 IDF_PATH 的值
otatool_dir = os.path.join(idf_path, "components", "app_update")  # otatool.py 位于
↪$IDF_PATH/components/app_update 下
```

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```
sys.path.append(otatool_dir)  # 使能 Python 寻找 otatool 模块
from otatool import *  # 导入 otatool 模块内的所有名称
```

要使用 OTA 工具的 Python API，第一步是创建 *OtatoolTarget*：

```
# 创建 partool.py 的目标设备，并将目标设备连接到串行端口 /dev/ttyUSB1
target = OtatoolTarget("/dev/ttyUSB1")
```

现在，可使用创建的 *OtatoolTarget* 在目标设备上完成操作：

```
# 擦除 otadata，将设备复位至工厂应用程序
target.erase_otadata()

# 擦除 OTA 应用程序分区 0
target.erase_ota_partition(0)

# 将启动分区切换至 OTA 应用程序分区 1
target.switch_ota_partition(1)

# 读取 OTA 分区 'ota_3'，将内容保存至文件 'ota_3.bin'
target.read_ota_partition("ota_3", "ota_3.bin")
```

要操作的 OTA 分区通过应用程序分区序号或分区名称指定。

更多关于 Python API 的信息，请查看 OTA 工具的代码注释。

命令行界面 *otatool.py* 的命令行界面具有如下结构：

```
otatool.py [command-args] [subcommand] [subcommand-args]

- command-args - 执行主命令 (otatool.py) 所需的实际参数，多与目标设备有关
- subcommand - 要执行的操作
- subcommand-args - 所选操作的实际参数
```

```
# 擦除 otadata，将设备复位至工厂应用程序
otatool.py --port "/dev/ttyUSB1" erase_otadata

# 擦除 OTA 应用程序分区 0
otatool.py --port "/dev/ttyUSB1" erase_ota_partition --slot 0

# 将启动分区切换至 OTA 应用程序分区 1
otatool.py --port "/dev/ttyUSB1" switch_ota_partition --slot 1

# 读取 OTA 分区 'ota_3'，将内容保存至文件 'ota_3.bin'
otatool.py --port "/dev/ttyUSB1" read_ota_partition --name=ota_3 --output=ota_3.bin
```

更多信息可用 *–help* 指令查看：

```
# 显示可用的子命令和主命令描述
otatool.py --help

# 显示子命令的描述
otatool.py [subcommand] --help
```

相关文档

- 分区表
- *SPI Flash 和分区 API*
- *ESP HTTPS OTA*

应用程序示例

端对端的 OTA 固件升级示例请参考 system/ota。

API 参考

Header File

- app_update/include/esp_ota_ops.h

Functions

const *esp_app_desc_t* ***esp_ota_get_app_description**(void)

Return esp_app_desc structure. This structure includes app version.

Return description for running app.

Return Pointer to esp_app_desc structure.

int **esp_ota_get_app_elf_sha256**(char **dst*, size_t *size*)

Fill the provided buffer with SHA256 of the ELF file, formatted as hexadecimal, null-terminated. If the buffer size is not sufficient to fit the entire SHA256 in hex plus a null terminator, the largest possible number of bytes will be written followed by a null.

Return Number of bytes written to dst (including null terminator)

Parameters

- `dst`: Destination buffer
- `size`: Size of the buffer

esp_err_t **esp_ota_begin**(**const** *esp_partition_t* **partition*, size_t *image_size*, *esp_ota_handle_t* **out_handle*)

Commence an OTA update writing to the specified partition.

The specified partition is erased to the specified image size.

If image size is not yet known, pass OTA_SIZE_UNKNOWN which will cause the entire partition to be erased.

On success, this function allocates memory that remains in use until esp_ota_end() is called with the returned handle.

Note: If the rollback option is enabled and the running application has the ESP_OTA_IMG_PENDING_VERIFY state then it will lead to the ESP_ERR_OTA_ROLLBACK_INVALID_STATE error. Confirm the running app before to run download a new app, use esp_ota_mark_app_valid_cancel_rollback() function for it (this should be done as early as possible when you first download a new application).

Return

- ESP_OK: OTA operation commenced successfully.
- ESP_ERR_INVALID_ARG: partition or out_handle arguments were NULL, or partition doesn't point to an OTA app partition.
- ESP_ERR_NO_MEM: Cannot allocate memory for OTA operation.
- ESP_ERR_OTA_PARTITION_CONFLICT: Partition holds the currently running firmware, cannot update in place.
- ESP_ERR_NOT_FOUND: Partition argument not found in partition table.
- ESP_ERR_OTA_SELECT_INFO_INVALID: The OTA data partition contains invalid data.
- ESP_ERR_INVALID_SIZE: Partition doesn't fit in configured flash size.
- ESP_ERR_FLASH_OP_TIMEOUT or ESP_ERR_FLASH_OP_FAIL: Flash write failed.
- ESP_ERR_OTA_ROLLBACK_INVALID_STATE: If the running app has not confirmed state. Before performing an update, the application must be valid.

Parameters

- `partition`: Pointer to info for partition which will receive the OTA update. Required.
- `image_size`: Size of new OTA app image. Partition will be erased in order to receive this size of image. If 0 or OTA_SIZE_UNKNOWN, the entire partition is erased.

- `out_handle`: On success, returns a handle which should be used for subsequent esp_ota_write() and esp_ota_end() calls.

esp_err_t **esp_ota_write**(*esp_ota_handle_t handle*, **const** void **data*, size_t *size*)

Write OTA update data to partition.

This function can be called multiple times as data is received during the OTA operation. Data is written sequentially to the partition.

Return

- ESP_OK: Data was written to flash successfully.
- ESP_ERR_INVALID_ARG: handle is invalid.
- ESP_ERR_OTA_VALIDATE_FAILED: First byte of image contains invalid app image magic byte.
- ESP_ERR_FLASH_OP_TIMEOUT or ESP_ERR_FLASH_OP_FAIL: Flash write failed.
- ESP_ERR_OTA_SELECT_INFO_INVALID: OTA data partition has invalid contents

Parameters

- `handle`: Handle obtained from esp_ota_begin
- `data`: Data buffer to write
- `size`: Size of data buffer in bytes.

esp_err_t **esp_ota_write_with_offset**(*esp_ota_handle_t handle*, **const** void **data*, size_t *size*, uint32_t *offset*)

Write OTA update data to partition.

This function can write data in non contiguous manner. If flash encryption is enabled, data should be 16 byte aligned.

Note While performing OTA, if the packets arrive out of order, esp_ota_write_with_offset() can be used to write data in non contiguous manner. Use of esp_ota_write_with_offset() in combination with esp_ota_write() is not recommended.

Return

- ESP_OK: Data was written to flash successfully.
- ESP_ERR_INVALID_ARG: handle is invalid.
- ESP_ERR_OTA_VALIDATE_FAILED: First byte of image contains invalid app image magic byte.
- ESP_ERR_FLASH_OP_TIMEOUT or ESP_ERR_FLASH_OP_FAIL: Flash write failed.
- ESP_ERR_OTA_SELECT_INFO_INVALID: OTA data partition has invalid contents

Parameters

- `handle`: Handle obtained from esp_ota_begin
- `data`: Data buffer to write
- `size`: Size of data buffer in bytes
- `offset`: Offset in flash partition

esp_err_t **esp_ota_end**(*esp_ota_handle_t handle*)

Finish OTA update and validate newly written app image.

Note After calling esp_ota_end(), the handle is no longer valid and any memory associated with it is freed (regardless of result).

Return

- ESP_OK: Newly written OTA app image is valid.
- ESP_ERR_NOT_FOUND: OTA handle was not found.
- ESP_ERR_INVALID_ARG: Handle was never written to.
- ESP_ERR_OTA_VALIDATE_FAILED: OTA image is invalid (either not a valid app image, or - if secure boot is enabled - signature failed to verify.)
- ESP_ERR_INVALID_STATE: If flash encryption is enabled, this result indicates an internal error writing the final encrypted bytes to flash.

Parameters

- `handle`: Handle obtained from esp_ota_begin().

esp_err_t **esp_ota_set_boot_partition**(**const** *esp_partition_t *partition*)

Configure OTA data for a new boot partition.

Note If this function returns ESP_OK, calling esp_restart() will boot the newly configured app partition.

Return

- ESP_OK: OTA data updated, next reboot will use specified partition.
- ESP_ERR_INVALID_ARG: partition argument was NULL or didn't point to a valid OTA partition of type "app".
- ESP_ERR_OTA_VALIDATE_FAILED: Partition contained invalid app image. Also returned if secure boot is enabled and signature validation failed.
- ESP_ERR_NOT_FOUND: OTA data partition not found.
- ESP_ERR_FLASH_OP_TIMEOUT or ESP_ERR_FLASH_OP_FAIL: Flash erase or write failed.

Parameters

- `partition`: Pointer to info for partition containing app image to boot.

const *esp_partition_t* ***esp_ota_get_boot_partition**(void)

Get partition info of currently configured boot app.

If esp_ota_set_boot_partition() has been called, the partition which was set by that function will be returned.

If esp_ota_set_boot_partition() has not been called, the result is usually the same as esp_ota_get_running_partition(). The two results are not equal if the configured boot partition does not contain a valid app (meaning that the running partition will be an app that the bootloader chose via fallback).

If the OTA data partition is not present or not valid then the result is the first app partition found in the partition table. In priority order, this means: the factory app, the first OTA app slot, or the test app partition.

Note that there is no guarantee the returned partition is a valid app. Use esp_image_verify(ESP_IMAGE_VERIFY, ···) to verify if the returned partition contains a bootable image.

Return Pointer to info for partition structure, or NULL if partition table is invalid or a flash read operation failed. Any returned pointer is valid for the lifetime of the application.

const *esp_partition_t* ***esp_ota_get_running_partition**(void)

Get partition info of currently running app.

This function is different to esp_ota_get_boot_partition() in that it ignores any change of selected boot partition caused by esp_ota_set_boot_partition(). Only the app whose code is currently running will have its partition information returned.

The partition returned by this function may also differ from esp_ota_get_boot_partition() if the configured boot partition is somehow invalid, and the bootloader fell back to a different app partition at boot.

Return Pointer to info for partition structure, or NULL if no partition is found or flash read operation failed. Returned pointer is valid for the lifetime of the application.

const *esp_partition_t* ***esp_ota_get_next_update_partition**(**const** *esp_partition_t* **start_from*)

Return the next OTA app partition which should be written with a new firmware.

Call this function to find an OTA app partition which can be passed to esp_ota_begin().

Finds next partition round-robin, starting from the current running partition.

Return Pointer to info for partition which should be updated next. NULL result indicates invalid OTA data partition, or that no eligible OTA app slot partition was found.

Parameters

- `start_from`: If set, treat this partition info as describing the current running partition. Can be NULL, in which case esp_ota_get_running_partition() is used to find the currently running partition. The result of this function is never the same as this argument.

esp_err_t **esp_ota_get_partition_description**(**const** *esp_partition_t* **partition*, *esp_app_desc_t* **app_desc*)

Returns esp_app_desc structure for app partition. This structure includes app version.

Returns a description for the requested app partition.

Return

- ESP_OK Successful.
- ESP_ERR_NOT_FOUND app_desc structure is not found. Magic word is incorrect.

- ESP_ERR_NOT_SUPPORTED Partition is not application.
- ESP_ERR_INVALID_ARG Arguments is NULL or if partition’ s offset exceeds partition size.
- ESP_ERR_INVALID_SIZE Read would go out of bounds of the partition.
- or one of error codes from lower-level flash driver.

Parameters

- `[in] partition`: Pointer to app partition. (only app partition)
- `[out] app_desc`: Structure of info about app.

esp_err_t **esp_ota_mark_app_valid_cancel_rollback**(void)

This function is called to indicate that the running app is working well.

Return

- ESP_OK: if successful.

esp_err_t **esp_ota_mark_app_invalid_rollback_and_reboot**(void)

This function is called to roll back to the previously workable app with reboot.

If rollback is successful then device will reset else API will return with error code. Checks applications on a flash drive that can be booted in case of rollback. If the flash does not have at least one app (except the running app) then rollback is not possible.

Return

- ESP_FAIL: if not successful.
- ESP_ERR_OTA_ROLLBACK_FAILED: The rollback is not possible due to flash does not have any apps.

const *esp_partition_t* ***esp_ota_get_last_invalid_partition**(void)

Returns last partition with invalid state (ESP_OTA_IMG_INVALID or ESP_OTA_IMG_ABORTED).

Return partition.

esp_err_t **esp_ota_get_state_partition**(**const** *esp_partition_t* **partition*, esp_ota_img_states_t **ota_state*)

Returns state for given partition.

Return

- ESP_OK: Successful.
- ESP_ERR_INVALID_ARG: partition or ota_state arguments were NULL.
- ESP_ERR_NOT_SUPPORTED: partition is not ota.
- ESP_ERR_NOT_FOUND: Partition table does not have otadata or state was not found for given partition.

Parameters

- `[in] partition`: Pointer to partition.
- `[out] ota_state`: state of partition (if this partition has a record in otadata).

esp_err_t **esp_ota_erase_last_boot_app_partition**(void)

Erase previous boot app partition and corresponding otadata select for this partition.

When current app is marked to as valid then you can erase previous app partition.

Return

- ESP_OK: Successful, otherwise ESP_ERR.

bool **esp_ota_check_rollback_is_possible**(void)

Checks applications on the slots which can be booted in case of rollback.

These applications should be valid (marked in otadata as not UNDEFINED, INVALID or ABORTED and crc is good) and be able booted, and secure_version of app >= secure_version of efuse (if anti-rollback is enabled).

Return

- True: Returns true if the slots have at least one app (except the running app).
- False: The rollback is not possible.

esp_err_t **esp_ota_revoke_secure_boot_public_key**(*esp_ota_secure_boot_public_key_index_t* *index*)

Revokes the old signature digest. To be called in the application after the rollback logic.

Relevant for Secure boot v2 on ESP32-S2 where upto 3 key digests can be stored (Key N-1, Key N, Key N+1). When key N-1 used to sign an app is invalidated, an OTA update is to be sent with an app signed with key N-1 & Key N. After successfully booting the OTA app should call this function to revoke Key N-1.

Return

- ESP_OK: If revocation is successful.
- ESP_ERR_INVALID_ARG: If the index of the public key to be revoked is incorrect.
- ESP_FAIL: If secure boot v2 has not been enabled.

Parameters

- `index`: - The index of the signature block to be revoked

Macros

OTA_SIZE_UNKNOWN

Used for esp_ota_begin() if new image size is unknown

ESP_ERR_OTA_BASE

Base error code for ota_ops api

ESP_ERR_OTA_PARTITION_CONFLICT

Error if request was to write or erase the current running partition

ESP_ERR_OTA_SELECT_INFO_INVALID

Error if OTA data partition contains invalid content

ESP_ERR_OTA_VALIDATE_FAILED

Error if OTA app image is invalid

ESP_ERR_OTA_SMALL_SEC_VER

Error if the firmware has a secure version less than the running firmware.

ESP_ERR_OTA_ROLLBACK_FAILED

Error if flash does not have valid firmware in passive partition and hence rollback is not possible

ESP_ERR_OTA_ROLLBACK_INVALID_STATE

Error if current active firmware is still marked in pending validation state (ESP_OTA_IMG_PENDING_VERIFY), essentially first boot of firmware image post upgrade and hence firmware upgrade is not possible

Type Definitions

typedef uint32_t **esp_ota_handle_t**

Opaque handle for an application OTA update.

esp_ota_begin() returns a handle which is then used for subsequent calls to esp_ota_write() and esp_ota_end().

Enumerations

enum esp_ota_secure_boot_public_key_index_t

Secure Boot V2 public key indexes.

Values:

SECURE_BOOT_PUBLIC_KEY_INDEX_0

Points to the 0th index of the Secure Boot v2 public key

SECURE_BOOT_PUBLIC_KEY_INDEX_1

Points to the 1st index of the Secure Boot v2 public key

SECURE_BOOT_PUBLIC_KEY_INDEX_2

Points to the 2nd index of the Secure Boot v2 public key

2.6.19 Performance Monitor

The Performance Monitor component provides APIs to use ESP32-S2 internal performance counters to profile functions and applications.

Application Example

An example which combines performance monitor is provided in `examples/system/perfmon` directory. This example initializes the performance monitor structure and execute them with printing the statistics.

High level API Reference

Header Files

- perfmon/include/perfmon.h

API Reference

Header File

- perfmon/include/xtensa_perfmon_access.h

Functions

esp_err_t **xtensa_perfmon_init** (int *id*, uint16_t *select*, uint16_t *mask*, int *kernelcnt*, int *tracelevel*)

Init Performance Monitoor.

Initialize performance monitor register with define values

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if one of the arguments is not correct

Parameters
- `[in] id`: performance counter number
- `[in] select`: select value from PMCTRLx register
- `[in] mask`: mask value from PMCTRLx register
- `[in] kernelcnt`: kernelcnt value from PMCTRLx register
- `[in] tracelevel`: tracelevel value from PMCTRLx register

esp_err_t **xtensa_perfmon_reset** (int *id*)

Reset PM counter.

Reset PM counter. Writes 0 to the PMx register.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if id out of range

Parameters
- `[in] id`: performance counter number

void **xtensa_perfmon_start** (void)

Start PM counters.

Start all PM counters synchronously. Write 1 to the PGM register

void **xtensa_perfmon_stop** (void)

Stop PM counters.

Stop all PM counters synchronously. Write 0 to the PGM register

uint32_t **xtensa_perfmon_value**(int *id*)
Read PM counter.

Read value of defined PM counter.

Return
- Performance counter value

Parameters
- `[in] id`: performance counter number

esp_err_t **xtensa_perfmon_overflow**(int *id*)
Read PM overflow state.

Read overflow value of defined PM counter.

Return
- ESP_OK if there is no overflow (overflow = 0)
- ESP_FAIL if overflow occure (overflow = 1)

Parameters
- `[in] id`: performance counter number

void **xtensa_perfmon_dump**(void)
Dump PM values.

Dump all PM register to the console.

Header File

- perfmon/include/xtensa_perfmon_apis.h

Functions

esp_err_t **xtensa_perfmon_exec**(**const** *xtensa_perfmon_config_t* **config*)
Execute PM.

Execute performance counter for dedicated function with defined parameters

Return
- ESP_OK if no errors
- ESP_ERR_INVALID_ARG if one of the required parameters not defined
- ESP_FAIL - counter overflow

Parameters
- `[in] config`: pointer to the configuration structure

void **xtensa_perfmon_view_cb**(void **params*, uint32_t *select*, uint32_t *mask*, uint32_t *value*)
Dump PM results.

Callback to dump perfmon result to a FILE* stream specified in perfmon_config_t::callback_params. If callback_params is set to NULL, will print to stdout

Parameters
- `[in] params`: used parameters passed from configuration (callback_params). This parameter expected as FILE* hanle, where data will be stored. If this parameter NULL, then data will be stored to the stdout.
- `[in] select`: select value for current counter
- `[in] mask`: mask value for current counter
- `[in] value`: counter value for current counter

Structures

struct xtensa_perfmon_config
Performance monitor configuration structure.

Structure to configure performance counter to measure dedicated function

Public Members

int **repeat_count**
how much times function will be called before the calback will be repeated

float **max_deviation**
Difference between min and max counter number 0..1, 0 - no difference, 1 - not used

void ***call_params**
This pointer will be passed to the call_function as a parameter

void (***call_function**) (void *params)
pointer to the function that have to be called

void (***callback**) (void *params, uint32_t select, uint32_t mask, uint32_t value)
pointer to the function that will be called with result parameters

void ***callback_params**
parameter that will be passed to the callback

int **tracelevel**
trace level for all counters. In case of negative value, the filter will be ignored. If it's >=0, then the perfmon will count only when interrupt level > tracelevel. It's useful to monitor interrupts.

uint32_t **counters_size**
amount of counter in the list

const uint32_t ***select_mask**
list of the select/mask parameters

Type Definitions

typedef struct *xtensa_perfmon_config* **xtensa_perfmon_config_t**
Performance monitor configuration structure.

Structure to configure performance counter to measure dedicated function

2.6.20 电源管理

概述

ESP-IDF 中集成的电源管理算法可以根据应用程序组件的需求，调整外围总线 (APB) 频率、CPU 频率，并使芯片进入 Light-sleep 模式，尽可能减少运行应用程序的功耗。

应用程序组件可以通过创建和获取电源管理锁来控制功耗。

例如：

- 对于从 APB 获得时钟频率的外设，其驱动可以要求在使用该外设时，将 APB 频率设置为 80 MHz。
- RTOS 可以要求 CPU 在有任务准备开始运行时以最高配置频率工作。
- 一些外设可能需要中断才能启用，因此其驱动也会要求禁用 Light-sleep 模式。

因为请求较高的 APB 频率或 CPU 频率，以及禁用 Light-sleep 模式会增加功耗，请将组件使用的电源管理锁降到最少。

电源管理配置

编译时可使用*CONFIG_PM_ENABLE* 选项启用电源管理功能。

启用电源管理功能将会增加中断延迟。额外延迟与多个因素有关，例如：CPU 频率、单/双核模式、是否需要进行频率切换等。CPU 频率为 240 MHz 且未启用频率调节时，最小额外延迟为 0.2 us；如果启用频率调节，且在中断入口将频率由 40 MHz 调节至 80 MHz，则最大额外延迟为 40 us。

应用程序可以通过调用*esp_pm_configure()* 函数启用动态调频 (DFS) 功能和自动 Light-sleep 模式。此函数的参数为*esp_pm_config_esp32s2_t*，定义了频率调节的相关设置。在此参数结构中，需要初始化下面三个字段：

- max_freq_mhz：最大 CPU 频率 (MHz)，即获取 ESP_PM_CPU_FREQ_MAX 锁后所使用的频率。该字段通常设置为*CONFIG_ESP32S2_DEFAULT_CPU_FREQ_MHZ*。
- min_freq_mhz：最小 CPU 频率 (MHz)，即仅获取 ESP_PM_APB_FREQ_MAX 锁后所使用的频率。该字段可设置为晶振 (XTAL) 频率值，或者 XTAL 频率值除以整数。注意，10 MHz 是生成 1 MHz 的 REF_TICK 默认时钟所需的最小频率。
- light_sleep_enable：没有获取任何管理锁时，决定系统是否需要自动进入 Light-sleep 状态 (true/false)。

或者，如果在 menuconfig 中启用了*CONFIG_PM_DFS_INIT_AUTO* 选项，最大 CPU 频率将由*CONFIG_ESP32S2_DEFAULT_CPU_FREQ_MHZ* 设置决定，最小 CPU 频率将锁定为 XTAL 频率。

注解:

1. 自动 Light-sleep 模式基于 FreeRTOS Tickless Idle 功能，因此如果在 menuconfig 中没有启用*CONFIG_FREERTOS_USE_TICKLESS_IDLE* 选项，在请求自动 Light-sleep 时，*esp_pm_configure()* 将会返回 *ESP_ERR_NOT_SUPPORTED* 错误。
2. 在 Light-sleep 状态下，外设没有时钟门控，不会产生来自 GPIO 和内部外设的中断。*Sleep Modes* 文档中所提到的唤醒源可用于从 Light-sleep 状态触发唤醒。例如，EXT0 和 EXT1 唤醒源就可以通过 GPIO 唤醒芯片。

电源管理锁

应用程序可以通过获取或释放管理锁来控制电源管理算法。应用程序获取电源管理锁后，电源管理算法的操作将受到下面的限制。释放电源管理锁后，限制解除。

电源管理锁设有获取/释放计数器，如果已多次获取电源管理锁，则需要将电源管理锁释放相同次数以解除限制。

ESP32-S2 支持下表中所述的三种电源管理锁。

电源管理锁	描述
ESP_PM_CPU_FREQ_MAX	请求使用*esp_pm_configure()* 将 CPU 频率设置为最大值。ESP32-S2 可以将该值设置为 80 MHz、160 MHz 或 240 MHz。
ESP_PM_APB_FREQ_MAX	请求将 APB 频率设置为最大值，ESP32-S2 支持的最大频率为 80 MHz。
ESP_PM_NO_LIGHT_SLEEP	禁止自动切换至 Light-sleep 模式。

ESP32-S2 电源管理算法

下表列出了启用动态调频时如何切换 CPU 频率和 APB 频率。您可以使用*esp_pm_configure()* 或者*CONFIG_ESP32S2_DEFAULT_CPU_FREQ_MHZ* 指定 CPU 最大频率。

CPU 最高频率	电源管理锁获取情况	APB 频率和 CPU 频率
240	获取 ESP_PM_CPU_FREQ_MAX	CPU: 240 MHz APB: 80 MHz
	获取 ESP_PM_APB_FREQ_MAX，未获得 ESP_PM_CPU_FREQ_MAX	CPU: 80 MHz APB: 80 MHz
	无	使用 esp_pm_configure() 为二者设置最小值
160	获取 ESP_PM_CPU_FREQ_MAX	CPU: 160 MHz APB: 80 MHz
	获取 ESP_PM_APB_FREQ_MAX，未获得 ESP_PM_CPU_FREQ_MAX	CPU: 80 MHz APB: 80 MHz
	无	使用 esp_pm_configure() 为二者设置最小值
80	获取 ESP_PM_CPU_FREQ_MAX 或 ESP_PM_APB_FREQ_MAX	CPU: 80 MHz APB: 80 MHz
	无	使用 esp_pm_configure() 为二者设置最小值

如果没有获取任何管理锁，调用 *esp_pm_configure()* 将启动 Light-sleep 模式。Light-sleep 模式持续时间由以下因素决定：

- 处于阻塞状态的 FreeRTOS 任务数（有限超时）
- 高分辨率定时器 API 注册的计数器数量

您也可以设置 Light-sleep 模式在最近事件（任务解除阻塞，或计时器超时）之前持续多久才唤醒芯片。

动态调频和外设驱动

启用动态调频后，APB 频率可在一个 RTOS 滴答周期内多次更改。有些外设不受 APB 频率变更的影响，但有些外设可能会出现问题。例如，Timer Group 外设定时器会继续计数，但定时器计数的速度将随 APB 频率的变更而变更。

下面的外设不受 APB 频率变更的影响：

- **UART**：如果 REF_TICK 用作时钟源，则 UART 不受 APB 频率变更影响。请查看 *uart_config_t* 中的 *use_ref_tick*。
- **LEDC**：如果 REF_TICK 用作时钟源，则 LEDC 不受 APB 频率变更影响。请查看 *ledc_timer_config()* 函数。
- **RMT**：如果 REF_TICK 用作时钟源，则 RMT 不受 APB 频率变更影响。请查看 *rmt_config_t* 结构体中的 *flags* 成员以及 *RMT_CHANNEL_FLAGS_ALWAYS_ON* 宏。

目前以下外设驱动程序可感知动态调频，并在调频期间使用 ESP_PM_APB_FREQ_MAX 锁：

- SPI master
- I2C
- I2S（如果 APLL 锁在使用中，I2S 则会启用 ESP_PM_NO_LIGHT_SLEEP 锁）
- SDMMC

启用以下驱动程序时，将占用 ESP_PM_APB_FREQ_MAX 锁：

- **SPI slave**：从调用 *spi_slave_initialize()* 至 *spi_slave_free()* 期间。
- **Ethernet**：从调用 *esp_eth_driver_install()* 至 *esp_eth_driver_uninstall()* 期间。
- **WiFi**：从调用 *esp_wifi_start()* 至 *esp_wifi_stop()* 期间。如果启用了调制解调器睡眠模式，广播关闭时将释放此管理锁。
- **TWAI**：从调用 *twai_driver_install()* 至 *twai_driver_uninstall()* 期间。

以下外设驱动程序无法感知动态调频，应用程序需自己获取/释放管理锁：

- PCNT
- Sigma-delta
- Timer group

API 参考

Header File

- esp_common/include/esp_pm.h

Functions

esp_err_t **esp_pm_configure**(**const** void **config*)

Set implementation-specific power management configuration.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the configuration values are not correct
- ESP_ERR_NOT_SUPPORTED if certain combination of values is not supported, or if CONFIG_PM_ENABLE is not enabled in sdkconfig

Parameters

- `config`: pointer to implementation-specific configuration structure (e.g. esp_pm_config_esp32)

esp_err_t **esp_pm_get_configuration**(void **config*)

Get implementation-specific power management configuration.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the pointer is null

Parameters

- `config`: pointer to implementation-specific configuration structure (e.g. esp_pm_config_esp32)

esp_err_t **esp_pm_lock_create**(*esp_pm_lock_type_t lock_type*, int *arg*, **const** char **name*, *esp_pm_lock_handle_t *out_handle*)

Initialize a lock handle for certain power management parameter.

When lock is created, initially it is not taken. Call esp_pm_lock_acquire to take the lock.

This function must not be called from an ISR.

Return

- ESP_OK on success
- ESP_ERR_NO_MEM if the lock structure can not be allocated
- ESP_ERR_INVALID_ARG if out_handle is NULL or type argument is not valid
- ESP_ERR_NOT_SUPPORTED if CONFIG_PM_ENABLE is not enabled in sdkconfig

Parameters

- `lock_type`: Power management constraint which the lock should control
- `arg`: argument, value depends on lock_type, see esp_pm_lock_type_t

- `name`: arbitrary string identifying the lock (e.g. “wifi” or “spi”). Used by the esp_pm_dump_locks function to list existing locks. May be set to NULL. If not set to NULL, must point to a string which is valid for the lifetime of the lock.
- `[out] out_handle`: handle returned from this function. Use this handle when calling esp_pm_lock_delete, esp_pm_lock_acquire, esp_pm_lock_release. Must not be NULL.

esp_err_t **esp_pm_lock_acquire**(*esp_pm_lock_handle_t handle*)

Take a power management lock.

Once the lock is taken, power management algorithm will not switch to the mode specified in a call to esp_pm_lock_create, or any of the lower power modes (higher numeric values of ‘mode’).

The lock is recursive, in the sense that if esp_pm_lock_acquire is called a number of times, esp_pm_lock_release has to be called the same number of times in order to release the lock.

This function may be called from an ISR.

This function is not thread-safe w.r.t. calls to other esp_pm_lock_* functions for the same handle.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the handle is invalid
- ESP_ERR_NOT_SUPPORTED if CONFIG_PM_ENABLE is not enabled in sdkconfig

Parameters

- `handle`: handle obtained from esp_pm_lock_create function

esp_err_t **esp_pm_lock_release**(*esp_pm_lock_handle_t handle*)

Release the lock taken using esp_pm_lock_acquire.

Call to this functions removes power management restrictions placed when taking the lock.

Locks are recursive, so if esp_pm_lock_acquire is called a number of times, esp_pm_lock_release has to be called the same number of times in order to actually release the lock.

This function may be called from an ISR.

This function is not thread-safe w.r.t. calls to other esp_pm_lock_* functions for the same handle.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the handle is invalid
- ESP_ERR_INVALID_STATE if lock is not acquired
- ESP_ERR_NOT_SUPPORTED if CONFIG_PM_ENABLE is not enabled in sdkconfig

Parameters

- `handle`: handle obtained from esp_pm_lock_create function

esp_err_t **esp_pm_lock_delete**(*esp_pm_lock_handle_t handle*)

Delete a lock created using esp_pm_lock.

The lock must be released before calling this function.

This function must not be called from an ISR.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the handle argument is NULL
- ESP_ERR_INVALID_STATE if the lock is still acquired
- ESP_ERR_NOT_SUPPORTED if CONFIG_PM_ENABLE is not enabled in sdkconfig

Parameters

- `handle`: handle obtained from esp_pm_lock_create function

esp_err_t **esp_pm_dump_locks**(FILE **stream*)

Dump the list of all locks to stderr

This function dumps debugging information about locks created using esp_pm_lock_create to an output stream.

This function must not be called from an ISR. If esp_pm_lock_acquire/release are called while this function is running, inconsistent results may be reported.

Return

- ESP_OK on success
- ESP_ERR_NOT_SUPPORTED if CONFIG_PM_ENABLE is not enabled in sdkconfig

Parameters

- `stream`: stream to print information to; use stdout or stderr to print to the console; use fmemopen/open_memstream to print to a string buffer.

Type Definitions

typedef struct esp_pm_lock ***esp_pm_lock_handle_t**

Opaque handle to the power management lock.

Enumerations

enum esp_pm_lock_type_t

Power management constraints.

Values:

ESP_PM_CPU_FREQ_MAX

Require CPU frequency to be at the maximum value set via esp_pm_configure. Argument is unused and should be set to 0.

ESP_PM_APB_FREQ_MAX

Require APB frequency to be at the maximum value supported by the chip. Argument is unused and should be set to 0.

ESP_PM_NO_LIGHT_SLEEP

Prevent the system from going into light sleep. Argument is unused and should be set to 0.

Header File

- esp32s2/include/esp32s2/pm.h

Structures

struct esp_pm_config_esp32s2_t

Power management config for ESP32.

Pass a pointer to this structure as an argument to esp_pm_configure function.

Public Members

int **max_freq_mhz**

Maximum CPU frequency, in MHz

int **min_freq_mhz**

Minimum CPU frequency to use when no locks are taken, in MHz

bool **light_sleep_enable**

Enter light sleep when no locks are taken

2.6.21 Sleep Modes

Overview

ESP32-S2 is capable of light sleep and deep sleep power saving modes.

In light sleep mode, digital peripherals, most of the RAM, and CPUs are clock-gated, and supply voltage is reduced. Upon exit from light sleep, peripherals and CPUs resume operation, their internal state is preserved.

In deep sleep mode, CPUs, most of the RAM, and all the digital peripherals which are clocked from APB_CLK are powered off. The only parts of the chip which can still be powered on are: RTC controller, RTC peripherals (including ULP coprocessor), and RTC memories (slow and fast).

Wakeup from deep and light sleep modes can be done using several sources. These sources can be combined, in this case the chip will wake up when any one of the sources is triggered. Wakeup sources can be enabled using `esp_sleep_enable_X_wakeup` APIs and can be disabled using *`esp_sleep_disable_wakeup_source()`* API. Next section describes these APIs in detail. Wakeup sources can be configured at any moment before entering light or deep sleep mode.

Additionally, the application can force specific powerdown modes for the RTC peripherals and RTC memories using *`esp_sleep_pd_config()`* API.

Once wakeup sources are configured, application can enter sleep mode using *`esp_light_sleep_start()`* or *`esp_deep_sleep_start()`* APIs. At this point the hardware will be configured according to the requested wakeup sources, and RTC controller will either power down or power off the CPUs and digital peripherals.

WiFi/BT and sleep modes

In deep sleep and light sleep modes, wireless peripherals are powered down. Before entering light sleep modes, applications must disable WiFi and BT using appropriate calls (`esp_bluedroid_disable()`, `esp_bt_controller_disable()`, *`esp_wifi_stop()`*). WiFi and BT connections will not be maintained in deep sleep or light sleep, even if these functions are not called.

If WiFi connection needs to be maintained, enable WiFi modem sleep, and enable automatic light sleep feature (see *Power Management APIs*). This will allow the system to wake up from sleep automatically when required by WiFi driver, thereby maintaining connection to the AP.

Wakeup sources

Timer RTC controller has a built in timer which can be used to wake up the chip after a predefined amount of time. Time is specified at microsecond precision, but the actual resolution depends on the clock source selected for RTC SLOW_CLK. See chapter "Reset and Clock" of the ESP32-S2 Technical Reference Manual for details about RTC clock options.

This wakeup mode doesn't require RTC peripherals or RTC memories to be powered on during sleep.

`esp_sleep_enable_timer_wakeup()` function can be used to enable deep sleep wakeup using a timer.

Touch pad RTC IO module contains logic to trigger wakeup when a touch sensor interrupt occurs. You need to configure the touch pad interrupt before the chip starts deep sleep.

`esp_sleep_enable_touchpad_wakeup()` function can be used to enable this wakeup source.

External wakeup (ext0) RTC IO module contains logic to trigger wakeup when one of RTC GPIOs is set to a predefined logic level. RTC IO is part of RTC peripherals power domain, so RTC peripherals will be kept powered on during deep sleep if this wakeup source is requested.

Because RTC IO module is enabled in this mode, internal pullup or pulldown resistors can also be used. They need to be configured by the application using *`rtc_gpio_pullup_en()`* and *`rtc_gpio_pulldown_en()`* functions, before calling `esp_sleep_start()`.

`esp_sleep_enable_ext0_wakeup()` function can be used to enable this wakeup source.

警告: After wake up from sleep, IO pad used for wakeup will be configured as RTC IO. Before using this pad as digital GPIO, reconfigure it using `rtc_gpio_deinit(gpio_num)` function.

External wakeup (ext1) RTC controller contains logic to trigger wakeup using multiple RTC GPIOs. One of the two logic functions can be used to trigger wakeup:

- wake up if any of the selected pins is high (`ESP_EXT1_WAKEUP_ANY_HIGH`)
- wake up if all the selected pins are low (`ESP_EXT1_WAKEUP_ALL_LOW`)

This wakeup source is implemented by the RTC controller. As such, RTC peripherals and RTC memories can be powered down in this mode. However, if RTC peripherals are powered down, internal pullup and pulldown resistors will be disabled. To use internal pullup or pulldown resistors, request RTC peripherals power domain to be kept on during sleep, and configure pullup/pulldown resistors using `rtc_gpio_` functions, before entering sleep:

```
esp_sleep_pd_config(ESP_PD_DOMAIN_RTC_PERIPH, ESP_PD_OPTION_ON);
gpio_pullup_dis(gpio_num);
gpio_pulldown_en(gpio_num);
```

警告: After wake up from sleep, IO pad(s) used for wakeup will be configured as RTC IO. Before using these pads as digital GPIOs, reconfigure them using `rtc_gpio_deinit(gpio_num)` function.

esp_sleep_enable_ext1_wakeup() function can be used to enable this wakeup source.

ULP coprocessor wakeup ULP coprocessor can run while the chip is in sleep mode, and may be used to poll sensors, monitor ADC or touch sensor values, and wake up the chip when a specific event is detected. ULP coprocessor is part of RTC peripherals power domain, and it runs the program stored in RTC slow memory. RTC slow memory will be powered on during sleep if this wakeup mode is requested. RTC peripherals will be automatically powered on before ULP coprocessor starts running the program; once the program stops running, RTC peripherals are automatically powered down again.

esp_sleep_enable_ulp_wakeup() function can be used to enable this wakeup source.

GPIO wakeup (light sleep only) In addition to EXT0 and EXT1 wakeup sources described above, one more method of wakeup from external inputs is available in light sleep mode. With this wakeup source, each pin can be individually configured to trigger wakeup on high or low level using *gpio_wakeup_enable()* function. Unlike EXT0 and EXT1 wakeup sources, which can only be used with RTC IOs, this wakeup source can be used with any IO (RTC or digital).

esp_sleep_enable_gpio_wakeup() function can be used to enable this wakeup source.

UART wakeup (light sleep only) When ESP32-S2 receives UART input from external devices, it is often required to wake up the chip when input data is available. UART peripheral contains a feature which allows waking up the chip from light sleep when a certain number of positive edges on RX pin are seen. This number of positive edges can be set using *uart_set_wakeup_threshold()* function. Note that the character which triggers wakeup (and any characters before it) will not be received by the UART after wakeup. This means that the external device typically needs to send an extra character to the ESP32-S2 to trigger wakeup, before sending the data.

esp_sleep_enable_uart_wakeup() function can be used to enable this wakeup source.

Power-down of RTC peripherals and memories

By default, *esp_deep_sleep_start()* and *esp_light_sleep_start()* functions will power down all RTC power domains which are not needed by the enabled wakeup sources. To override this behaviour, *esp_sleep_pd_config()* function is provided.

If some variables in the program are placed into RTC slow memory (for example, using `RTC_DATA_ATTR` attribute), RTC slow memory will be kept powered on by default. This can be overridden using *esp_sleep_pd_config()* function, if desired.

Entering light sleep

esp_light_sleep_start() function can be used to enter light sleep once wakeup sources are configured. It is also possible to go into light sleep with no wakeup sources configured, in this case the chip will be in light sleep mode indefinitely, until external reset is applied.

Entering deep sleep

esp_deep_sleep_start() function can be used to enter deep sleep once wakeup sources are configured. It is also possible to go into deep sleep with no wakeup sources configured, in this case the chip will be in deep sleep mode indefinitely, until external reset is applied.

Configuring IOs

Some ESP32-S2 IOs have internal pullups or pulldowns, which are enabled by default. If an external circuit drives this pin in deep sleep mode, current consumption may increase due to current flowing through these pullups and pulldowns.

To isolate a pin, preventing extra current draw, call *rtc_gpio_isolate()* function.

For example, on ESP32-WROVER module, GPIO12 is pulled up externally. GPIO12 also has an internal pull-down in the ESP32 chip. This means that in deep sleep, some current will flow through these external and internal resistors, increasing deep sleep current above the minimal possible value. Add the following code before *esp_deep_sleep_start()* to remove this extra current:

```
rtc_gpio_isolate(GPIO_NUM_12);
```

UART output handling

Before entering sleep mode, *esp_deep_sleep_start()* will flush the contents of UART FIFOs.

When entering light sleep mode using *esp_light_sleep_start()*, UART FIFOs will not be flushed. Instead, UART output will be suspended, and remaining characters in the FIFO will be sent out after wakeup from light sleep.

Checking sleep wakeup cause

esp_sleep_get_wakeup_cause() function can be used to check which wakeup source has triggered wakeup from sleep mode.

For touch pad and ext1 wakeup sources, it is possible to identify pin or touch pad which has caused wakeup using *esp_sleep_get_touchpad_wakeup_status()* and *esp_sleep_get_ext1_wakeup_status()* functions.

Disable sleep wakeup source

Previously configured wakeup source can be disabled later using *esp_sleep_disable_wakeup_source()* API. This function deactivates trigger for the given wakeup source. Additionally it can disable all triggers if the argument is `ESP_SLEEP_WAKEUP_ALL`.

Application Example

Implementation of basic functionality of deep sleep is shown in protocols/sntp example, where ESP module is periodically waken up to retrieve time from NTP server.

More extensive example in system/deep_sleep illustrates usage of various deep sleep wakeup triggers and ULP co-processor programming.

API Reference

Header File

- esp32s2/include/esp_sleep.h

Functions

esp_err_t **esp_sleep_disable_wakeup_source**(*esp_sleep_source_t source*)

Disable wakeup source.

This function is used to deactivate wake up trigger for source defined as parameter of the function.

See docs/sleep-modes.rst for details.

Note This function does not modify wake up configuration in RTC. It will be performed in esp_sleep_start function.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if trigger was not active

Parameters

- `source`: - number of source to disable of type esp_sleep_source_t

esp_err_t **esp_sleep_enable_ulp_wakeup**(void)

Enable wakeup by ULP coprocessor.

Note In revisions 0 and 1 of the ESP32, ULP wakeup source can not be used when RTC_PERIPH power domain is forced to be powered on (ESP_PD_OPTION_ON) or when ext0 wakeup source is used.

Return

- ESP_OK on success
- ESP_ERR_NOT_SUPPORTED if additional current by touch (CONFIG_ESP32_RTC_EXT_CRYST_ADDIT_CURRENT) is enabled.
- ESP_ERR_INVALID_STATE if ULP co-processor is not enabled or if wakeup triggers conflict

esp_err_t **esp_sleep_enable_timer_wakeup**(uint64_t *time_in_us*)

Enable wakeup by timer.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if value is out of range (TBD)

Parameters

- `time_in_us`: time before wakeup, in microseconds

esp_err_t **esp_sleep_enable_touchpad_wakeup**(void)

Enable wakeup by touch sensor.

Note In revisions 0 and 1 of the ESP32, touch wakeup source can not be used when RTC_PERIPH power domain is forced to be powered on (ESP_PD_OPTION_ON) or when ext0 wakeup source is used.

Note The FSM mode of the touch button should be configured as the timer trigger mode.

Return

- ESP_OK on success
- ESP_ERR_NOT_SUPPORTED if additional current by touch (CONFIG_ESP32_RTC_EXT_CRYST_ADDIT_CURRENT) is enabled.
- ESP_ERR_INVALID_STATE if wakeup triggers conflict

touch_pad_t **esp_sleep_get_touchpad_wakeup_status**(void)

Get the touch pad which caused wakeup.

If wakeup was caused by another source, this function will return TOUCH_PAD_MAX;

Return touch pad which caused wakeup

esp_err_t **esp_sleep_enable_ext0_wakeup**(*gpio_num_t gpio_num*, int *level*)

Enable wakeup using a pin.

This function uses external wakeup feature of RTC_IO peripheral. It will work only if RTC peripherals are kept on during sleep.

This feature can monitor any pin which is an RTC IO. Once the pin transitions into the state given by level argument, the chip will be woken up.

Note This function does not modify pin configuration. The pin is configured in esp_sleep_start, immediately before entering sleep mode.

Note In revisions 0 and 1 of the ESP32, ext0 wakeup source can not be used together with touch or ULP wakeup sources.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if the selected GPIO is not an RTC GPIO, or the mode is invalid
- ESP_ERR_INVALID_STATE if wakeup triggers conflict

Parameters

- `gpio_num`: GPIO number used as wakeup source. Only GPIOs which are have RTC functionality can be used: 0,2,4,12-15,25-27,32-39.
- `level`: input level which will trigger wakeup (0=low, 1=high)

esp_err_t **esp_sleep_enable_ext1_wakeup**(uint64_t *mask*, *esp_sleep_ext1_wakeup_mode_t* *mode*)

Enable wakeup using multiple pins.

This function uses external wakeup feature of RTC controller. It will work even if RTC peripherals are shut down during sleep.

This feature can monitor any number of pins which are in RTC IOs. Once any of the selected pins goes into the state given by mode argument, the chip will be woken up.

Note This function does not modify pin configuration. The pins are configured in esp_sleep_start, immediately before entering sleep mode.

Note internal pullups and pulldowns don't work when RTC peripherals are shut down. In this case, external resistors need to be added. Alternatively, RTC peripherals (and pullups/pulldowns) may be kept enabled using esp_sleep_pd_config function.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if any of the selected GPIOs is not an RTC GPIO, or mode is invalid

Parameters

- `mask`: bit mask of GPIO numbers which will cause wakeup. Only GPIOs which are have RTC functionality can be used in this bit map: 0,2,4,12-15,25-27,32-39.
- `mode`: select logic function used to determine wakeup condition:
  - ESP_EXT1_WAKEUP_ALL_LOW: wake up when all selected GPIOs are low
  - ESP_EXT1_WAKEUP_ANY_HIGH: wake up when any of the selected GPIOs is high

esp_err_t **esp_sleep_enable_gpio_wakeup**(void)

Enable wakeup from light sleep using GPIOs.

Each GPIO supports wakeup function, which can be triggered on either low level or high level. Unlike EXT0 and EXT1 wakeup sources, this method can be used both for all IOs: RTC IOs and digital IOs. It can only be used to wakeup from light sleep though.

To enable wakeup, first call gpio_wakeup_enable, specifying gpio number and wakeup level, for each GPIO which is used for wakeup. Then call this function to enable wakeup feature.

Note In revisions 0 and 1 of the ESP32, GPIO wakeup source can not be used together with touch or ULP wakeup sources.

Return

- ESP_OK on success
- ESP_ERR_INVALID_STATE if wakeup triggers conflict

esp_err_t **esp_sleep_enable_uart_wakeup**(int *uart_num*)

Enable wakeup from light sleep using UART.

Use uart_set_wakeup_threshold function to configure UART wakeup threshold.

Wakeup from light sleep takes some time, so not every character sent to the UART can be received by the application.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if wakeup from given UART is not supported

Parameters
- `uart_num`: UART port to wake up from

uint64_t **esp_sleep_get_ext1_wakeup_status**(void)
Get the bit mask of GPIOs which caused wakeup (ext1)

If wakeup was caused by another source, this function will return 0.

Return bit mask, if GPIOn caused wakeup, BIT(n) will be set

esp_err_t **esp_sleep_enable_wifi_wakeup**(void)
Enable wakeup by WiFi MAC.

Return
- ESP_OK on success

esp_err_t **esp_sleep_pd_config**(*esp_sleep_pd_domain_t domain*, *esp_sleep_pd_option_t option*)
Set power down mode for an RTC power domain in sleep mode.

If not set set using this API, all power domains default to ESP_PD_OPTION_AUTO.

Return
- ESP_OK on success
- ESP_ERR_INVALID_ARG if either of the arguments is out of range

Parameters
- `domain`: power domain to configure
- `option`: power down option (ESP_PD_OPTION_OFF, ESP_PD_OPTION_ON, or ESP_PD_OPTION_AUTO)

void **esp_deep_sleep_start**(void)
Enter deep sleep with the configured wakeup options.

This function does not return.

esp_err_t **esp_light_sleep_start**(void)
Enter light sleep with the configured wakeup options.

Return
- ESP_OK on success (returned after wakeup)
- ESP_ERR_INVALID_STATE if WiFi or BT is not stopped

void **esp_deep_sleep**(uint64_t *time_in_us*)
Enter deep-sleep mode.

The device will automatically wake up after the deep-sleep time Upon waking up, the device calls deep sleep wake stub, and then proceeds to load application.

Call to this function is equivalent to a call to esp_deep_sleep_enable_timer_wakeup followed by a call to esp_deep_sleep_start.

esp_deep_sleep does not shut down WiFi, BT, and higher level protocol connections gracefully. Make sure relevant WiFi and BT stack functions are called to close any connections and deinitialize the peripherals. These include:

- esp_bluedroid_disable
- esp_bt_controller_disable
- esp_wifi_stop

This function does not return.

Note The device will wake up immediately if the deep-sleep time is set to 0
Parameters

- `time_in_us`: deep-sleep time, unit: microsecond

esp_sleep_wakeup_cause_t **esp_sleep_get_wakeup_cause** (void)
Get the wakeup source which caused wakeup from sleep.

Return cause of wake up from last sleep (deep sleep or light sleep)

void **esp_wake_deep_sleep** (void)
Default stub to run on wake from deep sleep.

Allows for executing code immediately on wake from sleep, before the software bootloader or ESP-IDF app has started up.

This function is weak-linked, so you can implement your own version to run code immediately when the chip wakes from sleep.

See docs/deep-sleep-stub.rst for details.

void **esp_set_deep_sleep_wake_stub** (*esp_deep_sleep_wake_stub_fn_t new_stub*)
Install a new stub at runtime to run on wake from deep sleep.

If implementing esp_wake_deep_sleep() then it is not necessary to call this function.

However, it is possible to call this function to substitute a different deep sleep stub. Any function used as a deep sleep stub must be marked RTC_IRAM_ATTR, and must obey the same rules given for esp_wake_deep_sleep().

esp_deep_sleep_wake_stub_fn_t **esp_get_deep_sleep_wake_stub** (void)
Get current wake from deep sleep stub.

Return Return current wake from deep sleep stub, or NULL if no stub is installed.

void **esp_default_wake_deep_sleep** (void)
The default esp-idf-provided esp_wake_deep_sleep() stub.

See docs/deep-sleep-stub.rst for details.

Type Definitions

typedef *esp_sleep_source_t* **esp_sleep_wakeup_cause_t**
typedef void (***esp_deep_sleep_wake_stub_fn_t**) (void)
Function type for stub to run on wake from sleep.

Enumerations

enum esp_sleep_ext1_wakeup_mode_t
Logic function used for EXT1 wakeup mode.

Values:

ESP_EXT1_WAKEUP_ALL_LOW = 0
Wake the chip when all selected GPIOs go low.

ESP_EXT1_WAKEUP_ANY_HIGH = 1
Wake the chip when any of the selected GPIOs go high.

enum esp_sleep_pd_domain_t
Power domains which can be powered down in sleep mode.

Values:

ESP_PD_DOMAIN_RTC_PERIPH
RTC IO, sensors and ULP co-processor.

ESP_PD_DOMAIN_RTC_SLOW_MEM
RTC slow memory.

ESP_PD_DOMAIN_RTC_FAST_MEM
RTC fast memory.

ESP_PD_DOMAIN_XTAL
XTAL oscillator.

ESP_PD_DOMAIN_MAX
Number of domains.

enum esp_sleep_pd_option_t
Power down options.

Values:

ESP_PD_OPTION_OFF
Power down the power domain in sleep mode.

ESP_PD_OPTION_ON
Keep power domain enabled during sleep mode.

ESP_PD_OPTION_AUTO
Keep power domain enabled in sleep mode, if it is needed by one of the wakeup options. Otherwise power it down.

enum esp_sleep_source_t
Sleep wakeup cause.

Values:

ESP_SLEEP_WAKEUP_UNDEFINED
In case of deep sleep, reset was not caused by exit from deep sleep.

ESP_SLEEP_WAKEUP_ALL
Not a wakeup cause, used to disable all wakeup sources with esp_sleep_disable_wakeup_source.

ESP_SLEEP_WAKEUP_EXT0
Wakeup caused by external signal using RTC_IO.

ESP_SLEEP_WAKEUP_EXT1
Wakeup caused by external signal using RTC_CNTL.

ESP_SLEEP_WAKEUP_TIMER
Wakeup caused by timer.

ESP_SLEEP_WAKEUP_TOUCHPAD
Wakeup caused by touchpad.

ESP_SLEEP_WAKEUP_ULP
Wakeup caused by ULP program.

ESP_SLEEP_WAKEUP_GPIO
Wakeup caused by GPIO (light sleep only)

ESP_SLEEP_WAKEUP_UART
Wakeup caused by UART (light sleep only)

ESP_SLEEP_WAKEUP_WIFI
Wakeup caused by WIFI (light sleep only)

2.6.22 Watchdogs

Overview

The ESP-IDF has support for two types of watchdogs: The Interrupt Watchdog Timer and the Task Watchdog Timer (TWDT). The Interrupt Watchdog Timer and the TWDT can both be enabled using *Project Configuration Menu*, however the TWDT can also be enabled during runtime. The Interrupt Watchdog is responsible for detecting instances where FreeRTOS task switching is blocked for a prolonged period of time. The TWDT is responsible for detecting instances of tasks running without yielding for a prolonged period.

Interrupt watchdog The interrupt watchdog makes sure the FreeRTOS task switching interrupt isn't blocked for a long time. This is bad because no other tasks, including potentially important ones like the WiFi task and the idle task, can't get any CPU runtime. A blocked task switching interrupt can happen because a program runs into an infinite loop with interrupts disabled or hangs in an interrupt.

The default action of the interrupt watchdog is to invoke the panic handler. causing a register dump and an opportunity for the programmer to find out, using either OpenOCD or gdbstub, what bit of code is stuck with interrupts disabled. Depending on the configuration of the panic handler, it can also blindly reset the CPU, which may be preferred in a production environment.

The interrupt watchdog is built around the hardware watchdog in timer group 1. If this watchdog for some reason cannot execute the NMI handler that invokes the panic handler (e.g. because IRAM is overwritten by garbage), it will hard-reset the SOC. If the panic handler executes, it will display the panic reason as "Interrupt wdt timeout on CPU0" or "Interrupt wdt timeout on CPU1" (as applicable).

Configuration The interrupt watchdog is enabled by default via the *CONFIG_ESP_INT_WDT* configuration flag. The timeout is configured by setting *CONFIG_ESP_INT_WDT_TIMEOUT_MS*. The default timeout is higher if PSRAM support is enabled, as a critical section or interrupt routine that accesses a large amount of PSRAM will take longer to complete in some circumstances. The INT WDT timeout should always be longer than the period between FreeRTOS ticks (see *CONFIG_FREERTOS_HZ*).

Tuning If you find the Interrupt watchdog timeout is triggering because an interrupt or critical section is running longer than the timeout period, consider rewriting the code: critical sections should be made as short as possible, with non-critical computation happening outside the critical section. Interrupt handlers should also perform the minimum possible amount of computation, consider pushing data into a queue from the ISR and processing it in a task instead. Neither critical sections or interrupt handlers should ever block waiting for another event to occur.

If changing the code to reduce the processing time is not possible or desirable, it's possible to increase the *CONFIG_ESP_INT_WDT_TIMEOUT_MS* setting instead.

Task Watchdog Timer The Task Watchdog Timer (TWDT) is responsible for detecting instances of tasks running for a prolonged period of time without yielding. This is a symptom of CPU starvation and is usually caused by a higher priority task looping without yielding to a lower-priority task thus starving the lower priority task from CPU time. This can be an indicator of poorly written code that spinloops on a peripheral, or a task that is stuck in an infinite loop.

By default the TWDT will watch the Idle task, however any task can subscribe to be watched by the TWDT. Each watched task must 'reset' the TWDT periodically to indicate that they have been allocated CPU time. If a task does not reset within the TWDT timeout period, a warning will be printed with information about which tasks failed to reset the TWDT in time and which tasks are currently running.

It is also possible to redefine the function *esp_task_wdt_isr_user_handler* in the user code, in order to receive the timeout event and handle it differently.

The TWDT is built around the Hardware Watchdog Timer in Timer Group 0. The TWDT can be initialized by calling `esp_task_wdt_init()` which will configure the hardware timer. A task can then subscribe to the TWDT using `esp_task_wdt_add()` in order to be watched. Each subscribed task must periodically call `esp_task_wdt_reset()` to reset the TWDT. Failure by any subscribed tasks to periodically call `esp_task_wdt_reset()` indicates that one or more tasks have been starved of CPU time or are stuck in a loop somewhere.

A watched task can be unsubscribed from the TWDT using `esp_task_wdt_delete()`. A task that has been unsubscribed should no longer call `esp_task_wdt_reset()`. Once all tasks have unsubscribed form the TWDT, the TWDT can be deinitialized by calling `esp_task_wdt_deinit()`.

The default timeout period for the TWDT is set using config item *CONFIG_ESP_TASK_WDT_TIMEOUT_S*. This should be set to at least as long as you expect any single task will need to monopolise the CPU (for example, if you expect the app will do a long intensive calculation and should not yield to other tasks). It is also possible to change this timeout at runtime by calling `esp_task_wdt_init()`.

The following config options control TWDT configuration at startup. They are all enabled by default:

- *CONFIG_ESP_TASK_WDT* - the TWDT is initialized automatically during startup. If this option is disabled, it is still possible to initialize the Task WDT at runtime by calling `esp_task_wdt_init()`.
- *CONFIG_ESP_TASK_WDT_CHECK_IDLE_TASK_CPU0* - Idle task is subscribed to the TWDT during startup. If this option is disabled, it is still possible to subscribe the idle task by calling `esp_task_wdt_add()` at any time.

JTAG and watchdogs While debugging using OpenOCD, the CPUs will be halted every time a breakpoint is reached. However if the watchdog timers continue to run when a breakpoint is encountered, they will eventually trigger a reset making it very difficult to debug code. Therefore OpenOCD will disable the hardware timers of both the interrupt and task watchdogs at every breakpoint. Moreover, OpenOCD will not reenable them upon leaving the breakpoint. This means that interrupt watchdog and task watchdog functionality will essentially be disabled. No warnings or panics from either watchdogs will be generated when the ESP32-S2 is connected to OpenOCD via JTAG.

Interrupt Watchdog API Reference

Header File

- esp_common/include/esp_int_wdt.h

Functions

void **esp_int_wdt_init**(void)

Initialize the non-CPU-specific parts of interrupt watchdog. This is called in the init code if the interrupt watchdog is enabled in menuconfig.

Task Watchdog API Reference

A full example using the Task Watchdog is available in esp-idf: system/task_watchdog

Header File

- esp_common/include/esp_task_wdt.h

Functions

esp_err_t **esp_task_wdt_init**(uint32_t *timeout*, bool *panic*)

Initialize the Task Watchdog Timer (TWDT)

This function configures and initializes the TWDT. If the TWDT is already initialized when this function is called, this function will update the TWDT's timeout period and panic configurations instead. After initializing the TWDT, any task can elect to be watched by the TWDT by subscribing to it using esp_task_wdt_add().

Return

- ESP_OK: Initialization was successful
- ESP_ERR_NO_MEM: Initialization failed due to lack of memory

Note esp_task_wdt_init() must only be called after the scheduler started

Parameters

- `[in] timeout`: Timeout period of TWDT in seconds
- `[in] panic`: Flag that controls whether the panic handler will be executed when the TWDT times out

esp_err_t **esp_task_wdt_deinit**(void)

Deinitialize the Task Watchdog Timer (TWDT)

This function will deinitialize the TWDT. Calling this function whilst tasks are still subscribed to the TWDT, or when the TWDT is already deinitialized, will result in an error code being returned.

Return

- ESP_OK: TWDT successfully deinitialized

- ESP_ERR_INVALID_STATE: Error, tasks are still subscribed to the TWDT
- ESP_ERR_NOT_FOUND: Error, TWDT has already been deinitialized

esp_err_t **esp_task_wdt_add**(*TaskHandle_t handle*)

Subscribe a task to the Task Watchdog Timer (TWDT)

This function subscribes a task to the TWDT. Each subscribed task must periodically call esp_task_wdt_reset() to prevent the TWDT from elapsing its timeout period. Failure to do so will result in a TWDT timeout. If the task being subscribed is one of the Idle Tasks, this function will automatically enable esp_task_wdt_reset() to called from the Idle Hook of the Idle Task. Calling this function whilst the TWDT is uninitialized or attempting to subscribe an already subscribed task will result in an error code being returned.

Return

- ESP_OK: Successfully subscribed the task to the TWDT
- ESP_ERR_INVALID_ARG: Error, the task is already subscribed
- ESP_ERR_NO_MEM: Error, could not subscribe the task due to lack of memory
- ESP_ERR_INVALID_STATE: Error, the TWDT has not been initialized yet

Parameters

- `[in] handle`: Handle of the task. Input NULL to subscribe the current running task to the TWDT

esp_err_t **esp_task_wdt_reset**(void)

Reset the Task Watchdog Timer (TWDT) on behalf of the currently running task.

This function will reset the TWDT on behalf of the currently running task. Each subscribed task must periodically call this function to prevent the TWDT from timing out. If one or more subscribed tasks fail to reset the TWDT on their own behalf, a TWDT timeout will occur. If the IDLE tasks have been subscribed to the TWDT, they will automatically call this function from their idle hooks. Calling this function from a task that has not subscribed to the TWDT, or when the TWDT is uninitialized will result in an error code being returned.

Return

- ESP_OK: Successfully reset the TWDT on behalf of the currently running task
- ESP_ERR_NOT_FOUND: Error, the current running task has not subscribed to the TWDT
- ESP_ERR_INVALID_STATE: Error, the TWDT has not been initialized yet

esp_err_t **esp_task_wdt_delete**(*TaskHandle_t handle*)

Unsubscribes a task from the Task Watchdog Timer (TWDT)

This function will unsubscribe a task from the TWDT. After being unsubscribed, the task should no longer call esp_task_wdt_reset(). If the task is an IDLE task, this function will automatically disable the calling of esp_task_wdt_reset() from the Idle Hook. Calling this function whilst the TWDT is uninitialized or attempting to unsubscribe an already unsubscribed task from the TWDT will result in an error code being returned.

Return

- ESP_OK: Successfully unsubscribed the task from the TWDT
- ESP_ERR_INVALID_ARG: Error, the task is already unsubscribed
- ESP_ERR_INVALID_STATE: Error, the TWDT has not been initialized yet

Parameters

- `[in] handle`: Handle of the task. Input NULL to unsubscribe the current running task.

esp_err_t **esp_task_wdt_status**(*TaskHandle_t handle*)

Query whether a task is subscribed to the Task Watchdog Timer (TWDT)

This function will query whether a task is currently subscribed to the TWDT, or whether the TWDT is initialized.

Return :

- ESP_OK: The task is currently subscribed to the TWDT
- ESP_ERR_NOT_FOUND: The task is currently not subscribed to the TWDT
- ESP_ERR_INVALID_STATE: The TWDT is not initialized, therefore no tasks can be subscribed

Parameters

- `[in] handle`: Handle of the task. Input NULL to query the current running task.

2.6.23 System Time

Overview

System time can be kept using either one time source or two time sources simultaneously. The choice depends on the application purpose and accuracy requirements for system time.

There are the following two time sources:

- **RTC timer**: Allows keeping the system time during any resets and sleep modes, only the power-up reset leads to resetting the RTC timer. The frequency deviation depends on an *RTC Clock Source* and affects accuracy only in sleep modes, in which case the time will be measured at 6.6667 us resolution.
- **High-resolution timer**: Not available during any reset and sleep modes. The reason for using this timer is to achieve greater accuracy. It uses the APB_CLK clock source (typically 80 MHz), which has a frequency deviation of less than ±10 ppm. Time will be measured at 1 us resolution.

The settings for the system time source are as follows:

- RTC and high-resolution timer (default)
- RTC
- High-resolution timer
- None

It is recommended to stick to the default setting which provides maximum accuracy. If you want to choose a different timer, configure *CONFIG_ESP32S2_TIME_SYSCALL* in project configuration.

RTC Clock Source

The RTC timer has the following clock sources:

- `Internal 90kHz RC oscillator` (default): Features lowest deep sleep current consumption and no dependence on any external components. However, as frequency stability is affected by temperature fluctuations, time may drift in both Deep and Light sleep modes.
- `External 32kHz crystal`: Requires a 32kHz crystal to be connected to the 32K_XP and 32K_XN pins. Provides better frequency stability at the expense of slightly higher (by 1 uA) Deep sleep current consumption.
- `External 32kHz oscillator at 32K_XN pin`: Allows using 32kHz clock generated by an external circuit. The external clock signal must be connected to the 32K_XN pin. The amplitude should be less than 1.2 V for sine wave signal and less than 1 V for square wave signal. Common mode voltage should be in the range of 0.1 < Vcm < 0.5xVamp, where Vamp is signal amplitude. Additionally, a 1 nF capacitor must be placed between the 32K_XP pin and ground. In this case, the 32K_XP pin cannot be used as a GPIO pin.
- `Internal 8.5MHz oscillator, divided by 256 (~33kHz)`: Provides better frequency stability than the `internal 90kHz RC oscillator` at the expense of higher (by 5 uA) deep sleep current consumption. It also does not require external components.

The choice depends on your requirements for system time accuracy and power consumption in sleep modes. To modify the RTC clock source, set *CONFIG_ESP32S2_RTC_CLK_SRC* in project configuration.

More details on wiring requirements for the `External 32kHz crystal` and `External 32kHz oscillator at 32K_XN pin` sources can be found in Section *Crystal Oscillator* of ESP32-S2 Hardware Design Guidelines.

Get Current Time

To get the current time, use the POSIX function `gettimeofday()`. Additionally, you can use the following standard C library functions to obtain time and manipulate it:

```
gettimeofday
time
asctime
```

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```
clock
ctime
difftime
gmtime
localtime
mktime
strftime
adjtime*
```

* –To stop smooth time adjustment and update the current time immediately, use the POSIX function `settimeofday()`.

If you need to obtain time with one second resolution, use the following method:

```
time_t now;
char strftime_buf[64];
struct tm timeinfo;

time(&now);
// Set timezone to China Standard Time
setenv("TZ", "CST-8", 1);
tzset();

localtime_r(&now, &timeinfo);
strftime(strftime_buf, sizeof(strftime_buf), "%c", &timeinfo);
ESP_LOGI(TAG, "The current date/time in Shanghai is: %s", strftime_buf);
```

If you need to obtain time with one microsecond resolution, use the code snippet below:

```
struct timeval tv_now;
gettimeofday(&tv_now, NULL);
int64_t time_us = (int64_t)tv_now.tv_sec * 1000000L + (int64_t)tv_now.tv_usec;
```

SNTP Time Synchronization

To set the current time, you can use the POSIX functions `settimeofday()` and `adjtime()`. They are used internally in the lwIP SNTP library to set current time when a response from the NTP server is received. These functions can also be used separately from the lwIP SNTP library.

A function to use inside the lwIP SNTP library depends on a sync mode for system time. Use the function *`sntp_set_sync_mode()`* to set one of the following sync modes:

- `SNTP_SYNC_MODE_IMMED` (default) updates system time immediately upon receiving a response from the SNTP server after using `settimeofday()`.
- `SNTP_SYNC_MODE_SMOOTH` updates time smoothly by gradually reducing time error using the funcion `adjtime()`. If the difference between the SNTP response time and system time is more than 35 minutes, update system time immediately by using `settimeofday()`.

The lwIP SNTP library has API functions for setting a callback function for a certain event. You might need the following functions:

- `sntp_set_time_sync_notification_cb()` - use it for setting a callback function that will notify of the time synchronization process
- `sntp_get_sync_status()` and `sntp_set_sync_status()` - use it to get/set time synchronization status

To start synchronization via SNTP, just call the following three functions.

```
sntp_setoperatingmode(SNTP_OPMODE_POLL);
sntp_setservername(0, "pool.ntp.org");
sntp_init();
```

An application with this initialization code will periodically synchronize the time. The time synchronization period is determined by `CONFIG_LWIP_SNTP_UPDATE_DELAY` (default value is one hour). To modify the variable, set *CONFIG_LWIP_SNTP_UPDATE_DELAY* in project configuration.

A code example that demonstrates the implementation of time synchronization based on the lwIP SNTP library is provided in protocols/sntp directory.

Timezones

To set local timezone, use the following POSIX functions:

1. Call `setenv()` to set the `TZ` environment variable to the correct value depending on the device location. The format of the time string is the same as described in the GNU libc documentation (although the implementation is different).
2. Call `tzset()` to update C library runtime data for the new time zone.

Once these steps are completed, call the standard C library function `localtime()`, and it will return correct local time taking into account the time zone offset and daylight saving time.

API Reference

Header File

- lwip/include/apps/sntp/sntp.h

Functions

void **sntp_sync_time**(**struct** timeval *tv)

This function updates the system time.

This is a weak-linked function. It is possible to replace all SNTP update functionality by placing a sntp_sync_time() function in the app firmware source. If the default implementation is used, calling sntp_set_sync_mode() allows the time synchronization mode to be changed to instant or smooth. If a callback function is registered via sntp_set_time_sync_notification_cb(), it will be called following time synchronization.

Parameters

- `tv`: Time received from SNTP server.

void **sntp_set_sync_mode**(*sntp_sync_mode_t sync_mode*)

Set the sync mode.

Allowable two mode: SNTP_SYNC_MODE_IMMED and SNTP_SYNC_MODE_SMOOTH.

Parameters

- `sync_mode`: Sync mode.

sntp_sync_mode_t **sntp_get_sync_mode**(void)

Get set sync mode.

Return SNTP_SYNC_MODE_IMMED: Update time immediately. SNTP_SYNC_MODE_SMOOTH: Smooth time updating.

sntp_sync_status_t **sntp_get_sync_status**(void)

Get status of time sync.

After the update is completed, the status will be returned as SNTP_SYNC_STATUS_COMPLETED. After that, the status will be reset to SNTP_SYNC_STATUS_RESET. If the update operation is not completed yet, the status will be SNTP_SYNC_STATUS_RESET. If a smooth mode was chosen and the synchronization is still continuing (adjtime works), then it will be SNTP_SYNC_STATUS_IN_PROGRESS.

Return SNTP_SYNC_STATUS_RESET: Reset status. SNTP_SYNC_STATUS_COMPLETED: Time is synchronized. SNTP_SYNC_STATUS_IN_PROGRESS: Smooth time sync in progress.

void **sntp_set_sync_status** (*sntp_sync_status_t sync_status*)
Set status of time sync.

Parameters
- `sync_status`: status of time sync (see sntp_sync_status_t)

void **sntp_set_time_sync_notification_cb** (*sntp_sync_time_cb_t callback*)
Set a callback function for time synchronization notification.

Parameters
- `callback`: a callback function

void **sntp_set_sync_interval** (uint32_t *interval_ms*)
Set the sync interval of SNTP operation.

Note: SNTPv4 RFC 4330 enforces a minimum sync interval of 15 seconds. This sync interval will be used in the next attempt update time throught SNTP. To apply the new sync interval call the sntp_restart() function, otherwise, it will be applied after the last interval expired.

Parameters
- `interval_ms`: The sync interval in ms. It cannot be lower than 15 seconds, otherwise 15 seconds will be set.

uint32_t **sntp_get_sync_interval** (void)
Get the sync interval of SNTP operation.

Return the sync interval

bool **sntp_restart** (void)
Restart SNTP.

Return True - Restart False - SNTP was not initialized yet

Type Definitions

typedef void (***sntp_sync_time_cb_t**)(**struct** timeval *tv)
SNTP callback function for notifying about time sync event.

Parameters
- `tv`: Time received from SNTP server.

Enumerations

enum sntp_sync_mode_t
SNTP time update mode.

Values:

SNTP_SYNC_MODE_IMMED
Update system time immediately when receiving a response from the SNTP server.

SNTP_SYNC_MODE_SMOOTH
Smooth time updating. Time error is gradually reduced using adjtime function. If the difference between SNTP response time and system time is large (more than 35 minutes) then update immediately.

enum sntp_sync_status_t
SNTP sync status.

Values:

SNTP_SYNC_STATUS_RESET

SNTP_SYNC_STATUS_COMPLETED

SNTP_SYNC_STATUS_IN_PROGRESS

Code examples for this API section are provided in the system directory of ESP-IDF examples.

2.7 Project Configuration

2.7.1 Introduction

ESP-IDF uses kconfiglib which is a Python-based extension to the Kconfig system which provides a compile-time project configuration mechanism. Kconfig is based around options of several types: integer, string, boolean. Kconfig files specify dependencies between options, default values of the options, the way the options are grouped together, etc.

For the complete list of available features please see Kconfig and kconfiglib extentions.

2.7.2 Project Configuration Menu

Application developers can open a terminal-based project configuration menu with the `idf.py menuconfig` build target.

After being updated, this configuration is saved inside `sdkconfig` file in the project root directory. Based on `sdkconfig`, application build targets will generate `sdkconfig.h` file in the build directory, and will make sdkconfig options available to the project build system and source files.

(For the legacy GNU Make build system, the project configuration menu is opened with `make menuconfig`.)

2.7.3 Using sdkconfig.defaults

In some cases, such as when `sdkconfig` file is under revision control, the fact that `sdkconfig` file gets changed by the build system may be inconvenient. The build system offers a way to avoid this, in the form of `sdkconfig.defaults` file. This file is never touched by the build system, and must be created manually. It can contain all the options which matter for the given application. The format is the same as that of the `sdkconfig` file. Once `sdkconfig.defaults` is created, `sdkconfig` can be deleted and added to the ignore list of the revision control system (e.g. `.gitignore` file for git). Project build targets will automatically create `sdkconfig` file, populated with the settings from `sdkconfig.defaults` file, and the rest of the settings will be set to their default values. Note that the build process will not override settings that are already in `sdkconfig` by ones from `sdkconfig.defaults`. For more information, see 自定义 *sdkconfig* 的默认值.

2.7.4 Kconfig Formatting Rules

The following attributes of `Kconfig` files are standardized:

- Within any menu, option names should have a consistent prefix. The prefix length is currently set to at least 3 characters.
- The indentation style is 4 characters created by spaces. All sub-items belonging to a parent item are indented by one level deeper. For example, `menu` is indented by 0 characters, the `config` inside of the menu by 4 characters, the help of the `config` by 8 characters and the text of the `help` by 12 characters.
- No trailing spaces are allowed at the end of the lines.
- The maximum length of options is set to 40 characters.
- The maximum length of lines is set to 120 characters.
- Lines cannot be wrapped by backslash (because there is a bug in earlier versions of `conf-idf` which causes that Windows line endings are not recognized after a backslash).

Format checker

`tools/check_kconfigs.py` is provided for checking the `Kconfig` formatting rules. The checker checks all `Kconfig` and `Kconfig.projbuild` files in the ESP-IDF directory and generates a new file with suffix `.new` with some recommendations how to fix issues (if there are any). Please note that the checker cannot correct all rules and the responsibility of the developer is to check and make final corrections in order to pass the tests. For

example, indentations will be corrected if there isn't some misleading previous formatting but it cannot come up with a common prefix for options inside a menu.

2.7.5 Backward Compatibility of Kconfig Options

The standard Kconfig tools ignore unknown options in `sdkconfig`. So if a developer has custom settings for options which are renamed in newer ESP-IDF releases then the given setting for the option would be silently ignored. Therefore, several features have been adopted to avoid this:

1. `confgen.py` is used by the tool chain to pre-process `sdkconfig` files before anything else, for example `menuconfig`, would read them. As the consequence, the settings for old options will be kept and not ignored.
2. `confgen.py` recursively finds all `sdkconfig.rename` files in ESP-IDF directory which contain old and new `Kconfig` option names. Old options are replaced by new ones in the `sdkconfig` file.
3. `confgen.py` post-processes `sdkconfig` files and generates all build outputs (`sdkconfig.h`, `sdkconfig.cmake`, `auto.conf`) by adding a list of compatibility statements, i.e. value of the old option is set the value of the new option (after modification). This is done in order to not break customer codes where old option might still be used.
4. *Deprecated options and their replacements* are automatically generated by `confgen.py`.

2.7.6 Configuration Options Reference

Subsequent sections contain the list of available ESP-IDF options, automatically generated from Kconfig files. Note that depending on the options selected, some options listed here may not be visible by default in the interface of menuconfig.

By convention, all option names are upper case with underscores. When Kconfig generates sdkconfig and sdkconfig.h files, option names are prefixed with `CONFIG_`. So if an option `ENABLE_FOO` is defined in a Kconfig file and selected in menuconfig, then sdkconfig and sdkconfig.h files will have `CONFIG_ENABLE_FOO` defined. In this reference, option names are also prefixed with `CONFIG_`, same as in the source code.

SDK tool configuration

Contains:

- *CONFIG_SDK_TOOLPREFIX*
- *CONFIG_SDK_PYTHON*
- *CONFIG_SDK_MAKE_WARN_UNDEFINED_VARIABLES*
- *CONFIG_SDK_TOOLCHAIN_SUPPORTS_TIME_WIDE_64_BITS*

CONFIG_SDK_TOOLPREFIX

Compiler toolchain path/prefix

Found in: SDK tool configuration

The prefix/path that is used to call the toolchain. The default setting assumes a crosstool-ng gcc setup that is in your PATH.

CONFIG_SDK_PYTHON

Python interpreter

Found in: SDK tool configuration

The executable name/path that is used to run python.

(Note: This option is used with the legacy GNU Make build system only.)

CONFIG_SDK_MAKE_WARN_UNDEFINED_VARIABLES

'make' warns on undefined variables

Found in: SDK tool configuration

Adds –warn-undefined-variables to MAKEFLAGS. This causes make to print a warning any time an undefined variable is referenced.

This option helps find places where a variable reference is misspelled or otherwise missing, but it can be unwanted if you have Makefiles which depend on undefined variables expanding to an empty string.

(Note: this option is used with the legacy GNU Make build system only.)

CONFIG_SDK_TOOLCHAIN_SUPPORTS_TIME_WIDE_64_BITS

Toolchain supports time_t wide 64-bits

Found in: SDK tool configuration

Enable this option in case you have a custom toolchain which supports time_t wide 64-bits. This option checks time_t is 64-bits and disables ROM time functions to use the time functions from the toolchain instead. This option allows resolving the Y2K38 problem. See "Setup Linux Toolchain from Scratch" to build a custom toolchain which supports 64-bits time_t.

Note: ESP-IDF does not currently come with any pre-compiled toolchain that supports 64-bit wide time_t. This will change in a future major release, but currently 64-bit time_t requires a custom built toolchain.

Build type

Contains:

- *CONFIG_APP_BUILD_TYPE*

CONFIG_APP_BUILD_TYPE

Application build type

Found in: Build type

Select the way the application is built.

By default, the application is built as a binary file in a format compatible with the ESP32 bootloader. In addition to this application, 2nd stage bootloader is also built. Application and bootloader binaries can be written into flash and loaded/executed from there.

Another option, useful for only very small and limited applications, is to only link the .elf file of the application, such that it can be loaded directly into RAM over JTAG. Note that since IRAM and DRAM sizes are very limited, it is not possible to build any complex application this way. However for kinds of testing and debugging, this option may provide faster iterations, since the application does not need to be written into flash. Note that at the moment, ESP-IDF does not contain all the startup code required to initialize the CPUs and ROM memory (data/bss). Therefore it is necessary to execute a bit of ROM code prior to executing the application. A gdbinit file may look as follows:

```
# Connect to a running instance of OpenOCD target remote :3333 # Reset and halt the target
mon reset halt # Run to a specific point in ROM code, # where most of initialization is
complete. thb *0x40007d54 c # Load the application into RAM load # Run till app_main tb
app_main c
```

Execute this gdbinit file as follows:

```
xtensa-esp32-elf-gdb build/app-name.elf -x gdbinit
```

Recommended sdkconfig.defaults for building loadable ELF files is as follows. CONFIG_APP_BUILD_TYPE_ELF_RAM is required, other options help reduce application memory footprint.

CONFIG_APP_BUILD_TYPE_ELF_RAM=y CONFIG_VFS_SUPPORT_TERMIOS=
CONFIG_NEWLIB_NANO_FORMAT=y CONFIG_ESP_SYSTEM_PANIC_PRINT_HALT=y
CONFIG_ESP_DEBUG_STUBS_ENABLE= CONFIG_ESP_ERR_TO_NAME_LOOKUP=

Available options:

- Default (binary application + 2nd stage bootloader) (APP_BUILD_TYPE_APP_2NDBOOT)
- ELF file, loadable into RAM (EXPERIMENTAL)) (APP_BUILD_TYPE_ELF_RAM)

Partition Table

Contains:

- *CONFIG_PARTITION_TABLE_TYPE*
- *CONFIG_PARTITION_TABLE_CUSTOM_FILENAME*
- *CONFIG_PARTITION_TABLE_OFFSET*
- *CONFIG_PARTITION_TABLE_MD5*

CONFIG_PARTITION_TABLE_TYPE

Partition Table

Found in: Partition Table

The partition table to flash to the ESP32. The partition table determines where apps, data and other resources are expected to be found.

The predefined partition table CSV descriptions can be found in the components/partition_table directory. Otherwise it's possible to create a new custom partition CSV for your application.

Available options:

- Single factory app, no OTA (PARTITION_TABLE_SINGLE_APP)
- Factory app, two OTA definitions (PARTITION_TABLE_TWO_OTA)
- Custom partition table CSV (PARTITION_TABLE_CUSTOM)

CONFIG_PARTITION_TABLE_CUSTOM_FILENAME

Custom partition CSV file

Found in: Partition Table

Name of the custom partition CSV filename. This path is evaluated relative to the project root directory.

CONFIG_PARTITION_TABLE_OFFSET

Offset of partition table

Found in: Partition Table

The address of partition table (by default 0x8000). Allows you to move the partition table, it gives more space for the bootloader. Note that the bootloader and app will both need to be compiled with the same PARTITION_TABLE_OFFSET value.

This number should be a multiple of 0x1000.

Note that partition offsets in the partition table CSV file may need to be changed if this value is set to a higher value. To have each partition offset adapt to the configured partition table offset, leave all partition offsets blank in the CSV file.

CONFIG_PARTITION_TABLE_MD5

Generate an MD5 checksum for the partition table

Found in: Partition Table

Generate an MD5 checksum for the partition table for protecting the integrity of the table. The generation should be turned off for legacy bootloaders which cannot recognize the MD5 checksum in the partition table.

Serial flasher config

Contains:

- *CONFIG_ESPTOOLPY_PORT*
- *CONFIG_ESPTOOLPY_BAUD*
- *CONFIG_ESPTOOLPY_BAUD_OTHER_VAL*
- *CONFIG_ESPTOOLPY_COMPRESSED*
- *CONFIG_ESPTOOLPY_FLASHMODE*
- *CONFIG_ESPTOOLPY_FLASHFREQ*
- *CONFIG_ESPTOOLPY_FLASHSIZE*
- *CONFIG_ESPTOOLPY_FLASHSIZE_DETECT*
- *CONFIG_ESPTOOLPY_BEFORE*
- *CONFIG_ESPTOOLPY_AFTER*
- *CONFIG_ESPTOOLPY_MONITOR_BAUD*
- *CONFIG_ESPTOOLPY_MONITOR_BAUD_OTHER_VAL*

CONFIG_ESPTOOLPY_PORT

Default serial port

Found in: Serial flasher config

The serial port that’s connected to the ESP chip. This can be overridden by setting the ESPPORT environment variable.

This value is ignored when using the CMake-based build system or idf.py.

CONFIG_ESPTOOLPY_BAUD

Default baud rate

Found in: Serial flasher config

Default baud rate to use while communicating with the ESP chip. Can be overridden by setting the ESPBAUD variable.

This value is ignored when using the CMake-based build system or idf.py.

Available options:

- 115200 baud (ESPTOOLPY_BAUD_115200B)
- 230400 baud (ESPTOOLPY_BAUD_230400B)
- 921600 baud (ESPTOOLPY_BAUD_921600B)
- 2Mbaud (ESPTOOLPY_BAUD_2MB)
- Other baud rate (ESPTOOLPY_BAUD_OTHER)

CONFIG_ESPTOOLPY_BAUD_OTHER_VAL

Other baud rate value

Found in: Serial flasher config

CONFIG_ESPTOOLPY_COMPRESSED

Use compressed upload

Found in: Serial flasher config

The flasher tool can send data compressed using zlib, letting the ROM on the ESP chip decompress it on the fly before flashing it. For most payloads, this should result in a speed increase.

CONFIG_ESPTOOLPY_FLASHMODE

Flash SPI mode

Found in: Serial flasher config

Mode the flash chip is flashed in, as well as the default mode for the binary to run in.

Available options:

- QIO (ESPTOOLPY_FLASHMODE_QIO)
- QOUT (ESPTOOLPY_FLASHMODE_QOUT)
- DIO (ESPTOOLPY_FLASHMODE_DIO)
- DOUT (ESPTOOLPY_FLASHMODE_DOUT)

CONFIG_ESPTOOLPY_FLASHFREQ

Flash SPI speed

Found in: Serial flasher config

The SPI flash frequency to be used.

Available options:

- 80 MHz (ESPTOOLPY_FLASHFREQ_80M)
- 40 MHz (ESPTOOLPY_FLASHFREQ_40M)
- 26 MHz (ESPTOOLPY_FLASHFREQ_26M)
- 20 MHz (ESPTOOLPY_FLASHFREQ_20M)

CONFIG_ESPTOOLPY_FLASHSIZE

Flash size

Found in: Serial flasher config

SPI flash size, in megabytes

Available options:

- 1 MB (ESPTOOLPY_FLASHSIZE_1MB)
- 2 MB (ESPTOOLPY_FLASHSIZE_2MB)
- 4 MB (ESPTOOLPY_FLASHSIZE_4MB)
- 8 MB (ESPTOOLPY_FLASHSIZE_8MB)
- 16 MB (ESPTOOLPY_FLASHSIZE_16MB)

CONFIG_ESPTOOLPY_FLASHSIZE_DETECT

Detect flash size when flashing bootloader

Found in: Serial flasher config

If this option is set, flashing the project will automatically detect the flash size of the target chip and update the bootloader image before it is flashed.

CONFIG_ESPTOOLPY_BEFORE

Before flashing

Found in: Serial flasher config

Configure whether esptool.py should reset the ESP32 before flashing.

Automatic resetting depends on the RTS & DTR signals being wired from the serial port to the ESP32. Most USB development boards do this internally.

Available options:

- Reset to bootloader (ESPTOOLPY_BEFORE_RESET)
- No reset (ESPTOOLPY_BEFORE_NORESET)

CONFIG_ESPTOOLPY_AFTER

After flashing

Found in: Serial flasher config

Configure whether esptool.py should reset the ESP32 after flashing.

Automatic resetting depends on the RTS & DTR signals being wired from the serial port to the ESP32. Most USB development boards do this internally.

Available options:

- Reset after flashing (ESPTOOLPY_AFTER_RESET)
- Stay in bootloader (ESPTOOLPY_AFTER_NORESET)

CONFIG_ESPTOOLPY_MONITOR_BAUD

‘idf.py monitor’ baud rate

Found in: Serial flasher config

Baud rate to use when running ‘idf.py monitor’ or ‘make monitor’ to view serial output from a running chip.

Can override by setting the MONITORBAUD environment variable.

Available options:

- 9600 bps (ESPTOOLPY_MONITOR_BAUD_9600B)
- 57600 bps (ESPTOOLPY_MONITOR_BAUD_57600B)
- 115200 bps (ESPTOOLPY_MONITOR_BAUD_115200B)
- 230400 bps (ESPTOOLPY_MONITOR_BAUD_230400B)
- 921600 bps (ESPTOOLPY_MONITOR_BAUD_921600B)
- 2 Mbps (ESPTOOLPY_MONITOR_BAUD_2MB)
- Custom baud rate (ESPTOOLPY_MONITOR_BAUD_OTHER)

CONFIG_ESPTOOLPY_MONITOR_BAUD_OTHER_VAL

Custom baud rate value

Found in: Serial flasher config

Bootloader config

Contains:

- *CONFIG_BOOTLOADER_COMPILER_OPTIMIZATION*
- *CONFIG_BOOTLOADER_LOG_LEVEL*
- *CONFIG_BOOTLOADER_VDDSDIO_BOOST*
- *CONFIG_BOOTLOADER_FACTORY_RESET*

- *CONFIG_BOOTLOADER_APP_TEST*
- *CONFIG_BOOTLOADER_HOLD_TIME_GPIO*
- *CONFIG_BOOTLOADER_WDT_ENABLE*
- *CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE*
- *CONFIG_BOOTLOADER_SKIP_VALIDATE_IN_DEEP_SLEEP*
- *CONFIG_BOOTLOADER_CUSTOM_RESERVE_RTC*
- *CONFIG_BOOTLOADER_FLASH_XMC_SUPPORT*

CONFIG_BOOTLOADER_COMPILER_OPTIMIZATION

Bootloader optimization Level

Found in: Bootloader config

This option sets compiler optimization level (gcc -O argument) for the bootloader.

- The default "Size" setting will add the -0s flag to CFLAGS.
- The "Debug" setting will add the -Og flag to CFLAGS.
- The "Performance" setting will add the -O2 flag to CFLAGS.
- The "None" setting will add the -O0 flag to CFLAGS.

Note that custom optimization levels may be unsupported.

Available options:

- Size (-Os) (BOOTLOADER_COMPILER_OPTIMIZATION_SIZE)
- Debug (-Og) (BOOTLOADER_COMPILER_OPTIMIZATION_DEBUG)
- Optimize for performance (-O2) (BOOTLOADER_COMPILER_OPTIMIZATION_PERF)
- Debug without optimization (-O0) (BOOTLOADER_COMPILER_OPTIMIZATION_NONE)

CONFIG_BOOTLOADER_LOG_LEVEL

Bootloader log verbosity

Found in: Bootloader config

Specify how much output to see in bootloader logs.

Available options:

- No output (BOOTLOADER_LOG_LEVEL_NONE)
- Error (BOOTLOADER_LOG_LEVEL_ERROR)
- Warning (BOOTLOADER_LOG_LEVEL_WARN)
- Info (BOOTLOADER_LOG_LEVEL_INFO)
- Debug (BOOTLOADER_LOG_LEVEL_DEBUG)
- Verbose (BOOTLOADER_LOG_LEVEL_VERBOSE)

CONFIG_BOOTLOADER_VDDSDIO_BOOST

VDDSDIO LDO voltage

Found in: Bootloader config

If this option is enabled, and VDDSDIO LDO is set to 1.8V (using eFuse or MTDI bootstrapping pin), bootloader will change LDO settings to output 1.9V instead. This helps prevent flash chip from browning out during flash programming operations.

This option has no effect if VDDSDIO is set to 3.3V, or if the internal VDDSDIO regulator is disabled via eFuse.

Available options:

- 1.8V (BOOTLOADER_VDDSDIO_BOOST_1_8V)
- 1.9V (BOOTLOADER_VDDSDIO_BOOST_1_9V)

CONFIG_BOOTLOADER_FACTORY_RESET

GPIO triggers factory reset

Found in: Bootloader config

Allows to reset the device to factory settings: - clear one or more data partitions; - boot from "factory" partition. The factory reset will occur if there is a GPIO input pulled low while device starts up. See settings below.

CONFIG_BOOTLOADER_NUM_PIN_FACTORY_RESET

Number of the GPIO input for factory reset

Found in: Bootloader config > CONFIG_BOOTLOADER_FACTORY_RESET

The selected GPIO will be configured as an input with internal pull-up enabled. To trigger a factory reset, this GPIO must be pulled low on reset. Note that GPIO34-39 do not have an internal pullup and an external one must be provided.

CONFIG_BOOTLOADER_OTA_DATA_ERASE

Clear OTA data on factory reset (select factory partition)

Found in: Bootloader config > CONFIG_BOOTLOADER_FACTORY_RESET

The device will boot from "factory" partition (or OTA slot 0 if no factory partition is present) after a factory reset.

CONFIG_BOOTLOADER_DATA_FACTORY_RESET

Comma-separated names of partitions to clear on factory reset

Found in: Bootloader config > CONFIG_BOOTLOADER_FACTORY_RESET

Allows customers to select which data partitions will be erased while factory reset.

Specify the names of partitions as a comma-delimited with optional spaces for readability. (Like this: "nvs, phy_init, ···") Make sure that the name specified in the partition table and here are the same. Partitions of type "app" cannot be specified here.

CONFIG_BOOTLOADER_APP_TEST

GPIO triggers boot from test app partition

Found in: Bootloader config

Allows to run the test app from "TEST" partition. A boot from "test" partition will occur if there is a GPIO input pulled low while device starts up. See settings below.

CONFIG_BOOTLOADER_NUM_PIN_APP_TEST

Number of the GPIO input to boot TEST partition

Found in: Bootloader config > CONFIG_BOOTLOADER_APP_TEST

The selected GPIO will be configured as an input with internal pull-up enabled. To trigger a test app, this GPIO must be pulled low on reset. After the GPIO input is deactivated and the device reboots, the old application will boot. (factory or OTA[x]). Note that GPIO34-39 do not have an internal pullup and an external one must be provided.

CONFIG_BOOTLOADER_HOLD_TIME_GPIO

Hold time of GPIO for reset/test mode (seconds)

Found in: Bootloader config

The GPIO must be held low continuously for this period of time after reset before a factory reset or test partition boot (as applicable) is performed.

CONFIG_BOOTLOADER_WDT_ENABLE

Use RTC watchdog in start code

Found in: Bootloader config

Tracks the execution time of startup code. If the execution time is exceeded, the RTC_WDT will restart system. It is also useful to prevent a lock up in start code caused by an unstable power source. NOTE: Tracks the execution time starts from the bootloader code - re-set timeout, while selecting the source for slow_clk - and ends calling app_main. Re-set timeout is needed due to WDT uses a SLOW_CLK clock source. After changing a frequency slow_clk a time of WDT needs to re-set for new frequency. slow_clk depends on ESP32_RTC_CLK_SRC (INTERNAL_RC or EXTERNAL_CRYSTAL).

CONFIG_BOOTLOADER_WDT_DISABLE_IN_USER_CODE

Allows RTC watchdog disable in user code

Found in: Bootloader config > CONFIG_BOOTLOADER_WDT_ENABLE

If it is set, the client must itself reset or disable rtc_wdt in their code (app_main()). Otherwise rtc_wdt will be disabled before calling app_main function. Use function rtc_wdt_feed() for resetting counter of rtc_wdt. Use function rtc_wdt_disable() for disabling rtc_wdt.

CONFIG_BOOTLOADER_WDT_TIME_MS

Timeout for RTC watchdog (ms)

Found in: Bootloader config > CONFIG_BOOTLOADER_WDT_ENABLE

Verify that this parameter is correct and more then the execution time. Pay attention to options such as reset to factory, trigger test partition and encryption on boot - these options can increase the execution time. Note: RTC_WDT will reset while encryption operations will be performed.

CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE

Enable app rollback support

Found in: Bootloader config

After updating the app, the bootloader runs a new app with the "ESP_OTA_IMG_PENDING_VERIFY" state set. This state prevents the re-run of this app. After the first boot of the new app in the user code, the function should be called to confirm the operability of the app or vice versa about its non-operability. If the app is working, then it is marked as valid. Otherwise, it is marked as not valid and rolls back to the previous working app. A reboot is performed, and the app is booted before the software update. Note: If during the first boot a new app the power goes out or the WDT works, then roll back will happen. Rollback is possible only between the apps with the same security versions.

CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK

Enable app anti-rollback support

Found in: Bootloader config > CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE

This option prevents rollback to previous firmware/application image with lower security version.

CONFIG_BOOTLOADER_APP_SECURE_VERSION

eFuse secure version of app

Found in: Bootloader config > CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE > CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK

The secure version is the sequence number stored in the header of each firmware. The security version is set in the bootloader, version is recorded in the eFuse field as the number of set ones. The allocated number of bits in the efuse field for storing the security version is limited (see BOOTLOADER_APP_SEC_VER_SIZE_EFUSE_FIELD option).

Bootloader: When bootloader selects an app to boot, an app is selected that has a security version greater or equal that recorded in eFuse field. The app is booted with a higher (or equal) secure version.

The security version is worth increasing if in previous versions there is a significant vulnerability and their use is not acceptable.

Your partition table should has a scheme with ota_0 + ota_1 (without factory).

CONFIG_BOOTLOADER_APP_SEC_VER_SIZE_EFUSE_FIELD

Size of the efuse secure version field

Found in: Bootloader config > CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE > CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK

The size of the efuse secure version field. Its length is limited to 32 bits for ESP32 and 16 bits for ESP32-S2. This determines how many times the security version can be increased.

CONFIG_BOOTLOADER_EFUSE_SECURE_VERSION_EMULATE

Emulate operations with efuse secure version(only test)

Found in: Bootloader config > CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE > CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK

This option allow emulate read/write operations with efuse secure version. It allow to test anti-rollback implemention without permanent write eFuse bits. In partition table should be exist this partition *emul_efuse, data, 5, , 0x2000*.

CONFIG_BOOTLOADER_SKIP_VALIDATE_IN_DEEP_SLEEP

Skip image validation when exiting deep sleep

Found in: Bootloader config

This option disables the normal validation of an image coming out of deep sleep (checksums, SHA256, and signature). This is a trade-off between wakeup performance from deep sleep, and image integrity checks.

Only enable this if you know what you are doing. It should not be used in conjunction with using deep_sleep() entry and changing the active OTA partition as this would skip the validation upon first load of the new OTA partition.

CONFIG_BOOTLOADER_CUSTOM_RESERVE_RTC

Reserve RTC FAST memory for custom purposes

Found in: Bootloader config

This option allows the customer to place data in the RTC FAST memory, this area remains valid when rebooted, except for power loss. This memory is located at a fixed address and is available for both the bootloader and the application. (The application and bootoloader must be compiled with the same option). The RTC FAST memory has access only through PRO_CPU.

CONFIG_BOOTLOADER_CUSTOM_RESERVE_RTC_SIZE

Size in bytes for custom purposes

Found in: Bootloader config > CONFIG_BOOTLOADER_CUSTOM_RESERVE_RTC

This option reserves in RTC FAST memory the area for custom purposes. If you want to create your own bootloader and save more information in this area of memory, you can increase it. It must be a multiple of 4 bytes. This area (rtc_retain_mem_t) is reserved and has access from the bootloader and an application.

CONFIG_BOOTLOADER_FLASH_XMC_SUPPORT

Enable the support for flash chips of XMC (READ HELP FIRST)

Found in: Bootloader config

Perform the startup flow recommended by XMC. Please consult XMC for the details of this flow. XMC chips will be forbidden to be used, when this option is disabled.

DON’T DISABLE THIS UNLESS YOU KNOW WHAT YOU ARE DOING.

Security features

Contains:

- *CONFIG_SECURE_SIGNED_APPS_NO_SECURE_BOOT*
- *CONFIG_SECURE_SIGNED_APPS_SCHEME*
- *CONFIG_SECURE_SIGNED_ON_BOOT_NO_SECURE_BOOT*
- *CONFIG_SECURE_SIGNED_ON_UPDATE_NO_SECURE_BOOT*
- *CONFIG_SECURE_BOOT*
- *CONFIG_SECURE_BOOT_BUILD_SIGNED_BINARIES*
- *CONFIG_SECURE_BOOT_VERIFICATION_KEY*
- *CONFIG_SECURE_BOOT_INSECURE*
- *CONFIG_SECURE_FLASH_ENC_ENABLED*
- *Potentially insecure options*
- *CONFIG_SECURE_DISABLE_ROM_DL_MODE*
- *CONFIG_SECURE_ENABLE_SECURE_ROM_DL_MODE*

CONFIG_SECURE_SIGNED_APPS_NO_SECURE_BOOT

Require signed app images

Found in: Security features

Require apps to be signed to verify their integrity.

This option uses the same app signature scheme as hardware secure boot, but unlike hardware secure boot it does not prevent the bootloader from being physically updated. This means that the device can be secured against remote network access, but not physical access. Compared to using hardware Secure Boot this option is much simpler to implement.

CONFIG_SECURE_SIGNED_APPS_SCHEME

App Signing Scheme

Found in: Security features

Select the Secure App signing scheme. Depends on the Chip Revision. There are two options: 1. ECDSA based secure boot scheme. (Only choice for Secure Boot V1) Supported in ESP32 and ESP32-ECO3. 2. The RSA based secure boot scheme. (Only choice for Secure Boot V2) Supported in ESP32-ECO3. (ESP32 Chip Revision 3 onwards)

Available options:

- ECDSA (SECURE_SIGNED_APPS_ECDSA_SCHEME)
  Embeds the ECDSA public key in the bootloader and signs the application with an ECDSA key.
  Refer to the documentation before enabling.
- RSA (SECURE_SIGNED_APPS_RSA_SCHEME)
  Appends the RSA-3072 based Signature block to the application. Refer to <Secure Boot Version 2 documentation link> before enabling.

CONFIG_SECURE_SIGNED_ON_BOOT_NO_SECURE_BOOT

Bootloader verifies app signatures

Found in: Security features

If this option is set, the bootloader will be compiled with code to verify that an app is signed before booting it.

If hardware secure boot is enabled, this option is always enabled and cannot be disabled. If hardware secure boot is not enabled, this option doesn't add significant security by itself so most users will want to leave it disabled.

CONFIG_SECURE_SIGNED_ON_UPDATE_NO_SECURE_BOOT

Verify app signature on update

Found in: Security features

If this option is set, any OTA updated apps will have the signature verified before being considered valid.

When enabled, the signature is automatically checked whenever the esp_ota_ops.h APIs are used for OTA updates, or esp_image_format.h APIs are used to verify apps.

If hardware secure boot is enabled, this option is always enabled and cannot be disabled. If hardware secure boot is not enabled, this option still adds significant security against network-based attackers by preventing spoofing of OTA updates.

CONFIG_SECURE_BOOT

Enable hardware Secure Boot in bootloader (READ DOCS FIRST)

Found in: Security features

Build a bootloader which enables Secure Boot on first boot.

Once enabled, Secure Boot will not boot a modified bootloader. The bootloader will only load a partition table or boot an app if the data has a verified digital signature. There are implications for reflashing updated apps once secure boot is enabled.

When enabling secure boot, JTAG and ROM BASIC Interpreter are permanently disabled by default.

CONFIG_SECURE_BOOT_VERSION

Select secure boot version

Found in: Security features > CONFIG_SECURE_BOOT

Select the Secure Boot Version. Depends on the Chip Revision. Secure Boot V2 is the new RSA based secure boot scheme. Supported in ESP32-ECO3. (ESP32 Chip Revision 3 onwards) Secure Boot V1 is the AES based secure boot scheme. Supported in ESP32 and ESP32-ECO3.

Available options:

- Enable Secure Boot version 1 (SECURE_BOOT_V1_ENABLED)
  Build a bootloader which enables secure boot version 1 on first boot. Refer to the Secure Boot section of the ESP-IDF Programmer's Guide for this version before enabling.
- Enable Secure Boot version 2 (SECURE_BOOT_V2_ENABLED)
  Build a bootloader which enables Secure Boot version 2 on first boot. Refer to Secure Boot V2 section of the ESP-IDF Programmer's Guide for this version before enabling.

CONFIG_SECURE_BOOT_BUILD_SIGNED_BINARIES

Sign binaries during build

Found in: Security features

Once secure boot or signed app requirement is enabled, app images are required to be signed.

If enabled (default), these binary files are signed as part of the build process. The file named in "Secure boot private signing key" will be used to sign the image.

If disabled, unsigned app/partition data will be built. They must be signed manually using espsecure.py. Version 1 to enable ECDSA Based Secure Boot and Version 2 to enable RSA based Secure Boot. (for example, on a remote signing server.)

CONFIG_SECURE_BOOT_SIGNING_KEY

Secure boot private signing key

Found in: Security features > CONFIG_SECURE_BOOT_BUILD_SIGNED_BINARIES

Path to the key file used to sign app images.

Key file is an ECDSA private key (NIST256p curve) in PEM format for Secure Boot V1. Key file is an RSA private key in PEM format for Secure Boot V2.

Path is evaluated relative to the project directory.

You can generate a new signing key by running the following command: espsecure.py generate_signing_key secure_boot_signing_key.pem

See the Secure Boot section of the ESP-IDF Programmer's Guide for this version for details.

CONFIG_SECURE_BOOT_VERIFICATION_KEY

Secure boot public signature verification key

Found in: Security features

Path to a public key file used to verify signed images. Secure Boot V1: This ECDSA public key is compiled into the bootloader and/or app, to verify app images. Secure Boot V2: This RSA public key is compiled into the signature block at the end of the bootloader/app.

Key file is in raw binary format, and can be extracted from a PEM formatted private key using the espsecure.py extract_public_key command.

Refer to the Secure Boot section of the ESP-IDF Programmer's Guide for this version before enabling.

CONFIG_SECURE_BOOT_INSECURE

Allow potentially insecure options

Found in: Security features

You can disable some of the default protections offered by secure boot, in order to enable testing or a custom combination of security features.

Only enable these options if you are very sure.

Refer to the Secure Boot section of the ESP-IDF Programmer's Guide for this version before enabling.

CONFIG_SECURE_FLASH_ENC_ENABLED

Enable flash encryption on boot (READ DOCS FIRST)

Found in: Security features

If this option is set, flash contents will be encrypted by the bootloader on first boot.

Note: After first boot, the system will be permanently encrypted. Re-flashing an encrypted system is complicated and not always possible.

Read *Flash 加密* before enabling.

CONFIG_SECURE_FLASH_ENCRYPTION_KEYSIZE

Size of generated AES-XTS key

Found in: Security features > CONFIG_SECURE_FLASH_ENC_ENABLED

Size of generated AES-XTS key.

AES-128 uses a 256-bit key (32 bytes) which occupies one Efuse key block. AES-256 uses a 512-bit key (64 bytes) which occupies two Efuse key blocks.

This setting is ignored if either type of key is already burned to Efuse before the first boot. In this case, the pre-burned key is used and no new key is generated.

Available options:

- AES-128 (256-bit key) (SECURE_FLASH_ENCRYPTION_AES128)
- AES-256 (512-bit key) (SECURE_FLASH_ENCRYPTION_AES256)

CONFIG_SECURE_FLASH_ENCRYPTION_MODE

Enable usage mode

Found in: Security features > CONFIG_SECURE_FLASH_ENC_ENABLED

By default Development mode is enabled which allows UART bootloader to perform flash encryption operations

Select Release mode only for production or manufacturing. Once enabled you can not reflash using UART bootloader

Refer to the Secure Boot section of the ESP-IDF Programmer's Guide for this version and *Flash 加密* for details.

Available options:

- Development(NOT SECURE) (SECURE_FLASH_ENCRYPTION_MODE_DEVELOPMENT)
- Release (SECURE_FLASH_ENCRYPTION_MODE_RELEASE)

Potentially insecure options Contains:

- *CONFIG_SECURE_BOOT_ALLOW_JTAG*
- *CONFIG_SECURE_BOOT_ALLOW_SHORT_APP_PARTITION*
- *CONFIG_SECURE_BOOT_V2_ALLOW_EFUSE_RD_DIS*
- *CONFIG_SECURE_INSECURE_ALLOW_DL_MODE*
- *CONFIG_SECURE_FLASH_UART_BOOTLOADER_ALLOW_ENC*
- *CONFIG_SECURE_FLASH_UART_BOOTLOADER_ALLOW_CACHE*
- *CONFIG_SECURE_FLASH_REQUIRE_ALREADY_ENABLED*

CONFIG_SECURE_BOOT_ALLOW_JTAG

Allow JTAG Debugging

Found in: Security features > Potentially insecure options

If not set (default), the bootloader will permanently disable JTAG (across entire chip) on first boot when either secure boot or flash encryption is enabled.

Setting this option leaves JTAG on for debugging, which negates all protections of flash encryption and some of the protections of secure boot.

Only set this option in testing environments.

CONFIG_SECURE_BOOT_ALLOW_SHORT_APP_PARTITION

Allow app partition length not 64KB aligned

Found in: Security features > Potentially insecure options

If not set (default), app partition size must be a multiple of 64KB. App images are padded to 64KB length, and the bootloader checks any trailing bytes after the signature (before the next 64KB boundary) have not been written. This is because flash cache maps entire 64KB pages into the address space. This prevents an attacker from appending unverified data after the app image in the flash, causing it to be mapped into the address space.

Setting this option allows the app partition length to be unaligned, and disables padding of the app image to this length. It is generally not recommended to set this option, unless you have a legacy partitioning scheme which doesn't support 64KB aligned partition lengths.

CONFIG_SECURE_BOOT_V2_ALLOW_EFUSE_RD_DIS

Allow additional read protecting of efuses

Found in: Security features > Potentially insecure options

If not set (default, recommended), on first boot the bootloader will burn the WR_DIS_RD_DIS efuse when Secure Boot is enabled. This prevents any more efuses from being read protected.

If this option is set, it will remain possible to write the EFUSE_RD_DIS efuse field after Secure Boot is enabled. This may allow an attacker to read-protect the BLK2 efuse (for ESP32) and BLOCK4-BLOCK10 (i.e. BLOCK_KEY0-BLOCK_KEY5)(for other chips) holding the public key digest, causing an immediate denial of service and possibly allowing an additional fault injection attack to bypass the signature protection.

NOTE: Once a BLOCK is read-protected, the application will read all zeros from that block

NOTE: If "UART ROM download mode (Permanently disabled (recommended))" or "UART ROM download mode (Permanently switch to Secure mode (recommended))" is set, then it is __NOT__ possible to read/write efuses using espefuse.py utility. However, efuse can be read/written from the application

CONFIG_SECURE_INSECURE_ALLOW_DL_MODE

Don't automatically restrict UART download mode

Found in: Security features > Potentially insecure options

By default, enabling either flash encryption in release mode or secure boot will automatically disable UART download mode on ESP32 ECO3, or enable secure download mode on newer chips. This is recommended to reduce the attack surface of the chip.

To allow the full UART download mode to stay enabled, enable this option and ensure the options SECURE_DISABLE_ROM_DL_MODE and SECURE_ENABLE_SECURE_ROM_DL_MODE are disabled as applicable. This is not recommended.

CONFIG_SECURE_FLASH_UART_BOOTLOADER_ALLOW_ENC

Leave UART bootloader encryption enabled

Found in: Security features > Potentially insecure options

If not set (default), the bootloader will permanently disable UART bootloader encryption access on first boot. If set, the UART bootloader will still be able to access hardware encryption.

It is recommended to only set this option in testing environments.

CONFIG_SECURE_FLASH_UART_BOOTLOADER_ALLOW_CACHE

Leave UART bootloader flash cache enabled

Found in: Security features > Potentially insecure options

If not set (default), the bootloader will permanently disable UART bootloader flash cache access on first boot. If set, the UART bootloader will still be able to access the flash cache.

Only set this option in testing environments.

CONFIG_SECURE_FLASH_REQUIRE_ALREADY_ENABLED

Require flash encryption to be already enabled

Found in: Security features > Potentially insecure options

If not set (default), and flash encryption is not yet enabled in eFuses, the 2nd stage bootloader will enable flash encryption: generate the flash encryption key and program eFuses. If this option is set, and flash encryption is not yet enabled, the bootloader will error out and reboot. If flash encryption is enabled in eFuses, this option does not change the bootloader behavior.

Only use this option in testing environments, to avoid accidentally enabling flash encryption on the wrong device. The device needs to have flash encryption already enabled using espefuse.py.

CONFIG_SECURE_DISABLE_ROM_DL_MODE

Permanently disable ROM Download Mode

Found in: Security features

If set, during startup the app will burn an eFuse bit to permanently disable the UART ROM Download Mode. This prevents any future use of esptool.py, espefuse.py and similar tools.

Once disabled, if the SoC is booted with strapping pins set for ROM Download Mode then an error is printed instead.

It is recommended to enable this option in any production application where Flash Encryption and/or Secure Boot is enabled and access to Download Mode is not required.

It is also possible to permanently disable Download Mode by calling esp_efuse_disable_rom_download_mode() at runtime.

CONFIG_SECURE_ENABLE_SECURE_ROM_DL_MODE

Permanently switch to ROM UART Secure Download mode

Found in: Security features

If set, during startup the app will burn an eFuse bit to permanently switch the UART ROM Download Mode into a separate Secure Download mode. This option can only work if Download Mode is not already disabled by eFuse.

Secure Download mode limits the use of Download Mode functions to simple flash read, write and erase operations, plus a command to return a summary of currently enabled security features.

Secure Download mode is not compatible with the esptool.py flasher stub feature, espefuse.py, read/writing memory or registers, encrypted download, or any other features that interact with unsupported Download Mode commands.

Secure Download mode should be enabled in any application where Flash Encryption and/or Secure Boot is enabled. Disabling this option does not immediately cancel the benefits of the security features, but it increases the potential "attack surface" for an attacker to try and bypass them with a successful physical attack.

It is also possible to enable secure download mode at runtime by calling esp_efuse_enable_rom_secure_download_mode()

Application manager

Contains:

- *CONFIG_APP_COMPILE_TIME_DATE*
- *CONFIG_APP_EXCLUDE_PROJECT_VER_VAR*
- *CONFIG_APP_EXCLUDE_PROJECT_NAME_VAR*
- *CONFIG_APP_PROJECT_VER_FROM_CONFIG*
- *CONFIG_APP_RETRIEVE_LEN_ELF_SHA*

CONFIG_APP_COMPILE_TIME_DATE

Use time/date stamp for app

Found in: Application manager

If set, then the app will be built with the current time/date stamp. It is stored in the app description structure. If not set, time/date stamp will be excluded from app image. This can be useful for getting the same binary image files made from the same source, but at different times.

CONFIG_APP_EXCLUDE_PROJECT_VER_VAR

Exclude PROJECT_VER from firmware image

Found in: Application manager

The PROJECT_VER variable from the build system will not affect the firmware image. This value will not be contained in the esp_app_desc structure.

CONFIG_APP_EXCLUDE_PROJECT_NAME_VAR

Exclude PROJECT_NAME from firmware image

Found in: Application manager

The PROJECT_NAME variable from the build system will not affect the firmware image. This value will not be contained in the esp_app_desc structure.

CONFIG_APP_PROJECT_VER_FROM_CONFIG

Get the project version from Kconfig

Found in: Application manager

If this is enabled, then config item APP_PROJECT_VER will be used for the variable PROJECT_VER. Other ways to set PROJECT_VER will be ignored.

CONFIG_APP_PROJECT_VER

Project version

Found in: Application manager > CONFIG_APP_PROJECT_VER_FROM_CONFIG

Project version

CONFIG_APP_RETRIEVE_LEN_ELF_SHA

The length of APP ELF SHA is stored in RAM(chars)

Found in: Application manager

At startup, the app will read this many hex characters from the embedded APP ELF SHA-256 hash value and store it in static RAM. This ensures the app ELF SHA-256 value is always available if it needs to be printed by the panic handler code. Changing this value will change the size of a static buffer, in bytes.

Compiler options

Contains:

- *CONFIG_COMPILER_OPTIMIZATION*
- *CONFIG_COMPILER_OPTIMIZATION_ASSERTION_LEVEL*
- *CONFIG_COMPILER_CXX_EXCEPTIONS*
- *CONFIG_COMPILER_CXX_RTTI*
- *CONFIG_COMPILER_STACK_CHECK_MODE*
- *CONFIG_COMPILER_WARN_WRITE_STRINGS*
- *CONFIG_COMPILER_DISABLE_GCC8_WARNINGS*

CONFIG_COMPILER_OPTIMIZATION

Optimization Level

Found in: Compiler options

This option sets compiler optimization level (gcc -O argument) for the app.

- The "Default" setting will add the -0g flag to CFLAGS.
- The "Size" setting will add the -0s flag to CFLAGS.
- The "Performance" setting will add the -O2 flag to CFLAGS.
- The "None" setting will add the -O0 flag to CFLAGS.

The "Size" setting cause the compiled code to be smaller and faster, but may lead to difficulties of correlating code addresses to source file lines when debugging.

The "Performance" setting causes the compiled code to be larger and faster, but will be easier to correlated code addresses to source file lines.

"None" with -O0 produces compiled code without optimization.

Note that custom optimization levels may be unsupported.

Compiler optimization for the IDF bootloader is set separately, see the BOOTLOADER_COMPILER_OPTIMIZATION setting.

Available options:

- Debug (-Og) (COMPILER_OPTIMIZATION_DEFAULT)
- Optimize for size (-Os) (COMPILER_OPTIMIZATION_SIZE)
- Optimize for performance (-O2) (COMPILER_OPTIMIZATION_PERF)
- Debug without optimization (-O0) (COMPILER_OPTIMIZATION_NONE)

CONFIG_COMPILER_OPTIMIZATION_ASSERTION_LEVEL

Assertion level

Found in: Compiler options

Assertions can be:

- Enabled. Failure will print verbose assertion details. This is the default.
- Set to "silent" to save code size (failed assertions will abort() but user needs to use the aborting address to find the line number with the failed assertion.)
- Disabled entirely (not recommended for most configurations.) -DNDEBUG is added to CPPFLAGS in this case.

Available options:

- Enabled (COMPILER_OPTIMIZATION_ASSERTIONS_ENABLE)
  Enable assertions. Assertion content and line number will be printed on failure.
- Silent (saves code size) (COMPILER_OPTIMIZATION_ASSERTIONS_SILENT)
  Enable silent assertions. Failed assertions will abort(), user needs to use the aborting address to find the line number with the failed assertion.
- Disabled (sets -DNDEBUG) (COMPILER_OPTIMIZATION_ASSERTIONS_DISABLE)
  If assertions are disabled, -DNDEBUG is added to CPPFLAGS.

CONFIG_COMPILER_CXX_EXCEPTIONS

Enable C++ exceptions

Found in: Compiler options

Enabling this option compiles all IDF C++ files with exception support enabled.

Disabling this option disables C++ exception support in all compiled files, and any libstdc++ code which throws an exception will abort instead.

Enabling this option currently adds an additional ~500 bytes of heap overhead when an exception is thrown in user code for the first time.

Contains:

- *CONFIG_COMPILER_CXX_EXCEPTIONS_EMG_POOL_SIZE*

CONFIG_COMPILER_CXX_EXCEPTIONS_EMG_POOL_SIZE

Emergency Pool Size

Found in: Compiler options > CONFIG_COMPILER_CXX_EXCEPTIONS

Size (in bytes) of the emergency memory pool for C++ exceptions. This pool will be used to allocate memory for thrown exceptions when there is not enough memory on the heap.

CONFIG_COMPILER_CXX_RTTI

Enable C++ run-time type info (RTTI)

Found in: Compiler options

Enabling this option compiles all C++ files with RTTI support enabled. This increases binary size (typically by tens of kB) but allows using dynamic_cast conversion and typeid operator.

CONFIG_COMPILER_STACK_CHECK_MODE

Stack smashing protection mode

Found in: Compiler options

Stack smashing protection mode. Emit extra code to check for buffer overflows, such as stack smashing attacks. This is done by adding a guard variable to functions with vulnerable objects. The guards are initialized when a function is entered and then checked when the function exits. If a guard check fails, program is halted. Protection has the following modes:

- In NORMAL mode (GCC flag: -fstack-protector) only functions that call alloca, and functions with buffers larger than 8 bytes are protected.
- STRONG mode (GCC flag: -fstack-protector-strong) is like NORMAL, but includes additional functions to be protected –those that have local array definitions, or have references to local frame addresses.
- In OVERALL mode (GCC flag: -fstack-protector-all) all functions are protected.

Modes have the following impact on code performance and coverage:

- performance: NORMAL > STRONG > OVERALL
- coverage: NORMAL < STRONG < OVERALL

Available options:

- None (COMPILER_STACK_CHECK_MODE_NONE)
- Normal (COMPILER_STACK_CHECK_MODE_NORM)
- Strong (COMPILER_STACK_CHECK_MODE_STRONG)
- Overall (COMPILER_STACK_CHECK_MODE_ALL)

CONFIG_COMPILER_WARN_WRITE_STRINGS

Enable -Wwrite-strings warning flag

Found in: Compiler options

Adds -Wwrite-strings flag for the C/C++ compilers.

For C, this gives string constants the type `const char[]` so that copying the address of one into a non-const `char \*` pointer produces a warning. This warning helps to find at compile time code that tries to write into a string constant.

For C++, this warns about the deprecated conversion from string literals to `char \*`.

CONFIG_COMPILER_DISABLE_GCC8_WARNINGS

Disable new warnings introduced in GCC 6 - 8

Found in: Compiler options

Enable this option if using GCC 6 or newer, and wanting to disable warnings which don't appear with GCC 5.

Component config

Contains:

- *TinyUSB*
- *OpenSSL*
- *mDNS*
- *libsodium*
- *jsmn*
- *Modbus configuration*
- *ADC-Calibration*
- *ESP32S2-specific*
- *Power Management*
- *CoAP Configuration*
- *Supplicant*
- *Wi-Fi Provisioning Manager*

- *Wear Levelling*
- *Virtual file system*
- *Unity unit testing library*
- *SPIFFS Configuration*
- *SPI Flash driver*
- *PThreads*
- *NVS*
- *Newlib*
- *ESP-MQTT Configurations*
- *mbedTLS*
- *LWIP*
- *Log output*
- *Heap memory debugging*
- *FreeRTOS*
- *FAT Filesystem support*
- *Core dump*
- *Wi-Fi*
- *PHY*
- *High resolution timer (esp_timer)*
- *ESP System Settings*
- *ESP NETIF Adapter*
- *ESP HTTPS server*
- *ESP HTTPS OTA*
- *HTTP Server*
- *ESP HTTP client*
- *GDB Stub*
- *Event Loop Library*
- *Ethernet*
- *Common ESP-related*
- *ESP-TLS*
- *eFuse Bit Manager*
- *Driver configurations*
- *Application Level Tracing*

TinyUSB Contains:

- *CONFIG_USB_ENABLED*
- *Descriptor configuration*

CONFIG_USB_ENABLED

Enable TinyUSB driver

Found in: Component config > TinyUSB

Adds support for TinyUSB

CONFIG_USB_DEBUG

Debug mode

Found in: Component config > TinyUSB > CONFIG_USB_ENABLED

Debug mode

Descriptor configuration Contains:

- *CONFIG_USB_DESC_USE_ESPRESSIF_VID*
- *CONFIG_USB_DESC_CUSTOM_VID*

- *CONFIG_USB_DESC_USE_DEFAULT_PID*
- *CONFIG_USB_DESC_CUSTOM_PID*
- *CONFIG_USB_DESC_BCDDEVICE*
- *CONFIG_USB_DESC_MANUFACTURER_STRING*
- *CONFIG_USB_DESC_PRODUCT_STRING*
- *CONFIG_USB_DESC_SERIAL_STRING*

CONFIG_USB_DESC_USE_ESPRESSIF_VID

VID: Use an Espressif ' s default value

Found in: Component config > TinyUSB > Descriptor configuration

Long description

CONFIG_USB_DESC_CUSTOM_VID

Custom VID value

Found in: Component config > TinyUSB > Descriptor configuration

Custom Vendor ID

CONFIG_USB_DESC_USE_DEFAULT_PID

PID: Use a default PID assigning

Found in: Component config > TinyUSB > Descriptor configuration

Default TinyUSB PID assigning uses values 0x4000···0x4007

CONFIG_USB_DESC_CUSTOM_PID

Custom PID value

Found in: Component config > TinyUSB > Descriptor configuration

Custom Product ID

CONFIG_USB_DESC_BCDDEVICE

bcdDevice

Found in: Component config > TinyUSB > Descriptor configuration

Version of the firmware of the USB device

CONFIG_USB_DESC_MANUFACTURER_STRING

Manufacturer

Found in: Component config > TinyUSB > Descriptor configuration

Name of the manufacturer of the USB device

CONFIG_USB_DESC_PRODUCT_STRING

Product

Found in: Component config > TinyUSB > Descriptor configuration

Name of the USB device

CONFIG_USB_DESC_SERIAL_STRING

Serial string

Found in: Component config > TinyUSB > Descriptor configuration

Specify serial number of the USB device

OpenSSL Contains:

- *CONFIG_OPENSSL_DEBUG*
- *CONFIG_OPENSSL_ASSERT*

CONFIG_OPENSSL_DEBUG

Enable OpenSSL debugging

Found in: Component config > OpenSSL

Enable OpenSSL debugging function.

If the option is enabled, "SSL_DEBUG" works.

CONFIG_OPENSSL_DEBUG_LEVEL

OpenSSL debugging level

Found in: Component config > OpenSSL > CONFIG_OPENSSL_DEBUG

OpenSSL debugging level.

Only function whose debugging level is higher than "OPENSSL_DEBUG_LEVEL" works.

For example: If OPENSSL_DEBUG_LEVEL = 2, you use function "SSL_DEBUG(1, "malloc failed")". Because 1 < 2, it will not print.

CONFIG_OPENSSL_LOWLEVEL_DEBUG

Enable OpenSSL low-level module debugging

Found in: Component config > OpenSSL > CONFIG_OPENSSL_DEBUG

If the option is enabled, low-level module debugging function of OpenSSL is enabled, e.g. mbedtls internal debugging function.

CONFIG_OPENSSL_ASSERT

Select OpenSSL assert function

Found in: Component config > OpenSSL

OpenSSL function needs "assert" function to check if input parameters are valid.

If you want to use assert debugging function, "OPENSSL_DEBUG" should be enabled.

Available options:

- Do nothing (OPENSSL_ASSERT_DO_NOTHING)
  Do nothing and "SSL_ASSERT" does not work.
- Check and exit (OPENSSL_ASSERT_EXIT)
  Enable assert exiting, it will check and return error code.
- Show debugging message (OPENSSL_ASSERT_DEBUG)
  Enable assert debugging, it will check and show debugging message.
- Show debugging message and exit (OPENSSL_ASSERT_DEBUG_EXIT)
  Enable assert debugging and exiting, it will check, show debugging message and return error code.

- Show debugging message and block (OPENSSL_ASSERT_DEBUG_BLOCK)
  Enable assert debugging and blocking, it will check, show debugging message and block by "while (1);" .

mDNS Contains:

- *CONFIG_MDNS_MAX_SERVICES*
- *CONFIG_MDNS_TASK_PRIORITY*
- *CONFIG_MDNS_TASK_STACK_SIZE*
- *CONFIG_MDNS_TASK_AFFINITY*
- *CONFIG_MDNS_SERVICE_ADD_TIMEOUT_MS*
- *CONFIG_MDNS_STRICT_MODE*
- *CONFIG_MDNS_TIMER_PERIOD_MS*

CONFIG_MDNS_MAX_SERVICES

Max number of services

Found in: Component config > mDNS

Services take up a certain amount of memory, and allowing fewer services to be open at the same time conserves memory. Specify the maximum amount of services here. The valid value is from 1 to 64.

CONFIG_MDNS_TASK_PRIORITY

mDNS task priority

Found in: Component config > mDNS

Allows setting mDNS task priority. Please do not set the task priority higher than priorities of system tasks. Compile time warning/error would be emitted if the chosen task priority were too high.

CONFIG_MDNS_TASK_STACK_SIZE

mDNS task stack size

Found in: Component config > mDNS

Allows setting mDNS task stacksize.

CONFIG_MDNS_TASK_AFFINITY

mDNS task affinity

Found in: Component config > mDNS

Allows setting mDNS tasks affinity, i.e. whether the task is pinned to CPU0, pinned to CPU1, or allowed to run on any CPU.

Available options:

- No affinity (MDNS_TASK_AFFINITY_NO_AFFINITY)
- CPU0 (MDNS_TASK_AFFINITY_CPU0)
- CPU1 (MDNS_TASK_AFFINITY_CPU1)

CONFIG_MDNS_SERVICE_ADD_TIMEOUT_MS

mDNS adding service timeout (ms)

Found in: Component config > mDNS

Configures timeout for adding a new mDNS service. Adding a service fails if could not be completed within this time.

CONFIG_MDNS_STRICT_MODE

mDNS strict mode

Found in: Component config > mDNS

Configures strict mode. Set this to 1 for the mDNS library to strictly follow the RFC6762: Currently the only strict feature: Do not repeat original questions in response packets (defined in RFC6762 sec. 6). Default configuration is 0, i.e. non-strict mode, since some implementations, such as lwIP mdns resolver (used by standard POSIX API like getaddrinfo, gethostbyname) could not correctly resolve advertised names.

CONFIG_MDNS_TIMER_PERIOD_MS

mDNS timer period (ms)

Found in: Component config > mDNS

Configures period of mDNS timer, which periodically transmits packets and schedules mDNS searches.

libsodium Contains:

- *CONFIG_LIBSODIUM_USE_MBEDTLS_SHA*

CONFIG_LIBSODIUM_USE_MBEDTLS_SHA

Use mbedTLS SHA256 & SHA512 implementations

Found in: Component config > libsodium

If this option is enabled, libsodium will use thin wrappers around mbedTLS for SHA256 & SHA512 operations.

This saves some code size if mbedTLS is also used. However it is incompatible with hardware SHA acceleration (due to the way libsodium's API manages SHA state).

jsmn Contains:

- *CONFIG_JSMN_PARENT_LINKS*
- *CONFIG_JSMN_STRICT*

CONFIG_JSMN_PARENT_LINKS

Enable parent links

Found in: Component config > jsmn

You can access to parent node of parsed json

CONFIG_JSMN_STRICT

Enable strict mode

Found in: Component config > jsmn

In strict mode primitives are: numbers and booleans

Modbus configuration Contains:

- *CONFIG_FMB_COMM_MODE_RTU_EN*
- *CONFIG_FMB_COMM_MODE_ASCII_EN*
- *CONFIG_FMB_MASTER_TIMEOUT_MS_RESPOND*
- *CONFIG_FMB_MASTER_DELAY_MS_CONVERT*
- *CONFIG_FMB_QUEUE_LENGTH*
- *CONFIG_FMB_SERIAL_TASK_STACK_SIZE*
- *CONFIG_FMB_SERIAL_BUF_SIZE*
- *CONFIG_FMB_SERIAL_ASCII_BITS_PER_SYMB*
- *CONFIG_FMB_SERIAL_ASCII_TIMEOUT_RESPOND_MS*
- *CONFIG_FMB_SERIAL_TASK_PRIO*
- *CONFIG_FMB_PORT_TASK_AFFINITY*
- *CONFIG_FMB_CONTROLLER_SLAVE_ID_SUPPORT*
- *CONFIG_FMB_CONTROLLER_NOTIFY_TIMEOUT*
- *CONFIG_FMB_CONTROLLER_NOTIFY_QUEUE_SIZE*
- *CONFIG_FMB_CONTROLLER_STACK_SIZE*
- *CONFIG_FMB_EVENT_QUEUE_TIMEOUT*
- *CONFIG_FMB_TIMER_PORT_ENABLED*
- *CONFIG_FMB_TIMER_GROUP*
- *CONFIG_FMB_TIMER_INDEX*
- *CONFIG_FMB_MASTER_TIMER_GROUP*
- *CONFIG_FMB_MASTER_TIMER_INDEX*
- *CONFIG_FMB_TIMER_ISR_IN_IRAM*

CONFIG_FMB_COMM_MODE_RTU_EN

Enable Modbus stack support for RTU mode

Found in: Component config > Modbus configuration

Enable RTU Modbus communication mode option for Modbus serial stack.

CONFIG_FMB_COMM_MODE_ASCII_EN

Enable Modbus stack support for ASCII mode

Found in: Component config > Modbus configuration

Enable ASCII Modbus communication mode option for Modbus serial stack.

CONFIG_FMB_MASTER_TIMEOUT_MS_RESPOND

Slave respond timeout (Milliseconds)

Found in: Component config > Modbus configuration

If master sends a frame which is not broadcast, it has to wait sometime for slave response. if slave is not respond in this time, the master will process timeout error.

CONFIG_FMB_MASTER_DELAY_MS_CONVERT

Slave conversion delay (Milliseconds)

Found in: Component config > Modbus configuration

If master sends a broadcast frame, it has to wait conversion time to delay, then master can send next frame.

CONFIG_FMB_QUEUE_LENGTH

Modbus serial task queue length

Found in: Component config > Modbus configuration

Modbus serial driver queue length. It is used by event queue task. See the serial driver API for more information.

CONFIG_FMB_SERIAL_TASK_STACK_SIZE

Modbus serial task stack size

Found in: Component config > Modbus configuration

Modbus serial task stack size for event queue task. It may be adjusted when debugging is enabled (for example).

CONFIG_FMB_SERIAL_BUF_SIZE

Modbus serial task RX/TX buffer size

Found in: Component config > Modbus configuration

Modbus serial task RX and TX buffer size for UART driver initialization. This buffer is used for modbus frame transfer. The Modbus protocol maximum frame size is 256 bytes. Bigger size can be used for non standard implementations.

CONFIG_FMB_SERIAL_ASCII_BITS_PER_SYMB

Number of data bits per ASCII character

Found in: Component config > Modbus configuration

This option defines the number of data bits per ASCII character.

CONFIG_FMB_SERIAL_ASCII_TIMEOUT_RESPOND_MS

Response timeout for ASCII communication mode (ms)

Found in: Component config > Modbus configuration

This option defines response timeout of slave in milliseconds for ASCII communication mode. Thus the timeout will expire and allow the master program to handle the error.

CONFIG_FMB_SERIAL_TASK_PRIO

Modbus serial task priority

Found in: Component config > Modbus configuration

Modbus UART driver event task priority. The priority of Modbus controller task is equal to (CONFIG_FMB_SERIAL_TASK_PRIO - 1).

CONFIG_FMB_PORT_TASK_AFFINITY

Modbus task affinity

Found in: Component config > Modbus configuration

Allows setting the core affinity of the Modbus controller task, i.e. whether the task is pinned to particular CPU, or allowed to run on any CPU.

Available options:

- No affinity (FMB_PORT_TASK_AFFINITY_NO_AFFINITY)

- CPU0 (FMB_PORT_TASK_AFFINITY_CPU0)
- CPU1 (FMB_PORT_TASK_AFFINITY_CPU1)

CONFIG_FMB_CONTROLLER_SLAVE_ID_SUPPORT

Modbus controller slave ID support

Found in: Component config > Modbus configuration

Modbus slave ID support enable. When enabled the Modbus <Report Slave ID> command is supported by stack.

CONFIG_FMB_CONTROLLER_SLAVE_ID

Modbus controller slave ID

Found in: Component config > Modbus configuration > CONFIG_FMB_CONTROLLER_SLAVE_ID_SUPPORT

Modbus slave ID value to identify modbus device in the network using <Report Slave ID> command. Most significant byte of ID is used as short device ID and other three bytes used as long ID.

CONFIG_FMB_CONTROLLER_NOTIFY_TIMEOUT

Modbus controller notification timeout (ms)

Found in: Component config > Modbus configuration

Modbus controller notification timeout in milliseconds. This timeout is used to send notification about accessed parameters.

CONFIG_FMB_CONTROLLER_NOTIFY_QUEUE_SIZE

Modbus controller notification queue size

Found in: Component config > Modbus configuration

Modbus controller notification queue size. The notification queue is used to get information about accessed parameters.

CONFIG_FMB_CONTROLLER_STACK_SIZE

Modbus controller stack size

Found in: Component config > Modbus configuration

Modbus controller task stack size. The Stack size may be adjusted when debug mode is used which requires more stack size (for example).

CONFIG_FMB_EVENT_QUEUE_TIMEOUT

Modbus stack event queue timeout (ms)

Found in: Component config > Modbus configuration

Modbus stack event queue timeout in milliseconds. This may help to optimize Modbus stack event processing time.

CONFIG_FMB_TIMER_PORT_ENABLED

Modbus stack use timer for 3.5T symbol time measurement

Found in: Component config > Modbus configuration

If this option is set the Modbus stack uses timer for T3.5 time measurement. Else the internal UART TOUT timeout is used for 3.5T symbol time measurement.

CONFIG_FMB_TIMER_GROUP

Slave Timer group number

Found in: Component config > Modbus configuration

Modbus slave Timer group number that is used for timeout measurement.

CONFIG_FMB_TIMER_INDEX

Slave Timer index in the group

Found in: Component config > Modbus configuration

Modbus slave Timer Index in the group that is used for timeout measurement.

CONFIG_FMB_MASTER_TIMER_GROUP

Master Timer group number

Found in: Component config > Modbus configuration

Modbus master Timer group number that is used for timeout measurement.

CONFIG_FMB_MASTER_TIMER_INDEX

Master Timer index

Found in: Component config > Modbus configuration

Modbus master Timer Index in the group that is used for timeout measurement. Note: Modbus master and slave should have different timer index to be able to work simultaneously.

CONFIG_FMB_TIMER_ISR_IN_IRAM

Place timer interrupt handler into IRAM

Found in: Component config > Modbus configuration

This option places Modbus timer IRQ handler into IRAM. This allows to avoid delays related to processing of non-IRAM-safe interrupts during a flash write operation (NVS updating a value, or some other flash API which has to perform an read/write operation and disable CPU cache). This option has dependency with the UART_ISR_IN_IRAM option which places UART interrupt handler into IRAM to prevent delays related to processing of UART events.

ADC-Calibration

ESP32S2-specific Contains:

- *CONFIG_ESP32S2_DEFAULT_CPU_FREQ_MHZ*
- *Memory protection*
- *Cache config*
- *CONFIG_ESP32S2_SPIRAM_SUPPORT*
- *CONFIG_ESP32S2_TRAX*
- *CONFIG_ESP32S2_UNIVERSAL_MAC_ADDRESSES*
- *CONFIG_ESP32S2_ULP_COPROC_ENABLED*
- *CONFIG_ESP32S2_DEBUG_OCDAWARE*
- *CONFIG_ESP32S2_DEBUG_STUBS_ENABLE*
- *CONFIG_ESP32S2_BROWNOUT_DET*
- *CONFIG_ESP32S2_TIME_SYSCALL*
- *CONFIG_ESP32S2_RTC_CLK_SRC*
- *CONFIG_ESP32S2_RTC_CLK_CAL_CYCLES*
- *CONFIG_ESP32S2_RTC_XTAL_CAL_RETRY*
- *CONFIG_ESP32S2_NO_BLOBS*
- *CONFIG_ESP32S2_RTCDATA_IN_FAST_MEM*
- *CONFIG_ESP32S2_ALLOW_RTC_FAST_MEM_AS_HEAP*

CONFIG_ESP32S2_DEFAULT_CPU_FREQ_MHZ

CPU frequency

Found in: Component config > ESP32S2-specific

CPU frequency to be set on application startup.

Available options:

- FPGA (ESP32S2_DEFAULT_CPU_FREQ_FPGA)
- 80 MHz (ESP32S2_DEFAULT_CPU_FREQ_80)
- 160 MHz (ESP32S2_DEFAULT_CPU_FREQ_160)
- 240 MHz (ESP32S2_DEFAULT_CPU_FREQ_240)

Memory protection Contains:

- *CONFIG_ESP32S2_MEMPROT_FEATURE*

CONFIG_ESP32S2_MEMPROT_FEATURE

Enable memory protection

Found in: Component config > ESP32S2-specific > Memory protection

If enabled, permission control module watches all memory access and fires panic handler if permission violation is detected. This feature automatically splits memory into data and instruction segments and sets Read/Execute permissions for instruction part (below splitting address) and Read/Write permissions for data part (above splitting address). The memory protection is effective on all access through IRAM0 and DRAM0 buses.

CONFIG_ESP32S2_MEMPROT_FEATURE_LOCK

Lock memory protection settings

Found in: Component config > ESP32S2-specific > Memory protection > CONFIG_ESP32S2_MEMPROT_FEATURE

Once locked, memory protection settings cannot be changed anymore. The lock is reset only on the chip startup.

Cache config Contains:

- *CONFIG_ESP32S2_INSTRUCTION_CACHE_SIZE*
- *CONFIG_ESP32S2_INSTRUCTION_CACHE_LINE_SIZE*
- *CONFIG_ESP32S2_DATA_CACHE_SIZE*
- *CONFIG_ESP32S2_DATA_CACHE_LINE_SIZE*
- *CONFIG_ESP32S2_INSTRUCTION_CACHE_WRAP*
- *CONFIG_ESP32S2_DATA_CACHE_WRAP*

CONFIG_ESP32S2_INSTRUCTION_CACHE_SIZE

Instruction cache size

Found in: Component config > ESP32S2-specific > Cache config

Instruction cache size to be set on application startup. If you use 8KB instruction cache rather than 16KB instruction cache, then the other 8KB will be added to the heap.

Available options:

- 8KB (ESP32S2_INSTRUCTION_CACHE_8KB)
- 16KB (ESP32S2_INSTRUCTION_CACHE_16KB)

CONFIG_ESP32S2_INSTRUCTION_CACHE_LINE_SIZE

Instruction cache line size

Found in: Component config > ESP32S2-specific > Cache config

Instruction cache line size to be set on application startup.

Available options:

- 16 Bytes (ESP32S2_INSTRUCTION_CACHE_LINE_16B)
- 32 Bytes (ESP32S2_INSTRUCTION_CACHE_LINE_32B)

CONFIG_ESP32S2_DATA_CACHE_SIZE

Data cache size

Found in: Component config > ESP32S2-specific > Cache config

Data cache size to be set on application startup. If you use 8KB data cache rather than 16KB data cache, the other 8KB will be added to the heap.

Available options:

- 0KB (ESP32S2_DATA_CACHE_0KB)
- 8KB (ESP32S2_DATA_CACHE_8KB)
- 16KB (ESP32S2_DATA_CACHE_16KB)

CONFIG_ESP32S2_DATA_CACHE_LINE_SIZE

Data cache line size

Found in: Component config > ESP32S2-specific > Cache config

Data cache line size to be set on application startup.

Available options:

- 16 Bytes (ESP32S2_DATA_CACHE_LINE_16B)
- 32 Bytes (ESP32S2_DATA_CACHE_LINE_32B)

CONFIG_ESP32S2_INSTRUCTION_CACHE_WRAP

Enable instruction cache wrap

Found in: Component config > ESP32S2-specific > Cache config

If enabled, instruction cache will use wrap mode to read spi flash (maybe spiram). The wrap length equals to INSTRUCTION_CACHE_LINE_SIZE. However, it depends on complex conditions.

CONFIG_ESP32S2_DATA_CACHE_WRAP

Enable data cache wrap

Found in: Component config > ESP32S2-specific > Cache config

If enabled, data cache will use wrap mode to read spiram (maybe spi flash). The wrap length equals to DATA_CACHE_LINE_SIZE. However, it depends on complex conditions.

CONFIG_ESP32S2_SPIRAM_SUPPORT

Support for external, SPI-connected RAM

Found in: Component config > ESP32S2-specific

This enables support for an external SPI RAM chip, connected in parallel with the main SPI flash chip.

SPI RAM config Contains:

- *CONFIG_SPIRAM_TYPE*
- *CONFIG_SPIRAM_FETCH_INSTRUCTIONS*
- *CONFIG_SPIRAM_RODATA*
- *CONFIG_SPIRAM_USE_AHB_DBUS3*
- *CONFIG_SPIRAM_SPEED*
- *CONFIG_SPIRAM_BOOT_INIT*
- *CONFIG_SPIRAM_USE*
- *CONFIG_SPIRAM_MEMTEST*
- *CONFIG_SPIRAM_MALLOC_ALWAYSINTERNAL*
- *CONFIG_SPIRAM_TRY_ALLOCATE_WIFI_LWIP*
- *CONFIG_SPIRAM_MALLOC_RESERVE_INTERNAL*

CONFIG_SPIRAM_TYPE

Type of SPI RAM chip in use

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

Available options:

- Auto-detect (SPIRAM_TYPE_AUTO)
- ESP-PSRAM16 or APS1604 (SPIRAM_TYPE_ESPPSRAM16)
- ESP-PSRAM32 or IS25WP032 (SPIRAM_TYPE_ESPPSRAM32)
- ESP-PSRAM64 or LY68L6400 (SPIRAM_TYPE_ESPPSRAM64)

CONFIG_SPIRAM_FETCH_INSTRUCTIONS

Cache fetch instructions from SPI RAM

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

If enabled, instruction in flash will be copied into SPIRAM. If SPIRAM_RODATA also enabled, you can run the instruction when erasing or programming the flash.

CONFIG_SPIRAM_RODATA

Cache load read only data from SPI RAM

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

If enabled, radata in flash will be copied into SPIRAM. If SPIRAM_FETCH_INSTRUCTIONS also enabled, you can run the instruction when erasing or programming the flash.

CONFIG_SPIRAM_USE_AHB_DBUS3

Enable AHB DBUS3 to access SPIRAM

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

If Enabled, if SPI_CONFIG_SIZE is bigger then 10MB+576KB, then you can have 4MB more space to map the SPIRAM. However, the AHB bus is slower than other data cache buses.

CONFIG_SPIRAM_SPEED

Set RAM clock speed

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

Select the speed for the SPI RAM chip. If SPI RAM is enabled, we only support three combinations of SPI speed mode we supported now:

1. Flash SPI running at 40Mhz and RAM SPI running at 40Mhz
2. Flash SPI running at 80Mhz and RAM SPI running at 40Mhz
3. Flash SPI running at 80Mhz and RAM SPI running at 80Mhz

Note: If the third mode(80Mhz+80Mhz) is enabled for SPI RAM of type 32MBit, one of the HSPI/VSPI host will be occupied by the system. Which SPI host to use can be selected by the config item SPIRAM_OCCUPY_SPI_HOST. Application code should never touch HSPI/VSPI hardware in this case. The option to select 80MHz will only be visible if the flash SPI speed is also 80MHz. (ESPTOOLPY_FLASHFREQ_80M is true)

Available options:

- 80MHz clock speed (SPIRAM_SPEED_80M)
- 40Mhz clock speed (SPIRAM_SPEED_40M)
- 26Mhz clock speed (SPIRAM_SPEED_26M)
- 20Mhz clock speed (SPIRAM_SPEED_20M)

CONFIG_SPIRAM_BOOT_INIT

Initialize SPI RAM during startup

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

If this is enabled, the SPI RAM will be enabled during initial boot. Unless you have specific requirements, you’ ll want to leave this enabled so memory allocated during boot-up can also be placed in SPI RAM.

CONFIG_SPIRAM_IGNORE_NOTFOUND

Ignore PSRAM when not found

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config > CONFIG_SPIRAM_BOOT_INIT

Normally, if psram initialization is enabled during compile time but not found at runtime, it is seen as an error making the CPU panic. If this is enabled, booting will complete but no PSRAM will be available.

CONFIG_SPIRAM_USE

SPI RAM access method

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

The SPI RAM can be accessed in multiple methods: by just having it available as an unmanaged memory region in the CPU's memory map, by integrating it in the heap as 'special' memory needing heap_caps_malloc to allocate, or by fully integrating it making malloc() also able to return SPI RAM pointers.

Available options:

- Integrate RAM into memory map (SPIRAM_USE_MEMMAP)
- Make RAM allocatable using heap_caps_malloc(···, MALLOC_CAP_SPIRAM) (SPIRAM_USE_CAPS_ALLOC)
- Make RAM allocatable using malloc() as well (SPIRAM_USE_MALLOC)

CONFIG_SPIRAM_MEMTEST

Run memory test on SPI RAM initialization

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

Runs a rudimentary memory test on initialization. Aborts when memory test fails. Disable this for slightly faster startup.

CONFIG_SPIRAM_MALLOC_ALWAYSINTERNAL

Maximum malloc() size, in bytes, to always put in internal memory

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

If malloc() is capable of also allocating SPI-connected ram, its allocation strategy will prefer to allocate chunks less than this size in internal memory, while allocations larger than this will be done from external RAM. If allocation from the preferred region fails, an attempt is made to allocate from the non-preferred region instead, so malloc() will not suddenly fail when either internal or external memory is full.

CONFIG_SPIRAM_TRY_ALLOCATE_WIFI_LWIP

Try to allocate memories of WiFi and LWIP in SPIRAM firstly. If failed, allocate internal memory

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

Try to allocate memories of WiFi and LWIP in SPIRAM firstly. If failed, try to allocate internal memory then.

CONFIG_SPIRAM_MALLOC_RESERVE_INTERNAL

Reserve this amount of bytes for data that specifically needs to be in DMA or internal memory

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_SPIRAM_SUPPORT > SPI RAM config

Because the external/internal RAM allocation strategy is not always perfect, it sometimes may happen that the internal memory is entirely filled up. This causes allocations that are specifically done in internal memory, for example the stack for new tasks or memory to service DMA or have memory that's also

available when SPI cache is down, to fail. This option reserves a pool specifically for requests like that; the memory in this pool is not given out when a normal malloc() is called.

Set this to 0 to disable this feature.

Note that because FreeRTOS stacks are forced to internal memory, they will also use this memory pool; be sure to keep this in mind when adjusting this value.

Note also that the DMA reserved pool may not be one single contiguous memory region, depending on the configured size and the static memory usage of the app.

CONFIG_ESP32S2_TRAX

Use TRAX tracing feature

Found in: Component config > ESP32S2-specific

The ESP32S2 contains a feature which allows you to trace the execution path the processor has taken through the program. This is stored in a chunk of 32K (16K for single-processor) of memory that can't be used for general purposes anymore. Disable this if you do not know what this is.

CONFIG_ESP32S2_UNIVERSAL_MAC_ADDRESSES

Number of universally administered (by IEEE) MAC address

Found in: Component config > ESP32S2-specific

Configure the number of universally administered (by IEEE) MAC addresses. During initialization, MAC addresses for each network interface are generated or derived from a single base MAC address. If the number of universal MAC addresses is Two, all interfaces (WiFi station, WiFi softap) receive a universally administered MAC address. They are generated sequentially by adding 0, and 1 (respectively) to the final octet of the base MAC address. If the number of universal MAC addresses is one, only WiFi station receives a universally administered MAC address. It's generated by adding 0 to the base MAC address. The WiFi softap receives local MAC addresses. It's derived from the universal WiFi station MAC addresses. When using the default (Espressif-assigned) base MAC address, either setting can be used. When using a custom universal MAC address range, the correct setting will depend on the allocation of MAC addresses in this range (either 1 or 2 per device.)

Available options:

- One (ESP32S2_UNIVERSAL_MAC_ADDRESSES_ONE)
- Two (ESP32S2_UNIVERSAL_MAC_ADDRESSES_TWO)

CONFIG_ESP32S2_ULP_COPROC_ENABLED

Enable Ultra Low Power (ULP) Coprocessor

Found in: Component config > ESP32S2-specific

Set to 'y' if you plan to load a firmware for the coprocessor.

If this option is enabled, further coprocessor configuration will appear in the Components menu.

CONFIG_ESP32S2_ULP_COPROC_RESERVE_MEM

RTC slow memory reserved for coprocessor

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_ULP_COPROC_ENABLED

Bytes of memory to reserve for ULP coprocessor firmware & data.

Data is reserved at the beginning of RTC slow memory.

CONFIG_ESP32S2_DEBUG_OCDAWARE

Make exception and panic handlers JTAG/OCD aware

Found in: Component config > ESP32S2-specific

The FreeRTOS panic and unhandled exception handers can detect a JTAG OCD debugger and instead of panicking, have the debugger stop on the offending instruction.

CONFIG_ESP32S2_DEBUG_STUBS_ENABLE

OpenOCD debug stubs

Found in: Component config > ESP32S2-specific

Debug stubs are used by OpenOCD to execute pre-compiled onboard code which does some useful debugging, e.g. GCOV data dump.

CONFIG_ESP32S2_BROWNOUT_DET

Hardware brownout detect & reset

Found in: Component config > ESP32S2-specific

The ESP32-S2 has a built-in brownout detector which can detect if the voltage is lower than a specific value. If this happens, it will reset the chip in order to prevent unintended behaviour.

CONFIG_ESP32S2_BROWNOUT_DET_LVL_SEL

Brownout voltage level

Found in: Component config > ESP32S2-specific > CONFIG_ESP32S2_BROWNOUT_DET

The brownout detector will reset the chip when the supply voltage is approximately below this level. Note that there may be some variation of brownout voltage level between each ESP3-S2 chip.

#The voltage levels here are estimates, more work needs to be done to figure out the exact voltages #of the brownout threshold levels.

Available options:

- 2.44V (ESP32S2_BROWNOUT_DET_LVL_SEL_7)
- 2.56V (ESP32S2_BROWNOUT_DET_LVL_SEL_6)
- 2.67V (ESP32S2_BROWNOUT_DET_LVL_SEL_5)
- 2.84V (ESP32S2_BROWNOUT_DET_LVL_SEL_4)
- 2.98V (ESP32S2_BROWNOUT_DET_LVL_SEL_3)
- 3.19V (ESP32S2_BROWNOUT_DET_LVL_SEL_2)
- 3.30V (ESP32S2_BROWNOUT_DET_LVL_SEL_1)

CONFIG_ESP32S2_TIME_SYSCALL

Timers used for gettimeofday function

Found in: Component config > ESP32S2-specific

This setting defines which hardware timers are used to implement 'gettimeofday' and 'time' functions in C library.

- If both high-resolution and RTC timers are used, timekeeping will continue in deep sleep. Time will be reported at 1 microsecond resolution. This is the default, and the recommended option.
- If only high-resolution timer is used, gettimeofday will provide time at microsecond resolution. Time will not be preserved when going into deep sleep mode.
- If only RTC timer is used, timekeeping will continue in deep sleep, but time will be measured at 6.(6) microsecond resolution. Also the gettimeofday function itself may take longer to run.
- If no timers are used, gettimeofday and time functions return -1 and set errno to ENOSYS.

- When RTC is used for timekeeping, two RTC_STORE registers are used to keep time in deep sleep mode.

Available options:

- RTC and high-resolution timer (ESP32S2_TIME_SYSCALL_USE_RTC_FRC1)
- RTC (ESP32S2_TIME_SYSCALL_USE_RTC)
- High-resolution timer (ESP32S2_TIME_SYSCALL_USE_FRC1)
- None (ESP32S2_TIME_SYSCALL_USE_NONE)

CONFIG_ESP32S2_RTC_CLK_SRC

RTC clock source

Found in: Component config > ESP32S2-specific

Choose which clock is used as RTC clock source.

- "Internal 90kHz oscillator" option provides lowest deep sleep current consumption, and does not require extra external components. However frequency stability with respect to temperature is poor, so time may drift in deep/light sleep modes.
- "External 32kHz crystal" provides better frequency stability, at the expense of slightly higher (1uA) deep sleep current consumption.
- "External 32kHz oscillator" allows using 32kHz clock generated by an external circuit. In this case, external clock signal must be connected to 32K_XN pin. Amplitude should be <1.2V in case of sine wave signal, and <1V in case of square wave signal. Common mode voltage should be 0.1 < Vcm < 0.5Vamp, where Vamp is the signal amplitude. Additionally, 1nF capacitor must be connected between 32K_XP pin and ground. 32K_XP pin can not be used as a GPIO in this case.
- "Internal 8MHz oscillator divided by 256" option results in higher deep sleep current (by 5uA) but has better frequency stability than the internal 90kHz oscillator. It does not require external components.

Available options:

- Internal 90kHz RC oscillator (ESP32S2_RTC_CLK_SRC_INT_RC)
- External 32kHz crystal (ESP32S2_RTC_CLK_SRC_EXT_CRYS)
- External 32kHz oscillator at 32K_XN pin (ESP32S2_RTC_CLK_SRC_EXT_OSC)
- Internal 8MHz oscillator, divided by 256 (~32kHz) (ESP32S2_RTC_CLK_SRC_INT_8MD256)

CONFIG_ESP32S2_RTC_CLK_CAL_CYCLES

Number of cycles for RTC_SLOW_CLK calibration

Found in: Component config > ESP32S2-specific

When the startup code initializes RTC_SLOW_CLK, it can perform calibration by comparing the RTC_SLOW_CLK frequency with main XTAL frequency. This option sets the number of RTC_SLOW_CLK cycles measured by the calibration routine. Higher numbers increase calibration precision, which may be important for applications which spend a lot of time in deep sleep. Lower numbers reduce startup time.

When this option is set to 0, clock calibration will not be performed at startup, and approximate clock frequencies will be assumed:

- 90000 Hz if internal RC oscillator is used as clock source. For this use value 1024.
- 32768 Hz if the 32k crystal oscillator is used. For this use value 3000 or more. In case more value will help improve the definition of the launch of the crystal. If the crystal could not start, it will be switched to internal RC.

CONFIG_ESP32S2_RTC_XTAL_CAL_RETRY

Number of attempts to repeat 32k XTAL calibration

Found in: Component config > ESP32S2-specific

Number of attempts to repeat 32k XTAL calibration before giving up and switching to the internal RC. Increase this option if the 32k crystal oscillator does not start and switches to internal RC.

CONFIG_ESP32S2_NO_BLOBS

No Binary Blobs

Found in: Component config > ESP32S2-specific

If enabled, this disables the linking of binary libraries in the application build. Note that after enabling this Wi-Fi/Bluetooth will not work.

CONFIG_ESP32S2_RTCDATA_IN_FAST_MEM

Place RTC_DATA_ATTR and RTC_RODATA_ATTR variables into RTC fast memory segment

Found in: Component config > ESP32S2-specific

This option allows to place .rtc_data and .rtc_rodata sections into RTC fast memory segment to free the slow memory region for ULP programs.

CONFIG_ESP32S2_ALLOW_RTC_FAST_MEM_AS_HEAP

Enable RTC fast memory for dynamic allocations

Found in: Component config > ESP32S2-specific

This config option allows to add RTC fast memory region to system heap with capability similar to that of DRAM region but without DMA. This memory will be consumed first per heap initialization order by early startup services and scheduler related code. Speed wise RTC fast memory operates on APB clock and hence does not have much performance impact.

Power Management

Contains:

- *CONFIG_PM_ENABLE*

CONFIG_PM_ENABLE

Support for power management

Found in: Component config > Power Management

If enabled, application is compiled with support for power management. This option has run-time overhead (increased interrupt latency, longer time to enter idle state), and it also reduces accuracy of RTOS ticks and timers used for timekeeping. Enable this option if application uses power management APIs.

CONFIG_PM_DFS_INIT_AUTO

Enable dynamic frequency scaling (DFS) at startup

Found in: Component config > Power Management > CONFIG_PM_ENABLE

If enabled, startup code configures dynamic frequency scaling. Max CPU frequency is set to CONFIG_ESP32S2_DEFAULT_CPU_FREQ_MHZ setting, min frequency is set to XTAL frequency. If disabled, DFS will not be active until the application configures it using esp_pm_configure function.

CONFIG_PM_USE_RTC_TIMER_REF

Use RTC timer to prevent time drift (EXPERIMENTAL)

Found in: Component config > Power Management > CONFIG_PM_ENABLE

When APB clock frequency changes, high-resolution timer (esp_timer) scale and base value need to be adjusted. Each adjustment may cause small error, and over time such small errors may cause time drift. If this option is enabled, RTC timer will be used as a reference to compensate for the drift. It is recommended that this option is only used if 32k XTAL is selected as RTC clock source.

CONFIG_PM_PROFILING

Enable profiling counters for PM locks

Found in: Component config > Power Management > CONFIG_PM_ENABLE

If enabled, esp_pm_* functions will keep track of the amount of time each of the power management locks has been held, and esp_pm_dump_locks function will print this information. This feature can be used to analyze which locks are preventing the chip from going into a lower power state, and see what time the chip spends in each power saving mode. This feature does incur some run-time overhead, so should typically be disabled in production builds.

CONFIG_PM_TRACE

Enable debug tracing of PM using GPIOs

Found in: Component config > Power Management > CONFIG_PM_ENABLE

If enabled, some GPIOs will be used to signal events such as RTOS ticks, frequency switching, entry/exit from idle state. Refer to pm_trace.c file for the list of GPIOs. This feature is intended to be used when analyzing/debugging behavior of power management implementation, and should be kept disabled in applications.

CoAP Configuration

Contains:

- *CONFIG_COAP_MBEDTLS_ENCRYPTION_MODE*
- *CONFIG_COAP_MBEDTLS_DEBUG*

CONFIG_COAP_MBEDTLS_ENCRYPTION_MODE

CoAP Encryption method

Found in: Component config > CoAP Configuration

If the CoAP information is to be encrypted, the encryption environment can be set up in one of two ways (default being Pre-Shared key mode)

- Encrypt using defined Pre-Shared Keys (PSK if uri includes coaps://)
- Encrypt using defined Public Key Infrastructure (PKI if uri includes coaps://)

Available options:

- Pre-Shared Keys (COAP_MBEDTLS_PSK)
- PKI Certificates (COAP_MBEDTLS_PKI)

CONFIG_COAP_MBEDTLS_DEBUG

Enable CoAP debugging

Found in: Component config > CoAP Configuration

Enable CoAP debugging functions at compile time for the example code.

If this option is enabled, call coap_set_log_level() at runtime in order to enable CoAP debug output via the ESP log mechanism.

CONFIG_COAP_MBEDTLS_DEBUG_LEVEL

Set CoAP debugging level

Found in: Component config > CoAP Configuration > CONFIG_COAP_MBEDTLS_DEBUG

Set CoAP debugging level

Available options:

- Emergency (COAP_LOG_EMERG)
- Alert (COAP_LOG_ALERT)
- Critical (COAP_LOG_CRIT)
- Error (COAP_LOG_ERROR)
- Warning (COAP_LOG_WARNING)
- Notice (COAP_LOG_NOTICE)
- Info (COAP_LOG_INFO)
- Debug (COAP_LOG_DEBUG)

Supplicant Contains:

- *CONFIG_WPA_MBEDTLS_CRYPTO*
- *CONFIG_WPA_DEBUG_PRINT*
- *CONFIG_WPA_TESTING_OPTIONS*
- *CONFIG_WPA_WPS_STRICT*

CONFIG_WPA_MBEDTLS_CRYPTO

Use MbedTLS crypto API’s

Found in: Component config > Supplicant

Select this option to use MbedTLS crypto API’s which utilize hardware acceleration.

CONFIG_WPA_DEBUG_PRINT

Print debug messages from WPA Supplicant

Found in: Component config > Supplicant

Select this option to print logging information from WPA supplicant, this includes handshake information and key hex dumps depending on the project logging level.

Enabling this could increase the build size ~60kb depending on the project logging level.

CONFIG_WPA_TESTING_OPTIONS

Add DPP testing code

Found in: Component config > Supplicant

Select this to enable unity test for DPP.

CONFIG_WPA_WPS_STRICT

Strictly validate all WPS attributes

Found in: Component config > Supplicant

Select this option to enable validate each WPS attribute rigorously. Disabling this add the workaorunds with various APs. Enabling this may cause inter operability issues with some APs.

Wi-Fi Provisioning Manager Contains:

- *CONFIG_WIFI_PROV_SCAN_MAX_ENTRIES*
- *CONFIG_WIFI_PROV_AUTOSTOP_TIMEOUT*

CONFIG_WIFI_PROV_SCAN_MAX_ENTRIES

Max Wi-Fi Scan Result Entries

Found in: Component config > Wi-Fi Provisioning Manager

This sets the maximum number of entries of Wi-Fi scan results that will be kept by the provisioning manager

CONFIG_WIFI_PROV_AUTOSTOP_TIMEOUT

Provisioning auto-stop timeout

Found in: Component config > Wi-Fi Provisioning Manager

Time (in seconds) after which the Wi-Fi provisioning manager will auto-stop after connecting to a Wi-Fi network successfully.

Wear Levelling Contains:

- *CONFIG_WL_SECTOR_SIZE*
- *CONFIG_WL_SECTOR_MODE*

CONFIG_WL_SECTOR_SIZE

Wear Levelling library sector size

Found in: Component config > Wear Levelling

Sector size used by wear levelling library. You can set default sector size or size that will fit to the flash device sector size.

With sector size set to 4096 bytes, wear levelling library is more efficient. However if FAT filesystem is used on top of wear levelling library, it will need more temporary storage: 4096 bytes for each mounted filesystem and 4096 bytes for each opened file.

With sector size set to 512 bytes, wear levelling library will perform more operations with flash memory, but less RAM will be used by FAT filesystem library (512 bytes for the filesystem and 512 bytes for each file opened).

Available options:

- 512 (WL_SECTOR_SIZE_512)
- 4096 (WL_SECTOR_SIZE_4096)

CONFIG_WL_SECTOR_MODE

Sector store mode

Found in: Component config > Wear Levelling

Specify the mode to store data into flash:

- In Performance mode a data will be stored to the RAM and then stored back to the flash. Compared to the Safety mode, this operation is faster, but if power will be lost when erase sector operation is in progress, then the data from complete flash device sector will be lost.
- In Safety mode data from complete flash device sector will be read from flash, modified, and then stored back to flash. Compared to the Performance mode, this operation is slower, but if power is lost during erase sector operation, then the data from full flash device sector will not be lost.

Available options:
- Perfomance (WL_SECTOR_MODE_PERF)
- Safety (WL_SECTOR_MODE_SAFE)

Virtual file system Contains:

- *CONFIG_VFS_SUPPORT_IO*

CONFIG_VFS_SUPPORT_IO

Provide basic I/O functions

Found in: Component config > Virtual file system

If enabled, the following functions are provided by the VFS component.

open, close, read, write, pread, pwrite, lseek, fstat, fsync, ioctl, fcntl

Filesystem drivers can then be registered to handle these functions for specific paths.

Disabling this option can save memory when the support for these functions is not required.

CONFIG_VFS_SUPPORT_DIR

Provide directory related functions

Found in: Component config > Virtual file system > CONFIG_VFS_SUPPORT_IO

If enabled, the following functions are provided by the VFS component.

stat, link, unlink, rename, utime, access, truncate, rmdir, mkdir, opendir, closedir, readdir, readdir_r, seekdir, telldir, rewinddir

Filesystem drivers can then be registered to handle these functions for specific paths.

Disabling this option can save memory when the support for these functions is not required.

CONFIG_VFS_SUPPORT_SELECT

Provide select function

Found in: Component config > Virtual file system > CONFIG_VFS_SUPPORT_IO

If enabled, select function is provided by the VFS component, and can be used on peripheral file descriptors (such as UART) and sockets at the same time.

If disabled, the default select implementation will be provided by LWIP for sockets only.

Disabling this option can reduce code size if support for "select" on UART file descriptors is not required.

CONFIG_VFS_SUPPRESS_SELECT_DEBUG_OUTPUT

Suppress select() related debug outputs

Found in: Component config > Virtual file system > CONFIG_VFS_SUPPORT_IO > CONFIG_VFS_SUPPORT_SELECT

Select() related functions might produce an unconveniently lot of debug outputs when one sets the default log level to DEBUG or higher. It is possible to suppress these debug outputs by enabling this option.

CONFIG_VFS_SUPPORT_TERMIOS

Provide termios.h functions

Found in: Component config > Virtual file system > CONFIG_VFS_SUPPORT_IO

Disabling this option can save memory when the support for termios.h is not required.

Host File System I/O (Semihosting) Contains:

- *CONFIG_VFS_SEMIHOSTFS_MAX_MOUNT_POINTS*
- *CONFIG_VFS_SEMIHOSTFS_HOST_PATH_MAX_LEN*

CONFIG_VFS_SEMIHOSTFS_MAX_MOUNT_POINTS

Host FS: Maximum number of the host filesystem mount points

Found in: Component config > Virtual file system > CONFIG_VFS_SUPPORT_IO > Host File System I/O (Semihosting)

Define maximum number of host filesystem mount points.

CONFIG_VFS_SEMIHOSTFS_HOST_PATH_MAX_LEN

Host FS: Maximum path length for the host base directory

Found in: Component config > Virtual file system > CONFIG_VFS_SUPPORT_IO > Host File System I/O (Semihosting)

Define maximum path length for the host base directory which is to be mounted. If host path passed to esp_vfs_semihost_register() is longer than this value it will be truncated.

Unity unit testing library Contains:

- *CONFIG_UNITY_ENABLE_FLOAT*
- *CONFIG_UNITY_ENABLE_DOUBLE*
- *CONFIG_UNITY_ENABLE_COLOR*
- *CONFIG_UNITY_ENABLE_IDF_TEST_RUNNER*
- *CONFIG_UNITY_ENABLE_FIXTURE*
- *CONFIG_UNITY_ENABLE_BACKTRACE_ON_FAIL*

CONFIG_UNITY_ENABLE_FLOAT

Support for float type

Found in: Component config > Unity unit testing library

If not set, assertions on float arguments will not be available.

CONFIG_UNITY_ENABLE_DOUBLE

Support for double type

Found in: Component config > Unity unit testing library

If not set, assertions on double arguments will not be available.

CONFIG_UNITY_ENABLE_COLOR

Colorize test output

Found in: Component config > Unity unit testing library

If set, Unity will colorize test results using console escape sequences.

CONFIG_UNITY_ENABLE_IDF_TEST_RUNNER

Include ESP-IDF test registration/running helpers

Found in: Component config > Unity unit testing library

If set, then the following features will be available:

- TEST_CASE macro which performs automatic registration of test functions
- Functions to run registered test functions: unity_run_all_tests, unity_run_tests_with_filter, unity_run_single_test_by_name.
- Interactive menu which lists test cases and allows choosing the tests to be run, available via unity_run_menu function.

Disable if a different test registration mechanism is used.

CONFIG_UNITY_ENABLE_FIXTURE

Include Unity test fixture

Found in: Component config > Unity unit testing library

If set, unity_fixture.h header file and associated source files are part of the build. These provide an optional set of macros and functions to implement test groups.

CONFIG_UNITY_ENABLE_BACKTRACE_ON_FAIL

Print a backtrace when a unit test fails

Found in: Component config > Unity unit testing library

If set, the unity framework will print the backtrace information before jumping back to the test menu. The jumping is usually occurs in assert functions such as TEST_ASSERT, TEST_FAIL etc.

SPIFFS Configuration

Contains:

- *CONFIG_SPIFFS_MAX_PARTITIONS*
- *SPIFFS Cache Configuration*
- *CONFIG_SPIFFS_PAGE_CHECK*
- *CONFIG_SPIFFS_GC_MAX_RUNS*
- *CONFIG_SPIFFS_GC_STATS*
- *CONFIG_SPIFFS_PAGE_SIZE*
- *CONFIG_SPIFFS_OBJ_NAME_LEN*
- *CONFIG_SPIFFS_FOLLOW_SYMLINKS*
- *CONFIG_SPIFFS_USE_MAGIC*
- *CONFIG_SPIFFS_META_LENGTH*
- *CONFIG_SPIFFS_USE_MTIME*
- *CONFIG_SPIFFS_MTIME_WIDE_64_BITS*
- *Debug Configuration*

CONFIG_SPIFFS_MAX_PARTITIONS

Maximum Number of Partitions

Found in: Component config > SPIFFS Configuration

Define maximum number of partitions that can be mounted.

SPIFFS Cache Configuration Contains:

- *CONFIG_SPIFFS_CACHE*

CONFIG_SPIFFS_CACHE

Enable SPIFFS Cache

Found in: Component config > SPIFFS Configuration > SPIFFS Cache Configuration

Enables/disable memory read caching of nucleus file system operations.

CONFIG_SPIFFS_CACHE_WR

Enable SPIFFS Write Caching

Found in: Component config > SPIFFS Configuration > SPIFFS Cache Configuration > CONFIG_SPIFFS_CACHE

Enables memory write caching for file descriptors in hydrogen.

CONFIG_SPIFFS_CACHE_STATS

Enable SPIFFS Cache Statistics

Found in: Component config > SPIFFS Configuration > SPIFFS Cache Configuration > CONFIG_SPIFFS_CACHE

Enable/disable statistics on caching. Debug/test purpose only.

CONFIG_SPIFFS_PAGE_CHECK

Enable SPIFFS Page Check

Found in: Component config > SPIFFS Configuration

Always check header of each accessed page to ensure consistent state. If enabled it will increase number of reads from flash, especially if cache is disabled.

CONFIG_SPIFFS_GC_MAX_RUNS

Set Maximum GC Runs

Found in: Component config > SPIFFS Configuration

Define maximum number of GC runs to perform to reach desired free pages.

CONFIG_SPIFFS_GC_STATS

Enable SPIFFS GC Statistics

Found in: Component config > SPIFFS Configuration

Enable/disable statistics on gc. Debug/test purpose only.

CONFIG_SPIFFS_PAGE_SIZE

SPIFFS logical page size

Found in: Component config > SPIFFS Configuration

Logical page size of SPIFFS partition, in bytes. Must be multiple of flash page size (which is usually 256 bytes). Larger page sizes reduce overhead when storing large files, and improve filesystem performance when reading large files. Smaller page sizes reduce overhead when storing small (< page size) files.

CONFIG_SPIFFS_OBJ_NAME_LEN

Set SPIFFS Maximum Name Length

Found in: Component config > SPIFFS Configuration

Object name maximum length. Note that this length include the zero-termination character, meaning maximum string of characters can at most be SPIFFS_OBJ_NAME_LEN - 1.

SPIFFS_OBJ_NAME_LEN + SPIFFS_META_LENGTH should not exceed SPIFFS_PAGE_SIZE - 64.

CONFIG_SPIFFS_FOLLOW_SYMLINKS

Enable symbolic links for image creation

Found in: Component config > SPIFFS Configuration

If this option is enabled, symbolic links are taken into account during partition image creation.

CONFIG_SPIFFS_USE_MAGIC

Enable SPIFFS Filesystem Magic

Found in: Component config > SPIFFS Configuration

Enable this to have an identifiable spiffs filesystem. This will look for a magic in all sectors to determine if this is a valid spiffs system or not at mount time.

CONFIG_SPIFFS_USE_MAGIC_LENGTH

Enable SPIFFS Filesystem Length Magic

Found in: Component config > SPIFFS Configuration > CONFIG_SPIFFS_USE_MAGIC

If this option is enabled, the magic will also be dependent on the length of the filesystem. For example, a filesystem configured and formatted for 4 megabytes will not be accepted for mounting with a configuration defining the filesystem as 2 megabytes.

CONFIG_SPIFFS_META_LENGTH

Size of per-file metadata field

Found in: Component config > SPIFFS Configuration

This option sets the number of extra bytes stored in the file header. These bytes can be used in an application-specific manner. Set this to at least 4 bytes to enable support for saving file modification time.

SPIFFS_OBJ_NAME_LEN + SPIFFS_META_LENGTH should not exceed SPIFFS_PAGE_SIZE - 64.

CONFIG_SPIFFS_USE_MTIME

Save file modification time

Found in: Component config > SPIFFS Configuration

If enabled, then the first 4 bytes of per-file metadata will be used to store file modification time (mtime), accessible through stat/fstat functions. Modification time is updated when the file is opened.

CONFIG_SPIFFS_MTIME_WIDE_64_BITS

The time field occupies 64 bits in the image instead of 32 bits

Found in: Component config > SPIFFS Configuration

If this option is not set, the time field is 32 bits (up to 2106 year), otherwise it is 64 bits and make sure it matches SPIFFS_META_LENGTH. If the chip already has the spiffs image with the time field = 32 bits then this option cannot be applied in this case. Erase it first before using this option. To resolve the Y2K38 problem for the spiffs, use a toolchain with support time_t 64 bits (see SDK_TOOLCHAIN_SUPPORTS_TIME_WIDE_64_BITS).

Debug Configuration Contains:

- *CONFIG_SPIFFS_DBG*
- *CONFIG_SPIFFS_API_DBG*
- *CONFIG_SPIFFS_GC_DBG*
- *CONFIG_SPIFFS_CACHE_DBG*
- *CONFIG_SPIFFS_CHECK_DBG*
- *CONFIG_SPIFFS_TEST_VISUALISATION*

CONFIG_SPIFFS_DBG

Enable general SPIFFS debug

Found in: Component config > SPIFFS Configuration > Debug Configuration

Enabling this option will print general debug mesages to the console.

CONFIG_SPIFFS_API_DBG

Enable SPIFFS API debug

Found in: Component config > SPIFFS Configuration > Debug Configuration

Enabling this option will print API debug mesages to the console.

CONFIG_SPIFFS_GC_DBG

Enable SPIFFS Garbage Cleaner debug

Found in: Component config > SPIFFS Configuration > Debug Configuration

Enabling this option will print GC debug mesages to the console.

CONFIG_SPIFFS_CACHE_DBG

Enable SPIFFS Cache debug

Found in: Component config > SPIFFS Configuration > Debug Configuration

Enabling this option will print cache debug mesages to the console.

CONFIG_SPIFFS_CHECK_DBG

Enable SPIFFS Filesystem Check debug

Found in: Component config > SPIFFS Configuration > Debug Configuration

Enabling this option will print Filesystem Check debug mesages to the console.

CONFIG_SPIFFS_TEST_VISUALISATION

Enable SPIFFS Filesystem Visualization

Found in: Component config > SPIFFS Configuration > Debug Configuration

Enable this option to enable SPIFFS_vis function in the API.

SPI Flash driver Contains:

- *CONFIG_SPI_FLASH_VERIFY_WRITE*
- *CONFIG_SPI_FLASH_ENABLE_COUNTERS*
- *CONFIG_SPI_FLASH_ROM_DRIVER_PATCH*
- *CONFIG_SPI_FLASH_DANGEROUS_WRITE*
- *CONFIG_SPI_FLASH_USE_LEGACY_IMPL*
- *CONFIG_SPI_FLASH_BYPASS_BLOCK_ERASE*
- *CONFIG_SPI_FLASH_YIELD_DURING_ERASE*
- *CONFIG_SPI_FLASH_SIZE_OVERRIDE*
- *Auto-detect flash chips*
- *CONFIG_SPI_FLASH_ENABLE_ENCRYPTED_READ_WRITE*

CONFIG_SPI_FLASH_VERIFY_WRITE

Verify SPI flash writes

Found in: Component config > SPI Flash driver

If this option is enabled, any time SPI flash is written then the data will be read back and verified. This can catch hardware problems with SPI flash, or flash which was not erased before verification.

CONFIG_SPI_FLASH_LOG_FAILED_WRITE

Log errors if verification fails

Found in: Component config > SPI Flash driver > CONFIG_SPI_FLASH_VERIFY_WRITE

If this option is enabled, if SPI flash write verification fails then a log error line will be written with the address, expected & actual values. This can be useful when debugging hardware SPI flash problems.

CONFIG_SPI_FLASH_WARN_SETTING_ZERO_TO_ONE

Log warning if writing zero bits to ones

Found in: Component config > SPI Flash driver > CONFIG_SPI_FLASH_VERIFY_WRITE

If this option is enabled, any SPI flash write which tries to set zero bits in the flash to ones will log a warning. Such writes will not result in the requested data appearing identically in flash once written, as SPI NOR flash can only set bits to one when an entire sector is erased. After erasing, individual bits can only be written from one to zero.

Note that some software (such as SPIFFS) which is aware of SPI NOR flash may write one bits as an optimisation, relying on the data in flash becoming a bitwise AND of the new data and any existing data. Such software will log spurious warnings if this option is enabled.

CONFIG_SPI_FLASH_ENABLE_COUNTERS

Enable operation counters

Found in: Component config > SPI Flash driver

This option enables the following APIs:

- spi_flash_reset_counters
- spi_flash_dump_counters
- spi_flash_get_counters

These APIs may be used to collect performance data for spi_flash APIs and to help understand behaviour of libraries which use SPI flash.

CONFIG_SPI_FLASH_ROM_DRIVER_PATCH

Enable SPI flash ROM driver patched functions

Found in: Component config > SPI Flash driver

Enable this flag to use patched versions of SPI flash ROM driver functions. This option should be enabled, if any one of the following is true: (1) need to write to flash on ESP32-D2WD; (2) main SPI flash is connected to non-default pins; (3) main SPI flash chip is manufactured by ISSI.

CONFIG_SPI_FLASH_DANGEROUS_WRITE

Writing to dangerous flash regions

Found in: Component config > SPI Flash driver

SPI flash APIs can optionally abort or return a failure code if erasing or writing addresses that fall at the beginning of flash (covering the bootloader and partition table) or that overlap the app partition that contains the running app.

It is not recommended to ever write to these regions from an IDF app, and this check prevents logic errors or corrupted firmware memory from damaging these regions.

Note that this feature *does not* check calls to the esp_rom_xxx SPI flash ROM functions. These functions should not be called directly from IDF applications.

Available options:

- Aborts (SPI_FLASH_DANGEROUS_WRITE_ABORTS)
- Fails (SPI_FLASH_DANGEROUS_WRITE_FAILS)
- Allowed (SPI_FLASH_DANGEROUS_WRITE_ALLOWED)

CONFIG_SPI_FLASH_USE_LEGACY_IMPL

Use the legacy implementation before IDF v4.0

Found in: Component config > SPI Flash driver

The implementation of SPI flash has been greatly changed in IDF v4.0. Enable this option to use the legacy implementation.

CONFIG_SPI_FLASH_BYPASS_BLOCK_ERASE

Bypass a block erase and always do sector erase

Found in: Component config > SPI Flash driver

Some flash chips can have very high "max" erase times, especially for block erase (32KB or 64KB). This option allows to bypass "block erase" and always do sector erase commands. This will be much slower overall in most cases, but improves latency for other code to run.

CONFIG_SPI_FLASH_YIELD_DURING_ERASE

Enables yield operation during flash erase

Found in: Component config > SPI Flash driver

This allows to yield the CPUs between erase commands. Prevents starvation of other tasks.

CONFIG_SPI_FLASH_ERASE_YIELD_DURATION_MS

Duration of erasing to yield CPUs (ms)

Found in: Component config > SPI Flash driver > CONFIG_SPI_FLASH_YIELD_DURING_ERASE

If a duration of one erase command is large then it will yield CPUs after finishing a current command.

CONFIG_SPI_FLASH_ERASE_YIELD_TICKS

CPU release time (tick)

Found in: Component config > SPI Flash driver > CONFIG_SPI_FLASH_YIELD_DURING_ERASE

Defines how many ticks will be before returning to continue a erasing.

CONFIG_SPI_FLASH_SIZE_OVERRIDE

Override flash size in bootloader header by ESPTOOLPY_FLASHSIZE

Found in: Component config > SPI Flash driver

SPI Flash driver uses the flash size configured in bootloader header by default. Enable this option to override flash size with latest ESPTOOLPY_FLASHSIZE value from the app header if the size in the bootloader header is incorrect.

Auto-detect flash chips Contains:

- *CONFIG_SPI_FLASH_SUPPORT_ISSI_CHIP*
- *CONFIG_SPI_FLASH_SUPPORT_MXIC_CHIP*
- *CONFIG_SPI_FLASH_SUPPORT_GD_CHIP*

CONFIG_SPI_FLASH_SUPPORT_ISSI_CHIP

ISSI

Found in: Component config > SPI Flash driver > Auto-detect flash chips

Enable this to support auto detection of ISSI chips if chip vendor not directly given by `chip\_drv` member of the chip struct. This adds support for variant chips, however will extend detecting time.

CONFIG_SPI_FLASH_SUPPORT_MXIC_CHIP

MXIC

Found in: Component config > SPI Flash driver > Auto-detect flash chips

Enable this to support auto detection of MXIC chips if chip vendor not directly given by `chip\_drv` member of the chip struct. This adds support for variant chips, however will extend detecting time.

CONFIG_SPI_FLASH_SUPPORT_GD_CHIP

GigaDevice

Found in: Component config > SPI Flash driver > Auto-detect flash chips

Enable this to support auto detection of GD (GigaDevice) chips if chip vendor not directly given by `chip\_drv` member of the chip struct. If you are using Wrover modules, please don't disable this, otherwise your flash may not work in 4-bit mode.

This adds support for variant chips, however will extend detecting time and image size. Note that the default chip driver supports the GD chips with product ID 60H.

CONFIG_SPI_FLASH_ENABLE_ENCRYPTED_READ_WRITE

Enable encrypted partition read/write operations

Found in: Component config > SPI Flash driver

This option enables flash read/write operations to encrypted partition/s. This option is kept enabled irrespective of state of flash encryption feature. However, in case application is not using flash encryption feature and is in need of some additional memory from IRAM region (~1KB) then this config can be disabled.

PThreads Contains:

- *CONFIG_PTHREAD_TASK_PRIO_DEFAULT*
- *CONFIG_PTHREAD_TASK_STACK_SIZE_DEFAULT*
- *CONFIG_PTHREAD_STACK_MIN*
- *CONFIG_PTHREAD_TASK_CORE_DEFAULT*
- *CONFIG_PTHREAD_TASK_NAME_DEFAULT*

CONFIG_PTHREAD_TASK_PRIO_DEFAULT

Default task priority

Found in: Component config > PThreads

Priority used to create new tasks with default pthread parameters.

CONFIG_PTHREAD_TASK_STACK_SIZE_DEFAULT

Default task stack size

Found in: Component config > PThreads

Stack size used to create new tasks with default pthread parameters.

CONFIG_PTHREAD_STACK_MIN

Minimum allowed pthread stack size

Found in: Component config > PThreads

Minimum allowed pthread stack size set in attributes passed to pthread_create

CONFIG_PTHREAD_TASK_CORE_DEFAULT

Default pthread core affinity

Found in: Component config > PThreads

The default core to which pthreads are pinned.

Available options:
- No affinity (PTHREAD_DEFAULT_CORE_NO_AFFINITY)
- Core 0 (PTHREAD_DEFAULT_CORE_0)
- Core 1 (PTHREAD_DEFAULT_CORE_1)

CONFIG_PTHREAD_TASK_NAME_DEFAULT

Default name of pthreads

Found in: Component config > PThreads

The default name of pthreads.

NVS Contains:

- *CONFIG_NVS_ENCRYPTION*

CONFIG_NVS_ENCRYPTION

Enable NVS encryption

Found in: Component config > NVS

This option enables encryption for NVS. When enabled, AES-XTS is used to encrypt the complete NVS data, except the page headers. It requires XTS encryption keys to be stored in an encrypted partition. This means enabling flash encryption is a pre-requisite for this feature.

Newlib Contains:

- *CONFIG_NEWLIB_STDOUT_LINE_ENDING*
- *CONFIG_NEWLIB_STDIN_LINE_ENDING*
- *CONFIG_NEWLIB_NANO_FORMAT*

CONFIG_NEWLIB_STDOUT_LINE_ENDING

Line ending for UART output

Found in: Component config > Newlib

This option allows configuring the desired line endings sent to UART when a newline ('n', LF) appears on stdout. Three options are possible:

CRLF: whenever LF is encountered, prepend it with CR

LF: no modification is applied, stdout is sent as is

CR: each occurence of LF is replaced with CR

This option doesn't affect behavior of the UART driver (drivers/uart.h).

Available options:
- CRLF (NEWLIB_STDOUT_LINE_ENDING_CRLF)
- LF (NEWLIB_STDOUT_LINE_ENDING_LF)
- CR (NEWLIB_STDOUT_LINE_ENDING_CR)

CONFIG_NEWLIB_STDIN_LINE_ENDING

Line ending for UART input

Found in: Component config > Newlib

This option allows configuring which input sequence on UART produces a newline ('n', LF) on stdin. Three options are possible:

CRLF: CRLF is converted to LF

LF: no modification is applied, input is sent to stdin as is

CR: each occurence of CR is replaced with LF

This option doesn't affect behavior of the UART driver (drivers/uart.h).

Available options:

- CRLF (NEWLIB_STDIN_LINE_ENDING_CRLF)
- LF (NEWLIB_STDIN_LINE_ENDING_LF)
- CR (NEWLIB_STDIN_LINE_ENDING_CR)

CONFIG_NEWLIB_NANO_FORMAT

Enable 'nano' formatting options for printf/scanf family

Found in: Component config > Newlib

ESP32 ROM contains parts of newlib C library, including printf/scanf family of functions. These functions have been compiled with so-called "nano" formatting option. This option doesn't support 64-bit integer formats and C99 features, such as positional arguments.

For more details about "nano" formatting option, please see newlib readme file, search for '–enable-newlib-nano-formatted-io': https://sourceware.org/newlib/README

If this option is enabled, build system will use functions available in ROM, reducing the application binary size. Functions available in ROM run faster than functions which run from flash. Functions available in ROM can also run when flash instruction cache is disabled.

If you need 64-bit integer formatting support or C99 features, keep this option disabled.

ESP-MQTT Configurations

Contains:

- *CONFIG_MQTT_PROTOCOL_311*
- *CONFIG_MQTT_TRANSPORT_SSL*
- *CONFIG_MQTT_TRANSPORT_WEBSOCKET*
- *CONFIG_MQTT_USE_CUSTOM_CONFIG*
- *CONFIG_MQTT_TASK_CORE_SELECTION_ENABLED*
- *CONFIG_MQTT_CUSTOM_OUTBOX*

CONFIG_MQTT_PROTOCOL_311

Enable MQTT protocol 3.1.1

Found in: Component config > ESP-MQTT Configurations

If not, this library will use MQTT protocol 3.1

CONFIG_MQTT_TRANSPORT_SSL

Enable MQTT over SSL

Found in: Component config > ESP-MQTT Configurations

Enable MQTT transport over SSL with mbedtls

CONFIG_MQTT_TRANSPORT_WEBSOCKET

Enable MQTT over Websocket

Found in: Component config > ESP-MQTT Configurations

Enable MQTT transport over Websocket.

CONFIG_MQTT_TRANSPORT_WEBSOCKET_SECURE

Enable MQTT over Websocket Secure

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_TRANSPORT_WEBSOCKET

Enable MQTT transport over Websocket Secure.

CONFIG_MQTT_USE_CUSTOM_CONFIG

MQTT Using custom configurations

Found in: Component config > ESP-MQTT Configurations

Custom MQTT configurations.

CONFIG_MQTT_TCP_DEFAULT_PORT

Default MQTT over TCP port

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

Default MQTT over TCP port

CONFIG_MQTT_SSL_DEFAULT_PORT

Default MQTT over SSL port

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

Default MQTT over SSL port

CONFIG_MQTT_WS_DEFAULT_PORT

Default MQTT over Websocket port

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

Default MQTT over Websocket port

CONFIG_MQTT_WSS_DEFAULT_PORT

Default MQTT over Websocket Secure port

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

Default MQTT over Websocket Secure port

CONFIG_MQTT_BUFFER_SIZE

Default MQTT Buffer Size

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

This buffer size using for both transmit and receive

CONFIG_MQTT_TASK_STACK_SIZE

MQTT task stack size

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

MQTT task stack size

CONFIG_MQTT_DISABLE_API_LOCKS

Disable API locks

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

Default config employs API locks to protect internal structures. It is possible to disable these locks if the user code doesn't access MQTT API from multiple concurrent tasks

CONFIG_MQTT_TASK_PRIORITY

MQTT task priority

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_USE_CUSTOM_CONFIG

MQTT task priority. Higher number denotes higher priority.

CONFIG_MQTT_TASK_CORE_SELECTION_ENABLED

Enable MQTT task core selection

Found in: Component config > ESP-MQTT Configurations

This will enable core selection

CONFIG_MQTT_TASK_CORE_SELECTION

Core to use ?

Found in: Component config > ESP-MQTT Configurations > CONFIG_MQTT_TASK_CORE_SELECTION_ENABLED

Available options:

- Core 0 (MQTT_USE_CORE_0)
- Core 1 (MQTT_USE_CORE_1)

CONFIG_MQTT_CUSTOM_OUTBOX

Enable custom outbox implementation

Found in: Component config > ESP-MQTT Configurations

Set to true if a specific implementation of message outbox is needed (e.g. persistant outbox in NVM or similar).

mbedTLS Contains:

- *CONFIG_MBEDTLS_MEM_ALLOC_MODE*
- *CONFIG_MBEDTLS_SSL_MAX_CONTENT_LEN*
- *CONFIG_MBEDTLS_ASYMMETRIC_CONTENT_LEN*
- *CONFIG_MBEDTLS_DYNAMIC_BUFFER*
- *CONFIG_MBEDTLS_DEBUG*
- *Certificate Bundle*
- *CONFIG_MBEDTLS_ECP_RESTARTABLE*
- *CONFIG_MBEDTLS_CMAC_C*
- *CONFIG_MBEDTLS_HARDWARE_AES*
- *CONFIG_MBEDTLS_HARDWARE_MPI*
- *CONFIG_MBEDTLS_HARDWARE_SHA*
- *CONFIG_MBEDTLS_ATCA_HW_ECDSA_SIGN*
- *CONFIG_MBEDTLS_ATCA_HW_ECDSA_VERIFY*
- *CONFIG_MBEDTLS_HAVE_TIME*
- *CONFIG_MBEDTLS_ECDSA_DETERMINISTIC*
- *CONFIG_MBEDTLS_SHA512_C*

- *CONFIG_MBEDTLS_TLS_MODE*
- *TLS Key Exchange Methods*
- *CONFIG_MBEDTLS_SSL_RENEGOTIATION*
- *CONFIG_MBEDTLS_SSL_PROTO_SSL3*
- *CONFIG_MBEDTLS_SSL_PROTO_TLS1*
- *CONFIG_MBEDTLS_SSL_PROTO_TLS1_1*
- *CONFIG_MBEDTLS_SSL_PROTO_TLS1_2*
- *CONFIG_MBEDTLS_SSL_PROTO_DTLS*
- *CONFIG_MBEDTLS_SSL_ALPN*
- *CONFIG_MBEDTLS_CLIENT_SSL_SESSION_TICKETS*
- *CONFIG_MBEDTLS_SERVER_SSL_SESSION_TICKETS*
- *Symmetric Ciphers*
- *CONFIG_MBEDTLS_RIPEMD160_C*
- *Certificates*
- *CONFIG_MBEDTLS_ECP_C*
- *CONFIG_MBEDTLS_POLY1305_C*
- *CONFIG_MBEDTLS_CHACHA20_C*
- *CONFIG_MBEDTLS_HKDF_C*
- *CONFIG_MBEDTLS_THREADING_C*
- *CONFIG_MBEDTLS_SECURITY_RISKS*

CONFIG_MBEDTLS_MEM_ALLOC_MODE

Memory allocation strategy

Found in: Component config > mbedTLS

Allocation strategy for mbedTLS, essentially provides ability to allocate all required dynamic allocations from,

- Internal DRAM memory only
- External SPIRAM memory only
- Either internal or external memory based on default malloc() behavior in ESP-IDF
- Custom allocation mode, by overwriting calloc()/free() using mbedtls_platform_set_calloc_free() function
- Internal IRAM memory wherever applicable else internal DRAM

Recommended mode here is always internal (*), since that is most preferred from security perspective. But if application requirement does not allow sufficient free internal memory then alternate mode can be selected.

(*) In case of ESP32-S2, hardware allows encryption of external SPIRAM contents provided hardware flash encryption feature is enabled. In that case, using external SPIRAM allocation strategy is also safe choice from security perspective.

Available options:

- Internal memory (MBEDTLS_INTERNAL_MEM_ALLOC)
- External SPIRAM (MBEDTLS_EXTERNAL_MEM_ALLOC)
- Default alloc mode (MBEDTLS_DEFAULT_MEM_ALLOC)
- Custom alloc mode (MBEDTLS_CUSTOM_MEM_ALLOC)
- Internal IRAM (MBEDTLS_IRAM_8BIT_MEM_ALLOC)
  Allows to use IRAM memory region as 8bit accessible region.
  TLS input and output buffers will be allocated in IRAM section which is 32bit aligned memory. Every unaligned (8bit or 16bit) access will result in an exception and incur penalty of certain clock cycles per unaligned read/write.

CONFIG_MBEDTLS_SSL_MAX_CONTENT_LEN

TLS maximum message content length

Found in: Component config > mbedTLS

Maximum TLS message length (in bytes) supported by mbedTLS.

16384 is the default and this value is required to comply fully with TLS standards.

However you can set a lower value in order to save RAM. This is safe if the other end of the connection supports Maximum Fragment Length Negotiation Extension (max_fragment_length, see RFC6066) or you know for certain that it will never send a message longer than a certain number of bytes.

If the value is set too low, symptoms are a failed TLS handshake or a return value of MBEDTLS_ERR_SSL_INVALID_RECORD (-0x7200).

CONFIG_MBEDTLS_ASYMMETRIC_CONTENT_LEN

Asymmetric in/out fragment length

Found in: Component config > mbedTLS

If enabled, this option allows customizing TLS in/out fragment length in asymmetric way. Please note that enabling this with default values saves 12KB of dynamic memory per TLS connection.

CONFIG_MBEDTLS_SSL_IN_CONTENT_LEN

TLS maximum incoming fragment length

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ASYMMETRIC_CONTENT_LEN

This defines maximum incoming fragment length, overriding default maximum content length (MBEDTLS_SSL_MAX_CONTENT_LEN).

CONFIG_MBEDTLS_SSL_OUT_CONTENT_LEN

TLS maximum outgoing fragment length

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ASYMMETRIC_CONTENT_LEN

This defines maximum outgoing fragment length, overriding default maximum content length (MBEDTLS_SSL_MAX_CONTENT_LEN).

CONFIG_MBEDTLS_DYNAMIC_BUFFER

Using dynamic TX/RX buffer

Found in: Component config > mbedTLS

Using dynamic TX/RX buffer. After enabling this option, mbedTLS will allocate TX buffer when need to send data and then free it if all data is sent, allocate RX buffer when need to receive data and then free it when all data is used or read by upper layer.

By default, when SSL is initialized, mbedTLS also allocate TX and RX buffer with the default value of “MBEDTLS_SSL_OUT_CONTENT_LEN” or “MBEDTLS_SSL_IN_CONTENT_LEN” , so to save more heap, users can set the options to be an appropriate value.

CONFIG_MBEDTLS_DYNAMIC_FREE_PEER_CERT

Free SSL peer certificate after its usage

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_DYNAMIC_BUFFER

Free peer certificate after its usage in handshake process.

CONFIG_MBEDTLS_DYNAMIC_FREE_CONFIG_DATA

Free private key and DHM data after its usage

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_DYNAMIC_BUFFER

Free private key and DHM data after its usage in handshake process.

The option will decrease heap cost when handshake, but also lead to problem:

Becasue all certificate, private key and DHM data are freed so users should register certificate and private key to ssl config object again.

CONFIG_MBEDTLS_DYNAMIC_FREE_CA_CERT

Free SSL CA certificate after its usage

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_DYNAMIC_BUFFER > CONFIG_MBEDTLS_DYNAMIC_FREE_CONFIG_DATA

Free CA certificate after its usage in the handshake process. This option will decrease the heap footprint for the TLS handshake, but may lead to a problem: If the respective ssl object needs to perform the TLS handshake again, the CA certificate should once again be registered to the ssl object.

CONFIG_MBEDTLS_DEBUG

Enable mbedTLS debugging

Found in: Component config > mbedTLS

Enable mbedTLS debugging functions at compile time.

If this option is enabled, you can include "mbedtls/esp_debug.h" and call mbedtls_esp_enable_debug_log() at runtime in order to enable mbedTLS debug output via the ESP log mechanism.

CONFIG_MBEDTLS_DEBUG_LEVEL

Set mbedTLS debugging level

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_DEBUG

Set mbedTLS debugging level

Available options:

- Warning (MBEDTLS_DEBUG_LEVEL_WARN)
- Info (MBEDTLS_DEBUG_LEVEL_INFO)
- Debug (MBEDTLS_DEBUG_LEVEL_DEBUG)
- Verbose (MBEDTLS_DEBUG_LEVEL_VERBOSE)

Certificate Bundle Contains:

- *CONFIG_MBEDTLS_CERTIFICATE_BUNDLE*

CONFIG_MBEDTLS_CERTIFICATE_BUNDLE

Enable trusted root certificate bundle

Found in: Component config > mbedTLS > Certificate Bundle

Enable support for large number of default root certificates

When enabled this option allows user to store default as well as customer specific root certificates in compressed format rather than storing full certificate. For the root certificates the public key and the subject name will be stored.

CONFIG_MBEDTLS_DEFAULT_CERTIFICATE_BUNDLE

Default certificate bundle options

Found in: Component config > mbedTLS > Certificate Bundle > CONFIG_MBEDTLS_CERTIFICATE_BUNDLE

Available options:

- Use the full default certificate bundle (MBEDTLS_CERTIFICATE_BUNDLE_DEFAULT_FULL)
- Use only the most common certificates from the default bundles (MBEDTLS_CERTIFICATE_BUNDLE_DEFAULT_CMN)
  Use only the most common certificates from the default bundles, reducing the size with 50%, while still having around 99% coverage.
- Do not use the default certificate bundle (MBEDTLS_CERTIFICATE_BUNDLE_DEFAULT_NONE)

CONFIG_MBEDTLS_CUSTOM_CERTIFICATE_BUNDLE

Add custom certificates to the default bundle

Found in: Component config > mbedTLS > Certificate Bundle > CONFIG_MBEDTLS_CERTIFICATE_BUNDLE

CONFIG_MBEDTLS_CUSTOM_CERTIFICATE_BUNDLE_PATH

Custom certificate bundle path

Found in: Component config > mbedTLS > Certificate Bundle > CONFIG_MBEDTLS_CERTIFICATE_BUNDLE > CONFIG_MBEDTLS_CUSTOM_CERTIFICATE_BUNDLE

Name of the custom certificate directory or file. This path is evaluated relative to the project root directory.

CONFIG_MBEDTLS_ECP_RESTARTABLE

Enable mbedTLS ecp restartable

Found in: Component config > mbedTLS

Enable "non-blocking" ECC operations that can return early and be resumed.

CONFIG_MBEDTLS_CMAC_C

Enable CMAC mode for block ciphers

Found in: Component config > mbedTLS

Enable the CMAC (Cipher-based Message Authentication Code) mode for block ciphers.

CONFIG_MBEDTLS_HARDWARE_AES

Enable hardware AES acceleration

Found in: Component config > mbedTLS

Enable hardware accelerated AES encryption & decryption.

Note that if the ESP32 CPU is running at 240MHz, hardware AES does not offer any speed boost over software AES.

CONFIG_MBEDTLS_AES_USE_INTERRUPT

Use interrupt for long AES operations

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_HARDWARE_AES

Use an interrupt to coordinate long AES operations.

This allows other code to run on the CPU while an AES operation is pending. Otherwise the CPU busy-waits.

CONFIG_MBEDTLS_HARDWARE_GCM

Enable partially hardware accelerated GCM

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_HARDWARE_AES

Enable partially hardware accelerated GCM. GHASH calculation is still done in software.

If MBEDTLS_HARDWARE_GCM is disabled and MBEDTLS_HARDWARE_AES is enabled then mbedTLS will still use the hardware accelerated AES block operation, but on a single block at a time.

CONFIG_MBEDTLS_HARDWARE_MPI

Enable hardware MPI (bignum) acceleration

Found in: Component config > mbedTLS

Enable hardware accelerated multiple precision integer operations.

Hardware accelerated multiplication, modulo multiplication, and modular exponentiation for up to 4096 bit results.

These operations are used by RSA.

CONFIG_MBEDTLS_HARDWARE_SHA

Enable hardware SHA acceleration

Found in: Component config > mbedTLS

Enable hardware accelerated SHA1, SHA256, SHA384 & SHA512 in mbedTLS.

Due to a hardware limitation, on the ESP32 hardware acceleration is only guaranteed if SHA digests are calculated one at a time. If more than one SHA digest is calculated at the same time, one will be calculated fully in hardware and the rest will be calculated (at least partially calculated) in software. This happens automatically.

SHA hardware acceleration is faster than software in some situations but slower in others. You should benchmark to find the best setting for you.

CONFIG_MBEDTLS_ATCA_HW_ECDSA_SIGN

Enable hardware ECDSA sign acceleration when using ATECC608A

Found in: Component config > mbedTLS

This option enables hardware acceleration for ECDSA sign function, only when using ATECC608A cryptoauth chip (integrated with ESP32-WROOM-32SE)

CONFIG_MBEDTLS_ATCA_HW_ECDSA_VERIFY

Enable hardware ECDSA verify acceleration when using ATECC608A

Found in: Component config > mbedTLS

This option enables hardware acceleration for ECDSA sign function, only when using ATECC608A cryptoauth chip (integrated with ESP32-WROOM-32SE)

CONFIG_MBEDTLS_HAVE_TIME

Enable mbedtls time

Found in: Component config > mbedTLS

System has time.h and time(). The time does not need to be correct, only time differences are used.

CONFIG_MBEDTLS_HAVE_TIME_DATE

Enable mbedtls certificate expiry check

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_HAVE_TIME

System has time.h and time(), gmtime() and the clock is correct. The time needs to be correct (not necesarily very accurate, but at least the date should be correct). This is used to verify the validity period of X.509 certificates.

It is suggested that you should get the real time by “SNTP”.

CONFIG_MBEDTLS_ECDSA_DETERMINISTIC

Enable deterministic ECDSA

Found in: Component config > mbedTLS

Standard ECDSA is “fragile” in the sense that lack of entropy when signing may result in a compromise of the long-term signing key.

CONFIG_MBEDTLS_SHA512_C

Enable the SHA-384 and SHA-512 cryptographic hash algorithms

Found in: Component config > mbedTLS

Enable MBEDTLS_SHA512_C adds support for SHA-384 and SHA-512.

CONFIG_MBEDTLS_TLS_MODE

TLS Protocol Role

Found in: Component config > mbedTLS

mbedTLS can be compiled with protocol support for the TLS server, TLS client, or both server and client.

Reducing the number of TLS roles supported saves code size.

Available options:

- Server & Client (MBEDTLS_TLS_SERVER_AND_CLIENT)
- Server (MBEDTLS_TLS_SERVER_ONLY)
- Client (MBEDTLS_TLS_CLIENT_ONLY)
- None (MBEDTLS_TLS_DISABLED)

TLS Key Exchange Methods Contains:

- *CONFIG_MBEDTLS_PSK_MODES*
- *CONFIG_MBEDTLS_KEY_EXCHANGE_RSA*
- *CONFIG_MBEDTLS_KEY_EXCHANGE_DHE_RSA*
- *CONFIG_MBEDTLS_KEY_EXCHANGE_ELLIPTIC_CURVE*
- *CONFIG_MBEDTLS_KEY_EXCHANGE_ECJPAKE*

CONFIG_MBEDTLS_PSK_MODES

Enable pre-shared-key ciphersuites

Found in: Component config > mbedTLS > TLS Key Exchange Methods

Enable to show configuration for different types of pre-shared-key TLS authentatication methods.

Leaving this options disabled will save code size if they are not used.

CONFIG_MBEDTLS_KEY_EXCHANGE_PSK

Enable PSK based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_PSK_MODES

Enable to support symmetric key PSK (pre-shared-key) TLS key exchange modes.

CONFIG_MBEDTLS_KEY_EXCHANGE_DHE_PSK

Enable DHE-PSK based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_PSK_MODES

Enable to support Diffie-Hellman PSK (pre-shared-key) TLS authentication modes.

CONFIG_MBEDTLS_KEY_EXCHANGE_ECDHE_PSK

Enable ECDHE-PSK based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_PSK_MODES

Enable to support Elliptic-Curve-Diffie-Hellman PSK (pre-shared-key) TLS authentication modes.

CONFIG_MBEDTLS_KEY_EXCHANGE_RSA_PSK

Enable RSA-PSK based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_PSK_MODES

Enable to support RSA PSK (pre-shared-key) TLS authentication modes.

CONFIG_MBEDTLS_KEY_EXCHANGE_RSA

Enable RSA-only based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods

Enable to support ciphersuites with prefix TLS-RSA-WITH-

CONFIG_MBEDTLS_KEY_EXCHANGE_DHE_RSA

Enable DHE-RSA based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods

Enable to support ciphersuites with prefix TLS-DHE-RSA-WITH-

CONFIG_MBEDTLS_KEY_EXCHANGE_ELLIPTIC_CURVE

Support Elliptic Curve based ciphersuites

Found in: Component config > mbedTLS > TLS Key Exchange Methods

Enable to show Elliptic Curve based ciphersuite mode options.

Disabling all Elliptic Curve ciphersuites saves code size and can give slightly faster TLS handshakes, provided the server supports RSA-only ciphersuite modes.

CONFIG_MBEDTLS_KEY_EXCHANGE_ECDHE_RSA

Enable ECDHE-RSA based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_KEY_EXCHANGE_ELLIPTIC_CURVE

Enable to support ciphersuites with prefix TLS-ECDHE-RSA-WITH-

CONFIG_MBEDTLS_KEY_EXCHANGE_ECDHE_ECDSA

Enable ECDHE-ECDSA based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_KEY_EXCHANGE_ELLIPTIC_CURVE

Enable to support ciphersuites with prefix TLS-ECDHE-RSA-WITH-

CONFIG_MBEDTLS_KEY_EXCHANGE_ECDH_ECDSA

Enable ECDH-ECDSA based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_KEY_EXCHANGE_ELLIPTIC_CURVE

Enable to support ciphersuites with prefix TLS-ECDHE-RSA-WITH-

CONFIG_MBEDTLS_KEY_EXCHANGE_ECDH_RSA

Enable ECDH-RSA based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods > CONFIG_MBEDTLS_KEY_EXCHANGE_ELLIPTIC_CURVE

Enable to support ciphersuites with prefix TLS-ECDHE-RSA-WITH-

CONFIG_MBEDTLS_KEY_EXCHANGE_ECJPAKE

Enable ECJPAKE based ciphersuite modes

Found in: Component config > mbedTLS > TLS Key Exchange Methods

Enable to support ciphersuites with prefix TLS-ECJPAKE-WITH-

CONFIG_MBEDTLS_SSL_RENEGOTIATION

Support TLS renegotiation

Found in: Component config > mbedTLS

The two main uses of renegotiation are (1) refresh keys on long-lived connections and (2) client authentication after the initial handshake. If you don't need renegotiation, disabling it will save code size and reduce the possibility of abuse/vulnerability.

CONFIG_MBEDTLS_SSL_PROTO_SSL3

Legacy SSL 3.0 support

Found in: Component config > mbedTLS

Support the legacy SSL 3.0 protocol. Most servers will speak a newer TLS protocol these days.

CONFIG_MBEDTLS_SSL_PROTO_TLS1

Support TLS 1.0 protocol

Found in: Component config > mbedTLS

CONFIG_MBEDTLS_SSL_PROTO_TLS1_1

Support TLS 1.1 protocol

Found in: Component config > mbedTLS

CONFIG_MBEDTLS_SSL_PROTO_TLS1_2

Support TLS 1.2 protocol

Found in: Component config > mbedTLS

CONFIG_MBEDTLS_SSL_PROTO_DTLS

Support DTLS protocol (all versions)

Found in: Component config > mbedTLS

Requires TLS 1.1 to be enabled for DTLS 1.0 Requires TLS 1.2 to be enabled for DTLS 1.2

CONFIG_MBEDTLS_SSL_ALPN

Support ALPN (Application Layer Protocol Negotiation)

Found in: Component config > mbedTLS

Disabling this option will save some code size if it is not needed.

CONFIG_MBEDTLS_CLIENT_SSL_SESSION_TICKETS

TLS: Client Support for RFC 5077 SSL session tickets

Found in: Component config > mbedTLS

Client support for RFC 5077 session tickets. See mbedTLS documentation for more details. Disabling this option will save some code size.

CONFIG_MBEDTLS_SERVER_SSL_SESSION_TICKETS

TLS: Server Support for RFC 5077 SSL session tickets

Found in: Component config > mbedTLS

Server support for RFC 5077 session tickets. See mbedTLS documentation for more details. Disabling this option will save some code size.

Symmetric Ciphers Contains:

- *CONFIG_MBEDTLS_AES_C*
- *CONFIG_MBEDTLS_CAMELLIA_C*
- *CONFIG_MBEDTLS_DES_C*
- *CONFIG_MBEDTLS_RC4_MODE*
- *CONFIG_MBEDTLS_BLOWFISH_C*
- *CONFIG_MBEDTLS_XTEA_C*
- *CONFIG_MBEDTLS_CCM_C*
- *CONFIG_MBEDTLS_GCM_C*

CONFIG_MBEDTLS_AES_C

AES block cipher

Found in: Component config > mbedTLS > Symmetric Ciphers

CONFIG_MBEDTLS_CAMELLIA_C

Camellia block cipher

Found in: Component config > mbedTLS > Symmetric Ciphers

CONFIG_MBEDTLS_DES_C

DES block cipher (legacy, insecure)

Found in: Component config > mbedTLS > Symmetric Ciphers

Enables the DES block cipher to support 3DES-based TLS ciphersuites.

3DES is vulnerable to the Sweet32 attack and should only be enabled if absolutely necessary.

CONFIG_MBEDTLS_RC4_MODE

RC4 Stream Cipher (legacy, insecure)

Found in: Component config > mbedTLS > Symmetric Ciphers

ARCFOUR (RC4) stream cipher can be disabled entirely, enabled but not added to default ciphersuites, or enabled completely.

Please consider the security implications before enabling RC4.

Available options:

- Disabled (MBEDTLS_RC4_DISABLED)
- Enabled, not in default ciphersuites (MBEDTLS_RC4_ENABLED_NO_DEFAULT)
- Enabled (MBEDTLS_RC4_ENABLED)

CONFIG_MBEDTLS_BLOWFISH_C

Blowfish block cipher (read help)

Found in: Component config > mbedTLS > Symmetric Ciphers

Enables the Blowfish block cipher (not used for TLS sessions.)

The Blowfish cipher is not used for mbedTLS TLS sessions but can be used for other purposes. Read up on the limitations of Blowfish (including Sweet32) before enabling.

CONFIG_MBEDTLS_XTEA_C

XTEA block cipher

Found in: Component config > mbedTLS > Symmetric Ciphers

Enables the XTEA block cipher.

CONFIG_MBEDTLS_CCM_C

CCM (Counter with CBC-MAC) block cipher modes

Found in: Component config > mbedTLS > Symmetric Ciphers

Enable Counter with CBC-MAC (CCM) modes for AES and/or Camellia ciphers.

Disabling this option saves some code size.

CONFIG_MBEDTLS_GCM_C

GCM (Galois/Counter) block cipher modes

Found in: Component config > mbedTLS > Symmetric Ciphers

Enable Galois/Counter Mode for AES and/or Camellia ciphers.

This option is generally faster than CCM.

CONFIG_MBEDTLS_RIPEMD160_C

Enable RIPEMD-160 hash algorithm

Found in: Component config > mbedTLS

Enable the RIPEMD-160 hash algorithm.

Certificates Contains:

- *CONFIG_MBEDTLS_PEM_PARSE_C*
- *CONFIG_MBEDTLS_PEM_WRITE_C*
- *CONFIG_MBEDTLS_X509_CRL_PARSE_C*
- *CONFIG_MBEDTLS_X509_CSR_PARSE_C*

CONFIG_MBEDTLS_PEM_PARSE_C

Read & Parse PEM formatted certificates

Found in: Component config > mbedTLS > Certificates

Enable decoding/parsing of PEM formatted certificates.

If your certificates are all in the simpler DER format, disabling this option will save some code size.

CONFIG_MBEDTLS_PEM_WRITE_C

Write PEM formatted certificates

Found in: Component config > mbedTLS > Certificates

Enable writing of PEM formatted certificates.

If writing certificate data only in DER format, disabling this option will save some code size.

CONFIG_MBEDTLS_X509_CRL_PARSE_C

X.509 CRL parsing

Found in: Component config > mbedTLS > Certificates

Support for parsing X.509 Certifificate Revocation Lists.

CONFIG_MBEDTLS_X509_CSR_PARSE_C

X.509 CSR parsing

Found in: Component config > mbedTLS > Certificates

Support for parsing X.509 Certifificate Signing Requests

CONFIG_MBEDTLS_ECP_C

Elliptic Curve Ciphers

Found in: Component config > mbedTLS

Contains:

- *CONFIG_MBEDTLS_ECDH_C*
- *CONFIG_MBEDTLS_ECJPAKE_C*
- *CONFIG_MBEDTLS_ECP_DP_SECP192R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP224R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP256R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP384R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP521R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP192K1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP224K1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_SECP256K1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_BP256R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_BP384R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_BP512R1_ENABLED*
- *CONFIG_MBEDTLS_ECP_DP_CURVE25519_ENABLED*
- *CONFIG_MBEDTLS_ECP_NIST_OPTIM*

CONFIG_MBEDTLS_ECDH_C

Elliptic Curve Diffie-Hellman (ECDH)

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable ECDH. Needed to use ECDHE-xxx TLS ciphersuites.

CONFIG_MBEDTLS_ECDSA_C

Elliptic Curve DSA

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C > CONFIG_MBEDTLS_ECDH_C

Enable ECDSA. Needed to use ECDSA-xxx TLS ciphersuites.

CONFIG_MBEDTLS_ECJPAKE_C

Elliptic curve J-PAKE

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable ECJPAKE. Needed to use ECJPAKE-xxx TLS ciphersuites.

CONFIG_MBEDTLS_ECP_DP_SECP192R1_ENABLED

Enable SECP192R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP192R1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP224R1_ENABLED

Enable SECP224R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP224R1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP256R1_ENABLED

Enable SECP256R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP256R1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP384R1_ENABLED

Enable SECP384R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP384R1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP521R1_ENABLED

Enable SECP521R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP521R1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP192K1_ENABLED

Enable SECP192K1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP192K1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP224K1_ENABLED

Enable SECP224K1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP224K1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_SECP256K1_ENABLED

Enable SECP256K1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for SECP256K1 Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_BP256R1_ENABLED

Enable BP256R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

support for DP Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_BP384R1_ENABLED

Enable BP384R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

support for DP Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_BP512R1_ENABLED

Enable BP512R1 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

support for DP Elliptic Curve.

CONFIG_MBEDTLS_ECP_DP_CURVE25519_ENABLED

Enable CURVE25519 curve

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

Enable support for CURVE25519 Elliptic Curve.

CONFIG_MBEDTLS_ECP_NIST_OPTIM

NIST ‘modulo p’ optimisations

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_ECP_C

NIST ‘modulo p’ optimisations increase Elliptic Curve operation performance.

Disabling this option saves some code size.

end of Elliptic Curve options

CONFIG_MBEDTLS_POLY1305_C

Poly1305 MAC algorithm

Found in: Component config > mbedTLS

Enable support for Poly1305 MAC algorithm.

CONFIG_MBEDTLS_CHACHA20_C

Chacha20 stream cipher

Found in: Component config > mbedTLS

Enable support for Chacha20 stream cipher.

CONFIG_MBEDTLS_CHACHAPOLY_C

ChaCha20-Poly1305 AEAD algorithm

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_CHACHA20_C

Enable support for ChaCha20-Poly1305 AEAD algorithm.

CONFIG_MBEDTLS_HKDF_C

HKDF algorithm (RFC 5869)

Found in: Component config > mbedTLS

Enable support for the Hashed Message Authentication Code (HMAC)-based key derivation function (HKDF).

CONFIG_MBEDTLS_THREADING_C

Enable the threading abstraction layer

Found in: Component config > mbedTLS

If you do intend to use contexts between threads, you will need to enable this layer to prevent race conditions.

CONFIG_MBEDTLS_THREADING_ALT

Enable threading alternate implementation

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_THREADING_C

Enable threading alt to allow your own alternate threading implementation.

CONFIG_MBEDTLS_THREADING_PTHREAD

Enable threading pthread implementation

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_THREADING_C

Enable the pthread wrapper layer for the threading layer.

CONFIG_MBEDTLS_SECURITY_RISKS

Show configurations with potential security risks

Found in: Component config > mbedTLS

Contains:

- *CONFIG_MBEDTLS_ALLOW_UNSUPPORTED_CRITICAL_EXT*

CONFIG_MBEDTLS_ALLOW_UNSUPPORTED_CRITICAL_EXT

X.509 CRT parsing with unsupported critical extensions

Found in: Component config > mbedTLS > CONFIG_MBEDTLS_SECURITY_RISKS

Allow the X.509 certificate parser to load certificates with unsupported critical extensions

LWIP Contains:

- *CONFIG_LWIP_LOCAL_HOSTNAME*
- *CONFIG_LWIP_DNS_SUPPORT_MDNS_QUERIES*
- *CONFIG_LWIP_L2_TO_L3_COPY*
- *CONFIG_LWIP_IRAM_OPTIMIZATION*
- *CONFIG_LWIP_TIMERS_ONDEMAND*
- *CONFIG_LWIP_MAX_SOCKETS*
- *CONFIG_LWIP_USE_ONLY_LWIP_SELECT*
- *CONFIG_LWIP_SO_LINGER*
- *CONFIG_LWIP_SO_REUSE*
- *CONFIG_LWIP_SO_RCVBUF*
- *CONFIG_LWIP_NETBUF_RECVINFO*
- *CONFIG_LWIP_IP4_FRAG*
- *CONFIG_LWIP_IP6_FRAG*
- *CONFIG_LWIP_IP4_REASSEMBLY*
- *CONFIG_LWIP_IP6_REASSEMBLY*
- *CONFIG_LWIP_IP_FORWARD*
- *CONFIG_LWIP_STATS*
- *CONFIG_LWIP_ETHARP_TRUST_IP_MAC*
- *CONFIG_LWIP_ESP_GRATUITOUS_ARP*
- *CONFIG_LWIP_TCPIP_RECVMBOX_SIZE*
- *CONFIG_LWIP_DHCP_DOES_ARP_CHECK*
- *CONFIG_LWIP_DHCP_DISABLE_CLIENT_ID*
- *CONFIG_LWIP_DHCP_RESTORE_LAST_IP*
- *CONFIG_LWIP_DHCP_COARSE_TIMER_SECS*
- *DHCP server*
- *CONFIG_LWIP_AUTOIP*
- *CONFIG_LWIP_IPV6_AUTOCONFIG*
- *CONFIG_LWIP_NETIF_LOOPBACK*
- *TCP*
- *UDP*
- *CONFIG_LWIP_TCPIP_TASK_STACK_SIZE*
- *CONFIG_LWIP_TCPIP_TASK_AFFINITY*
- *CONFIG_LWIP_PPP_SUPPORT*
- *CONFIG_LWIP_IPV6_MEMP_NUM_ND6_QUEUE*
- *CONFIG_LWIP_IPV6_ND6_NUM_NEIGHBORS*
- *CONFIG_LWIP_PPP_NOTIFY_PHASE_SUPPORT*
- *CONFIG_LWIP_PPP_PAP_SUPPORT*
- *CONFIG_LWIP_PPP_CHAP_SUPPORT*
- *CONFIG_LWIP_PPP_MSCHAP_SUPPORT*
- *CONFIG_LWIP_PPP_MPPE_SUPPORT*
- *CONFIG_LWIP_PPP_DEBUG_ON*

- *ICMP*
- *LWIP RAW API*
- *SNTP*
- *CONFIG_LWIP_ESP_LWIP_ASSERT*

CONFIG_LWIP_LOCAL_HOSTNAME

Local netif hostname

Found in: Component config > LWIP

The default name this device will report to other devices on the network. Could be updated at runtime with esp_netif_set_hostname()

CONFIG_LWIP_DNS_SUPPORT_MDNS_QUERIES

Enable mDNS queries in resolving host name

Found in: Component config > LWIP

If this feature is enabled, standard API such as gethostbyname support .local addresses by sending one shot multicast mDNS query

CONFIG_LWIP_L2_TO_L3_COPY

Enable copy between Layer2 and Layer3 packets

Found in: Component config > LWIP

If this feature is enabled, all traffic from layer2(WIFI Driver) will be copied to a new buffer before sending it to layer3(LWIP stack), freeing the layer2 buffer. Please be notified that the total layer2 receiving buffer is fixed and ESP32 currently supports 25 layer2 receiving buffer, when layer2 buffer runs out of memory, then the incoming packets will be dropped in hardware. The layer3 buffer is allocated from the heap, so the total layer3 receiving buffer depends on the available heap size, when heap runs out of memory, no copy will be sent to layer3 and packet will be dropped in layer2. Please make sure you fully understand the impact of this feature before enabling it.

CONFIG_LWIP_IRAM_OPTIMIZATION

Enable LWIP IRAM optimization

Found in: Component config > LWIP

If this feature is enabled, some functions relating to RX/TX in LWIP will be put into IRAM, it can improve UDP/TCP throughput by >10% for single core mode, it doesn't help too much for dual core mode. On the other hand, it needs about 10KB IRAM for these optimizations.

If this feature is disabled, all lwip functions will be put into FLASH.

CONFIG_LWIP_TIMERS_ONDEMAND

Enable LWIP Timers on demand

Found in: Component config > LWIP

If this feature is enabled, IGMP and MLD6 timers will be activated only when joining groups or receiving QUERY packets.

This feature will reduce the power consumption for applications which do not use IGMP and MLD6.

CONFIG_LWIP_MAX_SOCKETS

Max number of open sockets

Found in: Component config > LWIP

Sockets take up a certain amount of memory, and allowing fewer sockets to be open at the same time conserves memory. Specify the maximum amount of sockets here. The valid value is from 1 to 16.

CONFIG_LWIP_USE_ONLY_LWIP_SELECT

Support LWIP socket select() only (DEPRECATED)

Found in: Component config > LWIP

This option is deprecated. Use VFS_SUPPORT_SELECT instead, which is the inverse of this option.

The virtual filesystem layer of select() redirects sockets to lwip_select() and non-socket file descriptors to their respective driver implementations. If this option is enabled then all calls of select() will be redirected to lwip_select(), therefore, select can be used for sockets only.

CONFIG_LWIP_SO_LINGER

Enable SO_LINGER processing

Found in: Component config > LWIP

Enabling this option allows SO_LINGER processing. l_onoff = 1,l_linger can set the timeout.

If l_linger=0, When a connection is closed, TCP will terminate the connection. This means that TCP will discard any data packets stored in the socket send buffer and send an RST to the peer.

If l_linger!=0,Then closesocket() calls to block the process until the remaining data packets has been sent or timed out.

CONFIG_LWIP_SO_REUSE

Enable SO_REUSEADDR option

Found in: Component config > LWIP

Enabling this option allows binding to a port which remains in TIME_WAIT.

CONFIG_LWIP_SO_REUSE_RXTOALL

SO_REUSEADDR copies broadcast/multicast to all matches

Found in: Component config > LWIP > CONFIG_LWIP_SO_REUSE

Enabling this option means that any incoming broadcast or multicast packet will be copied to all of the local sockets that it matches (may be more than one if SO_REUSEADDR is set on the socket.)

This increases memory overhead as the packets need to be copied, however they are only copied per matching socket. You can safely disable it if you don't plan to receive broadcast or multicast traffic on more than one socket at a time.

CONFIG_LWIP_SO_RCVBUF

Enable SO_RCVBUF option

Found in: Component config > LWIP

Enabling this option allows checking for available data on a netconn.

CONFIG_LWIP_NETBUF_RECVINFO

Enable IP_PKTINFO option

Found in: Component config > LWIP

Enabling this option allows checking for the destination address of a received IPv4 Packet.

CONFIG_LWIP_IP4_FRAG

Enable fragment outgoing IP4 packets

Found in: Component config > LWIP

Enabling this option allows fragmenting outgoing IP4 packets if their size exceeds MTU.

CONFIG_LWIP_IP6_FRAG

Enable fragment outgoing IP6 packets

Found in: Component config > LWIP

Enabling this option allows fragmenting outgoing IP6 packets if their size exceeds MTU.

CONFIG_LWIP_IP4_REASSEMBLY

Enable reassembly incoming fragmented IP4 packets

Found in: Component config > LWIP

Enabling this option allows reassemblying incoming fragmented IP4 packets.

CONFIG_LWIP_IP6_REASSEMBLY

Enable reassembly incoming fragmented IP6 packets

Found in: Component config > LWIP

Enabling this option allows reassemblying incoming fragmented IP6 packets.

CONFIG_LWIP_IP_FORWARD

Enable IP forwarding

Found in: Component config > LWIP

Enabling this option allows packets forwarding across multiple interfaces.

CONFIG_LWIP_IPV4_NAPT

Enable NAT (new/experimental)

Found in: Component config > LWIP > CONFIG_LWIP_IP_FORWARD

Enabling this option allows Network Address and Port Translation.

CONFIG_LWIP_STATS

Enable LWIP statistics

Found in: Component config > LWIP

Enabling this option allows LWIP statistics

CONFIG_LWIP_ETHARP_TRUST_IP_MAC

Enable LWIP ARP trust

Found in: Component config > LWIP

Enabling this option allows ARP table to be updated.

If this option is enabled, the incoming IP packets cause the ARP table to be updated with the source MAC and IP addresses supplied in the packet. You may want to disable this if you do not trust LAN peers to have the correct addresses, or as a limited approach to attempt to handle spoofing. If disabled, lwIP will need to make a new ARP request if the peer is not already in the ARP table, adding a little latency. The peer *is* in the ARP table if it requested our address before. Also notice that this slows down input processing of every IP packet!

There are two known issues in real application if this feature is enabled: - The LAN peer may have bug to update the ARP table after the ARP entry is aged out. If the ARP entry on the LAN peer is aged out but failed to be updated, all IP packets sent from LWIP to the LAN peer will be dropped by LAN peer. - The LAN peer may not be trustful, the LAN peer may send IP packets to LWIP with two different MACs, but the same IP address. If this happens, the LWIP has problem to receive IP packets from LAN peer.

So the recommendation is to disable this option. Here the LAN peer means the other side to which the ESP station or soft-AP is connected.

CONFIG_LWIP_ESP_GRATUITOUS_ARP

Send gratuitous ARP periodically

Found in: Component config > LWIP

Enable this option allows to send gratuitous ARP periodically.

This option solve the compatibility issues.If the ARP table of the AP is old, and the AP doesn’t send ARP request to update it’s ARP table, this will lead to the STA sending IP packet fail. Thus we send gratuitous ARP periodically to let AP update it’s ARP table.

CONFIG_LWIP_GARP_TMR_INTERVAL

GARP timer interval(seconds)

Found in: Component config > LWIP > CONFIG_LWIP_ESP_GRATUITOUS_ARP

Set the timer interval for gratuitous ARP. The default value is 60s

CONFIG_LWIP_TCPIP_RECVMBOX_SIZE

TCPIP task receive mail box size

Found in: Component config > LWIP

Set TCPIP task receive mail box size. Generally bigger value means higher throughput but more memory. The value should be bigger than UDP/TCP mail box size.

CONFIG_LWIP_DHCP_DOES_ARP_CHECK

DHCP: Perform ARP check on any offered address

Found in: Component config > LWIP

Enabling this option performs a check (via ARP request) if the offered IP address is not already in use by another host on the network.

CONFIG_LWIP_DHCP_DISABLE_CLIENT_ID

DHCP: Disable Use of HW address as client identification

Found in: Component config > LWIP

This option could be used to disable DHCP client identification with its MAC address. (Client id is used by DHCP servers to uniquely identify clients and are included in the DHCP packets as an option 61) Set this option to "y" in order to exclude option 61 from DHCP packets.

CONFIG_LWIP_DHCP_RESTORE_LAST_IP

DHCP: Restore last IP obtained from DHCP server

Found in: Component config > LWIP

When this option is enabled, DHCP client tries to re-obtain last valid IP address obtained from DHCP server. Last valid DHCP configuration is stored in nvs and restored after reset/power-up. If IP is still available, there is no need for sending discovery message to DHCP server and save some time.

CONFIG_LWIP_DHCP_COARSE_TIMER_SECS

DHCP coarse timer interval(s)

Found in: Component config > LWIP

Set DHCP coarse interval in seconds. A higher value will be less precise but cost less power consumption.

DHCP server Contains:

- *CONFIG_LWIP_DHCPS_LEASE_UNIT*
- *CONFIG_LWIP_DHCPS_MAX_STATION_NUM*

CONFIG_LWIP_DHCPS_LEASE_UNIT

Multiplier for lease time, in seconds

Found in: Component config > LWIP > DHCP server

The DHCP server is calculating lease time multiplying the sent and received times by this number of seconds per unit. The default is 60, that equals one minute.

CONFIG_LWIP_DHCPS_MAX_STATION_NUM

Maximum number of stations

Found in: Component config > LWIP > DHCP server

The maximum number of DHCP clients that are connected to the server. After this number is exceeded, DHCP server removes of the oldest device from it's address pool, without notification.

CONFIG_LWIP_AUTOIP

Enable IPV4 Link-Local Addressing (AUTOIP)

Found in: Component config > LWIP

Enabling this option allows the device to self-assign an address in the 169.256/16 range if none is assigned statically or via DHCP.

See RFC 3927.

Contains:

- *CONFIG_LWIP_AUTOIP_TRIES*

- *CONFIG_LWIP_AUTOIP_MAX_CONFLICTS*
- *CONFIG_LWIP_AUTOIP_RATE_LIMIT_INTERVAL*

CONFIG_LWIP_AUTOIP_TRIES

DHCP Probes before self-assigning IPv4 LL address

Found in: Component config > LWIP > CONFIG_LWIP_AUTOIP

DHCP client will send this many probes before self-assigning a link local address.

From LWIP help: "This can be set as low as 1 to get an AutoIP address very quickly, but you should be prepared to handle a changing IP address when DHCP overrides AutoIP." (In the case of ESP-IDF, this means multiple SYSTEM_EVENT_STA_GOT_IP events.)

CONFIG_LWIP_AUTOIP_MAX_CONFLICTS

Max IP conflicts before rate limiting

Found in: Component config > LWIP > CONFIG_LWIP_AUTOIP

If the AUTOIP functionality detects this many IP conflicts while self-assigning an address, it will go into a rate limited mode.

CONFIG_LWIP_AUTOIP_RATE_LIMIT_INTERVAL

Rate limited interval (seconds)

Found in: Component config > LWIP > CONFIG_LWIP_AUTOIP

If rate limiting self-assignment requests, wait this long between each request.

CONFIG_LWIP_IPV6_AUTOCONFIG

Enable IPV6 stateless address autoconfiguration

Found in: Component config > LWIP

Enabling this option allows the devices to IPV6 stateless address autoconfiguration.

See RFC 4862.

CONFIG_LWIP_NETIF_LOOPBACK

Support per-interface loopback

Found in: Component config > LWIP

Enabling this option means that if a packet is sent with a destination address equal to the interface's own IP address, it will "loop back" and be received by this interface.

Contains:

- *CONFIG_LWIP_LOOPBACK_MAX_PBUFS*

CONFIG_LWIP_LOOPBACK_MAX_PBUFS

Max queued loopback packets per interface

Found in: Component config > LWIP > CONFIG_LWIP_NETIF_LOOPBACK

Configure the maximum number of packets which can be queued for loopback on a given interface. Reducing this number may cause packets to be dropped, but will avoid filling memory with queued packet data.

TCP Contains:

- *CONFIG_LWIP_TCP_ISN_HOOK*
- *CONFIG_LWIP_MAX_ACTIVE_TCP*
- *CONFIG_LWIP_MAX_LISTENING_TCP*
- *CONFIG_LWIP_TCP_HIGH_SPEED_RETRANSMISSION*
- *CONFIG_LWIP_TCP_MAXRTX*
- *CONFIG_LWIP_TCP_SYNMAXRTX*
- *CONFIG_LWIP_TCP_MSS*
- *CONFIG_LWIP_TCP_TMR_INTERVAL*
- *CONFIG_LWIP_TCP_MSL*
- *CONFIG_LWIP_TCP_FIN_WAIT_TIMEOUT*
- *CONFIG_LWIP_TCP_SND_BUF_DEFAULT*
- *CONFIG_LWIP_TCP_WND_DEFAULT*
- *CONFIG_LWIP_TCP_RECVMBOX_SIZE*
- *CONFIG_LWIP_TCP_QUEUE_OOSEQ*
- *CONFIG_LWIP_TCP_SACK_OUT*
- *CONFIG_LWIP_TCP_KEEP_CONNECTION_WHEN_IP_CHANGES*
- *CONFIG_LWIP_TCP_OVERSIZE*
- *CONFIG_LWIP_WND_SCALE*
- *CONFIG_LWIP_TCP_RTO_TIME*

CONFIG_LWIP_TCP_ISN_HOOK

Enable TCP ISN Hook

Found in: Component config > LWIP > TCP

Enables custom TCP ISN hook to randomize initial sequence number in TCP connection. This is recommended as default lwIP implementation (*tcp_next_iss*) is not very strong, as it does not take into consideration any platform specific entropy source.

CONFIG_LWIP_MAX_ACTIVE_TCP

Maximum active TCP Connections

Found in: Component config > LWIP > TCP

The maximum number of simultaneously active TCP connections. The practical maximum limit is determined by available heap memory at runtime.

Changing this value by itself does not substantially change the memory usage of LWIP, except for preventing new TCP connections after the limit is reached.

CONFIG_LWIP_MAX_LISTENING_TCP

Maximum listening TCP Connections

Found in: Component config > LWIP > TCP

The maximum number of simultaneously listening TCP connections. The practical maximum limit is determined by available heap memory at runtime.

Changing this value by itself does not substantially change the memory usage of LWIP, except for preventing new listening TCP connections after the limit is reached.

CONFIG_LWIP_TCP_HIGH_SPEED_RETRANSMISSION

TCP high speed retransmissions

Found in: Component config > LWIP > TCP

Speed up the TCP retransmission interval. If disabled, it is recommended to change the number of SYN retransmissions to 6, and TCP initial rto time to 3000.

CONFIG_LWIP_TCP_MAXRTX

Maximum number of retransmissions of data segments

Found in: Component config > LWIP > TCP

Set maximum number of retransmissions of data segments.

CONFIG_LWIP_TCP_SYNMAXRTX

Maximum number of retransmissions of SYN segments

Found in: Component config > LWIP > TCP

Set maximum number of retransmissions of SYN segments.

CONFIG_LWIP_TCP_MSS

Maximum Segment Size (MSS)

Found in: Component config > LWIP > TCP

Set maximum segment size for TCP transmission.

Can be set lower to save RAM, the default value 1460(ipv4)/1440(ipv6) will give best throughput. IPv4 TCP_MSS Range: 576 <= TCP_MSS <= 1460 IPv6 TCP_MSS Range: 1220<= TCP_mSS <= 1440

CONFIG_LWIP_TCP_TMR_INTERVAL

TCP timer interval(ms)

Found in: Component config > LWIP > TCP

Set TCP timer interval in milliseconds.

Can be used to speed connections on bad networks. A lower value will redeliver unacked packets faster.

CONFIG_LWIP_TCP_MSL

Maximum segment lifetime (MSL)

Found in: Component config > LWIP > TCP

Set maximum segment lifetime in milliseconds.

CONFIG_LWIP_TCP_FIN_WAIT_TIMEOUT

Maximum FIN segment lifetime

Found in: Component config > LWIP > TCP

Set maximum segment lifetime in milliseconds.

CONFIG_LWIP_TCP_SND_BUF_DEFAULT

Default send buffer size

Found in: Component config > LWIP > TCP

Set default send buffer size for new TCP sockets.

Per-socket send buffer size can be changed at runtime with lwip_setsockopt(s, TCP_SNDBUF, ⋯).

This value must be at least 2x the MSS size, and the default is 4x the default MSS size.

Setting a smaller default SNDBUF size can save some RAM, but will decrease performance.

CONFIG_LWIP_TCP_WND_DEFAULT

Default receive window size

Found in: Component config > LWIP > TCP

Set default TCP receive window size for new TCP sockets.

Per-socket receive window size can be changed at runtime with lwip_setsockopt(s, TCP_WINDOW, ⋯).

Setting a smaller default receive window size can save some RAM, but will significantly decrease performance.

CONFIG_LWIP_TCP_RECVMBOX_SIZE

Default TCP receive mail box size

Found in: Component config > LWIP > TCP

Set TCP receive mail box size. Generally bigger value means higher throughput but more memory. The recommended value is: LWIP_TCP_WND_DEFAULT/TCP_MSS + 2, e.g. if LWIP_TCP_WND_DEFAULT=14360, TCP_MSS=1436, then the recommended receive mail box size is (14360/1436 + 2) = 12.

TCP receive mail box is a per socket mail box, when the application receives packets from TCP socket, LWIP core firstly posts the packets to TCP receive mail box and the application then fetches the packets from mail box. It means LWIP can caches maximum LWIP_TCP_RECCVMBOX_SIZE packets for each TCP socket, so the maximum possible cached TCP packets for all TCP sockets is LWIP_TCP_RECCVMBOX_SIZE multiples the maximum TCP socket number. In other words, the bigger LWIP_TCP_RECVMBOX_SIZE means more memory. On the other hand, if the receiv mail box is too small, the mail box may be full. If the mail box is full, the LWIP drops the packets. So generally we need to make sure the TCP receive mail box is big enough to avoid packet drop between LWIP core and application.

CONFIG_LWIP_TCP_QUEUE_OOSEQ

Queue incoming out-of-order segments

Found in: Component config > LWIP > TCP

Queue incoming out-of-order segments for later use.

Disable this option to save some RAM during TCP sessions, at the expense of increased retransmissions if segments arrive out of order.

CONFIG_LWIP_TCP_SACK_OUT

Support sending selective acknowledgements

Found in: Component config > LWIP > TCP

TCP will support sending selective acknowledgements (SACKs).

CONFIG_LWIP_TCP_KEEP_CONNECTION_WHEN_IP_CHANGES

Keep TCP connections when IP changed

Found in: Component config > LWIP > TCP

This option is enabled when the following scenario happen: network dropped and reconnected, IP changes is like: 192.168.0.2->0.0.0.0->192.168.0.2

Disable this option to keep consistent with the original LWIP code behavior.

CONFIG_LWIP_TCP_OVERSIZE

Pre-allocate transmit PBUF size

Found in: Component config > LWIP > TCP

Allows enabling "oversize" allocation of TCP transmission pbufs ahead of time, which can reduce the length of pbuf chains used for transmission.

This will not make a difference to sockets where Nagle's algorithm is disabled.

Default value of MSS is fine for most applications, 25% MSS may save some RAM when only transmitting small amounts of data. Disabled will have worst performance and fragmentation characteristics, but uses least RAM overall.

Available options:

- MSS (LWIP_TCP_OVERSIZE_MSS)
- 25% MSS (LWIP_TCP_OVERSIZE_QUARTER_MSS)
- Disabled (LWIP_TCP_OVERSIZE_DISABLE)

CONFIG_LWIP_WND_SCALE

Support TCP window scale

Found in: Component config > LWIP > TCP

Enable this feature to support TCP window scaling.

CONFIG_LWIP_TCP_RCV_SCALE

Set TCP receiving window scaling factor

Found in: Component config > LWIP > TCP > CONFIG_LWIP_WND_SCALE

Enable this feature to support TCP window scaling.

CONFIG_LWIP_TCP_RTO_TIME

Default TCP rto time

Found in: Component config > LWIP > TCP

Set default TCP rto time for a reasonable initial rto. In bad network environment, recommend set value of rto time to 1500.

UDP Contains:

- *CONFIG_LWIP_MAX_UDP_PCBS*
- *CONFIG_LWIP_UDP_RECVMBOX_SIZE*

CONFIG_LWIP_MAX_UDP_PCBS

Maximum active UDP control blocks

Found in: Component config > LWIP > UDP

The maximum number of active UDP "connections" (ie UDP sockets sending/receiving data). The practical maximum limit is determined by available heap memory at runtime.

CONFIG_LWIP_UDP_RECVMBOX_SIZE

Default UDP receive mail box size

Found in: Component config > LWIP > UDP

Set UDP receive mail box size. The recommended value is 6.

UDP receive mail box is a per socket mail box, when the application receives packets from UDP socket, LWIP core firstly posts the packets to UDP receive mail box and the application then fetches the packets from mail box. It means LWIP can caches maximum UDP_RECCVMBOX_SIZE packets for each UDP socket, so the maximum possible cached UDP packets for all UDP sockets is UDP_RECCVMBOX_SIZE multiples the maximum UDP socket number. In other words, the bigger UDP_RECVMBOX_SIZE means more memory. On the other hand, if the receiv mail box is too small, the mail box may be full. If the mail box is full, the LWIP drops the packets. So generally we need to make sure the UDP receive mail box is big enough to avoid packet drop between LWIP core and application.

CONFIG_LWIP_TCPIP_TASK_STACK_SIZE

TCP/IP Task Stack Size

Found in: Component config > LWIP

Configure TCP/IP task stack size, used by LWIP to process multi-threaded TCP/IP operations. Setting this stack too small will result in stack overflow crashes.

CONFIG_LWIP_TCPIP_TASK_AFFINITY

TCP/IP task affinity

Found in: Component config > LWIP

Allows setting LwIP tasks affinity, i.e. whether the task is pinned to CPU0, pinned to CPU1, or allowed to run on any CPU. Currently this applies to "TCP/IP" task and "Ping" task.

Available options:

- No affinity (LWIP_TCPIP_TASK_AFFINITY_NO_AFFINITY)
- CPU0 (LWIP_TCPIP_TASK_AFFINITY_CPU0)
- CPU1 (LWIP_TCPIP_TASK_AFFINITY_CPU1)

CONFIG_LWIP_PPP_SUPPORT

Enable PPP support (new/experimental)

Found in: Component config > LWIP

Enable PPP stack. Now only PPP over serial is possible.

PPP over serial support is experimental and unsupported.

Contains:

- *CONFIG_LWIP_PPP_ENABLE_IPV6*

CONFIG_LWIP_PPP_ENABLE_IPV6

Enable IPV6 support for PPP connections (IPV6CP)

Found in: Component config > LWIP > CONFIG_LWIP_PPP_SUPPORT

Enable IPV6 support in PPP for the local link between the DTE (processor) and DCE (modem). There are some modems which do not support the IPV6 addressing in the local link. If they are requested for IPV6CP negotiation, they may time out. This would in turn fail the configuration for the whole link. If your modem is not responding correctly to PPP Phase Network, try to disable IPV6 support.

CONFIG_LWIP_IPV6_MEMP_NUM_ND6_QUEUE

Max number of IPv6 packets to queue during MAC resolution

Found in: Component config > LWIP

Config max number of IPv6 packets to queue during MAC resolution.

CONFIG_LWIP_IPV6_ND6_NUM_NEIGHBORS

Max number of entries in IPv6 neighbor cache

Found in: Component config > LWIP

Config max number of entries in IPv6 neighbor cache

CONFIG_LWIP_PPP_NOTIFY_PHASE_SUPPORT

Enable Notify Phase Callback

Found in: Component config > LWIP

Enable to set a callback which is called on change of the internal PPP state machine.

CONFIG_LWIP_PPP_PAP_SUPPORT

Enable PAP support

Found in: Component config > LWIP

Enable Password Authentication Protocol (PAP) support

CONFIG_LWIP_PPP_CHAP_SUPPORT

Enable CHAP support

Found in: Component config > LWIP

Enable Challenge Handshake Authentication Protocol (CHAP) support

CONFIG_LWIP_PPP_MSCHAP_SUPPORT

Enable MSCHAP support

Found in: Component config > LWIP

Enable Microsoft version of the Challenge-Handshake Authentication Protocol (MSCHAP) support

CONFIG_LWIP_PPP_MPPE_SUPPORT

Enable MPPE support

Found in: Component config > LWIP

Enable Microsoft Point-to-Point Encryption (MPPE) support

CONFIG_LWIP_PPP_DEBUG_ON

Enable PPP debug log output

Found in: Component config > LWIP

Enable PPP debug log output

ICMP Contains:

- *CONFIG_LWIP_MULTICAST_PING*
- *CONFIG_LWIP_BROADCAST_PING*

CONFIG_LWIP_MULTICAST_PING

Respond to multicast pings

Found in: Component config > LWIP > ICMP

CONFIG_LWIP_BROADCAST_PING

Respond to broadcast pings

Found in: Component config > LWIP > ICMP

LWIP RAW API Contains:

- *CONFIG_LWIP_MAX_RAW_PCBS*

CONFIG_LWIP_MAX_RAW_PCBS

Maximum LWIP RAW PCBs

Found in: Component config > LWIP > LWIP RAW API

The maximum number of simultaneously active LWIP RAW protocol control blocks. The practical maximum limit is determined by available heap memory at runtime.

SNTP Contains:

- *CONFIG_LWIP_DHCP_MAX_NTP_SERVERS*
- *CONFIG_LWIP_SNTP_UPDATE_DELAY*

CONFIG_LWIP_DHCP_MAX_NTP_SERVERS

Maximum number of NTP servers

Found in: Component config > LWIP > SNTP

Set maximum number of NTP servers used by LwIP SNTP module. First argument of sntp_setserver/sntp_setservername functions is limited to this value.

CONFIG_LWIP_SNTP_UPDATE_DELAY

Request interval to update time (ms)

Found in: Component config > LWIP > SNTP

This option allows you to set the time update period via SNTP. Default is 1 hour. Must not be below 15 seconds by specification. (SNTPv4 RFC 4330 enforces a minimum update time of 15 seconds).

CONFIG_LWIP_ESP_LWIP_ASSERT

Enable LWIP ASSERT checks

Found in: Component config > LWIP

Enable this option allows lwip to check assert. It is recommended to keep it open, do not close it.

Log output Contains:

- *CONFIG_LOG_DEFAULT_LEVEL*
- *CONFIG_LOG_COLORS*
- *CONFIG_LOG_TIMESTAMP_SOURCE*

CONFIG_LOG_DEFAULT_LEVEL

Default log verbosity

Found in: Component config > Log output

Specify how much output to see in logs by default. You can set lower verbosity level at runtime using esp_log_level_set function.

Note that this setting limits which log statements are compiled into the program. So setting this to, say, "Warning" would mean that changing log level to "Debug" at runtime will not be possible.

Available options:

- No output (LOG_DEFAULT_LEVEL_NONE)
- Error (LOG_DEFAULT_LEVEL_ERROR)
- Warning (LOG_DEFAULT_LEVEL_WARN)
- Info (LOG_DEFAULT_LEVEL_INFO)
- Debug (LOG_DEFAULT_LEVEL_DEBUG)
- Verbose (LOG_DEFAULT_LEVEL_VERBOSE)

CONFIG_LOG_COLORS

Use ANSI terminal colors in log output

Found in: Component config > Log output

Enable ANSI terminal color codes in bootloader output.

In order to view these, your terminal program must support ANSI color codes.

CONFIG_LOG_TIMESTAMP_SOURCE

Log Timestamps

Found in: Component config > Log output

Choose what sort of timestamp is displayed in the log output:

- Milliseconds since boot is calulated from the RTOS tick count multiplied by the tick period. This time will reset after a software reboot. e.g. (90000)

- System time is taken from POSIX time functions which use the ESP32's RTC and FRC1 timers to maintain an accurate time. The system time is initialized to 0 on startup, it can be set with an SNTP sync, or with POSIX time functions. This time will not reset after a software reboot. e.g. (00:01:30.000)
- NOTE: Currently this will not get used in logging from binary blobs (i.e WiFi & Bluetooth libraries), these will always print milliseconds since boot.

Available options:
- Milliseconds Since Boot (LOG_TIMESTAMP_SOURCE_RTOS)
- System Time (LOG_TIMESTAMP_SOURCE_SYSTEM)

Heap memory debugging Contains:

- *CONFIG_HEAP_CORRUPTION_DETECTION*
- *CONFIG_HEAP_TRACING_DEST*
- *CONFIG_HEAP_TRACING_STACK_DEPTH*
- *CONFIG_HEAP_TASK_TRACKING*
- *CONFIG_HEAP_ABORT_WHEN_ALLOCATION_FAILS*

CONFIG_HEAP_CORRUPTION_DETECTION

Heap corruption detection

Found in: Component config > Heap memory debugging

Enable heap poisoning features to detect heap corruption caused by out-of-bounds access to heap memory.

See the "Heap Memory Debugging" page of the IDF documentation for a description of each level of heap corruption detection.

Available options:
- Basic (no poisoning) (HEAP_POISONING_DISABLED)
- Light impact (HEAP_POISONING_LIGHT)
- Comprehensive (HEAP_POISONING_COMPREHENSIVE)

CONFIG_HEAP_TRACING_DEST

Heap tracing

Found in: Component config > Heap memory debugging

Enables the heap tracing API defined in esp_heap_trace.h.

This function causes a moderate increase in IRAM code side and a minor increase in heap function (malloc/free/realloc) CPU overhead, even when the tracing feature is not used. So it's best to keep it disabled unless tracing is being used.

Available options:
- Disabled (HEAP_TRACING_OFF)
- Standalone (HEAP_TRACING_STANDALONE)
- Host-based (HEAP_TRACING_TOHOST)

CONFIG_HEAP_TRACING_STACK_DEPTH

Heap tracing stack depth

Found in: Component config > Heap memory debugging

Number of stack frames to save when tracing heap operation callers.

More stack frames uses more memory in the heap trace buffer (and slows down allocation), but can provide useful information.

CONFIG_HEAP_TASK_TRACKING

Enable heap task tracking

Found in: Component config > Heap memory debugging

Enables tracking the task responsible for each heap allocation.

This function depends on heap poisoning being enabled and adds four more bytes of overhead for each block allocated.

CONFIG_HEAP_ABORT_WHEN_ALLOCATION_FAILS

Abort if memory allocation fails

Found in: Component config > Heap memory debugging

When enabled, if a memory allocation operation fails it will cause a system abort.

FreeRTOS Contains:

- *CONFIG_FREERTOS_UNICORE*
- *CONFIG_FREERTOS_CORETIMER*
- *CONFIG_FREERTOS_OPTIMIZED_SCHEDULER*
- *CONFIG_FREERTOS_HZ*
- *CONFIG_FREERTOS_ASSERT_ON_UNTESTED_FUNCTION*
- *CONFIG_FREERTOS_CHECK_STACKOVERFLOW*
- *CONFIG_FREERTOS_WATCHPOINT_END_OF_STACK*
- *CONFIG_FREERTOS_INTERRUPT_BACKTRACE*
- *CONFIG_FREERTOS_THREAD_LOCAL_STORAGE_POINTERS*
- *CONFIG_FREERTOS_ASSERT*
- *CONFIG_FREERTOS_IDLE_TASK_STACKSIZE*
- *CONFIG_FREERTOS_ISR_STACKSIZE*
- *CONFIG_FREERTOS_LEGACY_HOOKS*
- *CONFIG_FREERTOS_MAX_TASK_NAME_LEN*
- *CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION*
- *CONFIG_FREERTOS_TIMER_TASK_PRIORITY*
- *CONFIG_FREERTOS_TIMER_TASK_STACK_DEPTH*
- *CONFIG_FREERTOS_TIMER_QUEUE_LENGTH*
- *CONFIG_FREERTOS_QUEUE_REGISTRY_SIZE*
- *CONFIG_FREERTOS_USE_TRACE_FACILITY*
- *CONFIG_FREERTOS_GENERATE_RUN_TIME_STATS*
- *CONFIG_FREERTOS_USE_TICKLESS_IDLE*
- *CONFIG_FREERTOS_TASK_FUNCTION_WRAPPER*
- *CONFIG_FREERTOS_CHECK_MUTEX_GIVEN_BY_OWNER*
- *CONFIG_FREERTOS_CHECK_PORT_CRITICAL_COMPLIANCE*

CONFIG_FREERTOS_UNICORE

Run FreeRTOS only on first core

Found in: Component config > FreeRTOS

This version of FreeRTOS normally takes control of all cores of the CPU. Select this if you only want to start it on the first core. This is needed when e.g. another process needs complete control over the second core.

This invisible config value sets the value of tskNO_AFFINITY in task.h. # Intended to be used as a constant from other Kconfig files. # Value is (32-bit) INT_MAX.

CONFIG_FREERTOS_CORETIMER

Xtensa timer to use as the FreeRTOS tick source

Found in: Component config > FreeRTOS

FreeRTOS needs a timer with an associated interrupt to use as the main tick source to increase counters, run timers and do pre-emptive multitasking with. There are multiple timers available to do this, with different interrupt priorities. Check

Available options:

- Timer 0 (int 6, level 1) (FREERTOS_CORETIMER_0)
  Select this to use timer 0
- Timer 1 (int 15, level 3) (FREERTOS_CORETIMER_1)
  Select this to use timer 1

CONFIG_FREERTOS_OPTIMIZED_SCHEDULER

Enable FreeRTOS plátform optimized scheduler

Found in: Component config > FreeRTOS

On most platforms there are instructions can speedup the ready task searching. Enabling this option the FreeRTOS with this instructions support will be built.

CONFIG_FREERTOS_HZ

Tick rate (Hz)

Found in: Component config > FreeRTOS

Select the tick rate at which FreeRTOS does pre-emptive context switching.

CONFIG_FREERTOS_ASSERT_ON_UNTESTED_FUNCTION

Halt when an SMP-untested function is called

Found in: Component config > FreeRTOS

Some functions in FreeRTOS have not been thoroughly tested yet when moving to the SMP implementation of FreeRTOS. When this option is enabled, these fuctions will throw an assert().

CONFIG_FREERTOS_CHECK_STACKOVERFLOW

Check for stack overflow

Found in: Component config > FreeRTOS

FreeRTOS can check for stack overflows in threads and trigger an user function called vApplicationStackOverflowHook when this happens.

Available options:

- No checking (FREERTOS_CHECK_STACKOVERFLOW_NONE)
  Do not check for stack overflows (configCHECK_FOR_STACK_OVERFLOW=0)
- Check by stack pointer value (FREERTOS_CHECK_STACKOVERFLOW_PTRVAL)
  Check for stack overflows on each context switch by checking if the stack pointer is in a valid range. Quick but does not detect stack overflows that happened between context switches (configCHECK_FOR_STACK_OVERFLOW=1)
- Check using canary bytes (FREERTOS_CHECK_STACKOVERFLOW_CANARY)
  Places some magic bytes at the end of the stack area and on each context switch, check if these bytes are still intact. More thorough than just checking the pointer, but also slightly slower. (configCHECK_FOR_STACK_OVERFLOW=2)

CONFIG_FREERTOS_WATCHPOINT_END_OF_STACK

Set a debug watchpoint as a stack overflow check

Found in: Component config > FreeRTOS

FreeRTOS can check if a stack has overflown its bounds by checking either the value of the stack pointer or by checking the integrity of canary bytes. (See FREERTOS_CHECK_STACKOVERFLOW for more information.) These checks only happen on a context switch, and the situation that caused the stack overflow may already be long gone by then. This option will use the debug memory watchpoint 1 (the second one) to allow breaking into the debugger (or panic'ing) as soon as any of the last 32 bytes on the stack of a task are overwritten. The side effect is that using gdb, you effectively only have one watchpoint; the 2nd one is overwritten as soon as a task switch happens.

This check only triggers if the stack overflow writes within 4 bytes of the end of the stack, rather than overshooting further, so it is worth combining this approach with one of the other stack overflow check methods.

When this watchpoint is hit, gdb will stop with a SIGTRAP message. When no JTAG OCD is attached, esp-idf will panic on an unhandled debug exception.

CONFIG_FREERTOS_INTERRUPT_BACKTRACE

Enable backtrace from interrupt to task context

Found in: Component config > FreeRTOS

If this option is enabled, interrupt stack frame will be modified to point to the code of the interrupted task as its return address. This helps the debugger (or the panic handler) show a backtrace from the interrupt to the task which was interrupted. This also works for nested interrupts: higer level interrupt stack can be traced back to the lower level interrupt. This option adds 4 instructions to the interrupt dispatching code.

CONFIG_FREERTOS_THREAD_LOCAL_STORAGE_POINTERS

Number of thread local storage pointers

Found in: Component config > FreeRTOS

FreeRTOS has the ability to store per-thread pointers in the task control block. This controls the number of pointers available.

This value must be at least 1. Index 0 is reserved for use by the pthreads API thread-local-storage. Other indexes can be used for any desired purpose.

CONFIG_FREERTOS_ASSERT

FreeRTOS assertions

Found in: Component config > FreeRTOS

Failed FreeRTOS configASSERT() assertions can be configured to behave in different ways.

Available options:

- abort() on failed assertions (FREERTOS_ASSERT_FAIL_ABORT)
  If a FreeRTOS configASSERT() fails, FreeRTOS will abort() and halt execution. The panic handler can be configured to handle the outcome of an abort() in different ways.
- Print and continue failed assertions (FREERTOS_ASSERT_FAIL_PRINT_CONTINUE)
  If a FreeRTOS assertion fails, print it out and continue.
- Disable FreeRTOS assertions (FREERTOS_ASSERT_DISABLE)
  FreeRTOS configASSERT() will not be compiled into the binary.

CONFIG_FREERTOS_IDLE_TASK_STACKSIZE

Idle Task stack size

Found in: Component config > FreeRTOS

The idle task has its own stack, sized in bytes. The default size is enough for most uses. Size can be reduced to 768 bytes if no (or simple) FreeRTOS idle hooks are used and pthread local storage or FreeRTOS local storage cleanup callbacks are not used.

The stack size may need to be increased above the default if the app installs idle or thread local storage cleanup hooks that use a lot of stack memory.

CONFIG_FREERTOS_ISR_STACKSIZE

ISR stack size

Found in: Component config > FreeRTOS

The interrupt handlers have their own stack. The size of the stack can be defined here. Each processor has its own stack, so the total size occupied will be twice this.

CONFIG_FREERTOS_LEGACY_HOOKS

Use FreeRTOS legacy hooks

Found in: Component config > FreeRTOS

FreeRTOS offers a number of hooks/callback functions that are called when a timer tick happens, the idle thread runs etc. esp-idf replaces these by runtime registerable hooks using the esp_register_freertos_xxx_hook system, but for legacy reasons the old hooks can also still be enabled. Please enable this only if you have code that for some reason can't be migrated to the esp_register_freertos_xxx_hook system.

CONFIG_FREERTOS_MAX_TASK_NAME_LEN

Maximum task name length

Found in: Component config > FreeRTOS

Changes the maximum task name length. Each task allocated will include this many bytes for a task name. Using a shorter value saves a small amount of RAM, a longer value allows more complex names.

For most uses, the default of 16 is OK.

CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION

Enable FreeRTOS static allocation API

Found in: Component config > FreeRTOS

FreeRTOS gives the application writer the ability to instead provide the memory themselves, allowing the following objects to optionally be created without any memory being allocated dynamically:

- Tasks
- Software Timers (Daemon task is still dynamic. See documentation)
- Queues
- Event Groups
- Binary Semaphores
- Counting Semaphores
- Recursive Semaphores
- Mutexes

Whether it is preferable to use static or dynamic memory allocation is dependent on the application, and the preference of the application writer. Both methods have pros and cons, and both methods can be used within the same RTOS application.

Creating RTOS objects using statically allocated RAM has the benefit of providing the application writer with more control: RTOS objects can be placed at specific memory locations. The maximum RAM footprint can be determined at link time, rather than run time. The application writer does not need to concern themselves with graceful handling of memory allocation failures. It allows the RTOS to be used in applications that simply don't allow any dynamic memory allocation (although FreeRTOS includes allocation schemes that can overcome most objections).

CONFIG_FREERTOS_ENABLE_STATIC_TASK_CLEAN_UP

Enable static task clean up hook

Found in: Component config > FreeRTOS > CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION

Enable this option to make FreeRTOS call the static task clean up hook when a task is deleted.

Bear in mind that if this option is enabled you will need to implement the following function:

```
void vPortCleanUpTCB ( void \*pxTCB ) {
    // place clean up code here
}
```

CONFIG_FREERTOS_TIMER_TASK_PRIORITY

FreeRTOS timer task priority

Found in: Component config > FreeRTOS

The timer service task (primarily) makes use of existing FreeRTOS features, allowing timer functionality to be added to an application with minimal impact on the size of the application's executable binary.

Use this constant to define the priority that the timer task will run at.

CONFIG_FREERTOS_TIMER_TASK_STACK_DEPTH

FreeRTOS timer task stack size

Found in: Component config > FreeRTOS

The timer service task (primarily) makes use of existing FreeRTOS features, allowing timer functionality to be added to an application with minimal impact on the size of the application's executable binary.

Use this constant to define the size (in bytes) of the stack allocated for the timer task.

CONFIG_FREERTOS_TIMER_QUEUE_LENGTH

FreeRTOS timer queue length

Found in: Component config > FreeRTOS

FreeRTOS provides a set of timer related API functions. Many of these functions use a standard FreeRTOS queue to send commands to the timer service task. The queue used for this purpose is called the 'timer command queue'. The 'timer command queue' is private to the FreeRTOS timer implementation, and cannot be accessed directly.

For most uses the default value of 10 is OK.

CONFIG_FREERTOS_QUEUE_REGISTRY_SIZE

FreeRTOS queue registry size

Found in: Component config > FreeRTOS

FreeRTOS uses the queue registry as a means for kernel aware debuggers to locate queues, semaphores, and mutexes. The registry allows for a textual name to be associated with a queue for easy identification within a debugging GUI. A value of 0 will disable queue registry functionality, and a value larger than 0 will specify the number of queues/semaphores/mutexes that the registry can hold.

CONFIG_FREERTOS_USE_TRACE_FACILITY

Enable FreeRTOS trace facility

Found in: Component config > FreeRTOS

If enabled, configUSE_TRACE_FACILITY will be defined as 1 in FreeRTOS. This will allow the usage of trace facility functions such as uxTaskGetSystemState().

CONFIG_FREERTOS_USE_STATS_FORMATTING_FUNCTIONS

Enable FreeRTOS stats formatting functions

Found in: Component config > FreeRTOS > CONFIG_FREERTOS_USE_TRACE_FACILITY

If enabled, configUSE_STATS_FORMATTING_FUNCTIONS will be defined as 1 in FreeRTOS. This will allow the usage of stats formatting functions such as vTaskList().

CONFIG_FREERTOS_VTASKLIST_INCLUDE_COREID

Enable display of xCoreID in vTaskList

Found in: Component config > FreeRTOS > CONFIG_FREERTOS_USE_TRACE_FACILITY > CONFIG_FREERTOS_USE_STATS_FORMATTING_FUNCTIONS

If enabled, this will include an extra column when vTaskList is called to display the CoreID the task is pinned to (0,1) or -1 if not pinned.

CONFIG_FREERTOS_GENERATE_RUN_TIME_STATS

Enable FreeRTOS to collect run time stats

Found in: Component config > FreeRTOS

If enabled, configGENERATE_RUN_TIME_STATS will be defined as 1 in FreeRTOS. This will allow FreeRTOS to collect information regarding the usage of processor time amongst FreeRTOS tasks. Run time stats are generated using either the ESP Timer or the CPU Clock as the clock source (Note that run time stats are only valid until the clock source overflows). The function vTaskGetRunTimeStats() will also be available if FREERTOS_USE_STATS_FORMATTING_FUNCTIONS and FREERTOS_USE_TRACE_FACILITY are enabled. vTaskGetRunTimeStats() will display the run time of each task as a % of the total run time of all CPUs (task run time / no of CPUs) / (total run time / 100)

CONFIG_FREERTOS_RUN_TIME_STATS_CLK

Choose the clock source for run time stats

Found in: Component config > FreeRTOS > CONFIG_FREERTOS_GENERATE_RUN_TIME_STATS

Choose the clock source for FreeRTOS run time stats. Options are CPU0's CPU Clock or the ESP Timer. Both clock sources are 32 bits. The CPU Clock can run at a higher frequency hence provide a finer resolution but will overflow much quicker. Note that run time stats are only valid until the clock source overflows.

Available options:

- Use ESP TIMER for run time stats (FREERTOS_RUN_TIME_STATS_USING_ESP_TIMER)
  ESP Timer will be used as the clock source for FreeRTOS run time stats. The ESP Timer runs at a frequency of 1MHz regardless of Dynamic Frequency Scaling. Therefore the ESP Timer will overflow in approximately 4290 seconds.
- Use CPU Clock for run time stats (FREERTOS_RUN_TIME_STATS_USING_CPU_CLK)
  CPU Clock will be used as the clock source for the generation of run time stats. The CPU Clock has a frequency dependent on ESP32_DEFAULT_CPU_FREQ_MHZ and Dynamic Frequency Scaling (DFS). Therefore the CPU Clock frequency can fluctuate between 80 to 240MHz. Run time stats generated using the CPU Clock represents the number of CPU cycles each task is allocated and DOES NOT reflect the amount of time each task runs for (as CPU clock frequency can change). If the CPU clock consistently runs at the maximum frequency of 240MHz, it will overflow in approximately 17 seconds.

CONFIG_FREERTOS_USE_TICKLESS_IDLE

Tickless idle support

Found in: Component config > FreeRTOS

If power management support is enabled, FreeRTOS will be able to put the system into light sleep mode when no tasks need to run for a number of ticks. This number can be set using FREERTOS_IDLE_TIME_BEFORE_SLEEP option. This feature is also known as "automatic light sleep" .

Note that timers created using esp_timer APIs may prevent the system from entering sleep mode, even when no tasks need to run.

If disabled, automatic light sleep support will be disabled.

CONFIG_FREERTOS_IDLE_TIME_BEFORE_SLEEP

Minimum number of ticks to enter sleep mode for

Found in: Component config > FreeRTOS > CONFIG_FREERTOS_USE_TICKLESS_IDLE

FreeRTOS will enter light sleep mode if no tasks need to run for this number of ticks.

CONFIG_FREERTOS_TASK_FUNCTION_WRAPPER

Enclose all task functions in a wrapper function

Found in: Component config > FreeRTOS

If enabled, all FreeRTOS task functions will be enclosed in a wrapper function. If a task function mistakenly returns (i.e. does not delete), the call flow will return to the wrapper function. The wrapper function will then log an error and abort the application. This option is also required for GDB backtraces and C++ exceptions to work correctly inside top-level task functions.

CONFIG_FREERTOS_CHECK_MUTEX_GIVEN_BY_OWNER

Check that mutex semaphore is given by owner task

Found in: Component config > FreeRTOS

If enabled, assert that when a mutex semaphore is given, the task giving the semaphore is the task which is currently holding the mutex.

CONFIG_FREERTOS_CHECK_PORT_CRITICAL_COMPLIANCE

Tests compliance with Vanilla FreeRTOS port*_CRITICAL calls

Found in: Component config > FreeRTOS

If enabled, context of port*_CRITICAL calls (ISR or Non-ISR) would be checked to be in compliance with Vanilla FreeRTOS. e.g Calling port*_CRITICAL from ISR context would cause assert failure

FAT Filesystem support Contains:

- *CONFIG_FATFS_CHOOSE_CODEPAGE*
- *CONFIG_FATFS_LONG_FILENAMES*
- *CONFIG_FATFS_MAX_LFN*
- *CONFIG_FATFS_API_ENCODING*
- *CONFIG_FATFS_FS_LOCK*
- *CONFIG_FATFS_TIMEOUT_MS*
- *CONFIG_FATFS_PER_FILE_CACHE*
- *CONFIG_FATFS_ALLOC_PREFER_EXTRAM*

CONFIG_FATFS_CHOOSE_CODEPAGE

OEM Code Page

Found in: Component config > FAT Filesystem support

OEM code page used for file name encodings.

If "Dynamic" is selected, code page can be chosen at runtime using f_setcp function. Note that choosing this option will increase application size by ~480kB.

Available options:

- Dynamic (all code pages supported) (FATFS_CODEPAGE_DYNAMIC)
- US (CP437) (FATFS_CODEPAGE_437)
- Arabic (CP720) (FATFS_CODEPAGE_720)
- Greek (CP737) (FATFS_CODEPAGE_737)
- KBL (CP771) (FATFS_CODEPAGE_771)
- Baltic (CP775) (FATFS_CODEPAGE_775)
- Latin 1 (CP850) (FATFS_CODEPAGE_850)
- Latin 2 (CP852) (FATFS_CODEPAGE_852)
- Cyrillic (CP855) (FATFS_CODEPAGE_855)
- Turkish (CP857) (FATFS_CODEPAGE_857)
- Portugese (CP860) (FATFS_CODEPAGE_860)
- Icelandic (CP861) (FATFS_CODEPAGE_861)
- Hebrew (CP862) (FATFS_CODEPAGE_862)
- Canadian French (CP863) (FATFS_CODEPAGE_863)
- Arabic (CP864) (FATFS_CODEPAGE_864)
- Nordic (CP865) (FATFS_CODEPAGE_865)
- Russian (CP866) (FATFS_CODEPAGE_866)
- Greek 2 (CP869) (FATFS_CODEPAGE_869)
- Japanese (DBCS) (CP932) (FATFS_CODEPAGE_932)
- Simplified Chinese (DBCS) (CP936) (FATFS_CODEPAGE_936)
- Korean (DBCS) (CP949) (FATFS_CODEPAGE_949)
- Traditional Chinese (DBCS) (CP950) (FATFS_CODEPAGE_950)

CONFIG_FATFS_LONG_FILENAMES

Long filename support

Found in: Component config > FAT Filesystem support

Support long filenames in FAT. Long filename data increases memory usage. FATFS can be configured to store the buffer for long filename data in stack or heap.

Available options:

- No long filenames (FATFS_LFN_NONE)
- Long filename buffer in heap (FATFS_LFN_HEAP)
- Long filename buffer on stack (FATFS_LFN_STACK)

CONFIG_FATFS_MAX_LFN

Max long filename length

Found in: Component config > FAT Filesystem support

Maximum long filename length. Can be reduced to save RAM.

CONFIG_FATFS_API_ENCODING

API character encoding

Found in: Component config > FAT Filesystem support

Choose encoding for character and string arguments/returns when using FATFS APIs. The encoding of arguments will usually depend on text editor settings.

Available options:

- API uses ANSI/OEM encoding (FATFS_API_ENCODING_ANSI_OEM)
- API uses UTF-16 encoding (FATFS_API_ENCODING_UTF_16)
- API uses UTF-8 encoding (FATFS_API_ENCODING_UTF_8)

CONFIG_FATFS_FS_LOCK

Number of simultaneously open files protected by lock function

Found in: Component config > FAT Filesystem support

This option sets the FATFS configuration value _FS_LOCK. The option _FS_LOCK switches file lock function to control duplicated file open and illegal operation to open objects.

* 0: Disable file lock function. To avoid volume corruption, application should avoid illegal open, remove and rename to the open objects.

* >0: Enable file lock function. The value defines how many files/sub-directories can be opened simultaneously under file lock control.

Note that the file lock control is independent of re-entrancy.

CONFIG_FATFS_TIMEOUT_MS

Timeout for acquiring a file lock, ms

Found in: Component config > FAT Filesystem support

This option sets FATFS configuration value _FS_TIMEOUT, scaled to milliseconds. Sets the number of milliseconds FATFS will wait to acquire a mutex when operating on an open file. For example, if one task is performing a lenghty operation, another task will wait for the first task to release the lock, and time out after amount of time set by this option.

CONFIG_FATFS_PER_FILE_CACHE

Use separate cache for each file

Found in: Component config > FAT Filesystem support

This option affects FATFS configuration value _FS_TINY.

If this option is set, _FS_TINY is 0, and each open file has its own cache, size of the cache is equal to the _MAX_SS variable (512 or 4096 bytes). This option uses more RAM if more than 1 file is open, but needs less reads and writes to the storage for some operations.

If this option is not set, _FS_TINY is 1, and single cache is used for all open files, size is also equal to _MAX_SS variable. This reduces the amount of heap used when multiple files are open, but increases the number of read and write operations which FATFS needs to make.

CONFIG_FATFS_ALLOC_PREFER_EXTRAM

Perfer external RAM when allocating FATFS buffers

Found in: Component config > FAT Filesystem support

When the option is enabled, internal buffers used by FATFS will be allocated from external RAM. If the allocation from external RAM fails, the buffer will be allocated from the internal RAM. Disable this option if optimizing for performance. Enable this option if optimizing for internal memory size.

Core dump Contains:

- *CONFIG_ESP32_COREDUMP_TO_FLASH_OR_UART*
- *CONFIG_ESP32_COREDUMP_DATA_FORMAT*
- *CONFIG_ESP32_COREDUMP_CHECKSUM*
- *CONFIG_ESP32_CORE_DUMP_MAX_TASKS_NUM*
- *CONFIG_ESP32_CORE_DUMP_UART_DELAY*
- *CONFIG_ESP32_CORE_DUMP_STACK_SIZE*
- *CONFIG_ESP32_CORE_DUMP_DECODE*

CONFIG_ESP32_COREDUMP_TO_FLASH_OR_UART

Data destination

Found in: Component config > Core dump

Select place to store core dump: flash, uart or none (to disable core dumps generation).

Core dumps to Flash are not available if PSRAM is used for task stacks.

If core dump is configured to be stored in flash and custom partition table is used add corresponding entry to your CSV. For examples, please see predefined partition table CSV descriptions in the components/partition_table directory.

Available options:

- Flash (ESP32_ENABLE_COREDUMP_TO_FLASH)
- UART (ESP32_ENABLE_COREDUMP_TO_UART)
- None (ESP32_ENABLE_COREDUMP_TO_NONE)

CONFIG_ESP32_COREDUMP_DATA_FORMAT

Core dump data format

Found in: Component config > Core dump

Select the data format for core dump.

Available options:

- Binary format (ESP32_COREDUMP_DATA_FORMAT_BIN)

- ELF format (ESP32_COREDUMP_DATA_FORMAT_ELF)

CONFIG_ESP32_COREDUMP_CHECKSUM

Core dump data integrity check

Found in: Component config > Core dump

Select the integrity check for the core dump.

Available options:

- Use CRC32 for integrity verification (ESP32_COREDUMP_CHECKSUM_CRC32)
- Use SHA256 for integrity verification (ESP32_COREDUMP_CHECKSUM_SHA256)

CONFIG_ESP32_CORE_DUMP_MAX_TASKS_NUM

Maximum number of tasks

Found in: Component config > Core dump

Maximum number of tasks snapshots in core dump.

CONFIG_ESP32_CORE_DUMP_UART_DELAY

Delay before print to UART

Found in: Component config > Core dump

Config delay (in ms) before printing core dump to UART. Delay can be interrupted by pressing Enter key.

CONFIG_ESP32_CORE_DUMP_STACK_SIZE

Reserved stack size

Found in: Component config > Core dump

Size of the memory to be reserved for core dump stack. If 0 core dump process will run on the stack of crashed task/ISR, otherwise special stack will be allocated. To ensure that core dump itself will not overflow task/ISR stack set this to the value above 800. NOTE: It eats DRAM.

CONFIG_ESP32_CORE_DUMP_DECODE

Handling of UART core dumps in IDF Monitor

Found in: Component config > Core dump

Available options:

- Decode and show summary (info_corefile) (ESP32_CORE_DUMP_DECODE_INFO)
- Don’t decode (ESP32_CORE_DUMP_DECODE_DISABLE)

Wi-Fi Contains:

- *CONFIG_ESP32_WIFI_STATIC_RX_BUFFER_NUM*
- *CONFIG_ESP32_WIFI_DYNAMIC_RX_BUFFER_NUM*
- *CONFIG_ESP32_WIFI_TX_BUFFER*
- *CONFIG_ESP32_WIFI_STATIC_TX_BUFFER_NUM*
- *CONFIG_ESP32_WIFI_CACHE_TX_BUFFER_NUM*
- *CONFIG_ESP32_WIFI_DYNAMIC_TX_BUFFER_NUM*
- *CONFIG_ESP32_WIFI_CSI_ENABLED*
- *CONFIG_ESP32_WIFI_AMPDU_TX_ENABLED*
- *CONFIG_ESP32_WIFI_AMPDU_RX_ENABLED*

- *CONFIG_ESP32_WIFI_NVS_ENABLED*
- *CONFIG_ESP32_WIFI_TASK_CORE_ID*
- *CONFIG_ESP32_WIFI_SOFTAP_BEACON_MAX_LEN*
- *CONFIG_ESP32_WIFI_MGMT_SBUF_NUM*
- *CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE*
- *CONFIG_ESP32_WIFI_IRAM_OPT*
- *CONFIG_ESP32_WIFI_RX_IRAM_OPT*
- *CONFIG_ESP32_WIFI_ENABLE_WPA3_SAE*

CONFIG_ESP32_WIFI_STATIC_RX_BUFFER_NUM

Max number of WiFi static RX buffers

Found in: Component config > Wi-Fi

Set the number of WiFi static RX buffers. Each buffer takes approximately 1.6KB of RAM. The static rx buffers are allocated when esp_wifi_init is called, they are not freed until esp_wifi_deinit is called.

WiFi hardware use these buffers to receive all 802.11 frames. A higher number may allow higher throughput but increases memory use. If ESP32_WIFI_AMPDU_RX_ENABLED is enabled, this value is recommended to set equal or bigger than ESP32_WIFI_RX_BA_WIN in order to achieve better throughput and compatibility with both stations and APs.

CONFIG_ESP32_WIFI_DYNAMIC_RX_BUFFER_NUM

Max number of WiFi dynamic RX buffers

Found in: Component config > Wi-Fi

Set the number of WiFi dynamic RX buffers, 0 means unlimited RX buffers will be allocated (provided sufficient free RAM). The size of each dynamic RX buffer depends on the size of the received data frame.

For each received data frame, the WiFi driver makes a copy to an RX buffer and then delivers it to the high layer TCP/IP stack. The dynamic RX buffer is freed after the higher layer has successfully received the data frame.

For some applications, WiFi data frames may be received faster than the application can process them. In these cases we may run out of memory if RX buffer number is unlimited (0).

If a dynamic RX buffer limit is set, it should be at least the number of static RX buffers.

CONFIG_ESP32_WIFI_TX_BUFFER

Type of WiFi TX buffers

Found in: Component config > Wi-Fi

Select type of WiFi TX buffers:

If "Static" is selected, WiFi TX buffers are allocated when WiFi is initialized and released when WiFi is de-initialized. The size of each static TX buffer is fixed to about 1.6KB.

If "Dynamic" is selected, each WiFi TX buffer is allocated as needed when a data frame is delivered to the Wifi driver from the TCP/IP stack. The buffer is freed after the data frame has been sent by the WiFi driver. The size of each dynamic TX buffer depends on the length of each data frame sent by the TCP/IP layer.

If PSRAM is enabled, "Static" should be selected to guarantee enough WiFi TX buffers. If PSRAM is disabled, "Dynamic" should be selected to improve the utilization of RAM.

Available options:

- Static (ESP32_WIFI_STATIC_TX_BUFFER)
- Dynamic (ESP32_WIFI_DYNAMIC_TX_BUFFER)

CONFIG_ESP32_WIFI_STATIC_TX_BUFFER_NUM

Max number of WiFi static TX buffers

Found in: Component config > Wi-Fi

Set the number of WiFi static TX buffers. Each buffer takes approximately 1.6KB of RAM. The static RX buffers are allocated when esp_wifi_init() is called, they are not released until esp_wifi_deinit() is called.

For each transmitted data frame from the higher layer TCP/IP stack, the WiFi driver makes a copy of it in a TX buffer. For some applications especially UDP applications, the upper layer can deliver frames faster than WiFi layer can transmit. In these cases, we may run out of TX buffers.

CONFIG_ESP32_WIFI_CACHE_TX_BUFFER_NUM

Max number of WiFi cache TX buffers

Found in: Component config > Wi-Fi

Set the number of WiFi cache TX buffer number.

For each TX packet from uplayer, such as LWIP etc, WiFi driver needs to allocate a static TX buffer and makes a copy of uplayer packet. If WiFi driver fails to allocate the static TX buffer, it caches the uplayer packets to a dedicated buffer queue, this option is used to configure the size of the cached TX queue.

CONFIG_ESP32_WIFI_DYNAMIC_TX_BUFFER_NUM

Max number of WiFi dynamic TX buffers

Found in: Component config > Wi-Fi

Set the number of WiFi dynamic TX buffers. The size of each dynamic TX buffer is not fixed, it depends on the size of each transmitted data frame.

For each transmitted frame from the higher layer TCP/IP stack, the WiFi driver makes a copy of it in a TX buffer. For some applications, especially UDP applications, the upper layer can deliver frames faster than WiFi layer can transmit. In these cases, we may run out of TX buffers.

CONFIG_ESP32_WIFI_CSI_ENABLED

WiFi CSI(Channel State Information)

Found in: Component config > Wi-Fi

Select this option to enable CSI(Channel State Information) feature. CSI takes about CONFIG_ESP32_WIFI_STATIC_RX_BUFFER_NUM KB of RAM. If CSI is not used, it is better to disable this feature in order to save memory.

CONFIG_ESP32_WIFI_AMPDU_TX_ENABLED

WiFi AMPDU TX

Found in: Component config > Wi-Fi

Select this option to enable AMPDU TX feature

CONFIG_ESP32_WIFI_TX_BA_WIN

WiFi AMPDU TX BA window size

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_AMPDU_TX_ENABLED

Set the size of WiFi Block Ack TX window. Generally a bigger value means higher throughput but more memory. Most of time we should NOT change the default value unless special reason, e.g. test the maximum UDP TX throughput with iperf etc. For iperf test in shieldbox, the recommended value is 9~12.

CONFIG_ESP32_WIFI_AMPDU_RX_ENABLED

WiFi AMPDU RX

Found in: Component config > Wi-Fi

Select this option to enable AMPDU RX feature

CONFIG_ESP32_WIFI_RX_BA_WIN

WiFi AMPDU RX BA window size

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_AMPDU_RX_ENABLED

Set the size of WiFi Block Ack RX window. Generally a bigger value means higher throughput and better compatibility but more memory. Most of time we should NOT change the default value unless special reason, e.g. test the maximum UDP RX throughput with iperf etc. For iperf test in shieldbox, the recommended value is 9~12. If PSRAM is used and WiFi memory is prefered to allocat in PSRAM first, the default and minimum value should be 16 to achieve better throughput and compatibility with both stations and APs.

CONFIG_ESP32_WIFI_NVS_ENABLED

WiFi NVS flash

Found in: Component config > Wi-Fi

Select this option to enable WiFi NVS flash

CONFIG_ESP32_WIFI_TASK_CORE_ID

WiFi Task Core ID

Found in: Component config > Wi-Fi

Pinned WiFi task to core 0 or core 1.

Available options:

- Core 0 (ESP32_WIFI_TASK_PINNED_TO_CORE_0)
- Core 1 (ESP32_WIFI_TASK_PINNED_TO_CORE_1)

CONFIG_ESP32_WIFI_SOFTAP_BEACON_MAX_LEN

Max length of WiFi SoftAP Beacon

Found in: Component config > Wi-Fi

ESP-MESH utilizes beacon frames to detect and resolve root node conflicts (see documentation). However the default length of a beacon frame can simultaneously hold only five root node identifier structures, meaning that a root node conflict of up to five nodes can be detected at one time. In the occurence of more root nodes conflict involving more than five root nodes, the conflict resolution process will detect five of the root nodes, resolve the conflict, and re-detect more root nodes. This process will repeat until all root node conflicts are resolved. However this process can generally take a very long time.

To counter this situation, the beacon frame length can be increased such that more root nodes can be detected simultaneously. Each additional root node will require 36 bytes and should be added ontop of the default beacon frame length of 752 bytes. For example, if you want to detect 10 root nodes simultaneously, you need to set the beacon frame length as 932 (752+36*5).

Setting a longer beacon length also assists with debugging as the conflicting root nodes can be identified more quickly.

CONFIG_ESP32_WIFI_MGMT_SBUF_NUM

WiFi mgmt short buffer number

Found in: Component config > Wi-Fi

Set the number of WiFi management short buffer.

CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE

Enable WiFi debug log

Found in: Component config > Wi-Fi

Select this option to enable WiFi debug log

CONFIG_ESP32_WIFI_DEBUG_LOG_LEVEL

WiFi debug log level

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE

The WiFi log is divided into the following levels: ERROR,WARNING,INFO,DEBUG,VERBOSE. The ERROR,WARNING,INFO levels are enabled by default, and the DEBUG,VERBOSE levels can be enabled here.

Available options:

- WiFi Debug Log Debug (ESP32_WIFI_DEBUG_LOG_DEBUG)
- WiFi Debug Log Verbose (ESP32_WIFI_DEBUG_LOG_VERBOSE)

CONFIG_ESP32_WIFI_DEBUG_LOG_MODULE

WiFi debug log module

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE

The WiFi log module contains three parts: WIFI,COEX,MESH. The WIFI module indicates the logs related to WiFi, the COEX module indicates the logs related to WiFi and BT(or BLE) coexist, the MESH module indicates the logs related to Mesh. When ESP32_WIFI_LOG_MODULE_ALL is enabled, all modules are selected.

Available options:

- WiFi Debug Log Module All (ESP32_WIFI_DEBUG_LOG_MODULE_ALL)
- WiFi Debug Log Module WiFi (ESP32_WIFI_DEBUG_LOG_MODULE_WIFI)
- WiFi Debug Log Module Coex (ESP32_WIFI_DEBUG_LOG_MODULE_COEX)
- WiFi Debug Log Module Mesh (ESP32_WIFI_DEBUG_LOG_MODULE_MESH)

CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE

WiFi debug log submodule

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE

Enable this option to set the WiFi debug log submodule. Currently the log submodule contains the following parts: INIT,IOCTL,CONN,SCAN. The INIT submodule indicates the initialization process.The IOCTL submodule indicates the API calling process. The CONN submodule indicates the connecting process.The SCAN submodule indicates the scaning process.

CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE_ALL

WiFi Debug Log Submodule All

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE > CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE

When this option is enabled, all debug submodules are selected.

CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE_INIT

WiFi Debug Log Submodule Init

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE > CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE

CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE_IOCTL

WiFi Debug Log Submodule Ioctl

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE > CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE

CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE_CONN

WiFi Debug Log Submodule Conn

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE > CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE

CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE_SCAN

WiFi Debug Log Submodule Scan

Found in: Component config > Wi-Fi > CONFIG_ESP32_WIFI_DEBUG_LOG_ENABLE > CONFIG_ESP32_WIFI_DEBUG_LOG_SUBMODULE

CONFIG_ESP32_WIFI_IRAM_OPT

WiFi IRAM speed optimization

Found in: Component config > Wi-Fi

Select this option to place frequently called Wi-Fi library functions in IRAM. When this option is disabled, more than 10Kbytes of IRAM memory will be saved but Wi-Fi throughput will be reduced.

CONFIG_ESP32_WIFI_RX_IRAM_OPT

WiFi RX IRAM speed optimization

Found in: Component config > Wi-Fi

Select this option to place frequently called Wi-Fi library RX functions in IRAM. When this option is disabled, more than 17Kbytes of IRAM memory will be saved but Wi-Fi performance will be reduced.

CONFIG_ESP32_WIFI_ENABLE_WPA3_SAE

Enable WPA3-Personal

Found in: Component config > Wi-Fi

Select this option to allow the device to establish a WPA3-Personal connection with eligible AP's. PMF (Protected Management Frames) is a prerequisite feature for a WPA3 connection, it needs to be explicitly configured before attempting connection. Please refer to the Wi-Fi Driver API Guide for details.

PHY Contains:

- *CONFIG_ESP32_PHY_CALIBRATION_AND_DATA_STORAGE*
- *CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION*
- *CONFIG_ESP32_PHY_MAX_WIFI_TX_POWER*

CONFIG_ESP32_PHY_CALIBRATION_AND_DATA_STORAGE

Store phy calibration data in NVS

Found in: Component config > PHY

If this option is enabled, NVS will be initialized and calibration data will be loaded from there. PHY calibration will be skipped on deep sleep wakeup. If calibration data is not found, full calibration will be performed and stored in NVS. Normally, only partial calibration will be performed. If this option is disabled, full calibration will be performed.

If it's easy that your board calibrate bad data, choose 'n'. Two cases for example, you should choose 'n': 1.If your board is easy to be booted up with antenna disconnected. 2.Because of your board design, each time when you do calibration, the result are too unstable. If unsure, choose 'y'.

CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION

Use a partition to store PHY init data

Found in: Component config > PHY

If enabled, PHY init data will be loaded from a partition. When using a custom partition table, make sure that PHY data partition is included (type: 'data', subtype: 'phy'). With default partition tables, this is done automatically. If PHY init data is stored in a partition, it has to be flashed there, otherwise runtime error will occur.

If this option is not enabled, PHY init data will be embedded into the application binary.

If unsure, choose 'n'.

Contains:

- *CONFIG_ESP32_PHY_DEFAULT_INIT_IF_INVALID*
- *CONFIG_ESP32_SUPPORT_MULTIPLE_PHY_INIT_DATA_BIN*

CONFIG_ESP32_PHY_DEFAULT_INIT_IF_INVALID

Reset default PHY init data if invalid

Found in: Component config > PHY > CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION

If enabled, PHY init data will be restored to default if it cannot be verified successfully to avoid endless bootloops.

If unsure, choose 'n'.

CONFIG_ESP32_SUPPORT_MULTIPLE_PHY_INIT_DATA_BIN

Support multiple PHY init data bin

Found in: Component config > PHY > CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION

If enabled, the corresponding PHY init data type can be automatically switched according to the country code. China' s PHY init data bin is used by default. Can be modified by country information in API esp_wifi_set_country(). The priority of switching the PHY init data type is: 1. Country configured by API esp_wifi_set_country() and the parameter policy is WIFI_COUNTRY_POLICY_MANUAL. 2. Country notified by the connected AP. 3. Country configured by API esp_wifi_set_country() and the parameter policy is WIFI_COUNTRY_POLICY_AUTO.

CONFIG_ESP32_PHY_INIT_DATA_ERROR

Terminate operation when PHY init data error

Found in: Component config > PHY > CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION > CONFIG_ESP32_SUPPORT_MULTIPLE_PHY_INIT_DATA_BIN

If enabled, when an error occurs while the PHY init data is updated, the program will terminate and restart. If not enabled, the PHY init data will not be updated when an error occurs.

CONFIG_ESP32_PHY_MAX_WIFI_TX_POWER

Max WiFi TX power (dBm)

Found in: Component config > PHY

Set maximum transmit power for WiFi radio. Actual transmit power for high data rates may be lower than this setting.

High resolution timer (esp_timer)

Contains:

- *CONFIG_ESP_TIMER_PROFILING*
- *CONFIG_ESP_TIMER_TASK_STACK_SIZE*
- *CONFIG_ESP_TIMER_IMPL*

CONFIG_ESP_TIMER_PROFILING

Enable esp_timer profiling features

Found in: Component config > High resolution timer (esp_timer)

If enabled, esp_timer_dump will dump information such as number of times the timer was started, number of times the timer has triggered, and the total time it took for the callback to run. This option has some effect on timer performance and the amount of memory used for timer storage, and should only be used for debugging/testing purposes.

CONFIG_ESP_TIMER_TASK_STACK_SIZE

High-resolution timer task stack size

Found in: Component config > High resolution timer (esp_timer)

Configure the stack size of "timer_task" task. This task is used to dispatch callbacks of timers created using ets_timer and esp_timer APIs. If you are seing stack overflow errors in timer task, increase this value.

Note that this is not the same as FreeRTOS timer task. To configure FreeRTOS timer task size, see "FreeRTOS timer task stack size" option in "FreeRTOS" menu.

CONFIG_ESP_TIMER_IMPL

Hardware timer to use for esp_timer

Found in: Component config > High resolution timer (esp_timer)

esp_timer APIs can be implemented using different hardware timers.

- "FRC2 (legacy)" implementation has been used in ESP-IDF v2.x - v4.1.
- "LAC timer of Timer Group 0" implementation is simpler, and has smaller run time overhead because software handling of timer overflow is not needed.
- "SYSTIMER" implementation is similar to "LAC timer of Timer Group 0" but for ESP32-S2 chip.

Available options:

- FRC2 (legacy) timer (ESP_TIMER_IMPL_FRC2)
- LAC timer of Timer Group 0 (ESP_TIMER_IMPL_TG0_LAC)
- SYSTIMER (ESP_TIMER_IMPL_SYSTIMER)

ESP System Settings Contains:

- *CONFIG_ESP_SYSTEM_PANIC*

CONFIG_ESP_SYSTEM_PANIC

Panic handler behaviour

Found in: Component config > ESP System Settings

If FreeRTOS detects unexpected behaviour or an unhandled exception, the panic handler is invoked. Configure the panic handler's action here.

Available options:

- Print registers and halt (ESP_SYSTEM_PANIC_PRINT_HALT)
  Outputs the relevant registers over the serial port and halt the processor. Needs a manual reset to restart.
- Print registers and reboot (ESP_SYSTEM_PANIC_PRINT_REBOOT)
  Outputs the relevant registers over the serial port and immediately reset the processor.
- Silent reboot (ESP_SYSTEM_PANIC_SILENT_REBOOT)
  Just resets the processor without outputting anything
- Invoke GDBStub (ESP_SYSTEM_PANIC_GDBSTUB)
  Invoke gdbstub on the serial port, allowing for gdb to attach to it to do a postmortem of the crash.

ESP NETIF Adapter Contains:

- *CONFIG_ESP_NETIF_IP_LOST_TIMER_INTERVAL*
- *CONFIG_ESP_NETIF_USE_TCPIP_STACK_LIB*
- *CONFIG_ESP_NETIF_TCPIP_ADAPTER_COMPATIBLE_LAYER*

CONFIG_ESP_NETIF_IP_LOST_TIMER_INTERVAL

IP Address lost timer interval (seconds)

Found in: Component config > ESP NETIF Adapter

The value of 0 indicates the IP lost timer is disabled, otherwise the timer is enabled.

The IP address may be lost because of some reasons, e.g. when the station disconnects from soft-AP, or when DHCP IP renew fails etc. If the IP lost timer is enabled, it will be started everytime the IP is lost. Event SYSTEM_EVENT_STA_LOST_IP will be raised if the timer expires. The IP lost timer is stopped if the station get the IP again before the timer expires.

CONFIG_ESP_NETIF_USE_TCPIP_STACK_LIB

TCP/IP Stack Library

Found in: Component config > ESP NETIF Adapter

Choose the TCP/IP Stack to work, for example, LwIP, uIP, etc.

Available options:

- LwIP (ESP_NETIF_TCPIP_LWIP)
  lwIP is a small independent implementation of the TCP/IP protocol suite.
- Loopback (ESP_NETIF_LOOPBACK)
  Dummy implementation of esp-netif functionality which connects driver transmit to receive function. This option is for testing purpose only

CONFIG_ESP_NETIF_TCPIP_ADAPTER_COMPATIBLE_LAYER

Enable backward compatible tcpip_adapter interface

Found in: Component config > ESP NETIF Adapter

Backward compatible interface to tcpip_adapter is enabled by default to support legacy TCP/IP stack initialisation code. Disable this option to use only esp-netif interface.

ESP HTTPS server

Contains:

- *CONFIG_ESP_HTTPS_SERVER_ENABLE*

CONFIG_ESP_HTTPS_SERVER_ENABLE

Enable ESP_HTTPS_SERVER component

Found in: Component config > ESP HTTPS server

Enable ESP HTTPS server component

ESP HTTPS OTA

Contains:

- *CONFIG_OTA_ALLOW_HTTP*

CONFIG_OTA_ALLOW_HTTP

Allow HTTP for OTA (WARNING: ONLY FOR TESTING PURPOSE, READ HELP)

Found in: Component config > ESP HTTPS OTA

It is highly recommended to keep HTTPS (along with server certificate validation) enabled. Enabling this option comes with potential risk of: - Non-encrypted communication channel with server - Accepting firmware upgrade image from server with fake identity

HTTP Server

Contains:

- *CONFIG_HTTPD_MAX_REQ_HDR_LEN*
- *CONFIG_HTTPD_MAX_URI_LEN*
- *CONFIG_HTTPD_ERR_RESP_NO_DELAY*
- *CONFIG_HTTPD_PURGE_BUF_LEN*
- *CONFIG_HTTPD_LOG_PURGE_DATA*
- *CONFIG_HTTPD_WS_SUPPORT*

CONFIG_HTTPD_MAX_REQ_HDR_LEN

Max HTTP Request Header Length

Found in: Component config > HTTP Server

This sets the maximum supported size of headers section in HTTP request packet to be processed by the server

CONFIG_HTTPD_MAX_URI_LEN

Max HTTP URI Length

Found in: Component config > HTTP Server

This sets the maximum supported size of HTTP request URI to be processed by the server

CONFIG_HTTPD_ERR_RESP_NO_DELAY

Use TCP_NODELAY socket option when sending HTTP error responses

Found in: Component config > HTTP Server

Using TCP_NODEALY socket option ensures that HTTP error response reaches the client before the underlying socket is closed. Please note that turning this off may cause multiple test failures

CONFIG_HTTPD_PURGE_BUF_LEN

Length of temporary buffer for purging data

Found in: Component config > HTTP Server

This sets the size of the temporary buffer used to receive and discard any remaining data that is received from the HTTP client in the request, but not processed as part of the server HTTP request handler.

If the remaining data is larger than the available buffer size, the buffer will be filled in multiple iterations. The buffer should be small enough to fit on the stack, but large enough to avoid excessive iterations.

CONFIG_HTTPD_LOG_PURGE_DATA

Log purged content data at Debug level

Found in: Component config > HTTP Server

Enabling this will log discarded binary HTTP request data at Debug level. For large content data this may not be desirable as it will clutter the log.

CONFIG_HTTPD_WS_SUPPORT

WebSocket server support

Found in: Component config > HTTP Server

This sets the WebSocket server support.

ESP HTTP client Contains:

- *CONFIG_ESP_HTTP_CLIENT_ENABLE_HTTPS*
- *CONFIG_ESP_HTTP_CLIENT_ENABLE_BASIC_AUTH*

CONFIG_ESP_HTTP_CLIENT_ENABLE_HTTPS

Enable https

Found in: Component config > ESP HTTP client

This option will enable https protocol by linking mbedtls library and initializing SSL transport

CONFIG_ESP_HTTP_CLIENT_ENABLE_BASIC_AUTH

Enable HTTP Basic Authentication

Found in: Component config > ESP HTTP client

This option will enable HTTP Basic Authentication. It is disabled by default as Basic auth uses unencrypted encoding, so it introduces a vulnerability when not using TLS

GDB Stub Contains:

- *CONFIG_ESP_GDBSTUB_SUPPORT_TASKS*

CONFIG_ESP_GDBSTUB_SUPPORT_TASKS

Enable listing FreeRTOS tasks through GDB Stub

Found in: Component config > GDB Stub

If enabled, GDBStub can supply the list of FreeRTOS tasks to GDB. Thread list can be queried from GDB using ‘info threads’ command. Note that if GDB task lists were corrupted, this feature may not work. If GDBStub fails, try disabling this feature.

CONFIG_ESP_GDBSTUB_MAX_TASKS

Maximum number of tasks supported by GDB Stub

Found in: Component config > GDB Stub > CONFIG_ESP_GDBSTUB_SUPPORT_TASKS

Set the number of tasks which GDB Stub will support.

Event Loop Library Contains:

- *CONFIG_ESP_EVENT_LOOP_PROFILING*
- *CONFIG_ESP_EVENT_POST_FROM_ISR*

CONFIG_ESP_EVENT_LOOP_PROFILING

Enable event loop profiling

Found in: Component config > Event Loop Library

Enables collections of statistics in the event loop library such as the number of events posted to/recieved by an event loop, number of callbacks involved, number of events dropped to to a full event loop queue, run time of event handlers, and number of times/run time of each event handler.

CONFIG_ESP_EVENT_POST_FROM_ISR

Support posting events from ISRs

Found in: Component config > Event Loop Library

Enable posting events from interrupt handlers.

CONFIG_ESP_EVENT_POST_FROM_IRAM_ISR

Support posting events from ISRs placed in IRAM

Found in: Component config > Event Loop Library > CONFIG_ESP_EVENT_POST_FROM_ISR

Enable posting events from interrupt handlers placed in IRAM. Enabling this option places API functions esp_event_post and esp_event_post_to in IRAM.

Ethernet Contains:

- *CONFIG_ETH_USE_SPI_ETHERNET*
- *CONFIG_ETH_USE_OPENETH*

CONFIG_ETH_USE_SPI_ETHERNET

Support SPI to Ethernet Module

Found in: Component config > Ethernet

ESP-IDF can also support some SPI-Ethernet modules.

Contains:

- *CONFIG_ETH_SPI_ETHERNET_DM9051*

CONFIG_ETH_SPI_ETHERNET_DM9051

Use DM9051

Found in: Component config > Ethernet > CONFIG_ETH_USE_SPI_ETHERNET

DM9051 is a fast Ethernet controller with an SPI interface. It’ s also integrated with a 10/100M PHY and MAC. Select to enable DM9051 driver.

CONFIG_ETH_USE_OPENETH

Support OpenCores Ethernet MAC (for use with QEMU)

Found in: Component config > Ethernet

OpenCores Ethernet MAC driver can be used when an ESP-IDF application is executed in QEMU. This driver is not supported when running on a real chip.

Contains:

- *CONFIG_ETH_OPENETH_DMA_RX_BUFFER_NUM*
- *CONFIG_ETH_OPENETH_DMA_TX_BUFFER_NUM*

CONFIG_ETH_OPENETH_DMA_RX_BUFFER_NUM

Number of Ethernet DMA Rx buffers

Found in: Component config > Ethernet > CONFIG_ETH_USE_OPENETH

Number of DMA receive buffers, each buffer is 1600 bytes.

CONFIG_ETH_OPENETH_DMA_TX_BUFFER_NUM

Number of Ethernet DMA Tx buffers

Found in: Component config > Ethernet > CONFIG_ETH_USE_OPENETH

Number of DMA transmit buffers, each buffer is 1600 bytes.

Common ESP-related Contains:

- *CONFIG_ESP_ERR_TO_NAME_LOOKUP*
- *CONFIG_ESP_SYSTEM_EVENT_QUEUE_SIZE*
- *CONFIG_ESP_SYSTEM_EVENT_TASK_STACK_SIZE*
- *CONFIG_ESP_MAIN_TASK_STACK_SIZE*
- *CONFIG_ESP_IPC_TASK_STACK_SIZE*
- *CONFIG_ESP_IPC_USES_CALLERS_PRIORITY*
- *CONFIG_ESP_MINIMAL_SHARED_STACK_SIZE*
- *CONFIG_ESP_CONSOLE_UART*
- *CONFIG_ESP_CONSOLE_UART_NUM*
- *CONFIG_ESP_CONSOLE_UART_TX_GPIO*
- *CONFIG_ESP_CONSOLE_UART_RX_GPIO*
- *CONFIG_ESP_CONSOLE_UART_BAUDRATE*
- *CONFIG_ESP_INT_WDT*
- *CONFIG_ESP_TASK_WDT*
- *CONFIG_ESP_PANIC_HANDLER_IRAM*

CONFIG_ESP_ERR_TO_NAME_LOOKUP

Enable lookup of error code strings

Found in: Component config > Common ESP-related

Functions esp_err_to_name() and esp_err_to_name_r() return string representations of error codes from a pre-generated lookup table. This option can be used to turn off the use of the look-up table in order to save memory but this comes at the price of sacrificing distinguishable (meaningful) output string representations.

CONFIG_ESP_SYSTEM_EVENT_QUEUE_SIZE

System event queue size

Found in: Component config > Common ESP-related

Config system event queue size in different application.

CONFIG_ESP_SYSTEM_EVENT_TASK_STACK_SIZE

Event loop task stack size

Found in: Component config > Common ESP-related

Config system event task stack size in different application.

CONFIG_ESP_MAIN_TASK_STACK_SIZE

Main task stack size

Found in: Component config > Common ESP-related

Configure the "main task" stack size. This is the stack of the task which calls app_main(). If app_main() returns then this task is deleted and its stack memory is freed.

CONFIG_ESP_IPC_TASK_STACK_SIZE

Inter-Processor Call (IPC) task stack size

Found in: Component config > Common ESP-related

Configure the IPC tasks stack size. One IPC task runs on each core (in dual core mode), and allows for cross-core function calls.

See IPC documentation for more details.

The default stack size should be enough for most common use cases. It can be shrunk if you are sure that you do not use any custom IPC functionality.

CONFIG_ESP_IPC_USES_CALLERS_PRIORITY

IPC runs at caller’s priority

Found in: Component config > Common ESP-related

If this option is not enabled then the IPC task will keep behavior same as prior to that of ESP-IDF v4.0, and hence IPC task will run at (configMAX_PRIORITIES - 1) priority.

CONFIG_ESP_MINIMAL_SHARED_STACK_SIZE

Minimal allowed size for shared stack

Found in: Component config > Common ESP-related

Minimal value of size, in bytes, accepted to execute a expression with shared stack.

CONFIG_ESP_CONSOLE_UART

UART for console output

Found in: Component config > Common ESP-related

Select whether to use UART for console output (through stdout and stderr).

- Default is to use UART0 on pre-defined GPIOs.
- If “Custom” is selected, UART0 or UART1 can be chosen, and any pins can be selected.
- If “None” is selected, there will be no console output on any UART, except for initial output from ROM bootloader. This output can be further suppressed by bootstrapping GPIO13 pin to low logic level.

Available options:

- Default: UART0 (ESP_CONSOLE_UART_DEFAULT)
- Custom (ESP_CONSOLE_UART_CUSTOM)
- None (ESP_CONSOLE_UART_NONE)

CONFIG_ESP_CONSOLE_UART_NUM

UART peripheral to use for console output (0-1)

Found in: Component config > Common ESP-related

Due of a ROM bug, UART2 is not supported for console output via ets_printf.

Available options:

- UART0 (ESP_CONSOLE_UART_CUSTOM_NUM_0)
- UART1 (ESP_CONSOLE_UART_CUSTOM_NUM_1)

CONFIG_ESP_CONSOLE_UART_TX_GPIO

UART TX on GPIO#

Found in: Component config > Common ESP-related

CONFIG_ESP_CONSOLE_UART_RX_GPIO

UART RX on GPIO#

Found in: Component config > Common ESP-related

CONFIG_ESP_CONSOLE_UART_BAUDRATE

UART console baud rate

Found in: Component config > Common ESP-related

CONFIG_ESP_INT_WDT

Interrupt watchdog

Found in: Component config > Common ESP-related

This watchdog timer can detect if the FreeRTOS tick interrupt has not been called for a certain time, either because a task turned off interrupts and did not turn them on for a long time, or because an interrupt handler did not return. It will try to invoke the panic handler first and failing that reset the SoC.

CONFIG_ESP_INT_WDT_TIMEOUT_MS

Interrupt watchdog timeout (ms)

Found in: Component config > Common ESP-related > CONFIG_ESP_INT_WDT

The timeout of the watchdog, in miliseconds. Make this higher than the FreeRTOS tick rate.

CONFIG_ESP_INT_WDT_CHECK_CPU1

Also watch CPU1 tick interrupt

Found in: Component config > Common ESP-related > CONFIG_ESP_INT_WDT

Also detect if interrupts on CPU 1 are disabled for too long.

CONFIG_ESP_TASK_WDT

Initialize Task Watchdog Timer on startup

Found in: Component config > Common ESP-related

The Task Watchdog Timer can be used to make sure individual tasks are still running. Enabling this option will cause the Task Watchdog Timer to be initialized automatically at startup. The Task Watchdog timer can be initialized after startup as well (see Task Watchdog Timer API Reference)

CONFIG_ESP_TASK_WDT_PANIC

Invoke panic handler on Task Watchdog timeout

Found in: Component config > Common ESP-related > CONFIG_ESP_TASK_WDT

If this option is enabled, the Task Watchdog Timer will be configured to trigger the panic handler when it times out. This can also be configured at run time (see Task Watchdog Timer API Reference)

CONFIG_ESP_TASK_WDT_TIMEOUT_S

Task Watchdog timeout period (seconds)

Found in: Component config > Common ESP-related > CONFIG_ESP_TASK_WDT

Timeout period configuration for the Task Watchdog Timer in seconds. This is also configurable at run time (see Task Watchdog Timer API Reference)

CONFIG_ESP_TASK_WDT_CHECK_IDLE_TASK_CPU0

Watch CPU0 Idle Task

Found in: Component config > Common ESP-related > CONFIG_ESP_TASK_WDT

If this option is enabled, the Task Watchdog Timer will watch the CPU0 Idle Task. Having the Task Watchdog watch the Idle Task allows for detection of CPU starvation as the Idle Task not being called is usually a symptom of CPU starvation. Starvation of the Idle Task is detrimental as FreeRTOS household tasks depend on the Idle Task getting some runtime every now and then.

CONFIG_ESP_TASK_WDT_CHECK_IDLE_TASK_CPU1

Watch CPU1 Idle Task

Found in: Component config > Common ESP-related > CONFIG_ESP_TASK_WDT

If this option is enabled, the Task Wtachdog Timer will wach the CPU1 Idle Task.

CONFIG_ESP_PANIC_HANDLER_IRAM

Place panic handler code in IRAM

Found in: Component config > Common ESP-related

If this option is disabled (default), the panic handler code is placed in flash not IRAM. This means that if ESP-IDF crashes while flash cache is disabled, the panic handler will automatically re-enable flash cache before running GDB Stub or Core Dump. This adds some minor risk, if the flash cache status is also corrupted during the crash.

If this option is enabled, the panic handler code is placed in IRAM. This allows the panic handler to run without needing to re-enable cache first. This may be necessary to debug some complex issues with crashes while flash cache is disabled (for example, when writing to SPI flash.)

ESP-TLS

Contains:

- *CONFIG_ESP_TLS_LIBRARY_CHOOSE*
- *CONFIG_ESP_TLS_SERVER*
- *CONFIG_ESP_TLS_PSK_VERIFICATION*
- *CONFIG_ESP_WOLFSSL_SMALL_CERT_VERIFY*
- *CONFIG_ESP_DEBUG_WOLFSSL*

CONFIG_ESP_TLS_LIBRARY_CHOOSE

Choose SSL/TLS library for ESP-TLS (See help for more Info)

Found in: Component config > ESP-TLS

The ESP-TLS APIs support multiple backend TLS libraries. Currently mbedTLS and WolfSSL are supported. Different TLS libraries may support different features and have different resource usage. Consult the ESP-TLS documentation in ESP-IDF Programming guide for more details.

Available options:

- mbedTLS (ESP_TLS_USING_MBEDTLS)
- wolfSSL (License info in wolfSSL directory README) (ESP_TLS_USING_WOLFSSL)

CONFIG_ESP_TLS_SERVER

Enable ESP-TLS Server

Found in: Component config > ESP-TLS

Enable support for creating server side SSL/TLS session, available for mbedTLS as well as wolfSSL TLS library.

CONFIG_ESP_TLS_PSK_VERIFICATION

Enable PSK verification

Found in: Component config > ESP-TLS

Enable support for pre shared key ciphers, supported for both mbedTLS as well as wolfSSL TLS library.

CONFIG_ESP_WOLFSSL_SMALL_CERT_VERIFY

Enable SMALL_CERT_VERIFY

Found in: Component config > ESP-TLS

Enables server verification with Intermediate CA cert, does not authenticate full chain of trust upto the root CA cert (After Enabling this option client only needs to have Intermediate CA certificate of the server to authenticate server, root CA cert is not necessary).

CONFIG_ESP_DEBUG_WOLFSSL

Enable debug logs for wolfSSL

Found in: Component config > ESP-TLS

Enable detailed debug prints for wolfSSL SSL library.

eFuse Bit Manager Contains:

- *CONFIG_EFUSE_CUSTOM_TABLE*
- *CONFIG_EFUSE_VIRTUAL*

CONFIG_EFUSE_CUSTOM_TABLE

Use custom eFuse table

Found in: Component config > eFuse Bit Manager

Allows to generate a structure for eFuse from the CSV file.

CONFIG_EFUSE_CUSTOM_TABLE_FILENAME

Custom eFuse CSV file

Found in: Component config > eFuse Bit Manager > CONFIG_EFUSE_CUSTOM_TABLE

Name of the custom eFuse CSV filename. This path is evaluated relative to the project root directory.

CONFIG_EFUSE_VIRTUAL

Simulate eFuse operations in RAM

Found in: Component config > eFuse Bit Manager

All read and writes operations are redirected to RAM instead of eFuse registers. If this option is set, all permanent changes (via eFuse) are disabled. Log output will state changes which would be applied, but they will not be.

Driver configurations Contains:

- *ADC configuration*
- *SPI configuration*
- *TWAI configuration*
- *UART configuration*

ADC configuration Contains:

- *CONFIG_ADC_FORCE_XPD_FSM*
- *CONFIG_ADC_DISABLE_DAC*

CONFIG_ADC_FORCE_XPD_FSM

Use the FSM to control ADC power

Found in: Component config > Driver configurations > ADC configuration

ADC power can be controlled by the FSM instead of software. This allows the ADC to be shut off when it is not working leading to lower power consumption. However using the FSM control ADC power will increase the noise of ADC.

CONFIG_ADC_DISABLE_DAC

Disable DAC when ADC2 is used on GPIO 25 and 26

Found in: Component config > Driver configurations > ADC configuration

If this is set, the ADC2 driver will disable the output of the DAC corresponding to the specified channel. This is the default value.

For testing, disable this option so that we can measure the output of DAC by internal ADC.

SPI configuration Contains:

- *CONFIG_SPI_MASTER_IN_IRAM*
- *CONFIG_SPI_MASTER_ISR_IN_IRAM*
- *CONFIG_SPI_SLAVE_IN_IRAM*
- *CONFIG_SPI_SLAVE_ISR_IN_IRAM*

CONFIG_SPI_MASTER_IN_IRAM

Place transmitting functions of SPI master into IRAM

Found in: Component config > Driver configurations > SPI configuration

Normally only the ISR of SPI master is placed in the IRAM, so that it can work without the flash when interrupt is triggered. For other functions, there's some possibility that the flash cache miss when running inside and out of SPI functions, which may increase the interval of SPI transactions. Enable this to put `queue\_trans`, `get\_trans\_result` and `transmit` functions into the IRAM to avoid possible cache miss.

During unit test, this is enabled to measure the ideal case of api.

CONFIG_SPI_MASTER_ISR_IN_IRAM

Place SPI master ISR function into IRAM

Found in: Component config > Driver configurations > SPI configuration

Place the SPI master ISR in to IRAM to avoid possible cache miss.

Also you can forbid the ISR being disabled during flash writing access, by add ESP_INTR_FLAG_IRAM when initializing the driver.

CONFIG_SPI_SLAVE_IN_IRAM

Place transmitting functions of SPI slave into IRAM

Found in: Component config > Driver configurations > SPI configuration

Normally only the ISR of SPI slave is placed in the IRAM, so that it can work without the flash when interrupt is triggered. For other functions, there's some possibility that the flash cache miss when running inside and out of SPI functions, which may increase the interval of SPI transactions. Enable this to put `queue\_trans`, `get\_trans\_result` and `transmit` functions into the IRAM to avoid possible cache miss.

CONFIG_SPI_SLAVE_ISR_IN_IRAM

Place SPI slave ISR function into IRAM

Found in: Component config > Driver configurations > SPI configuration

Place the SPI slave ISR in to IRAM to avoid possible cache miss.

Also you can forbid the ISR being disabled during flash writing access, by add ESP_INTR_FLAG_IRAM when initializing the driver.

TWAI configuration Contains:

- *CONFIG_TWAI_ISR_IN_IRAM*
- *CONFIG_TWAI_ERRATA_FIX_LISTEN_ONLY_DOM*

CONFIG_TWAI_ISR_IN_IRAM

Place TWAI ISR function into IRAM

Found in: Component config > Driver configurations > TWAI configuration

Place the TWAI ISR in to IRAM. This will allow the ISR to avoid cache misses, and also be able to run whilst the cache is disabled (such as when writing to SPI Flash). Note that if this option is enabled: - Users should also set the ESP_INTR_FLAG_IRAM in the driver configuration structure when installing the driver (see docs for specifics). - Alert logging (i.e., setting of the TWAI_ALERT_AND_LOG flag) will have no effect.

CONFIG_TWAI_ERRATA_FIX_LISTEN_ONLY_DOM

Add SW workaround for listen only transmits dominant bit errata

Found in: Component config > Driver configurations > TWAI configuration

When in the listen only mode, the TWAI controller must not influence the TWAI bus (i.e., must not send any dominant bits). However, while in listen only mode on the ESP32/ESP32-S2/ESP32-S3/ESP32-C3, the TWAI controller will still transmit dominant bits when it detects an error (i.e., as part of an active error frame). Enabling this option will add a workaround that forces the TWAI controller into an error passive state on initialization, thus preventing any dominant bits from being sent.

UART configuration Contains:

- *CONFIG_UART_ISR_IN_IRAM*

CONFIG_UART_ISR_IN_IRAM

Place UART ISR function into IRAM

Found in: Component config > Driver configurations > UART configuration

If this option is not selected, UART interrupt will be disabled for a long time and may cause data lost when doing spi flash operation.

Application Level Tracing Contains:

- *CONFIG_APPTRACE_DESTINATION*
- *CONFIG_APPTRACE_ONPANIC_HOST_FLUSH_TMO*
- *CONFIG_APPTRACE_POSTMORTEM_FLUSH_THRESH*
- *CONFIG_APPTRACE_PENDING_DATA_SIZE_MAX*
- *FreeRTOS SystemView Tracing*
- *CONFIG_APPTRACE_GCOV_ENABLE*

CONFIG_APPTRACE_DESTINATION

Data Destination

Found in: Component config > Application Level Tracing

Select destination for application trace: trace memory or none (to disable).

Available options:

- Trace memory (APPTRACE_DEST_TRAX)
- None (APPTRACE_DEST_NONE)

CONFIG_APPTRACE_ONPANIC_HOST_FLUSH_TMO

Timeout for flushing last trace data to host on panic

Found in: Component config > Application Level Tracing

Timeout for flushing last trace data to host in case of panic. In ms. Use -1 to disable timeout and wait forever.

CONFIG_APPTRACE_POSTMORTEM_FLUSH_THRESH

Threshold for flushing last trace data to host on panic

Found in: Component config > Application Level Tracing

Threshold for flushing last trace data to host on panic in post-mortem mode. This is minimal amount of data needed to perform flush. In bytes.

CONFIG_APPTRACE_PENDING_DATA_SIZE_MAX

Size of the pending data buffer

Found in: Component config > Application Level Tracing

Size of the buffer for events in bytes. It is useful for buffering events from the time critical code (scheduler, ISRs etc). If this parameter is 0 then events will be discarded when main HW buffer is full.

FreeRTOS SystemView Tracing Contains:

- *CONFIG_SYSVIEW_ENABLE*

CONFIG_SYSVIEW_ENABLE

SystemView Tracing Enable

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing

Enables supporrt for SEGGER SystemView tracing functionality.

CONFIG_SYSVIEW_TS_SOURCE

Timer to use as timestamp source

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

SystemView needs to use a hardware timer as the source of timestamps when tracing. This option selects the timer for it.

Available options:

- CPU cycle counter (CCOUNT) (SYSVIEW_TS_SOURCE_CCOUNT)
- Timer 0, Group 0 (SYSVIEW_TS_SOURCE_TIMER_00)
- Timer 1, Group 0 (SYSVIEW_TS_SOURCE_TIMER_01)
- Timer 0, Group 1 (SYSVIEW_TS_SOURCE_TIMER_10)
- Timer 1, Group 1 (SYSVIEW_TS_SOURCE_TIMER_11)
- esp_timer high resolution timer (SYSVIEW_TS_SOURCE_ESP_TIMER)

CONFIG_SYSVIEW_MAX_TASKS

Maximum supported tasks

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Configures maximum supported tasks in sysview debug

CONFIG_SYSVIEW_BUF_WAIT_TMO

Trace buffer wait timeout

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Configures timeout (in us) to wait for free space in trace buffer. Set to -1 to wait forever and avoid lost events.

CONFIG_SYSVIEW_EVT_OVERFLOW_ENABLE

Trace Buffer Overflow Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Trace Buffer Overflow" event.

CONFIG_SYSVIEW_EVT_ISR_ENTER_ENABLE

ISR Enter Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "ISR Enter" event.

CONFIG_SYSVIEW_EVT_ISR_EXIT_ENABLE

ISR Exit Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "ISR Exit" event.

CONFIG_SYSVIEW_EVT_ISR_TO_SCHEDULER_ENABLE

ISR Exit to Scheduler Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "ISR to Scheduler" event.

CONFIG_SYSVIEW_EVT_TASK_START_EXEC_ENABLE

Task Start Execution Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Task Start Execution" event.

CONFIG_SYSVIEW_EVT_TASK_STOP_EXEC_ENABLE

Task Stop Execution Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Task Stop Execution" event.

CONFIG_SYSVIEW_EVT_TASK_START_READY_ENABLE

Task Start Ready State Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Task Start Ready State" event.

CONFIG_SYSVIEW_EVT_TASK_STOP_READY_ENABLE

Task Stop Ready State Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Task Stop Ready State" event.

CONFIG_SYSVIEW_EVT_TASK_CREATE_ENABLE

Task Create Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Task Create" event.

CONFIG_SYSVIEW_EVT_TASK_TERMINATE_ENABLE

Task Terminate Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Task Terminate" event.

CONFIG_SYSVIEW_EVT_IDLE_ENABLE

System Idle Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "System Idle" event.

CONFIG_SYSVIEW_EVT_TIMER_ENTER_ENABLE

Timer Enter Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Timer Enter" event.

CONFIG_SYSVIEW_EVT_TIMER_EXIT_ENABLE

Timer Exit Event

Found in: Component config > Application Level Tracing > FreeRTOS SystemView Tracing > CONFIG_SYSVIEW_ENABLE

Enables "Timer Exit" event.

CONFIG_APPTRACE_GCOV_ENABLE

GCOV to Host Enable

Found in: Component config > Application Level Tracing

Enables support for GCOV data transfer to host.

Compatibility options

Contains:

- *CONFIG_LEGACY_INCLUDE_COMMON_HEADERS*

CONFIG_LEGACY_INCLUDE_COMMON_HEADERS

Include headers across components as before IDF v4.0

Found in: Compatibility options

Soc, esp32, and driver components, the most common components. Some header of these components are included implicitly by headers of other components before IDF v4.0. It's not required for high-level components, but still included through long header chain everywhere.

This is harmful to the modularity. So it's changed in IDF v4.0.

You can still include these headers in a legacy way until it is totally deprecated by enable this option.

Deprecated options and their replacements

- CONFIG_ADC2_DISABLE_DAC (*CONFIG_ADC_DISABLE_DAC*)
- CONFIG_APP_ANTI_ROLLBACK (*CONFIG_BOOTLOADER_APP_ANTI_ROLLBACK*)
- CONFIG_APP_ROLLBACK_ENABLE (*CONFIG_BOOTLOADER_APP_ROLLBACK_ENABLE*)
- CONFIG_APP_SECURE_VERSION (*CONFIG_BOOTLOADER_APP_SECURE_VERSION*)
- CONFIG_APP_SECURE_VERSION_SIZE_EFUSE_FIELD (*CONFIG_BOOTLOADER_APP_SEC_VER_SIZE_EFUSE_FIEL*
- **CONFIG_CONSOLE_UART (*CONFIG_ESP_CONSOLE_UART*)**
  - CONFIG_CONSOLE_UART_DEFAULT
  - CONFIG_CONSOLE_UART_CUSTOM
  - CONFIG_CONSOLE_UART_NONE
- CONFIG_CONSOLE_UART_BAUDRATE (*CONFIG_ESP_CONSOLE_UART_BAUDRATE*)
- **CONFIG_CONSOLE_UART_NUM (*CONFIG_ESP_CONSOLE_UART_NUM*)**
  - CONFIG_CONSOLE_UART_CUSTOM_NUM_0
  - CONFIG_CONSOLE_UART_CUSTOM_NUM_1
- CONFIG_CONSOLE_UART_RX_GPIO (*CONFIG_ESP_CONSOLE_UART_RX_GPIO*)
- CONFIG_CONSOLE_UART_TX_GPIO (*CONFIG_ESP_CONSOLE_UART_TX_GPIO*)
- CONFIG_CXX_EXCEPTIONS (*CONFIG_COMPILER_CXX_EXCEPTIONS*)
- CONFIG_CXX_EXCEPTIONS_EMG_POOL_SIZE (*CONFIG_COMPILER_CXX_EXCEPTIONS_EMG_POOL_SIZE*)
- CONFIG_DISABLE_GCC8_WARNINGS (*CONFIG_COMPILER_DISABLE_GCC8_WARNINGS*)
- CONFIG_EFUSE_SECURE_VERSION_EMULATE (*CONFIG_BOOTLOADER_EFUSE_SECURE_VERSION_EMULATE*)
- CONFIG_ENABLE_STATIC_TASK_CLEAN_UP_HOOK (*CONFIG_FREERTOS_ENABLE_STATIC_TASK_CLEAN_UP*)
- **CONFIG_ESP32S2_PANIC (*CONFIG_ESP_SYSTEM_PANIC*)**
  - CONFIG_ESP32S2_PANIC_PRINT_HALT
  - CONFIG_ESP32S2_PANIC_PRINT_REBOOT
  - CONFIG_ESP32S2_PANIC_SILENT_REBOOT
  - CONFIG_ESP32S2_PANIC_GDBSTUB
- **CONFIG_ESP32_APPTRACE_DESTINATION (*CONFIG_APPTRACE_DESTINATION*)**
  - CONFIG_ESP32_APPTRACE_DEST_TRAX
  - CONFIG_ESP32_APPTRACE_DEST_NONE
- CONFIG_ESP32_APPTRACE_ONPANIC_HOST_FLUSH_TMO (*CONFIG_APPTRACE_ONPANIC_HOST_FLUSH_TMO*)
- CONFIG_ESP32_APPTRACE_PENDING_DATA_SIZE_MAX (*CONFIG_APPTRACE_PENDING_DATA_SIZE_MAX*)
- CONFIG_ESP32_APPTRACE_POSTMORTEM_FLUSH_TRAX_THRESH (*CONFIG_APPTRACE_POSTMORTEM_FLUSH_THRESH*)
- CONFIG_ESP32_GCOV_ENABLE (*CONFIG_APPTRACE_GCOV_ENABLE*)
- **CONFIG_ESP32_PANIC (*CONFIG_ESP_SYSTEM_PANIC*)**
  - CONFIG_ESP32S2_PANIC_PRINT_HALT
  - CONFIG_ESP32S2_PANIC_PRINT_REBOOT
  - CONFIG_ESP32S2_PANIC_SILENT_REBOOT
  - CONFIG_ESP32S2_PANIC_GDBSTUB
- CONFIG_ESP32_PTHREAD_STACK_MIN (*CONFIG_PTHREAD_STACK_MIN*)
- **CONFIG_ESP32_PTHREAD_TASK_CORE_DEFAULT (*CONFIG_PTHREAD_TASK_CORE_DEFAULT*)**
  - CONFIG_ESP32_DEFAULT_PTHREAD_CORE_NO_AFFINITY
  - CONFIG_ESP32_DEFAULT_PTHREAD_CORE_0
  - CONFIG_ESP32_DEFAULT_PTHREAD_CORE_1
- CONFIG_ESP32_PTHREAD_TASK_NAME_DEFAULT (*CONFIG_PTHREAD_TASK_NAME_DEFAULT*)
- CONFIG_ESP32_PTHREAD_TASK_PRIO_DEFAULT (*CONFIG_PTHREAD_TASK_PRIO_DEFAULT*)
- CONFIG_ESP32_PTHREAD_TASK_STACK_SIZE_DEFAULT (*CONFIG_PTHREAD_TASK_STACK_SIZE_DEFAULT*)
- CONFIG_ESP_GRATUITOUS_ARP (*CONFIG_LWIP_ESP_GRATUITOUS_ARP*)
- CONFIG_ESP_TCP_KEEP_CONNECTION_WHEN_IP_CHANGES (*CONFIG_LWIP_TCP_KEEP_CONNECTION_WHEN_IP_CHANGES*)
- CONFIG_EVENT_LOOP_PROFILING (*CONFIG_ESP_EVENT_LOOP_PROFILING*)
- CONFIG_FLASH_ENCRYPTION_ENABLED (*CONFIG_SECURE_FLASH_ENC_ENABLED*)
- CONFIG_FLASH_ENCRYPTION_UART_BOOTLOADER_ALLOW_CACHE (*CONFIG_SECURE_FLASH_UART_BOOTLOADER_ALLOW_CACHE*)
- CONFIG_FLASH_ENCRYPTION_UART_BOOTLOADER_ALLOW_ENCRYPT (*CON-*

FIG_SECURE_FLASH_UART_BOOTLOADER_ALLOW_ENC)
- CONFIG_GARP_TMR_INTERVAL (*CONFIG_LWIP_GARP_TMR_INTERVAL*)
- CONFIG_GDBSTUB_MAX_TASKS (*CONFIG_ESP_GDBSTUB_MAX_TASKS*)
- CONFIG_GDBSTUB_SUPPORT_TASKS (*CONFIG_ESP_GDBSTUB_SUPPORT_TASKS*)
- CONFIG_INT_WDT (*CONFIG_ESP_INT_WDT*)
- CONFIG_INT_WDT_CHECK_CPU1 (*CONFIG_ESP_INT_WDT_CHECK_CPU1*)
- CONFIG_INT_WDT_TIMEOUT_MS (*CONFIG_ESP_INT_WDT_TIMEOUT_MS*)
- CONFIG_IPC_TASK_STACK_SIZE (*CONFIG_ESP_IPC_TASK_STACK_SIZE*)
- CONFIG_L2_TO_L3_COPY (*CONFIG_LWIP_L2_TO_L3_COPY*)
- **CONFIG_LOG_BOOTLOADER_LEVEL (*CONFIG_BOOTLOADER_LOG_LEVEL*)**
  - CONFIG_LOG_BOOTLOADER_LEVEL_NONE
  - CONFIG_LOG_BOOTLOADER_LEVEL_ERROR
  - CONFIG_LOG_BOOTLOADER_LEVEL_WARN
  - CONFIG_LOG_BOOTLOADER_LEVEL_INFO
  - CONFIG_LOG_BOOTLOADER_LEVEL_DEBUG
  - CONFIG_LOG_BOOTLOADER_LEVEL_VERBOSE
- CONFIG_MAIN_TASK_STACK_SIZE (*CONFIG_ESP_MAIN_TASK_STACK_SIZE*)
- CONFIG_MAKE_WARN_UNDEFINED_VARIABLES (*CONFIG_SDK_MAKE_WARN_UNDEFINED_VARIABLES*)
- CONFIG_MB_CONTROLLER_NOTIFY_QUEUE_SIZE (*CONFIG_FMB_CONTROLLER_NOTIFY_QUEUE_SIZE*)
- CONFIG_MB_CONTROLLER_NOTIFY_TIMEOUT (*CONFIG_FMB_CONTROLLER_NOTIFY_TIMEOUT*)
- CONFIG_MB_CONTROLLER_SLAVE_ID (*CONFIG_FMB_CONTROLLER_SLAVE_ID*)
- CONFIG_MB_CONTROLLER_SLAVE_ID_SUPPORT (*CONFIG_FMB_CONTROLLER_SLAVE_ID_SUPPORT*)
- CONFIG_MB_CONTROLLER_STACK_SIZE (*CONFIG_FMB_CONTROLLER_STACK_SIZE*)
- CONFIG_MB_EVENT_QUEUE_TIMEOUT (*CONFIG_FMB_EVENT_QUEUE_TIMEOUT*)
- CONFIG_MB_MASTER_DELAY_MS_CONVERT (*CONFIG_FMB_MASTER_DELAY_MS_CONVERT*)
- CONFIG_MB_MASTER_TIMEOUT_MS_RESPOND (*CONFIG_FMB_MASTER_TIMEOUT_MS_RESPOND*)
- CONFIG_MB_QUEUE_LENGTH (*CONFIG_FMB_QUEUE_LENGTH*)
- CONFIG_MB_SERIAL_BUF_SIZE (*CONFIG_FMB_SERIAL_BUF_SIZE*)
- CONFIG_MB_SERIAL_TASK_PRIO (*CONFIG_FMB_SERIAL_TASK_PRIO*)
- CONFIG_MB_SERIAL_TASK_STACK_SIZE (*CONFIG_FMB_SERIAL_TASK_STACK_SIZE*)
- CONFIG_MB_TIMER_GROUP (*CONFIG_FMB_TIMER_GROUP*)
- CONFIG_MB_TIMER_INDEX (*CONFIG_FMB_TIMER_INDEX*)
- CONFIG_MB_TIMER_PORT_ENABLED (*CONFIG_FMB_TIMER_PORT_ENABLED*)
- **CONFIG_MONITOR_BAUD (*CONFIG_ESPTOOLPY_MONITOR_BAUD*)**
  - CONFIG_MONITOR_BAUD_9600B
  - CONFIG_MONITOR_BAUD_57600B
  - CONFIG_MONITOR_BAUD_115200B
  - CONFIG_MONITOR_BAUD_230400B
  - CONFIG_MONITOR_BAUD_921600B
  - CONFIG_MONITOR_BAUD_2MB
  - CONFIG_MONITOR_BAUD_OTHER
- CONFIG_MONITOR_BAUD_OTHER_VAL (*CONFIG_ESPTOOLPY_MONITOR_BAUD_OTHER_VAL*)
- **CONFIG_OPTIMIZATION_ASSERTION_LEVEL (*CONFIG_COMPILER_OPTIMIZATION_ASSERTION_LEVEL*)**
  - CONFIG_OPTIMIZATION_ASSERTIONS_ENABLED
  - CONFIG_OPTIMIZATION_ASSERTIONS_SILENT
  - CONFIG_OPTIMIZATION_ASSERTIONS_DISABLED
- **CONFIG_OPTIMIZATION_COMPILER (*CONFIG_COMPILER_OPTIMIZATION*)**
  - CONFIG_COMPILER_OPTIMIZATION_LEVEL_DEBUG
  - CONFIG_COMPILER_OPTIMIZATION_LEVEL_RELEASE
- CONFIG_POST_EVENTS_FROM_IRAM_ISR (*CONFIG_ESP_EVENT_POST_FROM_IRAM_ISR*)
- CONFIG_POST_EVENTS_FROM_ISR (*CONFIG_ESP_EVENT_POST_FROM_ISR*)
- CONFIG_PPP_CHAP_SUPPORT (*CONFIG_LWIP_PPP_CHAP_SUPPORT*)
- CONFIG_PPP_DEBUG_ON (*CONFIG_LWIP_PPP_DEBUG_ON*)
- CONFIG_PPP_MPPE_SUPPORT (*CONFIG_LWIP_PPP_MPPE_SUPPORT*)
- CONFIG_PPP_MSCHAP_SUPPORT (*CONFIG_LWIP_PPP_MSCHAP_SUPPORT*)
- CONFIG_PPP_NOTIFY_PHASE_SUPPORT (*CONFIG_LWIP_PPP_NOTIFY_PHASE_SUPPORT*)
- CONFIG_PPP_PAP_SUPPORT (*CONFIG_LWIP_PPP_PAP_SUPPORT*)

- CONFIG_PPP_SUPPORT (*CONFIG_LWIP_PPP_SUPPORT*)
- CONFIG_PYTHON (*CONFIG_SDK_PYTHON*)
- CONFIG_SEMIHOSTFS_HOST_PATH_MAX_LEN (*CONFIG_VFS_SEMIHOSTFS_HOST_PATH_MAX_LEN*)
- CONFIG_SEMIHOSTFS_MAX_MOUNT_POINTS (*CONFIG_VFS_SEMIHOSTFS_MAX_MOUNT_POINTS*)
- **CONFIG_SPI_FLASH_WRITING_DANGEROUS_REGIONS (*CONFIG_SPI_FLASH_DANGEROUS_WRITE*)**
  - CONFIG_SPI_FLASH_WRITING_DANGEROUS_REGIONS_ABORTS
  - CONFIG_SPI_FLASH_WRITING_DANGEROUS_REGIONS_FAILS
  - CONFIG_SPI_FLASH_WRITING_DANGEROUS_REGIONS_ALLOWED
- **CONFIG_STACK_CHECK_MODE (*CONFIG_COMPILER_STACK_CHECK_MODE*)**
  - CONFIG_STACK_CHECK_NONE
  - CONFIG_STACK_CHECK_NORM
  - CONFIG_STACK_CHECK_STRONG
  - CONFIG_STACK_CHECK_ALL
- CONFIG_SUPPORT_STATIC_ALLOCATION (*CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION*)
- CONFIG_SUPPORT_TERMIOS (*CONFIG_VFS_SUPPORT_TERMIOS*)
- CONFIG_SUPPRESS_SELECT_DEBUG_OUTPUT (*CONFIG_VFS_SUPPRESS_SELECT_DEBUG_OUTPUT*)
- CONFIG_SYSTEM_EVENT_QUEUE_SIZE (*CONFIG_ESP_SYSTEM_EVENT_QUEUE_SIZE*)
- CONFIG_SYSTEM_EVENT_TASK_STACK_SIZE (*CONFIG_ESP_SYSTEM_EVENT_TASK_STACK_SIZE*)
- CONFIG_TASK_WDT (*CONFIG_ESP_TASK_WDT*)
- CONFIG_TASK_WDT_CHECK_IDLE_TASK_CPU0 (*CONFIG_ESP_TASK_WDT_CHECK_IDLE_TASK_CPU0*)
- CONFIG_TASK_WDT_CHECK_IDLE_TASK_CPU1 (*CONFIG_ESP_TASK_WDT_CHECK_IDLE_TASK_CPU1*)
- CONFIG_TASK_WDT_PANIC (*CONFIG_ESP_TASK_WDT_PANIC*)
- CONFIG_TASK_WDT_TIMEOUT_S (*CONFIG_ESP_TASK_WDT_TIMEOUT_S*)
- CONFIG_TCPIP_RECVMBOX_SIZE (*CONFIG_LWIP_TCPIP_RECVMBOX_SIZE*)
- **CONFIG_TCPIP_TASK_AFFINITY (*CONFIG_LWIP_TCPIP_TASK_AFFINITY*)**
  - CONFIG_TCPIP_TASK_AFFINITY_NO_AFFINITY
  - CONFIG_TCPIP_TASK_AFFINITY_CPU0
  - CONFIG_TCPIP_TASK_AFFINITY_CPU1
- CONFIG_TCPIP_TASK_STACK_SIZE (*CONFIG_LWIP_TCPIP_TASK_STACK_SIZE*)
- CONFIG_TCP_MAXRTX (*CONFIG_LWIP_TCP_MAXRTX*)
- CONFIG_TCP_MSL (*CONFIG_LWIP_TCP_MSL*)
- CONFIG_TCP_MSS (*CONFIG_LWIP_TCP_MSS*)
- **CONFIG_TCP_OVERSIZE (*CONFIG_LWIP_TCP_OVERSIZE*)**
  - CONFIG_TCP_OVERSIZE_MSS
  - CONFIG_TCP_OVERSIZE_QUARTER_MSS
  - CONFIG_TCP_OVERSIZE_DISABLE
- CONFIG_TCP_QUEUE_OOSEQ (*CONFIG_LWIP_TCP_QUEUE_OOSEQ*)
- CONFIG_TCP_RECVMBOX_SIZE (*CONFIG_LWIP_TCP_RECVMBOX_SIZE*)
- CONFIG_TCP_SND_BUF_DEFAULT (*CONFIG_LWIP_TCP_SND_BUF_DEFAULT*)
- CONFIG_TCP_SYNMAXRTX (*CONFIG_LWIP_TCP_SYNMAXRTX*)
- CONFIG_TCP_WND_DEFAULT (*CONFIG_LWIP_TCP_WND_DEFAULT*)
- CONFIG_TIMER_QUEUE_LENGTH (*CONFIG_FREERTOS_TIMER_QUEUE_LENGTH*)
- CONFIG_TIMER_TASK_PRIORITY (*CONFIG_FREERTOS_TIMER_TASK_PRIORITY*)
- CONFIG_TIMER_TASK_STACK_DEPTH (*CONFIG_FREERTOS_TIMER_TASK_STACK_DEPTH*)
- CONFIG_TIMER_TASK_STACK_SIZE (*CONFIG_ESP_TIMER_TASK_STACK_SIZE*)
- CONFIG_TOOLPREFIX (*CONFIG_SDK_TOOLPREFIX*)
- CONFIG_UDP_RECVMBOX_SIZE (*CONFIG_LWIP_UDP_RECVMBOX_SIZE*)
- CONFIG_USE_ONLY_LWIP_SELECT (*CONFIG_LWIP_USE_ONLY_LWIP_SELECT*)
- CONFIG_WARN_WRITE_STRINGS (*CONFIG_COMPILER_WARN_WRITE_STRINGS*)

2.7.7 Customisations

Because IDF builds by default with 警告未定义的变量, when the Kconfig tool generates Makefiles (the `auto.conf` file) its behaviour has been customised. In normal Kconfig, a variable which is set to "no" is undefined. In IDF's version of Kconfig, this variable is defined in the Makefile but has an empty value.

(Note that `ifdef` and `ifndef` can still be used in Makefiles, because they test if a variable is defined *and has a*

non-empty value.)

When generating header files for C & C++, the behaviour is not customised - so `#ifdef` can be used to test if a boolean config item is set or not.

2.8 Error Codes Reference

This section lists various error code constants defined in ESP-IDF.

For general information about error codes in ESP-IDF, see *Error Handling*.

ESP_FAIL (-1): Generic esp_err_t code indicating failure

ESP_OK (0): esp_err_t value indicating success (no error)

ESP_ERR_NO_MEM **(0x101)**: Out of memory

ESP_ERR_INVALID_ARG **(0x102)**: Invalid argument

ESP_ERR_INVALID_STATE **(0x103)**: Invalid state

ESP_ERR_INVALID_SIZE **(0x104)**: Invalid size

ESP_ERR_NOT_FOUND **(0x105)**: Requested resource not found

ESP_ERR_NOT_SUPPORTED **(0x106)**: Operation or feature not supported

ESP_ERR_TIMEOUT **(0x107)**: Operation timed out

ESP_ERR_INVALID_RESPONSE **(0x108)**: Received response was invalid

ESP_ERR_INVALID_CRC **(0x109)**: CRC or checksum was invalid

ESP_ERR_INVALID_VERSION **(0x10a)**: Version was invalid

ESP_ERR_INVALID_MAC **(0x10b)**: MAC address was invalid

ESP_ERR_NOT_FINISHED **(0x201)**

ESP_ERR_NVS_BASE **(0x1100)**: Starting number of error codes

ESP_ERR_NVS_NOT_INITIALIZED **(0x1101)**: The storage driver is not initialized

ESP_ERR_NVS_NOT_FOUND **(0x1102)**: Id namespace doesn’t exist yet and mode is NVS_READONLY

ESP_ERR_NVS_TYPE_MISMATCH **(0x1103)**: The type of set or get operation doesn’t match the type of value stored in NVS

ESP_ERR_NVS_READ_ONLY **(0x1104)**: Storage handle was opened as read only

ESP_ERR_NVS_NOT_ENOUGH_SPACE **(0x1105)**: There is not enough space in the underlying storage to save the value

ESP_ERR_NVS_INVALID_NAME **(0x1106)**: Namespace name doesn’t satisfy constraints

ESP_ERR_NVS_INVALID_HANDLE **(0x1107)**: Handle has been closed or is NULL

ESP_ERR_NVS_REMOVE_FAILED **(0x1108)**: The value wasn’t updated because flash write operation has failed. The value was written however, and update will be finished after re-initialization of nvs, provided that flash operation doesn’t fail again.

ESP_ERR_NVS_KEY_TOO_LONG **(0x1109)**: Key name is too long

ESP_ERR_NVS_PAGE_FULL **(0x110a)**: Internal error; never returned by nvs API functions

ESP_ERR_NVS_INVALID_STATE **(0x110b)**: NVS is in an inconsistent state due to a previous error. Call nvs_flash_init and nvs_open again, then retry.

ESP_ERR_NVS_INVALID_LENGTH **(0x110c)**: String or blob length is not sufficient to store data

ESP_ERR_NVS_NO_FREE_PAGES **(0x110d)**: NVS partition doesn't contain any empty pages. This may happen if NVS partition was truncated. Erase the whole partition and call nvs_flash_init again.

ESP_ERR_NVS_VALUE_TOO_LONG **(0x110e)**: String or blob length is longer than supported by the implementation

ESP_ERR_NVS_PART_NOT_FOUND **(0x110f)**: Partition with specified name is not found in the partition table

ESP_ERR_NVS_NEW_VERSION_FOUND **(0x1110)**: NVS partition contains data in new format and cannot be recognized by this version of code

ESP_ERR_NVS_XTS_ENCR_FAILED **(0x1111)**: XTS encryption failed while writing NVS entry

ESP_ERR_NVS_XTS_DECR_FAILED **(0x1112)**: XTS decryption failed while reading NVS entry

ESP_ERR_NVS_XTS_CFG_FAILED **(0x1113)**: XTS configuration setting failed

ESP_ERR_NVS_XTS_CFG_NOT_FOUND **(0x1114)**: XTS configuration not found

ESP_ERR_NVS_ENCR_NOT_SUPPORTED **(0x1115)**: NVS encryption is not supported in this version

ESP_ERR_NVS_KEYS_NOT_INITIALIZED **(0x1116)**: NVS key partition is uninitialized

ESP_ERR_NVS_CORRUPT_KEY_PART **(0x1117)**: NVS key partition is corrupt

ESP_ERR_NVS_CONTENT_DIFFERS **(0x1118)**: Internal error; never returned by nvs API functions. NVS key is different in comparison

ESP_ERR_ULP_BASE **(0x1200)**: Offset for ULP-related error codes

ESP_ERR_ULP_SIZE_TOO_BIG **(0x1201)**: Program doesn't fit into RTC memory reserved for the ULP

ESP_ERR_ULP_INVALID_LOAD_ADDR **(0x1202)**: Load address is outside of RTC memory reserved for the ULP

ESP_ERR_ULP_DUPLICATE_LABEL **(0x1203)**: More than one label with the same number was defined

ESP_ERR_ULP_UNDEFINED_LABEL **(0x1204)**: Branch instructions references an undefined label

ESP_ERR_ULP_BRANCH_OUT_OF_RANGE **(0x1205)**: Branch target is out of range of B instruction (try replacing with BX)

ESP_ERR_OTA_BASE **(0x1500)**: Base error code for ota_ops api

ESP_ERR_OTA_PARTITION_CONFLICT **(0x1501)**: Error if request was to write or erase the current running partition

ESP_ERR_OTA_SELECT_INFO_INVALID **(0x1502)**: Error if OTA data partition contains invalid content

ESP_ERR_OTA_VALIDATE_FAILED **(0x1503)**: Error if OTA app image is invalid

ESP_ERR_OTA_SMALL_SEC_VER **(0x1504)**: Error if the firmware has a secure version less than the running firmware.

ESP_ERR_OTA_ROLLBACK_FAILED **(0x1505)**: Error if flash does not have valid firmware in passive partition and hence rollback is not possible

ESP_ERR_OTA_ROLLBACK_INVALID_STATE **(0x1506)**: Error if current active firmware is still marked in pending validation state (ESP_OTA_IMG_PENDING_VERIFY), essentially first boot of firmware image post upgrade and hence firmware upgrade is not possible

ESP_ERR_EFUSE **(0x1600)**: Base error code for efuse api.

ESP_OK_EFUSE_CNT **(0x1601)**: OK the required number of bits is set.

ESP_ERR_EFUSE_CNT_IS_FULL **(0x1602)**: Error field is full.

ESP_ERR_EFUSE_REPEATED_PROG **(0x1603)**: Error repeated programming of programmed bits is strictly forbidden.

ESP_ERR_CODING **(0x1604)**: Error while a encoding operation.

ESP_ERR_IMAGE_BASE **(0x2000)**

ESP_ERR_IMAGE_FLASH_FAIL **(0x2001)**

ESP_ERR_IMAGE_INVALID **(0x2002)**

ESP_ERR_WIFI_BASE **(0x3000)**: Starting number of WiFi error codes

ESP_ERR_WIFI_NOT_INIT **(0x3001)**: WiFi driver was not installed by esp_wifi_init

ESP_ERR_WIFI_NOT_STARTED **(0x3002)**: WiFi driver was not started by esp_wifi_start

ESP_ERR_WIFI_NOT_STOPPED **(0x3003)**: WiFi driver was not stopped by esp_wifi_stop

ESP_ERR_WIFI_IF **(0x3004)**: WiFi interface error

ESP_ERR_WIFI_MODE **(0x3005)**: WiFi mode error

ESP_ERR_WIFI_STATE **(0x3006)**: WiFi internal state error

ESP_ERR_WIFI_CONN **(0x3007)**: WiFi internal control block of station or soft-AP error

ESP_ERR_WIFI_NVS **(0x3008)**: WiFi internal NVS module error

ESP_ERR_WIFI_MAC **(0x3009)**: MAC address is invalid

ESP_ERR_WIFI_SSID **(0x300a)**: SSID is invalid

ESP_ERR_WIFI_PASSWORD **(0x300b)**: Password is invalid

ESP_ERR_WIFI_TIMEOUT **(0x300c)**: Timeout error

ESP_ERR_WIFI_WAKE_FAIL **(0x300d)**: WiFi is in sleep state(RF closed) and wakeup fail

ESP_ERR_WIFI_WOULD_BLOCK **(0x300e)**: The caller would block

ESP_ERR_WIFI_NOT_CONNECT **(0x300f)**: Station still in disconnect status

ESP_ERR_WIFI_POST **(0x3012)**: Failed to post the event to WiFi task

ESP_ERR_WIFI_INIT_STATE **(0x3013)**: Invalid WiFi state when init/deinit is called

ESP_ERR_WIFI_STOP_STATE **(0x3014)**: Returned when WiFi is stopping

ESP_ERR_WIFI_NOT_ASSOC **(0x3015)**: The WiFi connection is not associated

ESP_ERR_WIFI_TX_DISALLOW **(0x3016)**: The WiFi TX is disallowed

ESP_ERR_WIFI_REGISTRAR **(0x3033)**: WPS registrar is not supported

ESP_ERR_WIFI_WPS_TYPE **(0x3034)**: WPS type error

ESP_ERR_WIFI_WPS_SM **(0x3035)**: WPS state machine is not initialized

ESP_ERR_ESPNOW_BASE **(0x3064)**: ESPNOW error number base.

ESP_ERR_ESPNOW_NOT_INIT **(0x3065)**: ESPNOW is not initialized.

ESP_ERR_ESPNOW_ARG **(0x3066)**: Invalid argument

ESP_ERR_ESPNOW_NO_MEM **(0x3067)**: Out of memory

ESP_ERR_ESPNOW_FULL **(0x3068)**: ESPNOW peer list is full

ESP_ERR_ESPNOW_NOT_FOUND **(0x3069)**: ESPNOW peer is not found

ESP_ERR_ESPNOW_INTERNAL **(0x306a)**: Internal error

ESP_ERR_ESPNOW_EXIST **(0x306b)**: ESPNOW peer has existed

ESP_ERR_ESPNOW_IF **(0x306c)**: Interface error

ESP_ERR_MESH_BASE **(0x4000)**: Starting number of MESH error codes

ESP_ERR_MESH_WIFI_NOT_START **(0x4001)**

ESP_ERR_MESH_NOT_INIT **(0x4002)**

ESP_ERR_MESH_NOT_CONFIG **(0x4003)**

ESP_ERR_MESH_NOT_START **(0x4004)**

ESP_ERR_MESH_NOT_SUPPORT **(0x4005)**

ESP_ERR_MESH_NOT_ALLOWED **(0x4006)**

ESP_ERR_MESH_NO_MEMORY **(0x4007)**

ESP_ERR_MESH_ARGUMENT **(0x4008)**

ESP_ERR_MESH_EXCEED_MTU **(0x4009)**

ESP_ERR_MESH_TIMEOUT **(0x400a)**

ESP_ERR_MESH_DISCONNECTED **(0x400b)**

ESP_ERR_MESH_QUEUE_FAIL **(0x400c)**

ESP_ERR_MESH_QUEUE_FULL **(0x400d)**

ESP_ERR_MESH_NO_PARENT_FOUND **(0x400e)**

ESP_ERR_MESH_NO_ROUTE_FOUND **(0x400f)**

ESP_ERR_MESH_OPTION_NULL **(0x4010)**

ESP_ERR_MESH_OPTION_UNKNOWN **(0x4011)**

ESP_ERR_MESH_XON_NO_WINDOW **(0x4012)**

ESP_ERR_MESH_INTERFACE **(0x4013)**

ESP_ERR_MESH_DISCARD_DUPLICATE **(0x4014)**

ESP_ERR_MESH_DISCARD **(0x4015)**

ESP_ERR_MESH_VOTING **(0x4016)**

ESP_ERR_MESH_XMIT **(0x4017)**

ESP_ERR_MESH_QUEUE_READ **(0x4018)**

ESP_ERR_MESH_PS **(0x4019)**

ESP_ERR_MESH_RECV_RELEASE **(0x401a)**

ESP_ERR_ESP_NETIF_BASE **(0x5000)**

ESP_ERR_ESP_NETIF_INVALID_PARAMS **(0x5001)**

ESP_ERR_ESP_NETIF_IF_NOT_READY **(0x5002)**

ESP_ERR_ESP_NETIF_DHCPC_START_FAILED **(0x5003)**

ESP_ERR_ESP_NETIF_DHCP_ALREADY_STARTED **(0x5004)**

ESP_ERR_ESP_NETIF_DHCP_ALREADY_STOPPED **(0x5005)**

ESP_ERR_ESP_NETIF_NO_MEM **(0x5006)**

ESP_ERR_ESP_NETIF_DHCP_NOT_STOPPED **(0x5007)**

ESP_ERR_ESP_NETIF_DRIVER_ATTACH_FAILED **(0x5008)**

ESP_ERR_ESP_NETIF_INIT_FAILED **(0x5009)**

ESP_ERR_ESP_NETIF_DNS_NOT_CONFIGURED **(0x500a)**

ESP_ERR_FLASH_BASE **(0x6000)**: Starting number of flash error codes

ESP_ERR_FLASH_OP_FAIL **(0x6001)**

ESP_ERR_FLASH_OP_TIMEOUT **(0x6002)**

ESP_ERR_FLASH_NOT_INITIALISED **(0x6003)**

ESP_ERR_FLASH_UNSUPPORTED_HOST **(0x6004)**

ESP_ERR_FLASH_UNSUPPORTED_CHIP **(0x6005)**

ESP_ERR_FLASH_PROTECTED **(0x6006)**

ESP_ERR_HTTP_BASE **(0x7000)**: Starting number of HTTP error codes

ESP_ERR_HTTP_MAX_REDIRECT **(0x7001)**: The error exceeds the number of HTTP redirects

ESP_ERR_HTTP_CONNECT **(0x7002)**: Error open the HTTP connection

ESP_ERR_HTTP_WRITE_DATA **(0x7003)**: Error write HTTP data

ESP_ERR_HTTP_FETCH_HEADER **(0x7004)**: Error read HTTP header from server

ESP_ERR_HTTP_INVALID_TRANSPORT **(0x7005)**: There are no transport support for the input scheme

ESP_ERR_HTTP_CONNECTING **(0x7006)**: HTTP connection hasn't been established yet

ESP_ERR_HTTP_EAGAIN **(0x7007)**: Mapping of errno EAGAIN to esp_err_t

ESP_ERR_ESP_TLS_BASE **(0x8000)**: Starting number of ESP-TLS error codes

ESP_ERR_ESP_TLS_CANNOT_RESOLVE_HOSTNAME **(0x8001)**: Error if hostname couldn't be resolved upon tls connection

ESP_ERR_ESP_TLS_CANNOT_CREATE_SOCKET **(0x8002)**: Failed to create socket

ESP_ERR_ESP_TLS_UNSUPPORTED_PROTOCOL_FAMILY **(0x8003)**: Unsupported protocol family

ESP_ERR_ESP_TLS_FAILED_CONNECT_TO_HOST **(0x8004)**: Failed to connect to host

ESP_ERR_ESP_TLS_SOCKET_SETOPT_FAILED **(0x8005)**: failed to set socket option

ESP_ERR_MBEDTLS_CERT_PARTLY_OK **(0x8006)**: mbedtls parse certificates was partly successful

ESP_ERR_MBEDTLS_CTR_DRBG_SEED_FAILED **(0x8007)**: mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_SET_HOSTNAME_FAILED **(0x8008)**: mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_CONFIG_DEFAULTS_FAILED **(0x8009)**: mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_CONF_ALPN_PROTOCOLS_FAILED **(0x800a)**: mbedtls api returned error

ESP_ERR_MBEDTLS_X509_CRT_PARSE_FAILED **(0x800b)**: mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_CONF_OWN_CERT_FAILED **(0x800c)**: mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_SETUP_FAILED **(0x800d)**: mbedtls api returned error

ESP_ERR_MBEDTLS_SSL_WRITE_FAILED **(0x800e)**: mbedtls api returned error

ESP_ERR_MBEDTLS_PK_PARSE_KEY_FAILED **(0x800f)**: mbedtls api returned failed

ESP_ERR_MBEDTLS_SSL_HANDSHAKE_FAILED **(0x8010)**: mbedtls api returned failed

ESP_ERR_MBEDTLS_SSL_CONF_PSK_FAILED **(0x8011)**: mbedtls api returned failed

ESP_ERR_ESP_TLS_CONNECTION_TIMEOUT **(0x8012)**: new connection in esp_tls_low_level_conn connection timeouted

ESP_ERR_WOLFSSL_SSL_SET_HOSTNAME_FAILED **(0x8013)**: wolfSSL api returned error

ESP_ERR_WOLFSSL_SSL_CONF_ALPN_PROTOCOLS_FAILED **(0x8014)**: wolfSSL api returned error

ESP_ERR_WOLFSSL_CERT_VERIFY_SETUP_FAILED **(0x8015)**: wolfSSL api returned error

ESP_ERR_WOLFSSL_KEY_VERIFY_SETUP_FAILED **(0x8016)**: wolfSSL api returned error

ESP_ERR_WOLFSSL_SSL_HANDSHAKE_FAILED **(0x8017)**: wolfSSL api returned failed

ESP_ERR_WOLFSSL_CTX_SETUP_FAILED **(0x8018)**: wolfSSL api returned failed

ESP_ERR_WOLFSSL_SSL_SETUP_FAILED **(0x8019)**: wolfSSL api returned failed

ESP_ERR_WOLFSSL_SSL_WRITE_FAILED **(0x801a)**: wolfSSL api returned failed

ESP_ERR_ESP_TLS_SE_FAILED **(0x801b)**

ESP_ERR_HTTPS_OTA_BASE **(0x9000)**

ESP_ERR_HTTPS_OTA_IN_PROGRESS **(0x9001)**

ESP_ERR_PING_BASE **(0xa000)**

ESP_ERR_PING_INVALID_PARAMS **(0xa001)**

ESP_ERR_PING_NO_MEM **(0xa002)**

ESP_ERR_HTTPD_BASE **(0xb000)**: Starting number of HTTPD error codes

ESP_ERR_HTTPD_HANDLERS_FULL **(0xb001)**: All slots for registering URI handlers have been consumed

ESP_ERR_HTTPD_HANDLER_EXISTS **(0xb002)**: URI handler with same method and target URI already registered

ESP_ERR_HTTPD_INVALID_REQ **(0xb003)**: Invalid request pointer

ESP_ERR_HTTPD_RESULT_TRUNC **(0xb004)**: Result string truncated

ESP_ERR_HTTPD_RESP_HDR **(0xb005)**: Response header field larger than supported

ESP_ERR_HTTPD_RESP_SEND **(0xb006)**: Error occured while sending response packet

ESP_ERR_HTTPD_ALLOC_MEM **(0xb007)**: Failed to dynamically allocate memory for resource

ESP_ERR_HTTPD_TASK **(0xb008)**: Failed to launch server task/thread

ESP_ERR_HW_CRYPTO_DS_BASE **(0xc000)**: Starting number of HW cryptography module error codes

ESP_ERR_HW_CRYPTO_DS_HMAC_FAIL **(0xc001)**: HMAC peripheral problem

ESP_ERR_HW_CRYPTO_DS_INVALID_KEY **(0xc002)**

ESP_ERR_HW_CRYPTO_DS_INVALID_DIGEST **(0xc004)**

ESP_ERR_HW_CRYPTO_DS_INVALID_PADDING **(0xc005)**

Chapter 3

ESP32-S2 H/W 硬件参考

3.1 ESP32-S2 系列模组和开发板

乐鑫设计并提供多种模组和开发板以供用户体验 ESP32-S2 系列芯片的强大功能。

本文档主要介绍了当前乐鑫所提供的各种模组和开发板。

注解： 如需了解较早版本或已停产的模组和开发板，请参考*ESP32-S2 模组与开发板（历史版本）*。

3.1.1 模组

ESP32-S2 系列模组集成了晶振、天线匹配电路等重要组件，可直接集成到终端产品中。如果再结合一些其他组件，例如编程接口、自举电阻和排针，您就可以体验 ESP32-S2 的强大功能了。

下表总结了上述模组的主要特点，详细信息见后文。

模组	芯片	Flash (MB)	PSRAM (MB)	天线	尺寸 (mm)
ESP32-S2-WROOM-32	ESP32-S2	2	N/A	MIFA	16 x 23 x 3

- MIFA - 蛇形倒 F 天线
- U.FL - U.FL / IPEX 天线连接器

3.1.2 开发板

ESP32-S2-Kaluga-1 套件 v1.3

ESP32-S2-Kaluga-1 是一款来自乐鑫的开发套件，包含 1 个主板和若干个扩展板，主要用于为用户提供基于 ESP32-S2 的人机交互应用开发工具。

相关文档

- *ESP32-S2-Kaluga-1 套件 v1.3*

较早版本

- *ESP32-S2-Kaluga-1 套件 v1.2*

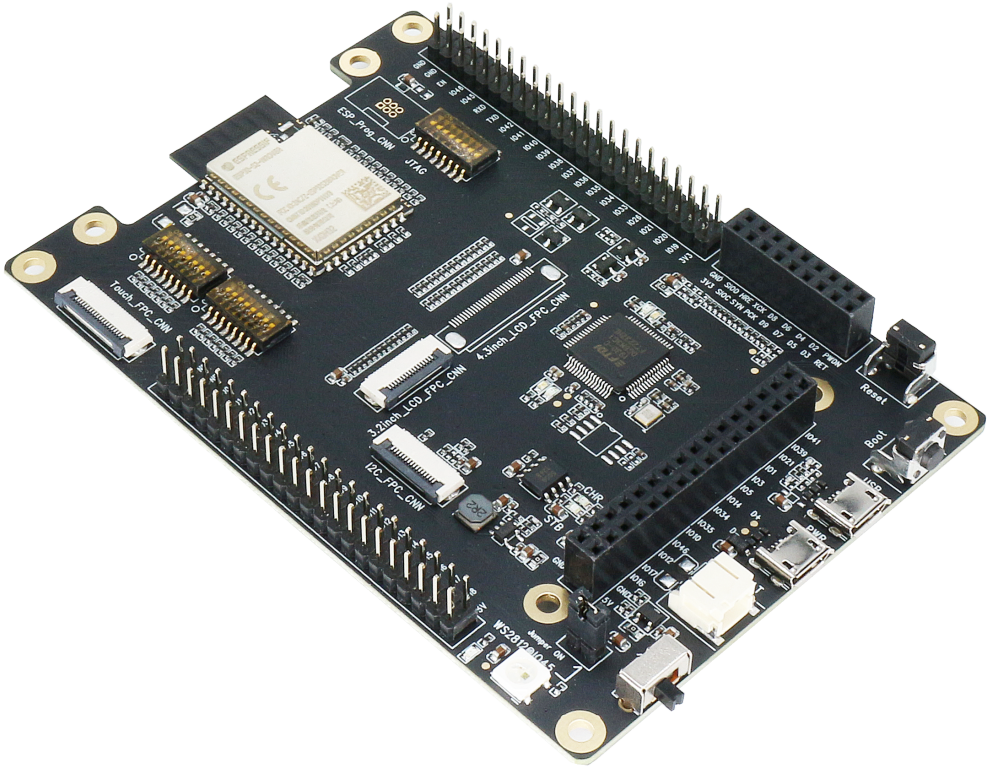

图 1: ESP32-S2-Kaluga-1（点击放大）

3.1.3 相关文档

- *ESP32-S2 模组与开发板（历史版本）*

3.2 ESP32-S2 模组与开发板（历史版本）

本节列出了旧版或已停产 ESP32-S2 模组和开发板的概述和文档链接，便于有意购买和使用旧版模组和开发板的用户参考。

3.2.1 模组

到目前为止没有模组停止更新或停产。

3.2.2 开发板

最新版本的开发板，请见章节*ESP32-S2 系列模组和开发板*。

ESP32-S2-Kaluga-1 套件 v1.2

ESP32-S2-Kaluga-1 是一款来自乐鑫的开发套件，包含 1 个主板和若干个扩展板，主要用于为用户提供基于 ESP32-S2 的人机交互应用开发工具。

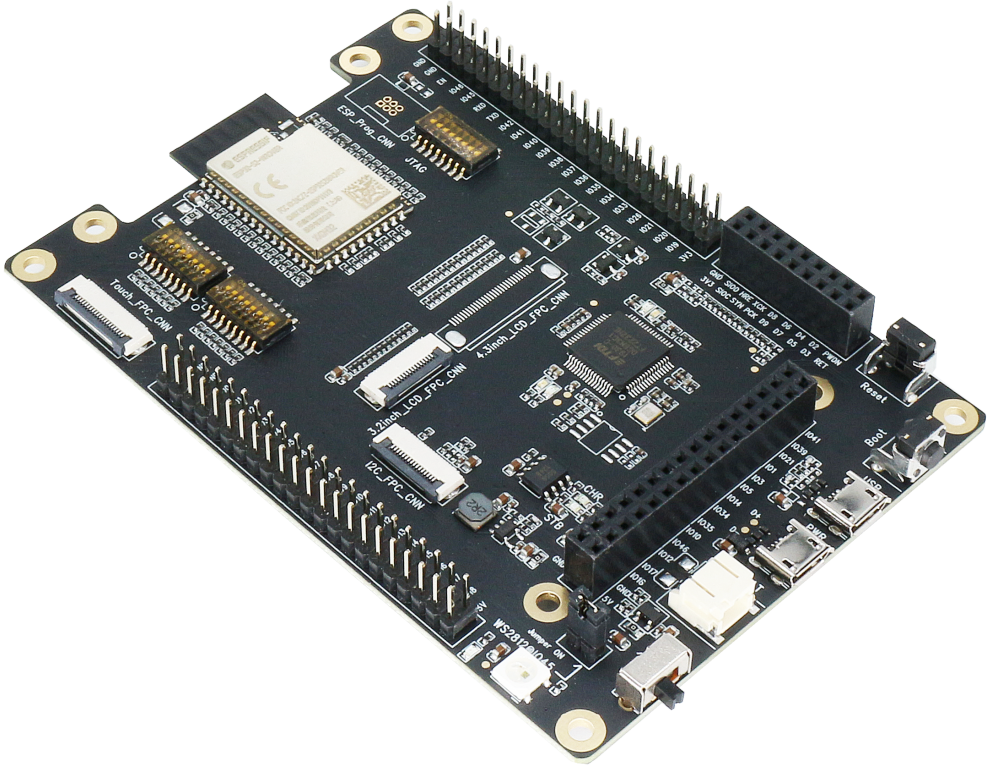

图 2: ESP32-S2-Kaluga-1 (click to enlarge)

相关文档

- *ESP32-S2-Kaluga-1 套件 v1.2*

3.2.3 相关文档

- *ESP32-S2 系列模组和开发板*

Chapter 4

API 指南

4.1 ESP-IDF 编程注意事项

4.1.1 应用程序的启动流程

本文将会介绍 ESP32-S2 从上电到运行 `app_main` 函数中间所经历的步骤（即启动流程）。

宏观上，该启动流程可以分为如下 3 个步骤：

1. 一级引导程序被固化在了 ESP32-S2 内部的 ROM 中，它会从 Flash 的 `0x1000` 偏移地址处加载二级引导程序至 RAM(IRAM & DRAM) 中。
2. 二级引导程序从 Flash 中加载分区表和主程序镜像至内存中，主程序中包含了 RAM 段和通过 Flash 高速缓存映射的只读段。
3. 主程序运行，这时第二个 CPU 和 RTOS 的调度器可以开始运行。

下面会对上述过程进行更为详细的阐述。

一级引导程序

SoC 复位后，PRO CPU 会立即开始运行，执行复位向量代码，而 APP CPU 仍然保持复位状态。在启动过程中，PRO CPU 会执行所有的初始化操作。APP CPU 的复位状态会在应用程序启动代码的 `call_start_cpu0` 函数中失效。复位向量代码位于 ESP32-S2 芯片掩膜 ROM 的 `0x40000400` 地址处，该地址不能被修改。

复位向量调用的启动代码会根据 `GPIO_STRAP_REG` 寄存器的值来确定 ESP32-S2 的工作模式，该寄存器保存着复位后 bootstrap 引脚的电平状态。根据不同的复位原因，程序会执行不同的操作：

1. 从深度睡眠模式复位：如果 `RTC_CNTL_STORE6_REG` 寄存器的值非零，并且 `RTC_CNTL_STORE7_REG` 寄存器中的 RTC 内存的 CRC 校验值有效，那么程序会使用 `RTC_CNTL_STORE6_REG` 寄存器的值作为入口地址，并立即跳转到该地址运行。如果 `RTC_CNTL_STORE6_REG` 的值为零，或者 `RTC_CNTL_STORE7_REG` 中的 CRC 校验值无效，又或者跳转到 `RTC_CNTL_STORE6_REG` 地址处运行的程序返回，那么将会执行上电复位的相关操作。**注意**：如果想在这里运行自定义的代码，可以参考深度睡眠 文档里面介绍的方法。
2. 上电复位、软件 SoC 复位、看门狗 SoC 复位：检查 `GPIO_STRAP_REG` 寄存器，判断是否 UART 或 SDIO 请求进入下载模式。如果是，则配置好 UART 或者 SDIO，然后等待下载代码。否则程序将会执行软件 CPU 复位的相关操作。
3. 软件 CPU 复位、看门狗 CPU 复位：根据 EFUSE 中的值配置 SPI Flash，然后尝试从 Flash 中加载代码，这部分的内存将会在后面一小节详细介绍。如果从 Flash 中加载代码失败，就会将 BASIC 解析器加压缩到 RAM 中启动。需要注意的是，此时 RTC 看门狗还在使能状态，如果在几百毫秒内没有任何输入事件，那么看门狗会再次复位 SoC，重复整个过程。如果解析器收到了来自 UART 的输入，程序会关闭看门狗。

应用程序的二进制镜像会从 Flash 的 `0x1000` 地址处加载。Flash 的第一个 4kB 扇区用于存储安全引导程序和应用程序镜像的签名。有关详细信息，请查看安全启动文档。

二级引导程序

在 ESP-IDF 中，存放在 Flash 的 0x1000 偏移地址处的二进制镜像就是二级引导程序。二级引导程序的源码可以在 ESP-IDF 的 components/bootloader 目录下找到。请注意，对于 ESP32-S2 芯片来说，这并不是唯一的安排程序镜像的方式。事实上用户完全可以把一个功能齐全的应用程序烧写到 Flash 的 0x1000 偏移地址处运行，但这超出本文档的范围。ESP-IDF 使用二级引导程序可以增加 Flash 分区的灵活性（使用分区表），并且方便实现 Flash 加密，安全引导和空中升级（OTA）等功能。

当一级引导程序校验并加载完二级引导程序后，它会从二进制镜像的头部找到二级引导程序的入口点，并跳转过去运行。

二级引导程序从 Flash 的 0x8000 偏移地址处读取分区表。详细信息请参阅分区表文档分区表 。二级引导程序会寻找出厂分区和 OTA 分区，然后根据 OTA 信息分区的数据决引导哪个分区。

对于选定的分区，二级引导程序将映射到 IRAM 和 DRAM 的数据和代码段复制到它们的加载地址处。对于一些加载地址位于 DROM 和 IROM 区域的段，会通过配置 Flash MMU 为其提供正确的映射。请注意，二级引导程序会为 PRO CPU 和 APP CPU 都配置 Flash MMU，但它只使能了 PRO CPU 的 Flash MMU。这么做的原因在于二级引导程序的代码被加载到了 APP CPU 的高速缓存使用的内存区域，因此使能 APP CPU 高速缓存的任务就交给了应用程序。一旦代码加载完毕并且设置好 Flash MMU，二级引导程序会从应用程序二进制镜像文件的头部寻找入口地址，然后跳转到该地址处运行。

目前还不支持添加钩子函数到二级引导程序中以自定义应用程序分区选择的逻辑，但是可以通过别的途径实现这个需求，比如根据某个 GPIO 的不同状态来引导不同的应用程序镜像。此类自定义的功能将在未来添加到 ESP-IDF 中。目前，可以通过将 bootloader 组件复制到应用程序目录并在那里进行必要的更改来自定义引导程序。在这种情况下，ESP-IDF 的编译系统将编译应用程序目录中的组件而不是 ESP-IDF 组件目录。

应用程序启动阶段

ESP-IDF 应用程序的入口是 components/esp32s2/cpu_start.c 文件中的 call_start_cpu0 函数，该函数主要完成了两件事，一是启用堆分配器，二是使 APP CPU 跳转到其入口点——call_start_cpu1 函数。PRO CPU 上的代码会给 APP CPU 设置好入口地址，解除其复位状态，然后等待 APP CPU 上运行的代码设置一个全局标志，以表明 APP CPU 已经正常启动。完成后，PRO CPU 跳转到 start_cpu0 函数，APP CPU 跳转到 start_cpu1 函数。

start_cpu0 和 start_cpu1 这两个函数都是弱类型的，这意味着如果某些特定的应用程序需要修改初始化顺序，就可以通过重写这两个函数来实现。start_cpu0 默认的实现方式是初始化用户在 menuconfig 中选择的组件，具体实现步骤可以阅读 components/esp32s2/cpu_start.c 文件中的源码。请注意，此阶段会调用应用程序中存在的 C++ 全局构造函数。一旦所有必要的组件都初始化好，就会创建 *main task* ，并启动 FreeRTOS 的调度器。

当 PRO CPU 在 start_cpu0 函数中进行初始化的时候，APP CPU 在 start_cpu1 函数中自旋，等待 PRO CPU 上的调度器启动。一旦 PRO CPU 上的调度器启动后，APP CPU 上的代码也会启动调度器。

主任务是指运行 app_main 函数的任务，主任务的堆栈大小和优先级可以在 menuconfig 中进行配置。应用程序可以用此任务来完成用户程序相关的初始化设置，比如启动其他的任务。应用程序还可以将主任务用于事件循环和其他通用活动。如果 app_main 函数返回，那么主任务将会被删除。

4.1.2 应用程序的内存布局

ESP32-S2 芯片具有灵活的内存映射功能，本小节将介绍 ESP-IDF 默认使用这些功能的方式。

ESP-IDF 应用程序的代码可以放在以下内存区域之一。

IRAM（指令 RAM）

ESP-IDF 将内部 SRAM0 区域（在技术参考手册中有定义）的一部分分配为指令 RAM。除了开始的 64kB 用作 PRO CPU 和 APP CPU 的高速缓存外，剩余内存区域（从 0x40080000 至 0x400A0000 ）被用来存储应用程序中部分需要在 RAM 中运行的代码。

一些 ESP-IDF 的组件和 WiFi 协议栈的部分代码通过链接脚本文件被存放到了这块内存区域。

如果一些应用程序的代码需要放在 IRAM 中运行，可以使用 IRAM_ATTR 宏定义进行声明。

```
#include "esp_attr.h"

void IRAM_ATTR gpio_isr_handler(void* arg)
{
    // ...
}
```

下面列举了应用程序中可能或者应该放入 IRAM 中运行例子。

- 当注册中断处理程序的时候设置了 ESP_INTR_FLAG_IRAM ，那么中断处理程序就必须要放在 IRAM 中运行。这种情况下，ISR 只能调用存放在 IRAM 或者 ROM 中的函数。注意：目前所有 FreeRTOS 的 API 都已经存放到了 IRAM 中，所以在中断中调用 FreeRTOS 的中断专属 API 是安全的。如果将 ISR 放在 IRAM 中运行，那么必须使用宏定义 DRAM_ATTR 将该 ISR 用到所有常量数据和调用的函数（包括但不限于 const char 数组）放入 DRAM 中。
- 可以将一些时间关键的代码放在 IRAM 中，这样可以缩减从 Flash 加载代码所消耗的时间。ESP32-S2 是通过 32kB 的高速缓存来从外部 Flash 中读取代码和数据的，将函数放在 IRAM 中运行可以减少由高速缓存未命中引起的时间延迟。

IROM（代码从 Flash 中运行）

如果一个函数没有被显式地声明放在 IRAM 或者 RTC 内存中，则将其置于 Flash 中。Flash 技术参考手册中介绍了 Flash MMU 允许代码从 Flash 执行的机制。ESP-IDF 将从 Flash 中执行的代码放在 0x400D0000—0x40400000 区域的开始，在启动阶段，二级引导程序会初始化 Flash MMU，将代码在 Flash 中的位置映射到这个区域的开头。对这个区域的访问会被透明地缓存到 0x40070000 —0x40080000 范围内的两个 32kB 的块中。

请注意，使用 Window ABI CALLx 指令可能无法访问 0x40000000 —0x40400000 区域以外的代码，所以要特别留意应用程序是否使用了 0x40400000 —0x40800000 或者 0x40800000 —0x40C00000 区域，ESP-IDF 默认不会使用这两个区域。

RTC 快速内存

从深度睡眠模式唤醒后必须要运行的代码要放在 RTC 内存中，更多信息请查阅文档深度睡眠 。

DRAM（数据 RAM）

链接器将非常量静态数据和零初始化数据放入 0x3FFB0000 — 0x3FFF0000 这 256kB 的区域。注意，如果使用蓝牙堆栈，此区域会减少 64kB（通过将起始地址移至 0x3FFC0000 ）。如果使用了内存跟踪的功能，该区域的长度还要减少 16kB 或者 32kB。放置静态数据后，留在此区域中的剩余空间都用作运行时堆。

常量数据也可以放在 DRAM 中，例如，用在 ISR 中的常量数据（参见上面 IRAM 部分的介绍），为此需要使用 DRAM_ATTR 宏来声明。

```
DRAM_ATTR const char[] format_string = "%p %x";
char buffer[64];
sprintf(buffer, format_string, ptr, val);
```

毋庸置疑，不建议在 ISR 中使用 printf 和其余输出函数。出于调试的目的，可以在 ISR 中使用 ESP_EARLY_LOGx 来输出日志，不过要确保将 TAG 和格式字符串都放在了 DRAM 中。

宏 __NOINIT_ATTR 可以用来声明将数据放在 .noinit 段中，放在此段中的数据不会在启动时被初始化，并且在软件重启后会保留原来的值。

例子：

```
__NOINIT_ATTR uint32_t noinit_data;
```

DROM（数据存储在 Flash 中）

默认情况下，链接器将常量数据放入一个 4MB 区域 (0x3F400000 — 0x3F800000)，该区域用于通过 Flash MMU 和高速缓存来访问外部 Flash。一种特例情况是，字面量会被编译器嵌入到应用程序代码中。

RTC 慢速内存

从 RTC 内存运行的代码（例如深度睡眠模块的代码）使用的全局和静态变量必须要放在 RTC 慢速内存中。更多详细说明请查看文档深度睡眠 。

宏 RTC_NOINIT_ATTR 用来声明将数据放入 RTC 慢速内存中，该数据在深度睡眠唤醒后将保持不变。

例子：

```
RTC_NOINIT_ATTR uint32_t rtc_noinit_data;
```

4.1.3 DMA 能力要求

大多数的 DMA 控制器（比如 SPI，SDMMC 等）都要求发送/接收缓冲区放在 DRAM 中，并且按字对齐。我们建议将 DMA 缓冲区放在静态变量中而不是堆栈中。使用 DMA_ATTR 宏可以声明该全局/本地的静态变量具备 DMA 能力，例如：

```
DMA_ATTR uint8_t buffer[]="I want to send something";

void app_main()
{
    // 初始化代码...
    spi_transaction_t temp = {
        .tx_buffer = buffer,
        .length = 8*sizeof(buffer),
    };
    spi_device_transmit( spi, &temp );
    // 其他程序
}
```

或者：

```
void app_main()
{
    DMA_ATTR static uint8_t buffer[]="I want to send something";
    // 初始化代码...
    spi_transaction_t temp = {
        .tx_buffer = buffer,
        .length = 8*sizeof(buffer),
    };
    spi_device_transmit( spi, &temp );
    // 其他程序
}
```

在堆栈中放置 DMA 缓冲区仍然是允许的，但是你必须记住：

1. 如果堆栈在 pSRAM 中，切勿尝试这么做，因为堆栈在 pSRAM 中的话就要按照片外 *SRAM* 文档介绍的步骤来操作（至少要在 menuconfig 中使能 SPIRAM_ALLOW_STACK_EXTERNAL_MEMORY），所以请确保你的任务不在 pSRAM 中。
2. 在函数中使用 WORD_ALIGNED_ATTR 宏来修饰变量，将其放在适当的位置上，比如：

```
void app_main()
{
    uint8_t stuff;
    WORD_ALIGNED_ATTR uint8_t buffer[]="I want to send something";    //否则
buffer 数组会被存储在 stuff 变量的后面
    // 初始化代码...
    spi_transaction_t temp = {
        .tx_buffer = buffer,
        .length = 8*sizeof(buffer),
    };
    spi_device_transmit( spi, &temp );
    // 其他程序
}
```

4.2 构建系统（CMake 版）

本文档主要介绍 ESP-IDF 构建系统的实现原理以及 组件等相关概念。如需您想了解如何组织和构建新的 ESP-IDF 项目或组件，请阅读本文档。

4.2.1 概述

一个 ESP-IDF 项目可以看作是多个不同组件的集合，例如一个显示当前湿度的网页服务器会包含以下组件：

- ESP-IDF 基础库，包括 libc、ROM bindings 等
- Wi-Fi 驱动
- TCP/IP 协议栈
- FreeRTOS 操作系统
- 网页服务器
- 湿度传感器的驱动
- 负责将上述组件整合到一起的主程序

ESP-IDF 可以显式地指定和配置每个组件。在构建项目的时候，构建系统会前往 ESP-IDF 目录、项目目录和用户自定义组件目录（可选）中查找所有组件，允许用户通过文本菜单系统配置 ESP-IDF 项目中用到的每个组件。在所有组件配置结束后，构建系统开始编译整个项目。

概念

- 项目特指一个目录，其中包含了构建可执行应用程序所需的全部文件和配置，以及其他支持型文件，例如分区表、数据/文件系统分区和引导程序。
- 项目配置保存在项目根目录下名为 `sdkconfig` 的文件中，可以通过 `idf.py menuconfig` 进行修改，且一个项目只能包含一个项目配置。
- 应用程序是由 ESP-IDF 构建得到的可执行文件。一个项目通常会构建两个应用程序：项目应用程序（可执行的主文件，即用户自定义的固件）和引导程序（启动并初始化项目应用程序）。
- 组件是模块化且独立的代码，会被编译成静态库（.a 文件）并链接到应用程序。部分组件由 ESP-IDF 官方提供，其他组件则来源于其它开源项目。
- 目标特指运行构建后应用程序的硬件设备。ESP-IDF 当前仅支持 `esp32` 和 `esp32s2` 这三个硬件目标。

请注意，以下内容并不属于项目的组成部分：

- `ESP-IDF` 并不是项目的一部分，它独立于项目，通过 `IDF_PATH` 环境变量（保存 `esp-idf` 目录的路径）链接到项目，从而将 IDF 框架与项目分离。
- 交叉编译工具链并不是项目的组成部分，它应该被安装在系统 PATH 环境变量中。

4.2.2 使用构建系统

idf.py

idf.py 命令行工具提供了一个前端，可以帮助您轻松管理项目的构建过程，它管理了以下工具：

- CMake，配置待构建的项目
- 命令行构建工具（Ninja 或 *GNU Make*）
- esptool.py，烧录目标硬件设备

入门指南 简要介绍了如何设置 idf.py 用于配置、构建并烧录项目。

idf.py 应运行在 ESP-IDF 的 项目目录下，即包含 CMakeLists.txt 文件的目录。仅包含 Makefile 的老式项目并不支持 idf.py。

运行 idf.py --help 查看完整的命令列表。下面总结了最常用的命令：

- idf.py set-target <target> 会设置构建项目的目标（芯片）。请参考选择目标芯片。
- idf.py menuconfig 会运行 menuconfig 工具来配置项目。
- idf.py build 会构建在当前目录下找到的项目，它包括以下步骤：
  - 根据需要创建 build 构建目录，它用于保存构建过程的输出文件，可以使用 -B 选项修改默认的构建目录。
  - 根据需要运行 CMake 来配置项目，为主构建工具生成构建文件。
  - 运行主构建工具（Ninja 或 *GNU Make*）。默认情况下，构建工具会被自动检测，可以使用 -G 选项显式地指定构建工具。

  构建过程是增量式的，如果自上次构建以来源文件或项目配置没有发生改变，则不会执行任何操作。
- idf.py clean 会把构建输出的文件从构建目录中删除，从而清理整个项目。下次构建时会强制“重新完整构建”这个项目。清理时，不会删除 CMake 配置输出及其他文件。
- idf.py fullclean 会将整个 build 目录下的内容全部删除，包括所有 CMake 的配置输出文件。下次构建项目时，CMake 会从头开始配置项目。请注意，该命令会递归删除构建目录下的 所有文件，请谨慎使用。项目配置文件不会被删除。
- idf.py flash 会在必要时自动构建项目，并将生成的二进制程序烧录进目标 ESP32-S2 设备中。-p 和 -b 选项可分别设置串口的设备名和烧录时的波特率。
- idf.py monitor 用于显示目标 ESP32-S2 设备的串口输出。-p 选项可用于设置主机端串口的设备名，按下 Ctrl-] 可退出监视器。更多有关监视器的详情，请参阅*IDF 监视器*。

多个 idf.py 命令可合并成一个，例如，idf.py -p COM4 clean flash monitor 会依次清理源码树，构建项目，烧录进目标 ESP32-S2 设备，最后运行串口监视器。

对于 idf.py 不知道的指令，idf.py 会尝试将其作为构建系统的目标来执行。

注解： 环境变量 ESPPORT 和 ESPBAUD 可分别用来设置 -p 和 -b 选项的默认值。在命令行中，重新为这两个选项赋值，会覆盖其默认值。

高级命令

- idf.py app，idf.py bootloader，idf.py partition_table 仅可用于从适用的项目中构建应用程序、引导程序或分区表。
- idf.py app-flash 等匹配命令，仅将项目的特定部分烧录至 ESP32-S2。
- idf.py -p PORT erase_flash 会使用 esptool.py 擦除 ESP32-S2 的整个 Flash。
- idf.py size 会打印应用程序相关的大小信息，idf.py size-components 和 idf.py size-files 这两个命令相似，分别用于打印每个组件或源文件的详细信息。如果您在运行 CMake（或 idf.py）时定义了变量 -DOUTPUT_JSON=1，那么输出的格式会变成 JSON 而不是可读文本。
- idf.py reconfigure 命令会重新运行 CMake （即便无需重新运行）。正常使用时，并不需要运行此命令，但当源码树中添加/删除文件后或更改 CMake cache 变量时，此命令会非常有用，例如，idf.py -DNAME='VALUE' reconfigure 会将 CMake cache 中的变量 NAME 的值设置为 VALUE。

- `idf.py python-clean` 会从 IDF 目录中删除生成的 Python 字节码，Python 字节码可能会在切换 IDF 和 Python 版本时引发问题，因此建议在切换 Python 后运行该命令。

同时调用多个 `idf.py` 命令时，命令的输入顺序并不重要，它们会按照正确的顺序依次执行，并保证每一条命令都生效（即先构建后烧录，先擦除后烧录等）。

idf.py 选项 运行 `idf.py --help` 命令列出所有根级选项。运行 `idf.py <command> --help` 命令列出针对某一子命令的选项，如 `idf.py monitor --help`。下列是一些常用选项：

- `-C <dir>` 可用来从默认的当前工作目录覆盖项目目录。
- `-B <dir>` 可用来从项目目录默认的 `build` 子目录覆盖构建目录。
- `--ccache` 可用来在编译源文件时启用 CCache，安装了 CCache 工具后可极大缩短编译时间。

请注意，一些旧版本的 CCache 在某些平台上可能会出现 bug，因此如果文件没有按预期重新构建，请尝试禁用 CCache 后再次构建。通过设置环境变量 `IDF_CCACHE_ENABLE` 为非零值，可以默认启用 CCache。

- `-v` 可以让 `idf.py` 和编译系统产生详细的编译输出，对于调试编译问题会非常有用。
- `--cmake-warn-uninitialized`` （或 ``-w`）会让 CMake 在项目目录内打印未初始化的变量警告（不包括在项目目录外的目录），这一选项只能控制 CMake 内部的 CMake 变量警告，不包括其它类型的编译警告。可以通过设置环境变量 `IDF_CMAKE_WARN_UNINITIALIZED` 为非零值，从而永久设置该选项。

直接使用 CMake

为了方便，*idf.py* 已经封装了 CMake 命令，但是您愿意，也可以直接调用 CMake。

当 `idf.py` 在执行某些操作时，它会打印出其运行的每条命令以便参考。例如运行 `idf.py build` 命令与在 bash shell（或者 Windows Command Prompt）中运行以下命令是相同的：

```
mkdir -p build
cd build
cmake .. -G Ninja    # 或者 'Unix Makefiles'
ninja
```

在上面的命令列表中，`cmake` 命令对项目进行配置，并生成用于最终构建工具的构建文件。在这个例子中，最终构建工具是 Ninja: 运行 `ninja` 来构建项目。

没有必要多次运行 `cmake`。第一次构建后，往后每次只需运行 `ninja` 即可。如果项目需要重新配置，`ninja` 会自动重新调用 `cmake`。

若在 CMake 中使用 `ninja` 或 `make`，则多数 `idf.py` 子命令也会有其对应的目标，例如在构建目录下运行 `make menuconfig` 或 `ninja menuconfig` 与运行 `idf.py menuconfig` 是相同的。

注解： 如果您已经熟悉了 CMake，那么可能会发现 ESP-IDF 的 CMake 构建系统不同寻常，为了减少样板文件，该系统封装了 CMake 的许多功能。请参考*编写纯 CMake 组件* 以编写更多 “CMake 风格” 的组件。

使用 Ninja/Make 来烧录 您可以直接使用 ninja 或 make 运行如下命令来构建项目并烧录：

```
ninja flash
```

或：

```
make app-flash
```

可用的目标还包括：`flash`、`app-flash`（仅用于 app）、`bootloader-flash`（仅用于 bootloader）。

以这种方式烧录时，可以通过设置 `ESPPORT` 和 `ESPBAUD` 环境变量来指定串口设备和波特率。您可以在操作系统或 IDE 项目中设置该环境变量，或者直接在命令行中进行设置：

```
ESPPORT=/dev/ttyUSB0 ninja flash
```

注解： 在命令的开头为环境变量赋值属于 Bash shell 的语法，可在 Linux 、macOS 和 Windows 的类 Bash shell 中运行，但在 Windows Command Prompt 中无法运行。

或:

```
make -j3 app-flash ESPPORT=COM4 ESPBAUD=2000000
```

注解： 在命令末尾为变量赋值属于 `make` 的语法，适用于所有平台的 `make`。

在 IDE 中使用 CMake

您还可以使用集成了 CMake 的 IDE，仅需将项目 `CMakeLists.txt` 文件的路径告诉 IDE 即可。集成 CMake 的 IDE 通常会有自己的构建工具（CMake 称之为“生成器”），它是组成 IDE 的一部分，用来构建源文件。

向 IDE 中添加除 `build` 目标以外的自定义目标（如添加“Flash”目标到 IDE）时，建议调用 `idf.py` 命令来执行这些“特殊”的操作。

有关将 ESP-IDF 同 CMake 集成到 IDE 中的详细信息，请参阅*构建系统的元数据*。

设置 Python 解释器

ESP-IDF 适用于所有支持的 Python 版本。即使您系统中默认的 `python` 解释器仍是 Python 2.7，ESP-IDF 也可以使用，但建议您升级至 Python 3。

`idf.py` 和其他的 Python 脚本会使用默认的 Python 解释器运行，如 `python`。您可以通过 `python3 $IDF_PATH/tools/idf.py ...` 命令切换到别的 Python 解释器，或者您可以通过设置 shell 别名或其他脚本来简化该命令。

如果直接使用 CMake，运行 `cmake -D PYTHON=python3 ...`，CMake 会使用传入的值覆盖默认的 Python 解释器。

如果使用集成 CMake 的 IDE，可以在 IDE 的图形用户界面中给名为 `PYTHON` 的 CMake cache 变量设置新的值来覆盖默认的 Python 解释器。

如果想在命令行中更优雅地管理 Python 的各个版本，请查看 pyenv 或 virtualenv 工具，它们会帮助您更改默认的 python 版本。

潜在问题 使用 `idf.py` 可能会出现如下 `ImportError` 错误:

```
Traceback (most recent call last):
  File "/Users/user_name/e/esp-idf/tools/kconfig_new/confgen.py", line 27, in
↪<module>
    import kconfiglib
ImportError: bad magic number in 'kconfiglib': b'\x03\xf3\r\n'
```

该错误通常是由不同 Python 版本生成的 `.pyc` 文件引起的，可以通过运行以下命令解决该问题:

```
idf.py python-clean
```

4.2.3 示例项目

示例项目的目录树结构可能如下所示:

```
- myProject/
             - CMakeLists.txt
             - sdkconfig
             - components/ - component1/ - CMakeLists.txt
                                         - Kconfig
                                         - src1.c
                           - component2/ - CMakeLists.txt
                                         - Kconfig
                                         - src1.c
                                         - include/ - component2.h
             - main/       - CMakeLists.txt
                           - src1.c
                           - src2.c

             - build/
```

该示例项目 “myProject” 包含以下组成部分:

- 顶层项目 CMakeLists.txt 文件，这是 CMake 用于学习如何构建项目的主要文件，可以在这个文件中设置项目全局的 CMake 变量。顶层项目 CMakeLists.txt 文件会导入 /tools/cmake/project.cmake 文件，由它负责实现构建系统的其余部分。该文件最后会设置项目的名称，并定义该项目。
- “sdkconfig” 项目配置文件，执行 `idf.py menuconfig` 时会创建或更新此文件，文件中保存了项目中所有组件（包括 ESP-IDF 本身）的配置信息。`sdkconfig` 文件可能会也可能不会被添加到项目的源码管理系统中。
- 可选的 “components” 目录中包含了项目的部分自定义组件，并不是每个项目都需要这种自定义组件，但它有助于构建可复用的代码或者导入第三方（不属于 ESP-IDF）的组件。或者，您也可以在顶层 CMakeLists.txt 中设置 `EXTRA_COMPONENT_DIRS` 变量以查找其他指定位置处的组件。有关详细信息，请参阅*重命名 main 组件*。如果项目中源文件较多，建议将其归于组件中，而不是全部放在 “main” 中。
- “main” 目录是一个特殊的组件，它包含项目本身的源代码。” main” 是默认名称，CMake 变量 `COMPONENT_DIRS` 默认包含此组件，但您可以修改此变量。
- “build” 目录是存放构建输出的地方，如果没有此目录，`idf.py` 会自动创建。CMake 会配置项目，并在此目录下生成临时的构建文件。随后，在主构建进程的运行期间，该目录还会保存临时目标文件、库文件以及最终输出的二进制文件。此目录通常不会添加到项目的源码管理系统中，也不会随项目源码一同发布。

每个组件目录都包含一个 `CMakeLists.txt` 文件，里面会定义一些变量以控制该组件的构建过程，以及其与整个项目的集成。更多详细信息请参阅*组件 CMakeLists 文件*。

每个组件还可以包含一个 `Kconfig` 文件，它用于定义 `menuconfig` 时展示的*组件配置* 选项。某些组件可能还会包含 `Kconfig.projbuild` 和 `project_include.cmake` 特殊文件，它们用于*覆盖项目的部分设置*。

4.2.4 项目 CMakeLists 文件

每个项目都有一个顶层 `CMakeLists.txt` 文件，包含整个项目的构建设置。默认情况下，项目 CMakeLists 文件会非常小。

最小 CMakeLists 文件示例

最小项目:

```
cmake_minimum_required(VERSION 3.5)
include($ENV{IDF_PATH}/tools/cmake/project.cmake)
project(myProject)
```

必要部分

每个项目都要按照上面显示的顺序添加上述三行代码：

- `cmake_minimum_required(VERSION 3.5)` 必须放在 CMakeLists.txt 文件的第一行，它会告诉 CMake 构建该项目所需要的最小版本号。ESP-IDF 支持 CMake 3.5 或更高的版本。
- `include($ENV{IDF_PATH}/tools/cmake/project.cmake)` 会导入 CMake 的其余功能来完成配置项目、检索组件等任务。
- `project(myProject)` 会创建项目本身，并指定项目名称。该名称会作为最终输出的二进制文件的名字，即 `myProject.elf` 和 `myProject.bin`。每个 CMakeLists 文件只能定义一个项目。

可选的项目变量

以下这些变量都有默认值，用户可以覆盖这些变量值以自定义构建行为。更多实现细节，请参阅 /tools/cmake/project.cmake 文件。

- `COMPONENT_DIRS`、`COMPONENTS_DIRS`：组件的搜索目录，默认为 `IDF_PATH/components`、`PROJECT_DIR/components`、和 `EXTRA_COMPONENT_DIRS`。如果您不想在这些位置搜索组件，请覆盖此变量。
- `EXTRA_COMPONENT_DIRS`、`EXTRA_COMPONENTS_DIRS`：用于搜索组件的其它可选目录列表。路径可以是相对于项目目录的相对路径，也可以是绝对路径。
- `COMPONENTS`：要构建进项目中的组件名称列表，默认为 `COMPONENT_DIRS` 目录下检索到的所有组件。使用此变量可以“精简”项目以缩短构建时间。请注意，如果一个组件通过 `COMPONENT_REQUIRES` 指定了它依赖的另一个组件，则会自动将其添加到 `COMPONENTS` 中，所以 `COMPONENTS` 列表可能会非常短。

以上变量中的路径可以是绝对路径，或者是相对于项目目录的相对路径。

请使用 cmake 中的 set 命令 来设置这些变量，如 `set(VARIABLE "VALUE")`。请注意，`set()` 命令需放在 `include(...)` 之前，`cmake_minimum(...)` 之后。

重命名 `main` 组件

构建系统会对 `main` 组件进行特殊处理。假如 `main` 组件位于预期的位置（即 *${PROJECT_PATH}/main*），那么它会被自动添加到构建系统中。其他组件也会作为其依赖项被添加到构建系统中，这使用户免于处理依赖关系，并提供即时可用的构建功能。重命名 `main` 组件会减轻上述这些幕后工作量，但要求用户指定重命名后的组件位置，并手动为其添加依赖项。重命名 `main` 组件的步骤如下：

1. 重命名 `main` 目录。
2. 在项目 CMakeLists.txt 文件中设置 `EXTRA_COMPONENT_DIRS`，并添加重命名后的 `main` 目录。
3. 在组件的 CMakeLists.txt 文件中设置 `COMPONENT_REQUIRES` 或 `COMPONENT_PRIV_REQUIRES` 以指定依赖项。

4.2.5 组件 CMakeLists 文件

每个项目都包含一个或多个组件，这些组件可以是 ESP-IDF 的一部分，可以是项目自身组件目录的一部分，也可以从自定义组件目录添加（见上文）。

组件是 `COMPONENT_DIRS` 列表中包含 `CMakeLists.txt` 文件的任何目录。

搜索组件

搜索 `COMPONENT_DIRS` 中的目录列表以查找项目的组件，此列表中的目录可以是组件自身（即包含 *CMakeLists.txt* 文件的目录），也可以是子目录为组件的顶级目录。

当 CMake 运行项目配置时，它会记录本次构建包含的组件列表，它可用于调试某些组件的添加/排除。

同名组件

ESP-IDF 在搜索所有待构建的组件时，会按照 COMPONENT_DIRS 指定的顺序依次进行，这意味着在默认情况下，首先搜索 ESP-IDF 内部组件，然后是项目组件，最后是 EXTRA_COMPONENT_DIRS 中的组件。如果这些目录中的两个或者多个包含具有相同名字的组件，则使用搜索到的最后一个位置的组件。这就允许将组件复制到项目目录中再修改以覆盖 ESP-IDF 组件，如果使用这种方式，ESP-IDF 目录本身可以保持不变。

最小组件 CMakeLists 文件

最小组件 CMakeLists.txt 文件通过使用 idf_component_register 将组件添加到构建系统中。

idf_component_register(SRCS "foo.c" "bar.c" INCLUDE_DIRS "include" REQUIRES mbedtls)

- SRCS 是源文件列表（*.c、*.cpp、*.cc、*.S），里面所有的源文件都将会编译进组件库中。
- INCLUDE_DIRS 是目录列表，里面的路径会被添加到所有需要该组件的组件（包括 main 组件）全局 include 搜索路径中。
- REQUIRES 实际上并不是必需的，但通常需要它来声明该组件需要使用哪些其它组件，请参考组件依赖。

上述命令会构建生成与组件同名的库，并最终被链接到应用程序中。

上述目录通常设置为相对于 CMakeLists.txt 文件的相对路径，当然也可以设置为绝对路径。

还有其它参数可以传递给 idf_component_register，具体可参考*here*。

有关更完整的 CMakeLists.txt 示例，请参阅组件依赖示例 和组件 *CMakeLists* 示例。

预设的组件变量

以下专用于组件的变量可以在组件 CMakeLists 中使用，但不建议修改：

- COMPONENT_DIR：组件目录，即包含 CMakeLists.txt 文件的绝对路径，它与 CMAKE_CURRENT_SOURCE_DIR 变量一样，路径中不能包含空格。
- COMPONENT_NAME：组件名，与组件目录名相同。
- COMPONENT_ALIAS：库别名，由构建系统在内部为组件创建。
- COMPONENT_LIB：库名，由构建系统在内部为组件创建。

以下变量在项目级别中被设置，但可在组件 CMakeLists 中使用：

- CONFIG_*：项目配置中的每个值在 cmake 中都对应一个以 CONFIG_ 开头的变量。更多详细信息请参阅*Kconfig*。
- ESP_PLATFORM：ESP-IDF 构建系统处理 CMake 文件时，其值设为 1。

构建/项目变量

以下是可作为构建属性的构建/项目变量，可通过组件 CMakeLists.txt 中的 idf_build_get_property 查询其变量值。

- PROJECT_NAME：项目名，在项目 CMakeLists.txt 文件中设置。
- PROJECT_DIR：项目目录（包含项目 CMakeLists 文件）的绝对路径，与 CMAKE_SOURCE_DIR 变量相同。
- COMPONENTS：此次构建中包含的所有组件的名称，具体格式为用分号隔开的 CMake 列表。
- IDF_VER：ESP-IDF 的 git 版本号，由 git describe 命令生成。
- IDF_VERSION_MAJOR, IDF_VERSION_MINOR, IDF_VERSION_PATCH: ESP-IDF 的组件版本，可用于条件表达式。请注意这些信息的精确度不如 IDF_VER 变量，版本号 v4.0-dev-*，v4.0-beta1，v4.0-rc1 和 v4.0 对应的 IDF_VERSION_* 变量值是相同的，但是 IDF_VER 的值是不同的。
- IDF_TARGET：项目的硬件目标名称。
- PROJECT_VER：项目版本号。

  - 如果设置*CONFIG_APP_PROJECT_VER_FROM_CONFIG* 选项，将会使用*CONFIG_APP_PROJECT_VER* 的值。
  - 或者，如果在项目 CMakeLists.txt 文件中设置了 PROJECT_VER 变量，则该变量值可以使用。
  - 或者，如果 PROJECT_DIR/version.txt 文件存在，其内容会用作 PROJECT_VER 的值。
  - 或者，如果项目位于某个 Git 仓库中，则使用 git describe 命令的输出作为 PROJECT_VER 的值。
  - 否则，PROJECT_VER 的值为 1。

其它与构建属性有关的信息请参考这里。

组件编译控制

在编译特定组件的源文件时，可以使用 target_compile_options 命令来传递编译器选项:

```
target_compile_options(${COMPONENT_LIB} PRIVATE -Wno-unused-variable)
```

如果给单个源文件指定编译器标志，可以使用 CMake 的 set_source_files_properties 命令:

```
set_source_files_properties(mysrc.c
    PROPERTIES COMPILE_FLAGS
    -Wno-unused-variable
)
```

如果上游代码在编译的时候发出了警告，那这么做可能会很有效。

请注意，上述两条命令只能在组件 CMakeLists 文件的 idf_component_register 命令之后调用。

4.2.6 组件配置

每个组件都可以包含一个 Kconfig 文件，和 CMakeLists.txt 放在同一目录下。Kconfig 文件中包含要添加到该组件配置菜单中的一些配置设置信息。

运行 menuconfig 时，可以在 Component Settings 菜单栏下找到这些设置。

创建一个组件的 Kconfig 文件，最简单的方法就是使用 ESP-IDF 中现有的 Kconfig 文件作为模板，在这基础上进行修改。

有关示例请参阅添加条件配置。

4.2.7 预处理器定义

ESP-IDF 构建系统会在命令行中添加以下 C 预处理器定义：

- ESP_PLATFORM：可以用来检测在 ESP-IDF 内发生了构建行为。
- IDF_VER：定义 git 版本字符串，例如：v2.0 用于标记已发布的版本，v1.0-275-g0efaa4f 则用于标记任意某次的提交记录。

4.2.8 组件依赖

编译各个组件时，ESP-IDF 系统会递归评估其依赖项。这意味着每个组件都需要声明它所依赖的组件，即“requires”。

编写组件

```
idf_component_register(...
                       REQUIRES mbedtls
                       PRIV_REQUIRES console spiffs)
```

- REQUIRES 需要包含所有在当前组件的 公共头文件里 *#include* 的头文件所在的组件。
- PRIV_REQUIRES 需要包含被当前组件的源文件 *#include* 的头文件所在的组件（除非已经被设置在了 REQUIRES 中）。以及是当前组件正常工作必须要链接的组件。
- REQUIRES 和 PRIV_REQUIRES 的值不能依赖于任何配置选项 (CONFIG_xxx 宏)。这是因为在配置加载之前，依赖关系就已经被展开。其它组件变量（比如包含路径或源文件）可以依赖配置选择。
- 如果当前组件除了通用组件依赖项 中设置的通用组件（比如 RTOS、libc 等）外，并不依赖其它组件，那么对于上述两个 REQUIRES 变量，可以选择其中一个或是两个都不设置。

如果组件仅支持某些硬件目标（IDF_TARGET 的值），则可以在 idf_component_register 中指定 REQUIRED_IDF_TARGETS 来声明这个需求。在这种情况下，如果构建系统导入了不支持当前硬件目标的组件时就会报错。

注解： 在 CMake 中，REQUIRES 和 PRIV_REQUIRES 是 CMake 函数 target_link_libraries(... PUBLIC ...) 和 target_link_libraries(... PRIVATE ...) 的近似包装。

组件依赖示例

假设现在有一个 car 组件，它需要使用 engine 组件，而 engine 组件需要使用 spark_plug 组件：

```
- autoProject/
              - CMakeLists.txt
              - components/ - car/ - CMakeLists.txt
                                   - car.c
                                   - car.h
                            - engine/ - CMakeLists.txt
                                      - engine.c
                                      - include/ - engine.h
                            - spark_plug/  - CMakeLists.txt
                                           - plug.c
                                           - plug.h
```

Car 组件 car.h 头文件是 car 组件的公共接口。该头文件直接包含了 engine.h，这是因为它需要使用 engine.h 中的一些声明：

```
/* car.h */
#include "engine.h"

#ifdef ENGINE_IS_HYBRID
#define CAR_MODEL "Hybrid"
#endif
```

同时 car.c 也包含了 car.h:

```
/* car.c */
#include "car.h"
```

这代表文件 car/CMakeLists.txt 需要声明 car 需要 engine：

```
idf_component_register(SRCS "car.c"
                    INCLUDE_DIRS "."
                    REQUIRES engine)
```

- SRCS 提供 car 组件中源文件列表。
- INCLUDE_DIRS 提供该组件公共头文件目录列表，由于 car.h 是公共接口，所以这里列出了所有包含了 car.h 的目录。
- REQUIRES 给出该组件的公共接口所需的组件列表。由于 car.h 是一个公共头文件并且包含了来自 engine 的头文件，所以我们这里包含 engine。这样可以确保任何包含 car.h 的其他组件也能递归地包含所需的 engine.h。

Engine 组件 engine 组件也有一个公共头文件 include/engine.h，但这个头文件更为简单:

```
/* engine.h */
#define ENGINE_IS_HYBRID

void engine_start(void);
```

在 engine.c 中执行:

```
/* engine.c */
#include "engine.h"
#include "spark_plug.h"

...
```

在该组件中，engine 依赖于 spark_plug，但这是私有依赖关系。编译 engine.c 需要 spark_plug.h 但不需要包含 engine.h。

这代表文件 engine/CMakeLists.txt 可以使用 PRIV_REQUIRES:

```
idf_component_register(SRCS "engine.c"
                       INCLUDE_DIRS "include"
                       PRIV_REQUIRES spark_plug)
```

因此，car 组件中的源文件不需要在编译器搜索路径中添加 spark_plug include 目录。这可以加快编译速度，避免编译器命令行过于的冗长。

Spark Plug 组件 spark_plug 组件没有依赖项，它有一个公共头文件 spark_plug.h，但不包含其他组件的头文件。

这代表 spark_plug/CMakeLists.txt 文件不需要任何 REQUIRES 或 PRIV_REQUIRES:

```
idf_component_register(SRCS "spark_plug.c"
                       INCLUDE_DIRS ".")
```

源文件 Include 目录

每个组件的源文件都是用这些 Include 路径目录编译的，这些路径在传递给 idf_component_register 的参数中指定:

```
idf_component_register(..
                       INCLUDE_DIRS "include"
                       PRIV_INCLUDE_DIRS "other")
```

- 当前组件的 INCLUDE_DIRS 和 PRIV_INCLUDE_DIRS。
- REQUIRES 和 PRIV_REQUIRES 参数指定的所有其他组件（即当前组件的所有公共和私有依赖项）所设置的 INCLUDE_DIRS。
- 递归列出所有组件 REQUIRES 列表中 INCLUDE_DIRS 目录（如递归展开这个组件的所有公共依赖项）。

主要组件依赖项

main 组件比较特别，因为它在构建过程中自动依赖所有其他组件。所以不需要向这个组件传递 REQUIRES 或 PRIV_REQUIRES。有关不再使用 main 组件时需要更改哪些内容，请参考重命名 *main* 组件。

通用组件依赖项

为避免重复性工作，各组件都用自动依赖一些“通用”IDF 组件，即使它们没有被明确提及。这些组件的头文件会一直包含在构建系统中。

通用组件包括：freertos、newlib、heap、log、soc、esp_rom、esp_common、xtensa/riscv、cxx。

在构建中导入组件

- 默认情况下，每个组件都会包含在构建系统中。
- 如果将 COMPONENTS 变量设置为项目直接使用的最小组件列表，那么构建系统会扩展到包含所有组件。完整的组件列表为：
  - COMPONENTS 中明确提及的组件。
  - 这些组件的依赖项（以及递归运算后的组件）。
  - 每个组件都依赖的通用组件。
- 将 COMPONENTS 设置为所需组件的最小列表，可以显著减少项目的构建时间。

构建系统中依赖处理的实现细节

- 在 CMake 配置进程的早期阶段会运行 expand_requirements.cmake 脚本。该脚本会对所有组件的 CMakeLists.txt 文件进行局部的运算，得到一张组件依赖关系图（此图可能会有闭环）。此图用于在构建目录中生成 component_depends.cmake 文件。
- CMake 主进程会导入该文件，并以此来确定要包含到构建系统中的组件列表（内部使用的 BUILD_COMPONENTS 变量）。BUILD_COMPONENTS 变量已排好序，依赖组件会排在前面。由于组件依赖关系图中可能存在闭环，因此不能保证每个组件都满足该排序规则。如果给定相同的组件集和依赖关系，那么最终的排序结果应该是确定的。
- CMake 会将 BUILD_COMPONENTS 的值以“Component names:”的形式打印出来。
- 然后执行构建系统中包含的每个组件的配置。
- 每个组件都被正常包含在构建系统中，然后再次执行 CMakeLists.txt 文件，将组件库加入构建系统。

组件依赖顺序 BUILD_COMPONENTS 变量中组件的顺序决定了构建过程中的其它顺序，包括：

- 项目导入*project_include.cmake* 文件的顺序。
- 生成用于编译（通过 -I 参数）的头文件路径列表的顺序。请注意，对于给定组件的源文件，仅需将该组件的依赖组件的头文件路径告知编译器。

覆盖项目的部分设置

project_include.cmake 如果组件的某些构建行为需要在组件 CMakeLists 文件之前被执行，您可以在组件目录下创建名为 project_include.cmake 的文件，project.cmake 在运行过程中会导入此 CMake 文件。

project_include.cmake 文件在 ESP-IDF 内部使用，以定义项目范围内的构建功能，比如 esptool.py 的命令行参数和 bootloader 这个特殊的应用程序。

与组件 CMakeLists.txt 文件有所不同，在导入 “project_include.cmake“ 文件的时候，当前源文件目录（即 CMAKE_CURRENT_SOURCE_DIR``和工作目录）为项目目录。如果想获得当前组件的绝对路径，可以使用 ``COMPONENT_PATH 变量。

请注意，project_include.cmake 对于大多数常见的组件并不是必需的。例如给项目添加 include 搜索目录，给最终的链接步骤添加 LDFLAGS 选项等等都可以通过 CMakeLists.txt 文件来自定义。详细信息请参考可选的项目变量。

project_include.cmake 文件会按照 BUILD_COMPONENTS 变量中组件的顺序（由 CMake 记录）依次导入。即只有在当前组件所有依赖组件的 project_include.cmake 文件都被导入后，当前组件的 project_include.cmake 文件才会被导入，除非两个组件在同一个依赖闭环中。如果某个 project_include.cmake 文件依赖于另一组件设置的变量，则要特别注意上述情况。更多详情请参阅构建系统中依赖处理的实现细节。

在 project_include.cmake 文件中设置变量或目标时要格外小心，这些值被包含在项目的顶层 CMake 文件中，因此他们会影响或破坏所有组件的功能。

KConfig.projbuild 与 project_include.cmake 类似，也可以为组件定义一个 KConfig 文件以实现全局的组件配置。如果要在 menuconfig 的顶层添加配置选项，而不是在 “Component Configuration” 子菜单中，则可以在 CMakeLists.txt 文件所在目录的 KConfig.projbuild 文件中定义这些选项。

在此文件中添加配置时要小心，因为这些配置会包含在整个项目配置中。在可能的情况下，请为组件配置 创建 KConfig 文件。

project_include.cmake 文件在 ESP-IDF 内部使用，以定义项目范围内的构建功能，比如 esptool.py 的命令行参数和 bootloader 这个特殊的应用程序。

仅配置组件 仅配置组件是一类不包含源文件的特殊组件，仅包含 Kconfig.projbuild、KConfig 和 CMakeLists.txt 文件，该 CMakeLists.txt 文件仅有一行代码，调用了 idf_component_register() 函数。此函数会将组件导入到项目构建中，但不会构建任何库，也不会将头文件添加到任何 include 搜索路径中。

CMake 调试

请查看 CMake v3.5 官方文档 获取更多关于 CMake 和 CMake 命令的信息。

调试 ESP-IDF CMake 构建系统的一些技巧：

- CMake 运行时，会打印大量诊断信息，包括组件列表和组件路径。
- 运行 cmake -DDEBUG=1，IDF 构建系统会生成更详细的诊断输出。
- 运行 cmake 时指定 --trace 或 --trace-expand 选项会提供大量有关控制流信息。详情请参考 CMake 命令行文档。

当从项目 CMakeLists 文件导入时，project.cmake 文件会定义工具模块和全局变量，并在系统环境中没有设置 IDF_PATH 时设置 IDF_PATH。

同时还定义了一个自定义版本的内置 CMake project 函数，这个函数被覆盖，以添加所有 ESP-IDF 特定的项目功能。

警告未定义的变量 默认情况下，idf.py 在调用 CMake 时会给它传递 --warn-uninitialized 标志，如果在构建的过程中引用了未定义的变量，CMake 会打印警告。这对查找有错误的 CMake 文件非常有用。

如果您不想启用此功能，可以给 idf.py 传递 --no-warnings 标志。

更多信息，请参考文件 /tools/cmake/project.cmake 以及 /tools/cmake/ 中支持的函数。

4.2.9 组件 CMakeLists 示例

因为构建环境试图设置大多数情况都能工作的合理默认值，所以组件 CMakeLists.txt 文件可能非常小，甚至是空的，请参考最小组件 *CMakeLists* 文件。但有些功能往往需要覆盖预设的组件变量 才能实现。

以下是组件 CMakeLists 文件的更高级的示例。

添加条件配置

配置系统可用于根据项目配置中选择的选项有条件地编译某些文件。

Kconfig:

```
config FOO_ENABLE_BAR
    bool "Enable the BAR feature."
    help
        This enables the BAR feature of the FOO component.
```

CMakeLists.txt:

```
 set(srcs "foo.c" "more_foo.c")

 if(CONFIG_FOO_ENABLE_BAR)
     list(APPEND srcs "bar.c")
 endif()

idf_component_register(SRCS "${srcs}"
                       ...)
```

上述示例使用了 CMake 的 if 函数和 list APPEND 函数。

也可用于选择或删除某一实现，如下所示：

Kconfig:

```
config ENABLE_LCD_OUTPUT
    bool "Enable LCD output."
    help
        Select this if your board has a LCD.

config ENABLE_LCD_CONSOLE
    bool "Output console text to LCD"
    depends on ENABLE_LCD_OUTPUT
    help
        Select this to output debugging output to the lcd

config ENABLE_LCD_PLOT
    bool "Output temperature plots to LCD"
    depends on ENABLE_LCD_OUTPUT
    help
        Select this to output temperature plots
```

CMakeLists.txt:

```
if(CONFIG_ENABLE_LCD_OUTPUT)
   set(srcs lcd-real.c lcd-spi.c)
else()
   set(srcs lcd-dummy.c)
endif()

# 如果启用了控制台或绘图功能，则需要加入字体
if(CONFIG_ENABLE_LCD_CONSOLE OR CONFIG_ENABLE_LCD_PLOT)
   list(APPEND srcs "font.c")
endif()

idf_component_register(SRCS "${srcs}"
                       ...)
```

硬件目标的条件判断

CMake 文件可以使用 IDF_TARGET 变量来获取当前的硬件目标。

此外，如果当前的硬件目标是 xyz （即 IDF_TARGET=xyz），那么 Kconfig 变量 CONFIG_IDF_TARGET_XYZ 同样也会被设置。

请注意，组件可以依赖 IDF_TARGET 变量，但不能依赖这个 Kconfig 变量。同样也不可在 CMake 文件的 include 语句中使用 Kconfig 变量，在这种上下文中可以使用 IDF_TARGET。

生成源代码

有些组件的源文件可能并不是由组件本身提供，而必须从另外的文件生成。假设组件需要一个头文件，该文件由 BMP 文件转换后（使用 bmp2h 工具）的二进制数据组成，然后将头文件包含在名为 graphics_lib.c 的文件中:

```
add_custom_command(OUTPUT logo.h
     COMMAND bmp2h -i ${COMPONENT_DIR}/logo.bmp -o log.h
     DEPENDS ${COMPONENT_DIR}/logo.bmp
     VERBATIM)

add_custom_target(logo DEPENDS logo.h)
add_dependencies(${COMPONENT_LIB} logo)

set_property(DIRECTORY "${COMPONENT_DIR}" APPEND PROPERTY
     ADDITIONAL_MAKE_CLEAN_FILES logo.h)
```

这个示例改编自 CMake 的一则 FAQ，其中还包含了一些同样适用于 ESP-IDF 构建系统的示例。

这个示例会在当前目录（构建目录）中生成 logo.h 文件，而 logo.bmp 会随组件一起提供在组件目录中。因为 logo.h 是一个新生成的文件，一旦项目需要清理，该文件也应该要被清除。因此，要将该文件添加到 ADDITIONAL_MAKE_CLEAN_FILES 属性中。

注解: 如果需要生成文件作为项目 CMakeLists.txt 的一部分，而不是作为组件 CMakeLists.txt 的一部分，此时需要使用 ${PROJECT_PATH} 替代 ${COMPONENT_DIR}，使用 ${PROJECT_NAME}.elf 替代 ${COMPONENT_LIB}。

如果某个源文件是从其他组件中生成，且包含 logo.h 文件，则需要调用 add_dependencies，在这两个组件之间添加一个依赖项，以确保组件源文件按照正确顺序进行编译。

嵌入二进制数据

有时您的组件希望使用一个二进制文件或者文本文件，但是您又不希望将它们重新格式化为 C 源文件。

这时，您可以在组件注册中指定 EMBED_FILES 参数，用空格分隔要嵌入的文件名称:

```
idf_component_register(...
                       EMBED_FILES server_root_cert.der)
```

或者，如果文件是字符串，则可以使用 EMBED_TXTFILES 变量，把文件的内容转成以 null 结尾的字符串嵌入:

```
idf_component_register(...
                       EMBED_TXTFILES server_root_cert.pem)
```

文件的内容会被添加到 Flash 的.rodata 段，用户可以通过符号名来访问，如下所示:

```
extern const uint8_t server_root_cert_pem_start[] asm("_binary_server_root_cert_
↪pem_start");
extern const uint8_t server_root_cert_pem_end[]   asm("_binary_server_root_cert_
↪pem_end");
```

符号名会根据文件全名生成，如 EMBED_FILES 中所示，字符 /、. 等都会被下划线替代。符号名称中的 _binary 前缀由 objcopy 命令添加，对文本文件和二进制文件都是如此。

如果要将文件嵌入到项目中，而非组件中，可以调用 target_add_binary_data 函数:

```
target_add_binary_data(myproject.elf "main/data.bin" TEXT)
```

并将这行代码放在项目 CMakeLists.txt 的 `project()` 命令之后，修改 `myproject.elf` 为你自己的项目名。如果最后一个参数是 `TEXT`，那么构建系统会嵌入以 null 结尾的字符串，如果最后一个参数被设置为 `BINARY`，则将文件内容按照原样嵌入。

有关使用此技术的示例，请查看 file_serving 示例 protocols/http_server/file_serving/main/CMakeLists.txt 中的 main 组件，两个文件会在编译时加载并链接到固件中。

代码和数据的存放

ESP-IDF 还支持自动生成链接脚本，它允许组件通过链接片段文件定义其代码和数据在内存中的存放位置。构建系统会处理这些链接片段文件，并将处理后的结果扩充进链接脚本，从而指导应用程序二进制文件的链接过程。更多详细信息与快速上手指南，请参阅链接脚本生成机制。

完全覆盖组件的构建过程

当然，在有些情况下，上面提到的方法不一定够用。如果组件封装了另一个第三方组件，而这个第三方组件并不能直接在 ESP-IDF 的构建系统中工作，在这种情况下，就需要放弃 ESP-IDF 的构建系统，改为使用 CMake 的 ExternalProject 功能。组件 CMakeLists 示例如下:

```
# 用于 quirc 的外部构建过程，在源目录中运行
# 并生成 libquirc.a
externalproject_add(quirc_build
    PREFIX ${COMPONENT_DIR}
    SOURCE_DIR ${COMPONENT_DIR}/quirc
    CONFIGURE_COMMAND ""
    BUILD_IN_SOURCE 1
    BUILD_COMMAND make CC=${CMAKE_C_COMPILER} libquirc.a
    INSTALL_COMMAND ""
    )

# 将 libquirc.a 添加到构建系统中
add_library(quirc STATIC IMPORTED GLOBAL)
add_dependencies(quirc quirc_build)

set_target_properties(quirc PROPERTIES IMPORTED_LOCATION
    ${COMPONENT_DIR}/quirc/libquirc.a)
set_target_properties(quirc PROPERTIES INTERFACE_INCLUDE_DIRECTORIES
    ${COMPONENT_DIR}/quirc/lib)

set_directory_properties( PROPERTIES ADDITIONAL_MAKE_CLEAN_FILES
    "${COMPONENT_DIR}/quirc/libquirc.a")
```

（上述 CMakeLists.txt 可用于创建名为 `quirc` 的组件，该组件使用自己的 Makefile 构建 quirc 项目。）

- `externalproject_add` 定义了一个外部构建系统。
  - 设置 `SOURCE_DIR`、`CONFIGURE_COMMAND`、`BUILD_COMMAND` 和 `INSTALL_COMMAND`。如果外部构建系统没有配置这一步骤，可以将 `CONFIGURE_COMMAND` 设置为空字符串。在 ESP-IDF 的构建系统中，一般会将 `INSTALL_COMMAND` 变量设置为空。
  - 设置 `BUILD_IN_SOURCE`，即构建目录与源目录相同。否则，您也可以设置 `BUILD_DIR` 变量。
  - 有关 `externalproject_add()` 命令的详细信息，请参阅 ExternalProject。
- 第二组命令添加了一个目标库，指向外部构建系统生成的库文件。为了添加 include 目录，并告知 CMake 该文件的位置，需要再设置一些属性。
- 最后，生成的库被添加到 ADDITIONAL_MAKE_CLEAN_FILES 中。即执行 `make clean` 后会删除该库。请注意，构建系统中的其他目标文件不会被删除。

ExternalProject 的依赖与构建清理 对于外部项目的构建，CMake 会有一些不同寻常的行为：

- ADDITIONAL_MAKE_CLEAN_FILES 仅在使用 Make 构建系统时有效。如果使用 Ninja 或 IDE 自带的构建系统，执行项目清理时，这些文件不会被删除。
- ExternalProject 会在 clean 运行后自动重新运行配置和构建命令。
- 可以采用以下两种方法来配置外部构建命令：
  1. 将外部 `BUILD_COMMAND` 命令设置为对所有源代码完整的重新编译。如果传递给 `externalproject_add` 命令的 `DEPENDS` 的依赖项发生了改变，或者当前执行的是项目清理操作（即运行了 `idf.py clean`、`ninja clean` 或者 `make clean`），那么就会执行该命令。
  2. 将外部 `BUILD_COMMAND` 命令设置为增量式构建命令，并给 `externalproject_add` 传递 `BUILD_ALWAYS 1` 参数。即不管实际的依赖情况，每次构建时，都会构建外部项目。这种方式仅当外部构建系统具备增量式构建的能力，且运行时间不会很长时才推荐。

构建外部项目的最佳方法取决于项目本身、其构建系统，以及是否需要频繁重新编译项目。

4.2.10 自定义 sdkconfig 的默认值

对于示例工程或者其他您不想指定完整 sdkconfig 配置的项目，但是您确实希望覆盖 ESP-IDF 默认值中的某些键值，则可以在项目中创建 `sdkconfig.defaults` 文件。重新创建新配置时将会用到此文件，另外在 `sdkconfig` 没有设置新配置值时，上述文件也会被用到。

如若需要覆盖此文件的名称或指定多个文件，请设置 `SDKCONFIG_DEFAULTS` 环境变量或在顶层 CMakeLists.txt 文件中设置 `SDKCONFIG_DEFAULTS`。在指定多个文件时，使用分号作为分隔符。未指定完整路径的文件名将以当前项目的相对路径来解析。

依赖于硬件目标的 sdkconfig 默认值

除了 `sdkconfig.defaults` 之外，构建系统还将从 `sdkconfig.defaults.TARGET_NAME` 文件加载默认值，其中 `IDF_TARGET` 的值为 `TARGET_NAME`。例如，对于 `ESP32` 这个硬件目标，sdkconfig 的默认值会首先从 `sdkconfig.defaults` 获取，然后再从 `sdkconfig.defaults.esp32` 获取。

如果使用 `SDKCONFIG_DEFAULTS` 覆盖了 sdkconfig 默认文件的名称，则硬件目标的 sdkconfig 默认文件名也会从 `SDKCONFIG_DEFAULTS` 值中派生。

4.2.11 Flash 参数

有些情况下，我们希望在没有 IDF 时也能烧写目标板，为此，我们希望可以保存已构建的二进制文件、esptool.py 和 esptool write_flash 命令的参数。可以通过编写一段简单的脚本来保存二进制文件和 esptool.py。

运行项目构建之后，构建目录将包含项目二进制输出文件（`.bin` 文件），同时也包含以下烧录数据文件：

- `flash_project_args` 包含烧录整个项目的参数，包括应用程序 (app)、引导程序 (bootloader)、分区表，如果设置了 PHY 数据，也会包含此数据。
- `flash_app_args` 只包含烧录应用程序的参数。
- `flash_bootloader_args` 只包含烧录引导程序的参数。

您可以参照如下命令将任意烧录参数文件传递给 `esptool.py`:

```
python esptool.py --chip esp32s2 write_flash @build/flash_project_args
```

也可以手动复制参数文件中的数据到命令行中执行。

构建目录中还包含生成的 `flasher_args.json` 文件，此文件包含 JSON 格式的项目烧录信息，可用于 `idf.py` 和其它需要项目构建信息的工具。

4.2.12 构建 Bootloader

引导程序默认作为 `idf.py build` 的一部分被构建，也可以通过 `idf.py bootloader` 来单独构建。

引导程序是 /components/bootloader/subproject 内部独特的“子项目”，它有自己的项目 CMakeLists.txt 文件，能够构建独立于主项目的 .ELF 和 .BIN 文件，同时它又与主项目共享配置和构建目录。

子项目通过 /components/bootloader/project_include.cmake 文件作为外部项目插入到项目的顶层，主构建进程会运行子项目的 CMake，包括查找组件（主项目使用的组件的子集），生成引导程序专用的配置文件（从主 sdkconfig 文件中派生）。

4.2.13 选择目标芯片

ESP-IDF 支持多款芯片，它们通过在软件中使用不同的“目标”(target) 名进行区分，具体对应关系如下：

- esp32 —适用于 ESP32-D0WD、ESP32-D2WD、ESP32-S0WD (ESP-SOLO)、ESP32-U4WDH、ESP32-PICO-D4
- esp32s2—适用于 ESP32-S2

在构建项目前，请首先根据您的芯片选择正确的软件目标，具体命令为 idf.py set-target <target>。举例

```
idf.py set-target esp32s2
```

重要： 运行 idf.py set-target 命令将清除 build 文件夹的内容，并重新生成一个 sdkconfig 文件。之前的 sdkconfig 将另存为 sdkconfig.old。

注解： 运行 idf.py set-target 命令相当于分别运行以下几个命令：

1. 清除 build 文件夹 (idf.py fullclean)
2. 移除 sdkconfig 文件 (mv sdkconfig sdkconfig.old)
3. 根据选择的“目标”芯片配置项目 (idf.py -DIDF_TARGET=esp32 reconfigure)

您也可以将要用的 IDF_TARGET 设置为环境变量，比如：export IDF_TARGET=esp32s2；或设置为 CMake 变量，比如将 -DIDF_TARGET=esp32s2 以参数形式传递给 CMake 或 idf.py。如果您大多数时间仅使用一款芯片，则将 IDF_TARGET 配置为环境变量比较方便。

对于特定项目，您可以使用以下方式为 IDF_TARGET 配置 _default_ 值：把 CONFIG_IDF_TARGET 的值加入 sdkconfig.defaults。举例而言，配置 CONFIG_IDF_TARGET="esp32s2"。这样一来，除非特别设置（比如使用环境变量、CMake 变量或 idf.py set-target 命令），否则 IDF_TARGET 将默认采用 CONFIG_IDF_TARGET。

如果您从未通过以上述任何方式配置过“目标”芯片，则构建系统会默认将 esp32 设定为“目标”芯片。

4.2.14 编写纯 CMake 组件

ESP-IDF 构建系统用“组件”的概念“封装”了 CMake，并提供了很多帮助函数来自动将这些组件集成到项目构建当中。

然而，“组件”概念的背后是一个完整的 CMake 构建系统，因此可以制作纯 CMake 组件。

下面是使用纯 CMake 语法为 json 组件编写的最小 CMakeLists 文件的示例：

```
add_library(json STATIC
cJSON/cJSON.c
cJSON/cJSON_Utils.c)

target_include_directories(json PUBLIC cJSON)
```

- 这实际上与 IDF 中的 json 组件 是等效的。
- 因为组件中的源文件不多，所以这个 CMakeLists 文件非常简单。对于具有大量源文件的组件而言，ESP-IDF 支持的组件通配符，可以简化组件 CMakeLists 的样式。

- 每当组件中新增一个与组件同名的库目标时，ESP-IDF 构建系统会自动将其添加到构建中，并公开公共的 include 目录。如果组件想要添加一个与组件不同名的库目标，就需要使用 CMake 命令手动添加依赖关系。

4.2.15 组件中使用第三方 CMake 项目

CMake 在许多开源的 C/C++ 项目中广泛使用，用户可以在自己的应用程序中使用开源代码。CMake 构建系统的一大好处就是可以导入这些第三方的项目，有时候甚至不用做任何改动。这就允许用户使用当前 ESP-IDF 组件尚未提供的功能，或者使用其它库来实现相同的功能。

假设 `main` 组件需要导入一个假想库 `foo`，相应的组件 CMakeLists 文件如下所示:

```
# 注册组件
idf_component_register(...)

# 设置 `foo` 项目中的一些 CMake 变量，以控制 `foo` 的构建过程
set(FOO_BUILD_STATIC OFF)
set(FOO_BUILD_TESTS OFF)

# 创建并导入第三方库目标
add_subdirectory(foo)

# 将 `foo` 目标公开链接至 `main` 组件
target_link_libraries(main PUBLIC foo)
```

实际的案例请参考 build_system/cmake/import_lib。请注意，导入第三方库所需要做的工作可能会因库的不同而有所差异。建议仔细阅读第三方库的文档，了解如何将其导入到其它项目中。阅读第三方库的 CMakeLists.txt 文件以及构建结构也会有所帮助。

用这种方式还可以将第三方库封装成 ESP-IDF 的组件。例如 mbedtls 组件就是封装了 mbedtls 项目 得到的。详情请参考 mbedtls 组件的 CMakeLists.txt 文件 。

每当使用 ESP-IDF 构建系统时，CMake 变量 `ESP_PLATFORM` 都会被设置为 1。如果要在通用的 CMake 代码加入 IDF 特定的代码时，可以采用 `if (ESP_PLATFORM)` 的形式加以分隔。

外部库中使用 ESP-IDF 组件

上述示例中假设的是外部库 `foo` （或 `import_lib` 示例中的 `tinyxml` 库）除了常见的 API 如 libc、libstdc++ 等外不需要使用其它 ESP-IDF API。如果外部库需要使用其它 ESP-IDF 组件提供的 API，则需要在外部 CMakeLists.txt 文件中通过添加对库目标 `idf::<componentname>` 的依赖关系。

例如，在 `foo/CMakeLists.txt` 文件:

```
add_library(foo bar.c fizz.cpp buzz.cpp)

if(ESP_PLATFORM)
  # 在 ESP-IDF 中，bar.c 需要包含 spi_flash 组件中的 esp_spi_flash.h
  target_link_libraries(foo PRIVATE idf::spi_flash)
endif()
```

4.2.16 组件中使用预建库

还有一种情况是您有一个由其它构建过程生成预建静态库（`.a` 文件）。

ESP-IDF 构建系统为用户提供了一个实用函数 `add_prebuilt_library`，能够轻松导入并使用预建库:

```
add_prebuilt_library(target_name lib_path [REQUIRES req1 req2 ...] [PRIV_REQUIRES␣
↪req1 req2 ...])
```

其中:

- target_name- 用于引用导入库的名称，如链接到其它目标时
- lib_path- 预建库的路径，可以是绝对路径或是相对于组件目录的相对路径

可选参数 REQUIRES 和 PRIV_REQUIRES 指定对其它组件的依赖性。这些参数与 idf_component_register 的参数的意义相同。

注意预建库的编译目标需与目前的项目相同。预建库的相关参数也要匹配。如果不特别注意，这两个因素可能会导致应用程序中出现 bug。

请查看示例 build_system/cmake/import_prebuilt。

4.2.17 在自定义 CMake 项目中使用 ESP-IDF

ESP-IDF 提供了一个模板 CMake 项目，可以基于此轻松创建应用程序。然而在有些情况下，用户可能已有一个现成的 CMake 项目，或者想自己创建一个 CMake 项目，此时就希望将 IDF 中的组件以库的形式链接到用户目标（库/可执行文件）。

可以通过 tools/cmake/idf.cmake 提供的*build system APIs* 实现该目标。例如：

```
cmake_minimum_required(VERSION 3.5)
project(my_custom_app C)

# 导入提供 ESP-IDF CMake 构建系统 API 的 CMake 文件
include($ENV{IDF_PATH}/tools/cmake/idf.cmake)

# 在构建中导入 ESP-IDF 组件，可以视作等同 add_subdirectory()
# 但为 ESP-IDF 构建增加额外的构建过程
# 具体构建过程
idf_build_process(esp32)

# 创建项目可执行文件
# 使用其别名 idf::newlib 将其链接到 newlib 组件
add_executable(${CMAKE_PROJECT_NAME}.elf main.c)
target_link_libraries(${CMAKE_PROJECT_NAME}.elf idf::newlib)

# 让构建系统知道项目到可执行文件是什么，从而添加更多的目标以及依赖关系等
idf_build_executable(${CMAKE_PROJECT_NAME}.elf)
```

build_system/cmake/idf_as_lib 中的示例演示了如何在自定义的 CMake 项目创建一个类似于 Hello World 的应用程序。

4.2.18 ESP-IDF CMake 构建系统 API

idf 构建命令

```
idf_build_get_property(var property [GENERATOR_EXPRESSION])
```

检索一个构建属性 *property*，并将其存储在当前作用域可访问的 var 中。特定 *GENERATOR_EXPRESSION* 将检索该属性的生成器表达式字符串（不是实际值），它可与支持生成器表达式的 CMake 命令一起使用。

```
idf_build_set_property(property val [APPEND])
```

设置构建属性 *property* 的值为 *val*。特定 *APPEND* 将把指定的值附加到属性当前值之后。如果该属性之前不存在或当前为空，则指定的值将变为第一个元素/成员。

```
idf_build_component(component_dir)
```

向构建系统提交一个包含组件的 *component_dir* 目录。相对路径会被转换为相对于当前目录的绝对路径。所有对该命令的调用必须在 'idf_build_process'之前执行。

该命令并不保证组件在构建过程中会被处理（参见 *idf_build_process* 中 *COMPONENTS* 参数说明）

```
idf_build_process(target
                  [PROJECT_DIR project_dir]
                  [PROJECT_VER project_ver]
                  [PROJECT_NAME project_name]
                  [SDKCONFIG sdkconfig]
                  [SDKCONFIG_DEFAULTS sdkconfig_defaults]
                  [BUILD_DIR build_dir]
                  [COMPONENTS component1 component2 ...])
```

为导入 ESP-IDF 组件执行大量的幕后工作，包括组件配置、库创建、依赖性扩展和解析。在这些功能中，对于用户最重要的可能是通过调用每个组件的 `idf_component_register` 来创建库。该命令为每个组件创建库，这些库可以使用别名来访问，其形式为 idf::*component_name*。这些别名可以用来将组件链接到用户自己的目标、库或可执行文件上。

该调用要求用 *target* 参数指定目标芯片。调用的可选参数包括：

- PROJECT_DIR - 项目目录，默认为 CMAKE_SOURCE_DIR。
- PROJECT_NAME - 项目名称，默认为 CMAKE_PROJECT_NAME。
- PROJECT_VER - 项目的版本/版本号，默认为“1”。
- SDKCONFIG - 生成的 sdkconfig 文件的输出路径，根据是否设置 PROJECT_DIR，默认为 PROJECT_DIR/sdkconfig 或 CMAKE_SOURCE_DIR/sdkconfig。
- SDKCONFIG_DEFAULTS - 包含默认配置的文件列表（列表中必须包含完整的路径），默认为空；对于列表中的每个值 *filename*，如果存在的话，也会加载文件 *filename.target* 中的配置。对于列表中的 *filename* 的每一个值，也会加载文件 *filename.target* （如果存在的话）中的配置。
- BUILD_DIR - 用于放置 ESP-IDF 构建相关工具的目录，如生成的二进制文件、文本文件、组件；默认为 CMAKE_BINARY_DIR。
- COMPONENTS - 从构建系统已知的组件中选择要处理的组件（通过 `idf_build_component` 添加）。这个参数用于精简构建过程。如果在依赖链中需要其它组件，则会自动添加，即自动添加这个列表中组件的公共和私有依赖项，进而添加这些依赖项的公共和私有依赖，以此类推。如果不指定，则会处理构建系统已知的所有组件。

```
idf_build_executable(executable)
```

指定 ESP-IDF 构建的可执行文件 *executable*。这将添加额外的目标，如与 flash 相关的依赖关系，生成额外的二进制文件等。应在 `idf_build_process` 之后调用。

```
idf_build_get_config(var config [GENERATOR_EXPRESSION])
```

获取指定配置的值。就像构建属性一样，特定 *GENERATOR_EXPRESSION* 将检索该配置的生成器表达式字符串，而不是实际值，即可以与支持生成器表达式的 CMake 命令一起使用。然而，实际的配置值只有在调用 `idf_build_process` 后才能知道。

idf 构建属性

可以通过使用构建命令 `idf_build_get_property` 来获取构建属性的值。例如，以下命令可以获取构建过程中使用的 Python 解释器的相关信息。

```
idf_build_get_property(python PYTHON)
message(STATUS "The Python intepreter is: ${python}")
```

- BUILD_DIR - 构建目录；由 `idf_build_process` 的 BUILD_DIR 参数设置。
- BUILD_COMPONENTS - 包含在构建中的组件列表；由 `idf_build_process` 设置。
- BUILD_COMPONENT_ALIASES - 包含在构建中的组件的库别名列表；由 `idf_build_process` 设置。
- C_COMPILE_OPTIONS - 适用于所有组件的 C 源代码文件的编译选项。
- COMPILE_OPTIONS - 适用于所有组件的源文件（无论是 C 还是 C++）的编译选项。
- COMPILE_DEFINITIONS - 适用于所有组件源文件的编译定义。
- CXX_COMPILE_OPTIONS - 适用于所有组件的 C++ 源文件的编译选项。
- EXECUTABLE - 项目可执行文件；通过调用 `idf_build_executable` 设置。

- EXECUTABLE_NAME - 不含扩展名的项目可执行文件的名称；通过调用 `idf_build_executable` 设置。
- EXECUTABLE_DIR - 输出的可执行文件的路径
- IDF_PATH - ESP-IDF 路径；由 IDF_PATH 环境变量设置，或者从 `idf.cmake` 的位置推断。
- IDF_TARGET - 构建的目标芯片；由 `idf_build_process` 的目标参数设置。
- IDF_VER - ESP-IDF 版本；由版本文件或 IDF_PATH 仓库的 Git 版本设置。
- INCLUDE_DIRECTORIES - 包含所有组件源文件的目录。
- KCONFIGS - 构建过程中组件里的 Kconfig 文件的列表；由 `idf_build_process` 设置。
- KCONFIG_PROJBUILDS - 构建过程中组件中的 Kconfig.projbuild 文件的列表；由 `idf_build_process` 设置。
- PROJECT_NAME - 项目名称；由 `idf_build_process` 的 PROJECT_NAME 参数设置。
- PROJECT_DIR - 项目的目录；由 `idf_build_process` 的 PROJECT_DIR 参数设置。
- PROJECT_VER - 项目的版本；由 `idf_build_process` 的 PROJECT_VER 参数设置。
- PYTHON - 用于构建的 Python 解释器；如果有则从 PYTHON 环境变量中设置，如果没有，则使用“python”。
- SDKCONFIG - 输出的配置文件的完整路径；由 `idf_build_process` SDKCONFIG 参数设置。
- SDKCONFIG_DEFAULTS - 包含默认配置的文件列表；由 `idf_build_process` SDKCONFIG_DEFAULTS 参数设置。
- SDKCONFIG_HEADER - 包含组件配置的 C/C++ 头文件的完整路径；由 `idf_build_process` 设置。
- SDKCONFIG_CMAKE - 包含组件配置的 CMake 文件的完整路径；由 `idf_build_process` 设置。
- SDKCONFIG_JSON - 包含组件配置的 JSON 文件的完整路径；由 `idf_build_process` 设置。
- SDKCONFIG_JSON_MENUS - 包含配置菜单的 JSON 文件的完整路径；由 `idf_build_process` 设置。

idf 组件命令

```
idf_component_get_property(var component property [GENERATOR_EXPRESSION])
```

检索一个指定的 *component* 的组件属性 *property*，并将其存储在当前作用域可访问的 *var* 中。指定 *GENERATOR_EXPRESSION* 将检索该属性的生成器表达式字符串（不是实际值），它可以在支持生成器表达式的 CMake 命令中使用。

```
idf_component_set_property(component property val [APPEND])
```

设置指定的 *component* 的组件属性，*property* 的值为 *val*。特定 *APPEND* 将把指定的值追加到属性的当前值后。如果该属性之前不存在或当前为空，指定的值将成为第一个元素/成员。

```
idf_component_register([[SRCS src1 src2 ...] | [[SRC_DIRS dir1 dir2 ...] [EXCLUDE_
↪SRCS src1 src2 ...]]
                       [INCLUDE_DIRS dir1 dir2 ...]
                       [PRIV_INCLUDE_DIRS dir1 dir2 ...]
                       [REQUIRES component1 component2 ...]
                       [PRIV_REQUIRES component1 component2 ...]
                       [LDFRAGMENTS ldfragment1 ldfragment2 ...]
                       [REQUIRED_IDF_TARGETS target1 target2 ...]
                       [EMBED_FILES file1 file2 ...]
                       [EMBED_TXTFILES file1 file2 ...])
```

将一个组件注册到构建系统中。就像 `project()` CMake 命令一样，该命令应该直接从组件的 CMakeLists.txt 中调用（而不是通过函数或宏），且建议在其他命令之前调用该命令。下面是一些关于在 `idf_component_register` 之前不能调用哪些命令的指南：

- 在 CMake 脚本模式下无效的命令。
- 在 project_include.cmake 中定义的自定义命令。
- 除了 `idf_build_get_property` 之外，构建系统的 API 命令；但要考虑该属性是否有被设置。

对变量进行设置和操作的命令，一般可在 `idf_component_register` 之前调用。

`idf_component_register` 的参数包括：

- SRCS - 组件的源文件，用于为组件创建静态库；如果没有指定，组件将被视为仅配置组件，从而创建接口库。
- SRC_DIRS, EXCLUDE_SRCS - 用于通过指定目录来 glob 源文件 (.c、.cpp、.S)，而不是通过 SRCS 手动指定源文件。请注意，这受*CMake 中通配符的限制*。在 EXCLUDE_SRCS 中指定的源文件会从被 glob 的文件中移除。
- INCLUDE_DIRS - 相对于组件目录的路径，该路径将被添加到需要当前组件的所有其他组件的 include 搜索路径中。
- PRIV_INCLUDE_DIRS - 必须是相对于组件目录的目录路径，它仅被添加到这个组件源文件的 include 搜索路径中。
- REQUIRES - 组件的公共组件依赖项。
- PRIV_REQUIRES - 组件的私有组件依赖项；在仅用于配置的组件上会被忽略。
- LDFRAGMENTS - 组件链接器片段文件。
- REQUIRED_IDF_TARGETS - 指定该组件唯一支持的目标。

以下内容用于*将数据嵌入到组件中*，并在确定组件是否仅用于配置时被视为源文件。这意味着，即使组件没有指定源文件，如果组件指定了以下其中之一，仍然会在内部为组件创建一个静态库。

- EMBED_FILES - 嵌入组件的二进制文件
- EMBED_TXTFILES - 嵌入组件的文本文件

idf 组件属性

组件的属性值可以通过使用构建命令 `idf_component_get_property` 来获取。例如，以下命令可以获取 `freertos` 组件的目录。

```
idf_component_get_property(dir freertos COMPONENT_DIR)
message(STATUS "The 'freertos' component directory is: ${dir}")
```

- COMPONENT_ALIAS - COMPONENT_LIB 的别名，用于将组件链接到外部目标；由 `idf_build_component` 设置，别名库本身由 `idf_component_register` 创建。
- COMPONENT_DIR - 组件目录；由 `idf_build_component` 设置。
- COMPONENT_LIB - 所创建的组件静态/接口库的名称；由 `idf_build_component` 设置，库本身由 `idf_component_register` 创建。
- COMPONENT_NAME - 组件的名称；由 `idf_build_component` 根据组件的目录名设置。
- COMPONENT_TYPE - 组件的类型（LIBRARY 或 CONFIG_ONLY）。如果一个组件指定了源文件或嵌入了一个文件，那么它的类型就是 LIBRARY。
- EMBED_FILES - 要嵌入组件的文件列表；由 `idf_component_register` EMBED_FILES 参数设置。
- EMBED_TXTFILES - 要嵌入组件的文本文件列表；由 `idf_component_register` EMBED_TXTFILES 参数设置。
- INCLUDE_DIRS - 组件 include 目录列表；由 `idf_component_register` INCLUDE_DIRS 参数设置。
- KCONFIG - 组件 Kconfig 文件；由 `idf_build_component` 设置。
- KCONFIG_PROJBUILD - 组件 Kconfig.projbuild；由 `idf_build_component` 设置。
- LDFRAGMENTS - 组件链接器片段文件列表；由 `idf_component_register` LDFRAGMENTS 参数设置。
- PRIV_INCLUDE_DIRS - 组件私有 include 目录列表；在 LIBRARY 类型的组件 `idf_component_register` PRIV_INCLUDE_DIRS 参数中设置。
- PRIV_REQUIRES - 私有组件依赖关系列表；由 `idf_component_register` PRIV_REQUIRES 参数设置。
- REQUIRED_IDF_TARGETS - 组件支持的目标列表；由 `idf_component_register` EMBED_TXTFILES 参数设置。
- REQUIRES - 公共组件依赖关系列表；由 `idf_component_register` REQUIRES 参数设置。
- SRCS - 组件源文件列表；由 `idf_component_register` 的 SRCS 或 SRC_DIRS/EXCLUDE_SRCS 参数设置。

4.2.19 文件通配 & 增量构建

在 ESP-IDF 组件中添加源文件的首选方法是在 COMPONENT_SRCS 中手动列出它们:

```
idf_component_register(SRCS library/a.c library/b.c platform/platform.c
                       ...)
```

这是在 CMake 中手动列出源文件的 最佳实践。然而，当有许多源文件都需要添加到构建中时，这种方法就会很不方便。ESP-IDF 构建系统因此提供了另一种替代方法，即使用 SRC_DIRS 来指定源文件:

```
idf_component_register(SRC_DIRS library platform
                       ...)
```

后台会使用通配符在指定的目录中查找源文件。但是请注意，在使用这种方法的时候，如果组件中添加了一个新的源文件，CMake 并不知道重新运行配置，最终该文件也没有被加入构建中。

如果是自己添加的源文件，这种折衷还是可以接受的，因为用户可以触发一次干净的构建，或者运行 idf.py reconfigure 来手动重启 CMake。但是，如果你需要与其他使用 Git 等版本控制工具的开发人员共享项目时，问题就会变得更加困难，因为开发人员有可能会拉取新的版本。

ESP-IDF 中的组件使用了第三方的 Git CMake 集成模块(/tools/cmake/third_party/GetGitRevisionDescription.cmake)，任何时候源码仓库的提交记录发生了改变，该模块就会自动重新运行 CMake。即只要拉取了新的 ESP-IDF 版本，CMake 就会重新运行。

对于不属于 ESP-IDF 的项目组件，有以下几个选项供参考:

- 如果项目文件保存在 Git 中，ESP-IDF 会自动跟踪 Git 修订版本，并在它发生变化时重新运行 CMake。
- 如果一些组件保存在第三方 Git 仓库中（不在项目仓库或 ESP-IDF 仓库），则可以在组件 CMakeLists 文件中调用 git_describe 函数，以便在 Git 修订版本发生变化时自动重启 CMake。
- 如果没有使用 Git，请记住在源文件发生变化时手动运行 idf.py reconfigure。
- 使用 idf_component_register 的 SRCS 参数来列出项目组件中的所有源文件则可以完全避免这一问题。

具体选择哪一方式，就要取决于项目本身，以及项目用户。

4.2.20 构建系统的元数据

为了将 ESP-IDF 集成到 IDE 或者其它构建系统中，CMake 在构建的过程中会在 build/ 目录下生成大量元数据文件。运行 cmake 或 idf.py reconfigure （或任何其它 idf.py 构建命令），可以重新生成这些元数据文件。

- compile_commands.json 是标准格式的 JSON 文件，它描述了在项目中参与编译的每个源文件。CMake 其中的一个功能就是生成此文件，许多 IDE 都知道如何解析此文件。
- project_description.json 包含有关 ESP-IDF 项目、已配置路径等的一些常规信息。
- flasher_args.json 包含 esptool.py 工具用于烧录项目二进制文件的参数，此外还有 flash_*_args 文件，可直接与 esptool.py 一起使用。更多详细信息请参阅*Flash 参数*。
- CMakeCache.txt 是 CMake 的缓存文件，包含 CMake 进程、工具链等其它信息。
- config/sdkconfig.json 包含 JSON 格式的项目配置结果。
- config/kconfig_menus.json 是在 menuconfig 中显示菜单的 JSON 格式版本，用于外部 IDE 的 UI。

JSON 配置服务器

confserver.py 工具可以帮助 IDE 轻松地与配置系统的逻辑进行集成，它运行在后台，通过使用 stdin 和 stdout 读写 JSON 文件的方式与调用进程交互。

您可以通过 idf.py confserver 或 ninja confserver 从项目中运行 confserver.py，也可以使用不同的构建生成器来触发类似的目标。

有关 confserver.py 的更多信息，请参阅 tools/kconfig_new/README.md

4.2.21 构建系统内部

构建脚本

ESP-IDF 构建系统的列表文件位于 /tools/cmake 中。实现构建系统核心功能的模块如下

- build.cmake - 构建相关命令，即构建初始化、检索/设置构建属性、构建处理。
- component.cmake - 组件相关的命令，如添加组件、检索/设置组件属性、注册组件。
- kconfig.cmake - 从 Kconfig 文件中生成配置文件（sdkconfig、sdkconfig.h、sdkconfig.cmake 等）。
- ldgen.cmake - 从链接器片段文件生成最终链接器脚本。
- target.cmake - 设置构建目标和工具链文件。
- utilities.cmake - 其它帮助命令。

除了这些文件，还有两个重要的 CMake 脚本在 /tools/cmake 中：

- idf.cmake - 设置构建参数并导入上面列出的核心模块。之所以包括在 CMake 项目中，是为了方便访问 ESP-IDF 构建系统功能。
- project.cmake - 导入 `idf.cmake`，并提供了一个自定义的 “project()“命令，该命令负责处理建立可执行文件时所有的繁重工作。包含在标准 ESP-IDF 项目的顶层 CMakeLists.txt 中。

/tools/cmake 中的其它文件都是构建过程中的支持性文件或第三方脚本。

构建过程

本节介绍了标准的 ESP-IDF 应用构建过程。构建过程可以大致分为四个阶段：

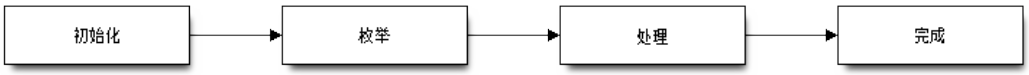

图 1: ESP-IDF Build System Process

初始化

该阶段为构建设置必要的参数。

- **在将 `idf.cmake` 导入 `project.cmake` 后，将执行以下步骤：**
  - 在环境变量中设置 `IDF_PATH` 或从顶层 CMakeLists.txt 中包含的 `project.cmake` 路径推断相对路径。
  - 将 /tools/cmake 添加到 `CMAKE_MODULE_PATH` 中，并导入核心模块和各种辅助/第三方脚本。
  - 设置构建工具/可执行文件，如默认的 Python 解释器。
  - 获取 ESP-IDF git 修订版，并存储为 `IDF_VER`。
  - 设置全局构建参数，即编译选项、编译定义、包括所有组件的 include 目录。
  - 将 components 中的组件添加到构建中。
- **自定义 `project()` 命令的初始部分执行以下步骤：**
  - 在环境变量或 CMake 缓存中设置 `IDF_TARGET` 以及设置相应要使用的 “CMAKE_TOOLCHAIN_FILE“。
  - 添加 `EXTRA_COMPONENTS_DIRS` 中的组件至构建中
  - 从 `COMPONENTS/EXCLUDE_COMPONENTS`、`SDKCONFIG`、`SDKCONFIG_DEFAULTS` 等变量中为调用命令 `idf_build_process()` 准备参数。

调用 `idf_build_process()` 命令标志着这个阶段的结束。

枚举

这个阶段会建立一个需要在构建过程中处理的组件列表，该阶段在 `idf_build_process()` 的前半部分进行。

- 检索每个组件的公共和私有依赖。创建一个子进程，以脚本模式执行每个组件的 CMakeLists.txt。`idf_component_register` REQUIRES 和 PRIV_REQUIRES 参数的值会返回给父进程。这就是所谓的早期扩展。在这一步中定义变量 `CMAKE_BUILD_EARLY_EXPANSION`。
- 根据公共和私有的依赖关系，递归地导入各个组件。

处理

该阶段处理构建中的组件，是 `idf_build_process()` 的后半部分。

- 从 sdkconfig 文件中加载项目配置，并生成 sdkconfig.cmake 和 sdkconfig.h 头文件。这两个文件分别定义了可以从构建脚本和 C/C++ 源文件/头文件中访问的配置变量/宏。
- 导入各组件的 `project_include.cmake`。
- 将每个组件添加为一个子目录，处理其 CMakeLists.txt。组件 CMakeLists.txt 调用注册命令 `idf_component_register` 添加源文件、导入目录、创建组件库、链接依赖关系等。

完成

该阶段是 `idf_build_process()` 剩余的步骤。

- 创建可执行文件并将其链接到组件库中。
- 生成 project_description.json 等项目元数据文件并且显示所建项目等相关信息。

请参考 /tools/cmake/project.cmake 获取更多信息。

4.2.22 从 ESP-IDF GNU Make 构建系统迁移到 CMake 构建系统

ESP-IDF CMake 构建系统与旧版的 GNU Make 构建系统在某些方面非常相似，开发者都需要提供 include 目录、源文件等。然而，有一个语法上的区别，即对于 ESP-IDF CMake 构建系统，开发者需要将这些作为参数传递给注册命令 `idf_component_register`。

自动转换工具

/tools/cmake/convert_to_cmake.py 中提供了一个项目自动转换工具。运行此命令时需要加上项目路径，如下所示:

```
$IDF_PATH/tools/cmake/convert_to_cmake.py /path/to/project_dir
```

项目目录必须包含 Makefile 文件，并确保主机已安装 GNU Make (`make`) 工具，并且被添加到了 PATH 环境变量中。

该工具会将项目 Makefile 文件和所有组件的 `component.mk` 文件转换为对应的 `CMakeLists.txt` 文件。

转换过程如下：该工具首先运行 `make` 来展开 ESP-IDF 构建系统设置的变量，然后创建相应的 CMakelists 文件来设置相同的变量。

重要： 当转换工具转换一个 `component.mk` 文件时，它并不能确定该组件依赖于哪些其他组件。这些信息需要通过编辑新的组件 `CMakeLists.txt` 文件并添加 `REQUIRES` 和/或 `PRIV_REQUIRES` 子句来手动添加。否则，组件中的源文件会因为找不到其他组件的头文件而编译失败。请参考组件依赖 获取更多信息。

转换工具并不能处理复杂的 Makefile 逻辑或异常的目标，这些需要手动转换。

CMake 中不可用的功能

有些功能已从 CMake 构建系统中移除，或者已经发生很大改变。GNU Make 构建系统中的以下变量已从 CMake 构建系统中删除：

- COMPONENT_BUILD_DIR：由 CMAKE_CURRENT_BINARY_DIR 替代。
- COMPONENT_LIBRARY：默认为 $(COMPONENT_NAME).a 但是库名可以被组件覆盖。在 CMake 构建系统中，组件库名称不可再被组件覆盖。
- CC、LD、AR、OBJCOPY：gcc xtensa 交叉工具链中每个工具的完整路径。CMake 使用 CMAKE_C_COMPILER、CMAKE_C_LINK_EXECUTABLE 和 CMAKE_OBJCOPY 进行替代。完整列表请参阅 CMake 语言变量。
- HOSTCC、HOSTLD、HOSTAR：宿主机本地工具链中每个工具的全名。CMake 系统不再提供此变量，外部项目需要手动检测所需的宿主机工具链。
- COMPONENT_ADD_LDFLAGS：用于覆盖链接标志。CMake 中使用 target_link_libraries 命令替代。
- COMPONENT_ADD_LINKER_DEPS：链接过程依赖的文件列表。target_link_libraries 通常会自动推断这些依赖。对于链接脚本，可以使用自定义的 CMake 函数 target_linker_scripts。
- COMPONENT_SUBMODULES：不再使用。CMake 会自动枚举 ESP-IDF 仓库中所有的子模块。
- COMPONENT_EXTRA_INCLUDES：曾是 COMPONENT_PRIV_INCLUDEDIRS 变量的替代版本，仅支持绝对路径。CMake 系统中统一使用 COMPONENT_PRIV_INCLUDEDIRS（可以是相对路径，也可以是绝对路径）。
- COMPONENT_OBJS：以前，可以以目标文件列表的方式指定组件源，现在，可以通过 COMPONENT_SRCS 以源文件列表的形式指定组件源。
- COMPONENT_OBJEXCLUDE：已被 COMPONENT_SRCEXCLUDE 替换。用于指定源文件（绝对路径或组件目录的相对路径）。
- COMPONENT_EXTRA_CLEAN：已被 ADDITIONAL_MAKE_CLEAN_FILES 属性取代，注意，*CMake 对此项功能有部分限制*。
- COMPONENT_OWNBUILDTARGET & COMPONENT_OWNCLEANTARGET：已被 CMake 外部项目 替代，详细内容请参阅*完全覆盖组件的构建过程*。
- COMPONENT_CONFIG_ONLY：已被 register_config_only_component() 函数替代，请参阅*仅配置组件*。
- CFLAGS、CPPFLAGS、CXXFLAGS：已被相应的 CMake 命令替代，请参阅*组件编译控制*。

无默认值的变量

以下变量不再具有默认值：

- 源目录（Make 中的 COMPONENT_SRCDIRS 变量，CMake 中 idf_component_register 的 SRC_DIRS 参数）
- include 目录（Make 中的 COMPONENT_ADD_INCLUDEDIRS 变量，CMake 中 idf_component_register 的 INCLUDE_DIRS 参数)

不再需要的变量

在 CMake 构建系统中，如果设置了 COMPONENT_SRCS，就不需要再设置 COMPONENT_SRCDIRS。实际上，CMake 构建系统中如果设置了 COMPONENT_SRCDIRS，那么 COMPONENT_SRCS 就会被忽略。

从 Make 中烧录

仍然可以使用 make flash 或者类似的目标来构建和烧录，但是项目 sdkconfig 不能再用来指定串口和波特率。可以使用环境变量来覆盖串口和波特率的设置，详情请参阅*使用 Ninja/Make 来烧录*。

4.3 错误处理

4.3.1 概述

在应用程序开发中，及时发现并处理在运行时期的错误，对于保证应用程序的健壮性非常重要。常见的运行时错误有如下几种：

- 可恢复的错误：
  - 通过函数的返回值（错误码）表示的错误
  - 使用 `throw` 关键字抛出的 C++ 异常
- 不可恢复（严重）的错误：
  - 断言失败（使用 `assert` 宏或者其它类似方法）或者直接调用 `abort()` 函数造成的错误
  - CPU 异常：访问受保护的内存区域、非法指令等
  - 系统级检查：看门狗超时、缓存访问错误、堆栈溢出、堆栈粉碎、堆栈损坏等

本文将介绍 ESP-IDF 中针对可恢复错误的错误处理机制，并提供一些常见错误的处理模式。

关于如何处理不可恢复的错误，请查阅*不可恢复错误*。

4.3.2 错误码

ESP-IDF 中大多数函数会返回*esp_err_t* 类型的错误码，*esp_err_t* 实质上是带符号的整型，`ESP_OK` 代表成功（没有错误），具体值定义为 0。

在 ESP-IDF 中，许多头文件都会使用预处理器，定义可能出现的错误代码。这些错误代码通常均以 `ESP_ERR_` 前缀开头，一些常见错误（比如内存不足、超时、无效参数等）的错误代码则已经在 `esp_err.h` 文件中定义好了。此外，ESP-IDF 中的各种组件 (component) 也都可以针对具体情况，自行定义更多错误代码。

完整错误代码列表，请见*错误代码参考* 中查看完整的错误列表。

4.3.3 错误码到错误消息

错误代码并不直观，因此 ESP-IDF 还可以使用*esp_err_to_name()* 或者*esp_err_to_name_r()* 函数，将错误代码转换为具体的错误消息。例如，我们可以向*esp_err_to_name()* 函数传递错误代码 `0x101`，可以得到返回字符串 “ESP_ERR_NO_MEM”。这样一来，我们可以在日志中输出更加直观的错误消息，而不是简单的错误码，从而帮助研发人员更快理解发生了何种错误。

此外，如果出现找不到匹配的 `ESP_ERR_` 值的情况，函数*esp_err_to_name_r()* 则会尝试将错误码作为一种 标准 POSIX 错误代码 进行解释。具体过程为：POSIX 错误代码（例如 `ENOENT`，`ENOMEM`）定义在 `errno.h` 文件中，可以通过 `errno` 变量获得，进而调用 `strerror_r` 函数实现。在 ESP-IDF 中，`errno` 是一个基于线程的局部变量，即每个 FreeRTOS 任务都有自己的 `errno` 副本，通过函数修改 `errno` 也只会作用于当前任务中的 `errno` 变量值。

该功能（即在无法匹配 `ESP_ERR_` 值时，尝试用标准 POSIX 解释错误码）默认启用。用户也可以禁用该功能，从而减小应用程序的二进制文件大小，详情可见*CONFIG_ESP_ERR_TO_NAME_LOOKUP*。注意，该功能对禁用并不影响*esp_err_to_name()* 和*esp_err_to_name_r()* 函数的定义，用户仍可调用这两个函数转化错误码。在这种情况下，*esp_err_to_name()* 函数在遇到无法匹配错误码的情况会返回 `UNKNOWN ERROR`，而*esp_err_to_name_r()* 函数会返回 `Unknown error 0xXXXX(YYYYY)`，其中 `0xXXXX` 和 `YYYYY` 分别代表错误代码的十六进制和十进制表示。

4.3.4 ESP_ERROR_CHECK 宏

宏 `ESP_ERROR_CHECK()` 的功能和 `assert` 类似，不同之处在于：这个宏会检查*esp_err_t* 的值，而非判断 `bool` 条件。如果传给 `ESP_ERROR_CHECK()` 的参数不等于*ESP_OK*，则会在控制台上打印错误消息，然后调用 `abort()` 函数。

错误消息通常如下所示：

```
ESP_ERROR_CHECK failed: esp_err_t 0x107 (ESP_ERR_TIMEOUT) at 0x400d1fdf

file: "/Users/user/esp/example/main/main.c" line 20
func: app_main
expression: sdmmc_card_init(host, &card)

Backtrace: 0x40086e7c:0x3ffb4ff0 0x40087328:0x3ffb5010 0x400d1fdf:0x3ffb5030␣
↪0x400d0816:0x3ffb5050
```

- 第一行打印错误代码的十六进制表示，及该错误在源代码中的标识符。这个标识符取决于*CONFIG_ESP_ERR_TO_NAME_LOOKUP* 选项的设定。最后，第一行还会打印程序中该错误发生的具体位置。
- 下面几行显示了程序中调用 ESP_ERROR_CHECK() 宏的具体位置，以及传递给该宏的参数。
- 最后一行打印回溯结果。对于所有不可恢复错误，这里在应急处理程序中打印的内容都是一样的。更多有关回溯结果的详细信息，请参阅不可恢复错误 。

注解： 如果使用*IDF monitor*, 则最后一行回溯结果中的地址将会被替换为相应的文件名和行号。

4.3.5 错误处理模式

1. 尝试恢复。根据具体情况不同，我们具体可以：
   - 在一段时间后，重新调用该函数；
   - 尝试删除该驱动，然后重新进行 “初始化”；
   - 采用其他带外机制，修改导致错误发生的条件（例如，对一直没有响应的外设进行复位等）。

   示例：

```
esp_err_t err;
do {
    err = sdio_slave_send_queue(addr, len, arg, timeout);
    // 如果发送队列已满就不断重试
} while (err == ESP_ERR_TIMEOUT);
if (err != ESP_OK) {
    // 处理其他错误
}
```

2. 将错误传递回调用程序。在某些中间件组件中，采用此类处理模式代表函数必须以相同的错误码退出，这样才能确保所有分配的资源都能得到释放。
   示例：

```
sdmmc_card_t* card = calloc(1, sizeof(sdmmc_card_t));
if (card == NULL) {
    return ESP_ERR_NO_MEM;
}
esp_err_t err = sdmmc_card_init(host, &card);
if (err != ESP_OK) {
    // 释放内存
    free(card);
    // 将错误码传递给上层（例如通知用户）
    // 或者，应用程序可以自定义错误代码并返回
    return err;
}
```

3. 转为不可恢复错误，比如使用 ESP_ERROR_CHECK。详情请见*ESP_ERROR_CHECK 宏* 章节。
   对于中间件组件而言，通常并不希望在发生错误时中止应用程序。不过，有时在应用程序级别，这种做法是可以接受的。在 ESP-IDF 的示例代码中，很多都会使用 ESP_ERROR_CHECK 来处理各种 API 引发的错误，虽然这不是应用程序的最佳做法，但可以让示例代码看起来更加简洁。
   示例：

```
ESP_ERROR_CHECK(spi_bus_initialize(host, bus_config, dma_chan));
```

4.3.6 C++ 异常

默认情况下，ESP-IDF 会禁用对 C++ 异常的支持，但是可以通过*CONFIG_COMPILER_CXX_EXCEPTIONS* 选项启用。

通常情况下，启用异常处理会让应用程序的二进制文件增加几 kB。此外，启用该功能时还应为异常事故池预留一定内存。当应用程序无法从堆中分配异常对象时，就可以使用这个池中的内存。该内存池的大小可以通过*CONFIG_COMPILER_CXX_EXCEPTIONS_EMG_POOL_SIZE* 来设定。

如果 C++ 程序抛出了异常，但是程序中并没有 `catch` 代码块来捕获该异常，那么程序的运行就会被 `abort` 函数中止，然后打印回溯信息。有关回溯的更多信息，请参阅 不可恢复错误 。

4.4 严重错误

4.4.1 概述

在某些情况下，程序并不会按照我们的预期运行，在 ESP-IDF 中，这些情况包括：

- CPU 异常：非法指令，加载/存储时的内存对齐错误，加载/存储时的访问权限错误，双重异常。
- 系统级检查错误：
  - 中断看门狗 超时
  - 任务看门狗 超时（只有开启*CONFIG_ESP_TASK_WDT_PANIC* 后才会触发严重错误）
  - 高速缓存访问错误
  - 掉电检测事件
  - 堆栈溢出
  - Stack 粉碎保护检查
  - Heap 完整性检查
- 使用 `assert`、`configASSERT` 等类似的宏断言失败。

本指南会介绍 ESP-IDF 中这类错误的处理流程，并给出对应的解决建议。

4.4.2 紧急处理程序

概述 中列举的所有错误都会由 紧急处理程序（*Panic Handler*） 负责处理。

紧急处理程序首先会将出错原因打印到控制台，例如 CPU 异常的错误信息通常会类似于:

```
Guru Meditation Error: Core 0 panic'ed (IllegalInstruction). Exception was␣
↪unhandled.
```

对于一些系统级检查错误（如中断看门狗超时，高速缓存访问错误等），错误信息会类似于:

```
Guru Meditation Error: Core 0 panic'ed (Cache disabled but cached memory region␣
↪accessed)
```

不管哪种情况，错误原因都会被打印在括号中。请参阅*Guru Meditation 错误* 以查看所有可能的出错原因。

紧急处理程序接下来的行为将取决于*CONFIG_ESP_SYSTEM_PANIC* 的设置，支持的选项包括：

- 打印 CPU 寄存器，然后重启（`CONFIG_ESP_SYSTEM_PANIC_PRINT_REBOOT`）- 默认选项
  打印系统发生异常时 CPU 寄存器的值，打印回溯，最后重启芯片。
- 打印 CPU 寄存器，然后暂停（`CONFIG_ESP_SYSTEM_PANIC_PRINT_HALT`）
  与上一个选项类似，但不会重启，而是选择暂停程序的运行。重启程序需要外部执行复位操作。
- 静默重启（`CONFIG_ESP_SYSTEM_PANIC_SILENT_REBOOT`）
  不打印 CPU 寄存器的值，也不打印回溯，立即重启芯片。

- 调用 GDB Stub（CONFIG_ESP_SYSTEM_PANIC_GDBSTUB）
  启动 GDB 服务器，通过控制台 UART 接口与 GDB 进行通信。详细信息请参阅*GDB Stub*。

紧急处理程序的行为还受到另外两个配置项的影响：

- 如果*CONFIG_ESP32S2_DEBUG_OCDAWARE* 被使能了（默认），紧急处理程序会检测 ESP32-S2 是否已经连接 JTAG 调试器。如果检测成功，程序会暂停运行，并将控制权交给调试器。在这种情况下，寄存器和回溯不会被打印到控制台，并且也不会使用 GDB Stub 和 Core Dump 的功能。
- 如果使能了*Core Dump* 功能（CONFIG_ESP32_ENABLE_COREDUMP_TO_FLASH 或者 CONFIG_ESP32_ENABLE_COREDUMP_TO_UART 选项），系统状态（任务堆栈和寄存器）会被转储到 Flash 或者 UART 以供后续分析。

下图展示了紧急处理程序的行为：

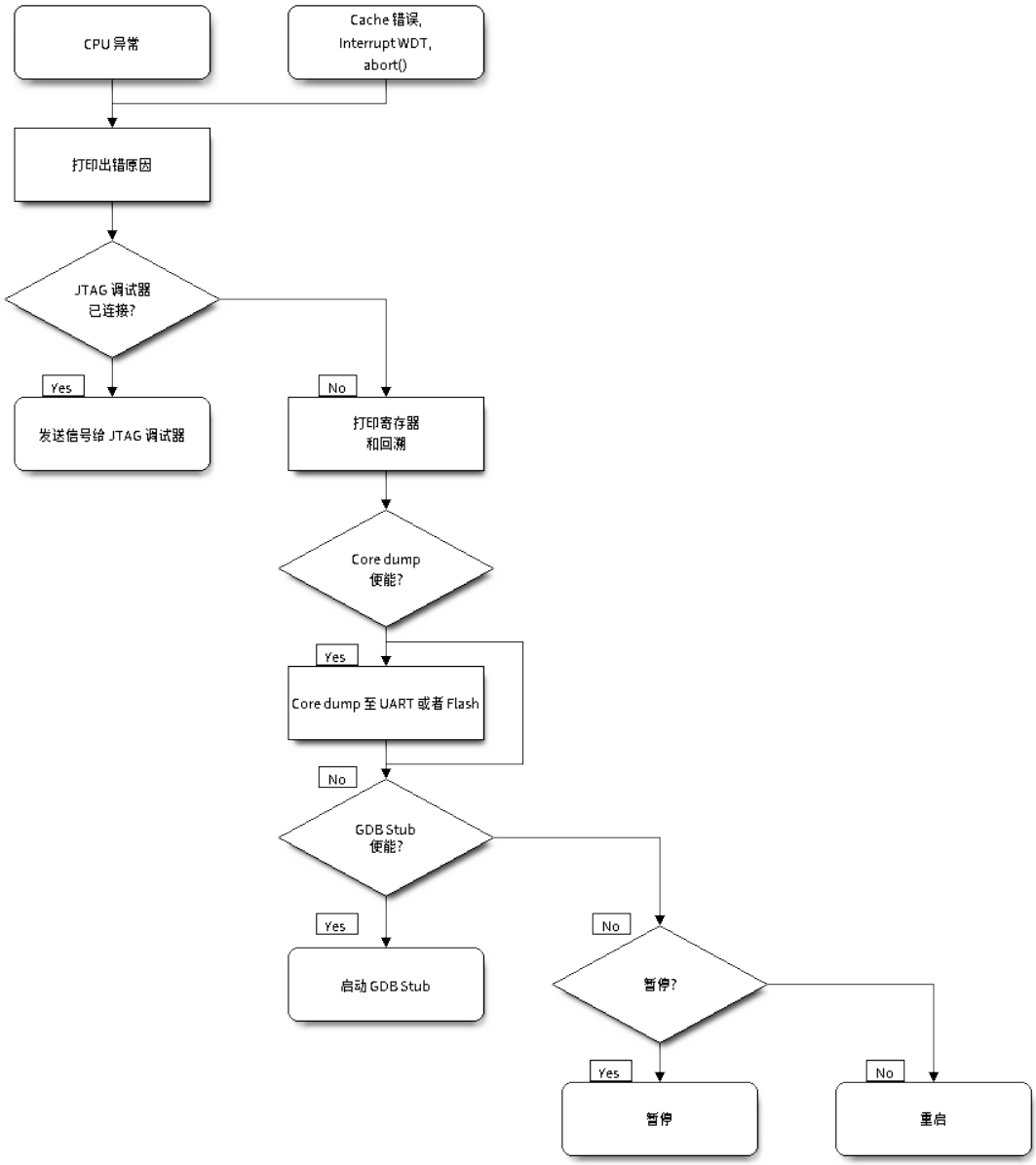

图 2: 紧急处理程序流程图（点击放大）

4.4.3 寄存器转储与回溯

除非启用了 CONFIG_ESP_SYSTEM_PANIC_SILENT_REBOOT 否则紧急处理程序会将 CPU 寄存器和回溯打印到控制台:

```
Core 0 register dump:
PC      : 0x400e14ed  PS      : 0x00060030  A0      : 0x800d0805  A1      :␣
↪0x3ffb5030
A2      : 0x00000000  A3      : 0x00000001  A4      : 0x00000001  A5      :␣
↪0x3ffb50dc
A6      : 0x00000000  A7      : 0x00000001  A8      : 0x00000000  A9      :␣
↪0x3ffb5000
A10     : 0x00000000  A11     : 0x3ffb2bac  A12     : 0x40082d1c  A13     :␣
↪0x06ff1ff8
A14     : 0x3ffb7078  A15     : 0x00000000  SAR     : 0x00000014  EXCCAUSE:␣
↪0x0000001d
EXCVADDR: 0x00000000  LBEG    : 0x4000c46c  LEND    : 0x4000c477  LCOUNT  :␣
↪0xffffffff

Backtrace: 0x400e14ed:0x3ffb5030 0x400d0802:0x3ffb5050
```

仅会打印异常帧中 CPU 寄存器的值，即引发 CPU 异常或者其它严重错误时刻的值。

紧急处理程序如果是因 abort() 而调用，则不会打印寄存器转储。

在某些情况下，例如中断看门狗超时，紧急处理程序会额外打印 CPU 寄存器（EPC1-EPC4）的值，以及另一个 CPU 的寄存器值和代码回溯。

回溯行包含了当前任务中每个堆栈帧的 PC:SP 对（PC 是程序计数器，SP 是堆栈指针）。如果在 ISR 中发生了严重错误，回溯会同时包括被中断任务的 PC:SP 对，以及 ISR 中的 PC:SP 对。

如果使用了*IDF 监视器*，该工具会将程序计数器的值转换为对应的代码位置（函数名，文件名，行号），并加以注释:

```
Core 0 register dump:
PC      : 0x400e14ed  PS      : 0x00060030  A0      : 0x800d0805  A1      :␣
↪0x3ffb5030
0x400e14ed: app_main at /Users/user/esp/example/main/main.cpp:36

A2      : 0x00000000  A3      : 0x00000001  A4      : 0x00000001  A5      :␣
↪0x3ffb50dc
A6      : 0x00000000  A7      : 0x00000001  A8      : 0x00000000  A9      :␣
↪0x3ffb5000
A10     : 0x00000000  A11     : 0x3ffb2bac  A12     : 0x40082d1c  A13     :␣
↪0x06ff1ff8
0x40082d1c: _calloc_r at /Users/user/esp/esp-idf/components/newlib/syscalls.c:51

A14     : 0x3ffb7078  A15     : 0x00000000  SAR     : 0x00000014  EXCCAUSE:␣
↪0x0000001d
EXCVADDR: 0x00000000  LBEG    : 0x4000c46c  LEND    : 0x4000c477  LCOUNT  :␣
↪0xffffffff

Backtrace: 0x400e14ed:0x3ffb5030 0x400d0802:0x3ffb5050
0x400e14ed: app_main at /Users/user/esp/example/main/main.cpp:36

0x400d0802: main_task at /Users/user/esp/esp-idf/components/esp32s2/cpu_start.c:470
```

若要查找发生严重错误的代码位置，请查看“Backtrace”的后面几行，发生严重错误的代码显示在顶行，后续几行显示的是调用堆栈。

4.4.4 GDB Stub

如果启用了 CONFIG_ESP_SYSTEM_PANIC_GDBSTUB 选项，在发生严重错误时，紧急处理程序不会复位芯片，相反，它将启动 GDB 远程协议服务器，通常称为 GDB Stub。发生这种情况时，可以让主机上运行的 GDB 实例通过 UART 端口连接到 ESP32。

如果使用了*IDF 监视器*，该工具会在 UART 端口检测到 GDB Stub 提示符后自动启动 GDB，输出会类似于:

```
Entering gdb stub now.
$T0b#e6GNU gdb (crosstool-NG crosstool-ng-1.22.0-80-gff1f415) 7.10
Copyright (C) 2015 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law.  Type "show copying"
and "show warranty" for details.
This GDB was configured as "--host=x86_64-build_apple-darwin16.3.0 --target=xtensa-
↪esp32s2-elf".
Type "show configuration" for configuration details.
For bug reporting instructions, please see:
<http://www.gnu.org/software/gdb/bugs/>.
Find the GDB manual and other documentation resources online at:
<http://www.gnu.org/software/gdb/documentation/>.
For help, type "help".
Type "apropos word" to search for commands related to "word"...
Reading symbols from /Users/user/esp/example/build/example.elf...done.
Remote debugging using /dev/cu.usbserial-31301
0x400e1b41 in app_main ()
    at /Users/user/esp/example/main/main.cpp:36
36      *((int*) 0) = 0;
(gdb)
```

在 GDB 会话中，我们可以检查 CPU 寄存器，本地和静态变量以及内存中任意位置的值。但是不支持设置断点，改变 PC 值或者恢复程序的运行。若要复位程序，请退出 GDB 会话，在 IDF 监视器中连续输入 Ctrl-T Ctrl-R，或者按下开发板上的复位按键也可以重新运行程序。

4.4.5 Guru Meditation 错误

本节将对打印在 Guru Meditation Error: Core panic'ed 后面括号中的致错原因进行逐一解释。

注解: 想要了解“Guru Meditation”的历史渊源，请参阅 维基百科 。

IllegalInstruction

此 CPU 异常表示当前执行的指令不是有效指令，引起此错误的常见原因包括:

- FreeRTOS 中的任务函数已返回。在 FreeRTOS 中，如果想终止任务函数，需要调用 *vTaskDelete()* 函数释放当前任务的资源，而不是直接返回。
- 无法从 SPI Flash 中加载下一条指令，这通常发生在:
  - 应用程序将 SPI Flash 的引脚重新配置为其它功能（如 GPIO，UART 等等）。有关 SPI Flash 引脚的详细信息，请参阅硬件设计指南和芯片/模组的数据手册。
  - 某些外部设备意外连接到 SPI Flash 的引脚上，干扰了 ESP32-S2 和 SPI Flash 之间的通信。

InstrFetchProhibited

此 CPU 异常表示 CPU 无法加载指令，因为指令的地址不在 IRAM 或者 IROM 中的有效区域中。

通常这意味着代码中调用了并不指向有效代码块的函数指针。这种情况下，可以查看 PC （程序计数器）寄存器的值并做进一步判断：若为 0 或者其它非法值（即只要不是 0x4xxxxxxx 的情况），则证实确实是该原因。

LoadProhibited, StoreProhibited

当应用程序尝试读取或写入无效的内存位置时，会发生此类 CPU 异常。此类无效内存地址可以在寄存器转储的 EXCVADDR 中找到。如果该地址为零，通常意味着应用程序正尝试解引用一个 NULL 指针。如果该地址接近于零，则通常意味着应用程序尝试访问某个结构体的成员，但是该结构体的指针为 NULL。如果该地址是其它非法值（不在 0x3fxxxxxx - 0x6xxxxxxx 的范围内），则可能意味着用于访问数据的指针未初始化或者已经损坏。

IntegerDivideByZero

应用程序尝试将整数除以零。

LoadStoreAlignment

应用程序尝试读取/写入的内存位置不符合加载/存储指令对字节对齐大小的要求，例如，32 位加载指令只能访问 4 字节对齐的内存地址，而 16 位加载指令只能访问 2 字节对齐的内存地址。

LoadStoreError

应用程序尝试从仅支持 32 位加载/存储的内存区域执行 8 位或 16 位加载/存储操作，例如，解引用一个指向指令内存区域的 char* 指针就会导致这样的错误。

Unhandled debug exception

这后面通常会再跟一条消息:

```
Debug exception reason: Stack canary watchpoint triggered (task_name)
```

此错误表示应用程序写入的位置越过了 task_name 任务堆栈的末尾，请注意，并非每次堆栈溢出都会触发此错误。任务有可能会绕过堆栈金丝雀（stack canary）的位置访问堆栈，在这种情况下，监视点就不会被触发。

Interrupt wdt timeout on CPU0 / CPU1

这表示发生了中断看门狗超时，详细信息请查阅*看门狗* 文档。

Cache disabled but cached memory region accessed

在某些情况下，ESP-IDF 会暂时禁止通过高速缓存访问外部 SPI Flash 和 SPI RAM，例如在使用 spi_flash API 读取/写入/擦除/映射 SPI Flash 的时候。在这些情况下，任务会被挂起，并且未使用 ESP_INTR_FLAG_IRAM 注册的中断处理程序会被禁用。请确保任何使用此标志注册的中断处理程序所访问的代码和数据分别位于 IRAM 和 DRAM 中。更多详细信息请参阅*SPI Flash API 文档*。

4.4.6 其它严重错误

欠压

ESP32-S2 内部集成掉电检测电路，并且会默认启用。如果电源电压低于安全值，掉电检测器可以触发系统复位。掉电检测器可以使用*CONFIG_ESP32S2_BROWNOUT_DET* 和*CONFIG_ESP32S2_BROWNOUT_DET_LVL_SEL* 这两个选项进行设置。

当掉电检测器被触发时，会打印如下信息:

```
Brownout detector was triggered
```

芯片会在该打印信息结束后复位。

请注意，如果电源电压快速下降，则只能在控制台上看到部分打印信息。

Heap 不完整

ESP-IDF 堆的实现包含许多运行时的堆结构检查，可以在 menuconfig 中开启额外的检查（“Heap Poisoning”）。如果其中的某项检查失败，则会打印类似如下信息:

```
CORRUPT HEAP: Bad tail at 0x3ffe270a. Expected 0xbaad5678 got 0xbaac5678
assertion "head != NULL" failed: file "/Users/user/esp/esp-idf/components/heap/
↪multi_heap_poisoning.c", line 201, function: multi_heap_free
abort() was called at PC 0x400dca43 on core 0
```

更多详细信息，请查阅堆内存调试 文档。

Stack 粉碎

Stack 粉碎保护（基于 GCC `-fstack-protector*` 标志）可以通过 ESP-IDF 中的*CONFIG_COMPILER_STACK_CHECK_MODE* 选项来开启。如果检测到 Stack 粉碎，则会打印类似如下的信息:

```
Stack smashing protect failure!

abort() was called at PC 0x400d2138 on core 0

Backtrace: 0x4008e6c0:0x3ffc1780 0x4008e8b7:0x3ffc17a0 0x400d2138:0x3ffc17c0␣
↪0x400e79d5:0x3ffc17e0 0x400e79a7:0x3ffc1840 0x400e79df:0x3ffc18a0␣
↪0x400e2235:0x3ffc18c0 0x400e1916:0x3ffc18f0 0x400e19cd:0x3ffc1910␣
↪0x400e1a11:0x3ffc1930 0x400e1bb2:0x3ffc1950 0x400d2c44:0x3ffc1a80
0
```

回溯信息会指明发生 Stack 粉碎的函数，建议检查函数中是否有代码访问本地数组时发生了越界。

4.5 Event Handling

Several ESP-IDF components use *events* to inform application about state changes, such as connection or disconnection. This document gives an overview of these event mechanisms.

4.5.1 Wi-Fi, Ethernet, and IP Events

Before the introduction of *esp_event library*, events from Wi-Fi driver, Ethernet driver, and TCP/IP stack were dispatched using the so-called *legacy event loop*. The following sections explain each of the methods.

esp_event Library Event Loop

esp_event library is designed to supersede the legacy event loop for the purposes of event handling in ESP-IDF. In the legacy event loop, all possible event types and event data structures had to be defined in `system_event_id_t` enumeration and `system_event_info_t` union, which made it impossible to send custom events to the event loop, and use the event loop for other kinds of events (e.g. Mesh). Legacy event loop also supported only one event handler function, therefore application components could not handle some of Wi-Fi or IP events themselves, and required application to forward these events from its event handler function.

See *esp_event library API reference* for general information on using this library. Wi-Fi, Ethernet, and IP events are sent to the *default event loop* provided by this library.

Legacy Event Loop

This event loop implementation is started using `esp_event_loop_init()` function. Application typically supplies an *event handler*, a function with the following signature:

```
esp_err_t event_handler(void *ctx, system_event_t *event)
{
}
```

Both the pointer to event handler function, and an arbitrary context pointer are passed to `esp_event_loop_init()`.

When Wi-Fi, Ethernet, or IP stack generate an event, this event is sent to a high-priority `event` task via a queue. Application-provided event handler function is called in the context of this task. Event task stack size and event queue size can be adjusted using *CONFIG_ESP_SYSTEM_EVENT_TASK_STACK_SIZE* and *CONFIG_ESP_SYSTEM_EVENT_QUEUE_SIZE* options, respectively.

Event handler receives a pointer to the event structure (`system_event_t`) which describes current event. This structure follows a *tagged union* pattern: `event_id` member indicates the type of event, and `event_info` member is a union of description structures. Application event handler will typically use `switch(event->event_id)` to handle different kinds of events.

If application event handler needs to relay the event to some other task, it is important to note that event pointer passed to the event handler is a pointer to temporary structure. To pass the event to another task, application has to make a copy of the entire structure.

Event IDs and Corresponding Data Structures

Event ID (legacy event ID)	Event data structure
Wi-Fi	
WIFI_EVENT_WIFI_READY (SYSTEM_EVENT_WIFI_READY)	n/a
WIFI_EVENT_SCAN_DONE (SYSTEM_EVENT_SCAN_DONE)	*wifi_event_sta_scan_done_t*
WIFI_EVENT_STA_START (SYSTEM_EVENT_STA_START)	n/a
WIFI_EVENT_STA_STOP (SYSTEM_EVENT_STA_STOP)	n/a
WIFI_EVENT_STA_CONNECTED (SYSTEM_EVENT_STA_CONNECTED)	*wifi_event_sta_connected_t*
WIFI_EVENT_STA_DISCONNECTED (SYSTEM_EVENT_STA_DISCONNECTED)	*wifi_event_sta_disconnected_t*
WIFI_EVENT_STA_AUTHMODE_CHANGE (SYSTEM_EVENT_STA_AUTHMODE_CHANGE)	*wifi_event_sta_authmode_change_t*
WIFI_EVENT_STA_WPS_ER_SUCCESS (SYSTEM_EVENT_STA_WPS_ER_SUCCESS)	n/a
WIFI_EVENT_STA_WPS_ER_FAILED (SYSTEM_EVENT_STA_WPS_ER_FAILED)	*wifi_event_sta_wps_fail_reason_t*
WIFI_EVENT_STA_WPS_ER_TIMEOUT (SYSTEM_EVENT_STA_WPS_ER_TIMEOUT)	n/a
WIFI_EVENT_STA_WPS_ER_PIN (SYSTEM_EVENT_STA_WPS_ER_PIN)	*wifi_event_sta_wps_er_pin_t*
WIFI_EVENT_AP_START (SYSTEM_EVENT_AP_START)	n/a
WIFI_EVENT_AP_STOP (SYSTEM_EVENT_AP_STOP)	n/a
WIFI_EVENT_AP_STACONNECTED (SYSTEM_EVENT_AP_STACONNECTED)	*wifi_event_ap_staconnected_t*
WIFI_EVENT_AP_STADISCONNECTED (SYSTEM_EVENT_AP_STADISCONNECTED)	*wifi_event_ap_stadisconnected_t*
WIFI_EVENT_AP_PROBEREQRECVED (SYSTEM_EVENT_AP_PROBEREQRECVED)	*wifi_event_ap_probe_req_rx_t*
Ethernet	
ETHERNET_EVENT_START (SYSTEM_EVENT_ETH_START)	n/a
ETHERNET_EVENT_STOP (SYSTEM_EVENT_ETH_STOP)	n/a
ETHERNET_EVENT_CONNECTED (SYSTEM_EVENT_ETH_CONNECTED)	n/a
ETHERNET_EVENT_DISCONNECTED (SYSTEM_EVENT_ETH_DISCONNECTED)	n/a
IP	
IP_EVENT_STA_GOT_IP (SYSTEM_EVENT_STA_GOT_IP)	`ip_event_got_ip_t`
IP_EVENT_STA_LOST_IP (SYSTEM_EVENT_STA_LOST_IP)	n/a
IP_EVENT_AP_STAIPASSIGNED (SYSTEM_EVENT_AP_STAIPASSIGNED)	n/a
IP_EVENT_GOT_IP6 (SYSTEM_EVENT_GOT_IP6)	`ip_event_got_ip6_t`
IP_EVENT_ETH_GOT_IP (SYSTEM_EVENT_ETH_GOT_IP)	`ip_event_got_ip_t`

4.5.2 Mesh Events

ESP-MESH uses a system similar to the *Legacy Event Loop* to deliver events to the application. See 系统事件 for details.

4.5.3 Bluetooth Events

Various modules of the Bluetooth stack deliver events to applications via dedicated callback functions. Callback functions receive the event type (enumerated value) and event data (union of structures for each event type). The

following list gives the registration API name, event enumeration type, and event parameters type.

- BLE GAP: `esp_ble_gap_register_callback()`, `esp_gap_ble_cb_event_t`, `esp_ble_gap_cb_param_t`.
- BT GAP: `esp_bt_gap_register_callback()`, `esp_bt_gap_cb_event_t`, `esp_bt_gap_cb_param_t`.
- GATTC: `esp_ble_gattc_register_callback()`, `esp_bt_gattc_cb_event_t`, `esp_bt_gattc_cb_param_t`.
- GATTS: `esp_ble_gatts_register_callback()`, `esp_bt_gatts_cb_event_t`, `esp_bt_gatts_cb_param_t`.
- SPP: `esp_spp_register_callback()`, `esp_spp_cb_event_t`, `esp_spp_cb_param_t`.
- Blufi: `esp_blufi_register_callbacks()`, `esp_blufi_cb_event_t`, `esp_blufi_cb_param_t`.
- A2DP: `esp_a2d_register_callback()`, `esp_a2d_cb_event_t`, `esp_a2d_cb_param_t`.
- AVRC: `esp_avrc_ct_register_callback()`, `esp_avrc_ct_cb_event_t`, `esp_avrc_ct_cb_param_t`.
- HFP Client: `esp_hf_client_register_callback()`, `esp_hf_client_cb_event_t`, `esp_hf_client_cb_param_t`.
- HFP AG: `esp_hf_ag_register_callback()`, `esp_hf_ag_cb_event_t`, `esp_hf_ag_cb_param_t`.

4.6 Deep Sleep Wake Stubs

ESP32-S2 supports running a "deep sleep wake stub" when coming out of deep sleep. This function runs immediately as soon as the chip wakes up - before any normal initialisation, bootloader, or ESP-IDF code has run. After the wake stub runs, the SoC can go back to sleep or continue to start ESP-IDF normally.

Deep sleep wake stub code is loaded into "RTC Fast Memory" and any data which it uses must also be loaded into RTC memory. RTC memory regions hold their contents during deep sleep.

4.6.1 Rules for Wake Stubs

Wake stub code must be carefully written:

- As the SoC has freshly woken from sleep, most of the peripherals are in reset states. The SPI flash is unmapped.
- The wake stub code can only call functions implemented in ROM or loaded into RTC Fast Memory (see below.)
- The wake stub code can only access data loaded in RTC memory. All other RAM will be unintiailised and have random contents. The wake stub can use other RAM for temporary storage, but the contents will be overwritten when the SoC goes back to sleep or starts ESP-IDF.
- RTC memory must include any read-only data (.rodata) used by the stub.
- Data in RTC memory is initialised whenever the SoC restarts, except when waking from deep sleep. When waking from deep sleep, the values which were present before going to sleep are kept.
- Wake stub code is a part of the main esp-idf app. During normal running of esp-idf, functions can call the wake stub functions or access RTC memory. It is as if these were regular parts of the app.

4.6.2 Implementing A Stub

The wake stub in esp-idf is called `esp_wake_deep_sleep()`. This function runs whenever the SoC wakes from deep sleep. There is a default version of this function provided in esp-idf, but the default function is weak-linked so if your app contains a function named `esp_wake_deep_sleep()` then this will override the default.

If supplying a custom wake stub, the first thing it does should be to call `esp_default_wake_deep_sleep()`.

It is not necessary to implement `esp_wake_deep_sleep()` in your app in order to use deep sleep. It is only necessary if you want to have special behaviour immediately on wake.

If you want to swap between different deep sleep stubs at runtime, it is also possible to do this by calling the `esp_set_deep_sleep_wake_stub()` function. This is not necessary if you only use the default `esp_wake_deep_sleep()` function.

All of these functions are declared in the `esp_sleep.h` header under components/esp32s2.

4.6.3 Loading Code Into RTC Memory

Wake stub code must be resident in RTC Fast Memory. This can be done in one of two ways.

The first way is to use the `RTC_IRAM_ATTR` attribute to place a function into RTC memory:

```
void RTC_IRAM_ATTR esp_wake_deep_sleep(void) {
    esp_default_wake_deep_sleep();
    // Add additional functionality here
}
```

The second way is to place the function into any source file whose name starts with `rtc_wake_stub`. Files names `rtc_wake_stub*` have their contents automatically put into RTC memory by the linker.

The first way is simpler for very short and simple code, or for source files where you want to mix "normal" and "RTC" code. The second way is simpler when you want to write longer pieces of code for RTC memory.

4.6.4 Loading Data Into RTC Memory

Data used by stub code must be resident in RTC memory. The data can be placed in RTC Fast memory or in RTC Slow memory which is also used by the ULP.

Specifying this data can be done in one of two ways:

The first way is to use the `RTC_DATA_ATTR` and `RTC_RODATA_ATTR` to specify any data (writeable or read-only, respectively) which should be loaded into RTC memory:

```
RTC_DATA_ATTR int wake_count;

void RTC_IRAM_ATTR esp_wake_deep_sleep(void) {
    esp_default_wake_deep_sleep();
    static RTC_RODATA_ATTR const char fmt_str[] = "Wake count %d\n";
    ets_printf(fmt_str, wake_count++);
}
```

The attributes `RTC_FAST_ATTR` and `RTC_SLOW_ATTR` can be used to specify data that will be force placed into RTC_FAST and RTC_SLOW memory respectively. Any access to data marked with `RTC_FAST_ATTR` is allowed by PRO_CPU only and it is responsibility of user to make sure about it.

Unfortunately, any string constants used in this way must be declared as arrays and marked with RTC_RODATA_ATTR, as shown in the example above.

The second way is to place the data into any source file whose name starts with `rtc_wake_stub`.

For example, the equivalent example in `rtc_wake_stub_counter.c`:

```
int wake_count;

void RTC_IRAM_ATTR esp_wake_deep_sleep(void) {
    esp_default_wake_deep_sleep();
    ets_printf("Wake count %d\n", wake_count++);
}
```

The second way is a better option if you need to use strings, or write other more complex code.

To reduce wake-up time use the *CONFIG_BOOTLOADER_SKIP_VALIDATE_IN_DEEP_SLEEP* Kconfig option, see more information in *Fast boot from Deep Sleep*.

4.7 Device Firmware Upgrade through USB

Device Firmware Upgrade (DFU) is a mechanism for upgrading the firmware of devices through Universal Serial Bus (USB). DFU is supported by ESP32-S2 chips. The necessary connections for the USB peripheral are shown in the following table.

GPIO	USB
20	D+ (green)
19	D- (white)
GND	GND (black)
+5V	+5V (red)

注解： The ESP32-S2 chip needs to be in bootloader mode for the detection as a DFU device and flashing. This can be achieved by pulling GPIO0 down (e.g. pressing the BOOT button), pulsing RESET down for a moment and releasing GPIO0.

警告： Some cables are wired up with non-standard colors and some drivers are able to work with swapped D+ and D- connections. Please try to swap the cables connecting to D+ and D- if your device is not detected.

The software requirements of DFU are included in 第一步：安装准备 of the Getting Started Guide.

Section *Building the DFU Image* describes how to build firmware for DFU with ESP-IDF and Section *Flashing the Chip with the DFU Image* deals with flashing the firmware.

4.7.1 Building the DFU Image

The DFU image can be created by running:

```
idf.py dfu
```

which creates `dfu.bin` in the build directory.

注解： Don't forget to set the target chip by `idf.py set-target` before running `idf.py dfu`. Otherwise, you might create an image for a different chip or receive an error message something like `unknown target 'dfu'`.

4.7.2 Flashing the Chip with the DFU Image

The DFU image is downloaded into the chip by running:

```
idf.py dfu-flash
```

which relies on dfu-util. Please see 第一步：安装准备 for installing `dfu-util`. `dfu-util` needs additional setup for *USB drivers (Windows only)* or setting up an *udev rule (Linux only)*. Mac OS users should be able to use `dfu-util` without further setup.

See *Common errors and known issues* and their solutions.

udev rule (Linux only)

udev is a device manager for the Linux kernel. It allows us to run `dfu-util` (and `idf.py dfu-flash`) without `sudo` for gaining access to the chip.

Create file `/etc/udev/rules.d/40-dfuse.rules` with the following content:

```
SUBSYSTEMS=="usb", ATTRS{idVendor}=="303a", ATTRS{idProduct}=="00??", GROUP=
↪"plugdev", MODE="0666"
```

注解： Please check the output of command `groups`. The user has to be a member of the *GROUP* specified above. You may use some other existing group for this purpose (e.g. *uucp* on some systems instead of *plugdev*) or create a new group for this purpose.

Restart your computer so the previous setting could take into affect or run `sudo udevadm trigger` to force manually udev to trigger your new rule.

USB drivers (Windows only)

`dfu-util` uses *libusb* to access the device. You have to register on Windows the device with the *WinUSB* driver. Please see the libusb wiki for more details.

The drivers can be installed by the Zadig tool. Please make sure that the device is in download mode before running the tool and that it detects the ESP32-S2 device before installing the drivers. The Zadig tool might detect several USB interfaces of ESP32-S2. Please install the WinUSB driver for only that interface for which there is no driver installed (probably it is Interface 2) and don't re-install the driver for the other interface.

警告： The manual installation of the driver in Device Manager of Windows is not recommended because the flashing might not work properly.

Common errors and known issues

- `dfu-util: command not found` might indicate that the tool hasn't been installed or is not available from the terminal. An easy way of checking the tool is running `dfu-util --version`. Please see 第一步：安装准备 for installing `dfu-util`.
- The reason for `No DFU capable USB device available` could be that the USB driver wasn't properly installed on Windows (see *USB drivers (Windows only)*), udev rule was not setup on Linux (see *udev rule (Linux only)*) or the device isn't in bootloader mode.
- Flashing with `dfu-util` on Windows fails on the first attempt with error `Lost device after RESET?`. Please retry the flashing and it should succeed the next time.

4.8 Core Dump

4.8.1 Overview

注解： The python utility does not currently support ESP32-S2

ESP-IDF provides support to generate core dumps on unrecoverable software errors. This useful technique allows post-mortem analysis of software state at the moment of failure. Upon the crash system enters panic state, prints some information and halts or reboots depending configuration. User can choose to generate core dump in order to analyse the reason of failure on PC later on. Core dump contains snapshots of all tasks in the system at the moment of failure.

Snapshots include tasks control blocks (TCB) and stacks. So it is possible to find out what task, at what instruction (line of code) and what callstack of that task lead to the crash. ESP-IDF provides special script *espcoredump.py* to help users to retrieve and analyse core dumps. This tool provides two commands for core dumps analysis:

- info_corefile - prints crashed task's registers, callstack, list of available tasks in the system, memory regions and contents of memory stored in core dump (TCBs and stacks)
- dbg_corefile - creates core dump ELF file and runs GDB debug session with this file. User can examine memory, variables and tasks states manually. Note that since not all memory is saved in core dump only values of variables allocated on stack will be meaningfull

For more information about core dump internals see the - Core dump internals

4.8.2 Configuration

There are a number of core dump related configuration options which user can choose in project configuration menu (*idf.py menuconfig*).

1. Core dump data destination (*Components -> Core dump -> Data destination*):

- Save core dump to Flash (Flash)
- Print core dump to UART (UART)
- Disable core dump generation (None)

2. Core dump data format (*Components -> Core dump -> Core dump data format*):

- ELF format (Executable and Linkable Format file for core dump)
- Binary format (Basic binary format for core dump)

The ELF format contains extended features and allow to save more information about broken tasks and crashed software but it requires more space in the flash memory. It also stores SHA256 of crashed application image. This format of core dump is recommended for new software designs and is flexible enough to extend saved information for future revisions. The Binary format is kept for compatibility standpoint, it uses less space in the memory to keep data and provides better performance.

3. Maximum number of tasks snapshots in core dump (*Components -> Core dump -> Maximum number of tasks*).
4. Delay before core dump is printed to UART (*Components -> Core dump -> Delay before print to UART*). Value is in ms.
5. Type of data integrity check for core dump (*Components -> Core dump -> Core dump data integrity check*).

- Use CRC32 for core dump integrity verification
- Use SHA256 for core dump integrity verification

The SHA256 hash algorithm provides greater probability of detecting corruption than a CRC32 with multiple bit errors. The CRC32 option provides better calculation performance and consumes less memory for storage.

4.8.3 Save core dump to flash

When this option is selected core dumps are saved to special partition on flash. When using default partition table files which are provided with ESP-IDF it automatically allocates necessary space on flash, But if user wants to use its own layout file together with core dump feature it should define separate partition for core dump as it is shown below:

```
# Name,   Type, SubType, Offset,  Size
# Note: if you have increased the bootloader size, make sure to update the offsets␣
↪to avoid overlap
nvs,      data, nvs,     0x9000,  0x6000
phy_init, data, phy,     0xf000,  0x1000
factory,  app,  factory, 0x10000, 1M
coredump, data, coredump,,        64K
```

There are no special requrements for partition name. It can be choosen according to the user application needs, but partition type should be 'data' and sub-type should be 'coredump'. Also when choosing partition size note that core dump data structure introduces constant overhead of 20 bytes and per-task overhead of 12 bytes. This overhead

does not include size of TCB and stack for every task. So partirion size should be at least 20 + max tasks number x (12 + TCB size + max task stack size) bytes.

The example of generic command to analyze core dump from flash is: *espcoredump.py -p </path/to/serial/port> info_corefile </path/to/program/elf/file>* or *espcoredump.py -p </path/to/serial/port> dbg_corefile </path/to/program/elf/file>*

4.8.4 Print core dump to UART

When this option is selected base64-encoded core dumps are printed on UART upon system panic. In this case user should save core dump text body to some file manually and then run the following command: *espcoredump.py info_corefile -t b64 -c </path/to/saved/base64/text> </path/to/program/elf/file>* or *espcoredump.py dbg_corefile -t b64 -c </path/to/saved/base64/text> </path/to/program/elf/file>*

Base64-encoded body of core dump will be between the following header and footer:

```
================ CORE DUMP START =================
<body of base64-encoded core dump, save it to file on disk>
================= CORE DUMP END ==================
```

The *CORE DUMP START* and *CORE DUMP END* lines must not be included in core dump text file.

4.8.5 ROM Functions in Backtraces

It is possible situation that at the moment of crash some tasks or/and crashed task itself have one or more ROM functions in their callstacks. Since ROM is not part of the program ELF it will be impossible for GDB to parse such callstacks, because it tries to analyse functions' prologues to acomplish that. In that case callstack printing will be broken with error message at the first ROM function. To overcome this issue you can use ROM ELF provided by Espressif (https://dl.espressif.com/dl/esp32_rom.elf) and pass it to 'espcoredump.py'.

4.8.6 Running 'espcoredump.py'

Generic command syntax:

espcoredump.py [options] command [args]

Script Options

- –chip,-c {auto,esp32}. Target chip type. Supported values are *auto* and *esp32*.
- –port,-p PORT. Serial port device.
- –baud,-b BAUD. Serial port baud rate used when flashing/reading.

Commands

- info_corefile. Retrieve core dump and print useful info.
- dbg_corefile. Retrieve core dump and start GDB session with it.

Command Arguments

- –debug,-d DEBUG. Log level (0..3).
- –gdb,-g GDB. Path to gdb to use for data retrieval.
- –core,-c CORE. Path to core dump file to use (if skipped core dump will be read from flash).
- –core-format,-t CORE_FORMAT. Specifies that file passed with "-c" is an ELF ("elf"), dumped raw binary ("raw") or base64-encoded ("b64") format.
- –off,-o OFF. Offset of coredump partition in flash (type *idf.py partition_table* to see it).
- –save-core,-s SAVE_CORE. Save core to file. Othwerwise temporary core file will be deleted. Ignored with "-c".
- –rom-elf,-r ROM_ELF. Path to ROM ELF file to use (if skipped "esp32_rom.elf" is used).
- –print-mem,-m Print memory dump. Used only with "info_corefile".
- <prog> Path to program ELF file.

4.9 Flash 加密

本文档将介绍 ESP32-S2 的 Flash 加密功能，并通过示例展示用户如何在开发及生产过程中使用此功能。本文档旨在引导用户快速入门如何测试及验证 Flash 加密的相关操作。有关 Flash 加密块的详细信息可参见 ESP32-S2 技术参考手册。

4.9.1 概述

Flash 加密功能用于加密与 ESP32-S2 搭载使用的 SPI Flash 中的内容。启用 Flash 加密功能后，物理读取 SPI Flash 便无法恢复大部分 Flash 内容。通过明文数据烧录 ESP32-S2 可应用加密功能，（若已启用加密功能）引导加载程序会在首次启动时对数据进行加密。

启用 Flash 加密后，系统将默认加密下列类型的 Flash 数据：

- 引导加载程序
- 分区表
- 所有“app”类型的分区

其他类型的 Flash 数据将视情况进行加密：

- 安全启动引导加载程序摘要（如果已启用安全启动）
- 分区表中标有“加密”标记的分区

重要： 启用 Flash 加密将限制后续 ESP32-S2 更新。请务必阅读本文档（包括*Flash 加密的局限性*）了解启用 Flash 加密的影响。

4.9.2 Flash 加密过程中使用的 eFuse

Flash 加密操作由 ESP32-S2 上的多个 eFuse 控制。以下是这些 eFuse 列表及其描述：

eFuse	描述	是否可锁定为读取/写入	默认值
编码方案	该 2 位宽 eFuse 控制 BLOCK1 中使用的实际位数，从而获得最终的 256 位 AES 密钥。编码方案值解码如下： 0: 256 bits 1: 192 bits 2: 128 bits 最终的 AES 密钥根据 FLASH_CRYPT_CONFIG 的值产生。	是	0
BLOCK1	存储 AES 密钥的 256 位宽 eFuse 块	是	x
FLASH_CRYPT_CONFIG	4 位宽 eFuse，控制 AES 加密进程	是	0xF
download_dis_encrypt	设置后，在 UART 下载模式运行时关闭 Flash 加密操作	是	0
download_dis_decrypt	设置后，在 UART 下载模式运行时关闭 Flash 解密操作	是	0

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```
FLASH_CRYPT_CNT        7 位 eFuse，在启动时启用/关闭加密功能                    是
↪           0

                       偶数位 (0, 2, 4, 6):
                       启动时加密 Flash
                       奇数位 (1, 3, 5, 7):
                       启动时不加密 Flash
```

对上述位的读写访问由 `efuse_wr_disable` 和 `efuse_rd_disable` 寄存器中的相应位控制。有关 ESP32-S2 eFuse 的详细信息可参见*eFuse 管理器*。

4.9.3 Flash 的加密过程

假设 eFuse 值处于默认状态，且第二阶段的引导加载程序编译为支持 Flash 加密，则 Flash 加密过程执行如下：

- 首次上电复位时，Flash 中的所有数据都是未加密形式（明文数据）。第一阶段加载器 (Rom) 将在 IRAM 中加载第二阶段加载器。
- 第二阶段引导加载程序将读取 flash_crypt_cnt (=00000000b) eFuse 值，因为该值为 0（偶数位），第二阶段引导加载程序将配置并启用 Flash 加密块，同时将 `FLASH_CRYPT_CFG` eFuse 的值编程为 0xF。
- Flash 加密块将生成 AES-256 位密钥，并将其储存于 BLOCK1 eFuse 中。该操作在硬件中执行，软件将无法访问此密钥。
- 接着，Flash 加密块将加密 Flash 的内容（根据分区表的标记值）。原地加密可能会有耗时（取决于大分区的耗时）。
- 随后，第二阶段引导加载程序将在 flash_crypt_cnt (=00000001b) 中设置第一个可用位，从而标记已加密的 Flash 内容（偶数位）。
- 在*释放模式*下，第二阶段引导加载程序将把 `download_dis_encrypt`、`download_dis_decrypt` 和 `download_dis_cache` 的 eFuse 位改写为 1，防止 UART 引导加载程序解密 Flash 的内容。同时也将写保护 `FLASH_CRYPT_CNT` 的 eFuse 位。
- 在*开发模式*下，第二阶段引导加载程序将仅改写 `download_dis_decrypt` 和 `download_dis_cache` 的 eFuse 位，从而允许 UART 引导加载程序重新烧录加密的二进制文件。同时将不会写保护 `FLASH_CRYPT_CNT` 的 eFuse 位。
- 然后，第二阶段引导加载程序重启设备并开始执行加密映像，同时将透明解密 Flash 的内容并将其加载至 IRAM。

在开发阶段常需编写不同的明文 Flash 映像，以及测试 Flash 的加密过程。这要求 UART 下载模式能够根据需求不断加载新的明文映像。但是，在量产和生产过程中，出于安全考虑，UART 下载模式不应有权限访问 Flash 内容。因此需要有两种不同的 ESP32-S2 配置：一种用于开发，另一种用于生产。以下章节介绍了 Flash 加密的*开发模式*和*释放模式*及其使用指南。

重要： 顾名思义，开发模式应 **仅开发过程**使用，因为该模式可以修改和回读加密的 Flash 内容。

4.9.4 设置 Flash 加密的步骤

开发模式

可使用 ESP32-S2 内部生成的密钥或外部主机生成的密钥在开发中运行 Flash 加密。

使用 ESP32-S2 生成的 Flash 加密密钥

正如上文所说，*开发模式*允许用户使用 UART 下载模式多次下载明文映像。需完成以下步骤测试 Flash 加密过程：

- 确保您的 ESP32-S2 设备有*Flash 加密过程中使用的 eFuse* 中所示的 Flash 加密 eFuse 的默认设置。

- 可在 `$IDF_PATH/examples/security/flash_encryption` 文件夹中找到 Flash 加密的示例应用程序。该示例应用程序中有显示 Flash 加密的状态（启用或关闭）以及 `FLASH_CRYPT_CNT` eFuse 值。
- 在第二阶段引导加载程序中启用 Flash 加密支持。请前往*Project Configuration Menu*，选择 “Security Features”。
- 选择*Enable flash encryption on boot*。
- 默认设置模式为 **开发模式**。
- 在引导加载程序 config 下选择适当详细程度的日志。
- 保存配置并退出。

构建并烧录完整的映像包括：引导加载程序、分区表和 app。这些分区最初以未加密形式写入 Flash。

```
idf.py flash monitor
```

一旦烧录完成，设备将重置，在下次启动时，第二阶段引导加载程序将加密 Flash 的 app 分区，然后重置该分区。现在，示例应用程序将在运行时解密并执行命令。以下是首次启用 Flash 加密后 ESP32-S2 启动时的样例输出。

```
--- idf_monitor on /dev/cu.SLAB_USBtoUART 115200 ---
--- Quit: Ctrl+] | Menu: Ctrl+T | Help: Ctrl+T followed by Ctrl+H ---
ets Jun  8 2016 00:22:57

rst:0x1 (POWERON_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
configsip: 0, SPIWP:0xee
clk_drv:0x00,q_drv:0x00,d_drv:0x00,cs0_drv:0x00,hd_drv:0x00,wp_drv:0x00
mode:DIO, clock div:2
load:0x3fff0018,len:4
load:0x3fff001c,len:8452
load:0x40078000,len:13608
load:0x40080400,len:6664
entry 0x40080764
I (28) boot: ESP-IDF v4.0-dev-850-gc4447462d-dirty 2nd stage bootloader
I (29) boot: compile time 15:37:14
I (30) boot: Enabling RNG early entropy source...
I (35) boot: SPI Speed      : 40MHz
I (39) boot: SPI Mode       : DIO
I (43) boot: SPI Flash Size : 4MB
I (47) boot: Partition Table:
I (51) boot: ## Label            Usage          Type ST Offset   Length
I (58) boot:  0 nvs              WiFi data        01 02 0000a000 00006000
I (66) boot:  1 phy_init         RF data          01 01 00010000 00001000
I (73) boot:  2 factory          factory app      00 00 00020000 00100000
I (81) boot: End of partition table
I (85) esp_image: segment 0: paddr=0x00020020 vaddr=0x3f400020␣
↪size=0x0808c ( 32908) map
I (105) esp_image: segment 1: paddr=0x000280b4 vaddr=0x3ffb0000␣
↪size=0x01ea4 (  7844) load
I (109) esp_image: segment 2: paddr=0x00029f60 vaddr=0x40080000␣
↪size=0x00400 (  1024) load
0x40080000: _WindowOverflow4 at esp-idf/esp-idf/components/freertos/
↪xtensa_vectors.S:1778

I (114) esp_image: segment 3: paddr=0x0002a368 vaddr=0x40080400␣
↪size=0x05ca8 ( 23720) load
I (132) esp_image: segment 4: paddr=0x00030018 vaddr=0x400d0018␣
↪size=0x126a8 ( 75432) map
0x400d0018: _flash_cache_start at ??:?

I (159) esp_image: segment 5: paddr=0x000426c8 vaddr=0x400860a8␣
↪size=0x01f4c (  8012) load
0x400860a8: prvAddNewTaskToReadyList at esp-idf/esp-idf/components/
↪freertos/tasks.c:4561
```

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```
I (168) boot: Loaded app from partition at offset 0x20000
I (168) boot: Checking flash encryption...
I (168) flash_encrypt: Generating new flash encryption key...
I (187) flash_encrypt: Read & write protecting new key...
I (187) flash_encrypt: Setting CRYPT_CONFIG efuse to 0xF
W (188) flash_encrypt: Not disabling UART bootloader encryption
I (195) flash_encrypt: Disable UART bootloader decryption...
I (201) flash_encrypt: Disable UART bootloader MMU cache...
I (208) flash_encrypt: Disable JTAG...
I (212) flash_encrypt: Disable ROM BASIC interpreter fallback...
I (219) esp_image: segment 0: paddr=0x00001020 vaddr=0x3fff0018␣
↪size=0x00004 (     4)
I (227) esp_image: segment 1: paddr=0x0000102c vaddr=0x3fff001c␣
↪size=0x02104 (  8452)
I (239) esp_image: segment 2: paddr=0x00003138 vaddr=0x40078000␣
↪size=0x03528 ( 13608)
I (249) esp_image: segment 3: paddr=0x00006668 vaddr=0x40080400␣
↪size=0x01a08 (  6664)
I (657) esp_image: segment 0: paddr=0x00020020 vaddr=0x3f400020␣
↪size=0x0808c ( 32908) map
I (669) esp_image: segment 1: paddr=0x000280b4 vaddr=0x3ffb0000␣
↪size=0x01ea4 (  7844)
I (672) esp_image: segment 2: paddr=0x00029f60 vaddr=0x40080000␣
↪size=0x00400 (  1024)
0x40080000: _WindowOverflow4 at esp-idf/esp-idf/components/freertos/
↪xtensa_vectors.S:1778

I (676) esp_image: segment 3: paddr=0x0002a368 vaddr=0x40080400␣
↪size=0x05ca8 ( 23720)
I (692) esp_image: segment 4: paddr=0x00030018 vaddr=0x400d0018␣
↪size=0x126a8 ( 75432) map
0x400d0018: _flash_cache_start at ??:?

I (719) esp_image: segment 5: paddr=0x000426c8 vaddr=0x400860a8␣
↪size=0x01f4c (  8012)
0x400860a8: prvAddNewTaskToReadyList at esp-idf/esp-idf/components/
↪freertos/tasks.c:4561

I (722) flash_encrypt: Encrypting partition 2 at offset 0x20000...
I (13229) flash_encrypt: Flash encryption completed
I (13229) boot: Resetting with flash encryption enabled...
```

启用 Flash 加密后，在下次启动时输出将显示已启用 Flash 加密。

```
rst:0x1 (POWERON_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
configsip: 0, SPIWP:0xee
clk_drv:0x00,q_drv:0x00,d_drv:0x00,cs0_drv:0x00,hd_drv:0x00,wp_drv:0x00
mode:DIO, clock div:2
load:0x3fff0018,len:4
load:0x3fff001c,len:8452
load:0x40078000,len:13652
ho 0 tail 12 room 4
load:0x40080400,len:6664
entry 0x40080764
I (30) boot: ESP-IDF v4.0-dev-850-gc4447462d-dirty 2nd stage bootloader
I (30) boot: compile time 16:32:53
I (31) boot: Enabling RNG early entropy source...
I (37) boot: SPI Speed      : 40MHz
I (41) boot: SPI Mode       : DIO
I (45) boot: SPI Flash Size : 4MB
```

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```
I (49) boot: Partition Table:
I (52) boot: ## Label            Usage          Type ST Offset   Length
I (60) boot:  0 nvs              WiFi data        01 02 0000a000 00006000
I (67) boot:  1 phy_init         RF data          01 01 00010000 00001000
I (75) boot:  2 factory          factory app      00 00 00020000 00100000
I (82) boot: End of partition table
I (86) esp_image: segment 0: paddr=0x00020020 vaddr=0x3f400020␣
↪size=0x0808c ( 32908) map
I (107) esp_image: segment 1: paddr=0x000280b4 vaddr=0x3ffb0000␣
↪size=0x01ea4 (  7844) load
I (111) esp_image: segment 2: paddr=0x00029f60 vaddr=0x40080000␣
↪size=0x00400 (  1024) load
0x40080000: _WindowOverflow4 at esp-idf/esp-idf/components/freertos/
↪xtensa_vectors.S:1778

I (116) esp_image: segment 3: paddr=0x0002a368 vaddr=0x40080400␣
↪size=0x05ca8 ( 23720) load
I (134) esp_image: segment 4: paddr=0x00030018 vaddr=0x400d0018␣
↪size=0x126a8 ( 75432) map
0x400d0018: _flash_cache_start at ??:?

I (162) esp_image: segment 5: paddr=0x000426c8 vaddr=0x400860a8␣
↪size=0x01f4c (  8012) load
0x400860a8: prvAddNewTaskToReadyList at esp-idf/esp-idf/components/
↪freertos/tasks.c:4561

I (171) boot: Loaded app from partition at offset 0x20000
I (171) boot: Checking flash encryption...
I (171) flash_encrypt: flash encryption is enabled (3 plaintext flashes␣
↪left)
I (178) boot: Disabling RNG early entropy source...
I (184) cpu_start: Pro cpu up.
I (188) cpu_start: Application information:
I (193) cpu_start: Project name:     flash-encryption
I (198) cpu_start: App version:      v4.0-dev-850-gc4447462d-dirty
I (205) cpu_start: Compile time:     Jun 17 2019 16:32:52
I (211) cpu_start: ELF file SHA256:  8770c886bdf561a7...
I (217) cpu_start: ESP-IDF:          v4.0-dev-850-gc4447462d-dirty
I (224) cpu_start: Starting app cpu, entry point is 0x40080e4c
0x40080e4c: call_start_cpu1 at esp-idf/esp-idf/components/esp32s2/cpu_
↪start.c:265

I (0) cpu_start: App cpu up.
I (235) heap_init: Initializing. RAM available for dynamic allocation:
I (241) heap_init: At 3FFAE6E0 len 00001920 (6 KiB): DRAM
I (247) heap_init: At 3FFB2EC8 len 0002D138 (180 KiB): DRAM
I (254) heap_init: At 3FFE0440 len 00003AE0 (14 KiB): D/IRAM
I (260) heap_init: At 3FFE4350 len 0001BCB0 (111 KiB): D/IRAM
I (266) heap_init: At 40087FF4 len 0001800C (96 KiB): IRAM
I (273) cpu_start: Pro cpu start user code
I (291) cpu_start: Starting scheduler on PRO CPU.
I (0) cpu_start: Starting scheduler on APP CPU.

Sample program to check Flash Encryption
This is ESP32-S2 chip with 2 CPU cores, WiFi/BT/BLE, silicon revision 1,␣
↪4MB external flash
Flash encryption feature is enabled
Flash encryption mode is DEVELOPMENT
Flash in encrypted mode with flash_crypt_cnt = 1
Halting...
```

在此阶段，如果用户希望以加密格式将已修改的明文应用程序映像更新到 Flash 中，可使用以下命令：

```
idf.py encrypted-app-flash monitor
```

加密多重分区

如果所有分区都需以加密格式更新，则可使用以下命令：

```
idf.py encrypted-flash monitor
```

使用主机生成的 Flash 加密密钥

可在主机中预生成 Flash 加密密钥，并将其烧录到 ESP32-S2 的 eFuse 密钥块中。这样，无需明文 Flash 更新便可以在主机上预加密数据并将其烧录到 ESP32-S2 中。该功能允许在开发模式 和释放模式 modes 两模式下加密烧录。

- 确保您的 ESP32-S2 设备有*Flash 加密过程中使用的 eFuse* 中所示 Flash 加密 eFuse 的默认设置。
- 使用 espsecure.py 随机生成一个密钥:

```
espsecure.py generate_flash_encryption_key my_flash_encryption_key.bin
```

- 将该密钥烧录到设备上（一次性）。**该步骤须在第一次加密启动前完成**，否则 ESP32-S2 将随机生成一个软件无权限访问或修改的密钥:

```
espefuse.py --port PORT burn_key flash_encryption my_flash_encryption_key.bin
```

- 在第二阶段引导加载程序中启用 Flash 加密支持。请前往*Project Configuration Menu*，选择 “Security Features”。
- 选择*Enable flash encryption on boot*。
- 模式默认设置为 **开发模式**。
- 在引导加载程序 config 下选择适当详细程度的日志。
- 保存配置并退出。

构建并烧录完整的映像包括：引导加载程序、分区表和 app。这些分区最初以未加密形式写入 Flash

```
idf.py flash monitor
```

下次启动时，第二阶段引导加载程序将加密 Flash 的 app 分区并重置该分区。现在，示例应用程序将在运行时解密并执行命令。

在此阶段，如果用户希望将新的明文应用程序映像更新到 Flash，应调用以下命令

```
idf.py encrypted-app-flash monitor
```

如何以加密格式重新编程所有分区，可参考加密多重分区。

释放模式

在释放模式下，UART 引导加载程序无法执行 Flash 加密操作，**只能**使用 OTA 方案下载新的明文映像，该方案将在写入 Flash 前加密明文映像。

- 确保您的 ESP32-S2 设备有*Flash 加密过程中使用的 eFuse* 中所示 Flash 加密 eFuse 的默认设置。
- 在第二阶段引导加载程序中启用 Flash 加密支持。请前往*Project Configuration Menu*，选择 “Security Features”。
- 选择*Enable flash encryption on boot*。
- 选择 **释放模式**，模式默认设置为 **开发模式**。请注意，**一旦选择了释放模式，“download_dis_encrypt“ 和 “download_dis_decrypt“ eFuse 位将被编程为禁止 UART 引导加载程序访问 Flash 的内容**。
- 在引导加载程序 config 下选择适当详细程度的日志。
- 保存配置并退出。

构建并烧录完整的映像包括：引导加载程序、分区表和 app。这些分区最初以未加密形式写入 Flash

```
idf.py flash monitor
```

下次启动时，第二阶段引导加载程序将加密 Flash app 分区并重置该分区。现在，示例应用程序应正确执行命令。

一旦在释放模式下启用 Flash 加密，引导加载程序将写保护 `FLASH_CRYPT_CNT` eFuse。

应使用 OTA 方案对字段中的明文进行后续更新。详情可参见*OTA*。

可能出现的错误

启用 Flash 加密后，如果 `FLASH_CRYPT_CNT` eFuse 值中有奇数位，则所有（标有加密标志的）分区都应包含加密密文。以下为 ESP32-S2 加载明文数据会产生的三种典型错误情况：

1. 如果通过明文引导加载程序映像重新更新了引导加载程序分区，则 ROM 加载器将无法加载引导加载程序，并会显示以下错误类型：

```
rst:0x3 (SW_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
flash read err, 1000
ets_main.c 371
ets Jun  8 2016 00:22:57

rst:0x7 (TG0WDT_SYS_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
flash read err, 1000
ets_main.c 371
ets Jun  8 2016 00:22:57

rst:0x7 (TG0WDT_SYS_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
flash read err, 1000
ets_main.c 371
ets Jun  8 2016 00:22:57

rst:0x7 (TG0WDT_SYS_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
flash read err, 1000
ets_main.c 371
ets Jun  8 2016 00:22:57

rst:0x7 (TG0WDT_SYS_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
flash read err, 1000
ets_main.c 371
ets Jun  8 2016 00:22:57
```

2. 如果引导加载程序已加密，但使用明文分区表映像重新更新了分区表，则引导加载程序将无法读取分区表，并会显示以下错误类型：

```
rst:0x3 (SW_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
configsip: 0, SPIWP:0xee
clk_drv:0x00,q_drv:0x00,d_drv:0x00,cs0_drv:0x00,hd_drv:0x00,wp_drv:0x00
mode:DIO, clock div:2
load:0x3fff0018,len:4
load:0x3fff001c,len:10464
ho 0 tail 12 room 4
load:0x40078000,len:19168
load:0x40080400,len:6664
entry 0x40080764
I (60) boot: ESP-IDF v4.0-dev-763-g2c55fae6c-dirty 2nd stage bootloader
I (60) boot: compile time 19:15:54
I (62) boot: Enabling RNG early entropy source...
I (67) boot: SPI Speed      : 40MHz
I (72) boot: SPI Mode       : DIO
```

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```
I (76) boot: SPI Flash Size : 4MB
E (80) flash_parts: partition 0 invalid magic number 0x94f6
E (86) boot: Failed to verify partition table
E (91) boot: load partition table error!
```

3. 如果引导加载程序和分区表已加密，但使用明文应用程序映像重新更新了应用程序，则引导加载程序将无法加载新的应用程序，并会显示以下错误类型：

```
rst:0x3 (SW_RESET),boot:0x13 (SPI_FAST_FLASH_BOOT)
configsip: 0, SPIWP:0xee
clk_drv:0x00,q_drv:0x00,d_drv:0x00,cs0_drv:0x00,hd_drv:0x00,wp_drv:0x00
mode:DIO, clock div:2
load:0x3fff0018,len:4
load:0x3fff001c,len:8452
load:0x40078000,len:13616
load:0x40080400,len:6664
entry 0x40080764
I (56) boot: ESP-IDF v4.0-dev-850-gc4447462d-dirty 2nd stage bootloader
I (56) boot: compile time 15:37:14
I (58) boot: Enabling RNG early entropy source...
I (64) boot: SPI Speed      : 40MHz
I (68) boot: SPI Mode       : DIO
I (72) boot: SPI Flash Size : 4MB
I (76) boot: Partition Table:
I (79) boot: ## Label            Usage          Type ST Offset   Length
I (87) boot:  0 nvs              WiFi data        01 02 0000a000 00006000
I (94) boot:  1 phy_init         RF data          01 01 00010000 00001000
I (102) boot:  2 factory          factory app      00 00 00020000 00100000
I (109) boot: End of partition table
E (113) esp_image: image at 0x20000 has invalid magic byte
W (120) esp_image: image at 0x20000 has invalid SPI mode 108
W (126) esp_image: image at 0x20000 has invalid SPI size 11
E (132) boot: Factory app partition is not bootable
E (138) boot: No bootable app partitions in the partition table
```

4.9.5 Flash 加密的要点

- 使用 AES-256 加密 Flash 的内容。Flash 加密密钥存储于 eFuse 内部的芯片中，并（默认）受保护防止软件访问。
- *flash 加密算法*采用的是 AES-256，其中密钥随着 Flash 的每个 32 字节块的偏移地址“调整”。这意味着，每个 32 字节块（2 个连续的 16 字节 AES 块）使用从 Flash 加密密钥中产生的一个特殊密钥进行加密。
- 通过 ESP32-S2 的 Flash 缓存映射功能，Flash 可支持透明访问——读取任何映射到地址空间的 Flash 区域时，都将透明解密该区域。

  为便于访问，某些数据分区最好保持未加密状态，或者也可使用对已加密数据无效的 Flash 友好型更新算法。由于 NVS 库无法与 Flash 加密直接兼容，因此无法加密非易失性存储器的 NVS 分区。详情可参见*NVS 加密*。
- 如果可能已启用 Flash 加密，则编写*使用加密 flash* 的代码时，编程人员须小心谨慎。
- 如果已启用安全启动，则重新烧录加密设备的引导加载程序则需要“可重新烧录”的安全启动摘要（可参见*Flash 加密与安全启动*）。

重要： 在首次启动加密过程中，请勿中断 ESP32-S2 的电源。如果电源中断，Flash 的内容将受到破坏，并需要重新烧录未加密数据。而这类重新烧录将不计入烧录限制次数。

4.9.6 使用加密的 Flash

ESP32-S2 app 代码可通过调用函数*esp_flash_encryption_enabled()* 来确认当前是否已启用 Flash 加密。同时，设备可通过调用函数*esp_get_flash_encryption_mode()* 来识别使用的 Flash 加密模式。

启用 Flash 加密后，使用代码访问 Flash 内容时需加注意。

Flash 加密的范围

只要 FLASH_CRYPT_CNT eFuse 设置为奇数位的值，所有通过 MMU 的 Flash 缓存访问的 Flash 内容都将被透明解密。包括：

- Flash 中可执行的应用程序代码 (IROM)。
- 所有存储于 Flash 中的只读数据 (DROM)。
- 通过函数 spi_flash_mmap() 访问的任意数据。
- ROM 引导加载程序读取的软件引导加载程序映像。

重要： MMU Flash 缓存将无条件解密所有数据。Flash 中未加密存储的数据将通过 Flash 缓存“被透明解密”，并在软件中存储为随机垃圾数据。

读取加密的 Flash

如在不使用 Flash 缓存 MMU 映射的情况下读取数据，推荐使用分区读取函数*esp_partition_read()*。使用该函数时，只有从加密分区读取的数据才会被解密。其他分区的数据将以未加密形式读取。这样，软件便能同样访问加密和未加密的 Flash。

通过其他 SPI 读取 API 读取的数据均未解密：

- 通过函数 spi_flash_read() 读取的数据均未解密。
- 通过 ROM 函数 SPIRead() 读取的数据均未解密（esp-idf app 不支持该函数）。
- 使用非易失性存储器 (NVS) API 存储的数据始终从 Flash 加密的角度进行存储和读取解密。如有需要，则由库提供加密功能。详情可参见*NVS 加密*。

写入加密的 Flash

在可能的情况下，推荐使用分区写入函数 esp_partition_write。使用该函数时，只有向加密分区写入的数据才会被加密。而写入其他分区的数据均未加密。这样，软件便可同样访问加密和未加密的 Flash。

当 write_encrypted 参数设置为“是”时，函数 esp_spi_flash_write 将写入数据。否则，数据将以未加密形式写入。

ROM 函数 esp_rom_spiflash_write_encrypted 将在 Flash 中写入加密数据，而 ROM 函数 SPI_Write 将在 Flash 中写入未加密数据（esp-idf app 不支持上述函数）。

由于数据均采用块加密方式，加密数据最小的写入大小为 16 字节（16 字节对齐）。

4.9.7 更新加密的 Flash

OTA 更新

只要使用了函数 esp_partition_write，则加密分区的 OTA 更新将自动以加密形式写入。

4.9.8 关闭 Flash 加密

若因某些原因意外启用了 Flash 加密，则接下来烧录明文数据时将使 ESP32-S2 软砖（设备不断重启，并报错 `flash read err, 1000`）。

可通过写入 `FLASH_CRYPT_CNT` eFuse 再次关闭 Flash 加密（仅适用于开发模式下）：

- 首先，前往*Project Configuration Menu*，在“安全性能”目录下关闭*启用 Flash 加密启动*。
- 退出 menuconfig 并保存最新配置。
- 再次运行 `idf.py menuconfig` 并复核是否确认已关闭该选项！如果该选项仍处于已启用状态，则引导加载程序会在启动后立即重新启用加密。
- 在未启用 Flash 加密的状态下，运行 `idf.py flash` 构建并烧录新的引导加载程序与 app。
- 运行 `espefuse.py`（`components/esptool_py/esptool` 中）以关闭 FLASH_CRYPT_CNT:

```
espefuse.py burn_efuse FLASH_CRYPT_CNT
```

重置 ESP32-S2，Flash 加密应处于关闭状态，引导加载程序将正常启动。

4.9.9 Flash 加密的局限性

Flash 加密可防止从加密 Flash 中读取明文，从而保护固件防止未经授权的读取与修改。了解 Flash 加密系统的局限之处亦十分重要：

- Flash 加密功能与密钥同样稳固。因而，推荐您首次启动设备时在设备上生成密钥（默认行为）。如果在设备外生成密钥，请确保遵循正确的后续步骤。
- 并非所有数据都是加密存储。因而在 Flash 上存储数据时，请检查您使用的存储方式（库、API 等）是否支持 Flash 加密。
- Flash 加密无法防止攻击者获取 Flash 的高层次布局信息。这是因为同一个 AES 密钥要用于每对相邻的 16 字节 AES 块。当这些相邻的 16 字节块中包含相同内容时（如空白或填充区域），这些字节块将加密以产生匹配的加密块对。这可能使得攻击者可在加密设备间进行高层次对比（例如，确认两设备是否可能在运行相同的固件版本）。
- 出于相同原因，攻击者始终可获知一对相邻的 16 字节块（32 字节对齐）何时包含相同内容。因此，在 Flash 上存储敏感数据时应牢记这点，并进行相关设置避免该情况发生（可使用计数器字节或每 16 字节设置不同的值即可）。

4.9.10 Flash 加密与安全启动

推荐搭配使用 Flash 加密与安全启动。但是，如果已启用安全启动，则重新烧录设备时会受到其他限制：

- *OTA 更新* 不受限制（如果新的 app 已使用安全启动签名密钥进行正确签名）。

4.9.11 使用无安全启动的 Flash 加密

尽管 Flash 加密与安全启动可独立使用，但强烈建议您将这二者 **搭配使用**以确保更高的安全性。

4.9.12 Flash 加密的高级功能

以下信息可帮助您使用 Flash 加密的高级功能：

加密分区标志

部分分区默认为已加密。除此之外，可将任意分区标记为需加密：

在*分区表* 文档对 CSV 文件的描述中有标志字段。

该字段通常保留为空白。如果在字段中写入”encrypted”，则这个分区将在分区表中标记为已加密，此处写入的数据也视为加密数据（app 分区同样适用）：

```
# Name,   Type, SubType, Offset,  Size, Flags
nvs,      data, nvs,     0x9000,  0x6000
phy_init, data, phy,     0xf000,  0x1000
factory,  app,  factory, 0x10000, 1M
secret_data, 0x40, 0x01, 0x20000, 256K, encrypted
```

- 默认分区表都不包含任何加密数据分区。
- “app” 分区一般都视为加密分区，因此无需将其标记为已加密。
- 如果未启用 Flash 加密，则” encrypted” 标记无效。
- 可将带有 phy_init 数据的可选 phy 分区标记为已加密，保护该数据防止物理访问读取或修改。
- nvs 分区无法标记为已加密。

启用 UART 引导加载程序加密/解密

默认情况下，首次启动 Flash 加密过程中将烧录 eFuse DISABLE_DL_ENCRYPT、DISABLE_DL_DECRYPT 和 DISABLE_DL_CACHE：

- DISABLE_DL_ENCRYPT 在 UART 引导加载程序启动模式下运行时，终止 Flash 加密操作。
- DISABLE_DL_DECRYPT 在 UART 引导加载程序模式下运行时，终止透明 Flash 解密（即使 FLASH_CRYPT_CNT 已设置为在正常操作中启用 Flash 透明解密）。
- DISABLE_DL_CACHE 在 UART 引导加载程序模式下运行时终止整个 MMU flash 缓存。

为了完整保存数据，可在首次启动前仅烧录部分 eFuse，并写保护其他部分（未设置值为 0）。例如:

```
espefuse.py --port PORT burn_efuse DISABLE_DL_DECRYPT
espefuse.py --port PORT write_protect_efuse DISABLE_DL_ENCRYPT
```

（请注意，一个写保护位即可关闭这 3 个 eFuse，因此，写保护一个 eFuse 将写保护上述所有 eFuse。所以，在写保护前须设置任意位）。

重要： 由于 esptool.py 不支持读取加密的 Flash，因此目前基本无法通过写保护这些 eFuse 来将其保持为未设状态。

重要： 如果保留 DISABLE_DL_DECRYPT 未设置（为 0），则实际上将使 Flash 加密无效，因为此时有物理访问权限的攻击者便可使用 UART 引导加载程序模式（使用自定义存根代码）读取 Flash 的内容。

设置 FLASH_CRYPT_CONFIG

FLASH_CRYPT_CONFIG eFuse 决定 Flash 加密密钥中随块偏移 “调整” 的位数。详情可参见*Flash 加密算法*。

首次启动引导加载程序时，该值始终设置为最大 *0xF*。

可手动写入这些 eFuse，并在首次启动前对其写保护，以便选择不同的调整值。但不推荐该操作。

当 FLASH_CRYPT_CONFIG 的值为 0 时，强烈建议始终不对其进行写保护。如果该 eFuse 设置为 0，则 Flash 加密密钥中无调整位，且 Flash 加密算法相当于 AES ECB 模式。

4.9.13 技术细节

下节将提供 Flash 加密操作的相关信息。

Flash 加密算法

- AES-256 在 16 字节的数据块上运行。Flash 加密引擎在 32 字节的数据块和 2 个串行 AES 块上加密或解密数据。
- Flash 加密的主密钥存储于 eFuse (BLOCK1) 中，默认受保护防止进一步写入或软件读取。
- AES-256 密钥大小为 256 位（32 字节），从 eFuse block 1 中读取。硬件 AES 引擎使用反字节序密钥于 eFuse 块中存储的字节序。
  - 如果 `CODING_SCHEME` eFuse 设置为 0（默认“无”编码方案），则 eFuse 密钥块为 256 位，且密钥按原方式存储（反字节序）。
  - 如果 `CODING_SCHEME` eFuse 设置为 1（3/4 编码），则 eFuse 密钥块为 192 位（反字节序），信息熵总量减少。硬件 Flash 加密仍在 256 字节密钥上运行，在读取后（字节序未反向），密钥扩展为 `key = key[0:255] + key[64:127]`。
- Flash 加密中使用了逆向 AES 算法，因此 Flash 加密的“加密”操作相当于 AES 解密，而其“解密”操作则相当于 AES 加密。这是为了优化性能，不会影响算法的有效性。
- 每个 32 字节块（2 个相邻的 16 字节 AES 块）都由一个特殊的密钥进行加密。该密钥由 eFuse 中 Flash 加密的主密钥产生，并随 Flash 中该字节块的偏移进行 XOR 运算（一次“密钥调整”）。
- 具体调整量取决于 `FLASH_CRYPT_CONFIG` eFuse 的设置。该 eFuse 共 4 位，每位可对特定范围的密钥位进行 XOR 运算：
  - Bit 1，对密钥的 0-66 位进行 XOR 运算。
  - Bit 2，对密钥的 67-131 位进行 XOR 运算。
  - Bit 3，对密钥的 132-194 位进行 XOR 运算。
  - Bit 4，对密钥的 195-256 位进行 XOR 运算。

  建议将 `FLASH_CRYPT_CONFIG` 的值始终保留为默认值 *0xF*，这样所有密钥位都随块偏移进行 XOR 运算。详情可参见设置 *FLASH_CRYPT_CONFIG*。
- 块偏移的 19 个高位（第 5-23 位）由 Flash 加密的主密钥进行 XOR 运算。选定该范围的原因为：Flash 的最大尺寸为 16MB（24 位），每个块大小为 32 字节，因而 5 个最低有效位始终为 0。
- 从 19 个块偏移位中每个位到 Flash 加密密钥的 256 位都有一个特殊的映射，以决定与哪个位进行 XOR 运算。有关完整映射可参见 `espsecure.py` 源代码中的变量 `_FLASH_ENCRYPTION_TWEAK_PATTERN`。
- 有关在 Python 中实现的完整 Flash 加密算法，可参见 `espsecure.py` 源代码中的函数 *_flash_encryption_operation()*。

4.10 ESP-IDF FreeRTOS SMP Changes

4.10.1 Overview

The ESP-IDF FreeRTOS is a modified version of vanilla FreeRTOS which supports symmetric multiprocessing (SMP). ESP-IDF FreeRTOS is based on the Xtensa port of FreeRTOS v8.2.0. This guide outlines the major differences between vanilla FreeRTOS and ESP-IDF FreeRTOS. The API reference for vanilla FreeRTOS can be found via https://www.freertos.org/a00106.html

For information regarding features that are exclusive to ESP-IDF FreeRTOS, see *ESP-IDF FreeRTOS Additions*.

Backported Features: Although ESP-IDF FreeRTOS is based on the Xtensa port of FreeRTOS v8.2.0, a number of FreeRTOS v9.0.0 features have been backported to ESP-IDF.

Task Deletion: Task deletion behavior has been backported from FreeRTOS v9.0.0 and modified to be SMP compatible. Task memory will be freed immediately when *`vTaskDelete()`* is called to delete a task that is not currently running and not pinned to the other core. Otherwise, freeing of task memory will still be delegated to the Idle Task.

Thread Local Storage Pointers & Deletion Callbacks: ESP-IDF FreeRTOS has backported the Thread Local Storage Pointers (TLSP) feature. However the extra feature of Deletion Callbacks has been added. Deletion callbacks are called automatically during task deletion and are used to free memory pointed to by TLSP. Call *`vTaskSetThreadLocalStoragePointerAndDelCallback()`* to set TLSP and Deletion Callbacks.

Configuring ESP-IDF FreeRTOS: Several aspects of ESP-IDF FreeRTOS can be set in the project configuration (`idf.py menuconfig`) such as running ESP-IDF in Unicore (single core) Mode, or configuring the number of Thread Local Storage Pointers each task will have.

4.10.2 Backported Features

The following features have been backported from FreeRTOS v9.0.0 to ESP-IDF.

Static Alocation

This feature has been backported from FreeRTOS v9.0.0 to ESP-IDF. The *CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION* option must be enabled in *menuconfig* in order for static allocation functions to be available. Once enabled, the following functions can be called…

- `xTaskCreateStatic()` (see *Backporting Notes* below)
- `xQueueCreateStatic`
- `xSemaphoreCreateBinaryStatic`
- `xSemaphoreCreateCountingStatic`
- `xSemaphoreCreateMutexStatic`
- `xSemaphoreCreateRecursiveMutexStatic`
- `xTimerCreateStatic()` (see *Backporting Notes* below)
- `xEventGroupCreateStatic()`

Other Features

- `vTaskSetThreadLocalStoragePointer()` (see *Backporting Notes* below)
- `pvTaskGetThreadLocalStoragePointer()` (see *Backporting Notes* below)
- `vTimerSetTimerID()`
- `xTimerGetPeriod()`
- `xTimerGetExpiryTime()`
- `pcQueueGetName()`
- `uxSemaphoreGetCount`

Backporting Notes

1) `xTaskCreateStatic()` has been made SMP compatible in a similar fashion to `xTaskCreate()` (see *Tasks and Task Creation*). Therefore `xTaskCreateStaticPinnedToCore()` can also be called.

2) Although vanilla FreeRTOS allows the Timer feature's daemon task to be statically allocated, the daemon task is always dynamically allocated in ESP-IDF. Therefore `vApplicationGetTimerTaskMemory` **does not** need to be defined when using statically allocated timers in ESP-IDF FreeRTOS.

3) The Thread Local Storage Pointer feature has been modified in ESP-IDF FreeRTOS to include Deletion Callbacks (see *Thread Local Storage Pointers & Deletion Callbacks*). Therefore the function `vTaskSetThreadLocalStoragePointerAndDelCallback()` can also be called.

4.10.3 Tasks and Task Creation

Tasks in ESP-IDF FreeRTOS are designed to run on a particular core, therefore two new task creation functions have been added to ESP-IDF FreeRTOS by appending `PinnedToCore` to the names of the task creation functions in vanilla FreeRTOS. The vanilla FreeRTOS functions of `xTaskCreate()` and `xTaskCreateStatic()` have led to the addition of `xTaskCreatePinnedToCore()` and `xTaskCreateStaticPinnedToCore()` in ESP-IDF FreeRTOS (see *Backported Features*).

For more details see freertos/tasks.c

The ESP-IDF FreeRTOS task creation functions are nearly identical to their vanilla counterparts with the exception of the extra parameter known as `xCoreID`. This parameter specifies the core on which the task should run on and can be one of the following values.

- `0` pins the task to **PRO_CPU**
- `1` pins the task to **APP_CPU**

- `tskNO_AFFINITY` allows the task to be run on both CPUs

For example `xTaskCreatePinnedToCore(tsk_callback, "APP_CPU Task", 1000, NULL, 10, NULL, 1)` creates a task of priority 10 that is pinned to **APP_CPU** with a stack size of 1000 bytes. It should be noted that the `uxStackDepth` parameter in vanilla FreeRTOS specifies a task's stack depth in terms of the number of words, whereas ESP-IDF FreeRTOS specifies the stack depth in terms of bytes.

Note that the vanilla FreeRTOS functions *`xTaskCreate()`* and *`xTaskCreateStatic()`* have been defined in ESP-IDF FreeRTOS as inline functions which call *`xTaskCreatePinnedToCore()`* and *`xTaskCreateStaticPinnedToCore()`* respectively with `tskNO_AFFINITY` as the `xCoreID` value.

Each Task Control Block (TCB) in ESP-IDF stores the `xCoreID` as a member. Hence when each core calls the scheduler to select a task to run, the `xCoreID` member will allow the scheduler to determine if a given task is permitted to run on the core that called it.

4.10.4 Scheduling

The vanilla FreeRTOS implements scheduling in the `vTaskSwitchContext()` function. This function is responsible for selecting the highest priority task to run from a list of tasks in the Ready state known as the Ready Tasks List (described in the next section). In ESP-IDF FreeRTOS, each core will call `vTaskSwitchContext()` independently to select a task to run from the Ready Tasks List which is shared between both cores. There are several differences in scheduling behavior between vanilla and ESP-IDF FreeRTOS such as differences in Round Robin scheduling, scheduler suspension, and tick interrupt synchronicity.

Round Robin Scheduling

Given multiple tasks in the Ready state and of the same priority, vanilla FreeRTOS implements Round Robin scheduling between each task. This will result in running those tasks in turn each time the scheduler is called (e.g. every tick interrupt). On the other hand, the ESP-IDF FreeRTOS scheduler may skip tasks when Round Robin scheduling multiple Ready state tasks of the same priority.

The issue of skipping tasks during Round Robin scheduling arises from the way the Ready Tasks List is implemented in FreeRTOS. In vanilla FreeRTOS, `pxReadyTasksList` is used to store a list of tasks that are in the Ready state. The list is implemented as an array of length `configMAX_PRIORITIES` where each element of the array is a linked list. Each linked list is of type `List_t` and contains TCBs of tasks of the same priority that are in the Ready state. The following diagram illustrates the `pxReadyTasksList` structure.

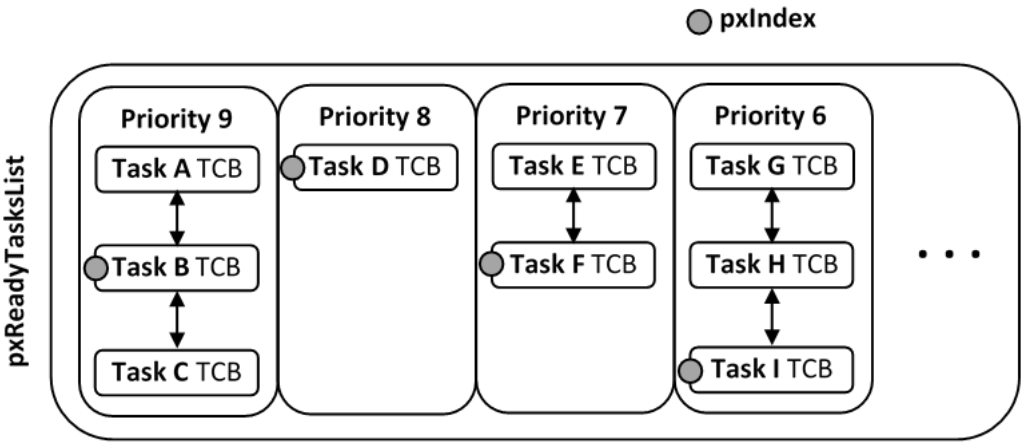

图 3: Illustration of FreeRTOS Ready Task List Data Structure

Each linked list also contains a `pxIndex` which points to the last TCB returned when the list was queried. This index allows the `vTaskSwitchContext()` to start traversing the list at the TCB immediately after `pxIndex` hence implementing Round Robin Scheduling between tasks of the same priority.

In ESP-IDF FreeRTOS, the Ready Tasks List is shared between cores hence `pxReadyTasksList` will contain tasks pinned to different cores. When a core calls the scheduler, it is able to look at the `xCoreID` member of each TCB in the list to determine if a task is allowed to run on calling the core. The ESP-IDF FreeRTOS `pxReadyTasksList` is illustrated below.

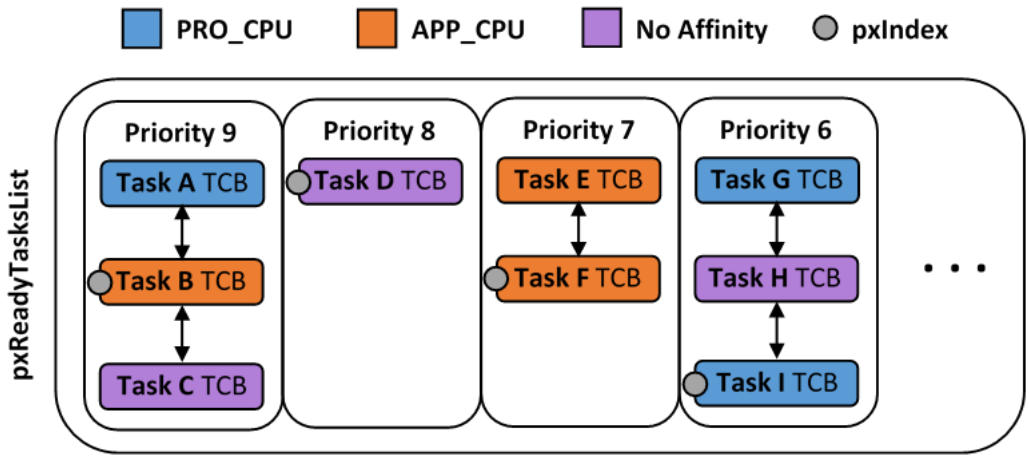

图 4: Illustration of FreeRTOS Ready Task List Data Structure in ESP-IDF

Therefore when **PRO_CPU** calls the scheduler, it will only consider the tasks in blue or purple. Whereas when **APP_CPU** calls the scheduler, it will only consider the tasks in orange or purple.

Although each TCB has an `xCoreID` in ESP-IDF FreeRTOS, the linked list of each priority only has a single `pxIndex`. Therefore when the scheduler is called from a particular core and traverses the linked list, it will skip all TCBs pinned to the other core and point the pxIndex at the selected task. If the other core then calls the scheduler, it will traverse the linked list starting at the TCB immediately after `pxIndex`. Therefore, TCBs skipped on the previous scheduler call from the other core would not be considered on the current scheduler call. This issue is demonstrated in the following illustration.

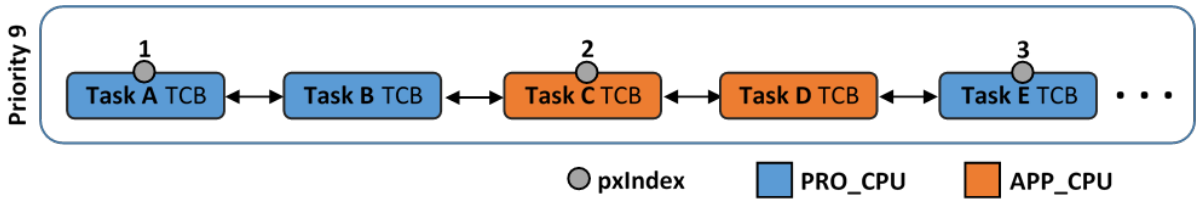

图 5: Illustration of pxIndex behavior in ESP-IDF FreeRTOS

Referring to the illustration above, assume that priority 9 is the highest priority, and none of the tasks in priority 9 will block hence will always be either in the running or Ready state.

1) **PRO_CPU** calls the scheduler and selects Task A to run, hence moves `pxIndex` to point to Task A

2) **APP_CPU** calls the scheduler and starts traversing from the task after `pxIndex` which is Task B. However Task B is not selected to run as it is not pinned to **APP_CPU** hence it is skipped and Task C is selected instead. `pxIndex` now points to Task C

3) **PRO_CPU** calls the scheduler and starts traversing from Task D. It skips Task D and selects Task E to run and points `pxIndex` to Task E. Notice that Task B isn't traversed because it was skipped the last time **APP_CPU** called the scheduler to traverse the list.

4) The same situation with Task D will occur if **APP_CPU** calls the scheduler again as `pxIndex` now points to Task E

One solution to the issue of task skipping is to ensure that every task will enter a blocked state so that they are removed from the Ready Task List. Another solution is to distribute tasks across multiple priorities such that a given priority will not be assigned multiple tasks that are pinned to different cores.

Scheduler Suspension

In vanilla FreeRTOS, suspending the scheduler via *vTaskSuspendAll()* will prevent calls of `vTaskSwitchContext` from context switching until the scheduler has been resumed with *xTaskResumeAll()*. However servicing ISRs are still permitted. Therefore any changes in task states as a result from the current running task or ISRSs will not be executed until the scheduler is resumed. Scheduler suspension in vanilla FreeRTOS is a common protection method against simultaneous access of data shared between tasks, whilst still allowing ISRs to be serviced.

In ESP-IDF FreeRTOS, `xTaskSuspendAll()` will only prevent calls of `vTaskSwitchContext()` from switching contexts on the core that called for the suspension. Hence if **PRO_CPU** calls *vTaskSuspendAll()*, **APP_CPU** will still be able to switch contexts. If data is shared between tasks that are pinned to different cores, scheduler suspension is **NOT** a valid method of protection against simultaneous access. Consider using critical sections (disables interrupts) or semaphores (does not disable interrupts) instead when protecting shared resources in ESP-IDF FreeRTOS.

In general, it's better to use other RTOS primitives like mutex semaphores to protect against data shared between tasks, rather than *vTaskSuspendAll()*.

Tick Interrupt Synchronicity

In ESP-IDF FreeRTOS, tasks on different cores that unblock on the same tick count might not run at exactly the same time due to the scheduler calls from each core being independent, and the tick interrupts to each core being unsynchronized.

In vanilla FreeRTOS the tick interrupt triggers a call to `xTaskIncrementTick()` which is responsible for incrementing the tick counter, checking if tasks which have called *vTaskDelay()* have fulfilled their delay period, and moving those tasks from the Delayed Task List to the Ready Task List. The tick interrupt will then call the scheduler if a context switch is necessary.

In ESP-IDF FreeRTOS, delayed tasks are unblocked with reference to the tick interrupt on PRO_CPU as PRO_CPU is responsible for incrementing the shared tick count. However tick interrupts to each core might not be synchronized (same frequency but out of phase) hence when PRO_CPU receives a tick interrupt, APP_CPU might not have received it yet. Therefore if multiple tasks of the same priority are unblocked on the same tick count, the task pinned to PRO_CPU will run immediately whereas the task pinned to APP_CPU must wait until APP_CPU receives its out of sync tick interrupt. Upon receiving the tick interrupt, APP_CPU will then call for a context switch and finally switches contexts to the newly unblocked task.

Therefore, task delays should **NOT** be used as a method of synchronization between tasks in ESP-IDF FreeRTOS. Instead, consider using a counting semaphore to unblock multiple tasks at the same time.

4.10.5 Critical Sections & Disabling Interrupts

Vanilla FreeRTOS implements critical sections in `vTaskEnterCritical` which disables the scheduler and calls `portDISABLE_INTERRUPTS`. This prevents context switches and servicing of ISRs during a critical section. Therefore, critical sections are used as a valid protection method against simultaneous access in vanilla FreeRTOS.

ESP-IDF contains some modifications to work with dual core concurrency, and the dual core API is used even on a single core only chip.

For this reason, ESP-IDF FreeRTOS implements critical sections using special mutexes, referred by portMUX_Type objects on top of specific spinlock component and calls to enter or exit a critical must provide a spinlock object that is associated with a shared resource requiring access protection. When entering a critical section in ESP-IDF FreeRTOS, the calling core will disable its scheduler and interrupts similar to the vanilla FreeRTOS implementation. However, the calling core will also take the locks whilst the other core is left unaffected during the critical section. If the other core attempts to take the spinlock, it will spin until the lock is released. Therefore, the ESP-IDF FreeRTOS implementation of critical sections allows a core to have protected access to a shared resource without disabling the other core. The other core will only be affected if it tries to concurrently access the same resource.

The ESP-IDF FreeRTOS critical section functions have been modified as follows···

- `taskENTER_CRITICAL(mux)`, `taskENTER_CRITICAL_ISR(mux)`, `portENTER_CRITICAL(mux)`, `portENTER_CRITICAL_ISR(mux)` are all macro defined to call `vTaskEnterCritical()`
- `taskEXIT_CRITICAL(mux)`, `taskEXIT_CRITICAL_ISR(mux)`, `portEXIT_CRITICAL(mux)`, `portEXIT_CRITICAL_ISR(mux)` are all macro defined to call `vTaskExitCritical()`
- `portENTER_CRITICAL_SAFE(mux)`, `portEXIT_CRITICAL_SAFE(mux)` macro identifies the context of execution, i.e ISR or Non-ISR, and calls appropriate critical section functions (`port*_CRITICAL` in Non-ISR and `port*_CRITICAL_ISR` in ISR) in order to be in compliance with Vanilla FreeRTOS.

For more details see soc/include/soc/spinlock.h and freertos/tasks.c

It should be noted that when modifying vanilla FreeRTOS code to be ESP-IDF FreeRTOS compatible, it is trivial to modify the type of critical section called as they are all defined to call the same function. As long as the same spinlock is provided upon entering and exiting, the type of call should not matter.

4.10.6 Task Deletion

FreeRTOS task deletion prior to v9.0.0 delegated the freeing of task memory entirely to the Idle Task. Currently, the freeing of task memory will occur immediately (within *`vTaskDelete()`*) if the task being deleted is not currently running or is not pinned to the other core (with respect to the core *`vTaskDelete()`* is called on). TLSP deletion callbacks will also run immediately if the same conditions are met.

However, calling *`vTaskDelete()`* to delete a task that is either currently running or pinned to the other core will still result in the freeing of memory being delegated to the Idle Task.

4.10.7 Thread Local Storage Pointers & Deletion Callbacks

Thread Local Storage Pointers (TLSP) are pointers stored directly in the TCB. TLSP allow each task to have its own unique set of pointers to data structures. However task deletion behavior in vanilla FreeRTOS does not automatically free the memory pointed to by TLSP. Therefore if the memory pointed to by TLSP is not explicitly freed by the user before task deletion, memory leak will occur.

ESP-IDF FreeRTOS provides the added feature of Deletion Callbacks. Deletion Callbacks are called automatically during task deletion to free memory pointed to by TLSP. Each TLSP can have its own Deletion Callback. Note that due to the to *Task Deletion* behavior, there can be instances where Deletion Callbacks are called in the context of the Idle Tasks. Therefore Deletion Callbacks **should never attempt to block** and critical sections should be kept as short as possible to minimize priority inversion.

Deletion callbacks are of type `void (*TlsDeleteCallbackFunction_t)( int, void * )` where the first parameter is the index number of the associated TLSP, and the second parameter is the TLSP itself.

Deletion callbacks are set alongside TLSP by calling *`vTaskSetThreadLocalStoragePointerAndDelCallback()`*. Calling the vanilla FreeRTOS function *`vTaskSetThreadLocalStoragePointer()`* will simply set the TLSP's associated Deletion Callback to *NULL* meaning that no callback will be called for that TLSP during task deletion. If a deletion callback is *NULL*, users should manually free the memory pointed to by the associated TLSP before task deletion in order to avoid memory leak.

CONFIG_FREERTOS_THREAD_LOCAL_STORAGE_POINTERS in menuconfig can be used to configure the number TLSP and Deletion Callbacks a TCB will have.

For more details see *FreeRTOS API reference*.

4.10.8 Configuring ESP-IDF FreeRTOS

The ESP-IDF FreeRTOS can be configured in the project configuration menu (`idf.py menuconfig`) under `Component Config/FreeRTOS`. The following section highlights some of the ESP-IDF FreeRTOS configuration options. For a full list of ESP-IDF FreeRTOS configurations, see *FreeRTOS*

As ESP32-S2 is a single core SoC, the config item *CONFIG_FREERTOS_UNICORE* is always set. This means ESP-IDF only runs on the single CPU. Note that this is **not equivalent to running vanilla FreeRTOS**. Behaviors of multiple components in ESP-IDF will be modified. For more details regarding the effects of running ESP-IDF FreeRTOS on a single core, search for occurences of `CONFIG_FREERTOS_UNICORE` in the ESP-IDF components.

CONFIG_FREERTOS_THREAD_LOCAL_STORAGE_POINTERS will define the number of Thread Local Storage Pointers each task will have in ESP-IDF FreeRTOS.

CONFIG_FREERTOS_SUPPORT_STATIC_ALLOCATION will enable the backported functionality of *xTaskCreateStaticPinnedToCore()* in ESP-IDF FreeRTOS

CONFIG_FREERTOS_ASSERT_ON_UNTESTED_FUNCTION will trigger a halt in particular functions in ESP-IDF FreeRTOS which have not been fully tested in an SMP context.

CONFIG_FREERTOS_TASK_FUNCTION_WRAPPER will enclose all task functions within a wrapper function. In the case that a task function mistakenly returns (i.e. does not call *vTaskDelete()*), the call flow will return to the wrapper function. The wrapper function will then log an error and abort the application, as illustrated below:

```
E (25) FreeRTOS: FreeRTOS task should not return. Aborting now!
abort() was called at PC 0x40085c53 on core 0
```

4.11 Thread Local Storage

4.11.1 Overview

Thread-local storage (TLS) is a mechanism by which variables are allocated such that there is one instance of the variable per extant thread. ESP-IDF provides three ways to make use of such variables:

- *FreeRTOS Native API*: ESP-IDF FreeRTOS native API.
- *Pthread API*: ESP-IDF's pthread API.
- *C11 Standard*: C11 standard introduces special keyword to declare variables as thread local.

4.11.2 FreeRTOS Native API

The ESP-IDF FreeRTOS provides the following API to manage thread local variables:

- *vTaskSetThreadLocalStoragePointer()*
- *pvTaskGetThreadLocalStoragePointer()*
- *vTaskSetThreadLocalStoragePointerAndDelCallback()*

In this case maximum number of variables that can be allocated is limited by `configNUM_THREAD_LOCAL_STORAGE_POINTERS` macro. Variables are kept in the task control block (TCB) and accessed by their index. Note that index 0 is reserved for ESP-IDF internal uses. Using that API user can allocate thread local variables of an arbitrary size and assign them to any number of tasks. Different tasks can have different sets of TLS variables. If size of the variable is more then 4 bytes then user is responsible for allocating/deallocating memory for it. Variable's deallocation is initiated by FreeRTOS when task is deleted, but user must provide function (callback) to do proper cleanup.

4.11.3 Pthread API

The ESP-IDF provides the following pthread API to manage thtread local variables:

- `pthread_key_create()`
- `pthread_key_delete()`
- `pthread_getspecific()`
- `pthread_setspecific()`

This API has all benefits of the one above, but eliminates some its limits. The number of variables is limited only by size of available memory on the heap. Due to the dynamic nature this API introduces additional performance overhead compared to the native one.

4.11.4 C11 Standard

The ESP-IDF FreeRTOS supports thread local variables according to C11 standard (ones specified with `__thread` keyword). For details on this GCC feature please see https://gcc.gnu.org/onlinedocs/gcc-5.5.0/gcc/Thread-Local.html#Thread-Local. Storage for that kind of variables is allocated on the task's stack. Note that area for all such variables in the program will be allocated on the stack of every task in the system even if that task does not use such variables at all. For example ESP-IDF system tasks (like `ipc`, `timer` tasks etc.) will also have that extra stack space allocated. So this feature should be used with care. There is a tradeoff: C11 thread local variables are quite handy to use in programming and can be accessed using just a few Xtensa instructions, but this benefit goes with the cost of additional stack usage for all tasks in the system. Due to static nature of variables allocation all tasks in the system have the same sets of C11 thread local variables.

4.12 High-Level Interrupts

The Xtensa architecture has support for 32 interrupts, divided over 8 levels, plus an assortment of exceptions. On the ESP32-S2, the interrupt mux allows most interrupt sources to be routed to these interrupts using the *interrupt allocator*. Normally, interrupts will be written in C, but ESP-IDF allows high-level interrupts to be written in assembly as well, allowing for very low interrupt latencies.

4.12.1 Interrupt Levels

Level	Symbol	Remark
1	N/A	Exception and level 0 interrupts. Handled by ESP-IDF
2-3	N/A	Medium level interrupts. Handled by ESP-IDF
4	xt_highint4	Normally used by ESP-IDF debug logic
5	xt_highint5	Free to use
NMI	xt_nmi	Free to use
dbg	xt_debugexception	Debug exception. Called on e.g. a BREAK instruction.

Using these symbols is done by creating an assembly file (suffix .S) and defining the named symbols, like this:

```
      .section .iram1,"ax"
      .global     xt_highint5
      .type       xt_highint5,@function
      .align      4
  xt_highint5:
      ... your code here
      rsr     a0, EXCSAVE_5
      rfi     5

For a real-life example, see the :component_file:`esp32s2/dport_panic_highint_hdl.
↪S` file; the panic handler interrupt is implemented there.
```

4.12.2 Notes

- Do not call C code from a high-level interrupt; because these interupts still run in critical sections, this can cause crashes. (The panic handler interrupt does call normal C code, but this is OK because there is no intention of returning to the normal code flow afterwards.)

- Make sure your assembly code gets linked in. If the interrupt handler symbol is the only symbol the rest of the code uses from this file, the linker will take the default ISR instead and not link the assembly file into the final project. To get around this, in the assembly file, define a symbol, like this:

```
    .global ld_include_my_isr_file
ld_include_my_isr_file:
```

(The symbol is called `ld_include_my_isr_file` here but can have any arbitrary name not defined anywhere else.) Then, in the component.mk, add this file as an unresolved symbol to the ld command line arguments:

```
COMPONENT_ADD_LDFLAGS := -u ld_include_my_isr_file
```

This should cause the linker to always include a file defining `ld_include_my_isr_file`, causing the ISR to always be linked in.

- High-level interrupts can be routed and handled using esp_intr_alloc and associated functions. The handler and handler arguments to esp_intr_alloc must be NULL, however.
- In theory, medium priority interrupts could also be handled in this way. For now, ESP-IDF does not support this.

4.13 JTAG 调试

本文将指导安装 ESP32-S2 的 OpenOCD 调试环境，并介绍如何使用 GDB 来调试 ESP32-S2 的应用程序。本文的组织结构如下：

引言 介绍本指南主旨。
工作原理 介绍 ESP32-S2，JTAG（Joint Test Action Group）接口，OpenOCD 和 GDB 是如何相互连接从而实现 ESP32-S2 的调试功能。
选择 *JTAG* 适配器 介绍有关 JTAG 硬件适配器的选择及参照标准。
安装 *OpenOCD* 介绍如何安装官方预编译好的 OpenOCD 软件包并验证是否安装成功。
配置 *ESP32-S2* 目标板 介绍如何设置 OpenOCD 软件并安装 JTAG 硬件适配器，这两者共同组成最终的调试目标。
启动调试器 介绍如何从*Eclipse 集成开发环境* 和*命令行终端* 启动 GDB 调试会话。
调试范例 如果你对 GDB 不太熟悉，本小节会分别针对*Eclipse 集成开发环境* 和*命令行终端* 来讲解调试的范例。
从源码构建 *OpenOCD* 介绍如何在*Windows*，*Linux* 和*MacOS* 操作系统上从源码构建 OpenOCD。
注意事项和补充内容 介绍使用 OpenOCD 和 GDB 通过 JTAG 接口调试 ESP32-S2 时的注意事项和补充内容。

4.13.1 引言

乐鑫已经为 ESP32-S2 处理器和多核 FreeRTOS 架构移植好了 OpenOCD，它将成为大多数 ESP32-S2 应用程序的基础。此外，乐鑫还提供了一些 OpenOCD 本身并不支持的工具来进一步丰富调试的功能。

本文将指导如何在 Linux，Windows 和 MacOS 环境下为 ESP32-S2 安装 OpenOCD，并使用 GDB 进行软件调试。除了个别操作系统的安装过程有所差别以外，软件用户界面和使用流程都是一样的。

注解： 本文使用的图片素材来自于 Ubuntu 16.04 LTS 上 Eclipse Neon 3 软件的截图，不同的操作系统（Windows，MacOS 或者 Linux）和 Eclipse 软件版本在用户界面上可能会有细微的差别。

4.13.2 工作原理

通过 JTAG（Joint Test Action Group）接口使用 OpenOCD 调试 ESP32-S2 时所需要的一些关键的软件和硬件包括 **xtensa-esp32s2-elf-gdb 调试器**，**OpenOCD 片上调试器**和连接到 **ESP32-S2** 目标的 **JTAG 适配器**。

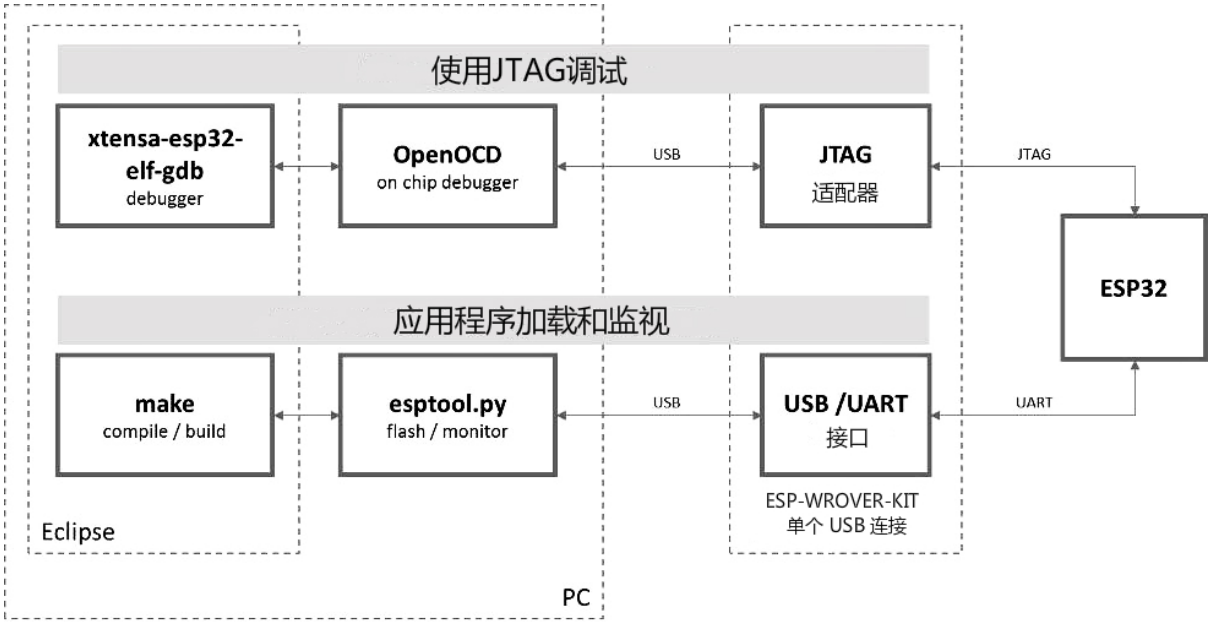

图 6: JTAG 调试 - 概述图

在“Application Loading and Monitoring”下还有另外一组软件和硬件，它们用来编译、构建和烧写应用程序到 ESP32-S2 上，以及监视来自 ESP32-S2 的运行诊断信息。

Eclipse 环境集成了 JTAG 调试和应用程序加载、监视的功能，它使得软件从编写、编译、加载到调试的迭代过程变得更加快速而简单。所有的软件均适用于 Windows，Linux 和 MacOS 平台。

如果你使用的是*ESP-S2-Kaluga-1*，得益于板载的 FT232H 芯片，PC 和 ESP32-S2 的连接仅仅需要一根 USB 线即可完成。FT232H 提供了两路 USB 通道，一路连接到 JTAG，另一路连接到 UART。

根据用户的喜好，除了使用 Eclipse 集成开发环境，还可以直接在命令行终端运行 *debugger* 和 *idf.py build*。

4.13.3 选择 JTAG 适配器

上手 JTAG 最快速便捷的方式是使用*ESP-S2-Kaluga-1*，因为它板载了 JTAG 调试接口，无需使用外部的 JTAG 硬件适配器和额外的线缆来连接 JTAG 与 ESP32-S2。ESP-S2-Kaluga-1 采用 FT2232H 提供的 JTAG 接口，可以稳定运行在 20 MHz 的时钟频率，外接的适配器很难达到这个速度。

如果你想使用单独的 JTAG 适配器，请确保其与 ESP32-S2 的电平电压和 OpenOCD 软件都兼容。ESP32-S2 使用的是业界标准的 JTAG 接口，它省略了（实际上也并不需要）TRST 信号脚。JTAG 使用的 IO 引脚由 VDD_3P3_RTC 电源引脚供电（通常连接到外部 3.3 V 的电源轨），因此 JTAG 硬件适配器的引脚需要能够在该电压范围内正常工作。

在软件方面，OpenOCD 支持相当多数量的 JTAG 适配器，可以参阅 OpenOCD 支持的适配器列表（尽管上面显示的器件不太完整），这个页面还列出了兼容 SWD 接口的适配器，但是请注意，ESP32-S2 目前并不支持 SWD。此外那些被硬编码为只支持特定产品线的 JTAG 适配器也不能在 ESP32-S2 上工作，比如用于 STM32 产品家族的 ST-LINK 适配器。

JTAG 正常工作至少需要连接的信号线有：TDI，TDO，TCK，TMS 和 GND。某些 JTAG 适配器还需要 ESP32-S2 提供一路电源到适配器的某个引脚上（比如 Vtar）用以设置适配器的工作电压。SRST 信号线是可选的，它可以连接到 ESP32-S2 的 CH_PD 引脚上，尽管目前 OpenOCD 对该信号线的支持还非常有限。

4.13.4 安装 OpenOCD

如果你已经按照快速入门一文中的介绍安装好了 ESP-IDF 及其 CMake 构建系统，那么 OpenOCD 已经被默认安装到了你的开发系统中。在设置开发环境结束后，你应该能够在终端中运行如下 OpenOCD 命令:

```
openocd --version
```

终端会输出以下信息（实际版本号可能会比这里列出的更新）：

```
Open On-Chip Debugger  v0.10.0-esp32-20190708 (2019-07-08-11:04)
Licensed under GNU GPL v2
For bug reports, read
    http://openocd.org/doc/doxygen/bugs.html
```

你还可以检查 `OPENOCD_SCRIPTS` 环境变量的值来确认 OpenOCD 配置文件的路径，Linux 和 macOS 用户可以在终端输入 `echo $OPENOCD_SCRIPTS`，Windows 用户需要输入 `echo %OPENOCD_SCRIPTS%`。如果终端打印了有效的路径，则表明 OpenOCD 已经被正确安装。

如果上述步骤没有成功执行，请返回快速入门手册，参考其中*设置安装工具* 章节的说明。

注解： 另外，我们还可以从源代码编译 OpenOCD 工具，相关详细信息请参阅*从源码构建 OpenOCD* 章节。

4.13.5 配置 ESP32-S2 目标板

安装好 OpenOCD 之后就可以配置 ESP32-S2 目标（即带 JTAG 接口的 ESP32-S2 板），具体可以通过以下三个步骤进行：

- 配置并连接 JTAG 接口
- 运行 OpenOCD
- 上传待调试的应用程序

配置并连接 JTAG 接口

此步骤取决于您使用的 JTAG 和 ESP32-S2 板，请参考以下两种情况。

配置 ESP-S2-Kaluga-1 上的 JTAG 接口 所有版本的 ESP-S2-Kaluga-1 板子都内置了 JTAG 调试功能，要使其正常工作，还需要设置相关跳帽来启用 JTAG 功能，设置 SPI 闪存电压和配置 USB 驱动程序。具体步骤请参考以下说明。

配置硬件

- 开箱即用，ESP32-S2-Kaluga-1 不需要任何其他硬件配置即可进行 JTAG 调试。但是，如果遇到问题，请检查“JTAG”DIP 开关的 2 ~ 5 号是否在“ON”的位置。
- 检查 ESP32-S2 上用于 JTAG 通信的引脚是否被接到了其它硬件上，这可能会影响 JTAG 的工作。

表 1: ESP32-S2 引脚和 JTAG 接口信号

ESP32-S2 引脚	JTAG 信号
MTDO / GPIO40	TDO
MTDI / GPIO41	TDI
MTCK / GPIO39	TCK
MTMS / GPIO42	TMS

配置 USB 驱动 安装和配置 USB 驱动，这样 OpenOCD 才能够与 ESP-S2-Kaluga-1 板上的 JTAG 接口通信，并且使用 UART 接口上传待烧写的镜像文件。请根据你的操作系统按照以下步骤进行安装配置。

注解： ESP-S2-Kaluga-1 使用了 FT2232 芯片实现了 JTAG 适配器，所以以下说明同样适用于其他基于 FT2232 的 JTAG 适配器。

Windows

1. 使用标准 USB A / micro USB B 线将 ESP-S2-Kaluga-1 与计算机相连接，并打开板子的电源。
2. 等待 Windows 识别出 ESP-S2-Kaluga-1 并且为其安装驱动。如果驱动没有被自动安装，请前往 官网 下载并手动安装。
3. 从 Zadig 官网 下载 Zadig 工具（Zadig_X.X.exe）并运行。
4. 在 Zadig 工具中，进入 “Options” 菜单中选中 “List All Devices”。
5. 检查设备列表，其中应该包含两条与 ESP-S2-Kaluga-1 相关的条目：“Dual RS232-HS (Interface 0)” 和 “Dual RS232-HS (Interface 1)”。驱动的名字应该是 “FTDIBUS (vxxxx)” 并且 USB ID 为：0403 6010。

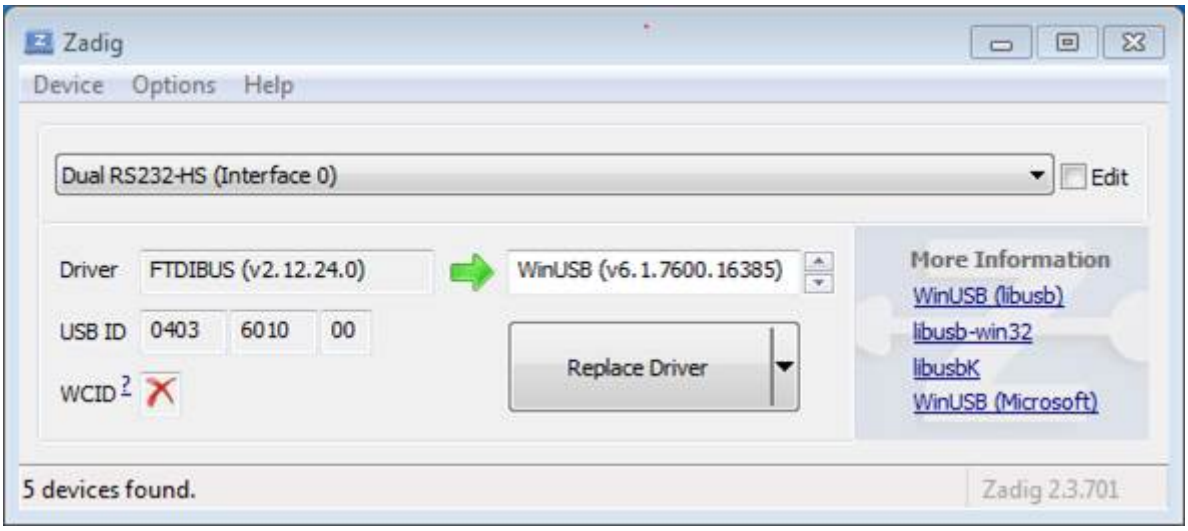

图 7: 在 Zadig 工具中配置 JTAG USB 驱动

6. 第一个设备“Dual RS232-HS（Interface 0）”连接到了 ESP32-S2 的 JTAG 端口，此设备原来的“FTDIBUS (vxxxx)” 驱动需要替换成 “WinUSB (v6xxxxx)”。为此，请选择 “Dual RS232-HS (Interface 0)” 并将驱动重新安装为 “WinUSB (v6xxxxx)”，具体可以参考上图。

注解： 请勿更改第二个设备 “Dual RS232-HS（Interface 1）” 的驱动，它被连接到 ESP32-S2 的串口（UART），用于上传应用程序映像给 ESP32-S2 进行烧写。

现在，ESP-S2-Kaluga-1 的 JTAG 接口应该可以被 OpenOCD 使用了，想要进一步设置调试环境，请前往*运行 OpenOCD* 章节。

Linux

1. 使用标准 USB A / micro USB B 线将 ESP-S2-Kaluga-1 与计算机相连接，并打开板子的电源。
2. 打开终端，输入 `ls -l /dev/ttyUSB*` 命令检查操作系统是否能够识别板子的 USB 端口。类似识别结果如下：

```
user-name@computer-name:~/esp$ ls -l /dev/ttyUSB*
crw-rw---- 1 root dialout 188, 0 Jul 10 19:04 /dev/ttyUSB0
crw-rw---- 1 root dialout 188, 1 Jul 10 19:04 /dev/ttyUSB1
```

3. 根据 OpenOCD README 文档 中 “Permissions delegation” 小节的介绍，设置这两个 USB 端口的访问权限。
4. 注销并重新登录 Linux 系统，然后重新插拔板子的电源使之前的改动生效。在终端再次输入 `ls -l /dev/ttyUSB*` 命令进行验证，查看这两个设备的组所有者是否已经从 `dialout` 更改为 `plugdev`:

```
user-name@computer-name:~/esp$ ls -l /dev/ttyUSB*
crw-rw-r-- 1 root plugdev 188, 0 Jul 10 19:07 /dev/ttyUSB0
crw-rw-r-- 1 root plugdev 188, 1 Jul 10 19:07 /dev/ttyUSB1
```

如果看到类似的输出结果，并且你也是 `plugdev` 组的成员，那么设置工作就完成了。
具有较低编号的 `/dev/ttyUSBn` 接口用于 JTAG 通信，另一路接口被连接到 ESP32-S2 的串口（UART），用于上传应用程序映像给 ESP32-S2 进行烧写。

现在，ESP-S2-Kaluga-1 的 JTAG 接口应该可以被 OpenOCD 使用了，想要进一步设置调试环境，请前往*运行 OpenOCD* 章节。

MacOS 在 macOS 上，同时使用 FT2232 的 JTAG 接口和串口还需另外进行其它操作。当操作系统加载 FTDI 串口驱动的时候，它会对 FT2232 芯片的两个通道做相同的操作。但是，这两个通道中只有一个是被用作串口，而另一个用于 JTAG，如果操作系统已经为用于 JTAG 的通道加载了 FTDI 串口驱动的话，OpenOCD 将无法连接到芯片。有两个方法可以解决这个问题：

1. 在启动 OpenOCD 之前手动卸载 FTDI 串口驱动程序，然后启动 OpenOCD，再加载串口驱动程序。
2. 修改 FTDI 驱动程序的配置，使其不会为 FT2232 芯片的通道 B 进行自我加载，该通道用于 ESP-S2-Kaluga-1 板上的 JTAG 通道。

手动卸载驱动程序

1. 从 FTDI 官网 安装驱动。
2. 使用 USB 线连接 ESP-S2-Kaluga-1。
3. 卸载串口驱动

```
sudo kextunload -b com.FTDI.driver.FTDIUSBSerialDriver
```

有时，您可能还需要卸载苹果的 FTDI 驱动:

```
sudo kextunload -b com.apple.driver.AppleUSBFTDI
```

4. 运行 OpenOCD:

```
.. include:: esp32s2.inc
    :start-after: run-openocd
    :end-before: ---
```

5. 在另一个终端窗口，再一次加载 FTDI 串口驱动:

```
sudo kextload -b com.FTDI.driver.FTDIUSBSerialDriver
```

注解： 如果你需要重启 OpenOCD，则无需再次卸载 FTDI 驱动程序，只需停止 OpenOCD 并再次启动它。只有在重新连接 ESP-S2-Kaluga-1 或者切换了电源的情况下才需要再次卸载驱动。

你也可以根据自身需求，将此过程包装进 shell 脚本中。

修改 FTDI 驱动 简而言之，这种方法需要修改 FTDI 驱动程序的配置文件，这样可以防止为 FT2232H 的通道 B 自动加载串口驱动。

注解： 其他板子可能将通道 A 用于 JTAG，因此请谨慎使用此选项。

警告： 此方法还需要操作系统禁止对驱动进行签名验证，因此可能无法被所有的用户所接受。

1. 使用文本编辑器打开 FTDI 驱动的配置文件（注意 `sudo`）：

```
sudo nano /Library/Extensions/FTDIUSBSerialDriver.kext/Contents/Info.plist
```

2. 找到并删除以下几行:

```
<key>FT2232H_B</key>
<dict>
    <key>CFBundleIdentifier</key>
    <string>com.FTDI.driver.FTDIUSBSerialDriver</string>
    <key>IOClass</key>
    <string>FTDIUSBSerialDriver</string>
    <key>IOProviderClass</key>
    <string>IOUSBInterface</string>
    <key>bConfigurationValue</key>
    <integer>1</integer>
    <key>bInterfaceNumber</key>
    <integer>1</integer>
    <key>bcdDevice</key>
    <integer>1792</integer>
    <key>idProduct</key>
    <integer>24592</integer>
    <key>idVendor</key>
    <integer>1027</integer>
</dict>
```

3. 保存并关闭文件
4. 禁用驱动的签名认证:
    1. 点击苹果的 logo，选择 “Restart…”
    2. 重启后当听到响铃时，立即按下键盘上的 CMD+R 组合键
    3. 进入恢复模式后，打开终端
    4. 运行命令:

    ```
    csrutil enable --without kext
    ```

    5. 再一次重启系统

完成这些步骤后，可以同时使用串口和 JTAG 接口了。

想要进一步设置调试环境，请前往运行 *OpenOCD* 章节。

配置其它 JTAG 接口 关于适配 OpenOCD 和 ESP32-S2 的 JTAG 接口选择问题，请参考选择 *JTAG* 适配器 章节，确保 JTAG 适配器能够与 OpenOCD 和 ESP32-S2 一同工作。然后按照以下三个步骤进行设置，使其正常工作。

配置硬件

1. 找到 JTAG 接口和 ESP32-S2 板上需要相互连接并建立通信的引脚/信号。

表 2: ESP32-S2 引脚和 JTAG 接口信号

ESP32-S2 引脚	JTAG 信号
MTDO / GPIO40	TDO
MTDI / GPIO41	TDI
MTCK / GPIO39	TCK
MTMS / GPIO42	TMS

2. 检查 ESP32-S2 上用于 JTAG 通信的的引脚是否被连接到了其它硬件上，这可能会影响 JTAG 的工作。
3. 连接 ESP32-S2 和 JTAG 接口上的引脚/信号。

配置驱动 你可能还需要安装软件驱动，才能使 JTAG 在计算机上正常工作，请参阅你所使用的 JTAG 适配器的有关文档，获取相关详细信息。

连接 将 JTAG 接口连接到计算机，打开 ESP32-S2 和 JTAG 接口板上的电源，然后检查计算机是否可以识别到 JTAG 接口。

要继续设置调试环境，请前往*运行 OpenOCD* 章节。

运行 OpenOCD

配置完目标并将其连接到电脑后，即可启动 OpenOCD。

打开终端，按照快速入门中的指南*设置好开发环境*，然后运行如下命令，启动 OpenOCD（该命令在 Windows，Linux，和 macOS 中通用）：

```
openocd -f board/esp32s2-kaluga-1.cfg
```

注解： 上述命令中 `-f` 选项后跟的配置文件专用于 ESP32-S2-Kaluga-1 开发板。您可能需要根据具体使用的硬件而选择或修改不同的配置文件，相关指导请参阅*Configuration of OpenOCD for specific target*。

现在应该可以看到如下输入（此日志来自 ESP32-S2-Kaluga-1 开发板）：

```
user-name@computer-name:~/esp/esp-idf$ openocd -f board/esp32s2-kaluga-1.cfg
Open On-Chip Debugger  v0.10.0-esp32-20200420 (2020-04-20-16:15)
Licensed under GNU GPL v2
For bug reports, read
        http://openocd.org/doc/doxygen/bugs.html
none separate
adapter speed: 20000 kHz
force hard breakpoints
Info : ftdi: if you experience problems at higher adapter clocks, try the command
↪"ftdi_tdo_sample_edge falling"
Info : clock speed 20000 kHz
Info : JTAG tap: esp32s2.cpu0 tap/device found: 0x120034e5 (mfg: 0x272 (Tensilica),
↪ part: 0x2003, ver: 0x1)
Info : esp32s2: Debug controller was reset (pwrstat=0x5F, after clear 0x0F).
Info : esp32s2: Core was reset (pwrstat=0x5F, after clear 0x0F).
```

- 如果出现指示权限问题的错误，请参阅 `~/esp/openocd-esp32` 目录下 OpenOCD README 文件中关于 “Permissions delegation” 的说明。
- 如果发现配置文件有错误，例如 `Can't find board/esp32s2-kaluga-1.cfg`，请检查 `-s` 后面的路径，OpenOCD 会根据此路径来查找 `-f` 指定的文件。此外，还需要检查配置文件是否确实位于该路径下。
- 如果看到 JTAG 错误（输出全是 1 或者全是 0），请检查硬件连接，除了 ESP32-S2 的引脚之外是否还有其他信号连接到了 JTAG，并查看是否所有器件都已经上电。

上传待调试的应用程序

您可以像往常一样构建并上传 ESP32-S2 应用程序，具体请参阅*第八步：编译工程* 章节。

除此以外，还支持使用 OpenOCD 通过 JTAG 接口将应用程序镜像烧写到闪存中，命令如下:

```
openocd -f board/esp32s2-kaluga-1.cfg -c "program_esp filename.bin 0x10000 verify␣
↪exit"
```

其中 OpenOCD 的烧写命令 `program_esp` 具有以下格式：

```
program_esp <image_file> <offset> [verify] [reset] [exit]
```

- `image_file` - 程序镜像文件存放的路径
- `offset` - 镜像烧写到闪存中的偏移地址
- `verify` - 烧写完成后校验闪存中的内容（可选）
- `reset` - 烧写完成后重启目标（可选）

- `exit` - 烧写完成后退出 OpenOCD（可选）

现在可以进行应用程序的调试了，请按照以下章节中讲解的步骤进行操作。

4.13.6 启动调试器

ESP32-S2 的工具链中带有 GNU 调试器（简称 GDB）`xtensa-esp32s2-elf-gdb`，它和其它工具链软件存放在同一个 bin 目录下。除了直接在命令行终端中调用并操作 GDB 外，还可以在 IDE (例如 Eclipse，Visual Studio Code 等）中调用它，在图形用户界面的帮助下间接操作 GDB，无需在终端中输入任何命令。

关于以上两种调试器的使用方法，详见以下链接。

- *在 Eclipse 中使用 GDB*
- *在命令行中使用 GDB*

建议首先检查调试器是否能在*命令行终端* 下正常工作，然后再转到使用 Eclipse 等*集成开发环境* 下进行调试工作。

4.13.7 调试范例

本节适用于不熟悉 GDB 的用户，将使用 get-started/blink 下简单的应用程序来演示*调试会话的工作流程*，同时会介绍以下常用的调试操作：

1. *浏览代码，查看堆栈和线程*
2. *设置和清除断点*
3. *手动暂停目标*
4. *单步执行代码*
5. *查看并设置内存*
6. *观察和设置程序变量*
7. *设置条件断点*

此外还会提供*在命令行终端进行调试* 的案例。

在演示之前，请设置好 ESP32-S2 目标板并加载 get-started/blink 至 ESP32-S2 中。

4.13.8 从源码构建 OpenOCD

请参阅以下文档，它们分别介绍了在各大操作系统平台上从源码构建 OpenOCD 的流程。

Windows 环境下从源码编译 OpenOCD

除了从 Espressif 官方 直接下载 OpenOCD 可执行文件，你还可以选择从源码编译得到 OpenOCD。如果想要快速设置 OpenOCD 而不是自行编译，请备份好当前文件，前往*安装 OpenOCD* 章节查阅。

下载 OpenOCD 源码 支持 ESP32-S2 的 OpenOCD 源代码可以从乐鑫官方的 GitHub 获得，网址为 https://github.com/espressif/openocd-esp32。请使用以下命令来下载源代码：

```
cd ~/esp
git clone --recursive https://github.com/espressif/openocd-esp32.git
```

克隆后的源代码被保存在 `~/esp/openocd-esp32` 目录中。

安装依赖的软件包 安装编译 OpenOCD 所需的软件包。

注解： 依次安装以下软件包，检查安装是否成功，然后继续下一个软件包的安装。在进行下一步操作之前，要先解决当前报告的问题。

```
pacman -S libtool
pacman -S autoconf
pacman -S automake
pacman -S texinfo
pacman -S mingw-w64-i686-libusb-compat-git
pacman -S pkg-config
```

注解: 安装 pkg-config 会破坏 esp-idf 的工具链，因而在 OpenOCD 构建完成后，应将其卸载。详见文末进一步说明。如果想要再次构建 OpenOCD，你需要再次运行 pacman -S pkg-config。此步骤安装的其他软件包（在 pkg-config 之前）并不会出现这一问题。

构建 OpenOCD 配置和构建 OpenOCD 的流程如下:

```
cd ~/esp/openocd-esp32
./bootstrap
./configure
make
```

你可以选择最后再执行 sudo make install，如果你已经安装过别的开发平台的 OpenOCD，请跳过这个步骤，因为它可能会覆盖掉原来的 OpenOCD。

注解:

- 如果发生错误，请解决后再次尝试编译，直到 make 成功为止。
- 如果 OpenOCD 存在子模块问题，请 cd 到 openocd-esp32 目录，并输入 git submodule update --init 命令。
- 如果 ./configure 成功运行，JTAG 被使能的信息会被打印在 OpenOCD configuration summary 下面。
- 如果您的设备信息未显示在日志中，请根据 ../openocd-esp32/doc/INSTALL.txt 文中的描述使用 ./configure 启用它。
- 有关编译 OpenOCD 的详细信息，请参阅 openocd-esp32/README.Windows。

一旦 make 过程成功完成，OpenOCD 的可执行文件会被保存到 ~/esp/openocd-esp32/src/openocd 目录中。

如安装依赖步骤所述，最后还需要移除 pkg-config 软件包:

```
pacman -Rs pkg-config
```

下一步 想要进一步配置调试环境，请前往配置 *ESP32-S2* 目标板 章节。

Linux 环境下从源码编译 OpenOCD

除了从 Espressif 官方 直接下载 OpenOCD 可执行文件，你还可以选择从源码编译得到 OpenOCD。如果想要快速设置 OpenOCD 而不是自行编译，请备份好当前文件，前往安装 *OpenOCD* 章节查阅。

下载 OpenOCD 源码 支持 ESP32-S2 的 OpenOCD 源代码可以从乐鑫官方的 GitHub 获得，网址为 https://github.com/espressif/openocd-esp32。请使用以下命令来下载源代码:

```
cd ~/esp
git clone --recursive https://github.com/espressif/openocd-esp32.git
```

克隆后的源代码被保存在 ~/esp/openocd-esp32 目录中。

安装依赖的软件包 安装编译 OpenOCD 所需的软件包。

注解: 依次安装以下软件包，检查安装是否成功，然后继续下一个软件包的安装。在进行下一步操作之前，要先解决当前报告的问题。

```
sudo apt-get install make
sudo apt-get install libtool
sudo apt-get install pkg-config
sudo apt-get install autoconf
sudo apt-get install automake
sudo apt-get install texinfo
sudo apt-get install libusb-1.0
```

注解:

- pkg-config 应为 0.2.3 或以上的版本。
- autoconf 应为 2.6.4 或以上的版本。
- automake 应为 1.9 或以上的版本。
- 当使用 USB-Blaster，ASIX Presto，OpenJTAG 和 FT2232 作为适配器时，需要下载安装 libFTDI 和 FTD2XX 的驱动。
- 当使用 CMSIS-DAP 时，需要安装 HIDAPI。

构建 OpenOCD 配置和构建 OpenOCD 的流程如下:

```
cd ~/esp/openocd-esp32
./bootstrap
./configure
make
```

你可以选择最后再执行 `sudo make install`，如果你已经安装过别的开发平台的 OpenOCD，请跳过这个步骤，因为它可能会覆盖掉原来的 OpenOCD。

注解:

- 如果发生错误，请解决后再次尝试编译，直到 `make` 成功为止。
- 如果 OpenOCD 存在子模块问题，请 `cd` 到 `openocd-esp32` 目录，并输入 `git submodule update --init` 命令。
- 如果 `./configure` 成功运行，JTAG 被使能的信息会被打印在 `OpenOCD configuration summary` 下面。
- 如果您的设备信息未显示在日志中，请根据 `../openocd-esp32/doc/INSTALL.txt` 文中的描述使用 `./configure` 启用它。
- 有关编译 OpenOCD 的详细信息，请参阅 `openocd-esp32/README`。

一旦 `make` 过程成功结束，OpenOCD 的可执行文件会被保存到 `~/openocd-esp32/bin` 目录中。

下一步 想要进一步配置调试环境，请前往*配置 ESP32-S2 目标板* 章节。

MacOS 环境下从源码编译 OpenOCD

除了从 Espressif 官方 直接下载 OpenOCD 可执行文件，你还可以选择从源码编译得到 OpenOCD。如果想要快速设置 OpenOCD 而不是自行编译，请备份好当前文件，前往*安装 OpenOCD* 章节查阅。

下载 OpenOCD 源码 支持 ESP32-S2 的 OpenOCD 源代码可以从乐鑫官方的 GitHub 获得，网址为 https://github.com/espressif/openocd-esp32。请使用以下命令来下载源代码:

```
cd ~/esp
git clone --recursive https://github.com/espressif/openocd-esp32.git
```

克隆后的源代码被保存在 ~/esp/openocd-esp32 目录中。

安装依赖的软件包 使用 Homebrew 安装编译 OpenOCD 所需的软件包:

```
brew install automake libtool libusb wget gcc@4.9 pkg-config
```

构建 OpenOCD 配置和构建 OpenOCD 的流程如下:

```
cd ~/esp/openocd-esp32
./bootstrap
./configure
make
```

你可以选择最后再执行 sudo make install，如果你已经安装过别的开发平台的 OpenOCD，请跳过这个步骤，因为它可能会覆盖掉原来的 OpenOCD。

注解:

- 如果发生错误，请解决后再次尝试编译，直到 make 成功为止。
- 如果 OpenOCD 存在子模块问题，请 cd 到 openocd-esp32 目录，并输入 git submodule update --init 命令。
- 如果 ./configure 成功运行，JTAG 被使能的信息会被打印在 OpenOCD configuration summary 下面。
- 如果您的设备信息未显示在日志中，请根据 ../openocd-esp32/doc/INSTALL.txt 文中的描述使用 ./configure 启用它。
- 有关编译 OpenOCD 的详细信息，请参阅 openocd-esp32/README.OSX。

一旦 make 过程成功结束，OpenOCD 的可执行文件会被保存到 ~/esp/openocd-esp32/src/openocd 目录中。

下一步 想要进一步配置调试环境，请前往*配置 ESP32-S2 目标板* 章节。

本文档演示所使用的 OpenOCD 是*安装 OpenOCD* 章节中介绍的预编译好的二进制发行版。

如果要使用本地从源代码编译的 OpenOCD 程序，需要将相应可执行文件的路径修改为 src/openocd，并设置 OPENOCD_SCRIPTS 环境变量，这样 OpenOCD 才能找到配置文件。Linux 和 macOS 用户可以执行:

```
cd ~/esp/openocd-esp32
export OPENOCD_SCRIPTS=$PWD/tcl
```

Windows 用户可以执行:

```
cd %USERPROFILE%\esp\openocd-esp32
set "OPENOCD_SCRIPTS=%CD%\tcl"
```

运行本地编译的 OpenOCD 的示例如下（Linux 和 macOS 用户）:

```
src/openocd -f board/esp32s2-kaluga-1.cfg
```

Windows 用户:

```
src\openocd -f board/esp32s2-kaluga-1.cfg
```

4.13.9 注意事项和补充内容

本节列出了本指南中提到的所有注意事项和补充内容的链接。

注意事项和补充内容

本节提供了本指南中各部分提到的一些注意事项和补充内容。

可用的断点和观察点 ESP32-S2 调试器支持 2 个硬件断点和 64 个软件断点。硬件断点是由 ESP32-S2 芯片内部的逻辑电路实现的，能够设置在代码的任何位置：闪存或者 IRAM 的代码区域。除此以外，OpenOCD 实现了两种软件断点：闪存断点（最多 32 个）和 IRAM 断点（最多 32 个）。目前 GDB 无法在闪存中设置软件断点，因此除非解决此限制，否则这些断点只能由 OpenOCD 模拟为硬件断点。（详细信息可以参阅下面）。ESP32-S2 还支持 2 个观察点，所以可以观察两个变量的变化或者通过 GDB 命令 `watch myVariable` 来读取变量的值。请注意 menuconfig 中的*CONFIG_FREERTOS_WATCHPOINT_END_OF_STACK* 选项会使用第二个观察点，如果你想在 OpenOCD 或者 GDB 中再次尝试使用这个观察点，可能不会得到预期的结果。详情请查看 menuconfig 中的帮助文档。

关于断点的补充知识 使用软件闪存模拟部分硬件断点的意思就是当使用 GDB 命令 `hb myFunction` 给某个函数设置硬件断点时，如果该函数位于闪存中，并且此时还有可用的硬件断点，那调试器就会使用硬件断点，否则就使用 32 个软件闪存断点中的一个来模拟。这个规则同样适用于 `b myFunction` 之类的命令，在这种情况下，GDB 会自己决定该使用哪种类型的断点。如果 `myFunction` 位于可写区域（IRAM），那就会使用软件 IRAM 断点，否则就会像处理 `hb` 命令一样使用硬件断点或者软件闪存断点。

闪存映射 vs 软件闪存断点 为了在闪存中设置或者清除软件断点，OpenOCD 需要知道它们在闪存中的地址。为了完成从 ESP32-S2 的地址空间到闪存地址的转换，OpenOCD 使用闪存中程序代码区域的映射。这些映射被保存在程序映像的头部，位于二进制数据（代码段和数据段）之前，并且特定于写入闪存的每一个应用程序的映像。因此，为了支持软件闪存断点，OpenOCD 需要知道待调试的应用程序映像在闪存中的位置。默认情况下，OpenOCD 会在 0x8000 处读取分区表并使用第一个找到的应用程序映像的映射，但是也可能会存在无法工作的情况，比如分区表不在标准的闪存位置，甚至可能有多个映像：一个出厂映像和两个 OTA 映像，你可能想要调试其中的任意一个。为了涵盖所有可能的调试情况，OpenOCD 支持特殊的命令，用于指定待调试的应用程序映像在闪存中的具体位置。该命令具有以下格式：

```
esp appimage_offset <offset>
```

偏移量应为十六进制格式，如果要恢复默认行为，可以将偏移地址设置为 `-1` 。

注解： 由于 GDB 在连接 OpenOCD 时仅仅请求一次内存映射，所以可以在 TCL 配置文件中指定该命令，或者通过命令行传递给 OpenOCD。对于后者，命令行示例如下：

```
openocd -f board/esp32s2-kaluga-1.cfg -c "init; halt; esp appimage_offset 0x210000"
```

另外还可以通过 OpenOCD 的 telnet 会话执行该命令，然后再连接 GDB，不过这种方式似乎没有那么便捷。

“next” 命令无法跳过子程序的原因 当使用 `next` 命令单步执行代码时，GDB 会在子程序的前面设置一个断点（两个中可用的一个），这样就可以跳过进入子程序内部的细节。如果这两个断点已经用在代码的其它位置，那么 `next` 命令将不起作用。在这种情况下，请删掉一个断点以使其中一个变得可用。当两个断点都已经被使用时，`next` 命令会像 `step` 命令一样工作，调试器就会进入子程序内部。

OpenOCD 支持的编译时的选项 ESP-IDF 有一些针对 OpenOCD 调试功能的选项可以在编译时进行设置：

- *CONFIG_ESP32S2_DEBUG_OCDAWARE* 默认会被使能。如果程序抛出了不可修复或者未处理的异常，并且此时已经连接上了 JTAG 调试器（即 OpenOCD 正在运行），那么 ESP-IDF 将会进入调试器工作模式。
- *CONFIG_FREERTOS_WATCHPOINT_END_OF_STACK* 默认没有使能。在所有任务堆栈的末尾设置观察点，从 1 号开始索引。这是调试任务堆栈溢出的最准确的方式。

更多有关设置编译时的选项的信息，请参阅项目配置菜单。

支持 FreeRTOS OpenOCD 完全支持 ESP-IDF 自带的 FreeRTOS 操作系统，GDB 会将 FreeRTOS 中的任务当做线程。使用 GDB 命令 `i threads` 可以查看所有的线程，使用命令 `thread n` 可以切换到某个具体任务的堆栈，其中 `n` 是线程的编号。检测 FreeRTOS 的功能可以在配置目标时被禁用。更多详细信息，请参阅*Configuration of OpenOCD for specific target*.

优化 JTAG 的速度 为了实现更高的数据通信速率同时最小化丢包数，建议优化 JTAG 时钟频率的设置，使其达到 JTAG 能稳定运行的最大值。为此，请参考以下建议。

1. 如果 CPU 以 80 MHz 运行，则 JTAG 时钟频率的上限为 20 MHz；如果 CPU 以 160 MHz 或者 240 MHz 运行，则上限为 26 MHz。
2. 根据特定的 JTAG 适配器和连接线缆的长度，你可能需要将 JTAG 的工作频率降低至 20 / 26 MHz 以下。
3. 在某些特殊情况下，如果你看到 DSR/DIR 错误（并且它并不是由 OpenOCD 试图从一个没有物理存储器映射的地址空间读取数据而导致的），请降低 JTAG 的工作频率。
4. ESP-WROVER-KIT 能够稳定运行在 20 / 26 MHz 频率下。

调试器的启动命令的含义 在启动时，调试器发出一系列命令来复位芯片并使其在特定的代码行停止运行。这个命令序列（如下所示）支持自定义，用户可以选择在最方便合适的代码行开始调试工作。

- `set remote hardware-watchpoint-limit 2` —限制 GDB 仅使用 ESP32-S2 支持的两个硬件观察点。更多详细信息，请查阅 GDB 配置远程目标 。
- `mon reset halt` —复位芯片并使 CPU 停止运行。
- `flushregs` —monitor (`mon`) 命令无法通知 GDB 目标状态已经更改，GDB 会假设在 `mon reset halt` 之前所有的任务堆栈仍然有效。实际上，复位后目标状态将发生变化。执行 `flushregs` 是一种强制 GDB 从目标获取最新状态的方法。
- `thb app_main` —在 `app_main` 处插入一个临时的硬件断点，如果有需要，可以将其替换为其他函数名。
- `c` —恢复程序运行，它将会在 `app_main` 的断点处停止运行。

Configuration of OpenOCD for specific target There are several kinds of OpenOCD configuration files (`*.cfg`). All configuration files are located in subdirectories of `share/openocd/scripts` directory of OpenOCD distribution (or `tcl/scripts` directory of the source repository). For the purposes of this guide, the most important ones are `board`, `interface` and `target`.

- `interface` configuration files describe the JTAG adapter. Examples of JTAG adapters are ESP-Prog and J-Link.
- `target` configuration files describe specific chips, or in some cases, modules.
- `board` configuration files are provided for development boards with a built-in JTAG adapter. Such files include an `interface` configuration file to choose the adapter, and `target` configuration file to choose the chip/module.

The following configuration files are available for ESP32-S2:

表 3: ESP32-S2 相关的 OpenOCD 配置文件

名字	描述
board/esp32s2-kaluga-1.cfg	ESP32-S2-Kaluga-1 开发板配置文件，内部已包含 ESP32-S2 目标配置和 JTAG 适配器配置
target/esp32s2.cfg	ESP32-S2 目标配置文件，可以和某个 interface/ 下的配置文件一同使用
interface/ftdi/esp32s2_kaluga_v1.cfg	适用于 ESP32-S2-Kaluga-1 开发板的 JTAG 适配器配置文件
interface/ftdi/esp32_devkitj_v1.cfg	适用于 ESP-Prog 板子的 JTAG 适配器配置文件

If you are using one of the boards which have a pre-defined configuration file, you only need to pass one `-f` argument to OpenOCD, specifying that file.

If you are using a board not listed here, you need to specify both the interface configuration file and target configuration file.

Custom configuration files OpenOCD configuration files are written in TCL, and include a variety of choices for customization and scripting. This can be useful for non-standard debugging situations. Please refer to OpenOCD Manual for the TCL scripting reference.

OpenOCD configuration variables The following variables can be optionally set before including the ESP-specific target configuration file. This can be done either in a custom configuration file, or from the command line.

The syntax for setting a variable in TCL is:

```
set VARIABLE_NAME value
```

To set a variable from the command line (replace the name of .cfg file with the correct file for your board):

```
openocd -c 'set VARIABLE_NAME value' -f board/esp-xxxxx-kit.cfg
```

It is important to set the variable before including the ESP-specific configuration file, otherwise the variable will not have effect. You can set multiple variables by repeating the `-c` option.

表 4: Common ESP-related OpenOCD variables

Variable	Description
ESP_RTOS	Set to `none` to disable RTOS support. In this case, thread list will not be available in GDB. Can be useful when debugging FreeRTOS itself, and stepping through the scheduler code.
ESP_FLASH_SIZE	Set to `0` to disable Flash breakpoints support.
ESP_SEMIHOST_BASEDIR	Set to the path (on the host) which will be the default directory for semihosting functions.

复位 ESP32-S2 通过在 GDB 中输入 `mon reset` 或者 `mon reset halt` 来复位板子。

不要将 JTAG 引脚用于其他功能 如果除了 ESP32-S2 模组和 JTAG 适配器之外的其他硬件也连接到了 JTAG 引脚，那么 JTAG 的操作可能会受到干扰。ESP32-S2 JTAG 使用以下引脚：

如果用户应用程序更改了 JTAG 引脚的配置，JTAG 通信可能会失败。如果 OpenOCD 正确初始化（检测到两个 Tensilica 内核），但在程序运行期间失去了同步并报出大量 DTR/DIR 错误，则应用程序可能将 JTAG 引脚重新配置为其他功能或者用户忘记将 Vtar 连接到 JTAG 适配器。

表 5: ESP32-S2 引脚和 JTAG 接口信号

ESP32-S2 引脚	JTAG 信号
MTDO / GPIO40	TDO
MTDI / GPIO41	TDI
MTCK / GPIO39	TCK
MTMS / GPIO42	TMS

下面是 GDB 在应用程序进入重新配置 MTDO/GPIO15 作为输入代码后报告的一系列错误摘录:

```
cpu0: xtensa_resume (line 431): DSR (FFFFFFFF) indicates target still busy!
cpu0: xtensa_resume (line 431): DSR (FFFFFFFF) indicates DIR instruction generated␣
↪an exception!
cpu0: xtensa_resume (line 431): DSR (FFFFFFFF) indicates DIR instruction generated␣
↪an overrun!
cpu1: xtensa_resume (line 431): DSR (FFFFFFFF) indicates target still busy!
cpu1: xtensa_resume (line 431): DSR (FFFFFFFF) indicates DIR instruction generated␣
↪an exception!
cpu1: xtensa_resume (line 431): DSR (FFFFFFFF) indicates DIR instruction generated␣
↪an overrun!
```

JTAG with Flash Encryption or Secure Boot By default, enabling Flash Encryption and/or Secure Boot will disable JTAG debugging. On first boot, the bootloader will burn an eFuse bit to permanently disable JTAG at the same time it enables the other features.

The project configuration option *CONFIG_SECURE_BOOT_ALLOW_JTAG* will keep JTAG enabled at this time, removing all physical security but allowing debugging. (Although the name suggests Secure Boot, this option can be applied even when only Flash Encryption is enabled).

However, OpenOCD may attempt to automatically read and write the flash in order to set *software breakpoints*. This has two problems:

- Software breakpoints are incompatible with Flash Encryption, OpenOCD currently has no support for encrypting or decrypting flash contents.
- If Secure Boot is enabled, setting a software breakpoint will change the digest of a signed app and make the signature invalid. This means if a software breakpoint is set and then a reset occurs, the signature verification will fail on boot.

To disable software breakpoints while using JTAG, add an extra argument `-c 'set ESP_FLASH_SIZE 0'` to the start of the OpenOCD command line. For example:

```
openocd -c 'set ESP_FLASH_SIZE 0' -f board/esp32-wrover-kit-3.3v.cfg
```

注解: For the same reason, the ESP-IDF app may fail bootloader verification of app signatures, when this option is enabled and a software breakpoint is set.

报告 OpenOCD / GDB 的问题 如果你遇到 OpenOCD 或者 GDB 程序本身的问题，并且在网上没有找到可用的解决方案，请前往 https://github.com/espressif/openocd-esp32/issues 新建一个议题。

1. 请在问题报告中提供你使用的配置的详细信息：
   a. JTAG 适配器类型。
   b. 用于编译和加载正在调试的应用程序的 ESP-IDF 版本号。
   c. 用于调试的操作系统的详细信息。
   d. 操作系统是在本地计算机运行还是在虚拟机上运行?
2. 创建一个能够演示问题的简单示例工程，描述复现该问题的步骤。且这个调试示例不能受到 Wi-Fi 协议栈引入的非确定性行为的影响，因而再次遇到同样问题时，更容易复现。

3. 在启动命令中添加额外的参数来输出调试日志。
   OpenOCD 端：

```
openocd -l openocd_log.txt -d3 -f board/esp32s2-kaluga-1.cfg
```

   这种方式会将日志输出到文件，但是它会阻止调试信息打印在终端上。当有大量信息需要输出的时候（比如调试等级提高到 -d3）这是个不错的选择。如果你仍然希望在屏幕上看到调试日志，请改用以下命令：

```
openocd -d3 -f board/esp32s2-kaluga-1.cfg 2>&1 | tee openocd.log
```

   Debugger 端：

```
xtensa-esp32s2-elf-gdb -ex "set remotelogfile gdb_log.txt" <all other options>
```

   也可以将命令 `remotelogfile gdb_log.txt` 添加到 `gdbinit` 文件中。
4. 请将 `openocd_log.txt` 和 `gdb_log.txt` 文件附在你的问题报告中。

4.13.10 相关文档

使用调试器

本节会在*Eclipse* 和命令行 中分别介绍配置和运行调试器的方法。我们建议你首先通过命令行 检查调试器是否正常工作，然后再转到使用*Eclipse* 平台。

在 Eclipse 中使用 GDB 标准的 Eclipse 安装流程默认安装调试功能，另外我们还可以使用插件来调试，比如 “GDB Hardware Debugging”。这个插件用起来非常方便，本指南会详细介绍该插件的使用方法。

首先，通过打开 Eclipse 并转到 “Help” > “Install New Software” 来安装 “GDB Hardware Debugging” 插件。

安装完成后，按照以下步骤配置调试会话。请注意，一些配置参数是通用的，有些则针对特定项目。我们会通过配置 “blink” 示例项目的调试环境来进行展示，请先按照使用 *Eclipse IDE* 编译和烧写 文章介绍的方法将该示例项目添加到 Eclipse 的工作空间。示例项目 get-started/blink 的源代码可以在 ESP-IDF 仓库的 examples 目录下找到。

1. 在 Eclipse 中，进入 *Run* > *Debug Configuration*，会出现一个新的窗口。在窗口的左侧窗格中，双击 “GDB Hardware Debugging”（或者选择 “GDB Hardware Debugging” 然后按下 “New” 按钮）来新建一个配置。
2. 在右边显示的表单中，“Name:” 一栏中输入配置的名称，例如： “Blink checking”。
3. 在下面的 “Main” 选项卡中，点击 “Project:” 边上的 “Browse” 按钮，然后选择当前的 “blink” 项目。
4. 在下一行的 “C/C++ Application:” 中，点击 “Browse” 按钮，选择 “blink.elf” 文件。如果 “blink.elf” 文件不存在，那么很有可能该项目还没有编译，请参考使用 *Eclipse IDE* 编辑和烧写 指南中的介绍。
5. 最后，在 “Build (if required) before launching” 下面点击 “Disable auto build”。
   上述步骤 1 - 5 的示例输入如下图所示。
6. 点击 “Debugger” 选项卡，在 “GDB Command” 栏中输入 `xtensa-esp32s2-elf-gdb` 来调用调试器。
7. 更改 “Remote host” 的默认配置，在 “Port number” 下面输入 `3333`。
   上述步骤 6 - 7 的示例输入如下图所示。
8. 最后一个需要更改默认配置的选项卡是 “Startup” 选项卡。在 “Initialization Commands” 下，取消选中 “Reset and Delay (seconds)” 和 “Halt”，然后在下面一栏中输入以下命令：

```
mon reset halt
flushregs
set remote hardware-watchpoint-limit 2
```

注解： 如果你想在启动新的调试会话之前自动更新闪存中的镜像，请在 “Initialization Commands” 文本框的开头添加以下命令行：

```
mon reset halt
mon program_esp ${workspace_loc:blink/build/blink.bin} 0x10000 verify
```

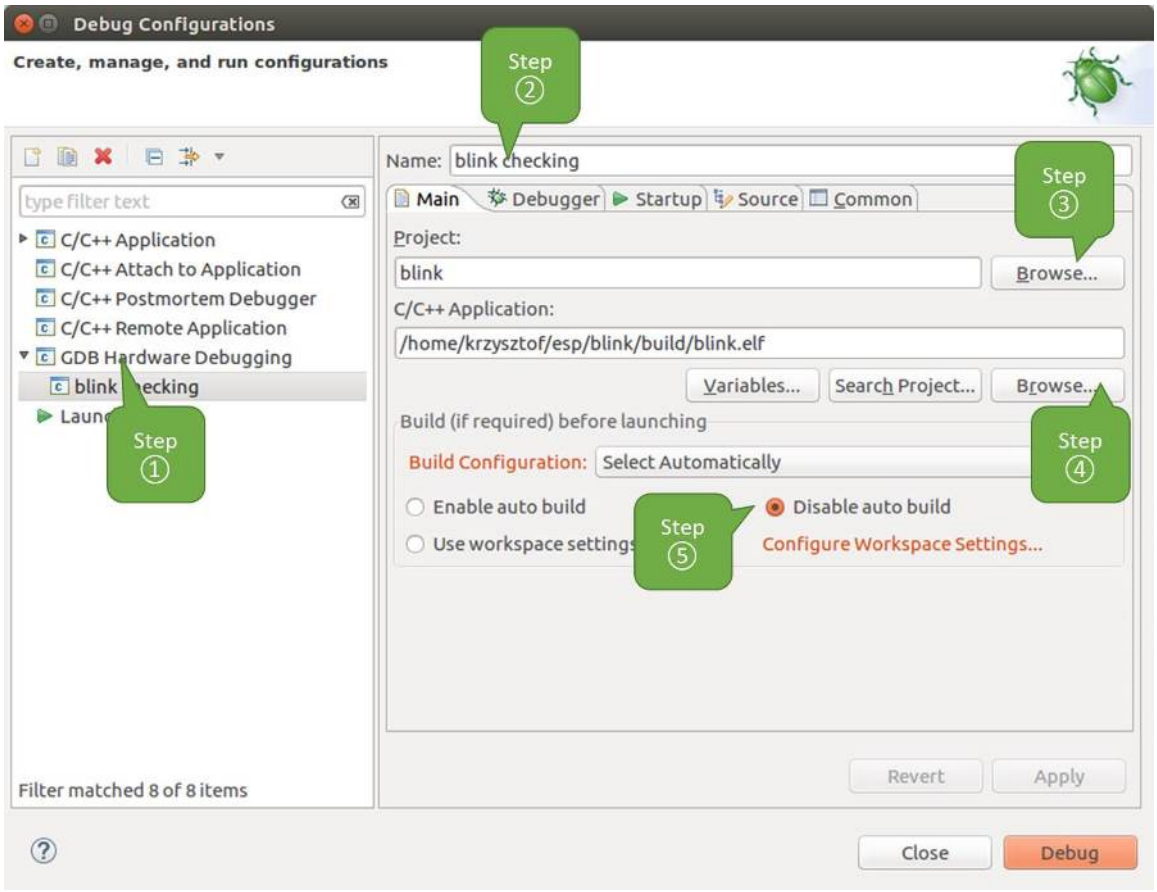

图 8: GDB 硬件调试的配置 - Main 选项卡

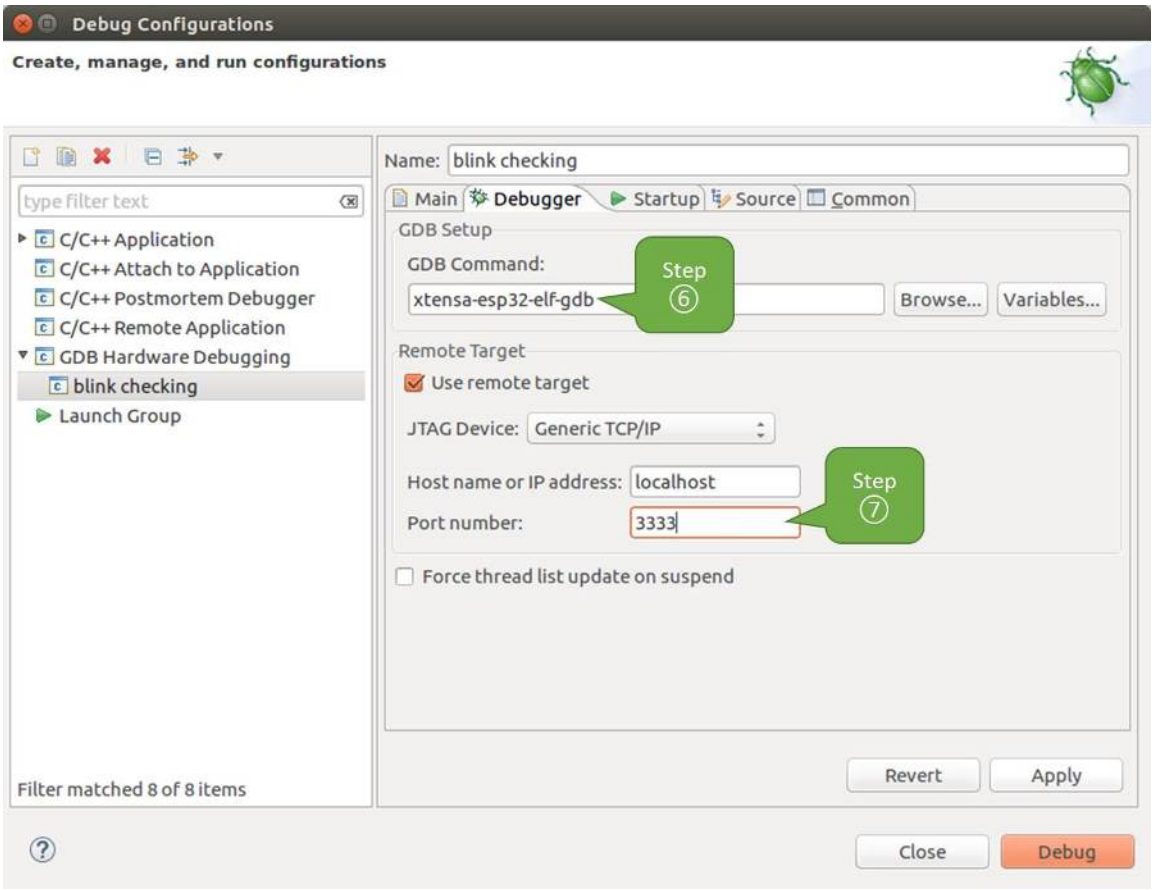

图 9: GDB 硬件调试的配置 - Debugger 选项卡

有关 `program_esp` 命令的说明请参考上传待调试的应用程序 章节。

9. 在 “Load Image and Symbols” 下，取消选中 “Load image” 选项。
10. 在同一个选项卡中继续往下浏览，建立一个初始断点用来在调试器复位后暂停 CPU。插件会根据 “Set break point at:” 一栏中输入的函数名，在该函数的开头设置断点。选中这一选项，并在相应的字段中输入 `app_main`。
11. 选中 “Resume” 选项，这会使得程序在每次调用步骤 8 中的 `mon reset halt` 之后恢复，然后在 `app_main` 的断点处停止。
    上述步骤 8 - 11 的示例输入如下图所示。

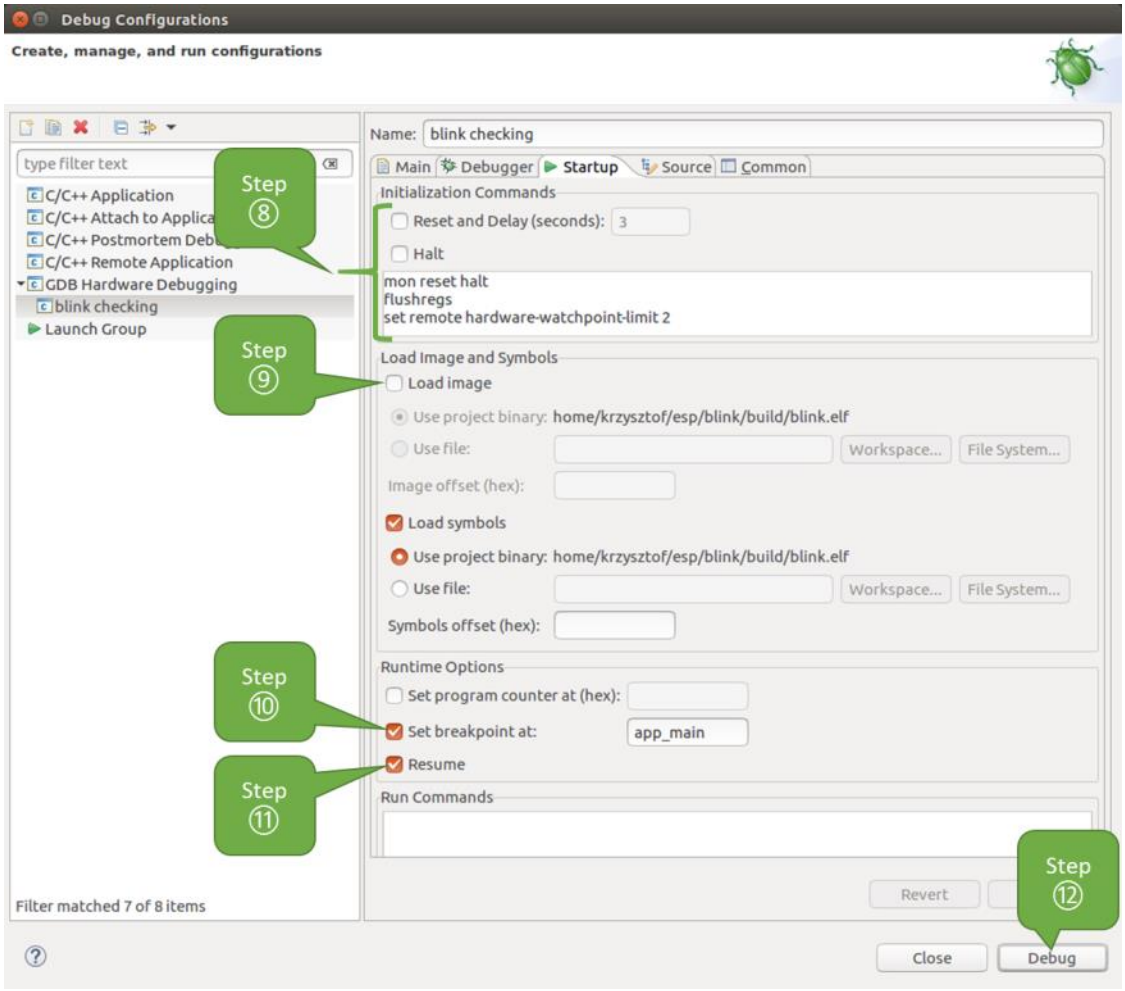

图 10: GDB 硬件调试的配置 - Startup 选项卡

上面的启动序列看起来有些复杂，如果你对其中的初始化命令不太熟悉，请查阅调试器的启动命令的含义 章节获取更多说明。

12. 如果你前面已经完成配置 *ESP32-S2* 目标板 中介绍的步骤，那么目标正在运行并准备与调试器进行对话。按下 “Debug” 按钮就可以直接调试。否则请按下 “Apply” 按钮保存配置，返回配置 *ESP32-S2* 目标板 章节进行配置，最后再回到这里开始调试。

一旦所有 1 - 12 的配置步骤都已经完成，Eclipse 就会打开 “Debug” 视图，如下图所示。

如果你不太了解 GDB 的常用方法，请查阅使用 *Eclipse* 的调试示例 文章中的调试示例章节调试范例。

在命令行中使用 GDB

1. 为了能够启动调试会话，需要先启动并运行目标，如果还没有完成，请按照配置 *ESP32-S2* 目标板 中的介绍进行操作。
2. 打开一个新的终端会话并前往待调试的项目目录，比如：

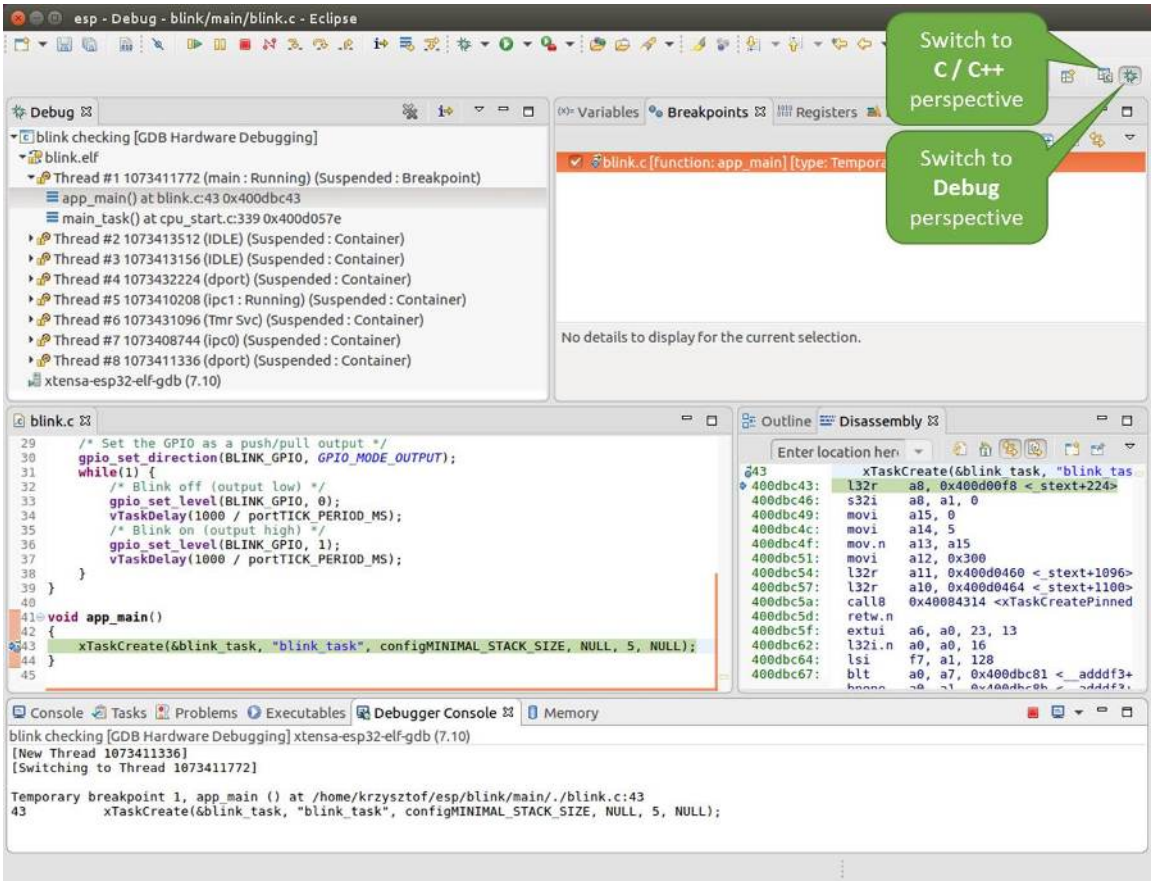

图 11: Eclipse 中的调试视图

```
cd ~/esp/blink
```

3. 当启动调试器时，通常需要提供几个配置参数和命令，为了避免每次都在命令行中逐行输入这些命令，我们可以新建一个配置文件，并将其命名为 gdbinit:

```
target remote :3333
set remote hardware-watchpoint-limit 2
mon reset halt
flushregs
thb app_main
c
```

将此文件保存在当前目录中。
有关 gdbinit 文件内部的更多详细信息，请参阅调试器的启动命令的含义 章节。

4. 准备好启动 GDB，请在终端中输入以下内容：

```
xtensa-esp32s2-elf-gdb -x gdbinit build/blink.elf
```

5. 如果前面的步骤已经正确完成，你会看到如下所示的输出日志，在日志的最后会出现 (gdb) 提示符：

```
user-name@computer-name:~/esp/blink$ xtensa-esp32s2-elf-gdb -x gdbinit build/
↪blink.elf
GNU gdb (crosstool-NG crosstool-ng-1.22.0-61-gab8375a) 7.10
Copyright (C) 2015 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law.  Type "show copying"
and "show warranty" for details.
This GDB was configured as "--host=x86_64-build_pc-linux-gnu --target=xtensa-
↪esp32s2-elf".
Type "show configuration" for configuration details.
For bug reporting instructions, please see:
<http://www.gnu.org/software/gdb/bugs/>.
Find the GDB manual and other documentation resources online at:
<http://www.gnu.org/software/gdb/documentation/>.
For help, type "help".
Type "apropos word" to search for commands related to "word"...
Reading symbols from build/blink.elf...done.
0x400d10d8 in esp_vApplicationIdleHook () at /home/user-name/esp/esp-idf/
↪components/esp32s2/./freertos_hooks.c:52
52          asm("waiti 0");
JTAG tap: esp32s2.cpu0 tap/device found: 0x120034e5 (mfg: 0x272 (Tensilica),␣
↪part: 0x2003, ver: 0x1)
JTAG tap: esp32s2.slave tap/device found: 0x120034e5 (mfg: 0x272 (Tensilica),␣
↪part: 0x2003, ver: 0x1)
esp32s2: Debug controller was reset (pwrstat=0x5F, after clear 0x0F).
esp32s2: Core was reset (pwrstat=0x5F, after clear 0x0F).
esp32s2 halted. PRO_CPU: PC=0x5000004B (active)    APP_CPU: PC=0x00000000
esp32s2: target state: halted
esp32s2: Core was reset (pwrstat=0x1F, after clear 0x0F).
Target halted. PRO_CPU: PC=0x40000400 (active)    APP_CPU: PC=0x40000400
esp32s2: target state: halted
Hardware assisted breakpoint 1 at 0x400db717: file /home/user-name/esp/blink/
↪main/./blink.c, line 43.
0x0:    0x00000000
Target halted. PRO_CPU: PC=0x400DB717 (active)    APP_CPU: PC=0x400D10D8
[New Thread 1073428656]
[New Thread 1073413708]
[New Thread 1073431316]
[New Thread 1073410672]
```

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```
[New Thread 1073408876]
[New Thread 1073432196]
[New Thread 1073411552]
[Switching to Thread 1073411996]

Temporary breakpoint 1, app_main () at /home/user-name/esp/blink/main/./blink.
↪c:43
43      xTaskCreate(&blink_task, "blink_task", 512, NULL, 5, NULL);
(gdb)
```

注意上面日志的倒数第三行显示了调试器已经在 `app_main()` 函数的断点处停止，该断点在 `gdbinit` 文件中设定。由于处理器已经暂停运行，LED 也不会闪烁。如果这也是你看到的现象，你可以开始调试了。

如果你不太了解 GDB 的常用方法，请查阅*使用命令行的调试示例* 文章中的调试示例章节*调试范例*。

使用 idf.py 进行调试 我们还可以使用 `idf.py` 更方便地执行上述提到的调试命令：

1. `idf.py openocd`
   在终端中运行 OpenOCD，其配置信息来源于环境变量或者命令行。默认会使用 `OPENOCD_SCRIPTS` 环境变量中指定的脚本路径，它是由 ESP-IDF 项目仓库中的导出脚本（`export.sh` or `export.bat`）添加到系统环境变量中的。当然，我们可以在命令行中通过 `--openocd-scripts` 来覆盖这个变量的值。
   你可以定义 `OPENOCD_COMMANDS` 环境变量来指定当前开发板的 JTAG 配置，或者通过 `--openocd-commands` 传递该参数。如果这两者都没有被定义，那么 OpenOCD 会使用 `-f board/esp32s2-kaluga-1.cfg` 参数来启动。
2. `idf.py gdb`
   根据当前项目的 elf 文件自动生成 gdb 启动脚本，然后会按照*在命令行中使用 GDB* 中所描述的步骤启动 GDB。
3. `idf.py gdbtui`
   和步骤 2 相同，但是会在启动 GDB 的时候传递 `tui` 参数，这样可以方便在调试过程中查看源代码。
4. `idf.py gdbgui`
   启动 gdbgui，在浏览器中打开调试器的前端界面。
   上述这些命令也可以合并到一起使用，`idf.py` 会自动将后台进程（比如 openocd）最先运行，交互式进程（比图 gdb，monitor）最后运行。
   常用的组合命令如下所示：

   ```
   idf.py openocd gdbgui monitor
   ```

   上述命令会将 OpenOCD 运行至后台，然后启动 gdbgui 打开一个浏览器窗口，显示调试器的前端界面，最后在活动终端打开串口监视器。

调试示例

本节将介绍如何在*Eclipse* 和*命令行* 中使用 GDB 进行调试的示例。

使用 Eclipse 的调试示例 请检查目标板是否已经准备好，并加载了 get-started/blink 示例代码，然后按照*在 Eclipse 中使用 GDB* 中介绍的步骤配置和启动调试器，最后选择让应用程序在 `app_main()` 建立的断点处停止。

本小节的示例

1. *浏览代码，查看堆栈和线程*
2. *设置和清除断点*
3. *手动暂停目标*
4. *单步执行代码*
5. *查看并设置内存*

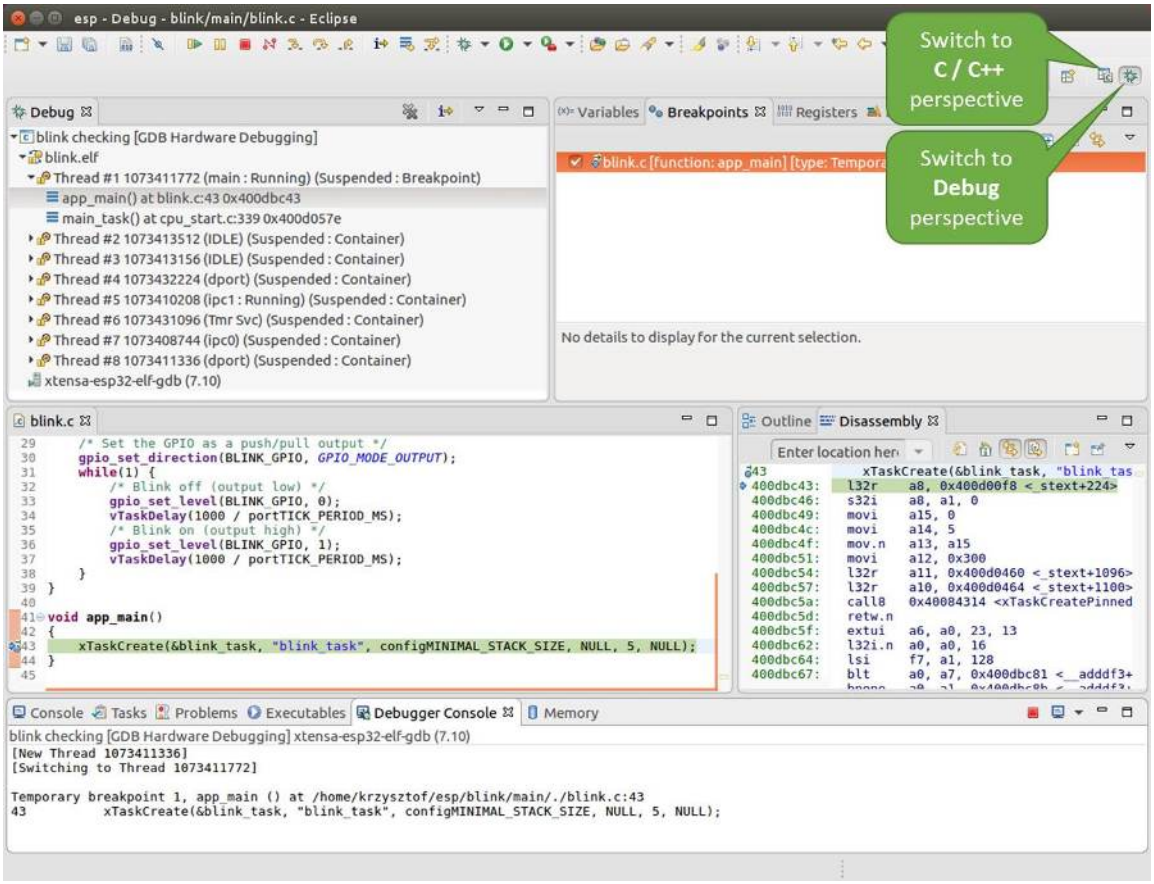

图 12: Eclipse 中的 Debug 视图

6. 观察和设置程序变量
7. 设置条件断点

浏览代码，查看堆栈和线程 当目标暂停时，调试器会在“Debug”窗口中显示线程的列表，程序暂停的代码行在下面的另一个窗口中被高亮显示，如下图所示。此时板子上的 LED 停止了闪烁。

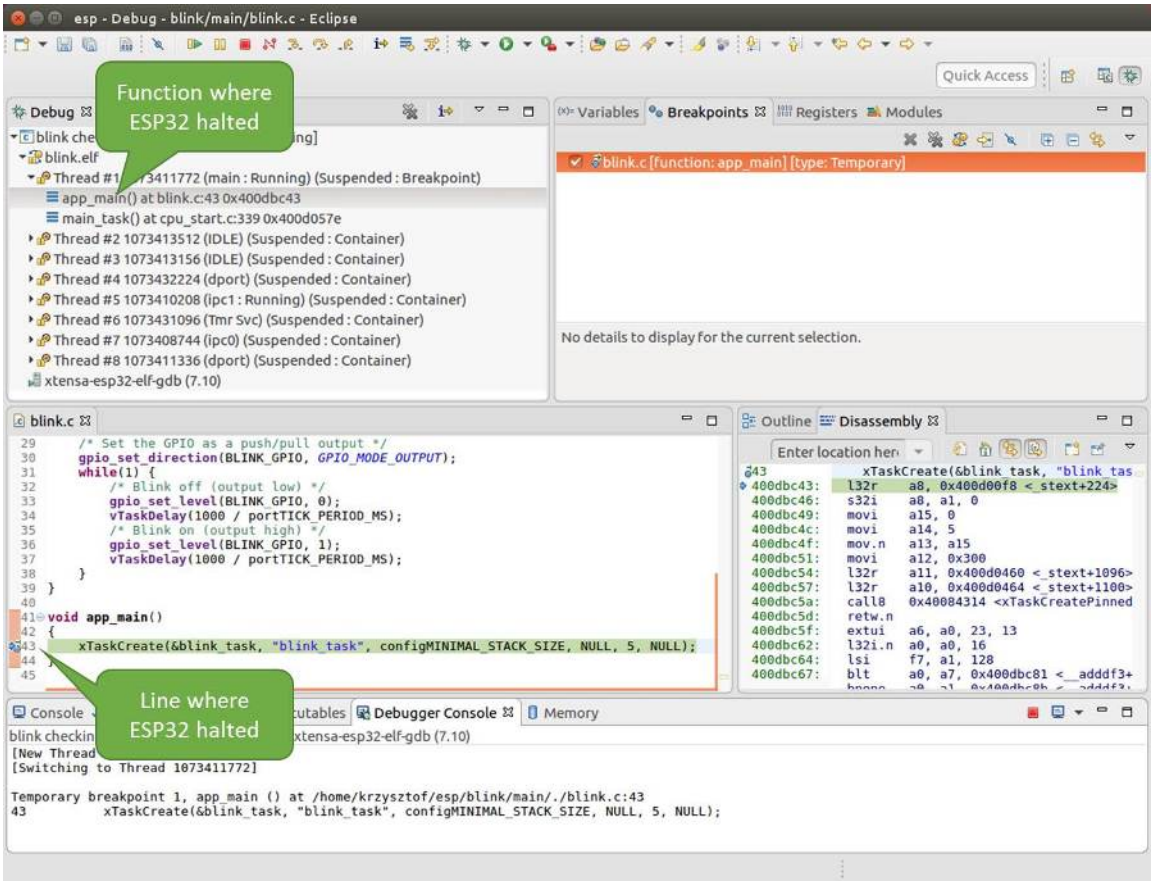

图 13: 调试时目标停止

暂停的程序所在线程也会被展开，显示函数调用的堆栈，它表示直到目标暂停所在代码行（下图高亮处）为止的相关函数的调用关系。1 号线程下函数调用堆栈的第一行包含了最后一个调用的函数 `app_main()`，根据下一行显示，它又是在函数 `main_task()` 中被调用的。堆栈的每一行还包含调用函数的文件名和行号。通过单击每个堆栈的条目，在下面的窗口中，你将看到此文件的内容。

通过展开线程，你可以浏览整个应用程序。展开 5 号线程，它包含了更长的函数调用堆栈，你可以看到函数调用旁边的数字，比如 `0x4000000c`，它们代表未以源码形式提供的二进制代码所在的内存地址。

无论项目是以源代码还是仅以二进制形式提供，在右边一个窗口中，都可以看到反汇编后的机器代码。

回到 1 号线程中的 `app_main()` 函数所在的 `blink.c` 源码文件，下面的示例将会以该文件为例介绍调试的常用功能。调试器可以轻松浏览整个应用程序的代码，这给单步调试代码和设置断点带来了很大的便利，下面将一一展开讨论。

设置和清除断点 在调试时，我们希望能够在关键的代码行停止应用程序，然后检查特定的变量、内存、寄存器和外设的状态。为此我们需要使用断点，以便在特定某行代码处快速访问和停止应用程序。

我们在控制 LED 状态发生变化的两处代码行分别设置一个断点。基于以上代码列表，这两处分别为第 33 和 36 代码行。按住键盘上的“Control”键，双击 blink.c 文件中的行号 33，并在弹出的对话框中点击“OK”按钮进行确定。如果你不想看到此对话框，双击行号即可。执行同样操作，在第 36 行设置另外一个断点。

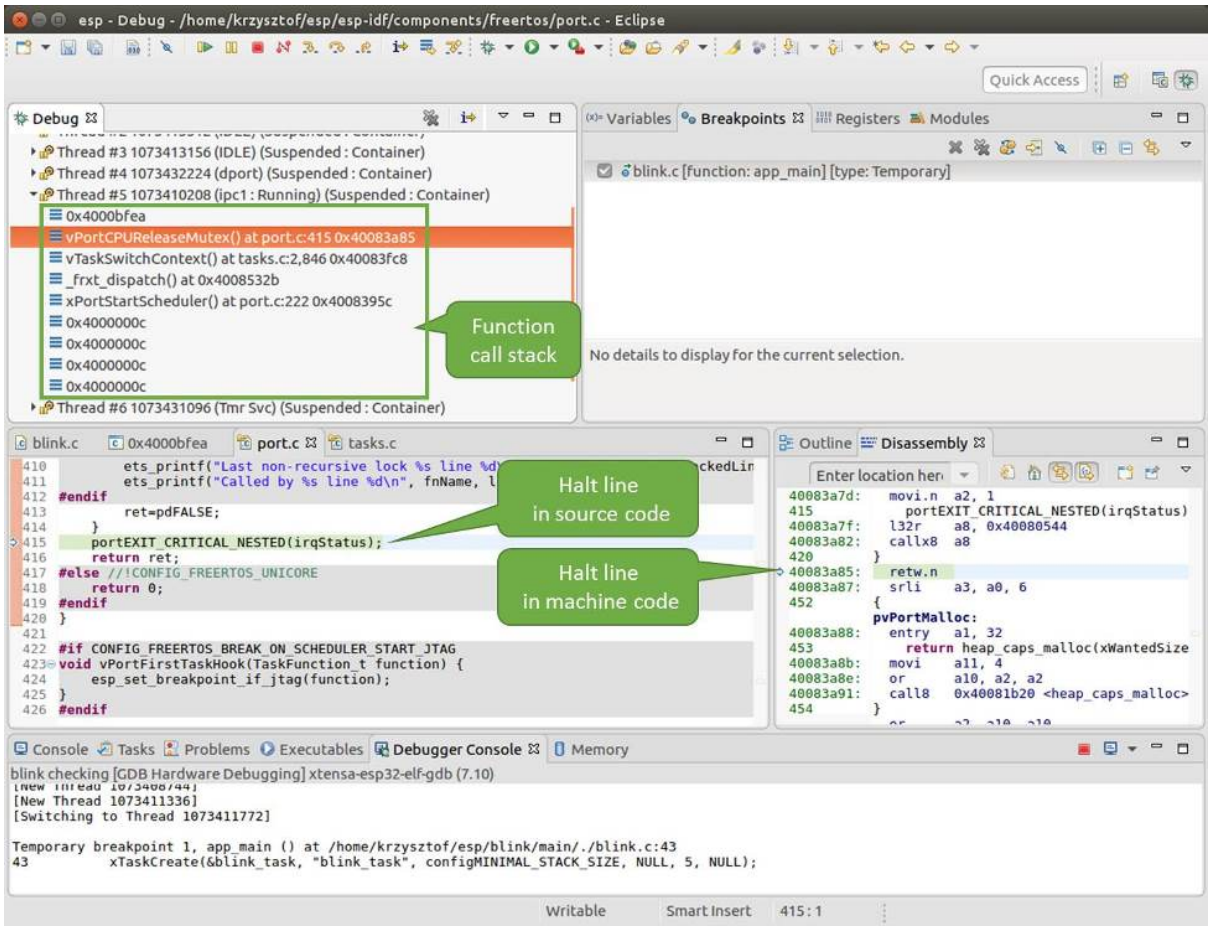

图 14: 浏览函数调用堆栈

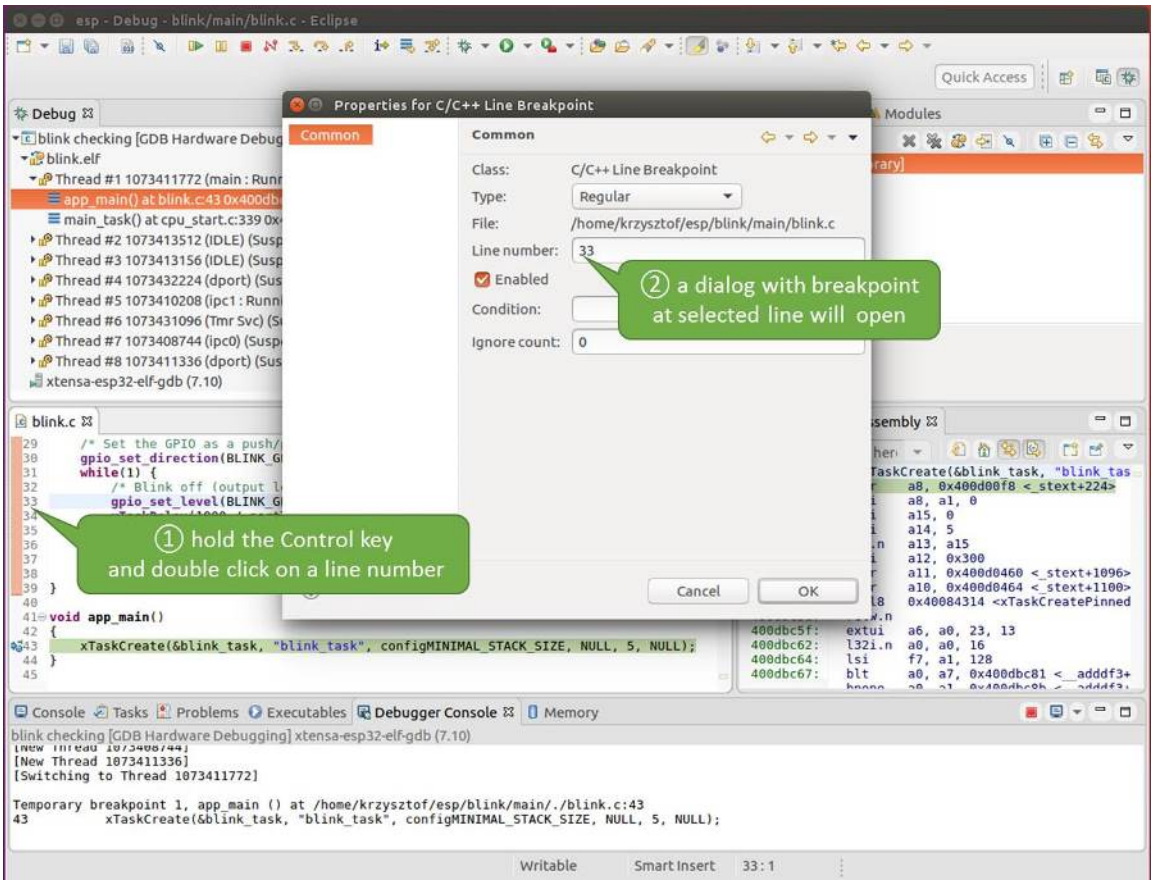
esp - Debug - blink/main/blink.c - Eclipse
Quick Access
Debug
Properties for C/C++ Line Breakpoint
Common
Class: C/C++ Line Breakpoint
Type: Regular
File: /home/krzysztof/esp/blink/main/blink.c
Line number: 33
Enabled
Condition:
Ignore count: 0
② a dialog with breakpoint at selected line will open
① hold the Control key and double click on a line number
Cancel
OK
blink.c
Console
Tasks
Problems
Executables
Debugger Console
Memory
[New Thread 1073411336]
[Switching to Thread 1073411772]
Temporary breakpoint 1, app_main () at /home/krzysztof/esp/blink/main/./blink.c:43
43 xTaskCreate(&blink_task, "blink_task", configMINIMAL_STACK_SIZE, NULL, 5, NULL);
Writable
Smart Insert
33 : 1

图 15: 设置断点

断点的数量和位置信息会显示在右上角的“断点”窗口中。单击“Show Breakpoints Supported by Selected Target”图标可以刷新此列表。除了刚才设置的两个断点外，列表中可能还包含在调试器启动时设置在 `app_main()` 函数处的临时断点。由于最多只允许设置两个断点（详细信息请参阅可用的断点和观察点），你需要将其删除，否则调试会失败。

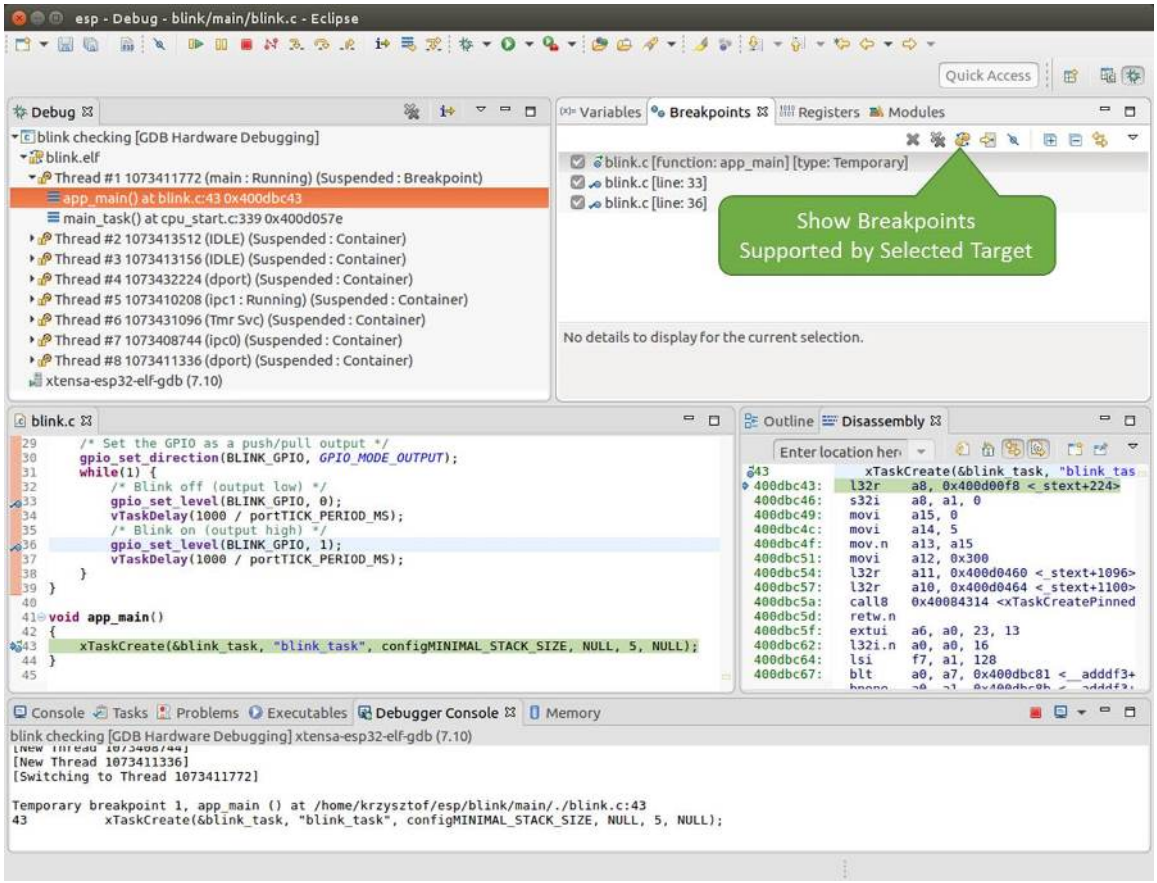

图 16: 设置了三个断点 / 最多允许两个断点

单击“Resume”（如果“Resume”按钮是灰色的，请先单击 8 号线程的 `blink_task()` 函数）后处理器将开始继续运行，并在断点处停止。再一次单击“Resume”按钮，使程序再次运行，然后停在第二个断点处，依次类推。

每次单击“Resume”按钮恢复程序运行后，都会看到 LED 切换状态。

更多关于断点的信息，请参阅可用的断点和观察点 和关于断点的补充知识。

手动暂停目标 在调试时，你可以恢复程序运行并输入代码等待某个事件发生或者保持无限循环而不设置任何断点。后者，如果想要返回调试模式，可以通过单击“Suspend”按钮来手动中断程序的运行。

在此之前，请删除所有的断点，然后单击“Resume”按钮。接着单击“Suspend”按钮，应用程序会停止在某个随机的位置，此时 LED 也将停止闪烁。调试器将展开线程并高亮显示停止的代码行。

在上图所示的情况中，应用程序已经在 `freertos_hooks.c` 文件的第 52 行暂停运行，现在你可以通过单击“Resume”按钮再次将其恢复运行或者进行下面要介绍的调试工作。

单步执行代码 我们还可以使用“Step Into (F5)”和“Step Over (F6)”命令单步执行代码，这两者之间的区别是执行“Step Into (F5)”命令会进入调用的子程序，而执行“Step Over (F6)”命令则会直接将子程序看成单个源码行，单步就能将其运行结束。

在继续演示此功能之前，请参照上文所述确保目前只在 `blink.c` 文件的第 36 行设置了一个断点。

按下 F8 键让程序继续运行然后在断点处停止运行，多次按下“Step Over (F6)”按钮，观察调试器是如何单步执行一行代码的。

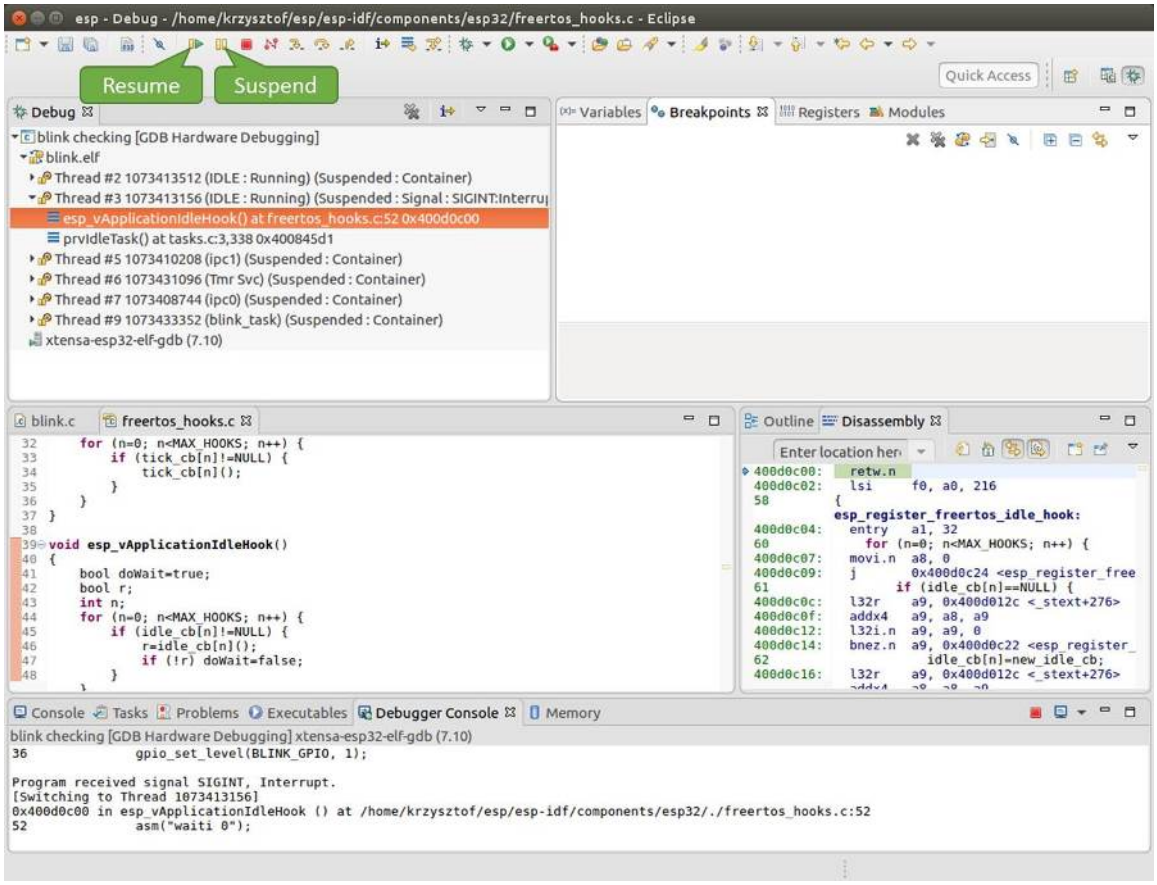
esp - Debug - /home/krzysztof/esp/esp-idf/components/esp32/freertos_hooks.c - Eclipse
Resume
Suspend
Quick Access
Debug
blink checking [GDB Hardware Debugging]
blink.elf
Thread #2 1073413512 (IDLE : Running) (Suspended : Container)
Thread #3 1073413156 (IDLE : Running) (Suspended : Signal : SIGINT:Interru
esp_vApplicationIdleHook() at freertos_hooks.c:52 0x400d0c00
prvIdleTask() at tasks.c:3,338 0x400845d1
Thread #5 1073410208 (ipc1) (Suspended : Container)
Thread #6 1073431096 (Tmr Svc) (Suspended : Container)
Thread #7 1073408744 (ipc0) (Suspended : Container)
Thread #9 1073433352 (blink_task) (Suspended : Container)
xtensa-esp32-elf-gdb (7.10)
Variables
Breakpoints
Registers
Modules
blink.c
freertos_hooks.c
Outline
Disassembly
Console
Tasks
Problems
Executables
Debugger Console
Memory
blink checking [GDB Hardware Debugging] xtensa-esp32-elf-gdb (7.10)
gpio_set_level(BLINK_GPIO, 1);
Program received signal SIGINT, Interrupt.
[Switching to Thread 1073413156]
0x400d0c00 in esp_vApplicationIdleHook () at /home/krzysztof/esp/esp-idf/components/esp32/./freertos_hooks.c:52
asm("waiti 0");

图 17: 手动暂停目标

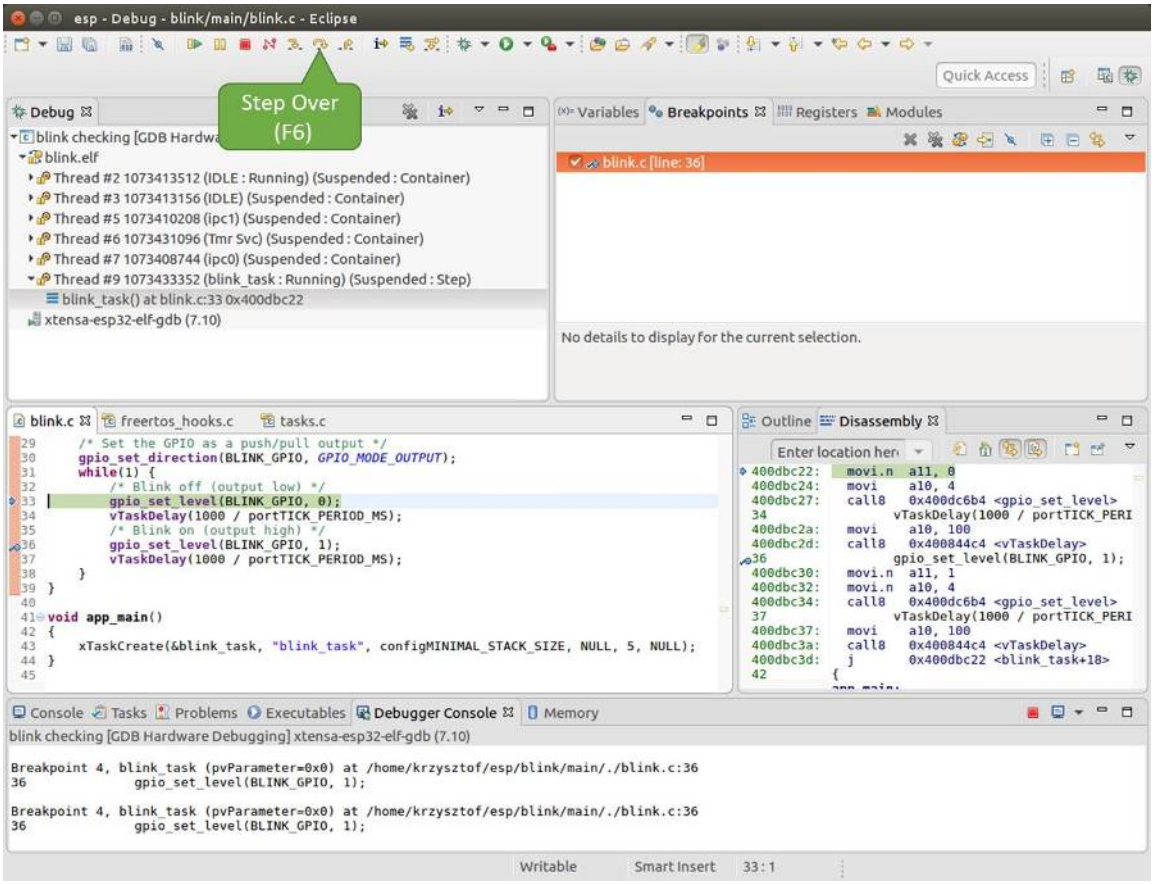

图 18: 使用“Step Over (F6)”单步执行代码

如果你改用“Step Into (F5)”，那么调试器将会进入调用的子程序内部。

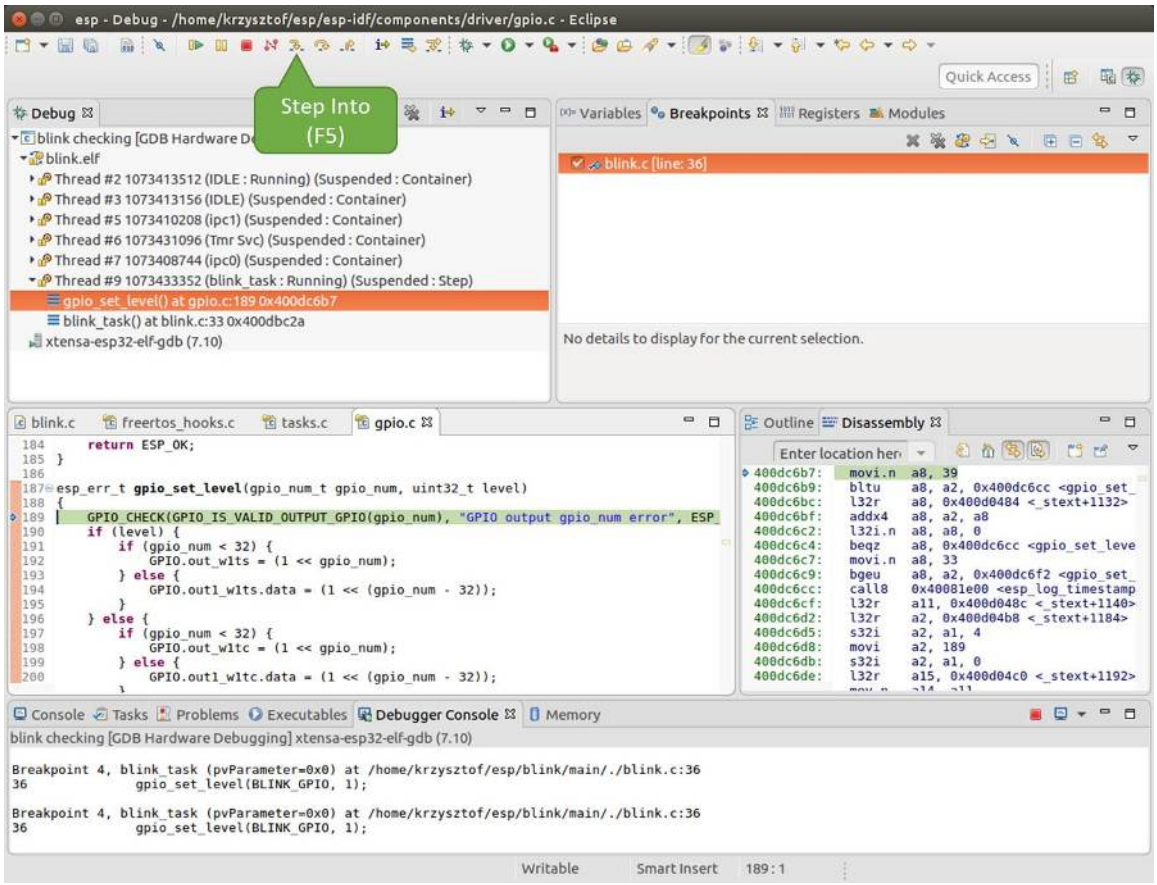

图 19: 使用“Step Into (F5)”单步执行代码

在上述例子中，调试器进入 `gpio_set_level(BLINK_GPIO, 0)` 代码内部，同时代码窗口快速切换到 `gpio.c` 驱动文件。

请参阅 *“next”命令无法跳过子程序的原因* 文档以了解 `next` 命令的潜在局限。

查看并设置内存 要显示或者设置内存的内容，请使用“调试”视图中位于底部的“Memory”选项卡。

在“Memory”选项卡下，我们将在内存地址 `0x3FF44004` 处读取和写入内容。该地址也是 `GPIO_OUT_REG` 寄存器的地址，可以用来控制（设置或者清除）某个 GPIO 的电平。

关于该寄存器的更多详细信息，请参阅 ESP32-S2 技术参考手册 中的 IO_MUX 和 GPIO Matrix 章节。

同样在 `blink.c` 项目文件中，在两个 `gpio_set_level` 语句的后面各设置一个断点，单击“Memory”选项卡，然后单击“Add Memory Monitor”按钮，在弹出的对话框中输入 `0x3FF44004`。

按下 F8 按键恢复程序运行，并观察“Monitor”选项卡。

每按一下 F8，你就会看到在内存 `0x3FF44004` 地址处的一个比特位被翻转（并且 LED 会改变状态）。

要修改内存的数值，请在“Monitor”选项卡中找到待修改的内存地址，如前面观察的结果一样，输入特定比特翻转后的值。当按下回车键后，将立即看到 LED 的状态发生了改变。

观察和设置程序变量 常见的调试任务是在程序运行期间检查程序中某个变量的值，为了演示这个功能，更新 `blink.c` 文件，在 `blink_task` 函数的上面添加一个全局变量的声明 `int i`，然后在 `while(1)` 里添加 `i++`，这样每次 LED 改变状态的时候，变量 `i` 都会增加 1。

退出调试器，这样就不会与新代码混淆，然后重新构建并烧写代码到 ESP32-S2 中，接着重启调试器。注意，这里不需要我们重启 OpenOCD。

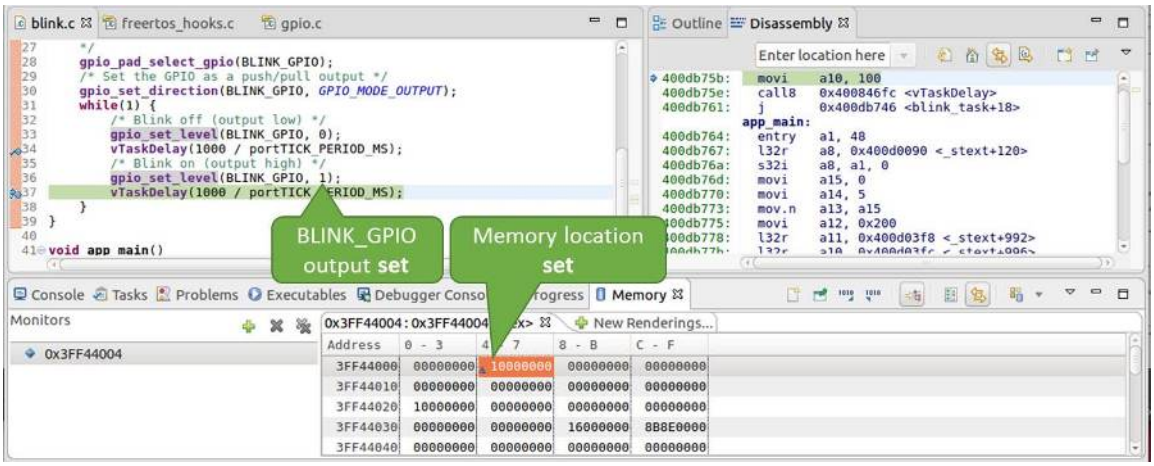

图 20: 观察内存地址 0x3FF44004 处的某个比特被置高

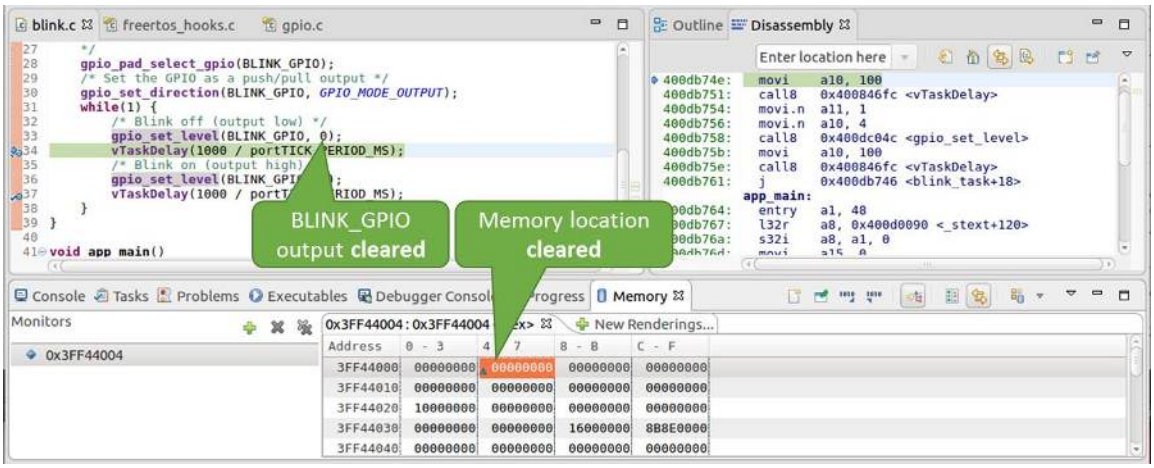

图 21: 观察内存地址 0x3FF44004 处的某个比特被置低

一旦程序停止运行，在代码 `i++` 处添加一个断点。

下一步，在 “Breakpoints” 所在的窗口中，选择 “Expressions” 选项卡。如果该选项卡不存在，请在顶部菜单栏的 Window > Show View > Expressions 中添加这一选项卡。然后在该选项卡中单击 “Add new expression”，并输入 `i`。

按下 F8 继续运行程序，每次程序停止时，都会看到变量 `i` 的值在递增。

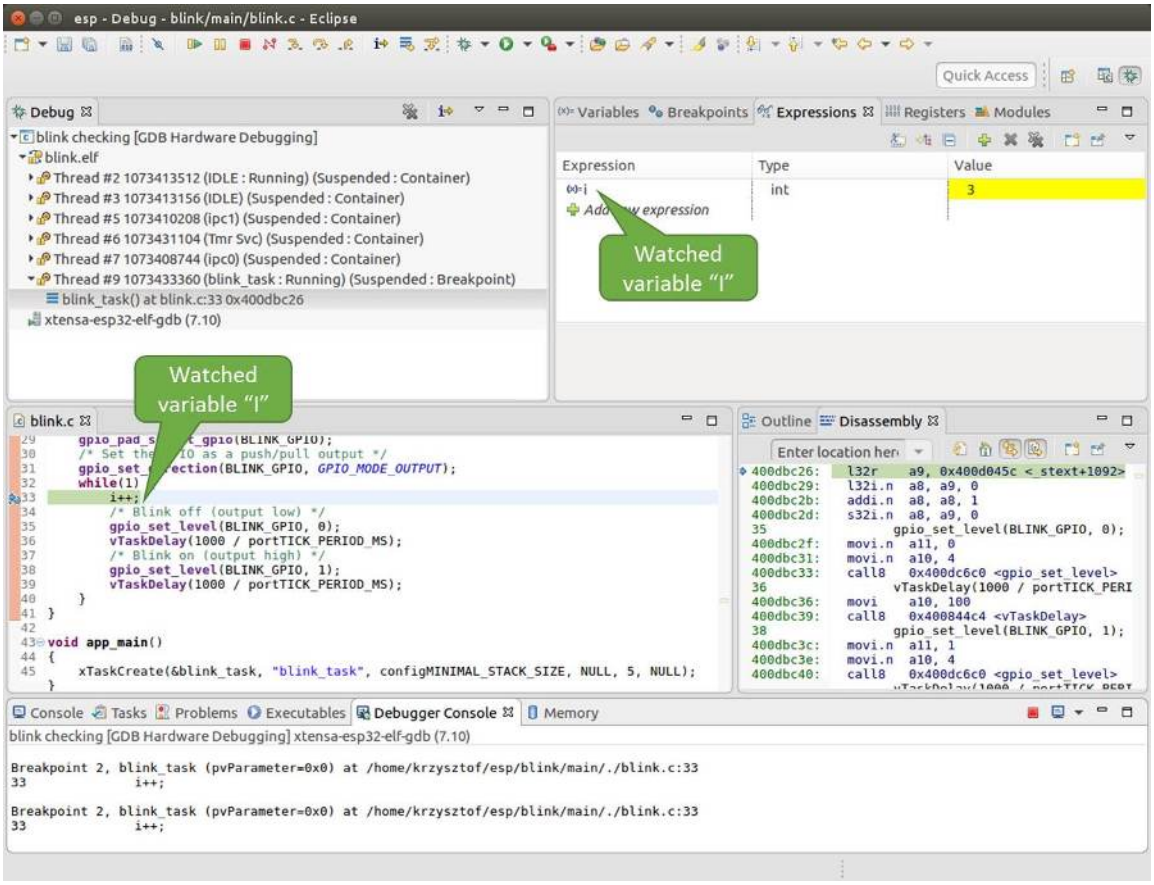

图 22: 观察程序变量 “i”

如想更改 `i` 的值，可以在 “Value” 一栏中输入新的数值。按下 “Resume (F8)” 后，程序将从新输入的数字开始递增 `i`。

设置条件断点 接下来的内容更为有趣，你可能想在一定条件满足的情况下设置断点，然后让程序停止运行。右击断点打开上下文菜单，选择 “Breakpoint Properties”，将 “Type:” 改选为 “Hardware” 然后在 “Condition:” 一栏中输入条件表达式，例如 `i == 2`。

如果当前 `i` 的值小于 2（如果有需要也可以更改这个阈值）并且程序被恢复运行，那么 LED 就会循环闪烁，直到 `i == 2` 条件成立，最后程序停止在该处。

使用命令行的调试示例 请检查您的目标板是否已经准备好，并加载了 get-started/blink 示例代码，然后按照*在命令行中使用 GDB* 中介绍的步骤配置和启动调试器，最后选择让应用程序在 `app_main()` 建立的断点处停止运行

```
Temporary breakpoint 1, app_main () at /home/user-name/esp/blink/main/./blink.c:43
43          xTaskCreate(&blink_task, "blink_task", configMINIMAL_STACK_SIZE, NULL,␣
↪5, NULL);
(gdb)
```

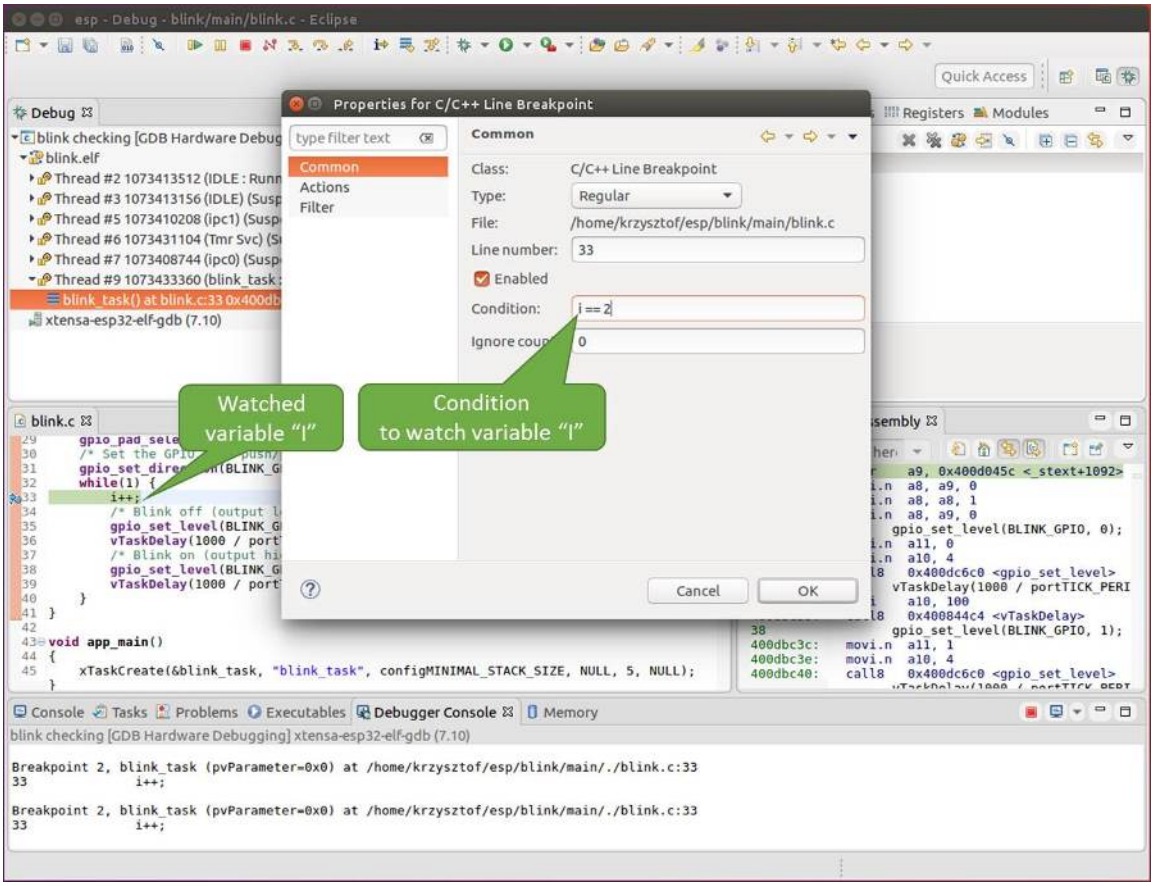
esp - Debug - blink/main/blink.c - Eclipse
Properties for C/C++ Line Breakpoint
Common
Class: C/C++ Line Breakpoint
Type: Regular
File: /home/krzysztof/esp/blink/main/blink.c
Line number: 33
Enabled
Condition: i == 2
Watched variable "i"
Condition to watch variable "i"
Cancel
OK
Debugger Console
Breakpoint 2, blink_task (pvParameter=0x0) at /home/krzysztof/esp/blink/main/./blink.c:33

图 23: 设置条件断点

本小节的示例

1. 浏览代码，查看堆栈和线程
2. 设置和清除断点
3. 暂停和恢复应用程序的运行
4. 单步执行代码
5. 查看并设置内存
6. 观察和设置程序变量
7. 设置条件断点

浏览代码，查看堆栈和线程 当看到 (gdb) 提示符的时候，应用程序已停止运行，LED 也停止闪烁。

要找到代码暂停的位置，输入 `l` 或者 `list` 命令，调试器会打印出停止点（`blink.c` 代码文件的第 43 行）附近的几行代码

```
(gdb) l
38          }
39      }
40
41      void app_main()
42      {
43          xTaskCreate(&blink_task, "blink_task", configMINIMAL_STACK_SIZE, NULL,␣
↪5, NULL);
44      }
(gdb)
```

也可以通过输入 `l 30, 40` 等命令来查看特定行号范围内的代码。

使用 `bt` 或者 `backtrace` 来查看哪些函数最终导致了此代码被调用:

```
(gdb) bt
#0  app_main () at /home/user-name/esp/blink/main/./blink.c:43
#1  0x400d057e in main_task (args=0x0) at /home/user-name/esp/esp-idf/components/
↪esp32s2/./cpu_start.c:339
(gdb)
```

输出的第 0 行表示应用程序暂停之前调用的最后一个函数，即我们之前列出的 `app_main ()`。`app_main ()` 又被位于 `cpu_start.c` 文件第 339 行的 `main_task` 函数调用。

想查看 `cpu_start.c` 文件中 `main_task` 函数的上下文，需要输入 `frame N`，其中 N = 1，因为根据前面的输出，`main_task` 位于 #1 下:

```
(gdb) frame 1
#1  0x400d057e in main_task (args=0x0) at /home/user-name/esp/esp-idf/components/
↪esp32s2/./cpu_start.c:339
339         app_main();
(gdb)
```

输入 `l` 将显示一段名为 `app_main()` 的代码（在第 339 行）：

```
(gdb) l
334             ;
335         }
336     #endif
337         //Enable allocation in region where the startup stacks were located.
338         heap_caps_enable_nonos_stack_heaps();
339         app_main();
340         vTaskDelete(NULL);
341     }
342
(gdb)
```

通过打印前面的一些行，你会看到我们一直在寻找的 `main_task` 函数:

```
(gdb) l 326, 341
326     static void main_task(void* args)
327     {
328         // Now that the application is about to start, disable boot watchdogs
329         REG_CLR_BIT(TIMG_WDTCONFIG0_REG(0), TIMG_WDT_FLASHBOOT_MOD_EN_S);
330         REG_CLR_BIT(RTC_CNTL_WDTCONFIG0_REG, RTC_CNTL_WDT_FLASHBOOT_MOD_EN);
331     #if !CONFIG_FREERTOS_UNICORE
332         // Wait for FreeRTOS initialization to finish on APP CPU, before␣
↪replacing its startup stack
333         while (port_xSchedulerRunning[1] == 0) {
334             ;
335         }
336     #endif
337         //Enable allocation in region where the startup stacks were located.
338         heap_caps_enable_nonos_stack_heaps();
339         app_main();
340         vTaskDelete(NULL);
341     }
(gdb)
```

如果要查看其他代码，可以输入 `i threads` 命令，则会输出目标板上运行的线程列表:

```
(gdb) i threads
  Id   Target Id         Frame
  8    Thread 1073411336 (dport) 0x400d0848 in dport_access_init_core (arg=
↪<optimized out>)
    at /home/user-name/esp/esp-idf/components/esp32s2/./dport_access.c:170
  7    Thread 1073408744 (ipc0) xQueueGenericReceive (xQueue=0x3ffae694,␣
↪pvBuffer=0x0, xTicksToWait=1644638200,
    xJustPeeking=0) at /home/user-name/esp/esp-idf/components/freertos/./queue.
↪c:1452
  6    Thread 1073431096 (Tmr Svc) prvTimerTask (pvParameters=0x0)
    at /home/user-name/esp/esp-idf/components/freertos/./timers.c:445
  5    Thread 1073410208 (ipc1 : Running) 0x4000bfea in ?? ()
  4    Thread 1073432224 (dport) dport_access_init_core (arg=0x0)
    at /home/user-name/esp/esp-idf/components/esp32s2/./dport_access.c:150
  3    Thread 1073413156 (IDLE) prvIdleTask (pvParameters=0x0)
    at /home/user-name/esp/esp-idf/components/freertos/./tasks.c:3282
  2    Thread 1073413512 (IDLE) prvIdleTask (pvParameters=0x0)
    at /home/user-name/esp/esp-idf/components/freertos/./tasks.c:3282
* 1    Thread 1073411772 (main : Running) app_main () at /home/user-name/esp/blink/
↪main/./blink.c:43
(gdb)
```

线程列表显示了每个线程最后一个被调用的函数以及所在的 C 源文件名（如果存在的话）。

您可以通过输入 `thread N` 进入特定的线程，其中 `N` 是线程 ID。我们进入 5 号线程来看一下它是如何工作的:

```
(gdb) thread 5
[Switching to thread 5 (Thread 1073410208)]
#0  0x4000bfea in ?? ()
(gdb)
```

然后查看回溯:

```
(gdb) bt
#0  0x4000bfea in ?? ()
#1  0x40083a85 in vPortCPUReleaseMutex (mux=<optimized out>) at /home/user-name/
↪esp/esp-idf/components/freertos/./port.c:415
#2  0x40083fc8 in vTaskSwitchContext () at /home/user-name/esp/esp-idf/components/
↪freertos/./tasks.c:2846
```

(下页继续)

(续上页)

```
#3  0x4008532b in _frxt_dispatch ()
#4  0x4008395c in xPortStartScheduler () at /home/user-name/esp/esp-idf/components/
↪freertos/./port.c:222
#5  0x4000000c in ?? ()
#6  0x4000000c in ?? ()
#7  0x4000000c in ?? ()
#8  0x4000000c in ?? ()
(gdb)
```

如上所示，回溯可能会包含多个条目，方便查看直至目标停止运行的函数调用顺序。如果找不到某个函数的源码文件，将会使用问号 `??` 替代，这表示该函数是以二进制格式提供的。像 `0x4000bfea` 这样的值是被调用函数所在的内存地址。

使用诸如 `bt`，`i threads`，`thread N` 和 `list` 命令可以浏览整个应用程序的代码。这给单步调试代码和设置断点带来很大的便利，下面将一一展开来讨论。

设置和清除断点 在调试时，我们希望能够在关键的代码行停止应用程序，然后检查特定的变量、内存、寄存器和外设的状态。为此我们需要使用断点，以便在特定某行代码处快速访问和停止应用程序。

我们在控制 LED 状态发生变化的两处代码行分别设置一个断点。基于以上代码列表，这两处分别为第 33 和 36 代码行。使用命令 `break M` 设置断点，其中 M 是具体的代码行:

```
(gdb) break 33
Breakpoint 2 at 0x400db6f6: file /home/user-name/esp/blink/main/./blink.c, line 33.
(gdb) break 36
Breakpoint 3 at 0x400db704: file /home/user-name/esp/blink/main/./blink.c, line 36.
```

输入命令 `c`，处理器将运行并在断点处停止。再次输入 `c` 将使其再次运行，并在第二个断点处停止，依此类推:

```
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB6F6 (active)    APP_CPU: PC=0x400D10D8

Breakpoint 2, blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:33
33          gpio_set_level(BLINK_GPIO, 0);
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB6F8 (active)    APP_CPU: PC=0x400D10D8
Target halted. PRO_CPU: PC=0x400DB704 (active)    APP_CPU: PC=0x400D10D8

Breakpoint 3, blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:36
36          gpio_set_level(BLINK_GPIO, 1);
(gdb)
```

只有在输入命令 `c` 恢复程序运行后才能看到 LED 改变状态。

查看已设置断点的数量和位置，请使用命令 `info break`:

```
(gdb) info break
Num     Type           Disp Enb Address    What
2       breakpoint     keep y   0x400db6f6 in blink_task at /home/user-name/esp/
↪blink/main/./blink.c:33
    breakpoint already hit 1 time
3       breakpoint     keep y   0x400db704 in blink_task at /home/user-name/esp/
↪blink/main/./blink.c:36
    breakpoint already hit 1 time
(gdb)
```

请注意，断点序号（在 Num 栏列出）从 2 开始，这是因为在调试器启动时执行 thb app_main 命令已经在 app_main() 函数处建立了第一个断点。由于它是一个临时断点，已经被自动删除，所以没有被列出。

要删除一个断点，请输入 delete N 命令（或者简写成 d N），其中 N 代表断点序号:

```
(gdb) delete 1
No breakpoint number 1.
(gdb) delete 2
(gdb)
```

更多关于断点的信息，请参阅可用的断点和观察点 和关于断点的补充知识。

暂停和恢复应用程序的运行 在调试时，可以恢复程序运行并输入代码等待某个事件发生或者保持无限循环而不设置任何断点。对于后者，想要返回调试模式，可以通过输入 Ctrl+C 手动中断程序的运行。

在此之前，请删除所有的断点，然后输入 c 恢复程序运行。接着输入 Ctrl+C，应用程序会停止在某个随机的位置，此时 LED 也将停止闪烁。调试器会打印如下信息:

```
(gdb) c
Continuing.
^CTarget halted. PRO_CPU: PC=0x400D0C00             APP_CPU: PC=0x400D0C00 (active)
[New Thread 1073433352]

Program received signal SIGINT, Interrupt.
[Switching to Thread 1073413512]
0x400d0c00 in esp_vApplicationIdleHook () at /home/user-name/esp/esp-idf/
↪components/esp32s2/./freertos_hooks.c:52
52              asm("waiti 0");
(gdb)
```

在上图所示的情况下，应用程序已经在 freertos_hooks.c 文件的第 52 行暂停运行，现在您可以通过输入 c 再次将其恢复运行或者进行如下所述的一些调试工作。

注解: 在 MSYS2 的 shell 中输入 Ctrl+C 并不会暂停目标的运行，而是会退出调试器。解决这个问题的方法可以通过使用 *Eclipse* 来调试 或者参考 http://www.mingw.org/wiki/Workaround_for_GDB_Ctrl_C_Interrupt 里的解决方案。

单步执行代码 我们还可以使用 step 和 next 命令（可以简写成 s 和 n）单步执行代码，这两者之间的区别是执行 “step” 命令会进入调用的子程序内部，而执行 “next” 命令则会直接将子程序看成单个源码行，单步就能将其运行结束。

在继续演示此功能之前，请使用前面介绍的 break 和 delete 命令，确保目前只在 blink.c 文件的第 36 行设置了一个断点:

```
(gdb) info break
Num     Type           Disp Enb Address    What
3       breakpoint     keep y   0x400db704 in blink_task at /home/user-name/esp/
↪blink/main/./blink.c:36
    breakpoint already hit 1 time
(gdb)
```

输入 c 恢复程序运行然后等它在断点处停止运行:

```
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB754 (active)    APP_CPU: PC=0x400D1128

Breakpoint 3, blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:36
```

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```
36          gpio_set_level(BLINK_GPIO, 1);
(gdb)
```

然后输入 n 多次，观察调试器是如何单步执行一行代码的:

```
(gdb) n
Target halted. PRO_CPU: PC=0x400DB756 (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB758 (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DC04C (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB75B (active)    APP_CPU: PC=0x400D1128
37          vTaskDelay(1000 / portTICK_PERIOD_MS);
(gdb) n
Target halted. PRO_CPU: PC=0x400DB75E (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400846FC (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB761 (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB746 (active)    APP_CPU: PC=0x400D1128
33          gpio_set_level(BLINK_GPIO, 0);
(gdb)
```

如果你输入 s，那么调试器将进入子程序:

```
(gdb) s
Target halted. PRO_CPU: PC=0x400DB748 (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB74B (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DC04C (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DC04F (active)    APP_CPU: PC=0x400D1128
gpio_set_level (gpio_num=GPIO_NUM_4, level=0) at /home/user-name/esp/esp-idf/
↪components/driver/./gpio.c:183
183     GPIO_CHECK(GPIO_IS_VALID_OUTPUT_GPIO(gpio_num), "GPIO output gpio_num error
↪", ESP_ERR_INVALID_ARG);
(gdb)
```

上述例子中，调试器进入 gpio_set_level(BLINK_GPIO, 0) 代码内部，同时代码窗口快速切换到 gpio.c 驱动文件。

请参阅 *"next" 命令无法跳过子程序的原因* 文档以了解 next 命令的潜在局限。

查看并设置内存 使用命令 x 可以显示内存的内容，配合其余参数还可以调整所显示内存位置的格式和数量。运行 help x 可以查看更多相关细节。与 x 命令配合使用的命令是 set，它允许你将值写入内存。

为了演示 x 和 set 的使用，我们将在内存地址 0x3FF44004 处读取和写入内容。该地址也是 GPIO_OUT_REG 寄存器的地址，可以用来控制（设置或者清除）某个 GPIO 的电平。关于该寄存器的更多详细信息，请参阅 ESP32-S2 技术参考手册 中的 IO_MUX 和 GPIO Matrix 章节。

同样在 blink.c 项目文件中，在两个 gpio_set_level 语句的后面各设置一个断点。输入两次 c 命令后停止在断点处，然后输入 x /1wx 0x3FF44004 来显示 GPIO_OUT_REG 寄存器的值:

```
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB75E (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB74E (active)    APP_CPU: PC=0x400D1128

Breakpoint 2, blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:34
34          vTaskDelay(1000 / portTICK_PERIOD_MS);
(gdb) x /1wx 0x3FF44004
0x3ff44004: 0x00000000
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB751 (active)    APP_CPU: PC=0x400D1128
Target halted. PRO_CPU: PC=0x400DB75B (active)    APP_CPU: PC=0x400D1128
```

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```
Breakpoint 3, blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:37
37          vTaskDelay(1000 / portTICK_PERIOD_MS);
(gdb) x /1wx 0x3FF44004
0x3ff44004: 0x00000010
(gdb)
```

如果闪烁的 LED 连接到了 GPIO4，那么每次 LED 改变状态时你会看到第 4 比特被翻转:

```
0x3ff44004: 0x00000000
...
0x3ff44004: 0x00000010
```

现在，当 LED 熄灭时，与之对应地会显示 `0x3ff44004: 0x00000000`，尝试使用 `set` 命令向相同的内存地址写入 `0x00000010` 来将该比特置高:

```
(gdb) x /1wx 0x3FF44004
0x3ff44004: 0x00000000
(gdb) set {unsigned int}0x3FF44004=0x000010
```

在输入 `set {unsigned int}0x3FF44004=0x000010` 命令后，你会立即看到 LED 亮起。

观察和设置程序变量 常见的调试任务是在程序运行期间检查程序中某个变量的值，为了能够演示这个功能，更新 `blink.c` 文件，在 `blink_task` 函数的上面添加一个全局变量的声明 `int i`，然后在 `while(1)` 里添加 `i++`，这样每次 LED 改变状态的时候，变量 `i` 都会增加 1。

退出调试器，这样就不会与新代码混淆，然后重新构建并烧写代码到 ESP32-S2 中，接着重启调试器。注意，这里不需要我们重启 OpenOCD。

一旦程序停止运行，输入命令 `watch i`:

```
(gdb) watch i
Hardware watchpoint 2: i
(gdb)
```

这会在所有变量 `i` 发生改变的代码处插入所谓的 “观察点”。现在输入 `continue` 命令来恢复应用程序的运行并观察它停止:

```
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB751 (active)    APP_CPU: PC=0x400D0811
[New Thread 1073432196]

Program received signal SIGTRAP, Trace/breakpoint trap.
[Switching to Thread 1073432196]
0x400db751 in blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:33
33          i++;
(gdb)
```

多次恢复程序运行后，变量 `i` 的值会增加，现在你可以输入 `print i`（简写 `p i`）来查看当前 `i` 的值:

```
(gdb) p i
$1 = 3
(gdb)
```

要修改 `i` 的值，请使用 `set` 命令，如下所示（可以将其打印输出来查看是否确已修改）:

```
(gdb) set var i = 0
(gdb) p i
$3 = 0
(gdb)
```

最多可以使用两个观察点，详细信息请参阅可用的断点和观察点。

设置条件断点 接下来的内容更为有趣，你可能想在一定条件满足的情况下设置断点。请先删除已有的断点，然后尝试如下命令:

```
(gdb) break blink.c:34 if (i == 2)
Breakpoint 3 at 0x400db753: file /home/user-name/esp/blink/main/./blink.c, line 34.
(gdb)
```

以上命令在 `blink.c` 文件的 `34` 处设置了一个条件断点，当 `i==2` 条件满足时，程序会停止运行。

如果当前 `i` 的值小于 `2` 并且程序被恢复运行，那么 LED 就会循环闪烁，直到 `i == 2` 条件成立，最后程序停止在该处:

```
(gdb) set var i = 0
(gdb) c
Continuing.
Target halted. PRO_CPU: PC=0x400DB755 (active)    APP_CPU: PC=0x400D112C
Target halted. PRO_CPU: PC=0x400DB753 (active)    APP_CPU: PC=0x400D112C
Target halted. PRO_CPU: PC=0x400DB755 (active)    APP_CPU: PC=0x400D112C
Target halted. PRO_CPU: PC=0x400DB753 (active)    APP_CPU: PC=0x400D112C

Breakpoint 3, blink_task (pvParameter=0x0) at /home/user-name/esp/blink/main/./
↪blink.c:34
34          gpio_set_level(BLINK_GPIO, 0);
(gdb)
```

获得命令的帮助信息 目前所介绍的都是些非常基础的命令，目的在于让您快速上手 JTAG 调试。如果想获得特定命令的语法和功能相关的信息，请在 `(gdb)` 提示符下输入 `help` 和命令名:

```
(gdb) help next
Step program, proceeding through subroutine calls.
Usage: next [N]
Unlike "step", if the current source line calls a subroutine,
this command does not enter the subroutine, but instead steps over
the call, in effect treating it as a single source line.
(gdb)
```

只需输入 `help` 命令，即可获得高级命令列表，帮助你了解更多详细信息。此外，还可以参考一些 GDB 命令速查表，比如 https://darkdust.net/files/GDB%20Cheat%20Sheet.pdf。虽然不是所有命令都适用于嵌入式环境，但还是会有所裨益。

结束调试会话 输入命令 `q` 可以退出调试器:

```
(gdb) q
A debugging session is active.

    Inferior 1 [Remote target] will be detached.

Quit anyway? (y or n) y
Detaching from program: /home/user-name/esp/blink/build/blink.elf, Remote target
Ending remote debugging.
user-name@computer-name:~/esp/blink$
```

应用层跟踪库

概述 为了分析应用程序的行为，IDF 提供了一个有用的功能：应用层跟踪。这个功能以库的形式提供，可以通过 menuconfig 开启。此功能使得用户可以在程序运行开销很小的前提下，通过 JTAG 接口在主机和 ESP32-S2 之间传输任意数据。

开发人员可以使用这个功能库将应用程序的运行状态发送给主机，在运行时接收来自主机的命令或者其他类型的信息。该库的主要使用场景有：

1. 收集应用程序特定的数据，具体请参阅*特定应用程序的跟踪*
2. 轻量级的日志记录，具体请参阅*记录日志到主机*
3. 系统行为分析，具体请参阅*基于 SEGGER SystemView 的系统行为分析*

使用 JTAG 接口的跟踪组件工作示意图：

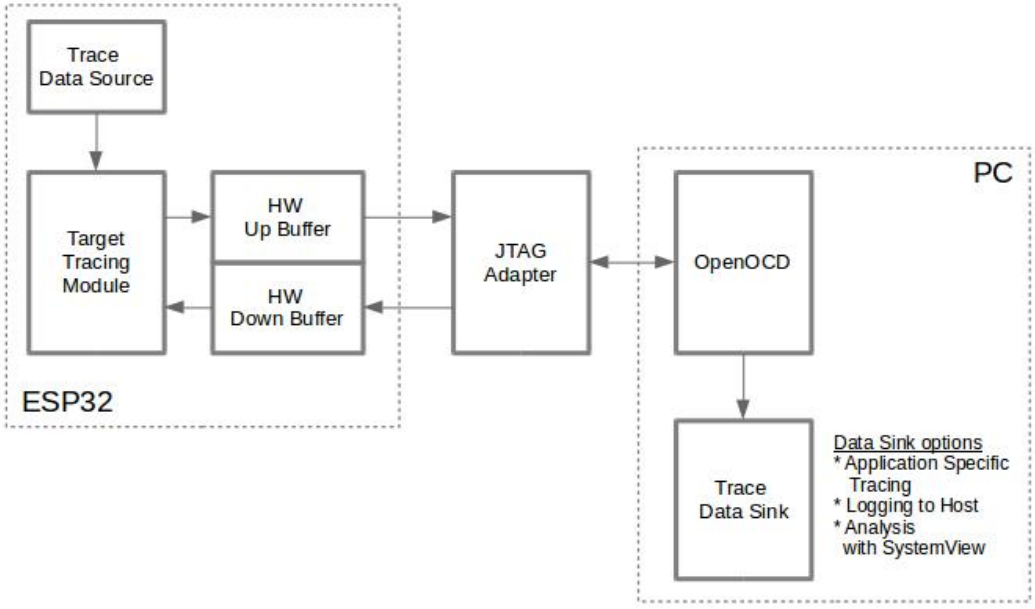

图 24: 使用 JTAG 接口的跟踪组件

运行模式 该库支持两种操作模式：

后验模式： 这是默认的模式，该模式不需要和主机进行交互。在这种模式下，跟踪模块不会检查主机是否已经从 *HW UP BUFFER* 缓冲区读走所有数据，而是直接使用新数据覆盖旧数据。该模式在用户仅对最新的跟踪数据感兴趣时会很有用，例如分析程序在崩溃之前的行为。主机可以稍后根据用户的请求来读取数据，例如通过特殊的 OpenOCD 命令（假如使用了 JTAG 接口）。

流模式： 当主机连接到 ESP32-S2 时，跟踪模块会进入此模式。在这种模式下，跟踪模块在新数据写入 *HW UP BUFFER* 之前会检查其中是否有足够的空间，并在必要的时候等待主机读取数据并释放足够的内存。用户会将最长的等待时间作为超时时间参数传递给相应的 API 函数，如果超时时间是个有限值，那么应用程序有可能会因为超时而将待写的数据丢弃。尤其需要注意，如果在讲究时效的代码中（如中断处理函数，操作系统调度等）指定了无限的超时时间，那么系统会产生故障。为了避免丢失此类关键数据，开发人员可以通过在 menuconfig 中开启*CONFIG_APPTRACE_PENDING_DATA_SIZE_MAX* 选项来启用额外的数据缓冲区。此宏还指定了在上述条件下可以缓冲的数据大小，它有助于缓解由于 USB 总线拥塞等原因导致的向主机传输数据间歇性减缓的状况。但是，当跟踪数据流的平均比特率超过硬件接口的能力时，它也无能为力。

配置选项与依赖项 使用此功能需要在主机端和目标端做相应的配置：

1. **主机端：** 应用程序跟踪是通过 JTAG 来完成的，因此需要在主机上安装并运行 OpenOCD。相关详细信息请参阅*JTAG Debugging*。
2. **目标端：** 在 menuconfig 中开启应用程序跟踪功能。*Component config* > *Application Level Tracing* 菜单允许选择跟踪数据的传输目标（具体用于传输的硬件接口），选择任一非 None 的目标都会自动开启 `CONFIG_APPTRACE_ENABLE` 这个选项。

注解： 为了实现更高的数据速率并降低丢包率，建议优化 JTAG 的时钟频率，使其达到能够稳定运行的最大值。详细信息请参阅*优化 JTAG 的速度*。

以下为前述未提及的另外两个 menuconfig 选项：

1. *Threshold for flushing last trace data to host on panic*（*CONFIG_APPTRACE_POSTMORTEM_FLUSH_THRESH*）。由于在 JTAG 上工作的性质，此选项是必选项。在该模式下，跟踪数据以 16 KB 数据块的形式曝露给主机。在后验模式中，当一个块被填充时，它会曝露给主机，而之前的块会变得不可用。换句话说，跟踪数据以 16 KB 的粒度进行覆盖。在发生 panic 的时候，当前输入块的最新数据将会被曝露给主机，主机可以读取它们以进行后续分析。如果系统发生 panic 的时候仍有少量数据还没来得及曝光给主机，那么之前收集的 16 KB 的数据将丢失，主机只能看到非常少的最新的跟踪部分，它可能不足以用来诊断问题所在。此 menuconfig 选项允许避免此类情况，它可以控制在发生 panic 时刷新数据的阈值，例如用户可以确定它需要不少于 512 字节的最新跟踪数据，所以如果在发生 panic 时待处理的数据少于 512 字节，它们不会被刷新，也不会覆盖之前的 16 KB。该选项仅在后验模式和 JTAG 工作时有意义。
2. *Timeout for flushing last trace data to host on panic*（*CONFIG_APPTRACE_ONPANIC_HOST_FLUSH_TMO*）。该选项仅在流模式下才起作用，它控制跟踪模块在发生 panic 时等待主机读取最新数据的最长时间。

如何使用这个库 该库提供了用于在主机和 ESP32-S2 之间传输任意数据的 API。当在 menuconfig 中启用时，目标应用程序的跟踪模块会在系统启动时自动初始化，因此用户需要做的就是调用相应的 API 来发送、接收或者刷新数据。

特定应用程序的跟踪 通常，用户需要决定在每个方向上待传输数据的类型以及如何解析（处理）这些数据。要想在目标和主机之间传输数据，用户必须要执行以下几个步骤。

1. 在目标端，用户需要实现将跟踪数据写入主机的算法，下面的代码片段展示了如何执行此操作。

```
#include "esp_app_trace.h"
...
char buf[] = "Hello World!";
esp_err_t res = esp_apptrace_write(ESP_APPTRACE_DEST_TRAX, buf, strlen(buf),␣
↪ESP_APPTRACE_TMO_INFINITE);
if (res != ESP_OK) {
    ESP_LOGE(TAG, "Failed to write data to host!");
    return res;
}
```

`esp_apptrace_write()` 函数使用 memcpy 把用户数据复制到内部缓存中。在某些情况下，使用 `esp_apptrace_buffer_get()` 和 `esp_apptrace_buffer_put()` 函数会更加理想，它们允许开发人员自行分配缓冲区并填充。下面的代码片段展示了如何执行此操作。

```
#include "esp_app_trace.h"
...
int number = 10;
char *ptr = (char *)esp_apptrace_buffer_get(ESP_APPTRACE_DEST_TRAX, 32, 100/
↪*tmo in us*/);
if (ptr == NULL) {
    ESP_LOGE(TAG, "Failed to get buffer!");
    return ESP_FAIL;
}
sprintf(ptr, "Here is the number %d", number);
```

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```
esp_err_t res = esp_apptrace_buffer_put(ESP_APPTRACE_DEST_TRAX, ptr, 100/*tmo␣
↪in us*/);
if (res != ESP_OK) {
    /* in case of error host tracing tool (e.g. OpenOCD) will report␣
↪incomplete user buffer */
    ESP_LOGE(TAG, "Failed to put buffer!");
    return res;
}
```

另外，根据实际项目的需要，用户可能希望从主机接收数据。下面的代码片段展示了如何执行此操作。

```
#include "esp_app_trace.h"
...
char buf[32];
char down_buf[32];
size_t sz = sizeof(buf);

/* config down buffer */
esp_apptrace_down_buffer_config(down_buf, sizeof(down_buf));
/* check for incoming data and read them if any */
esp_err_t res = esp_apptrace_read(ESP_APPTRACE_DEST_TRAX, buf, &sz, 0/*do not␣
↪wait*/);
if (res != ESP_OK) {
    ESP_LOGE(TAG, "Failed to read data from host!");
    return res;
}
if (sz > 0) {
    /* we have data, process them */
    ...
}
```

`esp_apptrace_read()` 函数使用 memcpy 来把主机端的数据复制到用户缓存区。在某些情况下，使用 `esp_apptrace_down_buffer_get()` 和 `esp_apptrace_down_buffer_put()` 函数可能更为理想。它们允许开发人员占用一块读缓冲区并就地进行有关处理操作。下面的代码片段展示了如何执行此操作。

```
#include "esp_app_trace.h"
...
char down_buf[32];
uint32_t *number;
size_t sz = 32;

/* config down buffer */
esp_apptrace_down_buffer_config(down_buf, sizeof(down_buf));
char *ptr = (char *)esp_apptrace_down_buffer_get(ESP_APPTRACE_DEST_TRAX, &sz,␣
↪100/*tmo in us*/);
if (ptr == NULL) {
    ESP_LOGE(TAG, "Failed to get buffer!");
    return ESP_FAIL;
}
if (sz > 4) {
    number = (uint32_t *)ptr;
    printf("Here is the number %d", *number);
} else {
    printf("No data");
}
esp_err_t res = esp_apptrace_down_buffer_put(ESP_APPTRACE_DEST_TRAX, ptr, 100/
↪*tmo in us*/);
if (res != ESP_OK) {
    /* in case of error host tracing tool (e.g. OpenOCD) will report␣
↪incomplete user buffer */
```

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```
    ESP_LOGE(TAG, "Failed to put buffer!");
    return res;
}
```

2. 下一步是编译应用程序的镜像并将其下载到目标板上，这一步可以参考文档构建并烧写。
3. 运行 OpenOCD（参见*JTAG 调试*）。
4. 连接到 OpenOCD 的 telnet 服务器，在终端执行如下命令 `telnet <oocd_host> 4444`。如果在运行 OpenOCD 的同一台机器上打开 telnet 会话，您可以使用 `localhost` 替换上面命令中的 `<oocd_host>`。
5. 使用特殊的 OpenOCD 命令开始收集待跟踪的命令，此命令将传输跟踪数据并将其重定向到指定的文件或套接字（当前仅支持文件作为跟踪数据目标）。相关命令的说明请参阅启动调试器 。
6. 最后一步是处理接收到的数据，由于数据格式由用户定义，因此处理阶段超出了本文档的范围。数据处理的范例可以参考位于 `$IDF_PATH/tools/esp_app_trace` 下的 Python 脚本 `app-trace_proc.py`（用于功能测试）和 `logtrace_proc.py`（请参阅记录日志到主机 章节中的详细信息）。

OpenOCD 应用程序跟踪命令 *HW UP BUFFER* 在用户数据块之间共享，并且会替 API 的调用者（在任务或者中断上下文中）填充分配到的内存。在多线程环境中，正在填充缓冲区的任务/中断可能会被另一个高优先级的任务/中断抢占，有可能发生主机读取还未准备好的用户数据的情况。为了处理这样的情况，跟踪模块在所有用户数据块之前添加一个数据头，其中包含有分配的用户缓冲区的大小（2 字节）和实际写入的数据长度（2 字节），也就是说数据头总共长 4 字节。负责读取跟踪数据的 OpenOCD 命令在读取到不完整的用户数据块时会报错，但是无论如何它都会将整个用户数据块（包括还未填充的区域）的内容放到输出文件中。

下面是 OpenOCD 应用程序跟踪命令的使用说明。

注解： 目前，OpenOCD 还不支持将任意用户数据发送到目标的命令。

命令用法：

```
esp32 apptrace [start <options>] | [stop] | [status] | [dump <cores_num>
<outfile>]
```

子命令：

`start` 开始跟踪（连续流模式）。
`stop` 停止跟踪。
`status` 获取跟踪状态。
`dump` 转储所有后验模式的数据。

Start 子命令的语法：

```
start <outfile> [poll_period [trace_size [stop_tmo [wait4halt
[skip_size]]]]
```

`outfile` 用于保存来自两个 CPU 的数据文件的路径，该参数需要具有以下格式：`file://path/to/file`。
`poll_period` 轮询跟踪数据的周期（单位：毫秒），如果大于 0 则以非阻塞模式运行。默认为 1 毫秒。
`trace_size` 最多要收集的数据量（单位：字节），接收到指定数量的数据后将会停止跟踪。默认情况下是 -1（禁用跟踪大小停止触发器）。
`stop_tmo` 空闲超时（单位：秒），如果指定的时间段内都没有数据就会停止跟踪。默认为 -1（禁用跟踪超时停止触发器）。还可以将其设置为比目标跟踪命令之间的最长暂停值更长的值（可选）。
`wait4halt` 如果设置为 0 则立即开始跟踪，否则命令等待目标停止（复位，打断点等），然后自动恢复它并开始跟踪。默认值为 0。
`skip_size` 开始时要跳过的字节数，默认为 0。

注解： 如果 `poll_period` 为 0，则在跟踪停止之前，OpenOCD 的 telnet 命令将不可用。必须通过复位电路板或者在 OpenOCD 的窗口中（不是 telnet 会话窗口）按下 Ctrl+C。另一种选择是设置 `trace_size`

并等待，当收集到指定数据量时，跟踪会自动停止。

命令使用示例：

1. 将 2048 个字节的跟踪数据收集到“trace.log”文件中，该文件将保存在“openocd-esp32”目录中。

   ```
   esp32 apptrace start file://trace.log 1 2048 5 0 0
   ```

   跟踪数据会被检索并以非阻塞的模式保存到文件中，如果收集满 2048 字节的数据或者在 5 秒内都没有新的数据，那么该过程就会停止。

   注解： 在将数据提供给 OpenOCD 之前，会对其进行缓冲。如果看到“Data timeout!”的消息，则目标可能在超时之前没有发送足够的数据给 OpenOCD 来清空缓冲区。增加超时时间或者使用函数 `esp_apptrace_flush()` 以特定间隔刷新数据都可以解决这个问题。

2. 在非阻塞模式下无限地检索跟踪数据。

   ```
   esp32 apptrace start file://trace.log 1 -1 -1 0 0
   ```

   对收集数据的大小没有限制，并且没有设置任何超时时间。可以通过在 OpenOCD 的 telnet 会话窗口中发送 `esp32 apptrace stop` 命令，或者在 OpenOCD 窗口中使用快捷键 Ctrl+C 来停止此过程。
3. 检索跟踪数据并无限期保存。

   ```
   esp32 apptrace start file://trace.log 0 -1 -1 0 0
   ```

   在跟踪停止之前，OpenOCD 的 telnet 会话窗口将不可用。要停止跟踪，请在 OpenOCD 的窗口中使用快捷键 Ctrl+C。
4. 等待目标停止，然后恢复目标的操作并开始检索数据。当收集满 2048 字节的数据后就停止：

   ```
   esp32 apptrace start file://trace.log 0 2048 -1 1 0
   ```

   想要复位后立即开始跟踪，请使用 OpenOCD 的 `reset halt` 命令。

记录日志到主机 记录日志到主机是 IDF 的一个非常实用的功能：通过应用层跟踪库将日志保存到主机端。某种程度上这也算是一种半主机（semihosting）机制，相较于调用 `ESP_LOGx` 将待打印的字符串发送到 UART 的日志记录方式，这个功能的优势在于它减少了本地的工作量，而将大部分工作转移到了主机端。

IDF 的日志库会默认使用类 vprintf 的函数将格式化的字符串输出到专用的 UART。一般来说，它涉及到以下几个步骤：

1. 解析格式字符串以获取每个参数的类型。
2. 根据其类型，将每个参数都转换为字符串。
3. 格式字符串与转换后的参数一起发送到 UART。

虽然可以将类 vprintf 函数优化到一定程度，但是上述步骤在任何情况下都是必须要执行的，并且每个步骤都会消耗一定的时间（尤其是步骤 3）。所以经常会发生以下这种情况：向程序中添加额外的打印信息以诊断问题，却改变了应用程序的行为，使得问题无法复现。在最差的情况下，程序会无法正常工作，最终导致报错甚至挂起。

解决此类问题的可能方法是使用更高的波特率或者其他更快的接口，并将字符串格式化的工作转移到主机端。

通过应用层跟踪库的 `esp_apptrace_vprintf` 函数，可以将日志信息发送到主机，该函数不执行格式字符串和参数的完全解析，而仅仅计算传递的参数的数量，并将它们与格式字符串地址一起发送给主机。主机端会通过一个特殊的 Python 脚本来处理并打印接收到的日志数据。

局限 目前通过 JTAG 实现记录日志还存在以下几点局限：

1. 不支持使用 `ESP_EARLY_LOGx` 宏进行跟踪。
2. 不支持大小超过 4 字节的 printf 参数（例如 `double` 和 `uint64_t`）。
3. 仅支持.rodata 段中的格式字符串和参数。

4. printf 参数最多 256 个。

如何使用 为了使用跟踪模块来记录日志，用户需要执行以下步骤：

1. 在目标端，需要安装特殊的类 vprintf 函数，正如前面提到过的，这个函数是 `esp_apptrace_vprintf`，它会负责将日志数据发送给主机。示例代码参见 system/app_trace_to_host 。
2. 按照特定应用程序的跟踪 章节中第 2-5 步骤中的说明进行操作。
3. 打印接收到的日志记录，请在终端运行以下命令：`$IDF_PATH/tools/esp_app_trace/logtrace_proc.py /path/to/trace/file /path/to/program/elf/file`。

Log Trace Processor 命令选项 命令用法：

```
logtrace_proc.py [-h] [--no-errors] <trace_file> <elf_file>
```

位置参数（必要）：

`trace_file` 日志跟踪文件的路径
`elf_file` 程序 ELF 文件的路径

可选参数：

`-h, --help` 显示此帮助信息并退出
`--no-errors, -n` 不打印错误信息

基于 SEGGER SystemView 的系统行为分析 IDF 中另一个基于应用层跟踪库的实用功能是系统级跟踪，它会生成与 SEGGER SystemView 工具 相兼容的跟踪信息。SEGGER SystemView 是一种实时记录和可视化工具，用来分析应用程序运行时的行为。

注解： 目前，基于 IDF 的应用程序能够以文件的形式生成与 SystemView 格式兼容的跟踪信息，并可以使用 SystemView 工具软件打开。但是还无法使用该工具控制跟踪的过程。

如何使用 若需使用这个功能，需要在 menuconfig 中开启*CONFIG_SYSVIEW_ENABLE* 选项，具体路径为：*Component config* > *Application Level Tracing* > *FreeRTOS SystemView Tracing* 。在同一个菜单栏下还开启了其他几个选项：

1. *ESP32-S2 timer to use as SystemView timestamp source*（*CONFIG_SYSVIEW_TS_SOURCE*）选择 SystemView 事件使用的时间戳来源。在单核模式下，使用 ESP32-S2 内部的循环计数器生成时间戳，其最大的工作频率是 240 MHz（时间戳粒度大约为 4 ns）。在双核模式下，使用工作在 40 MHz 的外部定时器，因此时间戳粒度为 25 ns。
2. 可以单独启用或禁用的 SystemView 事件集合（`CONFIG_SYSVIEW_EVT_XXX`）：
   - Trace Buffer Overflow Event
   - ISR Enter Event
   - ISR Exit Event
   - ISR Exit to Scheduler Event
   - Task Start Execution Event
   - Task Stop Execution Event
   - Task Start Ready State Event
   - Task Stop Ready State Event
   - Task Create Event
   - Task Terminate Event
   - System Idle Event
   - Timer Enter Event
   - Timer Exit Event

IDF 中已经包含了所有用于生成兼容 SystemView 跟踪信息的代码，用户只需配置必要的项目选项（如上所示），然后构建、烧写映像到目标板，接着参照前面的介绍，使用 OpenOCD 收集数据。

OpenOCD SystemView 跟踪命令选项 命令用法：

```
esp32 sysview [start <options>] | [stop] | [status]
```

自命令：

start 开启跟踪（连续流模式）。
stop 停止跟踪。
status 获取跟踪状态。

Start 子命令语法：

```
start <outfile1> [outfile2] [poll_period [trace_size [stop_tmo]]]
```

outfile1 保存 PRO CPU 数据的文件路径，此参数需要具有如下格式：`file://path/to/file`。
outfile2 保存 APP CPU 数据的文件路径，此参数需要具有如下格式：`file://path/to/file`。
poll_period 跟踪数据的轮询周期（单位：毫秒）。如果该值大于 0，则命令以非阻塞的模式运行。默认为 1 毫秒。
trace_size 最多要收集的数据量（单位：字节）。当收到指定数量的数据后，将停止跟踪。默认值是 -1（禁用跟踪大小停止触发器）。
stop_tmo 空闲超时（单位：秒）。如果指定的时间内没有数据，将停止跟踪。默认值是 -1（禁用跟踪超时停止触发器）。

注解： 如果 `poll_period` 为 0，则在跟踪停止之前，OpenOCD 的 telnet 命令行将不可用。你需要通过复位板卡或者在 OpenOCD 的窗口（不是 telnet 会话窗口）输入 Ctrl+C 命令来手动停止它。另一个办法是设置 `trace_size` 然后等到收集满指定数量的数据后自动停止跟踪。

命令使用示例：

1. 将 SystemView 跟踪数据收集到文件 “pro-cpu.SVDat” 和 “pro-cpu.SVDat” 中。这些文件会被保存在 “openocd-esp32” 目录中。

   ```
   esp32 sysview start file://pro-cpu.SVDat file://app-cpu.SVDat
   ```

   跟踪数据被检索并以非阻塞的方式保存，要停止此过程，需要在 OpenOCD 的 telnet 会话窗口输入 `esp32 apptrace stop` 命令，或者也可以在 OpenOCD 窗口中按下 Ctrl+C。
2. 检索跟踪数据并无限保存。

   ```
   esp32 sysview start file://pro-cpu.SVDat file://app-cpu.SVDat 0 -1 -1
   ```

   OpenOCD 的 telnet 命令行在跟踪停止前会无法使用，要停止跟踪，请在 OpenOCD 窗口按下 Ctrl+C。

数据可视化 收集到跟踪数据后，用户可以使用特殊的工具来可视化结果并分析程序的行为。遗憾的是，SystemView 不支持从多个核心进行跟踪。所以当追踪双核模式下的 ESP32 时会生成两个文件：一个用于 PRO CPU，另一个用于 APP CPU。用户可以将每个文件加载到工具中单独分析。

在工具中单独分析每个核的跟踪数据是比较棘手的，幸运的是，Eclipse 中有一款 *Impulse* 的插件可以加载多个跟踪文件，并且可以在同一个视图中检查来自两个内核的事件。此外，与免费版的 SystemView 相比，此插件没有 1,000,000 个事件的限制。

关于如何安装、配置 Impulse 并使用它可视化来自单个核心的跟踪数据，请参阅 官方教程 。

注解： IDF 使用自己的 SystemView FreeRTOS 事件 ID 映射，因此用户需要将 `$SYSVIEW_INSTALL_DIR/Description/SYSVIEW_FreeRTOS.txt` 替换成 `$IDF_PATH/docs/api-guides/SYSVIEW_FreeRTOS.txt`。在使用上述链接配置 SystemView 序列化程序时，也应该使用该 IDF 特定文件的内容。

配置 Impulse 实现双核跟踪 在安装好 Impulse 插件后，先确保它能够在单独的选项卡中成功加载每个核心的跟踪文件，然后用户可以添加特殊的 Multi Adapter 端口并将这两个文件加载到一个视图中。为此，用户需要在 Eclipse 中执行以下操作：

1. 打开 “Signal Ports” 视图，前往 Windows->Show View->Other 菜单，在 Impulse 文件夹中找到 “Signal Ports” 视图，然后双击它。
2. 在 “Signal Ports” 视图中，右键单击 “Ports” 并选择 “Add ···”，然后选择 New Multi Adapter Port。
3. 在打开的对话框中按下 “Add” 按钮，选择 “New Pipe/File”。
4. 在打开的对话框中选择 “SystemView Serializer” 并设置 PRO CPU 跟踪文件的路径，按下确定保存设置。
5. 对 APP CPU 的跟踪文件重复步骤 3 和 4。
6. 双击创建的端口，会打开此端口的视图。
7. 单击 Start/Stop Streaming 按钮，数据将会被加载。
8. 使用 “Zoom Out”，“Zoom In” 和 “Zoom Fit” 按钮来查看数据。
9. 有关设置测量光标和其他的功能，请参阅 Impulse 官方文档 。

注解： 如果您在可视化方面遇到了问题（未显示数据或者缩放操作很奇怪），您可以尝试删除当前的信号层次结构，再双击必要的文件或端口。Eclipse 会请求您创建新的信号层次结构。

- 使用调试器
- 调试示例
- 注意事项和补充内容
- 应用层跟踪库
- ESP-Prog 调试板介绍

4.14 引导加载程序（Bootloader）

引导加载程序（Bootloader）主要执行以下任务：

1. 内部模块的最小化初始配置；
2. 根据分区表和 ota_data（如果存在）选择需要引导的应用程序（app）分区；
3. 将此应用程序映像加载到 RAM（IRAM 和 DRAM）中，最后把控制权转交给应用程序。

引导加载程序位于 Flash 的 *0x1000* 偏移地址处。

4.14.1 恢复出厂设置

用户可以编写一个基本的工作固件，然后将其加载到工厂分区（factory）中。

接下来，通过 OTA（空中升级）更新固件，更新后的固件会被保存到某个 OTA app 分区中，OTA 数据分区也会做相应更新以指示从该分区引导应用程序。

如果你希望回滚到出厂固件并清除设置，则需要设置*CONFIG_BOOTLOADER_FACTORY_RESET*。

出厂重置机制允许将设备重置为出厂模式：

- 清除一个或多个数据分区。
- 从工厂分区启动。

CONFIG_BOOTLOADER_DATA_FACTORY_RESET - 允许用户选择在恢复出厂设置时需要删除的数据分区。可以通过逗号来分隔多个分区的名字，并适当增加空格以便阅读（例如 “nvs, phy_init, nvs_custom, ···”）。请确保此处指定的名称和分区表中的名称相同，且不含有 “app” 类型的分区。

CONFIG_BOOTLOADER_OTA_DATA_ERASE - 恢复出厂模式后，设备会从工厂分区启动，OTA 数据分区会被清除。

CONFIG_BOOTLOADER_NUM_PIN_FACTORY_RESET- 设置用于触发出厂重置的 GPIO 编号，必须在芯片复位时将此 GPIO 拉低才能触发出厂重置事件。

CONFIG_BOOTLOADER_HOLD_TIME_GPIO- 设置进入重置或测试模式所需要的保持时间（默认为 5 秒）。设备复位后，GPIO 必须在这段时间内持续保持低电平，然后才会执行出厂重置或引导测试分区。

示例分区表如下：

```
# Name,   Type, SubType, Offset,   Size, Flags
# 注意，如果你增大了引导加载程序的大小，请确保更新偏移量，避免和其它分区发生重叠
nvs,      data, nvs,     0x9000,   0x4000
otadata,  data, ota,     0xd000,   0x2000
phy_init, data, phy,     0xf000,   0x1000
factory,  0,    0,       0x10000,  1M
test,     0,    test,    ,         512K
ota_0,    0,    ota_0,   ,         512K
ota_1,    0,    ota_1,   ,         512K
```

4.14.2 从测试固件启动

用户可以编写在生产环境中测试用的特殊固件，然后在需要的时候运行。此时需要在分区表中专门申请一块分区用于保存该测试固件（详情请参阅分区表）。如果想要触发测试固件，还需要设置*CONFIG_BOOTLOADER_APP_TEST*。

CONFIG_BOOTLOADER_NUM_PIN_APP_TEST - 设置引导测试分区的 GPIO 管脚编号，该 GPIO 会被配置为输入模式，并且会使能内部上拉电阻。若想触发测试固件，该 GPIO 必须在芯片复位时拉低。设备重启时如果该 GPIO 没有被激活（即处于高电平状态），那么会加载常规配置的应用程序（可能位于工厂分区或者 OTA 分区）。

CONFIG_BOOTLOADER_HOLD_TIME_GPIO - 设置进入重置或测试模式所需要的保持时间（默认为 5 秒）。设备复位后，GPIO 必须在这段时间内持续保持低电平，然后才会执行出厂重置或引导测试分区。

4.14.3 自定义引导程序

用户可以重写当前的引导加载程序，具体做法是，复制 */esp-idf/components/bootloader* 文件夹到项目目录中，然后编辑 */your_project/components/bootloader/subproject/ain/bootloader_main.c* 文件。

在引导加载程序的代码中，用户不可以使用驱动和其他组件提供的函数，如果确实需要，请将该功能的实现部分放在 bootloader 目录中（注意，这会增加引导程序的大小）。监视生成的引导程序的大小是有必要的，因为它可能会与内存中的分区表发生重叠而损坏固件。目前，引导程序被限制在了分区表之前的区域（分区表位于 *0x8000* 地址处）。

4.15 分区表

4.15.1 概述

每片 ESP32-S2 的 flash 可以包含多个应用程序，以及多种不同类型的数据（例如校准数据、文件系统数据、参数存储器数据等）。因此，我们需要引入分区表的概念。

分区表中的每个条目都包括以下几个部分：Name（标签）、Type（app、data 等）、SubType 以及在 flash 中的偏移量（分区的加载地址）。

在使用分区表时，最简单的方法就是打开项目配置菜单（`idf.py menuconfig`），并在*CONFIG_PARTITION_TABLE_TYPE* 下选择一个预定义的分区表：

- “Single factory app, no OTA”
- “Factory app, two OTA definitions”

在以上两种选项中，出厂应用程序均将被烧录至 flash 的 0x10000 偏移地址处。这时，运行 *idf.py partition_table*，即可以打印当前使用分区表的信息摘要。

4.15.2 内置分区表

以下是“Single factory app, no OTA”选项的分区表信息摘要：

```
# ESP-IDF Partition Table # Name, Type, SubType, Offset, Size, Flags nvs, data, nvs, 0x9000, 0x6000,
phy_init, data, phy, 0xf000, 0x1000, factory, app, factory, 0x10000, 1M,
```

- flash 的 0x10000 (64KB) 偏移地址处存放一个标记为“factory”的二进制应用程序，且启动加载器将默认加载这个应用程序。
- 分区表中还定义了两个数据区域，分别用于存储 NVS 库专用分区和 PHY 初始化数据。

以下是“Factory app, two OTA definitions”选项的分区表信息摘要：

```
# ESP-IDF Partition Table # Name, Type, SubType, Offset, Size, Flags nvs, data, nvs, 0x9000, 0x4000,
otadata, data, ota, 0xd000, 0x2000, phy_init, data, phy, 0xf000, 0x1000, factory, app, factory, 0x10000,
1M, ota_0, app, ota_0, 0x110000, 1M, ota_1, app, ota_1, 0x210000, 1M,
```

- 分区表中定义了三个应用程序分区，这三个分区的类型都被设置为“app”，但具体 app 类型不同。其中，位于 0x10000 偏移地址处的为出厂应用程序（factory），其余两个为 OTA 应用程序（ota_0，ota_1）。
- 新增了一个名为“otadata”的数据分区，用于保存 OTA 升级时需要的数据。启动加载器会查询该分区的数据，以判断该从哪个 OTA 应用程序分区加载程序。如果“otadata”分区为空，则会执行出厂程序。

4.15.3 创建自定义分区表

如果在 `menuconfig` 中选择了“Custom partition table CSV”，则还需要输入该分区表的 CSV 文件在项目中的路径。CSV 文件可以根据需要，描述任意数量的分区信息。

CSV 文件的格式与上面摘要中打印的格式相同，但是在 CSV 文件中并非所有字段都是必需的。例如下面是一个自定义的 OTA 分区表的 CSV 文件：

```
# Name, Type, SubType, Offset, Size, Flags nvs, data, nvs, 0x9000, 0x4000 otadata, data, ota, 0xd000,
0x2000 phy_init, data, phy, 0xf000, 0x1000 factory, app, factory, 0x10000, 1M ota_0, app, ota_0, , 1M
ota_1, app, ota_1, , 1M nvs_key, data, nvs_keys, , 0x1000
```

- 字段之间的空格会被忽略，任何以 `#` 开头的行（注释）也会被忽略。
- CSV 文件中的每个非注释行均为一个分区定义。
- 每个分区的 `Offset` 字段可以为空，`gen_esp32part.py` 工具会从分区表位置的后面开始自动计算并填充该分区的偏移地址，同时确保每个分区的偏移地址正确对齐。

Name 字段

Name 字段可以是任何有意义的名称，但不能超过 16 个字符（之后的内容将被截断）。该字段对 ESP32-S2 并不是特别重要。

Type 字段

Type 字段可以指定为 app (0) 或者 data (1)，也可以直接使用数字 0-254（或者十六进制 0x00-0xFE）。注意，0x00-0x3F 不得使用（预留给 esp-idf 的核心功能）。

如果您的应用程序需要保存数据，请在 0x40-0xFE 内添加一个自定义分区类型。

注意，启动加载器将忽略 app (0) 和 data (1) 以外的其他分区类型。

SubType 字段

SubType 字段长度为 8 bit，内容与具体 Type 有关。目前，esp-idf 仅仅规定了“app”和“data”两种子类型。

- 当 Type 定义为 `app` 时，SubType 字段可以指定为 factory (0)，ota_0 (0x10) ···ota_15 (0x1F) 或者 test (0x20)。

  - factory (0) 是默认的 app 分区。启动加载器将默认加载该应用程序。但如果存在类型为 data/ota 分区，则启动加载器将加载 data/ota 分区中的数据，进而判断启动哪个 OTA 镜像文件。- OTA 升级永远都不会更新 factory 分区中的内容。- 如果您希望在 OTA 项目中预留更多 flash，可以删除 factory 分区，转而使用 ota_0 分区。
  - ota_0 (0x10) ··· ota_15 (0x1F) 为 OTA 应用程序分区，启动加载器将根据 OTA 数据分区中的数据来决定加载哪个 OTA 应用程序分区中的程序。在使用 OTA 功能时，应用程序应至少拥有 2 个 OTA 应用程序分区（ota_0 和 ota_1）。更多详细信息，请参考*OTA 文档*。
  - test (0x2) 为预留 app 子类型，用于工厂测试流程。如果没有其他有效 app 分区，test 将作为备选启动分区使用。也可以在每次启动时配置启动加载器读取 GPIO，如果 GPIO 被拉低则启动该分区。详细信息请查阅*从测试固件启动*。
- 当 Type 定义为 `data` 时，SubType 字段可以指定为 ota (0)，phy (1)，nvs (2) 或者 nvs_keys (4)。
  - ota (0) 即*OTA 数据分区*，用于存储当前所选的 OTA 应用程序的信息。这个分区的大小需要设定为 0x2000。更多详细信息，请参考*OTA 文档*。
  - phy (1) 分区用于存放 PHY 初始化数据，从而保证可以为每个设备单独配置 PHY，而非必须采用固件中的统一 PHY 初始化数据。
    * 默认配置下，phy 分区并不启用，而是直接将 phy 初始化数据编译至应用程序中，从而节省分区表空间（直接将此分区删掉）。
    * 如果需要从此分区加载 phy 初始化数据，请打开项目配置菜单（`idf.py menuconfig`），并且使能*CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION* 选项。此时，您还需要手动将 phy 初始化数据烧至设备 flash（esp-idf 编译系统并不会自动完成该操作）。
  - nvs (2) 是专门给*非易失性存储 (NVS) API* 使用的分区。
    * 用于存储每台设备的 PHY 校准数据（注意，并不是 PHY 初始化数据）。
    * 用于存储 Wi-Fi 数据（如果使用了*`esp_wifi_set_storage(WIFI_STORAGE_FLASH)`* 初始化函数）。
    * NVS API 还可以用于其他应用程序数据。
    * 强烈建议您应为 NVS 分区分配至少 0x3000 字节空间。
    * 如果使用 NVS API 存储大量数据，请增加 NVS 分区的大小（默认是 0x6000 字节）。
  - nvs_keys (4) 是 NVS 秘钥分区。详细信息，请参考*非易失性存储 (NVS) API* 文档。
    * 用于存储加密密钥（如果启用了 *NVS 加密功能*）。
    * 此分区应至少设定为 4096 字节。

其它数据子类型已预留给 esp-idf 未来使用。

Offset 和 Size 字段

分区若偏移地址为空，则会紧跟着前一个分区之后开始；若为首个分区，则将紧跟着分区表开始。

app 分区的偏移地址必须要与 0x10000 (64K) 对齐，如果将偏移字段留空，`gen_esp32part.py` 工具会自动计算得到一个满足对齐要求的偏移地址。如果 app 分区的偏移地址没有与 0x10000 (64K) 对齐，则该工具会报错。

app 分区的大小和偏移地址可以采用十进制数、以 0x 为前缀的十六进制数，且支持 K 或 M 的倍数单位（分别代表 1024 和 1024*1024 字节）。

如果您希望允许分区表中的分区采用任意起始偏移量 (*CONFIG_PARTITION_TABLE_OFFSET*)，请将分区表（CSV 文件）中所有分区的偏移字段都留空。注意，此时，如果您更改了分区表中任意分区的偏移地址，则其他分区的偏移地址也会跟着改变。这种情况下，如果您之前还曾设定某个分区采用固定偏移地址，则可能造成分区表冲突，从而导致报错。

Flags 字段

当前仅支持 `encrypted` 标记。如果 Flags 字段设置为 `encrypted`，且已启用*Flash Encryption* 功能，则该分区将会被加密。

注解： `app` 分区始终会被加密，不管 Flags 字段是否设置。

4.15.4 生成二进制分区表

烧写到 ESP32 中的分区表采用二进制格式，而不是 CSV 文件本身。此时，partition_table/gen_esp32part.py 工具可以实现 CSV 和二进制文件之间的转换。

如果您在项目配置菜单（`idf.py menuconfig`）中设置了分区表 CSV 文件的名称，然后构建项目或执行 `idf.py partition_table`。这时，转换将在编译过程中自动完成。

手动将 CSV 文件转换为二进制文件:

```
python gen_esp32part.py input_partitions.csv binary_partitions.bin
```

手动将二进制文件转换为 CSV 文件:

```
python gen_esp32part.py binary_partitions.bin input_partitions.csv
```

在标准输出（stdout）上，打印二进制分区表的内容（在运行 `idf.py partition_table` 时，我们正是这样打印上文展示的信息摘要的）：

```
python gen_esp32part.py binary_partitions.bin
```

MD5 校验和

二进制格式的分区表中含有一个 MD5 校验和。这个 MD5 校验和是根据分区表内容计算的，可在设备启动阶段，用于验证分区表的完整性。

注意，一些版本较老的启动加载器无法支持 MD5 校验，如果发现 MD5 校验和则将报错 `invalid magic number 0xebeb`。此时，用户可通过 `gen_esp32part.py` 的 `--disable-md5sum` 选项或者 `menuconfig` 的*CONFIG_PARTITION_TABLE_MD5* 选项关闭 MD5 校验。

4.15.5 烧写分区表

- `idf.py partition_table-flash`： 使用 esptool.py 工具烧写分区表。
- `idf.py flash`： 会烧写所有内容，包括分区表。

在执行 `idf.py partition_table` 命令时，手动烧写分区表的命令也将打印在终端上。

注解: 分区表的更新并不会擦除根据之前分区表存储的数据。此时，您可以使用 `idf.py erase_flash` 命令或者 `esptool.py erase_flash` 命令来擦除 flash 中的所有内容。

4.15.6 分区工具 (parttool.py)

partition_table 组件中有分区工具 parttool.py，可以在目标设备上完成分区相关操作。该工具有如下用途:

- 读取分区，将内容存储到文件中 (read_partition)
- 将文件中的内容写至分区 (write_partition)
- 擦除分区 (erase_partition)
- 检索特定分区的偏移和大小等信息 (get_partition_info)

用户若想通过编程方式完成相关操作，可从另一个 Python 脚本导入并使用分区工具，或者从 Shell 脚本调用分区工具。前者可使用工具的 Python API，后者可使用命令行界面。

Python API

首先请确保已导入 *parttool* 模块。

```
import sys
import os

idf_path = os.environ["IDF_PATH"]  # 从环境中获取 IDF_PATH 的值
parttool_dir = os.path.join(idf_path, "components", "partition_table")  # parttool.
↪py 位于 $IDF_PATH/components/partition_table 下

sys.path.append(parttool_dir)  # 使能 Python 寻找 parttool 模块
from parttool import *  # 导入 parttool 模块内的所有名称
```

要使用分区工具的 Python API，第一步是创建 *ParttoolTarget*：

```
# 创建 partool.py 的目标设备，并将目标设备连接到串行端口 /dev/ttyUSB1
target = ParttoolTarget("/dev/ttyUSB1")
```

现在，可使用创建的 *ParttoolTarget* 在目标设备上完成操作：

```
# 擦除名为 'storage' 的分区
target.erase_partition(PartitionName("storage"))

# 读取类型为 'data'、子类型为 'spiffs' 的分区，保存至文件 'spiffs.bin'
target.read_partition(PartitionType("data", "spiffs"), "spiffs.bin")

# 将 'factory.bin' 文件的内容写至 'factory' 分区
target.write_partition(PartitionName("factory"), "factory.bin")

# 打印默认启动分区的大小
storage = target.get_partition_info(PARTITION_BOOT_DEFAULT)
print(storage.size)
```

使用 *PartitionName*、*PartitionType* 或 PARTITION_BOOT_DEFAULT 指定要操作的分区。顾名思义，这三个参数可以指向拥有特定名称的分区、特定类型和子类型的分区或默认启动分区。

更多关于 Python API 的信息，请查看分区工具的代码注释。

命令行界面

parttool.py 的命令行界面具有如下结构：

```
parttool.py [command-args] [subcommand] [subcommand-args]

- command-args - 执行主命令 (parttool.py) 所需的实际参数，多与目标设备有关
- subcommand - 要执行的操作
- subcommand-args - 所选操作的实际参数
```

```
# 擦除名为 'storage' 的分区
parttool.py --port "/dev/ttyUSB1" erase_partition --partition-name=storage

# 读取类型为 'data'、子类型为 'spiffs' 的分区，保存到 'spiffs.bin' 文件
parttool.py --port "/dev/ttyUSB1" read_partition --partition-type=data --partition-
↪subtype=spiffs "spiffs.bin"

# 将 'factory.bin' 文件中的内容写入到 'factory' 分区
parttool.py --port "/dev/ttyUSB1" write_partition --partition-name=factory
↪"factory.bin"

# 打印默认启动分区的大小
parttool.py --port "/dev/ttyUSB1" get_partition_info --partition-boot-default --
↪info size
```

更多信息可用 *–help* 指令查看：

```
# 显示可用的子命令和主命令描述
parttool.py --help

# 显示子命令的描述
parttool.py [subcommand] --help
```

4.16 Secure Boot V2

重要: The references in this document are related to Secure Boot V2, the preferred scheme from ESP32-ECO3 onwards and in ESP32-S2.

Secure Boot V2 uses RSA based app and bootloader verification. This document can also be referred for signing apps with the RSA scheme without signing the bootloader.

4.16.1 Background

Secure Boot protects a device from running unsigned code (verification at time of load). A new RSA based secure boot verification scheme (Secure Boot V2) has been introduced for ESP32-S2 and ESP32 ECO3 onwards.

- The software bootloader's RSA-PSS signature is verified by the Mask ROM and it is executed post successful verification.
- The verified software bootloader verifies the RSA-PSS signature of the application image before it is executed.

4.16.2 Advantages

- The RSA public key is stored on the device. The corresponding RSA private key is kept secret on a server and is never accessed by the device.
  - Up to three public keys can be generated and stored in the chip during manufacturing.
  - ESP32-S2 provides the facility to permanently revoke individual public keys. This can be configured conservatively or aggressively.
  - Conservatively - The old key is revoked after the bootloader and application have successfully migrated to a new key. Aggressively - The key is revoked as soon as verification with this key fails.
- Same image format & signature verification is applied for applications & software bootloader.
- No secrets are stored on the device. Therefore immune to passive side-channel attacks (timing or power analysis, etc.)

4.16.3 Secure Boot V2 Process

This is an overview of the Secure Boot V2 Process, Step by step instructions are supplied under *How To Enable Secure Boot V2*.

1. Secure Boot V2 verifies the signature blocks appended to the bootloader and application binaries. The signature block contains the image binary signed by a RSA-3072 private key and its corresponding public key. More details on the *Signature Block Format*.
2. On startup, ROM code checks the Secure Boot V2 bit in eFuse.
3. If secure boot is enabled, ROM checks the SHA-256 of the public key in the signature block in the eFuse.
4. The ROM code validates the public key embedded in the software bootloader's signature block by matching the SHA-256 of its public key to the SHA-256 in eFuse as per the earlier step. Boot process will be aborted if a valid hash of the public key isn't found in the eFuse.
5. The ROM code verifies the signature of the bootloader with the pre-validated public key with the RSA-PSS Scheme. In depth information on *Verifying the signature Block*.

6. Software bootloader, reads the app partition and performs similar verification on the application. The application is verified on every boot up and OTA update. If selected OTA app partition fails verification, bootloader will fall back and look for another correctly signed partition.

4.16.4 Signature Block Format

The bootloader and application images are padded to the next 4096 byte boundary, thus the signature has a flash sector of its own. The signature is calculated over all bytes in the image including the padding bytes.

Each signature block contains the following:

- **Offset 0 (1 byte):** Magic byte (0xe7)
- **Offset 1 (1 byte):** Version number byte (currently 0x02), 0x01 is for Secure Boot V1.
- **Offset 2 (2 bytes):** Padding bytes, Reserved. Should be zero.
- **Offset 4 (32 bytes):** SHA-256 hash of only the image content, not including the signature block.
- **Offset 36 (384 bytes):** RSA Public Modulus used for signature verification. (value ‘n’ in RFC8017).
- **Offset 420 (4 bytes):** RSA Public Exponent used for signature verification (value ‘e’ in RFC8017).
- **Offset 424 (384 bytes):** Precalculated R, derived from ‘n’.
- **Offset 808 (4 bytes):** Precalculated M’, derived from ‘n’
- **Offset 812 (384 bytes):** RSA-PSS Signature result (section 8.1.1 of RFC8017) of image content, computed using following PSS parameters: SHA256 hash, MFG1 function, 0 length salt, default trailer field (0xBC).
- **Offset 1196:** CRC32 of the preceding 1095 bytes.
- **Offset 1200 (16 bytes):** Zero padding to length 1216 bytes.

注解： R and M’ are used for hardware-assisted Montgomery Multiplication.

The remainder of the signature sector is erased flash (0xFF) which allows writing other signature blocks after previous signature block.

4.16.5 Verifying the signature Block

A signature block is “valid” if the first byte is 0xe7 and a valid CRC32 is stored at offset 1196.

Upto 3 signature blocks can be appended to the bootloader or application image in ESP32-S2.

An image is “verified” if the public key stored in any signature block is valid for this device, and if the stored signature is valid for the image data read from flash.

1. The magic byte, signature block CRC is validated.
2. Public key digests are generated per signature block and compared with the digests from eFuse. If none of the digests match, the verification process is aborted.
3. The application image digest is generated and matched with the image digest in the signature blocks. The verification process is aborted is the digests don’t match.
4. The public key is used to verify the signature of the bootloader image, using RSA-PSS (section 8.1.2 of RFC8017) with the image digest calculated in step (3) for comparison.

- The application signing scheme is set to RSA for Secure Boot V2 and to ECDSA for Secure Boot V1.

重要： It is recommended to use Secure Boot V2 on the chip versions supporting them.

4.16.6 Bootloader Size

Enabling Secure boot and/or flash encryption will increase the size of bootloader, which might require updating partition table offset. See secure-boot-bootloader-size.

4.16.7 eFuse usage

- SECURE_BOOT_EN - Enables secure boot protection on boot.
- KEY_PURPOSE_X - Set the purpose of the key block on ESP32-S2 by programming SECURE_BOOT_DIGESTX (X = 0, 1, 2) into KEY_PURPOSE_X (X = 0, 1, 2, 3, 4, 5). Example: If KEY_PURPOSE_2 is set to SECURE_BOOT_DIGEST1, then BLOCK_KEY2 will have the Secure Boot V2 public key digest.
- BLOCK_KEYX - The block contains the data corresponding to its purpose programmed in KEY_PURPOSE_X. Stores the SHA-256 digest of the public key. SHA-256 hash of public key modulus, exponent, precalculated R & M' values (represented as 776 bytes –offsets 36 to 812 - as per the *Signature Block Format*) is written to an eFuse key block.
- KEY_REVOKEX - The revocation bits corresponding to each of the 3 key block. Ex. Setting KEY_REVOKE2 revokes the key block whose key purpose is SECURE_BOOT_DIGEST2.
- SECURE_BOOT_AGGRESSIVE_REVOKE - Enables aggressive revocation of keys. The key is revoked as soon as verification with this key fails.

4.16.8 How To Enable Secure Boot V2

1. Open the *Project Configuration Menu*, in "Security features" set "Enable hardware Secure Boot in bootloader" to enable Secure Boot.
2. The "Secure Boot V2" option will be selected and the "App Signing Scheme" would be set to RSA by default.
3. Specify the path to secure boot signing key, relative to the project directory.
4. Set other menuconfig options (as desired). Pay particular attention to the "Bootloader Config" options, as you can only flash the bootloader once. Then exit menuconfig and save your configuration.
5. The first time you run `make` or `idf.py build`, if the signing key is not found then an error message will be printed with a command to generate a signing key via `espsecure.py generate_signing_key`.

重要: A signing key generated this way will use the best random number source available to the OS and its Python installation (*/dev/urandom* on OSX/Linux and *CryptGenRandom()* on Windows). If this random number source is weak, then the private key will be weak.

重要: For production environments, we recommend generating the keypair using openssl or another industry standard encryption program. See *Generating Secure Boot Signing Key* for more details.

---

6. Run `idf.py bootloader` to build a secure boot enabled bootloader. The build output will include a prompt for a flashing command, using `esptool.py write_flash`.
7. When you're ready to flash the bootloader, run the specified command (you have to enter it yourself, this step is not performed by the build system) and then wait for flashing to complete.
8. Run `idf.py flash` to build and flash the partition table and the just-built app image. The app image will be signed using the signing key you generated in step 4.

注解: `idf.py flash` doesn't flash the bootloader if secure boot is enabled.

---

9. Reset the ESP32-S2 and it will boot the software bootloader you flashed. The software bootloader will enable secure boot on the chip, and then it verifies the app image signature and boots the app. You should watch the serial console output from the ESP32-S2 to verify that secure boot is enabled and no errors have occurred due to the build configuration.

注解: Secure boot won't be enabled until after a valid partition table and app image have been flashed. This is to prevent accidents before the system is fully configured.

注解： If the ESP32-S2 is reset or powered down during the first boot, it will start the process again on the next boot.

10. On subsequent boots, the secure boot hardware will verify the software bootloader has not changed and the software bootloader will verify the signed app image (using the validated public key portion of its appended signature block).

4.16.9 Restrictions after Secure Boot is enabled

- Any updated bootloader or app will need to be signed with a key matching the digest already stored in efuse.
- After Secure Boot is enabled, no further efuses can be read protected. (If *Flash 加密* is enabled then the bootloader will ensure that any flash encryption key generated on first boot will already be read protected.) If *CONFIG_SECURE_BOOT_INSECURE* is enabled then this behaviour can be disabled, but this is not recommended.

4.16.10 Generating Secure Boot Signing Key

The build system will prompt you with a command to generate a new signing key via `espsecure.py generate_signing_key`. The –version 2 parameter will generate the RSA 3072 private key for Secure Boot V2.

The strength of the signing key is proportional to (a) the random number source of the system, and (b) the correctness of the algorithm used. For production devices, we recommend generating signing keys from a system with a quality entropy source, and using the best available RSA key generation utilities.

For example, to generate a signing key using the openssl command line:

```
` openssl genrsa -out my_secure_boot_signing_key.pem 3072 `
```

Remember that the strength of the secure boot system depends on keeping the signing key private.

4.16.11 Remote Signing of Images

For production builds, it can be good practice to use a remote signing server rather than have the signing key on the build machine (which is the default esp-idf secure boot configuration). The espsecure.py command line program can be used to sign app images & partition table data for secure boot, on a remote system.

To use remote signing, disable the option "Sign binaries during build". The private signing key does not need to be present on the build system.

After the app image and partition table are built, the build system will print signing steps using espsecure.py:

```
espsecure.py sign_data --version 2 --keyfile PRIVATE_SIGNING_KEY BINARY_FILE
```

The above command appends the image signature to the existing binary. You can use the *–output* argument to write the signed binary to a separate file:

```
espsecure.py sign_data --version 2 --keyfile PRIVATE_SIGNING_KEY --output SIGNED_
↪BINARY_FILE BINARY_FILE
```

4.16.12 Secure Boot Best Practices

- Generate the signing key on a system with a quality source of entropy.
- Keep the signing key private at all times. A leak of this key will compromise the secure boot system.
- Do not allow any third party to observe any aspects of the key generation or signing process using espsecure.py. Both processes are vulnerable to timing or other side-channel attacks.
- Enable all secure boot options in the Secure Boot Configuration. These include flash encryption, disabling of JTAG, disabling BASIC ROM interpeter, and disabling the UART bootloader encrypted flash access.

- Use secure boot in combination with *flash encryption* to prevent local readout of the flash contents.

4.16.13 Key Management

- Between 1 and 3 RSA-3072 public keypairs (Keys #0, #1, #2) should be computed independently and stored separately.
- The KEY_DIGEST efuses should be write protected after being programmed.
- The unused KEY_DIGEST slots must have their corresponding KEY_REVOKE efuse burned to permanently disable them. This must happen before the device leaves the factory.
- The eFuses can either be written by the software bootloader during during first boot after enabling "Secure Boot V2" from menuconfig or can be done using *espefuse.py* which communicates with the serial bootloader program in ROM.
- The KEY_DIGESTs should be numbered sequentially beginning at key digest #0. (ie if key digest #1 is used, key digest #0 should be used. If key digest #2 is used, key digest #0 & #1 must be used.)
- The software bootloader (non OTA upgradeable) is signed using at least one, possibly all three, private keys and flashed in the factory.
- Apps should only be signed with a single private key (the others being stored securely elsewhere), however they may be signed with multiple private keys if some are being revoked (see Key Revocation, below).

4.16.14 Multiple Keys

- The bootloader should be signed with all the private key(s) that are needed for the life of the device, before it is flashed.
- The build system can sign with at most one private key, user has to run manual commands to append more signatures if necessary.
- **You can use the append functionality of `espsecure.py`, this command would also printed at the end of the Secure B**
  espsecure.py sign_data -k secure_boot_signing_key2.pem -v 2 –append_signatures -o signed_bootloader.bin build/bootloader/bootloader.bin
- While signing with multiple private keys, it is recommended that the private keys be signed independently, if possible on different servers and stored separately.
- **You can check the signatures attached to a binary using -** espsecure.py signature_info_v2 datafile.bin

4.16.15 Key Revocation

- Keys are processed in a linear order. (key #0, key #1, key #2).
- Applications should be signed with only one key at a time, to minimise the exposure of unused private keys.
- The bootloader can be signed with multiple keys from the factory.

Assuming a trusted private key (N-1) has been compromised, to update to new keypair (N).

1. Server sends an OTA update with an application signed with the new private key (#N).
2. The new OTA update is written to an unused OTA app partition.
3. The new application's signature block is validated. The public keys are checked against the digests programmed in the eFuse & the application is verified using the verified public key.
4. The active partition is set to the new OTA application's partition.
5. Device resets, loads the bootloader (verified with key #N-1) which then boots new app (verified with key #N).
6. The new app verifies bootloader with key #N (as a final check) and then runs code to revoke key #N-1 (sets KEY_REVOKE efuse bit).
7. The API *esp_ota_revoke_secure_boot_public_key()* can be used to revoke the key #N-1.

- A similiar approach can also be used to physically reflash with a new key. For physical reflashing, the bootloader content can also be changed at the same time.

4.16.16 Technical Details

The following sections contain low-level reference descriptions of various secure boot elements:

Manual Commands

Secure boot is integrated into the esp-idf build system, so `make` or `idf.py build` will sign an app image and `idf.py bootloader` will produce a signed bootloader if secure signed binaries on build is enabled.

However, it is possible to use the `espsecure.py` tool to make standalone signatures and digests.

To sign a binary image:

```
espsecure.py sign_data --version 2 --keyfile ./my_signing_key.pem --output ./image_
↪signed.bin image-unsigned.bin
```

Keyfile is the PEM file containing an RSA-3072 private signing key.

4.16.17 Secure Boot & Flash Encryption

If secure boot is used without *Flash Encryption*, it is possible to launch "time-of-check to time-of-use" attack, where flash contents are swapped after the image is verified and running. Therefore, it is recommended to use both the features together.

4.16.18 Advanced Features

JTAG Debugging

By default, when Secure Boot is enabled then JTAG debugging is disabled via eFuse. The bootloader does this on first boot, at the same time it enables Secure Boot.

See *JTAG with Flash Encryption or Secure Boot* for more information about using JTAG Debugging with either Secure Boot or signed app verification enabled.

4.17 ULP 协处理器编程

4.17.1 ESP32-S2 ULP coprocessor instruction set

This document provides details about the instructions used by ESP32-S2 ULP coprocessor assembler.

ULP coprocessor has 4 16-bit general purpose registers, labeled R0, R1, R2, R3. It also has an 8-bit counter register (stage_cnt) which can be used to implement loops. Stage count regiter is accessed using special instructions.

ULP coprocessor can access 8k bytes of RTC_SLOW_MEM memory region. Memory is addressed in 32-bit word units. It can also access peripheral registers in RTC_CNTL, RTC_IO, and SENS peripherals.

All instructions are 32-bit. Jump instructions, ALU instructions, peripheral register and memory access instructions are executed in 1 cycle. Instructions which work with peripherals (TSENS, ADC, I2C) take variable number of cycles, depending on peripheral operation.

The instruction syntax is case insensitive. Upper and lower case letters can be used and intermixed arbitrarily. This is true both for register names and instruction names.

Note about addressing

ESP32-S2 ULP coprocessor's JUMP, ST, LD instructions which take register as an argument (jump address, store/load base address) expect the argument to be expressed in 32-bit words.

Consider the following example program:

```
entry:
        NOP
        NOP
        NOP
        NOP
loop:
        MOVE R1, loop
        JUMP R1
```

When this program is assembled and linked, address of label `loop` will be equal to 16 (expressed in bytes). However *JUMP* instruction expects the address stored in register to be expressed in 32-bit words. To account for this common use case, assembler will convert the address of label *loop* from bytes to words, when generating `MOVE` instruction, so the code generated code will be equivalent to:

```
0000    NOP
0004    NOP
0008    NOP
000c    NOP
0010    MOVE R1, 4
0014    JUMP R1
```

The other case is when the argument of `MOVE` instruction is not a label but a constant. In this case assembler will use the value as is, without any conversion:

```
.set        val, 0x10
MOVE        R1, val
```

In this case, value loaded into R1 will be `0x10`.

Similar considerations apply to `LD` and `ST` instructions. Consider the following code:

```
        .global array
array:  .long 0
        .long 0
        .long 0
        .long 0

        MOVE R1, array
        MOVE R2, 0x1234
        ST R2, R1, 0        // write value of R2 into the first array element,
                            // i.e. array[0]

        ST R2, R1, 4        // write value of R2 into the second array element
                            // (4 byte offset), i.e. array[1]

        ADD R1, R1, 2       // this increments address by 2 words (8 bytes)
        ST R2, R1, 0        // write value of R2 into the third array element,
                            // i.e. array[2]
```

Note about instruction execution time

ULP coprocessor is clocked from RTC_FAST_CLK, which is normally derived from the internal 8MHz oscillator. Applications which need to know exact ULP clock frequency can calibrate it against the main XTAL clock:

```
#include "soc/rtc.h"

// calibrate 8M/256 clock against XTAL, get 8M/256 clock period
uint32_t rtc_8md256_period = rtc_clk_cal(RTC_CAL_8MD256, 100);
uint32_t rtc_fast_freq_hz = 1000000ULL * (1 << RTC_CLK_CAL_FRACT) * 256 / rtc_
↪8md256_period;
```

ULP coprocessor needs certain number of clock cycles to fetch each instruction, plus certain number of cycles to execute it, depending on the instruction. See description of each instruction below for details on the execution time.

Instruction fetch time is:

- 2 clock cycles —for instructions following ALU and branch instructions.
- 4 clock cycles —in other cases.

Note that when accessing RTC memories and RTC registers, ULP coprocessor has lower priority than the main CPUs. This means that ULP coprocessor execution may be suspended while the main CPUs access same memory region as the ULP.

Difference between ESP32 ULP and ESP32-S2 ULP Instruction sets

Compare to the ESP32 ULP coprocessor, the ESP-S2 ULP coprocessor has extended instruction set. The ESP32-S2 ULP is not binary compatible with ESP32 ULP, but the assembled program that was written for the ESP32 ULP will also work on the ESP32-S2 ULP after rebuild. The list of the new instructions that was added to the ESP32-S2 ULP is: LDL, LDH, STO, ST32, STI32. The detailed description of these commands please see below.

NOP - no operation

Syntax **NOP**
Operands None
Cycles 2 cycle to execute, 4 cycles to fetch next instruction
Description No operation is performed. Only the PC is incremented.

Example:

```
1:    NOP
```

ADD - Add to register

Syntax **ADD** *Rdst, Rsrc1, Rsrc2*
ADD *Rdst, Rsrc1, imm*
Operands

- *Rdst* - Register R[0..3]
- *Rsrc1* - Register R[0..3]
- *Rsrc2* - Register R[0..3]
- *Imm* - 16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction
Description The instruction adds source register to another source register or to a 16-bit signed value and stores result to the destination register.

Examples:

```
1:    ADD R1, R2, R3          //R1 = R2 + R3

2:    Add R1, R2, 0x1234      //R1 = R2 + 0x1234

3:    .set value1, 0x03       //constant value1=0x03
      Add R1, R2, value1      //R1 = R2 + value1


4:    .global label           //declaration of variable label
      Add R1, R2, label       //R1 = R2 + label
        ...
      label: nop              //definition of variable label
```

SUB - Subtract from register

Syntax **SUB** *Rdst, Rsrc1, Rsrc2*
SUB *Rdst, Rsrc1, imm*

Operands

- *Rdst* - Register R[0..3]
- *Rsrc1* - Register R[0..3]
- *Rsrc2* - Register R[0..3]
- *Imm* - 16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Description The instruction subtracts the source register from another source register or subtracts 16-bit signed value from a source register, and stores result to the destination register.

Examples:

```
1:          SUB R1, R2, R3              //R1 = R2 - R3

2:          sub R1, R2, 0x1234          //R1 = R2 - 0x1234

3:          .set value1, 0x03           //constant value1=0x03
            SUB R1, R2, value1          //R1 = R2 - value1
4:          .global label               //declaration of variable label
            SUB R1, R2, label           //R1 = R2 - label
              ....
  label:    nop                         //definition of variable label
```

AND - Logical AND of two operands

Syntax **AND** *Rdst, Rsrc1, Rsrc2*
AND *Rdst, Rsrc1, imm*

Operands

- *Rdst* - Register R[0..3]
- *Rsrc1* - Register R[0..3]
- *Rsrc2* - Register R[0..3]
- *Imm* - 16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Description The instruction does logical AND of a source register and another source register or 16-bit signed value and stores result to the destination register.

Examples:

```
1:          AND R1, R2, R3              //R1 = R2 & R3

2:          AND R1, R2, 0x1234          //R1 = R2 & 0x1234

3:          .set value1, 0x03           //constant value1=0x03
            AND R1, R2, value1          //R1 = R2 & value1

4:          .global label               //declaration of variable label
            AND R1, R2, label           //R1 = R2 & label
                ...
  label:    nop                         //definition of variable label
```

OR - Logical OR of two operands

Syntax **OR** *Rdst, Rsrc1, Rsrc2*
OR *Rdst, Rsrc1, imm*

Operands

- *Rdst* - Register R[0..3]

- *Rsrc1* - Register R[0..3]
- *Rsrc2* - Register R[0..3]
- *Imm* - 16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Description The instruction does logical OR of a source register and another source register or 16-bit signed value and stores result to the destination register.

Examples:

```
1:        OR R1, R2, R3             //R1 = R2 \| R3

2:        OR R1, R2, 0x1234         //R1 = R2 \| 0x1234

3:        .set value1, 0x03         //constant value1=0x03
          OR R1, R2, value1         //R1 = R2 \| value1

4:        .global label             //declaration of variable label
          OR R1, R2, label          //R1 = R2 \|label
          ...
  label: nop                        //definition of variable label
```

LSH - Logical Shift Left

Syntax **LSH** *Rdst, Rsrc1, Rsrc2*
LSH *Rdst, Rsrc1, imm*

Operands

- *Rdst* - Register R[0..3]
- *Rsrc1* - Register R[0..3]
- *Rsrc2* - Register R[0..3]
- *Imm* - 16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Description The instruction does logical shift to left of source register to number of bits from another source register or 16-bit signed value and store result to the destination register.

Examples:

```
1:        LSH R1, R2, R3              //R1 = R2 << R3

2:        LSH R1, R2, 0x03            //R1 = R2 << 0x03

3:        .set value1, 0x03           //constant value1=0x03
          LSH R1, R2, value1          //R1 = R2 << value1

4:        .global label               //declaration of variable label
          LSH R1, R2, label           //R1 = R2 << label
          ...
  label:  nop                          //definition of variable label
```

RSH - Logical Shift Right

Syntax **RSH** *Rdst, Rsrc1, Rsrc2*
RSH *Rdst, Rsrc1, imm*

Operands *Rdst* - Register R[0..3] *Rsrc1* - Register R[0..3] *Rsrc2* - Register R[0..3] *Imm* - 16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Description The instruction does logical shift to right of source register to number of bits from another source register or 16-bit signed value and store result to the destination register.

Examples:

```
1:        RSH R1, R2, R3              //R1 = R2 >> R3

2:        RSH R1, R2, 0x03            //R1 = R2 >> 0x03

3:        .set value1, 0x03           //constant value1=0x03
          RSH R1, R2, value1          //R1 = R2 >> value1

4:        .global label               //declaration of variable label
          RSH R1, R2, label           //R1 = R2 >> label
  label:  nop                         //definition of variable label
```

MOVE –Move to register

Syntax **MOVE** *Rdst, Rsrc*
MOVE *Rdst, imm*
Operands

- *Rdst* –Register R[0..3]
- *Rsrc* –Register R[0..3]
- *Imm* –16-bit signed value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction
Description The instruction move to destination register value from source register or 16-bit signed value.
Note that when a label is used as an immediate, the address of the label will be converted from bytes to words. This is because LD, ST, and JUMP instructions expect the address register value to be expressed in words rather than bytes. To avoid using an extra instruction

Examples:

```
1:        MOVE       R1, R2           //R1 = R2

2:        MOVE       R1, 0x03         //R1 = 0x03

3:        .set       value1, 0x03     //constant value1=0x03
          MOVE       R1, value1       //R1 = value1

4:        .global     label           //declaration of label
          MOVE        R1, label       //R1 = address_of(label) / 4
          ...
  label:  nop                         //definition of label
```

STL/ST –Store data to the low 16 bits of 32-bits memory

Syntax **ST** *Rsrc, Rdst, offset, Label* **STL** *Rsrc, Rdst, offset, Label*
Operands

- *Rsrc* –Register R[0..3], holds the 16-bit value to store
- *Rdst* –Register R[0..3], address of the destination, in 32-bit words
- *Offset* –11-bit signed value, offset in bytes
- *Label* –2-bit user defined unsigned value

Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction stores the 16-bit value of Rsrc to the lower half-word of memory with address Rdst+offset:

```
Mem[Rdst + offset / 4]{15:0} = {Rsrc[15:0]}
Mem[Rdst + offset / 4]{15:0} = {Label[1:0],Rsrc[13:0]}
```

The ST command introduced to make compatibility with previous versions of UPL core. The application can use higher 16 bits to determine which instruction in the ULP program has written any particular word into memory.

Examples:

```
1:          STL   R1, R2, 0x12          //MEM[R2+0x12] = R1

2:          .data                       //Data section definition
  Addr1:    .word     123               // Define label Addr1 16 bit
            .set      offs, 0x00        // Define constant offs
            .text                       //Text section definition
            MOVE      R1, 1             // R1 = 1
            MOVE      R2, Addr1         // R2 = Addr1
            STL       R1, R2, offs      // MEM[R2 +  0] = R1
                                        // MEM[Addr1 + 0] will be 32'hxxxx0001
3:
            MOVE      R1, 1                 // R1 = 1
            STL       R1, R2, 0x12,1        // MEM[R2+0x12] 0xxxxx4001
```

STH –Store data to the high 16 bits of 32-bits memory

Syntax STH *Rsrc, Rdst, offset, Label*
Operands

- *Rsrc* –Register R[0..3], holds the 16-bit value to store
- *Rdst* –Register R[0..3], address of the destination, in 32-bit words
- *Offset* –11-bit signed value, offset in bytes
- *Label* –2-bit user defined unsigned value

Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction stores the 16-bit value of Rsrc to the high half-word of memory with address Rdst+offset:

```
Mem[Rdst + offset / 4]{31:16} = {Rsrc[15:0]}
Mem[Rdst + offset / 4]{31:16} = {Label[1:0],Rsrc[13:0]}
```

Examples:

```
1:          STH   R1, R2, 0x12          //MEM[R2+0x12][31:16] = R1

2:          .data                       //Data section definition
  Addr1:    .word     123               // Define label Addr1 16 bit
            .set      offs, 0x00        // Define constant offs
            .text                       //Text section definition
            MOVE      R1, 1             // R1 = 1
            MOVE      R2, Addr1         // R2 = Addr1
            STH       R1, R2, offs      // MEM[R2 +  0] = R1
                                        // MEM[Addr1 + 0] will be 32'h0001xxxx
3:
            MOVE      R1, 1                 // R1 = 1
            STH       R1, R2, 0x12, 1       //MEM[R2+0x12] 0x4001xxxx
```

STO –Set offset value for auto increment operation

Syntax STO *offset*
Operands

- *Offset* –11-bit signed value, offset in bytes

Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction set 16-bit value to the offset register:

```
offset = value/ 4
```

Examples:

```
1:          STO   0x12                  // Offset = 0x12/4

2:          .data                       //Data section definition
  Addr1:    .word      123              // Define label Addr1 16 bit
            .set       offs, 0x00       // Define constant offs
            .text                       //Text section definition
            STO        offs             // Offset = 0x00
```

STI –Store data to the 32-bits memory with auto increment of predefined offset address

Syntax **STI** *Rsrc, Rdst, Label*

Operands

- *Rsrc* –Register R[0..3], holds the 16-bit value to store
- *Rdst* –Register R[0..3], address of the destination, in 32-bit words
- *Label* –2-bit user defined unsigned value

Cycles 4 cycles to execute, 4 cycles to fetch next instruction

Description The instruction stores the 16-bit value of Rsrc to the low and high half-word of memory with address Rdst+offset with auto increment of offset:

```
Mem[Rdst + offset / 4]{15:0/31:16} = {Rsrc[15:0]}
Mem[Rdst + offset / 4]{15:0/31:16} = {Label[1:0],Rsrc[13:0]}
```

Examples:

```
1:          STO   0                     // Set offset to 0
            STI   R1, R2, 0x12          //MEM[R2+0x12][15:0] = R1
            STI   R1, R2, 0x12          //MEM[R2+0x12][31:16] = R1

2:          .data                       //Data section definition
  Addr1:    .word      123              // Define label Addr1 16 bit
            .set       offs, 0x00       // Define constant offs
            .text                       //Text section definition
            STO        0                // Set offset to 0
            MOVE       R1, 1            // R1 = 1
            MOVE       R2, Addr1        // R2 = Addr1
            STI        R1, R2           // MEM[R2 +  0] = R1
                                        // MEM[Addr1 + 0] will be 32'hxxxx0001
            STIx        R1, R2           // MEM[R2 +  0] = R1
                                        // MEM[Addr1 + 0] will be 32'h00010001
3:
            STO        0                // Set offset to 0
            MOVE       R1, 1            // R1 = 1
            STI        R1, R2, 1        //MEM[R2+0x12] 0xxxxx4001
            STI        R1, R2, 1        //MEM[R2+0x12] 0x40014001
```

ST32 –Store 32-bits data to the 32-bits memory

Syntax **ST32** *Rsrc, Rdst, offset, Label*

Operands

- *Rsrc* –Register R[0..3], holds the 16-bit value to store
- *Rdst* –Register R[0..3], address of the destination, in 32-bit words
- *Offset* –11-bit signed value, offset in bytes
- *Label* –2-bit user defined unsigned value

Cycles 4 cycles to execute, 4 cycles to fetch next instruction

Description The instruction stores 11 bits of the PC value, label value and the 16-bit value of Rsrc to the 32-bits memory with address Rdst+offset:

```
Mem[Rdst + offset / 4]{31:0} = {PC[10:0],0[2:0],Label[1:0],Rsrc[15:0]}
```

Examples:

```
1:        ST32  R1, R2, 0x12, 0   //MEM[R2+0x12][31:0] = {PC[10:0],0[2:0],
↪Label[1:0],Rsrc[15:0]}

2:        .data                   //Data section definition
  Addr1:  .word     123           // Define label Addr1 16 bit
          .set      offs, 0x00    // Define constant offs
          .text                   //Text section definition
          MOVE      R1, 1         // R1 = 1
          MOVE      R2, Addr1     // R2 = Addr1
          ST32      R1, R2, offs,1// MEM[R2 +  0] = {PC[10:0],0[2:0],Label[1:0],
↪Rsrc[15:0]}
                                  // MEM[Addr1 + 0] will be 32'h00010001
```

STI32 –Store 32-bits data to the 32-bits memory with auto increment of adress offset

Syntax **STI32** *Rsrc, Rdst, Label*
Operands

- *Rsrc* –Register R[0..3], holds the 16-bit value to store
- *Rdst* –Register R[0..3], address of the destination, in 32-bit words
- *Label* –2-bit user defined unsigned value

Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction stores 11 bits of the PC value, label value and the 16-bit value of Rsrc to the 32-bits memory with address Rdst+offset:

```
Mem[Rdst + offset / 4]{31:0} = {PC[10:0],0[2:0],Label[1:0],Rsrc[15:0]}
```

Where offset value set by STO instruction

Examples:

```
1:        STO    0x12
          STI32  R1, R2, 0   //MEM[R2+0x12][31:0] = {PC[10:0],0[2:0],Label[1:0],
↪Rsrc[15:0]}
          STI32  R1, R2, 0   //MEM[R2+0x13][31:0] = {PC[10:0],0[2:0],Label[1:0],
↪Rsrc[15:0]}

2:        .data                   //Data section definition
  Addr1:  .word     123           // Define label Addr1 16 bit
          .set      offs, 0x00    // Define constant offs
          .text                   //Text section definition
          MOVE      R1, 1         // R1 = 1
          MOVE      R2, Addr1     // R2 = Addr1
          STO       offs
          STI32     R1, R2, 1// MEM[R2 +  0] = {PC[10:0],0[2:0],Label[1:0],
↪Rsrc[15:0]}
                                  // MEM[Addr1 + 0] will be 32'h00010001
          ST32      R1, R2, 1// MEM[R2 +  1] = {PC[10:0],0[2:0],Label[1:0],
↪Rsrc[15:0]}
                                  // MEM[Addr1 + 1] will be 32'h00010001
```

LDL/LD –Load data from low part of the 32-bits memory

Syntax **LD** *Rdst, Rsrc, offset* **LDL** *Rdst, Rsrc, offset*
Operands *Rdst* –Register R[0..3], destination
Rsrc –Register R[0..3], holds address of destination, in 32-bit words
Offset –10-bit signed value, offset in bytes

Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction loads lower 16-bit half-word from memory with address Rsrc+offset into the destination register Rdst:

```
Rdst[15:0] = Mem[Rsrc + offset / 4][15:0]
```

The LD command do the same as LDL, and included for compatibility with previous versions of ULP core.

Examples:

```
1:          LDL   R1, R2, 0x12              //R1 = MEM[R2+0x12]

2:          .data                          //Data section definition
  Addr1:    .word      123                 // Define label Addr1 16 bit
            .set       offs, 0x00          // Define constant offs
            .text                          //Text section definition
            MOVE       R1, 1               // R1 = 1
            MOVE       R2, Addr1           // R2 = Addr1 / 4 (address of label is␣
↪converted into words)
            LDL        R1, R2, offs        // R1 = MEM[R2 +  0]
                                           // R1 will be 123
```

LDH –Load data from high part of the 32-bits memory

Syntax **LDH** *Rdst, Rsrc, offset*
Operands *Rdst* –Register R[0..3], destination
Rsrc –Register R[0..3], holds address of destination, in 32-bit words
Offset –10-bit signed value, offset in bytes
Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction loads higher 16-bit half-word from memory with address Rsrc+offset into the destination register Rdst:

```
Rdst[15:0] = Mem[Rsrc + offset / 4][15:0]
```

The LD command do the same as LDL, and included for compatibility with previous versions of ULP core.

Examples:

```
1:          LDH   R1, R2, 0x12              //R1 = MEM[R2+0x12]

2:          .data                          //Data section definition
  Addr1:    .word      0x12345678                 // Define label Addr1 16 bit
            .set       offs, 0x00          // Define constant offs
            .text                          //Text section definition
            MOVE       R1, 1               // R1 = 1
            MOVE       R2, Addr1           // R2 = Addr1 / 4 (address of label is␣
↪converted into words)
            LDH        R1, R2, offs        // R1 = MEM[R2 +  0]
                                           // R1 will be 0x1234
```

JUMP –Jump to an absolute address

Syntax **JUMP** *Rdst*
JUMP *ImmAddr*
JUMP *Rdst, Condition*
JUMP *ImmAddr, Condition*
Operands

- *Rdst* –Register R[0..3] containing address to jump to (expressed in 32-bit words)
- *ImmAddr* –13 bits address (expressed in bytes), aligned to 4 bytes
- ***Condition*:**

- EQ –jump if last ALU operation result was zero
- OV –jump if last ALU has set overflow flag

Cycles 2 cycles to execute, 2 cycles to fetch next instruction

Description The instruction makes jump to the specified address. Jump can be either unconditional or based on an ALU flag.

Examples:

```
1:        JUMP       R1             // Jump to address in R1 (address in R1 is in␣
↪32-bit words)

2:        JUMP       0x120, EQ      // Jump to address 0x120 (in bytes) if ALU␣
↪result is zero

3:        JUMP       label          // Jump to label
          ...
  label:  nop                       // Definition of label

4:        .global    label          // Declaration of global label

          MOVE       R1, label      // R1 = label (value loaded into R1 is in words)
          JUMP       R1             // Jump to label
          ...
  label:  nop                       // Definition of label
```

JUMPR –Jump to a relative offset (condition based on R0)

Syntax **JUMPR** *Step, Threshold, Condition*

Operands

- *Step* –relative shift from current position, in bytes
- *Threshold* –threshold value for branch condition
- ***Condition*:**
  - *EQ* (equal) –jump if value in R0 == threshold
  - *LT* (less than) –jump if value in R0 < threshold
  - *LE* (less or equal) –jump if value in R0 <= threshold
  - *GT* (greater than) –jump if value in R0 > threshold
  - *GE* (greater or equal) –jump if value in R0 >= threshold

Cycles Conditions *EQ*, *GT* and *LT*: 2 cycles to execute, 2 cycles to fetch next instruction

Conditions *LE* and *GE* are implemented in the assembler using two **JUMPR** instructions:

```
// JUMPR target, threshold, LE is implemented as:

          JUMPR target, threshold, EQ
          JUMPR target, threshold, LT

// JUMPR target, threshold, GE is implemented as:

          JUMPR target, threshold, EQ
          JUMPR target, threshold, GT
```

Therefore the execution time will depend on the branches taken: either 2 cycles to execute + 2 cycles to fetch, or 4 cycles to execute + 4 cycles to fetch.

Description The instruction makes a jump to a relative address if condition is true. Condition is the result of comparison of R0 register value and the threshold value.

Examples:

```
1:pos:    JUMPR       16, 20, GE    // Jump to address (position + 16 bytes) if␣
↪value in R0 >= 20

2:        // Down counting loop using R0 register
```

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```
          MOVE        R0, 16        // load 16 into R0
  label:  SUB         R0, R0, 1     // R0--
          NOP                       // do something
          JUMPR       label, 1, GE // jump to label if R0 >= 1
```

JUMPS –Jump to a relative address (condition based on stage count)

Syntax **JUMPS** *Step, Threshold, Condition*

Operands

- *Step* –relative shift from current position, in bytes
- *Threshold* –threshold value for branch condition
- ***Condition*:**
  - *EQ* (equal) –jump if value in stage_cnt == threshold
  - *LT* (less than) –jump if value in stage_cnt < threshold
  - *LE* (less or equal) - jump if value in stage_cnt <= threshold
  - *GT* (greater than) –jump if value in stage_cnt > threshold
  - *GE* (greater or equal) —jump if value in stage_cnt >= threshold

Cycles 2 cycles to execute, 2 cycles to fetch next instruction:

```
// JUMPS target, threshold, EQ is implemented as:

          JUMPS next, threshold, LT
          JUMPS target, threshold, LE
next:

// JUMPS target, threshold, GT is implemented as:

          JUMPS next, threshold, LE
          JUMPS target, threshold, GE
next:
```

Therefore the execution time will depend on the branches taken: either 2 cycles to execute + 2 cycles to fetch, or 4 cycles to execute + 4 cycles to fetch.

Description The instruction makes a jump to a relative address if condition is true. Condition is the result of comparison of count register value and threshold value.

Examples:

```
1:pos:    JUMPS     16, 20, EQ     // Jump to (position + 16 bytes) if stage_cnt␣
↪== 20

2:        // Up counting loop using stage count register
          STAGE_RST                  // set stage_cnt to 0
  label:  STAGE_INC  1               // stage_cnt++
          NOP                        // do something
          JUMPS       label, 16, LT  // jump to label if stage_cnt < 16
```

STAGE_RST –Reset stage count register

Syntax **STAGE_RST**

Operands No operands

Description The instruction sets the stage count register to 0

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Examples:

```
1:        STAGE_RST      // Reset stage count register
```

STAGE_INC –Increment stage count register

Syntax **STAGE_INC** *Value*
Operands

- *Value* –8 bits value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction
Description The instruction increments stage count register by given value.

Examples:

```
1:        STAGE_INC      10          // stage_cnt += 10

2:        // Up counting loop example:
          STAGE_RST                  // set stage_cnt to 0
  label:  STAGE_INC  1               // stage_cnt++
          NOP                        // do something
          JUMPS      label, 16, LT   // jump to label if stage_cnt < 16
```

STAGE_DEC –Decrement stage count register

Syntax **STAGE_DEC** *Value*
Operands

- *Value* –8 bits value

Cycles 2 cycles to execute, 4 cycles to fetch next instruction
Description The instruction decrements stage count register by given value.

Examples:

```
1:        STAGE_DEC      10        // stage_cnt -= 10;

2:        // Down counting loop exaple
          STAGE_RST                // set stage_cnt to 0
          STAGE_INC  16            // increment stage_cnt to 16
  label:  STAGE_DEC  1             // stage_cnt--;
          NOP                      // do something
          JUMPS      label, 0, GT  // jump to label if stage_cnt > 0
```

HALT –End the program

Syntax **HALT**
Operands No operands
Cycles 2 cycles to execute
Description The instruction halts the ULP coprocessor and restarts ULP wakeup timer, if it is enabled.

Examples:

```
1:       HALT      // Halt the coprocessor
```

WAKE –Wake up the chip

Syntax **WAKE**
Operands No operands
Cycles 2 cycles to execute, 4 cycles to fetch next instruction
Description The instruction sends an interrupt from ULP to RTC controller.

- If the SoC is in deep sleep mode, and ULP wakeup is enabled, this causes the SoC to wake up.
- If the SoC is not in deep sleep mode, and ULP interrupt bit (RTC_CNTL_ULP_CP_INT_ENA) is set in RTC_CNTL_INT_ENA_REG register, RTC interrupt will be triggered.

Note that before using WAKE instruction, ULP program may needs to wait until RTC controller is ready to wake up the main CPU. This is indicated using RTC_CNTL_RDY_FOR_WAKEUP bit of RTC_CNTL_LOW_POWER_ST_REG register. If WAKE instruction is executed while RTC_CNTL_RDY_FOR_WAKEUP is zero, it has no effect (wake up does not occur).

Examples:

```
1: is_rdy_for_wakeup:                    // Read RTC_CNTL_RDY_FOR_WAKEUP bit
          READ_RTC_FIELD(RTC_CNTL_LOW_POWER_ST_REG, RTC_CNTL_RDY_FOR_WAKEUP)
          AND r0, r0, 1
          JUMP is_rdy_for_wakeup, eq     // Retry until the bit is set
          WAKE                           // Trigger wake up
          REG_WR 0x006, 24, 24, 0        // Stop ULP timer (clear RTC_CNTL_ULP_CP_
↪SLP_TIMER_EN)
          HALT                           // Stop the ULP program
          // After these instructions, SoC will wake up,
          // and ULP will not run again until started by the main program.
```

SLEEP –set ULP wakeup timer period

Syntax **SLEEP** *sleep_reg*

Operands

- *sleep_reg* –0..4, selects one of `SENS_ULP_CP_SLEEP_CYCx_REG` registers.

Cycles 2 cycles to execute, 4 cycles to fetch next instruction

Description The instruction selects which of the `SENS_ULP_CP_SLEEP_CYCx_REG` (x = 0..4) register values is to be used by the ULP wakeup timer as wakeup period. By default, the value from `SENS_ULP_CP_SLEEP_CYC0_REG` is used.

Examples:

```
1:        SLEEP     1          // Use period set in SENS_ULP_CP_SLEEP_CYC1_REG

2:        .set sleep_reg, 4    // Set constant
          SLEEP  sleep_reg     // Use period set in SENS_ULP_CP_SLEEP_CYC4_REG
```

WAIT –wait some number of cycles

Syntax **WAIT** *Cycles*

Operands

- *Cycles* –number of cycles for wait

Cycles 2 + *Cycles* cycles to execute, 4 cycles to fetch next instruction

Description The instruction delays for given number of cycles.

Examples:

```
1:        WAIT     10          // Do nothing for 10 cycles

2:        .set  wait_cnt, 10   // Set a constant
          WAIT  wait_cnt       // wait for 10 cycles
```

TSENS –do measurement with temperature sensor

Syntax

- **TSENS** *Rdst, Wait_Delay*

Operands

- *Rdst* –Destination Register R[0..3], result will be stored to this register
- *Wait_Delay* –number of cycles used to perform the measurement

Cycles 2 + *Wait_Delay* + 3 * TSENS_CLK to execute, 4 cycles to fetch next instruction

Description The instruction performs measurement using TSENS and stores the result into a general purpose register.

Examples:

```
1:        TSENS     R1, 1000      // Measure temperature sensor for 1000 cycles,
                                  // and store result to R1
```

ADC –do measurement with ADC

Syntax

- **ADC** *Rdst, Sar_sel, Mux*
- **ADC** *Rdst, Sar_sel, Mux, 0* —deprecated form

Operands

- *Rdst* –Destination Register R[0..3], result will be stored to this register
- *Sar_sel* –Select ADC: 0 = SARADC1, 1 = SARADC2
- *Mux* - selected PAD, SARADC Pad[Mux+1] is enabled

Cycles `23 + max(1, SAR_AMP_WAIT1) + max(1, SAR_AMP_WAIT2) + max(1, SAR_AMP_WAIT3) + SARx_SAMPLE_CYCLE + SARx_SAMPLE_BIT` cycles to execute, 4 cycles to fetch next instruction

Description The instruction makes measurements from ADC.

Examples:

```
1:        ADC       R1, 0, 1      // Measure value using ADC1 pad 2 and store␣
↪result into R1
```

I2C_RD - read single byte from I2C slave

Syntax

- **I2C_RD** *Sub_addr, High, Low, Slave_sel*

Operands

- *Sub_addr* –Address within the I2C slave to read.
- *High*, *Low* —Define range of bits to read. Bits outside of [High, Low] range are masked.
- *Slave_sel* - Index of I2C slave address to use.

Cycles Execution time mostly depends on I2C communication time. 4 cycles to fetch next instruction.

Description `I2C_RD` instruction reads one byte from I2C slave with index `Slave_sel`. Slave address (in 7-bit format) has to be set in advance into *SENS_I2C_SLAVE_ADDRx* register field, where `x == Slave_sel`. 8 bits of read result is stored into *R0* register.

Examples:

```
1:        I2C_RD      0x10, 7, 0, 0      // Read byte from sub-address 0x10 of␣
↪slave with address set in SENS_I2C_SLAVE_ADDR0
```

I2C_WR - write single byte to I2C slave

Syntax

- **I2C_WR** *Sub_addr, Value, High, Low, Slave_sel*

Operands

- *Sub_addr* –Address within the I2C slave to write.
- *Value* –8-bit value to be written.
- *High*, *Low* —Define range of bits to write. Bits outside of [High, Low] range are masked.
- *Slave_sel* - Index of I2C slave address to use.

Cycles Execution time mostly depends on I2C communication time. 4 cycles to fetch next instruction.

Description `I2C_WR` instruction writes one byte to I2C slave with index `Slave_sel`. Slave address (in 7-bit format) has to be set in advance into *SENS_I2C_SLAVE_ADDRx* register field, where `x == Slave_sel`.

Examples:

```
1:        I2C_WR      0x20, 0x33, 7, 0, 1      // Write byte 0x33 to sub-address␣
↪0x20 of slave with address set in SENS_I2C_SLAVE_ADDR1.
```

REG_RD –read from peripheral register

Syntax **REG_RD** *Addr, High, Low*
Operands

- *Addr* –Register address, in 32-bit words
- *High* –Register end bit number
- *Low* –Register start bit number

Cycles 4 cycles to execute, 4 cycles to fetch next instruction
Description The instruction reads up to 16 bits from a peripheral register into a general purpose register: `R0 = REG[Addr][High:Low]`.
This instruction can access registers in RTC_CNTL, RTC_IO, SENS, and RTC_I2C peripherals. Address of the the register, as seen from the ULP, can be calculated from the address of the same register on the PeriBUS1 as follows:

```
addr_ulp = (addr_peribus1 - DR_REG_RTCCNTL_BASE) / 4
```

Examples:

```
1:        REG_RD      0x120, 7, 4     // load 4 bits: R0 = {12'b0, REG[0x120][7:4]}
```

REG_WR –write to peripheral register

Syntax **REG_WR** *Addr, High, Low, Data*
Operands

- *Addr* –Register address, in 32-bit words.
- *High* –Register end bit number
- *Low* –Register start bit number
- *Data* –Value to write, 8 bits

Cycles 8 cycles to execute, 4 cycles to fetch next instruction
Description The instruction writes up to 8 bits from an immediate data value into a peripheral register: `REG[Addr][High:Low] = data`.
This instruction can access registers in RTC_CNTL, RTC_IO, SENS, and RTC_I2C peripherals. Address of the the register, as seen from the ULP, can be calculated from the address of the same register on the PeriBUS1 as follows:

```
addr_ulp = (addr_peribus1 - DR_REG_RTCCNTL_BASE) / 4
```

Examples:

```
1:        REG_WR      0x120, 7, 0, 0x10   // set 8 bits: REG[0x120][7:0] = 0x10
```

Convenience macros for peripheral registers access

ULP source files are passed through C preprocessor before the assembler. This allows certain macros to be used to facilitate access to peripheral registers.

Some existing macros are defined in `soc/soc_ulp.h` header file. These macros allow access to the fields of peripheral registers by their names. Peripheral registers names which can be used with these macros are the ones defined in `soc/rtc_cntl_reg.h`, `soc/rtc_io_reg.h`, `soc/sens_reg.h`, and `soc/rtc_i2c_reg.h`.

READ_RTC_REG(rtc_reg, low_bit, bit_width) Read up to 16 bits from rtc_reg[low_bit + bit_width - 1 : low_bit] into R0. For example:

```
#include "soc/soc_ulp.h"
#include "soc/rtc_cntl_reg.h"

/* Read 16 lower bits of RTC_CNTL_TIME0_REG into R0 */
READ_RTC_REG(RTC_CNTL_TIME0_REG, 0, 16)
```

READ_RTC_FIELD(rtc_reg, field) Read from a field in rtc_reg into R0, up to 16 bits. For example:

```
#include "soc/soc_ulp.h"
#include "soc/sens_reg.h"

/* Read 8-bit SENS_TSENS_OUT field of SENS_SAR_SLAVE_ADDR3_REG into R0 */
READ_RTC_FIELD(SENS_SAR_SLAVE_ADDR3_REG, SENS_TSENS_OUT)
```

WRITE_RTC_REG(rtc_reg, low_bit, bit_width, value) Write immediate value into rtc_reg[low_bit + bit_width - 1 : low_bit], bit_width <= 8. For example:

```
#include "soc/soc_ulp.h"
#include "soc/rtc_io_reg.h"

/* Set BIT(2) of RTC_GPIO_OUT_DATA_W1TS field in RTC_GPIO_OUT_W1TS_REG */
WRITE_RTC_REG(RTC_GPIO_OUT_W1TS_REG, RTC_GPIO_OUT_DATA_W1TS_S + 2, 1, 1)
```

WRITE_RTC_FIELD(rtc_reg, field, value) Write immediate value into a field in rtc_reg, up to 8 bits. For example:

```
#include "soc/soc_ulp.h"
#include "soc/rtc_cntl_reg.h"

/* Set RTC_CNTL_ULP_CP_SLP_TIMER_EN field of RTC_CNTL_STATE0_REG to 0 */
WRITE_RTC_FIELD(RTC_CNTL_STATE0_REG, RTC_CNTL_ULP_CP_SLP_TIMER_EN, 0)
```

4.17.2 Programming ULP coprocessor using C macros (legacy)

In addition to the existing binutils port for the ESP32 ULP coprocessor, it is possible to generate programs for the ULP by embedding assembly-like macros into an ESP32 application. Here is an example how this can be done:

```
const ulp_insn_t program[] = {
    I_MOVI(R3, 16),          // R3 <- 16
    I_LD(R0, R3, 0),         // R0 <- RTC_SLOW_MEM[R3 + 0]
    I_LD(R1, R3, 1),         // R1 <- RTC_SLOW_MEM[R3 + 1]
    I_ADDR(R2, R0, R1),      // R2 <- R0 + R1
    I_ST(R2, R3, 2),         // R2 -> RTC_SLOW_MEM[R2 + 2]
    I_HALT()
};
size_t load_addr = 0;
size_t size = sizeof(program)/sizeof(ulp_insn_t);
ulp_process_macros_and_load(load_addr, program, &size);
ulp_run(load_addr);
```

The `program` array is an array of `ulp_insn_t`, i.e. ULP coprocessor instructions. Each `I_XXX` preprocessor define translates into a single 32-bit instruction. Arguments of these preprocessor defines can be register numbers (`R0` —`R3`) and literal constants. See *ULP coprocessor instruction defines* section for descriptions of instructions and arguments they take.

注解: Because some of the instruction macros expand to inline function calls, defining such array in global scope will cause the compiler to produce an "initializer element is not constant" error. To fix this error, move the definition of instructions array into local scope.

Load and store instructions use addresses expressed in 32-bit words. Address 0 corresponds to the first word of `RTC_SLOW_MEM` (which is address 0x50000000 as seen by the main CPUs).

To generate branch instructions, special `M_` preprocessor defines are used. `M_LABEL` define can be used to define a branch target. Label identifier is a 16-bit integer. `M_Bxxx` defines can be used to generate branch instructions with target set to a particular label.

Implementation note: these `M_` preprocessor defines will be translated into two `ulp_insn_t` values: one is a token value which contains label number, and the other is the actual instruction. `ulp_process_macros_and_load` function resolves the label number to the address, modifies the branch instruction to use the correct address, and removes the the extra `ulp_insn_t` token which contains the label numer.

Here is an example of using labels and branches:

```
const ulp_insn_t program[] = {
    I_MOVI(R0, 34),          // R0 <- 34
    M_LABEL(1),              // label_1
    I_MOVI(R1, 32),          // R1 <- 32
    I_LD(R1, R1, 0),         // R1 <- RTC_SLOW_MEM[R1]
    I_MOVI(R2, 33),          // R2 <- 33
    I_LD(R2, R2, 0),         // R2 <- RTC_SLOW_MEM[R2]
    I_SUBR(R3, R1, R2),      // R3 <- R1 - R2
    I_ST(R3, R0, 0),         // R3 -> RTC_SLOW_MEM[R0 + 0]
    I_ADDI(R0, R0, 1),       // R0++
    M_BL(1, 64),             // if (R0 < 64) goto label_1
    I_HALT(),
};
RTC_SLOW_MEM[32] = 42;
RTC_SLOW_MEM[33] = 18;
size_t load_addr = 0;
size_t size = sizeof(program)/sizeof(ulp_insn_t);
ulp_process_macros_and_load(load_addr, program, &size);
ulp_run(load_addr);
```

Application Example

Demonstration of entering into deep sleep mode and waking up using several wake up sources: system/deep_sleep.

API Reference

Header File

- ulp/include/esp32s2/ulp.h

Functions

esp_err_t **ulp_process_macros_and_load**(uint32_t *load_addr*, **const** ulp_insn_t **program*, size_t **psize*)

Resolve all macro references in a program and load it into RTC memory.

Return

- ESP_OK on success
- ESP_ERR_NO_MEM if auxiliary temporary structure can not be allocated
- one of ESP_ERR_ULP_xxx if program is not valid or can not be loaded

Parameters

- `load_addr`: address where the program should be loaded, expressed in 32-bit words
- `program`: ulp_insn_t array with the program
- `psize`: size of the program, expressed in 32-bit words

esp_err_t **ulp_run**(uint32_t *entry_point*)

Run the program loaded into RTC memory.

Return ESP_OK on success

Parameters

- `entry_point`: entry point, expressed in 32-bit words

Error codes

ESP_ERR_ULP_BASE
Offset for ULP-related error codes

ESP_ERR_ULP_SIZE_TOO_BIG
Program doesn't fit into RTC memory reserved for the ULP

ESP_ERR_ULP_INVALID_LOAD_ADDR
Load address is outside of RTC memory reserved for the ULP

ESP_ERR_ULP_DUPLICATE_LABEL
More than one label with the same number was defined

ESP_ERR_ULP_UNDEFINED_LABEL
Branch instructions references an undefined label

ESP_ERR_ULP_BRANCH_OUT_OF_RANGE
Branch target is out of range of B instruction (try replacing with BX)

ULP coprocessor registers

ULP co-processor has 4 16-bit general purpose registers. All registers have same functionality, with one exception. R0 register is used by some of the compare-and-branch instructions as a source register.

These definitions can be used for all instructions which require a register.

R0
general purpose register 0

R1
general purpose register 1

R2
general purpose register 2

R3
general purpose register 3

ULP coprocessor instruction defines

I_DELAY(cycles_)
Delay (nop) for a given number of cycles

I_HALT()
Halt the coprocessor.

This instruction halts the coprocessor, but keeps ULP timer active. As such, ULP program will be restarted again by timer. To stop the program and prevent the timer from restarting the program, use I_END(0) instruction.

I_END()
Stop ULP program timer.

This is a convenience macro which disables the ULP program timer. Once this instruction is used, ULP program will not be restarted anymore until ulp_run function is called.

ULP program will continue running after this instruction. To stop the currently running program, use I_HALT().

I_ST(reg_val, reg_addr, offset_)
Store value from register reg_val into RTC memory.

The value is written to an offset calculated by adding value of reg_addr register and offset_ field (this offset is expressed in 32-bit words). 32 bits written to RTC memory are built as follows:

- bits [31:21] hold the PC of current instruction, expressed in 32-bit words
- bits [20:16] = 5’ b1
- bits [15:0] are assigned the contents of reg_val

RTC_SLOW_MEM[addr + offset_] = { 5’ b0, insn_PC[10:0], val[15:0] }

I_LD (reg_dest, reg_addr, offset_)
Load value from RTC memory into reg_dest register.

Loads 16 LSBs from RTC memory word given by the sum of value in reg_addr and value of offset_.

I_WR_REG (reg, low_bit, high_bit, val)
Write literal value to a peripheral register

reg[high_bit : low_bit] = val This instruction can access RTC_CNTL_, RTC_IO_, SENS_, and RTC_I2C peripheral registers.

I_RD_REG (reg, low_bit, high_bit)
Read from peripheral register into R0

R0 = reg[high_bit : low_bit] This instruction can access RTC_CNTL_, RTC_IO_, SENS_, and RTC_I2C peripheral registers.

I_BL (pc_offset, imm_value)
Branch relative if R0 less than immediate value.

pc_offset is expressed in words, and can be from -127 to 127 imm_value is a 16-bit value to compare R0 against

I_BGE (pc_offset, imm_value)
Branch relative if R0 greater or equal than immediate value.

pc_offset is expressed in words, and can be from -127 to 127 imm_value is a 16-bit value to compare R0 against

I_BXR (reg_pc)
Unconditional branch to absolute PC, address in register.

reg_pc is the register which contains address to jump to. Address is expressed in 32-bit words.

I_BXI (imm_pc)
Unconditional branch to absolute PC, immediate address.

Address imm_pc is expressed in 32-bit words.

I_BXZR (reg_pc)
Branch to absolute PC if ALU result is zero, address in register.

reg_pc is the register which contains address to jump to. Address is expressed in 32-bit words.

I_BXZI (imm_pc)
Branch to absolute PC if ALU result is zero, immediate address.

Address imm_pc is expressed in 32-bit words.

I_BXFR (reg_pc)
Branch to absolute PC if ALU overflow, address in register

reg_pc is the register which contains address to jump to. Address is expressed in 32-bit words.

I_BXFI (imm_pc)
Branch to absolute PC if ALU overflow, immediate address

Address imm_pc is expressed in 32-bit words.

I_ADDR (reg_dest, reg_src1, reg_src2)
Addition: dest = src1 + src2

I_SUBR (reg_dest, reg_src1, reg_src2)
Subtraction: dest = src1 - src2

I_ANDR (reg_dest, reg_src1, reg_src2)
Logical AND: dest = src1 & src2

I_ORR (reg_dest, reg_src1, reg_src2)
Logical OR: dest = src1 | src2

I_MOVR (reg_dest, reg_src)
Copy: dest = src

I_LSHR (reg_dest, reg_src, reg_shift)
Logical shift left: dest = src « shift

I_RSHR (reg_dest, reg_src, reg_shift)
Logical shift right: dest = src » shift

I_ADDI (reg_dest, reg_src, imm_)
Add register and an immediate value: dest = src1 + imm

I_SUBI (reg_dest, reg_src, imm_)
Subtract register and an immediate value: dest = src - imm

I_ANDI (reg_dest, reg_src, imm_)
Logical AND register and an immediate value: dest = src & imm

I_ORI (reg_dest, reg_src, imm_)
Logical OR register and an immediate value: dest = src | imm

I_MOVI (reg_dest, imm_)
Copy an immediate value into register: dest = imm

I_LSHI (reg_dest, reg_src, imm_)
Logical shift left register value by an immediate: dest = src « imm

I_RSHI (reg_dest, reg_src, imm_)
Logical shift right register value by an immediate: dest = val » imm

M_LABEL (label_num)
Define a label with number label_num.

This is a macro which doesn’ t generate a real instruction. The token generated by this macro is removed by ulp_process_macros_and_load function. Label defined using this macro can be used in branch macros defined below.

M_BL (label_num, imm_value)
Macro: branch to label label_num if R0 is less than immediate value.

This macro generates two ulp_insn_t values separated by a comma, and should be used when defining contents of ulp_insn_t arrays. First value is not a real instruction; it is a token which is removed by ulp_process_macros_and_load function.

M_BGE (label_num, imm_value)
Macro: branch to label label_num if R0 is greater or equal than immediate value

This macro generates two ulp_insn_t values separated by a comma, and should be used when defining contents of ulp_insn_t arrays. First value is not a real instruction; it is a token which is removed by ulp_process_macros_and_load function.

M_BX (label_num)
Macro: unconditional branch to label

This macro generates two ulp_insn_t values separated by a comma, and should be used when defining contents of ulp_insn_t arrays. First value is not a real instruction; it is a token which is removed by ulp_process_macros_and_load function.

M_BXZ (label_num)
Macro: branch to label if ALU result is zero

This macro generates two ulp_insn_t values separated by a comma, and should be used when defining contents of ulp_insn_t arrays. First value is not a real instruction; it is a token which is removed by ulp_process_macros_and_load function.

M_BXF(label_num)
Macro: branch to label if ALU overflow

This macro generates two ulp_insn_t values separated by a comma, and should be used when defining contents of ulp_insn_t arrays. First value is not a real instruction; it is a token which is removed by ulp_process_macros_and_load function.

Defines

RTC_SLOW_MEM
RTC slow memory, 8k size

ULP（Ultra Low Power 超低功耗）协处理器是一种简单的有限状态机 (FSM)，可以在主处理器处于深度睡眠模式时，使用 ADC、温度传感器和外部 I2C 传感器执行测量操作。ULP 协处理器可以访问 RTC_SLOW_MEM 内存区域及 RTC_CNTL、RTC_IO、SARADC 外设中的寄存器。ULP 协处理器使用 32 位固定宽度的指令，32 位内存寻址，配备 4 个 16 位通用寄存器。

4.17.3 安装工具链

ULP 协处理器代码是用汇编语言编写的，并使用 binutils-esp32ulp 工具链 进行编译。

如果你已经按照快速入门指南 中的介绍安装好了 ESP-IDF 及其 CMake 构建系统，那么 ULP 工具链已经被默认安装到了你的开发环境中。

4.17.4 编译 ULP 代码

若需要将 ULP 代码编译为某组件的一部分，则必须执行以下步骤：

1. 用汇编语言编写的 ULP 代码必须导入到一个或多个 *.S* 扩展文件中，且这些文件必须放在组件目录中一个独立的目录中，例如 *ulp/*。
2. 注册后从组件 CMakeLists.txt 中调用 `ulp_embed_binary` 示例如下：

```
...
idf_component_register()

set(ulp_app_name ulp_${COMPONENT_NAME})
set(ulp_s_sources ulp/ulp_assembly_source_file.S)
set(ulp_exp_dep_srcs "ulp_c_source_file.c")

ulp_embed_binary(${ulp_app_name} "${ulp_s_sources}" "${ulp_exp_dep_srcs}")
```

`ulp_embed_binary` 的第一个参数为 ULP 二进制文件命名。指定的此名称也用于生成的其他文件，如：ELF 文件、.map 文件、头文件和链接器导出文件。第二个参数指定 ULP 程序集源文件。最后，第三个参数指定组件源文件列表，其中包括被生成的头文件。此列表用以建立正确的依赖项，并确保在编译这些文件之前先创建生成的头文件。有关 ULP 应用程序生成的头文件等相关概念，请参考下文。

3. 使用常规方法（例如 *idf.py app*）编译应用程序
   在内部，构建系统将按照以下步骤编译 ULP 程序：
   1. **通过 C 预处理器运行每个程序集文件 (foo.S)**。此步骤在组件编译目录中生成预处理的程序集文件 (foo.ulp.S)，同时生成依赖文件 (foo.ulp.d)。
   2. **通过汇编器运行预处理过的汇编源码**。此步骤会生成目标文件 (foo.ulp.o) 和清单 (foo.ulp.lst)。清单文件仅用于调试，不用于编译进程的后续步骤。
   3. **通过 C 预处理器运行链接器脚本模板**。模板位于 `components/ulp/ld` 目录中。
   4. **将目标文件链接到 ELF 输出文件** (`ulp_app_name.elf`)。此步骤生成的.map 文件 (`ulp_app_name.map`) 默认用于调试。

5. **将 ELF 文件中的内容转储为二进制文件** (`ulp_app_name.bin`)，以便嵌入到应用程序中。
6. 使用 `esp32ulp-elf-nm` 在 ELF 文件中 **生成全局符号列表** (`ulp_app_name.sym`)。
7. **创建 LD 导出脚本和头文件** (`ulp_app_name.ld` 和 `ulp_app_name.h`)，包含来自 `ulp_app_name.sym` 的符号。此步骤可借助 `esp32ulp_mapgen.py` 工具来完成。
8. **将生成的二进制文件添加到要嵌入应用程序的二进制文件列表中**。

4.17.5 访问 ULP 程序变量

在 ULP 程序中定义的全局符号也可以在主程序中使用。

例如，ULP 程序可以定义 `measurement_count` 变量，此变量可以定义程序从深度睡眠中唤醒芯片之前需要进行的 ADC 测量的次数:

```
                        .global measurement_count
measurement_count:      .long 0

                        /* later, use measurement_count */
                        move r3, measurement_count
                        ld r3, r3, 0
```

主程序需要在启动 ULP 程序之前初始化 `measurement_count` 变量，构建系统通过生成定义 ULP 编程中全局符号的 `${ULP_APP_NAME}.h` 和 `${ULP_APP_NAME}.ld` 文件实现上述操作。这些文件包含了在 ULP 程序中定义的所有全局符号，文件以 `ulp_` 开头。

头文件包含对此类符号的声明:

```
extern uint32_t ulp_measurement_count;
```

注意，所有符号（包括变量、数组、函数）均被声明为 `uint32_t`。对于函数和数组，先获取符号地址，然后转换为适当的类型。

生成的链接器脚本文件定义了 RTC_SLOW_MEM 中的符号位置:

```
PROVIDE ( ulp_measurement_count = 0x50000060 );
```

如果要从主程序访问 ULP 程序变量，应先使用 `include` 语句包含生成的头文件，这样，就可以像访问常规变量一样访问 ulp 程序变量。操作如下:

```
#include "ulp_app_name.h"

// later
void init_ulp_vars() {
    ulp_measurement_count = 64;
}
```

注意，ULP 程序在 RTC 内存中只能使用 32 位字的低 16 位，因为寄存器是 16 位的，并且不具备从字的高位加载的指令。

同样，ULP 储存指令将寄存器值写入 32 位字的低 16 位中。高 16 位写入的值取决于储存指令的地址，因此在读取 ULP 写的变量时，主应用程序需要屏蔽高 16 位，例如:

```
printf("Last measurement value: %d\n", ulp_last_measurement & UINT16_MAX);
```

4.17.6 启动 ULP 程序

要运行 ULP 程序，主应用程序需要调用 `ulp_load_binary` 函数将 ULP 程序加载到 RTC 内存中，然后调用 `ulp_run` 函数，启动 ULP 程序。

注意，在 menuconfig 中必须启用 “Enable Ultra Low Power (ULP) Coprocessor” 选项，以便为 ULP 预留内存。”RTC slow memory reserved for coprocessor” 选项设置的值必须足够储存 ULP 代码和数据。如果应用程序组件包含多个 ULP 程序，则 RTC 内存必须足以容纳最大的程序。

每个 ULP 程序均以二进制 BLOB 的形式嵌入到 ESP-IDF 应用程序中。应用程序可以引用此 BLOB，并以下面的方式加载此 BLOB（假设 ULP_APP_NAME 已被定义为 `ulp_app_name`）：

```
extern const uint8_t bin_start[] asm("_binary_ulp_app_name_bin_start");
extern const uint8_t bin_end[]   asm("_binary_ulp_app_name_bin_end");

void start_ulp_program() {
    ESP_ERROR_CHECK( ulp_load_binary(
        0 /* load address, set to 0 when using default linker scripts */,
        bin_start,
        (bin_end - bin_start) / sizeof(uint32_t)) );
}
```

esp_err_t **ulp_load_binary**(uint32_t *load_addr*, **const** uint8_t **program_binary*, size_t *program_size*)

Load ULP program binary into RTC memory.

ULP program binary should have the following format (all values little-endian):

1. MAGIC, (value 0x00706c75, 4 bytes)
2. TEXT_OFFSET, offset of .text section from binary start (2 bytes)
3. TEXT_SIZE, size of .text section (2 bytes)
4. DATA_SIZE, size of .data section (2 bytes)
5. BSS_SIZE, size of .bss section (2 bytes)
6. (TEXT_OFFSET - 12) bytes of arbitrary data (will not be loaded into RTC memory)
7. .text section
8. .data section

Linker script in components/ulp/ld/esp32.ulp.ld produces ELF files which correspond to this format. This linker script produces binaries with load_addr == 0.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if load_addr is out of range
- ESP_ERR_INVALID_SIZE if program_size doesn’ t match (TEXT_OFFSET + TEXT_SIZE + DATA_SIZE)
- ESP_ERR_NOT_SUPPORTED if the magic number is incorrect

Parameters

- `load_addr`: address where the program should be loaded, expressed in 32-bit words
- `program_binary`: pointer to program binary
- `program_size`: size of the program binary

一旦上述程序加载到 RTC 内存后，应用程序即可启动此程序，并将入口点的地址传递给 `ulp_run` 函数:

```
ESP_ERROR_CHECK( ulp_run(&ulp_entry - RTC_SLOW_MEM) );
```

esp_err_t **ulp_run**(uint32_t *entry_point*)

Run the program loaded into RTC memory.

Return ESP_OK on success

Parameters

- `entry_point`: entry point, expressed in 32-bit words

上述生成的头文件 `${ULP_APP_NAME}.h` 声明了入口点符号。在 ULP 应用程序的汇编源代码中，此符号必须标记为 `.global`:

```
        .global entry
entry:
        /* code starts here */
```

4.17.7 ULP 程序流

ULP 协处理器由定时器启动，而调用 ulp_run 则可启动此定时器。定时器为 RTC_SLOW_CLK 的 Tick 事件计数（默认情况下，Tick 由内部 150 KHz 晶振器生成）。使用 SENS_ULP_CP_SLEEP_CYCx_REG 寄存器 (x = 0..4) 设置 Tick 数值。第一次启动 ULP 时，使用 SENS_ULP_CP_SLEEP_CYC0_REG 设置定时器 Tick 数值，之后，ULP 程序可以使用 sleep 指令来另外选择 SENS_ULP_CP_SLEEP_CYCx_REG 寄存器。

此应用程序可以调用 ulp_set_wakeup_period 函数来设置 ULP 定时器周期值 (SENS_ULP_CP_SLEEP_CYCx_REG, x = 0..4)。

esp_err_t **ulp_set_wakeup_period**(size_t *period_index*, uint32_t *period_us*)

Set one of ULP wakeup period values.

ULP coprocessor starts running the program when the wakeup timer counts up to a given value (called period). There are 5 period values which can be programmed into SENS_ULP_CP_SLEEP_CYCx_REG registers, x = 0..4. By default, wakeup timer will use the period set into SENS_ULP_CP_SLEEP_CYC0_REG, i.e. period number 0. ULP program code can use SLEEP instruction to select which of the SENS_ULP_CP_SLEEP_CYCx_REG should be used for subsequent wakeups.

However, please note that SLEEP instruction issued (from ULP program) while the system is in deep sleep mode does not have effect, and sleep cycle count 0 is used.

Note The ULP FSM requires two clock cycles to wakeup before being able to run the program. Then additional 16 cycles are reserved after wakeup waiting until the 8M clock is stable. The FSM also requires two more clock cycles to go to sleep after the program execution is halted. The minimum wakeup period that may be set up for the ULP is equal to the total number of cycles spent on the above internal tasks. For a default configuration of the ULP running at 150kHz it makes about 133us.

Return

- ESP_OK on success
- ESP_ERR_INVALID_ARG if period_index is out of range

Parameters

- period_index: wakeup period setting number (0 - 4)
- period_us: wakeup period, us

一旦定时器为所选的 SENS_ULP_CP_SLEEP_CYCx_REG 寄存器的 Tick 事件计数，ULP 协处理器就会启动，并调用 ulp_run 的入口点开始运行程序。

程序保持运行，直到遇到 halt 指令或非法指令。一旦程序停止，ULP 协处理器电源关闭，定时器再次启动。

如果想禁用定时器（有效防止 ULP 程序再次运行），请清除 RTC_CNTL_STATE0_REG 寄存器中的 RTC_CNTL_ULP_CP_SLP_TIMER_EN 位，可在 ULP 代码或主程序中进行以上操作。

4.18 ESP32-S2 中的单元测试

ESP-IDF 中附带了一个基于 Unity 的单元测试应用程序框架，且所有的单元测试用例分别保存在 ESP-IDF 仓库中每个组件的 test 子目录中。

4.18.1 添加常规测试用例

单元测试被添加在相应组件的 test 子目录中，测试用例写在 C 文件中，一个 C 文件可以包含多个测试用例。测试文件的名字要以“test”开头。

测试文件需要包含 unity.h 头文件，此外还需要包含待测试 C 模块需要的头文件。

测试用例需要通过 C 文件中特定的函数来添加，如下所示：

```
TEST_CASE("test name", "[module name]")
{
 // 在这里添加测试用例
}
```

- 第一个参数是此测试的描述性名称。
- 第二个参数用方括号中的标识符来表示，标识符用来对相关测试或具有特定属性的测试进行分组。

注解： 没有必要在每个测试用例中使用 UNITY_BEGIN() 和 UNITY_END() 来声明主函数的区域，unity_platform.c 会自动调用 UNITY_BEGIN()，然后运行测试用例，最后调用 UNITY_END()。

test 子目录应包含:ref: 组件 *CMakeLists.txt <component-directories>*，因为他们本身就是一种组件。ESP-IDF 使用了 unity 测试框架，需要将其指定为组件的依赖项。通常，组件:ref: 需要手动指定待编译的源文件 *<cmake-file-globbing>*；但是，对于测试组件来说，这个要求被放宽为仅建议将参数 *"SRC_DIRS"* 用于 idf_component_register。

总的来说，test 子目录下最小的 CMakeLists.txt 文件可能如下所示:

```
idf_component_register(SRC_DIRS "."
                       INCLUDE_DIRS "."
                       REQUIRES unity)
```

更多关于如何在 Unity 下编写测试用例的信息，请查阅 http://www.throwtheswitch.org/unity。

4.18.2 添加多设备测试用例

常规测试用例会在一个 DUT（Device Under Test，在试设备）上执行。但是，那些需要互相通信的组件（比如 GPIO、SPI）需要其他设备与其通信，因此不能使用常规测试用例进行测试。多设备测试用例包括写入多个测试函数，并在多个 DUT 进行运行测试。

以下是一个多设备测试用例：

```
void gpio_master_test()
{
    gpio_config_t slave_config = {
            .pin_bit_mask = 1 << MASTER_GPIO_PIN,
            .mode = GPIO_MODE_INPUT,
    };
    gpio_config(&slave_config);
    unity_wait_for_signal("output high level");
    TEST_ASSERT(gpio_get_level(MASTER_GPIO_PIN) == 1);
}

void gpio_slave_test()
{
    gpio_config_t master_config = {
            .pin_bit_mask = 1 << SLAVE_GPIO_PIN,
            .mode = GPIO_MODE_OUTPUT,
    };
    gpio_config(&master_config);
    gpio_set_level(SLAVE_GPIO_PIN, 1);
    unity_send_signal("output high level");
}

TEST_CASE_MULTIPLE_DEVICES("gpio multiple devices test example", "[driver]", gpio_
↪master_test, gpio_slave_test);
```

宏 TEST_CASE_MULTIPLE_DEVICES 用来声明多设备测试用例，

- 第一个参数指定测试用例的名字。

- 第二个参数是测试用例的描述。
- 从第三个参数开始，可以指定最多 5 个测试函数，每个函数都是单独运行在一个 DUT 上的测试入口点。

在不同的 DUT 上运行的测试用例，通常会要求它们之间进行同步。我们提供 `unity_wait_for_signal` 和 `unity_send_signal` 这两个函数来使用 UART 去支持同步操作。如上例中的场景，slave 应该在 master 设置好 GPIO 电平后再去读取 GPIO 电平，DUT 的 UART 终端会打印提示信息，并要求用户进行交互。

DUT1（master）终端:

```
Waiting for signal: [output high level]!
Please press "Enter" key once any board send this signal.
```

DUT2（slave）终端:

```
Send signal: [output high level]!
```

一旦 DUT2 发送了该信号，您需要在 DUT1 的终端按回车键，然后 DUT1 会从 `unity_wait_for_signal` 函数中解除阻塞，并开始更改 GPIO 的电平。

4.18.3 添加多阶段测试用例

常规的测试用例无需重启就会结束（或者仅需要检查是否发生了重启），可有些时候我们想在某些特定类型的重启事件后运行指定的测试代码，例如，我们想在深度睡眠唤醒后检查复位的原因是否正确。首先我们需要触发深度睡眠复位事件，然后检查复位的原因。为了实现这一点，我们可以定义多阶段测试用例来将这些测试函数组合在一起:

```
static void trigger_deepsleep(void)
{
    esp_sleep_enable_timer_wakeup(2000);
    esp_deep_sleep_start();
}

void check_deepsleep_reset_reason()
{
    RESET_REASON reason = rtc_get_reset_reason(0);
    TEST_ASSERT(reason == DEEPSLEEP_RESET);
}

TEST_CASE_MULTIPLE_STAGES("reset reason check for deepsleep", "[esp32s2]", trigger_
↪deepsleep, check_deepsleep_reset_reason);
```

多阶段测试用例向用户呈现了一组测试函数，它需要用户进行交互（选择用例并选择不同的阶段）来运行。

4.18.4 应用于不同芯片的单元测试

某些测试（尤其与硬件相关的）无法在所有的芯片上执行。请参照本节让你的单元测试只在其中一部分芯片上执行。

1. 使用宏 `!(TEMPORARY_)DISABLED_FOR_TARGETS()` 包装你的测试代码，并将其放于原始的测试文件中，或将代码分成按功能分组的文件。但请确保所有这些文件都会由编译器处理。例:

```
#if !TEMPORARY_DISABLED_FOR_TARGETS(ESP32, ESP8266)
TEST_CASE("a test that is not ready for esp32 and esp8266 yet", "[]")
{
}
#endif //!TEMPORARY_DISABLED_FOR_TARGETS(ESP32, ESP8266)
```

一旦你需要其中一个测试在某个芯片上被编译，只需要修改禁止的芯片列表。我们更鼓励使用一些通用的概念（能在 soc_caps.h 中被清楚描述）来禁止某些单元测试。如果你已经这样做，但有一些测试还没有在新的芯片版本中被调试通过，请同时使用上述两种方法，当调试完成后再移除 !(TEMPORARY_)DISABLED_FOR_TARGETS() 。例：

```
#if SOC_SDIO_SLAVE_SUPPORTED
#if !TEMPORARY_DISABLED_FOR_TARGETS(ESP64)
TEST_CASE("a sdio slave tests that is not ready for esp64 yet", "[sdio_slave]")
{
    //available for esp32 now, and will be available for esp64 in the future
}
#endif //!TEMPORARY_DISABLED_FOR_TARGETS(ESP64)
#endif //SOC_SDIO_SLAVE_SUPPORTED
```

2. 对于某些你确定不会支持的测试（例如，芯片根本没有该外设），使用 DISABLED_FOR_TARGETS 来禁止该测试；对于其他只是临时性需要关闭的（例如，没有 runner 资源等），使用 TEMPORARY_DISABLED_FOR_TARGETS 来暂时关闭该测试。

一些禁用目标芯片测试用例的旧方法，由于它们具有明显的缺陷，已经被废弃，请勿继续使用：

- 请勿将测试代码放在 test/芯片版本目录下面，然后用 CMakeLists.txt 来选择其中一个进行编译。这是因为测试代码比实现代码更容易被复用。如果你将一些代码放在 test/esp32 目录下来避免 esp32s2 芯片执行它，一旦你需要在新的芯片（比如 esp32s3 ）中启用该测试，你会发现这种结构非常难以保持代码的整洁。
- 请勿继续使用 CONFIG_IDF_TARGET_xxx 宏来禁止某些测试在一些芯片上编译。这种方法会让被禁止的测试项目难以追踪和重新打开。并且，相比于白名单式的 #if CONFIG_IDF_TARGET_xxx ，黑名单式的 #if !disabled 能避免新芯片引入时，这些测试被自动关闭。但对于用于测试的一些实现，#if CONFIG_IDF_TARGET_xxx 仍可用于给不同芯片版本选择实现代码。测试项目和测试实现区分如下：
  - 测试项目：某些你会在一些芯片上执行，而在另外一些上跳过的项目，例如：
    有三个测试项目 SD 1-bit 、SD 4-bit 和 SDSPI 。对于不支持 SD Host 外设的 ESP32-S2 芯片，只有 SDSPI 一个项目需要被执行。
  - 测试实现：某些代码始终会发生，但采取不同的做法。例如：
    ESP8266 芯片没有 SDIO_PKT_LEN 寄存器。如果在测试过程中需要获取从设备的数据长度，你可以用不同方式读取的 #if CONFIG_IDF_TARGET_ 宏来保护不同的实现代码。
    但请注意避免使用 #else 宏。这样当新芯片被引入时，测试就会在编译阶段失败，提示维护者去显示选择一个正确的测试实现。

4.18.5 编译单元测试程序

按照 esp-idf 顶层目录的 README 文件中的说明进行操作，请确保 IDF_PATH 环境变量已经被设置指向了 esp-idf 的顶层目录。

切换到 tools/unit-test-app 目录下进行配置和编译：

- idf.py menuconfig - 配置单元测试程序。
- idf.py -T all build - 编译单元测试程序，测试每个组件 test 子目录下的用例。
- idf.py -T "xxx yyy" build - 编译单元测试程序，测试指定的组件。(如 idf.py -T heap build - 仅对 heap 组件目录下的单元测试程序进行编译)
- idf.py -T all -E "xxx yyy" build - 编译单元测试程序，测试除指定组件之外的所有组件。(例如 idf.py -T all -E "ulp mbedtls" build - 编译所有的单元测试，不包括 ulp 和 mbedtls 组件。)

当编译完成时，它会打印出烧写芯片的指令。您只需要运行 idf.py flash 即可烧写所有编译输出的文件。

您还可以运行 idf.py -T all flash 或者 idf.py -T xxx flash 来编译并烧写，所有需要的文件都会在烧写之前自动重新编译。

使用 menuconfig 可以设置烧写测试程序所使用的串口。

4.18.6 运行单元测试

烧写完成后重启 ESP32-S2，它将启动单元测试程序。

当单元测试应用程序空闲时，输入回车键，它会打印出测试菜单，其中包含所有的测试项目:

```
Here's the test menu, pick your combo:
(1)     "esp_ota_begin() verifies arguments" [ota]
(2)     "esp_ota_get_next_update_partition logic" [ota]
(3)     "Verify bootloader image in flash" [bootloader_support]
(4)     "Verify unit test app image" [bootloader_support]
(5)     "can use new and delete" [cxx]
(6)     "can call virtual functions" [cxx]
(7)     "can use static initializers for non-POD types" [cxx]
(8)     "can use std::vector" [cxx]
(9)     "static initialization guards work as expected" [cxx]
(10)    "global initializers run in the correct order" [cxx]
(11)    "before scheduler has started, static initializers work correctly" [cxx]
(12)    "adc2 work with wifi" [adc]
(13)    "gpio master/slave test example" [ignore][misc][test_env=UT_T2_1][multi_
↪device]
        (1)     "gpio_master_test"
        (2)     "gpio_slave_test"
(14)    "SPI Master clockdiv calculation routines" [spi]
(15)    "SPI Master test" [spi][ignore]
(16)    "SPI Master test, interaction of multiple devs" [spi][ignore]
(17)    "SPI Master no response when switch from host1 (HSPI) to host2 (VSPI)"␣
↪[spi]
(18)    "SPI Master DMA test, TX and RX in different regions" [spi]
(19)    "SPI Master DMA test: length, start, not aligned" [spi]
(20)    "reset reason check for deepsleep" [esp32s2][test_env=UT_T2_1][multi_stage]
        (1)     "trigger_deepsleep"
        (2)     "check_deepsleep_reset_reason"
```

常规测试用例会打印用例名字和描述，主从测试用例还会打印子菜单（已注册的测试函数的名字）。

可以输入以下任意一项来运行测试用例:

- 引号中写入测试用例的名字，运行单个测试用例。
- 测试用例的序号，运行单个测试用例。
- 方括号中的模块名字，运行指定模块所有的测试用例。
- 星号，运行所有测试用例。

[multi_device] 和 [multi_stage]``标签告诉测试运行者该用例是多设备测试还是多阶段测试。这些标签由 ``TEST_CASE_MULTIPLE_STAGES 和 TEST_CASE_MULTIPLE_DEVICES 宏自动生成。

一旦选择了多设备测试用例，它会打印一个子菜单:

```
Running gpio master/slave test example...
gpio master/slave test example
        (1)     "gpio_master_test"
        (2)     "gpio_slave_test"
```

您需要输入数字以选择在 DUT 上运行的测试。

与多设备测试用例相似，多阶段测试用例也会打印子菜单:

```
Running reset reason check for deepsleep...
reset reason check for deepsleep
        (1)     "trigger_deepsleep"
        (2)     "check_deepsleep_reset_reason"
```

第一次执行此用例时，输入 1 来运行第一阶段（触发深度睡眠）。在重启 DUT 并再次选择运行此用例后，输入 2 来运行第二阶段。只有在最后一个阶段通过并且之前所有的阶段都成功触发了复位的情况下，该测试才算通过。

4.18.7 带缓存补偿定时器的定时代码

存储在外部存储器（如 SPI Flash 和 SPI RAM）中的指令和数据是通过 CPU 的统一指令和数据缓存来访问的。当代码或数据在缓存中时，访问速度会非常快（即缓存命中）。

然而，如果指令或数据不在缓存中，则需要从外部内存中获取（即缓存缺失）。访问外部存储器的速度明显较慢，因为 CPU 在等待从外部存储器获取指令或数据时会陷入停滞。这导致整体代码执行速度会依据缓存命中或缓存缺失的次数而变化。

在不同的编译中，代码和数据的位置可能会有所不同，一些可能会更有利于缓存访问（即，最大限度地减少缓存缺失）。理论上说这会影响执行速度，但这些因素通常却是无关紧要，因为它们的影响会在设备的运行过程中“平均化”。

然而，高速缓存对执行速度的影响可能与基准测试场景（尤其是微基准测试）有关。每次运行和构建时的测量时间可能会有所差异，消除部分差异的方法之一是将代码和数据分别放在指令或数据 RAM（IRAM/DRAM）中。CPU 可以直接访问 IRAM 和 DRAM，从而消除了高速缓存的影响因素。然而，由于 IRAM 和 DRAM 容量有限，该方法并不总是可行。

缓存补偿定时器是将要基准测试的代码/数据放置在 IRAM/DRAM 中的替代方法，该计时器使用处理器的内部事件计数器来确定在发生高速缓存未命中时等待代码/数据所花费的时间，然后从记录的实时时间中减去该时间。

```
// Start the timer
ccomp_timer_start();

// Function to time
func_code_to_time();

// Stop the timer, and return the elapsed time in microseconds relative to
// ccomp_timer_start
int64_t t = ccomp_timer_stop();
```

缓存补偿定时器的限制之一是基准功能必须固定在一个内核上。这是由于每个内核都有自己的事件计数器，这些事件计数器彼此独立。例如，如果在一个内核上调用 `ccomp_timer_start`，使调度器进入睡眠状态，唤醒并在在另一个内核上重新调度，那么对应的 `ccomp_timer_stop` 将无效。

4.19 ESP32-S2 ROM console

When an ESP32-S2 is unable to boot from flash ROM (and the fuse disabling it hasn ' t been blown), it boots into a rom console. The console is based on TinyBasic, and statements entered should be in the form of BASIC statements. As is common in the BASIC language, without a preceeding line number, commands entered are executed immediately; lines with a prefixed line number are stored as part of a program.

4.19.1 Full list of supported statements and functions

System

- BYE - *exits Basic, reboots and retries booting from flash*
- END - *stops execution from the program, also “STOP”*
- MEM - *displays memory usage statistics*
- NEW - *clears the current program*
- RUN - *executes the current program*

IO, Documentation

- PEEK(address) - *get a 32-bit value from a memory address*
- POKE - *write a 32-bit value to memory*
- USR(addr, arg1, ..) - *Execute a machine language function*

- PRINT expression - *print out the expression, also "?"*
- PHEX expression - *print expression as a hex number*
- REM stuff - *remark/comment, also "'"*

Expressions, Math

- A=V, LET A=V - *assign value to a variable*
- +, -, *, / - *Math*
- <,<=,=,<>,!=,>=,> - *Comparisons*
- ABS(expression) - *returns the absolute value of the expression*
- RSEED(v) - *sets the random seed to v*
- RND(m) - *returns a random number from 0 to m*
- A=1234 - * Assign a decimal value*
- A=&h1A2 - * Assign a hex value*
- A=&b1001 - *Assign a binary value*

Control

- IF expression statement - *perform statement if expression is true*
- FOR variable = start TO end - *start for block*
- FOR variable = start TO end STEP value - *start for block with step*
- NEXT - *end of for block*
- GOTO linenumber - *continue execution at this line number*
- GOSUB linenumber - *call a subroutine at this line number*
- RETURN - *return from a subroutine*
- DELAY - *Delay a given number of milliseconds*

Pin IO

- IODIR - *Set a GPIO-pin as an output (1) or input (0)*
- IOSET - *Set a GPIO-pin, configured as output, to high (1) or low (0)*
- IOGET - *Get the value of a GPIO-pin*

4.19.2 Example programs

Here are a few example commands and programs to get you started…

Read UART_DATE register of uart0

```
> PHEX PEEK(&h3FF40078)
15122500
```

Set GPIO2 using memory writes to GPIO_OUT_REG

Note: you can do this easier with the IOSET command

```
> POKE &h3FF44004,PEEK(&h3FF44004) OR &b100
```

Get value of GPIO0

```
> IODIR 0,0
> PRINT IOGET(0)
0
```

Blink LED

Hook up an LED between GPIO2 and ground. When running the program, the LED should blink 10 times.

```
10 IODIR 2,1
20 FOR A=1 TO 10
30 IOSET 2,1
40 DELAY 250
50 IOSET 2,0
60 DELAY 250
70 NEXT A
RUN
```

4.19.3 Credits

The ROM console is based on "TinyBasicPlus" by Mike Field and Scott Lawrence, which is based on "68000 TinyBasic" by Gordon Brandly

4.20 RF calibration

ESP32-S2 supports three RF calibration methods during RF initialization:

1. Partial calibration
2. Full calibration
3. No calibration

4.20.1 Partial calibration

During RF initialization, the partial calibration method is used by default for RF calibration. It is done based on the full calibration data which is stored in the NVS. To use this method, please go to `menuconfig` and enable *CONFIG_ESP32_PHY_CALIBRATION_AND_DATA_STORAGE*.

4.20.2 Full calibration

Full calibration is triggered in the follwing conditions:

1. NVS does not exist.
2. The NVS partition to store calibration data is erased.
3. Hardware MAC address is changed.
4. PHY library version is changed.
5. The RF calibration data loaded from the NVS partition is broken.

It takes about 100ms more than partial calibration. If boot duration is not critical, it is suggested to use the full calibration method. To switch to the full calibration method, go to `menuconfig` and disable *CONFIG_ESP32_PHY_CALIBRATION_AND_DATA_STORAGE*. If you use the default method of RF calibration, there are two ways to add the function of triggering full calibration as a last-resort remedy.

1. Erase the NVS partition if you don't mind all of the data stored in the NVS partition is erased. That is indeed the easiest way.

2. Call API `esp_phy_erase_cal_data_in_nvs()` before initializing WiFi and BT/BLE based on some conditions (e.g. an option provided in some diagnostic mode). In this case, only phy namespace of the NVS partition is erased.

4.20.3 No calibration

No calibration method is only used when the device wakes up from deep sleep.

4.20.4 PHY initialization data

The PHY initialization data is used for RF calibration. There are two ways to get the PHY initialization data.

One is the default initialization data which is located in the header file components/esp_wifi/esp32/include/phy_init_data.h.

It is embedded into the application binary after compiling and then stored into read-only memory (DROM). To use the default initialization data, please go to `menuconfig` and disable *CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION*.

Another is the initialization data which is stored in a partition. When using a custom partition table, make sure that PHY data partition is included (type: *data*, subtype: *phy*). With default partition table, this is done automatically. If initialization data is stored in a partition, it has to be flashed there, otherwise runtime error will occur. To switch to the initialization data stored in a partition, go to `menuconfig` and enable *CONFIG_ESP32_PHY_INIT_DATA_IN_PARTITION*.

4.21 Wi-Fi Driver

4.21.1 ESP32-S2 Wi-Fi Feature List

- Support 4 virtual WiFi interfaces, which are STA, AP, Sniffer and reserved.
- Support station-only mode, AP-only mode, station/AP-coexistence mode
- Support IEEE 802.11b, IEEE 802.11g, IEEE 802.11n, and APIs to configure the protocol mode
- Support WPA/WPA2/WPA2-Enterprise and WPS
- Support AMPDU, HT40, QoS, and other key features
- Support Modem-sleep
- Support an Espressif-specific protocol which, in turn, supports up to **1 km** of data traffic
- Up to 20 MBit/sec TCP throughput and 30 MBit/sec UDP throughput over the air
- Support Sniffer
- Support set fast_crypto algorithm and normal algorithm switch which used in wifi connect
- Support both fast scan and all channel scan feature
- Support multiple antennas
- Support channel state information

4.21.2 How To Write a Wi-Fi Application

Preparation

Generally, the most effective way to begin your own Wi-Fi application is to select an example which is similar to your own application, and port the useful part into your project. It is not a MUST but it is strongly recommended that you take some time to read this article first, especially if you want to program a robust Wi-Fi application. This article is supplementary to the Wi-Fi APIs/Examples. It describes the principles of using the Wi-Fi APIs, the limitations of the current Wi-Fi API implementation, and the most common pitfalls in using Wi-Fi. This article also reveals some design details of the Wi-Fi driver. We recommend that you become familiar at least with the following sections: <*ESP32-S2 Wi-Fi API Error Code*>, <*ESP32-S2 Wi-Fi Programming Model*>, and <*ESP32-S2 Wi-Fi Event Description*>.

Setting Wi-Fi Compile-time Options

Refer to <*Wi-Fi Menuconfig*>

Init Wi-Fi

Refer to <*ESP32-S2 Wi-Fi Station General Scenario*>, <*ESP32-S2 Wi-Fi AP General Scenario*>.

Start/Connect Wi-Fi

Refer to <*ESP32-S2 Wi-Fi Station General Scenario*>, <*ESP32-S2 Wi-Fi AP General Scenario*>.

Event-Handling

Generally, it is easy to write code in "sunny-day" scenarios, such as <*WIFI_EVENT_STA_START*>, <*WIFI_EVENT_STA_CONNECTED*> etc. The hard part is to write routines in "rainy-day" scenarios, such as <*WIFI_EVENT_STA_DISCONNECTED*> etc. Good handling of "rainy-day" scenarios is fundamental to robust Wi-Fi applications. Refer to <*ESP32-S2 Wi-Fi Event Description*>, <*ESP32-S2 Wi-Fi Station General Scenario*>, <*ESP32-S2 Wi-Fi AP General Scenario*>. See also *an overview of event handling in ESP-IDF*.

Write Error-Recovery Routines Correctly at All Times

Just like the handling of "rainy-day" scenarios, a good error-recovery routine is also fundamental to robust Wi-Fi applications. Refer to <*ESP32-S2 Wi-Fi API Error Code*>

4.21.3 ESP32-S2 Wi-Fi API Error Code

All of the ESP32-S2 Wi-Fi APIs have well-defined return values, namely, the error code. The error code can be categorized i

- No errors, e.g. ESP_OK means that the API returns successfully
- Recoverable errors, such as ESP_ERR_NO_MEM, etc.
- Non-recoverable, non-critical errors
- Non-recoverable, critical errors

Whether the error is critical or not depends on the API and the application scenario, and it is defined by the API user.

The primary principle to write a robust application with Wi-Fi API is to always check the error code and write the error-handling code. Generally, the error-handling code can be used:

- For recoverable errors, in which case you can write a recoverable-error code. For example, when `esp_wifi_start()` returns ESP_ERR_NO_MEM, the recoverable-error code vTaskDelay can be called in order to get a microseconds' delay for another try.
- For non-recoverable, yet non-critical errors, in which case printing the error code is a good method for error handling.
- For non-recoverable and also critical errors, in which case "assert" may be a good method for error handling. For example, if `esp_wifi_set_mode()` returns ESP_ERR_WIFI_NOT_INIT, it means that the Wi-Fi driver is not initialized by `esp_wifi_init()` successfully. You can detect this kind of error very quickly in the application development phase.

In esp_err.h, ESP_ERROR_CHECK checks the return values. It is a rather commonplace error-handling code and can be used as the default error-handling code in the application development phase. However, we strongly recommend that API users write their own error-handling code.

4.21.4 ESP32-S2 Wi-Fi API Parameter Initialization

When initializing struct parameters for the API, one of two approaches should be followed: - explicitly set all fields of the parameter or - use get API to get current configuration first, then set application specific fields

Initializing or getting the entire structure is very important because most of the time the value 0 indicates the default value is used. More fields may be added to the struct in the future and initializing these to zero ensures the application will still work correctly after IDF is updated to a new release.

4.21.5 ESP32-S2 Wi-Fi Programming Model

The ESP32-S2 Wi-Fi programming model is depicted as follows:

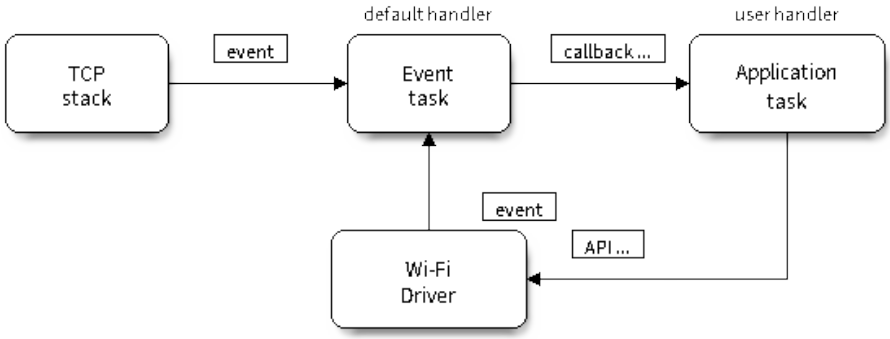


图 25: Wi-Fi Programming Model

The Wi-Fi driver can be considered a black box that knows nothing about high-layer code, such as the TCP/IP stack, application task, and event task. The application task (code) generally calls *Wi-Fi driver APIs* to initialize Wi-Fi and handles Wi-Fi events when necessary. Wi-Fi driver receives API calls, handles them, and posts events to the application.

Wi-Fi event handling is based on the *esp_event library*. Events are sent by the Wi-Fi driver to the *default event loop*. Application may handle these events in callbacks registered using `esp_event_handler_register()`. Wi-Fi events are also handled by *esp_netif component* to provide a set of default behaviors. For example, when Wi-Fi station connects to an AP, esp_netif will automatically start the DHCP client by default.

4.21.6 ESP32-S2 Wi-Fi Event Description

WIFI_EVENT_WIFI_READY

The Wi-Fi driver will never generate this event, which, as a result, can be ignored by the application event callback. This event may be removed in future releases.

WIFI_EVENT_SCAN_DONE

The scan-done event is triggered by esp_wifi_scan_start() and will arise in the following scenarios:

- The scan is completed, e.g., the target AP is found successfully, or all channels have been scanned.
- The scan is stopped by esp_wifi_scan_stop().
- The esp_wifi_scan_start() is called before the scan is completed. A new scan will override the current scan and a scan-done event will be generated.

The scan-done event will not arise in the following scenarios:

- It is a blocked scan.
- The scan is caused by esp_wifi_connect().

Upon receiving this event, the event task does nothing. The application event callback needs to call esp_wifi_scan_get_ap_num() and esp_wifi_scan_get_ap_records() to fetch the scanned AP list and trigger the Wi-Fi driver to free the internal memory which is allocated during the scan **(do not forget to do this)**! Refer to 'ESP32-S2 Wi-Fi Scan' for a more detailed description.

WIFI_EVENT_STA_START

If esp_wifi_start() returns ESP_OK and the current Wi-Fi mode is Station or AP+Station, then this event will arise. Upon receiving this event, the event task will initialize the LwIP network interface (netif). Generally, the application event callback needs to call esp_wifi_connect() to connect to the configured AP.

WIFI_EVENT_STA_STOP

If esp_wifi_stop() returns ESP_OK and the current Wi-Fi mode is Station or AP+Station, then this event will arise. Upon receiving this event, the event task will release the station's IP address, stop the DHCP client, remove TCP/UDP-related connections and clear the LwIP station netif, etc. The application event callback generally does not need to do anything.

WIFI_EVENT_STA_CONNECTED

If esp_wifi_connect() returns ESP_OK and the station successfully connects to the target AP, the connection event will arise. Upon receiving this event, the event task starts the DHCP client and begins the DHCP process of getting the IP address. Then, the Wi-Fi driver is ready for sending and receiving data. This moment is good for beginning the application work, provided that the application does not depend on LwIP, namely the IP address. However, if the application is LwIP-based, then you need to wait until the *got ip* event comes in.

WIFI_EVENT_STA_DISCONNECTED

This event can be generated in the following scenarios:

- When `esp_wifi_disconnect()`, or `esp_wifi_stop()`, or `esp_wifi_deinit()` is called and the station is already connected to the AP.
- When `esp_wifi_connect()` is called, but the Wi-Fi driver fails to set up a connection with the AP due to certain reasons, e.g. the scan fails to find the target AP, authentication times out, etc. If there are more than one AP with the same SSID, the disconnected event is raised after the station fails to connect all of the found APs.
- When the Wi-Fi connection is disrupted because of specific reasons, e.g., the station continuously loses N beacons, the AP kicks off the station, the AP's authentication mode is changed, etc.

Upon receiving this event, the default behavior of the event task is: - Shuts down the station's LwIP netif. - Notifies the LwIP task to clear the UDP/TCP connections which cause the wrong status to all sockets. For socket-based applications, the application callback can choose to close all sockets and re-create them, if necessary, upon receiving this event.

The most common event handle code for this event in application is to call esp_wifi_connect() to reconnect the Wi-Fi. However, if the event is raised because esp_wifi_disconnect() is called, the application should not call esp_wifi_connect() to reconnect. It's application's responsibility to distinguish whether the event is caused by esp_wifi_disconnect() or other reasons. Sometimes a better reconnect strategy is required, refer to <*Wi-Fi Reconnect*> and <*Scan When Wi-Fi Is Connecting*>.

Another thing deserves our attention is that the default behavior of LwIP is to abort all TCP socket connections on receiving the disconnect. Most of time it is not a problem. However, for some special application, this may not be what they want, consider following scenarios:

- The application creates a TCP connection to maintain the application-level keep-alive data that is sent out every 60 seconds.
- Due to certain reasons, the Wi-Fi connection is cut off, and the <*WIFI_EVENT_STA_DISCONNECTED*> is raised. According to the current implementation, all TCP connections will be removed and the keep-alive socket will be in a wrong status. However, since the application designer believes that the network layer should NOT care about this error at the Wi-Fi layer, the application does not close the socket.
- Five seconds later, the Wi-Fi connection is restored because esp_wifi_connect() is called in the application event callback function. **Moreover, the station connects to the same AP and gets the same IPV4 address as before**.
- Sixty seconds later, when the application sends out data with the keep-alive socket, the socket returns an error and the application closes the socket and re-creates it when necessary.

In above scenario, ideally, the application sockets and the network layer should not be affected, since the Wi-Fi connection only fails temporarily and recovers very quickly. The application can enable "Keep TCP connections when IP changed" via LwIP menuconfig.

IP_EVENT_STA_GOT_IP

This event arises when the DHCP client successfully gets the IPV4 address from the DHCP server, or when the IPV4 address is changed. The event means that everything is ready and the application can begin its tasks (e.g., creating sockets).

The IPV4 may be changed because of the following reasons:

- The DHCP client fails to renew/rebind the IPV4 address, and the station's IPV4 is reset to 0.
- The DHCP client rebinds to a different address.
- The static-configured IPV4 address is changed.

Whether the IPV4 address is changed or NOT is indicated by field `ip_change` of `ip_event_got_ip_t`.

The socket is based on the IPV4 address, which means that, if the IPV4 changes, all sockets relating to this IPV4 will become abnormal. Upon receiving this event, the application needs to close all sockets and recreate the application when the IPV4 changes to a valid one.

IP_EVENT_GOT_IP6

This event arises when the IPV6 SLAAC support auto-configures an address for the ESP32-S2, or when this address changes. The event means that everything is ready and the application can begin its tasks (e.g., creating sockets).

IP_STA_LOST_IP

This event arises when the IPV4 address become invalid.

IP_STA_LOST_IP doesn't arise immediately after the WiFi disconnects, instead it starts an IPV4 address lost timer, if the IPV4 address is got before ip lost timer expires, IP_EVENT_STA_LOST_IP doesn't happen. Otherwise, the event arises when IPV4 address lost timer expires.

Generally the application don't need to care about this event, it is just a debug event to let the application know that the IPV4 address is lost.

WIFI_EVENT_AP_START

Similar to <*WIFI_EVENT_STA_START*>.

WIFI_EVENT_AP_STOP

Similar to <*WIFI_EVENT_STA_STOP*>.

WIFI_EVENT_AP_STACONNECTED

Every time a station is connected to ESP32-S2 AP, the <*WIFI_EVENT_AP_STACONNECTED*> will arise. Upon receiving this event, the event task will do nothing, and the application callback can also ignore it. However, you may want to do something, for example, to get the info of the connected STA, etc.

WIFI_EVENT_AP_STADISCONNECTED

This event can happen in the following scenarios:

- The application calls `esp_wifi_disconnect()`, or `esp_wifi_deauth_sta()`, to manually disconnect the station.
- The Wi-Fi driver kicks off the station, e.g., because the AP has not received any packets in the past five minutes. The time can be modified by `esp_wifi_set_inactive_time()`.
- The station kicks off the AP.

When this event happens, the event task will do nothing, but the application event callback needs to do something, e.g., close the socket which is related to this station, etc.

WIFI_EVENT_AP_PROBEREQRECVED

This event is disabled by default. The application can enable it via API esp_wifi_set_event_mask(). When this event is enabled, it will be raised each time the AP receives a probe request.

4.21.7 ESP32-S2 Wi-Fi Station General Scenario

Below is a "big scenario" which describes some small scenarios in Station mode:

1. Wi-Fi/LwIP Init Phase

- s1.1: The main task calls `esp_netif_init()` to create an LwIP core task and initialize LwIP-related work.
- s1.2: The main task calls `esp_event_loop_init()` to create a system Event task and initialize an application event's callback function. In the scenario above, the application event's callback function does nothing but relaying the event to the application task.
- s1.3: The main task calls `esp_netif_create_default_wifi_ap()` or `esp_netif_create_default_wifi_sta()` to create default network interface instance binding station or AP with TCP/IP stack.
- s1.4: The main task calls `esp_wifi_init()` to create the Wi-Fi driver task and initialize the Wi-Fi driver.
- s1.5: The main task calls OS API to create the application task.

Step 1.1~1.5 is a recommended sequence that initializes a Wi-Fi-/LwIP-based application. However, it is **NOT** a must-follow sequence, which means that you can create the application task in step 1.1 and put all other initializations in the application task. Moreover, you may not want to create the application task in the initialization phase if the application task depends on the sockets. Rather, you can defer the task creation until the IP is obtained.

2. Wi-Fi Configuration Phase

Once the Wi-Fi driver is initialized, you can start configuring the Wi-Fi driver. In this scenario, the mode is Station, so you may need to call esp_wifi_set_mode(WIFI_MODE_STA) to configure the Wi-Fi mode as Station. You can call other esp_wifi_set_xxx APIs to configure more settings, such as the protocol mode, country code, bandwidth, etc. Refer to <*ESP32-S2 Wi-Fi Configuration*>.

Generally, the Wi-Fi driver should be configured before the Wi-Fi connection is set up. But this is **NOT** mandatory, which means that you can configure the Wi-Fi connection anytime, provided that the Wi-Fi driver is initialized successfully. However, if the configuration does not need to change after the Wi-Fi connection is set up, you should

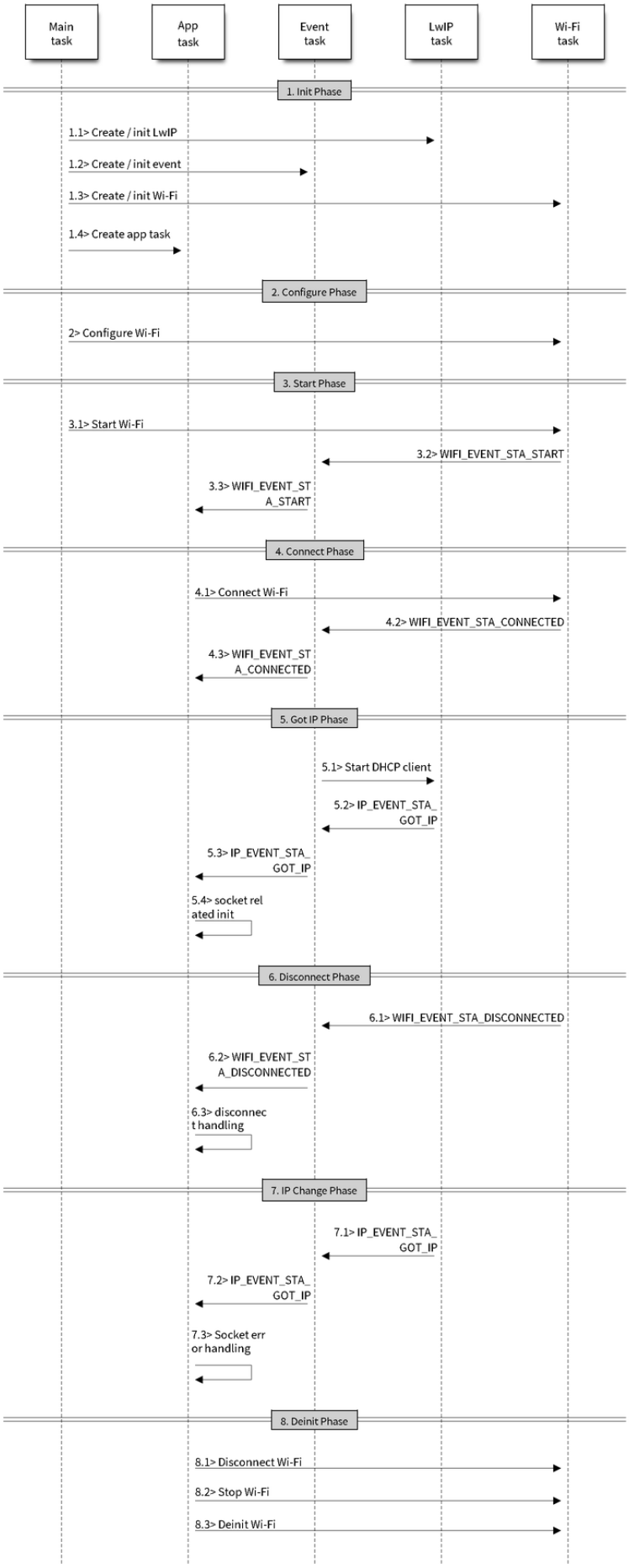


图 26: Sample Wi-Fi Event Scenarios in Station Mode

configure the Wi-Fi driver at this stage, because the configuration APIs (such as `esp_wifi_set_protocol()`) will cause the Wi-Fi to reconnect, which may not be desirable.

If the Wi-Fi NVS flash is enabled by menuconfig, all Wi-Fi configuration in this phase, or later phases, will be stored into flash. When the board powers on/reboots, you do not need to configure the Wi-Fi driver from scratch. You only need to call esp_wifi_get_xxx APIs to fetch the configuration stored in flash previously. You can also configure the Wi-Fi driver if the previous configuration is not what you want.

3. Wi-Fi Start Phase

- s3.1: Call esp_wifi_start to start the Wi-Fi driver.
- s3.2: The Wi-Fi driver posts <*WIFI_EVENT_STA_START*> to the event task; then, the event task will do some common things and will call the application event callback function.
- s3.3: The application event callback function relays the <*WIFI_EVENT_STA_START*> to the application task. We recommend that you call esp_wifi_connect(). However, you can also call esp_wifi_connect() in other phrases after the <*WIFI_EVENT_STA_START*> arises.

4. Wi-Fi Connect Phase

- s4.1: Once esp_wifi_connect() is called, the Wi-Fi driver will start the internal scan/connection process.
- s4.2: If the internal scan/connection process is successful, the <*WIFI_EVENT_STA_CONNECTED*> will be generated. In the event task, it starts the DHCP client, which will finally trigger the DHCP process.
- s4.3: In the above-mentioned scenario, the application event callback will relay the event to the application task. Generally, the application needs to do nothing, and you can do whatever you want, e.g., print a log, etc.

In step 4.2, the Wi-Fi connection may fail because, for example, the password is wrong, the AP is not found, etc. In a case like this, <*WIFI_EVENT_STA_DISCONNECTED*> will arise and the reason for such a failure will be provided. For handling events that disrupt Wi-Fi connection, please refer to phase 6.

5. Wi-Fi ‘Got IP’ Phase

- s5.1: Once the DHCP client is initialized in step 4.2, the *got IP* phase will begin.
- s5.2: If the IP address is successfully received from the DHCP server, then <*IP_EVENT_STA_GOT_IP*> will arise and the event task will perform common handling.
- s5.3: In the application event callback, <*IP_EVENT_STA_GOT_IP*> is relayed to the application task. For LwIP-based applications, this event is very special and means that everything is ready for the application to begin its tasks, e.g. creating the TCP/UDP socket, etc. A very common mistake is to initialize the socket before <*IP_EVENT_STA_GOT_IP*> is received. **DO NOT start the socket-related work before the IP is received.**

6. Wi-Fi Disconnect Phase

- s6.1: When the Wi-Fi connection is disrupted, e.g. because the AP is powered off, the RSSI is poor, etc., <*WIFI_EVENT_STA_DISCONNECTED*> will arise. This event may also arise in phase 3. Here, the event task will notify the LwIP task to clear/remove all UDP/TCP connections. Then, all application sockets will be in a wrong status. In other words, no socket can work properly when this event happens.
- s6.2: In the scenario described above, the application event callback function relays <*WIFI_EVENT_STA_DISCONNECTED*> to the application task. We recommend that esp_wifi_connect() be called to reconnect the Wi-Fi, close all sockets and re-create them if necessary. Refer to <*WIFI_EVENT_STA_DISCONNECTED*>.

7. Wi-Fi IP Change Phase

- s7.1: If the IP address is changed, the <*IP_EVENT_STA_GOT_IP*> will arise with “ip_change” set to true.

- s7.2: **This event is important to the application. When it occurs, the timing is good for closing all created sockets and recreating them.**

8. Wi-Fi Deinit Phase

- s8.1: Call esp_wifi_disconnect() to disconnect the Wi-Fi connectivity.
- s8.2: Call esp_wifi_stop() to stop the Wi-Fi driver.
- s8.3: Call esp_wifi_deinit() to unload the Wi-Fi driver.

4.21.8 ESP32-S2 Wi-Fi AP General Scenario

Below is a "big scenario" which describes some small scenarios in AP mode:

4.21.9 ESP32-S2 Wi-Fi Scan

Currently, the esp_wifi_scan_start() API is supported only in Station or Station+AP mode.

Scan Type

Mode	Description
Active Scan	Scan by sending a probe request. The default scan is an active scan.
Passive Scan	No probe request is sent out. Just switch to the specific channel and wait for a beacon. Application can enable it via the scan_type field of wifi_scan_config_t.
Foreground Scan	This scan is applicable when there is no Wi-Fi connection in Station mode. Foreground or background scanning is controlled by the Wi-Fi driver and cannot be configured by the application.
Background Scan	This scan is applicable when there is a Wi-Fi connection in Station mode or in Station+AP mode. Whether it is a foreground scan or background scan depends on the Wi-Fi driver and cannot be configured by the application.
All-Channel Scan	It scans all of the channels. If the channel field of wifi_scan_config_t is set to 0, it is an all-channel scan.
Specific Channel Scan	It scans specific channels only. If the channel field of wifi_scan_config_t set to 1, it is a specific-channel scan.

The scan modes in above table can be combined arbitrarily, so we totally have 8 different scans:

- All-Channel Background Active Scan
- All-Channel Background Passive Scan
- All-Channel Foreground Active Scan
- All-Channel Foreground Passive Scan
- Specific-Channel Background Active Scan
- Specific-Channel Background Passive Scan
- Specific-Channel Foreground Active Scan
- Specific-Channel Foreground Passive Scan

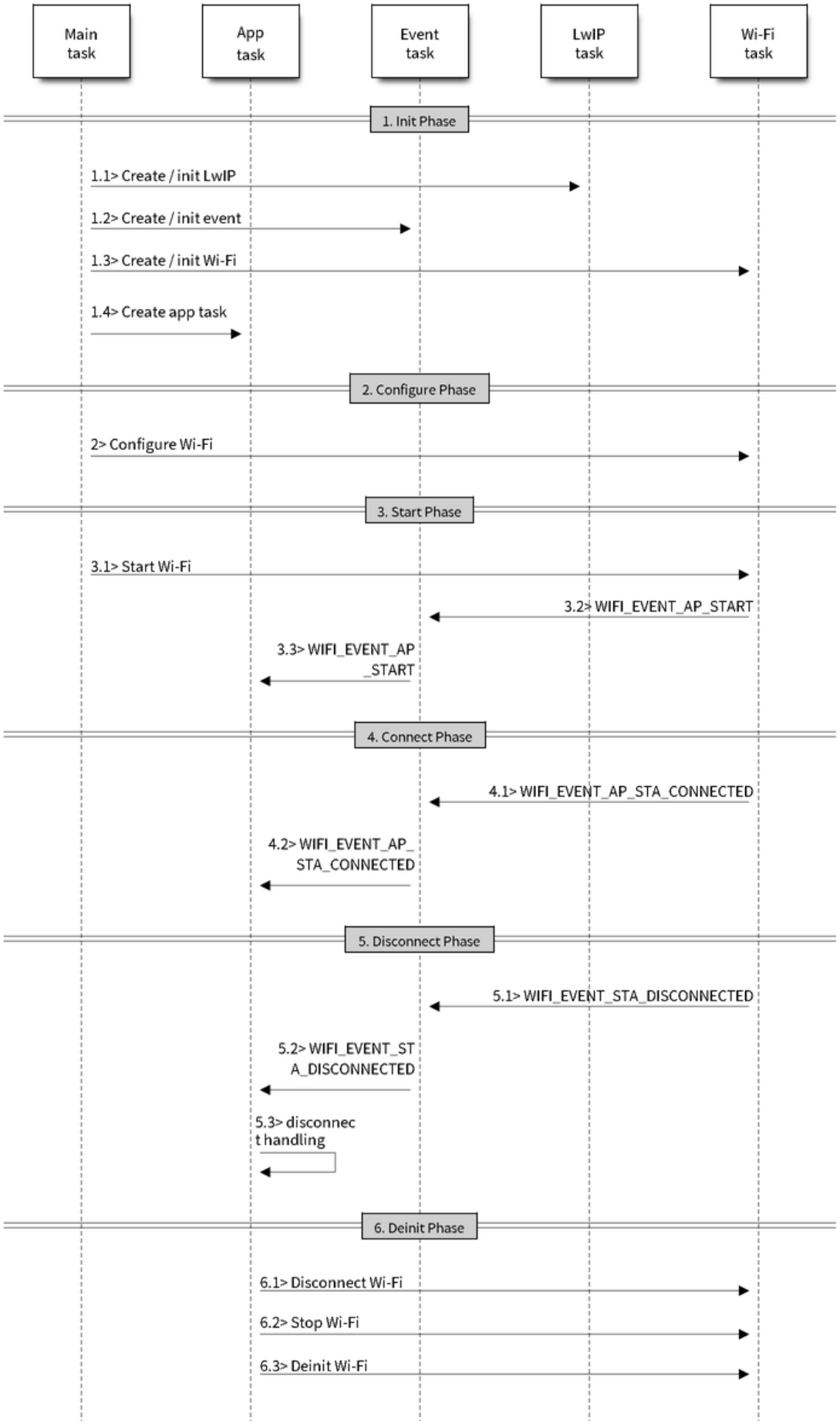


图 27: Sample Wi-Fi Event Scenarios in AP Mode

Scan Configuration

The scan type and other per-scan attributes are configured by esp_wifi_scan_start. The table below provides a detailed description of wifi_scan_config_t.

Field	Description
ssid	If the SSID is not NULL, it is only the AP with the same SSID that can be scanned.
bssid	If the BSSID is not NULL, it is only the AP with the same BSSID that can be scanned.
channel	If "channel" is 0, there will be an all-channel scan; otherwise, there will be a specific-channel scan.
show_hidden	If "show_hidden" is 0, the scan ignores the AP with a hidden SSID; otherwise, the scan considers the hidden AP a normal one.
scan_type	If "scan_type" is WIFI_SCAN_TYPE_ACTIVE, the scan is "active" ; otherwise, it is a "passive" one.
scan_time	This field is used to control how long the scan dwells on each channel. For passive scans, scan_time.passive designates the dwell time for each channel. For active scans, dwell times for each channel are listed in the table below. Here, min is short for scan time.active.min and max is short for scan_time.active.max. • min=0, max=0: scan dwells on each channel for 120 ms. • min>0, max=0: scan dwells on each channel for 120 ms. • min=0, max>0: scan dwells on each channel for `max` ms. • min>0, max>0: the minimum time the scan dwells on each channel is `min` ms. If no AP is found during this time frame, the scan switches to the next channel. Otherwise, the scan dwells on the channel for `max` ms. If you want to improve the performance of the the scan, you can try to modify these two parameters.

There also some global scan attributes which is configured by API esp_wifi_set_config, refer to *Station Basic Configuration*

Scan All APs In All Channels(foreground)

Scenario:

The scenario above describes an all-channel, foreground scan. The foreground scan can only occur in Station mode where the station does not connect to any AP. Whether it is a foreground or background scan is totally determined by the Wi-Fi driver, and cannot be configured by the application.

Detailed scenario description:

Scan Configuration Phase

- s1.1: Call esp_wifi_set_country() to set the country info if the default country info is not what you want, refer to <*Wi-Fi Country Code*>.

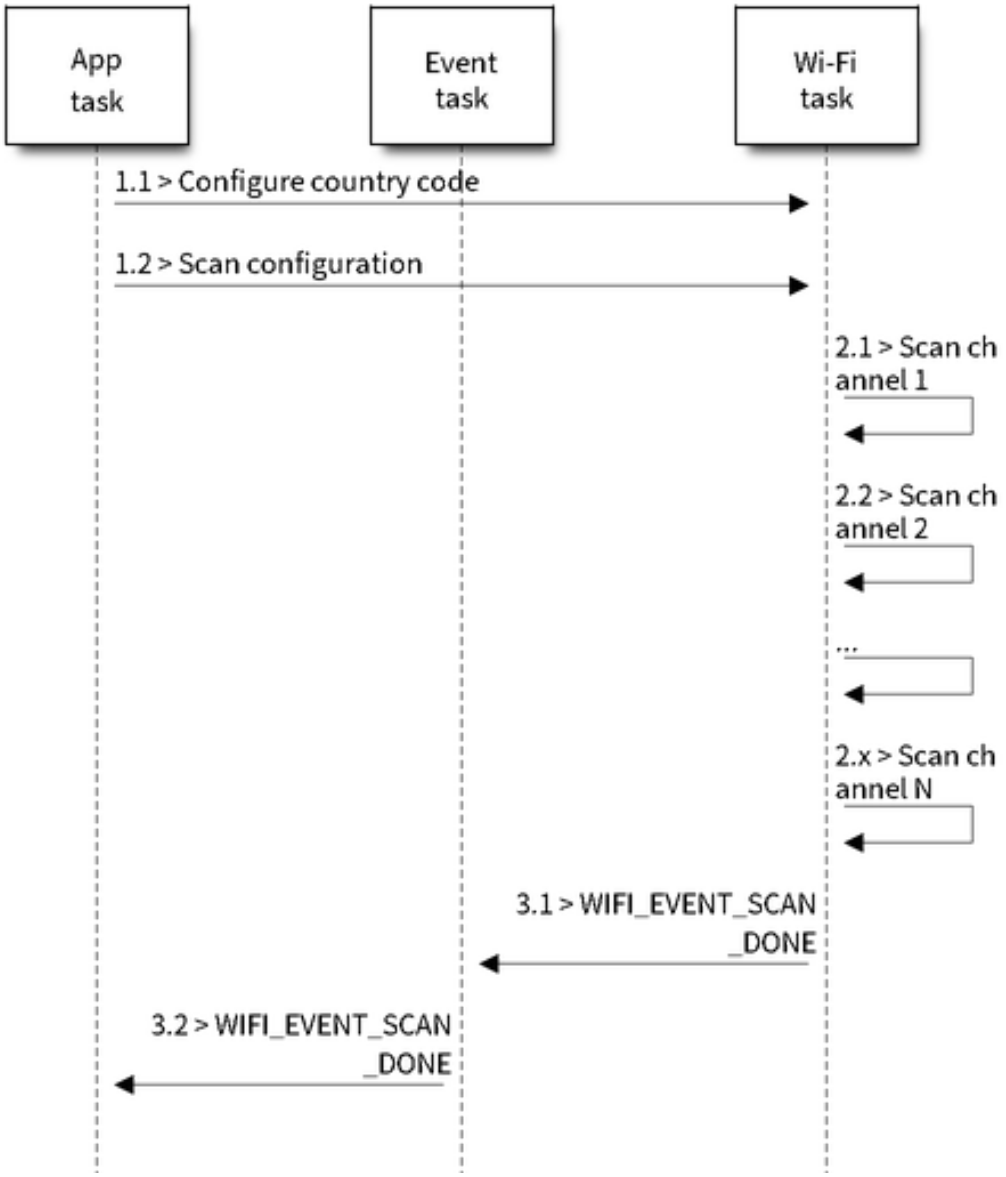


图 28: Foreground Scan of all Wi-Fi Channels

- s1.2: Call esp_wifi_scan_start() to configure the scan. To do so, you can refer to <*Scan Configuration*>. Since this is an all-channel scan, just set the SSID/BSSID/channel to 0.

Wi-Fi Driver’ s Internal Scan Phase

- s2.1: The Wi-Fi driver switches to channel 1, in case the scan type is WIFI_SCAN_TYPE_ACTIVE, and broadcasts a probe request. Otherwise, the Wi-Fi will wait for a beacon from the APs. The Wi-Fi driver will stay in channel 1 for some time. The dwell time is configured in min/max time, with default value being 120 ms.
- s2.2: The Wi-Fi driver switches to channel 2 and performs the same operation as in step 2.1.
- s2.3: The Wi-Fi driver scans the last channel N, where N is determined by the country code which is configured in step 1.1.

Scan-Done Event Handling Phase

- s3.1: When all channels are scanned, <*WIFI_EVENT_SCAN_DONE*> will arise.
- s3.2: The application’ s event callback function notifies the application task that <*WIFI_EVENT_SCAN_DONE*> is received. esp_wifi_scan_get_ap_num() is called to get the number of APs that have been found in this scan. Then, it allocates enough entries and calls esp_wifi_scan_get_ap_records() to get the AP records. Please note that the AP records in the Wi-Fi driver will be freed, once esp_wifi_scan_get_ap_records() is called. Do not call esp_wifi_scan_get_ap_records() twice for a single scan-done event. If esp_wifi_scan_get_ap_records() is not called when the scan-done event occurs, the AP records allocated by the Wi-Fi driver will not be freed. So, make sure you call esp_wifi_scan_get_ap_records(), yet only once.

Scan All APs on All Channels(background)

Scenario:

The scenario above is an all-channel background scan. Compared to *Scan All APs In All Channels(foreground)* , the difference in the all-channel background scan is that the Wi-Fi driver will scan the back-to-home channel for 30 ms before it switches to the next channel to give the Wi-Fi connection a chance to transmit/receive data.

Scan for a Specific AP in All Channels

Scenario:

This scan is similar to *Scan All APs In All Channels(foreground)*. The differences are:

- s1.1: In step 1.2, the target AP will be configured to SSID/BSSID.
- s2.1~s2.N: Each time the Wi-Fi driver scans an AP, it will check whether it is a target AP or not. If the scan is WIFI_FAST_SCAN scan and the target AP is found, then the scan-done event will arise and scanning will end; otherwise, the scan will continue. Please note that the first scanned channel may not be channel 1, because the Wi-Fi driver optimizes the scanning sequence.

If there are multiple APs which match the target AP info, for example, if we happen to scan two APs whose SSID is “ap” . If the scan is WIFI_FAST_SCAN, then only the first scanned “ap” will be found, if the scan is WIFI_ALL_CHANNEL_SCAN, both “ap” will be found and the station will connect the “ap” according to the configured strategy, refer to *Station Basic Configuration*.

You can scan a specific AP, or all of them, in any given channel. These two scenarios are very similar.

Scan in Wi-Fi Connect

When esp_wifi_connect() is called, then the Wi-Fi driver will try to scan the configured AP first. The scan in “Wi-Fi Connect” is the same as *Scan for a Specific AP In All Channels*, except that no scan-done event will be generated when the scan is completed. If the target AP is found, then the Wi-Fi driver will start the Wi-Fi connection; otherwise, <*WIFI_EVENT_STA_DISCONNECTED*> will be generated. Refer to *Scan for a Specific AP in All Channels*

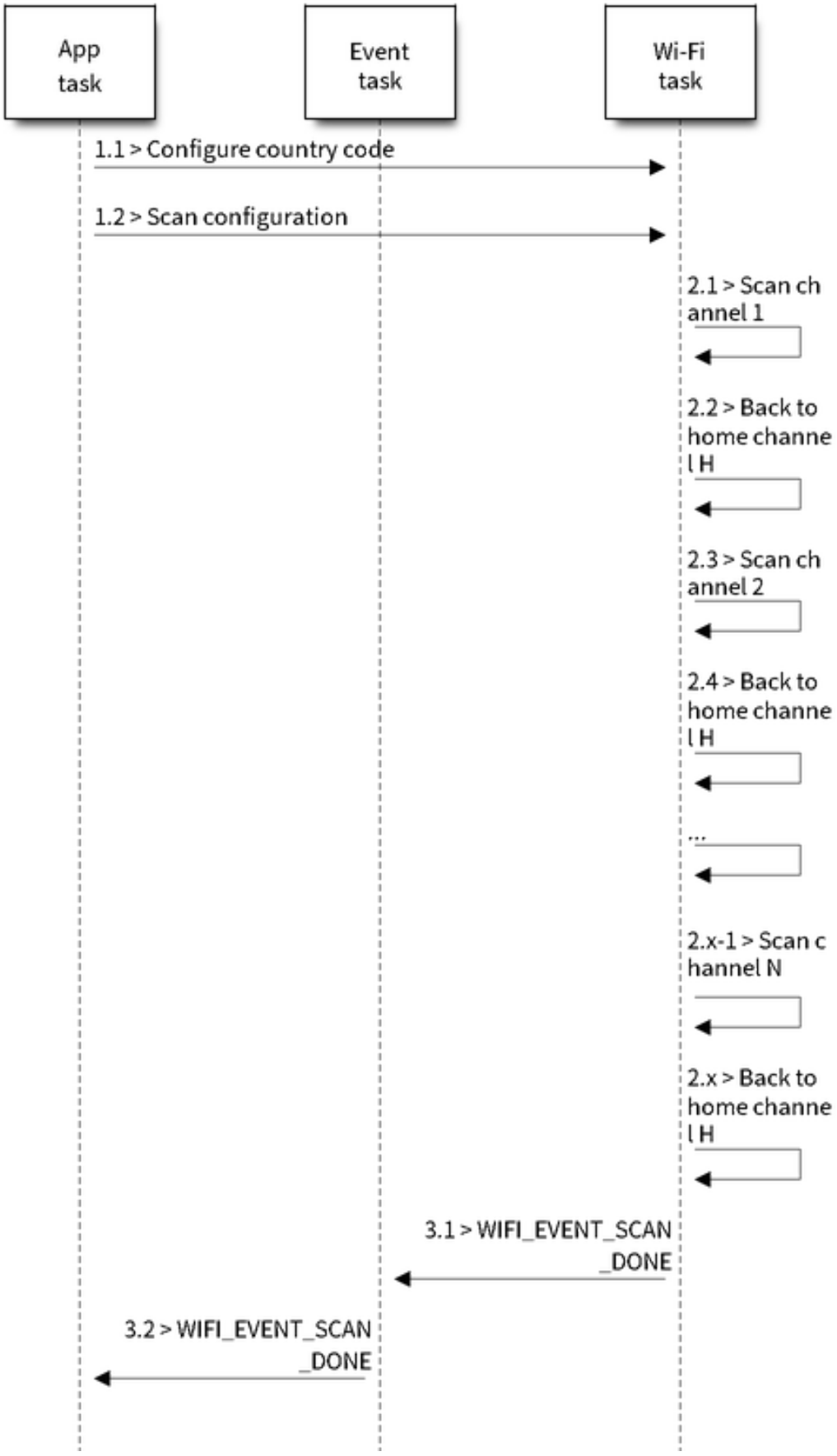


图 29: Background Scan of all Wi-Fi Channels

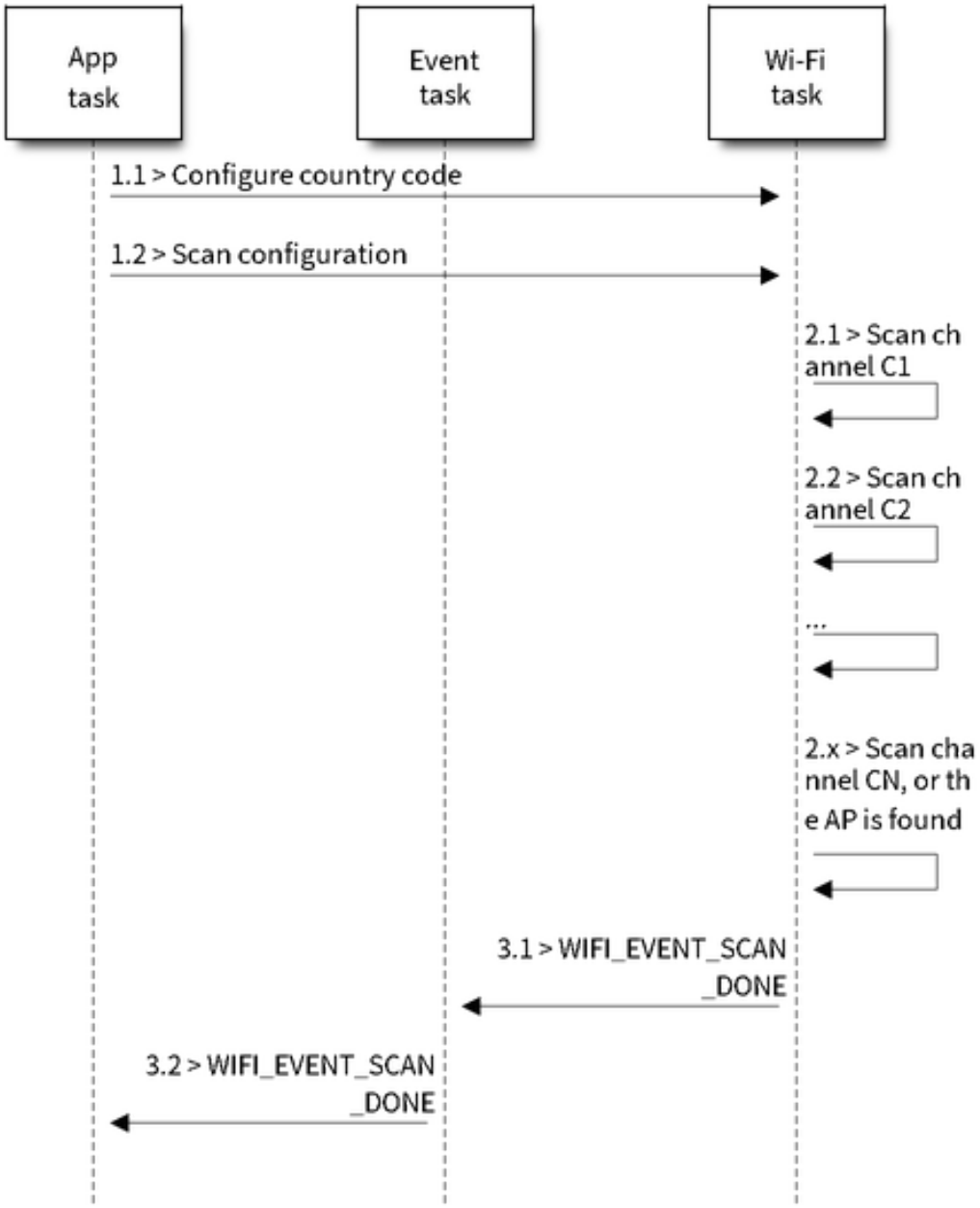


图 30: Scan of specific Wi-Fi Channels

Scan In Blocked Mode

If the block parameter of esp_wifi_scan_start() is true, then the scan is a blocked one, and the application task will be blocked until the scan is done. The blocked scan is similar to an unblocked one, except that no scan-done event will arise when the blocked scan is completed.

Parallel Scan

Two application tasks may call esp_wifi_scan_start() at the same time, or the same application task calls esp_wifi_scan_start() before it gets a scan-done event. Both scenarios can happen. **However, the Wi-Fi driver does not support multiple concurrent scans adequately. As a result, concurrent scans should be avoided.** Support for concurrent scan will be enhanced in future releases, as the ESP32-S2' s Wi-Fi functionality improves continuously.

Scan When Wi-Fi Is Connecting

The esp_wifi_scan_start() fails immediately if the Wi-Fi is in connecting process because the connecting has higher priority than the scan. If scan fails because of connecting, the recommended strategy is to delay sometime and retry scan again, the scan will succeed once the connecting is completed.

However, the retry/delay strategy may not work all the time. Considering following scenario: - The station is connecting a non-existed AP or if the station connects the existed AP with a wrong password, it always raises the event <*WIFI_EVENT_STA_DISCONNECTED*>. - The application call esp_wifi_connect() to do reconnection on receiving the disconnect event. - Another application task, e.g. the console task, call esp_wifi_scan_start() to do scan, the scan always fails immediately because the station is keeping connecting. - When scan fails, the application simply delay sometime and retry the scan.

In above scenario the scan will never succeed because the connecting is in process. So if the application supports similar scenario, it needs to implement a better reconnect strategy. E.g. - The application can choose to define a maximum continuous reconnect counter, stop reconnect once the reconnect reaches the max counter. - The application can choose to do reconnect immediately in the first N continous reconnect, then give a delay sometime and reconnect again.

The application can define its own reconnect strategy to avoid the scan starve to death. Refer to <*Wi-Fi Reconnect*>.

4.21.10 ESP32-S2 Wi-Fi Station Connecting Scenario

This scenario only depicts the case when there is only one target AP are found in scan phase, for the scenario that more than one AP with the same SSID are found, refer to <*ESP32-S2 Wi-Fi Station Connecting When Multiple APs Are Found*>.

Generally, the application does not need to care about the connecting process. Below is a brief introduction to the process for those who are really interested.

Scenario:

Scan Phase

- s1.1, The Wi-Fi driver begins scanning in "Wi-Fi Connect". Refer to <*Scan in Wi-Fi Connect*> for more details.
- s1.2, If the scan fails to find the target AP, <*WIFI_EVENT_STA_DISCONNECTED*> will arise and the reason-code will be WIFI_REASON_NO_AP_FOUND. Refer to <*Wi-Fi Reason Code*>.

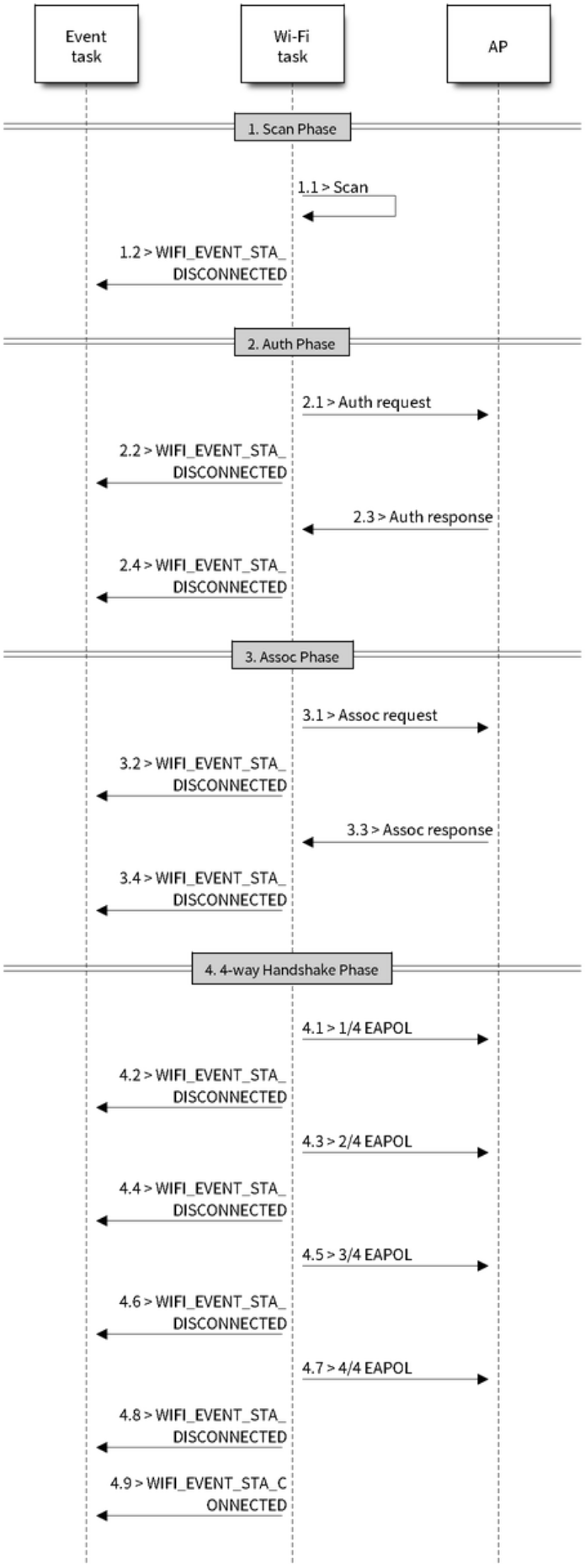


图 31: Wi-Fi Station Connecting Process

Auth Phase

- s2.1, The authentication request packet is sent and the auth timer is enabled.
- s2.2, If the authentication response packet is not received before the authentication timer times out, <*WIFI_EVENT_STA_DISCONNECTED*> will arise and the reason-code will be WIFI_REASON_AUTH_EXPIRE. Refer to <*Wi-Fi Reason Code*>.
- s2.3, The auth-response packet is received and the auth-timer is stopped.
- s2.4, The AP rejects authentication in the response and <*WIFI_EVENT_STA_DISCONNECTED*> arises, while the reason-code is WIFI_REASON_AUTH_FAIL or the reasons specified by the AP. Refer to <*Wi-Fi Reason Code*>.

Association Phase

- s3.1, The association request is sent and the association timer is enabled.
- s3.2, If the association response is not received before the association timer times out, <*WIFI_EVENT_STA_DISCONNECTED*> will arise and the reason-code will be WIFI_REASON_ASSOC_EXPIRE. Refer to <*Wi-Fi Reason Code*>.
- s3.3, The association response is received and the association timer is stopped.
- s3.4, The AP rejects the association in the response and <*WIFI_EVENT_STA_DISCONNECTED*> arises, while the reason-code is the one specified in the association response. Refer to <*Wi-Fi Reason Code*>.

Four-way Handshake Phase

- s4.1, The four-way handshake is sent out and the association timer is enabled.
- s4.2, If the association response is not received before the association timer times out, <*WIFI_EVENT_STA_DISCONNECTED*> will arise and the reason-code will be WIFI_REASON_ASSOC_EXPIRE. Refer to <*Wi-Fi Reason Code*>.
- s4.3, The association response is received and the association timer is stopped.
- s4.4, The AP rejects the association in the response and <*WIFI_EVENT_STA_DISCONNECTED*> arises and the reason-code will be the one specified in the association response. Refer to <*Wi-Fi Reason Code*>.

Wi-Fi Reason Code

The table below shows the reason-code defined in ESP32-S2. The first column is the macro name defined in esp_wifi_types.h. The common prefix *WIFI_REASON* is removed, which means that *UNSPECIFIED* actually stands for *WIFI_REASON_UNSPECIFIED* and so on. The second column is the value of the reason. The third column is the standard value to which this reason is mapped in section 8.4.1.7 of ieee802.11-2012. (For more information, refer to the standard mentioned above.) The last column is a description of the reason.

Reason code	Value	Mapped To	Description
UNSPECIFIED	1	1	Generally, it means an internal failure, e.g., the memory runs out, the internal TX fails, or the reason is received from the remote side, etc.
AUTH_EXPIRE	2	2	The previous authentication is no longer valid. For the ESP Station, this reason is reported when: • auth is timed out • the reason is received from the AP. For the ESP AP, this reason is reported when: • the AP has not received any packets from the station in the past five minutes. • the AP is stopped by calling esp_wifi_stop(). • the station is deauthed by calling esp_wifi_deauth_sta()
AUTH_LEAVE	3	3	De-authenticated, because the sending STA is leaving (or has left). For the ESP Station, this reason is reported when: • it is received from the AP.
ASSOC_EXPIRE	4	4	Disassociated due to inactivity. For the ESP Station, this reason is reported when: • it is received from the AP. For the ESP AP, this reason is reported when: • the AP has not received any packets from the station in the past five minutes. • the AP is stopped by calling esp_wifi_stop(). • the station is deauthed by calling esp_wifi_deauth_sta()
ASSOC_TOOMANY	5	5	Disassociated, because the AP is unable to handle all currently associated STAs at the same time. For the ESP Station, this reason is reported when: • it is received from

4.21.11 ESP32-S2 Wi-Fi Station Connecting When Multiple APs Are Found

This scenario is similar as <*ESP32-S2 Wi-Fi Station Connecting Scenario*>, the difference is the station will not raise the event <*WIFI_EVENT_STA_DISCONNECTED*> unless it fails to connect all of the found APs.

4.21.12 Wi-Fi Reconnect

The station may disconnect due to many reasons, e.g. the connected AP is restarted etc. It's the application's responsibility to do the reconnect. The recommended reconnect strategy is to call esp_wifi_connect() on receiving event <*WIFI_EVENT_STA_DISCONNECTED*>.

Sometimes the application needs more complex reconnect strategy: - If the disconnect event is raised because the esp_wifi_disconnect() is called, the application may not want to do reconnect. - If the esp_wifi_scan_start() may be called at anytime, a better reconnect strategy is necessary, refer to <*Scan When Wi-Fi Is Connecting*>.

Another thing we need to consider is the reconnect may not connect the same AP if there are more than one APs with the same SSID. The reconnect always select current best APs to connect.

4.21.13 Wi-Fi Beacon Timeout

The beacon timeout mechanism is used by ESP32-S2 station to detect whether the AP is alive or not. If the station continuously loses 60 beacons of the connected AP, the beacon timeout happens.

After the beacon timeout happens, the station sends 5 probe requests to AP, it disconnects the AP and raises the event <*WIFI_EVENT_STA_DISCONNECTED*> if still no probe response or beacon is received from AP.

4.21.14 ESP32-S2 Wi-Fi Configuration

All configurations will be stored into flash when the Wi-Fi NVS is enabled; otherwise, refer to <*Wi-Fi NVS Flash*>.

Wi-Fi Mode

Call esp_wifi_set_mode() to set the Wi-Fi mode.

Mode	Description
WIFI_MODE_NULL	Null mode: in this mode, the internal data struct is not allocated to the station and the AP, while both the station and AP interfaces are not initialized for RX/TX Wi-Fi data. Generally, this mode is used for Sniffer, or when you only want to stop both the STA and the AP without calling esp_wifi_deinit() to unload the whole Wi-Fi driver.
WIFI_MODE_STA	Station mode: in this mode, esp_wifi_start() will init the internal station data, while the station's interface is ready for the RX and TX Wi-Fi data. After esp_wifi_connect() is called, the STA will connect to the target AP.
WIFI_MODE_AP	AP mode: in this mode, esp_wifi_start() will init the internal AP data, while the AP's interface is ready for RX/TX Wi-Fi data. Then, the Wi-Fi driver starts broad- casting beacons, and the AP is ready to get connected to other stations.
WIFI_MODE_APSTA	Station-AP coexistence mode: in this mode, esp_wifi_start() will simultaneously init both the station and the AP. This is done in station mode and AP mode. Please note that the channel of the external AP, which the ESP Station is connected to, has higher priority over the ESP AP channel.

Station Basic Configuration

API esp_wifi_set_config() can be used to configure the station. The table below describes the fields in detail.

Field	Description
ssid	This is the SSID of the target AP, to which the station wants to connect to.
password	Password of the target AP
scan_method	For WIFI_FAST_SCAN scan, the scan ends when the first matched AP is found, for WIFI_ALL_CHANNEL_SCAN, the scan finds all matched APs in all channels. The default scan is WIFI_FAST_SCAN.
bssid_set	If bssid_set is 0, the station connects to the AP whose SSID is the same as the field "ssid", while the field "bssid" is ignored. In all other cases, the station connects to the AP whose SSID is the same as the "ssid" field, while its BSSID is the same the "bssid" field .
bssid	This is valid only when bssid_set is 1; see field "bssid_set".
channel	If the channel is 0, the station scans the channel 1~N to search for the target AP; otherwise, the station starts by scanning the channel whose value is the same as that of the "channel" field, and then scans others to find the target AP. If you do not know which channel the target AP is running on, set it to 0.
sort_method	This field is only for WIFI_ALL_CHANNEL_SCAN If the sort_method is WIFI_CONNECT_AP_BY_SIGNAL, all matched APs are sorted by signal, for AP with best signal will be connected firstly. E.g. if the station want to connect AP whose ssid is "apxx", the scan finds two AP whose ssid equals to "apxx", the first AP's signal is -90dBm, the second AP's signal is -30dBm, the station connects the second AP firstly, it doesn't connect the first one unless it fails to connect the second one. If the sort_method is WIFI_CONNECT_AP_BY_SECURITY, all matched APs are sorted by security. E.g. if the station wants to connect AP whose ssid is "apxx", the scan finds two AP whose ssid is "apxx", the security of the first found AP is open while the second one is WPA2, the stations connects to the second AP firstly, it doesn't connect the second one unless it fails to connect the first one.
threshold	The threshold is used to filter the found AP, if the RSSI or security mode is less than the configured threshold, the AP will be discard. If the RSSI set to 0, it means default threshold, the default RSSI threshold is -127dBm. If the authmode threshold is set to 0, it means default threshold, the default authmode threshold is open.

注意: WEP/WPA security modes are deprecated in IEEE802.11-2016 specifications and are recommended not to be used. These modes can be rejected using authmode threshold by setting threshold as WPA2 by threshold.authmode as WIFI_AUTH_WPA2_PSK.

AP Basic Configuration

API esp_wifi_set_config() can be used to configure the AP. The table below describes the fields in detail.

Field	Description
ssid	SSID of AP; if the ssid[0] is 0xFF and ssid[1] is 0xFF, the AP defaults the SSID to ESP_aabbcc, where "aabbcc" is the last three bytes of the AP MAC.
password	Password of AP; if the auth mode is WIFI_AUTH_OPEN, this field will be ignored.
ssid_len	Length of SSID; if ssid_len is 0, check the SSID until there is a termination character. If ssid_len > 32, change it to 32; otherwise, set the SSID length according to ssid_len.
channel	Channel of AP; if the channel is out of range, the Wi-Fi driver defaults the channel to channel 1. So, please make sure the channel is within the required range. For more details, refer to <*Wi-Fi Country Code*>.
authmode	Auth mode of ESP AP; currently, ESP Wi-Fi does not support AUTH_WEP. If the authmode is an invalid value, AP defaults the value to WIFI_AUTH_OPEN.
ssid_hidden	If ssid_hidden is 1, AP does not broadcast the SSID; otherwise, it does broadcast the SSID.
max_connection	Currently, ESP Wi-Fi supports up to 10 Wi-Fi connections. If max_connection > 10, AP defaults the value to 10.
beacon_interval	Beacon interval; the value is 100 ~ 60000 ms, with default value being 100 ms. If the value is out of range, AP defaults it to 100 ms.

Wi-Fi Protocol Mode

Currently, the IDF supports the following protocol modes:

Protocol Mode	Description
802.11 B	Call esp_wifi_set_protocol(ifx, WIFI_PROTOCOL_11B) to set the station/AP to 802.11B-only mode.
802.11 BG	Call esp_wifi_set_protocol(ifx, WIFI_PROTOCOL_11B|WIFI_ PROTOCOL_11G) to set the station/AP to 802.11BG mode.
802.11 BGN	Call esp_wifi_set_protocol(ifx, WIFI_PROTOCOL_11B| WIFI_PROTOCOL_11G|WIFI_PROTOCOL_11N) to set the station/ AP to BGN mode.
802.11 BGNLR	Call esp_wifi_set_protocol(ifx, WIFI_PROTOCOL_11B| WIFI_PROTOCOL_11G|WIFI_PROTOCOL_11N|WIFI_PROT to set the station/AP to BGN and the LR mode.
802.11 LR	Call esp_wifi_set_protocol(ifx, WIFI_PROTOCOL_LR) to set the station/AP only to the LR mode. **This mode is an Espressif-patented mode which can achieve a one-kilometer line of sight range. Please, make sure both the station and the AP are connected to an ESP device**

Long Range (LR)

Long Range (LR) mode is an Espressif-patented Wi-Fi mode which can achieve a one-kilometer line of sight range. It has better reception sensitivity, stronger anti-interference ability and longer transmission distance than the traditional 802.11B mode.

LR Compitability Since LR is Espressif unique Wi-Fi mode, only ESP32-S2 devices can transmit and receive the LR data. In other words, the ESP32-S2 device should NOT transmit the data in LR data rate if the connected device doesn’ t support LR. The application can achieve this by configuring suitable Wi-Fi mode. If the negotiated mode supports LR, the ESP32-S2 may transmit data in LR rate, otherwise, ESP32-S2 will transmit all data in traditional Wi-Fi data rate.

Following table depicts the Wi-Fi mode negotiation:

APSTA	BGN	BG	B	BGNLR	BGLR	BLR	LR
BGN	BGN	BG	B	BGN	BG	B	•
BG	BG	BG	B	BG	BG	B	•
B	B	B	B	B	B	B	•
BGNLR	•	•	•	BGNLR	BGLR	BLR	LR
BGLR	•	•	•	BGLR	BGLR	BLR	LR
BLR	•	•	•	BLR	BLR	BLR	LR
LR	•	•	•	LR	LR	LR	LR

In above table, the row is the Wi-Fi mode of AP and the column is the Wi-Fi mode of station. The “-” indicates Wi-Fi mode of the AP and station are not compatible.

According to the table, we can conclude that:

- For LR enabled in ESP32-S2 AP, it's incompatible with traditional 802.11 mode because the beacon is sent in LR mode.
- For LR enabled in ESP32-S2 station and the mode is NOT LR only mode, it's compatible with traditional 802.11 mode.
- If both station and AP are ESP32-S2 devices and both of them enable LR mode, the negotiated mode supports LR.

If the negotiated Wi-Fi mode supports both traditional 802.11 mode and LR mode, it's the WiFi driver's responsibility to automatically select the best data rate in different Wi-Fi mode and the application don't need to care about it.

LR Impacts to Traditional Wi-Fi device

The data transmission in LR rate has no impacts on the traditional Wi-Fi device because:

- The CCA and backoff process in LR mode are consistent with 802.11 specification.
- The traditional Wi-Fi device can detect the LR signal via CCA and do backoff.

In other words, the impact transmission in LR mode is similar as the impact in 802.11B mode.

LR Transmission Distance The reception sensitivity of LR has about 4 dB gain than the traditional 802.11 B mode, theoretically the transmission distance is about 2 to 2.5 times the distance of 11B.

LR Throughput The LR rate has very limited throughput, because the raw PHY data rates are 1/2 Mbps and 1/4 Mbps.

When to Use LR The general conditions for using LR are:

- Both the AP and station are Espressif devices.
- Long distance Wi-Fi connection and data transmission is required.
- Data throughput requirements are very small, such as remote device control, etc.

Wi-Fi Country Code

Call `esp_wifi_set_country()` to set the country info. The table below describes the fields in detail, please consult local 2.4 GHz RF operating regulations before configuring these fields.

Field	Description
cc[3]	Country code string, this attributes identify the country or noncountry entity in which the station/AP is operating. If it's a country, the first two octets of this string is the two character country info as described in document ISO/IEC3166-1. The third octect is one of the following: • an ASCII space character, if the regulations under which the station/AP is operating encompass all environments for the current frequency band in the country. • an ASCII 'O' character if the regulations under which the station/AP is operating are for an outdoor environment only. • an ASCII 'I' character if the regulations under which the station/AP is operating are for an indoor environment only. • an ASCII 'X' character if the station/AP is operating under a noncountry entity. The first two octets of the noncountry entity is two ASCII 'XX' characters. • the binary representation of the Operating Class table number currently in use. Refer to Annex E, IEEE Std 802.11-2020.
schan	Start channel, it's the minimum channel number of the regulations under which the station/AP can operate.
nchan	Total number of channels as per the regulations, e.g. if the schan=1, nchan=13, it means the station/AP can send data from channel 1 to 13.
policy	Country policy, this field control which country info will be used if the configured country info is conflict with the connected AP's. More description about policy is provided in following section.

The default country info is:

```
wifi_country_t config = {
    .cc = "CN",
    .schan = 1,
    .nchan = 13,
    .policy = WIFI_COUNTRY_POLICY_AUTO,
};
```

If the Wi-Fi Mode is station/AP coexist mode, they share the same configured country info. Sometimes, the country info of AP, to which the station is connected, is different from the country info of configured. For example, the configured station has country info:

```
wifi_country_t config = {
    .cc = "JP",
    .schan = 1,
    .nchan = 14,
    .policy = WIFI_COUNTRY_POLICY_AUTO,
};
```

but the connected AP has country info:

```
wifi_country_t config = {
    .cc = "CN",
    .schan = 1,
    .nchan = 13,
};
```

then country info of connected AP's is used.

Following table depicts which country info is used in different Wi-Fi Mode and different country policy, also describe the impact to active scan.

Wi-Fi Mode	Policy	Description
Station	WIFI_COUNTRY_POLICY_AUTO	If the connected AP has country IE in its beacon, the country info equals to the country info in beacon. Otherwise, use the default country info. For scan: Use active scan from 1 to 11 and use passive scan from 12 to 14. Always keep in mind that if an AP with hidden SSID and station is set to a passive scan channel, the passive scan will not find it. In other words, if the application hopes to find the AP with hidden SSID in every channel, the policy of country info should be configured to WIFI_COUNTRY_POLICY_MANUAL.
Station	WIFI_COUNTRY_POLICY_MANUAL	Always use the configured country info. For scan: Use active scan from schan to schan+nchan-1.
AP	WIFI_COUNTRY_POLICY_AUTO	Always use the configured country info.
AP	WIFI_COUNTRY_POLICY_MANUAL	Always use the configured country info.
Station/AP-coexistence	WIFI_COUNTRY_POLICY_AUTO	Station: Same as station mode with policy WIFI_COUNTRY_POLICY_AUTO. AP: If the station does not connect to any external AP, the AP uses the configured country info. If the station connects to an external AP, the AP has the same country info as the station.
Station/AP-coexistence	WIFI_COUNTRY_POLICY_MANUAL	Station: Same as station mode with policy WIFI_COUNTRY_POLICY_MANUAL. AP: Same as AP mode with policy WIFI_COUNTRY_POLICY_MANUAL.

Home Channel In AP mode, the home channel is defined as that of the AP channel. In Station mode, the home channel is defined as the channel of the AP to which the station is connected. In Station+AP mode, the home channel of AP and station must be the same. If the home channels of Station and AP are different, the station ' s home channel is always in priority. Take the following as an example: at the beginning, the AP is on channel 6, then the station connects to an AP whose channel is 9. Since the station ' s home channel has a higher priority, the AP needs to switch its channel from 6 to 9 to make sure that both station and AP have the same home channel. While switching channel, the ESP32-S2 in SoftAP mode will notify the connected stations about the channel migration using a Channel Switch Announcement (CSA). Stations that support channel switching will transition smoothly whereas stations who do not will disconnect and reconnect to the SoftAP.

Wi-Fi Vendor IE Configuration

By default, all Wi-Fi management frames are processed by the Wi-Fi driver, and the application does not need to care about them. Some applications, however, may have to handle the beacon, probe request, probe response and other management frames. For example, if you insert some vendor-specific IE into the management frames, it is only the management frames which contain this vendor-specific IE that will be processed. In ESP32-S2, esp_wifi_set_vendor_ie() and esp_wifi_set_vendor_ie_cb() are responsible for this kind of tasks.

4.21.15 Wi-Fi Security

In addition to traditional security methods (WEP/WPA-TKIP/WPA2-CCMP), ESP32-S2 Wi-Fi now supports state-of-the-art security protocols, namely Protected Management Frames based on 802.11w standard and Wi-Fi Protected Access 3 (WPA3-Personal). Together, PMF and WPA3 provide better privacy and robustness against known attacks in traditional modes.

Protected Management Frames (PMF)

In Wi-Fi, management frames such as beacons, probes, (de)authentication, (dis)association are used by non-AP stations to scan and connect to an AP. Unlike data frames, these frames are sent unencrypted. An attacker can use eavesdropping and packet injection to send spoofed (de)authentication/(dis)association frames at the right time, leading to following attacks in case of unprotected management frame exchanges.

- DOS attack on one or all clients in the range of the attacker.
- Tearing down existing association on AP side by sending association request.
- Forcing a client to perform 4-way handshake again in case PSK is compromised in order to get PTK.
- Getting SSID of hidden network from association request.
- Launching man-in-the-middle attack by forcing clients to deauth from legitimate AP and associating to a rogue one.

PMF provides protection against these attacks by encrypting unicast management frames and providing integrity checks for broadcast management frames. These include deauthentication, disassociation and robust management frames. It also provides Secure Association (SA) teardown mechanism to prevent spoofed association/authentication frames from disconnecting already connected clients.

There are 3 types of PMF configuration modes on both Station and AP side -

- PMF Optional
- PMF Required
- PMF Disabled

Depending on the PMF configuration on Station and AP side, the resulting connection will behave differently. Below table summarises all possible outcomes.

STA Setting	AP Setting	Outcome
PMF Optional	PMF Optional/Required	Mgmt Frames Protected
PMF Optional	PMF Disabled	Mgmt Frames Not Protected
PMF Required	PMF Optional/Required	Mgmt Frames Protected
PMF Required	PMF Disabled	STA refuses Connection
PMF Disabled	PMF Optional/Disabled	Mgmt Frames Not Protected
PMF Disabled	PMF Required	AP refuses Connection

ESP32-S2 supports PMF only in Station mode. Station defaults to PMF Optional mode and disabling PMF is not possible. For even higher security, PMF Required mode can be enabled by setting the `required` flag in *pmf_cfg* while using the *`esp_wifi_set_config()`* API. This will result in Station only connecting to a PMF enabled AP and rejecting all other AP's.

注意: `capable` flag in *pmf_cfg* is deprecated and set to true internally. This is to take the additional security benefit of PMF whenever possible.

WPA3-Personal

Wi-Fi Protected Access-3 (WPA3) is a set of enhancements to Wi-Fi access security intended to replace the current WPA2 standard. In order to provide more robust authentication, WPA3 uses Simultaneous Authentication of Equals (SAE), which is password-authenticated key agreement method based on Diffie-Hellman key exchange. Unlike WPA2, the technology is resistant to offline-dictionary attack, where the attacker attempts to determine shared password based on captured 4-way handshake without any further network interaction. WPA3 also provides forward secrecy, which means the captured data cannot be decrypted even if password is compromised after data transmission. Please refer to Security section of Wi-Fi Alliance's official website for further details.

In order to enable WPA3-Personal, "Enable WPA3-Personal" should be selected in menuconfig. If enabled, ESP32-S2 uses SAE for authentication if supported by the AP. Since PMF is a mandatory requirement for WPA3, PMF capability should be at least set to "PMF capable, but not required" for ESP32-S2 to use WPA3 mode. Application developers need not worry about the underlying security mode as highest available is chosen from security standpoint.

Note that Wi-Fi stack size requirement will increase approximately by 3k when WPA3 is used. Currently, WPA3 is supported only in Station mode.

4.21.16 ESP32-S2 Wi-Fi Power-saving Mode

Station Sleep

Currently, ESP32-S2 Wi-Fi supports the Modem-sleep mode which refers to the legacy power-saving mode in the IEEE 802.11 protocol. Modem-sleep mode works in Station-only mode and the station must connect to the AP first. If the Modem-sleep mode is enabled, station will switch between active and sleep state periodically. In sleep state, RF, PHY and BB are turned off in order to reduce power consumption. Station can keep connection with AP in modem-sleep mode.

Modem-sleep mode includes minimum and maximum power save modes. In minimum power save mode, station wakes up every DTIM to receive beacon. Broadcast data will not be lost because it is transmitted after DTIM. However, it can not save much more power if DTIM is short for DTIM is determined by AP.

In maximum power-saving mode, station wakes up in every listen interval to receive beacon. This listen interval can be set to be longer than the AP DTIM period. Broadcast data may be lost because station may be in sleep state at DTIM time. If listen interval is longer, more power is saved, but broadcast data is more easy to lose. Listen interval can be configured by calling API *esp_wifi_set_config()* before connecting to AP.

Call `esp_wifi_set_ps(WIFI_PS_MIN_MODEM)` to enable Modem-sleep minimum power-saving mode or `esp_wifi_set_ps(WIFI_PS_MAX_MODEM)` to enable Modem-sleep maximum power-saving mode after calling *esp_wifi_init()*. When station connects to AP, Modem-sleep will start. When station disconnects from AP, Modem-sleep will stop.

Call `esp_wifi_set_ps(WIFI_PS_NONE)` to disable modem sleep entirely. This has much higher power consumption, but provides minimum latency for receiving Wi-Fi data in real time. When modem sleep is enabled, received Wi-Fi data can be delayed for as long as the DTIM period (minimum power save mode) or the listen interval (maximum power save mode). Disabling modem sleep entirely is not possible for Wi-Fi and Bluetooth coexist mode.

The default Modem-sleep mode is WIFI_PS_MIN_MODEM.

AP Sleep

Currently ESP32-S2 AP doesn't support all of the power save feature defined in Wi-Fi specification. To be specific, the AP only caches unicast data for the stations connect to this AP, but doesn't cache the multicast data for the stations. If stations connected to the ESP32-S2 AP are power save enabled, they may experience multicast packet loss.

In future, all power save features will be supported on ESP32-S2 AP.

4.21.17 ESP32-S2 Wi-Fi Connect Crypto

Now ESP32-S2 have two group crypto functions can be used when do wifi connect, one is the original functions, the other is optimized by ESP hardware: 1. Original functions which is the source code used in the folder components/wpa_supplicant/src/crypto function; 2. The optimized functions is in the folder components/wpa_supplicant/src/fast_crypto, these function used the hardware crypto to make it faster than origin one, the type of function's name add *fast_* to distinguish with the original one. For example, the API aes_wrap() is used to encrypt frame information when do 4 way handshake, the fast_aes_wrap() has the same result but can be faster.

Two groups of crypto function can be used when register in the wpa_crypto_funcs_t, wpa2_crypto_funcs_t and wps_crypto_funcs_t structure, also we have given the recommend functions to register in the fast_crypto_ops.c, you can register the function as the way you need, however what should make action is that the crypto_hash_xxx function and crypto_cipher_xxx function need to register with the same function to operation. For example, if you register crypto_hash_init() function to initialize the esp_crypto_hash structure,

you need use the crypto_hash_update() and crypto_hash_finish() function to finish the operation, rather than fast_crypto_hash_update() or fast_crypto_hash_finish().

4.21.18 ESP32-S2 Wi-Fi Throughput

The table below shows the best throughput results we got in Espressif ' s lab and in a shield box.

Type/Throughput	Air In Lab	Shield-box	Test Tool	IDF Version (commit ID)
Raw 802.11 Packet RX	N/A	**130 MBit/sec**	Internal tool	NA
Raw 802.11 Packet TX	N/A	**130 MBit/sec**	Internal tool	NA
UDP RX	30 MBit/sec	90 MBit/sec	iperf example	05838641
UDP TX	30 MBit/sec	60 MBit/sec	iperf example	05838641
TCP RX	20 MBit/sec	50 MBit/sec	iperf example	05838641
TCP TX	20 MBit/sec	50 MBit/sec	iperf example	05838641

When the throughput is tested by iperf example, the sdkconfig is examples/wifi/iperf/sdkconfig.ci.99

4.21.19 Wi-Fi 80211 Packet Send

The *esp_wifi_80211_tx()* API can be used to:

- Send the beacon, probe request, probe response, action frame.
- Send the non-QoS data frame.

It cannot be used for sending encrypted or QoS frames.

Preconditions of Using esp_wifi_80211_tx()

- The Wi-Fi mode is station, or AP, or station/AP.
- Either esp_wifi_set_promiscuous(true), or *esp_wifi_start()*, or both of these APIs return ESP_OK. This is because Wi-Fi hardware must be initialized before *esp_wifi_80211_tx()* is called. In ESP32-S2, both esp_wifi_set_promiscuous(true) and *esp_wifi_start()* can trigger the initialization of Wi-Fi hardware.
- The parameters of *esp_wifi_80211_tx()* are hereby correctly provided.

Data rate

- If there is no WiFi connection, the data rate is 1Mbps.
- If there is WiFi connection and the packet is from station to AP or from AP to station, the data rate is same as the WiFi connection. Otherwise the data rate is 1Mbps.

Side-Effects to Avoid in Different Scenarios

Theoretically, if we do not consider the side-effects the API imposes on the Wi-Fi driver or other stations/APs, we can send a raw 802.11 packet over the air, with any destination MAC, any source MAC, any BSSID, or any other type of packet. However,robust/useful applications should avoid such side-effects. The table below provides some tips/recommendations on how to avoid the side-effects of esp_wifi_80211_tx in different scenarios.

Scenario	Description
No WiFi connection	In this scenario, no Wi-Fi connection is set up, so there are no side-effects on the Wi-Fi driver. If en_sys_seq==true, the Wi-Fi driver is responsible for the sequence control. If en_sys_seq==false, the application needs to ensure that the buffer has the correct sequence. Theoretically, the MAC address can be any address. However, this may impact other stations/APs with the same MAC/BSSID. Side-effect example#1 The application calls esp_wifi_80211_tx to send a beacon with BSSID == mac_x in AP mode, but the mac_x is not the MAC of the AP interface. Moreover, there is another AP, say "other-AP" , whose bssid is mac_x. If this happens, an "unexpected behavior" may occur, because the stations which connect to the "other-AP" cannot figure out whether the beacon is from the "other-AP" or the esp_wifi_80211_tx. To avoid the above-mentioned side-effects, we recommend that: • If esp_wifi_80211_tx is called in Station mode, the first MAC should be a multicast MAC or the exact target-device's MAC, while the second MAC should be that of the station interface. • If esp_wifi_80211_tx is called in AP mode, the first MAC should be a multicast MAC or the exact target-device's MAC, while the second MAC should be that of the AP interface. The recommendations above are only for avoiding side-effects and can be ignored when there are good reasons for doing this.
Have WiFi connection	When the Wi-Fi connection is already set up, and the sequence is controlled by the application, the latter may impact the sequence control of the Wi-Fi connection, as a whole. So, the en_sys_seq need to be true, otherwise ESP_ERR_WIFI_ARG is returned. The MAC-address recommendations in the "No WiFi connection" scenario also apply to this scenario. If the WiFi mode is station mode and the MAC address1 is the MAC of AP to which the station is connected, the MAC address2 is the MAC of station interface, we say the packets is from the station to AP. On the other hand, if the WiFi mode is AP mode and the MAC address1 is the MAC of the station who connects to this AP, the MAC address2 is the MAC of AP interface, we say the packet is from the AP to station. To avoid conflicting with WiFi connections, the following checks are applied: • If the packet type is data and is from the station to AP, the ToDS bit in ieee80211 frame control should be 1, the FromDS bit should be 0, otherwise the packet will be discarded by WiFi driver. • If the packet type is data and is from the AP to station, the ToDS bit in ieee80211 frame control should be 0, the FromDS bit should be 1, otherwise the packet will be discarded by WiFi driver. • If the packet is from station to AP or from AP to station, the Power Management, More Data, Re-Transmission bits should be 0, otherwise the packet will be discarded by WiFi driver. ESP_ERR_WIFI_ARG is returned if any check fails.

4.21.20 Wi-Fi Sniffer Mode

The Wi-Fi sniffer mode can be enabled by esp_wifi_set_promiscuous(). If the sniffer mode is enabled, the following packets **can** be dumped to the application:

- 802.11 Management frame
- 802.11 Data frame, including MPDU, AMPDU, AMSDU, etc.
- 802.11 MIMO frame, for MIMO frame, the sniffer only dumps the length of the frame.
- 802.11 Control frame

The following packets will **NOT** be dumped to the application:

- 802.11 error frame, such as the frame with a CRC error, etc.

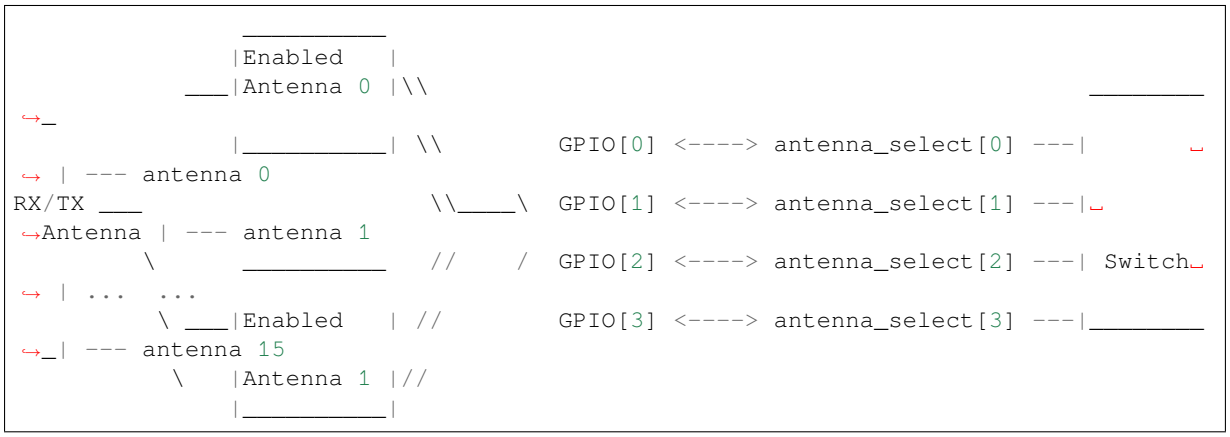
For frames that the sniffer **can** dump, the application can additionally decide which specific type of packets can be filtered to the application by using esp_wifi_set_promiscuous_filter() and esp_wifi_set_promiscuous_ctrl_filter(). By default, it will filter all 802.11 data and management frames to the application.

The Wi-Fi sniffer mode can be enabled in the Wi-Fi mode of WIFI_MODE_NULL, or WIFI_MODE_STA, or WIFI_MODE_AP, or WIFI_MODE_APSTA. In other words, the sniffer mode is active when the station is connected to the AP, or when the AP has a Wi-Fi connection. Please note that the sniffer has a **great impact** on the throughput of the station or AP Wi-Fi connection. Generally, we should **NOT** enable the sniffer, when the station/AP Wi-Fi connection experiences heavy traffic unless we have special reasons.

Another noteworthy issue about the sniffer is the callback wifi_promiscuous_cb_t. The callback will be called directly in the Wi-Fi driver task, so if the application has a lot of work to do for each filtered packet, the recommendation is to post an event to the application task in the callback and defer the real work to the application task.

4.21.21 Wi-Fi Multiple Antennas

The Wi-Fi multiple antennas selecting can be depicted as following picture:



ESP32-S2 supports up to sixteen antennas through external antenna switch. The antenna switch can be controlled by up to four address pins - antenna_select[0:3]. Different input value of antenna_select[0:3] means selecting different antenna. E.g. the value ‘0b1011’ means the antenna 11 is selected. The default value of antenna_select[3:0] is ‘0b0000’, it means the antenna 0 is selected by default.

Up to four GPIOs are connected to the four active high antenna_select pins. ESP32-S2 can select the antenna by control the GPIO[0:3]. The API *esp_wifi_set_ant_gpio()* is used to configure which GPIOs are connected to antenna_selects. If GPIO[x] is connected to antenna_select[x], then gpio_config->gpio_cfg[x].gpio_select should be set to 1 and gpio_config->gpio_cfg[x].gpio_num should be provided.

For the specific implementation of the antenna switch, there may be illegal values in *antenna_select[0:3]*. It means that ESP32-S2 may support less than sixteen antennas through the switch. For example, ESP32-WROOM-DA which uses RTC6603SP as the antenna switch, supports two antennas. Two GPIOs are connected to two active high antenna selection inputs. The value ‘0b01’ means the antenna 0 is selected, the value ‘0b10’ means the antenna 1 is selected. Values ‘0b00’ and ‘0b11’ are illegal.

Although up to sixteen antennas are supported, only one or two antennas can be simultaneously enabled for RX/TX. The API *esp_wifi_set_ant()* is used to configure which antennas are enabled.

The enabled antennas selecting algorithm is also configured by *esp_wifi_set_ant()*. The RX/TX antenna mode can be WIFI_ANT_MODE_ANT0, WIFI_ANT_MODE_ANT1 or WIFI_ANT_MODE_AUTO. If the antenna mode is WIFI_ANT_MODE_ANT0, the enabled antenna 0 is selected for RX/TX data. If the antenna mode is WIFI_ANT_MODE_ANT1, the enabled antenna 1 is selected for RX/TX data. Otherwise, WiFi automatically selects the antenna that has better signal from the enabled antennas.

If the RX antenna mode is WIFI_ANT_MODE_AUTO, the default antenna mode also needs to be set. Because the RX antenna switching only happens when some conditions are met, e.g. the RX antenna starts to switch if the RSSI is lower than -65dBm and if another antenna has better signal etc, RX uses the default antenna if the conditions are not met. If the default antenna mode is WIFI_ANT_MODE_ANT1, the enabled antenna 1 is used as the default RX antenna, otherwise the enabled antenna 0 is used as the default RX antenna.

Some limitations need to be considered:

- The TX antenna can be set to WIFI_ANT_MODE_AUTO only if the RX antenna mode is WIFI_ANT_MODE_AUTO because TX antenna selecting algorithm is based on RX antenna in WIFI_ANT_MODE_AUTO type.
- Currently BT doesn't support the multiple antennas feature, please don't use multiple antennas related APIs.

Following is the recommended scenarios to use the multiple antennas:

- In Wi-Fi mode WIFI_MODE_STA, both RX/TX antenna modes are configured to WIFI_ANT_MODE_AUTO. The WiFi driver selects the better RX/TX antenna automatically.
- The RX antenna mode is configured to WIFI_ANT_MODE_AUTO. The TX antenna mode is configured to WIFI_ANT_MODE_ANT0 or WIFI_ANT_MODE_ANT1. The applications can choose to always select a specified antenna for TX, or implement their own TX antenna selecting algorithm, e.g. selecting the TX antenna mode based on the channel switch information etc.
- Both RX/TX antenna modes are configured to WIFI_ANT_MODE_ANT0 or WIFI_ANT_MODE_ANT1.

Wi-Fi Multiple Antennas Configuration

Generally, following steps can be taken to configure the multiple antennas:

- Configure which GPIOs are connected to the antenna_selects, for example, if four antennas are supported and GPIO20/GPIO21 are connected to antenna_select[0]/antenna_select[1], the configurations look like:

```
wifi_ant_gpio_config_t config = {
    { .gpio_select = 1, .gpio_num = 20 },
    { .gpio_select = 1, .gpio_num = 21 }
};
```

- Configure which antennas are enabled and how RX/TX use the enabled antennas, for example, if antenna1 and antenna3 are enabled, the RX needs to select the better antenna automatically and uses antenna1 as its default antenna, the TX always selects the antenna3. The configuration looks like:

```
wifi_ant_config_t config = {
    .rx_ant_mode = WIFI_ANT_MODE_AUTO,
    .rx_ant_default = WIFI_ANT_ANT0,
    .tx_ant_mode = WIFI_ANT_MODE_ANT1,
    .enabled_ant0 = 1,
    .enabled_ant1 = 3
};
```

4.21.22 Wi-Fi Channel State Information

Channel state information (CSI) refers to the channel information of a Wi-Fi connection. In ESP32-S2, this information consists of channel frequency responses of sub-carriers and is estimated when packets are received from the transmitter. Each channel frequency response of sub-carrier is recorded by two bytes of signed characters. The first

one is imaginary part and the second one is real part. There are up to three fields of channel frequency responses according to the type of received packet. They are legacy long training field (LLTF), high throughput LTF (HT-LTF) and space time block code HT-LTF (STBC-HT-LTF). For different types of packets which are received on channels with different state, the sub-carrier index and total bytes of signed characters of CSI is shown in the following table.

channel	secondary channel	none			below					above				
packet information	signal mode	non HT	HT		non HT	HT				non HT	HT			
	channel bandwidth	20MHz	20MHz		20MHz	20MHz		40MHz		20MHz	20MHz		40MHz	
	STBC	non STBC	non STBC	STBC	non STBC	non STBC	STBC	non STBC	STBC	non STBC	non STBC	STBC	non STBC	STBC
sub-carrier index	LLTF	0~31,-32~-1	0~31,-32~-1	0~31,-32~-1	0~63	0~63	0~63	0~63	0~63	-64~-1	-64~-1	-64~-1	-64~-1	-64~-1
	HT-LTF	•	0~31,-32~-1	0~31,-32~-1	•	0~63	0~62	0~63,-64~-1	0~60,-60~-1	•	-64~-1	-62~-1	0~63,-64~-1	0~60,-60~-1
	STBC-HT-LTF	•	•	0~31,-32~-1	•	•	0~62	•	0~60,-60~-1	•	•	-62~-1	•	0~60,-60~-1
total bytes		128	256	384	128	256	380	384	612	128	256	376	384	612

All of the information in the table can be found in the structure wifi_csi_info_t.

- Secondary channel refers to secondary_channel field of rx_ctrl field.
- Signal mode of packet refers to sig_mode field of rx_ctrl field.
- Channel bandwidth refers to cwb field of rx_ctrl field.
- STBC refers to stbc field of rx_ctrl field.
- Total bytes refers to len field.
- The CSI data corresponding to each Long Training Field(LTF) type is stored in a buffer starting from the buf field. Each item is stored as two bytes: imaginary part followed by real part. The order of each item is the same as the sub-carrier in the table. The order of LTF is: LLTF, HT-LTF, STBC-HT-LTF. However all 3 LTFs may not be present, depending on the channel and packet information (see above).
- If first_word_invalid field of wifi_csi_info_t is true, it means that the first four bytes of CSI data is invalid due to a hardware limitation in ESP32-S2.
- More information like RSSI, noise floor of RF, receiving time and antenna is in the rx_ctrl field.

When imaginary part and real part data of sub-carrier are used, please refer to the table below.

PHY standard	Sub-carrier range	Pilot sub-carrier	Sub-carrier(total/data)
802.11a/g	-26 to +26	-21, -7, +7, +21	52 total, 48 usable
802.11n, 20MHz	-28 to +28	-21, -7, +7, +21	56 total, 52 usable
802.11n, 40MHz	-57 to +57	-53, -25, -11, +11, +25, +53	114 total, 108 usable

注解:

- For STBC packet, CSI is provided for every space-time stream without CSD (cyclic shift delay). As each cyclic shift on the additional chains shall be -200 ns, only the CSD angle of first space-time stream is recorded in sub-carrier 0 of HT-LTF and STBC-HT-LTF for there is no channel frequency response in sub-carrier 0. CSD[10:0] is 11 bits, ranging from -pi to pi.
- If LLTF, HT-LTF, or STBC-HT-LTF is not enabled by calling API `esp_wifi_set_csi_config()`, the total bytes of CSI data will be fewer than that in the table. For example, if LLTF and HT-LTF is not enabled and STBC-HT-LTF is enabled, when a packet is received with the condition above/HT/40MHz/STBC, the

total bytes of CSI data is 244 ((61 + 60) * 2 + 2 = 244. The result is aligned to four bytes, and the last two bytes are invalid).

4.21.23 Wi-Fi Channel State Information Configure

To use Wi-Fi CSI, the following steps need to be done.

- Select Wi-Fi CSI in menuconfig. Go to `Menuconfig > Components config > Wi-Fi > Wi-Fi CSI (Channel State Information)`.
- Set CSI receiving callback function by calling API *`esp_wifi_set_csi_rx_cb()`*.
- Configure CSI by calling API *`esp_wifi_set_csi_config()`*.
- Enable CSI by calling API *`esp_wifi_set_csi()`*.

The CSI receiving callback function runs from Wi-Fi task. So, do not do lengthy operations in the callback function. Instead, post necessary data to a queue and handle it from a lower priority task. Because station does not receive any packet when it is disconnected and only receives packets from AP when it is connected, it is suggested to enable sniffer mode to receive more CSI data by calling *`esp_wifi_set_promiscuous()`*.

4.21.24 Wi-Fi HT20/40

ESP32-S2 supports Wi-Fi bandwidth HT20 or HT40 and does not support HT20/40 coexist. *`esp_wifi_set_bandwidth()`* can be used to change the default bandwidth of station or AP. The default bandwidth for ESP32-S2 station and AP is HT40.

In station mode, the actual bandwidth is firstly negotiated during the Wi-Fi connection. It is HT40 only if both the station and the connected AP support HT40, otherwise it's HT20. If the bandwidth of connected AP is changes, the actual bandwidth is negotiated again without Wi-Fi disconnecting.

Similarly, in AP mode, the actual bandwidth is negotiated between AP and the stations that connect to the AP. It's HT40 if the AP and one of the stations support HT40, otherwise it's HT20.

In station/AP coexist mode, the station/AP can configure HT20/40 seperately. If both station and AP are negotiated to HT40, the HT40 channel should be the channel of station because the station always has higher priority than AP in ESP32-S2. E.g. the configured bandwidth of AP is HT40, the configured primary channel is 6 and the configured secondary channel is 10. The station is connected to an router whose primary channel is 6 and secondary channel is 2, then the actual channel of AP is changed to primary 6 and secondary 2 automatically.

Theoretically the HT40 can gain better throughput because the maximum raw physicial (PHY) data rate for HT40 is 150Mbps while it's 72Mbps for HT20. However, if the device is used in some special environment, e.g. there are too many other Wi-Fi devices around the ESP32-S2 device, the performance of HT40 may be degraded. So if the applications need to support same or similar scenarios, it's recommended that the bandwidth is always configured to HT20.

4.21.25 Wi-Fi QoS

ESP32-S2 supports all the mandatory features required in WFA Wi-Fi QoS Certification.

Four ACs(Access Category) are defined in Wi-Fi specification, each AC has a its own priority to access the Wi-Fi channel. Moreover a map rule is defined to map the QoS priority of other protocol, such as 802.11D or TCP/IP precedence to Wi-Fi AC.

Below is a table describes how the IP Precedences are mapped to Wi-Fi ACs in ESP32-S2, it also indicates whether the AMPDU is supported for this AC. The table is sorted with priority descending order, namely, the AC_VO has highest priority.

IP Precedence	Wi-Fi AC	Support AMPDU?
6, 7	AC_VO (Voice)	No
4, 5	AC_VI (Video)	Yes
3, 0	AC_BE (Best Effort)	Yes
1, 2	AC_BK (Background)	Yes

The application can make use of the QoS feature by configuring the IP precedence via socket option IP_TOS. Here is an example to make the socket to use VI queue:

```
const int ip_precedence_vi = 4;
const int ip_precedence_offset = 5;
int priority = (ip_precedence_vi << ip_precedence_offset);
setsockopt(socket_id, IPPROTO_IP, IP_TOS, &priority, sizeof(priority));
```

Theoretically the higher priority AC has better performance than the low priority AC, however, it's not always be true, here

- For some really important application traffic, can put it into AC_VO queue. Avoid sending big traffic via AC_VO queue. On one hand, the AC_VO queue doesn't support AMPDU and it can't get better performance than other queue if the traffic is big, on the other hand, it may impact the the management frames that also use AC_VO queue.
- Avoid using more than two different AMPDU supported precedences, e.g. socket A uses precedence 0, socket B uses precedence 1, socket C uses precedence 2, this is a bad design because it may need much more memory. To be detailed, the Wi-Fi driver may generate a Block Ack session for each precedence and it needs more memory if the Block Ack session is setup.

4.21.26 Wi-Fi AMSDU

ESP32-S2 supports receiving AMSDU but doesn't support transmitting AMSDU. The transmitting AMSDU is not necessary since ESP32-S2 has transmitting AMPDU.

4.21.27 Wi-Fi Fragment

ESP32-S2 supports Wi-Fi receiving fragment, but doesn't support Wi-Fi transmitting fragment. The Wi-Fi transmitting fragment will be supported in future release.

4.21.28 WPS Enrolle

ESP32-S2 supports WPS enrollee feature in Wi-Fi mode WIFI_MODE_STA or WIFI_MODE_APSTA. Currently ESP32-S2 supports WPS enrollee type PBC and PIN.

4.21.29 Wi-Fi Buffer Usage

This section is only about the dynamic buffer configuration.

Why Buffer Configuration Is Important

In order to get a robust, high-performance system, we need to consider the memory usage/configuration very carefully, becau

- the available memory in ESP32-S2 is limited.
- currently, the default type of buffer in LwIP and Wi-Fi drivers is "dynamic", **which means that both the LwIP and Wi-Fi share memory with the application**. Programmers should always keep this in mind; otherwise, they will face a memory issue, such as "running out of heap memory".

- it is very dangerous to run out of heap memory, as this will cause ESP32-S2 an "undefined behavior". Thus, enough heap memory should be reserved for the application, so that it never runs out of it.
- the Wi-Fi throughput heavily depends on memory-related configurations, such as the TCP window size, Wi-Fi RX/TX dynamic buffer number, etc.
- the peak heap memory that the ESP32-S2 LwIP/Wi-Fi may consume depends on a number of factors, such as the maximum TCP/UDP connections that the application may have, etc.
- the total memory that the application requires is also an important factor when considering memory configuration.

Due to these reasons, there is not a good-for-all application configuration. Rather, we have to consider memory configurations separately for every different application.

Dynamic vs. Static Buffer

The default type of buffer in Wi-Fi drivers is "dynamic". Most of the time the dynamic buffer can significantly save memory. However, it makes the application programming a little more difficult, because in this case the application needs to consider memory usage in Wi-Fi.

lwIP also allocates buffers at the TCP/IP layer, and this buffer allocation is also dynamic. See *lwIP documentation section about memory use and performance*.

Peak Wi-Fi Dynamic Buffer

The Wi-Fi driver supports several types of buffer (refer to *Wi-Fi Buffer Configure*). However, this section is about the usage of the dynamic Wi-Fi buffer only. The peak heap memory that Wi-Fi consumes is the **theoretically-maximum memory** that the Wi-Fi driver consumes. Generally, the peak memory depends on:

- the number of dynamic rx buffers that are configured: wifi_rx_dynamic_buf_num
- the number of dynamic tx buffers that are configured: wifi_tx_dynamic_buf_num
- the maximum packet size that the Wi-Fi driver can receive: wifi_rx_pkt_size_max
- the maximum packet size that the Wi-Fi driver can send: wifi_tx_pkt_size_max

So, the peak memory that the Wi-Fi driver consumes can be calculated with the following formula:

wifi_dynamic_peek_memory = (wifi_rx_dynamic_buf_num * wifi_rx_pkt_size_max) + (wifi_tx_dynamic_buf_num * wifi_tx_pkt_size_max)

Generally, we do not need to care about the dynamic tx long buffers and dynamic tx long long buffers, because they are management frames which only have a small impact on the system.

4.21.30 Wi-Fi Menuconfig

Wi-Fi Buffer Configure

If you are going to modify the default number or type of buffer, it would be helpful to also have an overview of how the buffer is allocated/freed in the data path. The following diagram shows this process in the TX direction:

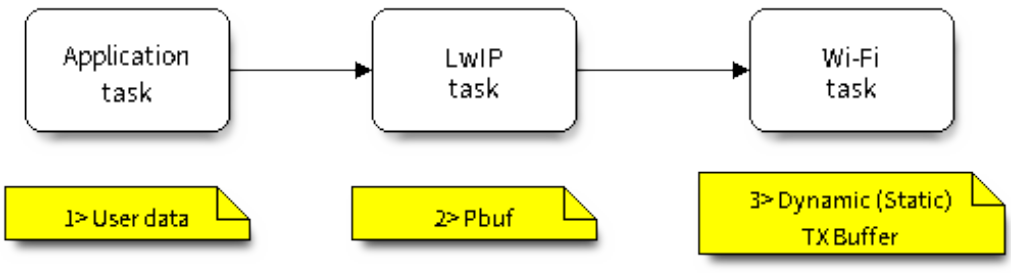


图 32: TX Buffer Allocation

Description:

- The application allocates the data which needs to be sent out.
- The application calls TCPIP-/Socket-related APIs to send the user data. These APIs will allocate a PBUF used in LwIP, and make a copy of the user data.
- When LwIP calls a Wi-Fi API to send the PBUF, the Wi-Fi API will allocate a "Dynamic Tx Buffer" or "Static Tx Buffer", make a copy of the LwIP PBUF, and finally send the data.

The following diagram shows how buffer is allocated/freed in the RX direction:

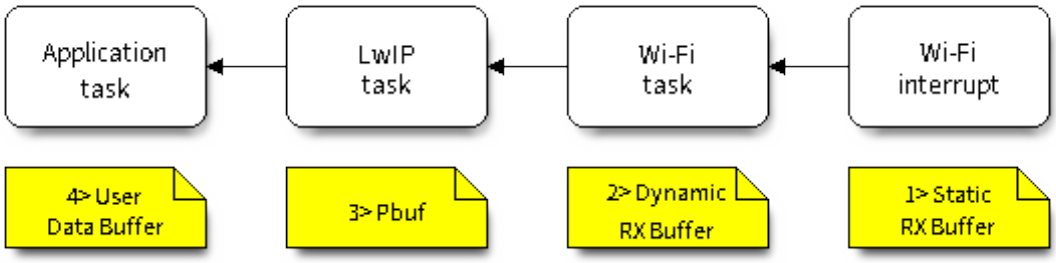


图 33: RX Buffer Allocation

Description:

- The Wi-Fi hardware receives a packet over the air and puts the packet content to the "Static Rx Buffer", which is also called "RX DMA Buffer".
- The Wi-Fi driver allocates a "Dynamic Rx Buffer", makes a copy of the "Static Rx Buffer", and returns the "Static Rx Buffer" to hardware.
- The Wi-Fi driver delivers the packet to the upper-layer (LwIP), and allocates a PBUF for holding the "Dynamic Rx Buffer".
- The application receives data from LwIP.

The diagram shows the configuration of the Wi-Fi internal buffer.

Buffer Type	Alloc Type	Default	Configurable	Description
Static RX Buffer (Hardware RX Buffer)	Static	10 * 1600 Bytes	Yes	This is a kind of DMA memory. It is initialized in esp_wifi_init() and freed in esp_wifi_deinit(). The 'Static Rx Buffer' forms the hardware receiving list. Upon receiving a frame over the air, hardware writes the frame into the buffer and raises an interrupt to the CPU. Then, the Wi-Fi driver reads the content from the buffer and returns the buffer back to the list. If the application want to reduce the the memory statically allocated by Wi-Fi, they can reduce this value from 10 to 6 to save 6400 Bytes memory. It's not recommended to reduce the configuration to a value less than 6 unless the AMPDU feature is disabled.
Dynamic RX Buffer	Dynamic	32	Yes	The buffer length is variable and it depends on the received frames' length. When the Wi-Fi driver receives a frame from the 'Hardware Rx Buffer', the 'Dynamic Rx Buffer' needs to be allocated from the heap. The number of the Dynamic Rx Buffer, configured in the menuconfig, is used to limit the total un-freed Dynamic Rx Buffer number.
Dynamic TX Buffer	Dynamic	32	Yes	This is a kind of DMA memory. It is allocated to the heap. When the upper- layer (LwIP) sends packets to the Wi-Fi driver, it firstly allocates

Wi-Fi NVS Flash

If the Wi-Fi NVS flash is enabled, all Wi-Fi configurations set via the Wi-Fi APIs will be stored into flash, and the Wi-Fi driver will start up with these configurations next time it powers on/reboots. However, the application can choose to disable the Wi-Fi NVS flash if it does not need to store the configurations into persistent memory, or has its own persistent storage, or simply due to debugging reasons, etc.

Wi-Fi AMPDU

ESP32-S2 supports both receiving and transmitting AMPDU, the AMPDU can greatly improve the Wi-Fi throughput.

Generally, the AMPDU should be enabled. Disabling AMPDU is usually for debugging purposes.

4.21.31 Troubleshooting

Please refer to a separate document with *乐鑫 Wireshark 使用指南*.

乐鑫 Wireshark 使用指南

1. 概述

1.1 什么是 Wireshark？ Wireshark（原称 Ethereal）是一个网络封包分析软件。网络封包分析软件的功能是撷取网络封包，并尽可能显示出最为详细的网络封包资料。Wireshark 使用 WinPCAP 作为接口，直接与网卡进行数据报文交换。

网络封包分析软件的功能可想像成“电工技师使用电表来量测电流、电压、电阻”的工作，只是将场景移植到网络上，并将电线替换成网线。

在过去，网络封包分析软件是非常昂贵，或是专门属于营利用的软件。Wireshark 的出现改变了这一切。

在 GNU GPL 通用许可证的保障范围下，使用者可以以免费的代价取得软件与其源代码，并拥有针对其源代码修改及客制化的权利。

Wireshark 是目前全世界最广泛的网络封包分析软件之一。

1.2 Wireshark 的主要应用 下面是 Wireshark 一些应用的举例：

- 网络管理员用来解决网络问题
- 网络安全工程师用来检测安全隐患
- 开发人员用来测试协议执行情况
- 用来学习网络协议

除了上面提到的，Wireshark 还可以用在其它许多场合。

1.3 Wireshark 的特性

- 支持 UNIX 和 Windows 平台
- 在接口实时捕捉包
- 能详细显示包的详细协议信息
- 可以打开/保存捕捉的包
- 可以导入导出其他捕捉程序支持的包数据格式
- 可以通过多种方式过滤包
- 多种方式查找包
- 通过过滤以多种色彩显示包
- 创建多种统计分析
- 等等

1.4 Wireshark 的“能”与“不能”？

- **捕捉多种网络接口**
  Wireshark 可以捕捉多种网络接口类型的包，哪怕是无线局域网接口。
- **支持多种其它程序捕捉的文件**
  Wireshark 可以打开多种网络分析软件捕捉的包。
- **支持多格式输出**
  Wireshark 可以将捕捉文件输出为多种其他捕捉软件支持的格式。
- **对多种协议解码提供支持**
  Wireshark 可以支持许多协议的解码。
- **Wireshark 不是入侵检测系统**
  如果您的网络中存在任何可疑活动，Wireshark 并不会主动发出警告。不过，当您希望对这些可疑活动一探究竟时，Wireshark 可以发挥作用。
- **Wireshark 不会处理网络事务，它仅仅是“测量”（监视）网络**
  Wireshark 不会发送网络包或做其它交互性的事情（名称解析除外，但您也可以禁止解析）。

2. 如何获取 Wireshark 官网链接：https://www.wireshark.org/download.html

Wireshark 支持多种操作系统，请在下载安装文件时，注意选择与您所用操作系统匹配的安装文件。

3. 使用步骤 **本文档仅以 Linux 系统下的 Wireshark（版本号：2.2.6）为例。**

1) 启动 Wireshark

Linux 下，可编写一个 Shell 脚本，运行该文件即可启动 Wireshark 配置抓包网卡和信道。Shell 脚本如下：

```
ifconfig $1 down
iwconfig $1 mode monitor
iwconfig $1 channel $2
ifconfig $1 up
Wireshark&
```

脚本中有两个参数：`$1` 和 `$2`，分别表示网卡和信道，例如，`./xxx.sh wlan0 6`（此处，`wlan0` 即为抓包使用的网卡，后面的数字 `6` 即为 AP 或 soft-AP 所在的 channel）。

2) 运行 Shell 脚本打开 Wireshark，会出现 Wireshark 抓包开始界面

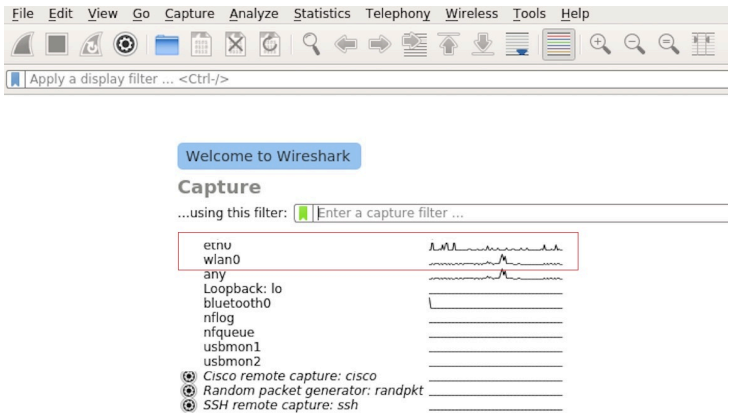


图 34: Wireshark 抓包界面

3) 选择接口，开始抓包

从上图红色框中可以看到有多个接口，第一个为本地网卡，第二个为无线网络。

可根据自己的需求选取相应的网卡，本文是以利用无线网卡抓取空中包为例进行简单说明。

双击 *wlan0* 即可开始抓包。

4) 设置过滤条件

抓包过程中会抓取到同信道所有的空中包，但其实很多都是我们不需要的，因此很多时候我们会设置抓包的过滤条件从而得到我们想要的包。

下图中红色框内即为设置 filter 的位置。

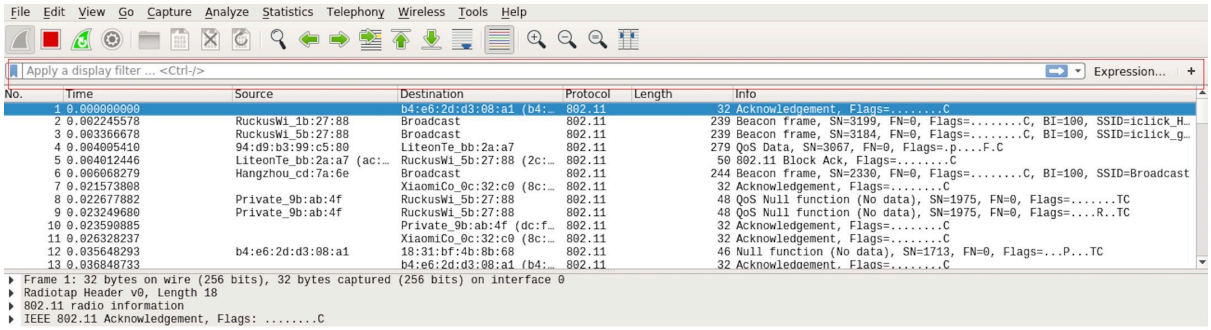


图 35: 设置 Wireshark 过滤条件

点击 *Filter* 按钮（下图的左上角蓝色按钮）会弹出 *display filter* 对话框。

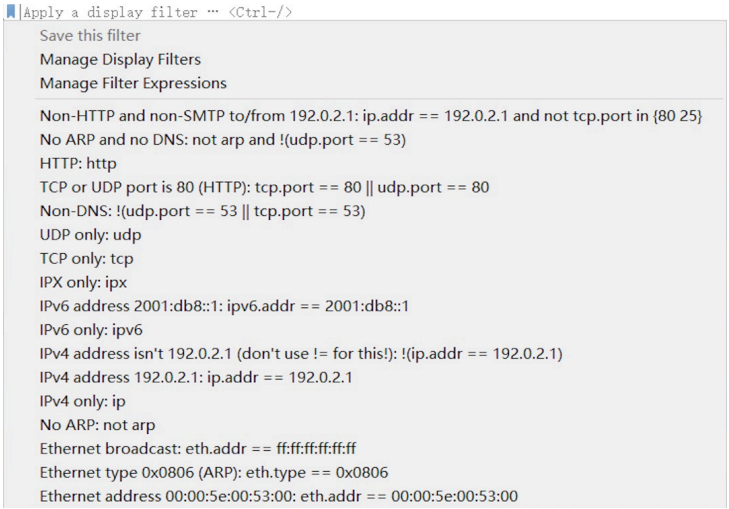


图 36: *Display Filter* 对话框

点击 *Expression* 按钮，会出现 *Filter Expression* 对话框，在此你可以根据需求进行 filter 的设置。

最直接的方法：直接在工具栏上输入过滤条件。

点击在此区域输入或修改显示的过滤字符，在输入过程中会进行语法检查。如果您输入的格式不正确，或者未输入完成，则背景显示为红色。直到您输入合法的表达式，背景会变为绿色。你可以点击下拉列表选择您先前键入的过滤字符。列表会一直保留，即使您重新启动程序。

例如：下图所示，直接输入 2 个 MAC 作为过滤条件，点击 *Apply*（即图中的蓝色箭头），则表示只抓取 2 个此 MAC 地址之间的交互的包。

5) 封包列表

若想查看包的具体的信息只需要选中要查看的包，在界面的下方会显示出包的具体的格式和包的内容。

如上图所示，我要查看第 1 个包，选中此包，图中红色框中即为包的具体内容。

6) 停止/开始包的捕捉

若要停止当前抓包，点击下图的红色按钮即可。

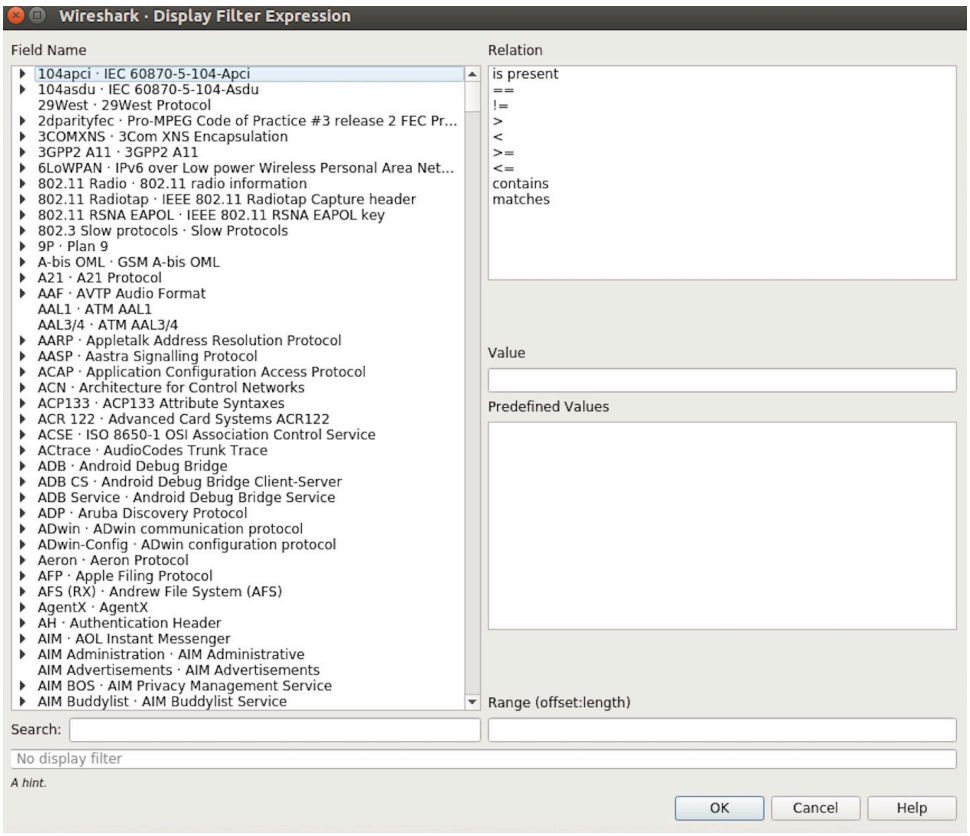


图 37: *Filter Expression* 对话框

图 38: 过滤条件工具栏

图 39: 在过滤条件工具栏中运用 MAC 地址过滤示例

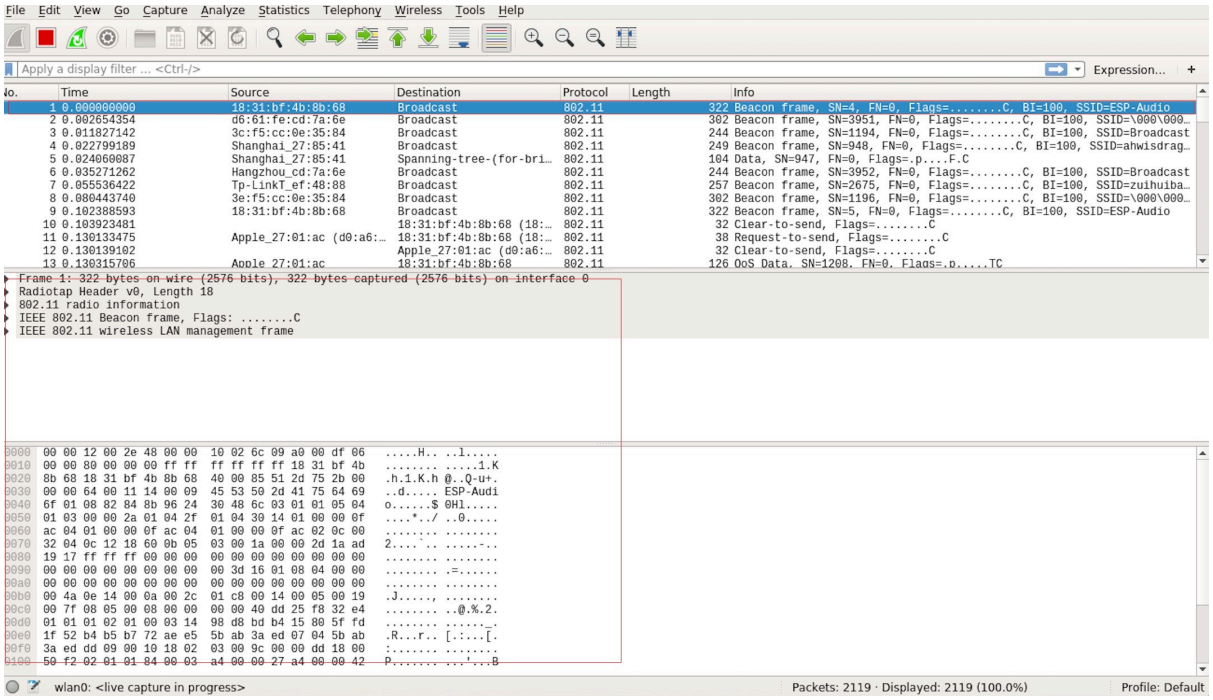


图 40: 封包列表具体信息示例

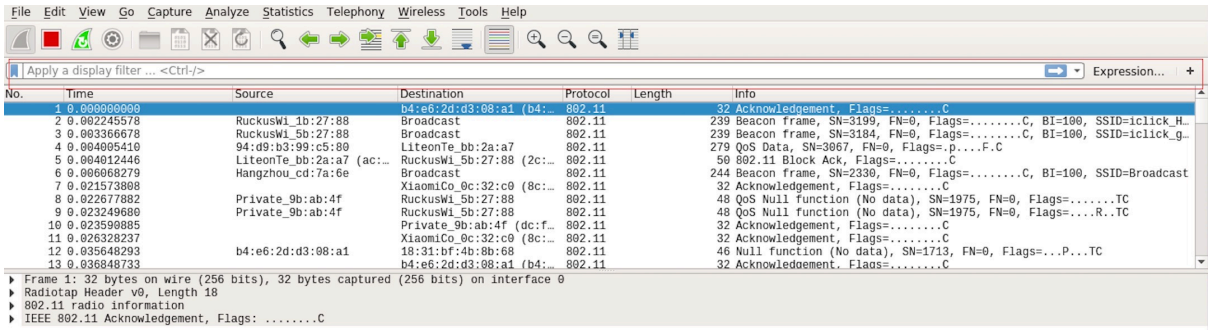


图 41: 停止包的捕捉

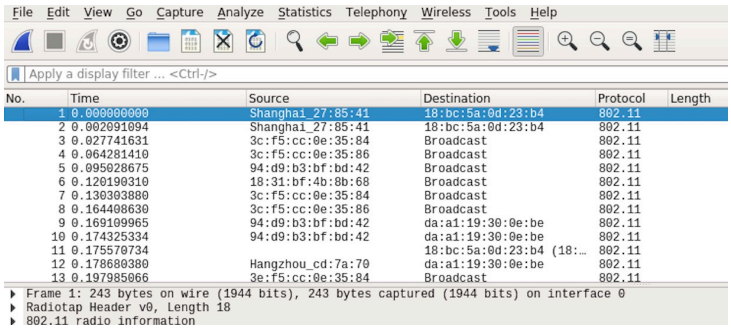


图 42: 开始或继续包的捕捉

若要重新开始抓包，点击下图左上角的蓝色按钮即可。

7) 保存当前捕捉包

Linux 下，可以通过依次点击“File” -> “Export Packet Dissections” -> “As Plain Text File”进行保存。

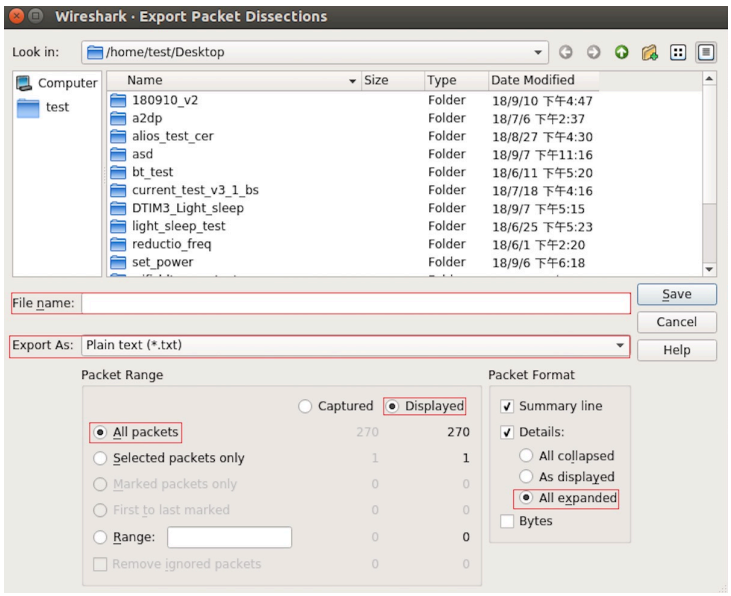


图 43: 保存捕捉包

上图中，需要注意的是，选择 *All packets*、*Displayed* 以及 *All expanded* 三项。

Wireshark 捕捉的包可以保存为其原生格式文件（libpcap），也可以保存为其他格式（如.txt 文件）供其他工具进行读取分析。

4.22 ESP-MESH

本指南提供有关 ESP-MESH 协议的介绍。更多有关 API 使用的信息，请见*MESH API 参考*。

4.22.1 概述

ESP-MESH 是一套建立在 Wi-Fi 协议之上的网络协议。ESP-MESH 允许分布在大范围区域内（室内和室外）的大量设备（下文称节点）在同一个 WLAN（无线局域网）中相互连接。ESP-MESH 具有自组网和自修复的特性，也就是说 mesh 网络可以自主地构建和维护。

本 ESP-MESH 指南分为以下几个部分：

1. 简介
2. *ESP-MESH* 概念
3. 建立网络
4. 管理网络
5. 数据传输
6. 信道切换
7. 性能
8. 更多注意事项

4.22.2 简介

传统基础设施 Wi-Fi 网络是一个“单点对多点”的网络。这种网络架构的中心节点为接入点 (AP)，其他节点 (station) 均与 AP 直接相连。其中，AP 负责各个 station 之间的仲裁和转发，一些 AP 还会通过路由

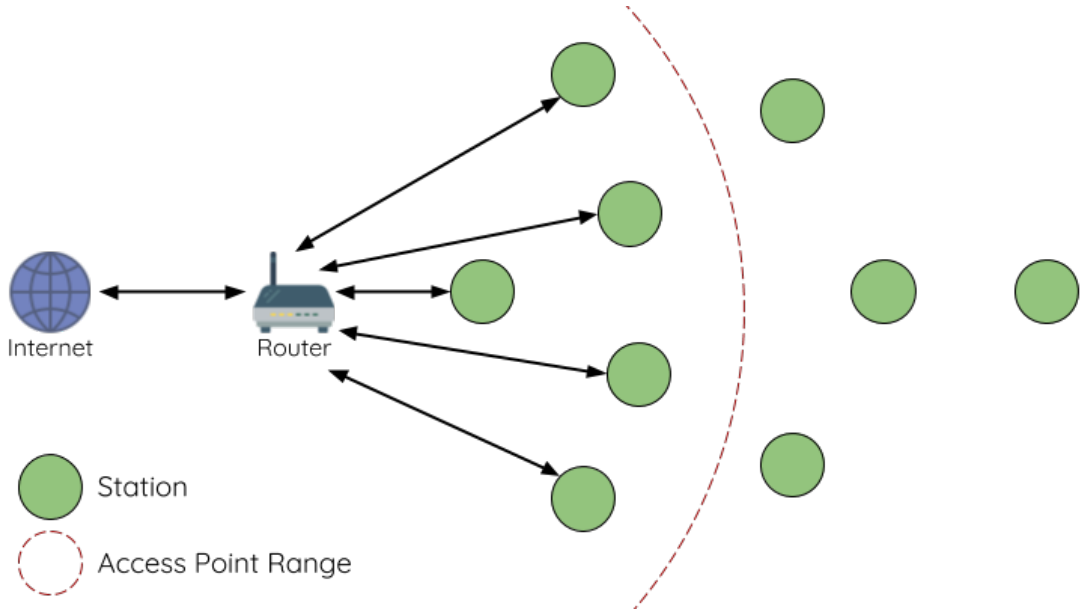


图 44: 传统 Wi-Fi 网络架构

器与外部 IP 网络交换数据。在传统 Wi-Fi 网络架构中，1）由于所有 station 均需与 AP 直接相连，不能距离 AP 太远，因此覆盖区域相对有限；2）受到 AP 容量的限制，因此网络中允许的 station 数量相对有限，很容易超载。

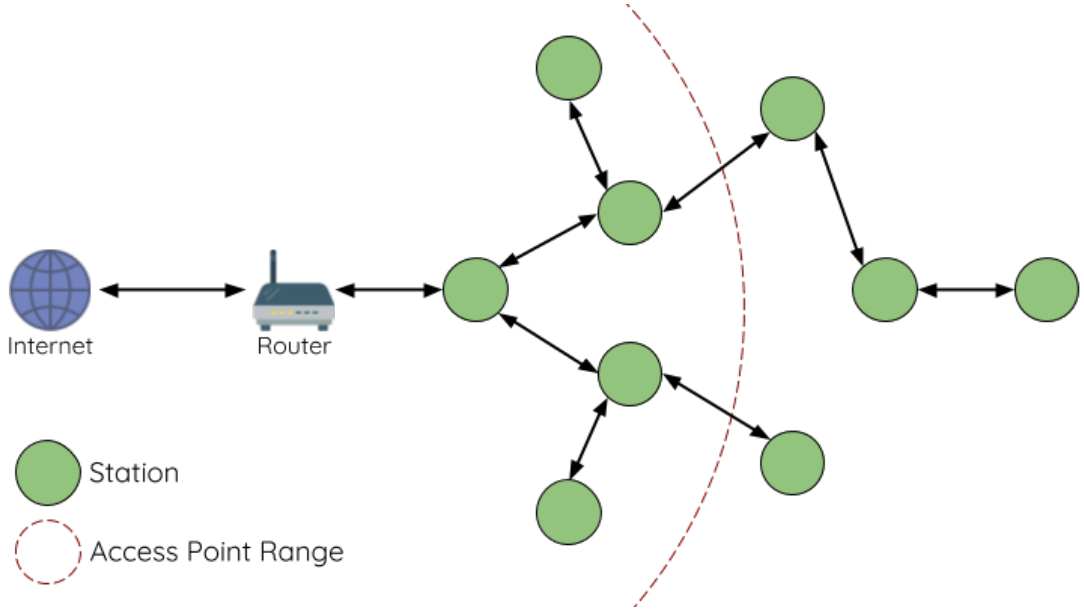


图 45: ESP-MESH 网络架构示意图

ESP-MESH 与传统 Wi-Fi 网络的不同之处在于：网络中的节点不需要连接到中心节点，而是可以与相邻节点连接。各节点均负责相连节点的数据中继。由于无需受限于距离中心节点的位置，所有节点仍可互连，因此 ESP-MESH 网络的覆盖区域更广。类似地，由于不再受限于中心节点的容量限制，ESP-MESH 允许更多节点接入，也不易于超载。

4.22.3 ESP-MESH 概念

术语

术语	描述
节点	任何 **属于**或 **可以成为** ESP-MESH 网络一部分的设备
根节点	网络顶部的节点
子节点	如节点 X 连接至节点 Y，且 X 相较 Y 与根节点的距离更远（跨越的连接数量更多），则称 X 为 Y 的子节点。
父节点	与子节点对应的概念
后裔节点	任何可以从根节点追溯到的节点
兄弟节点	连接至同一个父节点的所有节点
连接	AP 和 station 之间的传统 Wi-Fi 关联。ESP-MESH 中的节点使用 station 接口与另一个节点的 SoftAP 接口产生关联，进而形成连接。连接包括 Wi-Fi 网络中的身份验证和关联过程。
上行连接	从节点到其父节点的连接
下行连接	从父节点到其一个子节点的连接
无线 hop	源节点和目标节点间无线连接路径中的一部分。**单跳**指遍历单个连接的数据包，**多跳**指遍历多个连接的数据包。
子网	子网指 ESP-MESH 网络的一部分，包括一个节点及其所有后代节点。因此，根节点的子网包括 ESP-MESH 网络中的所有节点。
MAC 地址	在 ESP-MESH 网络中用于区别每个节点或路由器的唯一地址
DS	分布式系统（外部 IP 网络）

树型拓扑

ESP-MESH 建立在传统 Wi-Fi 协议之上，可被视为一种将多个独立 Wi-Fi 网络组合为一个单一 WLAN 网络的组网协议。在 Wi-Fi 网络中，station 在任何时候都仅限于与 AP 建立单个连接（上行连接），而 AP 则可以同时连接到多个 station（下行连接）。然而，ESP-MESH 网络则允许节点同时充当 station 和 AP。因此，ESP-MESH 中的节点可以使用 **其 SoftAP 接口建立多个下行连接**，同时使用 **其 station 接口建立一个上行连接**。这将自然产生一个由多层父子结构组成的树型网络拓扑结构。

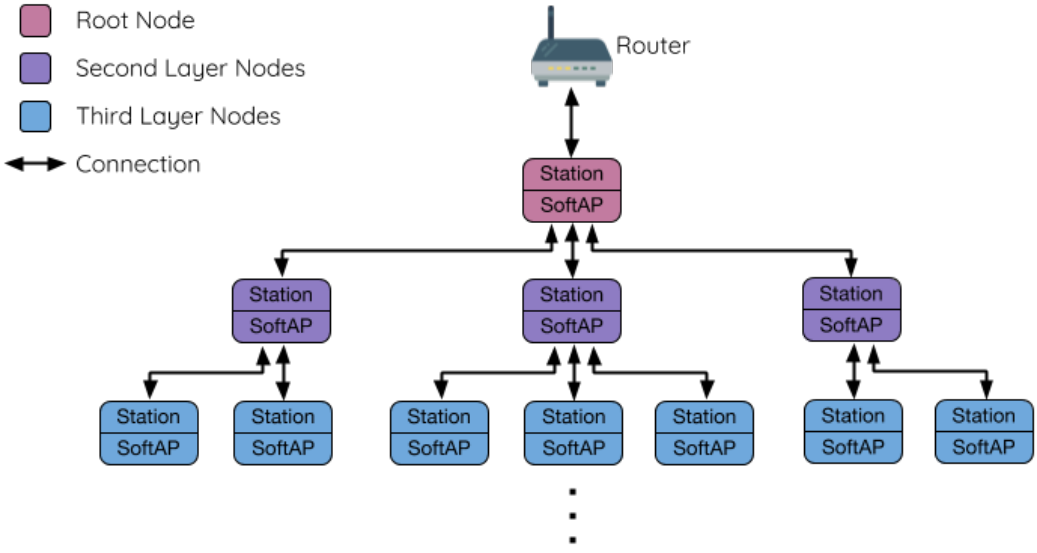


图 46: ESP-MESH 树型拓扑

ESP-MESH 是一个多跳网络，也就是说网络中的节点可以通过单跳或多跳向网络中的其他节点传送数据包。因此，ESP-MESH 中的节点不仅传输自己的数据包，而且同时充当其他节点的中继。假设 ESP-MESH 网络中的任意两个节点存在物理层上连接（通过单跳或多跳），则这两个节点可以进行通信。

注解： ESP-MESH 网络中的大小（节点总数）取决于网络中允许的最大层级，以及每个节点可以具有的最大下行连接数。因此，这两个变量可用于配置 ESP-MESH 网络的大小。

节点类型

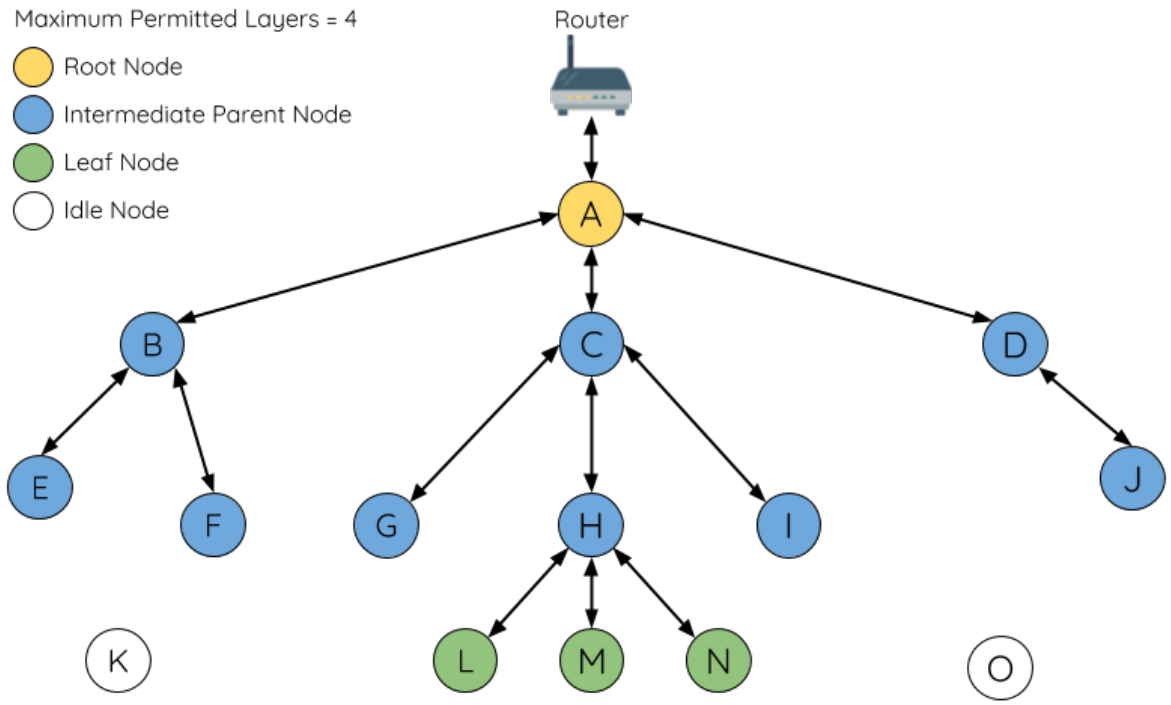


图 47: ESP-MESH 节点类型

根节点：指网络顶部的节点，是 ESP-MESH 网络和外部 IP 网络之间的唯一接口。根节点直接连接至传统的 Wi-Fi 路由器，并在 ESP-MESH 网络的节点和外部 IP 网络之间中继数据包。**ESP-MESH 网络中只能有一个根节点**，且根节点的上行连接只能是路由器。如上图所示，节点 A 即为该 ESP-MESH 网络的根节点。

叶子节点：指不允许拥有任何子节点（即无下行连接）的节点。因此，叶子节点只能传输或接收自己的数据包，但不能转发其他节点的数据包。如果节点处于 ESP-MESH 网络的最大允许层级，则该节点将成为叶子节点。叶子节点不回再产生下行连接，这可以防止节点继续生成下行连接，从而确保网络层级不会超出限制。由于建立下行连接必须使用 SoftAP 接口，因此一些没有 SoftAP 接口的节点（仅有 station 接口）也将被分配为叶子节点。如上图所示，位于网络最外层的 L/M/N 节点即为叶子节点。

中间父节点：既不是属于根节点也不属于叶子节点的节点即为中间父节点。中间父节点必须有且仅有一个上行连接（即一个父节点），但可以具有 0 个或多个下行连接（即 0 个或多个子节点）。因此，中间父节点可以发送和接收自己的数据包，也可以转发其上行和下行连接的数据包。如上图所示，节点 B 到 J 即为中间父节点。**注意，E/F/G/I/J 等没有下行连接的中间父节点并不等同于叶子节点**，原因在于这些节点仍允许形成下行连接。

空闲节点：尚未加入网络的节点即为空闲节点。空闲节点将尝试与中间父节点形成上行连接，或者在有条件的情况下（参见自动根节点选择）成为一个根节点。如上图所示，K 和 O 节点即为空闲节点。

信标帧和 RSSI 阈值

ESP-MESH 中能够形成下行连接的每个节点（即具有 SoftAP 接口）都会定期传输 Wi-Fi 信标帧。节点可以通过信标帧让其他节点检测自己的存在和状态。空闲节点将侦听信标帧以生成一个潜在父节点列表，

并与其中一个潜在父节点形成上行连接。ESP-MESH 使用“供应商信息元素”来存储元数据，例如：

- 节点类型（根节点、中间父节点、叶子节点、空闲节点）
- 节点当前所处的层级
- 网络中允许的最大层级
- 当前子节点数量
- 可接受的最大下行连接数量

潜在上行连接的信号强度可由潜在父节点信标帧的 RSSI 表示。为了防止节点形成弱上行连接，ESP-MESH 采用了针对信标帧的 RSSI 阈值控制机制。如果节点检测到某节点的信标帧 RSSI 过低（即低于预设阈值），则会在尝试形成上行连接时忽略该节点。

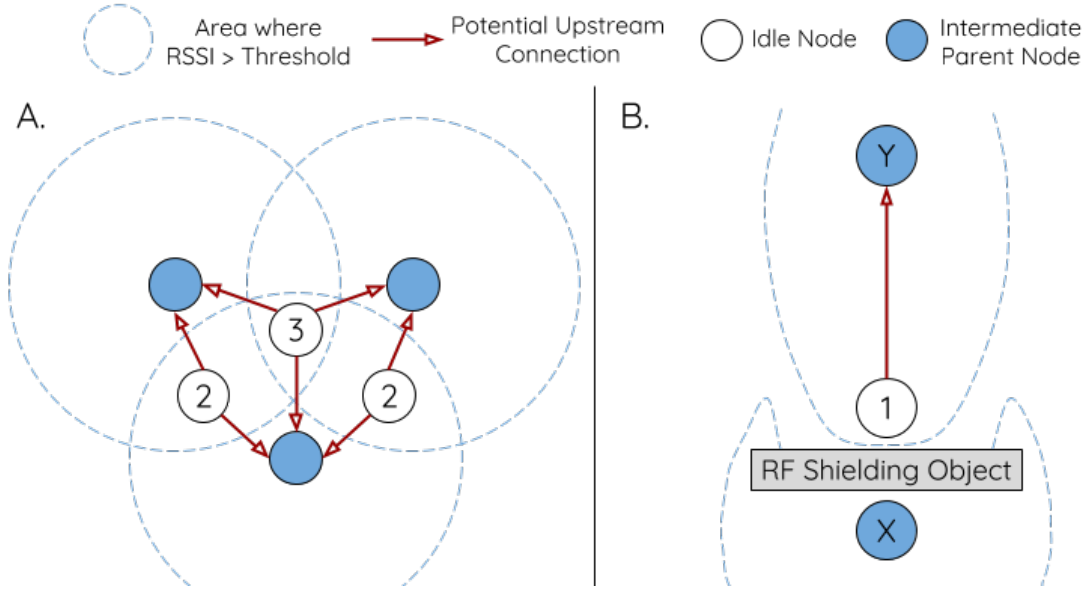


图 48: RSSI 阈值的影响

上图（A 侧） 展示了 RSSI 阈值将如何影响空闲节点的候选父节点数量。

上图（B 侧） 展示了 RF 屏蔽物将如何降低潜在父节点的 RSSI。由于存在 RF 屏蔽物，节点 X 的 RSSI 高于阈值的区域显著减小。这会导致空闲节点忽略节点 X，即使从地理位置上看 X 就在空闲节点附近。相反，该空闲节点将从更远的地方找到一个 RSSI 更强的节点 Y 形成上行连接。

注解： 事实上，ESP-MESH 网络中的节点在 MAC 层仍可以接收所有的信标帧，但 RSSI 阈值控制功能可以过滤掉所有 RSSI 低于预设阈值的信标帧。

首选父节点

当一个空闲节点有多个候选父节点（潜在父节点）时，空闲节点将与其中的 **首选父节点**形成上行连接。首选父节点基于以下条件确定：

- 候选父节点所处的层级
- 候选父节点当前具有的下行连接（子节点）数量

在网络中所处层级较浅的候选父节点（包括根节点）将优先成为首选父节点。这有助于在形成上行连接时控制 ESP-MESH 网络中的总层级使之最小。例如，在位于第二层和第三层的候选父节点间选择时，位于第二层的候选父节点将始终优先成为首选父节点。

如果同一层上存在多个候选父节点，则子节点最少的候选父节点将优先成为首选父节点。这有助于平衡同一层节点的下行连接数量。

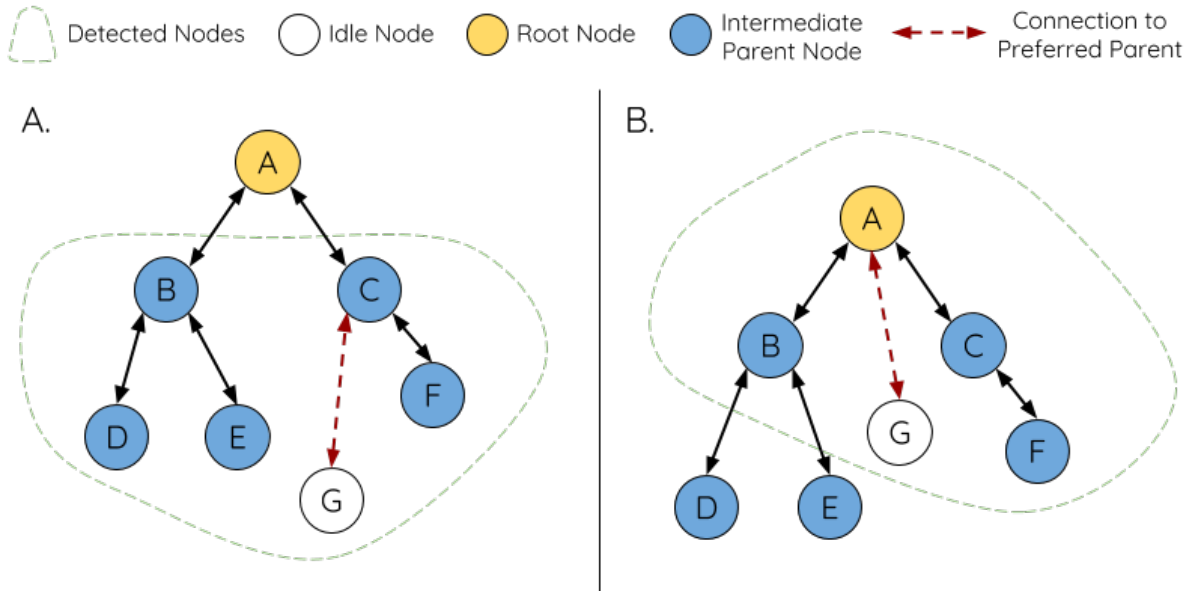


图 49: 首选父节点选择

上图（A 侧） 展示了空闲节点 G 如何在 B/C/D/E/F 五个候选父节点中选择首选父节点：首先，B/C 节点优于 D/E/F 节点，因为这两个节点所处的层级更浅。其次，C 节点优于 B 节点，因为 C 节点的下行连接数量（子节点数量）更少。

上图（B 侧） 展示了空闲节点 G 如何在根节点 A 和其他候选父节点中选择首选父节点，此时根节点 A 处于空闲节点 G 范围之内（即空闲节点 G 接收到的根节点 A 信标帧 RSSI 强度高于预设阈值）：由于根节点 A 处于网络中最浅的层，因此将成为首选父节点。

注解： 用户还可以自行定义首选父节点的选择规则，也可以直接指定某个节点为首选父节点（见 Mesh 手动配网示例 ）。

路由表

ESP-MESH 网络中的每个节点均会维护自己的路由表，并按路由表将数据包（请见*ESP-MESH 数据包*）沿正确的路线发送至正确的目标节点。某个特定节点的路由表将包含 **该节点的子网中所有节点的 MAC 地址**，也包括该节点自己的 MAC 地址。每个路由表会划分为多个子路由表，与每个子节点的子网对应。

以上图为例，节点 B 的路由表中将包含节点 B 到节点 I 的 MAC 地址（即相当于节点 B 的子网）。节点 B 的路由表可划分为节点 C 和 G 的子路由表，分别包含节点 C 到节点 F 的 MAC 地址、节点 G 到节点 I 的 MAC 地址。

ESP-MESH 利用路由表来使用以下规则进行转发，确定 ESP-MESH 数据包应根据向上行转发还是向下行转发。

1. 如果数据包的目标 MAC 地址处于当前节点的路由表中且不是当前节点本身，则选择包含目标 MAC 地址的子路由表，并将数据包向下转发给子路由表对应的子节点。

2. 如果数据包的目标 MAC 地址不在当前节点的路由表内，则将数据包向上转发给当前节点的父节点，并重复执行该操作直至数据包达到目标地址。此步骤可重复至根节点（根节点包含整个网络的全部节点）。

注解： 用户可以通过调用`esp_mesh_get_routing_table()` 获取一个节点的路由表，调用`esp_mesh_get_routing_table_size()` 获取一个路由表的大小，也可通过调用`esp_mesh_get_subnet_nodes_list()` 获取某个子节点的子路由表，调

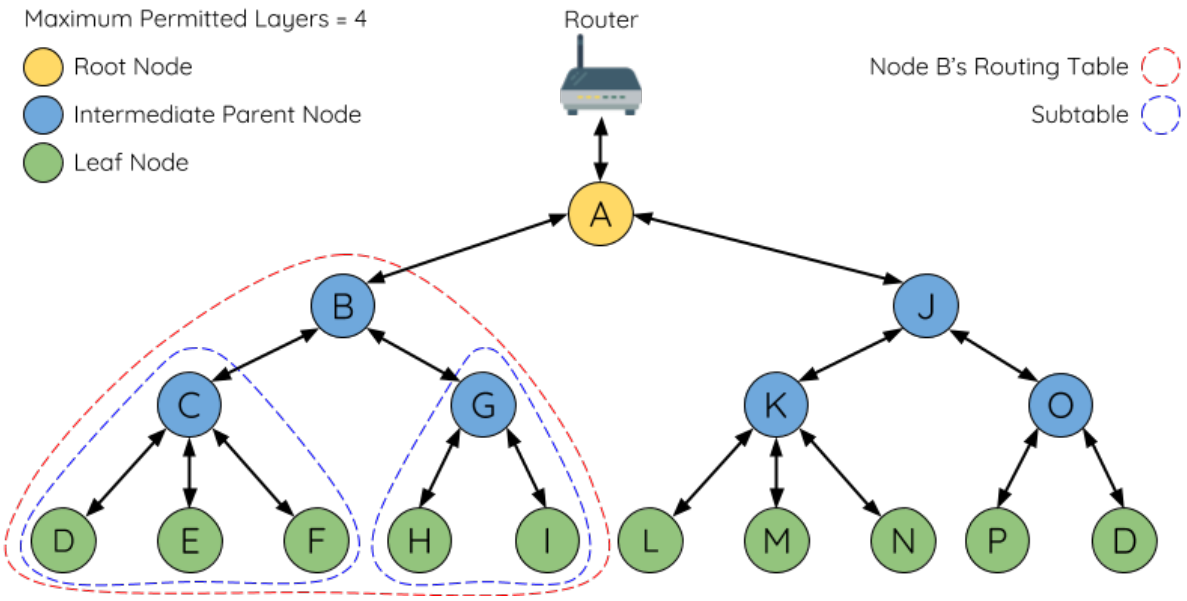


图 50: ESP-MESH 路由表示例

用*esp_mesh_get_subnet_nodes_num()* 获取子路由表的大小。

4.22.4 建立网络

一般过程

警告: ESP-MESH 正式开始构建网络前，必须确保网络中所有节点具有相同的配置（见*mesh_cfg_t*）。每个节点必须配置 **相同 MESH 网络 ID、路由器配置和 SoftAP 配置**。

ESP-MESH 网络将首先选择根节点，然后逐层形成下行连接，直到所有节点均加入网络。网络的布局可能取决于诸如根节点选择、父节点选择和异步上电复位等因素。但简单来说，一个 ESP-MESH 网络的构建过程可以概括为以下步骤：

1. 根节点选择 根节点直接进行指定（见*用户指定根节点*）或通过选举由信号强度最强的节点担任（见*自动根节点选择*）。一旦选定，根节点将与路由器连接，并开始允许下行连接形成。如上图所示，节点 A 被选为根节点，因此节点 A 上行连接到路由器。

2. 第二层形成 一旦根节点连接到路由器，根节点范围内的空闲节点将开始与根节点连接，从而形成第二层网络。一旦连接，第二层节点成为中间父节点（假设最大允许层级大于 2 层），并进而形成下一层。如上图所示，节点 B 到节点 D 都在根节点的连接范围内。因此，节点 B 到节点 D 将与根节点形成上行连接，并成为中间父节点。

3. 其余层形成 剩余的空闲节点将与所处范围内的中间父节点连接，并形成新的层。一旦连接，根据网络的最大允许层级，空闲节点成为中间父节点或叶子节点。此后重复该步骤，直到网络中的所有空闲节点均加入网络或达到网络最大允许层级。如上图所示，节点 E/F/G 分别与节点 B/C/D 连接，并成为中间父节点。

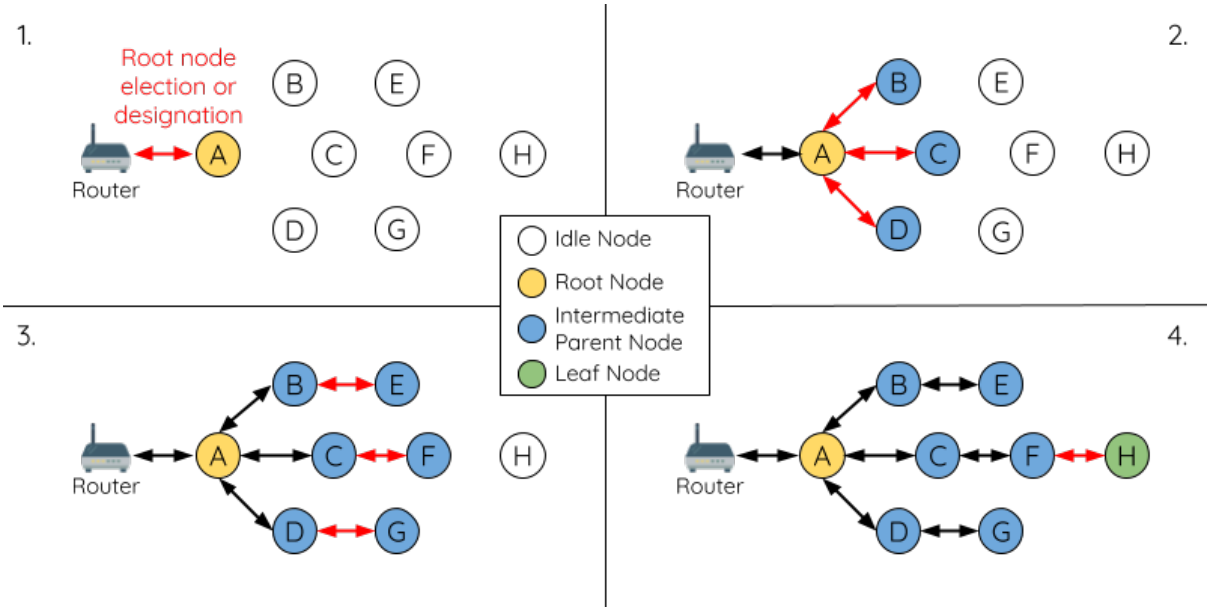


图 51: ESP-MESH 网络构建过程

4. 限制树深度 为了防止网络超过最大允许层级，最大允许层级上的节点将在完成连接后成为叶子节点。这样一来，其他空闲节点将无法与这些最大允许层上的叶子节点形成连接，因此不会超过最大允许层级。然而，如果空闲节点无法找到其他潜在父节点，则将无限期地保持空闲状态。如上图所示，网络的最大允许层级为四。因此，节点 H 在完成连接后将成为叶子节点，以防止任何下行连接的形成。

自动根节点选择

在自动模式下，根节点的选择取决于相对于路由器的信号强度。每个空闲节点将通过 Wi-Fi 信标帧发送自己的 MAC 地址和路由器 RSSI 值。**MAC 地址可以表示网络中的唯一节点**，而 **路由器 RSSI 值**代表相对于路由器的信号强度。

此后，每个节点将同时扫描来自其他空闲节点的信标帧。如果节点检测到具有更强的路由器 RSSI 的信标帧，则节点将开始传输该信标帧的内容（相当于为这个节点投票）。经过最小迭代次数（可预先设置，默认为 10 次）将选举出路由器 RSSI 值最强的信标帧。

在达到预设迭代次数后，每个节点将单独检查其 **得票百分比**（得票数/总票数）以确定它是否应该成为根节点。**如果节点的得票百分比大于预设的阈值（默认为 90%），则该节点将成为根节点**。

下图展示了在 ESP-MESH 网络中，根节点的自动选择过程。

1. 上电复位时，每个节点开始传输自己的信标帧（包括 MAC 地址和路由器 RSSI 值）。

2. 在多次传输和扫描迭代中，路由器 RSSI 最强的信标帧将在整个网络中传播。节点 C 具有最强的路由器 RSSI 值 (-10 dB)，因此它的信标帧将在整个网络中传播。所有参与选举的节点均给节点 C 投票，因此节点 C 的得票百分比为 100%。因此，节点 C 成为根节点，并与路由器连接。

3. 一旦节点 C 与路由器连接，节点 C 将成为节点 A/B/D/E 的首选父节点（即最浅的节点），并与这些节点连接。节点 A/B/D/E 将形成网络的第二层。

4. 节点 F 和节点 G 分别连接节点 D 和节点 E，并完成网络构建过程。

注解： 用户可以通过 `esp_mesh_set_attempts()` 配置选举的最小迭代次数。用户应根据网络内的节点数量配置迭代次数（即 mesh 网络越大，所需的迭代次数越高）。

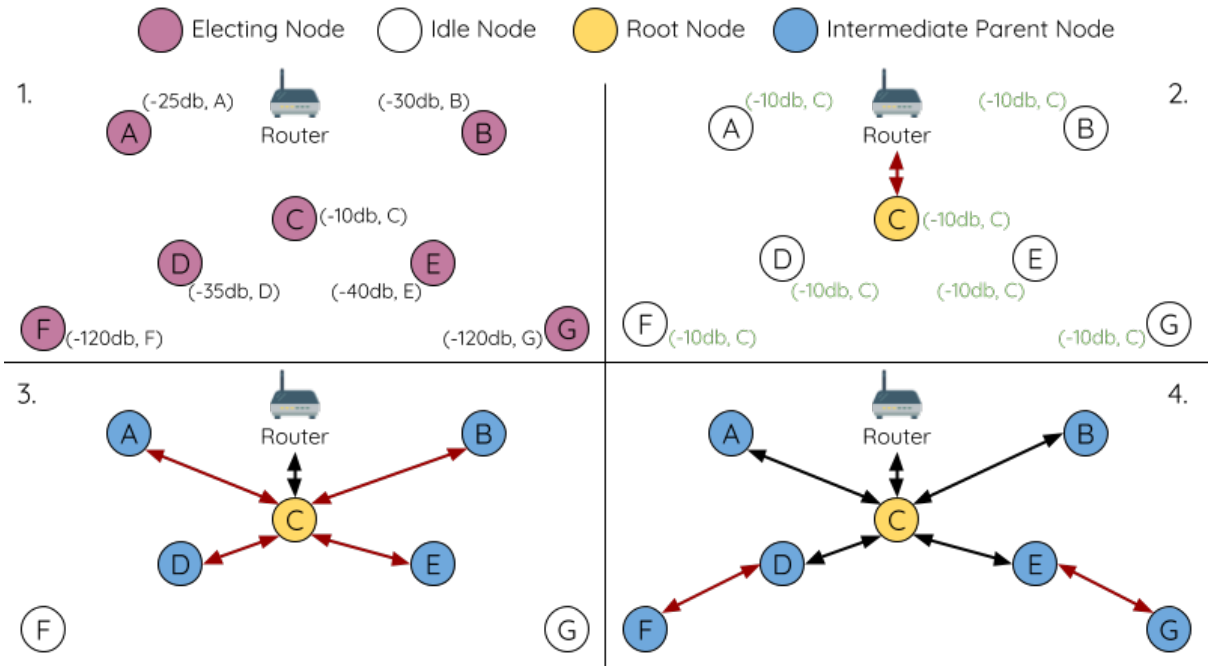


图 52: 根节点选举示例

警告： **得票百分比阈值**也可以使用`esp_mesh_set_vote_percentage()`进行配置。得票百分比阈值过低 **可能导致同一 mesh 网络中两个或多个节点成为根节点**，进而分化为多个 mesh 网络。如果发生这种情况，ESP-MESH 具有内部机制，可自主解决 **根节点冲突**。这些具有多个根节点的网络将围绕一个根节点形成一个网络。然而，两个或多个路由器 SSID 相同但路由器 BSSID 不同的根节点冲突尚无法解决。

用户指定根节点

根节点也可以由用户指定，即直接让指定的根节点与路由器连接，并放弃选举过程。当根节点指定后，网络内的所有其他节点也必须放弃选举过程，以防止根节点冲突的发生。下图展示了在 ESP-MESH 网络中，根节点的手动选择过程。

1. 节点 A 是由用户指定的根节点，因此直接与路由器连接。此时，所有其他节点放弃选举过程。

2. 节点 C 和节点 D 将节点 A 选为自己的首选父节点，并与其形成连接。这两个节点将形成网络的第二层。

3. 类似地，节点 B 和节点 E 将与节点 C 连接，节点 F 将与节点 D 连接。这三个节点将形成网络的第三层。

4. 节点 G 将与节点 E 连接，形成网络的第四层。然而，由于该网络的最大允许层级已配置为 4，因此节点 G 将成为叶子节点，以防止形成任何新层。

注解： 一旦指定根节点，该根节点应调用`esp_mesh_set_parent()`使其直接与路由器连接。类似地，所有其他节点都应该调用`esp_mesh_fix_root()`放弃选举过程。

选择父节点

默认情况下，ESP-MESH 具有可以自组网的特点，也就是每个节点都可以自主选择与其形成上行连接的潜在父节点。自主选择出的父节点被称为首选父节点。用于选择首选父节点的标准旨在减少 ESP-MESH 网络的层级，并平衡各个潜在父节点的下行连接数（参见首选父节点）。

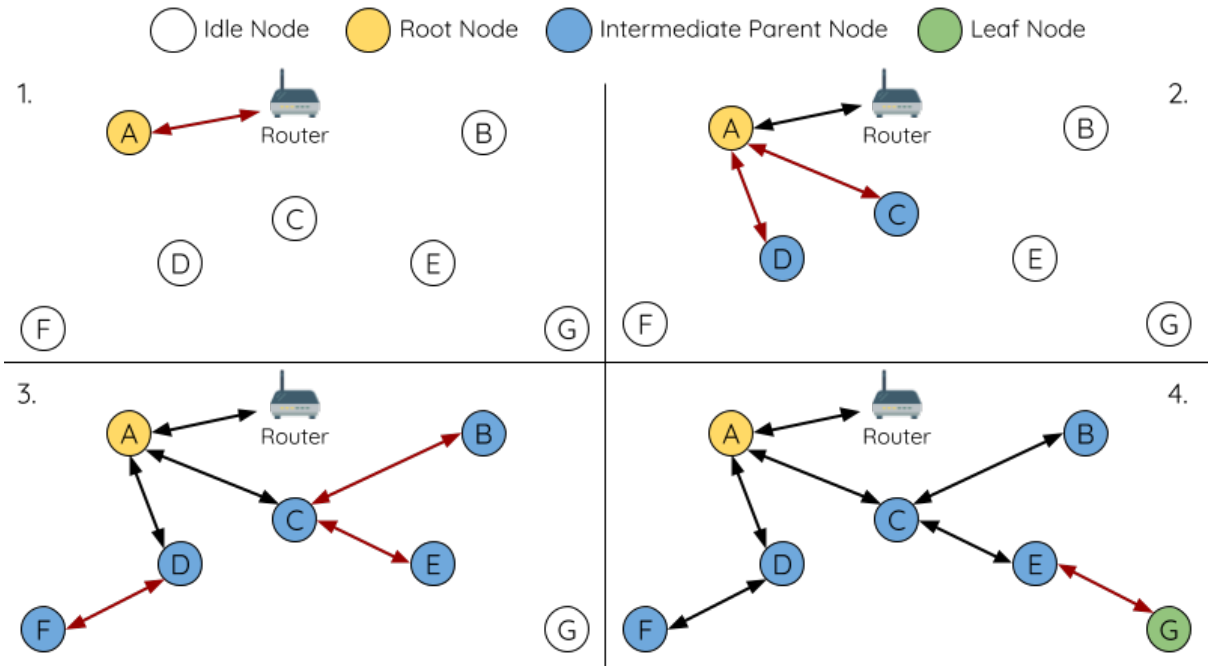


图 53: 根节点指定示例（根节点 = A，最大层级 = 4)

不过，ESP-MESH 也允许用户禁用自组网功能，即允许用户自己定义父节点选择标准，或直接指定某个节点为父节点（见：Mesh 手动组网示例 ）。

异步上电复位

ESP-MESH 网络构建可能会受到节点上电顺序的影响。如果网络中的某些节点为异步上电（即相隔几分钟上电），**网络的最终结构可能与所有节点同步上电时的理想情况不同**。延迟上电的节点将遵循以下规则：

规则 1：如果网络中已存在根节点，则延迟节点不会尝试选举成为新的根节点，即使自身的路由器 RSSI 更强。相反，延迟节点与任何其他空闲节点无异，将通过与首选父节点连接来加入网络。如果该延迟节点为用户指定的根节点，则网络中的所有其他节点将保持空闲状态，直到延迟节点完成上电。

规则 2：如果延迟节点形成上行连接，并成为中间父节点，则后续也可能成为其他节点（即其他更浅的节点）的新首选父节点。此时，其他节点切换上行连接至该延迟节点（见父节点切换）。

规则 3：如果空闲节点的指定父节点上电延迟了，则该空闲节点在没有找到指定父节点前不会尝试形成任何上行连接。空闲节点将无限期地保持空闲，直到其指定的父节点上电完成。

下方示例展示了异步上电对网络构建的影响。

1. 节点 A/C/D/F/G/H 同步上电，并通过广播其 MAC 地址和路由器 RSSI 开始选举根节点。节点 A 的 RSSI 最强，因此当选为根节点。

2. 一旦节点 A 成为根节点，其余的节点就开始与其首选父节点逐层形成上行连接，并最终形成一个具有五层的网络。

3. 节点 B/E 由于存在上电延迟，因此即使路由器 RSSI 比节点 A 更强（-20 dB 和 -10 dB）也不会尝试成为根节点。相反，这两个上电延迟节点均将与对应的首选父节点 A 和 C 形成上行连接。加入网络后，节点 B/E 均将成为中间父节点。

4. 节点 B 由于所处层级变化（现为第二层）而成为新的首选父节点，因此节点 D/G 将切换其上行连接从而选择新的首选父节点。由于切换的发生，最终的网络层级从原来的五层减少至三层。

同步上电：如果所有节点均同步上电，节点 E （-10 dB）由于路由器 RSSI 最强而成为根节点。此时形成的网络结构将与异步上电的情况截然不同。**但是，如果用户手动切换根节点，则仍可以达到同步上电的网络结构**（请见`esp_mesh_waive_root()`）。

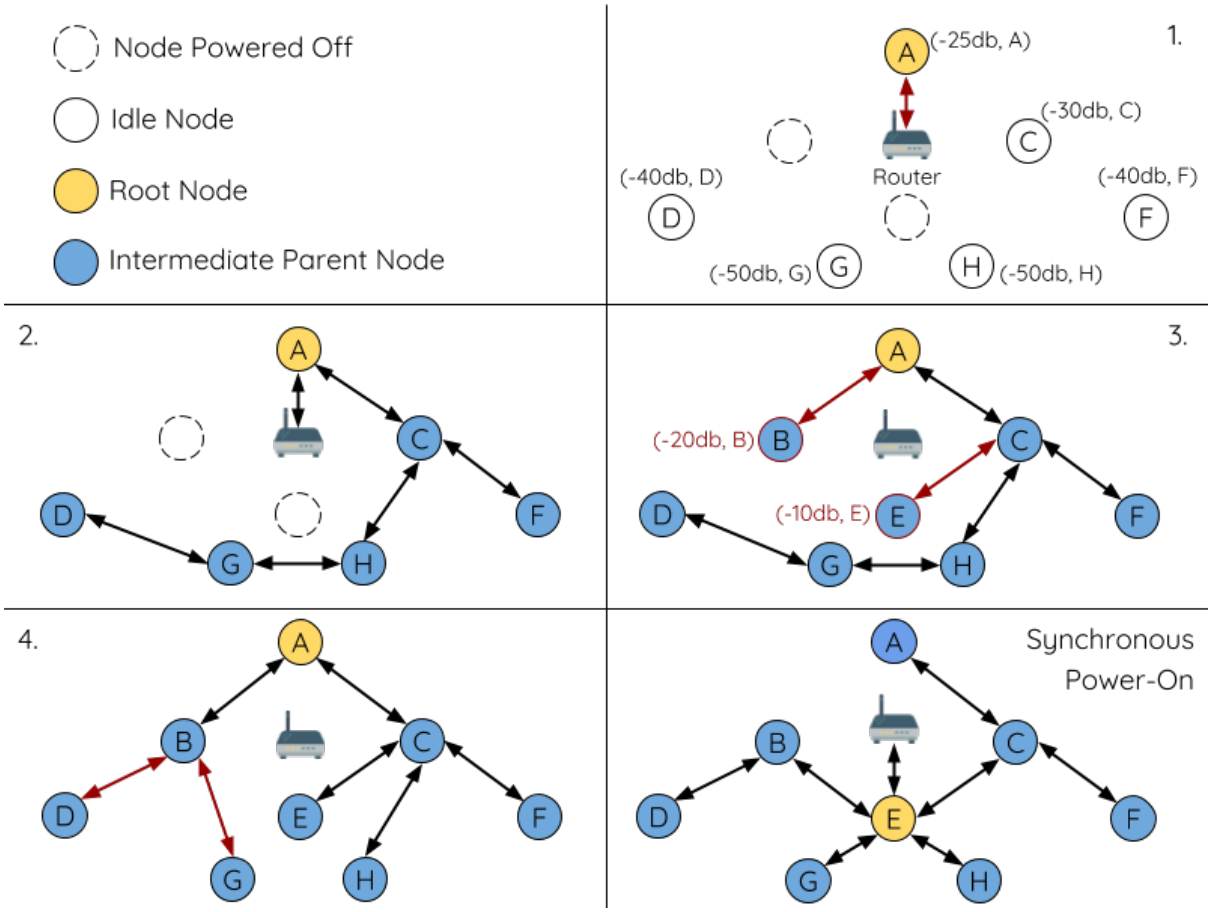


图 54: 网络构建（异步电源）示例

注解： 从某种程度上，ESP-MESH 可以自动修复部分因异步上电引起的父节点选择的偏差（请见父节点切换）

环路避免、检测和处理

环路是指特定节点与其后代节点（特定节点子网中的节点）形成上行连接的情况。因此产生的循环连接路径将打破 mesh 网络的树型拓扑结构。ESP-MESH 的节点在选择父节点时将主动排除路由表（见路由表）中的节点，从而避免与其子网中的节点建立上行连接并形成环路。

在存在环路的情况下，ESP-MESH 可利用路径验证机制和能量传递机制来检测环路的产生。因与子节点建立上行连接而导致环路形成的父节点将通知子节点环路的存在，并主动断开连接。

4.22.5 管理网络

作为一个自修复网络，ESP-MESH 可以检测并修正网络路由中的故障。当具有一个或多个子节点的父节点断开或父节点与其子节点之间的连接不稳定时，会发生故障。ESP-MESH 中的子节点将自主选择一个新的父节点，并与其形成上行连接，以维持网络互连。ESP-MESH 可以处理根节点故障和中间父节点故障。

根节点故障

如果根节点断开，则与其连接的节点（第二层节点）将及时检测到该根节点故障。第二层节点将主动尝试与根节点重连。但是在多次尝试失败后，第二层节点将启动新一轮的根节点选举。**第二层中 RSSI 最强的节点将当选为新的根节点**，而剩余的第二层节点将与新的根节点（如果不在范围内的话，也可与相邻父节点连接）形成上行连接。

如果根节点和下面多层的节点（例如根节点、第二层节点和第三层节点）同时断开，则位于最浅层的仍在正常工作的节点将发起根节点选举。下方示例展示了网络从根节点断开故障中进行自修复。

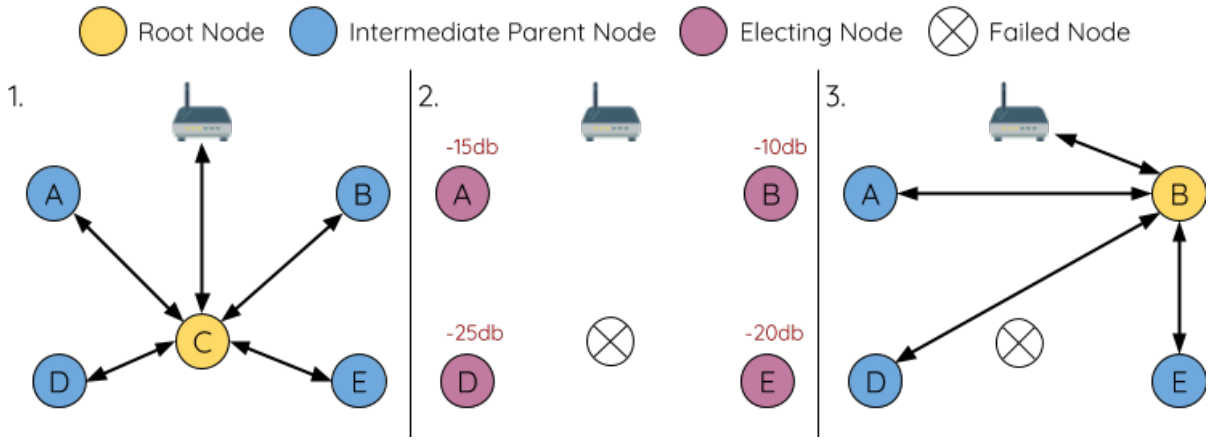


图 55: 根节点故障的自修复示意

1. 节点 C 是网络的根节点。节点 A/B/D/E 是连接到节点 C 的第二层节点。

2. 节点 C 断开。在多次重连尝试失败后，第二层节点开始通过广播其路由器 RSSI 开始新一轮的选举。此时，节点 B 的路由器 RSSI 最强。

3. 节点 B 被选为根节点，并开始接受下行连接。剩余的第二层节点 A/D/E 形成与节点 B 的上行连接，因此网络已经恢复，并且可以继续正常运行。

注解： 如果是手动指定的根节点断开，则无法进行自动修复。**任何节点不会在存在指定根节点的情况下开始选举过程**。

中间父节点故障

如果中间父节点断开，则与之断开的子节点将主动尝试与该父节点重连。在多次重连尝试失败后，每个子节点开始扫描潜在父节点（请见信标帧和 *RSSI* 阈值）。

如果存在其他可用的潜在父节点，每个子节点将分别给自己选择一个新的首选父节点（请见首选父节点），并与它形成上行连接。如果特定子节点没有其他潜在的父节点，则将无限期地保持空闲状态。

下方示例展示了网络从中间父节点断开故障中进行自修复。

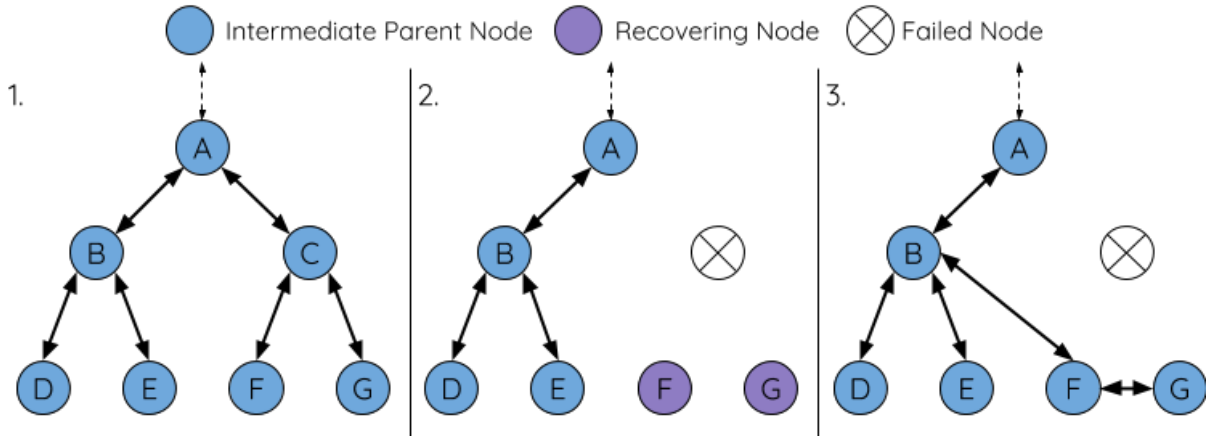


图 56: 中间父节点故障的自修复

1. 网络中存在节点 A 至 G。

2. 节点 C 断开。节点 F/G 检测到节点 C 的断开故障，并尝试与节点 C 重新连接。在多次重连尝试失败后，节点 F/G 将开始选择新的首选父节点。

3. 节点 G 因其范围内不存在任何父节点而暂时保持空闲。节点 F 的范围中有 B 和 E 两个节点，但节点 B 因为所处层级更浅而当选新的父节点。节点 F 将与节点 B 连接后，并成为一个中间父节点，节点 G 将于节点 F 相连。这样一来，网络已经恢复了，但结构发生了变化（网络层级增加了 1 层）。

注解： 如果子节点的父节点已被指定，则子节点不会尝试与其他潜在父节点连接。此时，该子节点将无限期地保持空闲状态。

根节点切换

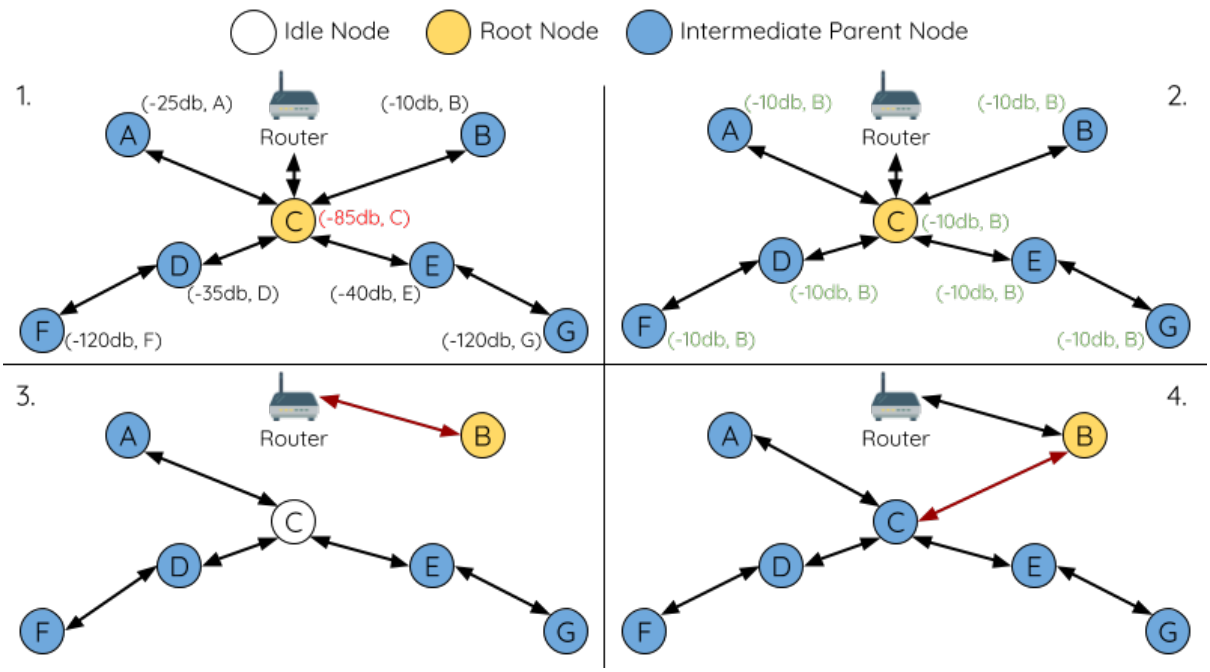
除非根节点断开，否则 ESP-MESH 不会自动切换根节点。即使根节点的路由器 RSSI 降低至必须断开的情况，根节点也将保持不变。根节点切换是指明确启动新选举过程的行为，即具有更强路由器 RSSI 的节点选为新的根节点。这可以用于应对根节点性能降低的情况。

要触发根节点切换，当前根节点必须明确调用`esp_mesh_waive_root()`以触发新的选举。当下根节点将指示网络中的所有节点开始发送并扫描信标帧（见自动根节点选择），**但与此同时一直保持联网（即不会变为空闲节点）**。如果另一个节点收到的票数超过当前根节点，则将启动根节点切换过程，**否则根节点将保持不变**。

新选出的根节点向当前的根节点发送 **切换请求**，而原先的根节点将返回一个应答通知，表示已经准备好切换。一旦接收到应答，新选出的根节点将与其父节点断开连接，并迅速与路由器形成上行连接，进而

成为网络的新根节点。原先的根节点将断开与路由器的连接，**并与此同时保持其所有下行连接**并进入空闲状态。之前的根节点将开始扫描潜在的父节点并选择首选父节点。

下图说明了根节点切换的示例。



切换根节点示例

1. 节点 C 是当前的根节点，但路由器 RSSI 值 (-85 dB) 降低至较低水平。此时，新的选举过程被触发了。所有节点开始传输和扫描信标帧（**此时仍保持连接**）。

2. 经过多轮传输和扫描后，节点 B 被选为新的根节点。节点 B 向节点 C 发送了一个 **切换请求**，节点 C 回复一个应答。

3. 节点 B 与其父节点断开连接，并与路由器连接，成为网络中的新根节点。节点 C 与路由器断开连接，进入空闲状态，并开始扫描并选择新的首选父节点。**节点 C 在整个过程中仍保持其所有的下行连接**。

4. 节点 C 选择节点 B 作为其的首选父节点，与之形成上行连接，并成为一个第二层节点。由于节点 C 仍保持相同的子网，因此根节点切换后的网络结构没有变化。然后，由于切换的发生，节点 C 子网中每个节点的所处层级均增加了一层。如果根节点切换过程中产生了新的根节点，则父节点切换 可以随后调整网络结构。

注解： 根节点切换必须要求选举，因此只有在使用自组网 ESP-MESH 网络时才支持。换句话说，如果使用指定的根节点，则不能进行根节点切换。

父节点切换

父节点切换是指一个子节点将其上行连接切换到更浅一层的另一个父节点。**父节点切换是自动的**，这意味着如果较浅层出现了可用的潜在父节点（因“异步上电复位”产生），子节点将自动更改其上行连接。

所有潜在的父节点将定期发送信标帧（参见信标帧和 *RSSI* 阈值），从而允许子节点扫描较浅层的父节点的可用性。由于父节点切换，自组网 ESP-MESH 网络可以动态调整其网络结构，以确保每个连接均具有良好的 RSSI 值，并且网络中的层级最小。

4.22.6 数据传输

ESP-MESH 数据包

ESP-MESH 网络使用 ESP-MESH 数据包传输数据。ESP-MESH 数据包 **完全包含在 Wi-Fi 数据帧**中。ESP-MESH 网络中的多跳数据传输将涉及通过不同 Wi-Fi 数据帧在每个无线跳上传输的单个 ESP-MESH 数据包。

下图显示了 ESP-MESH 数据包的结构及其与 Wi-Fi 数据帧的关系。

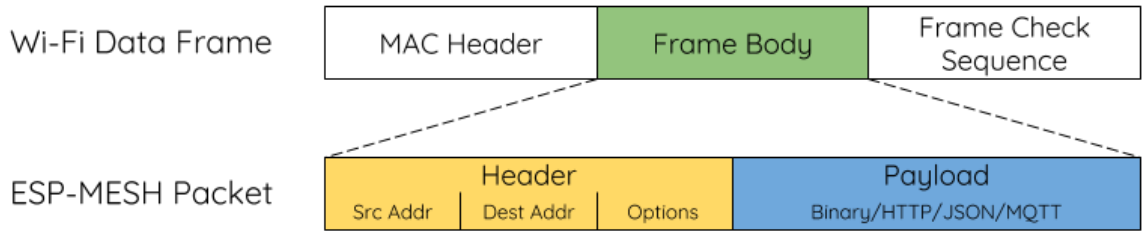


图 57: ESP-MESH 数据包

ESP-MESH 数据包的 **报头**包含源节点和目标节点的 MAC 地址。**选项 (option)** 字段包含有关特殊类型 ESP-MESH 数据包的信息，例如组传输或来自外部 IP 网络的数据包（请参阅*MESH_OPT_SEND_GROUP* 和*MESH_OPT_RECV_DS_ADDR*）。

ESP-MESH 数据包的 **有效载荷**包含实际的应用数据。该数据可以为原始二进制数据，也可以是使用 HTTP、MQTT 和 JSON 等应用层协议的编码数据（请见：*mesh_proto_t*）。

注解： 当向外部 IP 网络发送 ESP-MESH 数据包时，报头的目标地址字段将包含目标服务器的 IP 地址和端口号，而不是节点的 MAC 地址（请见：*mesh_addr_t*）。此外，根节点将处理外发 TCP/IP 数据包的形成。

组控制和组播

组播功能允许将单个 ESP-MESH 数据包同时发送给网络中的多个节点。ESP-MESH 中的组播可以通过“指定一个目标节点列表”或“预配置一个节点组”来实现。这两种组播方式均需调用*esp_mesh_send()* 实现。

如果通过“指定目标节点列表”实现组播，用户必须首先将 ESP-MESH 数据包的目标地址设置为 **组播组地址**（比如 `01:00:5E:xx:xx:xx`）。这表明 ESP-MESH 数据包是一个拥有一组地址的组播数据包，且该地址应该从报头选项中获得。然后，用户必须将目标节点的 MAC 地址列为选项（请见：*mesh_opt_t* 和*MESH_OPT_SEND_GROUP*）。这种组播方法不需要进行提前设置，但由于每个目标节点的 MAC 地址均需列为报头的选项字段，因此会产生大量开销数据。

分组组播允许 ESP-MESH 数据包被发送到一个预先配置的节点组。每个分组都有一个具有唯一性的 ID 标识。用户可通过*esp_mesh_set_group_id()* 将节点加入一个组。分组组播需要将 ESP-MESH 数据包的目标地址设置为目标组的 ID，还必须设置*MESH_DATA_GROUP* 标志位。分组组播产生的开销更小，但必须提前将节点加入分组中。

注解： 在组播期间，网络中的所有节点在 MAC 层都会收到 ESP-MESH 数据包。然而，不包括在 MAC 地址列表或目标组中的节点将简单地过滤掉这些数据包。

广播

广播功能允许将单个 ESP-MESH 数据包同时发送给网络中的所有节点。每个节点可以将一个广播包转发至其所有上行和下行连接，使得数据包尽可能快地在整个网络中传播。但是，ESP-MESH 利用以下方法

来避免在广播期间浪费带宽。

1. 当中间父节点收到来自其父节点的广播包时，它会将该数据包转发给自己的各个子节点，同时为自己保存一份数据包的副本。

2. 当中间父节点是广播的源节点时，它会将该数据包向上发送至其父节点，并向下发送给自己的各个子节点。

3. 当中间父节点接收到一个来自其子节点的广播包时，它会将该数据包转发给其父节点和其余子节点，同时为自己保存一份数据包的副本。

4. 当叶子节点是广播的源节点时，它会直接将该数据包发送至其父节点。

5. 当根节点是广播的源节点时，它会将该数据包发送至自己的所有子节点。

6. 当根节点收到来自其子节点的广播包时，它会将该数据包转发给其余子节点，同时为自己保存一份数据包的副本。

7. 当节点接收到一个源地址与自身 MAC 地址匹配的广播包时，它会将该广播包丢弃。

8. 当中间父节点收到一个来自其父节点的广播包时（该数据包最初来自该父节点的一个子节点），它会将该广播包丢弃。

上行流量控制

ESP-MESH 依赖父节点来控制其直接子节点的上行数据流。为了防止父节点的消息缓冲因上行传输过载而溢出，父节点将为每个子节点分配一个称为 **接收窗口** 的上行传输配额。**每个子节点均必须申请接收窗口才允许进行上行传输**。接收窗口的大小可以动态调整。完成从子节点到父节点的上行传输包括以下步骤：

1. 在每次传输之前，子节点向其父节点发送窗口请求。窗口请求中包括一个序号，与子节点的待传输数据包相对应。

2. 父节点接收窗口请求，并将序号与子节点发送的前一个数据包的序号进行比较，用于计算返回给子节点的接收窗口大小。

3. 子节点根据父节点指定的窗口大小发送数据包。如果子节点的接收窗口耗尽，它必须通过发送请求获得另一个接收窗口，然后才允许继续发送。

注解： ESP-MESH 不支持任何下行流量控制。

警告： 由于父节点切换，数据包可能会在上行传输期间丢失。

由于根节点是通向外部 IP 网络的唯一接口，因此下行节点必须了解根节点与外部 IP 网络的连接状态。否则，节点可能会尝试向一个已经与 IP 网络断开连接的根节点发送数据，从而造成不必要的传输和数据包丢失。ESP-MESH 可以基于监测根节点和外部 IP 网络的连接状态，提供一种稳定外发数据吞吐量的机制。根节点可以通过调用 `esp_mesh_post_toDS_state()` 将自身与外部 IP 网络的连接状态广播给所有其他节点。

双向数据流

下图展示了 ESP-MESH 双向数据流涉及的各种网络层。

由于使用路由表，**ESP-MESH 能够在 mesh 层中完全处理数据包的转发**。TCP/IP 层仅与 mesh 网络的根节点有关，可帮助根节点与外部 IP 网络的数据包传送。

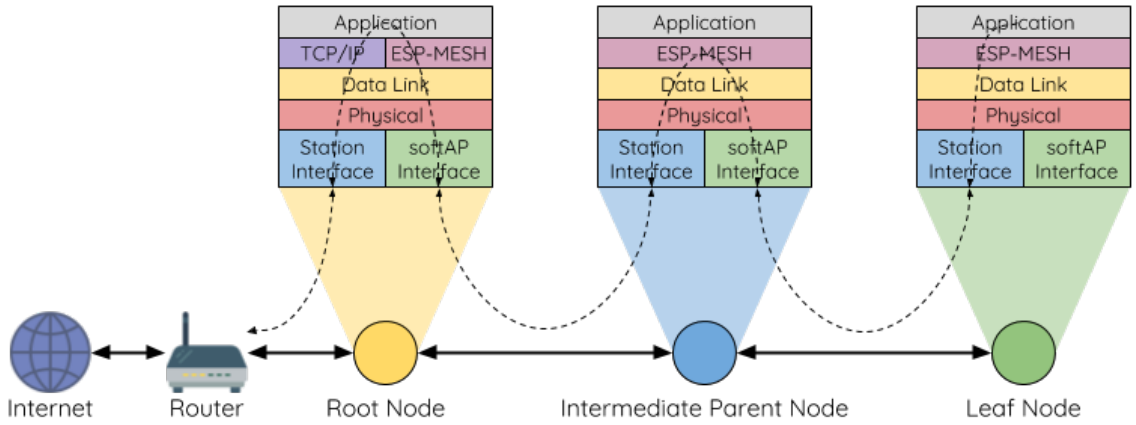


图 58: ESP-MESH 双向数据流

4.22.7 信道切换

背景

在传统的 Wi-Fi 网络中，**信道**代表预设的频率范围。在基础设施基本服务集 (BSS) 中，工作 AP 及与之相连的 station 必须处于传输信标的工作信道（1 到 14）中。物理上相邻的 BSS 使用相同的工作信道会导致干扰产生和性能下降。

为了允许 BSS 适应不断变化的物理层条件并保持性能，Wi-Fi 网络中增加了 **网络信道切换**的机制。网络信道切换是将 BSS 移至新的工作信道，并同时最大限度地减少期间对 BSS 的影响。然而，我们应该认识到，网络信道切换可能不会成功，无法将原信道中的所有 station 均移动至新的信道。

在基础设施 Wi-Fi 网络中，网络信道切换由 AP 触发，目的是将该 AP 及与之相连的所有 station 同步切换到新的信道。网络信道切换是通过在 AP 的周期性发送信标帧内嵌入一个 **信道切换公告 (CSA)** 元素来实现的。在网络信号切换前，该 CSA 元素用于向所有连接的 station 广播有关即将发生的网络信道切换，并且将包含在多个信标帧中。

一个 CSA 元素包含有关 **新信道号**和 **信道切换计数**的信息。其中，**信道切换计数**指示在网络信道切换之前剩余的信标帧间隔 (TBTT) 数量。因此，**信道切换计数**依每个信标帧递减，并且允许与之连接的 station 与 AP 同步进行信道切换。

ESP-MESH 网络信道切换

ESP-MESH 网络信道切换还利用包含 CSA 元素的信标帧。然而，ESP-MESH 作为一个多跳网络，其信标帧可能无法到达网络中的所有节点（这点与单跳网络不同），因此信道切换过程更加复杂。因此，ESP-MESH 网络依赖于通过节点转发 CSA 元素，从而实现在整个网络中的传播。

当具有一个或多个子节点的中间父节点接收到包含 CSA 元素的信标帧时，该节点会将该元素包含在其下一个发送的信标帧（即具有相同的 **新信道号**和 **信道切换计数**）中，从而实现该 CSA 元素的转发。鉴于 ESP-MESH 网络中的所有节点都接收到相同的 CSA 元素，这些节点可以使用 **信道切换计数**来同步其信道切换，但也会经历因 CSA 元素转发造成的延迟。

ESP-MESH 网络信道切换可以由路由器或根节点触发。

根节点触发 **由根节点触发的信道切换只能在 ESP-MESH 网络未连接到路由器**时才会发生。通过调用`esp_mesh_switch_channel()`，根节点将设置一个初始 **信道切换计数**值，并开始在其信标帧中包含 CSA 元素。接着，每个 CSA 元素将抵达第二层节点，并通过第二层节点自己的信标帧继续进行向下转发。

路由器触发 当 ESP-MESH 网络连接到路由器时，整个网络必须与路由器采用同一个信道。因此，**根节点在连接到路由器时无法触发信道切换**。

当根节点从路由器接收到包含 CSA 元素的信标帧时，**根节点将 CSA 元素中的信道切换计数值设置为自定义值，然后再通过信标帧继续向下转发**。此后，该 **信道切换计数**将依转发次数相对于自定义值依次递减。该自定义值可以基于诸如网络层级、当前节点数等因素。

ESP-MESH 网络及其路由器可能具有不同且变化的信标间隔，因此需要将 **信道切换计数**值设置为自定义值。也就是说，路由器提供的 **信道切换计数**值与 ESP-MESH 网络无关。通过使用自定义值，ESP-MESH 网络中的节点能够相对于 ESP-MESH 网络的信标间隔同步切换信道。也正因如此，ESP-MESH 网络也会出现信道与路由器及其连接 station 的信道切换不同步的情况。

网络信道切换的影响

- **由于 ESP-MESH 网络信道切换与路由器的信道切换不同步，ESP-MESH 网络和路由器之间会出现 临时信道差异**。
  - ESP-MESH 网络的信道切换时间取决于 ESP-MESH 网络的信标间隔和根节点的自定义 **信道切换计数**。
  - 在 ESP-MESH 网络切换期间，信道差异将阻止根节点和路由器之间的任何数据交换。
  - 在 ESP-MESH 网络中，根节点和中间父节点将请求与其连接的子节点停止传输，直至信道切换发生（通过将 CSA 元素的 **信道切换模式**字段置为 1）。
  - 频繁的路由器触发网络信道切换可能会降低 ESP-MESH 网络的性能。请注意，这可能是由 ESP-MESH 网络本身造成的（例如由于 ESP-MESH 网络的无线介质争用等原因）。此时，用户应该禁用路由器触发的自主信道切换，并直接指定一个信道。
- **当存在 临时信道差异时，根节点从技术上来说仍保持连接至路由器**。
  - 如果根节点经过一定数量信标间隔仍无法接到信标帧或探测来自路由器的响应，则会断开连接。
  - 断开连接时，根节点将自动重新扫描所有信道以确定是否存在路由器。
- **如果根节点无法接收任何路由器的 CSA 信标帧（例如短暂的路由器切换时间），则路由器将在没有 ESP-MESH 网络**
  - 在路由器切换信道后，根节点将不再能够接收路由器的信标帧和探测响应，并导致在一定数量的信标间隔后断开连接。
  - 在断开连接后，根节点将重新所有信道，寻找路由器。
  - 根节点将在整个过程中维护与之相连的下行连接。

注解： 虽然 ESP-MESH 网络信道切换的目的是将网络内的所有节点移动到新的工作信道，但也应该认识到，信道切换可能无法成功移动所有节点（比如由于节点故障等原因）。

信道和路由器切换配置

ESP-MESH 允许通过配置启用或禁用自主信道切换。同样，也可以通过配置启用或禁用自主路由器切换（即当根节点自主连接到另一个路由器时）。自主信道切换和自主路由器切换取决于以下配置参数和运行时间条件。

允许信道切换：本参数决定是否允许 ESP-MESH 网络进行自主信道切换，具体可通过*mesh_cfg_t* 结构体中的 `allow_channel_switch` 字段进行配置。

预设信道：ESP-MESH 网络可以将*mesh_cfg_t* 结构体中的 `channel` 字段设置为相应的信道号，而具备一个预设信道。如果未设置此字段，则 `allow_channel_switch` 的设置将被覆盖，即始终允许信道切换。

允许路由器切换：本参数决定是否允许 ESP-MESH 网络进行自主路由器切换，具体可通过*mesh_router_t* 结构体中的 `allow_router_switch` 字段进行配置。

预设路由器 BSSID：ESP-MESH 网络可以将*mesh_router_t* 结构体的 `bssid` 字段设置为目标路由器的 BSSID，而预设一个路由器。如果未设置此字段，则 `allow_router_switch` 的设置将被覆盖，即始终允许路由器切换。

存在根节点：根节点的存在也会影响是否允许信道或路由器切换。

下表说明了在不同参数/条件组合下是否允许信道切换和路由器切换。请注意，*X* 代表参数“不关心”。

预设信道	允许信道切换	预置路由器 BSSID	允许路由器切换	存在根节点	允许切换？
N	X	N	X	X	信道与路由器
N	X	Y	N	X	仅信道
N	X	Y	Y	X	信道与路由器
Y	Y	N	X	X	信道与路由器
Y	N	N	X	N	仅路由器
Y	N	N	X	Y	信道与路由器
Y	Y	Y	N	X	仅信道
Y	N	Y	N	N	无
Y	N	Y	N	Y	仅信道
Y	Y	Y	Y	X	信道与路由器
Y	N	Y	Y	N	仅路由器
Y	N	Y	Y	Y	信道与路由器

4.22.8 性能

ESP-MESH 网络的性能可以基于以下多个指标进行评估：

组网时长：从头开始构建 ESP-MESH 网络所需的总时长。

修复时间：从网络检测到节点断开到执行适当操作（例如生成新的根节点或形成新的连接等）以修复网络所需的时间。

每跳延迟：数据每经过一次无线 hop 而经历的延迟，即从父节点向子节点（或从子节点向父节点）发送一个数据包所需的时间。

网络节点容量：ESP-MESH 网络可以同时支持的节点总数。该指标取决于节点可以接受到的最大下行连接数和网络中允许的最大层级。

ESP-MESH 网络的常见性能指标如下表所示：

- 组网时长：< 60 秒
- **修复时间**
  - 根节点断开：< 10 秒
  - 子节点断开：< 5 秒
- 每条延迟：10 到 30 毫秒

注解： 上述性能指标的测试条件见下。

- 测试设备数量：**100**
- 最大允许下行连接数量：**6**
- 最大允许层级：**6**

注解： 吞吐量取决于数据包错误率和 hop 数量。

注解： 根节点访问外部 IP 网络的吞吐量直接受到 ESP-MESH 网络中节点数量和路由器带宽的影响。

注解： 用户应注意，ESP-MESH 网络的性能与网络配置和工作环境密切相关。

4.22.9 更多注意事项

- 数据传输使用 Wi-Fi WPA2-PSK 加密
- Mesh 网络 IE 使用 AES 加密

本文图片中使用的路由器与互联网图标来自 www.flaticon.com 的 Smashicons。

4.23 片外 RAM

4.23.1 简介

ESP32-S2 提供了 520 KB 的片上 SRAM，可以满足大部分需求。但有些场景可能需要更多 RAM，因此 ESP32-S2 另外提供了高达 4 MB 的片外 SPI RAM 存储器以供用户使用。片外 RAM 被添加到内存映射中，在某些范围内与片上 RAM 使用方式相同。

4.23.2 硬件

ESP32-S2 支持与 SPI Flash 芯片并联的 SPI PSRAM。ESP32-S2 支持多种类型的 RAM 芯片，但 ESP32-S2 SDK 当前仅支持 ESP-PSRAM32 芯片。

ESP-PSRAM32 芯片的工作电压为 1.8 V，只能与 1.8 V flash 并联使用。请确保在启动时将 MTDI 管脚设置为高电平，或者将 ESP32-S2 中的熔丝设置为始终使用 1.8 V 的 VDD_SIO 电平，否则有可能会损坏 PSRAM 和/或 flash 芯片。

要将 ESP-PSRAM 芯片连接到 ESP32D0W*，请连接以下信号：

- PSRAM /CE (pin 1) > ESP32 GPIO 16
- PSRAM SO (pin 2) > flash DO
- PSRAM SIO[2] (pin 3) > flash WP
- PSRAM SI (pin 5) > flash DI
- PSRAM SCLK (pin 6) > ESP32 GPIO 17
- PSRAM SIO[3] (pin 7) > flash HOLD
- PSRAM Vcc (pin 8) > ESP32 VCC_SDIO

ESP32D2W* 芯片的连接方式有待确定。

注解： 乐鑫同时提供 ESP32-WROVER 模组，内部搭载 ESP32 芯片，集成 1.8 V flash 和 ESP-PSRAM32，可直接用于终端产品 PCB 中。

4.23.3 配置片外 RAM

ESP-IDF 完全支持将外部存储器集成到您的应用程序中。您可以将 ESP-IDF 配置成启动并完成初始化后以多种方式处理片外 RAM：

- *集成片外 RAM 到 ESP32-S2 内存映射*
- *添加片外 RAM 到内存分配程序*
- *调用 malloc() 分配片外 RAM* （默认）

集成片外 RAM 到 ESP32-S2 内存映射

在*CONFIG_SPIRAM_USE* 中选择“Integrate RAM into ESP32-S2 memory map（集成片外 RAM 到 ESP32-S2 内存映射）”选项。

这是集成片外 RAM 最基础的设置选项，大多数用户需要用到其他更高级的选项。

ESP-IDF 启动过程中，片外 RAM 被映射到以 0x3F800000 起始的数据地址空间（字节可寻址），空间大小正好为 RAM 的大小 (4 MB)。

应用程序可以通过创建指向该区域的指针手动将数据放入片外存储器，同时应用程序全权负责管理片外 RAM，包括协调 Buffer 的使用、防止发生损坏等。

添加片外 RAM 到内存分配程序

在*CONFIG_SPIRAM_USE* 中选择“Make RAM allocatable using heap_caps_malloc(…, MALLOC_CAP_SPIRAM)”选项。

启用上述选项后，片外 RAM 被映射到地址 0x3F800000，并将这个区域添加到*内存分配程序* 里携带 `MALLOC_CAP_SPIRAM` 的标志

程序如果想从片外存储器分配存储空间，则需要调用 `heap_caps_malloc(size, MALLOC_CAP_SPIRAM)`，之后可以调用 `free()` 函数释放这部分存储空间。

调用 malloc() 分配片外 RAM

在*CONFIG_SPIRAM_USE* 中选择“Make RAM allocatable using malloc() as well”选项，该选项为默认选项。

启用此选项后，片外存储器将被添加到内存分配程序（与上一选项相同），同时也将被添加到由标准 `malloc()` 返回的 RAM 中。

这允许应用程序使用片外 RAM 而无需重写代码以使用 `heap_caps_malloc(..., MALLOC_CAP_SPIRAM)`。

如果某次内存分配偏向于片外存储器，您也可以使用*CONFIG_SPIRAM_MALLOC_ALWAYSINTERNAL* 设置分配空间的大小阈值，控制分配结果：

- 如果分配的空间小于阈值，分配程序将首先选择内部存储器。
- 如果分配的空间等于或大于阈值，分配程序将首先选择外部存储器。

如果优先考虑的内部或外部存储器中没有可用的存储块，分配程序则会选择其他类型存储。

由于有些 Buffer 仅可在内部存储器中分配，因此需要使用第二个配置项*CONFIG_SPIRAM_MALLOC_RESERVE_INTERNAL* 定义一个内部存储池，仅限显式的内部存储器分配使用（例如用于 DMA 的存储器）。常规 `malloc()` 将不会从该池中分配，但可以使用*MALLOC_CAP_DMA* 和 `MALLOC_CAP_INTERNAL` 旗标从该池中分配存储器。

4.23.4 片外 RAM 使用限制

使用片外 RAM 有下面一些限制：

- Flash cache 禁用时（比如，正在写入 flash），片外 RAM 将无法访问；同样，对片外 RAM 的读写操作也将导致 cache 访问异常。出于这个原因，ESP-IDF 不会在片外 RAM 中分配任务堆栈（详见下文）。
- 片外 RAM 不能用于储存 DMA 描述符，也不能用作 DMA 读写操作的缓冲区 (Buffer)。与 DMA 搭配使用的 Buffer 必须先使用 `heap_caps_malloc(size, MALLOC_CAP_DMA)` 进行分配，之后可以调用标准 `free()` 回调释放 Buffer。
- 片外 RAM 与片外 flash 使用相同的 cache 区域，即频繁在片外 RAM 访问的变量可以像在片上 RAM 中一样快速读取和修改。但访问大块数据时（大于 32 KB），cache 空间可能会不足，访问速度将回落到片外 RAM 访问速度。此外，访问大块数据可以“挤出”flash cache，可能会降低代码执行速度。
- 片外 RAM 不可用作任务堆栈存储器。因此*`xTaskCreate()`* 及类似函数将始终为堆栈和任务 TCB 分配片上储存器，而*`xTaskCreateStatic()`* 类型的函数将检查传递的 Buffer 是否属于片上存储器。
- 默认情况下，片外 RAM 初始化失败将终止 ESP-IDF 启动。如果想禁用此功能，可启用*CONFIG_SPIRAM_IGNORE_NOTFOUND* 配置选项。

4.24 链接脚本生成机制

4.24.1 概述

用于存放代码和数据的:ref:‘ 内存区域 <memory-layout>‘有多个。代码和只读数据默认存放在 flash 中，可写数据存放在 RAM 中。不过有时，我们必须更改默认存放区域，例如为了提高性能，将关键代码存放到 RAM 中，或者将代码存放到 RTC 存储器中以便在唤醒桩 和*ULP 协处理器* 中使用。

链接脚本生成机制可以让用户指定代码和数据在 ESP-IDF 组件中的存放区域。组件包含如何存放符号、目标或完整库的信息。在构建应用程序时，组件中的这些信息会被收集、解析并处理；生成的存放规则用于链接应用程序。

4.24.2 快速上手

本段将指导如何使用 ESP-IDF 的即用方案，快速将代码和数据放入 RAM 和 RTC 存储器中。

假设我们有:

```
- components/
                - my_component/
                                - CMakeLists.txt
                                - component.mk
                                - Kconfig
                                - src/
                                      - my_src1.c
                                      - my_src2.c
                                      - my_src3.c
                                - my_linker_fragment_file.lf
```

- 名为 my_component 的组件，在构建过程中存储为 libmy_component.a 库文件
- 库文件包含的三个源文件：my_src1.c、my_src2.c 和 my_src3.c，编译后分别为 my_src1.o、my_src2.o 和 my_src3.o，
- 在 my_src1.o 定义的 my_function1 功能；在 my_src2.o 定义的 my_function2 功能
- 存储在 my_component 下 Kconfig 中的布尔类型配置 PERFORMANCE_MODE (y/n) 和整数类型配置 PERFORMANCE_LEVEL （范围是 0-3）

创建和指定链接片段文件

首先，我们需要创建链接片段文件。链接片段文件是一个拓展名为 .lf 的文本文件，文件内写有想要存放的位置。文件创建成功后，需要将其呈现在构建系统中。ESP-IDF 支持的构建系统指南如下：

Make 在组件目录的 component.mk 文件中设置 COMPONENT_ADD_LDFRAGMENTS 变量的值，使其指向已创建的链接片段文件。路径可以为绝对路径，也可以为组件目录的相对路径。

```
COMPONENT_ADD_LDFRAGMENTS += my_linker_fragment_file.lf
```

CMake 在组件目录的 CMakeLists.txt 文件中，指定 idf_component_register 调用引数 LDFRAGMENTS 的值。LDFRAGMENTS 可以为绝对路径，也可为组件目录的相对路径，指向刚才创建的链接片段文件。

```
# 相对于组件的 CMakeLists.txt 的文件路径
idf_component_register(...
                       LDFRAGMENTS "path/to/linker_fragment_file.lf" "path/to/
↪another_linker_fragment_file.lf"
                       ...
                       )
```

指定存放区域

可以按照下列粒度指定存放区域：

- 目标文件（.obj 或 .o 文件）
- 符号（函数/变量）
- 库（.a 文件）

存放目标文件 假设整个 my_src1.o 目标文件对性能至关重要，所以最好把该文件放在 RAM 中。另外，my_src2.o 目标文件包含从深度睡眠唤醒所需的符号，因此需要将其存放到 RTC 存储器中。在链接片段文件中可以写入以下内容：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    my_src1 (noflash)      # 将所有 my_src1 代码和只读数据存放在 IRAM 和 DRAM 中
    my_src2 (rtc)          # 将所有 my_src2 代码、数据和只读数据存放到 RTC 快速 RAM 和 RTC␣
↪慢速 RAM 中
```

那么 my_src3.o 放在哪里呢？由于未指定存放区域，my_src3.o 会存放到默认区域。更多关于默认存放区域的信息，请查看这里。

存放符号 继续上文的例子，假设 object1.o 目标文件定义的功能中，只有 function1 影响到性能；object2.o 目标文件中只有 function2 需要在芯片从深度睡眠中唤醒后运行。可以在链接片段文件中写入以下内容实现：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    my_src1:my_function1 (noflash)
    my_src2:my_function2 (rtc)
```

my_src1.o 和 my_src2.o 中的其他函数以及整个 object3.o 目标文件会存放到默认区域。要指定数据的存放区域，仅需将上文的函数名替换为变量名即可，如:

```
my_src1:my_variable (noflash)
```

注意: 按照符号粒度存放代码和数据有一定的局限。为确保存放区域合适，您也可以将相关代码和数据集中在源文件中，参考使用目标文件的存放规则。

存放整个库 在这个例子中，假设整个组件库都需存放到 RAM 中，可以写入以下内容实现：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    * (noflash)
```

类似的，写入以下内容可以将整个组件存放到 RTC 存储器中：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    * (rtc)
```

根据具体配置存放 假设只有在某个条件为真时，比如 CONFIG_PERFORMANCE_MODE == y 时，整个组件库才有特定存放区域，可以写入以下内容实现：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    if PERFORMANCE_MODE = y:
        * (noflash)
    else:
        * (default)
```

来看一种更复杂的情况。假设 "CONFIG_PERFORMANCE_LEVEL == 1" 时，只有 object1.o 存放到 RAM 中；CONFIG_PERFORMANCE_LEVEL == 2 时，object1.o 和 object2.o 会存放到 RAM 中；CONFIG_PERFORMANCE_LEVEL == 3 时，库中的所有目标文件都会存放到 RAM 中。以上三个条件为假时，整个库会存放到 RTC 存储器中。虽然这种使用场景很罕见，不过，还是可以通过以下方式实现：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    if PERFORMANCE_LEVEL = 1:
        my_src1 (noflash)
    elif PERFORMANCE_LEVEL = 2:
        my_src1 (noflash)
        my_src2 (noflash)
    elif PERFORMANCE_LEVEL = 3:
        my_src1 (noflash)
        my_src2 (noflash)
        my_src3 (noflash)
    else:
        * (rtc)
```

也可以嵌套条件检查。以下内容与上述片段等效：

```
[mapping:my_component]
archive: libmy_component.a
entries:
    if PERFORMANCE_LEVEL <= 3 && PERFORMANCE_LEVEL > 0:
        if PERFORMANCE_LEVEL >= 1:
            object1 (noflash)
            if PERFORMANCE_LEVEL >= 2:
                object2 (noflash)
                if PERFORMANCE_LEVEL >= 3:
                    object2 (noflash)
    else:
        * (rtc)
```

默认存放区域

到目前为止，"默认存放区域" 在未指定 rtc 和 noflash 存放规则时才会使用，作为备选方案。需要注意的是，noflash 或者 rtc 标记不仅仅是关键字，实际上还是被称作片段的实体，确切地说是协议。

与 rtc 和 noflash 类似，还有一个 默认协议，定义了默认存放规则。顾名思义，该协议规定了代码和数据通常存放的区域，即代码和恒量存放在 flash 中，变量存放在 RAM 中。更多关于默认协议的信息，请见这里。

注解： 使用链接脚本生成机制的 IDF 组件示例，请参阅 freertos/CMakeLists.txt。为了提高性能，freertos 使用链接脚本生成机制，将其目标文件存放到 RAM 中。

快速入门指南到此结束，下文将详述这个机制的内核，有助于创建自定义存放区域或修改默认方式。

4.24.3 链接脚本生成机制内核

链接是将 C/C++ 源文件转换成可执行文件的最后一步。链接由工具链的链接器完成，接受指定代码和数据存放区域等信息的链接脚本。链接脚本生成机制的转换过程类似，区别在于传输给链接器的链接脚本根据 (1) 收集的链接片段文件 和 (2) 链接脚本模板 动态生成。

注解： 执行链接脚本生成机制的工具存放在 tools/ldgen 之下。

链接片段文件

如快速入门指南所述，片段文件是拓展名为 `.lf` 的简单文本文件，内含想要存放区域的信息。不过，这是对片段文件所包含内容的简化版描述。实际上，片段文件内包含的是“片段”。片段是实体，包含多条信息，这些信息放在一起组成了存放规则，说明目标文件各个段在二进制输出文件中的存放位置。片段一共有三种，分别是段、协议 和映射。

语法 三种片段类型使用同一种语法：

```
[type:name]
key: value
key:
    value
    value
    value
    ...
```

- 类型：片段类型，可以为 段、协议或 映射。
- 名称：片段名称，指定片段类型的片段名称应唯一。
- 键值：片段内容。每个片段类型可支持不同的键值和不同的键值语法。

注解： 多个片段的类型和名称相同时会引发异常。

注解： 片段名称和键值只能使用字母、数字和下划线。

条件检查

条件检查使得链接脚本生成机制可以感知配置。含有配置值的表达式是否为真，决定了使用哪些特定键值。检查使用的是 tools/kconfig_new/kconfiglib.py 脚本的 `eval_string`，遵循该脚本要求的语法和局限性，支持：

- **比较**
  - 小于 <
  - 小于等于 <=
  - 大于 >
  - 大于等于 >=
  - 等于 =
  - 不等于 !=
- **逻辑**
  - 或 ||
  - 和 &&
  - 否定？取反？ !
- **分组**
  - 圆括号 ()

条件检查和其他语言中的 `if...elseif/elif...else` 块作用一样。键值和完整片段都可以进行条件检查。以下两个示例效果相同：

```
# 键值取决于配置
[type:name]
key_1:
    if CONDITION = y:
        value_1
    else:
        value_2
key_2:
    if CONDITION = y:
        value_a
    else:
        value_b
```

```
# 完整片段的定义取决于配置
if CONDITION = y:
    [type:name]
    key_1:
        value_1
    key_2:
        value_b
else:
    [type:name]
    key_1:
        value_2
    key_2:
        value_b
```

注释

链接片段文件中的注释以 # 开头。和在其他语言中一样，注释提供了有用的描述和资料，在处理过程中会被忽略。

与 ESP-IDF v3.x 链接脚本片段文件兼容 ESP-IDF v4.0 变更了链接脚本片段文件使用的一些语法：

- 必须缩进，缩进不当的文件会产生解析异常；旧版本不强制缩进，但之前的文档和示例均遵循了正确的缩进语法
- 条件改用 `if...elif...else` 结构，可以嵌套检查，将完整片段置于条件内
- 映射片段和其他片段类型一样，需有名称

链接脚本生成器可解析 ESP-IDF v3.x 版本中缩进正确的链接片段文件（如 ESP-IDF v3.x 版本中的本文件所示），依然可以向后兼容此前的映射片段语法（可选名称和条件的旧语法），但是会有弃用警告。用户应换成本文档介绍的新语法，因为旧语法将在未来停用。

请注意，ESP-IDF v3.x 不支持使用 ESP-IDF v4.0 新语法的链接片段文件。

类型 段

段定义了 GCC 编译器输出的一系列目标文件段，可以是默认段（如 `.text`、`.data`），也可以是用户通过 `__attribute__` 关键字定义的段。

'+' 表示段列表开始，且当前段为列表中的第一个段。这种表达方式更加推荐。

```
[sections:name]
entries:
    .section+
    .section
    ...
```

示例：

```
# 不推荐的方式
[sections:text]
entries:
    .text
    .text.*
    .literal
    .literal.*

# 推荐的方式，效果与上面等同
[sections:text]
entries:
    .text+              # 即 .text 和 .text.*
    .literal+           # 即 .literal 和 .literal.*
```

协议

协议定义了每个段对应的 目标。

```
[scheme:name]
entries:
    sections -> target
    sections -> target
    ...
```

示例：

```
[scheme:noflash]
entries:
    text -> iram0_text          # text 段下的所有条目均归入 iram0_text
    rodata -> dram0_data        # rodata 段下的所有条目均归入 dram0_data
```

默认协议

注意，有一个 默认的协议很特殊，特殊在于包罗存放规则都是根据这个协议中的条目生成的。这意味着，如果该协议有一条条目是 `text -> flash_text`，则将为目标 `flash_text` 生成如下的存放规则:

```
*(.literal .literal.* .text .text.*)
```

这些生成的包罗规则将用于未指定映射规则的情况。

默认协议在 esp32s2/ld/esp32s2_fragments.lf 文件中定义，快速上手指南中提到的内置 `noflash` 协议和 `rtc` 协议也在该文件中定义。

映射

映射定义了可映射实体（即目标文件、函数名、变量名和库）对应的协议。

```
[mapping]
archive: archive                # 构建后输出的库文件名称（即 libxxx.a）
entries:
    object:symbol (scheme)      # 符号
    object (scheme)             # 目标
    * (scheme)                  # 库
```

有三种存放粒度：

- 符号：指定了目标文件名称和符号名称。符号名称可以是函数名或变量名。
- 目标：只指定目标文件名称。
- 库：指定 `*`，即某个库下面所有目标文件的简化表达法。

为了更好地理解条目的含义，我们看一个按目标存放的例子。

```
object (scheme)
```

根据条目定义，将这个协议展开：

```
object (sections -> target,
        sections -> target,
        ...)
```

再根据条目定义，将这个段展开：

```
object (.section,
        .section,
        ... -> target, # 根据目标文件将这里所列出的所有段放在该目标位置

        .section,
        .section,
        ... -> target, # 同样的方法指定其他段

        ...)           # 直至所有段均已展开
```

示例：

```
[mapping:map]
archive: libfreertos.a
entries:
    * (noflash)
```

按符号存放 按符号存放可通过编译器标志 `-ffunction-sections` 和 `-fdata-sections` 实现。ESP-IDF 默认用这些标志编译。用户若选择移除标志，便不能按符号存放。另外，即便有标志，也会其他限制，具体取决于编译器输出的段。

比如，使用 `-ffunction-sections`，针对每个功能会输出单独的段。段的名称可以预测，即 `.text.{func_name}` 和 `.literal.{func_name}`。但是功能内的字符串并非如此，因为字符串会进入字符串池，或者使用生成的段名称。

使用 `-fdata-sections`，对全局数据来说编译器可输出 `.data.{var_name}`、`.rodata.{var_name}` 或 `.bss.{var_name}`；因此类型 I 映射词条可以适用。但是，功能中声明的静态数据并非如此，生成的段名称是将变量名称和其他信息混合。

链接脚本模板

链接脚本模板是指定存放规则的存放位置的框架，与其他链接脚本没有本质区别，但带有特定的标记语法，可以指示存放生成的存放规则的位置。

如需引用一个目标标记下的所有存放规则，请使用以下语法：

```
mapping[target]
```

示例：

以下示例是某个链接脚本模板的摘录，定义了输出段 `.iram0.text`，该输出段包含一个引用目标 `iram0_text` 的标记。

```
.iram0.text :
{
    /* 标记 IRAM 空间不足 */
    _iram_text_start = ABSOLUTE(.);

    /* 引用 iram0_text */
    mapping[iram0_text]

    _iram_text_end = ABSOLUTE(.);
} > iram0_0_seg
```

假设链接脚本生成器收集到了以下片段定义：

```
[sections:text]
    .text+
    .literal+

[sections:iram]
    .iram1+

[scheme:default]
entries:
    text -> flash_text
    iram -> iram0_text

[scheme:noflash]
entries:
    text -> iram0_text

[mapping:freertos]
archive: libfreertos.a
entries:
    * (noflash)
```

则该脚本生成器生成的链接脚本文件，其摘录应如下所示：

```
.iram0.text :
{
    /* 标记 IRAM 空间不足 */
    _iram_text_start = ABSOLUTE(.);

    /* 处理片段生成的存放规则，存放在模板标记的位置处 */
    *(.iram1 .iram1.*)
    *libfreertos.a:(.literal .text .literal.* .text.*)

    _iram_text_end = ABSOLUTE(.);
} > iram0_0_seg
```

`*libfreertos.a:(.literal .text .literal.* .text.*)`

这是根据 `freertos` 映射的 `* (noflash)` 条目生成的规则。`libfreertos.a` 库下所有目标文件的所有 `text` 段会收集到 `iram0_text` 目标下（按照 `noflash` 协议），并放在模板中被 `iram0_text` 标记的地方。

`*(.iram1 .iram1.*)`

这是根据默认协议条目 `iram -> iram0_text` 生成的规则。默认协议指定了 `iram -> iram0_text` 条目，因此生成的规则同样也放在被 `iram0_text` 标记的地方。由于该规则是根据默认协议生成的，因此在同一目标下收集的所有规则下排在第一位。

目前使用的链接脚本模板是 esp32s2/ld/esp32s2.project.ld.in，由 `esp32s2` 组件指定，生成的脚本存放在构建目录下。

4.25 lwIP

ESP-IDF uses the open source lwIP lightweight TCP/IP stack. The ESP-IDF version of lwIP (esp-lwip) has some modifications and additions compared to the upstream project.

4.25.1 Supported APIs

ESP-IDF supports the following lwIP TCP/IP stack functions:

- *BSD Sockets API*

- *Netconn API* is enabled but not officially supported for ESP-IDF applications

Adapted APIs

Some common lwIP "app" APIs are supported indirectly by ESP-IDF:

- DHCP Server & Client are supported indirectly via the *ESP-NETIF* functionality
- Simple Network Time Protocol (SNTP) is supported via the lwip/include/apps/sntp/sntp.h lwip/lwip/src/include/lwip/apps/sntp.h functions (see also *SNTP Time Synchronization*)
- ICMP Ping is supported using a variation on the lwIP ping API. See *ICMP Echo*.
- NetBIOS lookup is available using the standard lwIP API. protocols/http_server/restful_server has an option to demonstrate using NetBIOS to look up a host on the LAN.
- mDNS uses a different implementation to the lwIP default mDNS (see *mDNS 服务*), but lwIP can look up mDNS hosts using standard APIs such as `gethostbyname()` and the convention `hostname.local`, provided the *CONFIG_LWIP_DNS_SUPPORT_MDNS_QUERIES* setting is enabled.

4.25.2 BSD Sockets API

The BSD Sockets API is a common cross-platform TCP/IP sockets API that originated in the Berkeley Standard Distribution of UNIX but is now standardized in a section of the POSIX specification. BSD Sockets are sometimes called POSIX Sockets or Berkeley Sockets.

As implemented in ESP-IDF, lwIP supports all of the common usages of the BSD Sockets API.

References

A wide range of BSD Sockets reference material is available, including:

- Single UNIX Specification BSD Sockets page
- Berkeley Sockets Wikipedia page

Examples

A number of ESP-IDF examples show how to use the BSD Sockets APIs:

- protocols/sockets/tcp_server
- protocols/sockets/tcp_client
- protocols/sockets/udp_server
- protocols/sockets/udp_client
- protocols/sockets/udp_multicast
- protocols/http_request (Note: this is a simplified example of using a TCP socket to send an HTTP request. The *ESP HTTP Client* is a much better option for sending HTTP requests.)

Supported functions

The following BSD socket API functions are supported. For full details see lwip/lwip/src/include/lwip/sockets.h.

- `socket()`
- `bind()`
- `accept()`
- `shutdown()`
- `getpeername()`
- `getsockopt()` & `setsockopt()` (see *Socket Options*)
- `close()` (via 虚拟文件系统组件)
- `read()`, `readv()`, `write()`, `writev()` (via 虚拟文件系统组件)
- `recv()`, `recvmsg()`, `recvfrom()`
- `send()`, `sendmsg()`, `sendto()`

- `select()` (via 虚拟文件系统组件)
- `poll()` (Note: on ESP-IDF, `poll()` is implemented by calling select internally, so using `select()` directly is recommended if a choice of methods is available.)
- `fcntl()` (see *fcntl*)

Non-standard functions:

- `ioctl()` (see *ioctls*)

注解: Some lwIP application sample code uses prefixed versions of BSD APIs, for example `lwip_socket()` instead of the standard `socket()`. Both forms can be used with ESP-IDF, but using standard names is recommended.

Socket Error Handling

BSD Socket error handling code is very important for robust socket applications. Normally the socket error handling involves the following aspects:

- Detecting the error.
- Geting the error reason code.
- Handle the error according to the reason code.

In lwIP, we have two different scenarios of handling socket errors:

- Socket API returns an error. For more information, see *Socket API Errors*.
- `select(int maxfdp1, fd_set *readset, fd_set *writeset, fd_set *exceptset, struct timeval *timeout)` has exception descriptor indicating that the socket has an error. For more information, see *select() Errors*.

Socket API Errors

The error detection

- We can know that the socket API fails according to its return value.

Get the error reason code

- When socket API fails, the return value doesn't contain the failure reason and the application can get the error reason code by accessing errno. Different values indicate different meanings. For more information, see <*Socket Error Reason Code*>.

Example:

```
int err;
int sockfd;

if (sockfd = socket(AF_INET,SOCK_STREAM,0) < 0) {
    // the error code is obtained from errno
    err = errno;
    return err;
}
```

select() Errors

The error detection

- Socket error when `select()` has exception descriptor

Get the error reason code

- If the `select` indicates that the socket fails, we can't get the error reason code by accessing errno, instead we should call `getsockopt()` to get the failure reason code. Because `select()` has exception descriptor, the error code will not be given to errno.

注解: `getsockopt` function prototype `int getsockopt(int s, int level, int optname, void *optval, socklen_t *optlen)`. Its function is to get the current value of the option of any type, any state socket, and store the result in optval. For example, when you get the error code on a socket, you can get it by `getsockopt(sockfd, SOL_SOCKET, SO_ERROR, &err, &optlen)`.

Example:

```
int err;

if (select(sockfd + 1, NULL, NULL, &exfds, &tval) <= 0) {
    err = errno;
    return err;
} else {
    if (FD_ISSET(sockfd, &exfds)) {
        // select() exception set using getsockopt()
        int optlen = sizeof(int);
        getsockopt(sockfd, SOL_SOCKET, SO_ERROR, &err, &optlen);
        return err;
    }
}
```

Socket Error Reason Code Below is a list of common error codes. For more detailed list of standard POSIX/C error codes, please see *newlib errno.h <https://github.com/espressif/newlib-esp32/blob/master/newlib/libc/include/sys/errno.h>* and the platform-specific extensions newlib/platform_include/errno.h

Error code	Description
ECONNREFUSED	Connection refused
EADDRINUSE	Address already in use
ECONNABORTED	Software caused connection abort
ENETUNREACH	Network is unreachable
ENETDOWN	Network interface is not configured
ETIMEDOUT	Connection timed out
EHOSTDOWN	Host is down
EHOSTUNREACH	Host is unreachable
EINPROGRESS	Connection already in progress
EALREADY	Socket already connected
EDESTADDRREQ	Destination address required
EPROTONOSUPPORT	Unknown protocol

Socket Options

The `getsockopt()` and `setsockopt()` functions allow getting/setting per-socket options.

Not all standard socket options are supported by lwIP in ESP-IDF. The following socket options are supported:

Common options Used with level argument `SOL_SOCKET`.

- `SO_REUSEADDR` (available if *CONFIG_LWIP_SO_REUSE* is set, behavior can be customized by setting *CONFIG_LWIP_SO_REUSE_RXTOALL*)
- `SO_KEEPALIVE`
- `SO_BROADCAST`
- `SO_ACCEPTCONN`
- `SO_RCVBUF` (available if *CONFIG_LWIP_SO_RCVBUF* is set)
- `SO_SNDTIMEO` / `SO_RCVTIMEO`

- SO_ERROR (this option is only used with select(), see *Socket Error Handling*)
- SO_TYPE
- SO_NO_CHECK (for UDP sockets only)

IP options Used with level argument IPPROTO_IP.

- IP_TOS
- IP_TTL
- IP_PKTINFO (available if *CONFIG_LWIP_NETBUF_RECVINFO* is set)

For multicast UDP sockets:

- IP_MULTICAST_IF
- IP_MULTICAST_LOOP
- IP_MULTICAST_TTL
- IP_ADD_MEMBERSHIP
- IP_DROP_MEMBERSHIP

TCP options TCP sockets only. Used with level argument IPPROTO_TCP.

- TCP_NODELAY

Options relating to TCP keepalive probes:

- TCP_KEEPALIVE (int value, TCP keepalive period in milliseconds)
- TCP_KEEPIDLE (same as TCP_KEEPALIVE, but the value is in seconds)
- TCP_KEEPINTVL (int value, interval between keepalive probes in seconds)
- TCP_KEEPCNT (int value, number of keepalive probes before timing out)

IPv6 options IPv6 sockets only. Used with level argument IPPROTO_IPV6

- IPV6_CHECKSUM
- IPV6_V6ONLY

For multicast IPv6 UDP sockets:

- IPV6_JOIN_GROUP / IPV6_ADD_MEMBERSHIP
- IPV6_LEAVE_GROUP / IPV6_DROP_MEMBERSHIP
- IPV6_MULTICAST_IF
- IPV6_MULTICAST_HOPS
- IPV6_MULTICAST_LOOP

fcntl

The fcntl() function is a standard API for manipulating options related to a file descriptor. In ESP-IDF, the 虚拟文件系统组件 layer is used to implement this function.

When the file descriptor is a socket, only the following fcntl() values are supported:

- O_NONBLOCK to set/clear non-blocking I/O mode. Also supports O_NDELAY, which is identical to O_NONBLOCK.
- O_RDONLY, O_WRONLY, O_RDWR flags for different read/write modes. These can read via F_GETFL only, they cannot be set using F_SETFL. A TCP socket will return a different mode depending on whether the connection has been closed at either end or is still open at both ends. UDP sockets always return O_RDWR.

ioctls

The ioctl() function provides a semi-standard way to access some internal features of the TCP/IP stack. In ESP-IDF, the 虚拟文件系统组件 layer is used to implement this function.

When the file descriptor is a socket, only the following ioctl() values are supported:

- `FIONREAD` returns the number of bytes of pending data already received in the socket's network buffer.
- `FIONBIO` is an alternative way to set/clear non-blocking I/O status for a socket, equivalent to `fcntl(fd, F_SETFL, O_NONBLOCK, ...)`.

4.25.3 Netconn API

lwIP supports two lower level APIs as well as the BSD Sockets API: the Netconn API and the Raw API.

The lwIP Raw API is designed for single threaded devices and is not supported in ESP-IDF.

The Netconn API is used to implement the BSD Sockets API inside lwIP, and it can also be called directly from ESP-IDF apps. This API has lower resource usage than the BSD Sockets API, in particular it can send and receive data without needing to first copy it into internal lwIP buffers.

重要： Espressif does not test the Netconn API in ESP-IDF. As such, this functionality is *enabled but not supported*. Some functionality may only work correctly when used from the BSD Sockets API.

For more information about the Netconn API, consult lwip/lwip/src/include/lwip/api.h and this wiki page which is part of the unofficial lwIP Application Developers Manual.

4.25.4 lwIP FreeRTOS Task

lwIP creates a dedicated TCP/IP FreeRTOS task to handle socket API requests from other tasks.

A number of configuration items are available to modify the task and the queues ("mailboxes") used to send data to/from the TCP/IP task:

- *CONFIG_LWIP_TCPIP_RECVMBOX_SIZE*
- *CONFIG_LWIP_TCPIP_TASK_STACK_SIZE*
- *CONFIG_LWIP_TCPIP_TASK_AFFINITY*

4.25.5 esp-lwip custom modifications

Additions

The following code is added which is not present in the upstream lwIP release:

Thread-safe sockets It is possible to `close()` a socket from a different thread to the one that created it. The `close()` call will block until any function calls currently using that socket from other tasks have returned.

It is, however, not possible to delete a task while it is actively waiting on `select()` or `poll()` APIs. It is always necessary that these APIs exit before destroying the task, as this might corrupt internal structures and cause subsequent crashes of the lwIP. (These APIs allocate globally referenced callback pointers on stack, so that when the task gets destroyed before unrolling the stack, the lwIP would still hold pointers to the deleted stack)

On demand timers lwIP IGMP and MLD6 features both initialize a timer in order to trigger timeout events at certain times.

The default lwIP implementation is to have these timers enabled all the time, even if no timeout events are active. This increases CPU usage and power consumption when using automatic light sleep mode. `esp-lwip` default behaviour is to set each timer "on demand" so it is only enabled when an event is pending.

To return to the default lwIP behaviour (always-on timers), disable *CONFIG_LWIP_TIMERS_ONDEMAND*.

Abort TCP connections when IP changes *CONFIG_LWIP_TCP_KEEP_CONNECTION_WHEN_IP_CHANGES* is disabled by default. This disables the default lwIP behaviour of keeping TCP connections open if an interface IP changes, in case the interface IP changes back (for example, if an interface connection goes down and comes back up). Enable this option to keep TCP connections open in this case, until they time out normally. This may increase the number of sockets in use if a network interface goes down temporarily.

Additional Socket Options

- Some standard IPV4 and IPV6 multicast socket options are implemented (see *Socket Options*).
- Possible to set IPV6-only UDP and TCP sockets with `IPV6_V6ONLY` socket option (normal lwIP is TCP only).

IP layer features

- IPV4 source based routing implementation is different.
- IPV4 mapped IPV6 addresses are supported.

Limitations

- Calling `send()` or `sendto()` repeatedly on a UDP socket may eventually fail with `errno` equal to `ENOMEM`. This is a limitation of buffer sizes in the lower layer network interface drivers. If all driver transmit buffers are full then UDP transmission will fail. Applications sending a high volume of UDP datagrams who don't wish for any to be dropped by the sender should check for this error code and re-send the datagram after a short delay.

4.25.6 Performance Optimization

TCP/IP performance is a complex subject, and performance can be optimized towards multiple goals. The default settings of ESP-IDF are tuned for a compromise between throughput, latency, and moderate memory usage.

Maximum throughput

Espressif tests ESP-IDF TCP/IP throughput using the wifi/iperf example in an RF sealed enclosure.

The wifi/iperf/sdkconfig.defaults file for the iperf example contains settings known to maximize TCP/IP throughput, usually at the expense of higher RAM usage. To get maximum TCP/IP throughput in an application at the expense of other factors then suggest applying settings from this file into the project sdkconfig.

重要： Suggest applying changes a few at a time and checking the performance each time with a particular application workload.

- If a lot of tasks are competing for CPU time on the system, consider that the lwIP task has configurable CPU affinity (*CONFIG_LWIP_TCPIP_TASK_AFFINITY*) and runs at fixed priority `ESP_TASK_TCPIP_PRIO` (18). Configure competing tasks to be pinned to a different core, or to run at a lower priority.
- If using `select()` function with socket arguments only, setting *CONFIG_LWIP_USE_ONLY_LWIP_SELECT* will make `select()` calls faster.

If using a Wi-Fi network interface, please also refer to *Wi-Fi Buffer Usage*.

Minimum latency

Except for increasing buffer sizes, most changes which increase throughput will also decrease latency by reducing the amount of CPU time spent in lwIP functions.

- For TCP sockets, lwIP supports setting the standard `TCP_NODELAY` flag to disable Nagle's algorithm.

Minimum RAM usage

Most lwIP RAM usage is on-demand, as RAM is allocated from the heap as needed. Therefore, changing lwIP settings to reduce RAM usage may not change RAM usage at idle but can change it at peak.

- Reducing *CONFIG_LWIP_MAX_SOCKETS* reduces the maximum number of sockets in the system. This will also cause TCP sockets in the `WAIT_CLOSE` state to be closed and recycled more rapidly (if needed to open a new socket), further reducing peak RAM usage.
- Reducing *CONFIG_LWIP_TCPIP_RECVMBOX_SIZE*, *CONFIG_LWIP_TCP_RECVMBOX_SIZE* and *CONFIG_LWIP_UDP_RECVMBOX_SIZE* reduce memory usage at the expense of throughput, depending on usage.
- Reducing *CONFIG_LWIP_TCP_MSL*, *CONFIG_LWIP_TCP_FIN_WAIT_TIMEOUT* reduces the maximum segment lifetime in the system. This will also cause TCP sockets in the `TIME_WAIT`, `FIN_WAIT_2` state to be closed and recycled more rapidly

If using Wi-Fi, please also refer to *Wi-Fi Buffer Usage*.

Peak Buffer Usage The peak heap memory that lwIP consumes is the **theoretically-maximum memory** that the lwIP driver consumes. Generally, the peak heap memory that lwIP consumes depends on:

- the memory required to create a UDP connection: lwip_udp_conn
- the memory required to create a TCP connection: lwip_tcp_conn
- the number of UDP connections that the application has: lwip_udp_con_num
- the number of TCP connections that the application has: lwip_tcp_con_num
- the TCP TX window size: lwip_tcp_tx_win_size
- the TCP RX window size: lwip_tcp_rx_win_size

So, the peak heap memory that the LwIP consumes can be calculated with the following formula:

lwip_dynamic_peek_memory = (lwip_udp_con_num * lwip_udp_conn) + (lwip_tcp_con_num * (lwip_tcp_tx_win_size + lwip_tcp_rx_win_size + lwip_tcp_conn))

Some TCP-based applications need only one TCP connection. However, they may choose to close this TCP connection and create a new one when an error (such as a sending failure) occurs. This may result in multiple TCP connections existing in the system simultaneously, because it may take a long time for a TCP connection to close, according to the TCP state machine (refer to RFC793).

4.26 工具

4.26.1 Downloadable Tools

ESP-IDF build process relies on a number of tools: cross-compiler toolchains, CMake build system, and others.

Installing the tools using an OS-specific package manager (like apt, yum, brew, etc.) is the preferred method when the required version of the tool is available. This recommendation is reflected in the Getting Started guide. For example, on Linux and macOS it is recommended to install CMake using an OS package manager.

However, some of the tools are IDF-specific and are not available in OS package repositories. Furthermore, different versions of ESP-IDF require different versions of the tools to operate correctly. To solve these two problems, ESP-IDF provides a set of scripts for downloading and installing the correct versions of tools, and exposing them in the environment.

The rest of the document refers to these downloadable tools simply as "tools". Other kinds of tools used in ESP-IDF are:

- Python scripts bundled with ESP-IDF (such as `idf.py`)
- Python packages installed from PyPI.

The following sections explain the installation method, and provide the list of tools installed on each platform.

注解: This document is provided for advanced users who need to customize their installation, users who wish to understand the installation process, and ESP-IDF developers.

If you are looking for instructions on how to install the tools, see the *Getting Started Guide*.

Tools metadata file

The list of tools and tool versions required for each platform is located in tools/tools.json. The schema of this file is defined by tools/tools_schema.json.

This file is used by tools/idf_tools.py script when installing the tools or setting up the environment variables.

Tools installation directory

`IDF_TOOLS_PATH` environment variable specifies the location where the tools are to be downloaded and installed. If not set, `IDF_TOOLS_PATH` defaults to `HOME/.espressif` on Linux and macOS, and `%USER_PROFILE%\.espressif` on Windows.

Inside `IDF_TOOLS_PATH`, the scripts performing tools installation create the following directories:

- `dist` —where the archives of the tools are downloaded.
- `tools` —where the tools are extracted. The tools are extracted into subdirectories: `tools/TOOL_NAME/VERSION/`. This arrangement allows different versions of tools to be installed side by side.

GitHub Assets Mirror

Most of the tools downloaded by the tools installer are GitHub Release Assets, which are files attached to a software release on GitHub.

If GitHub downloads are inaccessible or slow to access, it's possible to configure a GitHub assets mirror.

To use Espressif's download server, set the environment variable `IDF_GITHUB_ASSETS` to `dl.espressif.com/github_assets`. When the install process is downloading a tool from `github.com`, the URL will be rewritten to use this server instead.

Any mirror server can be used provided the URL matches the `github.com` download URL format: the install process will replace `https://github.com` with `https://${IDF_GITHUB_ASSETS}` for any GitHub asset URL that it downloads.

注解： The Espressif download server doesn't currently mirror everything from GitHub, it only mirrors files attached as Assets to some releases as well as source archives for some releases.

`idf_tools.py` script

tools/idf_tools.py script bundled with ESP-IDF performs several functions:

- `install`: Download the tool into `${IDF_TOOLS_PATH}/dist` directory, extract it into `${IDF_TOOLS_PATH}/tools/TOOL_NAME/VERSION`.
  `install` command accepts the list of tools to install, in `TOOL_NAME` or `TOOL_NAME@VERSION` format. If `all` is given, all the tools (required and optional ones) are installed. If no argument or `required` is given, only the required tools are installed.
- `download`: Similar to `install` but doesn't extract the tools. An optional `--platform` argument may be used to download the tools for the specific platform.
- `export`: Lists the environment variables which need to be set to use the installed tools. For most of the tools, setting `PATH` environment variable is sufficient, but some tools require extra environment variables.
  The environment variables can be listed in either of `shell` or `key-value` formats, set by `--format` parameter:
  - `shell` produces output suitable for evaluation in the shell. For example,

```
export PATH="/home/user/.espressif/tools/tool/v1.0.0/bin:$PATH"
```

on Linux and macOS, and

```
set "PATH=C:\Users\user\.espressif\tools\v1.0.0\bin;%PATH%"
```

on Windows.

注解： Exporting environment variables in Powershell format is not supported at the moment. `key-value` format may be used instead.

The output of this command may be used to update the environment variables, if the shell supports this. For example:

```
eval $($IDF_PATH/tools/idf_tools.py export)
```

- `key-value` produces output in *VARIABLE=VALUE* format, suitable for parsing by other scripts:

```
PATH=/home/user/.espressif/tools/tool/v1.0.0:$PATH
```

Note that the script consuming this output has to perform expansion of `$VAR` or `%VAR%` patterns found in the output.

- `list`: Lists the known versions of the tools, and indicates which ones are installed.
- `check`: For each tool, checks whether the tool is available in the system path and in `IDF_TOOLS_PATH`.

Install scripts

Shell-specific user-facing scripts are provided in the root of ESP-IDF repository to facilitate tools installation. These are:

- `install.bat` for Windows Command Prompt
- `install.ps1` for Powershell
- `install.sh` for Bash

Aside from downloading and installing the tools into `IDF_TOOLS_PATH`, these scripts prepare a Python virtual environment, and install the required packages into that environment.

Export scripts

Since the installed tools are not permanently added into the user or system `PATH` environment variable, an extra step is required to use them in the command line. The following scripts modify the environment variables in the current shell to make the correct versions of the tools available:

- `export.bat` for Windows Command Prompt
- `export.ps1` for Powershell
- `export.sh` for Bash

注解： To modify the shell environment in Bash, `export.sh` must be "sourced" : `. ./export.sh` (note the leading dot and space).

`export.sh` may be used with shells other than Bash (such as zsh). However in this case the `IDF_PATH` environment variable must be set before running the script. When used in Bash, the script will guess the `IDF_PATH` value from its own location.

In addition to calling `idf_tools.py`, these scripts list the directories which have been added to the `PATH`.

Other installation methods

Depending on the environment, more user-friendly wrappers for `idf_tools.py` are provided:

- *IDF Tools installer for Windows* can download and install the tools. Internally the installer uses `idf_tools.py`.
- *Eclipse plugin for ESP-IDF* includes a menu item to set up the tools. Internally the plugin calls `idf_tools.py`.
- Visual Studio Code extension for ESP-IDF includes an onboarding flow. This flow helps setting up the tools. Although the extension does not rely on `idf_tools.py`, the same installation method is used.

Custom installation

Although the methods above are recommended for ESP-IDF users, they are not a must for building ESP-IDF applications. ESP-IDF build system expects that all the necessary tools are installed somewhere, and made available in the `PATH`.

List of IDF Tools

xtensa-esp32-elf Toolchain for Xtensa (ESP32) based on GCC

License: GPL-3.0-with-GCC-exception

More info: https://github.com/espressif/crosstool-NG

Platform	Required	Download
linux-amd64	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-linux-amd64.tar.gz SHA256: 674080a12f9c5ebe5a3a5ce51c6deaeffe6dfb06d6416233df86f25b574e9279
linux-armel	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-linux-armel.tar.gz SHA256: 6771e011dffa2438ef84ff3474538b4a69df8f9d4cfae3b3707ca31c782ed7db
linux-i686	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-linux-i686.tar.gz SHA256: 076b7e05304e26aa6ec105c9e0dc74addca079bc2cae6e42ee7575c5ded29877
macos	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-macos.tar.gz SHA256: 6845f786303b26c4a55ede57487ba65bd25737232fe6104be03f25bb62f4631c
win32	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-win32.zip SHA256: 81cecd5493a3fcf2118977f3fd60bd0a13a4aeac8fe6760d912f96d2c34fab66
win64	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32-elf-gcc8_4_0-esp-2020r3-win64.zip SHA256: 58419852fefb7111fdec056ac2fd7c4bd09c1f4c17610a761a97b788413527cf

xtensa-esp32s2-elf Toolchain for Xtensa (ESP32-S2) based on GCC

License: GPL-3.0-with-GCC-exception

More info: https://github.com/espressif/crosstool-NG

Platform	Required	Download
linux-amd64	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32s2-elf-gcc8_4_0-esp-2020r3-linux-amd64.tar.gz SHA256: 40fafa47045167feda0cd07827db5207ebfeb4a3b6b24475957a921bc92805ed
linux-armel	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32s2-elf-gcc8_4_0-esp-2020r3-linux-armel.tar.gz SHA256: 6c1efec4c7829202279388ccb388e8a17a34464bc351d677c4f04d95ea4b4ce0
linux-i686	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32s2-elf-gcc8_4_0-esp-2020r3-linux-i686.tar.gz SHA256: bd3a91166206a1a7ff7c572e15389e1938c3cdce588032a5e915be677a945638
macos	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32s2-elf-gcc8_4_0-esp-2020r3-macos.tar.gz SHA256: fe19b0c873879d8d89ec040f4db04a3ab27d769d3fd5f55fe59a28b6b111d09c
win32	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32s2-elf-gcc8_4_0-esp-2020r3-win32.zip SHA256: d078d614ae864ae4a37fcb5b83323af0a5cfdbd8243607664becdd0f977a1e33
win64	required	https://github.com/espressif/crosstool-NG/releases/download/esp-2020r3/xtensa-esp32s2-elf-gcc8_4_0-esp-2020r3-win64.zip SHA256: 6ea78b72ec52330b69af91dbe6c77c22bdc827817f044aa30306270453032bb4

esp32ulp-elf Toolchain for ESP32 ULP coprocessor

License: GPL-2.0-or-later

More info: https://github.com/espressif/binutils-esp32ulp

Platform	Required	Download
linux-amd64	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32ulp-linux-amd64-2.28.51-esp-20191205.tar.gz SHA256: 3016c4fc551181175bd9979869bc1d1f28fa8efa25a0e29ad7f833fca4bc03d7
linux-armel	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32ulp-linux-armel-2.28.51-esp-20191205.tar.gz SHA256: 88967c086a6e71834282d9ae05841ee74dae1815f27807b25cdd3f7775a47101
macos	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32ulp-macos-2.28.51-esp-20191205.tar.gz SHA256: a35d9e7a0445247c7fc9dccd3fbc35682f0fafc28beeb10c94b59466317190c4
win32	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32ulp-win32-2.28.51-esp-20191205.zip SHA256: bade309353a9f0a4e5cc03bfe84845e33205f05502c4b199195e871ded271ab5
win64	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32ulp-win32-2.28.51-esp-20191205.zip SHA256: bade309353a9f0a4e5cc03bfe84845e33205f05502c4b199195e871ded271ab5

esp32s2ulp-elf Toolchain for ESP32-S2 ULP coprocessor

License: GPL-2.0-or-later

More info: https://github.com/espressif/binutils-esp32ulp

Platform	Required	Download
linux-amd64	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32s2ulp-linux-amd64-2.28.51-esp-20191205.tar.gz SHA256: df7b2ff6c7c718a7cbe3b4b6dbcd68180d835d164d1913bc4698fd3781b9a466
linux-armel	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32s2ulp-linux-armel-2.28.51-esp-20191205.tar.gz SHA256: 893b213c8f716d455a6efb2b08b6cf1bc34d08b78ee19c31e82ac44b1b45417e
macos	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32s2ulp-macos-2.28.51-esp-20191205.tar.gz SHA256: 5a9bb678a5246638cbda303f523d9bb8121a9a24dc01ecb22c21c46c41184155
win32	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32s2ulp-win32-2.28.51-esp-20191205.zip SHA256: 587de59fbb469a39f96168ae3eaa9f06b2601e6e0543c87eaf1bd97f23e5c4ca
win64	required	https://github.com/espressif/binutils-esp32ulp/releases/download/v2.28.51-esp-20191205/binutils-esp32s2ulp-win32-2.28.51-esp-20191205.zip SHA256: 587de59fbb469a39f96168ae3eaa9f06b2601e6e0543c87eaf1bd97f23e5c4ca

cmake CMake build system

On Linux and macOS, it is recommended to install CMake using the OS package manager. However, for convenience it is possible to install CMake using idf_tools.py along with the other tools.

License: BSD-3-Clause

More info: https://github.com/Kitware/CMake

Platform	Required	Download
linux-amd64	optional	https://github.com/Kitware/CMake/releases/download/v3.16.4/cmake-3.16.4-Linux-x86_64.tar.gz SHA256: 12a577aa04b6639766ae908f33cf70baefc11ac4499b8b1c8812d99f05fb6a02
linux-armel	optional	https://dl.espressif.com/dl/cmake/cmake-3.20.3-Linux-armv7l.tar.gz SHA256: f8bd050c2745f0dcc4b7cef9738bbfef775950a10f5bd377abb0062835e669dc
macos	optional	https://github.com/Kitware/CMake/releases/download/v3.16.4/cmake-3.16.4-Darwin-x86_64.tar.gz SHA256: f60e0ef96da48725cd8da7d6abe83cd9501167aa51625c90dd4d31081a631279
win32	required	https://github.com/Kitware/CMake/releases/download/v3.16.4/cmake-3.16.4-win64-x64.zip SHA256: f37963bcfcebdfe5864926a3623f6c21220c35790c39cd65e64bd521cbb39c55
win64	required	https://github.com/Kitware/CMake/releases/download/v3.16.4/cmake-3.16.4-win64-x64.zip SHA256: f37963bcfcebdfe5864926a3623f6c21220c35790c39cd65e64bd521cbb39c55

openocd-esp32 OpenOCD for ESP32

License: GPL-2.0-only

More info: https://github.com/espressif/openocd-esp32

Platform	Required	Download
linux-amd64	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-linux-amd64-0.11.0-esp32-20220706.tar.gz SHA256: 26f1f18dd93eb70a13203848d3fb1cc2e0de1fd6749c7dd771b2de8709735aed
linux-arm64	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-linux-arm64-0.11.0-esp32-20220706.tar.gz SHA256: f97792bc2852937ec0accb9f0eb2e49926c0f747a71f101a4e34aed75d2c6fcc
linux-armel	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-linux-armel-0.11.0-esp32-20220706.tar.gz SHA256: 27e4c628994cf342e7fc5b07f49ca5533ba892fd8a150683a382a08758c3dfbe
linux-armhf	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-linux-armhf-0.11.0-esp32-20220706.tar.gz SHA256: 7f3b57332104e8b8e6194553365a70a9d3754878cfc063d5dc5d839513a63de9
macos	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-macos-0.11.0-esp32-20220706.tar.gz SHA256: 333ee2ec3c9b5dc6ad4509faae55335cdea7f8bf83a56bfcf5327e4497c8538a
win32	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-win32-0.11.0-esp32-20220706.zip SHA256: c3d39eb4365a9947e71f1d3780ce031185bc6437f21186568a5c05f23f57a8d0
win64	required	https://github.com/espressif/openocd-esp32/releases/download/v0.11.0-esp32-20220706/openocd-esp32-win32-0.11.0-esp32-20220706.zip SHA256: c3d39eb4365a9947e71f1d3780ce031185bc6437f21186568a5c05f23f57a8d0

ninja Ninja build system

On Linux and macOS, it is recommended to install ninja using the OS package manager. However, for convenience it is possible to install ninja using idf_tools.py along with the other tools.

License: Apache-2.0

More info: https://github.com/ninja-build/ninja

Platform	Required	Download
linux-amd64	optional	https://dl.espressif.com/dl/ninja-1.10.0-linux64.tar.gz SHA256: 4b2ad65db438595813b981db921f675f5c775c302b34dc85710ffddf07ec9033
macos	optional	https://dl.espressif.com/dl/ninja-1.10.0-osx.tar.gz SHA256: 6cd22e2c5fc654282d26b485d9b3d68e242b09a22c2e73a253f2a4a7cfd6774c
win64	required	https://dl.espressif.com/dl/ninja-1.10.0-win64.zip SHA256: 919fd158c16bf135e8a850bb4046ec1ce28a7439ee08b977cd0b7f6b3463d178

idf-exe IDF wrapper tool for Windows

License: Apache-2.0

More info: tools/windows/idf_exe

Platform	Required	Download
win32	required	https://dl.espressif.com/dl/idf-exe-v1.0.1.zip SHA256: 53eb6aaaf034cc7ed1a97d5c577afa0f99815b7793905e9408e74012d357d04a
win64	required	https://dl.espressif.com/dl/idf-exe-v1.0.1.zip SHA256: 53eb6aaaf034cc7ed1a97d5c577afa0f99815b7793905e9408e74012d357d04a

ccache Ccache (compiler cache)

License: GPL-3.0-or-later

More info: https://github.com/ccache/ccache

Platform	Required	Download
win64	required	https://dl.espressif.com/dl/ccache-3.7-w64.zip SHA256: 37e833f3f354f1145503533e776c1bd44ec2e77ff8a2476a1d2039b0b10c78d6

dfu-util dfu-util (Device Firmware Upgrade Utilities)

License: GPL-2.0-only

More info: http://dfu-util.sourceforge.net/

Platform	Required	Download
win64	required	https://dl.espressif.com/dl/dfu-util-0.9-win64.zip SHA256: 5816d7ec68ef3ac07b5ac9fb9837c57d2efe45b6a80a2f2bbe6b40b1c15c470e

4.26.2 IDF Docker Image

IDF Docker image (`espressif/idf`) is intended for building applications and libraries with specific versions of ESP-IDF, when doing automated builds.

The image contains:

- Common utilities such as git, wget, curl, zip.
- Python 3.6 or newer.
- A copy of a specific version of ESP-IDF (see below for information about versions). `IDF_PATH` environment variable is set, and points to ESP-IDF location in the container.
- All the build tools required for the specific version of ESP-IDF: CMake, make, ninja, cross-compiler toolchains, etc.
- All Python packages required by ESP-IDF are installed in a virtual environment.

The image entrypoint sets up `PATH` environment variable to point to the correct version of tools, and activates the Python virtual environment. As a result, the environment is ready to use the ESP-IDF build system.

The image can also be used as a base for custom images, if additional utilities are required.

Tags

Multiple tags of this image are maintained:

- `latest`: tracks `master` branch of ESP-IDF
- `vX.Y`: corresponds to ESP-IDF release `vX.Y`
- `release-vX.Y`: tracks `release/vX.Y` branch of ESP-IDF

注解: Versions of ESP-IDF released before this feature was introduced do not have corresponding Docker image versions. You can check the up-to-date list of available tags at https://hub.docker.com/r/espressif/idf/tags.

Usage

Setting up Docker Before using the `espressif/idf` Docker image locally, make sure you have Docker installed. Follow the instructions at https://docs.docker.com/install/, if it is not installed yet.

If using the image in CI environment, consult the documentation of your CI service on how to specify the image used for the build process.

Building a project with CMake In the project directory, run:

```
docker run --rm -v $PWD:/project -w /project espressif/idf idf.py build
```

The above command explained:

- `docker run`: runs a Docker image. It is a shorter form of the command `docker container run`.
- `--rm`: removes the container when the build is finished
- `-v $PWD:/project`: mounts the current directory on the host (`$PWD`) as `/project` directory in the container
- `espressif/idf`: uses Docker image `espressif/idf` with tag `latest` (implicitly added by Docker when no tag is specified)
- `idf.py build`: runs this command inside the container

To build with a specific docker image tag, specify it as `espressif/idf:TAG`, for example:

```
docker run --rm -v $PWD:/project -w /project espressif/idf:release-v4.4 idf.py␣
↪build
```

You can check the up-to-date list of available tags at https://hub.docker.com/r/espressif/idf/tags.

Building a project with GNU Make Same as for CMake, except that the build command is different:

```
docker run --rm -v $PWD:/project -w /project espressif/idf make defconfig all -j4
```

注解: If the `sdkconfig` file does not exist, the default behavior of GNU Make build system is to open the menuconfig UI. This may be not desired in automated build environments. To ensure that the `sdkconfig` file exists, `defconfig` target is added before `all`.

If you intend to build the same project repeatedly, you may bind the `tools/kconfig` directory of ESP-IDF to a named volume. This will prevent Kconfig tools, located in ESP-IDF directory, from being rebuilt, causing a rebuild of the rest of the project:

```
docker run --rm -v $PWD:/project -v kconfig:/opt/esp/idf/tools/kconfig -w /project␣
↪espressif/idf make defconfig all -j4
```

If you need clean up the `kconfig` volume, run `docker volume rm kconfig`.

Binding the `tools/kconfig` directory to a volume is not necessary when using the CMake build system.

Using the image interactively It is also possible to do builds interactively, to debug build issues or test the automated build scripts. Start the container with *-i -t* flags:

```
docker run --rm -v $PWD:/project -w /project -it espressif/idf
```

Then inside the container, use `idf.py` as usual:

```
idf.py menuconfig
idf.py build
```

注解: Commands which communicate with the development board, such as `idf.py flash` and `idf.py monitor` will not work in the container unless the serial port is passed through into the container. However currently this is not possible with Docker for Windows (https://github.com/docker/for-win/issues/1018) and Docker for Mac (https://github.com/docker/for-mac/issues/900).

Building custom images

The Dockerfile in ESP-IDF repository provides several build arguments which can be used to customize the Docker image:

- `IDF_CLONE_URL`: URL of the repository to clone ESP-IDF from. Can be set to a custom URL when working with a fork of ESP-IDF. Default is `https://github.com/espressif/esp-idf.git`.
- `IDF_CLONE_BRANCH_OR_TAG`: Name of a git branch or tag use when cloning ESP-IDF. This value is passed to `git clone` command using the `--branch` argument. Default is `master`.
- `IDF_CHECKOUT_REF`: If this argument is set to a non-empty value, `git checkout $IDF_CHECKOUT_REF` command will be performed after cloning. This argument can be set to the SHA of the specific commit to check out, for example if some specific commit on a release branch is desired.
- `IDF_CLONE_SHALLOW`: If this argument is set to a non-empty value, `--depth=1 --shallow-submodules` arguments will be used when performing `git clone`. This significantly reduces the amount of data downloaded and the size of the resulting Docker image. However, if switching to a different branch in such a "shallow" repository is necessary, an additional `git fetch origin <branch>` command must be executed first.
- `IDF_INSTALL_TARGETS`: Comma-separated list of IDF targets to install toolchains for, or `all` to install toolchains for all targets. Selecting specific targets reduces the amount of data downloaded and the size of the resulting Docker image. Default is `all`.

To use these arguments, pass them via the `--build-arg` command line option. For example, the following command will build a Docker image with a shallow clone of ESP-IDF v4.4.1 and tools for ESP32-C3, only:

```
docker build -t idf-custom:v4.4.1-esp32c3 \
    --build-arg IDF_CLONE_BRANCH_OR_TAG=v4.4.1 \
    --build-arg IDF_CLONE_SHALLOW=1 \
    --build-arg IDF_INSTALL_TARGETS=esp32c3 \
    tools/docker
```

4.26.3 IDF Component Manager

The IDF Component manager is a tool that downloads dependencies for any ESP-IDF CMake project. The download happens automatically during a run of CMake. It can source components either from the component registry or from a git repository.

A list of components can be found on https://components.espressif.com/

Using with a project

Dependencies for each component in the project are defined in a separate manifest file named `idf_component.yml` placed in the root of the component. The manifest file template can be created for a component by running `idf.py create-manifest --component=my_component`. When a new manifest is added to one of the components in the project it's necessary to reconfigure it manually by running `idf.py reconfigure`. Then build will track changes in `idf_component.yml` manifests and automatically triggers CMake when necessary.

There is an example application: example:*build_system/cmake/component_manager* that uses components installed by the component manager.

It's not necessary to have a manifest for components that don't need any managed dependencies.

When CMake configures the project (e.g. `idf.py reconfigure`) component manager does a few things:

- Processes `idf_component.yml` manifests for every component in the project and recursively solves dependencies
- Creates a `dependencies.lock` file in the root of the project with a full list of dependencies
- Downloads all dependencies to the `managed_components` directory

The lock-file `dependencies.lock` and content of `managed_components` directory is not supposed to be modified by a user. When the component manager runs it always make sure they are up to date. If these files

were accidentally modified it’ s possible to re-run the component manager by triggering CMake with `idf.py reconfigure`

Defining dependencies in the manifest

```
dependencies:
  # Required IDF version
  idf: ">=4.1"
  # Defining a dependency from the registry:
  # https://components.espressif.com/component/example/cmp
  example/cmp: ">=1.0.0"

  # # Other ways to define dependencies
  #
  # # For components maintained by Espressif only name can be used.
  # # Same as `espressif/cmp`
  # component: "~1.0.0"
  #
  # # Or in a longer form with extra parameters
  # component2:
  #   version: ">=2.0.0"
  #
  #   # For transient dependencies `public` flag can be set.
  #   # `public` flag doesn't affect the `main` component.
  #   # All dependencies of `main` are public by default.
  #   public: true
  #
  #   # For components hosted on non-default registry:
  #   service_url: "https://componentregistry.company.com"
  #
  # # For components in git repository:
  # test_component:
  #   path: test_component
  #   git: ssh://git@gitlab.com/user/components.git
  #
  # # For test projects during component development
  # # components can be used from a local directory
  # # with relative or absolute path
  # some_local_component:
  #   path: ../../projects/component
```

Disabling the Component Manager

The component manager can be explicitly disabled by setting `IDF_COMPONENT_MANAGER` environment variable to `0`.

Chapter 5

Libraries and Frameworks

5.1 Cloud Frameworks

ESP32-S2 supports multiple cloud frameworks using agents built on top of ESP-IDF. Here are the pointers to various supported cloud frameworks' agents and examples:

5.1.1 AWS IoT

https://github.com/espressif/esp-aws-iot is an open source repository for ESP32-S2 based on Amazon Web Services' aws-iot-device-sdk-embedded-C.

5.1.2 Azure IoT

https://github.com/espressif/esp-azure is an open source repository for ESP32-S2 based on Microsoft Azure's azure-iot-sdk-c SDK.

5.1.3 Google IoT Core

https://github.com/espressif/esp-google-iot is an open source repository for ESP32-S2 based on Google's iot-device-sdk-embedded-c SDK.

5.1.4 Aliyun IoT

https://github.com/espressif/esp-aliyun is an open source repository for ESP32-S2 based on Aliyun's iotkit-embedded SDK.

5.1.5 Joylink IoT

https://github.com/espressif/esp-joylink is an open source repository for ESP32-S2 based on Joylink's joylink_dev_sdk SDK.

5.1.6 Tencent IoT

https://github.com/espressif/esp-welink is an open source repository for ESP32-S2 based on Tencent's welink SDK.

5.1.7 Tencentyun IoT

https://github.com/espressif/esp-qcloud is an open source repository for ESP32-S2 based on Tencentyun's qcloud-iot-sdk-embedded-c SDK.

5.1.8 Baidu IoT

https://github.com/espressif/esp-baidu-iot is an open source repository for ESP32-S2 based on Baidu's iot-sdk-c SDK.

Chapter 6

Contributions Guide

We welcome contributions to the esp-idf project!

6.1 How to Contribute

Contributions to esp-idf - fixing bugs, adding features, adding documentation - are welcome. We accept contributions via Github Pull Requests.

6.2 Before Contributing

Before sending us a Pull Request, please consider this list of points:

- Is the contribution entirely your own work, or already licensed under an Apache License 2.0 compatible Open Source License? If not then we unfortunately cannot accept it.
- Does any new code conform to the esp-idf *Style Guide*?
- Does the code documentation follow requirements in 编写代码文档?
- Is the code adequately commented for people to understand how it is structured?
- Is there documentation or examples that go with code contributions? There are additional suggestions for writing good examples in examples readme.
- Are comments and documentation written in clear English, with no spelling or grammar errors?
- Example contributions are also welcome. Please check the 创建示例项目 guide for these.
- If the contribution contains multiple commits, are they grouped together into logical changes (one major change per pull request)? Are any commits with names like "fixed typo" squashed into previous commits?
- If you're unsure about any of these points, please open the Pull Request anyhow and then ask us for feedback.

6.3 Pull Request Process

After you open the Pull Request, there will probably be some discussion in the comments field of the request itself.

Once the Pull Request is ready to merge, it will first be merged into our internal git system for in-house automated testing.

If this process passes, it will be merged onto the public github repository.

6.4 Legal Part

Before a contribution can be accepted, you will need to sign our *Contributor Agreement*. You will be prompted for this automatically as part of the Pull Request process.

6.5 Related Documents

6.5.1 Espressif IoT Development Framework Style Guide

About This Guide

Purpose of this style guide is to encourage use of common coding practices within the ESP-IDF.

Style guide is a set of rules which are aimed to help create readable, maintainable, and robust code. By writing code which looks the same way across the code base we help others read and comprehend the code. By using same conventions for spaces and newlines we reduce chances that future changes will produce huge unreadable diffs. By following common patterns for module structure and by using language features consistently we help others understand code behavior.

We try to keep rules simple enough, which means that they can not cover all potential cases. In some cases one has to bend these simple rules to achieve readability, maintainability, or robustness.

When doing modifications to third-party code used in ESP-IDF, follow the way that particular project is written. That will help propose useful changes for merging into upstream project.

C Code Formatting

Indentation Use 4 spaces for each indentation level. Don't use tabs for indentation. Configure the editor to emit 4 spaces each time you press tab key.

Vertical Space Place one empty line between functions. Don't begin or end a function with an empty line.

```
void function1()
{
    do_one_thing();
    do_another_thing();
                                    // INCORRECT, don't place empty line here
}
                                    // place empty line here
void function2()
{
                                    // INCORRECT, don't use an empty line here
    int var = 0;
    while (var < SOME_CONSTANT) {
        do_stuff(&var);
    }
}
```

The maximum line length is 120 characters as long as it doesn't seriously affect the readability.

Horizontal Space Always add single space after conditional and loop keywords:

```
if (condition) {     // correct
    // ...
}

switch (n) {         // correct
    case 0:
        // ...
}

for(int i = 0; i < CONST; ++i) {     // INCORRECT
    // ...
}
```

Add single space around binary operators. No space is necessary for unary operators. It is okay to drop space around multiply and divide operators:

```
const int y = y0 + (x - x0) * (y1 - y0) / (x1 - x0);     // correct

const int y = y0 + (x - x0)*(y1 - y0)/(x1 - x0);         // also okay

int y_cur = -y;                                          // correct
++y_cur;

const int y = y0+(x-x0)*(y1-y0)/(x1-x0);                 // INCORRECT
```

No space is necessary around `.` and `->` operators.

Sometimes adding horizontal space within a line can help make code more readable. For example, you can add space to align function arguments:

```
gpio_matrix_in(PIN_CAM_D6,   I2S0I_DATA_IN14_IDX, false);
gpio_matrix_in(PIN_CAM_D7,   I2S0I_DATA_IN15_IDX, false);
gpio_matrix_in(PIN_CAM_HREF, I2S0I_H_ENABLE_IDX,  false);
gpio_matrix_in(PIN_CAM_PCLK, I2S0I_DATA_IN15_IDX, false);
```

Note however that if someone goes to add new line with a longer identifier as first argument (e.g. `PIN_CAM_VSYNC`), it will not fit. So other lines would have to be realigned, adding meaningless changes to the commit.

Therefore, use horizontal alignment sparingly, especially if you expect new lines to be added to the list later.

Never use TAB characters for horizontal alignment.

Never add trailing whitespace at the end of the line.

Braces

- Function definition should have a brace on a separate line:

```
// This is correct:
void function(int arg)
{

}

// NOT like this:
void function(int arg) {

}
```

- Within a function, place opening brace on the same line with conditional and loop statements:

```
if (condition) {
    do_one();
} else if (other_condition) {
    do_two();
}
```

Comments Use `//` for single line comments. For multi-line comments it is okay to use either `//` on each line or a `/* */` block.

Although not directly related to formatting, here are a few notes about using comments effectively.

- Don't use single comments to disable some functionality:

```
void init_something()
{
    setup_dma();
    // load_resources();                    // WHY is this thing commented, asks␣
↪the reader?
    start_timer();
}
```

- If some code is no longer required, remove it completely. If you need it you can always look it up in git history of this file. If you disable some call because of temporary reasons, with an intention to restore it in the future, add explanation on the adjacent line:

```
void init_something()
{
    setup_dma();
    // TODO: we should load resources here, but loader is not fully integrated␣
↪yet.
    // load_resources();
    start_timer();
}
```

- Same goes for `#if 0 ... #endif` blocks. Remove code block completely if it is not used. Otherwise, add comment explaining why the block is disabled. Don't use `#if 0 ... #endif` or comments to store code snippets which you may need in the future.
- Don't add trivial comments about authorship and change date. You can always look up who modified any given line using git. E.g. this comment adds clutter to the code without adding any useful information:

```
void init_something()
{
    setup_dma();
    // XXX add 2016-09-01
    init_dma_list();
    fill_dma_item(0);
    // end XXX add
    start_timer();
}
```

Line Endings Commits should only contain files with LF (Unix style) endings.

Windows users can configure git to check out CRLF (Windows style) endings locally and commit LF endings by setting the `core.autocrlf` setting. *Github has a document about setting this option <github-line-endings>*. However because MSYS2 uses Unix-style line endings, it is often easier to configure your text editor to use LF (Unix style) endings when editing ESP-IDF source files.

If you accidentally have some commits in your branch that add LF endings, you can convert them to Unix by running this command in an MSYS2 or Unix terminal (change directory to the IDF working directory and check the correct branch is currently checked out, beforehand):

```
git rebase --exec 'git diff-tree --no-commit-id --name-only -r HEAD | xargs␣
↪dos2unix && git commit -a --amend --no-edit --allow-empty' master
```

(Note that this line rebases on master, change the branch name at the end to rebase on another branch.)

For updating a single commit, it's possible to run `dos2unix FILENAME` and then run `git commit --amend`

Formatting Your Code You can use `astyle` program to format your code according to the above recommendations.

If you are writing a file from scratch, or doing a complete rewrite, feel free to re-format the entire file. If you are changing a small portion of file, don't re-format the code you didn't change. This will help others when they review your changes.

To re-format a file, run:

```
tools/format.sh components/my_component/file.c
```

C++ Code Formatting

The same rules as for C apply. Where they are not enough, apply the following rules.

File Naming C++ Header files have the extension `.hpp`. C++ source files have the extension `.cpp`. The latter is important for the compiler to distiguish them from normal C source files.

Naming

- **Class and struct** names shall be written in `CamelCase` with a capital letter as beginning. Member variables and methods shall be in `snake_case`.
- **Namespaces** shall be in lower `snake_case`.
- **Templates** are specified in the line above the function declaration.
- Interfaces in terms of Object-Oriented Programming shall be named without the suffix `...Interface`. Later, this makes it easier to extract interfaces from normal classes and vice versa without making a breaking change.

Member Order in Classes In order of precedence:

- First put the public members, then the protected, then private ones. Omit public, protected or private sections without any members.
- First put constructors/destructors, then member functions, then member variables.

For example:

```
class ForExample {
public:
    // first constructors, then default constructor, then destructor
    ForExample(double example_factor_arg);
    ForExample();
    ~ForExample();

    // then remaining pubic methods
    set_example_factor(double example_factor_arg);

    // then public member variables
    uint32_t public_data_member;

private:
    // first private methods
    void internal_method();

    // then private member variables
    double example_factor;
};
```

Spacing

- Don't indent inside namespaces.
- Put `public`, `protected` and `private` labels at the same indentation level as the corresponding `class` label.

Simple Example

```
// file spaceship.h
#ifndef SPACESHIP_H_
#define SPACESHIP_H_
#include <cstdlib>

namespace spaceships {

class SpaceShip {
public:
    SpaceShip(size_t crew);
    size_t get_crew_size() const;

private:
    const size_t crew;
};

class SpaceShuttle : public SpaceShip {
public:
    SpaceShuttle();
};

class Sojuz : public SpaceShip {
public:
    Sojuz();
};

template <typename T>
class CargoShip {
public:
    CargoShip(const T &cargo);

private:
    T cargo;
};

} // namespace spaceships

#endif // SPACESHIP_H_

// file spaceship.cpp
#include "spaceship.h"

namespace spaceships {

// Putting the curly braces in the same line for constructors is OK if it only␣
↪initializes
// values in the initializer list
SpaceShip::SpaceShip(size_t crew) : crew(crew) { }

size_t SpaceShip::get_crew_size() const
{
    return crew;
}

SpaceShuttle::SpaceShuttle() : SpaceShip(7)
{
    // doing further initialization
}

Sojuz::Sojuz() : SpaceShip(3)
{
```

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```
    // doing further initialization
}

template <typename T>
CargoShip<T>::CargoShip(const T &cargo) : cargo(cargo) { }

} // namespace spaceships
```

CMake Code Style

- Indent with four spaces.
- Maximum line length 120 characters. When splitting lines, try to focus on readability where possible (for example, by pairing up keyword/argument pairs on individual lines).
- Don't put anything in the optional parentheses after `endforeach()`, `endif()`, etc.
- Use lowercase (`with_underscores`) for command, function, and macro names.
- For locally scoped variables, use lowercase (`with_underscores`).
- For globally scoped variables, use uppercase (`WITH_UNDERSCORES`).
- Otherwise follow the defaults of the cmake-lint project.

Configuring the Code Style for a Project Using EditorConfig

EditorConfig helps developers define and maintain consistent coding styles between different editors and IDEs. The EditorConfig project consists of a file format for defining coding styles and a collection of text editor plugins that enable editors to read the file format and adhere to defined styles. EditorConfig files are easily readable and they work nicely with version control systems.

For more information, see EditorConfig Website.

Documenting Code

Please see the guide here: 编写代码文档.

Naming

- Any variable or function which is only used in a single source file should be declared `static`.
- Public names (non-static variables and functions) should be namespaced with a per-component or per-unit prefix, to avoid naming collisions. ie `esp_vfs_register()` or `esp_console_run()`. Starting the prefix with `esp_` for Espressif-specific names is optional, but should be consistent with any other names in the same component.
- Static variables should be prefixed with `s_` for easy identification. For example, `static bool s_invert`.
- Avoid unnecessary abbreviations (ie shortening `data` to `dat`), unless the resulting name would otherwise be very long.

Structure

To be written.

Language Features

To be written.

6.5.2 编写代码文档

本文简要介绍了 espressif/esp-idf 项目库采用的文件风格以及如何在项目库中添加新文件。

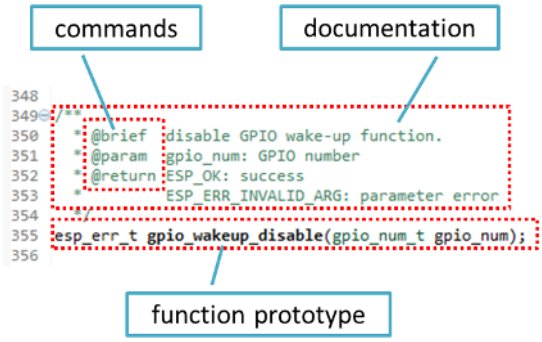
概述

在项目库内编写代码文档时，请遵循 Doxygen 代码注释风格。要采用这一风格，您可以将 `@param` 等特殊命令插入到标准注释块中，比如：

```
/**
 * @param ratio this is oxygen to air ratio
 */
```

Doxygen 会解析代码，提取命令和后续文本，生成代码文档。

注释块通常包含对功能的记述，如下所示。

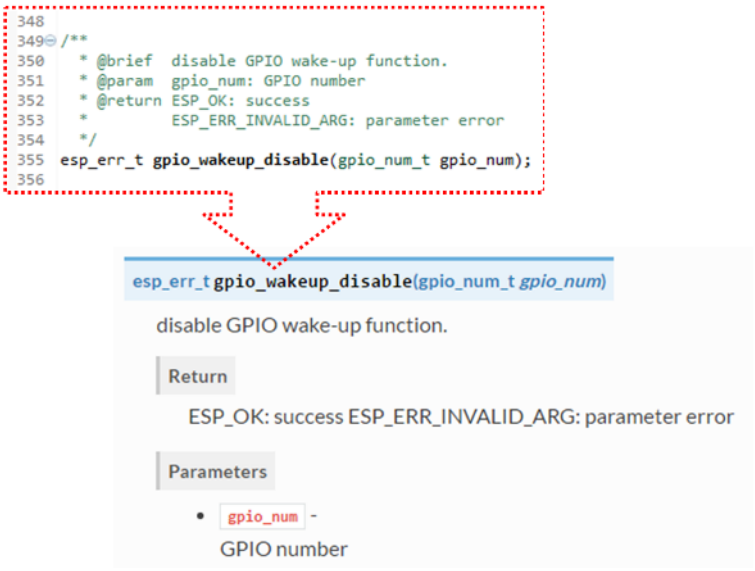


Doxygen 支持多种排版风格，对于文档中可以包含的细节非常灵活。请参考数据丰富、条理清晰的 Doxygen 手册 熟悉 Doxygen 特性。

为什么需要 Doxygen?

使用 Doxygen 的最终目的是确保所有代码编写风格一致，以便在代码变更时使用 Sphinx 和 Breathe 等工具协助筹备、自动更新 API 文档。

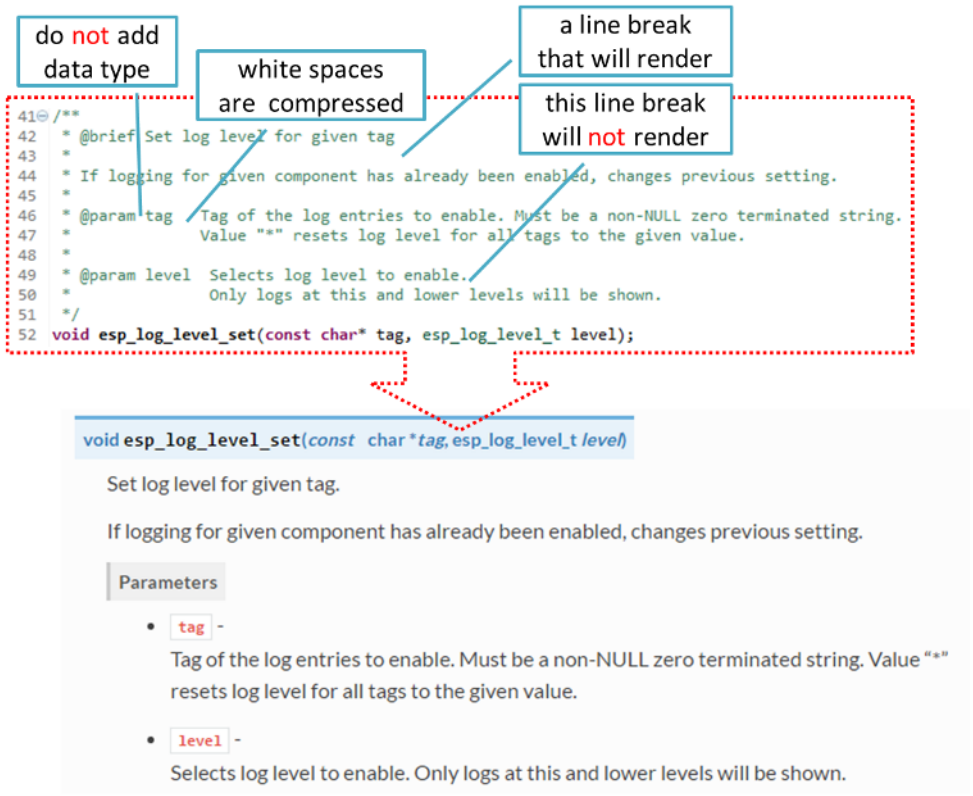
使用这类工具时，上文代码渲染后呈现效果如下：



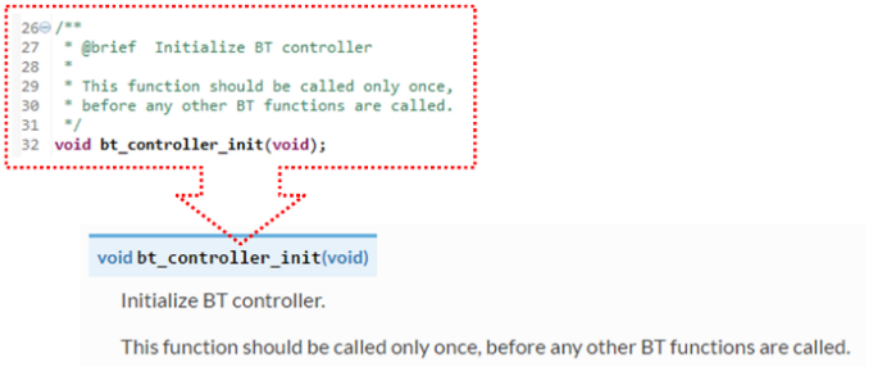
尝试一下！

在本项目库编写代码文档时，请遵守下列准则。

1. 写明代码的基本内容：函数、结构、类型定义、枚举、宏等。请详细说明代码的用途、功能和限制，因为在阅读他人的文档时你也想看到这些信息。
2. 函数文档需简述该函数的功能，并解释输入参数和返回值的含义。
3. 请不要在参数或除空格外的其他字符前面添加数据类型。所有空格和换行符都会压缩为一个空格。如需换行，请执行换行操作两次。



4. 如果函数没有输入参数或返回值，请跳过 `@param` 或 `@return`。



5. 为 `define`、`struct` 和 `enum` 的成员编写文档时，请在每一项后添加注释，如下所示。

6. 请在命令后换行（如下文中的 `@return`），呈现排版精美的列表。

```
*
* @return
*     - ESP_OK if erase operation was successful
*     - ESP_ERR_NVS_INVALID_HANDLE if handle has been closed or is NULL
*     - ESP_ERR_NVS_READ_ONLY if handle was opened as read only
*     - ESP_ERR_NVS_NOT_FOUND if the requested key doesn't exist
*     - other error codes from the underlying storage driver
*
```

7. 头文件的功能概览和库文件应当存在同一个项目库之下，放入单独的 `README.rst` 文件。如果目录下包含不同 API 的头文件，应将文件命名为 `apiname-readme.rst`。

进阶

以下小贴士可以帮助你进一步提高文档质量，增强可读性。

1. 添加代码片段举例说明。请在片段前后添加 `@code{c}` 和 `@endcode` 命令。

```
*
* @code{c}
* // Example of using nvs_get_i32:
* int32_t max_buffer_size = 4096; // default value
* esp_err_t err = nvs_get_i32(my_handle, "max_buffer_size", &max_buffer_size);
* assert(err == ESP_OK || err == ESP_ERR_NVS_NOT_FOUND);
* // if ESP_ERR_NVS_NOT_FOUND was returned, max_buffer_size will still
* // have its default value.
* @endcode
*
```

   代码片段应放入所介绍功能的注释块中。

2. 使用 `@attention` 或 `@note` 命令高亮显示重要信息。

```
*
* @attention
*     1. This API only impact WIFI_MODE_STA or WIFI_MODE_APSTA mode
*     2. If the ESP32 is connected to an AP, call esp_wifi_disconnect to␣
↪disconnect.
*
```

   上述例子介绍了如何使用编号列表。

3. 给相似的函数编写文档时，可在前后使用 `/**@{*/` 和 `/**@}*/` 标记命令。

```
/**@{*/
/**
 * @brief common description of similar functions
 *
 */
void  first_similar_function (void);
```

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```
void second_similar_function (void);
/**@}*/
```

示例请参照 nvs_flash/include/nvs.h。

4. 如果想跳过重复的宏定义、枚举项等代码，不添加描述，请在代码前后添加 `/** @cond */` 和 `/** @endcond */` 命令。示例请参照 driver/include/driver/gpio.h。
5. 使用 markdown 增强文档可读性，添加页眉、链接、表格及更多内容。

```
*
* [ESP32-S2 技术参考手册](https://www.espressif.com/sites/default/files/
↪documentation/esp32-s2_technical_reference_manual_cn.pdf)
*
```

注解： 代码片段、注释、链接等内容如没有附在所述对象对应的注释块中，将不会添加到文档中。

6. 准备一个或更多完整的代码示例和描述，将描述放入单独的 `README.md` 文件中，置于 examples 目录的特定文件夹中。

链接到示例

链接到 GitHub 上的示例时，请不要使用绝对 URLs 或硬编码 URLs。请使用 Docutils 自定义角色生成链接。自动生成的链接指向项目库中 git commit 编号（或标记）的 tree 或 blob。这种做法可以确保 master 分支上的文件移动或删除时，链接不会失效。

有如下角色可以选择：

- `:idf:`path`` - 指向 ESP-IDF 内的目录
- `:idf_file:`path`` - 指向 ESP-IDF 内的文件
- `:idf_raw:`path`` - 指向 ESP-IDF 内的原始格式文件
- `:component:`path`` - 指向 ESP-IDF components 内的文件夹
- `:component_file:`path`` - 指向 ESP-IDF components 内的文件
- `:component_raw:`path`` - 指向 ESP-IDF components 内的原始格式文件
- `:example:`path`` - 指向 ESP-IDF examples 内的文件夹
- `:example_file:`path`` - 指向 ESP-IDF examples 内的文件
- `:example_raw:`path`` - 指向 inside ESP-IDF examples 内的原始格式文件

示例:

```
* :example:`get-started/hello_world`
* :example:`Hello World! <get-started/hello_world>`
```

渲染效果：

- get-started/hello_world
- Hello World!

CI build 脚本中添加了检查功能，查找 RST 文件中的硬编码链接（通过 URL 的 tree/master、blob/master 或 raw/master 部分识别）。该功能可通过 `cd docs` 和 `make gh-linkcheck` 命令手动运行。

链接到其他语言文档

要切换不同语言的文档，可使用 `:link_to_translation:` 自定义角色。文档页面中的角色提供其他语言版本的链接。下文的例子说明了如何在文档中添加中英文版本的链接:

```
:link_to_translation:`zh_CN: 中文版 `
:link_to_translation:`en:English`
```

语言用 en 或 zh_CN 等标准简写表示。最后一个分号后的文本非标准化内容，可根据链接的位置自行输入，如:

```
:link_to_translation:`en:see description in English`
```

添加图例

请考虑使用图表和图片解释表述的概念。

相比于长篇的表述，图例有时可以更好地描述复杂的理念、数据结构或算法。本项目库使用 'blockdiag <http://blockdiag.com/en/index.html>'_ 工具包由简单的文本文件生成图表。

工具包支持下列图表类型:

- 框图
- 时序图
- 活动图
- 逻辑网络图

使用该工具包，可以将简单的文本（与 graphviz 的 DOT 格式类似）转换成美观的图片。图中内容自动排版。图标代码之后会转换为 ".png" 图片，在后台添加进 **Sphinx** 文档中。

要查看图表的渲染效果，可使用线上的 interactive shell 即时显示生成的图片。

下面是一些图表示例:

- 简单的 **框图** / `blockdiag` - Wi-Fi Buffer 配置
- 稍复杂的 **框图** - Wi-Fi 编程模型
- **时序图** / `seqdiag` - 在所有信道中扫描特定 AP
- **包图** / `packetdiag` - NVS 页面结构

尝试修改源代码，看看图表会发生什么变化。

注解: interactive shell 使用的字体和 esp-idf 文档使用的字体略有不同。

添加注释

写文档时，您可能需要:

- 留下建议，说明之后需添加会修改哪些内容。
- 提醒自己或其他人跟进。

这时，您可以使用 `.. todo::` 命令在 reST 文件中添加待做事项。如:

```
.. todo::

   Add a package diagram.
```

如果在 reST 文件中添加 `.. todolist::` 命令，整篇文档中的所有待做事项将会罗列成表。

默认情况下，文档生成器会忽视 `.. todo::` 和 `.. todolist::` 命令。如果您想在本地生成的文档中显示注释和注释列表，请执行下列步骤:

1. 打开本地的 `conf_common.py` 文件。
2. 找到 `todo_include_todos` 参数。
3. 将该参数的值由 `False` 改为 `True`。

将改动推送到远端分支之前，请把 `todo_include_todos` 的值重置为 `False`。

更多关于扩展的信息，请参阅 sphinx.ext.todo 的相关文档。

汇总文档

文档准备好后，请参照*API Documentation Template* 的要求创建一个文件，汇总所有准备好的文档。最后，在文件中添加链接指向 `/docs` 文件夹或子文件夹下 `index.rst` 文件的 `.. toctree::`。

Sphinx 新手怎么办

1. 不要担心。所有需要的软件均有详细文档，并且开源、免费。您可以先查看 Sphinx 文档。如果您不清楚如何用 rst markup 语言写作，请查看 reStructuredText Primer。您也可以使用 markdown (.md) 文件，查找更多在 Recommonmark parser' 文档页面 使用的特定 markdown 句法信息。
2. 查看本文档的源文件，了解本文档使用的代码。源文件存储于 GitHub espressif/esp-idf 项目库的 docs 文件夹下。您可以滑动到页面上方，点击右上角的链接，直接查看本页面的源文件。您也可以通过点击 `Raw` 按键打开源文件，在 GitHub 上查看文件的代码。
3. 想要查看在上传至 GitHub 前文档如何生成、呈现，有两种方式：
   - 安装 'Sphinx'_、Breathe、Blockdiag 和 Doxygen 本地生成文档，具体可查看下文。
   - 在 Read the Docs 建立账号，在云端生成文档。Read the Docs 免费提供文档生成和存储，且速度快、质量高。
4. 在生成文档前预览，可使用 Sublime Text 编辑器和 OmniMarkupPreviewer 插件。

搭建环境本地生成文档

您可以安装下列包，通过搭建环境在电脑上本地生成文档：

1. Doxygen - http://doxygen.nl/
2. Sphinx - https://github.com/sphinx-doc/sphinx/#readme-for-sphinx
3. Breathe - https://github.com/michaeljones/breathe#breathe
4. Document theme "sphinx_idf_theme" - https://github.com/rtfd/sphinx_idf_theme
5. Custom 404 page "sphinx-notfound-page" - https://github.com/rtfd/sphinx-notfound-page
6. Blockdiag - http://blockdiag.com/en/index.html
7. Recommonmark - https://github.com/rtfd/recommonmark

添加 "sphinx_idf_theme" 包之后，文档将与 ESP-IDF 编程指南 的风格保持一致。

不用担心需要安装太多包。除 Doxygen 和 sphinx_idf_theme 之外，其他包均使用纯 Python 语言，可一键安装。

Doxygen 的安装取决于操作系统：

Linux

```
sudo apt-get install doxygen
```

Windows - 在 MSYS2 控制台中安装

```
pacman -S doxygen
```

MacOS

```
brew install doxygen
```

注解： 如果您是在 Windows 系统上安装（Linux 和 MacOS 用户可以跳过此说明），在安装 **之前**，请完成以下两步。这是安装添加图例 提到的 "blockdiag" 依赖项的必须步骤。

1. 更新所有系统包：

   ```
   $ pacman -Syu
   ```

   该过程可能需要重启 MSYS2 MINGW32 控制台重复上述命令，直至更新完成。
2. 安装 *blockdiag* 的依赖项之一 *pillow*：

```
$ pacman -S mingw32/mingw-w64-i686-python-pillow
```

查看屏幕上的记录，确定 `mingw-w64-i686-python-pillow-4.3.0-1` 已安装。旧版本 *pillow* 无法运行。

Windows 安装 Doxygen 的缺点是 blockdiag pictures 字体不能正确加载，可能会存在乱码。在此问题解决之前，您可以使用 interactive shell 查看完整图片。

sphinx_idf_theme 编译 `sphinx_idf_theme` 需要同时使用 Python 和 JavaScript。因此，目前要进行本地编译还必须安装 node.js ，命令如下：

```
cd ~/esp
git clone https://github.com/espressif/sphinx_idf_theme.git
cd sphinx_idf_theme
npm install
python setup.py build
python setup.py install
```

我们计划在不久的将来支持安装预编译 `sphinx_idf_theme`，给您带来的暂时不便，敬请谅解。

其他所有应用都是 Python 包，可以按照下列步骤一键安装：

```
cd ~/esp/esp-idf/docs
pip install --user -r requirements.txt
```

注解： 安装步骤设定将 ESP-IDF 放在 `~/esp/esp-idf` 目录下，这是文档中使用的 ESP-IDF 默认地址。

更换到特定语言文件所在的目录:

```
cd en
```

现在可以调用如下命令生成文档:

```
make html
```

这一步骤需要几分钟时间。完成后，文档会放置在 `~/esp/esp-idf/docs/en/_build/html` 文件夹下。您可以在网页浏览器中打开 `index.html` 查看。

大功告成

我们喜欢可以做酷炫事情的好代码。但我们更喜欢有清晰文档的好代码，可以让读者快速上手，做酷炫的事情。

尝试一下，贡献你的代码和文档！

相关文档

- *API Documentation Template*
- *Documentation Add-ons and Extensions Reference*

6.5.3 Documentation Add-ons and Extensions Reference

This documentation is created using Sphinx application that renders text source files in reStructuredText (`.rst`) format located in docs directory. For some more details on that process, please refer to section 编写代码文档.

Besides Sphinx there are several other applications that help to provide nicely formatted and easy to navigate documentation. These applications are listed in section 搭建环境本地生成文档 with the installed version numbers provided in file docs/requirements.txt.

We build ESP-IDF documentation for two languages (English, Simplified Chinese) and for multiple chips. Therefore we don't run `sphinx` directly, there is a wrapper Python program `build_docs.py` that runs Sphinx.

On top of that we have created a couple of custom add-ons and extensions to help integrate documentation with underlining ESP-IDF repository and further improve navigation as well as maintenance of documentation.

The purpose of this section is to provide a quick reference to the add-ons and the extensions.

Documentation Folder Structure

- The ESP-IDF repository contains a dedicated documentation folder docs in the root.
- The `docs` folder contains localized documentation in docs/en (English) and docs/zh_CN (simplified Chinese) subfolders.
- Graphics files and fonts common to localized documentation are contained in docs/_static subfolder
- Remaining files in the root of `docs` as well as `docs/en` and `docs/zh_CN` provide configuration and scripts used to automate documentation processing including the add-ons and extensions.
- Sphinx extensions are provided in two directories, `extensions` and `idf_extensions`
- A `_build` directory is created in the `docs` folder by `build_docs.py`. This directory is not added to the ESP-IDF repository.

Add-ons and Extensions Reference

Config Files

docs/conf_common.py This file contains configuration common to each localized documentation (e.g. English, Chinese). The contents of this file is imported to standard Sphinx configuration file `conf.py` located in respective language folders (e.g. `docs/en`, `docs/zh_CN`) during build for each language.

docs/sphinx-known-warnings.txt There are couple of spurious Sphinx warnings that cannot be resolved without doing update to the Sphinx source code itself. For such specific cases respective warnings are documented in `sphinx-known-warnings.txt` file, that is checked during documentation build, to ignore the spurious warnings.

Scripts docs/build_docs.py

Top-level executable program which runs a Sphinx build for one or more language/target combinations. Run `build_docs.py --help` for full command line options.

When `build_docs.py` runs Sphinx it sets the `idf_target` configuration variable, sets a Sphinx tag with the same name as the configuration variable, and uses some environment variables to communicate paths to *IDF-Specific Extensions*.

docs/check_lang_folder_sync.sh To reduce potential discrepancies when maintaining concurrent language version, the structure and filenames of language folders `docs/en` and `docs/zh_CN` folders should be kept identical. The script `check_lang_folder_sync.sh` is run on each documentation build to verify if this condition is met.

注解: If a new content is provided in e.g. English, and there is no any translation yet, then the corresponding file in `zh_CN` folder should contain an `.. include::` directive pointing to the source file in English. This will automatically include the English version visible to Chinese readers. For example if a file `docs/zh_CN/contribute/documenting-code.rst` does not have a Chinese translation, then it should contain `.. include:: ../../en/contribute/documenting-code.rst` instead.

Non-Docs Scripts These scripts are used to build docs but also used for other purposes:

tools/gen_esp_err_to_name.py This script is traversing the ESP-IDF directory structure looking for error codes and messages in source code header files to generate an `.inc` file to include in documentation under *Error Codes Reference*.

tools/kconfig_new/confgen.py Options to configure ESP-IDF's components are contained in `Kconfig` files located inside directories of individual components, e.g. components/bt/Kconfig. This script is traversing the `component` directories to collect configuration options and generate an `.inc` file to include in documentation under *Configuration Options Reference*.

Generic Extensions These are Sphinx extensions developed for IDF that don't rely on any IDF-docs-specific behaviour or configuration:

docs/extensions/toctree_filter.py Sphinx extensions overrides the `:toctree:` directive to allow filtering entries based on whether a tag is set, as `:tagname: toctree_entry`. See the Python file for a more complete description.

docs/extensions/list_filter.py Sphinx extensions that provides a `.. list::` directive that allows filtering of entries in lists based on whether a tag is set, as `:tagname: - list content`. See the Python file for a more complete description.

docs/extensions/html_redirects.py During documentation lifetime some source files are moved between folders or renamed. This Sphinx extension adds a mechanism to redirect documentation pages that have changed URL by generating in the Sphinx output static HTML redirect pages. The script is used together with a redirection list `html_redirect_pages`. `conf_common.py` builds this list from docs/page_redirects.txt

Third Party Extensions

- `sphinxcontrib` extensions for blockdiag, seqdiag, actdiag, nwdiag, rackdiag & packetdiag diagrams.
- Sphinx selective exclude `eager_only` extension

IDF-Specific Extensions

Build System Integration docs/idf_extensions/build_system/

Python package implementing a Sphinx extension to pull IDF build system information into the docs build

- Creates a dummy CMake IDF project and runs CMake to generate metadata
- Registers some new configuration variables and emits a new Sphinx event, both for use by other extensions.

Configuration Variables

- `docs_root` - The absolute path of the $IDF_PATH/docs directory
- `idf_path` - The value of IDF_PATH variable, or the absolute path of IDF_PATH if environment unset
- `build_dir` - The build directory passed in by `build_docs.py`, default will be like `_build/<lang>/<target>`
- `idf_target` - The IDF_TARGET value. Expected that `build_docs.py` set this on the Sphinx command line

New Event `idf-info` event is emitted early in the build, after the dummy project CMake run is complete.

Arguments are `(app, project_description)` where `project_description` is a dict containing the values parsed from `project_description.json` in the CMake build directory.

Other IDF-specific extensions subscribe to this event and use it to set up some docs parameters based on build system info.

Other Extensions

docs/idf_extensions/include_build_file.py The `include-build-file` directive is like the built-in `include-file` directive, but file path is evaluated relative to `build_dir`.

docs/idf_extensions/kconfig_reference.py Subscribes to `idf-info` event and uses confgen to generate `kconfig.inc` from the components included in the default project build. This file is then included into *Project Configuration*.

docs/idf_extensions/link_roles.py This is an implementation of a custom Sphinx Roles to help linking from documentation to specific files and folders in ESP-IDF. For description of implemented roles please see 链接到示例 and 链接到其他语言文档.

docs/idf_extensions/esp_err_definitions.py Small wrapper extension that calls `gen_esp_err_to_name.py` and updates the included .rst file if it has changed.

docs/idf_extensions/gen_toolchain_links.py There couple of places in documentation that provide links to download the toolchain. To provide one source of this information and reduce effort to manually update several files, this script generates toolchain download links and toolchain unpacking code snippets based on information found in tools/toolchain_versions.mk.

docs/idf_extensions/gen_version_specific_includes.py Another extension to automatically generate reStructuredText Text `.inc` snippets with version-based content for this ESP-IDF version.

docs/idf_extensions/util.py A collection of utility functions useful primarily when building documentation locally (see 搭建环境本地生成文档) to reduce the time to generate documentation on a second and subsequent builds.

docs/idf_extensions/format_idf_target.py An extension for replacing generic target related names with the idf_target passed to the Sphinx command line. This is a {IDF_TARGET_NAME}, with /{IDF_TARGET_PATH_NAME}/soc.c, compiled with *xtensa-{IDF_TARGET_TOOLCHAIN_NAME}-elf-gcc* with *CONFIG_{IDF_TARGET_CFG_PREFIX}_MULTI_DOC* will, if the backspaces are removed, render as This is a ESP32-S2, with /esp32s2/soc.c, compiled with *xtensa-esp32s2-elf-gcc* with *CONFIG_ESP32S2_MULTI_DOC*.

Also supports markup for defining local (single .rst-file) substitions with the following syntax: {IDF_TARGET_TX_PIN:default=" IO3" ,esp32=" IO4" ,esp32s2=" IO5" }

This will define a replacement of the tag {IDF_TARGET_TX_PIN} in the current rst-file.

The extension also overrides the default `.. include::` directive in order to format any included content using the same rules.

These replacements cannot be used inside markup that rely on alignment of characters, e.g. tables.

docs/idf_extensions/latex_builder.py An extension for adding ESP-IDF specific functionality to the latex builder. Overrides the default Sphinx latex builder.

Creates and adds the espidf.sty latex package to the output directory, which contains some macros for run-time variables such as IDF-Target.

docs/idf_extensions/gen_defines.py Sphinx extension to integrate defines from IDF into the Sphinx build, runs after the IDF dummy project has been built.

Parses defines and adds them as sphinx tags.

Emits the new 'idf-defines-generated' event which has a dictionary of raw text define values that other extensions can use to generate relevant data.

docs/idf_extensions/exclude_docs.py Sphinx extension that updates the excluded documents according to the conditional_include_dict {tag:documents}. If the tag is set, then the list of documents will be included.

Also responsible for excluding documents when building with the config value `docs_to_build` set. In these cases all documents not listed in `docs_to_build` will be excluded.

Subscribes to `idf-defines-generated` as it relies on the sphinx tags to determine which documents to exclude

docs/idf_extensions/run_doxygen.py Subscribes to `idf-defines-generated` event and runs Doxygen (docs/Doxyfile) to generate XML files describing key headers, and then runs Breathe to convert these to `.inc` files which can be included directly into API reference pages.

Pushes a number of target-specific custom environment variables into Doxygen, including all macros defined in the project's default `sdkconfig.h` file and all macros defined in all `soc` component `xxx_caps.h` headers. This means that public API headers can depend on target-specific configuration options or `soc` capabilities headers options as `#ifdef` & `#if` preprocessor selections in the header.

This means we can generate different Doxygen files, depending on the target we are building docs for.

Please refer to 编写代码文档 and *API Documentation Template*, section **API Reference** for additional details on this process.

Related Documents

- 编写代码文档

6.5.4 创建示例项目

每个 ESP-IDF 的示例都是一个完整的项目，其他人可以将示例复制至本地，并根据实际情况进行一定修改。请注意，示例项目主要是为了展示 ESP-IDF 的功能。

示例项目结构

- `main` 目录需要包含一个名为 `(something)_example_main.c` 的源文件，里面包含示例项目的主要功能。
- 如果该示例项目的子任务比较多，请根据逻辑将其拆分为 `main` 目录下的多个 C 或者 C++ 源文件，并将对应的头文件也放在同一目录下。
- 如果该示例项目具有多种功能，可以考虑在项目中增加一个 `components` 子目录，通过库功能，将示例项目的不同功能划分为不同的组件。注意，如果该组件提供的功能相对完整，且具有一定的通用性，则应该将它们添加到 ESP-IDF 的 `components` 目录中，使其成为 ESP-IDF 的一部分。
- 示例项目需要包含一个 `README.md` 文件，建议使用 示例项目 README 模板 ，并根据项目实际情况进行修改。
- 示例项目需要包含一个 `example_test.py` 文件，用于进行自动化测试。如果在 GitHub 上初次提交 Pull Request 时，可以先不包含这个脚本文件。具体细节，请见有关 Pull Request 的相关内容。

一般准则

示例代码需要遵循《乐鑫物联网开发框架风格指南》。

检查清单

提交一个新的示例项目之前，需要检查以下内容：

- 示例项目的名字（包括 `Makefile` 和 `README.md` 中）应使用 `example`，而不要写 “demo”，“test” 等词汇。
- 每个示例项目只能有一个主要功能。如果某个示例项目有多个主要功能，请将其拆分为两个或更多示例项目。
- 每个示例项目应包含一个 `README.md` 文件，建议使用 示例项目 README 模板 。
- 示例项目中的函数和变量的命令要遵循命名规范 中的要求。对于仅在示例项目源文件中使用的非静态变量/函数，请使用 `example` 或其他类似的前缀。
- 示例项目中的所有代码结构良好，关键代码要有详细注释。
- 示例项目中所有不必要的代码（旧的调试日志，注释掉的代码等）都必须清除掉。
- 示例项目中使用的选项（比如网络名称，地址等）不得直接硬编码，应尽可能地使用配置项，或者定义为宏或常量。
- 配置项可见 `KConfig.projbuild` 文件，该文件中包含一个名为 “Example Configuration” 的菜单。具体情况，请查看现有示例项目。
- 所有的源代码都需要在文件开头指定许可信息（表示该代码是 `in the public domain CC0`）和免责声明。或者，源代码也可以应用 `Apache License 2.0` 许可条款。请查看现有示例项目的许可信息和免责声明，并根据实际情况进行修改。
- 任何第三方代码（无论是直接使用，还是进行了一些改进）均应保留原始代码中的许可信息，且这些代码的许可必须兼容 `Apache License 2.0` 协议。

6.5.5 API Documentation Template

注解： *INSTRUCTIONS*

1. Use this file (docs/en/api-reference/template.rst) as a template to document API.
2. Change the file name to the name of the header file that represents documented API.
3. Include respective files with descriptions from the API folder using `..include::`
   - README.rst
   - example.rst
   - ...
4. Optionally provide description right in this file.
5. Once done, remove all instructions like this one and any superfluous headers.

Overview

注解： *INSTRUCTIONS*

1. Provide overview where and how this API may be used.
2. Where applicable include code snippets to illustrate functionality of particular functions.
3. To distinguish between sections, use the following heading levels:
   - `#` with overline, for parts
   - `*` with overline, for chapters
   - =, for sections
   - −, for subsections
   - ^, for subsubsections
   - ", for paragraphs

Application Example

注解： *INSTRUCTIONS*

1. Prepare one or more practical examples to demonstrate functionality of this API.
2. Each example should follow pattern of projects located in `esp-idf/examples/` folder.
3. Place example in this folder complete with `README.md` file.
4. Provide overview of demonstrated functionality in `README.md`.
5. With good overview reader should be able to understand what example does without opening the source code.
6. Depending on complexity of example, break down description of code into parts and provide overview of functionality of each part.
7. Include flow diagram and screenshots of application output if applicable.
8. Finally add in this section synopsis of each example together with link to respective folder in `esp-idf/examples/`.

API Reference

注解： *INSTRUCTIONS*

1. This repository provides for automatic update of API reference documentation using *code markup retrieved by Doxygen from header files*.

1. Update is done on each documentation build by invoking Sphinx extension docs/idf_extensions/run_doxygen.py for all header files listed in the `INPUT` statement of docs/Doxyfile.

1. Each line of the `INPUT` statement (other than a comment that begins with `##`) contains a path to header file `*.h` that will be used to generate corresponding `*.inc` files:

```
##
## Wi-Fi - API Reference
##
../components/esp32/include/esp_wifi.h \
../components/esp32/include/esp_smartconfig.h \
```

1. When the headers are expanded, any macros defined by default in `sdkconfig.h` as well as any macros defined in SOC-specific `include/soc/*_caps.h` headers will be expanded. This allows the headers to include/exclude material based on the `IDF_TARGET` value.

1. The `*.inc` files contain formatted reference of API members generated automatically on each documentation build. All `*.inc` files are placed in Sphinx `_build` directory. To see directives generated for e.g. `esp_wifi.h`, run `python gen-dxd.py esp32/include/esp_wifi.h`.

1. To show contents of `*.inc` file in documentation, include it as follows:

   ```
   .. include-build-file:: inc/esp_wifi.inc
   ```

   For example see docs/en/api-reference/network/esp_wifi.rst

1. Optionally, rather that using `*.inc` files, you may want to describe API in you own way. See docs/en/api-guides/ulp.rst for example.
   Below is the list of common `.. doxygen...::` directives:
   - Functions - `.. doxygenfunction:: name_of_function`
   - Unions - `.. doxygenunion:: name_of_union`
   - Structures - `.. doxygenstruct:: name_of_structure` together with `:members:`
   - Macros - `.. doxygendefine:: name_of_define`
   - Type Definitions - `.. doxygentypedef:: name_of_type`
   - Enumerations - `.. doxygenenum:: name_of_enumeration`

   See Breathe documentation for additional information.
   To provide a link to header file, use the *link custom role* as follows:

   ```
   * :component_file:`path_to/header_file.h`
   ```

1. In any case, to generate API reference, the file docs/Doxyfile should be updated with paths to `*.h` headers that are being documented.

1. When changes are committed and documentation is build, check how this section has been rendered. *Correct annotations* in respective header files, if required.

6.5.6 Contributor Agreement

Individual Contributor Non-Exclusive License Agreement

including the Traditional Patent License OPTION

Thank you for your interest in contributing to Espressif IoT Development Framework (esp-idf) ("We" or "Us").

The purpose of this contributor agreement ("Agreement") is to clarify and document the rights granted by contributors to Us. To make this document effective, please follow the instructions at CONTRIBUTING.rst

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8.4 You agree to notify Us of any facts or circumstances of which you become aware that would make this Agreement inaccurate in any respect.

You

Date:	
Name:	
Title:	
Address:	

Us

Date:	
Name:	
Title:	
Address:	

Chapter 7

ESP-IDF 版本简介

ESP-IDF 的 GitHub 仓库时常更新，特别是用于开发新特性的 master 分支。

如有量产需求，请使用稳定版本。

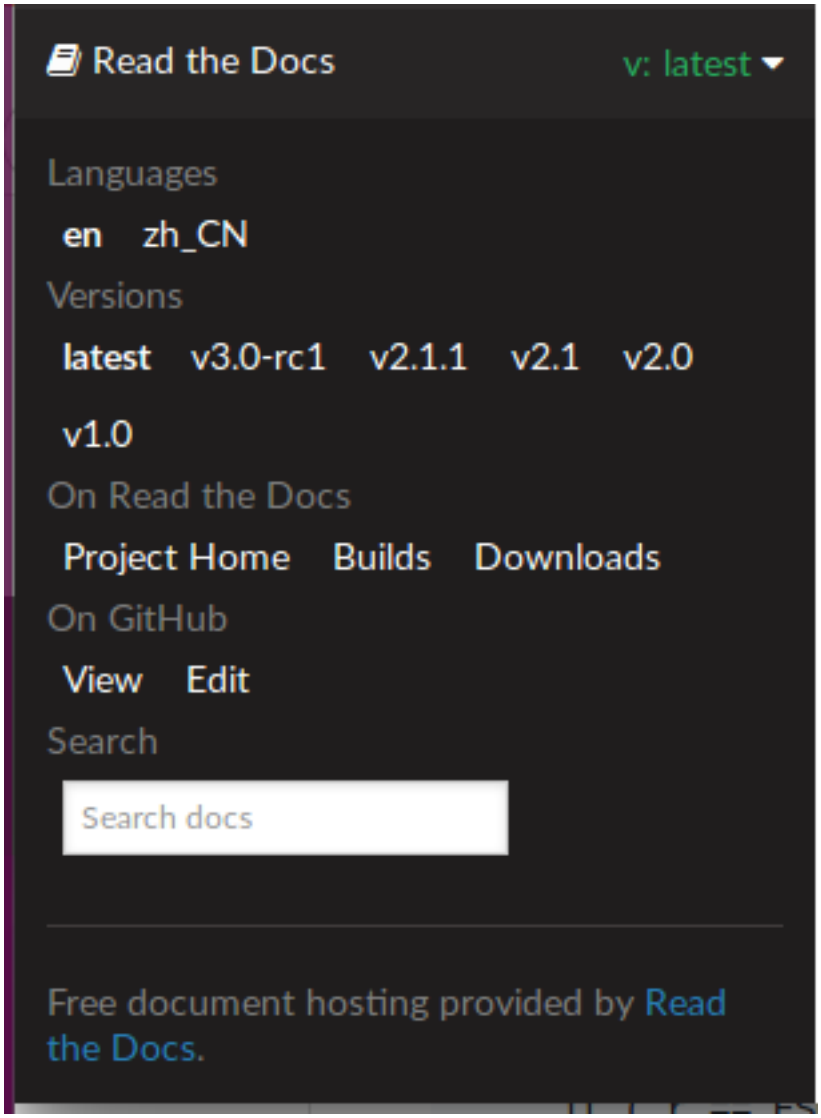
7.1 发布版本

您可以通过以下链接访问各个版本的配套文档：

- 最新稳定版 ESP-IDF：https://docs.espressif.com/projects/esp-idf/zh_CN/stable/
- 最新版 ESP-IDF（即 master 分支）：https://docs.espressif.com/projects/esp-idf/zh_CN/latest/

ESP-IDF 在 GitHub 平台上的完整发布历史请见 发布说明页面。您可以在该页面查看各个版本的发布说明、配套文档及相应获取方式。

此外，您还可以直接前往文档页面，查看部分 ESP-IDF 版本的配套文档，具体可通过点击页面左下角的小三角标志，在不同版本间切换。



7.2 我该选择哪个版本？

- 如有量产需求，请使用 最新稳定版本 。稳定版本已通过人工测试，后续更新仅修复 bug，主要特性不受影响（更多详情，请见 版本管理 ）。请访问 发布说明页面 界面查看每一个稳定发布版本。
- 如需尝试/测试 ESP-IDF 的最新特性，请使用 最新版本（在 master 分支上）。最新版本包含 ESP-IDF 的所有最新特性，已通过自动化测试，但尚未全部完成人工测试（因此存在一定风险）。
- 如需使用稳定版本中没有的新特性，但同时又不希望受到 master 分支更新的影响，您可以将一个最适合您的稳定版本更新至一个预发布版本 或更新至一个发布分支。

有关如何更新 ESP-IDF 本地副本的内容，请参考 *更新 ESP-IDF* 章节。

7.3 版本管理

ESP-IDF 采用了 语义版本管理方法，即您可以从字面含义理解每个版本的差异。其中

- 主要版本（例 `v3.0`）代表有重大更新，包括增加新特性、改变现有特性及移除已弃用的特性。升级至一个新的主要版本（例 `v2.1` 升级至 `v3.0`）意味着您可能需要更新您的工程代码，并重新测试您的工程，具体可参考 发布说明页面 的重大变更 (Breaking Change) 部分。
- 次要版本（例 `v3.1`）代表有新增特性和 bug 修复，但现有特性不受影响，公开 API 的使用也不受影响。

升级至一个新的次要版本（例 `v3.0` 升级至 `v3.1`）意味着您可能不需要更新您的工程代码，但需重新测试您的工程，特别是 发布说明页面 中专门提到的部分。

- Bugfix 版本（例 `v3.0.1`）仅修复 bug，并不增加任何新特性。
  升级至一个新的 Bugfix 版本（例 `v3.0` 升级至 `v3.0.1`）意味着您不需要更新您的工程代码，仅需测试与本次发布修复 bug（列表见 发布说明页面）直接相关的特性。

7.4 支持期限

ESP-IDF 的每个主要版本和次要版本都有相应的支持期限。支持期限满后，版本停止更新维护，将不再支持。

支持期限政策 对此有具体描述，并介绍了每个版本的支持期限是如何界定的。

发布说明页面 界面上的每一个发布版本都提供了该版本的支持期限信息。

一般而言：

- 如您刚开始一个新项目，建议使用最新稳定版本。
- 如您有 GitHub 账号，请点击 发布说明页面 界面右上角的 “Watch” 按键，并选中 “Releases only” 选项。GitHub 将会在新版本发布的时候通知您。当您所使用的版本有 bugfix 版本发布时，请做好升级至该 bugfix 版本的规划。
- 如可能，请定期（如每年一次）将项目的 IDF 版本升级至一个新的主要版本或次要版本。对于小版本更新，更新过程应该比较简单，但对于主要版本更新，可能需要细致查看发布说明并做对应的更新规划。
- 请确保您所使用的版本停止更新维护前，已做好升级至新版本的规划。

7.5 查看当前版本

查看 ESP-IDF 本地副本的版本，请使用 idf.py 命令：

```
idf.py --version
```

此外，由于 ESP-IDF 的版本也已编译至固件中，因此您也可以使用宏 `IDF_VER` 查看 ESP-IDF 的版本（以字符串的格式）。ESP-IDF 默认引导程序会在设备启动时打印 ESP-IDF 的版本。请注意，在 GitHub 仓库中的代码更新时，代码中的版本信息仅会在源代码重新编译或在清除编译时才会更新，因此打印出来的版本可能并不是最新的。

几个 ESP-IDF 版本的例子：

版本字符串	含义
`v3.2-dev-306-gbeb3611ca`	master 分支上的预发布版本。 - `v3.2-dev`：为 v3.2 进行的开发。 - `306`：v3.2 开发启动后的 commit 数量。 - `beb3611ca`：commit 标识符。
`v3.0.2`	稳定版本，标签为 `v3.0.2`。
`v3.1-beta1-75-g346d6b0ea`	v3.1 的 beta 测试版本（可参考更新至一个发布分支）。 - `v3.1-beta1` - 预发布标签。 - `75`：添加预发布 beta 标签后的 commit 数量。 - `346d6b0ea`：commit 标识符。
`v3.0.1-dirty`	稳定版本，标签为 `v3.0.1`。 - `dirty` 代表 ESP-IDF 的本地副本有修改。

7.6 Git 工作流

乐鑫 ESP-IDF 团队的 (Git) 开发工作流程如下：

- 新的改动总是在 master 分支（最新版本）上进行。master 分支上的 ESP-IDF 版本总带有 `-dev` 标签，表示“正在开发中”，例 `v3.1-dev`。
- 这些改动将首先在乐鑫的内部 Git 仓库进行代码审阅与测试，而后在自动化测试完成后推至 GitHub。
- 新版本一旦完成特性开发（在 `master` 分支上进行）并达到进入 beta 测试的标准，则将该版本签出至一个新分支（例 `release/v3.1`）。此外，该分支还打上预发布标签（例 `v3.1-beta1`）。您可以在 GitHub 平台上查看 ESP-IDF 的完整 分支列表 和 标签列表。Beta 预发布版本可能仍存在大量“已知问题”(Known Issue)。
- 随着对 beta 版本的不断测试，bug 修复将同时增加至该发布分支和 `master` 分支。而且，`master` 分支可能也已经开始为下个版本开发新特性了。
- 当测试快结束时，该发布分支上将增加一个 rc 标签，代表候选发布 (Release Candidate)，例 `v3.1-rc1`。此时，该分支仍属于预发布版本。
- 如果一直未发现或报告重大 bug，则该预发布版本将最终增加“主要版本”（例 `v4.0`）或“次要版本”标记（例 `v3.1`），成为正式发布版本，并体现在 发布说明页面 。
- 后续，该版本中发现的 bug 都将在该发布分支上进行修复。人工测试完成后，该分支将增加一个 bugfix 版本标签（例 `v3.1.1`），并体现在 发布说明页面 。

7.7 更新 ESP-IDF

请根据您的实际情况，对 ESP-IDF 进行更新。

- 如有量产用途，建议参考更新至一个稳定发布版本。
- 如需测试/研发/尝试最新特性，建议参考更新至 *master* 分支 。
- 两者折衷建议参考更新至一个发布分支。

注解： 在参考本指南时，请首先获得 ESP-IDF 的本地副本，具体步骤请参考入门指南 中的介绍。

7.7.1 更新至一个稳定发布版本

（推荐量产用户）如需更新至一个新的 ESP-IDF 发布版本，请参考以下步骤：

- 请定期查看 发布说明页面 ，了解最新发布情况。
- 如有新发布的 bugfix 版本（例 `v3.0.1` 或 `v3.0.2`）时，请将新的 bugfix 版本更新至您的 ESP-IDF 目录：

```
cd $IDF_PATH
git fetch
git checkout vX.Y.Z
git submodule update --init --recursive
```

- 如有主要版本或次要版本新发布时，请查看发布说明中的具体描述，并决定是否升级您的版本。具体命令与上方描述一致。

注解： 如果您之前在安装 ESP-IDF 时使用了 zip 文件包，而非通过 Git 命令，则您将无法使用 Git 命令进行版本升级，此属正常情况。这种情况下，请重新下载最新 zip 文件包，并替换掉之前 `IDF_PATH` 下的全部内容。

7.7.2 更新至一个预发布版本

您也可以将您的本地副本签出（命令 `git checkout`）至一个预发布版本或 rc 版本，具体方法请参考*更新至一个稳定发布版本* 中的描述。

预发布版本通常不体现在 发布说明页面。更多详情，请查看完整 标签列表。使用预发布版本的注意事项，请参考*更新至一个发布分支* 中的描述。

7.7.3 更新至 master 分支

注解： ESP-IDF 中 master 分支上的代码会时时更新，因此使用 master 分支相当在 “流血的边缘试探”，存在一定风险。

如需使用 ESP-IDF 的 master 分支，请参考以下步骤：

- 本地签出至 master 分支：

```
cd $IDF_PATH
git checkout master
git pull
git submodule update --init --recursive
```

- 此外，您还应在后续工作中不时使用 `git pull` 命令，将远端 master 上的更新同步到本地。注意，在更新 master 分支后，您可能需要更改项代码，也可能遇到新的 bug。
- 如需从 master 分支切换至一个发布分支或稳定版本，请使用 `git checkout` 命令。

重要： 强烈建议您定期使用 `git pull` 和 `git submodule update --init --recursive` 命令，确保本地副本的及时更新。旧的 master 分支相当于一个 “快照”，可能存在未记录的问题，且无法获得支持。对于半稳定版本，请参考*更新至一个发布分支*。

7.7.4 更新至一个发布分支

从稳定性来说，使用 “发布分支” 相当于在使用 master 分支和稳定版本之间进行折衷，包含一些 master 分支上的新特性，但也同时保证可通过 beta 测试且基本完成了 bug 修复。

更多详情，请前往 GitHub 查看完整 标签列表。

举例，您可以关注 ESP-IDF v3.1 分支，随时关注该分支上的 bugfix 版本发布（例 `v3.1.1` 等）：

```
cd $IDF_PATH
git fetch
git checkout release/v3.1
git pull
git submodule update --init --recursive
```

您每次在该分支上使用 `git pull` 时都相当于把最新的 bugfix 版本发布更新至您的本地副本中。

注解： 发布分支并不会有专门的配套文档，建议您使用与本分支最接近的版本。

Chapter 8

资源

8.1 PlatformIO

8.1.1 What is PlatformIO?

PlatformIO is a cross-platform embedded development environment with out-of-the-box support for ESP-IDF.

Since ESP-IDF support within PlatformIO is not maintained by the Espressif team, please report any issues with PlatformIO directly to its developers in the official PlatformIO repositories.

A detailed overview of the PlatformIO ecosystem and its philosophy can be found in the official PlatformIO documentation.

8.1.2 Installation

- PlatformIO IDE is a toolset for embedded C/C++ development available on Windows, macOS and Linux platforms
- PlatformIO Core (CLI) is a command-line tool that consists of multi-platform build system, platform and library managers and other integration components. It can be used with a variety of code development environments and allows integration with cloud platforms and web services

8.1.3 Configuration

Please go through the official PlatformIO configuration guide for ESP-IDF.

8.1.4 Tutorials

- ESP-IDF and ESP32-DevKitC: debugging, unit testing, project analysis

8.1.5 Project Examples

Please check ESP-IDF page in the official PlatformIO documentation

8.1.6 Next Steps

Here are some useful links for exploring the PlatformIO ecosystem:

- Learn more about integrations with other IDEs/Text Editors
- Get help from PlatformIO community

8.2 有用的链接

- 您可以在 ESP32 论坛 中提出您的问题，访问社区资源。
- 您可以通过 GitHub 的 Issues 版块提交 bug 或功能请求。在提交新 Issue 之前，请先查看现有的 Issues。
- 您可以在 ESP32 IoT Solution 库中找到基于 ESP-IDF 的 解决方案、应用实例、组件和驱动 等内容。
- 通过 Arduino 平台开发应用，请参考 ESP32 Wi-Fi 芯片的 Arduino 内核。
- 关于 ESP32 的书籍列表，请查看 乐鑫 网站。
- 如果您有兴趣参与到 ESP-IDF 的开发, 请查阅*Contributions Guide*。
- 关于 ESP32 的其它信息，请查看官网 文档 版块。
- 关于本文档的 PDF 和 HTML 格式下载（最新版本和早期版本），请点击 下载。

Chapter 9

Copyrights and Licenses

9.1 Software Copyrights

All original source code in this repository is Copyright (C) 2015-2019 Espressif Systems. This source code is licensed under the Apache License 2.0 as described in the file LICENSE.

Additional third party copyrighted code is included under the following licenses.

Where source code headers specify Copyright & License information, this information takes precedence over the summaries made here.

9.1.1 Firmware Components

These third party libraries can be included into the application (firmware) produced by ESP-IDF.

- Newlib is licensed under the BSD License and is Copyright of various parties, as described in COPYING.NEWLIB.
- Xtensa header files are Copyright (C) 2013 Tensilica Inc and are licensed under the MIT License as reproduced in the individual header files.
- Original parts of FreeRTOS (components/freertos) are Copyright (C) 2015 Real Time Engineers Ltd and is licensed under the GNU General Public License V2 with the FreeRTOS Linking Exception, as described in license.txt.
- Original parts of LWIP (components/lwip) are Copyright (C) 2001, 2002 Swedish Institute of Computer Science and are licensed under the BSD License as described in COPYING file.
- wpa_supplicant Copyright (c) 2003-2005 Jouni Malinen and licensed under the BSD license.
- FreeBSD net80211 Copyright (c) 2004-2008 Sam Leffler, Errno Consulting and licensed under the BSD license.
- JSMN JSON Parser (components/jsmn) Copyright (c) 2010 Serge A. Zaitsev and licensed under the MIT license.
- argtable3 argument parsing library Copyright (C) 1998-2001,2003-2011,2013 Stewart Heitmann and licensed under 3-clause BSD license.
- linenoise line editing library Copyright (c) 2010-2014 Salvatore Sanfilippo, Copyright (c) 2010-2013 Pieter Noordhuis, licensed under 2-clause BSD license.
- libcoap COAP library Copyright (c) 2010-2017 Olaf Bergmann and others, is licensed under 2-clause BSD license as described in LICENSE file and COPYING file .
- libexpat XML parsing library Copyright (c) 1998-2000 Thai Open Source Software Center Ltd and Clark Cooper, Copyright (c) 2001-2017 Expat maintainers, is licensed under MIT license as described in COPYING file .
- FatFS library, Copyright (C) 2017 ChaN, is licensed under a BSD-style license .
- cJSON library, Copyright (c) 2009-2017 Dave Gamble and cJSON contributors, is licensed under MIT license as described in LICENSE file .
- libsodium library, Copyright (c) 2013-2018 Frank Denis, is licensed under ISC license as described in LICENSE file .

- micro-ecc library, Copyright (c) 2014 Kenneth MacKay, is licensed under 2-clause BSD license.
- nghttp2 library, Copyright (c) 2012, 2014, 2015, 2016 Tatsuhiro Tsujikawa, Copyright (c) 2012, 2014, 2015, 2016 nghttp2 contributors, is licensed under MIT license as described in COPYING file .
- Mbed TLS library, Copyright (C) 2006-2018 ARM Limited, is licensed under Apache License 2.0 as described in LICENSE file .
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- SD/MMC driver is derived from OpenBSD SD/MMC driver, Copyright (c) 2006 Uwe Stuehler, and is licensed under BSD license.
- Asio , Copyright (c) 2003-2018 Christopher M. Kohlhoff is licensed under the Boost Software License as described in COPYING file.
- ESP-MQTT MQTT Package (contiki-mqtt) - Copyright (c) 2014, Stephen Robinson, MQTT-ESP - Tuan PM <tuanpm at live dot com> is licensed under Apache License 2,0 as described in LICENSE file .
- BLE Mesh is adapted from Zephyr Project, Copyright (c) 2017-2018 Intel Corporation and licensed under Apache License 2.0
- mynewt-nimble Apache Mynewt NimBLE, Copyright 2015-2018, The Apache Software Foundation, is licensed under Apache License 2.0 as described in LICENSE file.
- cryptoauthlib Microchip CryptoAuthentication Library - Copyright (c) 2015 - 2018 Microchip Technology Inc, is licensed under common Microchip software License as described in LICENSE file
- freemodbus Copyright (c) 2006-2013 Christian Walter, Armink and licensed under the BSD license.

9.1.2 Build Tools

This is the list of licenses for tools included in this repository, which are used to build applications. The tools do not become part of the application (firmware), so their license does not affect licensing of the application.

- esptool.py is Copyright (C) 2014-2016 Fredrik Ahlberg, Angus Gratton and is licensed under the GNU General Public License v2, as described in LICENSE file.
- KConfig is Copyright (C) 2002 Roman Zippel and others, and is licensed under the GNU General Public License V2.
- Kconfiglib is Copyright (C) 2011-2019, Ulf Magnusson, and is licensed under the ISC License.
- Menuconfig of Kconfiglib is Copyright (C) 2018-2019, Nordic Semiconductor ASA and Ulf Magnusson, and is licensed under the ISC License.

9.1.3 Documentation

- HTML version of the ESP-IDF Programming Guide uses the Sphinx theme sphinx_idf_theme, which is Copyright (c) 2013-2020 Dave Snider, Read the Docs, Inc. & contributors, and Espressif Systems (Shanghai) CO., LTD. It is based on sphinx_rtd_theme. Both are licensed under MIT license.

9.2 ROM Source Code Copyrights

ESP32 mask ROM hardware includes binaries compiled from portions of the following third party software:

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9.5 TJpgDec License

TJpgDec - Tiny JPEG Decompressor R0.01 (C)ChaN, 2011 The TJpgDec is a generic JPEG decompressor module for tiny embedded systems. This is a free software that opened for education, research and commercial developments under license policy of following terms.

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Chapter 10

关于本指南

本指南为 乐鑫 公司 ESP32 系列芯片 官方应用开发框架 ESP-IDF 的配套文档。

ESP32 芯片是一款 2.4 GHz Wi-Fi 和蓝牙双模芯片，内置 1 或 2 个 32 位处理器，运算能力最高可达 600 DMIPS。

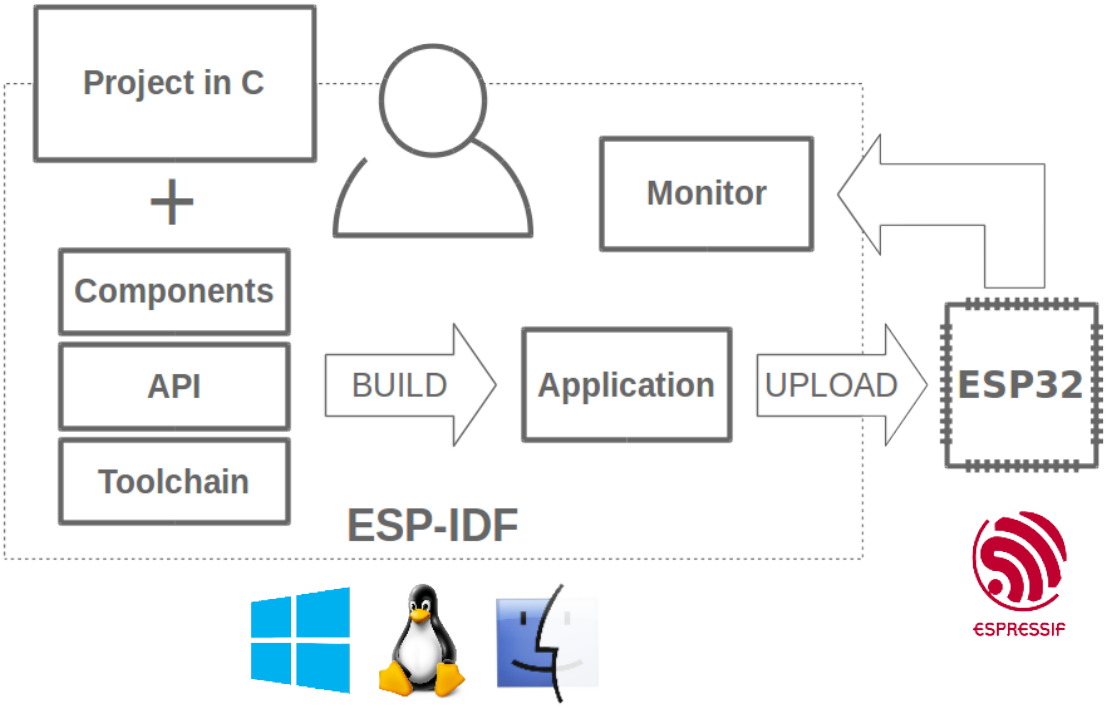

图 1: 乐鑫物联网综合开发框架

ESP-IDF 即乐鑫物联网开发框架，可为在 Windows、Linux 和 macOS 系统平台上开发 ESP32 应用程序提供工具链、API、组件和工作流的支持。

Chapter 11

Switch Between Languages/切换语言

The ESP-IDF Programming Manual is now available in two languages. Please refer to the English version if there is any discrepancy.

《ESP-IDF 编程手册》部分文档现在有两种语言的版本。如有出入请以英文版本为准。

- English/英文
- Chinese/中文

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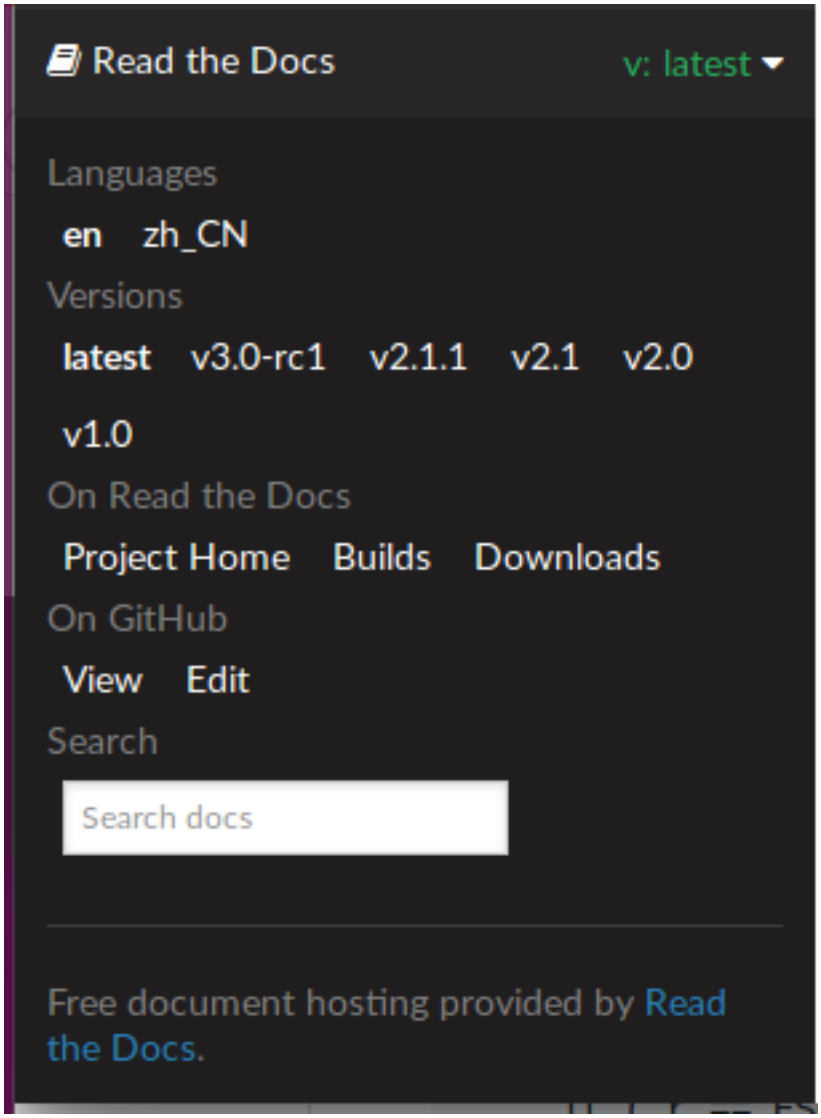
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